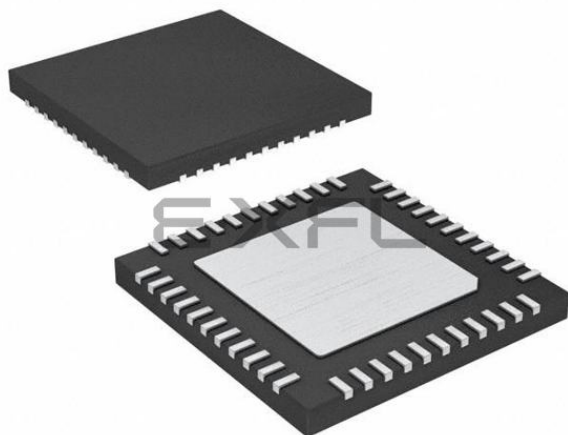




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Details

Product Status	Active
Core Processor	AVR
Core Size	8-Bit
Speed	10MHz
Connectivity	I ² C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	32
Program Memory Size	32KB (16K x 16)
Program Memory Type	FLASH
EEPROM Size	1K x 8
RAM Size	2K x 8
Voltage - Supply (Vcc/Vdd)	1.8V ~ 5.5V
Data Converters	A/D 8x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	44-VFQFN Exposed Pad
Supplier Device Package	44-VQFN (7x7)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/atmega324pv-10mur

5.6 Register Description

5.6.1 EEARH and EEARL – The EEPROM Address Register

Bit	15	14	13	12	11	10	9	8	
0x22 (0x42)	–	–	–	–	EEAR11	EEAR10	EEAR9	EEAR8	EEARH
0x21 (0x41)	EEAR7	EEAR6	EEAR5	EEAR4	EEAR3	EEAR2	EEAR1	EEAR0	EEARL
	7	6	5	4	3	2	1	0	
Read/Write	R	R	R	R	R/W	R/W	R/W	R/W	
	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	X	X	X	X	
	X	X	X	X	X	X	X	X	

- **Bits 15:12 – Res: Reserved Bits**

These bits are reserved bits in the ATmega164P/324P/644P and will always read as zero.

- **Bits 11:0 – EEAR8:0: EEPROM Address**

The EEPROM Address Registers – EEARH and EEARL specify the EEPROM address in the 4K bytes EEPROM space. The EEPROM data bytes are addressed linearly between 0 and 4096. The initial value of EEAR is undefined. A proper value must be written before the EEPROM may be accessed.

5.6.2 EEDR – The EEPROM Data Register

Bit	7	6	5	4	3	2	1	0	
0x20 (0x40)	MSB							LSB	EEDR
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bits 7:0 – EEDR7:0: EEPROM Data**

For the EEPROM write operation, the EEDR Register contains the data to be written to the EEPROM in the address given by the EEAR Register. For the EEPROM read operation, the EEDR contains the data read out from the EEPROM at the address given by EEAR.

5.6.3 EECR – The EEPROM Control Register

Bit	7	6	5	4	3	2	1	0	
0x1F (0x3F)	–	–	EEPM1	EEPM0	EERIE	EEMPE	EEPE	EERE	EECR
Read/Write	R	R	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	X	X	0	0	X	0	

- **Bits 7:6 – Res: Reserved Bits**

These bits are reserved bits in the ATmega164P/324P/644P and will always read as zero.

- **Bits 5:4 – EEPM1 and EEPM0: EEPROM Programming Mode Bits**

The EEPROM Programming mode bit setting defines which programming action that will be triggered when writing EEPE. It is possible to program data in one atomic operation (erase the old value and program the new value) or to split the Erase and Write operations in two different operations. The Programming times for the different modes are shown in Table 5-1 on page 24.

pletion of the currently executing instruction to generate an interrupt. If enabled, a level triggered interrupt will generate an interrupt request as long as the pin is held low. When changing the ISCn bit, an interrupt can occur. Therefore, it is recommended to first disable INTn by clearing its Interrupt Enable bit in the EIMSK Register. Then, the ISCn bit can be changed. Finally, the INTn interrupt flag should be cleared by writing a logical one to its Interrupt Flag bit (INTFn) in the EIFR Register before the interrupt is re-enabled.

Table 10-1. Interrupt Sense Control⁽¹⁾

ISCn1	ISCn0	Description
0	0	The low level of INTn generates an interrupt request.
0	1	Any edge of INTn generates asynchronously an interrupt request.
1	0	The falling edge of INTn generates asynchronously an interrupt request.
1	1	The rising edge of INTn generates asynchronously an interrupt request.

Note: 1. n = 2, 1 or 0.

When changing the ISCn1/ISCn0 bits, the interrupt must be disabled by clearing its Interrupt Enable bit in the EIMSK Register. Otherwise an interrupt can occur when the bits are changed.

10.2.2 EIMSK – External Interrupt Mask Register

Bit	7	6	5	4	3	2	1	0	
0x1D (0x3D)	–	–	–	–	–	INT2	INT1	IINT0	EIMSK
Read/Write	R	R	R	R	R	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

• Bits 2:0 – INT2:0: External Interrupt Request 2 - 0 Enable

When an INT2:0 bit is written to one and the I-bit in the Status Register (SREG) is set (one), the corresponding external pin interrupt is enabled. The Interrupt Sense Control bits in the External Interrupt Control Register, EICRA, defines whether the external interrupt is activated on rising or falling edge or level sensed. Activity on any of these pins will trigger an interrupt request even if the pin is enabled as an output. This provides a way of generating a software interrupt.

10.2.3 EIFR –External Interrupt Flag Register

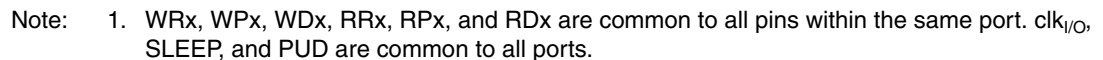
Bit	7	6	5	4	3	2	1	0	
0x1C (0x3C)	–	–	–	–	–	INTF2	INTF1	IINTF0	EIFR
Read/Write	R/W	R	R	R	R	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

• Bits 2:0 – INTF2:0: External Interrupt Flags 2 - 0

When an edge or logic change on the INT2:0 pin triggers an interrupt request, INTF2:0 becomes set (one). If the I-bit in SREG and the corresponding interrupt enable bit, INT2:0 in EIMSK, are set (one), the MCU will jump to the interrupt vector. The flag is cleared when the interrupt routine is executed. Alternatively, the flag can be cleared by writing a logical one to it. These flags are always cleared when INT2:0 are configured as level interrupt. Note that when entering sleep mode with the INT2:0 interrupts disabled, the input buffers on these pins will be disabled. This may cause a logic change in internal signals which will set the INTF2:0 flags. See "Digital Input Enable and Sleep Modes" on page 76 for more information.

