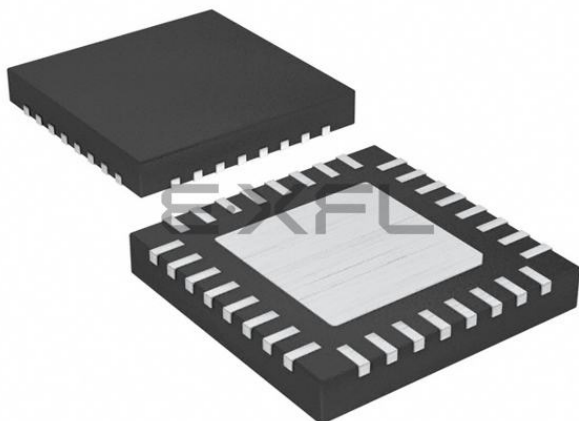




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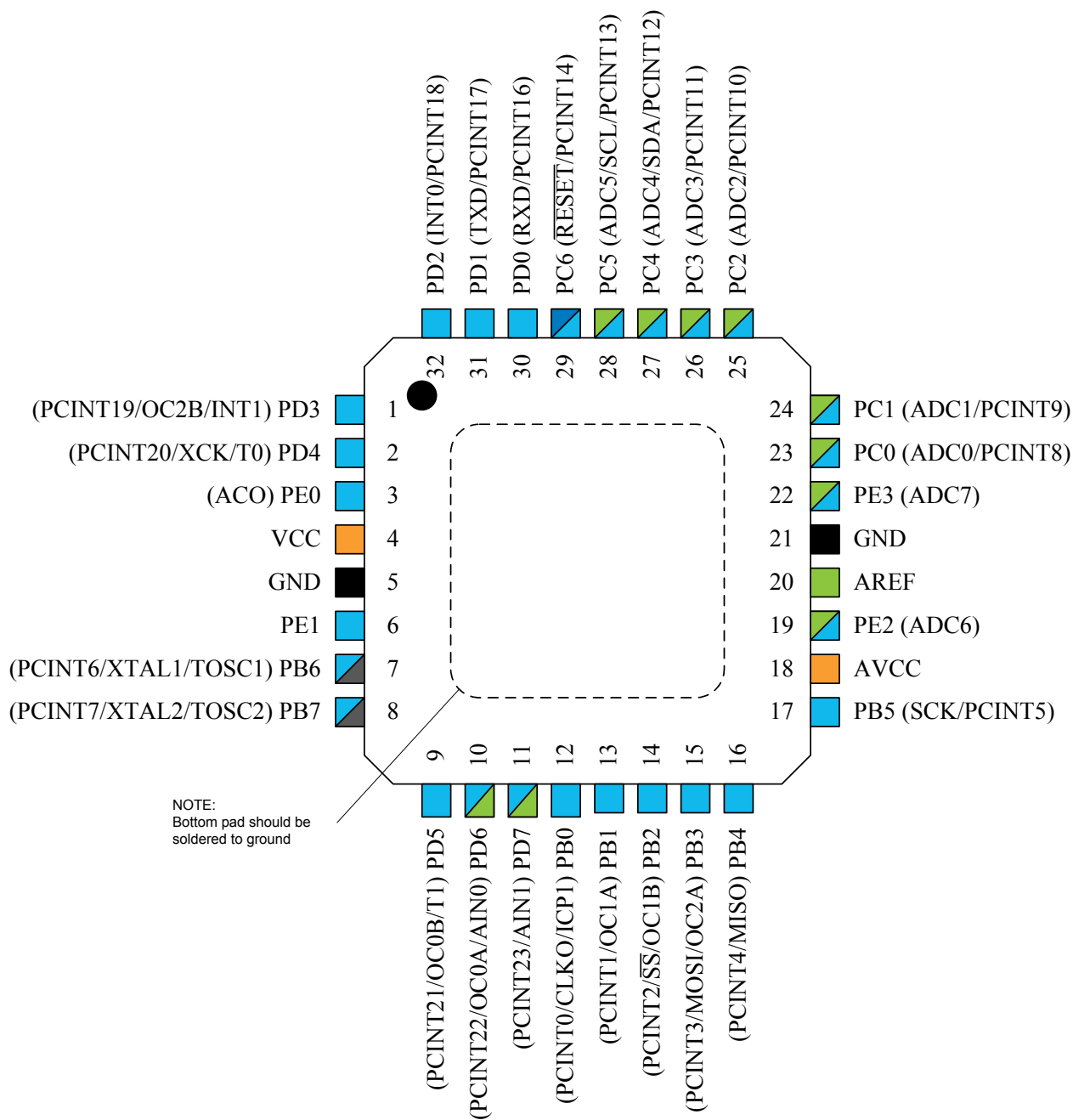
"[Embedded - Microcontrollers](#)" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

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Details

Product Status	Active
Core Processor	AVR
Core Size	8-Bit
Speed	20MHz
Connectivity	I ² C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	27
Program Memory Size	16KB (8K x 16)
Program Memory Type	FLASH
EEPROM Size	512 x 8
RAM Size	1K x 8
Voltage - Supply (Vcc/Vdd)	1.8V ~ 5.5V
Data Converters	A/D 8x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 105°C (TA)
Mounting Type	Surface Mount
Package / Case	32-VFQFN Exposed Pad
Supplier Device Package	32-VFQFN (5x5)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/atmega168pb-mn

Figure 5-2. 32 VQFN Pinout ATmega48PB/88PB/168PB



5.1. Pin Descriptions

5.1.1. VCC

Digital supply voltage.

5.1.2. GND

Ground.

