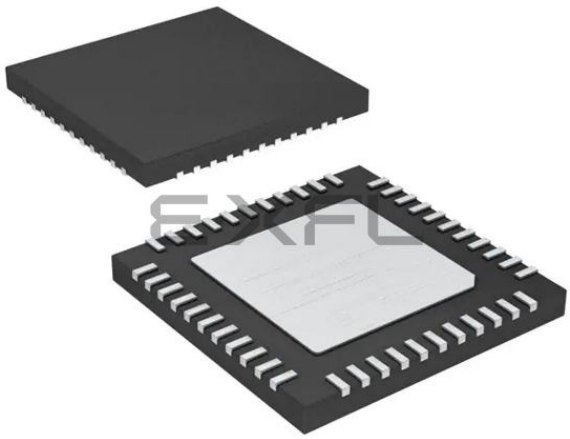

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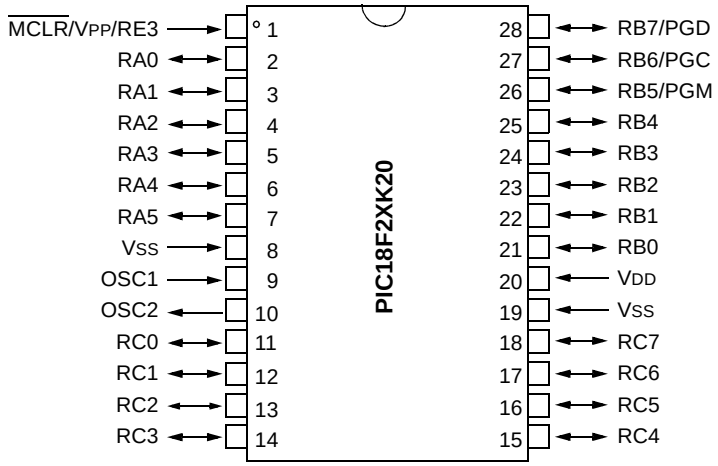
Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	64MHz
Connectivity	I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, HLVD, POR, PWM, WDT
Number of I/O	35
Program Memory Size	8KB (4K x 16)
Program Memory Type	FLASH
EEPROM Size	256 x 8
RAM Size	512 x 8
Voltage - Supply (Vcc/Vdd)	1.8V ~ 3.6V
Data Converters	A/D 14x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	44-VQFN Exposed Pad
Supplier Device Package	44-QFN (8x8)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18f43k20t-i-ml

FIGURE 2-1: 28-PIN SDIP, SSOP AND SOIC PIN DIAGRAMS

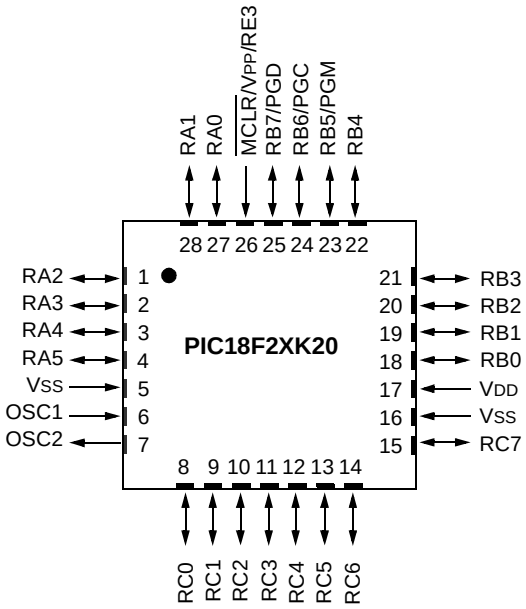
SDIP, SSOP, SOIC (300 MIL)



Note: The following devices are included in 28-pin SDIP, SSOP and SOIC parts: PIC18F23K20, PIC18F24K20, PIC18F25K20, PIC18F26K20.

FIGURE 2-2: 28-PIN QFN PIN DIAGRAMS

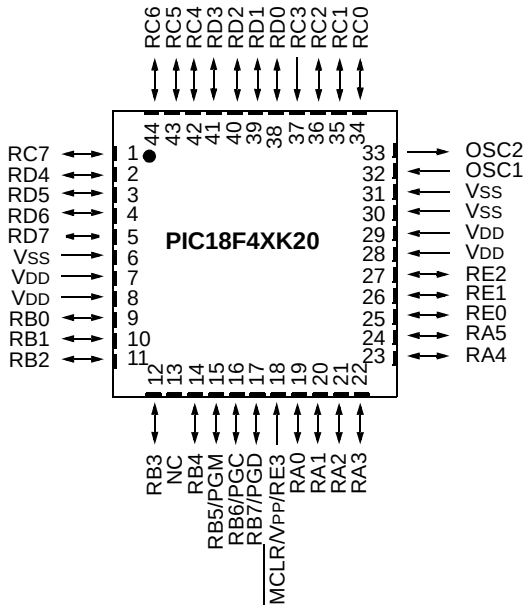
28-Pin QFN



Note: The following devices are included in 28-pin QFN parts: PIC18F23K20, PIC18F24K20, PIC18F25K20, PIC18F26K20.

FIGURE 2-5: 44-PIN QFN PIN DIAGRAMS

44-PIN QFN



Note: The following devices are included in 44-pin QFN parts: PIC18F43K20, PIC18F44K20, PIC18F45K20, PIC18F46K20.