11.2 Ports as General Digital I/O

Figure 11-2. General Digital I/O⁽¹⁾



If PORTxn is written logic one when the pin is configured as an input pin, the pull-up resistor is activated. To switch the pull-up resistor off, PORTxn has to be written logic zero or the pin has to be configured as an output pin. The port pins are tri-stated when reset condition becomes active, even if no clocks are running.

If PORTxn is written logic one when the pin is configured as an output pin, the port pin is driven high (one). If PORTxn is written logic zero when the pin is configured as an output pin, the port pin is driven low (zero).

11.2.2 Toggling the Pin

Writing a logic one to PINxn toggles the value of PORTxn, independent on the value of DDRxn. Note that the SBI instruction can be used to toggle one single bit in a port.

11.2.3 Switching Between Input and Output

When switching between tri-state ({DDxn, PORTxn} = 0b00) and output high ({DDxn, PORTxn} = 0b11), an intermediate state with either pull-up enabled {DDxn, PORTxn} = 0b01) or output low ({DDxn, PORTxn} = 0b10) must occur. Normally, the pull-up enabled state is fully acceptable, as a high-impedant environment will not notice the difference between a strong high driver and a pull-up. If this is not the case, the PUD bit in the MCUCR Register can be set to disable all pull-ups in all ports.

Switching between input with pull-up and output low generates the same problem. The user must use either the tri-state ({DDxn, PORTxn} = 0b00) or the output high state ({DDxn, PORTxn} = 0b11) as an intermediate step.

Table 11-1 summarizes the control signals for the pin value.

Table 11-1. Port Pin Configurations

DDxn	PORTxn	PUD (in MCUCR)	I/O	Pull-up	Comment
0	0	X	Input	No	Tri-state (Hi-Z)
0	1	0	Input	Yes	Pxn will source current if ext. pulled low.
0	1	1	Input	No	Tri-state (Hi-Z)
1	0	X	Output	No	Output Low (Sink)
1	1	X	Output	No	Output High (Source)

11.2.4 Reading the Pin Value

Independent of the setting of Data Direction bit DDxn, the port pin can be read through the PINxn Register bit. As shown in Figure 11-2, the PINxn Register bit and the preceding latch constitute a synchronizer. This is needed to avoid metastability if the physical pin changes value near the edge of the internal clock, but it also introduces a delay. Figure 11-3 shows a timing diagram of the synchronization when reading an externally applied pin value. The maximum and minimum propagation delays are denoted $t_{pd,max}$ and $t_{pd,min}$ respectively.

Table 11-11. Overriding Signals for Alternate Functions in PC3:PC0

Signal Name	PC3/TMS/ PCINT19	PC2/TCK/ PCINT18	PC1/SDA/ PCINT17	PC0/SCL/ PCINT16
PUOE	JTAGEN	JTAGEN	TWEN	TWEN
PUOV	1	1	PORTC1 • $\overline{\text{PUD}}$	PORTC0 • $\overline{\text{PUD}}$
DDOE	JTAGEN	JTAGEN	TWEN	TWEN
DDOV	0	0	0	0
PVOE	0	0	TWEN	TWEN
PVOV	0	0	SDA OUT	SCL OUT
DIEOE	JTAGEN + PCINT19 • PCIE2	JTAGEN + PCINT18 • PCIE2	PCINT17 • PCIE2	PCINT16 • PCIE2
DIEOV	$\overline{\text{JTAGEN}}$	$\overline{\text{JTAGEN}}$	1	1
DI	PCINT19 INPUT	PCINT18 INPUT	PCINT17 INPUT	PCINT16 INPUT
AIO	TMS INPUT	TCK INPUT	SDA INPUT	SCL INPUT

11.3.4 Alternate Functions of Port D

The Port D pins with alternate functions are shown in Table 11-12.

Table 11-12. Port D Pins Alternate Functions

Port Pin	Alternate Function
PD7	OC2A (Timer/Counter2 Output Compare Match A Output) PCINT31 (Pin Change Interrupt 31)
PD6	ICP1 (Timer/Counter1 Input Capture Trigger) OC2B (Timer/Counter2 Output Compare Match B Output) PCINT30 (Pin Change Interrupt 30)
PD5	OC1A (Timer/Counter1 Output Compare Match A Output) PCINT29 (Pin Change Interrupt 29)
PD4	OC1B (Timer/Counter1 Output Compare Match B Output) XCK1 (USART1 External Clock Input/Output) PCINT28 (Pin Change Interrupt 28)
PD3	INT1 (External Interrupt1 Input) TXD1 (USART1 Transmit Pin) PCINT27 (Pin Change Interrupt 27)
PD2	INT0 (External Interrupt0 Input) RXD1 (USART1 Receive Pin) PCINT26 (Pin Change Interrupt 26)
PD1	TXD0 (USART0 Transmit Pin) PCINT25 (Pin Change Interrupt 25)
PD0	RXD0 (USART0 Receive Pin) PCINT24 (Pin Change Interrupt 24)

Assembly Code Example⁽¹⁾⁽²⁾

```

USART_Transmit:
    ; Wait for empty transmit buffer
    sbis UCSRnA, UDREn
    rjmp USART_Transmit
    ; Copy 9th bit from r17 to TXB8
    cbi UCSRnB, TXB8
    sbrc r17, 0
    sbi UCSRnB, TXB8
    ; Put LSB data (r16) into buffer, sends the data
    out UDRn, r16
    ret

```

C Code Example⁽¹⁾⁽²⁾

```

void USART_Transmit( unsigned int data )
{
    /* Wait for empty transmit buffer */
    while ( !( UCSRnA & (1<<UDREn)) )
        ;
    /* Copy 9th bit to TXB8 */
    UCSRnB &= ~(1<<TXB8);
    if ( data & 0x0100 )
        UCSRnB |= (1<<TXB8);
    /* Put data into buffer, sends the data */
    UDRn = data;
}

```

- Notes:
1. These transmit functions are written to be general functions. They can be optimized if the contents of the UCSRnB is static. For example, only the TXB8 bit of the UCSRnB Register is used after initialization.
 2. See "About Code Examples" on page 8.

The ninth bit can be used for indicating an address frame when using multi processor communication mode or for other protocol handling as for example synchronization.

16.7.3 Transmitter Flags and Interrupts

The USART Transmitter has two flags that indicate its state: USART Data Register Empty (UDREn) and Transmit Complete (TXCn). Both flags can be used for generating interrupts.