Bit 0 – C: Carry Flag

The Carry Flag C indicates a carry in an arithmetic or logic operation. See the *Instruction Set Description* for detailed information.

12.4. General Purpose Register File

The Register File is optimized for the AVR Enhanced RISC instruction set. In order to achieve the required performance and flexibility, the following input/output schemes are supported by the Register File:

- One 8-bit output operand and one 8-bit result input
- Two 8-bit output operands and one 8-bit result input
- Two 8-bit output operands and one 16-bit result input
- One 16-bit output operand and one 16-bit result input

The figure shows the structure of the 32 general purpose working registers in the CPU.

Figure 12-2. AVR CPU General Purpose Working Registers

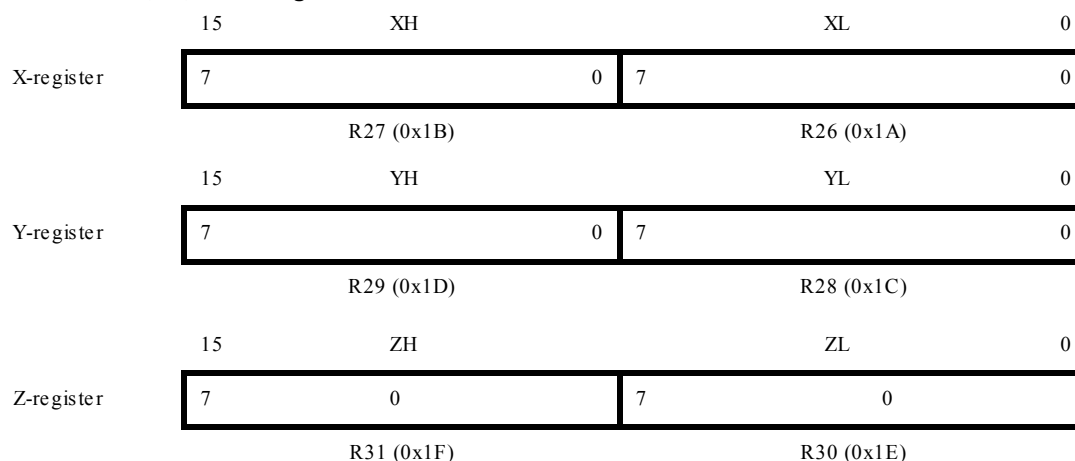
	7	0	Addr.	
General Purpose Working Registers	R0		0x00	
	R1		0x01	
	R2		0x02	
	...			
	R13		0x0D	
	R14		0x0E	
	R15		0x0F	
	R16		0x10	
	R17		0x11	
	...			
	R26		0x1A	X-register Low Byte
	R27		0x1B	X-register High Byte
	R28		0x1C	Y-register Low Byte
	R29		0x1D	Y-register High Byte
	R30		0x1E	Z-register Low Byte
	R31		0x1F	Z-register High Byte

Most of the instructions operating on the Register File have direct access to all registers, and most of them are single cycle instructions. As shown in the figure, each register is also assigned a data memory address, mapping them directly into the first 32 locations of the user Data Space. Although not being physically implemented as SRAM locations, this memory organization provides great flexibility in access of the registers, as the X-, Y-, and Z-pointer registers can be set to index any register in the file.

12.4.1. The X-register, Y-register, and Z-register

The registers R26...R31 have some added functions to their general purpose usage. These registers are 16-bit address pointers for indirect addressing of the data space. The three indirect address registers X, Y, and Z are defined as described in the figure.

Figure 12-3. The X-, Y-, and Z-registers



In the different addressing modes these address registers have functions as fixed displacement, automatic increment, and automatic decrement. See *Instruction Set Summary* for details.

12.5. Stack Pointer

The Stack is mainly used for storing temporary data, for storing local variables and for storing return addresses after interrupts and subroutine calls. The Stack is implemented as growing from higher to lower memory locations. The Stack Pointer Register always points to the top of the Stack.

The Stack Pointer points to the data SRAM Stack area where the Subroutine and Interrupt Stacks are located. A Stack PUSH command will decrease the Stack Pointer. The Stack in the data SRAM must be defined by the program before any subroutine calls are executed or interrupts are enabled. Initial Stack Pointer value equals the last address of the internal SRAM and the Stack Pointer must be set to point above start of the SRAM. See the table for Stack Pointer details.

Table 12-1. Stack Pointer Instructions

Instruction	Stack pointer	Description
PUSH	Decrement by 1	Data is pushed onto the stack
ICALL RCALL	Decrement by 2	Return address is pushed onto the stack with a subroutine call or interrupt
POP	Increment by 1	Data is popped from the stack
RET RETI	Increment by 2	Return address is popped from the stack with return from subroutine or return from interrupt

The AVR Stack Pointer is implemented as two 8-bit registers in the I/O space. The number of bits actually used is implementation dependent. Note that the data space in some implementations of the AVR architecture is so small that only SPL is needed. In this case, the SPH Register will not be present.

12.5.2. Stack Pointer Low Register

When addressing I/O Registers as data space using LD and ST instructions, the provided offset must be used. When using the I/O specific commands IN and OUT, the offset is reduced by 0x20, resulting in an I/O address offset within 0x00 - 0x3F.

Reset value of SPL is RAMEND.

