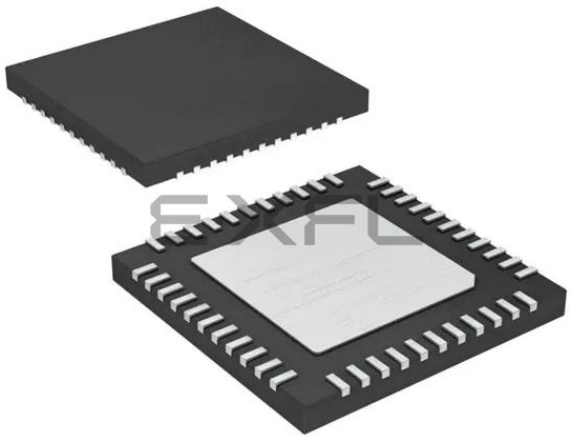

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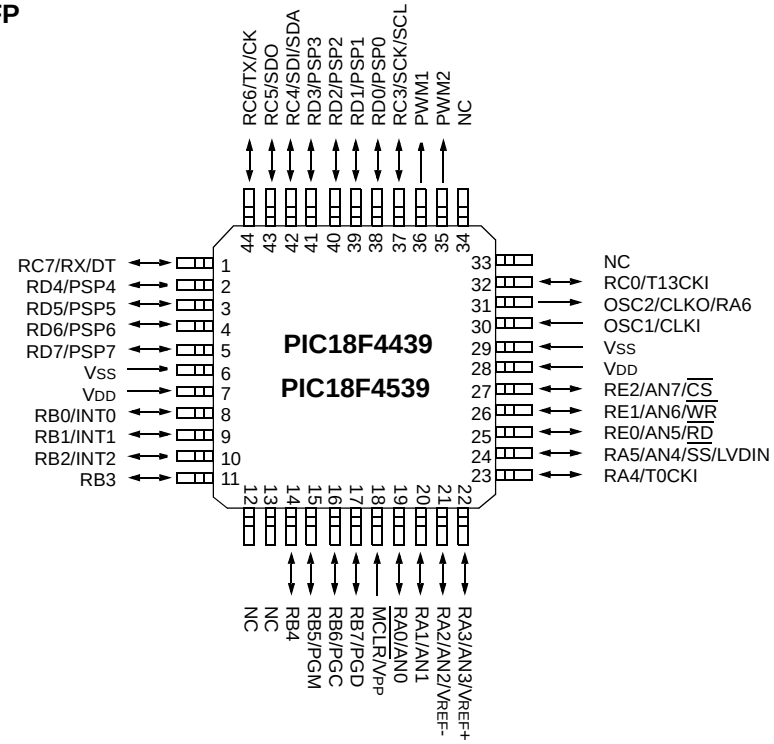
Applications of "Embedded - Microcontrollers"

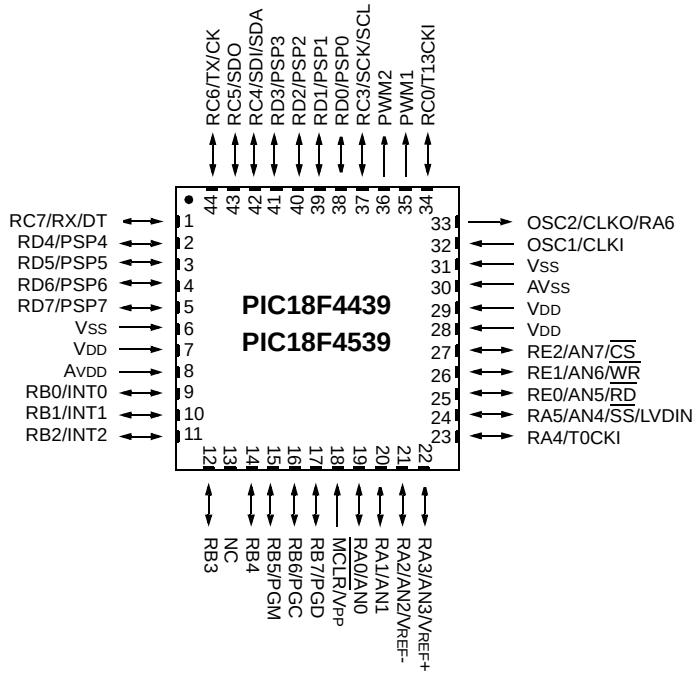
Details	
Product Status	Obsolete
Core Processor	PIC
Core Size	8-Bit
Speed	40MHz
Connectivity	I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, LVD, POR, PWM, WDT
Number of I/O	32
Program Memory Size	12KB (6K x 16)
Program Memory Type	FLASH
EEPROM Size	256 x 8
RAM Size	640 x 8
Voltage - Supply (Vcc/Vdd)	4.2V ~ 5.5V
Data Converters	A/D 8x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 125°C (TA)
Mounting Type	Surface Mount
Package / Case	44-VQFN Exposed Pad
Supplier Device Package	44-QFN (8x8)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18f4439-e-ml

Pin Diagrams

44-Pin TQFP



44-Pin QFN



4.6 Instruction Flow/Pipelining

An “Instruction Cycle” consists of four Q cycles (Q1, Q2, Q3 and Q4). The instruction fetch and execute are pipelined such that fetch takes one instruction cycle, while decode and execute takes another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the program counter to change (e.g., GOTO), then two cycles are required to complete the instruction (Example 4-1).

A fetch cycle begins with the program counter (PC) incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the “Instruction Register” (IR) in cycle Q1. This instruction is then decoded and executed during the Q2, Q3, and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).

EXAMPLE 4-1: INSTRUCTION PIPELINE FLOW

	TCY0	TCY1	TCY2	TCY3	TCY4	TCY5
1. MOVLW 55h	Fetch 1	Execute 1				
2. MOVWF PORTB		Fetch 2	Execute 2			
3. BRA SUB_1			Fetch 3	Execute 3		
4. BSF PORTA, BIT3 (Forced NOP)				Fetch 4	Flush (NOP)	
5. Instruction @ address SUB_1					Fetch SUB_1	Execute SUB_1

All instructions are single cycle, except for any program branches. These take two cycles since the fetch instruction is “flushed” from the pipeline while the new instruction is being fetched and then executed.

4.7 Instructions in Program Memory

The program memory is addressed in bytes. Instructions are stored as two bytes or four bytes in program memory. The Least Significant Byte of an instruction word is always stored in a program memory location with an even address (LSB = 0). Figure 4-4 shows an example of how instruction words are stored in the program memory. To maintain alignment with instruction boundaries, the PC increments in steps of 2 and the LSB will always read ‘0’ (see Section 4.4).

The CALL and GOTO instructions have an absolute program memory address embedded into the instruction. Since instructions are always stored on word boundaries, the data contained in the instruction is a word address. The word address is written to PC<20:1>, which accesses the desired byte address in program memory. Instruction #2 in Figure 4-4 shows how the instruction, ‘GOTO 000006h’, is encoded in the program memory. Program branch instructions, which encode a relative address offset, operate in the same manner. The offset value stored in a branch instruction represents the number of single word instructions that the PC will be offset by. Section 21.0 provides further details of the instruction set.

FIGURE 4-4: INSTRUCTIONS IN PROGRAM MEMORY

			Program Memory Byte Locations →	LSB = 1	LSB = 0	Word Address ↓
						000000h
						000002h
						000004h
						000006h
Instruction 1:	MOVLW	055h		0Fh	55h	000008h
Instruction 2:	GOTO	000006h		EFh	03h	00000Ah
				F0h	00h	00000Ch
Instruction 3:	MOVFF	123h, 456h		C1h	23h	00000Eh
				F4h	56h	000010h
						000012h
						000014h

4.13 STATUS Register

The STATUS register, shown in Register 4-2, contains the arithmetic status of the ALU. The STATUS register can be the destination for any instruction, as with any other register. If the STATUS register is the destination for an instruction that affects the Z, DC, C, OV, or N bits, then the write to these five bits is disabled. These bits are set or cleared according to the device logic. Therefore, the result of an instruction with the STATUS register as destination may be different than intended.

