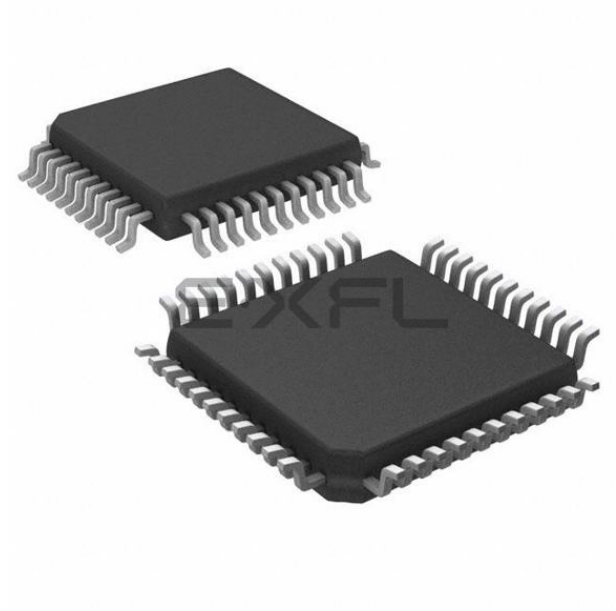

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Applications of "Embedded - Microcontrollers"

Details	
Product Status	Obsolete
Core Processor	HC08
Core Size	8-Bit
Speed	8MHz
Connectivity	SCI, SPI
Peripherals	LVD, POR, PWM
Number of I/O	36
Program Memory Size	16KB (16K x 8)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	512 x 8
Voltage - Supply (Vcc/Vdd)	2.7V ~ 5.5V
Data Converters	A/D 8x8b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	44-QFP
Supplier Device Package	44-QFP (10x10)
Purchase URL	https://www.e-xfl.com/product-detail/nxp-semiconductors/mc908gt16cfber

Chapter 2
Memory

2.1 Introduction

The CPU08 can address 64 Kbytes of memory space. The memory map, shown in Figure 2-1, includes:

- User Flash memory:
  - MC68HC908GT16 — 15,872 bytes
  - MC68HC908GT8 — 7,680 bytes
- 512 bytes of random-access memory (RAM)
- 720 bytes of Flash programming routines read-only memory (ROM)
- 36 bytes of user-defined vectors
- 304 bytes of monitor ROM

2.2 Unimplemented Memory Locations

Accessing an unimplemented location can cause an illegal address reset. In the memory map (Figure 2-1) and in register figures in this document, unimplemented locations are shaded.

2.3 Reserved Memory Locations

Accessing a reserved location can have unpredictable effects on MCU operation. In the Figure 2-1 and in register figures in this document, reserved locations are marked with the word Reserved or with the letter R.

2.4 Input/Output (I/O) Section

Most of the control, status, and data registers are in the zero page area of $0000–$003F. Additional I/O registers have these addresses:

- $FE00; SIM break status register, SBSR
- $FE01; SIM reset status register, SRSR
- $FE02; reserved, SUBAR
- $FE03; SIM break flag control register, SBFCR
- $FE04; interrupt status register 1, INT1
- $FE05; interrupt status register 2, INT2
- $FE06; interrupt status register 3, INT3
- $FE07; reserved
- $FE08; Flash control register, FLCR
- $FE09; break address register high, BRKH
- $FE0A; break address register low, BRKL
- $FE0B; break status and control register, BRKSCR

Addr.	Register Name		Bit 7	6	5	4	3	2	1	Bit 0
$001A	Keyboard Status and Control Register (INTKBSCR) See page 111.	Read:	0	0	0	0	KEYF	0	IMASKK	MODEK
		Write:						ACKK		
		Reset:	0	0	0	0	0	0	0	0
$001B	Keyboard Interrupt Enable Register (INTKBIER) See page 112.	Read:	KBIE7	KBIE6	KBIE5	KBIE4	KBIE3	KBIE2	KBIE1	KBIE0
		Write:								
		Reset:	0	0	0	0	0	0	0	0
$001C	Timebase Module Control Register (TBCR) See page 218.	Read:	TBIF	TBR2	TBR1	TBR0	0	TBIE	TBON	R
		Write:					TACK			
		Reset:	0	0	0	0	0	0	0	0
$001D	IRQ Status and Control Register (INTSCR) See page 104.	Read:	0	0	0	0	IRQF1	0	IMASK1	MODE1
		Write:						ACK1		
		Reset:	0	0	0	0	0	0	0	0
$001E	Configuration Register 2 (CONFIG2)† See page 58.	Read:	R	0	EXT-XTALEN	EXT-SLOW	EXT-CLKEN	0	OSCENIN-STOP	R
		Write:								
		Reset:	0	0	0	0	0	0	0	0
$001F	Configuration Register 1 (CONFIG1)† See page 58.	Read:	COPRS	LVISTOP	LVIRSTD	LVIPWRD	LVI5OR3(1)	SSREC	STOP	COPD
		Write:								
		Reset:	0	0	0	0	0	0	0	0

1. One-time writable register after each reset, except LVI5OR3 bit. LVI5OR3 bit is only reset via POR (power-on reset).

$0020	Timer 1 Status and Control Register (T1SC) See page 231.	Read:	TOF	TOIE	TSTOP	0	0	PS2	PS1	PS0
		Write:	0			TRST				
		Reset:	0	0	1	0	0	0	0	0
$0021	Timer 1 Counter Register High (T1CNTH) See page 232.	Read:	Bit 15	14	13	12	11	10	9	Bit 8
		Write:								
		Reset:	0	0	0	0	0	0	0	0
$0022	Timer 1 Counter Register Low (T1CNTL) See page 232.	Read:	Bit 7	6	5	4	3	2	1	Bit 0
		Write:								
		Reset:	0	0	0	0	0	0	0	0
$0023	Timer 1 Counter Modulo Register High (T1MODH) See page 233.	Read:	Bit 15	14	13	12	11	10	9	Bit 8
		Write:								
		Reset:	1	1	1	1	1	1	1	1
$0024	Timer 1 Counter Modulo Register Low (T1MODL) See page 233.	Read:	Bit 7	6	5	4	3	2	1	Bit 0
		Write:								
		Reset:	1	1	1	1	1	1	1	1
$0025	Timer 1 Channel 0 Status and Control Register (T1SC0) See page 233.	Read:	CH0F	CH0IE	MS0B	MS0A	ELS0B	ELS0A	TOV0	CH0MAX
		Write:	0							
		Reset:	0	0	0	0	0	0	0	0