Name: SPL

Offset: 0x5D

Reset: 0xFF

Property: When addressing as I/O Register: address offset is 0x3D

Bit	7	6	5	4	3	2	1	0
	SP7	SP6	SP5	SP4	SP3	SP2	SP1	SP0
Access	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
Reset	x	x	x	x	x	x	x	x

Bit 7 – SP7: Stack Pointer Address 7

Bit 6 – SP6: Stack Pointer Address 6

Bit 5 – SP5: Stack Pointer Address 5

Bit 4 – SP4: Stack Pointer Address 4

Bit 3 – SP3: Stack Pointer Address 3

Bit 2 – SP2: Stack Pointer Address 2

Bit 1 – SP1: Stack Pointer Address 1

Bit 0 – SP0: Stack Pointer Address 0

12.6. Instruction Execution Timing

This section describes the general access timing concepts for instruction execution. The AVR CPU is driven by the CPU clock clk_{CPU} , directly generated from the selected clock source for the chip. No internal clock division is used. The Figure below shows the parallel instruction fetches and instruction executions enabled by the Harvard architecture and the fast-access Register File concept. This is the basic pipelining concept to obtain up to 1 MIPS per MHz with the corresponding unique results for functions per cost, functions per clocks, and functions per power-unit.

12.7.1. Interrupt Response Time

The interrupt execution response for all the enabled AVR interrupts is four clock cycles minimum. After four clock cycles the program vector address for the actual interrupt handling routine is executed. During this four clock cycle period, the Program Counter is pushed onto the Stack. The vector is normally a jump to the interrupt routine, and this jump takes three clock cycles. If an interrupt occurs during execution of a multi-cycle instruction, this instruction is completed before the interrupt is served. If an interrupt occurs when the MCU is in sleep mode, the interrupt execution response time is increased by four clock cycles. This increase comes in addition to the start-up time from the selected sleep mode. A return from an interrupt handling routine takes four clock cycles. During these four clock cycles, the Program Counter (two bytes) is popped back from the Stack, the Stack Pointer is incremented by two, and the I-bit in SREG is set.

CLKPS[3:0]	Clock Division Factor
1010	Reserved
1011	Reserved
1100	Reserved
1101	Reserved
1110	Reserved
1111	Reserved

VectorNo.	Program Address ⁽²⁾	Source	Interrupt Definition
22	0x002A	ADC	ADC Conversion Complete
23	0x002C	EE READY	EEPROM Ready
24	0x002E	ANALOG COMP	Analog Comparator
25	0x0030	TWI	2-wire Serial Interface (I2C)
26	0x0032	SPM READY	Store Program Memory Ready
27	0x0034	USART, START	USART Start Edge Interrupt

Note:

1. When the BOOTRST Fuse is programmed, the device will jump to the Boot Loader address at reset, please refer to *Boot Loader Support – Read-While-Write Self-Programming* chapter.
2. When the IVSEL bit in MCUCR (MCUCR.IVSEL) is set, Interrupt Vectors will be moved to the start of the Boot Flash Section. The address of each Interrupt Vector will then be the address in this table added to the start address of the Boot Flash Section.

The following table shows reset and Interrupt Vectors placement for the various combinations of BOOTRST and IVSEL settings. If the program never enables an interrupt source, the Interrupt Vectors are not used, and regular program code can be placed at these locations. This is also the case if the Reset Vector is in the Application section while the Interrupt Vectors are in the Boot section or vice versa.

Table 17-5. Reset and Interrupt Vectors Placement in ATmega168PB

BOOTRST	IVSEL	Reset Address	Interrupt Vectors Start Address
1	0	0x000	0x002
1	1	0x000	Boot Reset Address + 0x0002
0	0	Boot Reset Address	0x002
0	1	Boot Reset Address	Boot Reset Address + 0x0002

Note: The Boot Reset Address is shown in Table. Boot Size Configuration, ATmega168PB in *ATmega168PB Boot Loader Parameters*. For the BOOTRST Fuse “1” means unprogrammed while “0” means programmed.

The most typical and general program setup for the Reset and Interrupt Vector Addresses in ATmega168PB is:

Address	Labels	Code Comments
0x0000	jmp RESET	; Reset Handler
0x0002	jmp EXT_INT0	; IRQ0 Handler
0x0004	jmp EXT_INT1	; IRQ1 Handler
0x0006	jmp PCINT0	; PCINT0 Handler
0x0008	jmp PCINT1	; PCINT1 Handler
0x000A	jmp PCINT2	; PCINT2 Handler
0x000C	jmp WDT	; Watchdog Timer Handler

Table 19-4. Overriding Signals for Alternate Functions in PB7...PB4

Signal Name	PB7/XTAL2/TOSC2/ PCINT7 ⁽¹⁾	PB6/XTAL1/TOSC1/ PCINT6 ⁽¹⁾	PB5/SCK0/XCK0/ PCINT5	PB4/MISO0/RXD1/ PCINT4
PUOE	$\overline{\text{INTRC}} \cdot \overline{\text{EXTCK}} + \text{AS2}$	$\text{INTRC} + \text{AS2}$	$\text{SPE0} \cdot \overline{\text{MSTR}}$	$\text{SPE0} \cdot \text{MSTR} + \text{RXEN1}$
PUOV	0	0	$\text{PORTB5} \cdot \overline{\text{PUD}}$	$\text{PORTB4} \cdot \overline{\text{PUD}}$
DDOE	$\overline{\text{INTRC}} \cdot \overline{\text{EXTCK}} + \text{AS2}$	$\text{INTRC} + \text{AS2}$	$\text{SPE0} \cdot \overline{\text{MSTR}}$	$\text{SPE0} \cdot \text{MSTR} + \text{RXEN1}$
DDOV	0	0	0	0
PVOE	0	0	$\text{SPE0} \cdot \text{MSTR}$	$\text{SPE0} \cdot \overline{\text{MSTR}}$
PVOV	0	0	SCK0 OUTPUT	SPI0 SLAVE OUTPUT
DIEOE	$\overline{\text{INTRC}} \cdot \overline{\text{EXTCK}} + \text{AS2} + \text{PCINT7} \cdot \text{PCIE0}$	$\overline{\text{INTRC}} + \text{AS2} + \text{PCINT6} \cdot \text{PCIE0}$	$\text{PCINT5} \cdot \text{PCIE0}$	$\text{PCINT4} \cdot \text{PCIE0}$
DIEOV	$(\text{INTRC} + \text{EXTCK}) \cdot \overline{\text{AS2}}$	$\text{INTRC} \cdot \overline{\text{AS2}}$	1	1
DI	PCINT7 INPUT	PCINT6 INPUT	PCINT5 INPUT SCK0 INPUT	PCINT4 INPUT SPI0 MSTR INPUT RXD1
AIO	Oscillator Output	Oscillator/Clock Input	—	—

