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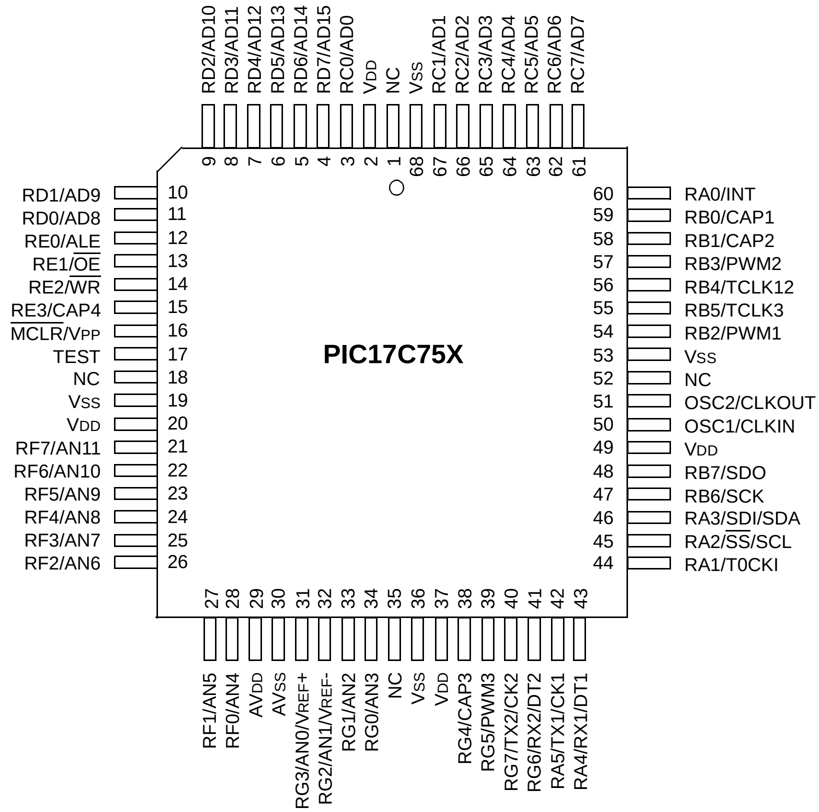
"Embedded - Microcontrollers" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

Applications of "Embedded - Microcontrollers"

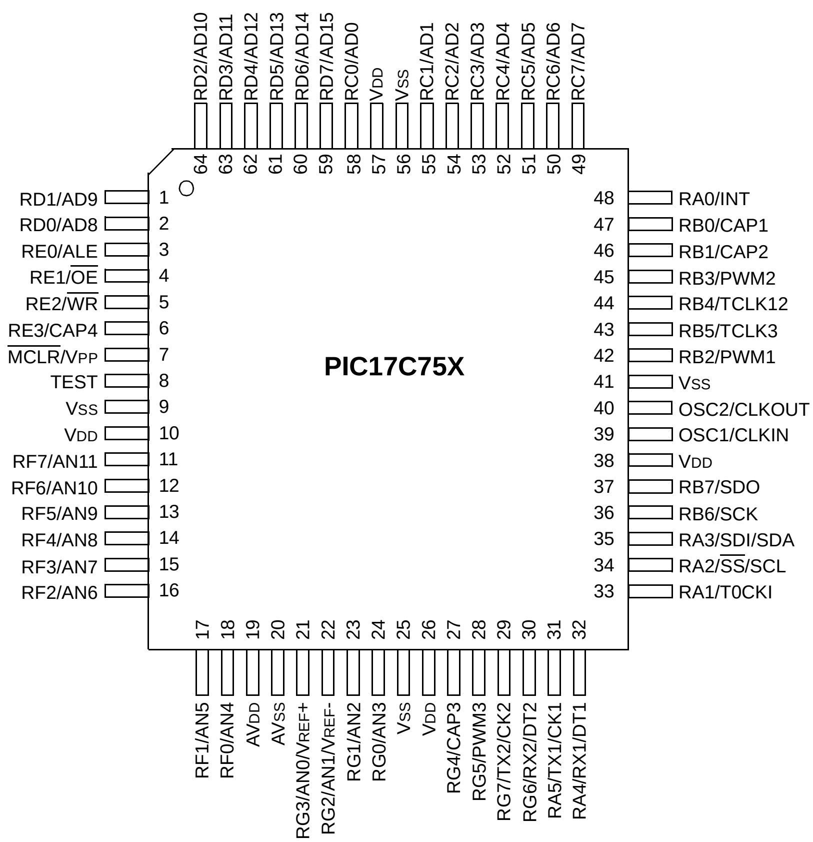
Details	
Product Status	Obsolete
Core Processor	PIC
Core Size	8-Bit
Speed	8MHz
Connectivity	I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	66
Program Memory Size	16KB (8K x 16)
Program Memory Type	OTP
EEPROM Size	-
RAM Size	678 x 8
Voltage - Supply (Vcc/Vdd)	3V ~ 5.5V
Data Converters	A/D 16x10b
Oscillator Type	External
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	84-LCC (J-Lead)
Supplier Device Package	84-PLCC (29.31x29.31)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic17lc762t-08i-l

Pin Diagrams cont.'d

68-Pin PLCC



64-Pin TQFP



2.0 DEVICE VARIETIES

Each device has a variety of frequency ranges and packaging options. Depending on application and production requirements, the proper device option can be selected using the information in the PIC17C7XX Product Selection System section at the end of this data sheet. When placing orders, please use the "PIC17C7XX Product Identification System" at the back of this data sheet to specify the correct part number. When discussing the functionality of the device, memory technology and voltage range does not matter.

There are two memory type options. These are specified in the middle characters of the part number.

1. **C**, as in PIC17**C**756A. These devices have EPROM type memory.
2. **CR**, as in PIC17**CR**756A. These devices have ROM type memory.

All these devices operate over the standard voltage range. Devices are also offered which operate over an extended voltage range (and reduced frequency range). Table 2-1 shows all possible memory types and voltage range designators for a particular device. These designators are in **bold** typeface.

TABLE 2-1: DEVICE MEMORY VARIETIES

Memory Type	Voltage Range	
	Standard	**Extended**
EPROM	PIC17**C**XXX	PIC17**LC**XXX
ROM	PIC17**CR**XXX	PIC17**LCR**XXX

Note: Not all memory technologies are available for a particular device.

2.1 UV Erasable Devices

The UV erasable version, offered in CERQUAD package, is optimal for prototype development and pilot programs.

The UV erasable version can be erased and reprogrammed to any of the configuration modes. Third party programmers also are available; refer to the *Third Party Guide* for a list of sources.

2.2 One-Time-Programmable (OTP) Devices

The availability of OTP devices is especially useful for customers expecting frequent code changes and updates.

The OTP devices, packaged in plastic packages, permit the user to program them once. In addition to the program memory, the configuration bits must be programmed.

2.3 Quick-Turnaround-Production (QTP) Devices

Microchip offers a QTP Programming Service for factory production orders. This service is made available for users who choose not to program a medium to high quantity of units and whose code patterns have stabilized. The devices are identical to the OTP devices but with all EPROM locations and configuration options already programmed by the factory. Certain code and prototype verification procedures apply before production shipments are available. Please contact your local Microchip Technology sales office for more details.

2.4 Serialized Quick-Turnaround Production (SQTP^{sm}) Devices

Microchip offers a unique programming service, where a few user defined locations in each device are programmed with different serial numbers. The serial numbers may be random, pseudo-random or sequential.

