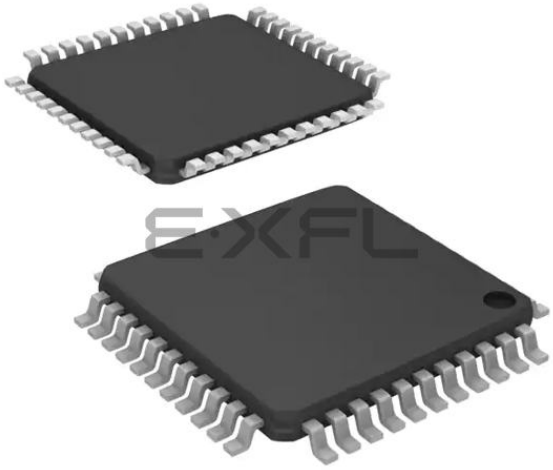

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Applications of "Embedded - Microcontrollers"

Details	
Product Status	Obsolete
Core Processor	PIC
Core Size	8-Bit
Speed	40MHz
Connectivity	I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, LVD, POR, PWM, WDT
Number of I/O	32
Program Memory Size	12KB (6K x 16)
Program Memory Type	FLASH
EEPROM Size	256 x 8
RAM Size	640 x 8
Voltage - Supply (Vcc/Vdd)	4.2V ~ 5.5V
Data Converters	A/D 8x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 125°C (TA)
Mounting Type	Surface Mount
Package / Case	44-TQFP
Supplier Device Package	44-TQFP (10x10)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18f4439t-e-pt

TABLE 1-3: PIC18F4X39 PINOUT I/O DESCRIPTIONS (CONTINUED)

Pin Name	Pin Number			Pin Type	Buffer Type	Description
	DIP	QFN	TQFP			
						PORTD is a bi-directional I/O port, or a Parallel Slave Port (PSP) for interfacing to a microprocessor port. These pins have TTL input buffers when PSP module is enabled.
RD0/PSP0	19	38	38	I/O		
RD0					ST	Digital I/O.
PSP0					TTL	Parallel Slave Port Data.
RD1/PSP1	20	39	39	I/O		
RD1					ST	Digital I/O.
PSP1					TTL	Parallel Slave Port Data.
RD2/PSP2	21	40	40	I/O		
RD2					ST	Digital I/O.
PSP2					TTL	Parallel Slave Port Data.
RD3/PSP3	22	41	41	I/O		
RD3					ST	Digital I/O.
PSP3					TTL	Parallel Slave Port Data.
RD4/PSP4	27	2	2	I/O		
RD4					ST	Digital I/O.
PSP4					TTL	Parallel Slave Port Data.
RD5/PSP5	28	3	3	I/O		
RD5					ST	Digital I/O.
PSP5					TTL	Parallel Slave Port Data.
RD6/PSP6	29	4	4	I/O		
RD6					ST	Digital I/O.
PSP6					TTL	Parallel Slave Port Data.
RD7/PSP7	30	5	5	I/O		
RD7					ST	Digital I/O.
PSP7					TTL	Parallel Slave Port Data.

Legend: TTL = TTL compatible input
CMOS = CMOS compatible input or output
ST = Schmitt Trigger input with CMOS levels
I = Input
O = Output
P = Power
OD = Open Drain (no P diode to V_{DD})

REGISTER 4-1: STKPTR REGISTER

R/C-0	R/C-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
STKFUL	STKUNF	—	SP4	SP3	SP2	SP1	SP0
bit 7							bit 0

bit 7(1) **STKFUL:** Stack Full Flag bit
1 = Stack became full or overflowed
0 = Stack has not become full or overflowed

bit 6(1) **STKUNF:** Stack Underflow Flag bit
1 = Stack underflow occurred
0 = Stack underflow did not occur

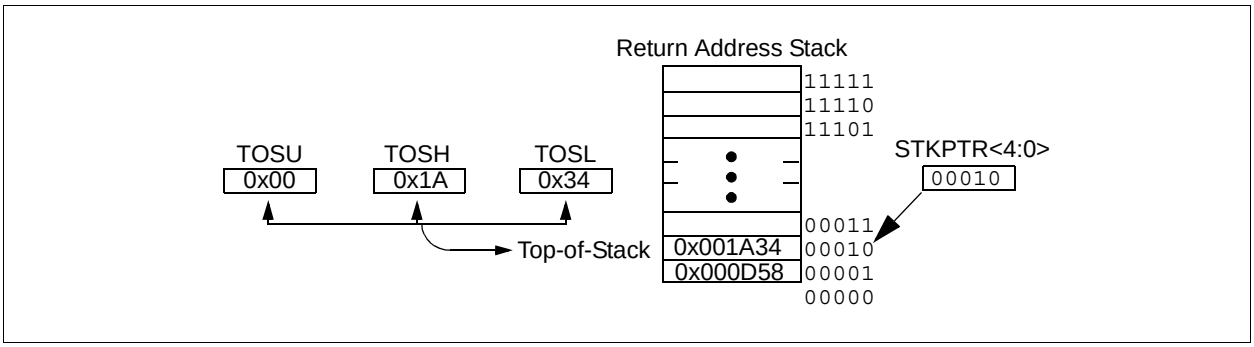
bit 5 **Unimplemented:** Read as '0'

bit 4-0 **SP4:SP0:** Stack Pointer Location bits

Note 1: Bit 7 and bit 6 can only be cleared in user software or by a POR.

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
- n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