= Unimplemented R = Reserved U = Unaffected

Figure 2-2. Control, Status, and Data Registers (Sheet 3 of 7)

Addr.	Register Name		Bit 7	6	5	4	3	2	1	Bit 0
$0033	Timer 2 Channel 1 Status and Control Register (T2SC1) See page 234.	Read:	CH1F	CH1IE	0	MS1A	ELS1B	ELS1A	TOV1	CH1MAX
		Write:	0							
		Reset:	0	0	0	0	0	0	0	0
$0034	Timer 2 Channel 1 Register High (T2CH1H) See page 236.	Read:	Bit 15	14	13	12	11	10	9	Bit 8
		Write:								
		Reset:	Indeterminate after reset							
$0035	Timer 2 Channel 1 Register Low (T2CH1L) See page 236.	Read:	Bit 7	6	5	4	3	2	1	Bit 0
		Write:								
		Reset:	Indeterminate after reset							
$0036	ICG Control Register (ICGCR) See page 98.	Read:	CMIE	CMF	CMON	CS	ICGON	ICGS	ECGON	ECGS
		Write:		0						
		Reset:	0	0	0	0	1	0	0	0
$0037	ICG Multiplier Register (ICGMR) See page 99.	Read:		N6	$\overline{N5}$	N4	N3	N2	N1	N0
		Write:								
		Reset:	0	0	0	1	0	1	0	1
$0038	ICG Trim Register (ICGTR) See page 100.	Read:	TRIM7	TRIM6	TRIM5	TRIM4	TRIM3	TRIM2	TRIM1	TRIM0
		Write:								
		Reset:	1	0	0	0	0	0	0	0
$0039	ICG Divider Control Register (ICGDVR) See page 100.	Read:					DDIV3	DDIV2	DDIV1	DDIV0
		Write:								
		Reset:	0	0	0	0	U	U	U	U
$003A	ICG DCO Stage Control Register (ICGDSR) See page 100.	Read:	DSTG7	DSTG6	DSTG5	DSTG4	DSTG3	DSTG2	DSTG1	DSTG0
		Write:	R	R	R	R	R	R	R	R
		Reset:	Unaffected by reset							
$003B	Reserved	Read:	R	R	R	R	R	R	R	R
		Write:								
		Reset:	Indeterminate after reset							
$003C	ADC Status and Control Register (ADSCR) See page 54.	Read:	COCO	AIEN	ADCO	ADCH4	ADCH3	ADCH2	ADCH1	ADCH0
		Write:	R							
		Reset:	0	0	0	1	1	1	1	1
$003D	ADC Data Register (ADR) See page 55.	Read:	AD7	AD6	AD5	AD4	AD3	AD2	AD1	AD0
		Write:								
		Reset:	0	0	0	0	0	0	0	0
$003E	ADC Clock Register (ADCLK) See page 56.	Read:	ADIV2	ADIV1	ADIV0	ADICLK	0	0	0	0
		Write:								
		Reset:	0	0	0	0	0	0	0	0
$003F	Unimplemented	Read:								
		Write:								
		Reset:								

= Unimplemented R = Reserved U = Unaffected

Figure 2-2. Control, Status, and Data Registers (Sheet 5 of 7)

7.3.1 Clock Enable Circuit

The clock enable circuit is used to enable the internal clock (ICLK) or external clock (ECLK) and the port logic which is shared with the oscillator pins (OSC1 and OSC2). The clock enable circuit generates an ICG stop (ICGSTOP) signal which stops all clocks (ICLK, ECLK, and the low-frequency base clock, IBASE). ICGSTOP is set and the ICG is disabled in stop mode if the oscillator enable stop bit (OSCENINSTOP) in the CONFIG2 register is clear. The ICG clocks will be enabled in stop mode if OSCENINSTOP is high.

The internal clock enable signal (ICGEN) turns on the internal clock generator which generates ICLK. ICGEN is set (active) whenever the ICGON bit is set and the ICGSTOP signal is clear. When ICGEN is clear, ICLK and IBASE are both low.

The external clock enable signal (ECGEN) turns on the external clock generator which generates ECLK. ECGEN is set (active) whenever the ECGON bit is set and the ICGSTOP signal is clear. ECGON cannot be set unless the external clock enable (EXTCLKEN) bit in the CONFIG2 register is set. when ECGEN is clear, ECLK is low.

The port E4 enable signal (PE4EN) turns on the port E4 logic. Since port E4 is on the same pin as OSC1, this signal is only active (set) when the external clock function is not desired. Therefore, PE4EN is clear when ECGON is set. PE4EN is not gated with ICGSTOP, which means that if the ECGON bit is set, the port E4 logic will remain disabled in stop mode.

The port E3 enable signal (PE3EN) turns on the port E3 logic. Since port E3 is on the same pin as OSC2, this signal is only active (set) when 2-pin oscillator function is not desired. Therefore, PE3EN is clear when ECGON and the external crystal enable (EXTXTALEN) bit in the CONFIG2 register are both set. PE3EN is not gated with ICGSTOP, which means that if ECGON and EXTXTALEN are set, the port E3 logic will remain disabled in stop mode.

7.3.2 Internal Clock Generator

The internal clock generator, shown in Figure 7-3, creates a low frequency base clock (IBASE), which operates at a nominal frequency (f_{NOM}) of 307.2 kHz $\pm$ 25 percent, and an internal clock (ICLK) which is an integer multiple of IBASE. This multiple is the ICG multiplier factor (N), which is programmed in the ICG multiplier register (ICGMR). The internal clock generator is turned off and the output clocks (IBASE and ICLK) are held low when the internal clock generator enable signal (ICGEN) is clear.

The internal clock generator contains:

- A digitally controlled oscillator
- A modulo N divider
- A frequency comparator, which contains voltage and current references, a frequency to voltage converter, and comparators
- A digital loop filter

Settling time depends primarily on how many corrections it takes to change the clock period and the period of each correction. Since the corrections require four periods of the low-frequency base clock ($4*\tau_{IBASE}$), and since ICLK is N (the ICG multiply factor for the desired frequency) times faster than IBASE, each correction takes $4*N*\tau_{ICLK}$. The period of ICLK, however, will vary as the corrections occur.

7.4.6.1 Settling to Within 15 Percent

When the error is greater than 15 percent, the filter takes eight corrections to double or halve the clock period. Due to how the DCO increases or decreases the clock period, the total period of these eight corrections is approximately 11 times the period of the fastest correction. (If the corrections were perfectly linear, the total period would be 11.5 times the minimum period; however, the ring must be slightly nonlinear.) Therefore, the total time it takes to double or halve the clock period is $44*N*\tau_{ICLKFAST}$.