For example, `CLRF STATUS` will clear the upper three bits and set the Z bit. This leaves the STATUS register as `000u u1uu` (where `u` = unchanged).

It is recommended, therefore, that only `BCF`, `BSF`, `SWAPF`, `MOVFF` and `MOVWF` instructions are used to alter the STATUS register, because these instructions do not affect the Z, C, DC, OV, or N bits from the STATUS register. For other instructions not affecting any status bits, see Table 21-2.

Note:	The C and DC bits operate as a $\overline{\text{borrow}}$ and $\overline{\text{digit borrow}}$ bit respectively, in subtraction.

REGISTER 4-2: STATUS REGISTER

U-0	U-0	U-0	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x
—	—	—	N	OV	Z	DC	C
bit 7							bit 0

bit 7-5 **Unimplemented:** Read as '0'

bit 4 **N:** Negative bit
This bit is used for signed arithmetic (2's complement). It indicates whether the result was negative (ALU MSB = 1).
`1` = Result was negative
`0` = Result was positive

bit 3 **OV:** Overflow bit
This bit is used for signed arithmetic (2's complement). It indicates an overflow of the 7-bit magnitude, which causes the sign bit (bit 7) to change state.
`1` = Overflow occurred for signed arithmetic (in this arithmetic operation)
`0` = No overflow occurred

bit 2 **Z:** Zero bit
`1` = The result of an arithmetic or logic operation is zero
`0` = The result of an arithmetic or logic operation is not zero

bit 1 **DC:** Digit carry/$\overline{\text{borrow}}$ bit
For `ADDWF`, `ADDLW`, `SUBLW`, and `SUBWF` instructions
`1` = A carry-out from the 4th low order bit of the result occurred
`0` = No carry-out from the 4th low order bit of the result

Note: For $\overline{\text{borrow}}$, the polarity is reversed. A subtraction is executed by adding the two's complement of the second operand. For rotate (`RRF`, `RLF`) instructions, this bit is loaded with either the bit 4 or bit 3 of the source register.

bit 0 **C:** Carry/$\overline{\text{borrow}}$ bit
For `ADDWF`, `ADDLW`, `SUBLW`, and `SUBWF` instructions
`1` = A carry-out from the Most Significant bit of the result occurred
`0` = No carry-out from the Most Significant bit of the result occurred

Note: For $\overline{\text{borrow}}$, the polarity is reversed. A subtraction is executed by adding the two's complement of the second operand. For rotate (`RRF`, `RLF`) instructions, this bit is loaded with either the high or low order bit of the source register.

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
- n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

Example 7-3 shows the sequence to do a 16 x 16 unsigned multiply. Equation 7-1 shows the algorithm that is used. The 32-bit result is stored in four registers, RES3:RES0.

EQUATION 7-1: 16 x 16 UNSIGNED MULTIPLICATION ALGORITHM

$$
\begin{aligned}
RES3:RES0 &= ARG1H:ARG1L \bullet ARG2H:ARG2L \\
&= (ARG1H \bullet ARG2H \bullet 2^{16}) + \\
&\quad (ARG1H \bullet ARG2L \bullet 2^{8}) + \\
&\quad (ARG1L \bullet ARG2H \bullet 2^{8}) + \\
&\quad (ARG1L \bullet ARG2L)
\end{aligned}
$$

EXAMPLE 7-3: 16 x 16 UNSIGNED MULTIPLY ROUTINE

```
    MOVF    ARG1L, W
    MULWF   ARG2L         ; ARG1L * ARG2L ->
                          ; PRODH:PRODL
    MOVFF   PRODH, RES1   ;
    MOVFF   PRODL, RES0   ;
;
    MOVF    ARG1H, W
    MULWF   ARG2H         ; ARG1H * ARG2H ->
                          ; PRODH:PRODL
    MOVFF   PRODH, RES3   ;
    MOVFF   PRODL, RES2   ;
;
    MOVF    ARG1L, W
    MULWF   ARG2H         ; ARG1L * ARG2H ->
                          ; PRODH:PRODL
    MOVF    PRODL, W      ;
    ADDWF   RES1,  F      ; Add cross
    MOVF    PRODH, W      ; products
    ADDWFC  RES2,  F      ;
    CLRF    WREG          ;
    ADDWFC  RES3,  F      ;
;
    MOVF    ARG1H, W      ;
    MULWF   ARG2L         ; ARG1H * ARG2L ->
                          ; PRODH:PRODL
    MOVF    PRODL, W      ;
    ADDWF   RES1,  F      ; Add cross
    MOVF    PRODH, W      ; products
    ADDWFC  RES2,  F      ;
    CLRF    WREG          ;
    ADDWFC  RES3,  F      ;
```

Example 7-4 shows the sequence to do a 16 x 16 signed multiply. Equation 7-2 shows the algorithm used. The 32-bit result is stored in four registers, RES3:RES0. To account for the sign bits of the arguments, each argument pairs Most Significant bit (MSb) is tested and the appropriate subtractions are done.

EQUATION 7-2: 16 x 16 SIGNED MULTIPLICATION ALGORITHM

$$
\begin{aligned}
RES3:RES0 &= ARG1H:ARG1L \bullet ARG2H:ARG2L \\
&= (ARG1H \bullet ARG2H \bullet 2^{16}) + \\
&\quad (ARG1H \bullet ARG2L \bullet 2^{8}) + \\
&\quad (ARG1L \bullet ARG2H \bullet 2^{8}) + \\
&\quad (ARG1L \bullet ARG2L) + \\
&\quad (-1 \bullet ARG2H\langle 7\rangle \bullet ARG1H:ARG1L \bullet 2^{16}) + \\
&\quad (-1 \bullet ARG1H\langle 7\rangle \bullet ARG2H:ARG2L \bullet 2^{16})
\end{aligned}
$$

EXAMPLE 7-4: 16 x 16 SIGNED MULTIPLY ROUTINE

```
    MOVF    ARG1L, W
    MULWF   ARG2L         ; ARG1L * ARG2L ->
                          ; PRODH:PRODL
    MOVFF   PRODH, RES1   ;
    MOVFF   PRODL, RES0   ;
;
    MOVF    ARG1H, W
    MULWF   ARG2H         ; ARG1H * ARG2H ->
                          ; PRODH:PRODL
    MOVFF   PRODH, RES3   ;
    MOVFF   PRODL, RES2   ;
;
    MOVF    ARG1L, W
    MULWF   ARG2H         ; ARG1L * ARG2H ->
                          ; PRODH:PRODL
    MOVF    PRODL, W      ;
    ADDWF   RES1,  F      ; Add cross
    MOVF    PRODH, W      ; products
    ADDWFC  RES2,  F      ;
    CLRF    WREG          ;
    ADDWFC  RES3,  F      ;
;
    MOVF    ARG1H, W      ;
    MULWF   ARG2L         ; ARG1H * ARG2L ->
                          ; PRODH:PRODL
    MOVF    PRODL, W      ;
    ADDWF   RES1,  F      ; Add cross
    MOVF    PRODH, W      ; products
    ADDWFC  RES2,  F      ;
    CLRF    WREG          ;
    ADDWFC  RES3,  F      ;
;
    BTFSS   ARG2H, 7      ; ARG2H:ARG2L neg?
    BRA     SIGN_ARG1     ; no, check ARG1
    MOVF    ARG1L, W      ;
    SUBWF   RES2          ;
    MOVF    ARG1H, W      ;
    SUBWFB  RES3
;
SIGN_ARG1
    BTFSS   ARG1H, 7      ; ARG1H:ARG1L neg?
    BRA     CONT_CODE     ; no, done
    MOVF    ARG2L, W      ;
    SUBWF   RES2          ;
    MOVF    ARG2H, W      ;
    SUBWFB  RES3
;
CONT_CODE
     :
```


