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#### Details

Product Status	Obsolete
Core Processor	ST7
Core Size	8-Bit
Speed	8MHz
Connectivity	I <sup>2</sup> C, SCI, SPI
Peripherals	LVD, POR, PWM, WDT
Number of I/O	32
Program Memory Size	16KB (16K x 8)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	512 x 8
Voltage - Supply (Vcc/Vdd)	3.8V ~ 5.5V
Data Converters	A/D 12x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	44-LQFP
Supplier Device Package	-
Purchase URL	https://www.e-xfl.com/product-detail/stmicroelectronics/st72f325j4t6tr

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# **3 REGISTER & MEMORY MAP**

As shown in Figure 6, the MCU is capable of addressing 64K bytes of memories and I/O registers.

The available memory locations consist of 128 bytes of register locations, up to 2Kbytes of RAM and up to 60Kbytes of user program memory. The RAM space includes up to 256 bytes for the stack from 0100h to 01FFh.

The highest address bytes contain the user reset and interrupt vectors.

**IMPORTANT:** Memory locations marked as "Reserved" must never be accessed. Accessing a reseved area can have unpredictable effects on the device.

## **Related Documentation**

AN 985: Executing Code in ST7 RAM

# Figure 6. Memory Map

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# **5 CENTRAL PROCESSING UNIT**

# **5.1 INTRODUCTION**

This CPU has a full 8-bit architecture and contains six internal registers allowing efficient 8-bit data manipulation.

# **5.2 MAIN FEATURES**

- Enable executing 63 basic instructions
- Fast 8-bit by 8-bit multiply
- 17 main addressing modes (with indirect addressing mode)
- Two 8-bit index registers
- 16-bit stack pointer
- Low power HALT and WAIT modes
- Priority maskable hardware interrupts
- Non-maskable software/hardware interrupts

# **5.3 CPU REGISTERS**

The six CPU registers shown in Figure 1 are not present in the memory mapping and are accessed by specific instructions.

## Accumulator (A)

The Accumulator is an 8-bit general purpose register used to hold operands and the results of the arithmetic and logic calculations and to manipulate data.

#### Index Registers (X and Y)

These 8-bit registers are used to create effective addresses or as temporary storage areas for data manipulation. (The Cross-Assembler generates a precede instruction (PRE) to indicate that the following instruction refers to the Y register.)

The Y register is not affected by the interrupt automatic procedures.

#### **Program Counter (PC)**

The program counter is a 16-bit register containing the address of the next instruction to be executed by the CPU. It is made of two 8-bit registers PCL (Program Counter Low which is the LSB) and PCH (Program Counter High which is the MSB).

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# Figure 10. CPU Registers

# CENTRAL PROCESSING UNIT (Cont'd)

## Condition Code Register (CC)

Read/Write

Reset Value: 111x1xxx



The 8-bit Condition Code register contains the interrupt masks and four flags representative of the result of the instruction just executed. This register can also be handled by the PUSH and POP instructions.

These bits can be individually tested and/or controlled by specific instructions.

#### **Arithmetic Management Bits**

Bit  $4 = \mathbf{H}$  Half carry.

This bit is set by hardware when a carry occurs between bits 3 and 4 of the ALU during an ADD or ADC instructions. It is reset by hardware during the same instructions.

0: No half carry has occurred. 1: A half carry has occurred.

This bit is tested using the JRH or JRNH instruction. The H bit is useful in BCD arithmetic subroutines.

#### Bit 2 = **N** *Negative*.

This bit is set and cleared by hardware. It is representative of the result sign of the last arithmetic, logical or data manipulation. It's a copy of the result 7<sup>th</sup> bit.

0: The result of the last operation is positive or null.

1: The result of the last operation is negative

(that is, the most significant bit is a logic 1).

This bit is accessed by the JRMI and JRPL instructions.

#### Bit 1 = **Z** Zero.

This bit is set and cleared by hardware. This bit indicates that the result of the last arithmetic, logical or data manipulation is zero.

