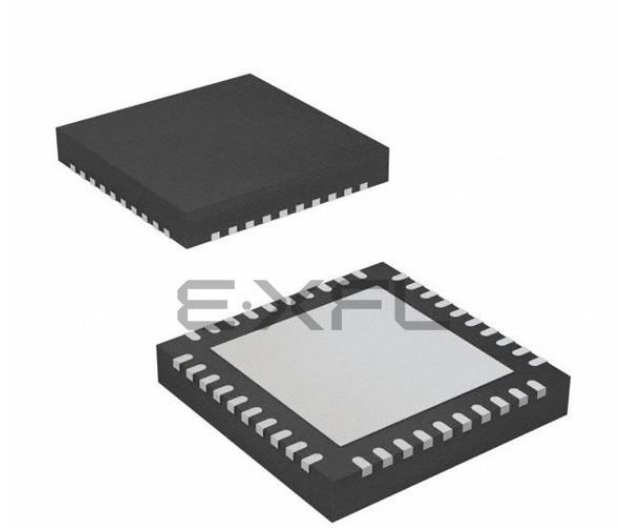

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Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	64MHz
Connectivity	I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, HLVD, POR, PWM, WDT
Number of I/O	35
Program Memory Size	64KB (32K x 16)
Program Memory Type	FLASH
EEPROM Size	1K x 8
RAM Size	3.8K x 8
Voltage - Supply (Vcc/Vdd)	1.8V ~ 3.6V
Data Converters	A/D 14x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	40-UFQFN Exposed Pad
Supplier Device Package	40-UQFN (5x5)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18f46k20t-i-mv

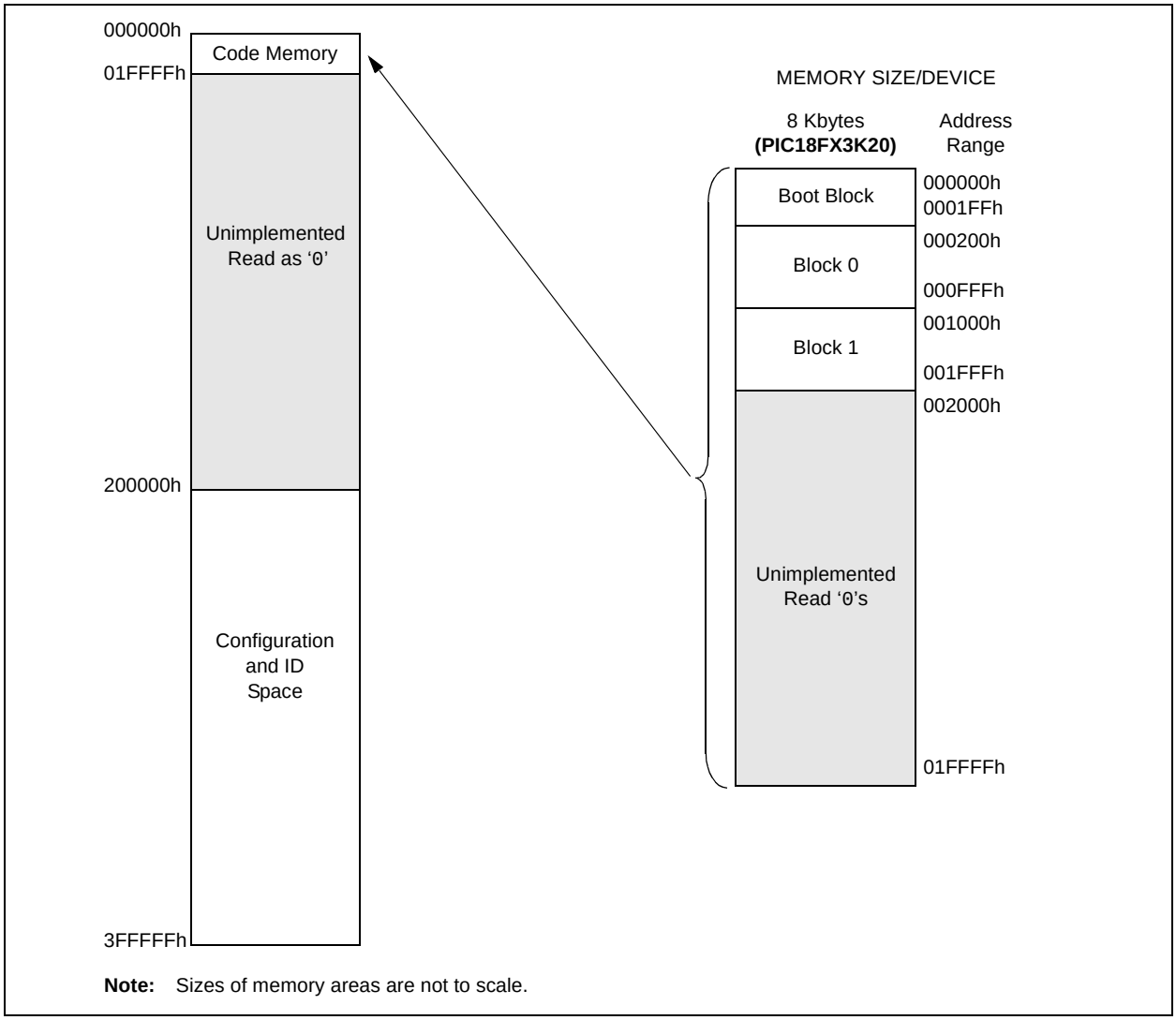
2.3 Memory Maps

For the PIC18FX3K20 devices, the code memory space extends from 0000h to 01FFFh (8 Kbytes) in two 4-Kbyte blocks. Addresses 0000h through 01FFh, however, define a “Boot Block” region that is treated separately from Block 0. All of these blocks define code protection boundaries within the code memory space.