REGISTER 8-7: PIE2: PERIPHERAL INTERRUPT ENABLE REGISTER 2

U-0	U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	U-0
—	—	—	EEIE	BCLIE	LVDIE	TMR3IE	—
bit 7							bit 0

bit 7-5 **Unimplemented:** Read as '0'

bit 4 **EEIE**: Data EEPROM/FLASH Write Operation Interrupt Enable bit
1 = Enabled
0 = Disabled

bit 3 **BCLIE**: Bus Collision Interrupt Enable bit
1 = Enabled
0 = Disabled

bit 2 **LVDIE**: Low Voltage Detect Interrupt Enable bit
1 = Enabled
0 = Disabled

bit 1 **TMR3IE**: TMR3 Overflow Interrupt Enable bit
1 = Enables the TMR3 overflow interrupt
0 = Disables the TMR3 overflow interrupt

bit 0 **Unimplemented:** Read as ‘0’

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as ‘0’	
- n = Value at POR	‘1’ = Bit is set	‘0’ = Bit is cleared	x = Bit is unknown

9.3 PORTC, TRISC and LATC Registers

PORTC is a 6-bit wide, bi-directional port. The corresponding Data Direction register is TRISC. Setting a TRISC bit (= 1) will make the corresponding PORTC pin an input (i.e., put the corresponding output driver in a High Impedance mode). Clearing a TRISC bit (= 0) will make the corresponding PORTC pin an output (i.e., put the contents of the output latch on the selected pin).

The Data Latch register (LATC) is also memory mapped. Read-modify-write operations on the LATC register reads and writes the latched output value for PORTC.

PORTC is multiplexed with the serial communication functions (Table 9-5). PORTC pins have Schmitt Trigger input buffers.

When enabling peripheral functions, care should be taken in defining TRIS bits for each PORTC pin. Some peripherals override the TRIS bit to make a pin an output, while other peripherals override the TRIS bit to make a pin an input. The user should refer to the corresponding peripheral section for the correct TRIS bit settings.

Note:	On a Power-on Reset, these pins are configured as digital inputs.

The pin override value is not loaded into the TRIS register. This allows read-modify-write of the TRIS register, without concern due to peripheral overrides.

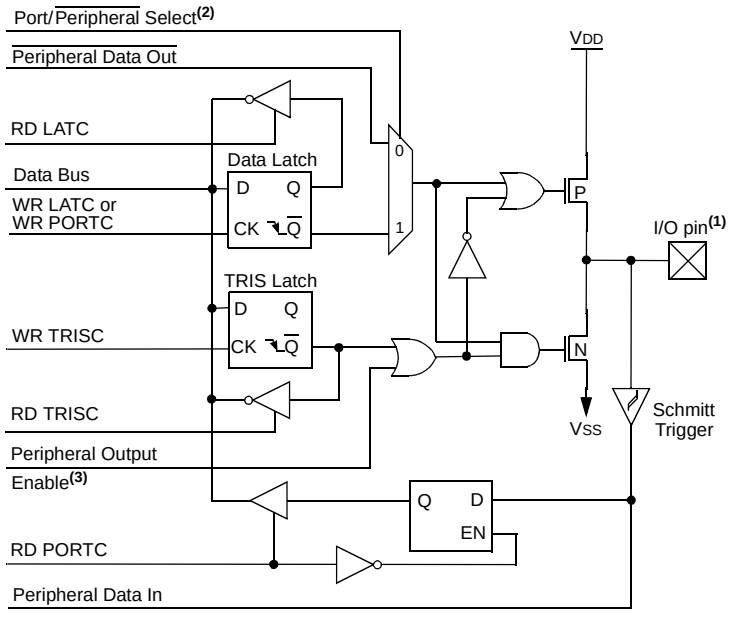
EXAMPLE 9-3: INITIALIZING PORTC

```
CLRF    PORTC    ; Initialize PORTC by
                 ; clearing output
                 ; data latches
CLRF    LATC     ; Alternate method
                 ; to clear output
                 ; data latches
MOVLW 0xC9       ; Value used to
                 ; initialize data
                 ; direction
MOVWF TRISC      ; Set RC<3>,RC<0> as inputs,
                 ; RC<5:4> as outputs, and
                 ; RC<7:6> as inputs
```

PIC18FXX39 devices differ from other PIC18 microcontrollers in allocation of PORTC pins. For most PIC18 devices, PORTC is an 8-bit-wide port. For the PIC18FXX39 family, two of the PORTC pins (RC1 and RC2) are re-allocated as PWM output only pins for use with the Motor Control kernel. To maintain pinout compatibility with other PIC® devices, the remaining PORTC pins are assigned in a manner consistent with other PIC18 devices. For this reason, PORTC has pins RC0 and RC3 through RC7, but not RC1 and RC2.

To maintain compatibility with PIC18FXX2 devices, the individual port and corresponding latch and direction bits for RC1 and RC2 are present in the appropriate registers, but are not available to the user. To avoid erratic device operation, the values of these bits should not be modified.

FIGURE 9-7: PORTC BLOCK DIAGRAM (PERIPHERAL OUTPUT OVERRIDE)



Note 1: I/O pins have diode protection to VDD and VSS.

2: Port/Peripheral Select signal selects between port data (input) and peripheral output.

3: Peripheral Output Enable is only active if Peripheral Select is active.

unsigned char ProMPT_GetBoostTime()

Resources used: 1 stack level

Range of values: 0 to 255

Description: Returns the time in seconds for Boost mode.

unsigned char ProMPT_GetDecelRate()

Resources used: 1 stack level

Range of values: 0 to 255

Description: Returns the current deceleration rate in Hz/second.

unsigned char ProMPT_GetFrequency(void)

Resources used: 1 stack level

Range of values: 0 to 127

Description: Returns the current output frequency in Hz. This may not be the frequency commanded due to Boost or Accel/Decel logic.

unsigned char ProMPT_GetModulation(void)

Resources used: Hardware Multiplier; 1 stack level

Range of values: 0 to 200

Description: Returns the current output modulation in %.

unsigned char ProMPT_GetParameter(unsigned char parameter)

Resources used: 1 stack level

Description: In addition to its pre-defined API methods, the ProMPT kernel allows the user to custom define up to 16 functions for control or communication purposes not covered by the ProMPT APIs. These parameters are used to communicate with motor control GUI evaluation tools, such as Microchip's DashDriveMP™. This method returns the current value of any one of the parameters.

unsigned char ProMPT_GetVFCurve(unsigned char point)

Resources used: Hardware Multiplier; 1 stack level

Description: This function returns one of the 17 modulation values (in %) of the V/F curve. Each point represents a frequency increment of 8 Hz, ranging from point 0 (0 Hz) to point 16 (128 Hz).

void ProMPT_Init(unsigned char PWMfrequency)

Resources used: 64 Bytes RAM; Timer2; PWM1 and PWM2; High Priority Interrupt Vector; Hardware Multiplier; fast call/return; FSR 0; TBLPTR; 2 stack levels

PWMfrequency values: 0 or 1

Description: This function must be called before all other ProMPT methods, and it must be called only once. This routine configures Timer2 and the PWM outputs.

When `PWMfrequency` is '0', the module's operating frequency is 9.75 kHz. When `PWMfrequency` is '1', the module's operating frequency is 19.53 kHz.

Note: Since the high priority interrupt is used, the fast call/return cannot be used by other routines.

16.4.4 CLOCK STRETCHING

Both 7- and 10-bit Slave modes implement automatic clock stretching during a transmit sequence.

The SEN bit (SSPCON2<0>) allows clock stretching to be enabled during receives. Setting SEN will cause the SCL pin to be held low at the end of each data receive sequence.

16.4.4.1 Clock Stretching for 7-bit Slave Receive Mode (SEN = 1)

In 7-bit Slave Receive mode, on the falling edge of the ninth clock at the end of the ACK sequence, if the BF bit is set, the CKP bit in the SSPCON1 register is automatically cleared, forcing the SCL output to be held low. The CKP being cleared to '0' will assert the SCL line low. The CKP bit must be set in the user's ISR before reception is allowed to continue. By holding the SCL line low, the user has time to service the ISR and read the contents of the SSPBUF before the master device can initiate another receive sequence. This will prevent buffer overruns from occurring (see Figure 16-13).