Serial programming allows each device to have a unique number which can serve as an entry code, password or ID number.

2.5 Read Only Memory (ROM) Devices

Microchip offers masked ROM versions of several of the highest volume parts, thus giving customers a low cost option for high volume, mature products.

ROM devices do not allow serialization information in the program memory space.

For information on submitting ROM code, please contact your regional sales office.

Note: Presently, NO ROM versions of the PIC17C7XX devices are available.

4.2 Clocking Scheme/Instruction Cycle

The clock input (from OSC1) is internally divided by four to generate four non-overlapping quadrature clocks, namely Q1, Q2, Q3 and Q4. Internally, the program counter (PC) is incremented every Q1 and the instruction is fetched from the program memory and latched into the instruction register in Q4. The instruction is decoded and executed during the following Q1 through Q4. The clocks and instruction execution flow are shown in Figure 4-8.

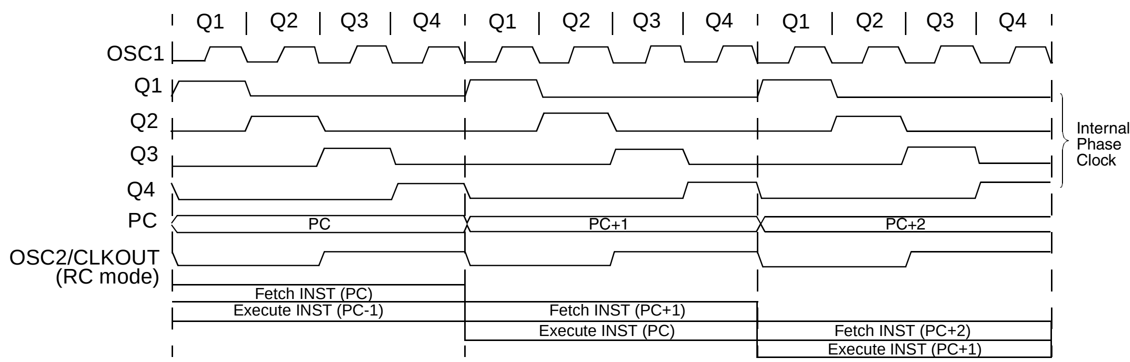
4.3 Instruction Flow/Pipelining

An “Instruction Cycle” consists of four Q cycles (Q1, Q2, Q3 and Q4). The instruction fetch and execute are pipelined such that fetch takes one instruction cycle, while decode and execute takes another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the program counter to change (e.g. `GOTO`), then two cycles are required to complete the instruction (Example 4-1).

A fetch cycle begins with the program counter incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the “Instruction Register (IR)” in cycle Q1. This instruction is then decoded and executed during the Q2, Q3 and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).

FIGURE 4-8: CLOCK/INSTRUCTION CYCLE



EXAMPLE 4-1: INSTRUCTION PIPELINE FLOW

	TCY0	TCY1	TCY2	TCY3	TCY4	TCY5
`1. MOVLW 55h`	Fetch 1	Execute 1				
`2. MOVWF PORTB`		Fetch 2	Execute 2			
`3. CALL SUB_1`			Fetch 3	Execute 3		
`4. BSF   PORTA, BIT3 (Forced NOP)`				Fetch 4	Flush	
`5. Instruction @ address SUB_1`					Fetch SUB_1	Execute SUB_1

All instructions are single cycle, except for any program branches. These take two cycles since the fetched instruction is “flushed” from the pipeline, while the new instruction is being fetched and then executed.

NOTES:

5.1 Power-on Reset (POR), Power-up Timer (PWRT), Oscillator Start-up Timer (OST) and Brown-out Reset (BOR)

5.1.1 POWER-ON RESET (POR)

The Power-on Reset circuit holds the device in RESET until VDD is above the trip point (in the range of 1.4V - 2.3V). The devices produce an internal RESET for both rising and falling VDD. To take advantage of the POR, just tie the MCLR/VPP pin directly (or through a resistor) to VDD. This will eliminate external RC components usually needed to create Power-on Reset. A minimum rise time for VDD is required. See Electrical Specifications for details.

Figure 5-2 and Figure 5-3 show two possible POR circuits.

FIGURE 5-2: USING ON-CHIP POR

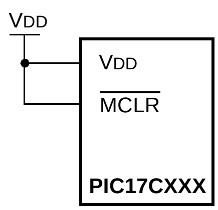
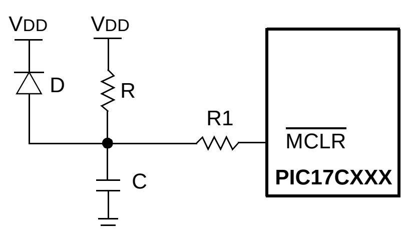


FIGURE 5-3: EXTERNAL POWER-ON RESET CIRCUIT (FOR SLOW VDD POWER-UP)



Note 1: An external Power-on Reset circuit is required only if VDD power-up time is too slow. The diode D helps discharge the capacitor quickly when VDD powers down.

2: R < 40 kΩ is recommended to ensure that the voltage drop across R does not exceed 0.2V (max. leakage current spec. on the MCLR/VPP pin is 5 μA). A larger voltage drop will degrade VIH level on the MCLR/VPP pin.

3: R1 = 100Ω to 1 kΩ will limit any current flowing into MCLR from external capacitor C in the event of MCLR/VPP pin breakdown due to Electrostatic Discharge (ESD) or Electrical Overstress (EOS).

5.1.2 POWER-UP TIMER (PWRT)

The Power-up Timer provides a fixed 96 ms time-out (nominal) on power-up. This occurs from the rising edge of the internal POR signal if VDD and MCLR are tied, or after the first rising edge of MCLR (detected high). The Power-up Timer operates on an internal RC oscillator. The chip is kept in RESET as long as the PWRT is active. In most cases, the PWRT delay allows VDD to rise to an acceptable level.

The power-up time delay will vary from chip to chip and with VDD and temperature. See DC parameters for details.

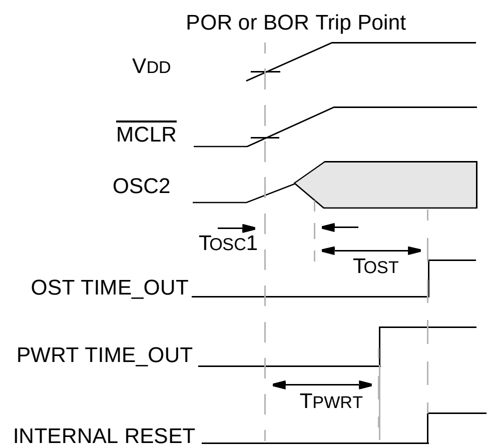
5.1.3 OSCILLATOR START-UP TIMER (OST)

The Oscillator Start-up Timer (OST) provides a 1024 oscillator cycle (1024TOSC) delay whenever the PWRT is invoked, or a wake-up from SLEEP event occurs in XT or LF mode. The PWRT and OST operate in parallel.

The OST counts the oscillator pulses on the OSC1/CLKIN pin. The counter only starts incrementing after the amplitude of the signal reaches the oscillator input thresholds. This delay allows the crystal oscillator or resonator to stabilize before the device exits RESET. The length of the time-out is a function of the crystal/resonator frequency.

Figure 5-4 shows the operation of the OST circuit. In this figure, the oscillator is of such a low frequency that although enabled simultaneously, the OST does not time-out until after the Power-up Timer time-out.

