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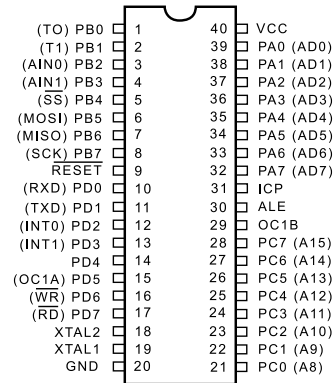
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Details

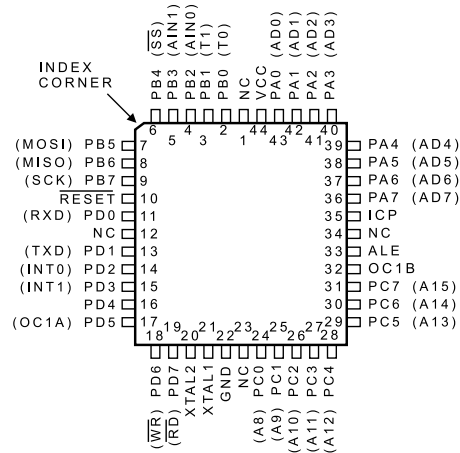
Product Status	Obsolete
Core Processor	AVR
Core Size	8-Bit
Speed	8MHz
Connectivity	SPI, UART/USART
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	32
Program Memory Size	8KB (4K x 16)
Program Memory Type	FLASH
EEPROM Size	512 x 8
RAM Size	512 x 8
Voltage - Supply (Vcc/Vdd)	4V ~ 6V
Data Converters	-
Oscillator Type	Internal
Operating Temperature	0°C ~ 70°C (TA)
Mounting Type	Through Hole
Package / Case	40-DIP (0.600", 15.24mm)
Supplier Device Package	40-PDIP
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/at90s8515a-8pc

Pin Configurations

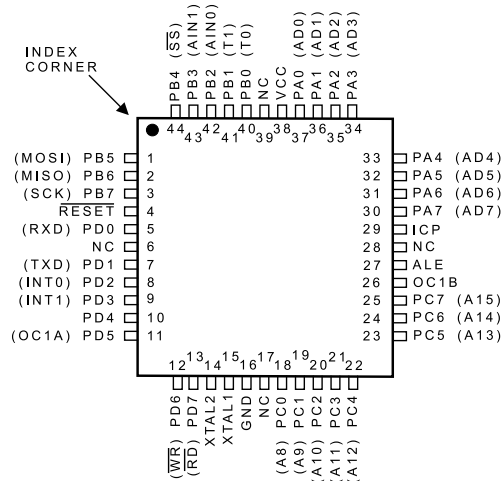
PDIP



PLCC



TQFP



Architectural Overview

The fast-access register file concept contains 32 x 8-bit general-purpose working registers with a single clock cycle access time. This means that during one single clock cycle, one ALU (Arithmetic Logic Unit) operation is executed. Two operands are output from the register file, the operation is executed and the result is stored back in the register file – in one clock cycle.

Six of the 32 registers can be used as three 16-bit indirect address register pointers for Data Space addressing, enabling efficient address calculations. One of the three address pointers is also used as the address pointer for the constant table look-up function. These added function registers are the 16-bit X-, Y-, and Z-register.

The ALU supports arithmetic and logic functions between registers or between a constant and a register. Single register operations are also executed in the ALU. Figure 4 shows the AT90S8515 AVR RISC microcontroller architecture.

In addition to the register operation, the conventional memory addressing modes can be used on the register file as well. This is enabled by the fact that the register file is assigned the 32 lowermost Data Space addresses (\$00 - \$1F), allowing them to be accessed as though they were ordinary memory locations.

The I/O memory space contains 64 addresses for CPU peripheral functions such as Control Registers, Timer/Counters, A/D converters and other I/O functions. The I/O memory can be accessed directly or as the Data Space locations following those of the register file, \$20 - \$5F.

The AVR uses a Harvard architecture concept – with separate memories and buses for program and data. The program memory is executed with a two-stage pipeline. While one instruction is being executed, the next instruction is pre-fetched from the program memory. This concept enables instructions to be executed in every clock cycle. The program memory is In-System Programmable Flash memory.

With the relative jump and call instructions, the whole 4K address space is directly accessed. Most AVR instructions have a single 16-bit word format. Every program memory address contains a 16- or 32-bit instruction.

During interrupts and subroutine calls, the return address Program Counter (PC) is stored on the stack. The stack is effectively allocated in the general data SRAM and consequently, the stack size is only limited by the total SRAM size and the usage of the SRAM. All user programs must initialize the SP in the reset routine (before subroutines or interrupts are executed). The 16-bit Stack Pointer (SP) is read/write-accessible in the I/O space.

The 512-byte data SRAM can be easily accessed through the five different addressing modes supported in the AVR architecture.

The memory spaces in the AVR architecture are all linear and regular memory maps.

In the different addressing modes these address registers have functions as fixed displacement, automatic increment and decrement (see the descriptions for the different instructions).

ALU – Arithmetic Logic Unit

The high-performance AVR ALU operates in direct connection with all the 32 general-purpose working registers. Within a single clock cycle, ALU operations between registers in the register file are executed. The ALU operations are divided into three main categories: arithmetic, logical and bit functions.

In-System Programmable Flash Program Memory

The AT90S8515 contains 8K bytes On-chip In-System Programmable Flash memory for program storage. Since all instructions are 16- or 32-bit words, the Flash is organized as 4K x 16. The Flash memory has an endurance of at least 1000 write/erase cycles. The AT90S8515 Program Counter (PC) is 12 bits wide, thus addressing the 4096 program memory addresses.

See page 86 for a detailed description of Flash data downloading.

See page 13 for the different program memory addressing modes.

- **Bit 6 – INTF0: External Interrupt Flag0**

When an edge on the INT0 pin triggers an interrupt request, the corresponding interrupt flag, INTF0, becomes set (one). If the I-bit in SREG and the corresponding interrupt enable bit, INT0 in GIMSK are set (one), the MCU will jump to the interrupt vector. The flag is cleared when the interrupt routine is executed. Alternatively, the flag is cleared by writing a logical “1” to it. This flag is always cleared when INT0 is configured as level interrupt.

- **Bits 5..0 – Res: Reserved Bits**

These bits are reserved bits in the AT90S8515 and always read as zero.