Note 1: If the user reads the contents of the SSPBUF before the falling edge of the ninth clock, thus clearing the BF bit, the CKP bit will not be cleared and clock stretching will not occur.

2: The CKP bit can be set in software, regardless of the state of the BF bit. The user should be careful to clear the BF bit in the ISR before the next receive sequence, in order to prevent an overflow condition.

16.4.4.2 Clock Stretching for 10-bit Slave Receive Mode (SEN = 1)

In 10-bit Slave Receive mode, during the address sequence, clock stretching automatically takes place but CKP is not cleared. During this time, if the UA bit is set after the ninth clock, clock stretching is initiated. The UA bit is set after receiving the upper byte of the 10-bit address, and following the receive of the second byte of the 10-bit address with the R/$\overline{W}$ bit cleared to '0'. The release of the clock line occurs upon updating SSPADD. Clock stretching will occur on each data receive sequence, as described in 7-bit mode.

Note: If the user polls the UA bit and clears it by updating the SSPADD register before the falling edge of the ninth clock occurs, and if the user hasn't cleared the BF bit by reading the SSPBUF register before that time, then the CKP bit will still NOT be asserted low. Clock stretching on the basis of the state of the BF bit only occurs during a data sequence, not an address sequence.

16.4.4.3 Clock Stretching for 7-bit Slave Transmit Mode

7-bit Slave Transmit mode implements clock stretching by clearing the CKP bit after the falling edge of the ninth clock, if the BF bit is clear. This occurs, regardless of the state of the SEN bit.

The user's ISR must set the CKP bit before transmission is allowed to continue. By holding the SCL line low, the user has time to service the ISR and load the contents of the SSPBUF before the master device can initiate another transmit sequence (see Figure 16-9).

Note 1: If the user loads the contents of SSPBUF, setting the BF bit before the falling edge of the ninth clock, the CKP bit will not be cleared and clock stretching will not occur.

2: The CKP bit can be set in software, regardless of the state of the BF bit.

16.4.4.4 Clock Stretching for 10-bit Slave Transmit Mode

In 10-bit Slave Transmit mode, clock stretching is controlled during the first two address sequences by the state of the UA bit, just as it is in 10-bit Slave Receive mode. The first two addresses are followed by a third address sequence, which contains the high order bits of the 10-bit address and the R/$\overline{W}$ bit set to '1'. After the third address sequence is performed, the UA bit is not set, the module is now configured in Transmit mode, and clock stretching is controlled by the BF flag, as in 7-bit Slave Transmit mode (see Figure 16-11).

16.4.9 I²C MASTER MODE REPEATED START CONDITION TIMING

A Repeated START condition occurs when the RSEN bit (SSPCON2<1>) is programmed high and the I²C logic module is in the IDLE state. When the RSEN bit is set, the SCL pin is asserted low. When the SCL pin is sampled low, the baud rate generator is loaded with the contents of SSPADD<5:0> and begins counting. The SDA pin is released (brought high) for one baud rate generator count (TBRG). When the baud rate generator times out, if SDA is sampled high, the SCL pin will be de-asserted (brought high). When SCL is sampled high, the baud rate generator is reloaded with the contents of SSPADD<6:0> and begins counting. SDA and SCL must be sampled high for one TBRG. This action is then followed by assertion of the SDA pin (SDA = 0) for one TBRG, while SCL is high. Following this, the RSEN bit (SSPCON2<1>) will be automatically cleared and the baud rate generator will not be reloaded, leaving the SDA pin held low. As soon as a START condition is detected on the SDA and SCL pins, the S bit (SSPSTAT<3>) will be set. The SSPIF bit will not be set until the baud rate generator has timed out.

Note 1: If RSEN is programmed while any other event is in progress, it will not take effect.

2: A bus collision during the Repeated START condition occurs if:

- SDA is sampled low when SCL goes from low to high.
- SCL goes low before SDA is asserted low. This may indicate that another master is attempting to transmit a data "1".

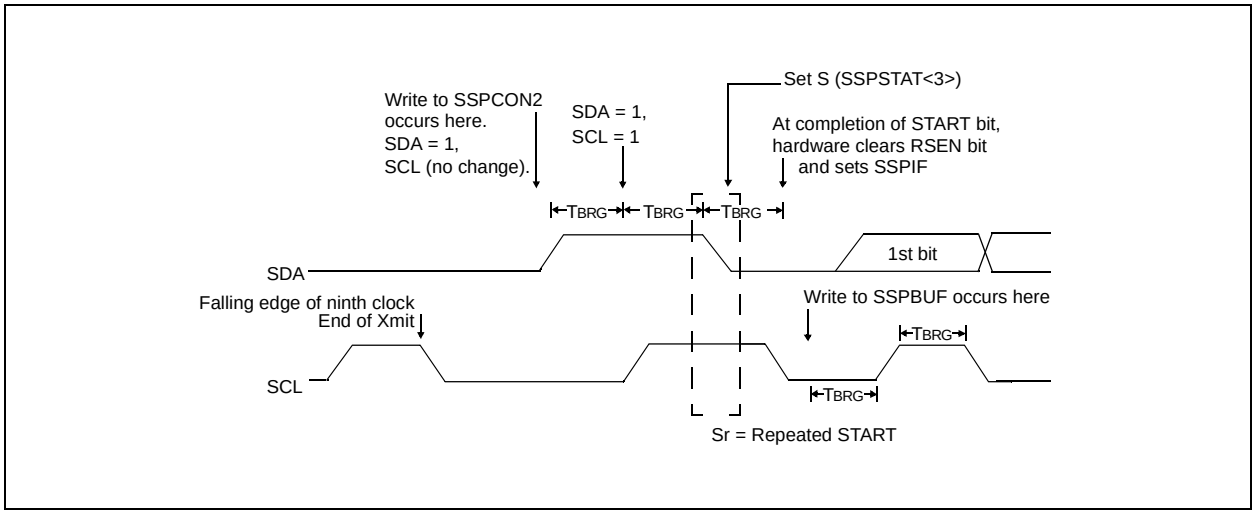
Immediately following the SSPIF bit getting set, the user may write the SSPBUF with the 7-bit address in 7-bit mode, or the default first address in 10-bit mode. After the first eight bits are transmitted and an ACK is received, the user may then transmit an additional eight bits of address (10-bit mode) or eight bits of data (7-bit mode).

16.4.9.1 WCOL Status Flag

If the user writes the SSPBUF when a Repeated START sequence is in progress, the WCOL is set and the contents of the buffer are unchanged (the write doesn't occur).

Note: Because queueing of events is not allowed, writing of the lower 5 bits of SSPCON2 is disabled until the Repeated START condition is complete.

FIGURE 16-20: REPEAT START CONDITION WAVEFORM



16.4.17.2 Bus Collision During a Repeated START Condition

During a Repeated START condition, a bus collision occurs if:

a) A low level is sampled on SDA when SCL goes from low level to high level.

b) SCL goes low before SDA is asserted low, indicating that another master is attempting to transmit a data ‘1’.

When the user de-asserts SDA and the pin is allowed to float high, the BRG is loaded with SSPADD<6:0> and counts down to ‘0’. The SCL pin is then de-asserted, and when sampled high, the SDA pin is sampled.

If SDA is low, a bus collision has occurred (i.e., another master is attempting to transmit a data ’0’, Figure 16-29). If SDA is sampled high, the BRG is reloaded and begins counting. If SDA goes from high to low before the BRG times out, no bus collision occurs because no two masters can assert SDA at exactly the same time.

If SCL goes from high to low before the BRG times out and SDA has not already been asserted, a bus collision occurs. In this case, another master is attempting to transmit a data ‘1’ during the Repeated START condition, see Figure 16-30.

If, at the end of the BRG time-out, both SCL and SDA are still high, the SDA pin is driven low and the BRG is reloaded and begins counting. At the end of the count, regardless of the status of the SCL pin, the SCL pin is driven low and the Repeated START condition is complete.