Notes: 1. INTRC means that one of the internal RC Oscillators are selected (by the CKSEL fuses), EXTCK means that external clock is selected (by the CKSEL fuses).

Table 19-5. Overriding Signals for Alternate Functions in PB3...PB0

Signal Name	PB3/MOSI0/TXD1/OC2A/PCINT3	PB2/SS0/OC1B/PCINT2	PB1/OC1A/PCINT1	PB0/ICP1/CLKO/ PCINT0
PUOE	$\text{SPE0} \cdot \overline{\text{MSTR}} + \text{TXEN1}$	$\text{SPE0} \cdot \overline{\text{MSTR}}$	0	0
PUOV	$\text{PORTB3} \cdot \overline{\text{PUD}}$	$\text{PORTB2} \cdot \overline{\text{PUD}}$	0	0
DDOE	$\text{SPE0} \cdot \overline{\text{MSTR}} + \text{TXEN1}$	$\text{SPE0} \cdot \overline{\text{MSTR}}$	0	0
DDOV	0	0	0	0
PVOE	$\text{SPE0} \cdot \text{MSTR} + \text{OC2A ENABLE}$	OC1B ENABLE	OC1A ENABLE	0
PVOV	SPI0 MSTR OUTPUT + OC2A + TXD1	OC1B	OC1A	0
DIEOE	$\text{PCINT3} \cdot \text{PCIE0}$	$\text{PCINT2} \cdot \text{PCIE0}$	$\text{PCINT1} \cdot \text{PCIE0}$	$\text{PCINT0} \cdot \text{PCIE0}$
DIEOV	1	1	1	1
DI	PCINT3 INPUT SPI0 SLAVE INPUT	PCINT2 INPUT SPI0 SS	PCINT1 INPUT	PCINT0 INPUT ICP1 INPUT
AIO	—	—	—	—

19.4.9. Port D Data Direction Register

When addressing I/O Registers as data space using LD and ST instructions, the provided offset must be used. When using the I/O specific commands IN and OUT, the offset is reduced by 0x20, resulting in an I/O address offset within 0x00 - 0x3F.

Name: DDRD

Offset: 0x2A

Reset: 0x00

Property: When addressing as I/O Register: address offset is 0x0A

Bit	7	6	5	4	3	2	1	0
	DDRD7	DDRD6	DDRD5	DDRD4	DDRD3	DDRD2	DDRD1	DDRD0
Access	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
Reset	0	0	0	0	0	0	0	0

Bits 7:0 – DDRDn: Port D Data Direction [n = 7:0]

Table 20-1. Definitions

Constant	Description
BOTTOM	The counter reaches the BOTTOM when it becomes zero (0x00 for 8-bit counters, or 0x0000 for 16-bit counters).
MAX	The counter reaches its Maximum when it becomes 0xFF (decimal 255, for 8-bit counters) or 0xFFFF (decimal 65535, for 16-bit counters).
TOP	The counter reaches the TOP when it becomes equal to the highest value in the count sequence. The TOP value can be assigned to be the fixed value MAX or the value stored in the OCR1A Register. The assignment is dependent on the mode of operation.

20.2.2. Registers

The Timer/Counter 0 register (TCNT0) and Output Compare TC0x registers (OCR0x) are 8-bit registers. Interrupt request (abbreviated to Int.Req. in the block diagram) signals are all visible in the Timer Interrupt Flag Register 0 (TIFR0). All interrupts are individually masked with the Timer Interrupt Mask Register 0 (TIMSK0). TIFR0 and TIMSK0 are not shown in the figure.

The TC can be clocked internally, via the prescaler, or by an external clock source on the T0 pin. The Clock Select logic block controls which clock source and edge is used by the Timer/Counter to increment (or decrement) its value. The TC is inactive when no clock source is selected. The output from the Clock Select logic is referred to as the timer clock (clk_{T0}).

The double buffered Output Compare Registers (OCR0A and OCR0B) are compared with the Timer/Counter value at all times. The result of the compare can be used by the Waveform Generator to generate a PWM or variable frequency output on the Output Compare pins (OC0A and OC0B). See [Output Compare Unit](#) for details. The compare match event will also set the Compare Flag (OCF0A or OCF0B) which can be used to generate an Output Compare interrupt request.