FIGURE 5-4: OSCILLATOR START-UP TIME (LOW FREQUENCY)



This figure shows in greater detail the timings involved with the oscillator start-up timer. In this example, the low frequency crystal start-up time is larger than power-up time (TPWRT).

TOSC1 = time for the crystal oscillator to react to an oscillation level detectable by the Oscillator Start-up Timer (OST).

TOST = 1024TOSC.

7.2 Data Memory Organization

Data memory is partitioned into two areas. The first is the General Purpose Registers (GPR) area, and the second is the Special Function Registers (SFR) area. The SFRs control and provide status of device operation.

Portions of data memory are banked, this occurs in both areas. The GPR area is banked to allow greater than 232 bytes of general purpose RAM.

Banking requires the use of control bits for bank selection. These control bits are located in the Bank Select Register (BSR). If an access is made to the unbanked region, the BSR bits are ignored. Figure 7-5 shows the data memory map organization.

Instructions `MOVPF` and `MOVFP` provide the means to move values from the peripheral area ("P") to any location in the register file ("F"), and vice-versa. The definition of the "P" range is from 0h to 1Fh, while the "F" range is 0h to FFh. The "P" range has six more locations than peripheral registers, which can be used as General Purpose Registers. This can be useful in some applications where variables need to be copied to other locations in the general purpose RAM (such as saving status information during an interrupt).

The entire data memory can be accessed either directly, or indirectly (through file select registers FSR0 and FSR1) (see Section 7.4). Indirect addressing uses the appropriate control bits of the BSR for access into the banked areas of data memory. The BSR is explained in greater detail in Section 7.8.

7.2.1 GENERAL PURPOSE REGISTER (GPR)

All devices have some amount of GPR area. The GPRs are 8-bits wide. When the GPR area is greater than 232, it must be banked to allow access to the additional memory space.

All the PIC17C7XX devices have banked memory in the GPR area. To facilitate switching between these banks, the `MOVLR bank` instruction has been added to the instruction set. GPRs are not initialized by a Power-on Reset and are unchanged on all other RESETS.

7.2.2 SPECIAL FUNCTION REGISTERS (SFR)

The SFRs are used by the CPU and peripheral functions to control the operation of the device (Figure 7-5). These registers are static RAM.

The SFRs can be classified into two sets, those associated with the "core" function and those related to the peripheral functions. Those registers related to the "core" are described here, while those related to a peripheral feature are described in the section for each peripheral feature.

The peripheral registers are in the banked portion of memory, while the core registers are in the unbanked region. To facilitate switching between the peripheral banks, the `MOVLB bank` instruction has been provided.

8.1 Table Writes to Internal Memory

A table write operation to internal memory causes a long write operation. The long write is necessary for programming the internal EPROM. Instruction execution is halted while in a long write cycle. The long write will be terminated by any enabled interrupt. To ensure that the EPROM location has been well programmed, a minimum programming time is required (see specification #D114). Having only one interrupt enabled to terminate the long write ensures that no unintentional interrupts will prematurely terminate the long write.

The sequence of events for programming an internal program memory location should be:

1. Disable all interrupt sources, except the source to terminate EPROM program write.
2. Raise $\overline{\text{MCLR}}$/VPP pin to the programming voltage.
3. Clear the WDT.
4. Do the table write. The interrupt will terminate the long write.
5. Verify the memory location (table read).

> **Note 1:** Programming requirements must be met. See timing specification in electrical specifications for the desired device. Violating these specifications (including temperature) may result in EPROM locations that are not fully programmed and may lose their state over time.
>
> **2:** If the VPP requirement is not met, the table write is a 2-cycle write and the program memory is unchanged.

8.1.1 TERMINATING LONG WRITES

An interrupt source or RESET are the only events that terminate a long write operation. Terminating the long write from an interrupt source requires that the interrupt enable and flag bits are set. The GLINTD bit only enables the vectoring to the interrupt address.

If the T0CKI, RA0/INT, or TMR0 interrupt source is used to terminate the long write, the interrupt flag of the highest priority enabled interrupt, will terminate the long write and automatically be cleared.

> **Note 1:** If an interrupt is pending, the `TABLWT` is aborted (a `NOP` is executed). The highest priority pending interrupt, from the T0CKI, RA0/INT, or TMR0 sources that is enabled, has its flag cleared.
>
> **2:** If the interrupt is not being used for the program write timing, the interrupt should be disabled. This will ensure that the interrupt is not lost, nor will it terminate the long write prematurely.

If a peripheral interrupt source is used to terminate the long write, the interrupt enable and flag bits must be set. The interrupt flag will not be automatically cleared upon the vectoring to the interrupt vector address.

The GLINTD bit determines whether the program will branch to the interrupt vector when the long write is terminated. If GLINTD is clear, the program will vector, if GLINTD is set, the program will not vector to the interrupt address.

TABLE 8-1: INTERRUPT - TABLE WRITE INTERACTION

Interrupt Source	GLINTD	Enable Bit	Flag Bit	Action
RA0/INT, TMR0, T0CKI	0	1	1	Terminate long table write (to internal program memory), branch to interrupt vector (branch clears flag bit).
	0	1	0	None.
	1	0	x	None.
	1	1	1	Terminate long table write, do not branch to interrupt vector (flag is automatically cleared).
Peripheral	0	1	1	Terminate long table write, branch to interrupt vector.
	0	1	0	None.
	1	0	x	None.
	1	1	1	Terminate long table write, do not branch to interrupt vector (flag remains set).

10.5 PORTE and DDRE Register

PORTE is a 4-bit bi-directional port. The corresponding data direction register is DDRE. A '1' in DDRE configures the corresponding port pin as an input. A '0' in the DDRE register configures the corresponding port pin as an output. Reading PORTE reads the status of the pins, whereas writing to PORTE will write to the port latch. PORTE is multiplexed with the system bus. When operating as the system bus, PORTE contains the control signals for the address/data bus (AD15:AD0). These control signals are Address Latch Enable (ALE), Output Enable ($\overline{OE}$) and Write ($\overline{WR}$). The control signals $\overline{OE}$ and $\overline{WR}$ are active low signals. The timing for the system bus is shown in the Electrical Specifications section.