TABLE 2-2: IMPLEMENTATION OF CODE MEMORY

Device	Code Memory Size (Bytes)
PIC18F23K20	000000h-001FFFh (8K)
PIC18F43K20	

FIGURE 2-6: MEMORY MAP AND THE CODE MEMORY SPACE FOR PIC18FX3K20 DEVICES



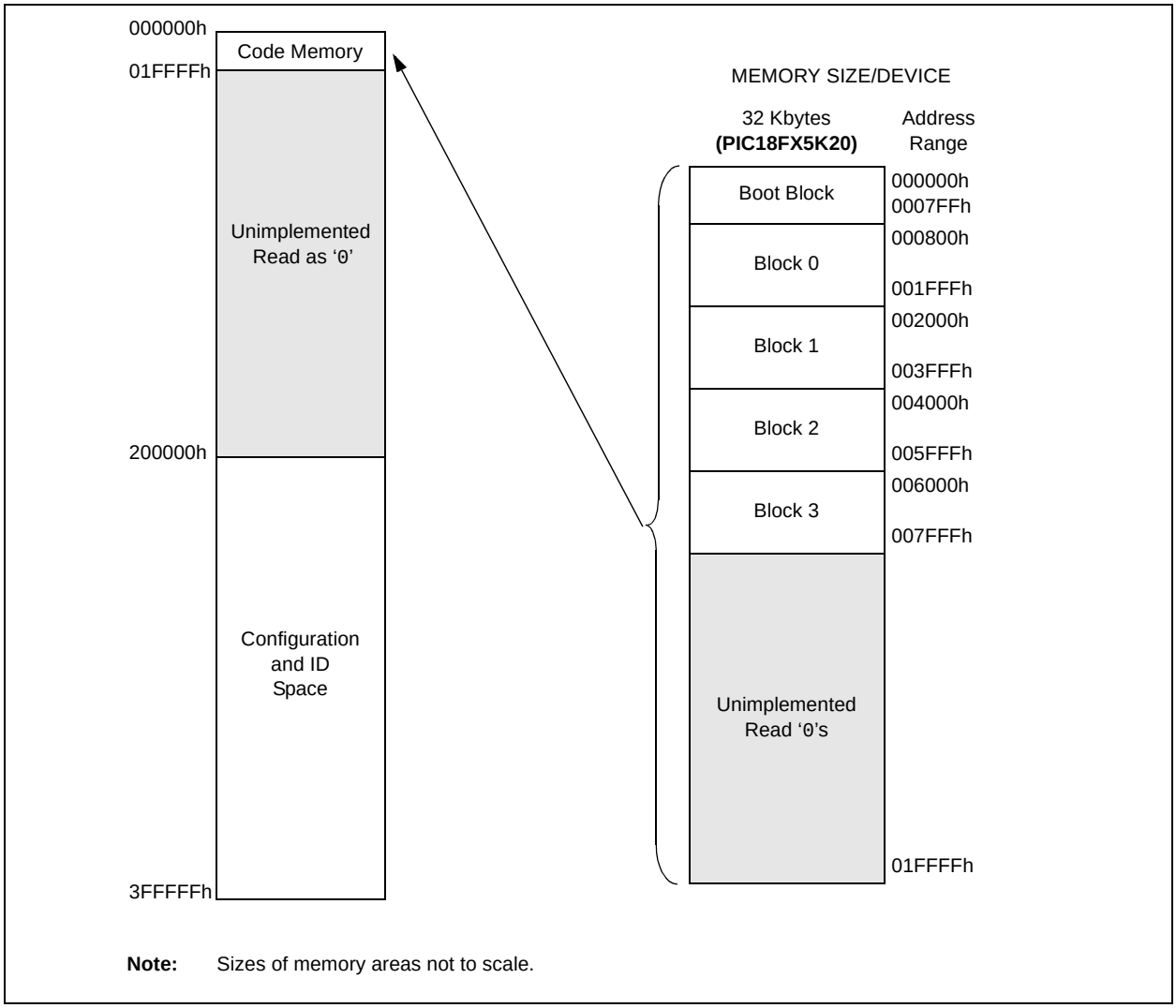
Note: Sizes of memory areas are not to scale.

For PIC18FX5K20 devices, the code memory space extends from 000000h to 007FFFh (32 Kbytes) in four 8-Kbyte blocks. Addresses 000000h through 0007FFh, however, define a “Boot Block” region that is treated separately from Block 0. All of these blocks define code protection boundaries within the code memory space.

TABLE 2-4: IMPLEMENTATION OF CODE MEMORY

Device	Code Memory Size (Bytes)
PIC18F25K20	000000h-007FFFh (32K)
PIC18F45K20	

FIGURE 2-8: MEMORY MAP AND THE CODE MEMORY SPACE FOR PIC18FX5K20 DEVICES



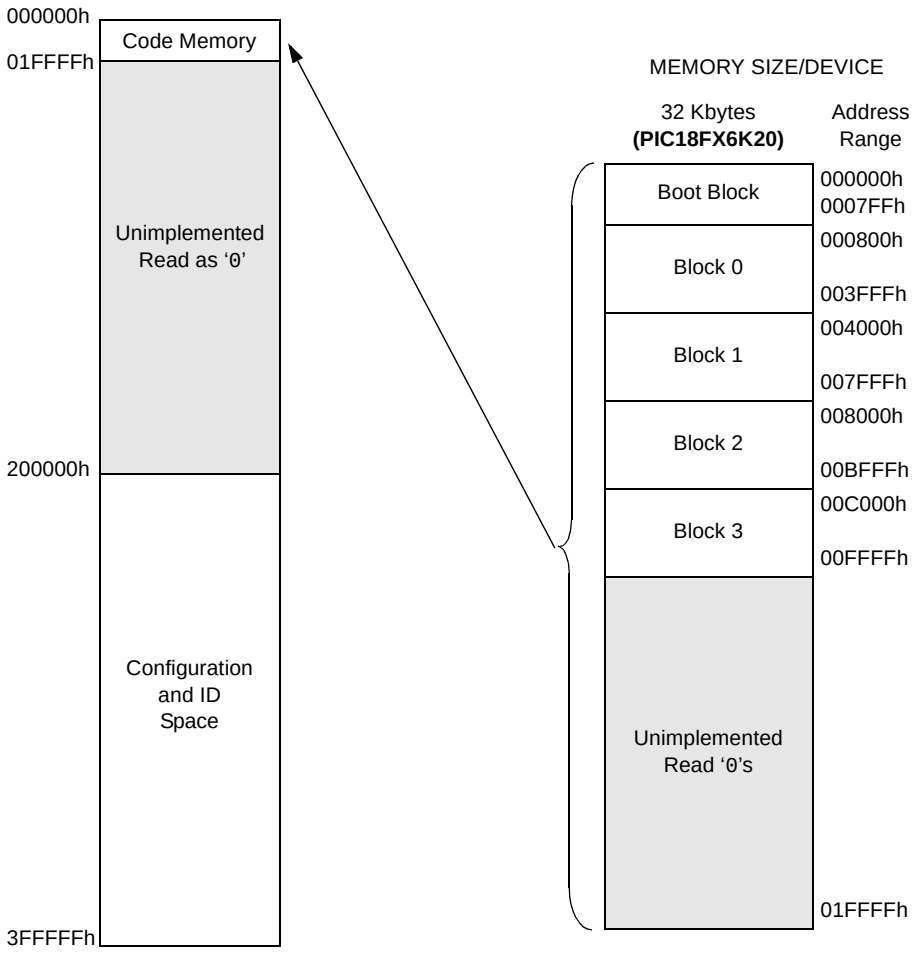
Note: Sizes of memory areas not to scale.

For PIC18FX6K20 devices, the code memory space extends from 000000h to 00FFFFh (64 Kbytes) in four 16-Kbyte blocks. Addresses 000000h through 0007FFh, however, define a “Boot Block” region that is treated separately from Block 0. All of these blocks define code protection boundaries within the code memory space.

TABLE 2-5: IMPLEMENTATION OF CODE MEMORY

Device	Code Memory Size (Bytes)
PIC18F26K20	000000h-00FFFFh (64K)
PIC18F46K20	

FIGURE 2-9: MEMORY MAP AND THE CODE MEMORY SPACE FOR PIC18FX6K20 DEVICES

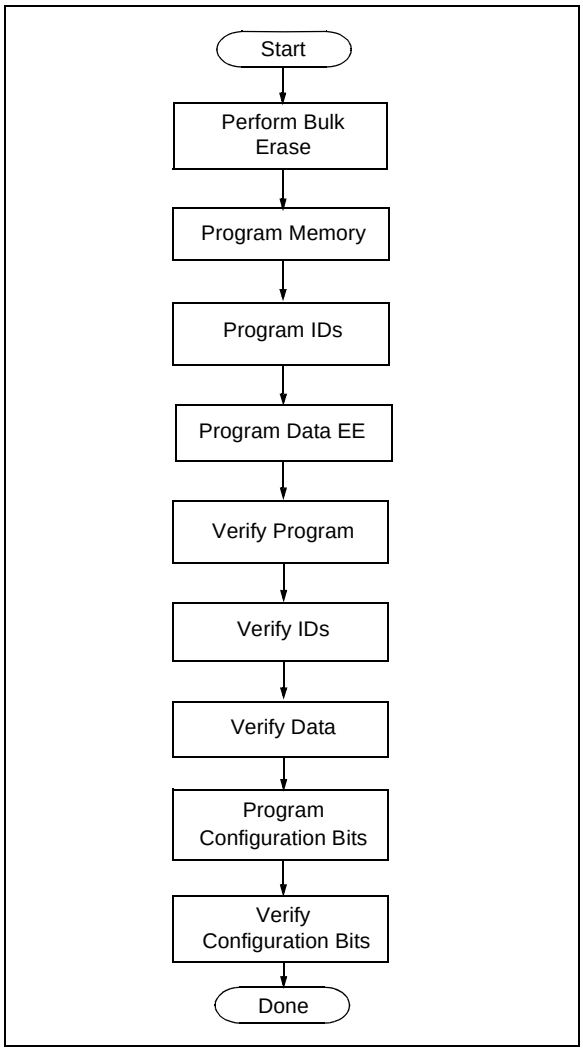


Note: Sizes of memory areas not to scale.

