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Understanding Embedded - Microprocessors

Embedded microprocessors are specialized computing chips designed to perform specific tasks within an embedded system. Unlike general-purpose microprocessors found in personal computers, embedded microprocessors are tailored for dedicated functions within larger systems, offering optimized performance, efficiency, and reliability. These microprocessors are integral to the operation of countless electronic devices, providing the computational power necessary for controlling processes, handling data, and managing communications.

Applications of Embedded - Microprocessors

Embedded microprocessors are utilized across a broad spectrum of applications, making them indispensable in

Details	
Product Status	Obsolete
Core Processor	68000
Number of Cores/Bus Width	1 Core, 32-Bit
Speed	25MHz
Co-Processors/DSP	-
RAM Controllers	-
Graphics Acceleration	No
Display & Interface Controllers	-
Ethernet	-
SATA	-
USB	-
Voltage - I/O	5.0V
Operating Temperature	-55°C ~ 125°C (TC)
Security Features	-
Package / Case	114-BCPGA
Supplier Device Package	114-CPGA (34.54x34.54)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/ts68020mr25

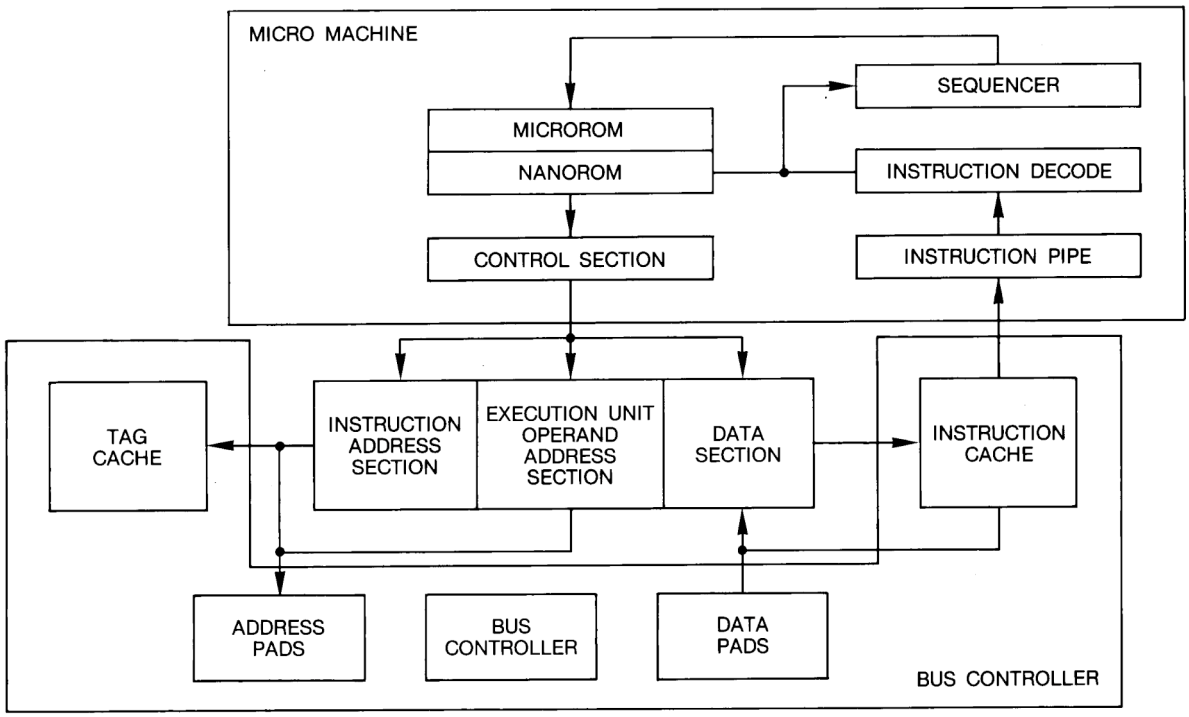
Introduction

The TS68020 is a high-performance 32-bit microprocessor. It is the first microprocessor to have evolved from a 16-bit machine to a full 32-bit machine that provides 32-bit address and data buses as well as 32-bit internal structures. Many techniques were utilized to improve performance and at the same time maintain compatibility with other processors of the TS68000 Family. Among the improvements are new addressing modes which better support high-level language structures, an expanded instruction set which provides 32-bit operations for the limited cases not supported by the TS68000 and several new instructions which support new data types. For special-purpose applications when a general-purpose processor alone is not adequate, a co-processor interface is provided.

The TS68020 is a high-performance microprocessor implemented in HCMOS, low power, small geometry process. This process allows CMOS and HMOS (high density NMOS) gates to be combined on the same device. CMOS structures are used where speed and low power is required, and HMOS structures are used where minimum silicon area is desired. This technology enables the TS68020 to be very fast while consuming less power (less than 1.5 watts) and still have a reasonably small die size. It utilizes about 190.000 transistors, 103.000 of which are actually implemented. The package is a pin-grid array (PGA) with 114 pins, arranged 13 pins on a side with a depopulated center and 132 pins ceramic quad flat pack.

Figure 1 is a block diagram of the TS68020. The processor can be divided into two main sections: the bus controller and the micromachine. This division reflects the autonomy with which the sections operate.

Figure 1. TS68020 Block Diagram



The bus controller consists of the address and data pads and multiplexers required to support dynamic bus sizing, a macro bus controller which schedules the bus cycles on the basis of priority with two state machines (one to control the bus cycles for operated accesses and the other to control the bus cycles for instruction accesses), and the instruction cache with its associated control.

The micromachine consists of an execution unit, nanorom and microrom storage, an instruction decoder, an instruction pipe, and associated control sections. The execution unit consists of an address section, an operand address section, and a data section. Microcode control is provided by a modified two-level store of microrom and nanorom. Programmed logical arrays (PLAs) are used to provide instruction decode and sequencing information. The instruction pipe and other individual control sections provide the secondary decode of instructions and generated the actual control signals that result in the decoding and interpretation of nanorom and micorom information.