FIGURE 4-2: RETURN ADDRESS STACK AND ASSOCIATED REGISTERS



4.2.3 PUSH AND POP INSTRUCTIONS

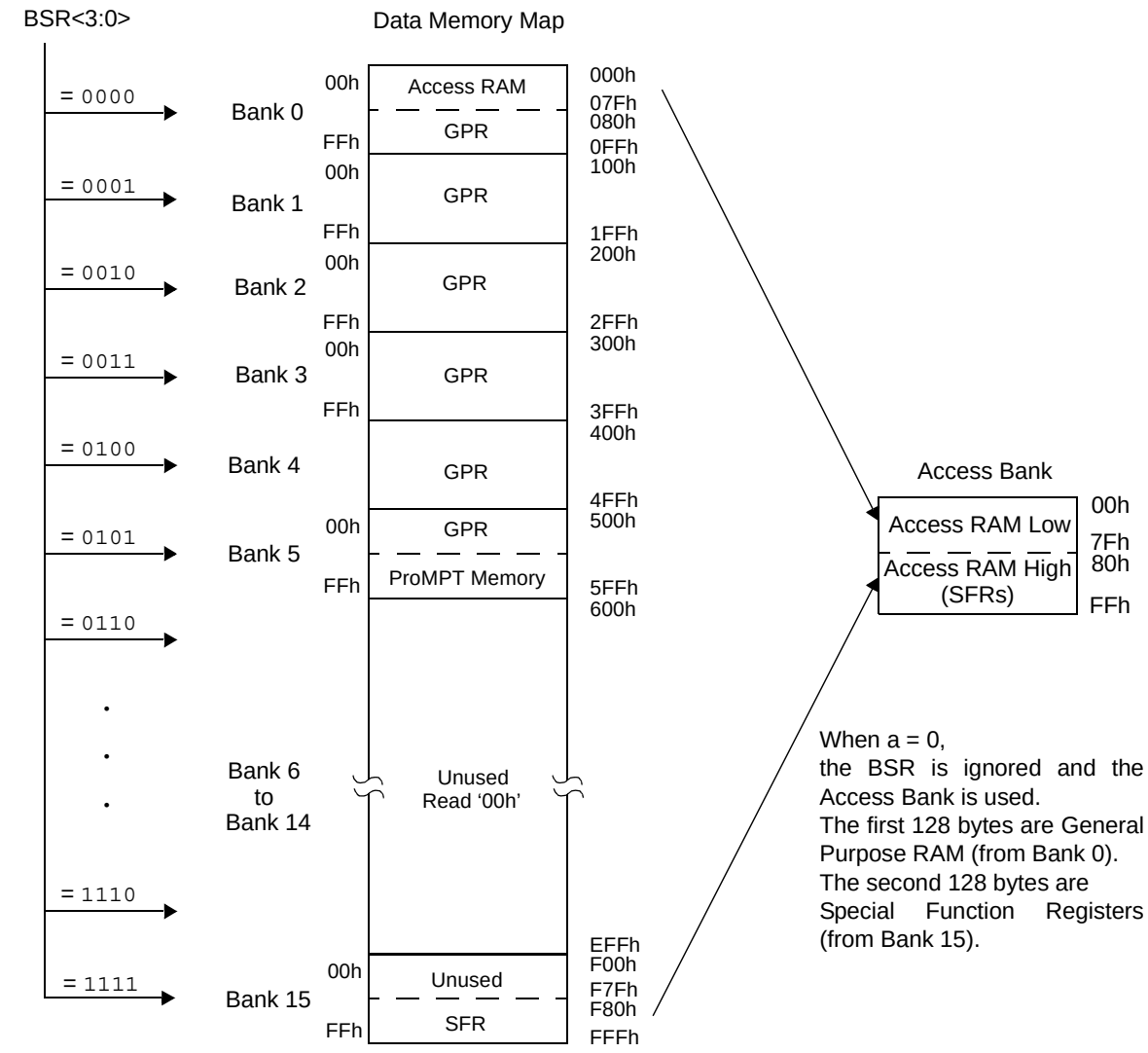
Since the Top-of-Stack (TOS) is readable and writable, the ability to push values onto the stack and pull values off the stack, without disturbing normal program execution, is a desirable option. To push the current PC value onto the stack, a PUSH instruction can be executed. This will increment the stack pointer and load the current PC value onto the stack. TOSU, TOSH and TOSL can then be modified to place a return address on the stack.

The ability to pull the TOS value off of the stack and replace it with the value that was previously pushed onto the stack, without disturbing normal execution, is achieved by using the POP instruction. The POP instruction discards the current TOS by decrementing the stack pointer. The previous value pushed onto the stack then becomes the TOS value.

4.2.4 STACK FULL/UNDERFLOW RESETS

These RESETS are enabled by programming the STVREN configuration bit. When the STVREN bit is disabled, a full or underflow condition will set the appropriate STKFUL or STKUNF bit, but not cause a device RESET. When the STVREN bit is enabled, a full or underflow condition will set the appropriate STKFUL or STKUNF bit and then cause a device RESET. The STKFUL or STKUNF bits are only cleared by the user software or a POR Reset.

FIGURE 4-6: DATA MEMORY MAP FOR PIC18FX539



REGISTER 5-1: EECON1 REGISTER (ADDRESS FA6h)

R/W-x	R/W-x	U-0	R/W-0	R/W-x	R/W-0	R/S-0	R/S-0
EEPGD	CFGS	—	FREE	WRERR	WREN	WR	RD
bit 7							bit 0

bit 7 **EEPGD:** FLASH Program or Data EEPROM Memory Select bit
1 = Access FLASH program memory
0 = Access data EEPROM memory

bit 6 **CFGS:** FLASH Program/Data EE or Configuration Select bit
1 = Access configuration registers
0 = Access FLASH program or data EEPROM memory

bit 5 **Unimplemented:** Read as '0'

bit 4 **FREE:** FLASH Row Erase Enable bit
1 = Erase the program memory row addressed by TBLPTR on the next WR command (cleared by completion of erase operation)
0 = Perform write only

bit 3 **WRERR:** FLASH Program/Data EE Error Flag bit
1 = A write operation is prematurely terminated (any RESET during self-timed programming in normal operation)
0 = The write operation completed

Note: When a WRERR occurs, the EEPGD and CFGS bits are not cleared. This allows tracing of the error condition.

bit 2 **WREN:** FLASH Program/Data EE Write Enable bit
1 = Allows write cycles
0 = Inhibits write to the EEPROM

bit 1 **WR:** Write Control bit
1 = Initiates a data EEPROM erase/write cycle or a program memory erase cycle or write cycle. (The operation is self-timed and the bit is cleared by hardware once write is complete. The WR bit can only be set (not cleared) in software.)
0 = Write cycle to the EEPROM is complete

bit 0 **RD:** Read Control bit
1 = Initiates an EEPROM read
(Read takes one cycle. RD is cleared in hardware. The RD bit can only be set (not cleared) in software. RD bit cannot be set when EEPGD = 1.)
0 = Does not initiate an EEPROM read

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as ‘0’	
- n = Value at POR	‘1’ = Bit is set	‘0’ = Bit is cleared	x = Bit is unknown

REGISTER 8-5: PIR2: PERIPHERAL INTERRUPT REQUEST (FLAG) REGISTER 2

U-0	U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	U-0
—	—	—	EEIF	BCLIF	LVDIF	TMR3IF	—
bit 7							bit 0

bit 7-5 **Unimplemented:** Read as '0'

bit 4 **EEIF**: Data EEPROM/FLASH Write Operation Interrupt Flag bit
1 = The write operation is complete (must be cleared in software)
0 = The write operation is not complete, or has not been started

bit 3 **BCLIF**: Bus Collision Interrupt Flag bit
1 = A bus collision occurred (must be cleared in software)
0 = No bus collision occurred

bit 2 **LVDIF**: Low Voltage Detect Interrupt Flag bit
1 = A low voltage condition occurred (must be cleared in software)
0 = The device voltage is above the Low Voltage Detect trip point

bit 1 **TMR3IF**: TMR3 Overflow Interrupt Flag bit
1 = TMR3 register overflowed (must be cleared in software)
0 = TMR3 register did not overflow

bit 0 **Unimplemented:** Read as '0'

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
- n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

8.3 PIE Registers

The PIE registers contain the individual enable bits for the peripheral interrupts. Due to the number of peripheral interrupt sources, there are two Peripheral Interrupt Enable registers (PIE1, PIE2). When IPEN = 0, the PEIE bit must be set to enable any of these peripheral interrupts.

REGISTER 8-6: PIE1: PERIPHERAL INTERRUPT ENABLE REGISTER 1

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	U-0	R/W-0	R/W-0
PSPIE(1)	ADIE	RCIE	TXIE	SSPIE	—	TMR2IE(2)	TMR1IE
bit 7							bit 0

bit 7 **PSPIE(1):** Parallel Slave Port Read/Write Interrupt Enable bit
1 = Enables the PSP read/write interrupt
0 = Disables the PSP read/write interrupt

bit 6 **ADIE**: A/D Converter Interrupt Enable bit
1 = Enables the A/D interrupt
0 = Disables the A/D interrupt

bit 5 **RCIE**: USART Receive Interrupt Enable bit
1 = Enables the USART receive interrupt
0 = Disables the USART receive interrupt

bit 4 **TXIE**: USART Transmit Interrupt Enable bit
1 = Enables the USART transmit interrupt
0 = Disables the USART transmit interrupt

bit 3 **SSPIE**: Master Synchronous Serial Port Interrupt Enable bit
1 = Enables the MSSP interrupt
0 = Disables the MSSP interrupt

bit 2 **Unimplemented:** Read as '0'

bit 1 **TMR2IE(2):** TMR2 to PR2 Match Interrupt Enable bit
1 = Enables the TMR2 to PR2 match interrupt
0 = Disables the TMR2 to PR2 match interrupt

bit 0 **TMR1IE**: TMR1 Overflow Interrupt Enable bit
1 = Enables the TMR1 overflow interrupt
0 = Disables the TMR1 overflow interrupt

Note 1: This bit is reserved on PIC18F2X39 devices; always maintain this bit clear.
2: This bit is reserved for use by the ProMPT kernel; do not alter its value.