2.4 High-Level Overview of the Programming Process

Figure 2-11 shows the high-level overview of the programming process. First, a Bulk Erase is performed. Next, the code memory, ID locations and data EEPROM are programmed. These memories are then verified to ensure that programming was successful. If no errors are detected, the Configuration bits are then programmed and verified.

FIGURE 2-11: HIGH-LEVEL PROGRAMMING FLOW



2.5 Entering and Exiting High-Voltage ICSP Program/Verify Mode

As shown in Figure 2-12, the High-Voltage ICSP Program/Verify mode is entered by holding PGC and PGD low and then raising $\overline{MCLR}/V_{PP}$/RE3 to V_{IHH} (high voltage). Once in this mode, the code memory, data EEPROM, ID locations and Configuration bits can be accessed and programmed in serial fashion. Figure 2-13 shows the exit sequence.

The sequence that enters the device into the Program/Verify mode places all unused I/Os in the high-impedance state.

FIGURE 2-12: ENTERING HIGH-VOLTAGE PROGRAM/VERIFY MODE

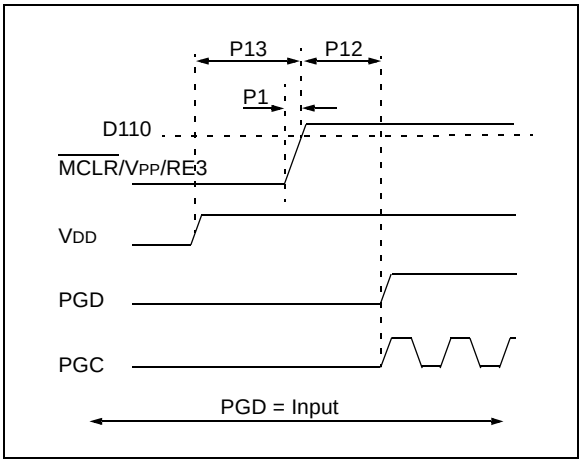


FIGURE 2-13: EXITING HIGH-VOLTAGE PROGRAM/VERIFY MODE

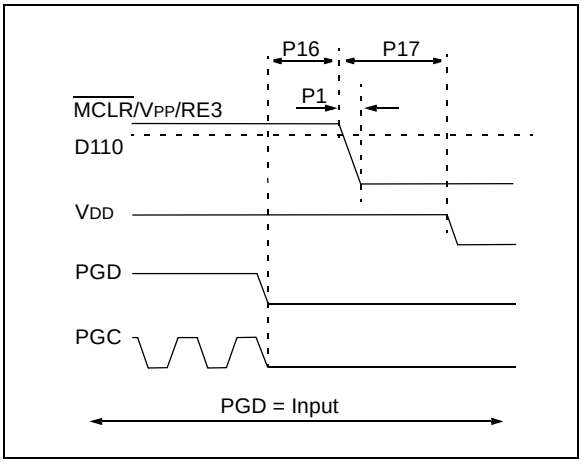


TABLE 2-7: SAMPLE COMMAND SEQUENCE

4-Bit Command	Data Payload	Core Instruction
1101	3C 40	Table Write, post-increment by 2

FIGURE 2-16: TABLE WRITE, POST-INCREMENT TIMING DIAGRAM (1101)

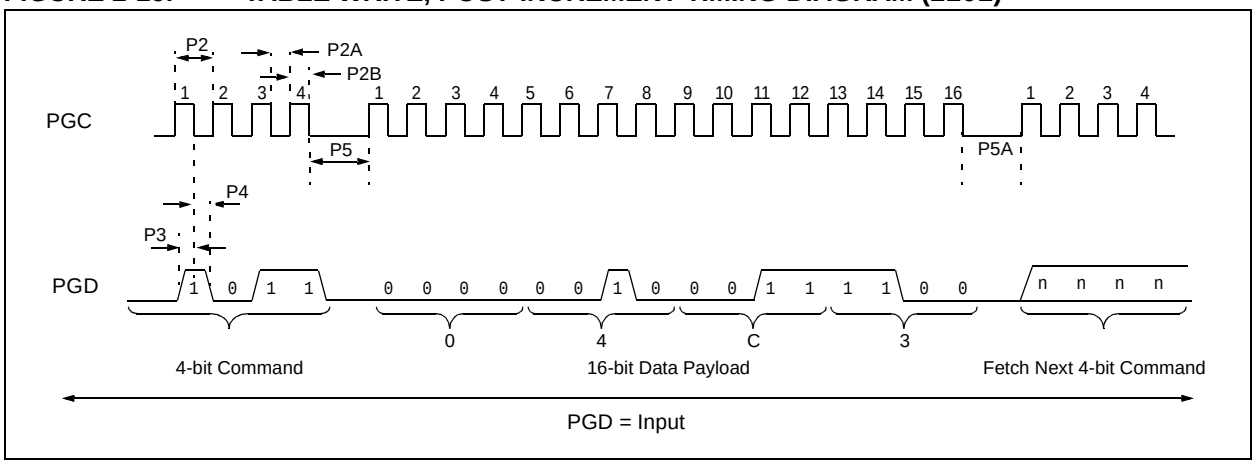
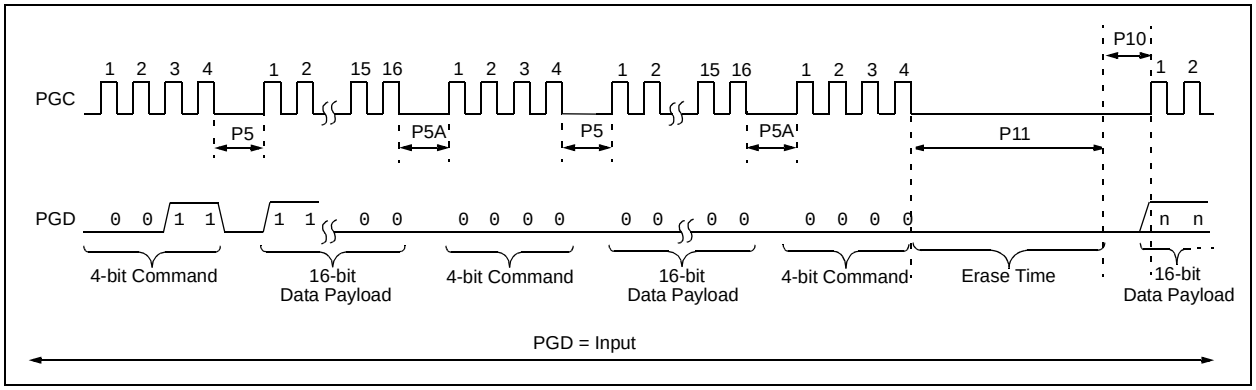


FIGURE 3-2: BULK ERASE TIMING DIAGRAM



3.1.2 LOW-VOLTAGE ICSP BULK ERASE

When using low-voltage ICSP, the part must be supplied by the voltage specified in parameter D111 if a Bulk Erase is to be executed. All other Bulk Erase details as described above apply.

If it is determined that a program memory erase must be performed at a supply voltage below the Bulk Erase limit, refer to the erase methodology described in **Section 3.1.3 "ICSP Row Erase"** and **Section 3.2.1 "Modifying Code Memory"**.

If it is determined that a data EEPROM erase must be performed at a supply voltage below the Bulk Erase limit, follow the methodology described in **Section 3.3 "Data EEPROM Programming"** and write '1's to the array.