- 0: The result of the last operation is different from zero.
- 1: The result of the last operation is zero.

This bit is accessed by the JREQ and JRNE test instructions.

#### Bit 0 = C Carry/borrow.

This bit is set and cleared by hardware and software. It indicates an overflow or an underflow has occurred during the last arithmetic operation.

0: No overflow or underflow has occurred.

1: An overflow or underflow has occurred.

This bit is driven by the SCF and RCF instructions and tested by the JRC and JRNC instructions. It is also affected by the "bit test and branch", shift and rotate instructions.

#### **Interrupt Management Bits**

Bit 5,3 = 11, 10 Interrupt

The combination of the I1 and I0 bits gives the current interrupt software priority.

Interrupt Software Priority	11	10
Level 0 (main)	1	0
Level 1	0	1
Level 2	0	0
Level 3 (= interrupt disable)	1	1

These two bits are set/cleared by hardware when entering in interrupt. The loaded value is given by the corresponding bits in the interrupt software priority registers (IxSPR). They can be also set/ cleared by software with the RIM, SIM, IRET, HALT, WFI and PUSH/POP instructions.

See the interrupt management chapter for more details.

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# CENTRAL PROCESSING UNIT (Cont'd)

## Stack Pointer (SP)

#### Read/Write

Reset Value: 01 FFh



The Stack Pointer is a 16-bit register which is always pointing to the next free location in the stack. It is then decremented after data has been pushed onto the stack and incremented before data is popped from the stack (see Figure 2).

Since the stack is 256 bytes deep, the 8 most significant bits are forced by hardware. Following an MCU Reset, or after a Reset Stack Pointer instruction (RSP), the Stack Pointer contains its reset value (the SP7 to SP0 bits are set) which is the stack higher address.

#### Figure 11. Stack Manipulation Example

The least significant byte of the Stack Pointer (called S) can be directly accessed by a LD instruction.

**Note:** When the lower limit is exceeded, the Stack Pointer wraps around to the stack upper limit, without indicating the stack overflow. The previously stored information is then overwritten and therefore lost. The stack also wraps in case of an underflow.

The stack is used to save the return address during a subroutine call and the CPU context during an interrupt. The user may also directly manipulate the stack by means of the PUSH and POP instructions. In the case of an interrupt, the PCL is stored at the first location pointed to by the SP. Then the other registers are stored in the next locations as shown in Figure 2.

- When an interrupt is received, the SP is decremented and the context is pushed on the stack.
- On return from interrupt, the SP is incremented and the context is popped from the stack.

A subroutine call occupies two locations and an interrupt five locations in the stack area.

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## INTERRUPTS (Cont'd)

# 7.5 INTERRUPT REGISTER DESCRIPTION

# CPU CC REGISTER INTERRUPT BITS

# Read/Write

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Reset Value: 111x 1010 (xAh)

7							0	
1	1	11	н	10	Ν	Z	С	

#### Bit 5, 3 = 11, 10 Software Interrupt Priority

These two bits indicate the current interrupt software priority.

Interrupt Software Priority	Level	11	10
Level 0 (main)	Low	1	0
Level 1		0	1
Level 2	▼	0	0
Level 3 (= interrupt disable*)	High	1	1

These two bits are set/cleared by hardware when entering in interrupt. The loaded value is given by the corresponding bits in the interrupt software priority registers (ISPRx).

They can be also set/cleared by software with the RIM, SIM, HALT, WFI, IRET and PUSH/POP instructions (see "Interrupt Dedicated Instruction Set" table).

\*Note: TRAP and RESET events can interrupt a level 3 program.