Figure 2. PGA Terminal Designation

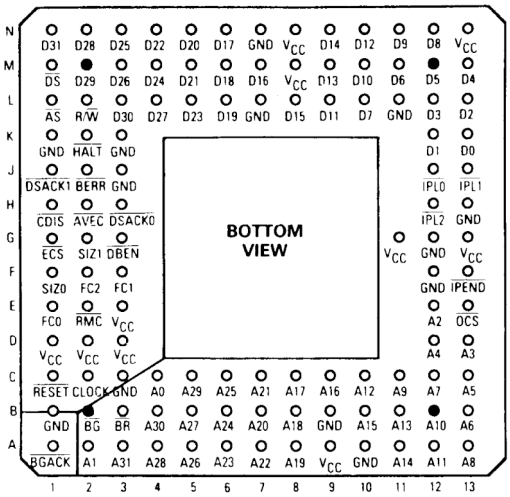


Figure 3. CQFP Terminal Designation

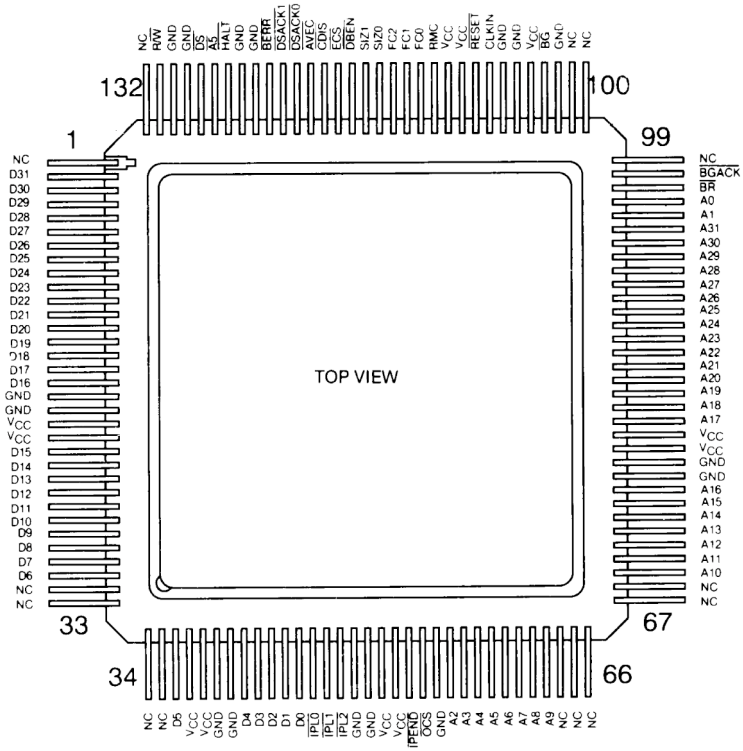


Table 1. Signal Index

Signal Name	Mnemonic	Function
Address Bus	A0-A31	32-bit Address Bus Used to address any of 4, 294, 967, 296 bytes.
Data Bus	D0-D31	32-bit Data Bus Used to Transfer 8, 16, 24 or 32 bits of Data Per Bus Cycle.
Function Codes	FC0-FC2	3-bit Function Case Used to Identify the Address Space of Each Bus Cycle.
Size	SIZ0/SIZ1	Indicates the Number of Bytes Remaining to be Transferred for this Cycle. These Signals, Together with A0 And A1, Define the Active Sections of the Data Bus.
Read-Modify-Write Cycle	$\overline{RMC}$	Provides an Indicator that the Current Bus Cycle is Part of an Indivisibleread-modify-write Operation.
External Cycle Start	$\overline{ECS}$	Provides an Indication that a Bus Cycle is Beginning.
Operand Cycle Start	$\overline{OCS}$	Identical Operation to that of ECS Except that OCS Is Asserted Only During the First Bus Cycle of an Operand Transfer.
Address Strobe	$\overline{AS}$	Indicates that a Valid Address is on The Bus.
Data Strobe	$\overline{DS}$	Indicates that Valid Data is to be Placed on the Data Bus by an External Device or has been Laced on the Data Bus by the TS68020.
Read/Write	R/$\overline{W}$	Defines the Bus Transfer as an MPU Read or Write.
Data Buffer Enable	$\overline{DBEN}$	Provides an Enable Signal for External Data Buffers.
Data Transfer and Size Acknowledge	$\overline{DSACK0}$/$\overline{DSACK1}$	Bus Response Signals that Indicate the Requested Data Transfer Operation is Completed. In Addition, these Two Lines Indicate the Size of the External Bus Port on a Cycle-by-cycle Basis.
Cache Disable	$\overline{CDIS}$	Dynamically Disables the On-chip Cache to Assist Emulator Support.
Interrupt Priority Level	$\overline{IPL0}$-$\overline{IPL2}$	Provides an Encoded Interrupt Level to the Processor.
Autovector	$\overline{AVEC}$	Requests an Autovector During an Interrupt Acknowledge Cycle.
Interrupt Pending	$\overline{IPEND}$	Indicates that an Interrupt is Pending.
Bus Request	$\overline{BR}$	Indicates that an External Device Requires Bus Mastership.
Bus Grant	$\overline{BG}$	Indicates that an External Device may Assume Bus Mastership.
Bus Grant Acknowledge	$\overline{BGACK}$	Indicates that an External Device has Assumed Bus Mastership.
Reset	$\overline{RESET}$	System Reset.
Halt	$\overline{HALT}$	Indicates that the Processor Should Suspend Bus Activity.
Bus Error	$\overline{BERR}$	Indicates an Invalid or Illegal Bus Operation is Being Attempted.
Clock	CLK	Clock Input to the Processor.
Power Supply	V_{CC}	+5-volt ± 10% Power Supply.
Ground	GND	Ground Connection.

Electrical Characteristics

Table 2. Absolute Maximum Ratings

Symbol	Parameter	Test Conditions	Min	Max	Unit
V_{CC}	Supply Voltage		-0.3	+7.0	V
V_I	Input Voltage		-0.5	+7.0	V
P_{dmax}	Max Power Dissipation	T_{case} = -55°C		2.0	W
		T_{case} = +125°C		1.9	W
T_{case}	Operating Temperature	M Suffix	-55	+125	°C
		V Suffix	-40	+85	°C
T_{stg}	Storage Temperature		-55	+150	°C
T_{leads}	Lead Temperature	Max 5 Sec. Soldering		+270	°C

Table 3. Recommended Condition of Use

Unless otherwise stated, all voltages are referenced to the reference terminal (see Table 1).

Symbol	Parameter		Min	Max	Unit
V_{CC}	Supply Voltage		4.5	5.5	V
V_{IL}	Low Level Input Voltage		-0.3	0.5	V
V_{IH}	High Level Input Voltage		2.4	5.25	V
T_{case}	Operating Temperature		-55	+125	°C
R_L	Value of Output Load Resistance		(1)		Ω
C_L	Output Loading Capacitance			(1)	pF
$t_r(c)$–$t_f(c)$	Clock Rise Time (See Figure 5)	68020-16		5	ns
		68020-20		5	
		68020-25		4	
f_c	Clock Frequency (See Figure 5)	68020-16	8	16.67	MHz
		68020-20	12.5	20	
		68020-25	12.5	25	
t_{cyc}	Cycle Time (see Figure 5)	68020-16	60	125	ns
		68020-20	50	80	
		68020-25	40	80	
$t_W(CL)$	Clock Pulse Width Low (See Figure 5)	68020-16	24	95	ns
		68020-20	20	54	
		68020-25	19	61	
$t_W(CH)$	Clock Pulse Width High (See Figure 5)	68020-16	24	95	ns
		68020-20	20	50	
		68020-25	19	61	

Note: 1. Load network number 1 to 4 as specified (Table 7) gives the maximum loading of the relevant output.