Timer/Counter Interrupt Mask Register – TIMSK

Bit	7	6	5	4	3	2	1	0	
\$39 (\$59)	TOIE1	OCIE1A	OCIE1B	–	TICIE1	–	TOIE0	–	TIMSK
Read/Write	R/W	R/W	R/W	R	R/W	R	R/W	R	
Initial Value	0	0	0	0	0	0	0	0	

- **Bit 7 – TOIE1: Timer/Counter1 Overflow Interrupt Enable**

When the TOIE1 bit is set (one) and the I-bit in the Status Register is set (one), the Timer/Counter1 Overflow interrupt is enabled. The corresponding interrupt (at vector \$006) is executed if an overflow in Timer/Counter1 occurs, i.e., when the TOV1 bit is set in the Timer/Counter Interrupt Flag Register (TIFR).

- **Bit 6 – OCE1A: Timer/Counter1 Output CompareA Match Interrupt Enable**

When the OCIE1A bit is set (one) and the I-bit in the Status Register is set (one), the Timer/Counter1 CompareA Match interrupt is enabled. The corresponding interrupt (at vector \$004) is executed if a CompareA match in Timer/Counter1 occurs, i.e., when the OCF1A bit is set in the Timer/Counter Interrupt Flag Register (TIFR).

- **Bit 5 – OCIE1B: Timer/Counter1 Output CompareB Match Interrupt Enable**

When the OCIE1B bit is set (one) and the I-bit in the Status Register is set (one), the Timer/Counter1 CompareB Match interrupt is enabled. The corresponding interrupt (at vector \$005) is executed if a CompareB match in Timer/Counter1 occurs, i.e., when the OCF1B bit is set in the Timer/Counter Interrupt Flag Register (TIFR).

- **Bit 4 – Res: Reserved Bit**

This bit is a reserved bit in the AT90S8515 and always reads zero.

- **Bit 3 – TICIE1: Timer/Counter1 Input Capture Interrupt Enable**

When the TICIE1 bit is set (one) and the I-bit in the Status Register is set (one), the Timer/Counter1 Input Capture Event interrupt is enabled. The corresponding interrupt (at vector \$003) is executed if a capture-triggering event occurs on pin 31, ICP, i.e., when the ICF1 bit is set in the Timer/Counter Interrupt Flag Register (TIFR).

- **Bit 2 – Res: Reserved Bit**

This bit is a reserved bit in the AT90S8515 and always reads zero.

- **Bit 1 – TOIE0: Timer/Counter0 Overflow Interrupt Enable**

When the TOIE0 bit is set (one) and the I-bit in the Status Register is set (one), the Timer/Counter0 Overflow interrupt is enabled. The corresponding interrupt (at vector \$007) is executed if an overflow in Timer/Counter0 occurs, i.e., when the TOV0 bit is set in the Timer/Counter Interrupt Flag Register (TIFR).

- **Bit 0 – Res: Reserved Bit**

This bit is a reserved bit in the AT90S8515 and always reads zero.

Timer/Counter Interrupt Flag Register – TIFR

Bit	7	6	5	4	3	2	1	0	
\$38 (\$58)	TOV1	OCF1A	OCF1B	–	ICF1	–	TOV0	–	TIFR
Read/Write	R/W	R/W	R/W	R	R/W	R	R/W	R	
Initial Value	0	0	0	0	0	0	0	0	

• Bit 7 – TOV1: Timer/Counter1 Overflow Flag

The TOV1 is set (one) when an overflow occurs in Timer/Counter1. TOV1 is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, TOV1 is cleared by writing a logical “1” to the flag. When the I-bit in SREG, TOIE1 (Timer/Counter1 Overflow Interrupt Enable) and TOV1 are set (one), the Timer/Counter1 Overflow interrupt is executed. In PWM mode, this bit is set when Timer/Counter1 changes counting direction at \$0000.

• Bit 6 – OCF1A: Output Compare Flag 1A

The OCF1A bit is set (one) when compare match occurs between the Timer/Counter1 and the data in OCR1A (Output Compare Register 1A). OCF1A is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, OCF1A is cleared by writing a logical “1” to the flag. When the I-bit in SREG, OCIE1A (Timer/Counter1 Compare Match InterruptA Enable) and the OCF1A are set (one), the Timer/Counter1 CompareA Match interrupt is executed.

• Bit 5 – OCF1B: Output Compare Flag 1B

The OCF1B bit is set (one) when compare match occurs between the Timer/Counter1 and the data in OCR1B (Output Compare Register 1B). OCF1B is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, OCF1B is cleared by writing a logical “1” to the flag. When the I-bit in SREG, OCIE1B (Timer/Counter1 Compare Match InterruptB Enable) and the OCF1B are set (one), the Timer/Counter1 CompareB Match interrupt is executed.

• Bit 4 – Res: Reserved Bit

This bit is a reserved bit in the AT90S8515 and always reads zero.

• Bit 3 – ICF1: Input Capture Flag 1

The ICF1 bit is set (one) to flag an input capture event, indicating that the Timer/Counter1 value has been transferred to the input capture register (ICR1). ICF1 is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, ICF1 is cleared by writing a logical “1” to the flag. When the SREG I-bit, TICIE1 (Timer/Counter1 Input Capture Interrupt Enable) and ICF1 are set (one), the Timer/Counter1 Capture interrupt is executed.

• Bit 2 – Res: Reserved Bit

This bit is a reserved bit in the AT90S8515 and always reads zero.

• Bit 1 – TOV: Timer/Counter0 Overflow Flag

The bit TOV0 is set (one) when an overflow occurs in Timer/Counter0. TOV0 is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, TOV0 is cleared by writing a logical “1” to the flag. When the SREG I-bit, TOIE0 (Timer/Counter0 Overflow Interrupt Enable) and TOV0 are set (one), the Timer/Counter0 Overflow interrupt is executed.

• Bit 0 – Res: Reserved Bit

This bit is a reserved bit in the AT90S8515 and always reads zero.

Sleep Modes

To enter the sleep modes, the SE bit in MCUCR must be set (one) and a SLEEP instruction must be executed. If an enabled interrupt occurs while the MCU is in a sleep mode, the MCU awakes, executes the interrupt routine and resumes execution from the instruction following SLEEP. The contents of the register file, SRAM and I/O memory are unaltered. If a reset occurs during Sleep Mode, the MCU wakes up and executes from the Reset vector.

