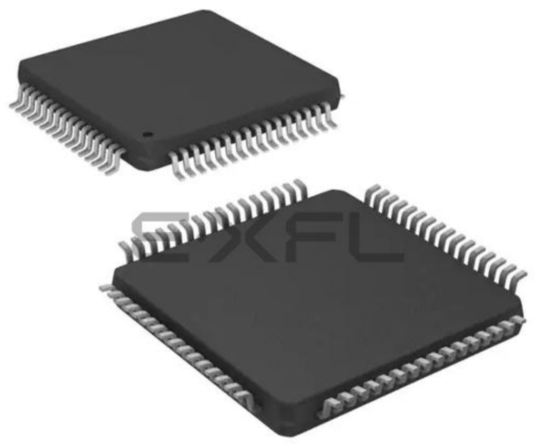
Microchip Technology - PIC17LC752T-08/PT Datasheet



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Applications of "Embedded - Microcontrollers"

Details	
Product Status	Obsolete
Core Processor	PIC
Core Size	8-Bit
Speed	8MHz
Connectivity	I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	50
Program Memory Size	16KB (8K x 16)
Program Memory Type	OTP
EEPROM Size	-
RAM Size	678 x 8
Voltage - Supply (Vcc/Vdd)	3V ~ 5.5V
Data Converters	A/D 12x10b
Oscillator Type	External
Operating Temperature	0°C ~ 70°C (TA)
Mounting Type	Surface Mount
Package / Case	64-TQFP
Supplier Device Package	64-TQFP (10x10)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic17lc752t-08-pt

TABLE 1-1: PIC17CXXX FAMILY OF DEVICES

Features		PIC17C42A	PIC17C43	PIC17C44	PIC17C752	PIC17C756A	PIC17C762	PIC17C766
Maximum Frequency of Operation		33 MHz	33 MHz	33 MHz	33 MHz	33 MHz	33 MHz	33 MHz
Operating Voltage Range		2.5 - 6.0V	2.5 - 6.0V	2.5 - 6.0V	3.0 - 5.5V	3.0 - 5.5V	3.0 - 5.5V	3.0 - 5.5V
Program Memory (x16)	(EPROM)	2 K	4 K	8 K	8 K	16 K	8 K	16 K
	(ROM)	—	—	—	—	—	—	—
Data Memory (bytes)		232	454	454	678	902	678	902
Hardware Multiplier (8 x 8)		Yes	Yes	Yes	Yes	Yes	Yes	Yes
Timer0 (16-bit + 8-bit postscaler)		Yes	Yes	Yes	Yes	Yes	Yes	Yes
Timer1 (8-bit)		Yes	Yes	Yes	Yes	Yes	Yes	Yes
Timer2 (8-bit)		Yes	Yes	Yes	Yes	Yes	Yes	Yes
Timer3 (16-bit)		Yes	Yes	Yes	Yes	Yes	Yes	Yes
Capture inputs (16-bit)		2	2	2	4	4	4	4
PWM outputs (up to 10-bit)		2	2	2	3	3	3	3
USART/SCI		1	1	1	2	2	2	2
A/D channels (10-bit)		—	—	—	12	12	16	16
SSP (SPI/I^{2}C w/Master mode)		—	—	—	Yes	Yes	Yes	Yes
Power-on Reset		Yes	Yes	Yes	Yes	Yes	Yes	Yes
Watchdog Timer		Yes	Yes	Yes	Yes	Yes	Yes	Yes
External Interrupts		Yes	Yes	Yes	Yes	Yes	Yes	Yes
Interrupt Sources		11	11	11	18	18	18	18
Code Protect		Yes	Yes	Yes	Yes	Yes	Yes	Yes
Brown-out Reset		—	—	—	Yes	Yes	Yes	Yes
In-Circuit Serial Programming		—	—	—	Yes	Yes	Yes	Yes
I/O Pins		33	33	33	50	50	66	66
I/O High Current Capability	Source	25 mA	25 mA	25 mA	25 mA	25 mA	25 mA	25 mA
	Sink	25 mA(1)	25 mA(1)	25 mA(1)	25 mA(1)	25 mA(1)	25 mA(1)	25 mA(1)
Package Types		40-pin DIP 44-pin PLCC 44-pin MQFP 44-pin TQFP	40-pin DIP 44-pin PLCC 44-pin MQFP 44-pin TQFP	40-pin DIP 44-pin PLCC 44-pin MQFP 44-pin TQFP	64-pin TQFP 68-pin PLCC	64-pin TQFP 68-pin PLCC	80-pin TQFP 84-pin PLCC	80-pin TQFP 84-pin PLCC

Note 1: Pins RA2 and RA3 can sink up to 60 mA.

NOTES:

6.0 INTERRUPTS

PIC17C7XX devices have 18 sources of interrupt:

- External interrupt from the RA0/INT pin
- Change on RB7:RB0 pins
- TMR0 Overflow
- TMR1 Overflow
- TMR2 Overflow
- TMR3 Overflow
- USART1 Transmit buffer empty
- USART1 Receive buffer full
- USART2 Transmit buffer empty
- USART2 Receive buffer full
- SSP Interrupt
- SSP I^2C bus collision interrupt
- A/D conversion complete
- Capture1
- Capture2
- Capture3
- Capture4
- T0CKI edge occurred

There are six registers used in the control and status of interrupts. These are:

- CPUSTA
- INTSTA
- PIE1
- PIR1
- PIE2
- PIR2

The CPUSTA register contains the GLINTD bit. This is the Global Interrupt Disable bit. When this bit is set, all interrupts are disabled. This bit is part of the controller core functionality and is described in the Section 6.4.

When an interrupt is responded to, the GLINTD bit is automatically set to disable any further interrupts, the return address is pushed onto the stack and the PC is loaded with the interrupt vector address. There are four interrupt vectors. Each vector address is for a specific interrupt source (except the peripheral interrupts, which all vector to the same address). These sources are:

- External interrupt from the RA0/INT pin
- TMR0 Overflow
- T0CKI edge occurred
- Any peripheral interrupt