FIGURE 16-29: BUS COLLISION DURING A REPEATED START CONDITION (CASE 1)

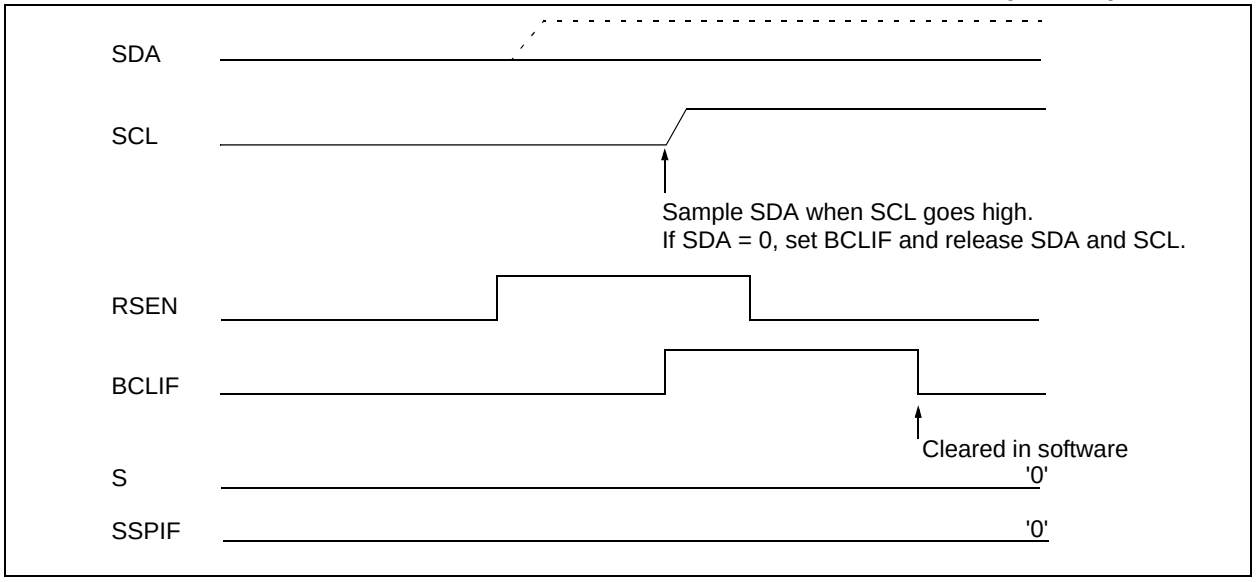
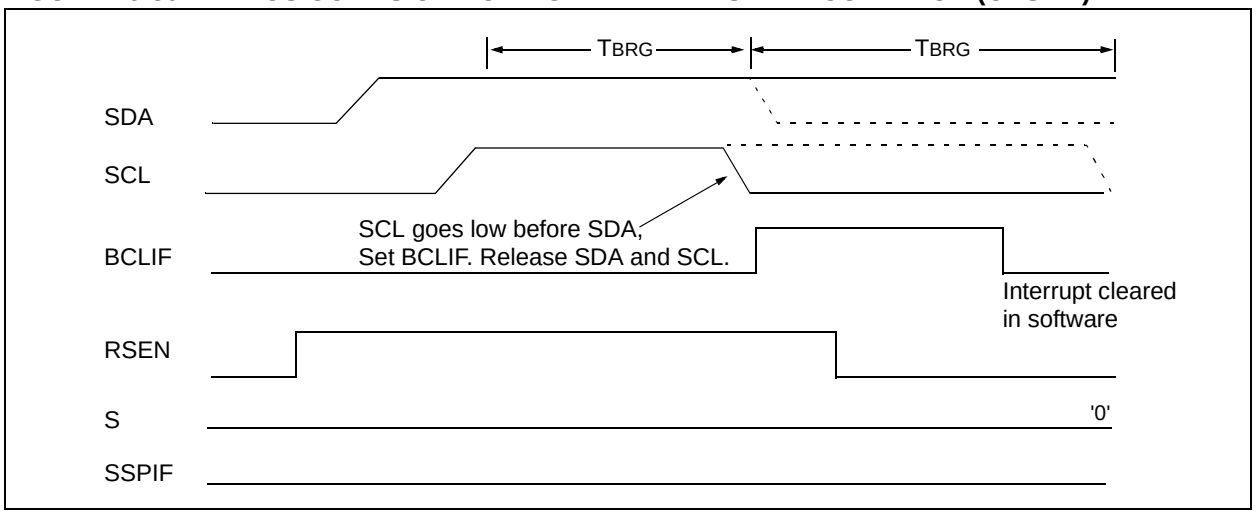


FIGURE 16-30: BUS COLLISION DURING REPEATED START CONDITION (CASE 2)



17.4.2 USART SYNCHRONOUS SLAVE RECEPTION

The operation of the Synchronous Master and Slave modes is identical, except in the case of the SLEEP mode and bit SREN, which is a “don't care” in Slave mode.

If receive is enabled by setting bit CREN prior to the SLEEP instruction, then a word may be received during SLEEP. On completely receiving the word, the RSR register will transfer the data to the RCREG register, and if enable bit RCIE bit is set, the interrupt generated will wake the chip from SLEEP. If the global interrupt is enabled, the program will branch to the interrupt vector.

To set up a Synchronous Slave Reception:

1. Enable the synchronous master serial port by setting bits SYNC and SPEN and clearing bit CSRC.
2. If interrupts are desired, set enable bit RCIE.
3. If 9-bit reception is desired, set bit RX9.
4. To enable reception, set enable bit CREN.
5. Flag bit RCIF will be set when reception is complete. An interrupt will be generated if enable bit RCIE was set.
6. Read the RCSTA register to get the ninth bit (if enabled) and determine if any error occurred during reception.
7. Read the 8-bit received data by reading the RCREG register.
8. If any error occurred, clear the error by clearing bit CREN.
9. If using interrupts, ensure that the GIE and PEIE bits in the INTCON register (INTCON<7:6>) are set.

TABLE 17-11: REGISTERS ASSOCIATED WITH SYNCHRONOUS SLAVE RECEPTION

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	Value on All Other RESETS
INTCON	GIE/ GIEH	PEIE/ GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF	0000 000x	0000 000u
PIR1	PSPIF(1)	ADIF	RCIF	TXIF	SSPIF	—	TMR2IF	TMR1IF	0000 0000	0000 0000
PIE1	PSPIE(1)	ADIE	RCIE	TXIE	SSPIE	—	TMR2IE	TMR1IE	0000 0000	0000 0000
IPR1	PSPIP(1)	ADIP	RCIP	TXIP	SSPIP	—	TMR2IP	TMR1IP	0000 0000	0000 0000
RCSTA	SPEN	RX9	SREN	CREN	ADDEN	FERR	OERR	RX9D	0000 -00x	0000 -00x
RCREG	USART Receive Register								0000 0000	0000 0000
TXSTA	CSRC	TX9	TXEN	SYNC	—	BRGH	TRMT	TX9D	0000 -010	0000 -010
SPBRG	Baud Rate Generator Register								0000 0000	0000 0000

Legend: x = unknown, - = unimplemented, read as '0'.
Shaded cells are not used for Synchronous Slave Reception.

Note 1: The PSPIF, PSPIE and PSPIP bits are reserved on the PIC18F2X39 devices; always maintain these bits clear.