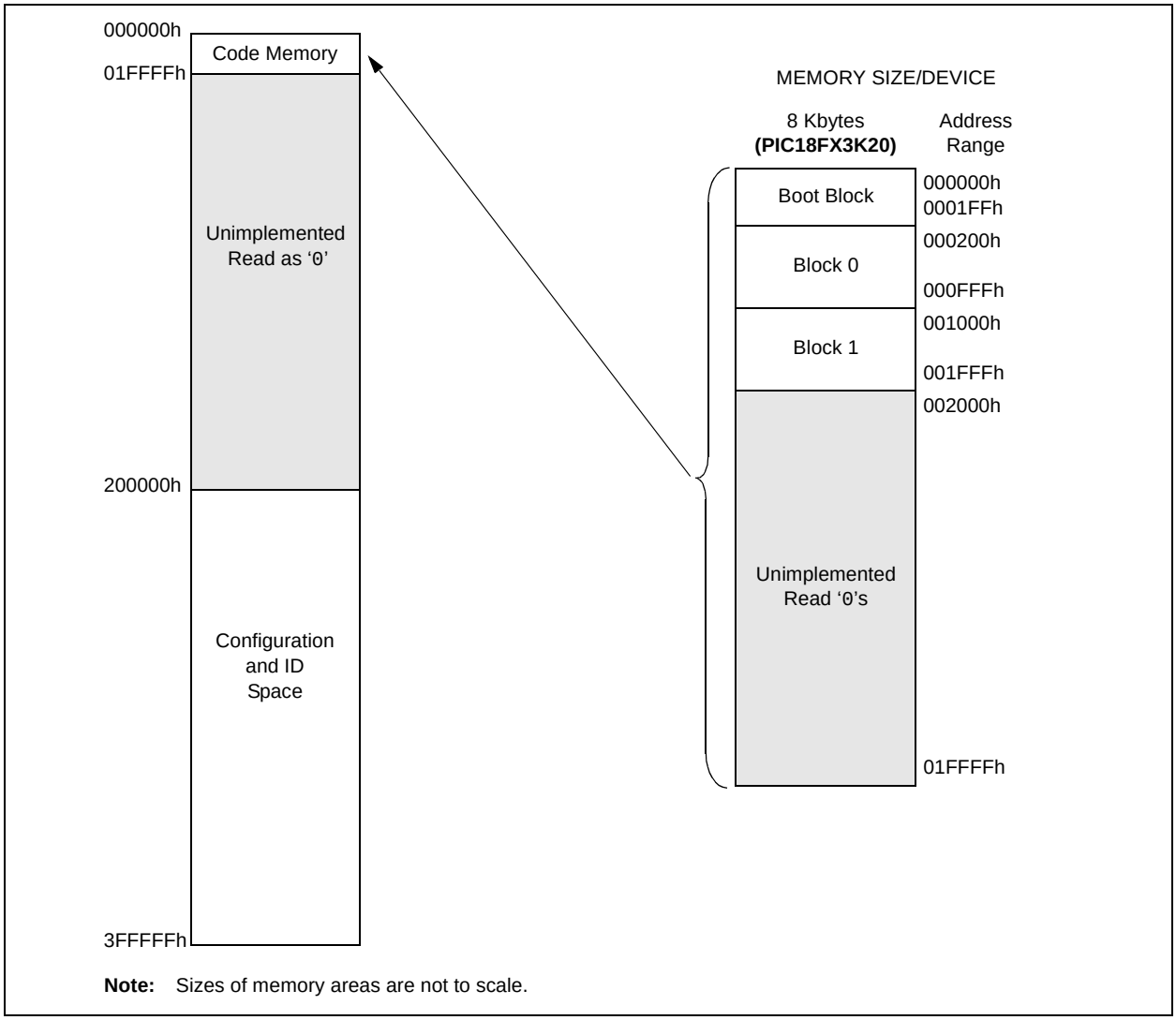
2.3 Memory Maps

For the PIC18FX3K20 devices, the code memory space extends from 0000h to 01FFFh (8 Kbytes) in two 4-Kbyte blocks. Addresses 0000h through 01FFh, however, define a “Boot Block” region that is treated separately from Block 0. All of these blocks define code protection boundaries within the code memory space.

TABLE 2-2: IMPLEMENTATION OF CODE MEMORY

Device	Code Memory Size (Bytes)
PIC18F23K20	000000h-001FFFh (8K)
PIC18F43K20	

FIGURE 2-6: MEMORY MAP AND THE CODE MEMORY SPACE FOR PIC18FX3K20 DEVICES



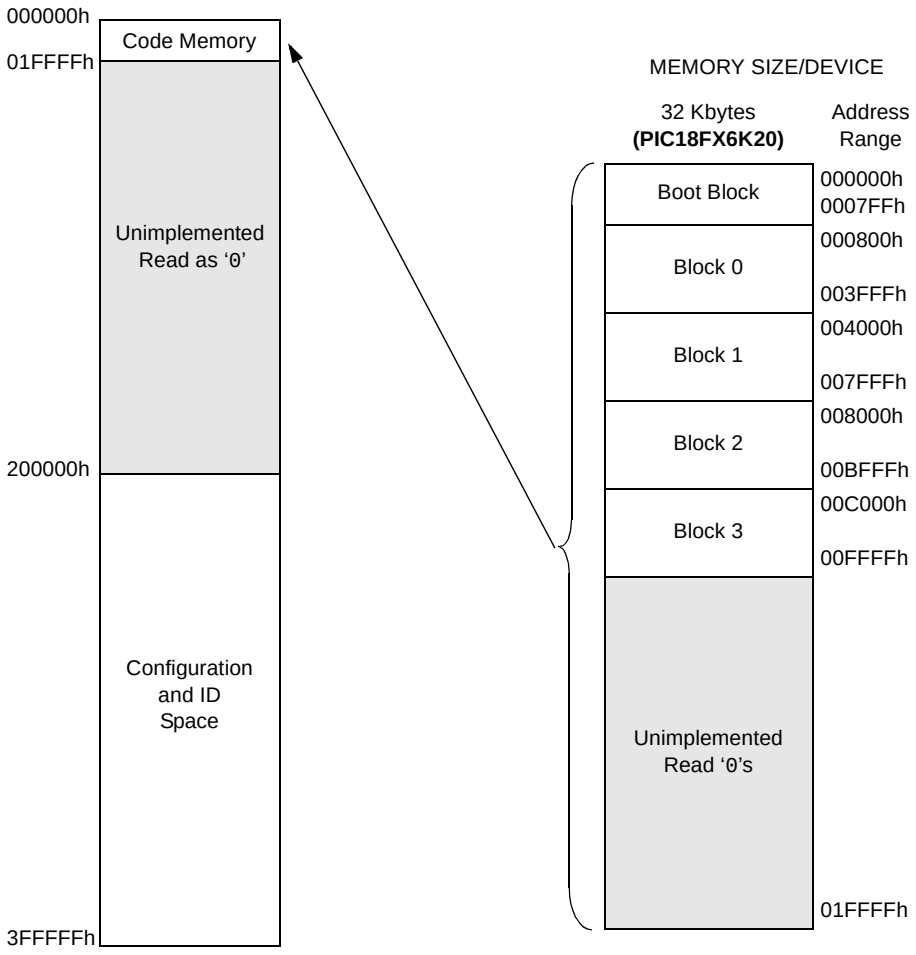
Note: Sizes of memory areas are not to scale.

For PIC18FX6K20 devices, the code memory space extends from 000000h to 00FFFFh (64 Kbytes) in four 16-Kbyte blocks. Addresses 000000h through 0007FFh, however, define a “Boot Block” region that is treated separately from Block 0. All of these blocks define code protection boundaries within the code memory space.