# INTERRUPT SOFTWARE PRIORITY REGISTERS (ISPRX)

Read/Write (bit 7:4 of **ISPR3** are read only) Reset Value: 1111 1111 (FFh)

	1							0
ISPR0	l1_3	10_3	l1_2	10_2	11_1	10_1	l1_0	10_0
ISPR1	11_7	10_7	l1_6	10_6	l1_5	10_5	11_4	10_4
ISPR2	11_11	10_11	11_10	10_10	l1_9	10_9	l1_8	10_8
ISPR3	1	1	1	1	11_13	10_13	11_12	10_12

These four registers contain the interrupt software priority of each interrupt vector.

 Each interrupt vector (except RESET and TRAP) has corresponding bits in these registers where its own software priority is stored. This correspondance is shown in the following table.

Vector address	ISPRx bits
FFFBh-FFFAh	11_0 and 10_0 bits*
FFF9h-FFF8h	I1_1 and I0_1 bits
FFE1h-FFE0h	I1_13 and I0_13 bits

Each I1\_x and I0\_x bit value in the ISPRx registers has the same meaning as the I1 and I0 bits in the CC register.

 Level 0 can not be written (l1\_x=1, l0\_x=0). In this case, the previously stored value is kept. (example: previous=CFh, write=64h, result=44h)

The TLI, RESET, and TRAP vectors have no software priorities. When one is serviced, the I1 and I0 bits of the CC register are both set.

\*Note: Bits in the ISPRx registers which correspond to the TLI can be read and written but they are not significant in the interrupt process management.

**Caution**: If the  $I1_x$  and  $I0_x$  bits are modified while the interrupt x is executed the following behaviour has to be considered: If the interrupt x is still pending (new interrupt or flag not cleared) and the new software priority is higher than the previous one, the interrupt x is re-entered. Otherwise, the software priority stays unchanged up to the next interrupt request (after the IRET of the interrupt x).

# POWER SAVING MODES (Cont'd)

#### 8.4.2 HALT MODE

The HALT mode is the lowest power consumption mode of the MCU. It is entered by executing the 'HALT' instruction when the OIE bit of the Main Clock Controller Status register (MCCSR) is cleared (see section 10.2 on page 61 for more details on the MCCSR register).

The MCU can exit HALT mode on reception of either a specific interrupt (see Table 9, "Interrupt Mapping," on page 41) or a RESET. When exiting HALT mode by means of a RESET or an interrupt, the oscillator is immediately turned on and the 256 or 4096 CPU cycle delay is used to stabilize the oscillator. After the start up delay, the CPU resumes operation by servicing the interrupt or by fetching the reset vector which woke it up (see Figure 32).

When entering HALT mode, the I[1:0] bits in the CC register are forced to '10b'to enable interrupts. Therefore, if an interrupt is pending, the MCU wakes up immediately.

In HALT mode, the main oscillator is turned off causing all internal processing to be stopped, including the operation of the on-chip peripherals. All peripherals are not clocked except the ones which get their clock supply from another clock generator (such as an external or auxiliary oscillator).

The compatibility of Watchdog operation with HALT mode is configured by the "WDGHALT" option bit of the option byte. The HALT instruction when executed while the Watchdog system is enabled, can generate a Watchdog RESET (see section 14.1 on page 181 for more details).

Figure 3	1. HAL	T Timing	Overview
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## Notes:

1. WDGHALT is an option bit. See option byte section for more details.

2. Peripheral clocked with an external clock source can still be active.

3. Only some specific interrupts can exit the MCU from HALT mode (such as external interrupt). Refer to Table 9, "Interrupt Mapping," on page 41 for more details.

4. Before servicing an interrupt, the CC register is pushed on the stack. The I[1:0] bits of the CC register are set to the current software priority level of the interrupt routine and recovered when the CC register is popped.

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# Figure 32. HALT Mode Flow-chart

# WATCHDOG TIMER (Cont'd)

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## 10.1.4 How to Program the Watchdog Timeout

Figure 2 shows the linear relationship between the 6-bit value to be loaded in the Watchdog Counter (CNT) and the resulting timeout duration in milliseconds. This can be used for a quick calculation without taking the timing variations into account. If

Figure 36. Approximate Timeout Duration

more precision is needed, use the formulae in Figure 3.