The total thermal resistance of a package (θ_{JA}) can be separated into two components, θ_{JC} and θ_{CA}, representing the barrier to heat flow from the semiconductor junction to the package (case), surface (θ_{JC}) and from the case to the outside ambient (θ_{CA}). These terms are related by the equation:

$$\theta_{JA} = \theta_{JC} = \theta_{CA} \quad (4)$$

θ_{JC} is device related and cannot be influenced by the user. However, θ_{CA} is user dependent and can be minimized by such thermal management techniques as heat sinks, ambient air cooling and thermal convection. Thus, good thermal management on the part of the user can significantly reduce θ_{CA} so that θ_{JA} approximately equals θ_{JC}. Substitution of θ_{JC} for θ_{JA} in equation **(1)** will result in a lower semiconductor junction temperature.

Mechanical and Environment

The microcircuits shall meet all mechanical environmental requirements of MIL-STD-883 for class B devices.

Marking

The document where are defined the marking are identified in the related reference documents. Each microcircuit are legible and permanently marked with the following information as minimum:

- ATMEL Logo
- Manufacturer's Part Number
- Class B Identification
- Date-code of Inspection Lot
- ESD Identifier if Available
- Country of Manufacturing

Quality Conformance Inspection

DESC/MIL-STD-883

Is in accordance with MIL-M-38510 and method 5005 of MIL-STD-883. Group A and B inspections are performed on each production lot. Group C and D inspections are performed on a periodical basis.

Electrical Characteristics

General Requirements

All static and dynamic electrical characteristics specified and the relevant measurement conditions are given below.

(last issue on request to our marketing services).

Table 5: Static electrical characteristics for all electrical variants.

Table 6: Dynamic electrical characteristics for 68020-16 (16.67 MHz), 68020-20 (20 MHz) and 68020-25 (25 MHz).

For static characteristics, test methods refer to "Test Conditions Specific to the Device" on page 14 hereafter of this specification (Table 7).

For dynamic characteristics (Table 6), test methods refer to IEC 748-2 method, where existing.

Indication of “min.” or “max.” in the column “test temperature” means minimum or maximum operating temperature.

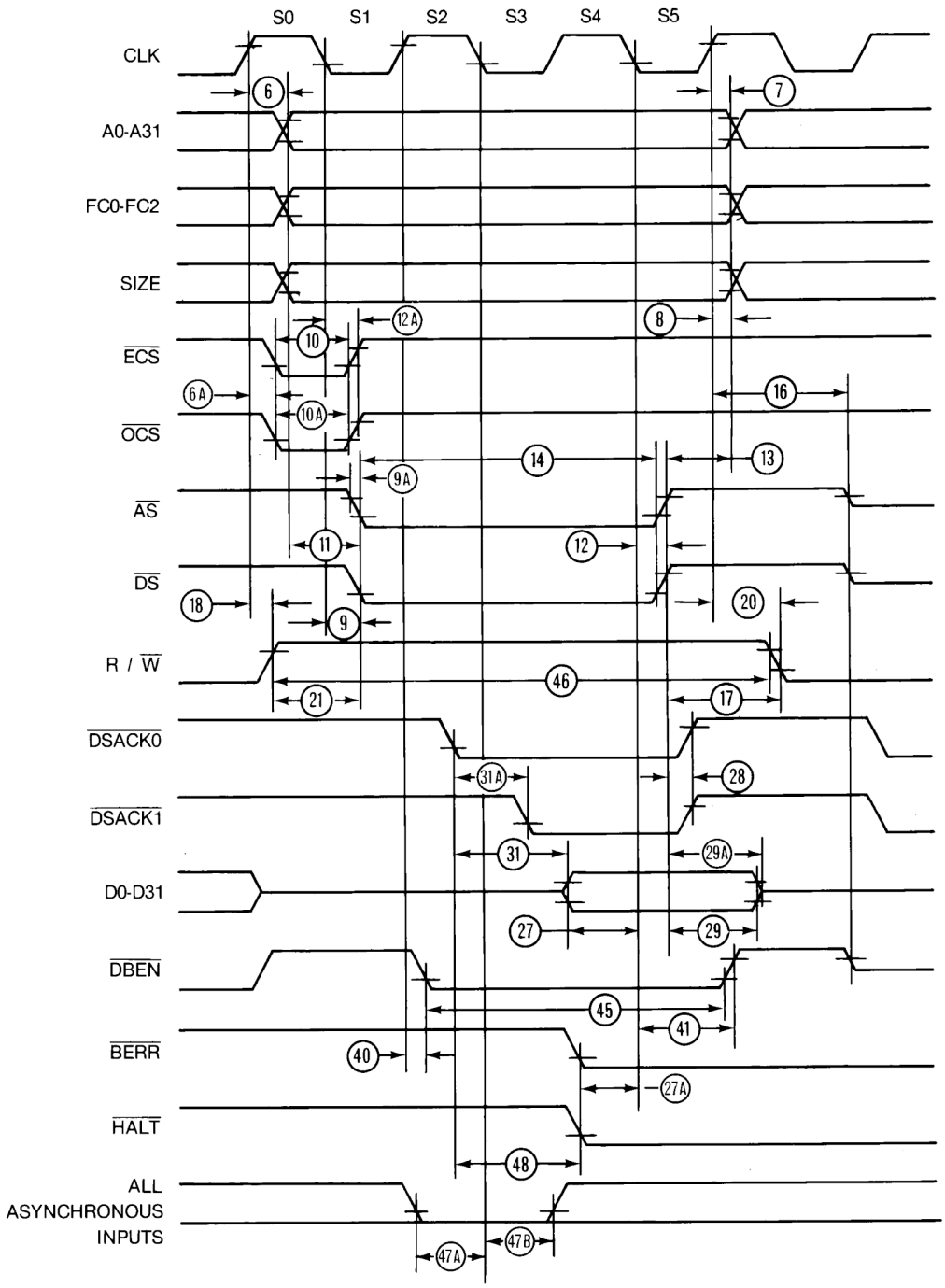
Table 5. Static Characteristics. V_{CC} = 5.0V_{DC} ± 10%; GND = 0V_{DC}; T_c = -55/+125°C or -40/+85°C (Figure 4 to Figure 8)