The Data Register Empty (UDREn) Flag indicates whether the transmit buffer is ready to receive new data. This bit is set when the transmit buffer is empty, and cleared when the transmit buffer contains data to be transmitted that has not yet been moved into the Shift Register. For compatibility with future devices, always write this bit to zero when writing the UCSRnA Register.

When the Data Register Empty Interrupt Enable (UDRIEn) bit in UCSRnB is written to one, the USART Data Register Empty Interrupt will be executed as long as UDREn is set (provided that global interrupts are enabled). UDREn is cleared by writing UDRn. When interrupt-driven data

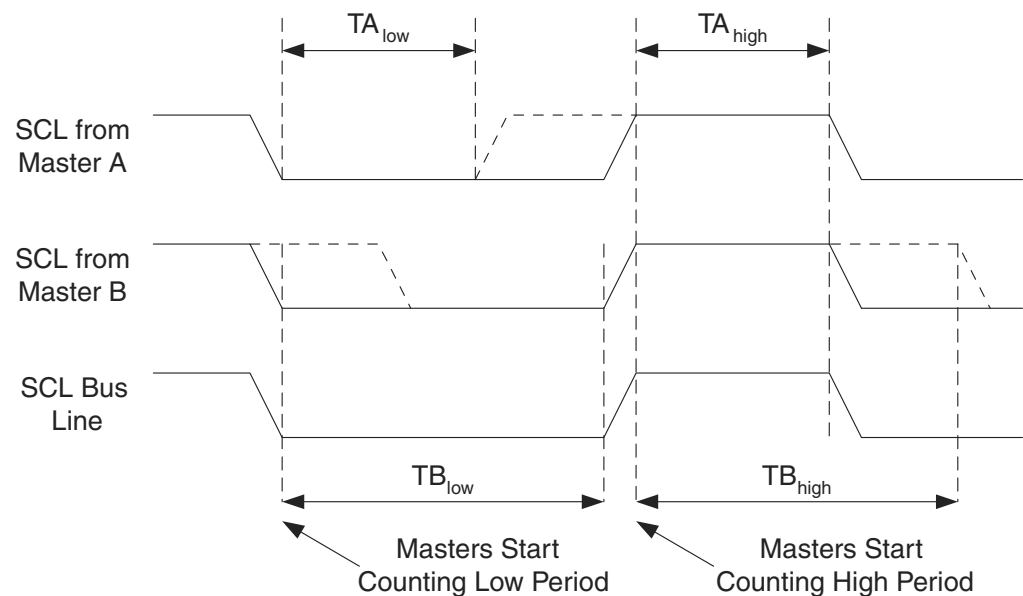
18.4 Multi-master Bus Systems, Arbitration and Synchronization

The TWI protocol allows bus systems with several masters. Special concerns have been taken in order to ensure that transmissions will proceed as normal, even if two or more masters initiate a transmission at the same time. Two problems arise in multi-master systems:

- An algorithm must be implemented allowing only one of the masters to complete the transmission. All other masters should cease transmission when they discover that they have lost the selection process. This selection process is called arbitration. When a contending master discovers that it has lost the arbitration process, it should immediately switch to Slave mode to check whether it is being addressed by the winning master. The fact that multiple masters have started transmission at the same time should not be detectable to the slaves, i.e. the data being transferred on the bus must not be corrupted.
- Different masters may use different SCL frequencies. A scheme must be devised to synchronize the serial clocks from all masters, in order to let the transmission proceed in a lockstep fashion. This will facilitate the arbitration process.

The wired-ANDing of the bus lines is used to solve both these problems. The serial clocks from all masters will be wired-ANDed, yielding a combined clock with a high period equal to the one from the Master with the shortest high period. The low period of the combined clock is equal to the low period of the Master with the longest low period. Note that all masters listen to the SCL line, effectively starting to count their SCL high and low time-out periods when the combined SCL line goes high or low, respectively.

Figure 18-7. SCL Synchronization Between Multiple Masters



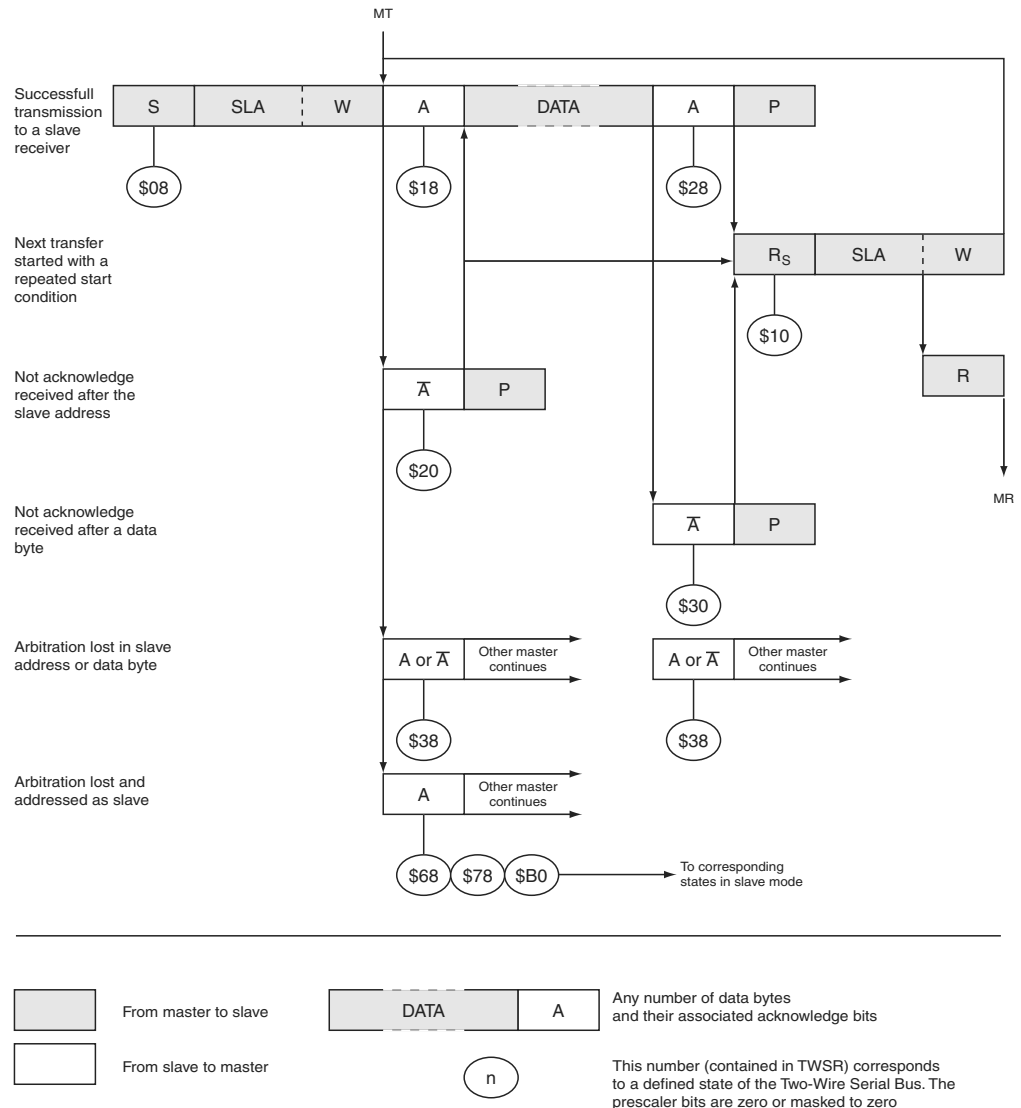
Arbitration is carried out by all masters continuously monitoring the SDA line after outputting data. If the value read from the SDA line does not match the value the Master had output, it has lost the arbitration. Note that a Master can only lose arbitration when it outputs a high SDA value while another Master outputs a low value. The losing Master should immediately go to Slave mode, checking if it is being addressed by the winning Master. The SDA line should be left high, but losing masters are allowed to generate a clock signal until the end of the current data or address packet. Arbitration will continue until only one Master remains, and this may take many