BNOV **Branch if Not Overflow**

Syntax: [*label*] BNOV n

Operands: $-128 \le n \le 127$

Operation: if overflow bit is '0'
(PC) + 2 + 2n → PC

Status Affected: None

Encoding:

1110	0101	nnnn	nnnn

Description: If the Overflow bit is '0', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC+2+2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BNOV Jump`

Before Instruction
PC = address `(HERE)`
After Instruction
If Overflow = 0;
PC = address `(Jump)`
If Overflow = 1;
PC = address `(HERE+2)`

BNZ **Branch if Not Zero**

Syntax: [*label*] BNZ n

Operands: $-128 \le n \le 127$

Operation: if zero bit is '0'
(PC) + 2 + 2n → PC

Status Affected: None

Encoding:

1110	0001	nnnn	nnnn

Description: If the Zero bit is '0', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC+2+2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BNZ Jump`

Before Instruction
PC = address `(HERE)`
After Instruction
If Zero = 0;
PC = address `(Jump)`
If Zero = 1;
PC = address `(HERE+2)`

BTG **Bit Toggle f**

Syntax: [*label*] BTG f,b[,a]

Operands: $0 \le f \le 255$
$0 \le b \le 7$
$a \in [0,1]$

Operation: $(\overline{f<b>}) \rightarrow f<b>$

Status Affected: None

Encoding:

0111	bbba	ffff	ffff

Description: Bit 'b' in data memory location 'f' is inverted. If 'a' is 0, the Access Bank will be selected, overriding the BSR value. If 'a' = 1, then the bank will be selected as per the BSR value (default).

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: `BTG PORTC, 4, 0`

Before Instruction:
PORTC = 0111 0101 [0x75]
After Instruction:
PORTC = 0110 0101 [0x65]

BOV **Branch if Overflow**

Syntax: [*label*] BOV n

Operands: $-128 \le n \le 127$

Operation: if overflow bit is '1'
$(PC) + 2 + 2n \rightarrow PC$

Status Affected: None

Encoding:

1110	0100	nnnn	nnnn

Description: If the Overflow bit is '1', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC+2+2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BOV Jump`

Before Instruction
PC = address (HERE)
After Instruction
If Overflow = 1;
PC = address (Jump)
If Overflow = 0;
PC = address (HERE+2)

INCFSZ **Increment f, skip if 0**

Syntax: [*label*] INCFSZ f [,d [,a]

Operands: $0 \le f \le 255$
$d \in [0,1]$
$a \in [0,1]$

Operation: (f) + 1 → dest,
skip if result = 0

Status Affected: None

Encoding:

0011	11da	ffff	ffff

Description: The contents of register 'f' are incremented. If 'd' is 0, the result is placed in W. If 'd' is 1, the result is placed back in register 'f'. (default) If the result is 0, the next instruction, which is already fetched, is discarded, and a NOP is executed instead, making it a two-cycle instruction. If 'a' is 0, the Access Bank will be selected, overriding the BSR value. If 'a' = 1, then the bank will be selected as per the BSR value (default).

Words: 1

Cycles: 1(2)
Note: 3 cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:
```
HERE    INCFSZ   CNT, 1, 0
NZERO   :
ZERO    :
```

Before Instruction
PC = Address (HERE)

After Instruction
CNT = CNT + 1
If CNT = 0;
PC = Address (ZERO)
If CNT ≠ 0;
PC = Address (NZERO)

INFSNZ **Increment f, skip if not 0**

Syntax: [*label*] INFSNZ f [,d [,a]

Operands: $0 \le f \le 255$
$d \in [0,1]$
$a \in [0,1]$

Operation: (f) + 1 → dest,
skip if result ≠ 0

Status Affected: None

Encoding:

0100	10da	ffff	ffff

Description: The contents of register 'f' are incremented. If 'd' is 0, the result is placed in W. If 'd' is 1, the result is placed back in register 'f' (default). If the result is not 0, the next instruction, which is already fetched, is discarded, and a NOP is executed instead, making it a two-cycle instruction. If 'a' is 0, the Access Bank will be selected, overriding the BSR value. If 'a' = 1, then the bank will be selected as per the BSR value (default).

Words: 1

Cycles: 1(2)
Note: 3 cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:
```
HERE    INFSNZ  REG, 1, 0
ZERO
NZERO
```

Before Instruction
PC = Address (HERE)

After Instruction
REG = REG + 1
If REG ≠ 0;
PC = Address (NZERO)
If REG = 0;
PC = Address (ZERO)

MULLW **Multiply Literal with W**

Syntax: [*label*] MULLW k

Operands: $0 \le k \le 255$

Operation: (W) x k → PRODH:PRODL

Status Affected: None

Encoding:

0000	1101	kkkk	kkkk

Description: An unsigned multiplication is carried out between the contents of W and the 8-bit literal 'k'. The 16-bit result is placed in PRODH:PRODL register pair. PRODH contains the high byte. W is unchanged.
None of the status flags are affected.
Note that neither overflow nor carry is possible in this operation. A zero result is possible but not detected.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write registers PRODH: PRODL

Example: `MULLW 0xC4`

Before Instruction
W = 0xE2
PRODH = ?
PRODL = ?

After Instruction
W = 0xE2
PRODH = 0xAD
PRODL = 0x08

MULWF **Multiply W with f**

Syntax: [*label*] MULWF f [,a]

Operands: $0 \le f \le 255$
$a \in [0,1]$

Operation: (W) x (f) → PRODH:PRODL

Status Affected: None

Encoding:

0000	001a	ffff	ffff

Description: An unsigned multiplication is carried out between the contents of W and the register file location 'f'. The 16-bit result is stored in the PRODH:PRODL register pair. PRODH contains the high byte. Both W and 'f' are unchanged.
None of the status flags are affected.
Note that neither overflow nor carry is possible in this operation. A zero result is possible but not detected. If 'a' is 0, the Access Bank will be selected, overriding the BSR value. If 'a' = 1, then the bank will be selected as per the BSR value (default).

Words: 1

Cycles: 1

Q Cycle Activity:

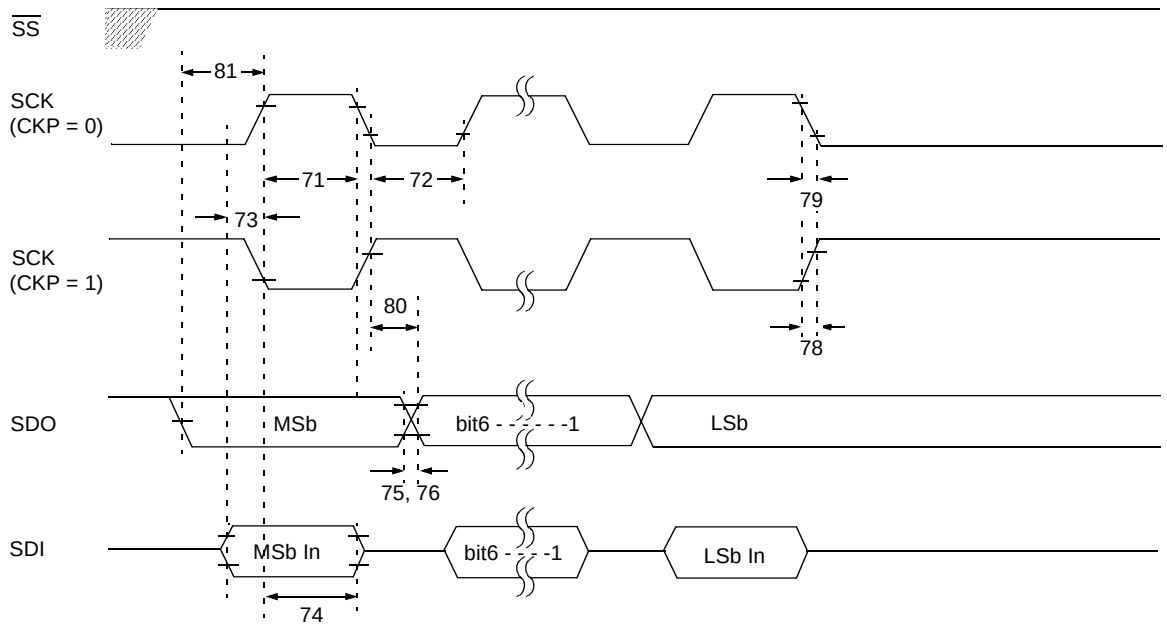
Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write registers PRODH: PRODL

Example: `MULWF REG, 1`

Before Instruction
W = 0xC4
REG = 0xB5
PRODH = ?
PRODL = ?