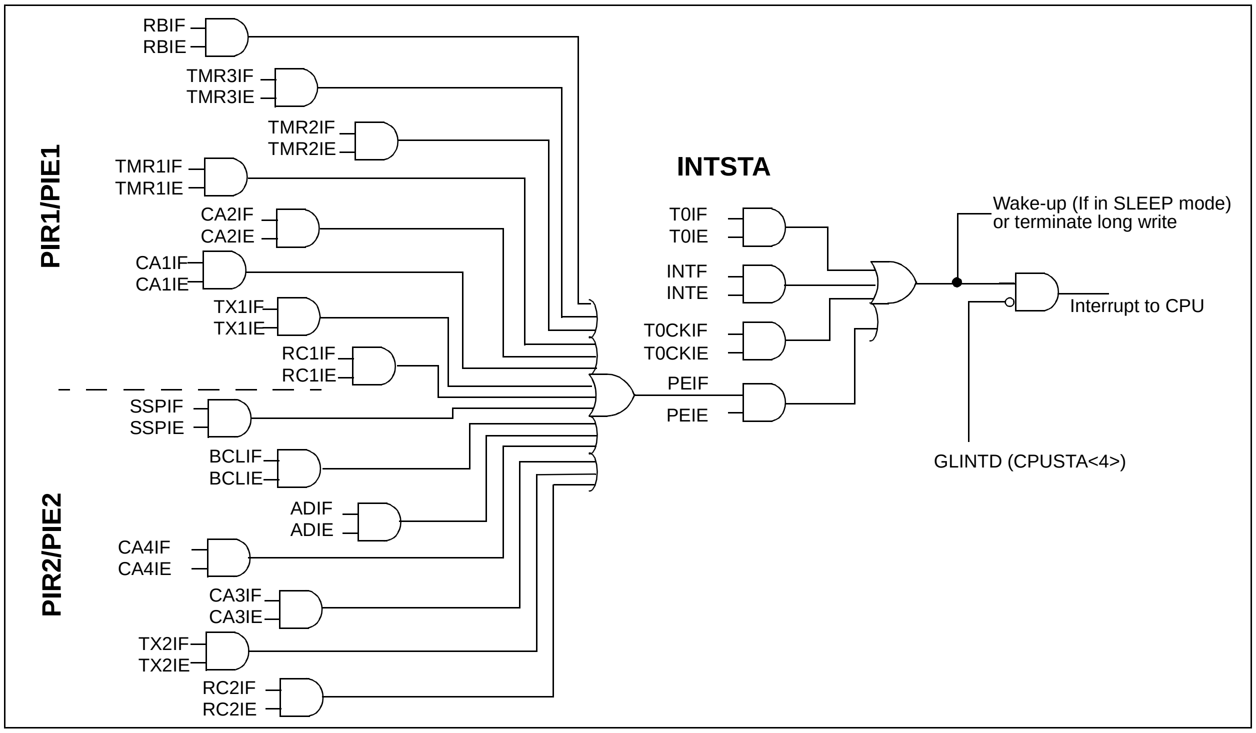
When program execution vectors to one of these interrupt vector addresses (except for the peripheral interrupts), the interrupt flag bit is automatically cleared. Vectoring to the peripheral interrupt vector address does not automatically clear the source of the interrupt. In the peripheral Interrupt Service Routine, the source(s) of the interrupt can be determined by testing the interrupt flag bits. The interrupt flag bit(s) must be cleared in software before re-enabling interrupts to avoid infinite interrupt requests.

When an interrupt condition is met, that individual interrupt flag bit will be set, regardless of the status of its corresponding mask bit or the GLINTD bit.

For external interrupt events, there will be an interrupt latency. For two-cycle instructions, the latency could be one instruction cycle longer.

The “return from interrupt” instruction, `RETFIE`, can be used to mark the end of the Interrupt Service Routine. When this instruction is executed, the stack is “POPed” and the GLINTD bit is cleared (to re-enable interrupts).

FIGURE 6-1: INTERRUPT LOGIC



8.1 Table Writes to Internal Memory

A table write operation to internal memory causes a long write operation. The long write is necessary for programming the internal EPROM. Instruction execution is halted while in a long write cycle. The long write will be terminated by any enabled interrupt. To ensure that the EPROM location has been well programmed, a minimum programming time is required (see specification #D114). Having only one interrupt enabled to terminate the long write ensures that no unintentional interrupts will prematurely terminate the long write.

The sequence of events for programming an internal program memory location should be:

1. Disable all interrupt sources, except the source to terminate EPROM program write.
2. Raise $\overline{\text{MCLR}}$/VPP pin to the programming voltage.
3. Clear the WDT.
4. Do the table write. The interrupt will terminate the long write.
5. Verify the memory location (table read).

Note 1: Programming requirements must be met. See timing specification in electrical specifications for the desired device. Violating these specifications (including temperature) may result in EPROM locations that are not fully programmed and may lose their state over time.

2: If the VPP requirement is not met, the table write is a 2-cycle write and the program memory is unchanged.

8.1.1 TERMINATING LONG WRITES

An interrupt source or RESET are the only events that terminate a long write operation. Terminating the long write from an interrupt source requires that the interrupt enable and flag bits are set. The GLINTD bit only enables the vectoring to the interrupt address.

If the T0CKI, RA0/INT, or TMR0 interrupt source is used to terminate the long write, the interrupt flag of the highest priority enabled interrupt, will terminate the long write and automatically be cleared.

Note 1: If an interrupt is pending, the `TABLWT` is aborted (a `NOP` is executed). The highest priority pending interrupt, from the T0CKI, RA0/INT, or TMR0 sources that is enabled, has its flag cleared.

2: If the interrupt is not being used for the program write timing, the interrupt should be disabled. This will ensure that the interrupt is not lost, nor will it terminate the long write prematurely.

If a peripheral interrupt source is used to terminate the long write, the interrupt enable and flag bits must be set. The interrupt flag will not be automatically cleared upon the vectoring to the interrupt vector address.

The GLINTD bit determines whether the program will branch to the interrupt vector when the long write is terminated. If GLINTD is clear, the program will vector, if GLINTD is set, the program will not vector to the interrupt address.

TABLE 8-1: INTERRUPT - TABLE WRITE INTERACTION

Interrupt Source	GLINTD	Enable Bit	Flag Bit	Action
RA0/INT, TMR0, T0CKI	0	1	1	Terminate long table write (to internal program memory), branch to interrupt vector (branch clears flag bit).
	0	1	0	None.
	1	0	x	None.
	1	1	1	Terminate long table write, do not branch to interrupt vector (flag is automatically cleared).
Peripheral	0	1	1	Terminate long table write, branch to interrupt vector.
	0	1	0	None.
	1	0	x	None.
	1	1	1	Terminate long table write, do not branch to interrupt vector (flag remains set).

10.0 I/O PORTS

PIC17C75X devices have seven I/O ports, PORTA through PORTG. PIC17C76X devices have nine I/O ports, PORTA through PORTJ. PORTB through PORTJ have a corresponding Data Direction Register (DDR), which is used to configure the port pins as inputs or outputs. Some of these ports pins are multiplexed with alternate functions.

PORTC, PORTD, and PORTE are multiplexed with the system bus. These pins are configured as the system bus when the device's configuration bits are selected to Microprocessor or Extended Microcontroller modes. In the two other microcontroller modes, these pins are general purpose I/O.

PORTA, PORTB, PORTE<3>, PORTF, PORTG and the upper four bits of PORTH are multiplexed with the peripheral features of the device. These peripheral features are:

- Timer Modules
- Capture Modules
- PWM Modules
- USART/SCI Modules
- SSP Module
- A/D Module
- External Interrupt pin

When some of these peripheral modules are turned on, the port pin will automatically configure to the alternate function. The modules that do this are:

- PWM Module
- SSP Module
- USART/SCI Module

When a pin is automatically configured as an output by a peripheral module, the pins data direction (DDR) bit is unknown. After disabling the peripheral module, the user should re-initialize the DDR bit to the desired configuration.