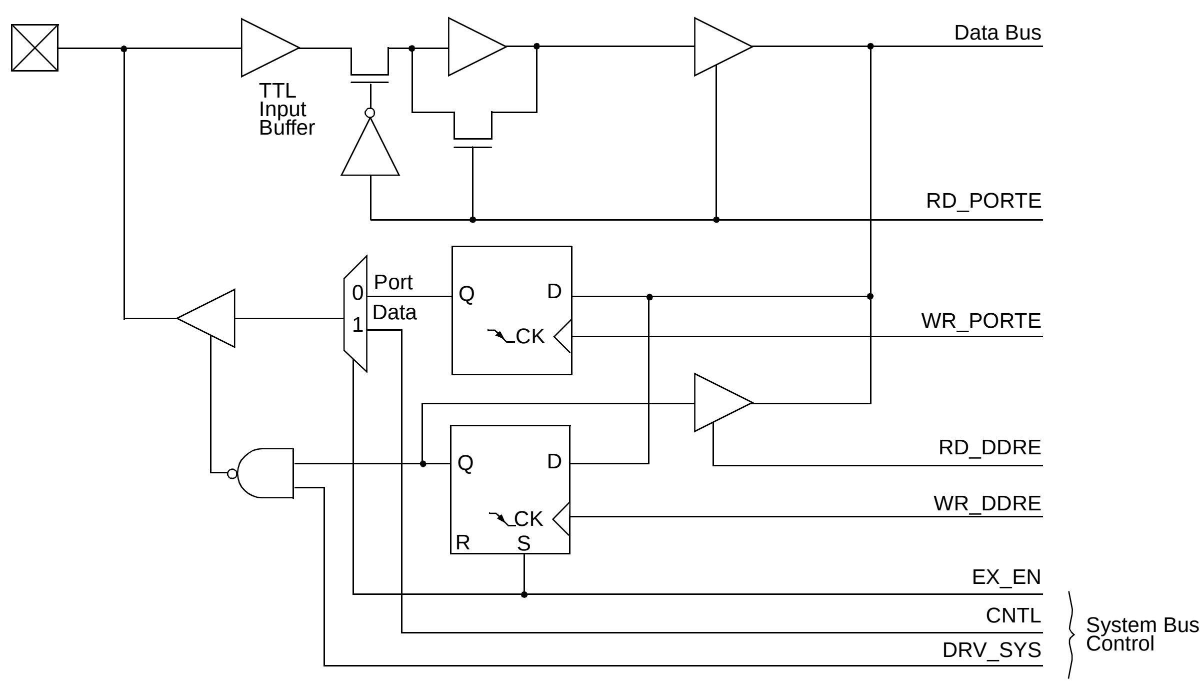
Note:	Three pins of this port are configured as the system bus when the device's configuration bits are selected to Microprocessor or Extended Microcontroller modes. The other pin is a general purpose I/O or Capture4 pin. In the two other microcontroller modes, RE2:RE0 are general purpose I/O pins.

Example 10-5 shows an instruction sequence to initialize PORTE. The Bank Select Register (BSR) must be selected to Bank 1 for the port to be initialized. The following example uses the `MOVLB` instruction to load the BSR register for bank selection.

EXAMPLE 10-5: INITIALIZING PORTE

```
MOVLB  1          ; Select Bank 1
CLRF   PORTE, F   ; Initialize PORTE data
                  ; latches before setting
                  ; the data direction
                  ; register
MOVLW   0x03      ; Value used to initialize
                  ; data direction
MOVWF   DDRE      ; Set RE<1:0> as inputs
                  ; RE<3:2> as outputs
                  ; RE<7:4> are always
                  ; read as '0'
```

FIGURE 10-11: BLOCK DIAGRAM OF RE2:RE0 (IN I/O PORT MODE)



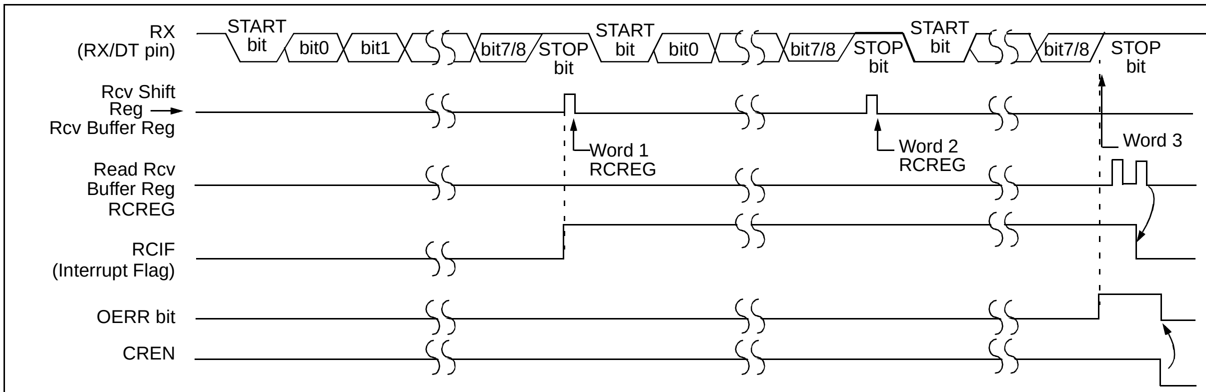
Note: I/O pins have protection diodes to VDD and VSS.

Steps to follow when setting up an Asynchronous Reception:

1. Initialize the SPBRG register for the appropriate baud rate.
2. Enable the asynchronous serial port by clearing the SYNC bit and setting the SPEN bit.
3. If interrupts are desired, then set the RCIE bit.
4. If 9-bit reception is desired, then set the RX9 bit.
5. Enable the reception by setting the CREN bit.
6. The RCIF bit will be set when reception completes and an interrupt will be generated if the RCIE bit was set.
7. Read RCSTA to get the ninth bit (if enabled) and FERR bit to determine if any error occurred during reception.
8. Read RCREG for the 8-bit received data.
9. If an overrun error occurred, clear the error by clearing the OERR bit.

Note: To terminate a reception, either clear the SREN and CREN bits, or the SPEN bit. This will reset the receive logic, so that it will be in the proper state when receive is re-enabled.

FIGURE 14-7: ASYNCHRONOUS RECEPTION



Note: This timing diagram shows three words appearing on the RX input. The RCREG (receive buffer) is read after the third word, causing the OERR (overrun) bit to be set.

TABLE 14-7: REGISTERS ASSOCIATED WITH ASYNCHRONOUS RECEPTION

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	$\overline{\text{MCLR}}$, WDT
16h, Bank 1	PIR1	RBIF	TMR3IF	TMR2IF	TMR1IF	CA2IF	CA1IF	TX1IF	RC1IF	x000 0010	u000 0010
17h, Bank 1	PIE1	RBIE	TMR3IE	TMR2IE	TMR1IE	CA2IE	CA1IE	TX1IE	RC1IE	0000 0000	0000 0000
13h, Bank 0	RCSTA1	SPEN	RX9	SREN	CREN	—	FERR	OERR	RX9D	0000 -00x	0000 -00u
14h, Bank 0	RCREG1	RX7	RX6	RX5	RX4	RX3	RX2	RX1	RX0	xxxx xxxx	uuuu uuuu
15h, Bank 0	TXSTA1	CSRC	TX9	TXEN	SYNC	—	—	TRMT	TX9D	0000 --1x	0000 --1u
17h, Bank 0	SPBRG1	Baud Rate Generator Register								0000 0000	0000 0000
10h, Bank 4	PIR2	SSPIF	BCLIF	ADIF	—	CA4IF	CA3IF	TX2IF	RC2IF	000- 0010	000- 0010
11h, Bank 4	PIE2	SSPIE	BCLIE	ADIE	—	CA4IE	CA3IE	TX2IE	RC2IE	000- 0000	000- 0000
13h, Bank 4	RCSTA2	SPEN	RX9	SREN	CREN	—	FERR	OERR	RX9D	0000 -00x	0000 -00u
14h, Bank 4	RCREG2	RX7	RX6	RX5	RX4	RX3	RX2	RX1	RX0	xxxx xxxx	uuuu uuuu
15h, Bank 4	TXSTA2	CSRC	TX9	TXEN	SYNC	—	—	TRMT	TX9D	0000 --1x	0000 --1u
17h, Bank 4	SPBRG2	Baud Rate Generator Register								0000 0000	0000 0000

Legend: x = unknown, u = unchanged, - = unimplemented, read as a '0'. Shaded cells are not used for asynchronous reception.