TABLE 2-5: IMPLEMENTATION OF CODE MEMORY

Device	Code Memory Size (Bytes)
PIC18F26K20	000000h-00FFFFh (64K)
PIC18F46K20	

FIGURE 2-9: MEMORY MAP AND THE CODE MEMORY SPACE FOR PIC18FX6K20 DEVICES



Note: Sizes of memory areas not to scale.

In addition to the code memory space, there are three blocks in the configuration and ID space that are accessible to the user through table reads and table writes. Their locations in the memory map are shown in Figure 2-10.

Users may store identification information (ID) in eight ID registers. These ID registers are mapped in addresses 200000h through 200007h. The ID locations read out normally, even after code protection is applied.

Locations 300000h through 30000Dh are reserved for the Configuration bits. These bits select various device options and are described in **Section 5.0 "Configuration Word"**. These Configuration bits read out normally, even after code protection.

Locations 3FFFFEh and 3FFFFFh are reserved for the device ID bits. These bits may be used by the programmer to identify what device type is being programmed and are described in **Section 5.0 "Configuration Word"**. These device ID bits read out normally, even after code protection.

2.3.1 MEMORY ADDRESS POINTER

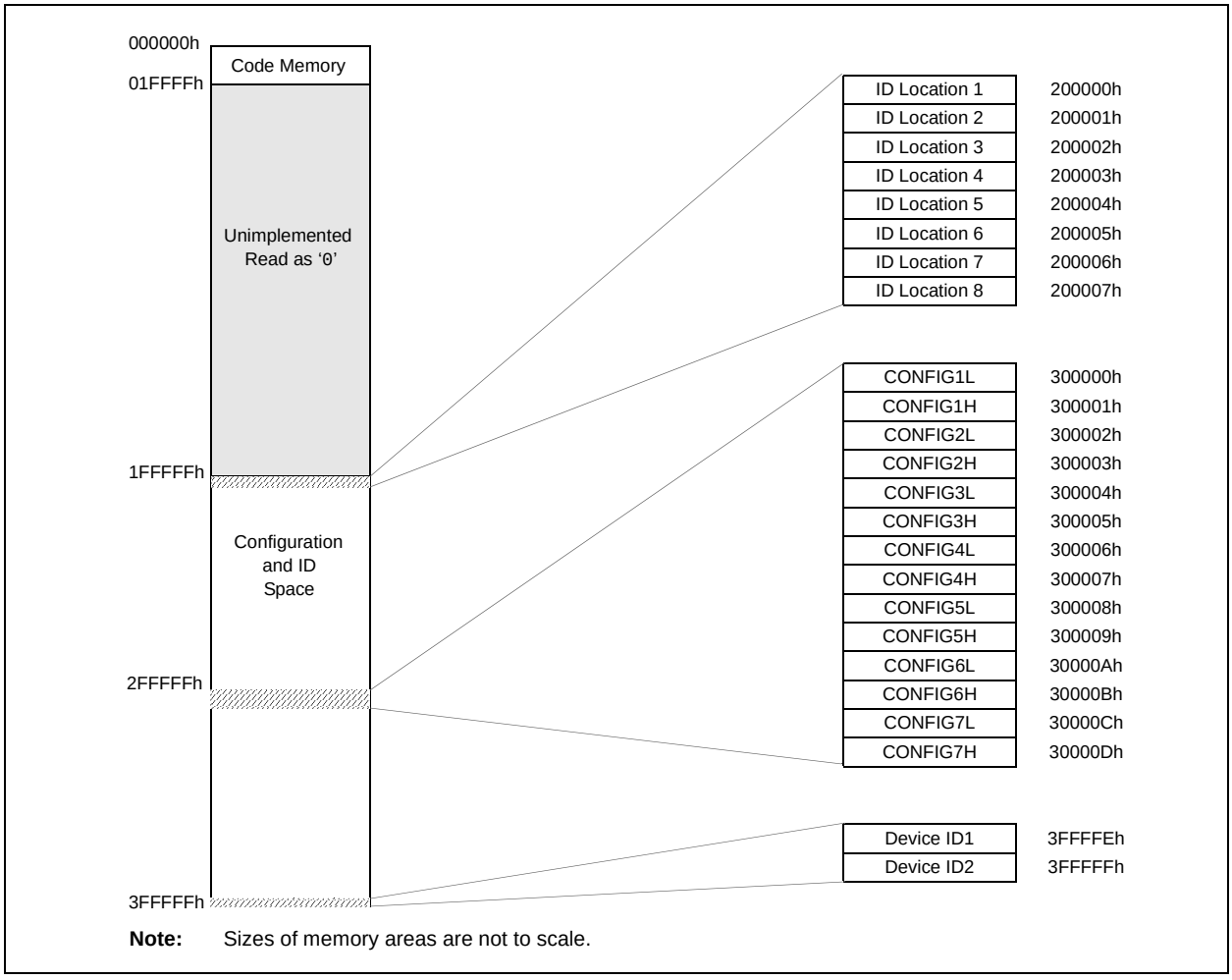
Memory in the address space, 0000000h to 3FFFFFh, is addressed via the Table Pointer register, which is comprised of three Pointer registers:

- TBLPTRU, at RAM address 0FF8h
- TBLPTRH, at RAM address 0FF7h
- TBLPTRL, at RAM address 0FF6h

TBLPTRU	TBLPTRH	TBLPTRL
Addr[21:16]	Addr[15:8]	Addr[7:0]

The 4-bit command, '0000' (core instruction), is used to load the Table Pointer prior to using any read or write operations.

FIGURE 2-10: CONFIGURATION AND ID LOCATIONS FOR PIC18F2XK20/4XK20 DEVICES



Note: Sizes of memory areas are not to scale.

TABLE 2-7: SAMPLE COMMAND SEQUENCE

4-Bit Command	Data Payload	Core Instruction
1101	3C 40	Table Write, post-increment by 2

FIGURE 2-16: TABLE WRITE, POST-INCREMENT TIMING DIAGRAM (1101)