The other peripheral modules (which require an input) must have their data direction bits configured appropriately.

Note: A pin that is a peripheral input, can be configured as an output (DDRx<y> is cleared). The peripheral events will be determined by the action output on the port pin.

When the device enters the "RESET state", the Data Direction registers (DDR) are forced set, which will make the I/O hi-impedance inputs. The RESET state of some peripheral modules may force the I/O to other operations, such as analog inputs or the system bus.

TABLE 10-3: PORTB FUNCTIONS

Name	Bit	Buffer Type	Function
RB0/CAP1	bit0	ST	Input/output or the Capture1 input pin. Software programmable weak pull-up and interrupt-on-change features.
RB1/CAP2	bit1	ST	Input/output or the Capture2 input pin. Software programmable weak pull-up and interrupt-on-change features.
RB2/PWM1	bit2	ST	Input/output or the PWM1 output pin. Software programmable weak pull-up and interrupt-on-change features.
RB3/PWM2	bit3	ST	Input/output or the PWM2 output pin. Software programmable weak pull-up and interrupt-on-change features.
RB4/TCLK12	bit4	ST	Input/output or the external clock input to Timer1 and Timer2. Software programmable weak pull-up and interrupt-on-change features.
RB5/TCLK3	bit5	ST	Input/output or the external clock input to Timer3. Software programmable weak pull-up and interrupt-on-change features.
RB6/SCK	bit6	ST	Input/output or the Master/Slave clock for the SPI. Software programmable weak pull-up and interrupt-on-change features.
RB7/SDO	bit7	ST	Input/output or data output for the SPI. Software programmable weak pull-up and interrupt-on-change features.

Legend: ST = Schmitt Trigger input

TABLE 10-4: REGISTERS/BITS ASSOCIATED WITH PORTB

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	$\overline{\text{MCLR}}$, WDT
12h, Bank 0	PORTB	RB7/ SDO	RB6/ SCK	RB5/ TCLK3	RB4/ TCLK12	RB3/ PWM2	RB2/ PWM1	RB1/ CAP2	RB0/ CAP1	`xxxx xxxx`	`uuuu uuuu`
11h, Bank 0	DDRB	Data Direction Register for PORTB								`1111 1111`	`1111 1111`
10h, Bank 0	PORTA	$\overline{\text{RBPU}}$	—	RA5/ TX1/CK1	RA4/ RX1/DT1	RA3/ SDI/SDA	RA2/ $\overline{\text{SS}}$/SCL	RA1/T0CKI	RA0/INT	`0-xx 11xx`	`0-uu 11uu`
06h, Unbanked	CPUSTA	—	—	STKAV	GLINTD	TO	PD	$\overline{\text{POR}}$	$\overline{\text{BOR}}$	`--11 11qq`	`--11 qquu`
07h, Unbanked	INTSTA	PEIF	T0CKIF	T0IF	INTF	PEIE	T0CKIE	T0IE	INTE	`0000 0000`	`0000 0000`
16h, Bank 1	PIR1	RBIF	TMR3IF	TMR2IF	TMR1IF	CA2IF	CA1IF	TX1IF	RC1IF	`x000 0010`	`u000 0010`
17h, Bank 1	PIE1	RBIE	TMR3IE	TMR2IE	TMR1IE	CA2IE	CA1IE	TX1IE	RC1IE	`0000 0000`	`0000 0000`
16h, Bank 3	TCON1	CA2ED1	CA2ED0	CA1ED1	CA1ED0	T16	TMR3CS	TMR2CS	TMR1CS	`0000 0000`	`0000 0000`
17h, Bank 3	TCON2	CA2OVF	CA1OVF	PWM2ON	PWM1ON	CA1/$\overline{\text{PR3}}$	TMR3ON	TMR2ON	TMR1ON	`0000 0000`	`0000 0000`

Legend: `x` = unknown, `u` = unchanged, `-` = unimplemented, read as '0', q = value depends on condition. Shaded cells are not used by PORTB.

TABLE 10-5: PORTC FUNCTIONS

Name	Bit	Buffer Type	Function
RC0/AD0	bit0	TTL	Input/output or system bus address/data pin.
RC1/AD1	bit1	TTL	Input/output or system bus address/data pin.
RC2/AD2	bit2	TTL	Input/output or system bus address/data pin.
RC3/AD3	bit3	TTL	Input/output or system bus address/data pin.
RC4/AD4	bit4	TTL	Input/output or system bus address/data pin.
RC5/AD5	bit5	TTL	Input/output or system bus address/data pin.
RC6/AD6	bit6	TTL	Input/output or system bus address/data pin.
RC7/AD7	bit7	TTL	Input/output or system bus address/data pin.

Legend: TTL = TTL input

TABLE 10-6: REGISTERS/BITS ASSOCIATED WITH PORTC

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	$\overline{\text{MCLR}}$, WDT
11h, Bank 1	PORTC	RC7/ AD7	RC6/ AD6	RC5/ AD5	RC4/ AD4	RC3/ AD3	RC2/ AD2	RC1/ AD1	RC0/ AD0	`xxxx xxxx`	`uuuu uuuu`
10h, Bank 1	DDRC	Data Direction Register for PORTC								`1111 1111`	`1111 1111`

Legend: `x` = unknown, `u` = unchanged

10.10.2 SUCCESSIVE OPERATIONS ON I/O PORTS

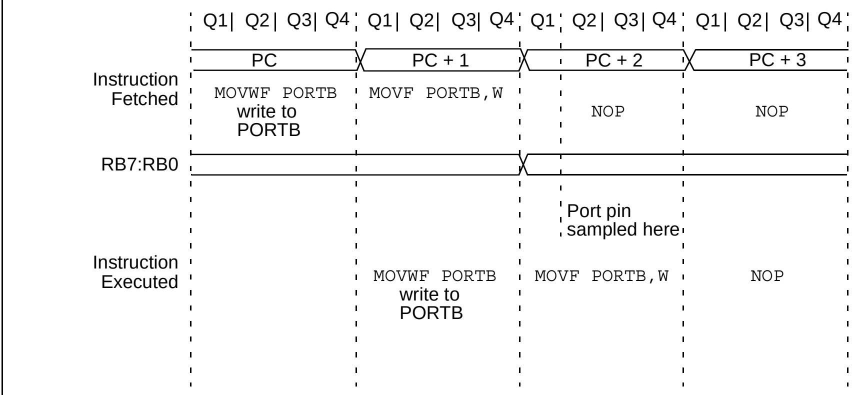
The actual write to an I/O port happens at the end of an instruction cycle, whereas for reading, the data must be valid at the beginning of the instruction cycle (Figure 10-20). Therefore, care must be exercised if a write followed by a read operation is carried out on the same I/O port. The sequence of instructions should be such to allow the pin voltage to stabilize (load dependent) before executing the instruction that reads the values on that I/O port. Otherwise, the previous state of that pin may be read into the CPU, rather than the "new" state. When in doubt, it is better to separate these instructions with a `NOP`, or another instruction not accessing this I/O port.