3.1.3 ICSP ROW ERASE

Regardless of whether high or low-voltage ICSP is used, it is possible to erase one row (64 bytes of data), provided the block is not code or write-protected. Rows are located at static boundaries beginning at program memory address 000000h, extending to the internal program memory limit (see **Section 2.3 "Memory Maps"**).

The Row Erase duration is self-timed. After the WR bit in EECON1 is set, two NOPs are issued. Erase starts upon the 4th PGC of the second NOP. It ends when the WR bit is cleared by hardware.

The code sequence to Row Erase a PIC18F2XK20/4XK20 device is shown in Table 3-3. The flowchart shown in Figure 3-3 depicts the logic necessary to completely erase a PIC18F2XK20/4XK20 device. The timing diagram for Row Erase is identical to the data EEPROM write timing shown in Figure 3-7.

Note:	The TBLPTR register can point at any byte within the row intended for erase.

TABLE 3-3: ERASE CODE MEMORY CODE SEQUENCE

4-bit Command	Data Payload	Core Instruction
Step 1: Direct access to code memory and enable writes.		
0000	8E A6	BSF EECON1, EEPGD
0000	9C A6	BCF EECON1, CFGS
0000	84 A6	BSF EECON1, WREN
Step 2: Point to first row in code memory.		
0000	6A F8	CLRF TBLPTRU
0000	6A F7	CLRF TBLPTRH
0000	6A F6	CLRF TBLPTRL
Step 3: Enable erase and erase single row.		
0000	88 A6	BSF EECON1, FREE
0000	82 A6	BSF EECON1, WR
0000	00 00	NOP
0000	00 00	NOP Erase starts on the 4th clock of this instruction
Step 4: Poll WR bit. Repeat until bit is clear.		
0000	50 A6	MOVF EECON1, W, 0
0000	6E F5	MOVWF TABLAT
0000	00 00	NOP
0010	<MSB><LSB>	Shift out data(1)
Step 5: Hold PGC low for time P10.		
Step 6: Repeat step 3 with Address Pointer incremented by 64 until all rows are erased.		
Step 7: Disable writes.		
0000	94 A6	BCF EECON1, WREN

Note 1: See Figure 4-4 for details on shift out data timing.

3.2.1 MODIFYING CODE MEMORY

The previous programming example assumed that the device has been Bulk Erased prior to programming (see **Section 3.1.1 "High-Voltage ICSP Bulk Erase"**). It may be the case, however, that the user wishes to modify only a section of an already programmed device.

The appropriate number of bytes required for the erase buffer must be read out of code memory (as described in **Section 4.2 "Verify Code Memory and ID Locations"**) and buffered. Modifications can be made on this buffer. Then, the block of code memory that was read out must be erased and rewritten with the modified data.

The WREN bit must be set if the WR bit in EECON1 is used to initiate a write sequence.

TABLE 3-6: MODIFYING CODE MEMORY

4-bit Command	Data Payload	Core Instruction
Step 1: Direct access to code memory.		
0000	8E A6	BSF EECON1, EEPGD
0000	9C A6	BCF EECON1, CFGS
Step 2: Read code memory into buffer (**Section 4.1 "Read Code Memory, ID Locations and Configuration Bits"**).		
Step 3: Set the Table Pointer for the block to be erased.		
0000	0E <Addr[21:16]>	MOVLW <Addr[21:16]>
0000	6E F8	MOVWF TBLPTRU
0000	0E <Addr[8:15]>	MOVLW <Addr[8:15]>
0000	6E F7	MOVWF TBLPTRH
0000	0E <Addr[7:0]>	MOVLW <Addr[7:0]>
0000	6E F6	MOVWF TBLPTRL
Step 4: Enable memory writes and setup an erase.		
0000	84 A6	BSF EECON1, WREN
0000	88 A6	BSF EECON1, FREE
Step 5: Initiate erase.		
0000	88 A6	BSF EECON1, FREE
0000	82 A6	BSF EECON1, WR
0000	00 00	NOP
0000	00 00	NOP Erase starts on the 4th clock of this instruction
Step 6: Poll WR bit. Repeat until bit is clear.		
0000	50 A6	MOVF EECON1, W, 0
0000	6E F5	MOVWF TABLAT
0000	00 00	NOP
0000	<MSB><LSB>	Shift out data(1)
Step 7: Load write buffer. The correct bytes will be selected based on the Table Pointer.		
0000	0E <Addr[21:16]>	MOVLW <Addr[21:16]>
0000	6E F8	MOVWF TBLPTRU
0000	0E <Addr[8:15]>	MOVLW <Addr[8:15]>
0000	6E F7	MOVWF TBLPTRH
0000	0E <Addr[7:0]>	MOVLW <Addr[7:0]>
0000	6E F6	MOVWF TBLPTRL
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
•	•	
•	•	Repeat as many times as necessary to fill the write buffer
•	•	Write 2 bytes and start programming.
1111	<MSB><LSB>	NOP - hold PGC high for time P9 and low for time P10.
0000	00 00	
To continue modifying data, repeat Steps 2 through 6, where the Address Pointer is incremented by the appropriate number of bytes (see Table 3-4) at each iteration of the loop. The write cycle must be repeated enough times to completely rewrite the contents of the erase buffer.		
Step 8: Disable writes.		
0000	94 A6	BCF EECON1, WREN

3.3 Data EEPROM Programming

Data EEPROM is accessed one byte at a time via an Address Pointer (register pair EEADRH:EEADR) and a data latch (EEDATA). Data EEPROM is written by loading EEADRH:EEADR with the desired memory location, EEDATA with the data to be written and initiating a memory write by appropriately configuring the EECON1 register. A byte write automatically erases the location and writes the new data (erase-before-write).

When using the EECON1 register to perform a data EEPROM write, both the EEPGD and CFGS bits must be cleared (EECON1<7:6> = 00). The WREN bit must be set (EECON1<2> = 1) to enable writes of any sort and this must be done prior to initiating a write sequence. The write sequence is initiated by setting the WR bit (EECON1<1> = 1).

The write begins on the falling edge of the 24th PGC after the WR bit is set. It ends when the WR bit is cleared by hardware.

After the programming sequence terminates, PGC must be held low for the time specified by parameter P10 to allow high-voltage discharge of the memory array.