15.0 MASTER SYNCHRONOUS SERIAL PORT (MSSP) MODULE

The Master Synchronous Serial Port (MSSP) module is a serial interface useful for communicating with other peripheral or microcontroller devices. These peripheral devices may be serial EEPROMs, shift registers, display drivers, A/D converters, etc. The MSSP module can operate in one of two modes:

- Serial Peripheral Interface (SPI)
- Inter-Integrated Circuit™ (I^2C)

Figure 15-1 shows a block diagram for the SPI mode, while Figure 15-2 and Figure 15-3 show the block diagrams for the two different I^2C modes of operation.

FIGURE 15-1: SPI MODE BLOCK DIAGRAM

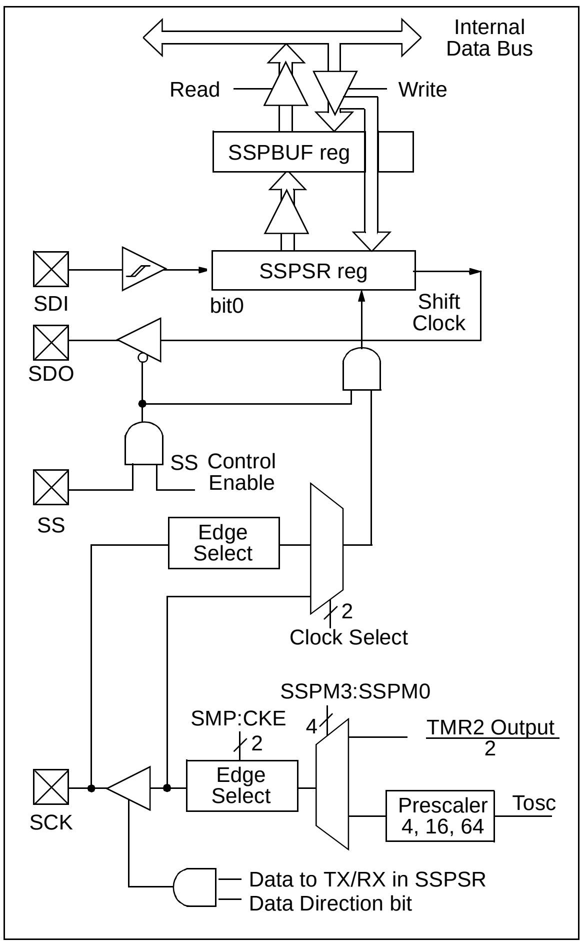


FIGURE 15-2: I^2C SLAVE MODE BLOCK DIAGRAM

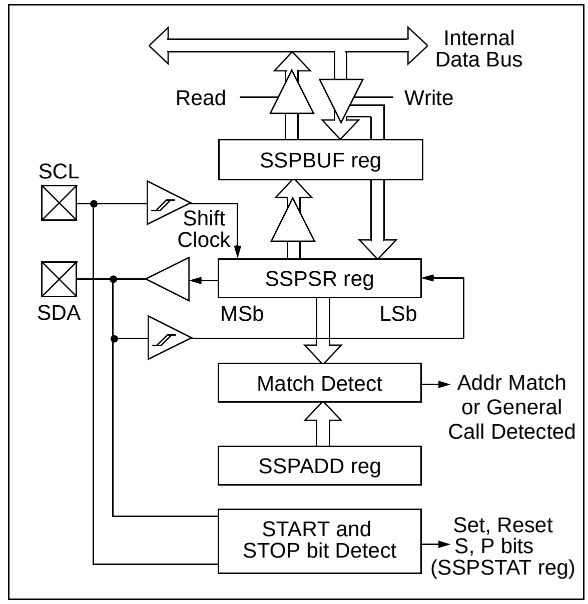
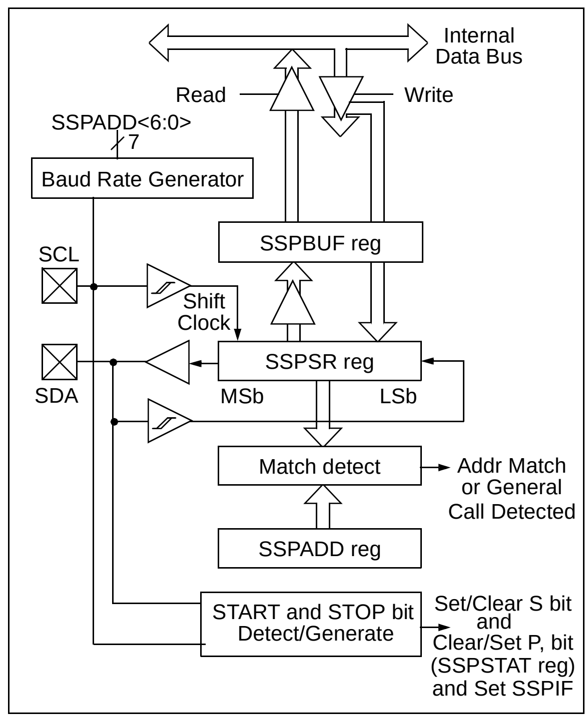


FIGURE 15-3: I^2C MASTER MODE BLOCK DIAGRAM



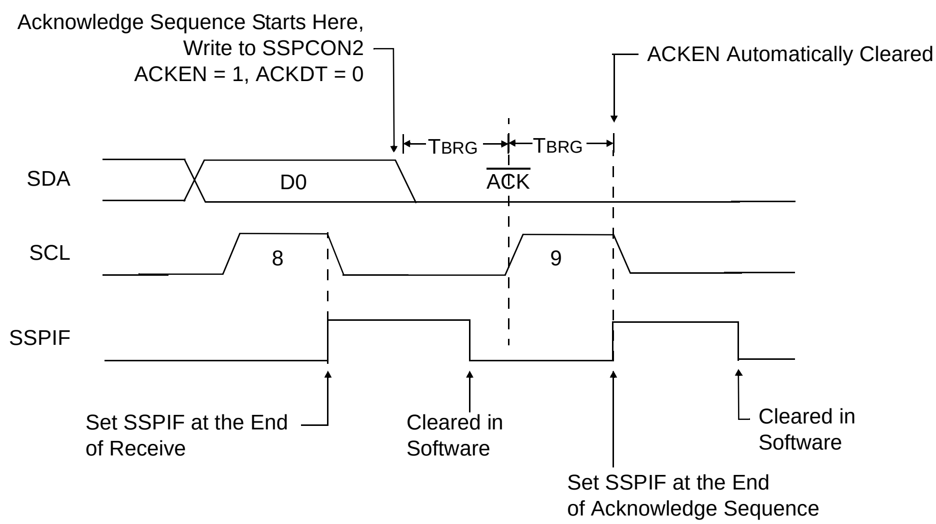
15.2.13 ACKNOWLEDGE SEQUENCE TIMING

An acknowledge sequence is enabled by setting the acknowledge sequence enable bit, ACKEN (SSPCON2<4>). When this bit is set, the SCL pin is pulled low and the contents of the acknowledge data bit is presented on the SDA pin. If the user wishes to generate an acknowledge, then the ACKDT bit should be cleared. If not, the user should set the ACKDT bit before starting an acknowledge sequence. The baud rate generator then counts for one rollover period (T_{BRG}), and the SCL pin is de-asserted (pulled high). When the SCL pin is sampled high (clock arbitration), the baud rate generator counts for T_{BRG}. The SCL pin is then pulled low. Following this, the ACKEN bit is automatically cleared, the baud rate generator is turned off and the SSP module then goes into IDLE mode (Figure 15-29).