Figure 10-21 shows the I/O model which causes this situation. As the effective capacitance (C) becomes larger, the rise/fall time of the I/O pin increases. As the device frequency increases, or the effective capacitance increases, the possibility of this subsequent PORTx read-modify-write instruction issue increases. This effective capacitance includes the effects of the board traces.

The best way to address this is to add a series resistor at the I/O pin. This resistor allows the I/O pin to get to the desired level before the next instruction.

The use of `NOP` instructions between the subsequent PORTx read-modify-write instructions, is a lower cost solution, but has the issue that the number of `NOP` instructions is dependent on the effective capacitance C and the frequency of the device.

FIGURE 10-20: SUCCESSIVE I/O OPERATION



Note:

This example shows a write to PORTB, followed by a read from PORTB.

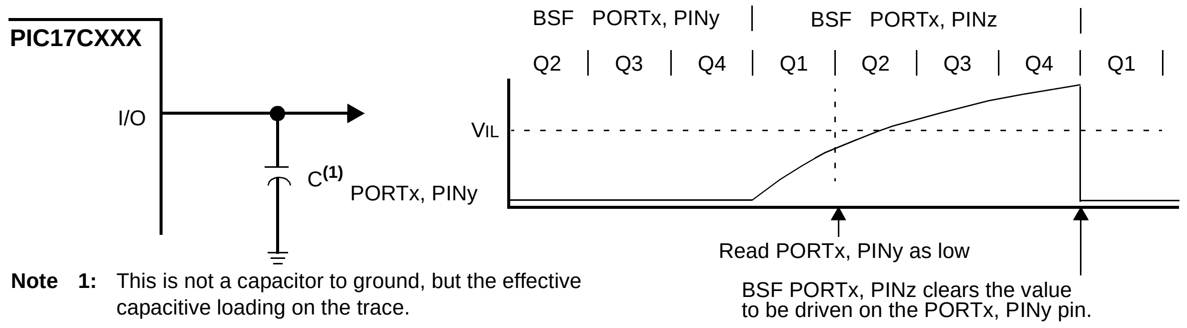
Note that:

data setup time = (0.25TCY - TPD)

where TCY = instruction cycle
TPD = propagation delay

Therefore, at higher clock frequencies, a write followed by a read may be problematic.

FIGURE 10-21: I/O CONNECTION ISSUES



Note 1: This is not a capacitor to ground, but the effective capacitive loading on the trace.

12.1 Timer0 Operation

When the T0CS (T0STA<5>) bit is set, TMR0 increments on the internal clock. When T0CS is clear, TMR0 increments on the external clock (RA1/T0CKI pin). The external clock edge can be selected in software. When the T0SE (T0STA<6>) bit is set, the timer will increment on the rising edge of the RA1/T0CKI pin. When T0SE is clear, the timer will increment on the falling edge of the RA1/T0CKI pin. The prescaler can be programmed to introduce a prescale of 1:1 to 1:256. The timer increments from 0000h to FFFFh and rolls over to 0000h. On overflow, the TMR0 Interrupt Flag bit (T0IF) is set. The TMR0 interrupt can be masked by clearing the corresponding TMR0 Interrupt Enable bit (T0IE). The TMR0 Interrupt Flag bit (T0IF) is automatically cleared when vectoring to the TMR0 interrupt vector.

12.2 Using Timer0 with External Clock

When an external clock input is used for Timer0, it is synchronized with the internal phase clocks. Figure 12-2 shows the synchronization of the external clock. This synchronization is done after the prescaler. The output of the prescaler (PSOUT) is sampled twice in every instruction cycle to detect a rising or a falling edge. The timing requirements for the external clock are detailed in the electrical specification section.

12.2.1 DELAY FROM EXTERNAL CLOCK EDGE

Since the prescaler output is synchronized with the internal clocks, there is a small delay from the time the external clock edge occurs to the time TMR0 is actually incremented. Figure 12-2 shows that this delay is between 3TOSC and 7TOSC. Thus, for example, measuring the interval between two edges (e.g. period) will be accurate within ±4TOSC (±121 ns @ 33 MHz).

FIGURE 12-1: TIMER0 MODULE BLOCK DIAGRAM

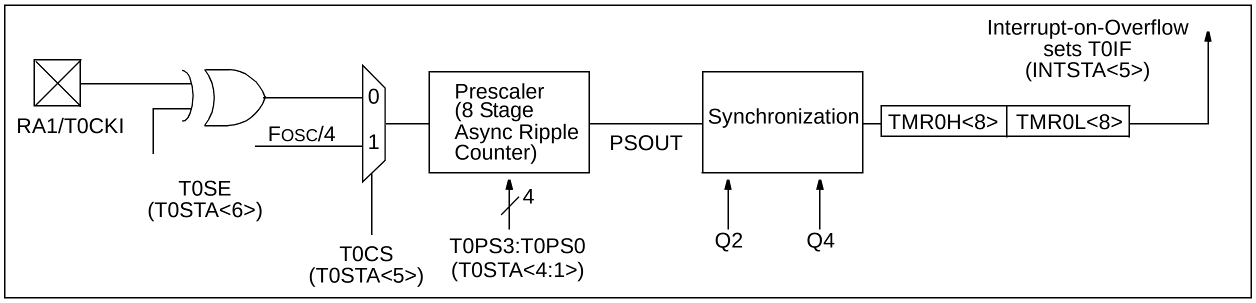
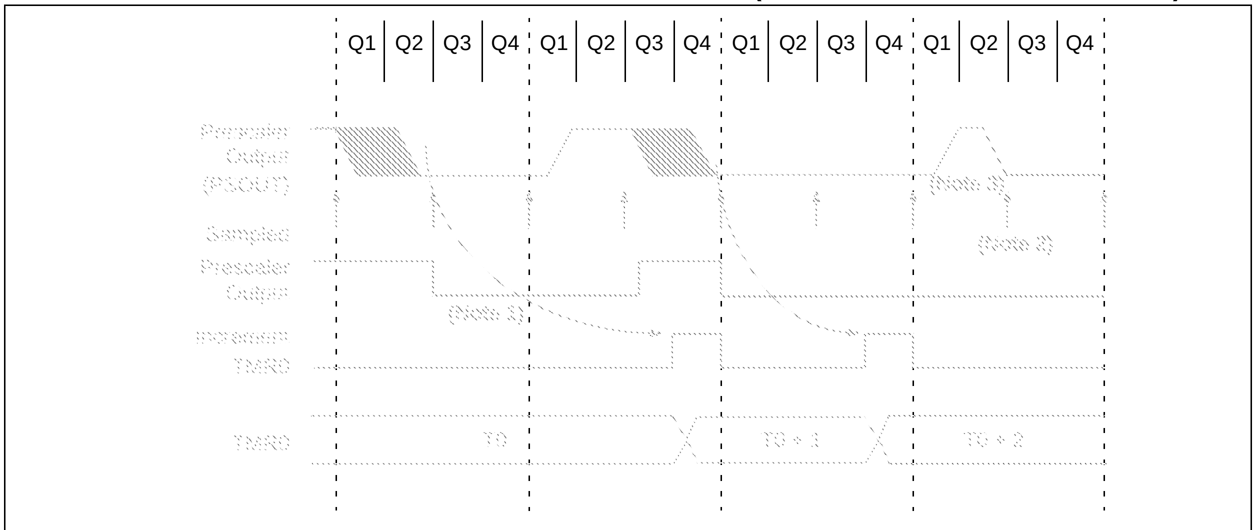