If the clock period needs more than doubled or halved, the same relationship applies, only for each time the clock period needs doubled, the total number of cycles doubles. That is, when transitioning from fast to slow, going from the initial speed to half speed takes $44*N*\tau_{ICLKFAST}$; from half speed to quarter speed takes $88*N*\tau_{ICLKFAST}$; going from quarter speed to eighth speed takes $176*N*\tau_{ICLKFAST}$; and so on. This series can be expressed as $(2^X–1)*44*N*\tau_{ICLKFAST}$, where x is the number of times the speed needs doubled or halved. Since 2^X happens to be equal to $\tau_{ICLKSLOW}/\tau_{ICLKFAST}$, the equation reduces to $44*N*(\tau_{ICLKSLOW}–\tau_{ICLKFAST})$.

Note that increasing speed takes much longer than decreasing speed since N is higher. This can be expressed in terms of the initial clock period (τ_1) minus the final clock period (τ_2) as such:

$$\tau_{15} = \text{abs}[44N(\tau_1 - \tau_2)]$$

7.4.6.2 Settling to Within 5 Percent

Once the clock period is within 15 percent of the desired clock period, the filter starts making smaller adjustments. When between 15 percent and 5 percent error, each correction will adjust the clock period between 1.61 percent and 2.94 percent. In this mode, a maximum of eight corrections will be required to get to less than 5 percent error. Since the clock period is relatively close to desired, each correction takes approximately the same period of time, or $4*\tau_{IBASE}$. At this point, the internal clock stable bit (ICGS) will be set and the clock frequency is usable, although the error will be as high as 5 percent. The total time to this point is:

$$\tau_5 = \text{abs}[44N(\tau_1 - \tau_2)] + 32\tau_{IBASE}$$

7.4.6.3 Total Settling Time

Once the clock period is within 5 percent of the desired clock period, the filter starts making minimum adjustments. In this mode, each correction will adjust the frequency between 0.202 percent and 0.368 percent. A maximum of 24 corrections will be required to get to the minimum error. Each correction takes approximately the same period of time, or $4*\tau_{IBASE}$. Added to the corrections for 15 percent to 5 percent, this makes 32 corrections ($128*\tau_{IBASE}$) to get from 15 percent to the minimum error. The total time to the minimum error is:

$$\tau_{tot} = \text{abs}[44N(\tau_1 - \tau_2)] + 128\tau_{IBASE}$$

The equations for τ_{15}, τ_5, and τ_{tot} are dependent on the actual initial and final clock periods τ_1 and τ_2, not the nominal. This means the variability in the ICLK frequency due to process, temperature, and voltage must be considered. Additionally, other process factors and noise can affect the actual tolerances of the points at which the filter changes modes. This means a worst case adjustment of up to 35 percent (ICLK

9.5.2 Stop Mode

The keyboard module remains active in stop mode. Clearing the IMASKK bit in the keyboard status and control register enables keyboard interrupt requests to bring the MCU out of stop mode.

9.6 Keyboard Module During Break Interrupts

The system integration module (SIM) controls whether the keyboard interrupt latch can be cleared during the break state. The BCFE bit in the SIM break flag control register (SBFCR) enables software to clear status bits during the break state.

To allow software to clear the keyboard interrupt latch during a break interrupt, write a 1 to the BCFE bit. If a latch is cleared during the break state, it remains cleared when the MCU exits the break state.

To protect the latch during the break state, write a 0 to the BCFE bit. With BCFE at 0 (its default state), writing to the keyboard acknowledge bit (ACKK) in the keyboard status and control register during the break state has no effect. See 9.7.1 Keyboard Status and Control Register.

9.7 I/O Registers

These registers control and monitor operation of the keyboard module:

- Keyboard status and control register (INTKBSCR)
- Keyboard interrupt enable register (INTKBIER)

9.7.1 Keyboard Status and Control Register

The keyboard status and control register:

- Flags keyboard interrupt requests
- Acknowledges keyboard interrupt requests
- Masks keyboard interrupt requests
- Controls keyboard interrupt triggering sensitivity

Address: $001A

	Bit 7	6	5	4	3	2	1	Bit 0
Read:	0	0	0	0	KEYF	0	IMASKK	MODEK
Write:						ACKK		
Reset:	0	0	0	0	0	0	0	0

= Unimplemented

Figure 9-4. Keyboard Status and Control Register (INTKBSCR)

Bits 7–4 — Not used

These read-only bits always read as 0s.

KEYF — Keyboard Flag Bit

This read-only bit is set when a keyboard interrupt is pending. Reset clears the KEYF bit.

1 = Keyboard interrupt pending
0 = No keyboard interrupt pending

Chapter 11
Low-Power Modes (MODES)

11.1 Introduction

The microcontroller (MCU) may enter two low-power modes: wait mode and stop mode. They are common to all HC08 MCUs and are entered through instruction execution. This section describes how each module acts in the low-power modes.

11.1.1 Wait Mode

The WAIT instruction puts the MCU in a low-power standby mode in which the central processor unit (CPU) clock is disabled but the bus clock continues to run. Power consumption can be further reduced by disabling the LVI module and/or the timebase module through bits in the CONFIG1 register. (See Chapter 4 Configuration Register (CONFIG).)

11.1.2 Stop Mode

Stop mode is entered when a STOP instruction is executed. The CPU clock is disabled and the bus clock is disabled if the OSCENINSTOP bit in the CONFIG2 register is at a 0. (See Chapter 4 Configuration Register (CONFIG).)

11.2 Analog-to-Digital Converter (ADC)

11.2.1 Wait Mode

The analog-to-digital converter (ADC) continues normal operation during wait mode. Any enabled CPU interrupt request from the ADC can bring the MCU out of wait mode. If the ADC is not required to bring the MCU out of wait mode, power down the ADC by setting ADCH4–ADCH0 bits in the ADC status and control register before executing the WAIT instruction.

11.2.2 Stop Mode

The ADC module is inactive after the execution of a STOP instruction. Any pending conversion is aborted. ADC conversions resume when the MCU exits stop mode after an external interrupt. Allow one conversion cycle to stabilize the analog circuitry.

11.3 Break Module (BRK)

11.3.1 Wait Mode

If enabled, the break (BRK) module is active in wait mode. In the break routine, the user can subtract one from the return address on the stack if the SBSW bit in the break status register is set.

11.3.2 Stop Mode

The break module is inactive in stop mode. A break interrupt causes exit from stop mode and sets the SBSW bit in the break status register. The STOP instruction does not affect break module register states.