15.2.13.1 WCOL Status Flag

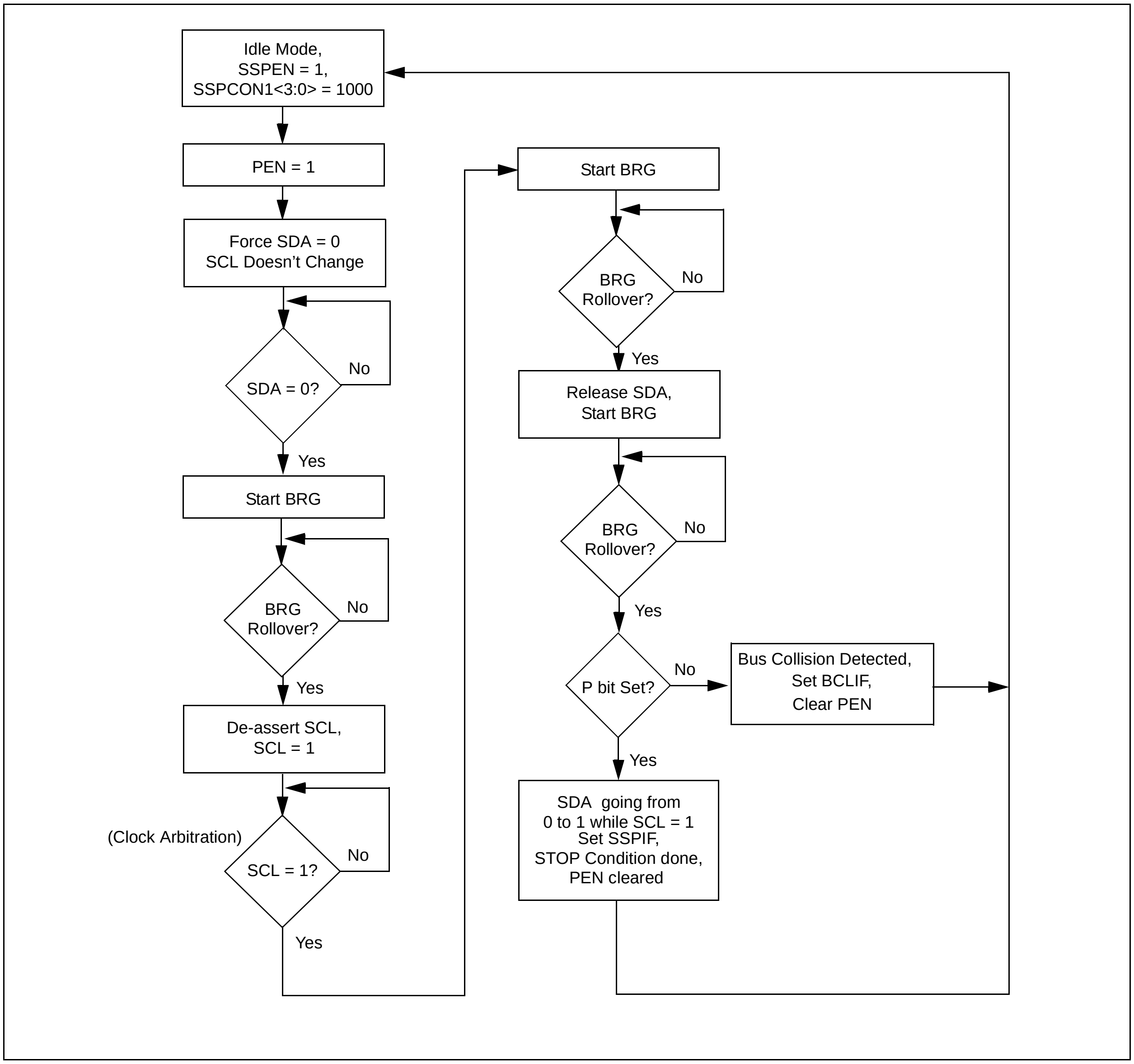
If the user writes the SSPBUF when an acknowledge sequence is in progress, then WCOL is set and the contents of the buffer are unchanged (the write doesn't occur).

FIGURE 15-29: ACKNOWLEDGE SEQUENCE WAVEFORM



Note: T_{BRG} = one baud rate generator period.

FIGURE 15-32: STOP CONDITION FLOW CHART



15.2.18 MULTI -MASTER COMMUNICATION, BUS COLLISION AND BUS ARBITRATION

Multi-Master mode support is achieved by bus arbitration. When the master outputs address/data bits onto the SDA pin, arbitration takes place when the master outputs a '1' on SDA, by letting SDA float high and another master asserts a '0'. When the SCL pin floats high, data should be stable. If the expected data on SDA is a '1' and the data sampled on the SDA pin = '0', then a bus collision has taken place. The master will set the Bus Collision Interrupt Flag, BCLIF and reset the I^{2}C port to its IDLE state (Figure 15-34).

If a transmit was in progress when the bus collision occurred, the transmission is halted, the BF flag is cleared, the SDA and SCL lines are de-asserted and the SSPBUF can be written to. When the user services the bus collision Interrupt Service Routine and if the I^{2}C bus is free, the user can resume communication by asserting a START condition.

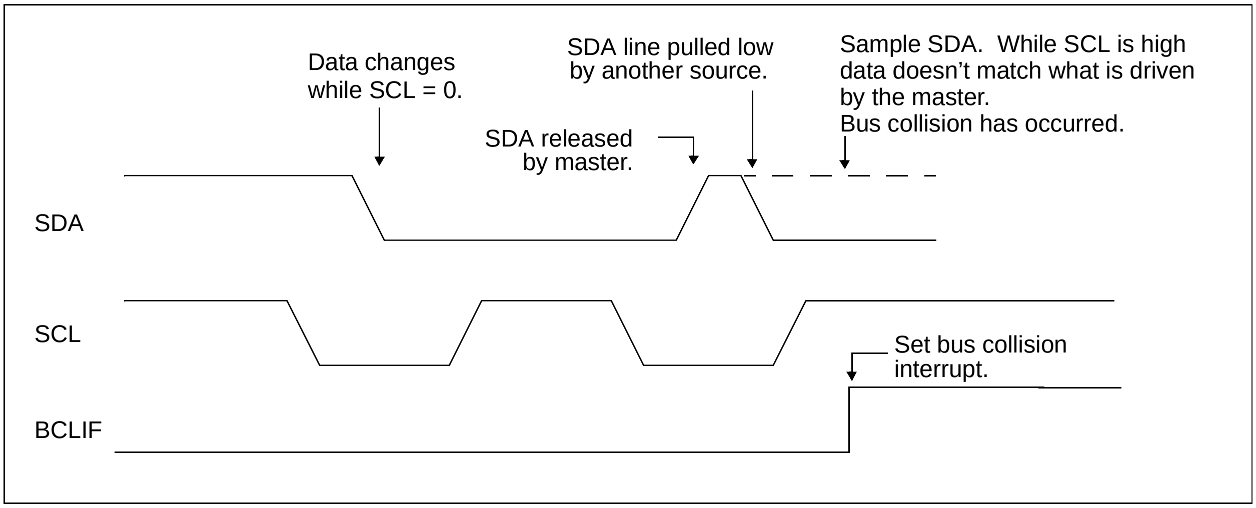
If a START, Repeated Start, STOP, or Acknowledge condition was in progress when the bus collision occurred, the condition is aborted, the SDA and SCL lines are de-asserted and the respective control bits in the SSPCON2 register are cleared. When the user services the bus collision Interrupt Service Routine, and if the I^{2}C bus is free, the user can resume communication by asserting a START condition.