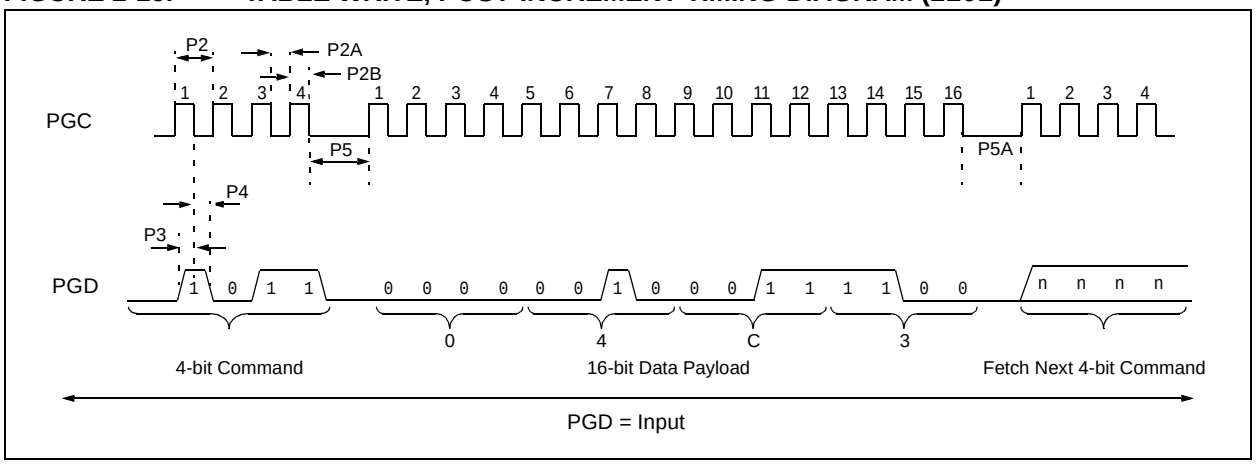


TABLE 3-3: ERASE CODE MEMORY CODE SEQUENCE

4-bit Command	Data Payload	Core Instruction
Step 1: Direct access to code memory and enable writes.		
0000	8E A6	BSF EECON1, EEPGD
0000	9C A6	BCF EECON1, CFGS
0000	84 A6	BSF EECON1, WREN
Step 2: Point to first row in code memory.		
0000	6A F8	CLRF TBLPTRU
0000	6A F7	CLRF TBLPTRH
0000	6A F6	CLRF TBLPTRL
Step 3: Enable erase and erase single row.		
0000	88 A6	BSF EECON1, FREE
0000	82 A6	BSF EECON1, WR
0000	00 00	NOP
0000	00 00	NOP Erase starts on the 4th clock of this instruction
Step 4: Poll WR bit. Repeat until bit is clear.		
0000	50 A6	MOVF EECON1, W, 0
0000	6E F5	MOVWF TABLAT
0000	00 00	NOP
0010	<MSB><LSB>	Shift out data(1)
Step 5: Hold PGC low for time P10.		
Step 6: Repeat step 3 with Address Pointer incremented by 64 until all rows are erased.		
Step 7: Disable writes.		
0000	94 A6	BCF EECON1, WREN

Note 1: See Figure 4-4 for details on shift out data timing.

3.2 Code Memory Programming

Programming code memory is accomplished by first loading data into the write buffer and then initiating a programming sequence. The write and erase buffer sizes shown in Table 3-4 can be mapped to any location of the same size beginning at 000000h. The actual memory write sequence takes the contents of this buffer and programs the proper amount of code memory that contains the Table Pointer.

The programming duration is externally timed and is controlled by PGC. After a Start Programming command is issued (4-bit command, '1111'), a NOP is issued, where the 4th PGC is held high for the duration of the programming time, P9.

After PGC is brought low, the programming sequence is terminated. PGC must be held low for the time specified by parameter P10 to allow high-voltage discharge of the memory array.

The code sequence to program a PIC18F2XK20/4XK20 device is shown in Table 3-5. The flowchart shown in Figure 3-4 depicts the logic necessary to completely write a PIC18F2XK20/4XK20 device. The timing diagram that details the Start Programming command and parameters P9 and P10 is shown in Figure 3-5.

Note:	The TBLPTR register must point to the same region when initiating the programming sequence as it did when the write buffers were loaded.

TABLE 3-4: WRITE AND ERASE BUFFER SIZES

Devices (Arranged by Family)	Write Buffer Size (bytes)	Erase Size (bytes)
PIC18F26K20, PIC18F46K20	64	64
PIC18F24K20, PIC18F25K20, PIC18F44K20, PIC18F45K20	32	64
PIC18F23K20, PIC18F43K20	16	64

TABLE 3-5: WRITE CODE MEMORY CODE SEQUENCE

4-bit Command	Data Payload	Core Instruction
Step 1: Direct access to code memory.		
0000 0000 0000	8E A6 9C A6 84 A6	BSF EECON1, EEPGD BCF EECON1, CFGS BSF EECON1, WREN
Step 2: Point to row to write.		
0000 0000 0000 0000 0000 0000	0E <Addr[21:16]> 6E F8 0E <Addr[15:8]> 6E F7 0E <Addr[7:0]> 6E F6	MOVLW <Addr[21:16]> MOVWF TBLPTRU MOVLW <Addr[15:8]> MOVWF TBLPTRH MOVLW <Addr[7:0]> MOVWF TBLPTRL
Step 3: Load write buffer. Repeat for all but the last two bytes.		
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
Step 4: Load write buffer for last two bytes and start programming.		
1111 0000	<MSB><LSB> 00 00	Write 2 bytes and start programming. NOP - hold PGC high for time P9 and low for time P10.
To continue writing data, repeat steps 2 through 4, where the Address Pointer is incremented by 2 at each iteration of the loop.		

3.3 Data EEPROM Programming

Data EEPROM is accessed one byte at a time via an Address Pointer (register pair EEADRH:EEADR) and a data latch (EEDATA). Data EEPROM is written by loading EEADRH:EEADR with the desired memory location, EEDATA with the data to be written and initiating a memory write by appropriately configuring the EECON1 register. A byte write automatically erases the location and writes the new data (erase-before-write).

When using the EECON1 register to perform a data EEPROM write, both the EEPGD and CFGS bits must be cleared (EECON1<7:6> = 00). The WREN bit must be set (EECON1<2> = 1) to enable writes of any sort and this must be done prior to initiating a write sequence. The write sequence is initiated by setting the WR bit (EECON1<1> = 1).

The write begins on the falling edge of the 24th PGC after the WR bit is set. It ends when the WR bit is cleared by hardware.

After the programming sequence terminates, PGC must be held low for the time specified by parameter P10 to allow high-voltage discharge of the memory array.