Symbol	Parameter	Condition	Min	Max	Units
I_{CC}	Maximum Supply Current	V_{CC} = 5.5V T_{case} -55°C to +25°C		333	mA
I_{CC}	Maximum Supply Current	V_{CC} = 5.5V T_{case} = 125°C		207	mA
V_{IH}	High Level Input Voltage	V_O = 0.5V or 2.5 V_{CC} = 4.5V to 5.5V	2.0	V_{CC}	V
V_{IL}	Low Level Input Voltage	V_O = 0.5V or 2.4V V_{CC} = 4.5V to 5.5V	-0.5	0.8	V
V_{OH}	High Level Output Voltage All Outputs	I_{OH} = 400 µA	2.4		V
V_{OL}	Low Level Output Voltage Outputs A0-A31, FC0-FC2, D0-D31, SIZ0-SIZ1, $\overline{BG}$	I_{OL} = 3.2 mA Load Circuit as Figure 8 R = 1.22 kΩ		0.5	V
V_{OL}	Low Level Output Voltage Outputs $\overline{AS}$, $\overline{DS}$, $\overline{RMC}$, R/$\overline{W}$, $\overline{DBEN}$, $\overline{IPEND}$	I_{OL} = 5.3 mA Load Circuit as Figure 8 R = 740Ω		0.5	V
V_{OL}	Low Level Output Voltage Outputs $\overline{ECS}$, $\overline{OCS}$	I_{OL} = 2.0 mA Load Circuit as Figure 8 R = 2 kΩ		0.5	V
V_{OL}	Low Level Output Voltage Outputs $\overline{HALT}$, $\overline{RESET}$	I_{OL} = 10.7 mA Load Circuit as Figure 6 and Figure 7		0.5	V
$\lvert I_{IN} \rvert$	Input Leakage Current (High and Low State)	-0.5V ≤ V_{IN} ≤ V_{CC} (Max)		2.5	µA
$\lvert I_{OHZ} \rvert$	High level leakage current at three-state outputs Outputs A0-A31, $\overline{AS}$, $\overline{DBEN}$, $\overline{DS}$, D0-D31, R/$\overline{W}$, FC0-FC2, $\overline{RMC}$, SIZ0-SIZ1	V_{OH} = 2.4V		2.5	µA
$\lvert I_{OLZ} \rvert$	Low Level Leakage Current at Three-state Outputs Outputs A0-A31, $\overline{AS}$, $\overline{DBEN}$, $\overline{DS}$, D0-D31 R/$\overline{W}$, FC0-FC2, $\overline{RMC}$, SIZ0-SIZ1	V_{OL} = 0.5V		2.5	µA
I_{OS}	Output Short-circuit Current (Any Output)	V_{CC} = 5.5V V_O = 0V (Pulsed. Duration 1 ms Duty Cycle 10:1)		200	mA

Time Definitions

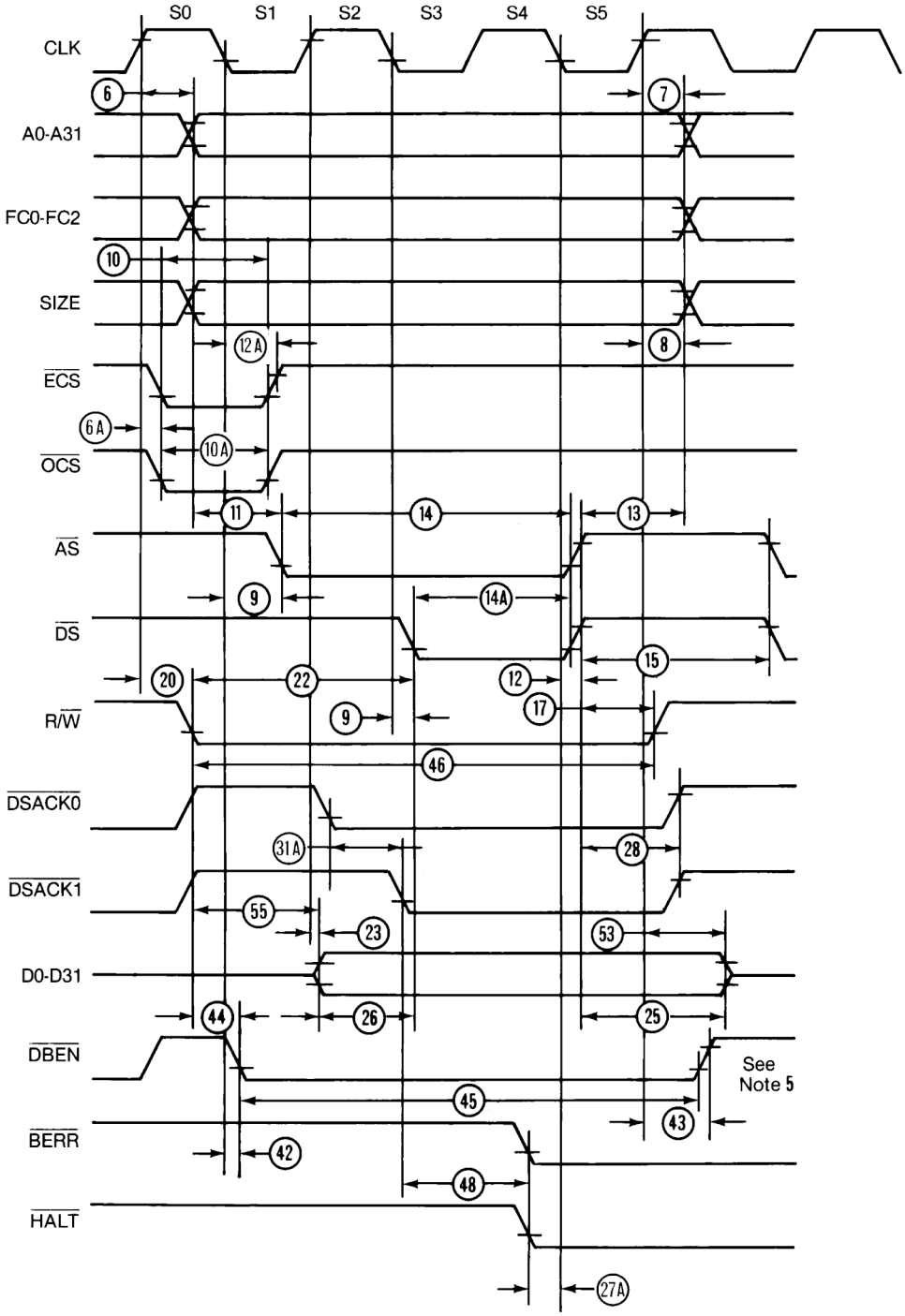
The times specified in Table 6 as dynamic characteristics are defined in Figure 9 below, by a reference number given the column “interval N°” of the tables together with the relevant figure number.