Idle Mode

When the SM bit is cleared (zero), the SLEEP instruction forces the MCU into the Idle Mode, stopping the CPU but allowing Timer/Counters, Watchdog and the interrupt system to continue operating. This enables the MCU to wake up from external triggered interrupts as well as internal ones like Timer Overflow interrupt and Watchdog reset. If wake-up from the Analog Comparator interrupt is not required, the Analog Comparator can be powered down by setting the ACD-bit in the Analog Comparator Control and Status Register (ACSR). This will reduce power consumption in Idle Mode. When the MCU wakes up from Idle Mode, the CPU starts program execution immediately.

Power-down Mode

When the SM bit is set (one), the SLEEP instruction forces the MCU into the Power-down mode. In this mode, the external oscillator is stopped, while the external interrupts and the Watchdog (if enabled) continue operating. Only an external reset, a Watchdog reset (if enabled), or an external level interrupt on INT0 or INT1 can wake up the MCU.

Note that when a level-triggered interrupt is used for wake-up from power-down, the low level must be held for a time longer than the reset delay Time-out period t_{TOU} . Otherwise, the MCU will fail to wake up.

1. In the same operation, write a logical “1” to WDTOE and WDE. A logical “1” must be written to WDE even though it is set to one before the disable operation starts.
2. Within the next four clock cycles, write a logical “0” to WDE. This disables the Watchdog.

• **Bits 2..0 – WDP2, WDP1, WDP0: Watchdog Timer Prescaler 2, 1 and 0**

The WDP2, WDP1 and WDP0 bits determine the Watchdog Timer prescaling when the Watchdog Timer is enabled. The different prescaling values and their corresponding Time-out periods are shown in Table 14.

Table 14. Watchdog Timer Prescale Select

WDP2	WDP1	WDP0	Number of WDT Oscillator Cycles	Typical Time-out at $V_{CC} = 3.0V$	Typical Time-out at $V_{CC} = 5.0V$
0	0	0	16K cycles	47.0 ms	15.0 ms
0	0	1	32K cycles	94.0 ms	30.0 ms
0	1	0	64K cycles	0.19 s	60.0 ms
0	1	1	128K cycles	0.38 s	0.12 s
1	0	0	256K cycles	0.75 s	0.24 s
1	0	1	512K cycles	1.5 s	0.49 s
1	1	0	1,024K cycles	3.0 s	0.97 s
1	1	1	2,048K cycles	6.0 s	1.9 s

Note: The frequency of the Watchdog oscillator is voltage-dependent as shown in the Electrical Characteristics section.

The WDR (Watchdog Reset) instruction should always be executed before the Watchdog Timer is enabled. This ensures that the reset period will be in accordance with the Watchdog Timer prescale settings. If the Watchdog Timer is enabled without reset, the Watchdog Timer may not start to count from zero.

To avoid unintentional MCU reset, the Watchdog Timer should be disabled or reset before changing the Watchdog Timer Prescale Select.

Prevent EEPROM Corruption

During periods of low V_{CC} , the EEPROM data can be corrupted because the supply voltage is too low for the CPU and the EEPROM to operate properly. These issues are the same as for board level systems using the EEPROM and the same design solutions should be applied.

An EEPROM data corruption can be caused by two situations when the voltage is too low. First, a regular write sequence to the EEPROM requires a minimum voltage to operate correctly. Second, the CPU itself can execute instructions incorrectly if the supply voltage for executing instructions is too low.

EEPROM data corruption can easily be avoided by following these design recommendations (one is sufficient):

1. Keep the AVR RESET active (low) during periods of insufficient power supply voltage. This is best done by an external low V_{CC} Reset Protection circuit, often referred to as a Brown-out Detector (BOD). Please refer to application note AVR 180 for design considerations regarding power-on reset and low-voltage detection.
2. Keep the AVR core in Power-down Sleep mode during periods of low V_{CC} . This will prevent the CPU from attempting to decode and execute instructions, effectively protecting the EEPROM registers from unintentional writes.
3. Store constants in Flash memory if the ability to change memory contents from software is not required. Flash memory cannot be updated by the CPU and will not be subject to corruption.

UART Control

UART I/O Data Register – UDR

Bit	7	6	5	4	3	2	1	0	
\$0C (\$2C)	MSB							LSB	UDR
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

The UDR register is actually two physically separate registers sharing the same I/O address. When writing to the register, the UART Transmit Data register is written. When reading from UDR, the UART Receive Data register is read.

UART Status Register – USR

Bit	7	6	5	4	3	2	1	0	
\$0B (\$2B)	RXC	TXC	UDRE	FE	OR	–	–	–	USR
Read/Write	R	R/W	R	R	R	R	R	R	
Initial Value	0	0	1	0	0	0	0	0	

The USR register is a read-only register providing information on the UART status.

- **Bit 7 – RXC: UART Receive Complete**

This bit is set (one) when a received character is transferred from the Receiver Shift register to UDR. The bit is set regardless of any detected framing errors. When the RXCIE bit in UCR is set, the UART Receive Complete interrupt will be executed when RXC is set (one). RXC is cleared by reading UDR. When interrupt-driven data reception is used, the UART Receive Complete Interrupt routine must read UDR in order to clear RXC, otherwise a new interrupt will occur once the interrupt routine terminates.

- **Bit 6 – TXC: UART Transmit Complete**

This bit is set (one) when the entire character (including the stop bit) in the Transmit Shift register has been shifted out and no new data has been written to UDR. This flag is especially useful in half-duplex communications interfaces, where a transmitting application must enter receive mode and free the communications bus immediately after completing the transmission.

When the TXCIE bit in UCR is set, setting of TXC causes the UART Transmit Complete interrupt to be executed. TXC is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, the TXC bit is cleared (zero) by writing a logical “1” to the bit.

- **Bit 5 – UDRE: UART Data Register Empty**

This bit is set (one) when a character written to UDR is transferred to the Transmit Shift register. Setting of this bit indicates that the transmitter is ready to receive a new character for transmission.

When the UDRIE bit in UCR is set, the UART Transmit Complete interrupt to be executed as long as UDRE is set. UDRE is cleared by writing UDR. When interrupt-driven data transmittal is used, the UART Data Register Empty Interrupt routine must write UDR in order to clear UDRE, otherwise a new interrupt will occur once the interrupt routine terminates.

UDRE is set (one) during reset to indicate that the transmitter is ready.

- **Bit 4 – FE: Framing Error**

This bit is set if a Framing Error condition is detected, i.e., when the stop bit of an incoming character is zero.