11.4 Central Processor Unit (CPU)

11.4.1 Wait Mode

The WAIT instruction:

- Clears the interrupt mask (I bit) in the condition code register, enabling interrupts. After exit from wait mode by interrupt, the I bit remains clear. After exit by reset, the I bit is set.
- Disables the CPU clock

11.4.2 Stop Mode

The STOP instruction:

- Clears the interrupt mask (I bit) in the condition code register, enabling external interrupts. After exit from stop mode by external interrupt, the I bit remains clear. After exit by reset, the I bit is set.
- Disables the CPU clock

After exiting stop mode, the CPU clock begins running after the oscillator stabilization delay.

11.5 Internal Clock Generator Module (ICG)

11.5.1 Wait Mode

The internal clock generator (ICG) module remains active in wait mode. If enabled, the ICG interrupt to the CPU can bring the MCU out of wait mode.

In some applications, low power-consumption is desired in wait mode and a high-frequency clock is not needed. In these applications, reduce power consumption by either selecting a low-frequency external clock and turn the internal clock generator off or reduce the bus frequency by minimizing the ICG multiplier factor (N) before executing the WAIT instruction.

11.5.2 Stop Mode

The value of the oscillator enable in stop (OSCENINSTOP) bit in the CONFIG2 register determines the behavior of the ICG in stop mode. If OSCENINSTOP is low, the ICG is disabled in stop and, upon execution of the STOP instruction, all ICG activity will cease and the output clocks (CGMXCLK, CGMOUT, COPCLK, and TBMCLK) will be held low. Power consumption will be minimal.

If OSCENINSTOP is high, the ICG is enabled in stop and activity will continue. This is useful if the timebase module (TBM) is required to bring the MCU out of stop mode. ICG interrupts will not bring the MCU out of stop mode in this case.

During stop mode, if OSCENINSTOP is low, several functions in the ICG are affected. The stable bits (ECGS and ICGS) are cleared, which will enable the external clock stabilization divider upon recovery. The clock monitor is disabled (CMON = 0) which will also clear the clock monitor interrupt enable (CMIE) and clock monitor flag (CMF) bits. The CS, ICGON, ECGON, N, TRIM, DDIV, and DSTG bits are unaffected.

Address:	$000C							
	Bit 7	6	5	4	3	2	1	Bit 0
Read:	0	0	0	DDRE4	DDRE3	DDRE2	DDRE1	DDRE0
Write:								
Reset:	0	0	0	0	0	0	0	0

= Unimplemented

Figure 12-18. Data Direction Register E (DDRE)

DDRE4–DDRE0 — Data Direction Register E Bits

These read/write bits control port E data direction. Reset clears DDRE4–DDRE0, configuring all port E pins as inputs.

1 = Corresponding port E pin configured as output
0 = Corresponding port E pin configured as input

NOTE

Avoid glitches on port E pins by writing to the port E data register before changing data direction register E bits from 0 to 1.

Figure 12-19 shows the port E I/O logic.

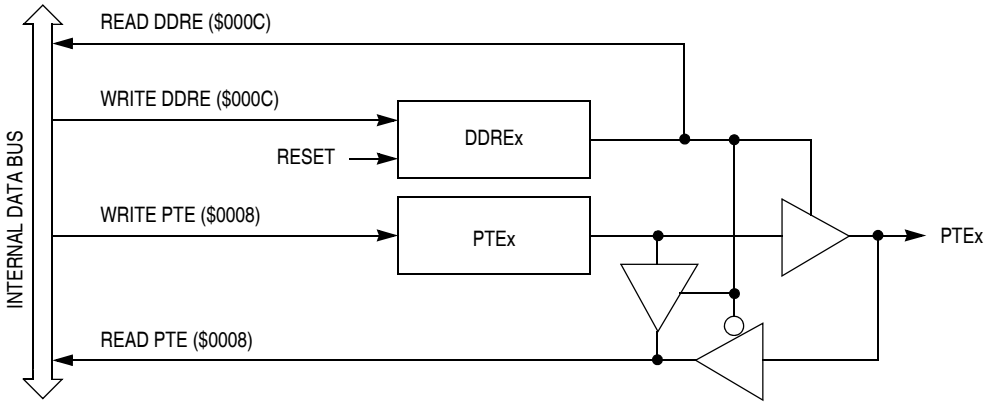


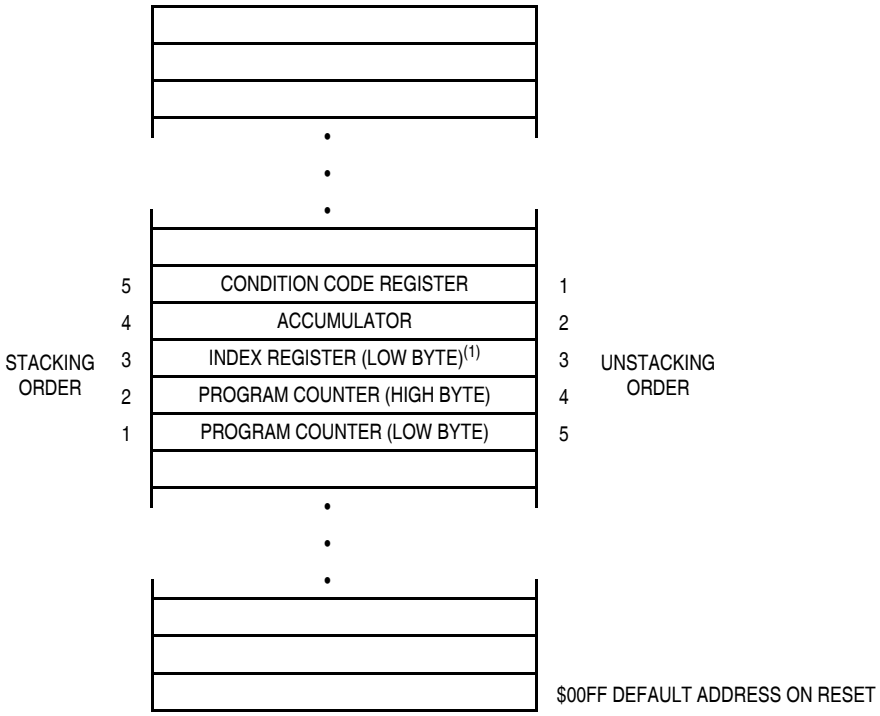
Figure 12-19. Port E I/O Circuit

When bit DDREx is a 1, reading address $0008 reads the PTEx data latch. When bit DDREx is a 0, reading address $0008 reads the voltage level on the pin. The data latch can always be written, regardless of the state of its data direction bit. Table 12-6 summarizes the operation of the port E pins.

