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### What is "[Embedded - Microcontrollers](#)"?

"[Embedded - Microcontrollers](#)" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

### Applications of "[Embedded - Microcontrollers](#)"

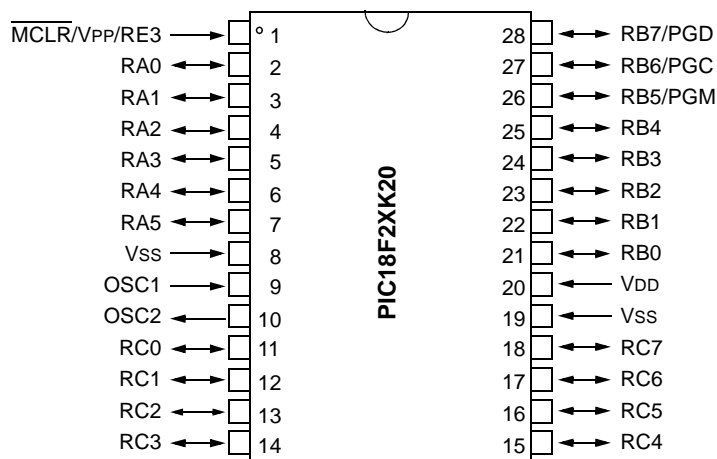
#### Details

Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	64MHz
Connectivity	I <sup>2</sup> C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, HLVD, POR, PWM, WDT
Number of I/O	24
Program Memory Size	8KB (4K x 16)
Program Memory Type	FLASH
EEPROM Size	256 x 8
RAM Size	512 x 8
Voltage - Supply (Vcc/Vdd)	1.8V ~ 3.6V
Data Converters	A/D 11x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	28-SSOP (0.209", 5.30mm Width)
Supplier Device Package	28-SSOP
Purchase URL	<a href="https://www.e-xfl.com/product-detail/microchip-technology/pic18f23k20-i-ss">https://www.e-xfl.com/product-detail/microchip-technology/pic18f23k20-i-ss</a>

# PIC18F2XK20/4XK20

**FIGURE 2-1: 28-PIN SDIP, SSOP AND SOIC PIN DIAGRAMS**

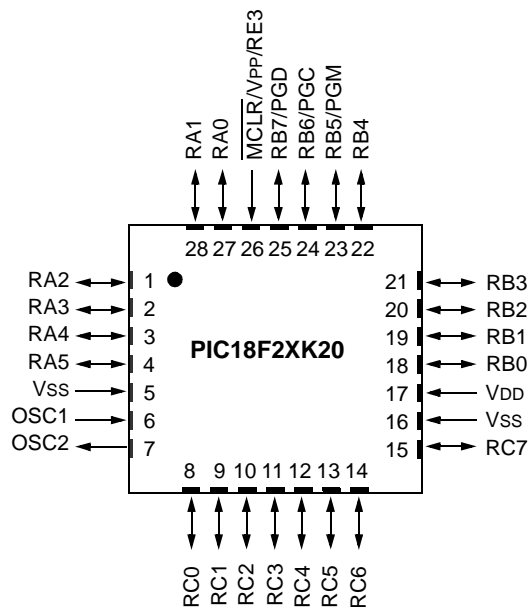
SDIP, SSOP, SOIC (300 MIL)



**Note:** The following devices are included in 28-pin SDIP, SSOP and SOIC parts: PIC18F23K20, PIC18F24K20, PIC18F25K20, PIC18F26K20.