FIGURE 3-6: PROGRAM DATA FLOW

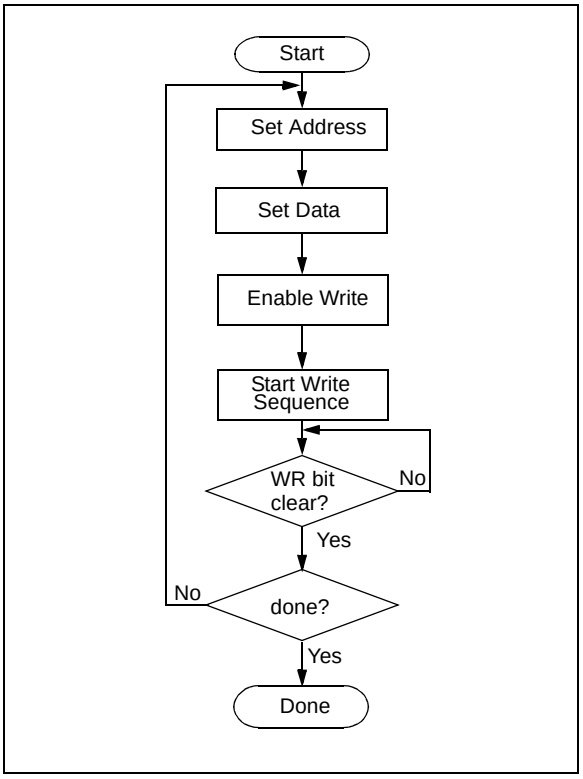
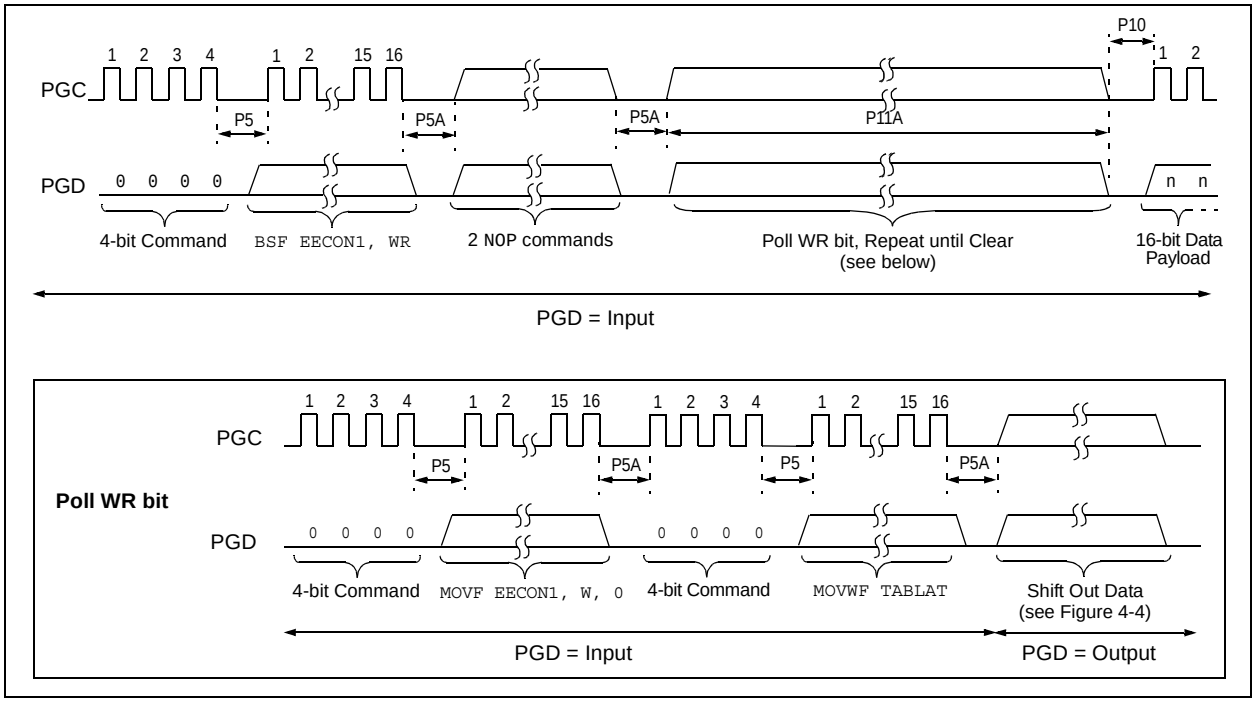


FIGURE 3-7: DATA EEPROM WRITE TIMING DIAGRAM



3.4 ID Location Programming

The ID locations are programmed much like the code memory. The ID registers are mapped in addresses 200000h through 200007h. These locations read out normally even after code protection.

Note:	The user only needs to fill the first 8 bytes of the write buffer in order to write the ID locations.

Table 3-8 demonstrates the code sequence required to write the ID locations.

In order to modify the ID locations, refer to the methodology described in **Section 3.2.1 "Modifying Code Memory"**. As with code memory, the ID locations must be erased before being modified.

When VDD is below the minimum for Bulk Erase operation, ID locations can be cleared with the Row Erase method described in **Section 3.1.3 "ICSP Row Erase"**.

TABLE 3-8: WRITE ID SEQUENCE

4-bit Command	Data Payload	Core Instruction
Step 1: Direct access to code memory.		
0000	8E A6	BSF EECON1, EEPGD
0000	9C A6	BCF EECON1, CFGS
0000	84 A6	BSF EECON1, WREN
Step 2: Set Table Pointer to ID. Load write buffer with 8 bytes and write.		
0000	0E 20	MOVLW 20h
0000	6E F8	MOVWF TBLPTRU
0000	0E 00	MOVLW 00h
0000	6E F7	MOVWF TBLPTRH
0000	0E 00	MOVLW 00h
0000	6E F6	MOVWF TBLPTRL
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
1111	<MSB><LSB>	Write 2 bytes and start programming.
0000	00 00	NOP - hold PGC high for time P9 and low for time P10.

4.0 READING THE DEVICE

4.1 Read Code Memory, ID Locations and Configuration Bits

Code memory is accessed one byte at a time via the 4-bit command, '1001' (table read, post-increment). The contents of memory pointed to by the Table Pointer (TBLPTRU:TBLPTRH:TBLPTRL) are serially output on PGD.

The 4-bit command is shifted in LSb first. The read is executed during the next 8 clocks, then shifted out on PGD during the last 8 clocks, LSb to MSb. A delay of P6 must be introduced after the falling edge of the 8th PGC of the operand to allow PGD to transition from an input to an output. During this time, PGC must be held low (see Figure 4-1). This operation also increments the Table Pointer by one, pointing to the next byte in code memory for the next read.