Table 12-6. Port E Pin Functions

DDRE Bit	PTE Bit	I/O Pin Mode	Accesses to DDRE	Accesses to PTE	
			Read/Write	Read	Write
0	X[1]	Input, Hi-Z[2]	DDRE4–DDRE0	Pin	PTE4–PTE0[3]
1	X	Output	DDRE4–DDRE0	PTE4–PTE0	PTE4–PTE0

1. X = Don't care
2. Hi-Z = High impedance
3. Writing affects data register, but does not affect input.



1. High byte of index register is not stacked.

Figure 13-4. Interrupt Stacking Order

After every instruction, the CPU checks all pending interrupts if the I bit is not set. If more than one interrupt is pending when an instruction is done, the highest priority interrupt is serviced first. In the example shown in Figure 13-5, if an interrupt is pending upon exit from the interrupt service routine, the pending interrupt is serviced before the LDA instruction is executed.

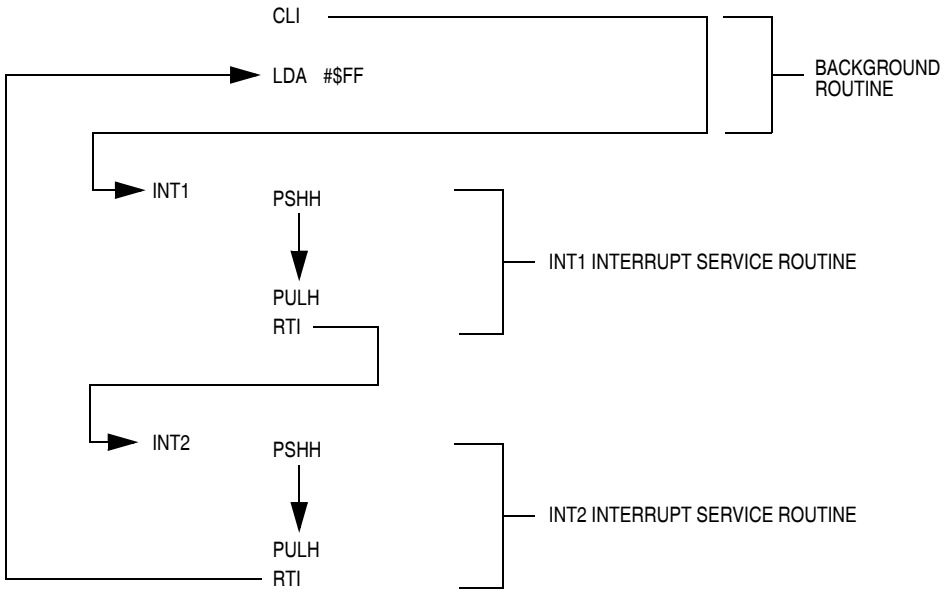


Figure 13-5. Interrupt Recognition Example

Table 14-3. Data Bit Recovery

RT8, RT9, and RT10 Samples	Data Bit Determination	Noise Flag
000	0	0
001	0	1
010	0	1
011	1	1
100	0	1
101	1	1
110	1	1
111	1	0

NOTE

The RT8, RT9, and RT10 samples do not affect start bit verification. If any or all of the RT8, RT9, and RT10 start bit samples are 1s following a successful start bit verification, the noise flag (NF) is set and the receiver assumes that the bit is a start bit.

To verify a stop bit and to detect noise, recovery logic takes samples at RT8, RT9, and RT10. Table 14-4 summarizes the results of the stop bit samples.

Table 14-4. Stop Bit Recovery

RT8, RT9, and RT10 Samples	Framing Error Flag	Noise Flag
000	1	0
001	1	1
010	1	1
011	0	1
100	1	1
101	0	1
110	0	1
111	0	0

14.4.3.4 Framing Errors

If the data recovery logic does not detect a 1 where the stop bit should be in an incoming character, it sets the framing error bit, FE, in SCS1. A break character also sets the FE bit because a break character has no stop bit. The FE bit is set at the same time that the SCRF bit is set.

14.4.3.5 Baud Rate Tolerance

A transmitting device may be operating at a baud rate below or above the receiver baud rate. Accumulated bit time misalignment can cause one of the three stop bit data samples to fall outside the actual stop bit. Then a noise error occurs. If more than one of the samples is outside the stop bit, a framing error occurs. In most applications, the baud rate tolerance is much more than the degree of misalignment that is likely to occur.

As the receiver samples an incoming character, it resynchronizes the RT clock on any valid falling edge within the character. Resynchronization within characters corrects misalignments between transmitter bit times and receiver bit times.

For an 8-bit character, data sampling of the stop bit takes the receiver 9 bit times × 16 RT cycles + 10 RT cycles = 154 RT cycles.

With the misaligned character shown in Figure 14-9, the receiver counts 154 RT cycles at the point when the count of the transmitting device is 10 bit times × 16 RT cycles = 160 RT cycles.

The maximum percent difference between the receiver count and the transmitter count of a fast 8-bit character with no errors is

$$\left|\frac{154 - 160}{154}\right| \times 100 = 3.90\%.$$

For a 9-bit character, data sampling of the stop bit takes the receiver 10 bit times × 16 RT cycles + 10 RT cycles = 170 RT cycles.

With the misaligned character shown in Figure 14-9, the receiver counts 170 RT cycles at the point when the count of the transmitting device is 11 bit times × 16 RT cycles = 176 RT cycles.

The maximum percent difference between the receiver count and the transmitter count of a fast 9-bit character with no errors is:

$$\left|\frac{170 - 176}{170}\right| \times 100 = 3.53\%.$$

14.4.3.6 Receiver Wakeup

So that the MCU can ignore transmissions intended only for other receivers in multiple-receiver systems, the receiver can be put into a standby state. Setting the receiver wakeup bit, RWU, in SCC2 puts the receiver into a standby state during which receiver interrupts are disabled.

Depending on the state of the WAKE bit in SCC1, either of two conditions on the RxD pin can bring the receiver out of the standby state:

1. Address mark — An address mark is a 1 in the MSB position of a received character. When the WAKE bit is set, an address mark wakes the receiver from the standby state by clearing the RWU bit. The address mark also sets the ESCI receiver full bit, SCRF. Software can then compare the character containing the address mark to the user-defined address of the receiver. If they are the same, the receiver remains awake and processes the characters that follow. If they are not the same, software can set the RWU bit and put the receiver back into the standby state.
2. Idle input line condition — When the WAKE bit is clear, an idle character on the RxD pin wakes the receiver from the standby state by clearing the RWU bit. The idle character that wakes the receiver does not set the receiver idle bit, IDLE, or the ESCI receiver full bit, SCRF. The idle line type bit, ILTY, determines whether the receiver begins counting 1s as idle character bits after the start bit or after the stop bit.