- When the TWINT Flag is set, the user must update all TWI Registers with the value relevant for the next TWI bus cycle. As an example, TWDR must be loaded with the value to be transmitted in the next bus cycle.
- After all TWI Register updates and other pending application software tasks have been completed, TWCR is written. When writing TWCR, the TWINT bit should be set. Writing a one to TWINT clears the flag. The TWI will then commence executing whatever operation was specified by the TWCR setting.

In the following an assembly and C implementation of the example is given. Note that the code below assumes that several definitions have been made, for example by using include-files.

	Assembly Code Example	C Example	Comments
1	<pre>ldi r16, (1<<TWINT) (1<<TWSTA) (1<<TWEN) out TWCR, r16</pre>	<pre>TWCR = (1<<TWINT) (1<<TWSTA) (1<<TWEN)</pre>	Send START condition
2	<pre>wait1: in r16,TWCR sbrs r16,TWINT rjmp wait1</pre>	<pre>while (!(TWCR & (1<<TWINT))) ;</pre>	Wait for TWINT Flag set. This indicates that the START condition has been transmitted
3	<pre>in r16,TWSR andi r16, 0xF8 cpi r16, START brne ERROR</pre>	<pre>if ((TWSR & 0xF8) != START) ERROR();</pre>	Check value of TWI Status Register. Mask prescaler bits. If status different from START go to ERROR
	<pre>ldi r16, SLA_W out TWDR, r16 ldi r16, (1<<TWINT) (1<<TWEN) out TWCR, r16</pre>	<pre>TWDR = SLA_W; TWCR = (1<<TWINT) (1<<TWEN);</pre>	Load SLA_W into TWDR Register. Clear TWINT bit in TWCR to start transmission of address
4	<pre>wait2: in r16,TWCR sbrs r16,TWINT rjmp wait2</pre>	<pre>while (!(TWCR & (1<<TWINT))) ;</pre>	Wait for TWINT Flag set. This indicates that the SLA+W has been transmitted, and ACK/NACK has been received.
5	<pre>in r16,TWSR andi r16, 0xF8 cpi r16, MT_SLA_ACK brne ERROR</pre>	<pre>if ((TWSR & 0xF8) != MT_SLA_ACK) ERROR();</pre>	Check value of TWI Status Register. Mask prescaler bits. If status different from MT_SLA_ACK go to ERROR
	<pre>ldi r16, DATA out TWDR, r16 ldi r16, (1<<TWINT) (1<<TWEN) out TWCR, r16</pre>	<pre>TWDR = DATA; TWCR = (1<<TWINT) (1<<TWEN);</pre>	Load DATA into TWDR Register. Clear TWINT bit in TWCR to start transmission of data

Figure 18-12. Formats and States in the Master Transmitter Mode



18.7.2 Master Receiver Mode

In the Master Receiver mode, a number of data bytes are received from a Slave Transmitter (Slave see Figure 18-13 on page 222). In order to enter a Master mode, a START condition must be transmitted. The format of the following address packet determines whether Master Transmitter or Master Receiver mode is to be entered. If SLA+W is transmitted, MT mode is entered, if SLA+R is transmitted, MR mode is entered. All the status codes mentioned in this section assume that the prescaler bits are zero or are masked to zero.

Table 18-4. Status Codes for Slave Receiver Mode

Status Code (TWSR) Prescaler Bits are 0	Status of the 2-wire Serial Bus and 2-wire Serial Interface Hardware	Application Software Response					Next Action Taken by TWI Hardware
		To/from TWDR	To TWCR				
			STA	STO	TWINT	TWEA	
0x60	Own SLA+W has been received; ACK has been returned	No TWDR action or	X	0	1	0	Data byte will be received and NOT ACK will be returned
		No TWDR action	X	0	1	1	Data byte will be received and ACK will be returned
0x68	Arbitration lost in SLA+R/W as Master; own SLA+W has been received; ACK has been returned	No TWDR action or	X	0	1	0	Data byte will be received and NOT ACK will be returned
		No TWDR action	X	0	1	1	Data byte will be received and ACK will be returned
0x70	General call address has been received; ACK has been returned	No TWDR action or	X	0	1	0	Data byte will be received and NOT ACK will be returned
		No TWDR action	X	0	1	1	Data byte will be received and ACK will be returned
0x78	Arbitration lost in SLA+R/W as Master; General call address has been received; ACK has been returned	No TWDR action or	X	0	1	0	Data byte will be received and NOT ACK will be returned
		No TWDR action	X	0	1	1	Data byte will be received and ACK will be returned
0x80	Previously addressed with own SLA+W; data has been received; ACK has been returned	Read data byte or	X	0	1	0	Data byte will be received and NOT ACK will be returned
		Read data byte	X	0	1	1	Data byte will be received and ACK will be returned
0x88	Previously addressed with own SLA+W; data has been received; NOT ACK has been returned	Read data byte or	0	0	1	0	Switched to the not addressed Slave mode; no recognition of own SLA or GCA
		Read data byte or	0	0	1	1	Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1"
		Read data byte or	1	0	1	0	Switched to the not addressed Slave mode; no recognition of own SLA or GCA; a START condition will be transmitted when the bus becomes free
		Read data byte	1	0	1	1	Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1"; a START condition will be transmitted when the bus becomes free
0x90	Previously addressed with general call; data has been re- ceived; ACK has been returned	Read data byte or	X	0	1	0	Data byte will be received and NOT ACK will be returned
		Read data byte	X	0	1	1	Data byte will be received and ACK will be returned
0x98	Previously addressed with general call; data has been received; NOT ACK has been returned	Read data byte or	0	0	1	0	Switched to the not addressed Slave mode; no recognition of own SLA or GCA
		Read data byte or	0	0	1	1	Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1"
		Read data byte or	1	0	1	0	Switched to the not addressed Slave mode; no recognition of own SLA or GCA; a START condition will be transmitted when the bus becomes free
		Read data byte	1	0	1	1	Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1"; a START condition will be transmitted when the bus becomes free
0xA0	A STOP condition or repeated START condition has been received while still addressed as Slave	No action	0	0	1	0	Switched to the not addressed Slave mode; no recognition of own SLA or GCA
			0	0	1	1	Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1"
			1	0	1	0	Switched to the not addressed Slave mode; no recognition of own SLA or GCA; a START condition will be transmitted when the bus becomes free
			1	0	1	1	Switched to the not addressed Slave mode; own SLA will be recognized; GCA will be recognized if TWGCE = "1"; a START condition will be transmitted when the bus becomes free

