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Details

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Product Status	Active
Core Processor	HCS12
Core Size	16-Bit
Speed	16MHz
Connectivity	CANbus, EBI/EMI, SCI, SPI
Peripherals	POR, PWM, WDT
Number of I/O	31
Program Memory Size	32KB (32K x 8)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	2K x 8
Voltage - Supply (Vcc/Vdd)	2.35V ~ 5.5V
Data Converters	A/D 8x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 105°C (TA)
Mounting Type	Surface Mount
Package / Case	48-LQFP
Supplier Device Package	48-LQFP (7x7)
Purchase URL	https://www.e-xfl.com/pro/item?MUrl=&PartUrl=mc9s12c32vfae16

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Chapter 1 MC9S12C and MC9S12GC Device Overview (MC9S12C128)

- Memory options:
 - 16K or 32Kbyte Flash EEPROM (erasable in 512-byte sectors)
 - 64K, 96K, or 128Kbyte Flash EEPROM (erasable in 1024-byte sectors)
 - 1K, 2K or 4K Byte RAM
- Analog-to-digital converters:
 - One 8-channel module with 10-bit resolution
 - External conversion trigger capability
- Available on MC9S12C Family:
 - One 1M bit per second, CAN 2.0 A, B software compatible module
 - Five receive and three transmit buffers
 - Flexible identifier filter programmable as 2 x 32 bit, 4 x 16 bit, or 8 x 8 bit
 - Four separate interrupt channels for Rx, Tx, error, and wake-up
 - Low-pass filter wake-up function
 - Loop-back for self test operation
- Timer module (TIM):
 - 8-channel timer
 - Each channel configurable as either input capture or output compare
 - Simple PWM mode
 - Modulo reset of timer counter
 - 16-bit pulse accumulator
 - External event counting
 - Gated time accumulation
- PWM module:
 - Programmable period and duty cycle
 - 8-bit 6-channel or 16-bit 3-channel
 - Separate control for each pulse width and duty cycle
 - Center-aligned or left-aligned outputs
 - Programmable clock select logic with a wide range of frequencies
 - Fast emergency shutdown input
- Serial interfaces:
 - One asynchronous serial communications interface (SCI)
 - One synchronous serial peripheral interface (SPI)
- CRG (clock reset generator module)
 - Windowed COP watchdog
 - Real time interrupt
 - Clock monitor
 - Pierce or low current Colpitts oscillator
 - Phase-locked loop clock frequency multiplier
 - Limp home mode in absence of external clock
 - Low power 0.5MHz to 16MHz crystal oscillator reference clock



1.3.4.6 PA[7:0] / ADDR[15:8] / DATA[15:8] — Port A I/O Pins

PA7–PA0 are general purpose input or output pins,. In MCU expanded modes of operation, these pins are used for the multiplexed external address and data bus. PA[7:1] pins are not available in the 48-pin package version. PA[7:3] are not available in the 52-pin package version.

1.3.4.7 PB[7:0] / ADDR[7:0] / DATA[7:0] — Port B I/O Pins

PB7–PB0 are general purpose input or output pins. In MCU expanded modes of operation, these pins are used for the multiplexed external address and data bus. PB[7:5] and PB[3:0] pins are not available in the 48-pin nor 52-pin package version.

1.3.4.8 PE7 / NOACC / XCLKS — Port E I/O Pin 7

PE7 is a general purpose input or output pin. During MCU expanded modes of operation, the NOACC signal, when enabled, is used to indicate that the current bus cycle is an unused or "free" cycle. This signal will assert when the CPU is not using the bus. The $\overline{\text{XCLKS}}$ is an input signal which controls whether a crystal in combination with the internal Colpitts (low power) oscillator is used or whether Pierce oscillator/external clock circuitry is used. The state of this pin is latched at the rising edge of $\overline{\text{RESET}}$. If the input is a logic low the EXTAL pin is configured for an external clock drive or a Pierce oscillator. If input is a logic high a Colpitts oscillator circuit is configured on EXTAL and XTAL. Since this pin is an input with a pull-up device during reset, if the pin is left floating, the default configuration is a Colpitts oscillator circuit on EXTAL and XTAL.