**Caution:** When writing to the WDGCR register, always write 1 in the T6 bit to avoid generating an immediate reset.



## **ON-CHIP PERIPHERALS** (Cont'd)

#### Input capture function

This mode allows the measurement of external signal pulse widths through ARTICRx registers.

Each input capture can generate an interrupt independently on a selected input signal transition. This event is flagged by a set of the corresponding CFx bits of the Input Capture Control/Status register (ARTICCSR).

These input capture interrupts are enabled through the CIEx bits of the ARTICCSR register.

The active transition (falling or rising edge) is software programmable through the CSx bits of the ARTICCSR register.

The read only input capture registers (ARTICRx) are used to latch the auto-reload counter value when a transition is detected on the ARTICx pin (CFx bit set in ARTICCSR register). After fetching the interrupt vector, the CFx flags can be read to identify the interrupt source.

**Note**: After a capture detection, data transfer in the ARTICRx register is inhibited until it is read (clearing the CFx bit).

The timer interrupt remains pending while the CFx flag is set when the interrupt is enabled (CIEx bit set). This means, the ARTICRx register has to be read at each capture event to clear the CFx flag.

The timing resolution is given by auto-reload counter cycle time  $(1/f_{COUNTER})$ .

**Note:** During HALT mode, if both input capture and external clock are enabled, the ARTICRx register value is not guaranteed if the input capture pin and the external clock change simultaneously.

#### Figure 44. Input Capture Timing Diagram

## External interrupt capability

This mode allows the Input capture capabilities to be used as external interrupt sources. The interrupts are generated on the edge of the ARTICx signal.

The edge sensitivity of the external interrupts is programmable (CSx bit of ARTICCSR register) and they are independently enabled through CIEx bits of the ARTICCSR register. After fetching the interrupt vector, the CFx flags can be read to identify the interrupt source.

During HALT mode, the external interrupts can be used to wake up the micro (if the CIEx bit is set).



# 16-BIT TIMER (Cont'd)

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# 16-BIT TIMER (Cont'd)

Table 19. 16-	-Bit Timer	Register	Мар	and	Reset	Values
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Address (Hex.)	Register Label	7	6	5	4	3	2	1	0
Timer A: 32	CR1	ICIE	OCIE	TOIE	FOLV2	FOLV1	OLVL2	IEDG1	OLVL1
Timer B: 42	Reset Value	0	0	0	0	0	0	0	0
Timer A: 31	CR2	OC1E	OC2E	OPM	PWM	CC1	CC0	IEDG2	EXEDG
Timer B: 41	Reset Value	0	0	0	0	0	0	0	0
Timer A: 33	CSR	ICF1	OCF1	TOF	ICF2	OCF2	TIMD	-	-
Timer B: 43	Reset Value	х	х	х	х	х	0	х	х
Timer A: 34	IC1HR	MSB							LSB
Timer B: 44	Reset Value	х	х	Х	х	х	х	х	х
Timer A: 35	IC1LR	MSB							LSB
Timer B: 45	Reset Value	х	х	х	х	х	х	х	х
Timer A: 36	OC1HR	MSB							LSB
Timer B: 46	Reset Value	1	0	0	0	0	0	0	0
Timer A: 37	OC1LR	MSB							LSB
Timer B: 47	Reset Value	0	0	0	0	0	0	0	0
Timer A: 3E	OC2HR	MSB							LSB
Timer B: 4E	Reset Value	1	0	0	0	0	0	0	0
Timer A: 3F	OC2LR	MSB							LSB
Timer B: 4F	Reset Value	0	0	0	0	0	0	0	0
Timer A: 38	CHR	MSB							LSB
Timer B: 48	Reset Value	1	1	1	1	1	1	1	1
Timer A: 39	CLR	MSB							LSB
Timer B: 49	Reset Value	1	1	1	1	1	1	0	0
Timer A: 3A	ACHR	MSB							LSB
Timer B: 4A	Reset Value	1	1	1	1	1	1	1	1
Timer A: 3B	ACLR	MSB							LSB
Timer B: 4B	Reset Value	1	1	1	1	1	1	0	0
Timer A: 3C	IC2HR	MSB							LSB
Timer B: 4C	Reset Value	х	х	Х	х	х	х	х	х
Timer A: 3D	IC2LR	MSB							LSB
Timer B: 4D	Reset Value	Х	Х	Х	х	х	х	х	х