The master will continue to monitor the SDA and SCL pins and if a STOP condition occurs, the SSPIF bit will be set.

A write to the SSPBUF will start the transmission of data at the first data bit, regardless of where the transmitter left off when bus collision occurred.

In Multi-Master mode, the interrupt generation on the detection of START and STOP conditions allows the determination of when the bus is free. Control of the I^{2}C bus can be taken when the P bit is set in the SSPSTAT register, or the bus is idle and the S and P bits are cleared.

FIGURE 15-34: BUS COLLISION TIMING FOR TRANSMIT AND ACKNOWLEDGE



EXAMPLE 15-2: INTERFACING TO A 24LC01B SERIAL EEPROM (USING MPLAB C17)

```
void ACKPoll(void)
{
        StartI2C();                     // Send start bit
        IdleI2C();                      // Wait for idle condition
        WriteI2C(CONTROL);              // Send control byte
        IdleI2C();                      // Wait for idle condition
        // Poll the ACK bit coming from the 24LC01B
        // Loop as long as the 24LC01B NACKs
        while (SSPCON2bits.ACKSTAT)
        {
                RestartI2C();           // Send a restart bit
                IdleI2C();              // Wait for idle condition
                WriteI2C(CONTROL);      // Send control byte
                IdleI2C();              // Wait for idle condition
        }
        IdleI2C();                      // Wait for idle condition
        StopI2C();                      // Send stop bit
        IdleI2C();                      // Wait for idle condition
        return;
}
```

DECF **Decrement f**

Syntax: [*label*] DECF f,d

Operands: $0 \leq f \leq 255$
$d \in [0,1]$

Operation: (f) – 1 → (dest)

Status Affected: OV, C, DC, Z

Encoding:

0000	011d	ffff	ffff

Description: Decrement register 'f'. If 'd' is 0, the result is stored in WREG. If 'd' is 1, the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `DECF CNT, 1`

Before Instruction
CNT = 0x01
Z = 0

After Instruction
CNT = 0x00
Z = 1

DECFSZ **Decrement f, skip if 0**

Syntax: [*label*] DECFSZ f,d

Operands: $0 \leq f \leq 255$
$d \in [0,1]$

Operation: (f) – 1 → (dest);
skip if result = 0

Status Affected: None

Encoding:

0001	011d	ffff	ffff

Description: The contents of register 'f' are decremented. If 'd' is 0, the result is placed in WREG. If 'd' is 1, the result is placed back in register 'f'.
If the result is 0, the next instruction, which is already fetched is discarded and a NOP is executed instead, making it a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

Example:

```
HERE      DECFSZ   CNT, 1
          GOTO     HERE
NZERO
ZERO
```

Before Instruction
PC = Address (HERE)

After Instruction
CNT = CNT - 1
If CNT = 0;
PC = Address (HERE)
If CNT ≠ 0;
PC = Address (NZERO)

RETFIE — Return from Interrupt

Syntax:	[*label*] RETFIE
Operands:	None
Operation:	TOS → (PC); 0 → GLINTD; PCLATH is unchanged.
Status Affected:	GLINTD
Encoding:	`0000 0000 0000 0101`
Description:	Return from Interrupt. Stack is POP'ed and Top-of-Stack (TOS) is loaded in the PC. Interrupts are enabled by clearing the GLINTD bit. GLINTD is the global interrupt disable bit (CPUSTA<4>).
Words:	1
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	Clear GLINTD	POP PC from stack
No operation	No operation	No operation	No operation

<u>Example</u>: `RETFIE`

After Interrupt
PC = TOS
GLINTD = 0

RETLW — Return Literal to WREG

Syntax:	[*label*] RETLW k
Operands:	$0 \le k \le 255$
Operation:	k → (WREG); TOS → (PC); PCLATH is unchanged
Status Affected:	None
Encoding:	`1011 0110 kkkk kkkk`
Description:	WREG is loaded with the eight-bit literal 'k'. The program counter is loaded from the top of the stack (the return address). The high address latch (PCLATH) remains unchanged.
Words:	1
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	POP PC from stack, Write to WREG
No operation	No operation	No operation	No operation

<u>Example</u>:

```
    CALL TABLE ; WREG contains table
               ;  offset value
               ;  WREG now has
               ;  table value
    :
TABLE
    ADDWF PC   ; WREG = offset
    RETLW k0   ; Begin table
    RETLW k1   ;
    :
    :
    RETLW kn   ; End of table
```

Before Instruction
WREG = 0x07

After Instruction
WREG = value of k7

RRNCF — Rotate Right f (no carry)

Syntax: [*label*] RRNCF f,d

Operands: $0 \le f \le 255$
$d \in [0,1]$

Operation: f<n> → d<n-1>;
f<0> → d<7>

Status Affected: None

Encoding:

0010	000d	ffff	ffff

Description: The contents of register 'f' are rotated one bit to the right. If 'd' is 0, the result is placed in WREG. If 'd' is 1, the result is placed back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example 1: `RRNCF REG, 1`

Before Instruction
WREG = ?
REG = 1101 0111

After Instruction
WREG = 0
REG = 1110 1011

Example 2: `RRNCF REG, 0`

Before Instruction
WREG = ?
REG = 1101 0111

After Instruction
WREG = 1110 1011
REG = 1101 0111

SETF — Set f

Syntax: [*label*] SETF f,s

Operands: $0 \le f \le 255$
$s \in [0,1]$

Operation: FFh → f;
FFh → d

Status Affected: None

Encoding:

0010	101s	ffff	ffff

Description: If 's' is 0, both the data memory location 'f' and WREG are set to FFh. If 's' is 1, only the data memory location 'f' is set to FFh.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f' and other specified register