1. Due to the nature of a translated ground Colpitts oscillator a DC voltage bias is applied to the crystal. Please contact the crystal manufacturer for crystal DC.

Figure 1-11. Colpitts Oscillator Connections (PE7 = 1)



1. RS can be zero (shorted) when used with higher frequency crystals, refer to manufacturer's data.





2.3.2.1 Port T Registers

2.3.2.1.1 Port T I/O Register (PTT)

Module Base + 0x0000

	7	6	5	4	3	2	1	0
R W	PTT7	PTT6	PTT5	PTT4	PTT3	PTT2	PTT1	PTT0
ТІМ	IOC7	IOC6	IOC5	IOC4	IOC3	IOC2	IOC1	IOC0
PWM				PWM4	PWM3	PWM2	PWM1	PWM0
Reset	0	0	0	0	0	0	0	0
		= Unimplemer	nted or Reserve	ed	•			

Figure 2-3. Port T I/O Register (PTT)

Read: Anytime.

Write: Anytime.

If the data direction bits of the associated I/O pins are set to 1, a read returns the value of the port register, otherwise the value at the pins is read.

If a TIM-channel is defined as output, the related port T is assigned to IOC function.

In addition to the possible timer functionality of port T pins PWM channels can be routed to port T. For this the Module Routing Register (MODRR) needs to be configured.

MODRR[x]	PWME[x]	TIMEN[x] (2)	Port T[x] Output
0	0	0	General Purpose I/O
0	0	1	Timer
0	1	0	General Purpose I/O
0	1	1	Timer
1	0	0	General Purpose I/O
1	0	1	Timer
1	1	0	PWM
1	1	1	PWM

 Table 2-3. Port T[4:0] Pin Functionality Configurations⁽¹⁾

1. All fields in the that are not shaded are standard use cases.

 TIMEN[x] means that the timer is enabled (TSCR1[7]), the related channel is configured for output compare function (TIOS[x] or special output on a timer overflow event — configurable in TTOV[x]) and the timer output is routed to the port pin (TCTL1/TCTL2).



Chapter 3 Module Mapping Control (MMCV4) Block Description

3.3.2.2 Initialization of Internal Registers Position Register (INITRG)

Module Base + 0x0011

Starting address location affected by INITRG register setting.



Figure 3-4. Initialization of Internal Registers Position Register (INITRG)

Read: Anytime

Write: Once in normal and emulation modes and anytime in special modes

This register initializes the position of the internal registers within the on-chip system memory map. The registers occupy either a 1K byte or 2K byte space and can be mapped to any 2K byte space within the first 32K bytes of the system's address space.

Table 3-3. INITRG Field Descriptions

Description
rnal Register Map Position — These four bits in combination with the leading zero supplied by bit 7 of
RG determine the upper five bits of the base address for the system's internal registers (i.e., the minimum address is 0x0000 and the maximum is 0x7EEE)
r r R



The PPAGE register holds the page select value for the program page window. The value of the PPAGE register can be manipulated by normal read and write (some devices don't allow writes in some modes) instructions as well as the CALL and RTC instructions.

Control registers, vector spaces, and a portion of on-chip memory are located in unpaged portions of the 64K byte physical address space. The stack and I/O addresses should also be in unpaged memory to make them accessible from any page.

The starting address of a service routine must be located in unpaged memory because the 16-bit exception vectors cannot point to addresses in paged memory. However, a service routine can call other routines that are in paged memory. The upper 16K byte block of memory space (0xC000–0xFFFF) is unpaged. It is recommended that all reset and interrupt vectors point to locations in this area.