FIGURE 12-2: TMR0 TIMING WITH EXTERNAL CLOCK (INCREMENT ON FALLING EDGE)



Note 1: The delay from the T0CKI edge to the TMR0 increment is 3Tosc to 7Tosc.
2: ↑ = PSOUT is sampled here
3: The PSOUT high time is too short and is missed by the sampling circuit.

15.2.1.3 Slave Transmission

When the R/$\overline{W}$ bit of the incoming address byte is set and an address match occurs, the R/$\overline{W}$ bit of the SSPSTAT register is set. The received address is loaded into the SSPBUF register. The $\overline{ACK}$ pulse will be sent on the ninth bit, and the SCL pin is held low. The transmit data must be loaded into the SSPBUF register, which also loads the SSPSR register. Then SCL pin should be enabled by setting bit CKP (SSPCON1<4>). The master must monitor the SCL pin prior to asserting another clock pulse. The slave devices may be holding off the master by stretching the clock. The eight data bits are shifted out on the falling edge of the SCL input. This ensures that the SDA signal is valid during the SCL high time (Figure 15-13).

An SSP interrupt is generated for each data transfer byte. The SSPIF flag bit must be cleared in software, and the SSPSTAT register is used to determine the status of the byte transfer. The SSPIF flag bit is set on the falling edge of the ninth clock pulse.

As a slave-transmitter, the $\overline{ACK}$ pulse from the master-receiver is latched on the rising edge of the ninth SCL input pulse. If the SDA line was high (not $\overline{ACK}$), then the data transfer is complete. When the not $\overline{ACK}$ is latched by the slave, the slave logic is reset and the slave then monitors for another occurrence of the START bit. If the SDA line was low ($\overline{ACK}$), the transmit data must be loaded into the SSPBUF register, which also loads the SSPSR register. Then, the SCL pin should be enabled by setting the CKP bit.

FIGURE 15-12: I^{2}C WAVEFORMS FOR RECEPTION (7-BIT ADDRESS)

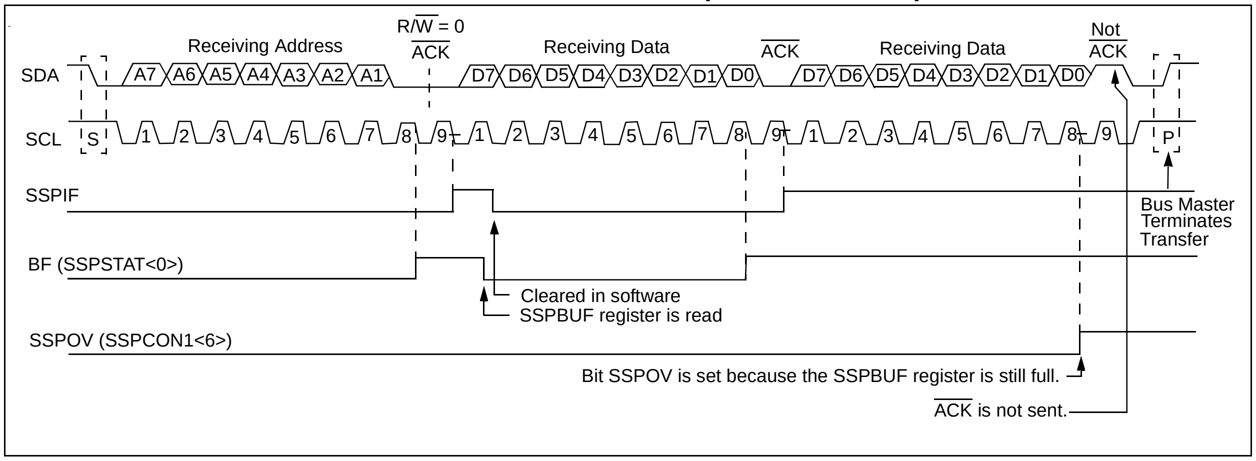
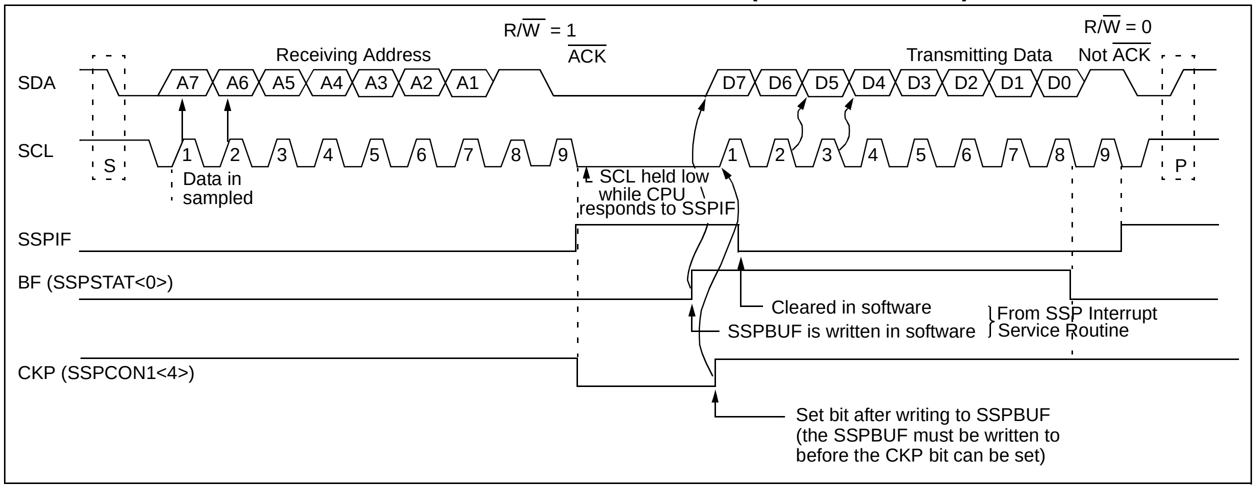


FIGURE 15-13: I^{2}C WAVEFORMS FOR TRANSMISSION (7-BIT ADDRESS)