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
- n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

TABLE 9-1: PORTA FUNCTIONS

Name	Bit#	Buffer	Function
RA0/AN0	bit0	TTL	Input/output or analog input.
RA1/AN1	bit1	TTL	Input/output or analog input.
RA2/AN2/VREF-	bit2	TTL	Input/output or analog input or VREF-.
RA3/AN3/VREF+	bit3	TTL	Input/output or analog input or VREF+.
RA4/T0CKI	bit4	ST	Input/output or external clock input for Timer0. Output is open drain type.
RA5/$\overline{SS}$/AN4/LVDIN	bit5	TTL	Input/output or slave select input for synchronous serial port or analog input, or low voltage detect input.
OSC2/CLKO/RA6	bit6	TTL	OSC2 or clock output or I/O pin.

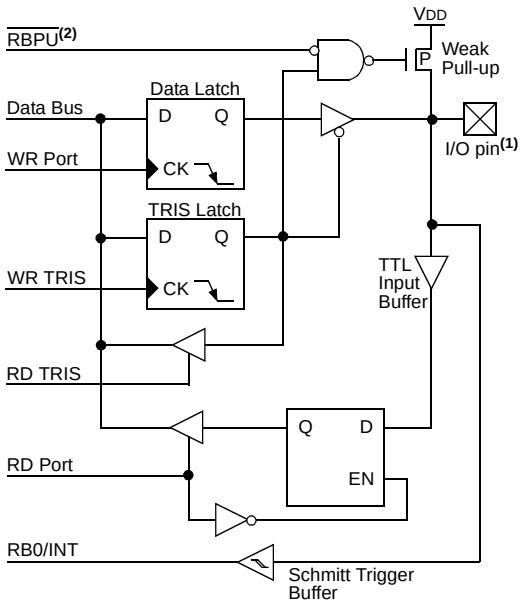
Legend: TTL = TTL input, ST = Schmitt Trigger input

TABLE 9-2: SUMMARY OF REGISTERS ASSOCIATED WITH PORTA

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	Value on All Other RESETS
PORTA	—	RA6	RA5	RA4	RA3	RA2	RA1	RA0	-x0x 0000	-u0u 0000
LATA	—	LATA Data Output Register							-xxx xxxx	-uuu uuuu
TRISA	—	PORTA Data Direction Register							-111 1111	-111 1111
ADCON1	ADFM	ADCS2	—	—	PCFG3	PCFG2	PCFG1	PCFG0	00-- 0000	00-- 0000

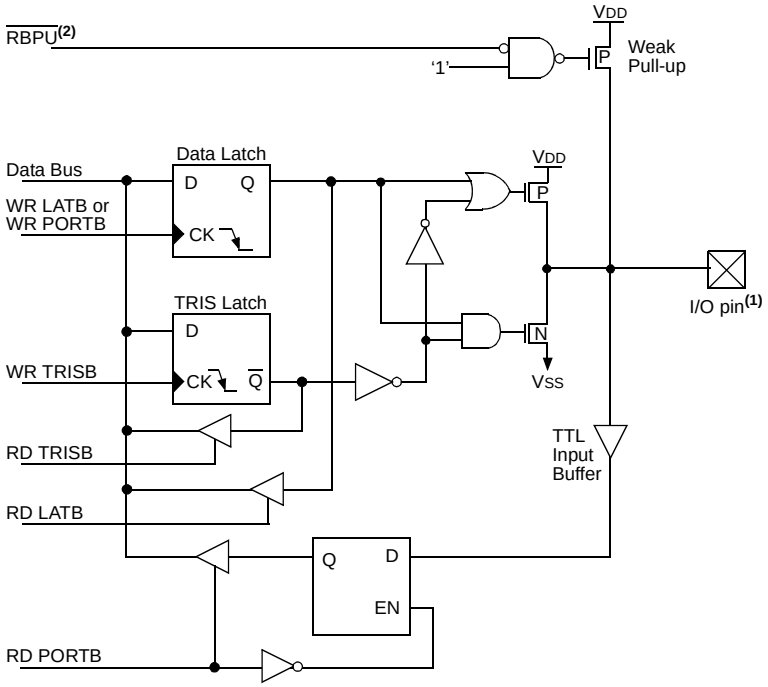
Legend: x = unknown, u = unchanged, - = unimplemented locations read as '0'. Shaded cells are not used by PORTA.

FIGURE 9-5: BLOCK DIAGRAM OF RB2:RB0 PINS



Note 1: I/O pins have diode protection to VDD and VSS.
2: To enable weak pull-ups, set the appropriate TRIS bit(s) and clear the $\overline{\text{RBPU}}$ bit (OPTION_REG<7>).

FIGURE 9-6: BLOCK DIAGRAM OF RB3 PIN



Note 1: I/O pin has diode protection to VDD and VSS.
2: To enable weak pull-ups, set the appropriate DDR bit(s) and clear the $\overline{\text{RBPU}}$ bit (INTCON2<7>).

15.1.2 PWM DUTY CYCLE

The PWM duty cycle is set by the Motor Control module when it writes a 10-bit value to the CCPR1L and CCP1CON registers, where CCPR1L contains the eight Most Significant bits and CCP1CON<5:4> contains the two Least Significant bits. The duty cycle time is given by the equation:

$$\text{PWM duty cycle} = (\text{10-bit CCP register value}) \bullet T_{OSC} \bullet (\text{TMR2 prescale value})$$

where TOSC and the duty cycle are in the same unit of time.

The CCPR1H register and a 2-bit internal latch are used to double-buffer the PWM duty cycle. This buffering is essential for glitchless PWM operation. At the same time, the value of TMR2 is concatenated with either an internal 2-bit Q clock, or 2 bits of the TMR2 prescaler. When the CCPR1H:latch pair value matches that of the TMR2:latch pair, the PWM1 pin is cleared.

The maximum PWM resolution (bits) for a given PWM frequency is given by the equation:

$$\text{PWM Resolution (max)} = \frac{\log\left(\frac{F_{OSC}}{F_{PWM}}\right)}{\log(2)} \text{bits}$$

where FPWM is the PWM frequency, or (1/PWM period).

Note:	If the PWM duty cycle value is longer than the PWM period, the PWM1 pin will not be cleared.