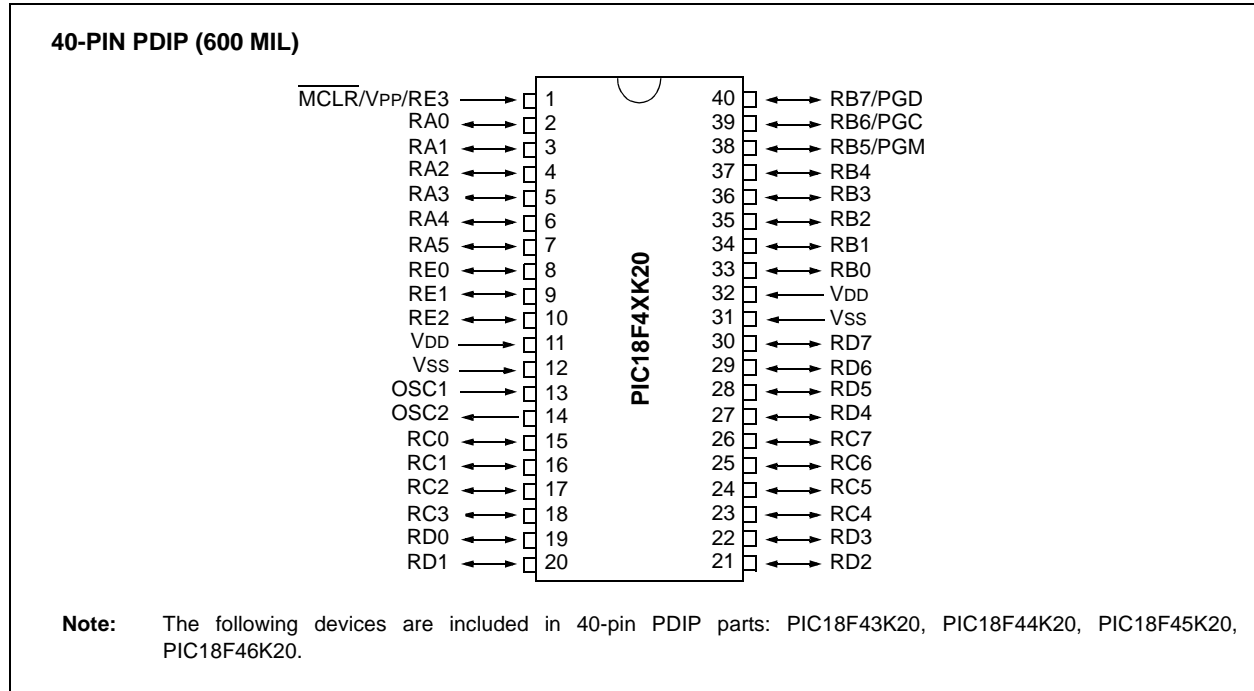
**FIGURE 2-2: 28-PIN QFN PIN DIAGRAMS**

28-Pin QFN

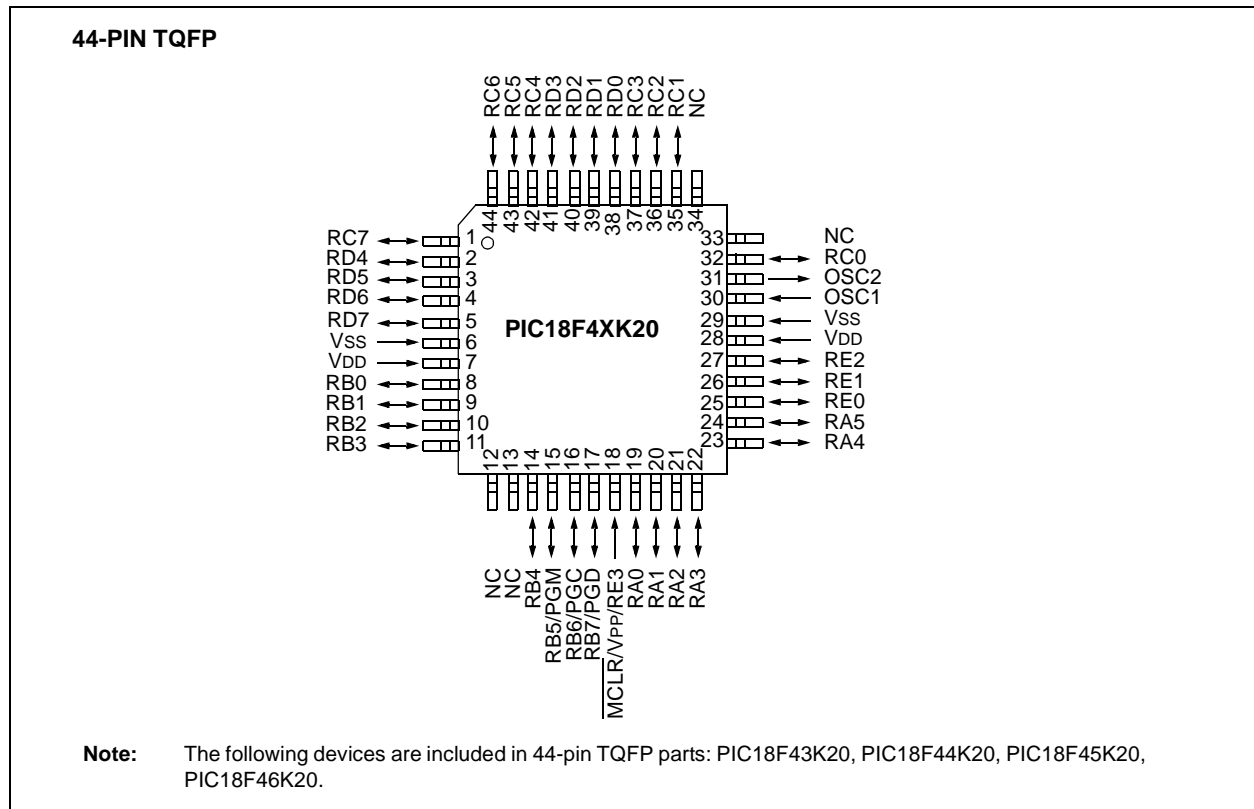


**Note:** The following devices are included in 28-pin QFN parts: PIC18F23K20, PIC18F24K20, PIC18F25K20, PIC18F26K20.

**FIGURE 2-3: 40-PIN PDIP PIN DIAGRAMS**

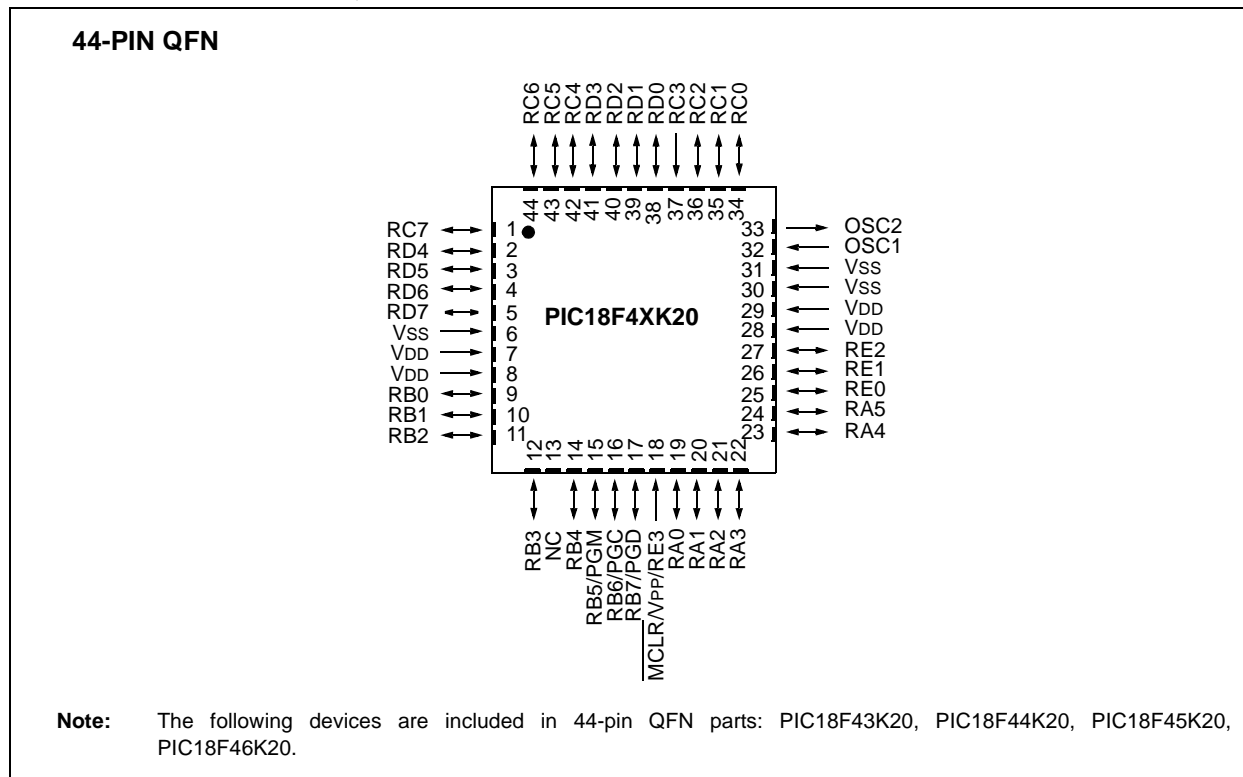


**FIGURE 2-4: 44-PIN TQFP PIN DIAGRAMS**



# PIC18F2XK20/4XK20

FIGURE 2-5: 44-PIN QFN PIN DIAGRAMS



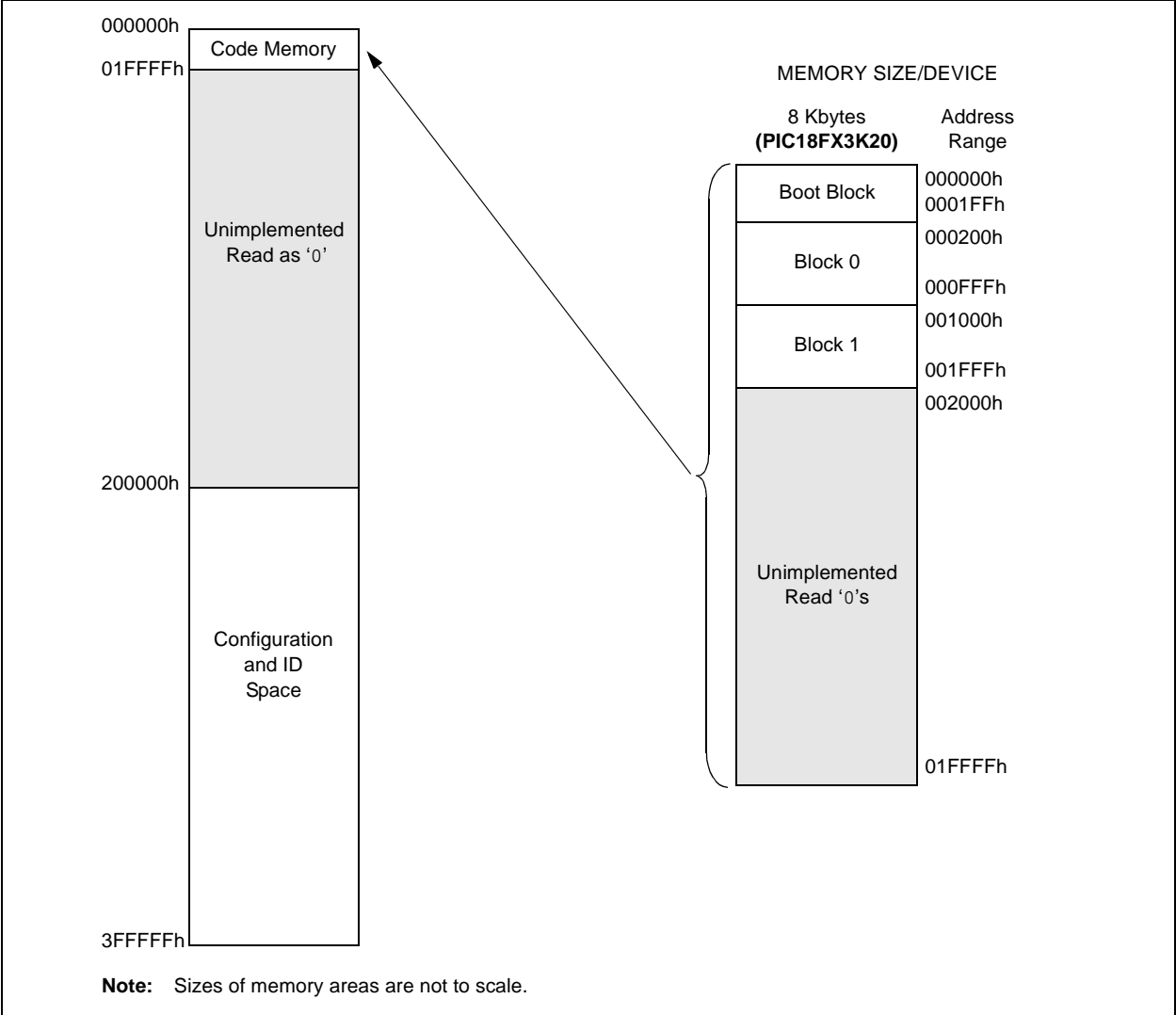
2.3 Memory Maps

For the PIC18FX3K20 devices, the code memory space extends from 0000h to 01FFFh (8 Kbytes) in two 4-Kbyte blocks. Addresses 0000h through 01FFFh, however, define a “Boot Block” region that is treated separately from Block 0. All of these blocks define code protection boundaries within the code memory space.

TABLE 2-2: IMPLEMENTATION OF CODE MEMORY

Device	Code Memory Size (Bytes)
PIC18F23K20	000000h-001FFFh (8K)
PIC18F43K20	

FIGURE 2-6: MEMORY MAP AND THE CODE MEMORY SPACE FOR PIC18FX3K20 DEVICES



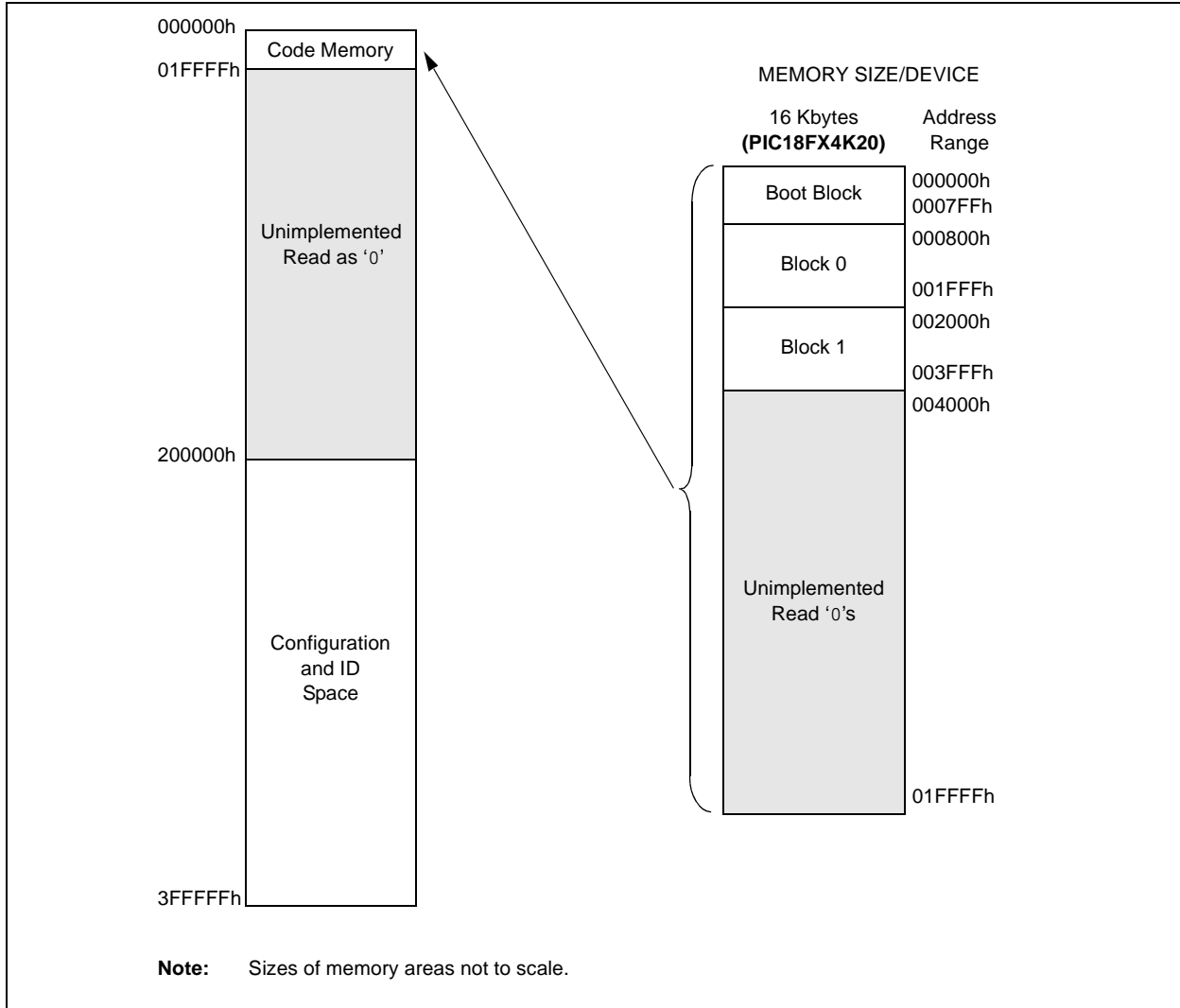
# PIC18F2XK20/4XK20

For PIC18FX4K20 devices, the code memory space extends from 000000h to 003FFFh (16 Kbytes) in two 8-Kbyte blocks. Addresses 000000h through 0007FFh, however, define a “Boot Block” region that is treated separately from Block 0. All of these blocks define code protection boundaries within the code memory space.

**TABLE 2-3: IMPLEMENTATION OF CODE MEMORY**

Device	Code Memory Size (Bytes)
PIC18F24K20	000000h-003FFFh (16K)
PIC18F44K20	

**FIGURE 2-7: MEMORY MAP AND THE CODE MEMORY SPACE FOR PIC18FX4K20 DEVICES**



In addition to the code memory space, there are three blocks in the configuration and ID space that are accessible to the user through table reads and table writes. Their locations in the memory map are shown in Figure 2-10.

Users may store identification information (ID) in eight ID registers. These ID registers are mapped in addresses 200000h through 200007h. The ID locations read out normally, even after code protection is applied.

Locations 300000h through 30000Dh are reserved for the Configuration bits. These bits select various device options and are described in **Section 5.0 “Configuration Word”**. These Configuration bits read out normally, even after code protection.

Locations 3FFFEh and 3FFFFh are reserved for the device ID bits. These bits may be used by the programmer to identify what device type is being programmed and are described in **Section 5.0 “Configuration Word”**. These device ID bits read out normally, even after code protection.