Figure 9. Read Cycle Timing Diagram



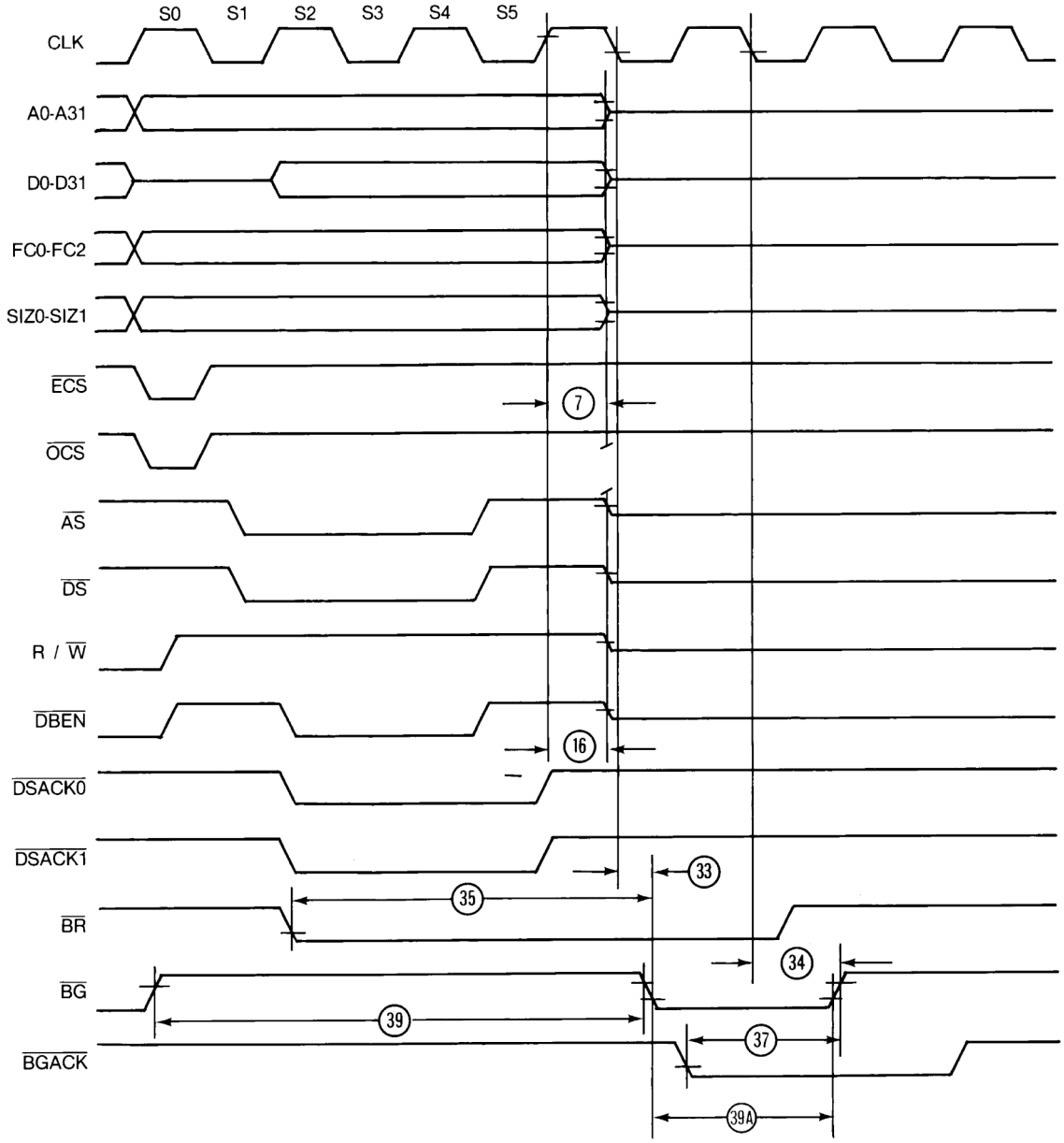
Note: Timing measurements are referenced to and from a low voltage of 0.8V and a high voltage of 2.0V, unless otherwise noted. The voltage swing through this range should start outside and pass through the range such that the rise or fall will be linear between 0.8V and 2.0V.

Figure 10. Write Cycle Timing Diagram (Continued)



Note: Timing measurements are referenced to and from a low voltage of 0.8V and a high voltage of 2.0V, unless otherwise noted. The voltage swing thorough this range should start outside and pass through the range such that the rise or fall will be linear between 0.8V and 2.0V.

Figure 11. Bus Arbitration Timing Diagram



Note: Timing measurements are referenced to and from a low voltage of 0.8V and a high voltage of 2.0V, unless otherwise noted. The voltage swing thorough this range should start outside and pass through the range such that the rise or fall will be linear between 0.8V and 2.0V.

Instruction Set Overview

The TS68020 instruction set is shown in Table 9. Special emphasis has been given to the instruction set's support of structured high-level languages and sophisticated operating systems. Each instruction, with few exceptions, operates on bytes, words, and long words and most instructions can use any of the 18 addressing modes. Many instruction extensions have been made on the TS68020 to take advantage of the full 32-bit operation where, on the earlier 68000 Family members, only 8 and 16 bits values were used. The TS68020 is upward source- and object-level code compatible with the family because it supports all of the instructions that previous family members offer. Additional instructions are now provided by the TS68020 in support of its advanced features.

Table 9. Instruction Set

Mnemonic	Description
ABCD	Add Decimal with Extend
ADD	Add
ADDA	Add Address
ADDI	Add Immediate
ADDQ	Add Quick
ADDX	Add with Extend
AND	Logical AND
ANDI	Logical AND Immediate
ASL, ASR	Arithmetic Shift Left and Right
Bcc	Branch Conditionally
BCHG	Test Bit and Change
BCLR	Test Bit and Clear
BFCHG	Test Bit Field and Change
BFCLR	Test Bit Field and Clear
BFEXTS	Signed Bit Field Extract
BFEXTU	Unsigned Bit Field Extract
BFFFO	Bit Field Find First One
BFINS	Bit Field Insert
BFSET	Test Bit Field and Set
BFTST	Test Bit Field
BKPT	Breakpoint
BRA	Branch
BSET	Test Bit and Set
BSR	Branch to Subroutine
BTST	Test Bit

Table 9. Instruction Set (Continued)