## **Related Documentation**

AN 973: SCI software communications using 16bit timer

AN 974: Real Time Clock with ST7 Timer Output Compare

AN 976: Driving a buzzer through the ST7 Timer PWM function

AN1041: Using ST7 PWM signal to generate analog input (sinusoid)

AN1046: UART emulation software

AN1078: PWM duty cycle switch implementing true 0 or 100 per cent duty cycle

AN1504: Starting a PWM signal directly at high level using the ST7 16-Bit timer

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# SERIAL PERIPHERAL INTERFACE (Cont'd)

# 10.5.6 Low Power Modes

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Mode	Description
WAIT	No effect on SPI. SPI interrupt events cause the device to exit from WAIT mode.
HALT	SPI registers are frozen. In HALT mode, the SPI is inactive. SPI oper- ation resumes when the MCU is woken up by an interrupt with "exit from HALT mode" ca- pability. The data received is subsequently read from the SPIDR register when the soft- ware is running (interrupt vector fetching). If several data are received before the wake- up event, then an overrun error is generated. This error can be detected after the fetch of the interrupt routine that woke up the device.

# 10.5.6.1 Using the SPI to wakeup the MCU from Halt mode

In slave configuration, the SPI is able to wakeup the ST7 device from HALT mode through a SPIF interrupt. The data received is subsequently read from the SPIDR register when the software is running (interrupt vector fetch). If multiple data transfers have been performed before software clears the SPIF bit, then the OVR bit is set by hardware. **Note:** When waking up from Halt mode, if the SPI remains in Slave mode, it is recommended to perform an extra communications cycle to bring the SPI from Halt mode state to normal state. If the SPI exits from Slave mode, it returns to normal state immediately.

**Caution:** The SPI can wake up the ST7 from Halt mode only if the Slave Select signal (external SS pin or the SSI bit in the SPICSR register) is low when the ST7 enters Halt mode. So if Slave selection is configured as external (see Section 10.5.3.2), make sure the master drives a low level on the SS pin when the slave enters Halt mode.

## 10.5.7 Interrupts

Interrupt Event	Event Flag	Enable Control Bit	Exit from Wait	Exit from Halt
SPI End of Transfer Event	SPIF		Yes	Yes
Master Mode Fault Event	MODF	SPIE	Yes	No
Overrun Error	OVR		Yes	No

**Note**: The SPI interrupt events are connected to the same interrupt vector (see Interrupts chapter). They generate an interrupt if the corresponding Enable Control Bit is set and the interrupt mask in

# SERIAL COMMUNICATIONS INTERFACE (Cont'd) CONTROL REGISTER 2 (SCICR2)

Read/Write

Reset Value: 0000 0000 (00h)

7							0
TIE	TCIE	RIE	ILIE	TE	RE	RWU	SBK

Bit 7 = **TIE** *Transmitter interrupt enable.* This bit is set and cleared by software.

0: Interrupt is inhibited

1: An SCI interrupt is generated whenever TDRE=1 in the SCISR register

Bit 6 = TCIE *Transmission complete interrupt enable* 

This bit is set and cleared by software.

0: Interrupt is inhibited

1: An SCI interrupt is generated whenever TC=1 in the SCISR register

Bit 5 = **RIE** *Receiver interrupt enable.* 

This bit is set and cleared by software.

0: Interrupt is inhibited

1: An SCI interrupt is generated whenever OR=1 or RDRF=1 in the SCISR register

Bit 4 = **ILIE** *Idle line interrupt enable.* 

This bit is set and cleared by software.