3.4.3.1 CALL and Return from Call Instructions

CALL and RTC are uninterruptable instructions that automate page switching in the program expansion window. CALL is similar to a JSR instruction, but the subroutine that is called can be located anywhere in the normal 64K byte address space or on any page of program expansion memory. CALL calculates and stacks a return address, stacks the current PPAGE value, and writes a new instruction-supplied value to PPAGE. The PPAGE value controls which of the 64 possible pages is visible through the 16K byte expansion window in the 64K byte memory map. Execution then begins at the address of the called subroutine.

During the execution of a CALL instruction, the CPU:

- Writes the old PPAGE value into an internal temporary register and writes the new instructionsupplied PPAGE value into the PPAGE register.
- Calculates the address of the next instruction after the CALL instruction (the return address), and pushes this 16-bit value onto the stack.
- Pushes the old PPAGE value onto the stack.
- Calculates the effective address of the subroutine, refills the queue, and begins execution at the new address on the selected page of the expansion window.

This sequence is uninterruptable; there is no need to inhibit interrupts during CALL execution. A CALL can be performed from any address in memory to any other address.

The PPAGE value supplied by the instruction is part of the effective address. For all addressing mode variations except indexed-indirect modes, the new page value is provided by an immediate operand in the instruction. In indexed-indirect variations of CALL, a pointer specifies memory locations where the new page value and the address of the called subroutine are stored. Using indirect addressing for both the new page value and the address within the page allows values calculated at run time rather than immediate values that must be known at the time of assembly.

The RTC instruction terminates subroutines invoked by a CALL instruction. RTC unstacks the PPAGE value and the return address and refills the queue. Execution resumes with the next instruction after the CALL.



Chapter 4 Multiplexed External Bus Interface (MEBIV3)

4.1 Introduction

This section describes the functionality of the multiplexed external bus interface (MEBI) sub-block of the S12 core platform. The functionality of the module is closely coupled with the S12 CPU and the memory map controller (MMC) sub-blocks.

Figure 4-1 is a block diagram of the MEBI. In Figure 4-1, the signals on the right hand side represent pins that are accessible externally. On some chips, these may not all be bonded out.

The MEBI sub-block of the core serves to provide access and/or visibility to internal core data manipulation operations including timing reference information at the external boundary of the core and/or system. Depending upon the system operating mode and the state of bits within the control registers of the MEBI, the internal 16-bit read and write data operations will be represented in 8-bit or 16-bit accesses externally. Using control information from other blocks within the system, the MEBI will determine the appropriate type of data access to be generated.

4.1.1 Features

The block name includes these distinctive features:

- External bus controller with four 8-bit ports A,B, E, and K
- Data and data direction registers for ports A, B, E, and K when used as general-purpose I/O
- Control register to enable/disable alternate functions on ports E and K
- Mode control register
- Control register to enable/disable pull resistors on ports A, B, E, and K
- Control register to enable/disable reduced output drive on ports A, B, E, and K
- Control register to configure external clock behavior
- Control register to configure IRQ pin operation
- Logic to capture and synchronize external interrupt pin inputs



4.3.2.8 Port E Assignment Register (PEAR)

Module Base + 0x000A

Starting address location affected by INITRG register setting.

	7	6	5	4	3	2	1	0
R		0					0	0
W	NOACCE		FIFUE	NEGLK	LOINE	RDWE		
Reset								
Special Single Chip	0	0	0	0	0	0	0	0
Special Test	0	0	1	0	1	1	0	0
Peripheral	0	0	0	0	0	0	0	0
Emulation Expanded Narrow	1	0	1	0	1	1	0	0
Emulation Expanded Wide	1	0	1	0	1	1	0	0
Normal Single Chip	0	0	0	1	0	0	0	0
Normal Expanded Narrow	0	0	0	0	0	0	0	0
Normal Expanded Wide	0	0	0	0	0	0	0	0
		= Unimplem	ented or Res	served				

Figure 4-12. Port E Assignment Register (PEAR)

Read: Anytime (provided this register is in the map).