Mnemonic	Description
CALLM	Call Module
CAS	Compare and Swap Operands
CAS2	Compare and Swap Dual Operands
CHK	Check Register Against Bound
CHK2	Check Register Against Upper and Lower Bounds
CLR	Clear
CMP	Compare
CMPA	Compare Address
CMPI	Compare Immediate
CMPM	Compare Memory to Memory
CMP2	Compare Register Against Upper and Lower Bounds
DBcc	Test Condition, Decrement and Branch
DIVS, DIVSL	Signed Divide
DIVU, DIVUL	Unsigned Divide
EOR	Logical Exclusive OR
EORI	Logical Exclusive OR Immediate
EXG	Exchange Registers
EXT, EXTB	Sign Extend
ILLEGAL	Take Illegal Instruction Tape
JMP	Jump
JSR	Jump to Subroutine
LEA	Load Effective Address
LINK	Link and Allocate
LSL, LSR	Logical Shift Left and Right
MOVE	Move
MOVEA	Move Address
MOVE CCR	Move Condition Code Register
MOVE SR	Move Status Register
MOVE USP	Move User Stack Pointer
MOVEC	Move Control Register
MOVEM	Move Multiple Registers
MOVEP	Move Peripheral
MOVEQ	Move Quick
MOVES	Move Alternate Address Space
MULS	Signed Multiply
MULU	Unsigned Multiply
NBCD	Negate Decimal with Extend
NEG	Negate
NEGX	Negate with Extend
NOP	No Operation
NOT	Logical Complement

Table 9. Instruction Set (Continued)

Mnemonic	Description
OR	Logical Inclusive OR
ORI	Logical Inclusive OR Immediate
PACK	Pack BCD
PEA	Push Effective Address
RESET	Reset External Devices
ROL, ROR	Rotate Left and Right
ROXL, ROXR	Rotate with Extend Left and Right
RTD	Return and Deallocate
RTE	Return and Exception
RTM	Return from Module
RTR	Return and Restore Codes
RTS	Return from Subroutine
SBCD	Subtract Decimal with Extend
Scc	Set Conditionally
STOP	Stop
SUB	Subtract
SUBA	Subtract Address
SUBI	Subtract Immediate
SUBQ	Subtract Quick
SUBX	Subtract with Extend
SWAP	Swap Register Words
TAS	Test Operand and Set
TRAP	Trap
TRAPcc	Trap Conditionally
TRAPV	Trap on Overflow
TST	Test Operand
UNLK	Unlink
UNPK	Unpack BCD
Co-processor Instructions	
cpBCC	Branch Conditionally
cpDBcc	Test Co-processor Condition, Decrement and Branch
cpGEN	Co-processor General Instruction
cpRESTORE	Restore Internal State of Co-processor
cpSAVE	Save Internal State of Co-processor
cpScc	Set Conditionally
cpTRAPcc	Trap Conditionally

Multi-processing

To further support multi-processing with the TS68020, a compare and swap instruction, CAS, has been added. This instruction makes use of the read-modify-write cycle to compare two operands and swap a third operand pending the results of the compare. A variant of this instruction, CAS2, performs similarly comparing dual operand pairs, and updating two operands.

These multi-processing operations are useful when using common memory to share or pass data between multiple processing elements. The read-modify-write cycle is an indivisible operand that allows reading and updating a “lock” operand used to control access to the common memory elements. The CAS2 instruction is more powerful since dual operands allow the “lock” to the checked and two values (i.e., both pointers in a doubly-linked list) to be updated according to the lock’s status, all in a single operation.

Module Support

The TS68020 includes support for modules with the call module (CALLM) and return from module (RTM) instructions. The CALLM instruction references a module descriptor. This descriptor contains control information for entry into the associated module. The CALLM instruction creates a module stack frame and stores the module state in that frame. The RTM instruction recovers the previous module state from the stack frame and returns to the calling module.

The module interface also provides a mechanism for finer resolution of access control by external hardware. Although the TS68020 does not interrupt the access control information, it does communicate with external hardware when the access control is to be changed, and relies on the external hardware to verify that the changes are legal.

CALLM and RTM, when used as subroutine calls and returns with proper descriptor formats, cause the TS68020 to perform the necessary actions to verify legitimate access to modules.

Virtual Memory/Machine Concepts

The full addressing range of the TS68020 is 4-Gbyte (4, 294, 967, 296). However, most TS68020 systems implement a smaller physical memory. Nonetheless, by using virtual memory techniques, the system can be made to appear to have a full 4-Gbyte of physical memory available to each user program. These techniques have been used for many years in large mainframe computers and minicomputers. With the TS68020 (as with the TS68010 and TS68012), virtual memory can be fully supported in microprocessor-based systems.

In a virtual memory system, a user program can be written as though it has a large amount of memory available to it when actually only a smaller amount of memory is physically present in the system. In a similar fashion, a system provides user programs access to other devices that are not physically present in the system, such as tape drives, disk drives, printers, or terminals. With proper software emulation, a physical system can be made to appear to a user program as any other 68000 computer system and the program may be given full access to all of the resources of that emulated system. Such an emulator system is called a virtual machine.

Virtual Memory

The basic mechanism for supporting virtual memory is to provides a limited amount of high-speed physical memory that can be accessed directly by the processor while maintaining of a much larger “virtual” memory on secondary storage devices such as large capacity disk drives. When the processor attempts to access a location in the virtual memory map that is not resident in the physical memory (referred to as a page fault), the access to that location is temporarily suspended while the necessary data is fetched from secondary storage and placed in physical memory; the suspended access is then either restarted or continued.

The TS68020 will always transfer the maximum amount of data on all bus cycles; i.e., it always assumes the port is 32-bit wide when beginning the bus cycle. In addition, the TS68020 has no restrictions concerning alignment of operands in memory; long word operands need not be aligned on long word address boundaries. When misaligned data requires multiple bus cycles, the TS68020 aligned data requires multiple bus cycles, the TS68020 automatically runs the minimum number of bus cycles.

The Co-processor Concept

The co-processor interface is a mechanism for extending the instruction set of the TS68000 Family. Examples of these extensions are the addition of specialized data operands for the existing data types or, for the case of the floating point, the inclusion of new data types and operations for them as implemented by the TS68881 and TS68882 floating point co-processors.

The programmer's model for the TS68000 Family of microprocessors is based on sequential, non-concurrent instruction execution. This means each instruction is completely executed prior to the beginning of the next instruction. Hence, instructions do not operate concurrently in the programmer's model. Most microprocessors implement the sequential model which greatly simplifies the programmer responsibilities since sequencing control is automatic and discrete.

The TS68000 co-processor interface is designed to extend the programmer's model and it provides full support for the sequential, non-concurrent instruction execution model. Hence, instruction execution by the co-processor is assumed to not overlap with instruction execution with the main microprocessor. Yet, the TS68000 co-processor interface does allow concurrent operation when concurrency can be properly accommodated. For example, the TS68881 or TS68882 floating-point co-processor will allow the TS68020 to proceed executing instruction while the co-processor continues a floating-point operation, up to the point that the TS68020 sends another request to the co-processor. Adhering to the sequential execution model, the request to the co-processor continues a floating-point operation, up to the co-processor completes each TS68881 and TS68882 instruction before it starts the next, and the TS68020 is allowed to proceed as it can in a concurrent fashion.

co-processors are divided into two types by their bus utilization characteristics. A co-processor is a DMA co-processor if it can control the bus independent of the main processor. A co-processor is a non-DMA co-processor if it does not have the capability of controlling the bus. Both co-processor types utilize the same protocol and main processor resources. Implementation of a co-processor as a DMA or non-DMA type is based primarily on bus bandwidth of the co-processor, performance, and cost issues.