Table 20-5. ADC Prescaler Selections (Continued)

ADPS2	ADPS1	ADPS0	Division Factor
1	0	0	16
1	0	1	32
1	1	0	64
1	1	1	128

20.9.3 ADCL and ADCH – The ADC Data Register

ADLAR = 0

Bit	15	14	13	12	11	10	9	8	
(0x79)	–	–	–	–	–	–	ADC9	ADC8	ADCH
(0x78)	ADC7	ADC6	ADC5	ADC4	ADC3	ADC2	ADC1	ADC0	ADCL
	7	6	5	4	3	2	1	0	
Read/Write	R	R	R	R	R	R	R	R	
	R	R	R	R	R	R	R	R	
Initial Value	0	0	0	0	0	0	0	0	
	0	0	0	0	0	0	0	0	

ADLAR = 1

Bit	15	14	13	12	11	10	9	8	
(0x79)	ADC9	ADC8	ADC7	ADC6	ADC5	ADC4	ADC3	ADC2	ADCH
(0x78)	ADC1	ADC0	–	–	–	–	–	–	ADCL
	7	6	5	4	3	2	1	0	
Read/Write	R	R	R	R	R	R	R	R	
	R	R	R	R	R	R	R	R	
Initial Value	0	0	0	0	0	0	0	0	
	0	0	0	0	0	0	0	0	

When an ADC conversion is complete, the result is found in these two registers. If differential channels are used, the result is presented in two's complement form.

When ADCL is read, the ADC Data Register is not updated until ADCH is read. Consequently, if the result is left adjusted and no more than 8-bit precision is required, it is sufficient to read ADCH. Otherwise, ADCL must be read first, then ADCH.

The ADLAR bit in ADMUX, and the MUXn bits in ADMUX affect the way the result is read from the registers. If ADLAR is set, the result is left adjusted. If ADLAR is cleared (default), the result is right adjusted.

• ADC9:0: ADC Conversion Result

These bits represent the result from the conversion, as detailed in "ADC Conversion Result" on page 253.

20.9.4 ADCSRB – ADC Control and Status Register B

Bit	7	6	5	4	3	2	1	0	
(0x7B)	–	ACME	–	–	–	ADTS2	ADTS1	ADTS0	ADCSRB
Read/Write	R	R/W	R	R	R	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

The JTAG programming capability supports:

- Flash programming and verifying.
- EEPROM programming and verifying.
- Fuse programming and verifying.
- Lock bit programming and verifying.

The Lock bit security is exactly as in parallel programming mode. If the Lock bits LB1 or LB2 are programmed, the OCDEN Fuse cannot be programmed unless first doing a chip erase. This is a security feature that ensures no back-door exists for reading out the content of a secured device.

The details on programming through the JTAG interface and programming specific JTAG instructions are given in the section "Programming via the JTAG Interface" on page 312.

21.9 Bibliography

For more information about general Boundary-scan, the following literature can be consulted:

- IEEE: IEEE Std. 1149.1-1990. IEEE Standard Test Access Port and Boundary-scan Architecture, IEEE, 1993.
- Colin Maunder: The Board Designers Guide to Testable Logic Circuits, Addison-Wesley, 1992.

21.10 Register Description

21.10.1 OCDR – On-chip Debug Register

Bit	7	6	5	4	3	2	1	0	
0x31 (0x51)	MSB/IDRD							LSB	OCDR
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

The OCDR Register provides a communication channel from the running program in the micro-controller to the debugger. The CPU can transfer a byte to the debugger by writing to this location. At the same time, an internal flag; I/O Debug Register Dirty – IDRD – is set to indicate to the debugger that the register has been written. When the CPU reads the OCDR Register the 7 LSB will be from the OCDR Register, while the MSB is the IDRD bit. The debugger clears the IDRD bit when it has read the information.

In some AVR devices, this register is shared with a standard I/O location. In this case, the OCDR Register can only be accessed if the OCDEN Fuse is programmed, and the debugger enables access to the OCDR Register. In all other cases, the standard I/O location is accessed.

Refer to the debugger documentation for further information on how to use this register.

```
    sbiw loophi:looplo, 1          ;use subi for PAGESIZEB<=256
    brne Rdloop

    ; return to RWW section
    ; verify that RWW section is safe to read
Return:
    in    temp1, SPMCSR
    sbrs temp1, RWWSB             ; If RWWSB is set, the RWW section is not ready yet
    ret
    ; re-enable the RWW section
    ldi   spmcval, (1<<RWWSRE) | (1<<SPMEN)
    call Do_spm
    rjmp Return

Do_spm:
    ; check for previous SPM complete
Wait_spm:
    in    temp1, SPMCSR
    sbrc temp1, SPMEN
    rjmp Wait_spm
    ; input: spmcval determines SPM action
    ; disable interrupts if enabled, store status
    in    temp2, SREG
    cli
    ; check that no EEPROM write access is present
Wait_ee:
    sbic EECR, EEPE
    rjmp Wait_ee
    ; SPM timed sequence
    out   SPMCSR, spmcval
    spm
    ; restore SREG (to enable interrupts if originally enabled)
    out   SREG, temp2
    ret
```

23.9 Register Description

23.9.1 SPMCSR – Store Program Memory Control and Status Register

The Store Program Memory Control and Status Register contains the control bits needed to control the Boot Loader operations.