Table 19. DDAn Effects on Port A Pins

Note: n: 7,6...0, pin number.

Note that all port pins are synchronized. The synchronization latch is, however, not shown in the figure.

WP: WRITE PORTA
WD: WRITE DDRA
RL: READ PORTA LATCH
RP: READ PORTA PIN
RD: READ DDRA
SRE: EXT. SRAM ENABLE
A: ADDRESS
W: WRITE
R: READ
n: 0-7

Port B

Port B is an 8-bit bi-directional I/O port.

Three I/O memory address locations are allocated for the Port B, one each for the Data Register – PORTB, \$18(\$38), Data Direction Register – DDRB, \$17(\$37) and the Port B Input Pins – PINB, \$16(\$36). The Port B Input Pins address is read-only, while the Data Register and the Data Direction Register are read/write.

All port pins have individually selectable pull-up resistors. The Port B output buffers can sink 20 mA and thus drive LED displays directly. When pins PB0 to PB7 are used as inputs and are externally pulled low, they will source current if the internal pull-up resistors are activated.

The Port B pins with alternate functions are shown in Table 20.

Table 20. Port B Pin Alternate Functions

Port Pin	Alternate Functions
PB0	T0 (Timer/Counter 0 External Counter Input)
PB1	T1 (Timer/Counter 1 External Counter Input)
PB2	AIN0 (Analog Comparator positive input)
PB3	AIN1 (Analog Comparator negative input)
PB4	$\overline{SS}$ (SPI Slave Select Input)
PB5	MOSI (SPI Bus Master Output/Slave Input)
PB6	MISO (SPI Bus Master Input/Slave Output)
PB7	SCK (SPI Bus Serial Clock)

When the pins are used for the alternate function, the DDRB and PORTB registers have to be set according to the alternate function description.

Port B Data Register – PORTB

Bit	7	6	5	4	3	2	1	0	
\$18 (\$38)	PORTB7	PORTB6	PORTB5	PORTB4	PORTB3	PORTB2	PORTB1	PORTB0	PORTB
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

Port B Data Direction Register – DDRB

Bit	7	6	5	4	3	2	1	0	
\$17 (\$37)	DDB7	DDB6	DDB5	DDB4	DDB3	DDB2	DDB1	DDB0	DDRB
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

Port B Input Pins Address – PINB

Bit	7	6	5	4	3	2	1	0	
\$16 (\$36)	PINB7	PINB6	PINB5	PINB4	PINB3	PINB2	PINB1	PINB0	PINB
Read/Write	R	R	R	R	R	R	R	R	
Initial Value	N/A	N/A	N/A	N/A	N/A	N/A	N/A	N/A	

The Port B Input Pins address (PINB) is not a register; this address enables access to the physical value on each Port B pin. When reading PORTB, the Port B Data Latch is read and when reading PINB, the logical values present on the pins are read.

Port C Data Direction Register – DDRC

Bit	7	6	5	4	3	2	1	0	
\$14 (\$34)	DDC7	DDC6	DDC5	DDC4	DDC3	DDC2	DDC1	DDC0	DDRC
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

Port C Input Pins Address – PINC

Bit	7	6	5	4	3	2	1	0	
\$13 (\$33)	PINC7	PINC6	PINC5	PINC4	PINC3	PINC2	PINC1	PINC0	PINC
Read/Write	R	R	R	R	R	R	R	R	
Initial Value	N/A	N/A	N/A	N/A	N/A	N/A	N/A	N/A	

The Port C Input Pins address (PINC) is not a register; this address enables access to the physical value on each Port C pin. When reading PORTC, the Port C Data Latch is read and when reading PINC, the logical values present on the pins are read.

Port C as General Digital I/O

All eight pins in Port C have equal functionality when used as digital I/O pins.

PCn, general I/O pin: The DDCn bit in the DDRC register selects the direction of this pin. If DDCn is set (one), PCn is configured as an output pin. If DDCn is cleared (zero), PCn is configured as an input pin. If PORTCn is set (one) when the pin is configured as an input pin, the MOS pull-up resistor is activated. To switch the pull-up resistor off, PORTCn has to be cleared (zero) or the pin has to be configured as an output pin. The Port C pins are tri-stated when a reset condition becomes active, even if the clock is not active.

Table 22. DDCn Effects on Port C Pins

DDCn	PORTCn	I/O	Pull-up	Comment
0	0	Input	No	Tri-state (high-Z)
0	1	Input	Yes	PCn will source current if ext. pulled low.
1	0	Output	No	Push-pull Zero Output
1	1	Output	No	Push-pull One Output

Note: n: 7...0, pin number

Table 27. Pin Name Mapping

Signal Name in Programming Mode	Pin Name	I/O	Function
RDY/ $\overline{\text{BSY}}$	PD1	O	0: Device is busy programming, 1: Device is ready for new command
$\overline{\text{OE}}$	PD2	I	Output Enable (Active low)
$\overline{\text{WR}}$	PD3	I	Write Pulse (Active low)
BS	PD4	I	Byte Select ("0" selects low byte, "1" selects high byte)
XA0	PD5	I	XTAL Action Bit 0
XA1	PD6	I	XTAL Action Bit 1
DATA	PB7-0	I/O	Bi-directional Data Bus (Output when $\overline{\text{OE}}$ is low)

Table 28. XA1 and XA0 Coding

XA1	XA0	Action when XTAL1 is Pulsed
0	0	Load Flash or EEPROM Address (High or low address byte determined by BS)
0	1	Load Data (High or low data byte for Flash determined by BS)
1	0	Load Command
1	1	No Action, Idle