R/W = 1
ACK
Receiving Address
Transmitting Data
R/W = 0
Not ACK
SDA
A7 A6 A5 A4 A3 A2 A1
D7 D6 D5 D4 D3 D2 D1 D0
SCL
S
1 2 3 4 5 6 7 8 9
P
Data in sampled
SCL held low while CPU responds to SSPIF
SSPIF
BF (SSPSTAT<0>)
Cleared in software
SSPBUF is written in software
From SSP Interrupt Service Routine
CKP (SSPCON1<4>)
Set bit after writing to SSPBUF (the SSPBUF must be written to before the CKP bit can be set)

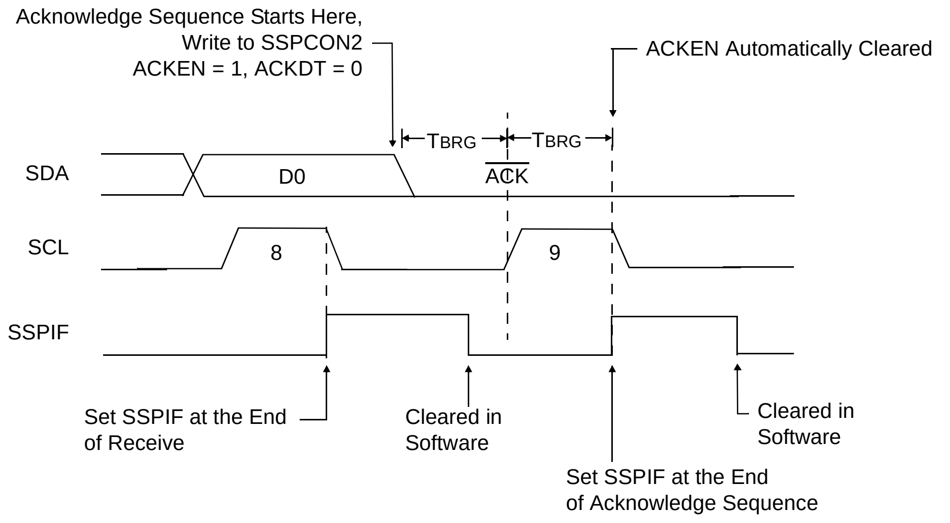
15.2.13 ACKNOWLEDGE SEQUENCE TIMING

An acknowledge sequence is enabled by setting the acknowledge sequence enable bit, ACKEN (SSPCON2<4>). When this bit is set, the SCL pin is pulled low and the contents of the acknowledge data bit is presented on the SDA pin. If the user wishes to generate an acknowledge, then the ACKDT bit should be cleared. If not, the user should set the ACKDT bit before starting an acknowledge sequence. The baud rate generator then counts for one rollover period (TBRG), and the SCL pin is de-asserted (pulled high). When the SCL pin is sampled high (clock arbitration), the baud rate generator counts for TBRG. The SCL pin is then pulled low. Following this, the ACKEN bit is automatically cleared, the baud rate generator is turned off and the SSP module then goes into IDLE mode (Figure 15-29).

15.2.13.1 WCOL Status Flag

If the user writes the SSPBUF when an acknowledge sequence is in progress, then WCOL is set and the contents of the buffer are unchanged (the write doesn't occur).

FIGURE 15-29: ACKNOWLEDGE SEQUENCE WAVEFORM


Acknowledge Sequence Starts Here,
Write to SSPCON2
ACKEN = 1, ACKDT = 0
ACKEN Automatically Cleared
TBRG
TBRG
SDA
D0
$\overline{ACK}$
SCL
8
9
SSPIF
Set SSPIF at the End of Receive
Cleared in Software
Cleared in Software
Set SSPIF at the End of Acknowledge Sequence

Note: TBRG = one baud rate generator period.

15.2.15 CLOCK ARBITRATION

Clock arbitration occurs when the master, during any receive, transmit, or Repeated Start/Stop condition, de-asserts the SCL pin (SCL allowed to float high). When the SCL pin is allowed to float high, the baud rate generator (BRG) is suspended from counting until the SCL pin is actually sampled high. When the SCL pin is sampled high, the baud rate generator is reloaded with the contents of SSPADD<6:0> and begins counting. This ensures that the SCL high time will always be at least one BRG rollover count, in the event that the clock is held low by an external device (Figure 15-33).

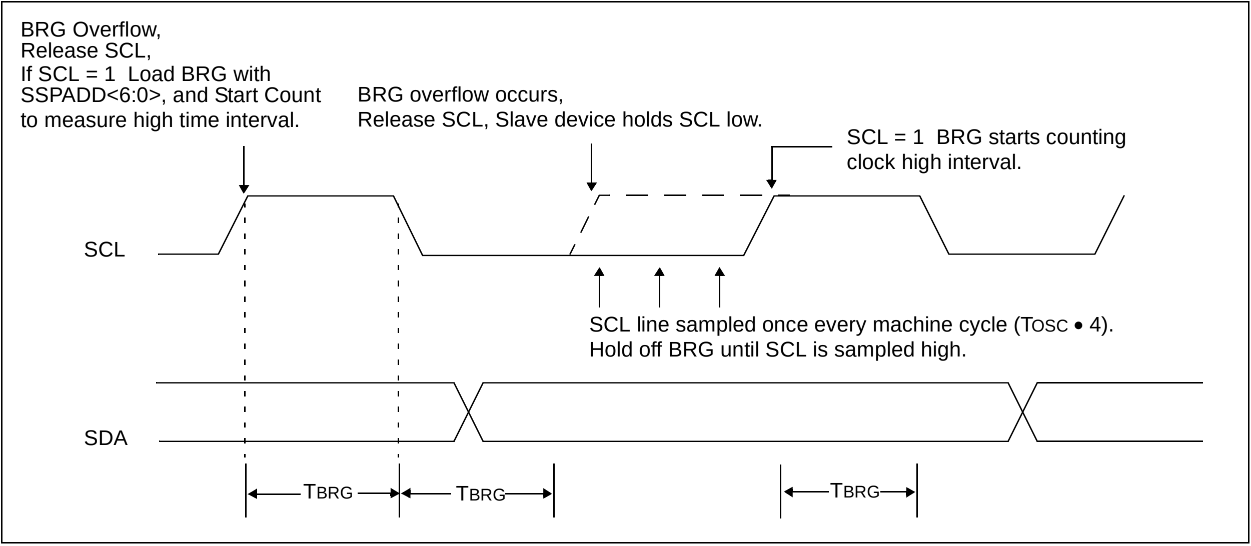
15.2.16 SLEEP OPERATION

While in SLEEP mode, the I^2C module can receive addresses or data and when an address match or complete byte transfer occurs, wake the processor from SLEEP (if the SSP interrupt is enabled).

15.2.17 EFFECTS OF A RESET

A RESET disables the SSP module and terminates the current transfer.