Write: Each bit has specific write conditions. Please refer to the descriptions of each bit on the following pages.

Port E serves as general-purpose I/O or as system and bus control signals. The PEAR register is used to choose between the general-purpose I/O function and the alternate control functions. When an alternate control function is selected, the associated DDRE bits are overridden.

The reset condition of this register depends on the mode of operation because bus control signals are needed immediately after reset in some modes. In normal single-chip mode, no external bus control signals are needed so all of port E is configured for general-purpose I/O. In normal expanded modes, only the E clock is configured for its alternate bus control function and the other bits of port E are configured for general-purpose I/O. As the reset vector is located in external memory, the E clock is required for this access. R/W is only needed by the system when there are external writable resources. If the normal expanded system needs any other bus control signals, PEAR would need to be written before any access that needed the additional signals. In special test and emulation modes, IPIPE1, IPIPE0, E, LSTRB, and R/W are configured out of reset as bus control signals.

This register is not in the on-chip memory map in expanded and special peripheral modes. Therefore, these accesses will be echoed externally.



DBGC2 being logic 1 or logic 0, respectively. BDM requests have a higher priority than SWI requests. No data breakpoints are allowed in this mode.

TAGAB in DBGC2 selects whether the breakpoint mode is forced or tagged. The BKxMBH:L bits in DBGC3 select whether or not the breakpoint is matched exactly or is a range breakpoint. They also select whether the address is matched on the high byte, low byte, both bytes, and/or memory expansion. The RWx and RWxEN bits in DBGC3 select whether the type of bus cycle to match is a read, write, or read/write when performing forced breakpoints.

7.4.1.2 Full Breakpoint Mode

Full breakpoint mode requires a match on address and data for a breakpoint to occur. Upon a successful match, the system will enter background debug mode or initiate a software interrupt based upon the state of BDM in DBGC2 being logic 1 or logic 0, respectively. BDM requests have a higher priority than SWI requests. R/W matches are also allowed in this mode.

TAGAB in DBGC2 selects whether the breakpoint mode is forced or tagged. When TAGAB is set in DBGC2, only addresses are compared and data is ignored. The BKAMBH:L bits in DBGC3 select whether or not the breakpoint is matched exactly, is a range breakpoint, or is in page space. The BKBMBH:L bits in DBGC3 select whether the data is matched on the high byte, low byte, or both bytes. RWA and RWAEN bits in DBGC2 select whether the type of bus cycle to match is a read or a write when performing forced breakpoints. RWB and RWBEN bits in DBGC2 are not used in full breakpoint mode.

NOTE

The full trigger mode is designed to be used for either a word access or a byte access, but not both at the same time. Confusing trigger operation (seemingly false triggers or no trigger) can occur if the trigger address occurs in the user program as both byte and word accesses.

7.4.1.3 Breakpoint Priority

Breakpoint operation is first determined by the state of the BDM module. If the BDM module is already active, meaning the CPU is executing out of BDM firmware, breakpoints are not allowed. In addition, while executing a BDM TRACE command, tagging into BDM is not allowed. If BDM is not active, the breakpoint will give priority to BDM requests over SWI requests. This condition applies to both forced and tagged breakpoints.

In all cases, BDM related breakpoints will have priority over those generated by the Breakpoint sub-block. This priority includes breakpoints enabled by the TAGLO and TAGHI external pins of the system that interface with the BDM directly and whose signal information passes through and is used by the breakpoint sub-block.