Table 29. Command Byte Bit Coding

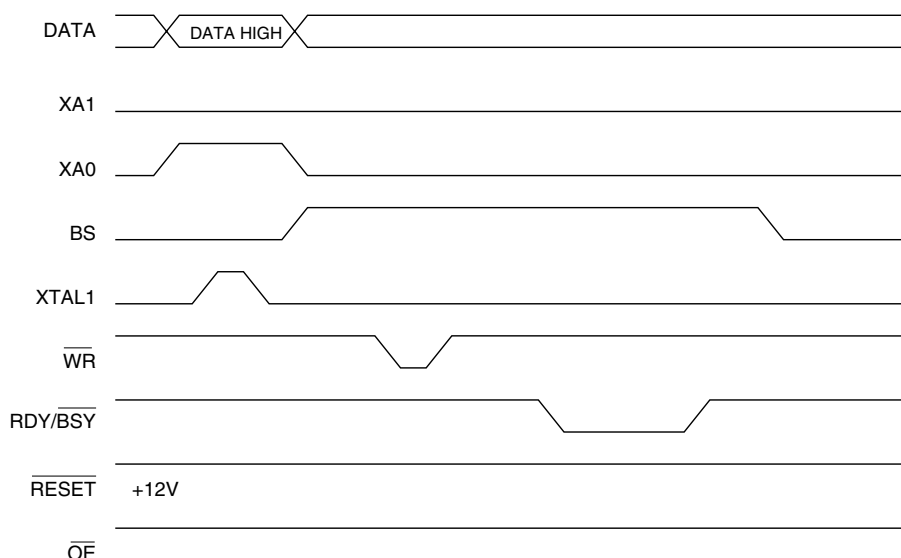
Command Byte	Command Executed
1000 0000	Chip Erase
0100 0000	Write Fuse Bits
0010 0000	Write Lock Bits
0001 0000	Write Flash
0001 0001	Write EEPROM
0000 1000	Read Signature Bytes
0000 0100	Read Lock and Fuse Bits
0000 0010	Read Flash
0000 0011	Read EEPROM

Enter Programming Mode

The following algorithm puts the device in Parallel Programming Mode:

1. Apply supply voltage according to Table 26, between V_{CC} and GND.
2. Set the $\overline{\text{RESET}}$ and BS pin to "0" and wait at least 100 ns.
3. Apply 11.5 - 12.5V to $\overline{\text{RESET}}$. Any activity on BS within 100 ns after +12V has been applied to $\overline{\text{RESET}}$ will cause the device to fail entering programming mode.

Figure 62. Programming the Flash Waveforms (Continued)



Reading the Flash

The algorithm for reading the Flash memory is as follows (refer to “Programming the Flash” for details on command and address loading):

1. A: Load Command “0000 0010”.
2. B: Load Address High Byte (\$00 - \$0F).
3. C: Load Address Low Byte (\$00 - \$FF).
4. Set $\overline{OE}$ to “0”, and BS to “0”. The Flash word low byte can now be read at DATA.
5. Set BS to “1”. The Flash word high byte can now be read from DATA.
6. Set $\overline{OE}$ to “1”.

Programming the EEPROM

The programming algorithm for the EEPROM data memory is as follows (refer to “Programming the Flash” for details on command, address and data loading):

1. A: Load Command “0001 0001”.
2. (AT90S8515 only) B: Load Address High Byte (\$00 - \$01).
3. C: Load Address Low Byte (\$00 - \$FF).
4. D: Load Data Low Byte (\$00 - \$FF).
5. E: Write Data Low Byte.

Reading the EEPROM

The algorithm for reading the EEPROM memory is as follows (refer to “Programming the Flash” for details on command and address loading):

1. A: Load Command “0000 0011”.
2. (AT90S8515 only) B: Load Address High Byte (\$00 - \$01).
3. C: Load Address Low Byte (\$00 - \$FF).
4. Set $\overline{OE}$ to “0”, and BS to “0”. The EEPROM data byte can now be read at DATA.
5. Set $\overline{OE}$ to “1”.

Programming the Fuse Bits

The algorithm for programming the Fuse bits is as follows (refer to “Programming the Flash” for details on command and data loading):

1. A: Load Command “0100 0000”.
2. D: Load Data Low Byte. Bit n = “0” programs and bit n = “1” erases the Fuse bit.

Parallel Programming Characteristics

Figure 63. Parallel Programming Timing

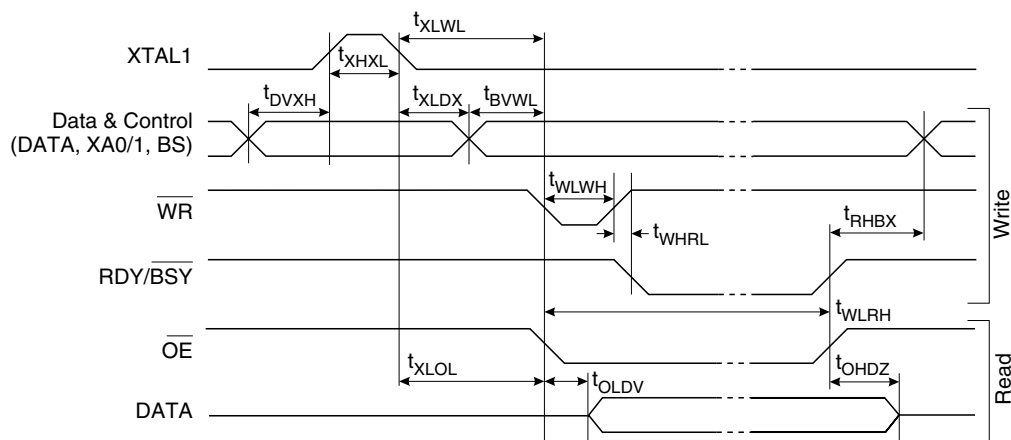


Table 30. Parallel Programming Characteristics, $T_A = 25^\circ\text{C} \pm 10\%$, $V_{CC} = 5\text{V} \pm 10\%$