The communication protocol between the main processor and the co-processor necessary to execute a co-processor instruction is based on a group of co-processor interface registers (Table 10) which are defined for the TS68000 Family co-processor interface. The TS68020 hardware uses standard TS68000 asynchronous bus cycles to access the registers. Thus, the co-processor doesn't require a special bus hardware; the bus interface implemented by a co-processor for its interface register set must only satisfy the TS68020 address, data, and control signal timing to guarantee proper communication with the main processor. The TS68020 implements the communication protocol with all co-processors in hardware (and microcode) and handles all operations automatically so the programmer is only concerned with the instructions and data types provided by the co-processor as extensions to the TS68020 instruction set and data types.

Table 11. Co-processor Primitives (Continued)

General Operand Transfer
Evaluate and Pass (Ea.)
Evaluate (Ea.) and Transfer Data
Write to Previously Evaluated (Ea.)
Take Address and Transfer Data
Transfer to/from Top of Stack
Register Transfer
Transfer CPU Register
Transfer CPU Control Register
Transfer Multiple CPU Registers
Transfer Multiple Co-processor Registers
Transfer CPU SR and/or ScanPC

Up to eight processors are supported in a single system with a system-unique co-processor identifier encoded in the co-processor instruction. When accessing a co-processor, the TS68020 executes standard read and write bus cycle in CPU address space, as encoded by the function codes, and places the co-processor identifier on the address bus to be used by chip-select logic to select the particular co-processor. Since standard bus cycle are used to access the co-processor, the co-processor may be located according to system design requirements, whether it be located on the micro-processor local bus, on another board on the system bus, or any other place where the chip-select and co-processor protocol using standard TS68000 bus cycles can be supported.

Co-processor Protocol

Interprocessor transfers are all initiated by the main processor during co-processor instruction execution. During the processing of a co-processor instruction, the main processor transfers instruction information and data to the associated co-processor, and receives data, requests, and status information from the co-processor. These transfers are all based on the TS68000 bus cycles.

The typical co-processor protocol which the main processor follows is:

a) The main processor initiates the communications by writing command information to a location in the co-processor interface.

b) The main processor reads the co-processor response to that information.

1) The response may indicate that the co-processor is busy, and the main processor should again query the co-processor. This allows the main processor and co-processor to synchronize their concurrent operations.

2) The response may indicate some exception condition; the main processor acknowledges the exception and begins exception processing.

3) The response may indicate that the co-processor needs the main processor to perform some service such as transferring data to or from the co-processor. The co-processor may also request that the main processor query the co-processor again after the service is complete.

4) The response may indicate that the main processor is not needed for further processing of the instruction. The communication is terminated, and the main processor is free to begin execution of the next instruction. At this point in the co-processor protocol, as the main processor continues to execute the instruction stream, the main processor may operate concurrently with the co-processor.

When the main processor encounters the next co-processor instruction, the main processor queries the co-processor until the co-processor is ready; meanwhile, the main processor can go on to service interrupts and do a context switch to execute other tasks, for example.

Each co-processor instruction type has specific requirements based on this simplified protocol. The co-processor interface may use as many extension words as requires to implement a co-processor instruction.

Primitives/Response

The response register is the means by which the co-processor communicates service requests to the main processor. The content of the co-processor response register is a primitive instruction to the main processor which is read during co-processor communication by the main processor. The main processor “executes” this primitive, thereby providing the services requires by the co-processor. Table 11 summarizes the co-processor primitives that the TS68020 accepts.

Exceptions

Kinds of Exceptions

Exception can be generated by either internal or external causes. The externally generated exceptions are the interrupts, the bus error, and reset requests. The interrupts are requests from peripheral devices for processor action while the bus error and reset pins are used for access control and processor restart. The internally generated exceptions come from instructions, address errors, tracing, or breakpoints. The TRAP, TRAPcc, TRAPV, cpTRAPcc, CHK, CHK2, and DIV instructions can all generate exceptions as part of their execution. Tracing behaves like a very high priority, internally generated interrupt whenever it is processed. The other internally generated exceptions are caused by illegal instructions, instruction fetches from odd addresses, and privilege violations.

Exception Processing Sequence

Exception processing occurs in four steps. During the first step, an internal copy is made of the status register. After the copy is made, the special processor state bits in the status register are changed. The S bit is set, putting the processor into supervisor privilege state. Also, the T1 and T0 bits are negated, allowing the exception handler to execute unhindered by tracing. For the reset and interrupt exceptions, the interrupt priority mask is also updated.

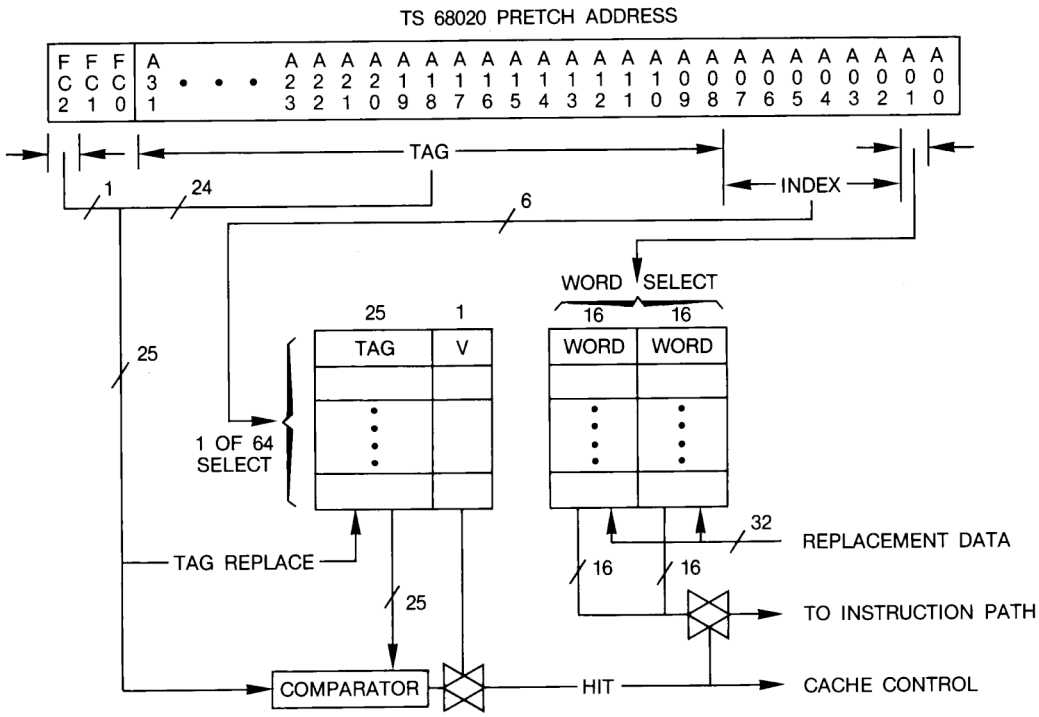
In the second step, the vector number of the exception is determined. For interrupts, the vector number is obtained by a processor read that is classified as an interrupt acknowledge cycle. For co-processor detected exceptions, the victor number is included in the co-processor exception primitive response. For all other exceptions, internal logic provides the vector number. This vector number is then used to generate the address of the exception vector.