Chapter 8 Analog-to-Digital Converter (ATD10B8C) Block Description

Address	Name		Bit 7	6	5	4	3	2	1	Bit 0
Left Justifi	ed Result Data									
0x0010	ATDDR0H	R	BIT 9 MSB BIT 7 MSB	BIT 8 BIT 6	BIT 7 BIT 5	BIT 6 BIT 4	BIT 5 BIT 3	BIT 4 BIT 2	BIT 3 BIT 1	BIT 2 BIT 0
		W								
0x0011	ATDDR0L	R	BIT 1 u	BIT 0 u	0 0	0 0	0 0	0 0	0 0	0 0
		W								
0x0012	ATDDR1H	R	BIT 9 MSB BIT 7 MSB	BIT 8 BIT 6	BIT 7 BIT 5	BIT 6 BIT 4	BIT 5 BIT 3	BIT 4 BIT 2	BIT 3 BIT 1	BIT 2 BIT 0
		W								
0x0013	ATDDR1L	R	BIT 1 u	BIT 0 u	0 0	0 0	0 0	0 0	0 0	0 0
		W								
0x0014	ATDDR2H	R	BIT 9 MSB BIT 7 MSB	BIT 8 BIT 6	BIT 7 BIT 5	BIT 6 BIT 4	BIT 5 BIT 3	BIT 4 BIT 2	BIT 3 BIT 1	BIT 2 BIT 0
		W								
0x0015	ATDDR2L	R	BIT 1 u	BIT 0 u	0 0	0 0	0 0	0 0	0 0	0 0
		W								
0x0016	ATDDR3H	R	BIT 9 MSB BIT 7 MSB	BIT 8 BIT 6	BIT 7 BIT 5	BIT 6 BIT 4	BIT 5 BIT 3	BIT 4 BIT 2	BIT 3 BIT 1	BIT 2 BIT 0
		W								
0x0017	ATDDR3L	R	BIT 1 u	BIT 0 u	0 0	0 0	0 0	0 0	0 0	0 0
		W								
0x0018	ATDDR4H	R	BIT 9 MSB BIT 7 MSB	BIT 8 BIT 6	BIT 7 BIT 5	BIT 6 BIT 4	BIT 5 BIT 3	BIT 4 BIT 2	BIT 3 BIT 1	BIT 2 BIT 0
		W								
0x0019	ATDDR4L	R	BIT 1 u	BIT 0 u	0	0	0	0	0	0
		W								
0x001A	ATDDR5H	R	BIT 9 MSB BIT 7 MSB	BIT 8 BIT 6	BIT 7 BIT 5	BIT 6 BIT 4	BIT 5 BIT 3	BIT 4 BIT 2	BIT 3 BIT 1	BIT 2 BIT 0
		W			-					
0x001B	ATDDR5L	н	BII 1 u	BIL 0	0	0	0	0	0	0
		W								D / T 0
0x001C	ATDDR6H	R	BIT 9 MSB BIT 7 MSB	BIT 8 BIT 6	BIT 7 BIT 5	BIT 6 BIT 4	BIT 5 BIT 3	BIT 4 BIT 2	BIT 3 BIT 1	BIT 2 BIT 0
		W								
0x001D	ATDDR6L	R	BIT 1 u	BIT 0 u	0 0	0	0 0	0 0	0 0	0 0
		W								
				= Unimpler	mented or F	Reserved				

Figure 8-2. ATD Register Summary (Sheet 2 of 4)



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Chapter 8 Analog-to-Digital Converter (ATD10B8C) Block Description
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8.5.1.3 Step 3

Configure how many conversions you want to perform in one sequence and define other settings in ATDCTL3.

Example: Write S4C=1 to do 4 conversions per sequence.

8.5.1.4 Step 4

Configure resolution, sampling time and ATD clock speed in ATDCTL4.

Example: Use default for resolution and sampling time by leaving SRES8, SMP1 and SMP0 clear. For a bus clock of 40MHz write 9 to PR4-0, this gives an ATD clock of 0.5*40MHz/(9+1) = 2MHz which is within the allowed range for f_{ATDCLK} .

8.5.1.5 Step 5

Configure starting channel, single/multiple channel, continuous or single sequence and result data format in ATDCTL5. Writing ATDCTL5 will start the conversion, so make sure your write ATDCTL5 in the last step.