0: Interrupt is inhibited

1: An SCI interrupt is generated whenever IDLE=1 in the SCISR register.

Bit 3 = **TE** *Transmitter enable*.

This bit enables the transmitter. It is set and cleared by software. 0: Transmitter is disabled 1: Transmitter is enabled

#### Notes:

- During transmission, a "0" pulse on the TE bit ("0" followed by "1") sends a preamble (idle line) after the current word.
- When TE is set there is a 1 bit-time delay before the transmission starts.

**CAUTION:** The TDO pin is free for general purpose I/O only when the TE and RE bits are both cleared (or if TE is never set).

#### Bit 2 = **RE** Receiver enable.

This bit enables the receiver. It is set and cleared by software.

- 0: Receiver is disabled
- 1: Receiver is enabled and begins searching for a start bit

Bit 1 = RWU Receiver wake-up.

This bit determines if the SCI is in mute mode or not. It is set and cleared by software and can be cleared by hardware when a wake-up sequence is recognized.

0: Receiver in Active mode

1: Receiver in Mute mode

**Note:** Before selecting Mute mode (setting the RWU bit), the SCI must receive some data first, otherwise it cannot function in Mute mode with wake-up by idle line detection.

#### Bit 0 = **SBK** Send break.

This bit set is used to send break characters. It is set and cleared by software.

0: No break character is transmitted

1: Break characters are transmitted

**Note:** If the SBK bit is set to "1" and then to "0", the transmitter sends a BREAK word at the end of the current word.

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# SERIAL COMMUNICATION INTERFACE (Cont'd)

Table 24. SCI Register Map and Reset Values

Address (Hex.)	Register Label	7	6	5	4	3	2	1	0
00504	SCISR	TDRE	TC	RDRF	IDLE	OVR	NF	FE	PE
005011	Reset Value	1	1	0	0	0	0	0	0
0051h	SCIDR	MSB							LSB
003111	Reset Value	х	х	х	х	х	х	х	х
0050h	SCIBRR	SCP1	SCP0	SCT2	SCT1	SCT0	SCR2	SCR1	SCR0
005211	Reset Value	0	0	0	0	0	0	0	0
0053h	SCICR1	R8	T8	SCID	М	WAKE	PCE	PS	PIE
	Reset Value	х	0	0	0	0	0	0	0
0054b	SCICR2	TIE	TCIE	RIE	ILIE	TE	RE	RWU	SBK
003411	Reset Value	0	0	0	0	0	0	0	0
0055h	SCIERPR	MSB							LSB
	Reset Value	0	0	0	0	0	0	0	0
00576	SCIPETPR	MSB							LSB
005711	Reset Value	0	0	0	0	0	0	0	0



# I<sup>2</sup>C BUS INTERFACE (Cont'd)

#### Bit 1 = **M/SL** *Master/Slave*.

This bit is set by hardware as soon as the interface is in Master mode (writing START=1). It is cleared by hardware after detecting a Stop condition on the bus or a loss of arbitration (ARLO=1). It is also cleared when the interface is disabled (PE=0). 0: Slave mode

1: Master mode

#### Bit 0 = SB Start bit (Master mode).

This bit is set by hardware as soon as the Start condition is generated (following a write START=1). An interrupt is generated if ITE=1. It is cleared by software reading SR1 register followed by writing the address byte in DR register. It is also cleared by hardware when the interface is disabled (PE=0).

0: No Start condition

1: Start condition generated

# I<sup>2</sup>C STATUS REGISTER 2 (SR2)

Read Only

Reset Value: 0000 0000 (00h)

7	7						
0	0	0	AF	STOPF	ARLO	BERR	GCAL

Bit 7:5 = Reserved. Forced to 0 by hardware.