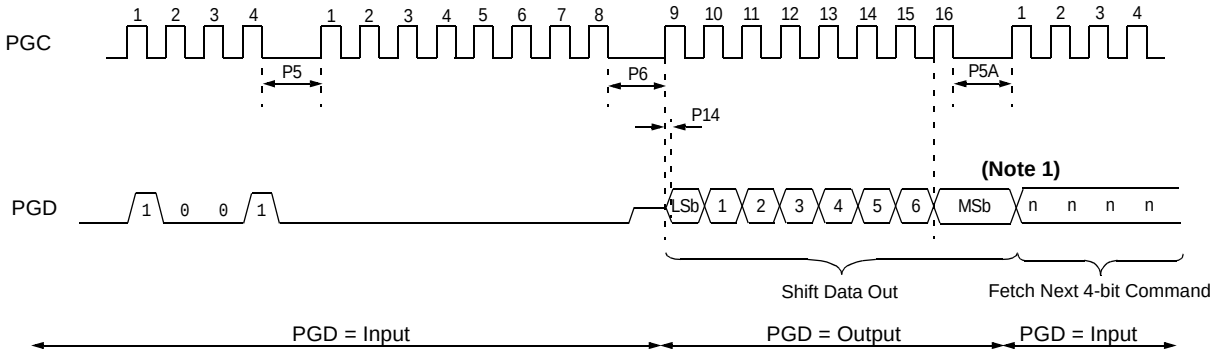
This technique will work to read any memory in the 000000h to 3FFFFFh address space, so it also applies to the reading of the ID and Configuration registers.

Note:	When table read protection is enabled, the first read access to a protected block should be discarded and the read repeated to retrieve valid data. Subsequent reads of the same block can be performed normally.

TABLE 4-1: READ CODE MEMORY SEQUENCE

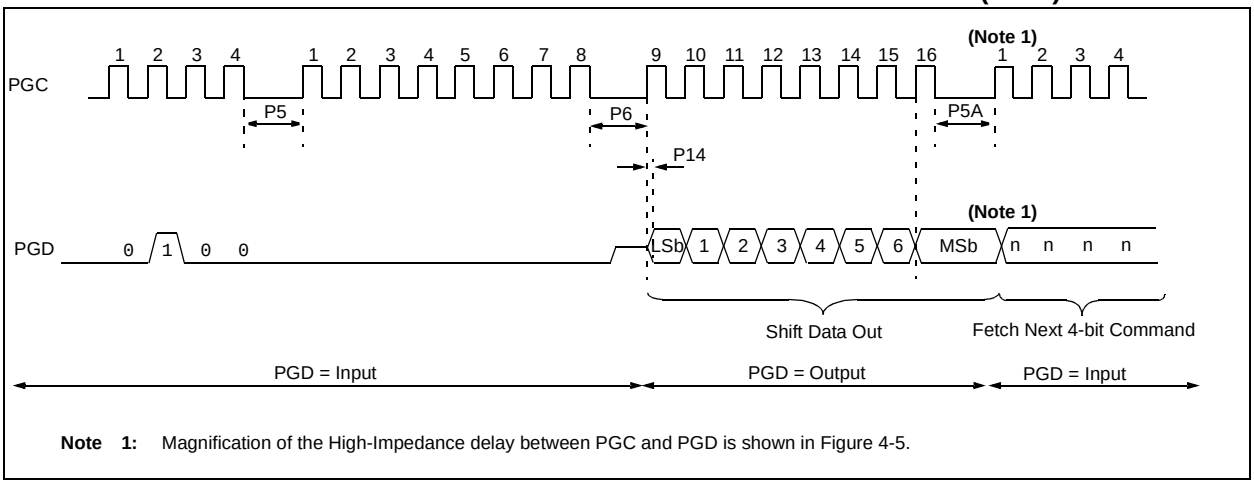
4-bit Command	Data Payload	Core Instruction
Step 1: Set Table Pointer		
0000	0E <Addr[21:16]>	MOVLW Addr[21:16]
0000	6E F8	MOVWF TBLPTRU
0000	0E <Addr[15:8]>	MOVLW <Addr[15:8]>
0000	6E F7	MOVWF TBLPTRH
0000	0E <Addr[7:0]>	MOVLW <Addr[7:0]>
0000	6E F6	MOVWF TBLPTRL
Step 2: Read memory and then shift out on PGD, LSb to MSb		
1001	00 00	TBLRD *+

FIGURE 4-1: TABLE READ POST-INCREMENT INSTRUCTION TIMING DIAGRAM (1001)



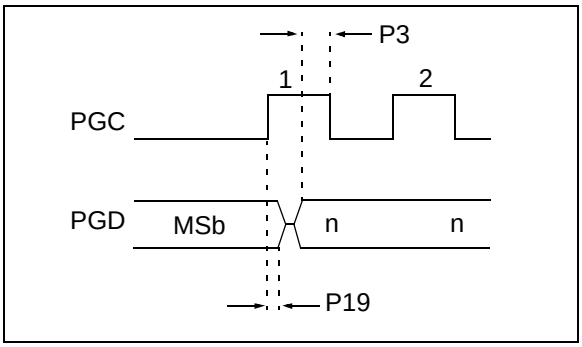
Note 1: Magnification of the high-impedance delay between PGC and PGD is shown in Figure 4-6.

FIGURE 4-4: SHIFT OUT DATA HOLDING REGISTER TIMING DIAGRAM (0010)



Note 1: Magnification of the High-Impedance delay between PGC and PGD is shown in Figure 4-5.

FIGURE 4-5: HIGH-IMPEDANCE DELAY



4.5 Verify Data EEPROM

A data EEPROM address may be read via a sequence of core instructions (4-bit command, '0000') and then output on PGD via the 4-bit command, '0010' (TABLAT register). The result may then be immediately compared to the appropriate data in the programmer's memory for verification. Refer to **Section 4.4 "Read Data EEPROM Memory"** for implementation details of reading data EEPROM.

4.6 Blank Check

The term "Blank Check" means to verify that the device has no programmed memory cells. All memories must be verified: code memory, data EEPROM, ID locations and Configuration bits. The device ID registers (3FFFFEh:3FFFFFh) should be ignored.

A "blank" or "erased" memory cell will read as a '1'. Therefore, Blank Checking a device merely means to verify that all bytes read as FFh except the Configuration bits. Unused (reserved) Configuration bits will read '0' (programmed). Refer to Table 5-1 for blank configuration expect data for the various PIC18F2XK20/4XK20 devices.

Given that Blank Checking is merely code and data EEPROM verification with FFh expect data, refer to **Section 4.4 "Read Data EEPROM Memory"** and **Section 4.2 "Verify Code Memory and ID Locations"** for implementation details.