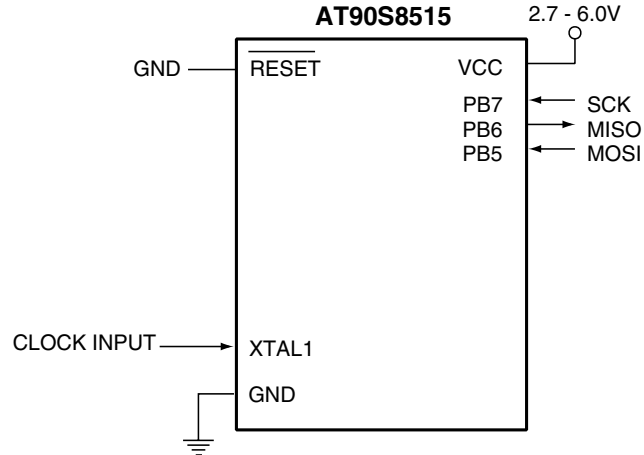
Symbol	Parameter	Min	Typ	Max	Units
V_{PP}	Programming Enable Voltage	11.5		12.5	V
I_{PP}	Programming Enable Current			250.0	μA
t_{DVXH}	Data and Control Setup before XTAL1 High	67.0			ns
t_{XHXL}	XTAL1 Pulse Width High	67.0			ns
t_{XLDX}	Data and Control Hold after XTAL1 Low	67.0			ns
t_{XLWL}	XTAL1 Low to $\overline{WR}$ Low	67.0			ns
t_{BVWL}	BS Valid to $\overline{WR}$ Low	67.0			ns
t_{RHBX}	BS Hold after RDY/ $\overline{BSY}$ High	67.0			ns
t_{WLWH}	$\overline{WR}$ Pulse Width Low ⁽¹⁾	67.0			ns
t_{WHRL}	$\overline{WR}$ High to RDY/ $\overline{BSY}$ Low ⁽²⁾		20.0		ns
t_{WLRH}	$\overline{WR}$ Low to RDY/ $\overline{BSY}$ High ⁽²⁾	0.5	0.7	0.9	ms
t_{XLOL}	XTAL1 Low to $\overline{OE}$ Low	67.0			ns
t_{OLDV}	$\overline{OE}$ Low to DATA Valid		20.0		ns
t_{OHDZ}	$\overline{OE}$ High to DATA Tri-stated			20.0	ns
t_{WLWH_CE}	$\overline{WR}$ Pulse Width Low for Chip Erase	5.0	10.0	15.0	ms
t_{WLWH_PFB}	$\overline{WR}$ Pulse Width Low for Programming the Fuse Bits	1.0	1.5	1.8	ms

Notes: 1. Use t_{WLWH_CE} for Chip Erase and t_{WLWH_PFB} for programming the Fuse bits.
2. If t_{WLWH} is held longer than t_{WLRH} , no RDY/ $\overline{BSY}$ pulse will be seen.

Serial Downloading

Both the program and data memory arrays can be programmed using the SPI bus while $\overline{\text{RESET}}$ is pulled to GND. The serial interface consists of pins SCK, MOSI (input) and MISO (output). See Figure 64. After $\overline{\text{RESET}}$ is set low, the Programming Enable instruction needs to be executed first before program/erase instructions can be executed.

Figure 64. Serial Programming and Verify



For the EEPROM, an auto-erase cycle is provided within the self-timed Write instruction and there is no need to first execute the Chip Erase instruction. The Chip Erase instruction turns the content of every memory location in both the program and EEPROM arrays into \$FF.

The program and EEPROM memory arrays have separate address spaces: \$0000 to \$0FFF (AT90S8515) for program memory and \$0000 to \$01FF (AT90S8515) for EEPROM memory.

Either an external clock is supplied at pin XTAL1 or a crystal needs to be connected across pins XTAL1 and XTAL2. The minimum low and high periods for the serial clock (SCK) input are defined as follows:

Low: > 2 XTAL1 clock cycles

High: > 2 XTAL1 clock cycles

Serial Programming Algorithm

When writing serial data to the AT90S8515, data is clocked on the rising edge of SCK.

When reading data from the AT90S8515, data is clocked on the falling edge of SCK. See Figure 65, Figure 66 and Table 33 on page 89 for timing details.

To program and verify the AT90S8515 in the Serial Programming Mode, the following sequence is recommended (see 4-byte instruction formats in Table 32):

1. Power-up sequence:

Apply power between V_{CC} and GND while $\overline{\text{RESET}}$ and SCK are set to "0". If a crystal is not connected across pins XTAL1 and XTAL2, apply a clock signal to the XTAL1 pin. In some systems, the programmer cannot guarantee that SCK is held low during power-up. In this case, $\overline{\text{RESET}}$ must be given a positive pulse of at least two XTAL1 cycles duration after SCK has been set to "0".

2. Wait for at least 20 ms and enable serial programming by sending the Programming Enable serial instruction to the MOSI (PB5) pin.
3. The serial programming instructions will not work if the communication is out of synchronization. When in sync, the second byte (\$53) will echo back when issu-

- Notes:
1. "Max" means the highest value where the pin is guaranteed to be read as low.
 2. "Min" means the lowest value where the pin is guaranteed to be read as high.
 3. Although each I/O port can sink more than the test conditions (20 mA at $V_{CC} = 5V$, 10 mA at $V_{CC} = 3V$) under steady state conditions (non-transient), the following must be observed:
 - 1] The sum of all I_{OL} , for all ports, should not exceed 200 mA.
 - 2] The sum of all I_{OL} , for ports B0 - B7, D0 - D7 and XTAL2, should not exceed 100 mA.
 - 3] The sum of all I_{OL} , for ports A0 - A7, ALE, OC1B and C0 - C7 should not exceed 100 mA.
 If I_{OL} exceeds the test condition, V_{OL} may exceed the related specification. Pins are not guaranteed to sink current greater than the listed test condition.
 4. Although each I/O port can source more than the test conditions (3 mA at $V_{CC} = 5V$, 1.5 mA at $V_{CC} = 3V$) under steady state conditions (non-transient), the following must be observed:
 - 1] The sum of all I_{OH} , for all ports, should not exceed 200 mA.
 - 2] The sum of all I_{OH} , for ports B0 - B7, D0 - D7 and XTAL2, should not exceed 100 mA.
 - 3] The sum of all I_{OH} , for ports A0 - A7, ALE, OC1B and C0 - C7 should not exceed 100 mA.
 If I_{OH} exceeds the test condition, V_{OH} may exceed the related specification. Pins are not guaranteed to source current greater than the listed test condition.
 5. Minimum V_{CC} for power-down is 2V.