NOTE

With the WAKE bit clear, setting the RWU bit after the RxD pin has been idle will cause the receiver to wake up.

a valid character that sets the SCRF bit before an idle condition can set the IDLE bit. Also, after the IDLE bit has been cleared, a valid character must again set the SCRF bit before an idle condition can set the IDLE bit. Reset clears the IDLE bit.

1 = Receiver input idle

0 = Receiver input active (or idle since the IDLE bit was cleared)

OR — Receiver Overrun Bit

This clearable, read-only bit is set when software fails to read the SCDR before the receive shift register receives the next character. The OR bit generates an ESCI error CPU interrupt request if the ORIE bit in SCC3 is also set. The data in the shift register is lost, but the data already in the SCDR is not affected. Clear the OR bit by reading SCS1 with OR set and then reading the SCDR. Reset clears the OR bit.

1 = Receive shift register full and SCRF = 1

0 = No receiver overrun

Software latency may allow an overrun to occur between reads of SCS1 and SCDR in the flag-clearing sequence. Figure 14-14 shows the normal flag-clearing sequence and an example of an overrun caused by a delayed flag-clearing sequence. The delayed read of SCDR does not clear the OR bit because OR was not set when SCS1 was read. Byte 2 caused the overrun and is lost. The next flag-clearing sequence reads byte 3 in the SCDR instead of byte 2.

In applications that are subject to software latency or in which it is important to know which byte is lost due to an overrun, the flag-clearing routine can check the OR bit in a second read of SCS1 after reading the data register.

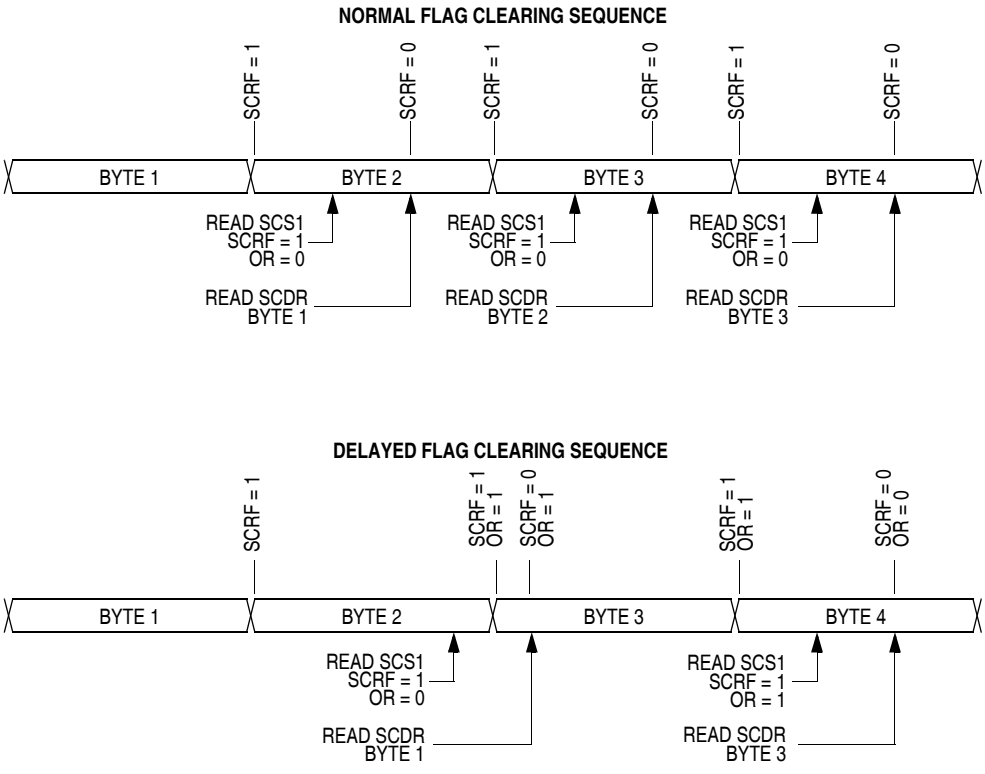
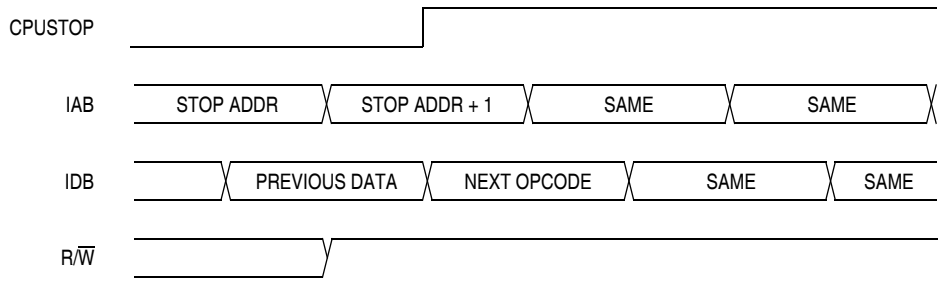


Figure 14-14. Flag Clearing Sequence



Note: Previous data can be operand data or the STOP opcode, depending on the last instruction.

Figure 15-18. Stop Mode Entry Timing

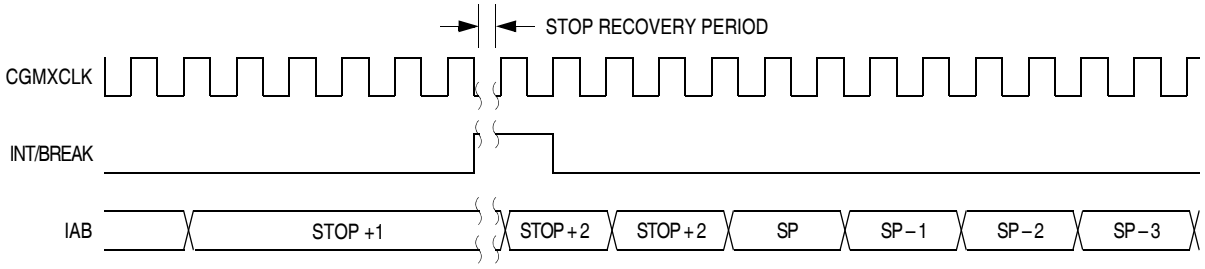


Figure 15-19. Stop Mode Recovery from Interrupt

15.7 SIM Registers

The SIM has three memory-mapped registers. Table 15-4 shows the mapping of these registers.