FIGURE 4-6: BLANK CHECK FLOW

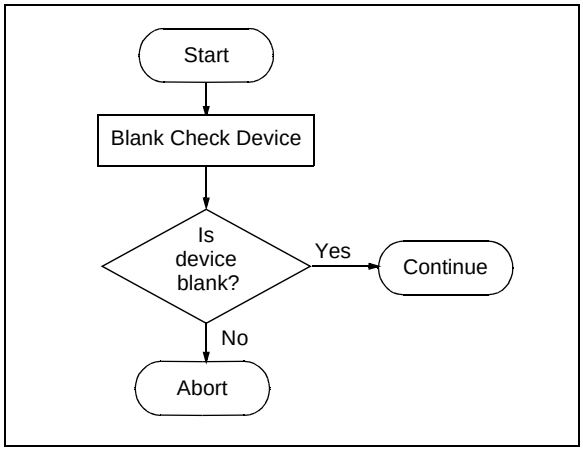


TABLE 5-3: PIC18F2XK20/4XK20 BIT DESCRIPTIONS (CONTINUED)

Bit Name	Configuration Words	Description
WDTEN	CONFIG2H	Watchdog Timer Enable bit 1 = WDT enabled 0 = WDT disabled (control is placed on SWDTEN bit)
MCLRE	CONFIG3H	$\overline{\text{MCLR}}$ Pin Enable bit 1 = $\overline{\text{MCLR}}$ pin enabled, RE3 input pin disabled 0 = RE3 input pin enabled, $\overline{\text{MCLR}}$ pin disabled
HFOFST	CONFIG3H	HFINTOSC Fast Start 1 = HFINTOSC output is not delayed 0 = HFINTOSC output is delayed until oscillator is stable (IOFS = 1)
LPT1OSC	CONFIG3H	Low-Power Timer1 Oscillator Enable bit 1 = Timer1 configured for low-power operation 0 = Timer1 configured for higher power operation
PBADEN	CONFIG3H	PORTB A/D Enable bit 1 = PORTB A/D<4:0> pins are configured as analog input channels on Reset 0 = PORTB A/D<4:0> pins are configured as digital I/O on Reset
CCP2MX	CONFIG3H	CCP2 MUX bit 1 = CCP2 input/output is multiplexed with RC1 0 = CCP2 input/output is multiplexed with RB3
DEBUG	CONFIG4L	Background Debugger Enable bit 1 = Background debugger disabled, RB6 and RB7 configured as general purpose I/O pins 0 = Background debugger enabled, RB6 and RB7 are dedicated to In-Circuit Debug
XINST	CONFIG4L	Extended Instruction Set Enable bit 1 = Instruction set extension and Indexed Addressing mode enabled 0 = Instruction set extension and Indexed Addressing mode disabled (Legacy mode)
LVP	CONFIG4L	Low-Voltage Programming Enable bit 1 = Low-Voltage Programming enabled, RB5 is the PGM pin 0 = Low-Voltage Programming disabled, RB5 is an I/O pin
STVREN	CONFIG4L	Stack Overflow/Underflow Reset Enable bit 1 = Reset on stack overflow/underflow enabled 0 = Reset on stack overflow/underflow disabled

.

TABLE 5-3: PIC18F2XK20/4XK20 BIT DESCRIPTIONS (CONTINUED)

Bit Name	Configuration Words	Description
CP3	CONFIG5L	Code Protection bits (Block 3 code memory area) 1 = Block 3 is not code-protected 0 = Block 3 is code-protected
CP2	CONFIG5L	Code Protection bits (Block 2 code memory area) 1 = Block 2 is not code-protected 0 = Block 2 is code-protected
CP1	CONFIG5L	Code Protection bits (Block 1 code memory area) 1 = Block 1 is not code-protected 0 = Block 1 is code-protected
CP0	CONFIG5L	Code Protection bits (Block 0 code memory area) 1 = Block 0 is not code-protected 0 = Block 0 is code-protected
CPD	CONFIG5H	Code Protection bits (Data EEPROM) 1 = Data EEPROM is not code-protected 0 = Data EEPROM is code-protected
CPB	CONFIG5H	Code Protection bits (Boot Block memory area) 1 = Boot Block is not code-protected 0 = Boot Block is code-protected
WRT3	CONFIG6L	Write Protection bits (Block 3 code memory area) 1 = Block 3 is not write-protected 0 = Block 3 is write-protected
WRT2	CONFIG6L	Write Protection bits (Block 2 code memory area) 1 = Block 2 is not write-protected 0 = Block 2 is write-protected
WRT1	CONFIG6L	Write Protection bits (Block 1 code memory area) 1 = Block 1 is not write-protected 0 = Block 1 is write-protected
WRT0	CONFIG6L	Write Protection bits (Block 0 code memory area) 1 = Block 0 is not write-protected 0 = Block 0 is write-protected
WRTD	CONFIG6H	Write Protection bit (Data EEPROM) 1 = Data EEPROM is not write-protected 0 = Data EEPROM is write-protected
WRTB	CONFIG6H	Write Protection bit (Boot Block memory area) 1 = Boot Block is not write-protected 0 = Boot Block is write-protected
WRTC	CONFIG6H	Write Protection bit (Configuration registers) 1 = Configuration registers are not write-protected 0 = Configuration registers are write-protected

.

5.3 Single-Supply ICSP Programming

The LVP bit in Configuration register, CONFIG4L, enables Single-Supply (Low-Voltage) ICSP Programming. The LVP bit defaults to a '1' (enabled) from the factory.

If Single-Supply Programming mode is not used, the LVP bit can be programmed to a '0' and RB5/PGM becomes a digital I/O pin. However, the LVP bit may only be programmed by entering the High-Voltage ICSP mode, where MCLR/VPP/RE3 is raised to VIHH. Once the LVP bit is programmed to a '0', only the High-Voltage ICSP mode is available and only the High-Voltage ICSP mode can be used to program the device.

Note 1: The High-Voltage ICSP mode is always available, regardless of the state of the LVP bit, by applying VIHH to the MCLR/VPP/RE3 pin.

2: While in Low-Voltage ICSP mode, the RB5 pin can no longer be used as a general purpose I/O.

5.4 Embedding Configuration Word Information in the HEX File

To allow portability of code, a PIC18F2XK20/4XK20 programmer is required to read the Configuration Word locations from the hex file. If Configuration Word information is not present in the hex file, then a simple warning message should be issued. Similarly, while saving a hex file, all Configuration Word information must be included. An option to not include the Configuration Word information may be provided. When embedding Configuration Word information in the hex file, it should start at address 300000h.

Microchip Technology Inc. feels strongly that this feature is important for the benefit of the end customer.

5.5 Embedding Data EEPROM Information In the HEX File

To allow portability of code, a PIC18F2XK20/4XK20 programmer is required to read the data EEPROM information from the hex file. If data EEPROM information is not present, a simple warning message should be issued. Similarly, when saving a hex file, all data EEPROM information must be included. An option to not include the data EEPROM information may be provided. When embedding data EEPROM information in the hex file, it should start at address F00000h.