#### Bit 4 = **AF** Acknowledge failure.

This bit is set by hardware when no acknowledge is returned. An interrupt is generated if ITE=1. It is cleared by software reading SR2 register or by hardware when the interface is disabled (PE=0).

The SCL line is not held low while AF=1 but by other flags (SB or BTF) that are set at the same time.

0: No acknowledge failure

1: Acknowledge failure

Note:

– When an AF event occurs, the SCL line is not held low; however, the SDA line can remain low if the last bits transmitted are all 0. It is then necessary to release both lines by software.

Bit 3 = **STOPF** *Stop detection (Slave mode).* This bit is set by hardware when a Stop condition is detected on the bus after an acknowledge (if ACK=1). An interrupt is generated if ITE=1. It is cleared by software reading SR2 register or by

hardware when the interface is disabled (PE=0).

The SCL line is not held low while STOPF=1.

- 0: No Stop condition detected
- 1: Stop condition detected

#### Bit 2 = ARLO Arbitration lost.

This bit is set by hardware when the interface loses the arbitration of the bus to another master. An interrupt is generated if ITE=1. It is cleared by software reading SR2 register or by hardware when the interface is disabled (PE=0).

After an ARLO event the interface switches back automatically to Slave mode (M/SL=0).

The SCL line is not held low while ARLO=1.

0: No arbitration lost detected

1: Arbitration lost detected

Note:

- In a Multimaster environment, when the interface is configured in Master Receive mode it does not perform arbitration during the reception of the Acknowledge Bit. Mishandling of the ARLO bit from the I2CSR2 register may occur when a second master simultaneously requests the same data from the same slave and the I<sup>2</sup>C master does not acknowledge the data. The ARLO bit is then left at 0 instead of being set.

# Bit 1 = **BERR** Bus error.

This bit is set by hardware when the interface detects a misplaced Start or Stop condition. An interrupt is generated if ITE=1. It is cleared by software reading SR2 register or by hardware when the interface is disabled (PE=0).

The SCL line is not held low while BERR=1.

0: No misplaced Start or Stop condition

1: Misplaced Start or Stop condition

Note:

 If a Bus Error occurs, a Stop or a repeated Start condition should be generated by the Master to re-synchronize communication, get the transmission acknowledged and the bus released for further communication

Bit 0 = GCAL General Call (Slave mode).

This bit is set by hardware when a general call address is detected on the bus while ENGC=1. It is cleared by hardware detecting a Stop condition (STOPF=1) or when the interface is disabled (PE=0).

- 0: No general call address detected on bus
- 1: general call address detected on bus

# CLOCK CHARACTERISTICS (Cont'd) 12.5.4 RC Oscillators

Symbol	Parameter	Conditions	Min	Тур	Max	Unit
fosc (RCINT)	Internal RC oscillator frequency	T₄=25°C, V <sub>DD</sub> =5V	2	3.5	5.6	MHz
	See Figure 77	A 900				

# Figure 77. Typical f<sub>OSC(RCINT)</sub> vs T<sub>A</sub>



Note: To reduce disturbance to the RC oscillator, it is recommended to place decoupling capacitors between  $V_{DD}$  and  $V_{SS}$  as shown in Figure 97

# CLOCK CHARACTERISTICS (Cont'd)

# 12.5.5 Clock Security System (CSS)

Symbol	Parameter	Conditions	Min	Тур	Max	Unit
f <sub>SFOSC</sub>	Safe Oscillator Frequency <sup>1)</sup>			3		MHz

#### Note:

1. Data based on characterization results.

## 12.5.6 PLL Characteristics

Symbol	Parameter	Conditions	Min	Тур	Max	Unit
f <sub>OSC</sub>	PLL input frequency range		2		4	MHz
$\Delta f_{\rm CPU} / f_{\rm CPU}$	Instantaneous PLL jitter <sup>1)</sup>	f <sub>OSC</sub> = 4 MHz.		0.7	2	%

#### Note:

1. Data characterized but not tested.

The user must take the PLL jitter into account in the application (for example in serial communication or sampling of high frequency signals). The PLL jitter is a periodic effect, which is integrated over several CPU cycles. Therefore the longer the period of the application signal, the less it will be impacted by the PLL jitter.