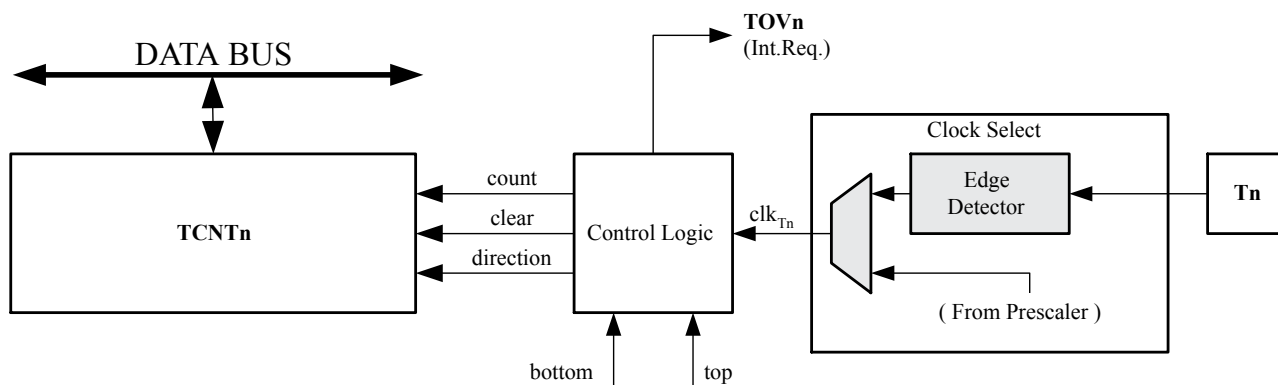
20.3. Timer/Counter Clock Sources

The TC can be clocked by an internal or an external clock source. The clock source is selected by writing to the Clock Select (CS0[2:0]) bits in the Timer/Counter Control Register (TCCR0B).

20.4. Counter Unit

The main part of the 8-bit Timer/Counter is the programmable bi-directional counter unit. Below is the block diagram of the counter and its surroundings.

Figure 20-2. Counter Unit Block Diagram



2. BOTTOM = 0x00

Table 20-10. Clock Select Bit Description

CA02	CA01	CS00	Description
0	0	0	No clock source (Timer/Counter stopped).
0	0	1	clk _{I/O} /1 (No prescaling)
0	1	0	clk _{I/O} /8 (From prescaler)
0	1	1	clk _{I/O} /64 (From prescaler)
1	0	0	clk _{I/O} /256 (From prescaler)
1	0	1	clk _{I/O} /1024 (From prescaler)
1	1	0	External clock source on T0 pin. Clock on falling edge.
1	1	1	External clock source on T0 pin. Clock on rising edge.

If external pin modes are used for the Timer/Counter0, transitions on the T0 pin will clock the counter even if the pin is configured as an output. This feature allows software control of the counting.

21.12.1. TC1 Control Register A

Name: TCCR1A

Offset: 0x80

Reset: 0x00

Property: -

Bit	7	6	5	4	3	2	1	0
	COM1A1	COM1A0	COM1B1	COM1B0			WGM11	WGM10
Access	R/W	R/W	R/W	R/W			R/W	R/W
Reset	0	0	0	0			0	0

Bits 7:6 – COM1An: Compare Output Mode for Channel A [n = 1:0]

Bits 5:4 – COM1Bn: Compare Output Mode for Channel B [n = 1:0]

The COM1A[1:0] and COM1B[1:0] control the Output Compare pins (OC1A and OC1B respectively) behavior. If one or both of the COM1A[1:0] bits are written to one, the OC1A output overrides the normal port functionality of the I/O pin it is connected to. If one or both of the COM1B[1:0] bit are written to one, the OC1B output overrides the normal port functionality of the I/O pin it is connected to. However, note that the Data Direction Register (DDR) bit corresponding to the OC1A or OC1B pin must be set in order to enable the output driver.

When the OC1A or OC1B is connected to the pin, the function of the COM1x[1:0] bits is dependent of the WGM1[3:0] bits setting. The table below shows the COM1x[1:0] bit functionality when the WGM1[3:0] bits are set to a Normal or a CTC mode (non-PWM).

Table 21-3. Compare Output Mode, non-PWM

COM1A1/COM1B1	COM1A0/COM1B0	Description
0	0	Normal port operation, OC1A/OC1B disconnected.
0	1	Toggle OC1A/OC1B on Compare Match.
1	0	Clear OC1A/OC1B on Compare Match (Set output to low level).
1	1	Set OC1A/OC1B on Compare Match (Set output to high level).

The table below shows the COM1x[1:0] bit functionality when the WGM1[3:0] bits are set to the fast PWM mode.

Table 21-4. Compare Output Mode, Fast PWM

COM1A1/COM1B1	COM1A0/COM1B0	Description
0	0	Normal port operation, OC1A/OC1B disconnected.
0	1	WGM1[3:0] = 14 or 15: Toggle OC1A on Compare Match, OC1B disconnected (normal port operation). For all other WGM1 settings, normal port operation, OC1A/OC1B disconnected.

24.5.1. SPI Control Register 0

When addressing I/O Registers as data space using LD and ST instructions, the provided offset must be used. When using the I/O specific commands IN and OUT, the offset is reduced by 0x20, resulting in an I/O address offset within 0x00 - 0x3F.

Name: SPCR

Offset: 0x4C

Reset: 0x00

Property: When addressing as I/O Register: address offset is 0x2C

Bit	7	6	5	4	3	2	1	0
	SPIE	SPE	DORD	MSTR	CPOL	CPHA	SPR1	SPR0
Access	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
Reset	0	0	0	0	0	0	0	0

Bit 7 – SPIE: SPI0 Interrupt Enable

This bit causes the SPI interrupt to be executed if SPIF bit in the SPSR Register is set and if the Global Interrupt Enable bit in SREG is set.

Bit 6 – SPE: SPI0 Enable

When the SPE bit is written to one, the SPI is enabled. This bit must be set to enable any SPI operations.

Bit 5 – DORD: Data0 Order

When the DORD bit is written to one, the LSB of the data word is transmitted first.

When the DORD bit is written to zero, the MSB of the data word is transmitted first.