TABLE 15-1: REGISTERS ASSOCIATED WITH PWM AND TIMER2

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	Value on All Other RESETS
INTCON	GIE/GIEH	PEIE/GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF	0000 000x	0000 000u
PIR1	PSPIF(1)	ADIF	RCIF	TXIF	SSPIF	—	TMR2IF	TMR1IF	0000 0000	0000 0000
PIE1	PSPIE(1)	ADIE	RCIE	TXIE	SSPIE	—	TMR2IE	TMR1IE	0000 0000	0000 0000
IPR1	PSPIP(1)	ADIP	RCIP	TXIP	SSPIP	—	TMR2IP	TMR1IP	0000 0000	0000 0000
TMR2*	*	*	*	*	*	*	*	*	0000 0000	0000 0000
PR2*	*	*	*	*	*	*	*	*	1111 1111	1111 1111
T2CON*	*	*	*	*	*	*	*	*	-000 0000	-000 0000
CCPR1L*	*	*	*	*	*	*	*	*	xxxx xxxx	uuuu uuuu
CCPR1H	PWM Register1 (MSB) (read-only)								xxxx xxxx	uuuu uuuu
CCP1CON*	—	—	*	*	*	*	*	*	--00 0000	--00 0000
CCPR2L*	*	*	*	*	*	*	*	*	xxxx xxxx	uuuu uuuu
CCPR2H*	PWM Register2 (MSB) (read-only)								xxxx xxxx	uuuu uuuu
CCP2CON*	—	—	*	*	*	*	*	*	--00 0000	--00 0000

Legend: x = unknown, u = unchanged, - = unimplemented, read as '0' unless otherwise noted. Shaded cells are not used by PWM and Timer2.

* These registers are retained to maintain compatibility with PIC18FXX2 devices; however, the indicated bits are reserved in PIC18FXX39 devices. Users should not alter the values of these bits.

Note 1: The PSPIF, PSPIE and PSPIP bits are reserved on the PIC18F2X39 devices; always maintain these bits clear.

16.0 MASTER SYNCHRONOUS SERIAL PORT (MSSP) MODULE

16.1 Master SSP (MSSP) Module Overview

The Master Synchronous Serial Port (MSSP) module is a serial interface useful for communicating with other peripheral or microcontroller devices. These peripheral devices may be serial EEPROMs, shift registers, display drivers, A/D converters, etc. The MSSP module can operate in one of two modes:

- Serial Peripheral Interface (SPI)
- Inter-Integrated Circuit (I^2C)
  - Full Master mode
  - Slave mode (with general address call)

The I^2C interface supports the following modes in hardware:

- Master mode
- Multi-Master mode
- Slave mode

16.2 Control Registers

The MSSP module has three associated registers. These include a status register (SSPSTAT) and two control registers (SSPCON1 and SSPCON2). The use of these registers and their individual configuration bits differ significantly, depending on whether the MSSP module is operated in SPI or I^2C mode.

Additional details are provided under the individual sections.

16.3 SPI Mode

The SPI mode allows 8 bits of data to be synchronously transmitted and received, simultaneously. All four modes of SPI are supported. To accomplish communication, typically three pins are used:

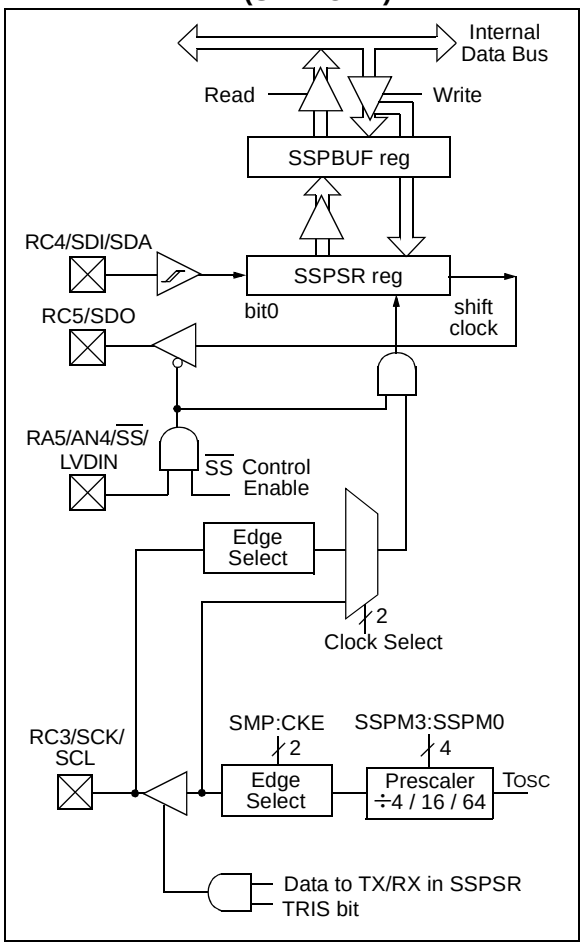
- Serial Data Out (SDO) - RC5/SDO
- Serial Data In (SDI) - RC4/SDI/SDA
- Serial Clock (SCK) - RC3/SCK/SCL

Additionally, a fourth pin may be used when in a Slave mode of operation:

- Slave Select ($\overline{SS}$) - RA5/AN4/$\overline{SS}$/LVDIN

Figure 16-1 shows the block diagram of the MSSP module when operating in SPI mode.

FIGURE 16-1: MSSP BLOCK DIAGRAM (SPI MODE)



16.4.9 I²C MASTER MODE REPEATED START CONDITION TIMING

A Repeated START condition occurs when the RSEN bit (SSPCON2<1>) is programmed high and the I²C logic module is in the IDLE state. When the RSEN bit is set, the SCL pin is asserted low. When the SCL pin is sampled low, the baud rate generator is loaded with the contents of SSPADD<5:0> and begins counting. The SDA pin is released (brought high) for one baud rate generator count (TBRG). When the baud rate generator times out, if SDA is sampled high, the SCL pin will be de-asserted (brought high). When SCL is sampled high, the baud rate generator is reloaded with the contents of SSPADD<6:0> and begins counting. SDA and SCL must be sampled high for one TBRG. This action is then followed by assertion of the SDA pin (SDA = 0) for one TBRG, while SCL is high. Following this, the RSEN bit (SSPCON2<1>) will be automatically cleared and the baud rate generator will not be reloaded, leaving the SDA pin held low. As soon as a START condition is detected on the SDA and SCL pins, the S bit (SSPSTAT<3>) will be set. The SSPIF bit will not be set until the baud rate generator has timed out.

Note 1: If RSEN is programmed while any other event is in progress, it will not take effect.

2: A bus collision during the Repeated START condition occurs if:

- SDA is sampled low when SCL goes from low to high.
- SCL goes low before SDA is asserted low. This may indicate that another master is attempting to transmit a data "1".