## 2.3.1 MEMORY ADDRESS POINTER

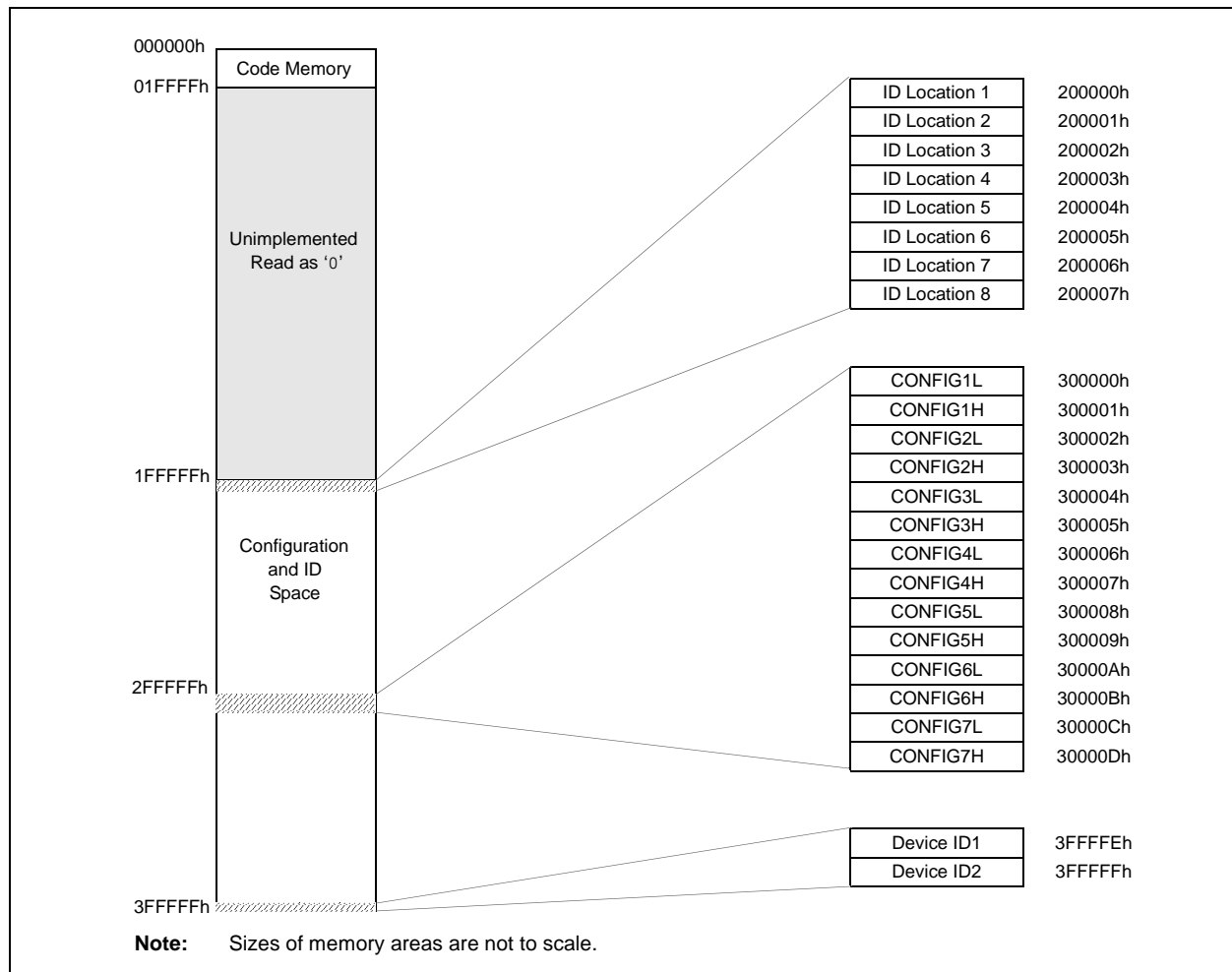
Memory in the address space, 000000h to 3FFFFFFh, is addressed via the Table Pointer register, which is comprised of three Pointer registers:

- TBLPTRU, at RAM address 0FF8h
- TBLPTRH, at RAM address 0FF7h
- TBLPTRL, at RAM address 0FF6h

TBLPTRU	TBLPTRH	TBLPTRL
Addr[21:16]	Addr[15:8]	Addr[7:0]

The 4-bit command, '0000' (core instruction), is used to load the Table Pointer prior to using any read or write operations.

**FIGURE 2-10: CONFIGURATION AND ID LOCATIONS FOR PIC18F2XK20/4XK20 DEVICES**

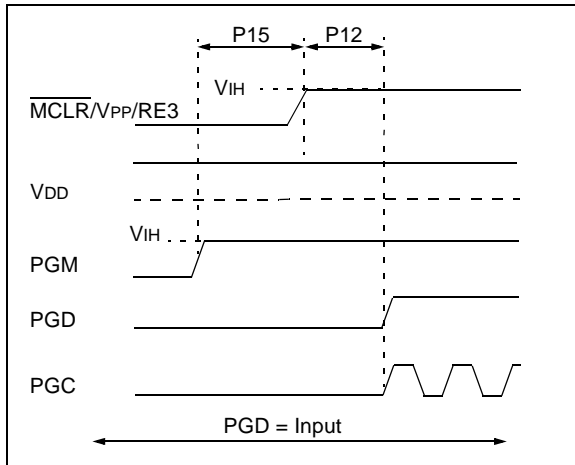


## 2.6 Entering and Exiting Low-Voltage ICSP Program/Verify Mode

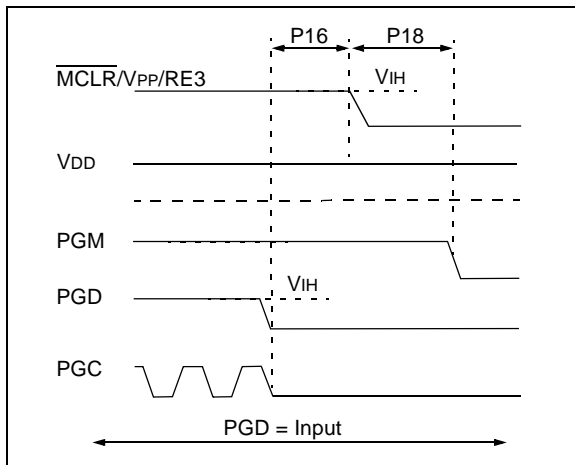
When the LVP Configuration bit is '1' (see **Section 5.3 “Single-Supply ICSP Programming”**), the Low-Voltage ICSP mode is enabled. As shown in Figure 2-14, Low-Voltage ICSP Program/Verify mode is entered by holding PGC and PGD low, placing a logic high on PGM and then raising MCLR/VPP/RE3 to  $V_{IH}$ . In this mode, the RB5/PGM pin is dedicated to the programming function and ceases to be a general purpose I/O pin. Figure 2-15 shows the exit sequence.

The sequence that enters the device into the Program/Verify mode places all unused I/Os in the high-impedance state.

**FIGURE 2-14: ENTERING LOW-VOLTAGE PROGRAM/VERIFY MODE**



**FIGURE 2-15: EXITING LOW-VOLTAGE PROGRAM/VERIFY MODE**



## 2.7 Serial Program/Verify Operation

The PGC pin is used as a clock input pin and the PGD pin is used for entering command bits and data input/output during serial operation. Commands and data are transmitted on the rising edge of PGC, latched on the falling edge of PGC and are Least Significant bit (LSb) first.

### 2.7.1 4-BIT COMMANDS

All instructions are 20 bits, consisting of a leading 4-bit command followed by a 16-bit operand, which depends on the type of command being executed. To input a command, PGC is cycled four times. The commands needed for programming and verification are shown in Table 2-6.

Depending on the 4-bit command, the 16-bit operand represents 16 bits of input data or 8 bits of input data and 8 bits of output data.

Throughout this specification, commands and data are presented as illustrated in Table 2-7. The 4-bit command is shown Most Significant bit (MSb) first. The command operand, or “Data Payload”, is shown <MSB><LSB>. Figure 2-16 demonstrates how to serially present a 20-bit command/operand to the device.

### 2.7.2 CORE INSTRUCTION

The core instruction passes a 16-bit instruction to the CPU core for execution. This is needed to set up registers as appropriate for use with other commands.