Bit	7	6	5	4	3	2	1	0	
0x37 (0x57)	SPMIE	RWWSB	SIGRD	RWWSRE	BLBSET	PGWRT	PGERS	SPMEN	SPMCSR
Read/Write	R/W	R	R/W	R/W	R/W	R/W	R/W	R/W	
Initial Value	0	0	0	0	0	0	0	0	

- **Bit 7 – SPMIE: SPM Interrupt Enable**

When the SPMIE bit is written to one, and the I-bit in the Status Register is set (one), the SPM ready interrupt will be enabled. The SPM ready Interrupt will be executed as long as the SPMEN bit in the SPMCSR Register is cleared.

- **Bit 6 – RWWSB: Read-While-Write Section Busy**

When a Self-Programming (Page Erase or Page Write) operation to the RWW section is initiated, the RWWSB will be set (one) by hardware. When the RWWSB bit is set, the RWW section cannot be accessed. The RWWSB bit will be cleared if the RWWSRE bit is written to one after a Self-Programming operation is completed. Alternatively the RWWSB bit will automatically be cleared if a page load operation is initiated.

- **Bit 5 – SIGRD: Signature Row Read**

If this bit is written to one at the same time as SPMEN, the next LPM instruction within three clock cycles will read a byte from the signature row into the destination register. see "Reading the Signature Row from Software" on page 284 for details. An SPM instruction within four cycles after SIGRD and SPMEN are set will have no effect. This operation is reserved for future use and should not be used.

- **Bit 4 – RWWSRE: Read-While-Write Section Read Enable**

When programming (Page Erase or Page Write) to the RWW section, the RWW section is blocked for reading (the RWWSB will be set by hardware). To re-enable the RWW section, the user software must wait until the programming is completed (SPMEN will be cleared). Then, if the RWWSRE bit is written to one at the same time as SPMEN, the next SPM instruction within four clock cycles re-enables the RWW section. The RWW section cannot be re-enabled while the Flash is busy with a Page Erase or a Page Write (SPMEN is set). If the RWWSRE bit is written while the Flash is being loaded, the Flash load operation will abort and the data loaded will be lost.

- **Bit 3 – BLBSET: Boot Lock Bit Set**

If this bit is written to one at the same time as SPMEN, the next SPM instruction within four clock cycles sets Boot Lock bits, according to the data in R0. The data in R1 and the address in the Z-pointer are ignored. The BLBSET bit will automatically be cleared upon completion of the Lock bit set, or if no SPM instruction is executed within four clock cycles.

An (E)LPM instruction within three cycles after BLBSET and SPMEN are set in the SPMCSR Register, will read either the Lock bits or the Fuse bits (depending on Z0 in the Z-pointer) into the destination register. See "Reading the Fuse and Lock Bits from Software" on page 284 for details.

3. Set DATA = Address extended high byte (0x00 - 0xFF).
4. Give XTAL1 a positive pulse. This loads the address high byte.

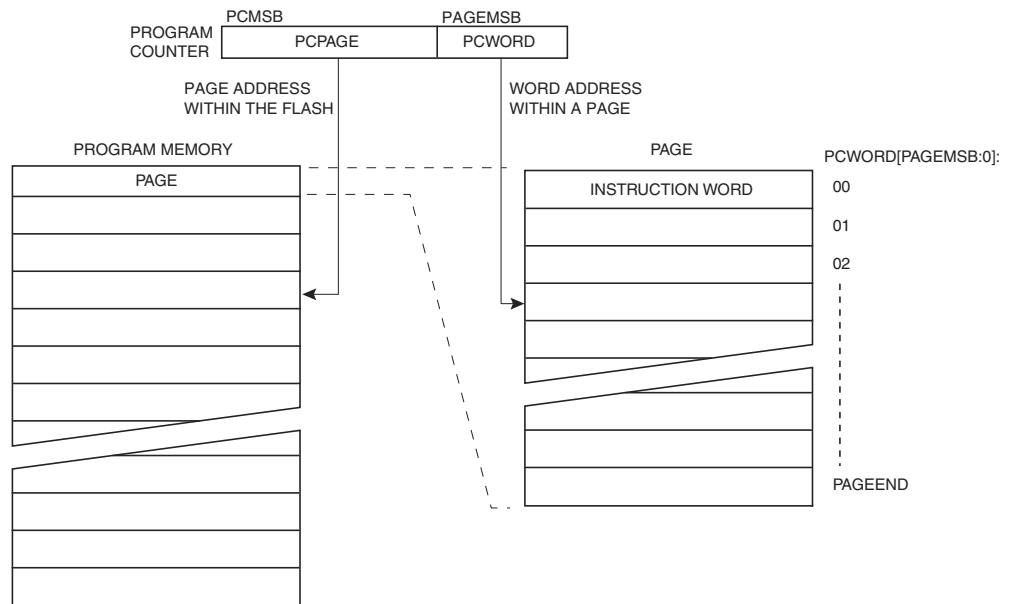
I. Program Page

1. Set BS2, BS1 to "00"
2. Give $\overline{WR}$ a negative pulse. This starts programming of the entire page of data. RDY/ $\overline{BSY}$ goes low.
3. Wait until RDY/ $\overline{BSY}$ goes high (See Figure 24-3 on page 302 for signal waveforms).
- J. Repeat B through I until the entire Flash is programmed or until all data has been programmed.

K. End Page Programming

1. Set XA1, XA0 to "10". This enables command loading.
2. Set DATA to "0000 0000". This is the command for No Operation.
3. Give XTAL1 a positive pulse. This loads the command, and the internal write signals are reset.

Figure 24-2. Addressing the Flash Which is Organized in Pages⁽¹⁾



Note: 1. PCPAGE and PCWORD are listed in Table 24-7 on page 296.

Table 24-16. Minimum Wait Delay Before Writing the Next Flash or EEPROM Location

Symbol	Minimum Wait Delay
t_{WD_FLASH}	4.5 ms
t_{WD_EEPROM}	9.0 ms
t_{WD_ERASE}	9.0 ms

24.9 Serial Programming Instruction set

Table 24-17 on page 310 and Figure 24-11 on page 311 describes the Instruction set.