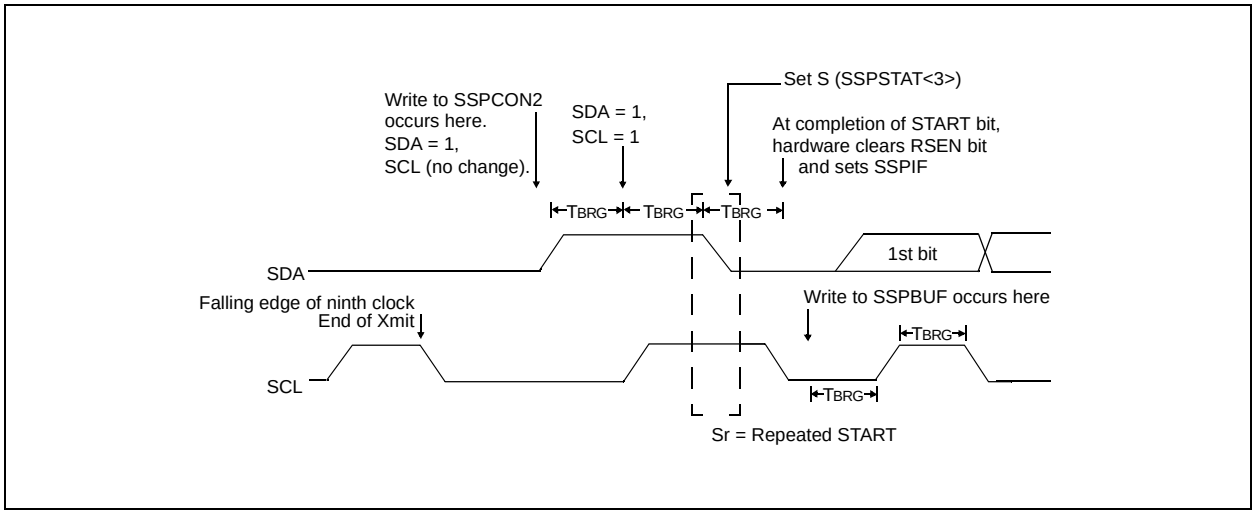
Immediately following the SSPIF bit getting set, the user may write the SSPBUF with the 7-bit address in 7-bit mode, or the default first address in 10-bit mode. After the first eight bits are transmitted and an ACK is received, the user may then transmit an additional eight bits of address (10-bit mode) or eight bits of data (7-bit mode).

16.4.9.1 WCOL Status Flag

If the user writes the SSPBUF when a Repeated START sequence is in progress, the WCOL is set and the contents of the buffer are unchanged (the write doesn't occur).

Note: Because queueing of events is not allowed, writing of the lower 5 bits of SSPCON2 is disabled until the Repeated START condition is complete.

FIGURE 16-20: REPEAT START CONDITION WAVEFORM



REGISTER 17-2: RCSTA: RECEIVE STATUS AND CONTROL REGISTER

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R-0	R-0	R-x
SPEN	RX9	SREN	CREN	ADDEN	FERR	OERR	RX9D
bit 7							bit 0

bit 7 **SPEN:** Serial Port Enable bit
1 = Serial port enabled (configures RX/DT and TX/CK pins as serial port pins)
0 = Serial port disabled

bit 6 **RX9**: 9-bit Receive Enable bit
1 = Selects 9-bit reception
0 = Selects 8-bit reception

bit 5 **SREN**: Single Receive Enable bit
Asynchronous mode:
Don't care
Synchronous mode - Master:
1 = Enables single receive
0 = Disables single receive
This bit is cleared after reception is complete.
Synchronous mode - Slave:
Don't care

bit 4 **CREN**: Continuous Receive Enable bit
Asynchronous mode:
1 = Enables receiver
0 = Disables receiver
Synchronous mode:
1 = Enables continuous receive until enable bit CREN is cleared (CREN overrides SREN)
0 = Disables continuous receive

bit 3 **ADDEN**: Address Detect Enable bit
Asynchronous mode 9-bit (RX9 = 1):
1 = Enables address detection, enables interrupt and load of the receive buffer when RSR<8> is set
0 = Disables address detection, all bytes are received, and ninth bit can be used as parity bit

bit 2 **FERR**: Framing Error bit
1 = Framing error (can be updated by reading RCREG register and receive next valid byte)
0 = No framing error

bit 1 **OERR**: Overrun Error bit
1 = Overrun error (can be cleared by clearing bit CREN)
0 = No overrun error

bit 0 **RX9D:** 9th bit of Received Data
This can be Address/Data bit or a parity bit, and must be calculated by user firmware

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
- n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

REGISTER 20-2: CONFIG2L: CONFIGURATION REGISTER 2 LOW (BYTE ADDRESS 300002h)

U-0	U-0	U-0	U-0	R/P-1	R/P-1	R/P-1	R/P-1
—	—	—	—	BORV1	BORV0	BOREN	$\overline{\text{PWRTEN}}$
bit 7							bit 0

bit 7-4 **Unimplemented:** Read as '0'

bit 3-2 **BORV1:BORV0:** Brown-out Reset Voltage bits
11 = V_{BOR} set to 2.5V
10 = V_{BOR} set to 2.7V
01 = V_{BOR} set to 4.2V
00 = V_{BOR} set to 4.5V

bit 1 **BOREN:** Brown-out Reset Enable bit
1 = Brown-out Reset enabled
0 = Brown-out Reset disabled

bit 0 **$\overline{\text{PWRTEN}}$:** Power-up Timer Enable bit
1 = PWRT disabled
0 = PWRT enabled

Legend:		
R = Readable bit	P = Programmable bit	U = Unimplemented bit, read as '0'
- n = Value when device is unprogrammed		u = Unchanged from programmed state

REGISTER 20-3: CONFIG2H: CONFIGURATION REGISTER 2 HIGH (BYTE ADDRESS 300003h)

U-0	U-0	U-0	U-0	R/P-1	R/P-1	R/P-1	R/P-1
—	—	—	—	WDTPS2	WDTPS1	WDTPS0	WDTEN
bit 7							bit 0

bit 7-4 **Unimplemented:** Read as '0'

bit 3-1 **WDTPS2:WDTPS0:** Watchdog Timer Postscale Select bits
111 = 1:128
110 = 1:64
101 = 1:32
100 = 1:16
011 = 1:8
010 = 1:4
001 = 1:2
000 = 1:1

bit 0 **WDTEN:** Watchdog Timer Enable bit
1 = WDT enabled
0 = WDT disabled (control is placed on the SWDTEN bit)

Legend:		
R = Readable bit	P = Programmable bit	U = Unimplemented bit, read as '0'
- n = Value when device is unprogrammed		u = Unchanged from programmed state

20.3 Power-down Mode (SLEEP)

Power-down mode is entered by executing a SLEEP instruction.

If enabled, the Watchdog Timer will be cleared, but keeps running, the $\overline{PD}$ bit (RCON<3>) is cleared, the $\overline{TO}$ (RCON<4>) bit is set, and the oscillator driver is turned off. The I/O ports maintain the status they had before the SLEEP instruction was executed (driving high, low or hi-impedance).

For lowest current consumption in this mode, place all I/O pins at either V_{DD} or V_{SS}, ensure no external circuitry is drawing current from the I/O pin, power-down the A/D and disable external clocks. Pull all I/O pins that are hi-impedance inputs, high or low externally, to avoid switching currents caused by floating inputs. The T0CKI input should also be at V_{DD} or V_{SS} for lowest current consumption. The contribution from on-chip pull-ups on PORTB should be considered.

The $\overline{MCLR}$ pin must be at a logic high level (V_{IHMC}).

20.3.1 WAKE-UP FROM SLEEP

The device can wake-up from SLEEP through one of the following events:

1. External RESET input on $\overline{MCLR}$ pin.
2. Watchdog Timer Wake-up (if WDT was enabled).
3. Interrupt from INT pin, RB port change or a Peripheral Interrupt.

The following peripheral interrupts can wake the device from SLEEP:

1. PSP read or write.
2. TMR1 interrupt. Timer1 must be operating as an asynchronous counter.
3. TMR3 interrupt. Timer3 must be operating as an asynchronous counter.
4. CCP Capture mode interrupt.
5. Special event trigger (Timer1 in Asynchronous mode using an external clock).
6. MSSP (START/STOP) bit detect interrupt.
7. MSSP transmit or receive in Slave mode (SPI/I^2C).
8. USART RX or TX (Synchronous Slave mode).
9. A/D conversion (when A/D clock source is RC).
10. EEPROM write operation complete.
11. LVD interrupt.