FIGURE 3-6: PROGRAM DATA FLOW

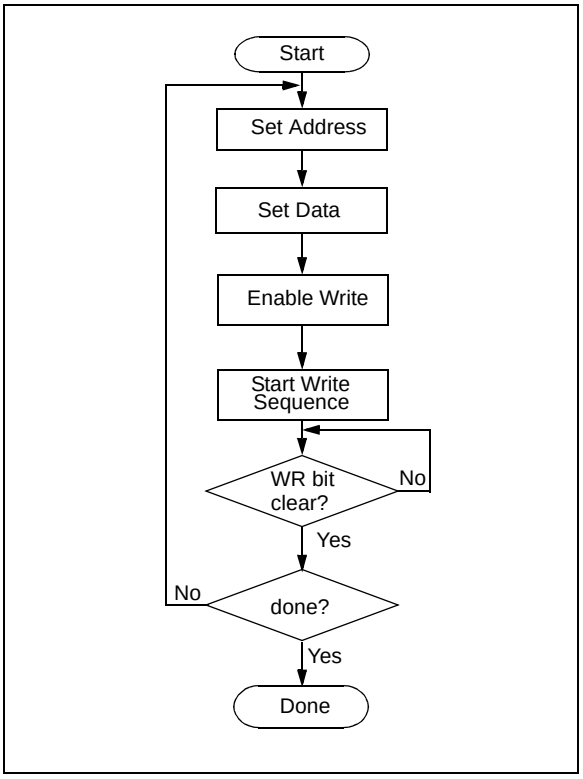
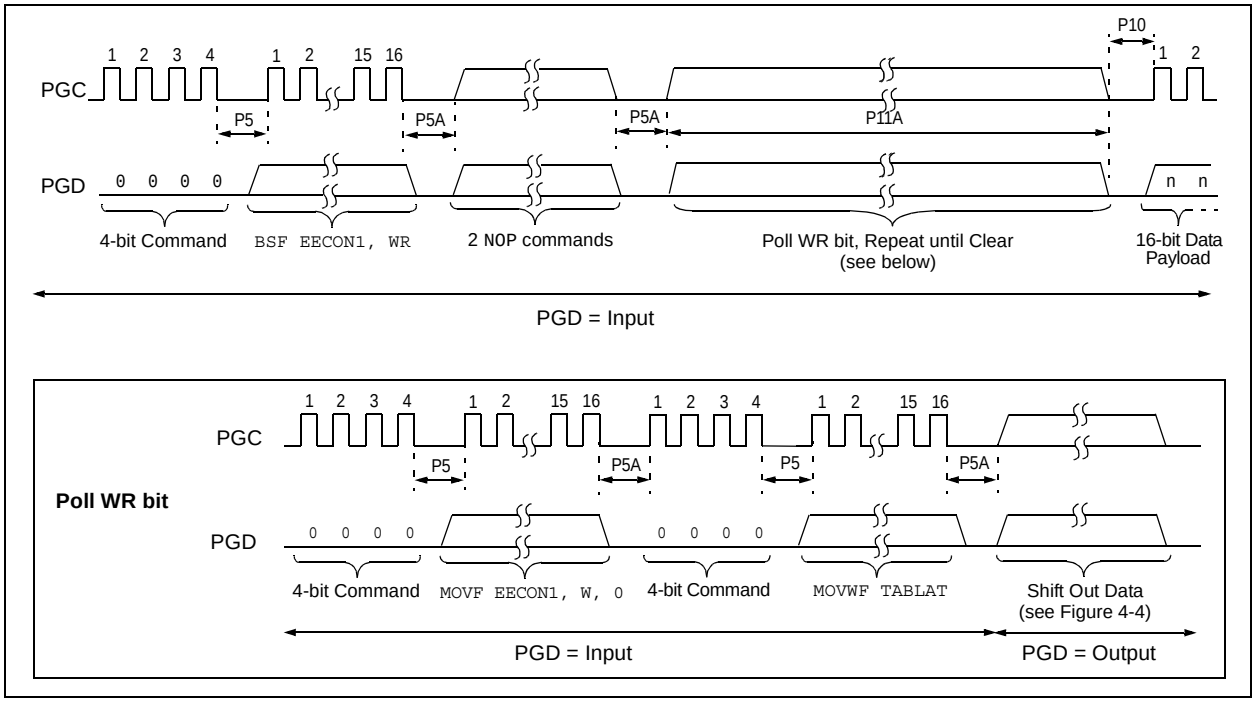


FIGURE 3-7: DATA EEPROM WRITE TIMING DIAGRAM



3.4 ID Location Programming

The ID locations are programmed much like the code memory. The ID registers are mapped in addresses 200000h through 200007h. These locations read out normally even after code protection.

Note:	The user only needs to fill the first 8 bytes of the write buffer in order to write the ID locations.

Table 3-8 demonstrates the code sequence required to write the ID locations.

In order to modify the ID locations, refer to the methodology described in **Section 3.2.1 "Modifying Code Memory"**. As with code memory, the ID locations must be erased before being modified.

When VDD is below the minimum for Bulk Erase operation, ID locations can be cleared with the Row Erase method described in **Section 3.1.3 "ICSP Row Erase"**.

TABLE 3-8: WRITE ID SEQUENCE

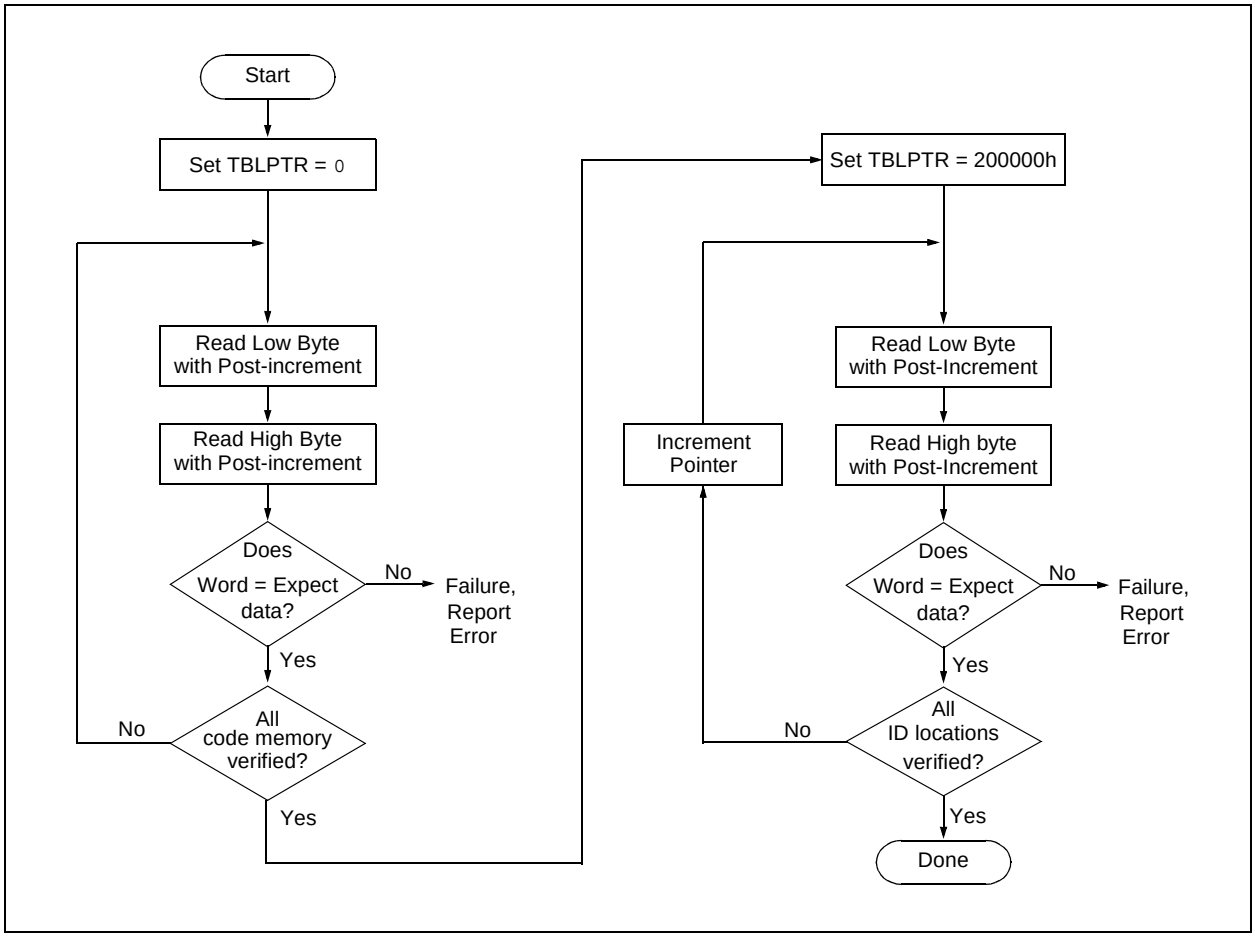
4-bit Command	Data Payload	Core Instruction
Step 1: Direct access to code memory.		
0000	8E A6	BSF EECON1, EEPGD
0000	9C A6	BCF EECON1, CFGS
0000	84 A6	BSF EECON1, WREN
Step 2: Set Table Pointer to ID. Load write buffer with 8 bytes and write.		
0000	0E 20	MOVLW 20h
0000	6E F8	MOVWF TBLPTRU
0000	0E 00	MOVLW 00h
0000	6E F7	MOVWF TBLPTRH
0000	0E 00	MOVLW 00h
0000	6E F6	MOVWF TBLPTRL
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
1101	<MSB><LSB>	Write 2 bytes and post-increment address by 2.
1111	<MSB><LSB>	Write 2 bytes and start programming.
0000	00 00	NOP - hold PGC high for time P9 and low for time P10.