Bit 4 – MSTR: Master/Slave0 Select

This bit selects Master SPI mode when written to one, and Slave SPI mode when written logic zero. If SS is configured as an input and is driven low while MSTR is set, MSTR will be cleared, and SPIF in SPSR will become set. The user will then have to set MSTR to re-enable SPI Master mode.

Bit 3 – CPOL: Clock0 Polarity

When this bit is written to one, SCK is high when idle. When CPOL is written to zero, SCK is low when idle. Refer to [Figure 24-3 SPI Transfer Format with CPHA = 0](#) and [Figure 24-4 SPI Transfer Format with CPHA = 1](#) for an example. The CPOL functionality is summarized below:

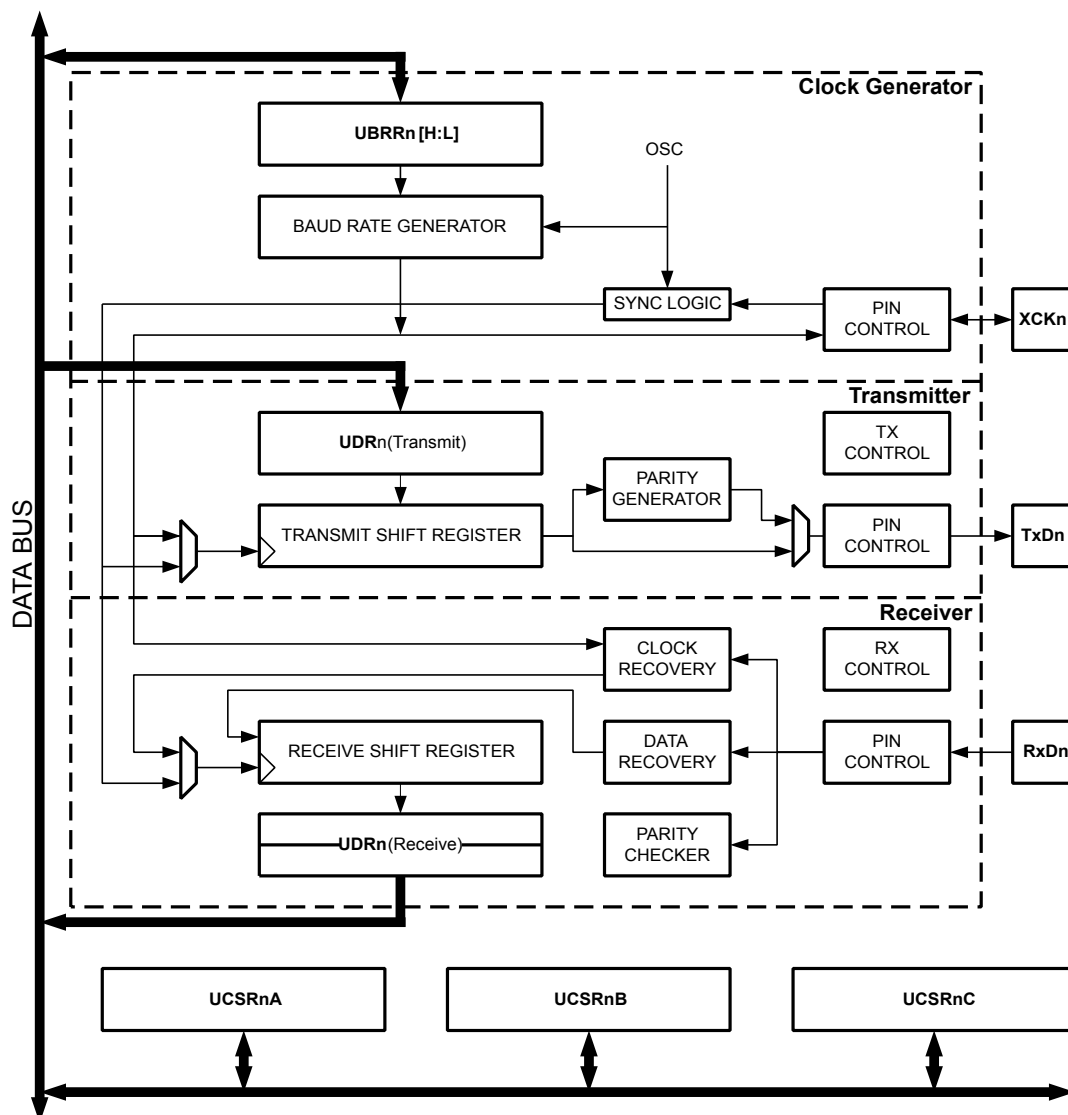
Table 24-3. CPOL0 Functionality

CPOL	Leading Edge	Trailing Edge
0	Rising	Falling
1	Falling	Rising

Bit 2 – CPHA: Clock0 Phase

The settings of the Clock Phase bit (CPHA) determine if data is sampled on the leading (first) or trailing (last) edge of SCK. Refer to [Figure 24-3 SPI Transfer Format with CPHA = 0](#) and [Figure 24-4 SPI Transfer Format with CPHA = 1](#) for an example. The CPHA functionality is summarized below:

Figure 25-1. USART Block Diagram



Note: Refer to the *Pin Configurations* and the *I/O-Ports* description for USART pin placement.

Related Links

[Pin Configurations](#) on page 14

[I/O-Ports](#) on page 105

[USART in SPI Mode](#) on page 272

25.4. Clock Generation

The Clock Generation logic generates the base clock for the Transmitter and Receiver. The USART supports four modes of clock operation: Normal asynchronous, Double Speed asynchronous, Master synchronous and Slave synchronous mode. The USART Mode Select bit 0 in the USART Control and Status Register n C (UCSRnC.UMSELn0) selects between asynchronous and synchronous operation. Double Speed (asynchronous mode only) is controlled by the U2Xn found in the UCSRnA Register. When using synchronous mode (UMSELn0=1), the Data Direction Register for the XCKn pin (DDR_XCKn) controls whether the clock source is internal (Master mode) or external (Slave mode). The XCKn pin is only active when using synchronous mode.