Other peripherals cannot generate interrupts, since during SLEEP, no on-chip clocks are present.

External $\overline{MCLR}$ Reset will cause a device RESET. All other events are considered a continuation of program execution and will cause a “wake-up”. The $\overline{TO}$ and $\overline{PD}$ bits in the RCON register can be used to determine the cause of the device RESET. The $\overline{PD}$ bit, which is set on power-up, is cleared when SLEEP is invoked. The $\overline{TO}$ bit is cleared, if a WDT time-out occurred (and caused wake-up).

When the SLEEP instruction is being executed, the next instruction (PC + 2) is pre-fetched. For the device to wake-up through an interrupt event, the corresponding interrupt enable bit must be set (enabled). Wake-up is regardless of the state of the GIE bit. If the GIE bit is clear (disabled), the device continues execution at the instruction after the SLEEP instruction. If the GIE bit is set (enabled), the device executes the instruction after the SLEEP instruction and then branches to the interrupt address. In cases where the execution of the instruction following SLEEP is not desirable, the user should have a NOP after the SLEEP instruction.

20.3.2 WAKE-UP USING INTERRUPTS

When global interrupts are disabled (GIE cleared) and any interrupt source has both its interrupt enable bit and interrupt flag bit set, one of the following will occur:

- If an interrupt condition (interrupt flag bit and interrupt enable bits are set) occurs **before** the execution of a SLEEP instruction, the SLEEP instruction will complete as a NOP. Therefore, the WDT and WDT postscaler will not be cleared, the $\overline{TO}$ bit will not be set and $\overline{PD}$ bits will not be cleared.
- If the interrupt condition occurs **during or after** the execution of a SLEEP instruction, the device will immediately wake-up from SLEEP. The SLEEP instruction will be completely executed before the wake-up. Therefore, the WDT and WDT postscaler will be cleared, the $\overline{TO}$ bit will be set and the $\overline{PD}$ bit will be cleared.

Even if the flag bits were checked before executing a SLEEP instruction, it may be possible for flag bits to become set before the SLEEP instruction completes. To determine whether a SLEEP instruction executed, test the $\overline{PD}$ bit. If the $\overline{PD}$ bit is set, the SLEEP instruction was executed as a NOP.

To ensure that the WDT is cleared, a CLRWDT instruction should be executed before a SLEEP instruction.

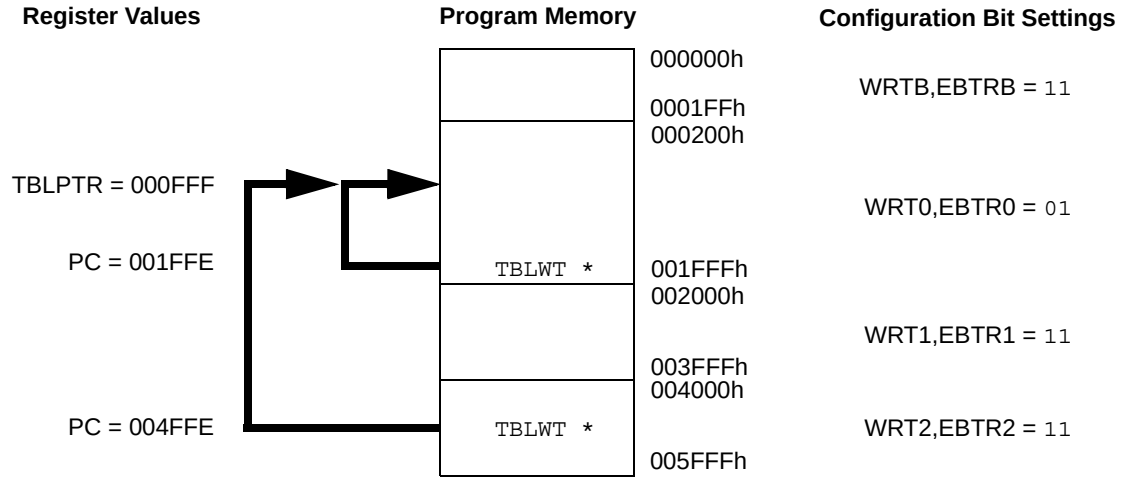
20.4.1 PROGRAM MEMORY CODE PROTECTION

The user memory may be read to, or written from, any location using the Table Read and Table Write instructions. The device ID may be read with Table Reads. The configuration registers may be read and written with the Table Read and Table Write instructions.

In User mode, the CPn bits have no direct effect. CPn bits inhibit external reads and writes. A block of user memory may be protected from Table Writes if the WRTn configuration bit is '0'. The EBTRn bits control Table Reads. For a block of user memory with the EBTRn bit set to '0', a Table Read instruction that executes from within that block is allowed to read. A Table Read instruction that executes from a location outside of that block is not allowed to read, and will result in reading '0's. Figures 20-4 through 20-6 illustrate Table Write and Table Read protection.

Note:	Code protection bits may only be written to a '0' from a '1' state. It is not possible to write a '1' to a bit in the '0' state. Code protection bits are only set to '1' by a block erase function. The block erase function can only be initiated via ICSP or an external programmer.

FIGURE 20-4: TABLE WRITE (WRTn) DISALLOWED



Results: All Table Writes disabled to Blockn whenever WRTn = 0.