4.2 Verify Code Memory and ID Locations

The verify step involves reading back the code memory space and comparing it against the copy held in the programmer's buffer. Memory reads occur a single byte at a time, so two bytes must be read to compare against the word in the programmer's buffer. Refer to **Section 4.1 "Read Code Memory, ID Locations and Configuration Bits"** for implementation details of reading code memory.

The Table Pointer must be manually set to 200000h (base address of the ID locations) once the code memory has been verified. The post-increment feature of the table read 4-bit command can not be used to increment the Table Pointer beyond the code memory space. In a 64-Kbyte device, for example, a post-increment read of address FFFFh will wrap the Table Pointer back to 000000h, rather than point to unimplemented address 010000h.

FIGURE 4-2: VERIFY CODE MEMORY FLOW



4.3 Verify Configuration Bits

A configuration address may be read and output on PGD via the 4-bit command, '1001'. Configuration data is read and written in a byte-wise fashion, so it is not necessary to merge two bytes into a word prior to a compare. The result may then be immediately compared to the appropriate configuration data in the programmer's memory for verification. Refer to **Section 4.1 "Read Code Memory, ID Locations and Configuration Bits"** for implementation details of reading configuration data.

4.4 Read Data EEPROM Memory

Data EEPROM is accessed one byte at a time via an Address Pointer (register pair EEADRH:EEADR) and a data latch (EEDATA). Data EEPROM is read by loading EEADRH:EEADR with the desired memory location and initiating a memory read by appropriately configuring the EECON1 register. The data will be loaded into EEDATA, where it may be serially output on PGD via the 4-bit command, '0010' (Shift Out Data Holding register). A delay of P6 must be introduced after the falling edge of the 8th PGC of the operand to allow PGD to transition from an input to an output. During this time, PGC must be held low (see Figure 4-4).

The command sequence to read a single byte of data is shown in Table 4-2.

FIGURE 4-3: READ DATA EEPROM FLOW

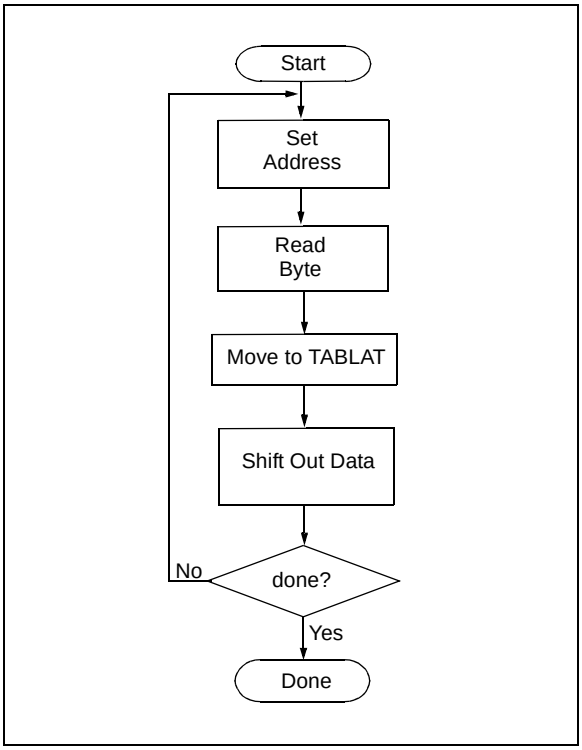
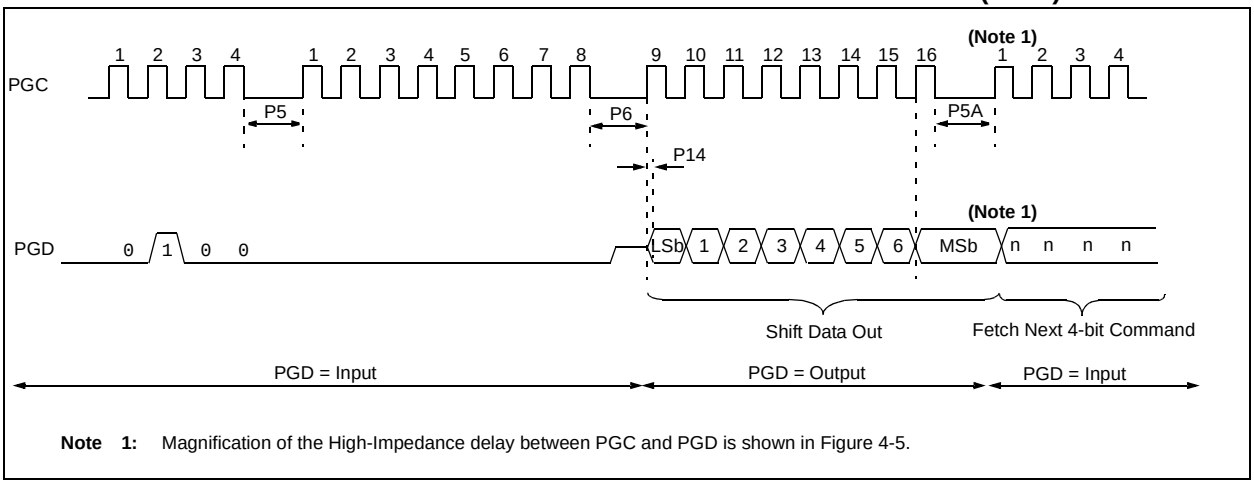


TABLE 4-2: READ DATA EEPROM MEMORY

4-bit Command	Data Payload	Core Instruction
Step 1: Direct access to data EEPROM.		
0000	9E A6	BCF EECON1, EEPGD
0000	9C A6	BCF EECON1, CFGS
Step 2: Set the data EEPROM Address Pointer.		
0000	0E <Addr>	MOVLW <Addr>
0000	6E A9	MOVWF EEADR
0000	OE <AddrH>	MOVLW <AddrH>
0000	6E AA	MOVWF EEADRH
Step 3: Initiate a memory read.		
0000	80 A6	BSF EECON1, RD
Step 4: Load data into the Serial Data Holding register.		
0000	50 A8	MOVF EEDATA, W, 0
0000	6E F5	MOVWF TABLAT
0000	00 00	NOP
0010	<MSB><LSB>	Shift Out Data(1)

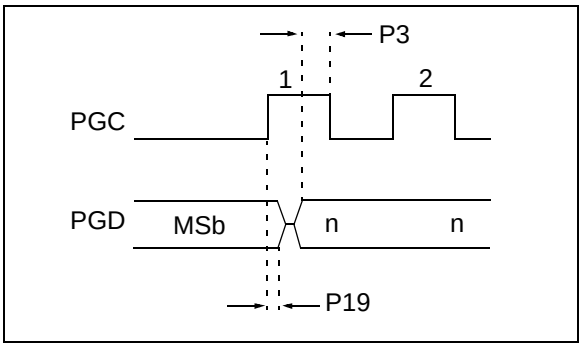
Note 1: The <LSB> is undefined. The <MSB> is the data.

FIGURE 4-4: SHIFT OUT DATA HOLDING REGISTER TIMING DIAGRAM (0010)



Note 1: Magnification of the High-Impedance delay between PGC and PGD is shown in Figure 4-5.

FIGURE 4-5: HIGH-IMPEDANCE DELAY



4.5 Verify Data EEPROM

A data EEPROM address may be read via a sequence of core instructions (4-bit command, '0000') and then output on PGD via the 4-bit command, '0010' (TABLAT register). The result may then be immediately compared to the appropriate data in the programmer's memory for verification. Refer to **Section 4.4 "Read Data EEPROM Memory"** for implementation details of reading data EEPROM.