Microchip Technology Inc. believes that this feature is important for the benefit of the end customer.

5.6 Checksum Computation

The checksum is calculated by summing the following:

- The contents of all code memory locations
- The Configuration Word, appropriately masked
- ID locations (Only if any portion of program memory is code-protected)

The Least Significant 16 bits of this sum are the checksum.

Code protection limits access to program memory by both external programmer (code-protect) and code execution (table read protect). The ID locations, when included in a code protected checksum, contain the checksum of an unprotected part. The unprotected checksum is distributed: one nibble per ID location. Each nibble is right justified.

Table 5-4 describes how to calculate the checksum for each device.

Note: The checksum calculation differs depending on the code-protect setting. Since the code memory locations read out differently depending on the code-protect setting, the table describes how to manipulate the actual code memory values to simulate the values that would be read from a protected device. When calculating a checksum by reading a device, the entire code memory can simply be read and summed. The Configuration Word and ID locations can always be read.

TABLE 5-4: CHECKSUM COMPUTATION (CONTINUED)

Device	Code-Protect	Checksum	Blank Value	0xAA at 0 and Max Address
PIC18FX5K20	None	SUM[0000:07FF]+SUM[0800:1FFF]+SUM[2000:3FFF]+ SUM[4000:5FFF]+SUM[6000:7FFF]+(CONFIG1L & 00h)+ (CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+ (CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+ (CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+ (CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+ (CONFIG7H & 40h)	8362h	82B8h
	Boot Block	SUM[0800:1FFF]+SUM[2000:3FFF]+SUM[4000:5FFF]+SUM[6000:7FFF]+ (CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+ (CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+ (CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+ (CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+ (CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	8B35h	8AEAh
	Boot/ Block 0/ Block 1	SUM[4000:5FFF]+SUM[6000:7FFF]+(CONFIG1L & 00h)+ (CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+ (CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+ (CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+ (CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+ (CONFIG7H & 40h)+SUM_ID	C332h	C2E7h
	All	(CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+ (CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+ (CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+ (CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+ (CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	0326h	0330h

Legend:

Item		Description
CONFIGx	=	Configuration Word
SUM[a:b]	=	Sum of locations, a to b inclusive
SUM_ID	=	Byte-wise sum of lower four bits of all customer ID locations
+	=	Addition
&	=	Bit-wise AND

TABLE 5-4: CHECKSUM COMPUTATION (CONTINUED)

Device	Code-Protect	Checksum	Blank Value	0xAA at 0 and Max Address
PIC18FX6K20	None	SUM[0000:07FF]+SUM[0800:3FFF]+SUM[4000:7FFF]+ SUM[8000:BFFF]+SUM[C000:FFFF]+(CONFIG1L & 00h)+ (CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+ (CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+ (CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+ (CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+ (CONFIG7H & 40h)	0362h	02B8h
	Boot Block	SUM[0800:3FFF]+SUM[4000:7FFF]+SUM[8000:BFFF]+SUM[C000:FFFF]+ (CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+ (CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+ (CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+ (CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+ (CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	0B2Dh	0AE2h
	Boot/ Block 0/ Block 1	SUM[3000:BFFF]+SUM[C000:FFFF]+(CONFIG1L & 00h)+ (CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+ (CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+ (CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+ (CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+ (CONFIG7H & 40h)+SUM_ID	832Ah	82DFh
	All	(CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+ (CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+ (CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+ (CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+ (CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	031Eh	0328h

Legend:

Item		Description
CONFIGx	=	Configuration Word
SUM[a:b]	=	Sum of locations, a to b inclusive
SUM_ID	=	Byte-wise sum of lower four bits of all customer ID locations
+	=	Addition
&	=	Bit-wise AND

6.0 AC/DC CHARACTERISTICS TIMING REQUIREMENTS FOR PROGRAM/ VERIFY TEST MODE (CONTINUED)

Standard Operating Conditions
Operating Temperature: 25°C is recommended

Param No.	Sym.	Characteristic	Min.	Max.	Units	Conditions
P12	THLD2	Input Data Hold Time from $\overline{MCLR}$/VPP/RE3 ↑	2	—	µs	
P13	TSET2	VDD ↑ Setup Time to $\overline{MCLR}$/VPP/RE3 ↑	100	—	ns	
P14	TVALID	Data Out Valid from PGC ↑	10	—	ns	
P15	TSET3	PGM ↑ Setup Time to $\overline{MCLR}$/VPP/RE3 ↑	2	—	µs	
P16	TDLY8	Delay between Last PGC ↓ and $\overline{MCLR}$/VPP/RE3 ↓	0	—	s	
P17	THLD3	$\overline{MCLR}$/VPP/RE3 ↓ to VDD ↓	—	100	ns	
P18	THLD4	$\overline{MCLR}$/VPP/RE3 ↓ to PGM ↓	0	—	s	
P19	THIZ	Delay from PGC ↑ to PGD High-Z	3	10	nS	
P20	TPPDP	Hold time after VPP changes	5	—	µs	

Note 1: Do not allow excess time when transitioning $\overline{MCLR}$ between VIL and VIHH; this can cause spurious program executions to occur. The maximum transition time is:
1 TCY + TPWRT (if enabled) + 1024 TOSC (for LP, HS, HS/PLL and XT modes only) + 2 ms (for HS/PLL mode only) + 1.5 µs (for EC mode only) where TCY is the instruction cycle time, TPWRT is the Power-up Timer period and TOSC is the oscillator period. For specific values, refer to the Electrical Characteristics section of the device data sheet for the particular device.

NOTES:

Note the following details of the code protection feature on Microchip devices:

- Microchip products meet the specification contained in their particular Microchip Data Sheet.
- Microchip believes that its family of products is one of the most secure families of its kind on the market today, when used in the intended manner and under normal conditions.
- There are dishonest and possibly illegal methods used to breach the code protection feature. All of these methods, to our knowledge, require using the Microchip products in a manner outside the operating specifications contained in Microchip's Data Sheets. Most likely, the person doing so is engaged in theft of intellectual property.
- Microchip is willing to work with the customer who is concerned about the integrity of their code.
- Neither Microchip nor any other semiconductor manufacturer can guarantee the security of their code. Code protection does not mean that we are guaranteeing the product as "unbreakable."

Code protection is constantly evolving. We at Microchip are committed to continuously improving the code protection features of our products. Attempts to break Microchip's code protection feature may be a violation of the Digital Millennium Copyright Act. If such acts allow unauthorized access to your software or other copyrighted work, you may have a right to sue for relief under that Act.

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