**TABLE 2-6: COMMANDS FOR PROGRAMMING**

Description	4-Bit Command
Core Instruction (Shift in 16-bit instruction)	0000
Shift out TABLAT register	0010
Table Read	1000
Table Read, post-increment	1001
Table Read, post-decrement	1010
Table Read, pre-increment	1011
Table Write	1100
Table Write, post-increment by 2	1101
Table Write, start programming, post-increment by 2	1110
Table Write, start programming	1111

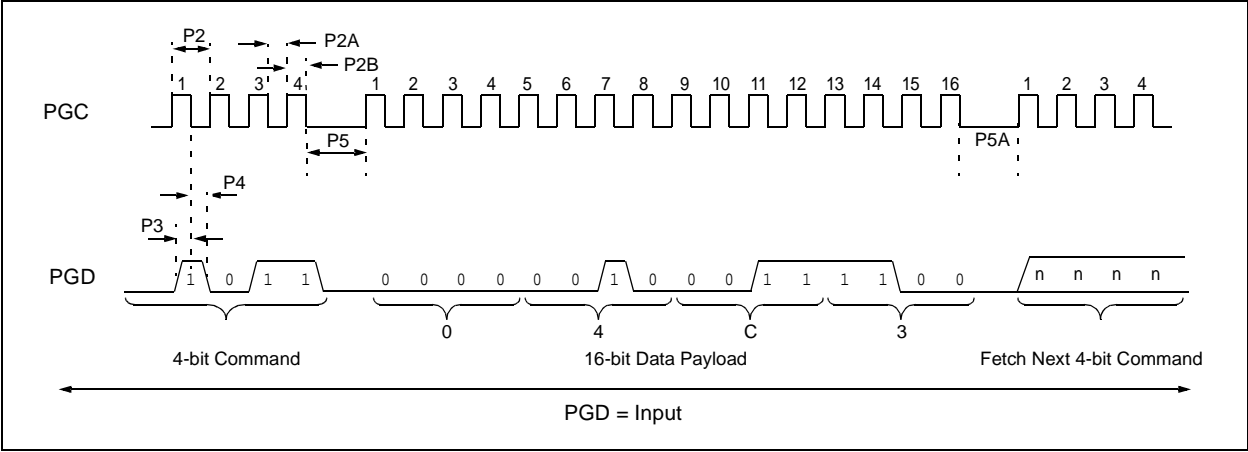


# PIC18F2XK20/4XK20

TABLE 2-7: SAMPLE COMMAND SEQUENCE

4-Bit Command	Data Payload	Core Instruction
1101	3C 40	Table Write, post-increment by 2

FIGURE 2-16: TABLE WRITE, POST-INCREMENT TIMING DIAGRAM (1101)



## 3.0 DEVICE PROGRAMMING

Programming includes the ability to erase or write the various memory regions within the device.

In all cases, except high-voltage ICSP Bulk Erase, the EECON1 register must be configured in order to operate on a particular memory region.

When using the EECON1 register to act on code memory, the EEPGD bit must be set (EECON1<7> = 1) and the CFGS bit must be cleared (EECON1<6> = 0). The WREN bit must be set (EECON1<2> = 1) to enable writes of any sort (e.g., erases) and this must be done prior to initiating a write sequence. The FREE bit must be set (EECON1<4> = 1) in order to erase the program space being pointed to by the Table Pointer. The erase or write sequence is initiated by setting the WR bit (EECON1<1> = 1). It is strongly recommended that the WREN bit only be set immediately prior to a program or erase.

### 3.1 ICSP Erase

#### 3.1.1 HIGH-VOLTAGE ICSP BULK ERASE

Erasing code or data EEPROM is accomplished by configuring two Bulk Erase Control registers located at 3C0004h and 3C0005h. Code memory may be erased portions at a time, or the user may erase the entire device in one action. Bulk Erase operations will also clear any code-protect settings associated with the memory block erased. Erase options are detailed in Table 3-1. If data EEPROM is code-protected (CPD = 0), the user must request an erase of data EEPROM (e.g., 0084h as shown in Table 3-1).

**TABLE 3-1: BULK ERASE OPTIONS**

Description	Data (3C0005h:3C0004h)
Chip Erase	0F8Fh
Erase User ID	0088h
Erase Data EEPROM	0084h
Erase Boot Block	0081h
Erase Config Bits	0082h
Erase Code EEPROM Block 0	0180h
Erase Code EEPROM Block 1	0280h
Erase Code EEPROM Block 2	0480h
Erase Code EEPROM Block 3	0880h

The actual Bulk Erase function is a self-timed operation. Once the erase has started (falling edge of the 4th PGC after the NOP command), serial execution will cease until the erase completes (parameter P11). During this time, PGC may continue to toggle but PGD must be held low.

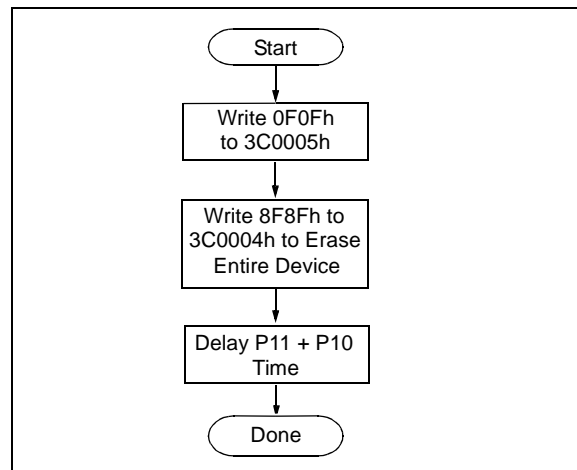
The code sequence to erase the entire device is shown in Table 3-2 and the flowchart is shown in Figure 3-1.

**Note:** A Bulk Erase is the only way to reprogram code-protect bits from an “on” state to an “off” state.

**TABLE 3-2: BULK ERASE COMMAND SEQUENCE**

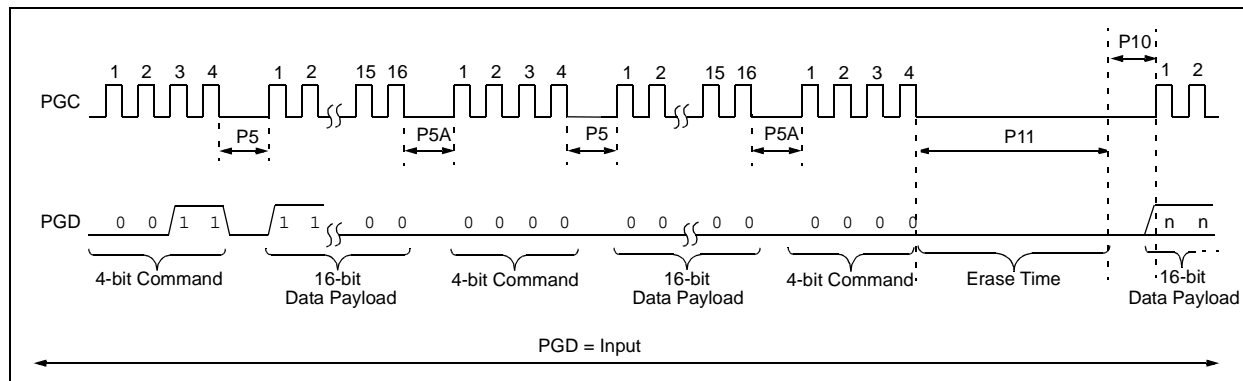
4-Bit Command	Data Payload	Core Instruction
0000	0E 3C	MOVLW 3Ch
0000	6E F8	MOVWF TBLPTRU
0000	0E 00	MOVLW 00h
0000	6E F7	MOVWF TBLPTRH
0000	0E 05	MOVLW 05h
0000	6E F6	MOVWF TBLPTRL
1100	0F 0F	Write 0Fh to 3C0005h
0000	0E 3C	MOVLW 3Ch
0000	6E F8	MOVWF TBLPTRU
0000	0E 00	MOVLW 00h
0000	6E F7	MOVWF TBLPTRH
0000	0E 04	MOVLW 04h
0000	6E F6	MOVWF TBLPTRL
1100	8F 8F	Write 8F8Fh TO 3C0004h to erase entire device.
0000	00 00	NOP
0000	00 00	Hold PGD low until erase completes.

**FIGURE 3-1: BULK ERASE FLOW**



# PIC18F2XK20/4XK20

**FIGURE 3-2: BULK ERASE TIMING DIAGRAM**



## 3.1.2 LOW-VOLTAGE ICSP BULK ERASE

When using low-voltage ICSP, the part must be supplied by the voltage specified in parameter D111 if a Bulk Erase is to be executed. All other Bulk Erase details as described above apply.