FIGURE 15-33: CLOCK ARBITRATION TIMING IN MASTER TRANSMIT MODE


BRG Overflow,
Release SCL,
If SCL = 1 Load BRG with
SSPADD<6:0>, and Start Count
to measure high time interval.
BRG overflow occurs,
Release SCL, Slave device holds SCL low.
SCL = 1 BRG starts counting
clock high interval.
SCL
SCL line sampled once every machine cycle (TOSC • 4).
Hold off BRG until SCL is sampled high.
SDA
TBRG
TBRG
TBRG

15.2.18.3 Bus Collision During a STOP Condition

Bus collision occurs during a STOP condition if:

a) After the SDA pin has been de-asserted and allowed to float high, SDA is sampled low after the BRG has timed out.

b) After the SCL pin is de-asserted, SCL is sampled low before SDA goes high.

The STOP condition begins with SDA asserted low. When SDA is sampled low, the SCL pin is allowed to float. When the pin is sampled high (clock arbitration), the baud rate generator is loaded with SSPADD<6:0> and counts down to ‘0’. After the BRG times out, SDA is sampled. If SDA is sampled low, a bus collision has occurred. This is due to another master attempting to drive a data '0'. If the SCL pin is sampled low before SDA is allowed to float high, a bus collision occurs. This is another case of another master attempting to drive a data '0' (Figure 15-40).

FIGURE 15-40: BUS COLLISION DURING A STOP CONDITION (CASE 1)

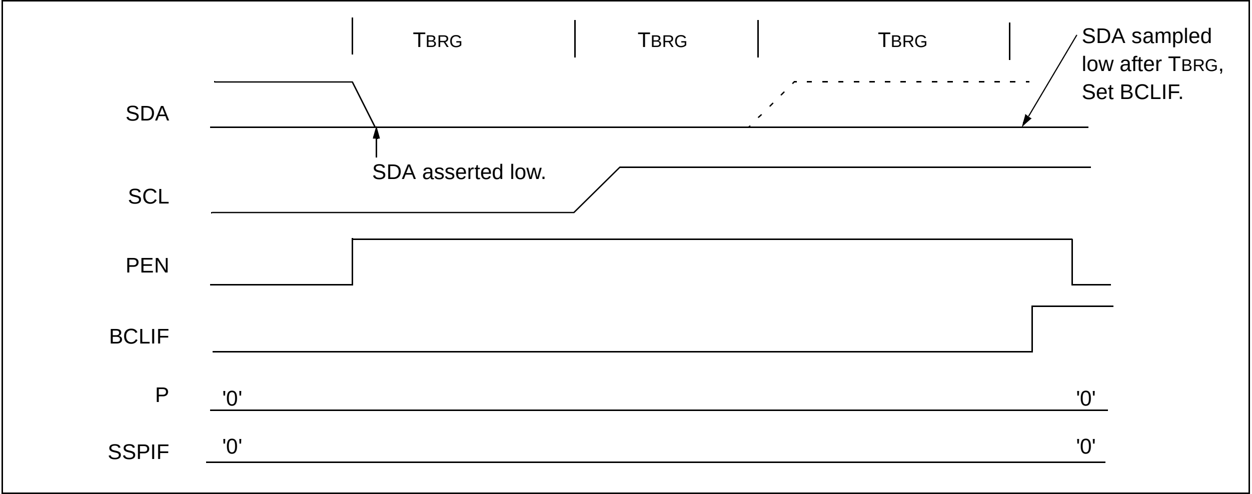
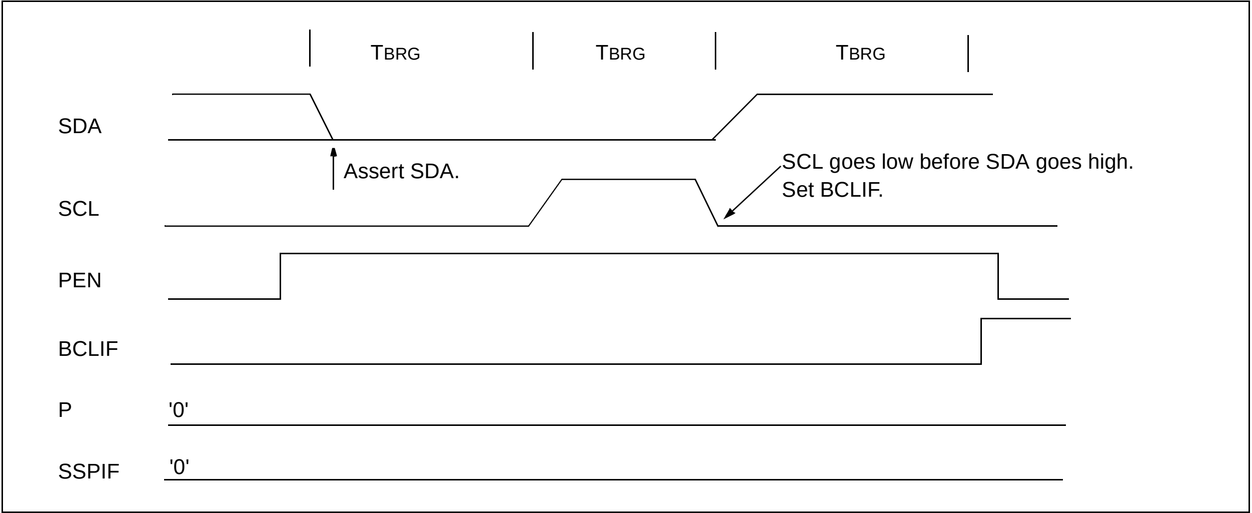


FIGURE 15-41: BUS COLLISION DURING A STOP CONDITION (CASE 2)



15.4 Example Program

Example 15-2 shows MPLAB® C17 ’C’ code for using the I^{2}C module in Master mode to communicate with a 24LC01B serial EEPROM. This example uses the PIC® MCU ‘C’ libraries included with MPLAB C17.

EXAMPLE 15-2: INTERFACING TO A 24LC01B SERIAL EEPROM (USING MPLAB C17)

```
// Include necessary header files
#include <p17c756.h>     // Processor header file
#include <delays.h>      // Delay routines header file
#include <stdlib.h>      // Standard Library header file
#include <i2c16.h>       // I2C routines header file

#define CONTROL 0xa0     // Control byte definition for 24LC01B

// Function declarations
void main(void);
void WritePORTD(static unsigned char data);
void ByteWrite(static unsigned char address,static unsigned char data);
unsigned char ByteRead(static unsigned char address);
void ACKPoll(void);

// Main program
void main(void)
{
static unsigned char address;    // I2C address of 24LC01B
static unsigned char datao;      // Data written to 24LC01B
static unsigned char datai;      // Data read from 24LC01B

    address = 0;                 // Preset address to 0
    OpenI2C(MASTER,SLEW_ON);     // Configure I2C Module Master mode, Slew rate control on
    SSPADD = 39;                 // Configure clock for 100KHz

    while(address<128)           // Loop 128 times, 24LC01B is 128x8
    {
        datao = PORTB;
        do
        {
            ByteWrite(address,datao);   // Write data to EEPROM
            ACKPoll();                  // Poll the 24LC01B for state
            datai = ByteRead(address);  // Read data from EEPROM into SSPBUF
        } while(datai != datao);        // Loop as long as data not correctly
                                        //    written to 24LC01B

        address++;                      // Increment address
    }
    while(1)                            // Done writing 128 bytes to 24LC01B, Loop forever
    {
        Nop();
    }
}
```