After Instruction
W = 0xC4
REG = 0xB5
PRODH = 0x8A
PRODL = 0x94

FIGURE 23-13: EXAMPLE SPI MASTER MODE TIMING (CKE = 1)



Note: Refer to Figure 23-4 for load conditions.

TABLE 23-12: EXAMPLE SPI MODE REQUIREMENTS (MASTER MODE, CKE = 1)

Param. No.	Symbol	Characteristic		Min	Max	Units	Conditions
71	TscH	SCK input high time (Slave mode)	Continuous	1.25 TCY + 30	—	ns	
71A			Single Byte	40	—	ns	**(Note 1)**
72	TscL	SCK input low time (Slave mode)	Continuous	1.25 TCY + 30	—	ns	
72A			Single Byte	40	—	ns	**(Note 1)**
73	TdiV2scH, TdiV2scL	Setup time of SDI data input to SCK edge		100	—	ns	
73A	TB2B	Last clock edge of Byte 1 to the 1st clock edge of Byte 2		1.5 TCY + 40	—	ns	**(Note 2)**
74	TscH2diL, TscL2diL	Hold time of SDI data input to SCK edge		100	—	ns	
75	TdoR	SDO data output rise time	PIC18**F**XXXX	—	25	ns	
			PIC18**LF**XXXX	—	60	ns	VDD = 2V
76	TdoF	SDO data output fall time	PIC18**F**XXXX	—	25	ns	
			PIC18**LF**XXXX	—	60	ns	VDD = 2V
78	TscR	SCK output rise time (Master mode)	PIC18**F**XXXX	—	25	ns	
			PIC18**LF**XXXX	—	60	ns	VDD = 2V
79	TscF	SCK output fall time (Master mode)	PIC18**F**XXXX	—	25	ns	
			PIC18**LF**XXXX	—	60	ns	VDD = 2V
80	TscH2doV, TscL2doV	SDO data output valid after SCK edge	PIC18**F**XXXX	—	50	ns	
			PIC18**LF**XXXX	—	150	ns	VDD = 2V
81	TdoV2scH, TdoV2scL	SDO data output setup to SCK edge		TCY	—	ns	

Note 1: Requires the use of Parameter # 73A.
2: Only if Parameter # 71A and # 72A are used.

FIGURE 24-11: ΔILVD vs. VDD OVER TEMPERATURE (LVD ENABLED, VLVD = 4.5 - 4.78V)

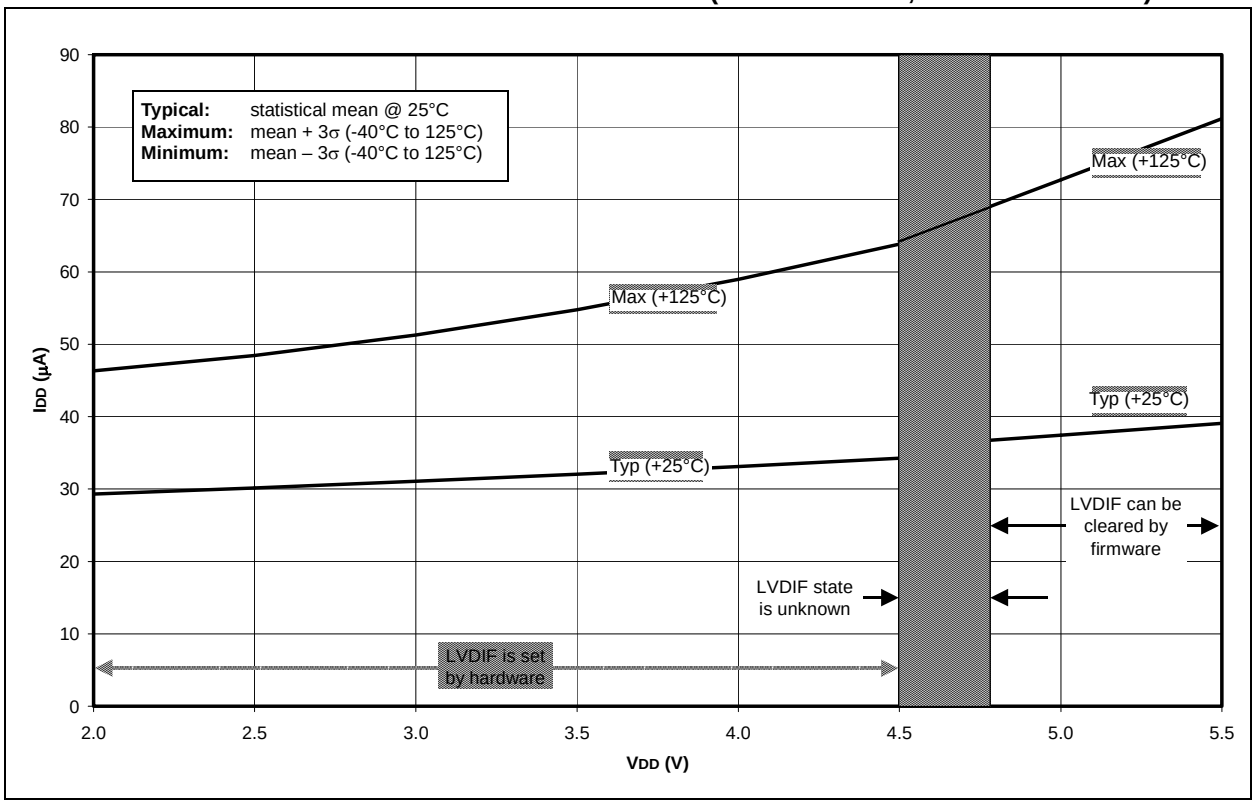


FIGURE 24-12: TYPICAL, MINIMUM AND MAXIMUM VOH vs. IOH (VDD = 5V, -40°C TO +125°C)

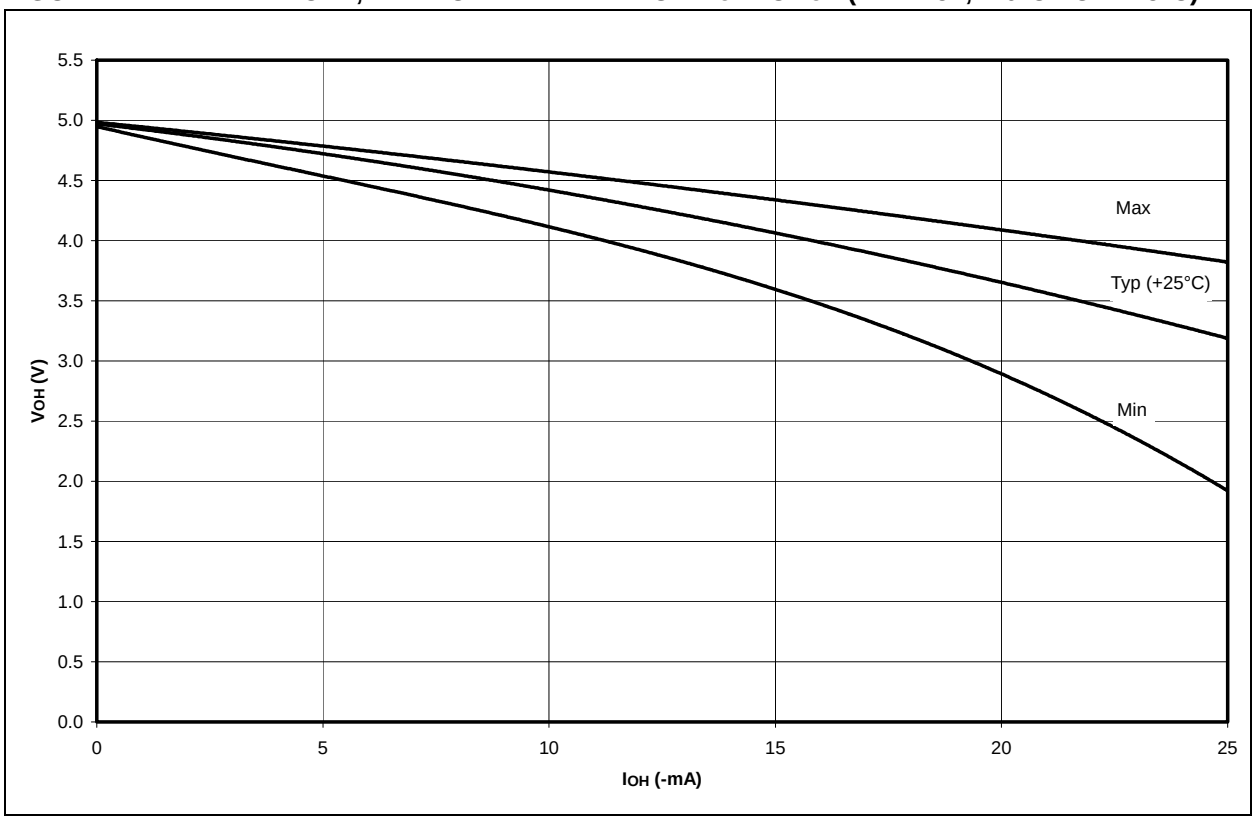


FIGURE 24-17: MINIMUM AND MAXIMUM VIN vs. VDD (TTL INPUT, -40°C TO +125°C)