If it is determined that a program memory erase must be performed at a supply voltage below the Bulk Erase limit, refer to the erase methodology described in **Section 3.1.3 “ICSP Row Erase”** and **Section 3.2.1 “Modifying Code Memory”**.

If it is determined that a data EEPROM erase must be performed at a supply voltage below the Bulk Erase limit, follow the methodology described in **Section 3.3 “Data EEPROM Programming”** and write ‘1’s to the array.

## 3.1.3 ICSP ROW ERASE

Regardless of whether high or low-voltage ICSP is used, it is possible to erase one row (64 bytes of data), provided the block is not code or write-protected. Rows are located at static boundaries beginning at program memory address 000000h, extending to the internal program memory limit (see **Section 2.3 “Memory Maps”**).

The Row Erase duration is self-timed. After the WR bit in EECON1 is set, two NOPs are issued. Erase starts upon the 4th PGC of the second NOP. It ends when the WR bit is cleared by hardware.

The code sequence to Row Erase a PIC18F2XK20/4XK20 device is shown in Table 3-3. The flowchart shown in Figure 3-3 depicts the logic necessary to completely erase a PIC18F2XK20/4XK20 device. The timing diagram for Row Erase is identical to the data EEPROM write timing shown in Figure 3-7.

**Note:** The TBLPTR register can point at any byte within the row intended for erase.

**TABLE 3-3: ERASE CODE MEMORY CODE SEQUENCE**

4-bit Command	Data Payload	Core Instruction
Step 1: Direct access to code memory and enable writes.		
0000	8E A6	BSF EECON1, EEPGD
0000	9C A6	BCF EECON1, CFGS
0000	84 A6	BSF EECON1, WREN
Step 2: Point to first row in code memory.		
0000	6A F8	CLRF TBLPTRU
0000	6A F7	CLRF TBLPTRH
0000	6A F6	CLRF TBLPTRL
Step 3: Enable erase and erase single row.		
0000	88 A6	BSF EECON1, FREE
0000	82 A6	BSF EECON1, WR
0000	00 00	NOP
0000	00 00	NOP Erase starts on the 4th clock of this instruction
Step 4: Poll WR bit. Repeat until bit is clear.		
0000	50 A6	MOVF EECON1, W, 0
0000	6E F5	MOVWF TABLAT
0000	00 00	NOP
0010	<MSB><LSB>	Shift out data <sup>(1)</sup>
Step 5: Hold PGC low for time P10.		
Step 6: Repeat step 3 with Address Pointer incremented by 64 until all rows are erased.		
Step 7: Disable writes.		
0000	94 A6	BCF EECON1, WREN

**Note 1:** See Figure 4-4 for details on shift out data timing.

---

```

graph TD
    Start([Start]) -- "N = 1  
LoopCount = 0" --> Configure[Configure Device for Writes]
    Configure --> Load[Load 2 Bytes to Write Buffer at <Addr>]
    Load --> AllBytesWritten{All bytes written?}
    AllBytesWritten -- No --> Nplus1[N = N + 1]
    Nplus1 --> Load
    AllBytesWritten -- Yes --> StartWrite[Start Write Sequence and Hold PGC High until Done and Wait P9]
    StartWrite --> HoldPGCLow[Hold PGC Low for Time P10]
    HoldPGCLow --> AllLocationsDone{All locations done?}
    AllLocationsDone -- No --> LoopCountInc[N = 1  
LoopCount = LoopCount + 1]
    LoopCountInc --> Configure
    AllLocationsDone -- Yes --> Done([Done])
  
```

The diagram illustrates the timing relationship between the PGC (Pulse Generator Clock) and PGD (Program Data) signals during a programming operation. The PGC signal is a periodic square wave, with specific pulses numbered 1 through 16. The PGD signal is a data bus that carries the programming data. The sequence of events is as follows:

- PGC Pulses:** The PGC signal consists of a series of pulses. Pulses 1 through 4 are followed by a period P5. Pulses 1 through 4 are followed by a period P5A. Pulses 1 through 3 are followed by a period P9(1). Pulses 1 through 3 are followed by a period P10.
- PGD Data:** The PGD signal carries the programming data. It starts with a 4-bit Command (1111), followed by a 16-bit Data Payload (n's), a 4-bit Command (0000), a Programming Time interval, and a 16-bit Data Payload (0000).
- PGD = Input:** A horizontal arrow at the bottom indicates that the PGD signal is the input to the device.

**Note 1:** Use P9A for User ID and Configuration Word programming.

# PIC18F2XK20/4XK20

## 3.4 ID Location Programming

The ID locations are programmed much like the code memory. The ID registers are mapped in addresses 200000h through 200007h. These locations read out normally even after code protection.

**Note:** The user only needs to fill the first 8 bytes of the write buffer in order to write the ID locations.

Table 3-8 demonstrates the code sequence required to write the ID locations.

In order to modify the ID locations, refer to the methodology described in **Section 3.2.1 “Modifying Code Memory”**. As with code memory, the ID locations must be erased before being modified.

When VDD is below the minimum for Bulk Erase operation, ID locations can be cleared with the Row Erase method described in **Section 3.1.3 “ICSP Row Erase”**.

**TABLE 3-8: WRITE ID SEQUENCE**

4-bit Command	Data Payload	Core Instruction
Step 1: Direct access to code memory.		
0000	8E A6	BSF EECON1, EEPGD
0000	9C A6	BCF EECON1, CFGS
0000	84 A6	BSF EECON1, WREN
Step 2: Set Table Pointer to ID. Load write buffer with 8 bytes and write.		
0000	0E 20	MOVLW 20h
0000	6E F8	MOVWF TBLPTRU
0000	0E 00	MOVLW 00h
0000	6E F7	MOVWF TBLPTRH
0000	0E 00	MOVLW 00h
0000	6E F6	MOVWF TBLPTRL
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
1111	<MSB><LSB>	Write 2 bytes and start programming.
0000	00 00	NOP - hold PGC high for time P9 and low for time P10.

# PIC18F2XK20/4XK20

## 4.0 READING THE DEVICE

### 4.1 Read Code Memory, ID Locations and Configuration Bits

Code memory is accessed one byte at a time via the 4-bit command, '1001' (table read, post-increment). The contents of memory pointed to by the Table Pointer (TBLPTRU:TBLPTRH:TBLPTRL) are serially output on PGD.

The 4-bit command is shifted in LSb first. The read is executed during the next 8 clocks, then shifted out on PGD during the last 8 clocks, LSb to MSb. A delay of P6 must be introduced after the falling edge of the 8th

PGC of the operand to allow PGD to transition from an input to an output. During this time, PGC must be held low (see Figure 4-1). This operation also increments the Table Pointer by one, pointing to the next byte in code memory for the next read.

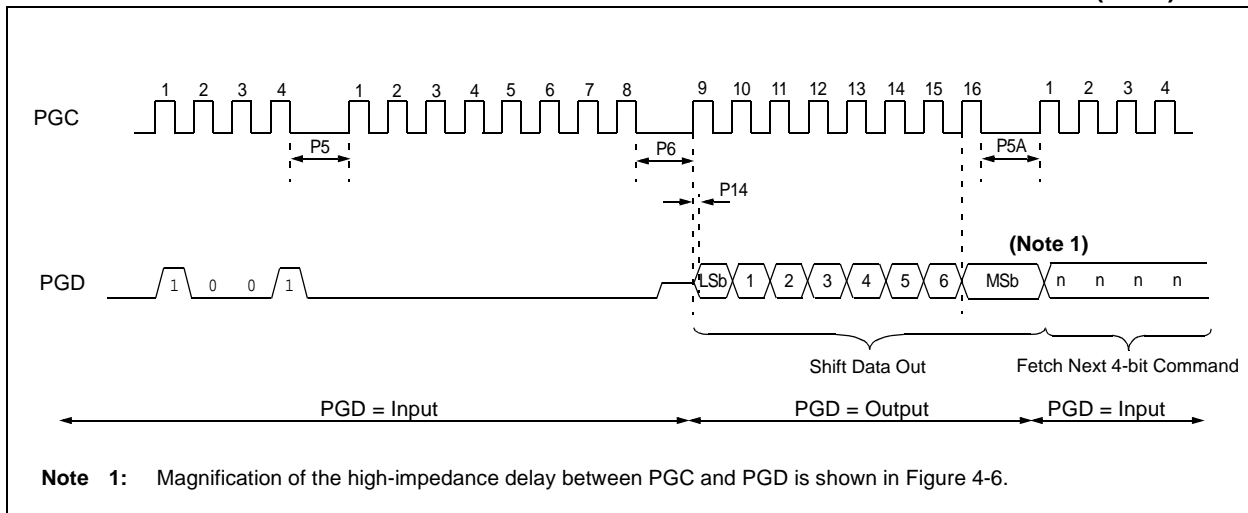
This technique will work to read any memory in the 000000h to 3FFFFFFh address space, so it also applies to the reading of the ID and Configuration registers.