Table 15-4. SIM Registers

Address	Register	Access Mode
$FE00	SBSR	User
$FE01	SRSR	User
$FE03	SBFCR	User

15.7.1 SIM Break Status Register

The SIM break status register (SBSR) contains a flag to indicate that a break caused an exit from wait mode. This register is only used in emulation mode.

Address: $FE00

	Bit 7	6	5	4	3	2	1	Bit 0
Read:	R	R	R	R	R	R	SBSW	R
Write:							Note(1)	
Reset:	0	0	0	0	0	0	0	0

R = Reserved

Note: 1. Writing a 0 clears SBSW.

Figure 15-20. SIM Break Status Register (SBSR)

SPR1 and SPR0 — SPI Baud Rate Select Bits

In master mode, these read/write bits select one of four baud rates as shown in Table 16-3. SPR1 and SPR0 have no effect in slave mode. Reset clears SPR1 and SPR0.

Table 16-3. SPI Master Baud Rate Selection

SPR1 and SPR0	Baud Rate Divisor (BD)
00	2
01	8
10	32
11	128

Use this formula to calculate the SPI baud rate:

$$\text{Baud rate} = \frac{\text{BUSCLK}}{\text{BD}}$$

16.12.3 SPI Data Register

The SPI data register consists of the read-only receive data register and the write-only transmit data register. Writing to the SPI data register writes data into the transmit data register. Reading the SPI data register reads data from the receive data register. The transmit data and receive data registers are separate registers that can contain different values. See Figure 16-2.

Address: $0012

	Bit 7	6	5	4	3	2	1	Bit 0
Read:	R7	R6	R5	R4	R3	R2	R1	R0
Write:	T7	T6	T5	T4	T3	T2	T1	T0
Reset:	Unaffected by reset							

Figure 16-16. SPI Data Register (SPDR)

R7–R0/T7–T0 — Receive/Transmit Data Bits

NOTE

Do not use read-modify-write instructions on the SPI data register since the register read is not the same as the register written.

Addr.	Register Name		Bit 7	6	5	4	3	2	1	Bit 0
$002C	Timer 2 Counter Register High (T2CNTH) See page 232.	Read:	Bit 15	14	13	12	11	10	9	Bit 8
		Write:								
		Reset:	0	0	0	0	0	0	0	0
$002D	Timer 2 Counter Register Low (T2CNTL) See page 232.	Read:	Bit 7	6	5	4	3	2	1	Bit 0
		Write:								
		Reset:	0	0	0	0	0	0	0	0
$002E	Timer 2 Counter Modulo Register High (T2MODH) See page 233.	Read:	Bit 15	14	13	12	11	10	9	Bit 8
		Write:								
		Reset:	1	1	1	1	1	1	1	1
$002F	Timer 2 Counter Modulo Register Low (T2MODL) See page 233.	Read:	Bit 7	6	5	4	3	2	1	Bit 0
		Write:								
		Reset:	1	1	1	1	1	1	1	1
$0030	Timer 2 Channel 0 Status and Control Register (T2SC0) See page 233.	Read:	CH0F	CH0IE	MS0B	MS0A	ELS0B	ELS0A	TOV0	CH0MAX
		Write:	0							
		Reset:	0	0	0	0	0	0	0	0
$0031	Timer 2 Channel 0 Register High (T2CH0H) See page 236.	Read:	Bit 15	14	13	12	11	10	9	Bit 8
		Write:								
		Reset:	Indeterminate after reset							
$0032	Timer 2 Channel 0 Register Low (T2CH0L) See page 236.	Read:	Bit 7	6	5	4	3	2	1	Bit 0
		Write:								
		Reset:	Indeterminate after reset							
$0033	Timer 2 Channel 1 Status and Control Register (T2SC1) See page 234.	Read:	CH1F	CH1IE	0	MS1A	ELS1B	ELS1A	TOV1	CH1MAX
		Write:	0							
		Reset:	0	0	0	0	0	0	0	0
$0034	Timer 2 Channel 1 Register High (T2CH1H) See page 236.	Read:	Bit 15	14	13	12	11	10	9	Bit 8
		Write:								
		Reset:	Indeterminate after reset							
$0035	Timer 2 Channel 1 Register Low (T2CH1L) See page 236.	Read:	Bit 7	6	5	4	3	2	1	Bit 0
		Write:								
		Reset:	Indeterminate after reset							

= Unimplemented

Figure 18-3. TIM I/O Register Summary (Sheet 2 of 2)

18.4.1 TIM Counter Prescaler

The TIM clock source can be one of the seven prescaler outputs. The prescaler generates seven clock rates from the internal bus clock. The prescaler select bits, PS[2:0], in the TIM status and control register select the TIM clock source.

18.4.2 Input Capture

With the input capture function, the TIM can capture the time at which an external event occurs. When an active edge occurs on the pin of an input capture channel, the TIM latches the contents of the TIM counter

18.4.4 Pulse Width Modulation (PWM)

By using the toggle-on-overflow feature with an output compare channel, the TIM can generate a PWM signal. The value in the TIM counter modulo registers determines the period of the PWM signal. The channel pin toggles when the counter reaches the value in the TIM counter modulo registers. The time between overflows is the period of the PWM signal.

As Figure 18-4 shows, the output compare value in the TIM channel registers determines the pulse width of the PWM signal. The time between overflow and output compare is the pulse width. Program the TIM to clear the channel pin on output compare if the state of the PWM pulse is logic 1. Program the TIM to set the pin if the state of the PWM pulse is logic 0.

The value in the TIM counter modulo registers and the selected prescaler output determines the frequency of the PWM output. The frequency of an 8-bit PWM signal is variable in 256 increments. Writing $00FF (255) to the TIM counter modulo registers produces a PWM period of 256 times the internal bus clock period if the prescaler select value is $000. See 18.9.1 TIM Status and Control Register.

The value in the TIM channel registers determines the pulse width of the PWM output. The pulse width of an 8-bit PWM signal is variable in 256 increments. Writing $0080 (128) to the TIM channel registers produces a duty cycle of 128/256 or 50%.