TABLE 21-2: PIC18FXXX INSTRUCTION SET

Mnemonic, Operands		Description	Cycles	16-Bit Instruction Word				Status Affected	Notes
				MSb			LSb		
BYTE-ORIENTED FILE REGISTER OPERATIONS									
ADDWF	f, d, a	Add WREG and f	1	0010	01da0	ffff	ffff	C, DC, Z, OV, N	1, 2
ADDWFC	f, d, a	Add WREG and Carry bit to f	1	0010	0da	ffff	ffff	C, DC, Z, OV, N	1, 2
ANDWF	f, d, a	AND WREG with f	1	0001	01da	ffff	ffff	Z, N	1,2
CLRF	f, a	Clear f	1	0110	101a	ffff	ffff	Z	2
COMF	f, d, a	Complement f	1	0001	11da	ffff	ffff	Z, N	1, 2
CPFSEQ	f, a	Compare f with WREG, skip =	1 (2 or 3)	0110	001a	ffff	ffff	None	4
CPFSGT	f, a	Compare f with WREG, skip >	1 (2 or 3)	0110	010a	ffff	ffff	None	4
CPFSLT	f, a	Compare f with WREG, skip <	1 (2 or 3)	0110	000a	ffff	ffff	None	1, 2
DECF	f, d, a	Decrement f	1	0000	01da	ffff	ffff	C, DC, Z, OV, N	1, 2, 3, 4
DECFSZ	f, d, a	Decrement f, Skip if 0	1 (2 or 3)	0010	11da	ffff	ffff	None	1, 2, 3, 4
DCFSNZ	f, d, a	Decrement f, Skip if Not 0	1 (2 or 3)	0100	11da	ffff	ffff	None	1, 2
INCF	f, d, a	Increment f	1	0010	10da	ffff	ffff	C, DC, Z, OV, N	1, 2, 3, 4
INCFSZ	f, d, a	Increment f, Skip if 0	1 (2 or 3)	0011	11da	ffff	ffff	None	4
INFSNZ	f, d, a	Increment f, Skip if Not 0	1 (2 or 3)	0100	10da	ffff	ffff	None	1, 2
IORWF	f, d, a	Inclusive OR WREG with f	1	0001	00da	ffff	ffff	Z, N	1, 2
MOVF	f, d, a	Move f	1	0101	00da	ffff	ffff	Z, N	1
MOVFF	f_s, f_d	Move f_s (source) to 1st word f_d (destination) 2nd word	2	1100 1111	ffff ffff	ffff ffff	ffff ffff	None	
MOVWF	f, a	Move WREG to f	1	0110	111a	ffff	ffff	None	
MULWF	f, a	Multiply WREG with f	1	0000	001a	ffff	ffff	None	
NEGF	f, a	Negate f	1	0110	110a	ffff	ffff	C, DC, Z, OV, N	1, 2
RLCF	f, d, a	Rotate Left f through Carry	1	0011	01da	ffff	ffff	C, Z, N	
RLNCF	f, d, a	Rotate Left f (No Carry)	1	0100	01da	ffff	ffff	Z, N	1, 2
RRCF	f, d, a	Rotate Right f through Carry	1	0011	00da	ffff	ffff	C, Z, N	
RRNCF	f, d, a	Rotate Right f (No Carry)	1	0100	00da	ffff	ffff	Z, N	
SETF	f, a	Set f	1	0110	100a	ffff	ffff	None	
SUBFWB	f, d, a	Subtract f from WREG with borrow	1	0101	01da	ffff	ffff	C, DC, Z, OV, N	1, 2
SUBWF	f, d, a	Subtract WREG from f	1	0101	11da	ffff	ffff	C, DC, Z, OV, N	
SUBWFB	f, d, a	Subtract WREG from f with borrow	1	0101	10da	ffff	ffff	C, DC, Z, OV, N	1, 2
SWAPF	f, d, a	Swap nibbles in f	1	0011	10da	ffff	ffff	None	4
TSTFSZ	f, a	Test f, skip if 0	1 (2 or 3)	0110	011a	ffff	ffff	None	1, 2
XORWF	f, d, a	Exclusive OR WREG with f	1	0001	10da	ffff	ffff	Z, N	
BIT-ORIENTED FILE REGISTER OPERATIONS									
BCF	f, b, a	Bit Clear f	1	1001	bbba	ffff	ffff	None	1, 2
BSF	f, b, a	Bit Set f	1	1000	bbba	ffff	ffff	None	1, 2
BTFSC	f, b, a	Bit Test f, Skip if Clear	1 (2 or 3)	1011	bbba	ffff	ffff	None	3, 4
BTFSS	f, b, a	Bit Test f, Skip if Set	1 (2 or 3)	1010	bbba	ffff	ffff	None	3, 4
BTG	f, d, a	Bit Toggle f	1	0111	bbba	ffff	ffff	None	1, 2

Note 1: When a PORT register is modified as a function of itself (e.g., `MOVF PORTB, 1, 0`), the value used will be that value present on the pins themselves. For example, if the data latch is '1' for a pin configured as input and is driven low by an external device, the data will be written back with a '0'.

2: If this instruction is executed on the TMR0 register (and, where applicable, d = 1), the prescaler will be cleared if assigned.

3: If Program Counter (PC) is modified or a conditional test is true, the instruction requires two cycles. The second cycle is executed as a `NOP`.

4: Some instructions are 2-word instructions. The second word of these instructions will be executed as a `NOP`, unless the first word of the instruction retrieves the information embedded in these 16-bits. This ensures that all program memory locations have a valid instruction.

5: If the Table Write starts the write cycle to internal memory, the write will continue until terminated.

BCF **Bit Clear f**

Syntax: [*label*] BCF f,b[,a]

Operands: $0 \le f \le 255$
$0 \le b \le 7$
$a \in [0,1]$

Operation: $0 \rightarrow f<b>$

Status Affected: None

Encoding:

1001	bbba	ffff	ffff

Description: Bit 'b' in register 'f' is cleared. If 'a' is 0, the Access Bank will be selected, overriding the BSR value. If 'a' = 1, then the bank will be selected as per the BSR value (default).

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: `BCF FLAG_REG, 7, 0`

Before Instruction
FLAG_REG = 0xC7
After Instruction
FLAG_REG = 0x47

BN **Branch if Negative**

Syntax: [*label*] BN n

Operands: $-128 \le n \le 127$

Operation: if negative bit is '1'
$(PC) + 2 + 2n \rightarrow PC$

Status Affected: None

Encoding:

1110	0110	nnnn	nnnn

Description: If the Negative bit is '1', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC+2+2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example:
```
HERE    BN   Jump
```
Before Instruction
PC = address (HERE)
After Instruction
If Negative = 1;
PC = address (Jump)
If Negative = 0;
PC = address (HERE+2)

BNOV **Branch if Not Overflow**

Syntax: [*label*] BNOV n

Operands: $-128 \le n \le 127$

Operation: if overflow bit is '0'
(PC) + 2 + 2n → PC

Status Affected: None

Encoding:

1110	0101	nnnn	nnnn

Description: If the Overflow bit is '0', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC+2+2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BNOV Jump`

Before Instruction
PC = address `(HERE)`
After Instruction
If Overflow = 0;
PC = address `(Jump)`
If Overflow = 1;
PC = address `(HERE+2)`

BNZ **Branch if Not Zero**

Syntax: [*label*] BNZ n

Operands: $-128 \le n \le 127$

Operation: if zero bit is '0'
(PC) + 2 + 2n → PC

Status Affected: None

Encoding:

1110	0001	nnnn	nnnn

Description: If the Zero bit is '0', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC+2+2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BNZ Jump`

Before Instruction
PC = address `(HERE)`
After Instruction
If Zero = 0;
PC = address `(Jump)`
If Zero = 1;
PC = address `(HERE+2)`

FIGURE 23-15: EXAMPLE SPI SLAVE MODE TIMING (CKE = 1)