**Note:** When table read protection is enabled, the first read access to a protected block should be discarded and the read repeated to retrieve valid data. Subsequent reads of the same block can be performed normally.

**TABLE 4-1: READ CODE MEMORY SEQUENCE**

4-bit Command	Data Payload	Core Instruction
Step 1: Set Table Pointer		
0000	0E <Addr[21:16]>	MOVLW Addr[21:16]
0000	6E F8	MOVWF TBLPTRU
0000	0E <Addr[15:8]>	MOVLW <Addr[15:8]>
0000	6E F7	MOVWF TBLPTRH
0000	0E <Addr[7:0]>	MOVLW <Addr[7:0]>
0000	6E F6	MOVWF TBLPTRL
Step 2: Read memory and then shift out on PGD, LSb to MSb		
1001	00 00	TBLRD *+

**FIGURE 4-1: TABLE READ POST-INCREMENT INSTRUCTION TIMING DIAGRAM (1001)**

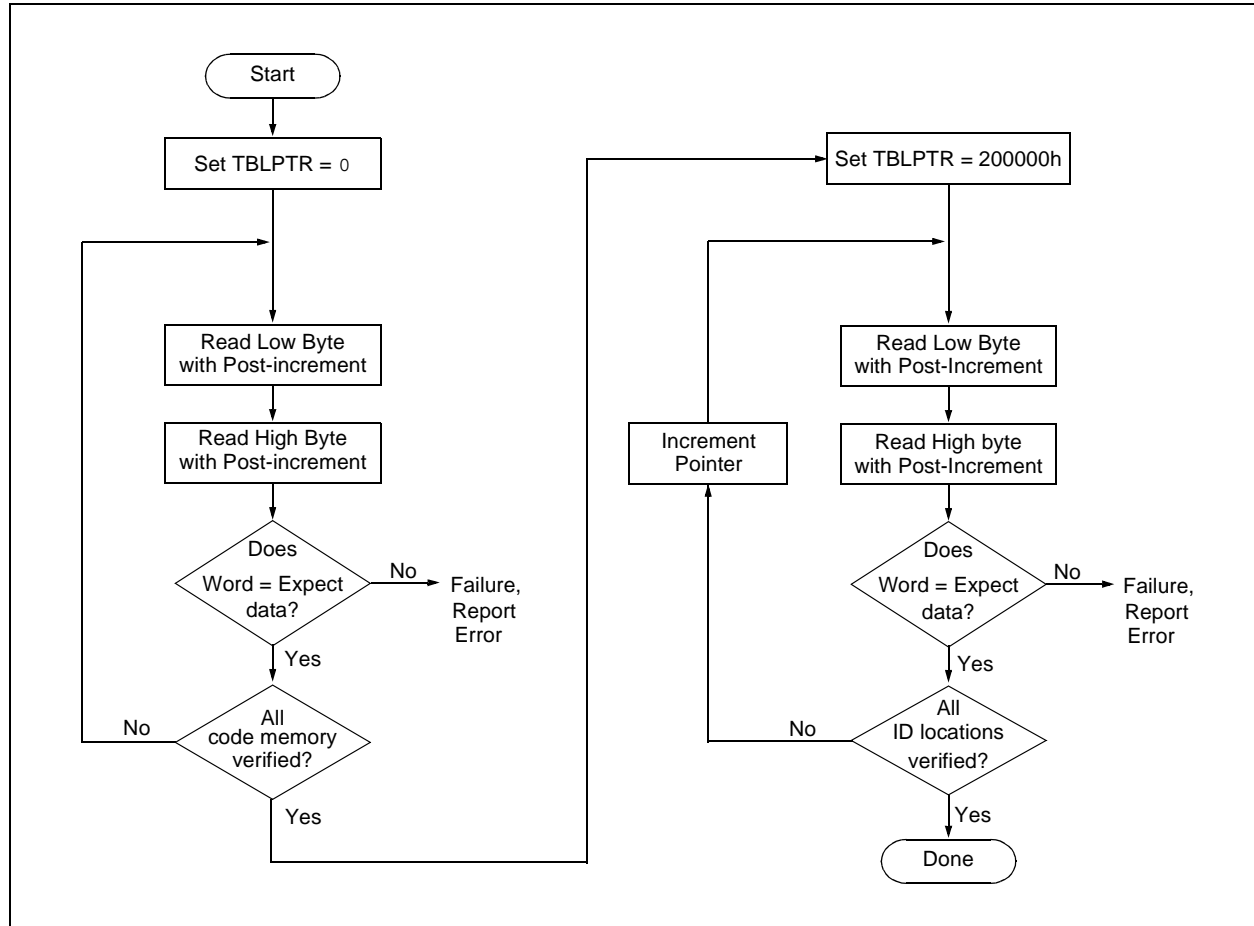


## 4.2 Verify Code Memory and ID Locations

The verify step involves reading back the code memory space and comparing it against the copy held in the programmer's buffer. Memory reads occur a single byte at a time, so two bytes must be read to compare against the word in the programmer's buffer. Refer to **Section 4.1 "Read Code Memory, ID Locations and Configuration Bits"** for implementation details of reading code memory.

The Table Pointer must be manually set to 200000h (base address of the ID locations) once the code memory has been verified. The post-increment feature of the table read 4-bit command can not be used to increment the Table Pointer beyond the code memory space. In a 64-Kbyte device, for example, a post-increment read of address FFFFh will wrap the Table Pointer back to 000000h, rather than point to unimplemented address 010000h.

**FIGURE 4-2: VERIFY CODE MEMORY FLOW**





**TABLE 5-3: PIC18F2XK20/4XK20 BIT DESCRIPTIONS (CONTINUED)**

Bit Name	Configuration Words	Description
WDTEN	CONFIG2H	Watchdog Timer Enable bit 1 = WDT enabled 0 = WDT disabled (control is placed on SWDTEN bit)
MCLRE	CONFIG3H	MCLR Pin Enable bit 1 = MCLR pin enabled, RE3 input pin disabled 0 = RE3 input pin enabled, MCLR pin disabled
HFOFST	CONFIG3H	HFINTOSC Fast Start 1 = HFINTOSC output is not delayed 0 = HFINTOSC output is delayed until oscillator is stable (IOFS = 1)
LPT1OSC	CONFIG3H	Low-Power Timer1 Oscillator Enable bit 1 = Timer1 configured for low-power operation 0 = Timer1 configured for higher power operation
PBADEN	CONFIG3H	PORTB A/D Enable bit 1 = PORTB A/D<4:0> pins are configured as analog input channels on Reset 0 = PORTB A/D<4:0> pins are configured as digital I/O on Reset
CCP2MX	CONFIG3H	CCP2 MUX bit 1 = CCP2 input/output is multiplexed with RC1 0 = CCP2 input/output is multiplexed with RB3
DEBUG	CONFIG4L	Background Debugger Enable bit 1 = Background debugger disabled, RB6 and RB7 configured as general purpose I/O pins 0 = Background debugger enabled, RB6 and RB7 are dedicated to In-Circuit Debug
XINST	CONFIG4L	Extended Instruction Set Enable bit 1 = Instruction set extension and Indexed Addressing mode enabled 0 = Instruction set extension and Indexed Addressing mode disabled (Legacy mode)
LVP	CONFIG4L	Low-Voltage Programming Enable bit 1 = Low-Voltage Programming enabled, RB5 is the PGM pin 0 = Low-Voltage Programming disabled, RB5 is an I/O pin
STVREN	CONFIG4L	Stack Overflow/Underflow Reset Enable bit 1 = Reset on stack overflow/underflow enabled 0 = Reset on stack overflow/underflow disabled