Table 24-17. Serial Programming Instruction Set (Hexadecimal values)

Instruction/Operation	Instruction Format			
	Byte 1	Byte 2	Byte 3	Byte4
Programming Enable	\$AC	\$53	\$00	\$00
Chip Erase (Program Memory/EEPROM)	\$AC	\$80	\$00	\$00
Poll RDY/ $\overline{BSY}$	\$F0	\$00	\$00	data byte out
Load Instructions				
Load Extended Address byte ⁽¹⁾	\$4D	\$00	Extended adr	\$00
Load Program Memory Page, High byte	\$48	\$00	adr LSB	high data byte in
Load Program Memory Page, Low byte	\$40	\$00	adr LSB	low data byte in
Load EEPROM Memory Page (page access)	\$C1	\$00	0000 000aa	data byte in
Read Instructions				
Read Program Memory, High byte	\$28	adr MSB	adr LSB	high data byte out
Read Program Memory, Low byte	\$20	adr MSB	adr LSB	low data byte out
Read EEPROM Memory	\$A0	0000 00aa	aaaa aaaa	data byte out
Read Lock bits	\$58	\$00	\$00	data byte out
Read Signature Byte	\$30	\$00	0000 000aa	data byte out
Read Fuse bits	\$50	\$00	\$00	data byte out
Read Fuse High bits	\$58	\$08	\$00	data byte out
Read Extended Fuse Bits	\$50	\$08	\$00	data byte out
Read Calibration Byte	\$38	\$00	\$00	data byte out
Write Instructions⁽⁶⁾				
Write Program Memory Page	\$4C	adr MSB	adr LSB	\$00
Write EEPROM Memory	\$C0	0000 00aa	aaaa aaaa	data byte in
Write EEPROM Memory Page (page access)	\$C2	0000 00aa	aaaa aa00	\$00
Write Lock bits	\$AC	\$E0	\$00	data byte in
Write Fuse bits	\$AC	\$A0	\$00	data byte in
Write Fuse High bits	\$AC	\$A8	\$00	data byte in
Write Extended Fuse Bits	\$AC	\$A4	\$00	data byte in

Table 24-18. JTAG Programming Instruction (Continued)

Set (Continued) **a** = address high bits, **b** = address low bits, **c** = address extended bits, **H** = 0 - Low byte, 1 - High Byte,
o = data out, **i** = data in, **x** = don't care

Instruction	TDI Sequence	TDO Sequence	Notes
5a. Enter EEPROM Read	0100011_00000011	xxxxxxx_xxxxxxxx	
5b. Load Address High Byte	0000111_aaaaaaaa	xxxxxxx_xxxxxxxx	(10)
5c. Load Address Low Byte	0000011_bbbbbbbb	xxxxxxx_xxxxxxxx	
5d. Read Data Byte	0110011_bbbbbbbb 0110010_00000000 0110011_00000000	xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_00000000	
6a. Enter Fuse Write	0100011_01000000	xxxxxxx_xxxxxxxx	
6b. Load Data Low Byte ⁽⁶⁾	0010011_iiiiiiii	xxxxxxx_xxxxxxxx	(3)
6c. Write Fuse Extended Byte	0111011_00000000 0111001_00000000 0111011_00000000 0111011_00000000	xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx	(1)
6d. Poll for Fuse Write Complete	0110111_00000000	xxxxxox_xxxxxxxx	(2)
6e. Load Data Low Byte ⁽⁷⁾	0010011_iiiiiiii	xxxxxxx_xxxxxxxx	(3)
6f. Write Fuse High Byte	0110111_00000000 0110101_00000000 0110111_00000000 0110111_00000000	xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx	(1)
6g. Poll for Fuse Write Complete	0110111_00000000	xxxxxox_xxxxxxxx	(2)
6h. Load Data Low Byte ⁽⁷⁾	0010011_iiiiiiii	xxxxxxx_xxxxxxxx	(3)
6i. Write Fuse Low Byte	0110011_00000000 0110001_00000000 0110011_00000000 0110011_00000000	xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx	(1)
6j. Poll for Fuse Write Complete	0110011_00000000	xxxxxox_xxxxxxxx	(2)
7a. Enter Lock Bit Write	0100011_00100000	xxxxxxx_xxxxxxxx	
7b. Load Data Byte ⁽⁹⁾	0010011_11iiiiii	xxxxxxx_xxxxxxxx	(4)
7c. Write Lock Bits	0110011_00000000 0110001_00000000 0110011_00000000 0110011_00000000	xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx xxxxxxx_xxxxxxxx	(1)
7d. Poll for Lock Bit Write complete	0110011_00000000	xxxxxox_xxxxxxxx	(2)
8a. Enter Fuse/Lock Bit Read	0100011_00000100	xxxxxxx_xxxxxxxx	
8b. Read Extended Fuse Byte ⁽⁶⁾	0111010_00000000 0111011_00000000	xxxxxxx_xxxxxxxx xxxxxxx_00000000	
8c. Read Fuse High Byte ⁽⁷⁾	0111110_00000000 0111111_00000000	xxxxxxx_xxxxxxxx xxxxxxx_00000000	
8d. Read Fuse Low Byte ⁽⁸⁾	0110010_00000000 0110011_00000000	xxxxxxx_xxxxxxxx xxxxxxx_00000000	

Figure 26-32. BOD Threshold vs. Temperature (Bodlevel is 2.7V).

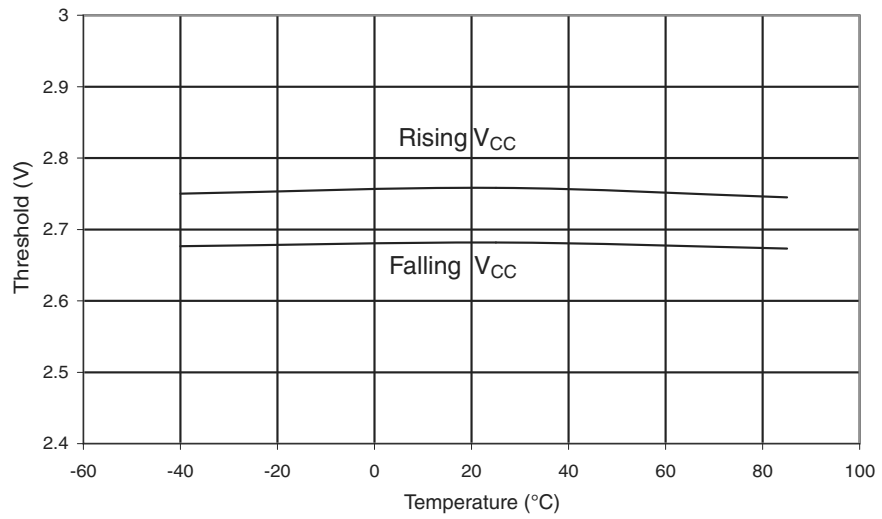
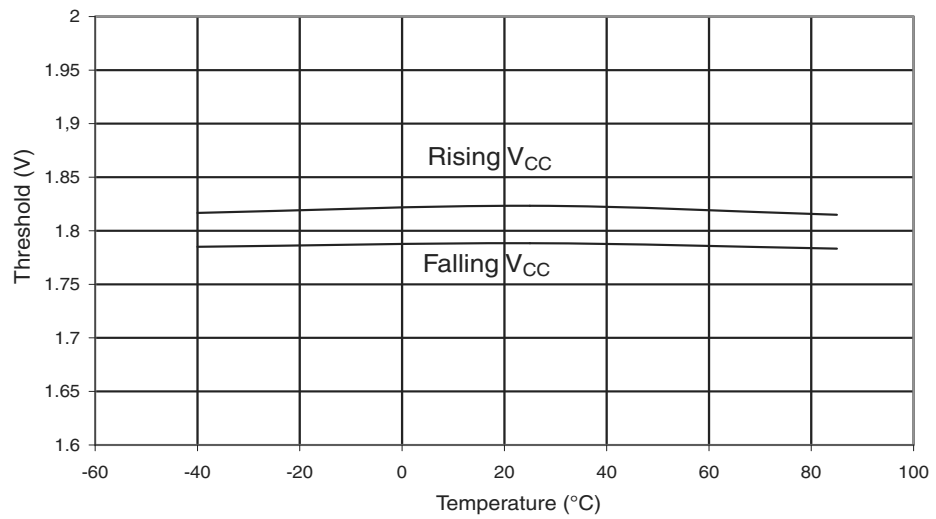