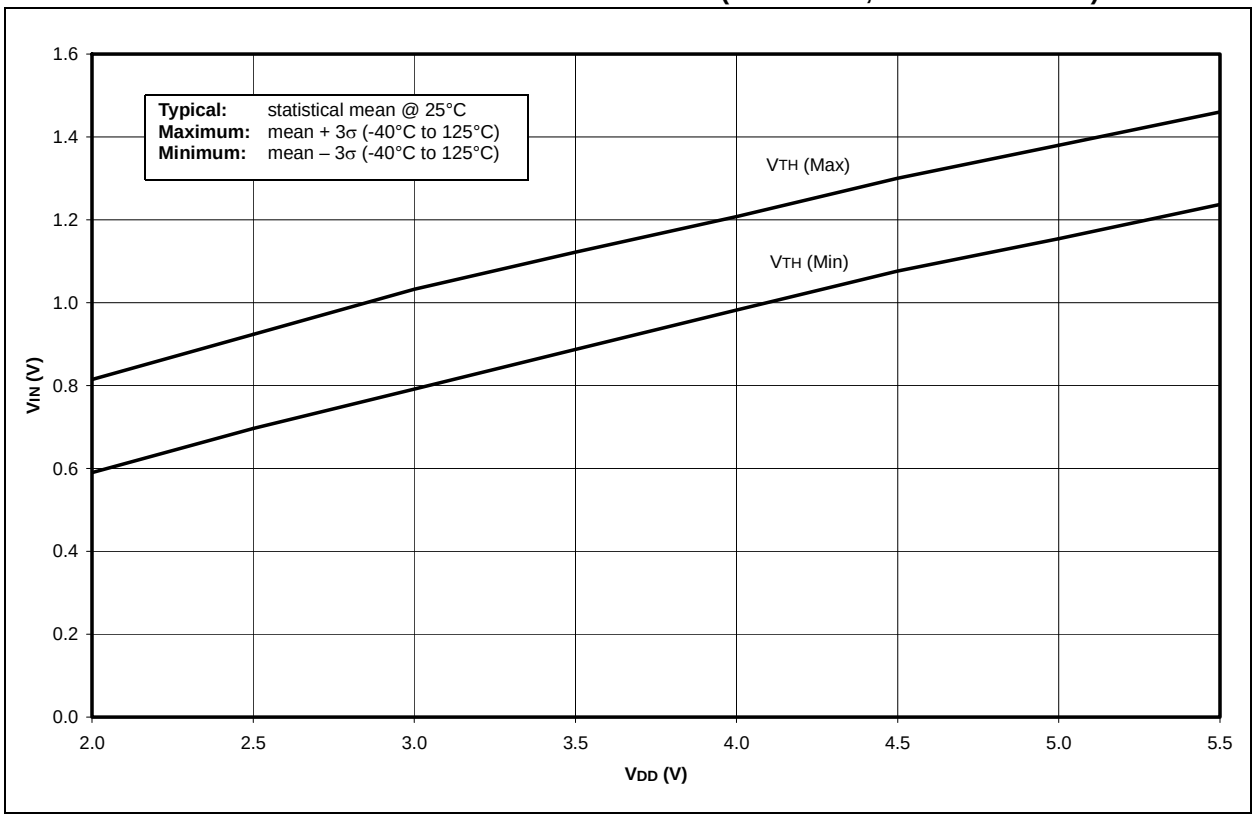
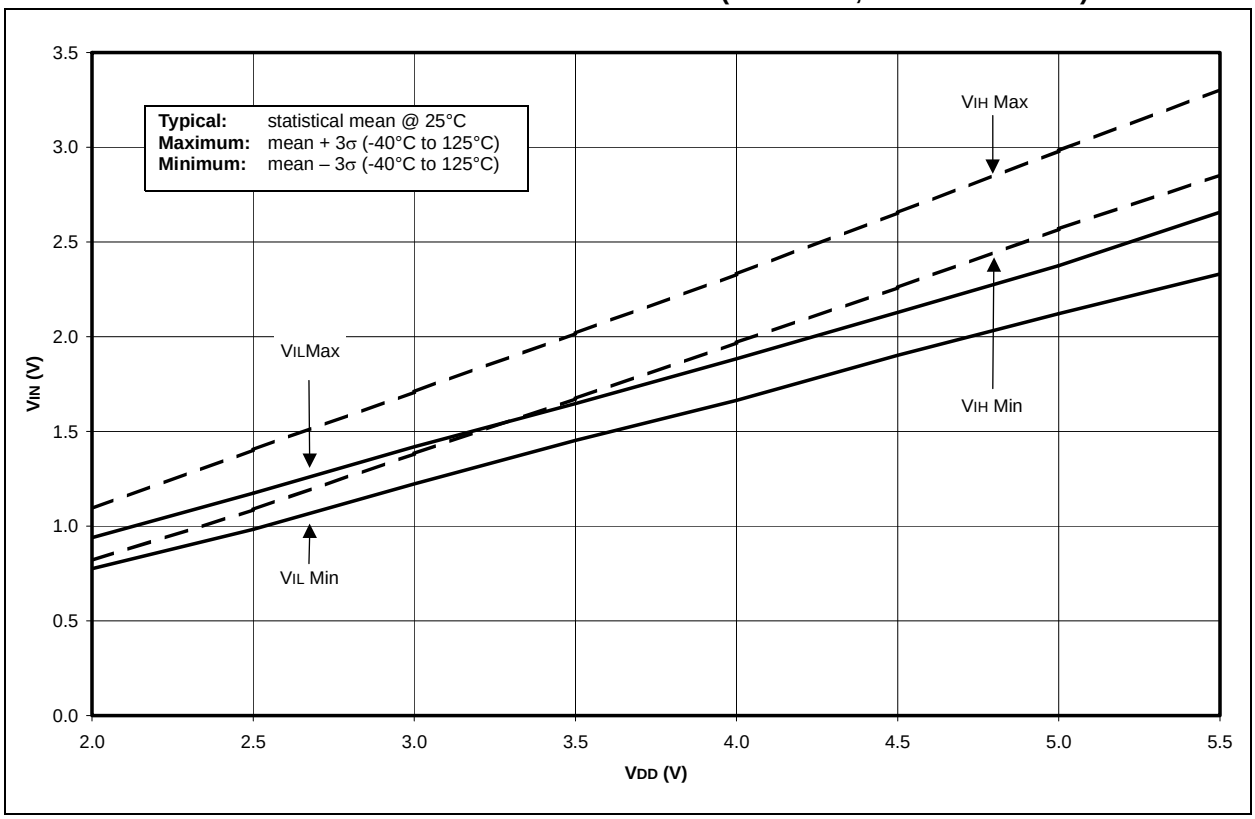


FIGURE 24-18: MINIMUM AND MAXIMUM VIN vs. VDD (I^2C INPUT, -40°C TO +125°C)



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