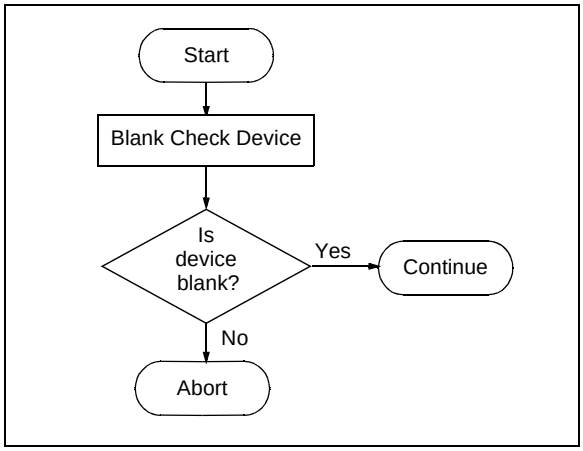
4.6 Blank Check

The term "Blank Check" means to verify that the device has no programmed memory cells. All memories must be verified: code memory, data EEPROM, ID locations and Configuration bits. The device ID registers (3FFFFEh:3FFFFFh) should be ignored.

A "blank" or "erased" memory cell will read as a '1'. Therefore, Blank Checking a device merely means to verify that all bytes read as FFh except the Configuration bits. Unused (reserved) Configuration bits will read '0' (programmed). Refer to Table 5-1 for blank configuration expect data for the various PIC18F2XK20/4XK20 devices.

Given that Blank Checking is merely code and data EEPROM verification with FFh expect data, refer to **Section 4.4 "Read Data EEPROM Memory"** and **Section 4.2 "Verify Code Memory and ID Locations"** for implementation details.

FIGURE 4-6: BLANK CHECK FLOW



5.0 CONFIGURATION WORD

The PIC18F2XK20/4XK20 devices have several Configuration Words. These bits can be set or cleared to select various device configurations. All other memory areas should be programmed and verified prior to setting Configuration Words. These bits may be read out normally, even after read or code protection. See Table 5-1 for a list of Configuration bits and device IDs and Table 5-3 for the Configuration bit descriptions.

5.1 User ID Locations

A user may store identification information (ID) in eight ID locations mapped in 200000h:200007h. It is recommended that the Most Significant nibble of each ID be Fh. In doing so, if the user code inadvertently tries to execute from the ID space, the ID data will execute as a NOP.

5.2 Device ID Word

The device ID word for the PIC18F2XK20/4XK20 devices is located at 3FFFFEh:3FFFFFh. These bits may be used by the programmer to identify what device type is being programmed and read out normally, even after code or read protection. See Table 5-2 for a complete list of device ID values.

FIGURE 5-1: READ DEVICE ID WORD FLOW

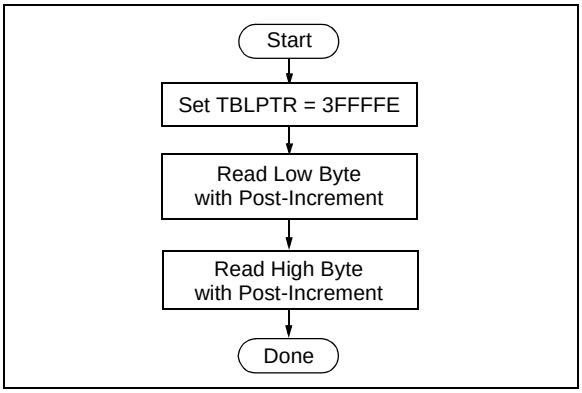


TABLE 5-1: CONFIGURATION BITS AND DEVICE IDs

File Name		Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Default/ Unprogrammed Value
300001h	CONFIG1H	IESO	FCMEN	—	—	FOSC3	FOSC2	FOSC1	FOSC0	00-- 0111
300002h	CONFIG2L	—	—	—	BORV1	BORV0	BOREN1	BOREN0	$\overline{\text{PWRTEN}}$	---1 1111
300003h	CONFIG2H	—	—	—	WDTPS3	WDTPS2	WDTPS1	WDTPS0	WDTEN	---1 1111
300005h	CONFIG3H	MCLRE	—	—	—	HFOFST	LPT1OSC	PBADEN	CCP2MX	1--- 1011
300006h	CONFIG4L	DEBUG	XINST	—	—	—	LVP	—	STVREN	10-- -1-1
300008h	CONFIG5L	—	—	—	—	CP3(1)	CP2(1)	CP1	CP0	---- 1111
300009h	CONFIG5H	CPD	CPB	—	—	—	—	—	—	11-- ----
30000Ah	CONFIG6L	—	—	—	—	WRT3(1)	WRT2(1)	WRT1	WRT0	---- 1111
30000Bh	CONFIG6H	WRTD	WRTB	WRTC	—	—	—	—	—	111- ----
30000Ch	CONFIG7L	—	—	—	—	EBTR3(1)	EBTR2(1)	EBTR1	EBTR0	---- 1111
30000Dh	CONFIG7H	—	EBTRB	—	—	—	—	—	—	-1-- ----
3FFFFEh	DEVID1(2)	DEV2	DEV1	DEV0	REV4	REV3	REV2	REV1	REV0	See Table 5-2
3FFFFFh	DEVID2(2)	DEV10	DEV9	DEV8	DEV7	DEV6	DEV5	DEV4	DEV3	See Table 5-2

Legend: x = unknown, u = unchanged, – = unimplemented. Shaded cells are unimplemented, read as '0'.

Note 1: These bits are only implemented on specific devices. Refer to **Section 2.3 "Memory Maps"** to determine which bits apply based on available memory.

2: DEVID registers are read-only and cannot be programmed by the user.

TABLE 5-2: DEVICE ID VALUE

Device	Device ID Value	
	DEVID2	DEVID1
PIC18F23K20	20h	111x xxxx
PIC18F24K20	20h	101x xxxx
PIC18F25K20	20h	011x xxxx
PIC18F26K20	20h	001x xxxx
PIC18F43K20	20h	110x xxxx
PIC18F44K20	20h	100x xxxx
PIC18F45K20	20h	010x xxxx
PIC18F46K20	20h	000x xxxx

Note: The ‘x’s in DEVID1 contain the device revision code.

TABLE 5-3: PIC18F2XK20/4XK20 BIT DESCRIPTIONS (CONTINUED)

Bit Name	Configuration Words	Description
WDTEN	CONFIG2H	Watchdog Timer Enable bit 1 = WDT enabled 0 = WDT disabled (control is placed on SWDTEN bit)
MCLRE	CONFIG3H	$\overline{\text{MCLR}}$ Pin Enable bit 1 = $\overline{\text{MCLR}}$ pin enabled, RE3 input pin disabled 0 = RE3 input pin enabled, $\overline{\text{MCLR}}$ pin disabled
HFOFST	CONFIG3H	HFINTOSC Fast Start 1 = HFINTOSC output is not delayed 0 = HFINTOSC output is delayed until oscillator is stable (IOFS = 1)
LPT1OSC	CONFIG3H	Low-Power Timer1 Oscillator Enable bit 1 = Timer1 configured for low-power operation 0 = Timer1 configured for higher power operation
PBADEN	CONFIG3H	PORTB A/D Enable bit 1 = PORTB A/D<4:0> pins are configured as analog input channels on Reset 0 = PORTB A/D<4:0> pins are configured as digital I/O on Reset
CCP2MX	CONFIG3H	CCP2 MUX bit 1 = CCP2 input/output is multiplexed with RC1 0 = CCP2 input/output is multiplexed with RB3
DEBUG	CONFIG4L	Background Debugger Enable bit 1 = Background debugger disabled, RB6 and RB7 configured as general purpose I/O pins 0 = Background debugger enabled, RB6 and RB7 are dedicated to In-Circuit Debug
XINST	CONFIG4L	Extended Instruction Set Enable bit 1 = Instruction set extension and Indexed Addressing mode enabled 0 = Instruction set extension and Indexed Addressing mode disabled (Legacy mode)
LVP	CONFIG4L	Low-Voltage Programming Enable bit 1 = Low-Voltage Programming enabled, RB5 is the PGM pin 0 = Low-Voltage Programming disabled, RB5 is an I/O pin
STVREN	CONFIG4L	Stack Overflow/Underflow Reset Enable bit 1 = Reset on stack overflow/underflow enabled 0 = Reset on stack overflow/underflow disabled

.