Figure 26-33. BOD Threshold vs. Temperature (Bodlevel is 1.8V).



26.2.4 Power-down Supply Current

Figure 26-58. Power-down Supply Current vs. V_{CC} (Watchdog Timer Disabled).

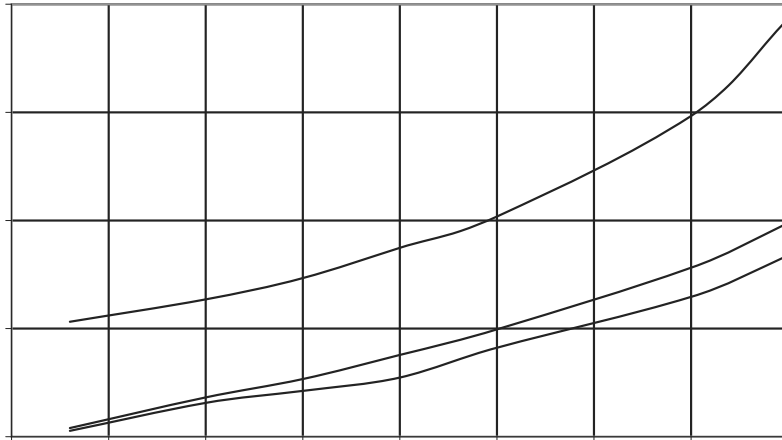


Figure 26-59. Power-down Supply Current vs. V_{CC} (Watchdog Timer Enabled).

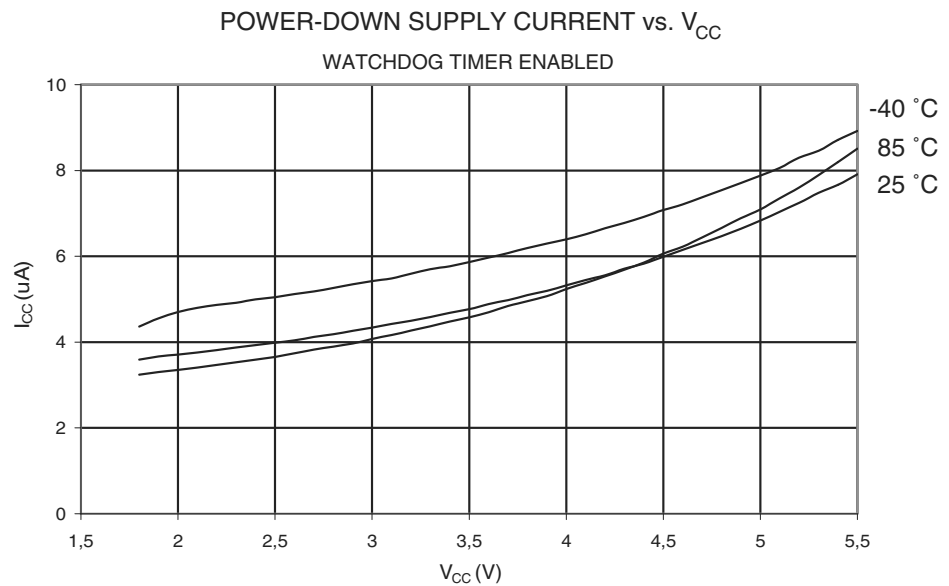


Figure 26-66. Reset Pull-up Resistor Current vs. Reset Pin Voltage ($V_{CC} = 2.7V$).

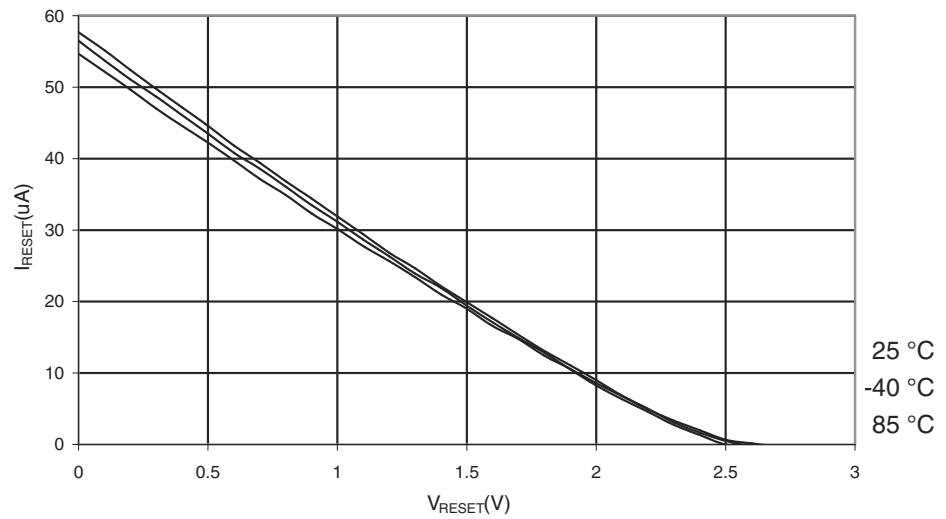


Figure 26-67. Reset Pull-up Resistor Current vs. Reset Pin Voltage ($V_{CC} = 5V$).