# PIC18F2XK20/4XK20

**TABLE 5-3: PIC18F2XK20/4XK20 BIT DESCRIPTIONS (CONTINUED)**

Bit Name	Configuration Words	Description
CP3	CONFIG5L	Code Protection bits (Block 3 code memory area) 1 = Block 3 is not code-protected 0 = Block 3 is code-protected
CP2	CONFIG5L	Code Protection bits (Block 2 code memory area) 1 = Block 2 is not code-protected 0 = Block 2 is code-protected
CP1	CONFIG5L	Code Protection bits (Block 1 code memory area) 1 = Block 1 is not code-protected 0 = Block 1 is code-protected
CP0	CONFIG5L	Code Protection bits (Block 0 code memory area) 1 = Block 0 is not code-protected 0 = Block 0 is code-protected
CPD	CONFIG5H	Code Protection bits (Data EEPROM) 1 = Data EEPROM is not code-protected 0 = Data EEPROM is code-protected
CPB	CONFIG5H	Code Protection bits (Boot Block memory area) 1 = Boot Block is not code-protected 0 = Boot Block is code-protected
WRT3	CONFIG6L	Write Protection bits (Block 3 code memory area) 1 = Block 3 is not write-protected 0 = Block 3 is write-protected
WRT2	CONFIG6L	Write Protection bits (Block 2 code memory area) 1 = Block 2 is not write-protected 0 = Block 2 is write-protected
WRT1	CONFIG6L	Write Protection bits (Block 1 code memory area) 1 = Block 1 is not write-protected 0 = Block 1 is write-protected
WRT0	CONFIG6L	Write Protection bits (Block 0 code memory area) 1 = Block 0 is not write-protected 0 = Block 0 is write-protected
WRTD	CONFIG6H	Write Protection bit (Data EEPROM) 1 = Data EEPROM is not write-protected 0 = Data EEPROM is write-protected
WRTB	CONFIG6H	Write Protection bit (Boot Block memory area) 1 = Boot Block is not write-protected 0 = Boot Block is write-protected
WRTC	CONFIG6H	Write Protection bit (Configuration registers) 1 = Configuration registers are not write-protected 0 = Configuration registers are write-protected

**TABLE 5-3: PIC18F2XK20/4XK20 BIT DESCRIPTIONS (CONTINUED)**

Bit Name	Configuration Words	Description
EBTR3	CONFIG7L	Table Read Protection bit (Block 3 code memory area) 1 = Block 3 is not protected from table reads executed in other blocks 0 = Block 3 is protected from table reads executed in other blocks
EBTR2	CONFIG7L	Table Read Protection bit (Block 2 code memory area) 1 = Block 2 is not protected from table reads executed in other blocks 0 = Block 2 is protected from table reads executed in other blocks
EBTR1	CONFIG7L	Table Read Protection bit (Block 1 code memory area) 1 = Block 1 is not protected from table reads executed in other blocks 0 = Block 1 is protected from table reads executed in other blocks
EBTR0	CONFIG7L	Table Read Protection bit (Block 0 code memory area) 1 = Block 0 is not protected from table reads executed in other blocks 0 = Block 0 is protected from table reads executed in other blocks
EBTRB	CONFIG7H	Table Read Protection bit (Boot Block memory area) 1 = Boot Block is not protected from table reads executed in other blocks 0 = Boot Block is protected from table reads executed in other blocks
DEV<10:3>	DEVID2	Device ID bits These bits are used with the DEV<2:0> bits in the DEVID1 register to identify part number.
DEV<2:0>	DEVID1	Device ID bits These bits are used with the DEV<10:3> bits in the DEVID2 register to identify part number.
REV<4:0>	DEVID1	Revision ID bits These bits are used to indicate the revision of the device.

## 5.3 Single-Supply ICSP Programming

The LVP bit in Configuration register, CONFIG4L, enables Single-Supply (Low-Voltage) ICSP Programming. The LVP bit defaults to a '1' (enabled) from the factory.

If Single-Supply Programming mode is not used, the LVP bit can be programmed to a '0' and RB5/PGM becomes a digital I/O pin. However, the LVP bit may only be programmed by entering the High-Voltage ICSP mode, where MCLR/VPP/RE3 is raised to  $V_{IH}$ . Once the LVP bit is programmed to a '0', only the High-Voltage ICSP mode is available and only the High-Voltage ICSP mode can be used to program the device.

**Note 1:** The High-Voltage ICSP mode is always available, regardless of the state of the LVP bit, by applying  $V_{IH}$  to the MCLR/VPP/RE3 pin.

**2:** While in Low-Voltage ICSP mode, the RB5 pin can no longer be used as a general purpose I/O.

## 5.4 Embedding Configuration Word Information in the HEX File

To allow portability of code, a PIC18F2XK20/4XK20 programmer is required to read the Configuration Word locations from the hex file. If Configuration Word information is not present in the hex file, then a simple warning message should be issued. Similarly, while saving a hex file, all Configuration Word information must be included. An option to not include the Configuration Word information may be provided. When embedding Configuration Word information in the hex file, it should start at address 300000h.

Microchip Technology Inc. feels strongly that this feature is important for the benefit of the end customer.

## 5.5 Embedding Data EEPROM Information In the HEX File

To allow portability of code, a PIC18F2XK20/4XK20 programmer is required to read the data EEPROM information from the hex file. If data EEPROM information is not present, a simple warning message should be issued. Similarly, when saving a hex file, all data EEPROM information must be included. An option to not include the data EEPROM information may be provided. When embedding data EEPROM information in the hex file, it should start at address F00000h.

Microchip Technology Inc. believes that this feature is important for the benefit of the end customer.

## 5.6 Checksum Computation

The checksum is calculated by summing the following:

- The contents of all code memory locations
- The Configuration Word, appropriately masked
- ID locations (Only if any portion of program memory is code-protected)

The Least Significant 16 bits of this sum are the checksum.

Code protection limits access to program memory by both external programmer (code-protect) and code execution (table read protect). The ID locations, when included in a code protected checksum, contain the checksum of an unprotected part. The unprotected checksum is distributed: one nibble per ID location. Each nibble is right justified.

Table 5-4 describes how to calculate the checksum for each device.

**Note:** The checksum calculation differs depending on the code-protect setting. Since the code memory locations read out differently depending on the code-protect setting, the table describes how to manipulate the actual code memory values to simulate the values that would be read from a protected device. When calculating a checksum by reading a device, the entire code memory can simply be read and summed. The Configuration Word and ID locations can always be read.

# PIC18F2XK20/4XK20

**TABLE 5-4: CHECKSUM COMPUTATION (CONTINUED)**

Device	Code-Protect	Checksum	Blank Value	0xAA at 0 and Max Address
PIC18FX5K20	None	SUM[0000:07FF]+SUM[0800:1FFF]+SUM[2000:3FFF]+SUM[4000:5FFF]+SUM[6000:7FFF]+(CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+(CONFIG7H & 40h)	8362h	82B8h
	Boot Block	SUM[0800:1FFF]+SUM[2000:3FFF]+SUM[4000:5FFF]+SUM[6000:7FFF]+(CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	8B35h	8AEAh
	Boot/Block 0/Block 1	SUM[4000:5FFF]+SUM[6000:7FFF]+(CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	C332h	C2E7h
	All	(CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	0326h	0330h

**Legend:**

Item	Description
CONFIGx	Configuration Word
SUM[a:b]	Sum of locations, a to b inclusive
SUM_ID	Byte-wise sum of lower four bits of all customer ID locations
+	Addition
&	Bit-wise AND