5.3 Single-Supply ICSP Programming

The LVP bit in Configuration register, CONFIG4L, enables Single-Supply (Low-Voltage) ICSP Programming. The LVP bit defaults to a '1' (enabled) from the factory.

If Single-Supply Programming mode is not used, the LVP bit can be programmed to a '0' and RB5/PGM becomes a digital I/O pin. However, the LVP bit may only be programmed by entering the High-Voltage ICSP mode, where MCLR/VPP/RE3 is raised to VIHH. Once the LVP bit is programmed to a '0', only the High-Voltage ICSP mode is available and only the High-Voltage ICSP mode can be used to program the device.

Note 1: The High-Voltage ICSP mode is always available, regardless of the state of the LVP bit, by applying VIHH to the MCLR/VPP/RE3 pin.

2: While in Low-Voltage ICSP mode, the RB5 pin can no longer be used as a general purpose I/O.

5.4 Embedding Configuration Word Information in the HEX File

To allow portability of code, a PIC18F2XK20/4XK20 programmer is required to read the Configuration Word locations from the hex file. If Configuration Word information is not present in the hex file, then a simple warning message should be issued. Similarly, while saving a hex file, all Configuration Word information must be included. An option to not include the Configuration Word information may be provided. When embedding Configuration Word information in the hex file, it should start at address 300000h.

Microchip Technology Inc. feels strongly that this feature is important for the benefit of the end customer.

5.5 Embedding Data EEPROM Information In the HEX File

To allow portability of code, a PIC18F2XK20/4XK20 programmer is required to read the data EEPROM information from the hex file. If data EEPROM information is not present, a simple warning message should be issued. Similarly, when saving a hex file, all data EEPROM information must be included. An option to not include the data EEPROM information may be provided. When embedding data EEPROM information in the hex file, it should start at address F00000h.

Microchip Technology Inc. believes that this feature is important for the benefit of the end customer.

5.6 Checksum Computation

The checksum is calculated by summing the following:

- The contents of all code memory locations
- The Configuration Word, appropriately masked
- ID locations (Only if any portion of program memory is code-protected)

The Least Significant 16 bits of this sum are the checksum.

Code protection limits access to program memory by both external programmer (code-protect) and code execution (table read protect). The ID locations, when included in a code protected checksum, contain the checksum of an unprotected part. The unprotected checksum is distributed: one nibble per ID location. Each nibble is right justified.

Table 5-4 describes how to calculate the checksum for each device.

Note: The checksum calculation differs depending on the code-protect setting. Since the code memory locations read out differently depending on the code-protect setting, the table describes how to manipulate the actual code memory values to simulate the values that would be read from a protected device. When calculating a checksum by reading a device, the entire code memory can simply be read and summed. The Configuration Word and ID locations can always be read.

TABLE 5-4: CHECKSUM COMPUTATION (CONTINUED)

Device	Code-Protect	Checksum	Blank Value	0xAA at 0 and Max Address
PIC18FX5K20	None	SUM[0000:07FF]+SUM[0800:1FFF]+SUM[2000:3FFF]+ SUM[4000:5FFF]+SUM[6000:7FFF]+(CONFIG1L & 00h)+ (CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+ (CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+ (CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+ (CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+ (CONFIG7H & 40h)	8362h	82B8h
	Boot Block	SUM[0800:1FFF]+SUM[2000:3FFF]+SUM[4000:5FFF]+SUM[6000:7FFF]+ (CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+ (CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+ (CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+ (CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+ (CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	8B35h	8AEAh
	Boot/ Block 0/ Block 1	SUM[4000:5FFF]+SUM[6000:7FFF]+(CONFIG1L & 00h)+ (CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+ (CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+ (CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+ (CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+ (CONFIG7H & 40h)+SUM_ID	C332h	C2E7h
	All	(CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+ (CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+ (CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+ (CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+ (CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	0326h	0330h

Legend:

Item		Description
CONFIGx	=	Configuration Word
SUM[a:b]	=	Sum of locations, a to b inclusive
SUM_ID	=	Byte-wise sum of lower four bits of all customer ID locations
+	=	Addition
&	=	Bit-wise AND

TABLE 5-4: CHECKSUM COMPUTATION (CONTINUED)

Device	Code-Protect	Checksum	Blank Value	0xAA at 0 and Max Address
PIC18FX6K20	None	SUM[0000:07FF]+SUM[0800:3FFF]+SUM[4000:7FFF]+ SUM[8000:BFFF]+SUM[C000:FFFF]+(CONFIG1L & 00h)+ (CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+ (CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+ (CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+ (CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+ (CONFIG7H & 40h)	0362h	02B8h
	Boot Block	SUM[0800:3FFF]+SUM[4000:7FFF]+SUM[8000:BFFF]+SUM[C000:FFFF]+ (CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+ (CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+ (CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+ (CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+ (CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	0B2Dh	0AE2h
	Boot/ Block 0/ Block 1	SUM[3000:BFFF]+SUM[C000:FFFF]+(CONFIG1L & 00h)+ (CONFIG1H & CFh)+(CONFIG2L & 1Fh)+(CONFIG2H & 1F)+ (CONFIG3L & 00h)+(CONFIG3H & 8Fh)+(CONFIG4L & C5h)+ (CONFIG4H & 00h)+(CONFIG5L & 0Fh)+(CONFIG5H & C0h)+ (CONFIG6L & 0Fh)+(CONFIG6H & E0h)+(CONFIG7L & 0Fh)+ (CONFIG7H & 40h)+SUM_ID	832Ah	82DFh
	All	(CONFIG1L & 00h)+(CONFIG1H & CFh)+(CONFIG2L & 1Fh)+ (CONFIG2H & 1F)+(CONFIG3L & 00h)+(CONFIG3H & 8Fh)+ (CONFIG4L & C5h)+(CONFIG4H & 00h)+(CONFIG5L & 0Fh)+ (CONFIG5H & C0h)+(CONFIG6L & 0Fh)+(CONFIG6H & E0h)+ (CONFIG7L & 0Fh)+(CONFIG7H & 40h)+SUM_ID	031Eh	0328h

Legend:

Item		Description
CONFIGx	=	Configuration Word
SUM[a:b]	=	Sum of locations, a to b inclusive
SUM_ID	=	Byte-wise sum of lower four bits of all customer ID locations
+	=	Addition
&	=	Bit-wise AND

Note the following details of the code protection feature on Microchip devices:

- Microchip products meet the specification contained in their particular Microchip Data Sheet.
- Microchip believes that its family of products is one of the most secure families of its kind on the market today, when used in the intended manner and under normal conditions.
- There are dishonest and possibly illegal methods used to breach the code protection feature. All of these methods, to our knowledge, require using the Microchip products in a manner outside the operating specifications contained in Microchip's Data Sheets. Most likely, the person doing so is engaged in theft of intellectual property.
- Microchip is willing to work with the customer who is concerned about the integrity of their code.
- Neither Microchip nor any other semiconductor manufacturer can guarantee the security of their code. Code protection does not mean that we are guaranteeing the product as "unbreakable."

Code protection is constantly evolving. We at Microchip are committed to continuously improving the code protection features of our products. Attempts to break Microchip's code protection feature may be a violation of the Digital Millennium Copyright Act. If such acts allow unauthorized access to your software or other copyrighted work, you may have a right to sue for relief under that Act.

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