Figure 78 shows the PLL jitter integrated on application signals in the range 125kHz to 4MHz. At frequencies of less than 125KHz, the jitter is negligible.

# Figure 78. Integrated PLL Jitter vs signal frequency<sup>1</sup>



Note 1: Measurement conditions: f<sub>CPU</sub> = 8MHz.



# **13 PACKAGE CHARACTERISTICS**

# **13.1 PACKAGE MECHANICAL DATA**

# Figure 99. 64-Pin Low Profile Quad Flat Package (14x14)



# PACKAGE MECHANICAL DATA (Cont'd)

# Figure 105. 32-Pin Low Profile Quad Flat Package





# **15 KNOWN LIMITATIONS**

## **15.1 ALL DEVICES**

## **15.1.1 Unexpected Reset Fetch**

If an interrupt request occurs while a "POP CC" instruction is executed, the interrupt controller does not recognise the source of the interrupt and, by default, passes the RESET vector address to the CPU.

#### Workaround

To solve this issue, a "POP CC" instruction must always be preceded by a "SIM" instruction.

#### 15.1.2 External interrupt missed

To avoid any risk if generating a parasitic interrupt, the edge detector is automatically disabled for one clock cycle during an access to either DDR and OR. Any input signal edge during this period will not be detected and will not generate an interrupt.

This case can typically occur if the application refreshes the port configuration registers at intervals during runtime.

#### Workaround

The workaround is based on software checking the level on the interrupt pin before and after writing to the PxOR or PxDDR registers. If there is a level change (depending on the sensitivity programmed for this pin) the interrupt routine is invoked using the call instruction with three extra PUSH instructions before executing the interrupt routine (this is to make the call compatible with the IRET instruction at the end of the interrupt service routine).

But detection of the level change does not make sure that edge occurs during the critical 1 cycle duration and the interrupt has been missed. This may lead to occurrence of same interrupt twice (one hardware and another with software call).

To avoid this, a semaphore is set to '1' before checking the level change. The semaphore is changed to level '0' inside the interrupt routine. When a level change is detected, the semaphore status is checked and if it is '1' this means that the last interrupt has been missed. In this case, the interrupt routine is invoked with the call instruction.

There is another possible case i.e. if writing to PxOR or PxDDR is done with global interrupts disabled (interrupt mask bit set). In this case, the semaphore is changed to '1' when the level change is detected. Detecting a missed interrupt is done after the global interrupts are enabled (interrupt mask bit reset) and by checking the status of the semaphore. If it is '1' this means that the last interrupt was missed and the interrupt routine is invoked with the call instruction.

To implement the workaround, the following software sequence is to be followed for writing into the PxOR/PxDDR registers. The example is for for Port PF1 with falling edge interrupt sensitivity. The software sequence is given for both cases (global interrupt disabled/enabled).

**Case 1:** Writing to PxOR or PxDDR with Global Interrupts Enabled:

LD A,#01

LD sema,A ; set the semaphore to '1'

LD A, PFDR

AND A,#02

LD X,A ; store the level before writing to PxOR/PxDDR

LD A,#\$90

LD PFDDR,A ; Write to PFDDR

LD A,#\$ff

LD PFOR,A ; Write to PFOR

LD A, PFDR

AND A,#02

LD Y,A ; store the level after writing to PxOR/PxDDR

LD A,X ; check for falling edge

cp A,#02

jrne OUT

TNZ Y

jrne OUT

LD A,sema ; check the semaphore status if edge is detected

CP A,#01

jrne OUT

call call\_routine; call the interrupt routine

OUT:LD A,#00

LD sema,A

.call\_routine ; entry to call\_routine

PUSH A

PUSH X

PUSH CC

.ext1\_rt ; entry to interrupt routine

LD A,#00