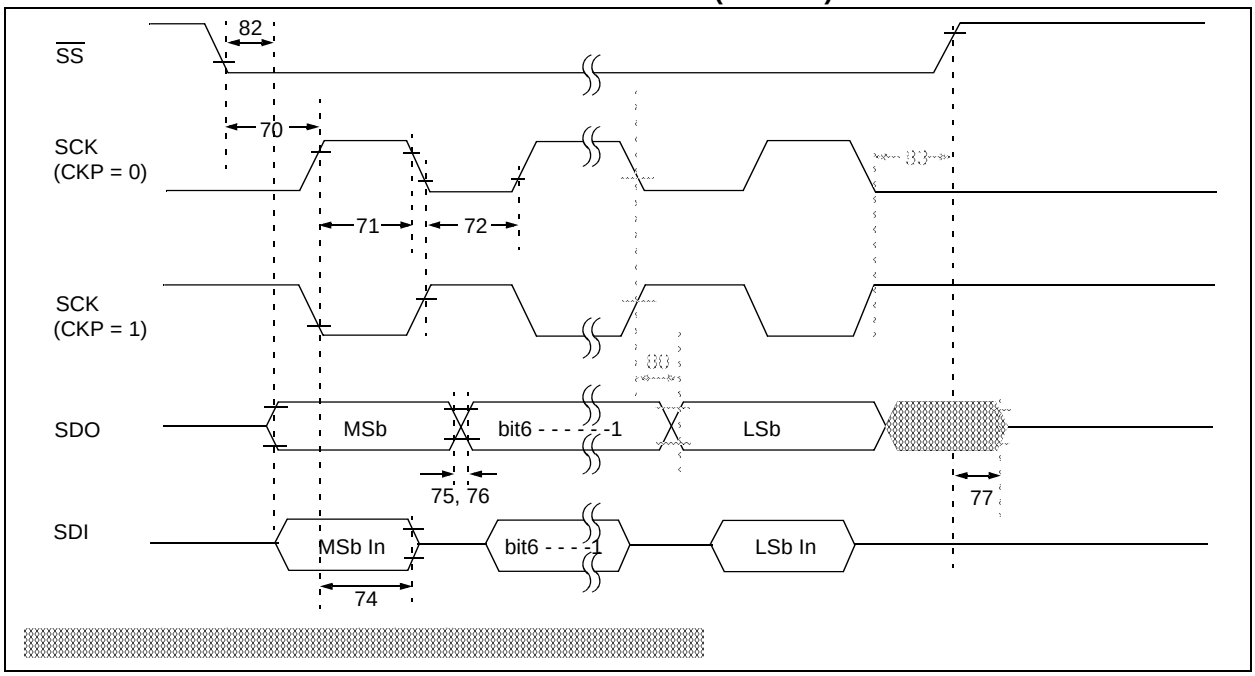


TABLE 23-14: EXAMPLE SPI SLAVE MODE REQUIREMENTS (CKE = 1)

Param. No.	Symbol	Characteristic		Min	Max	Units	Conditions
70	TssL2scH, TssL2scL	$\overline{SS}\downarrow$ to SCK$\downarrow$ or SCK$\uparrow$ input		TCY	—	ns	
71	TscH	SCK input high time (Slave mode)	Continuous	1.25 TCY + 30	—	ns	
71A			Single Byte	40	—	ns	**(Note 1)**
72	TscL	SCK input low time (Slave mode)	Continuous	1.25 TCY + 30	—	ns	
72A			Single Byte	40	—	ns	**(Note 1)**
73A	TB2B	Last clock edge of Byte 1 to the first clock edge of Byte 2		1.5 TCY + 40	—	ns	**(Note 2)**
74	TscH2diL, TscL2diL	Hold time of SDI data input to SCK edge		100	—	ns	
75	TdoR	SDO data output rise time	PIC18**F**XXXX	—	25	ns	
			PIC18**LF**XXXX	—	60	ns	VDD = 2V
76	TdoF	SDO data output fall time	PIC18**F**XXXX	—	25	ns	
			PIC18**LF**XXXX	—	60	ns	VDD = 2V
77	TssH2doZ	$\overline{SS}\uparrow$ to SDO output hi-impedance		10	50	ns	
78	TscR	SCK output rise time (Master mode)	PIC18**F**XXXX	—	25	ns	
			PIC18**LF**XXXX	—	60	ns	VDD = 2V
79	TscF	SCK output fall time (Master mode)	PIC18**F**XXXX	—	25	ns	
			PIC18**LF**XXXX	—	60	ns	VDD = 2V
80	TscH2doV, TscL2doV	SDO data output valid after SCK edge	PIC18**F**XXXX	—	50	ns	
			PIC18**LF**XXXX	—	150	ns	VDD = 2V
82	TssL2doV	SDO data output valid after $\overline{SS}\downarrow$ edge	PIC18**F**XXXX	—	50	ns	
			PIC18**LF**XXXX	—	150	ns	VDD = 2V
83	TscH2ssH, TscL2ssH	$\overline{SS}\uparrow$ after SCK edge		1.5 TCY + 40	—	ns	

Note 1: Requires the use of Parameter # 73A.
2: Only if Parameter # 71A and # 72A are used.

TABLE 23-21: A/D CONVERTER CHARACTERISTICS: PIC18FXX39 (INDUSTRIAL, EXTENDED) PIC18LFXX39 (INDUSTRIAL)

Param No.	Symbol	Characteristic	Min	Typ	Max	Units	Conditions
A01	NR	Resolution	—	—	10	bit	
A03	EIL	Integral linearity error	—	—	<±1	LSb	VREF = VDD = 5.0V
A04	EDL	Differential linearity error	—	—	<±1	LSb	VREF = VDD = 5.0V
A05	EG	Gain error	—	—	<±1	LSb	VREF = VDD = 5.0V
A06	EOFF	Offset error	—	—	<±1.5	LSb	VREF = VDD = 5.0V
A10	—	Monotonicity	guaranteed(2)			—	VSS ≤ VAIN ≤ VREF
A20 A20A	VREF	Reference Voltage (VREFH – VREFL)	1.8V 3V	— —	— —	V V	VDD < 3.0V VDD ≥ 3.0V
A21	VREFH	Reference voltage High	AVSS	—	AVDD + 0.3V	V	
A22	VREFL	Reference voltage Low	AVSS – 0.3V	—	VREFH	V	
A25	VAIN	Analog input voltage	AVSS – 0.3V	—	AVDD + 0.3V	V	VDD ≥ 2.5V **(Note 3)**
A30	ZAIN	Recommended impedance of analog voltage source	—	—	2.5	kΩ	**(Note 4)**
A50	IREF	VREF input current **(Note 1)**	— —	— —	5 150	μA μA	During VAIN acquisition During A/D conversion cycle

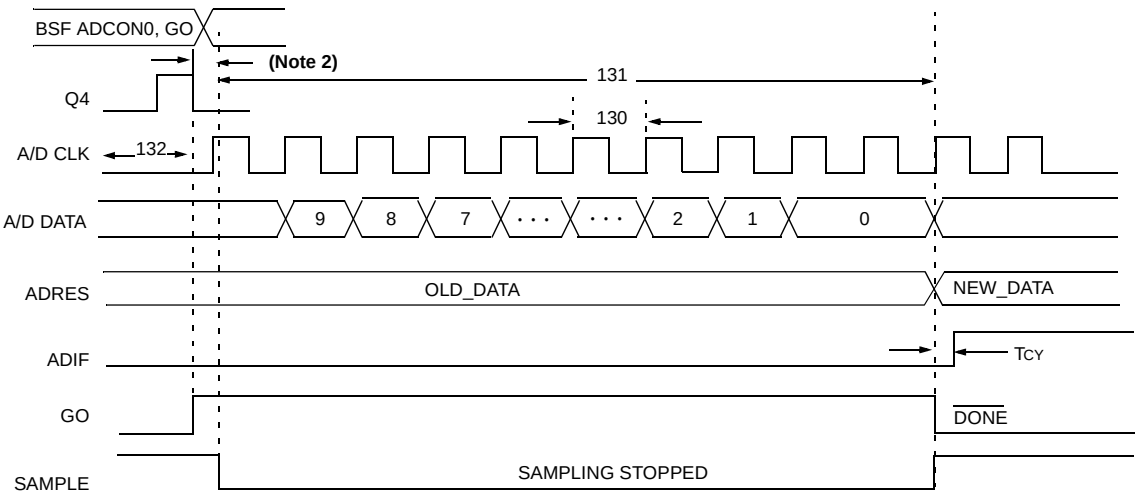
Note 1: VSS ≤ VAIN ≤ VREF

2: The A/D conversion result never decreases with an increase in the Input Voltage, and has no missing codes.

3: For VDD < 2.5V, VAIN should be limited to < .5 VDD.

4: Maximum allowed impedance for analog voltage source is 10 kΩ. This requires higher acquisition times.

FIGURE 23-22: A/D CONVERSION TIMING



Note 1: If the A/D clock source is selected as RC, a time of TCY is added before the A/D clock starts. This allows the SLEEP instruction to be executed.

2: This is a minimal RC delay (typically 100 ns), which also disconnects the holding capacitor from the analog input.