External Clock Drive Waveforms

Figure 67. External Clock

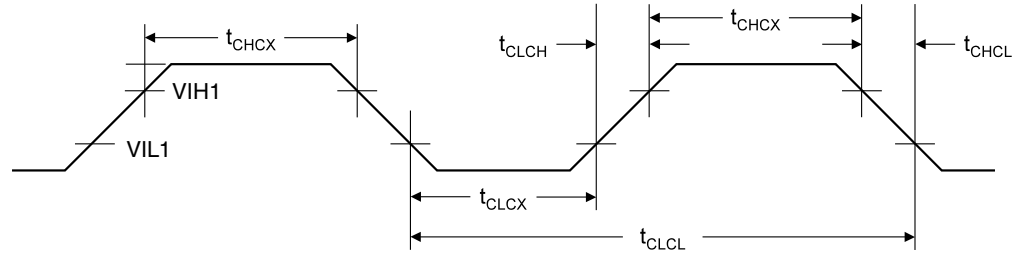
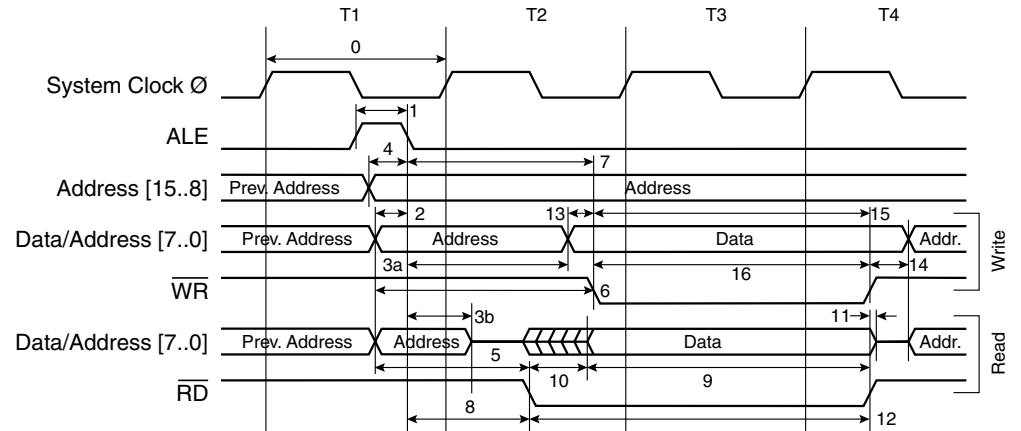


Table 36. External Clock Drive

Symbol	Parameter	V _{CC} = 2.7V to 4.0V		V _{CC} = 4.0V to 6.0V		Units
		Min	Max	Min	Max	
1/t _{CLCL}	Oscillator Frequency	0	4.0	0	8.0	MHz
t _{CLCL}	Clock Period	250.0		125.0		ns
t _{CHCX}	High Time	100.0		50.0		ns
t _{CLCX}	Low Time	100.0		50.0		ns
t _{CLCH}	Rise Time		1.6		0.5	μs
t _{CHCL}	Fall Time		1.6		0.5	μs

Note: See “External Data Memory Timing” for a description of how the duty cycle influences the timing for the external data memory.

Figure 68. External RAM Timing



Note: Clock cycle T3 is only present when external SRAM wait state is enabled.

Figure 82. I/O Pin Sink Current vs. Output Voltage

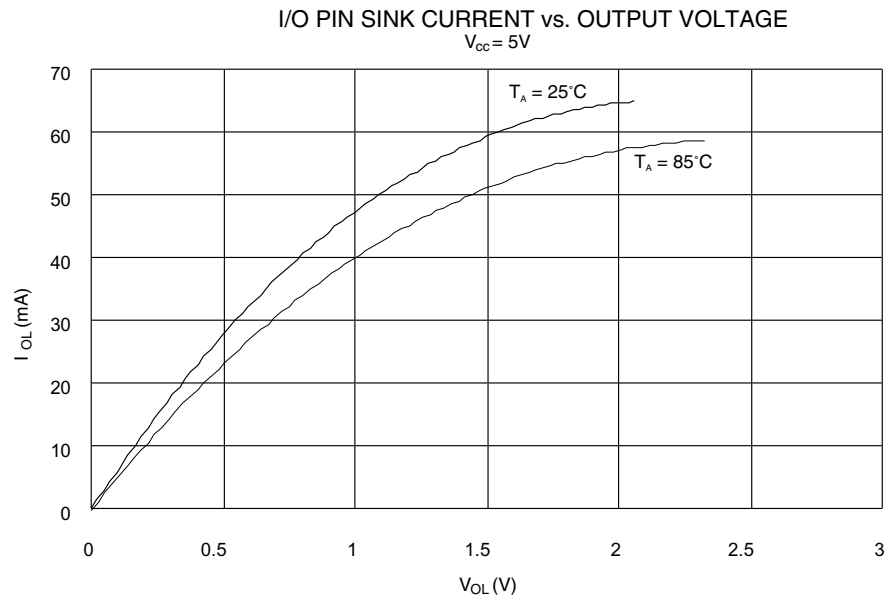
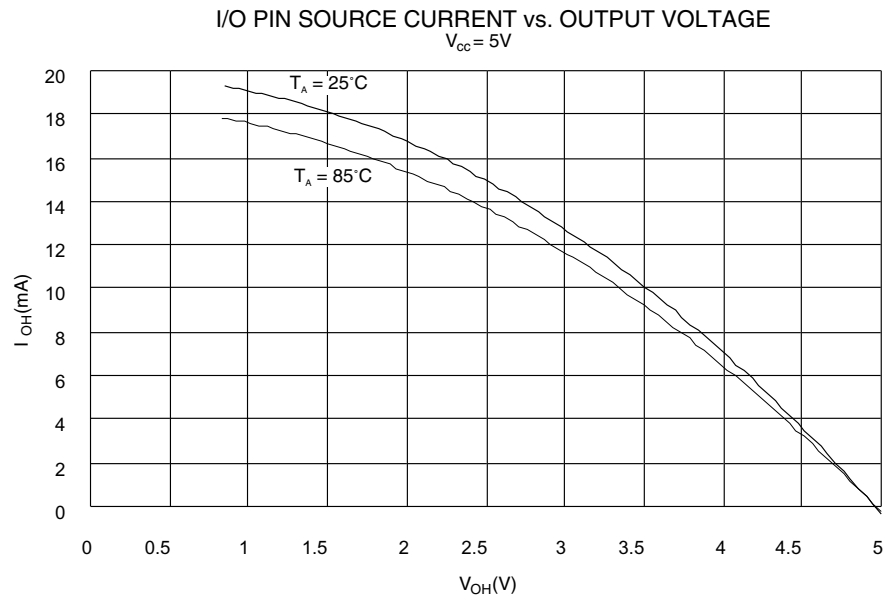


Figure 83. I/O Pin Source Current vs. Output Voltage



Instruction Set Summary (Continued)