27. TWI - 2-wire Serial Interface

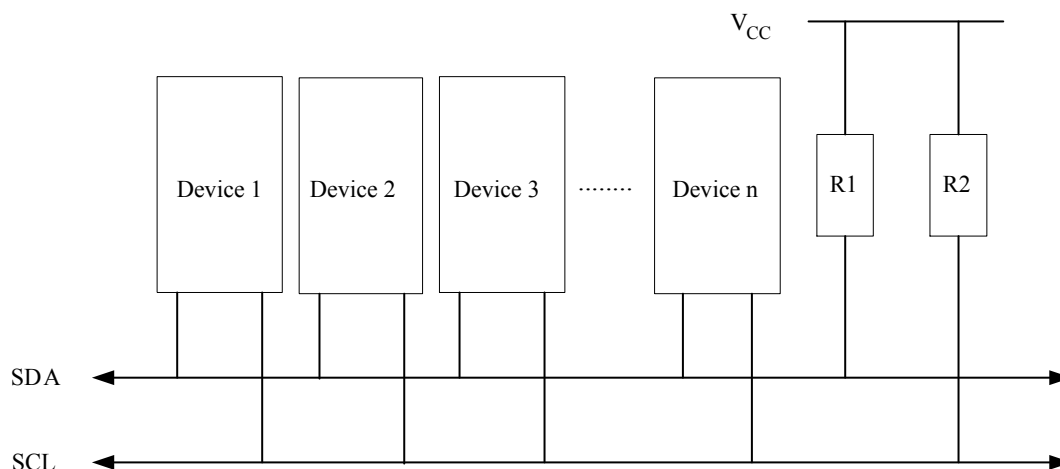
27.1. Features

- Simple, yet Powerful and Flexible Communication Interface, only two Bus Lines Needed
- Both Master and Slave Operation Supported
- Device can Operate as Transmitter or Receiver
- 7-bit Address Space Allows up to 128 Different Slave Addresses
- Multi-master Arbitration Support
- Up to 400kHz Data Transfer Speed
- Slew-rate Limited Output Drivers
- Noise Suppression Circuitry Rejects Spikes on Bus Lines
- Fully Programmable Slave Address with General Call Support
- Address Recognition Causes Wake-up When AVR is in Sleep Mode
- Compatible with Philips' I²C protocol

27.2. Two-Wire Serial Interface Bus Definition

The Two-Wire Serial Interface (TWI) is ideally suited for typical microcontroller applications. The TWI protocol allows the systems designer to interconnect up to 128 different devices using only two bi-directional bus lines: one for clock (SCL) and one for data (SDA). The only external hardware needed to implement the bus is a single pull-up resistor for each of the TWI bus lines. All devices connected to the bus have individual addresses, and mechanisms for resolving bus contention are inherent in the TWI protocol.

Figure 27-1. TWI Bus Interconnection



27.2.1. TWI Terminology

The following definitions are frequently encountered in this section.

28.3.2. Analog Comparator Control and Status Register C

The Store Program Memory Control and Status Register contains the control bits needed to control the Boot Loader operations.

When addressing I/O Registers as data space using LD and ST instructions, the provided offset must be used. When using the I/O specific commands IN and OUT, the offset is reduced by 0x20, resulting in an I/O address offset within 0x00 - 0x3F.

Name: ACSR0

Offset: 0x4F

Reset: 0x00

Property: When addressing as I/O Register: address offset is 0x2F

Bit	7	6	5	4	3	2	1	0
								ACOE
Access								R/W
Reset								0

Bit 0 – ACOE: Analog Comparator Output Enable

When this bit is set, the analog comparator output is connected to the ACO pin.

28.3.3. Analog Comparator Control and Status Register

When addressing I/O Registers as data space using LD and ST instructions, the provided offset must be used. When using the I/O specific commands IN and OUT, the offset is reduced by 0x20, resulting in an I/O address offset within 0x00 - 0x3F.

Name: ACSR

Offset: 0x50

Reset: N/A

Property: When addressing as I/O Register: address offset is 0x30

Bit	7	6	5	4	3	2	1	0
	ACD	ACBG	ACO	ACI	ACIE	ACIC	ACIS1	ACIS0
Access	R/W	R/W	R	R/W	R/W	R/W	R/W	R/W
Reset	0	0	a	0	0	0	0	0

Bit 7 – ACD: Analog Comparator Disable

When this bit is written logic one, the power to the Analog Comparator is switched off. This bit can be set at any time to turn off the Analog Comparator. This will reduce power consumption in Active and Idle mode. When changing the ACD bit, the Analog Comparator Interrupt must be disabled by clearing the ACIE bit in ACSR. Otherwise an interrupt can occur when the bit is changed.

Bit 6 – ACBG: Analog Comparator Bandgap Select

When this bit is set, a fixed bandgap reference voltage replaces the positive input to the Analog Comparator. When this bit is cleared, AIN0 is applied to the positive input of the Analog Comparator. When the bandgap reference is used as input to the Analog Comparator, it will take a certain time for the voltage to stabilize. If not stabilized, the first conversion may give a wrong value.

Bit 5 – ACO: Analog Comparator Output

The output of the Analog Comparator is synchronized and then directly connected to ACO. The synchronization introduces a delay of 1 - 2 clock cycles.