BSF **Bit Set f**

Syntax:	[*label*] BSF f,b
Operands:	0 ≤ f ≤ 255 0 ≤ b ≤ 7
Operation:	1 → (f<b>)
Status Affected:	None
Encoding:	`1000` `0bbb` `ffff` `ffff`
Description:	Bit 'b' in register 'f' is set.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: `BSF FLAG_REG, 7`

Before Instruction
FLAG_REG = 0x0A

After Instruction
FLAG_REG = 0x8A

BTFSC **Bit Test, skip if Clear**

Syntax:	[*label*] BTFSC f,b
Operands:	0 ≤ f ≤ 255 0 ≤ b ≤ 7
Operation:	skip if (f<b>) = 0
Status Affected:	None
Encoding:	`1001` `1bbb` `ffff` `ffff`
Description:	If bit 'b' in register 'f' is 0, then the next instruction is skipped. If bit 'b' is 0, then the next instruction fetched during the current instruction execution is discarded and a `NOP` is executed instead, making this a two-cycle instruction.
Words:	1
Cycles:	1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

Example:

```
HERE    BTFSC   FLAG,1
FALSE   :
TRUE    :
```

Before Instruction
PC = address `(HERE)`

After Instruction
If FLAG<1> = 0;
PC = address `(TRUE)`
If FLAG<1> = 1;
PC = address `(FALSE)`

INCF — Increment f

Syntax:	[*label*] INCF f,d
Operands:	$0 \le f \le 255$ $d \in [0,1]$
Operation:	(f) + 1 $\rightarrow$ (dest)
Status Affected:	OV, C, DC, Z
Encoding:	`0001` `010d` `ffff` `ffff`
Description:	The contents of register 'f' are incremented. If 'd' is 0, the result is placed in WREG. If 'd' is 1, the result is placed back in register 'f'.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `INCF CNT, 1`

Before Instruction

CNT	=	0xFF
Z	=	0
C	=	?

After Instruction

CNT	=	0x00
Z	=	1
C	=	1

INCFSZ — Increment f, skip if 0

Syntax:	[*label*] INCFSZ f,d
Operands:	$0 \le f \le 255$ $d \in [0,1]$
Operation:	(f) + 1 $\rightarrow$ (dest) skip if result = 0
Status Affected:	None
Encoding:	`0001` `111d` `ffff` `ffff`
Description:	The contents of register 'f' are incremented. If 'd' is 0, the result is placed in WREG. If 'd' is 1, the result is placed back in register 'f'. If the result is 0, the next instruction, which is already fetched is discarded and a `NOP` is executed instead, making it a two-cycle instruction.
Words:	1
Cycles:	1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

Example:

```
HERE     INCFSZ   CNT, 1
NZERO    :
ZERO     :
```

Before Instruction

PC	=	Address `(HERE)`

After Instruction

CNT	=	CNT + 1
If CNT	=	0;
PC	=	Address `(ZERO)`
If CNT	≠	0;
PC	=	Address `(NZERO)`

TABLRD **Table Read**

Example1: `TABLRD 1, 1, REG ;`

Before Instruction

REG	=	0x53
TBLATH	=	0xAA
TBLATL	=	0x55
TBLPTR	=	0xA356
MEMORY(TBLPTR)	=	0x1234

After Instruction (table write completion)

REG	=	0xAA
TBLATH	=	0x12
TBLATL	=	0x34
TBLPTR	=	0xA357
MEMORY(TBLPTR)	=	0x5678

Example2: `TABLRD 0, 0, REG ;`

Before Instruction

REG	=	0x53
TBLATH	=	0xAA
TBLATL	=	0x55
TBLPTR	=	0xA356
MEMORY(TBLPTR)	=	0x1234

After Instruction (table write completion)

REG	=	0x55
TBLATH	=	0x12
TBLATL	=	0x34
TBLPTR	=	0xA356
MEMORY(TBLPTR)	=	0x1234

TABLWT **Table Write**

Syntax: [*label*] TABLWT t,i,f

Operands: $0 \le f \le 255$
$i \in [0,1]$
$t \in [0,1]$

Operation: If t = 0,
f → TBLATL;
If t = 1,
f → TBLATH;
TBLAT → Prog Mem (TBLPTR);
If i = 1,
TBLPTR + 1 → TBLPTR
If i = 0,
TBLPTR is unchanged

Status Affected: None

Encoding:

`1010`	`11ti`	`ffff`	`ffff`

Description:

1. Load value in ’f’ into 16-bit table latch (TBLAT)
   If t = 1: load into high byte;
   If t = 0: load into low byte
2. The contents of TBLAT are written to the program memory location pointed to by TBLPTR.
   If TBLPTR points to external program memory location, then the instruction takes two-cycle.
   If TBLPTR points to an internal EPROM location, then the instruction is terminated when an interrupt is received.

Note: The $\overline{\text{MCLR}}$/VPP pin must be at the programming voltage for successful programming of internal memory.
If $\overline{\text{MCLR}}$/VPP = VDD
the programming sequence of internal memory will be interrupted. A short write will occur (2 TCY). The internal memory location will not be affected.

3. The TBLPTR can be automatically incremented
   If i = 1; TBLPTR is not incremented
   If i = 0; TBLPTR is incremented

Words: 1

Cycles: 2 (many if write is to on-chip EPROM program memory)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register TBLATH or TBLATL
No operation	No operation (Table Pointer on Address bus)	No operation	No operation (Table Latch on Address bus, $\overline{\text{WR}}$ goes low)

XORLW Exclusive OR Literal with WREG

Syntax:	[*label*] XORLW k
Operands:	$0 \le k \le 255$
Operation:	(WREG) .XOR. k $\rightarrow$ (WREG)
Status Affected:	Z
Encoding:	`1011` `0100` `kkkk` `kkkk`
Description:	The contents of WREG are XOR'ed with the 8-bit literal 'k'. The result is placed in WREG.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to WREG

Example: `XORLW 0xAF`

Before Instruction

WREG	=	0xB5

After Instruction

WREG	=	0x1A

XORWF Exclusive OR WREG with f

Syntax:	[*label*] XORWF f,d
Operands:	$0 \le f \le 255$ $d \in [0,1]$
Operation:	(WREG) .XOR. (f) $\rightarrow$ (dest)
Status Affected:	Z
Encoding:	`0000` `110d` `ffff` `ffff`
Description:	Exclusive OR the contents of WREG with register 'f'. If 'd' is 0, the result is stored in WREG. If 'd' is 1, the result is stored back in the register 'f'.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `XORWF REG, 1`

Before Instruction

REG	=	0xAF	1010 1111
WREG	=	0xB5	1011 0101

After Instruction

REG	=	0x1A	0001 1010
WREG	=	0xB5	

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