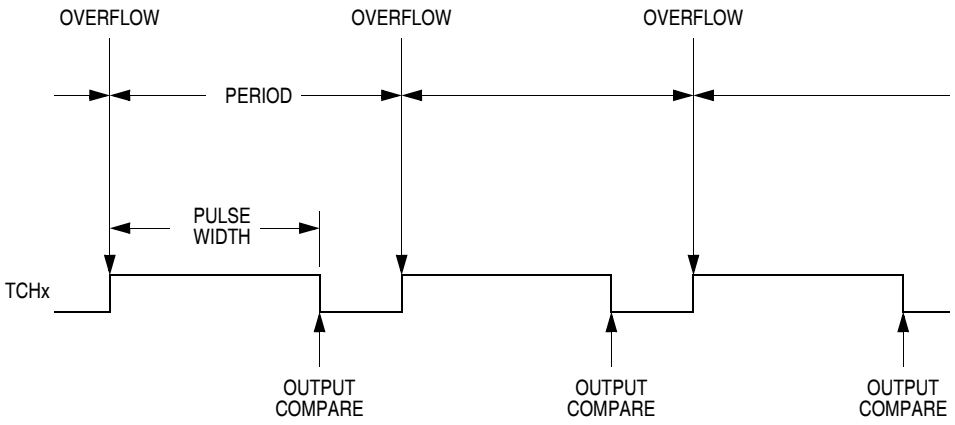


Figure 18-4. PWM Period and Pulse Width

18.4.4.1 Unbuffered PWM Signal Generation

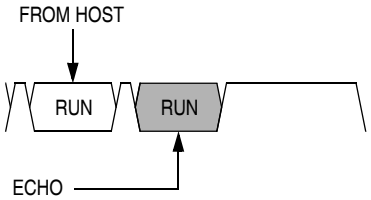
Any output compare channel can generate unbuffered PWM pulses as described in 18.4.4 Pulse Width Modulation (PWM). The pulses are unbuffered because changing the pulse width requires writing the new pulse width value over the old value currently in the TIM channel registers.

An unsynchronized write to the TIM channel registers to change a pulse width value could cause incorrect operation for up to two PWM periods. For example, writing a new value before the counter reaches the old value but after the counter reaches the new value prevents any compare during that PWM period. Also, using a TIM overflow interrupt routine to write a new, smaller pulse width value may cause the compare to be missed. The TIM may pass the new value before it is written.

Use the following methods to synchronize unbuffered changes in the PWM pulse width on channel x:

- When changing to a shorter pulse width, enable channel x output compare interrupts and write the new value in the output compare interrupt routine. The output compare interrupt occurs at the end of the current pulse. The interrupt routine has until the end of the PWM period to write the new value.

Table 19-7. RUN (Run User Program) Command

Description	Executes PULH and RTI instructions
Operand	None
Data Returned	None
Opcode	$28
Command Sequence	

The MCU executes the SWI and PSHH instructions when it enters monitor mode. The RUN command tells the MCU to execute the PULH and RTI instructions. Before sending the RUN command, the host can modify the stacked CPU registers to prepare to run the host program. The READSP command returns the incremented stack pointer value, SP + 1. The high and low bytes of the program counter are at addresses SP + 5 and SP + 6.

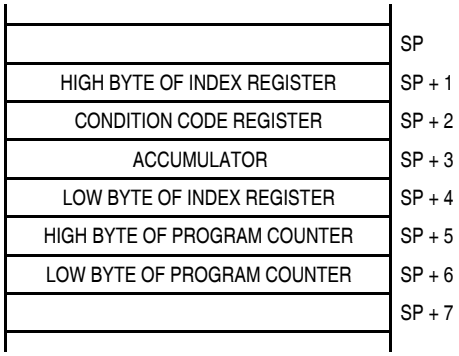


Figure 19-17. Stack Pointer at Monitor Mode Entry

19.3.2 Security

A security feature discourages unauthorized reading of Flash locations while in monitor mode. The host can bypass the security feature at monitor mode entry by sending eight security bytes that match the bytes at locations $FFF6–$FFFD. Locations $FFF6–$FFFD contain user-defined data.

NOTE

Do not leave locations $FFF6–$FFFD blank. For security reasons, program locations $FFF6–$FFFD even if they are not used for vectors.

During monitor mode entry, the MCU waits after the power-on reset for the host to send the eight security bytes on pin PTA0. If the received bytes match those at locations $FFF6–$FFFD, the host bypasses the security feature and can read all Flash locations and execute code from Flash. Security remains bypassed until a power-on reset occurs. If the reset was not a power-on reset, security remains bypassed and security code entry is not required. See Figure 19-18.

20.16 ADC Characteristics

Characteristic[1]	Symbol	Min	Max	Unit	Comments
Supply voltage	V_{DDA}	2.7 (V_{DD} min)	5.5 (V_{DD} max)	V	V_{DDA} should be tied to the same potential as V_{DD} via separate traces.
Input voltages	V_{ADIN}	0	V_{DDA}	V	
Resolution	B_{AD}	8	8	Bits	
Absolute accuracy (V_{REFL} = 0 V, V_{REFH} = V_{DDA} = 5 V ± 10%)	A_{AD}	—	± 1	LSB	Includes quantization
ADC internal clock	f_{ADIC}	0.5	1.048	MHz	$t_{AIC} = 1/f_{ADIC}$, tested only at 1 MHz
Conversion range	R_{AD}	V_{REFL}	V_{REFH}	V	$V_{SSA} \le V_{ADIN} \le V_{DDA}$
Power-up time	t_{ADPU}	16		t_{AIC} cycles	
ADC voltage reference high	V_{REFH}	V_{SSA} – 0.1	V_{DDA} + 0.1	V	$V_{REFL} \le V_{REFH}$
ADC voltage reference low	V_{REFL}	V_{SSA} – 0.1	V_{DDA} + 0.1	V	$V_{REFL} \le V_{REFH}$
Conversion time	t_{ADC}	16	17	t_{AIC} cycles	
Sample time[2]	t_{ADS}	5	—	t_{AIC} cycles	
Zero input reading[3]	Z_{ADI}	00	01	Hex	$V_{IN} = V_{REFL}$
Full-scale reading[3]	F_{ADI}	FE	FF	Hex	$V_{IN} = V_{REFH}$
Input capacitance	C_{ADI}	—	8	pF	Not tested
Input leakage[4] Port B	—	—	± 1	µA	

1. V_{DD} = 5.0 Vdc ± 10%, V_{SS} = 0 Vdc, V_{DDA} = 5.0 Vdc ± 10%, V_{SSA} = 0 Vdc, V_{REFH} = 5.0 Vdc ± 10%, V_{REFL} = 0
2. Source impedances greater than 10 kΩ adversely affect internal RC charging time during input sampling.
3. Zero-input/full-scale reading requires sufficient decoupling measures for accurate conversions.
4. The external system error caused by input leakage current is approximately equal to the product of R source and input current.