Example1: `SETF REG, 0`

Before Instruction
REG = 0xDA
WREG = 0x05

After Instruction
REG = 0xFF
WREG = 0xFF

Example2: `SETF REG, 1`

Before Instruction
REG = 0xDA
WREG = 0x05

After Instruction
REG = 0xFF
WREG = 0x05

FIGURE 21-2: TYPICAL RC OSCILLATOR FREQUENCY vs. VDD

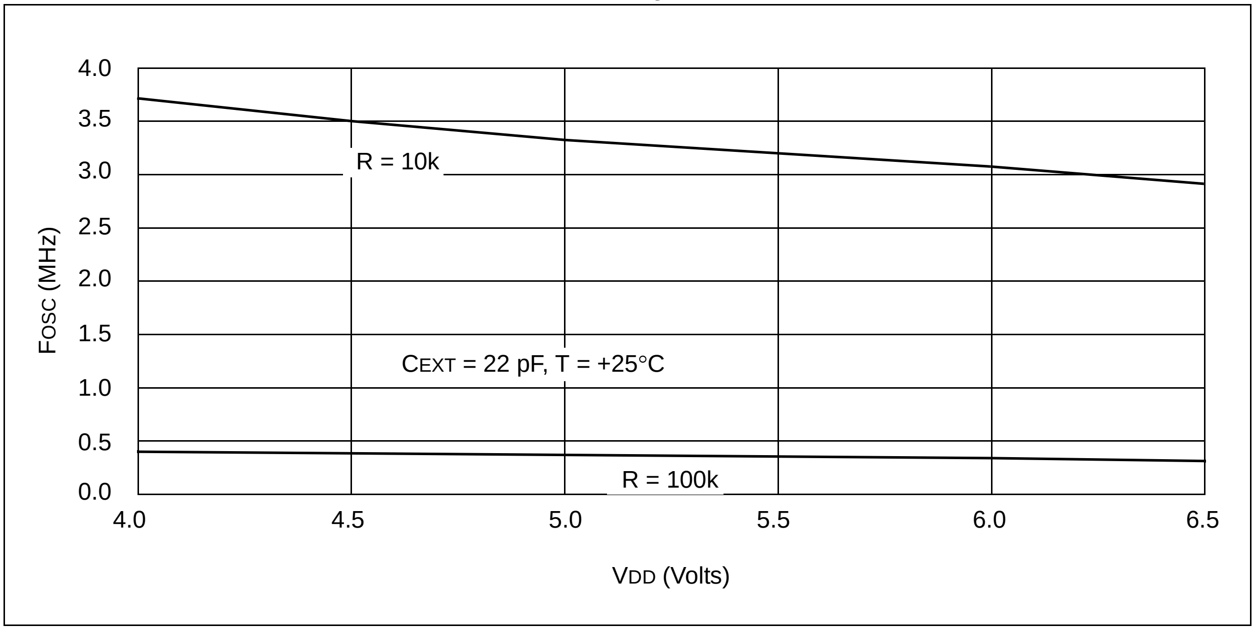


FIGURE 21-3: TYPICAL RC OSCILLATOR FREQUENCY vs. VDD

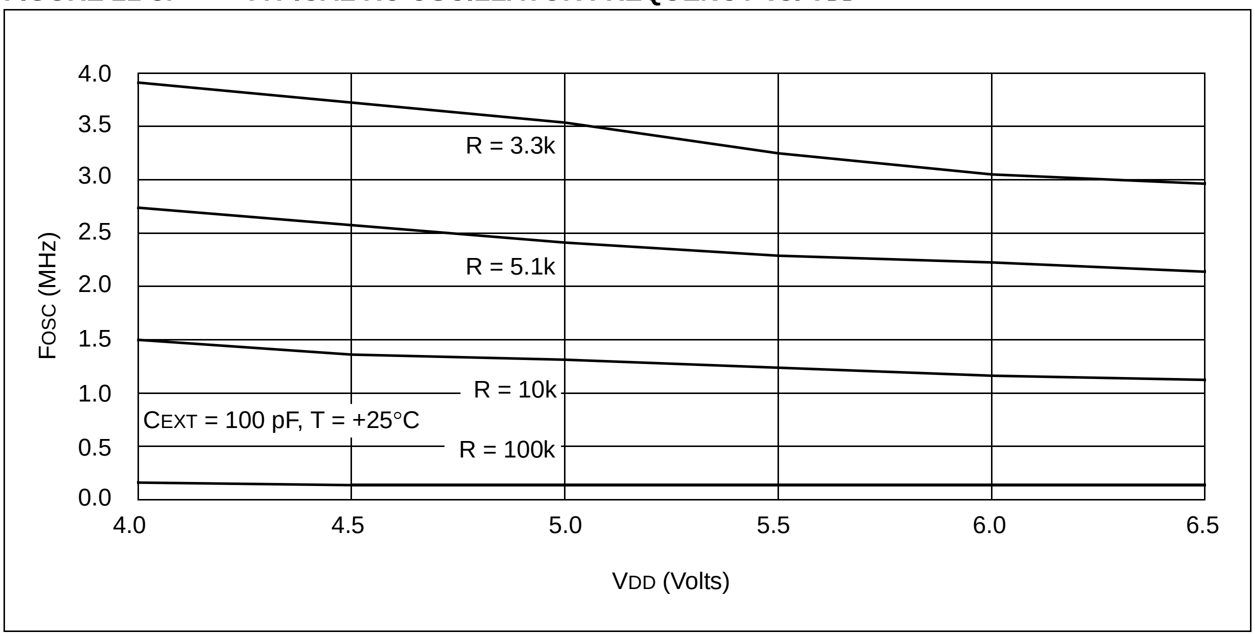


FIGURE 21-13: TYPICAL AND MAXIMUM ΔIPD vs. VDD (SLEEP MODE, WDT ENABLED, -40°C to +125°C)

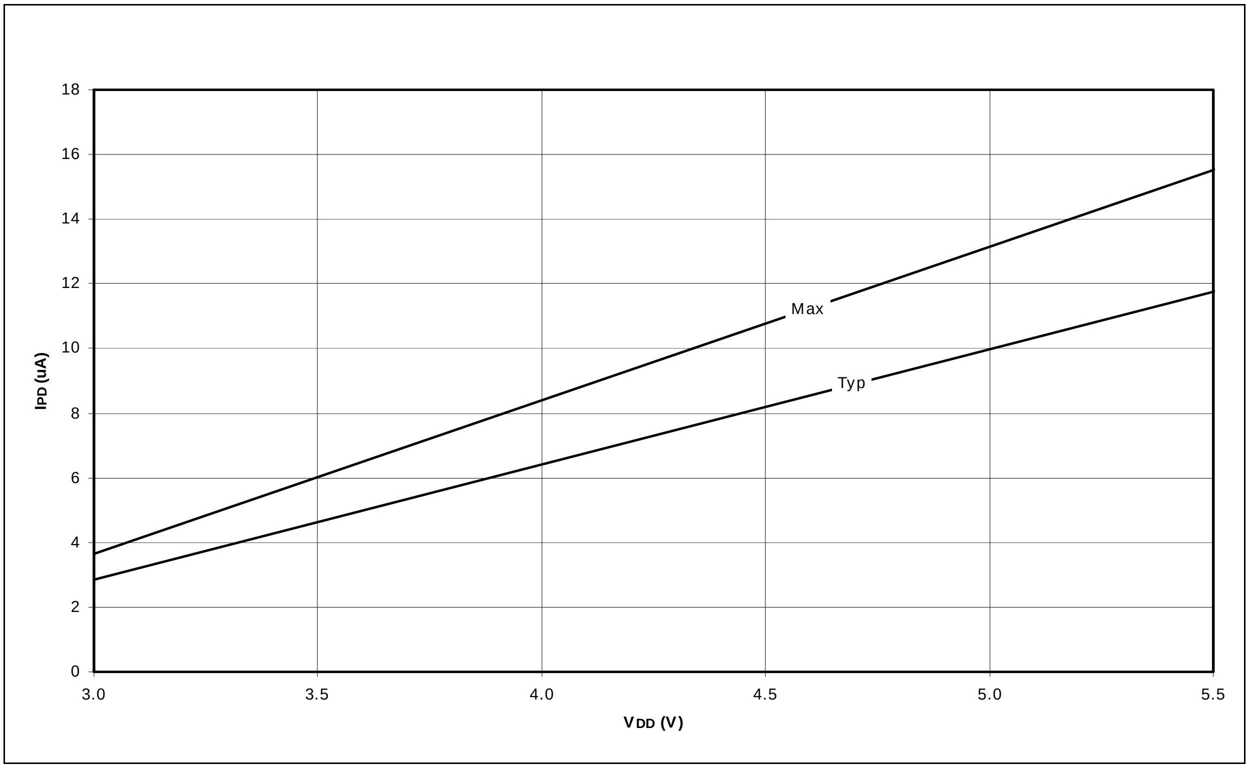


FIGURE 21-14: TYPICAL AND MAXIMUM ΔIRBPU vs. VDD (MEASURED PER INPUT PIN, -40°C TO +125°C)