The third step is to save the current processor status. The exception stack frame is created and filled on the supervisor stack. In order to minimize the amount of machine state that is saved, various stack frame sizes are used to contain the processor state depending on the type of exception and where it occurred during instruction execution. If the exception is an interrupt and the M bit is on, the M bit is forced off, and a short four word exception stack frame is saved on the master stack which indicates that the exception is saved on the interrupt stack. If the exception is a reset, the M bit is simply forced off, and the reset vector is accessed.

Second, and probably the most important benefit of the cache, is that it allows instruction stream fetches and operand accesses to proceed in parallel. For example, if the TS68020 requires both an instruction stream access and an operand access, and the instruction is resident in the cache, the operand access will proceed unimpeded rather than being queued behind the instruction fetch. Similarly, the TS68020 is fully capable of executing several internal instructions (instructions that do not require the bus) while completing an operand access for another instruction.

The TS68020 instruction cache is a 256-byte direct mapped cache organized as 64 long word entries. Each cache entry consists of a tag field made up of the upper 24 address bits, the FC2 (user/supervisor) value, one valid bit, and 32-bit of instruction data (Figure 22).

Figure 22. TS68020 On-chip Cache Organization



The TS68020 employs a 32-bit data bus and fetches instructions on long word address boundaries. Hence, each 32-bit instruction fetch brings in two 16-bit instruction words which are then written into the on-chip cache. When the cache is enabled, the subsequent prefetch will find the next 16-bit instruction word is already present in the cache and the related bus cycle is saved. If the cache were not enabled, the subsequent prefetch will find the bus controller still holds the full 32-bit and can satisfy the prefetch and again save the related bus cycle. So, even when the on-chip instruction cache is not enabled, the bus controller provides an instruction “cache hit” rate up to 50%.

Preparation for Delivery

Certificate of Compliance

Atmel offers a certificate of compliance with each shipment of parts, affirming the products are in compliance with MIL-STD-883 and guaranteeing the parameters are tested at extreme temperatures for the entire temperature range.

Handling

MOS devices must be handled with certain precautions to avoid damage due to accumulation of static charge. Input protection devices have been designed in the chip to minimize the effect of this static buildup. However, the following handling practices are recommended:

a) Device should be handled on benches with conductive and grounded surface.

b) Ground test equipment, tools and operator

c) Do not handle devices by the leads.

d) Store devices in conductive foam or carriers.

e) Avoid use of plastic, rubber, or silk in MOS areas.

f) Maintain relative humidity above 50%, if practical.

Mass

PGA 114 - 6 grams typically
CQFP 132 - 14 grams typically

Terminal Connections

114-lead - Ceramic Pin Grid Array

See Figure 2.

132-lead - Ceramic Quad Flat Pack

See Figure 3.

Ordering Information

Hi-REL Product

Commercial Atmel Part-Number	Norms	Package	Temperature Range T_c (°C)	Frequency (MHz)	Drawing Number
TS68020MRB/C16	MIL-STD-883	PGA 114	-55/+125	16.67	-
TS68020MR1B/C16	MIL-STD-883	PGA 114/tin	-55/+125	16.67	-
TS68020MRB/C20	MIL-STD-883	PGA 114	-55/+125	20	-
TS68020MR1B/C20	MIL-STD-883	PGA 114/tin	-55/+125	20	-
TS68020MRB/C25	MIL-STD-883	PGA 114	-55/+125	25	-
TS68020MR1B/C25	MIL-STD-883	PGA 114/tin	-55/+125	25	-
TS68020MFB/C16	MIL-STD-883	CQFP 132	-55/+125	16.67	-
TS68020MF1B/C16	MIL-STD-883	CQFP 132/tin	-55/+125	16.67	-
TS68020MFB/C20	MIL-STD-883	CQFP 132	-55/+125	20	-
TS68020MF1B/C20	MIL-STD-883	CQFP 132/tin	-55/+125	20	-
TS68020MFB/C25	MIL-STD-883	CQFP 132	-55/+125	25	-
TS68020MF1B/C25	MIL-STD-883	CQFP 132/tin	-55/+125	25	-
TS68020DESC02XA	DESC	PGA 114/tin	-55/+125	16.67	5962-8603202XA
TS68020DESC03XA	DESC	PGA 114/tin	-55/+125	20	5962-8603203XA
TS68020DESC04XA	DESC	PGA 114/tin	-55/+125	25	5962-8603204XA
TS68020DESC02XC	DESC	PGA 114	-55/+125	16.67	5962-8603202XC
TS68020DESC03XC	DESC	PGA 114	-55/+125	20	5962-8603203XC
TS68020DESC04XC	DESC	PGA 114	-55/+125	25	5962-8603204XC
TS68020DESC02YA	DESC	CQFP 132/tin	-55/+125	16.67	5962-8603202YA
TS68020DESC03YA	DESC	CQFP 132/tin	-55/+125	20	5962-8603203YA
TS68020DESC04YA	DESC	CQFP 132/tin	-55/+125	25	5962-8603204YA
TS68020DESC02YC	DESC	CQFP 132	-55/+125	16.67	5962-8603202YC
TS68020DESC03YC	DESC	CQFP 132	-55/+125	20	5962-8603203YC
TS68020DESC04YC	DESC	CQFP 132	-55/+125	25	5962-8603204YC

Standard Product

Commercial Atmel Part-Number	Norms	Package	Temperature Range T_c (°C)	Frequency (MHz)	Drawing Number
TS68020VR16	Internal Standard	PGA 114	-40/+85	16.67	Internal
TS68020VR20	Internal Standard	PGA 114	-40/+85	20	Internal
TS68020VR25	Internal Standard	PGA 114	-40/+85	25	Internal
TS68020MR16	Internal Standard	PGA 114	-55/+125	16.67	Internal
TS68020MR20	Internal Standard	PGA 114	-55/+125	20	Internal
TS68020MR25	Internal Standard	PGA 114	-55/+125	25	Internal