Bit 4 – ACI: Analog Comparator Interrupt Flag

This bit is set by hardware when a comparator output event triggers the interrupt mode defined by ACIS1 and ACIS0. The Analog Comparator interrupt routine is executed if the ACIE bit is set and the I-bit in SREG is set. ACI is cleared by hardware when executing the corresponding interrupt handling vector. Alternatively, ACI is cleared by writing a logic one to the flag.

Bit 3 – ACIE: Analog Comparator Interrupt Enable

When the ACIE bit is written logic one and the I-bit in the Status Register is set, the Analog Comparator interrupt is activated. When written logic zero, the interrupt is disabled.

Bit 2 – ACIC: Analog Comparator Input Capture Enable

When written logic one, this bit enables the input capture function in Timer/Counter1 to be triggered by the Analog Comparator. The comparator output is in this case directly connected to the input capture front-end logic, making the comparator utilize the noise canceler and edge select features of the Timer/Counter1 Input Capture interrupt. When written logic zero, no connection between the Analog Comparator and the input capture function exists. To make the comparator trigger the Timer/Counter1 Input Capture interrupt, the ICIE1 bit in the Timer Interrupt Mask Register (TIMSK1) must be set.

- To protect only the Boot Loader Flash section from a software update by the MCU
- To protect only the Application Flash section from a software update by the MCU
- Allow software update in the entire Flash

The Boot Lock bits can be set in software and in Serial or Parallel Programming mode, but they can be cleared by a Chip Erase command only. The general Write Lock (Lock Bit mode 2) does not control the programming of the Flash memory by SPM instruction. Similarly, the general Read/Write Lock (Lock Bit mode 1) does not control reading nor writing by LPM/SPM, if it is attempted.

Table 32-2. Boot Lock Bit0 Protection Modes (Application Section)

BLB0 Mode	BLB02	BLB01	Protection
1	1	1	No restrictions for SPM or LPM accessing the Application section.
2	1	0	SPM is not allowed to write to the Application section.
3	0	0	SPM is not allowed to write to the Application section, and LPM executing from the Boot Loader section is not allowed to read from the Application section. If Interrupt Vectors are placed in the Boot Loader section, interrupts are disabled while executing from the Application section.
4	0	1	LPM executing from the Boot Loader section is not allowed to read from the Application section. If Interrupt Vectors are placed in the Boot Loader section, interrupts are disabled while executing from the Application section.

Note: “1” means unprogrammed, “0” means programmed.

Table 32-3. Boot Lock Bit1 Protection Modes (Boot Loader Section)

BLB1 Mode	BLB12	BLB11	Protection
1	1	1	No restrictions for SPM or LPM accessing the Boot Loader section.
2	1	0	SPM is not allowed to write to the Boot Loader section.
3	0	0	SPM is not allowed to write to the Boot Loader section, and LPM executing from the Application section is not allowed to read from the Boot Loader section. If Interrupt Vectors are placed in the Application section, interrupts are disabled while executing from the Boot Loader section.
4	0	1	LPM executing from the Application section is not allowed to read from the Boot Loader section. If Interrupt Vectors are placed in the Application section, interrupts are disabled while executing from the Boot Loader section.

Note: “1” means unprogrammed, “0” means programmed.

32.6. Entering the Boot Loader Program

Entering the Boot Loader takes place by a jump or call from the application program. This may be initiated by a trigger such as a command received via USART, or SPI interface. Alternatively, the Boot Reset Fuse can be programmed so that the Reset Vector is pointing to the Boot Flash start address after a reset. In this case, the Boot Loader is started after a reset. After the application code is loaded, the program can

Note: The EEPROM memory is preserved during Chip Erase if the EESAVE Fuse is programmed.

Load Command "Chip Erase":

1. Set XA1, XA0 to "10". This enables command loading.
2. Set BS1 to "0".
3. Set DATA to "1000 0000". This is the command for Chip Erase.
4. Give XTAL1 a positive pulse. This loads the command.
5. Give $\overline{WR}$ a negative pulse. This starts the Chip Erase. RDY/ $\overline{BSY}$ goes low.
6. Wait until RDY/ $\overline{BSY}$ goes high before loading a new command.

33.7.4. Programming the Flash

The Flash is organized in pages as number of Words in a Page and number of Pages in the Flash. When programming the Flash, the program data is latched into a page buffer. This allows one page of program data to be programmed simultaneously. The following procedure describes how to program the entire Flash memory:

Step A. Load Command "Write Flash"

1. Set XA1, XA0 to "10". This enables command loading.
2. Set BS1 to "0".
3. Set DATA to "0001 0000". This is the command for Write Flash.
4. Give XTAL1 a positive pulse. This loads the command.

Step B. Load Address Low Byte

1. Set XA1, XA0 to "00". This enables address loading.
2. Set BS1 to "0". This selects low address.
3. Set DATA = Address low byte (0x00 - 0xFF).
4. Give XTAL1 a positive pulse. This loads the address low byte.

Step C. Load Data Low Byte

1. Set XA1, XA0 to "01". This enables data loading.
2. Set DATA = Data low byte (0x00 - 0xFF).
3. Give XTAL1 a positive pulse. This loads the data byte.

Step D. Load Data High Byte

1. Set BS1 to "1". This selects high data byte.
2. Set XA1, XA0 to "01". This enables data loading.
3. Set DATA = Data high byte (0x00 - 0xFF).
4. Give XTAL1 a positive pulse. This loads the data byte.

Step E. Latch Data

1. Set BS1 to "1". This selects high data byte.
2. Give PAGESL a positive pulse. This latches the data bytes. (Please refer to the figure, Programming the Flash Waveforms, in this section for signal waveforms)