Mnemonic	Operands	Description	Operation	Flags	# Clocks
DATA TRANSFER INSTRUCTIONS					
MOV	Rd, Rr	Move between Registers	$Rd \leftarrow Rr$	None	1
LDI	Rd, K	Load Immediate	$Rd \leftarrow K$	None	1
LD	Rd, X	Load Indirect	$Rd \leftarrow (X)$	None	2
LD	Rd, X+	Load Indirect and Post-inc.	$Rd \leftarrow (X), X \leftarrow X + 1$	None	2
LD	Rd, -X	Load Indirect and Pre-dec.	$X \leftarrow X - 1, Rd \leftarrow (X)$	None	2
LD	Rd, Y	Load Indirect	$Rd \leftarrow (Y)$	None	2
LD	Rd, Y+	Load Indirect and Post-inc.	$Rd \leftarrow (Y), Y \leftarrow Y + 1$	None	2
LD	Rd, -Y	Load Indirect and Pre-dec.	$Y \leftarrow Y - 1, Rd \leftarrow (Y)$	None	2
LDD	Rd, Y+q	Load Indirect with Displacement	$Rd \leftarrow (Y + q)$	None	2
LD	Rd, Z	Load Indirect	$Rd \leftarrow (Z)$	None	2
LD	Rd, Z+	Load Indirect and Post-inc.	$Rd \leftarrow (Z), Z \leftarrow Z + 1$	None	2
LD	Rd, -Z	Load Indirect and Pre-dec.	$Z \leftarrow Z - 1, Rd \leftarrow (Z)$	None	2
LDD	Rd, Z+q	Load Indirect with Displacement	$Rd \leftarrow (Z + q)$	None	2
LDS	Rd, k	Load Direct from SRAM	$Rd \leftarrow (k)$	None	2
ST	X, Rr	Store Indirect	$(X) \leftarrow Rr$	None	2
ST	X+, Rr	Store Indirect and Post-inc.	$(X) \leftarrow Rr, X \leftarrow X + 1$	None	2
ST	-X, Rr	Store Indirect and Pre-dec.	$X \leftarrow X - 1, (X) \leftarrow Rr$	None	2
ST	Y, Rr	Store Indirect	$(Y) \leftarrow Rr$	None	2
ST	Y+, Rr	Store Indirect and Post-inc.	$(Y) \leftarrow Rr, Y \leftarrow Y + 1$	None	2
ST	-Y, Rr	Store Indirect and Pre-dec.	$Y \leftarrow Y - 1, (Y) \leftarrow Rr$	None	2
STD	Y+q, Rr	Store Indirect with Displacement	$(Y + q) \leftarrow Rr$	None	2
ST	Z, Rr	Store Indirect	$(Z) \leftarrow Rr$	None	2
ST	Z+, Rr	Store Indirect and Post-inc.	$(Z) \leftarrow Rr, Z \leftarrow Z + 1$	None	2
ST	-Z, Rr	Store Indirect and Pre-dec.	$Z \leftarrow Z - 1, (Z) \leftarrow Rr$	None	2
STD	Z+q, Rr	Store Indirect with Displacement	$(Z + q) \leftarrow Rr$	None	2
STS	k, Rr	Store Direct to SRAM	$(k) \leftarrow Rr$	None	2
LPM		Load Program Memory	$R0 \leftarrow (Z)$	None	3
IN	Rd, P	In Port	$Rd \leftarrow P$	None	1
OUT	P, Rr	Out Port	$P \leftarrow Rr$	None	1
PUSH	Rr	Push Register on Stack	$STACK \leftarrow Rr$	None	2
POP	Rd	Pop Register from Stack	$Rd \leftarrow STACK$	None	2
BIT AND BIT-TEST INSTRUCTIONS					
SBI	P, b	Set Bit in I/O Register	$I/O(P,b) \leftarrow 1$	None	2
CBI	P, b	Clear Bit in I/O Register	$I/O(P,b) \leftarrow 0$	None	2
LSL	Rd	Logical Shift Left	$Rd(n+1) \leftarrow Rd(n), Rd(0) \leftarrow 0$	Z,C,N,V	1
LSR	Rd	Logical Shift Right	$Rd(n) \leftarrow Rd(n+1), Rd(7) \leftarrow 0$	Z,C,N,V	1
ROL	Rd	Rotate Left through Carry	$Rd(0) \leftarrow C, Rd(n+1) \leftarrow Rd(n), C \leftarrow Rd(7)$	Z,C,N,V	1
ROR	Rd	Rotate Right through Carry	$Rd(7) \leftarrow C, Rd(n) \leftarrow Rd(n+1), C \leftarrow Rd(0)$	Z,C,N,V	1
ASR	Rd	Arithmetic Shift Right	$Rd(n) \leftarrow Rd(n+1), n = 0..6$	Z,C,N,V	1
SWAP	Rd	Swap Nibbles	$Rd(3..0) \leftarrow Rd(7..4), Rd(7..4) \leftarrow Rd(3..0)$	None	1
BSET	s	Flag Set	$SREG(s) \leftarrow 1$	SREG(s)	1
BCLR	s	Flag Clear	$SREG(s) \leftarrow 0$	SREG(s)	1
BST	Rr, b	Bit Store from Register to T	$T \leftarrow Rr(b)$	T	1
BLD	Rd, b	Bit Load from T to Register	$Rd(b) \leftarrow T$	None	1
SEC		Set Carry	$C \leftarrow 1$	C	1
CLC		Clear Carry	$C \leftarrow 0$	C	1
SEN		Set Negative Flag	$N \leftarrow 1$	N	1
CLN		Clear Negative Flag	$N \leftarrow 0$	N	1
SEZ		Set Zero Flag	$Z \leftarrow 1$	Z	1
CLZ		Clear Zero Flag	$Z \leftarrow 0$	Z	1
SEI		Global Interrupt Enable	$I \leftarrow 1$	I	1
CLI		Global Interrupt Disable	$I \leftarrow 0$	I	1
SES		Set Signed Test Flag	$S \leftarrow 1$	S	1
CLS		Clear Signed Test Flag	$S \leftarrow 0$	S	1
SEV		Set Two's Complement Overflow	$V \leftarrow 1$	V	1
CLV		Clear Two's Complement Overflow	$V \leftarrow 0$	V	1
SET		Set T in SREG	$T \leftarrow 1$	T	1
CLT		Clear T in SREG	$T \leftarrow 0$	T	1
SEH		Set Half-carry Flag in SREG	$H \leftarrow 1$	H	1
CLH		Clear Half-carry Flag in SREG	$H \leftarrow 0$	H	1
NOP		No Operation		None	1
SLEEP		Sleep	(see specific descr. for Sleep function)	None	1
WDR		Watchdog Reset	(see specific descr. for WDR/timer)	None	1