Example: Leave CC,CB,CA clear to start on channel AN0. Write MULT=1 to convert channel AN0 to AN3 in a sequence (4 conversion per sequence selected in ATDCTL3).

8.5.2 Aborting an A/D conversion

8.5.2.1 Step 1

Disable the ATD Interrupt by writing ASCIE=0 in ATDCTL2. This will also abort any ongoing conversion sequence.

It is important to clear the interrupt enable at this point, prior to step 3, as depending on the device clock gating it may not always be possible to clear it or the SCF flag once the module is disabled (ADPU=0).

8.5.2.2 Step 2

Clear the SCF flag by writing a 1 in ATDSTAT0.

(Remaining flags will be cleared with the next start of a conversions, but SCF flag should be cleared to avoid SCF interrupt.)

8.5.2.3 Step 3

Power down ATD by writing ADPU=0 in ATDCTL2.

8.6 Resets

At reset the ATD10B8C is in a power down state. The reset state of each individual bit is listed within Section 8.3.2, "Register Descriptions" which details the registers and their bit-field.



Chapter 12 Pulse-Width Modulator (PWM8B6CV1) Block Description



Read: anytime

Write: anytime (any value written causes PWM counter to be reset to 0x0000).

12.3.2.13 PWM Channel Period Registers (PWMPERx)

There is a dedicated period register for each channel. The value in this register determines the period of the associated PWM channel.

The period registers for each channel are double buffered so that if they change while the channel is enabled, the change will NOT take effect until one of the following occurs:

- The effective period ends
- The counter is written (counter resets to 0x0000)



13.4.5 Single-Wire Operation

Normally, the SCI uses two pins for transmitting and receiving. In single-wire operation, the RXD pin is disconnected from the SCI. The SCI uses the TXD pin for both receiving and transmitting.



Figure 13-22. Single-Wire Operation (LOOPS = 1, RSRC = 1)

Enable single-wire operation by setting the LOOPS bit and the receiver source bit, RSRC, in SCI control register 1 (SCICR1). Setting the LOOPS bit disables the path from the **Rx Input** signal to the receiver. Setting the RSRC bit connects the receiver input to the output of the TXD pin driver. Both the transmitter and receiver must be enabled (TE = 1 and RE = 1). The TXDIR bit (SCISR2[1]) determines whether the TXD pin is going to be used as an input (TXDIR = 0) or an output (TXDIR = 1) in this mode of operation.

13.4.6 Loop Operation

In loop operation the transmitter output goes to the receiver input. The **Rx Input** signal is disconnected from the SCI



Figure 13-23. Loop Operation (LOOPS = 1, RSRC = 0)

Enable loop operation by setting the LOOPS bit and clearing the RSRC bit in SCI control register 1 (SCICR1). Setting the LOOPS bit disables the path from the **Rx Input** signal to the receiver. Clearing the RSRC bit connects the transmitter output to the receiver input. Both the transmitter and receiver must be enabled (TE = 1 and RE = 1).

13.5 Initialization Information

13.5.1 Reset Initialization

The reset state of each individual bit is listed in Section 13.3, "Memory Map and Registers" which details the registers and their bit fields. All special functions or modes which are initialized during or just following reset are described within this section.





Read: Anytime

Write: Anytime

Table 15-9. TCTL1/TCTL2 Field Descriptions

Field	Description
7:0 OMx	Output Mode — These eight pairs of control bits are encoded to specify the output action to be taken as a result of a successful OCx compare. When either OMx or OLx is 1, the pin associated with OCx becomes an output tied to OCx. Note: To enable output action by OMx bits on timer port, the corresponding bit in OC7M should be cleared.
7:0 OLx	Output Level — These eight pairs of control bits are encoded to specify the output action to be taken as a result of a successful OCx compare. When either OMx or OLx is 1, the pin associated with OCx becomes an output tied to OCx. Note: To enable output action by OLx bits on timer port, the corresponding bit in OC7M should be cleared.

Table 15-10. Compare Result Output Action

ОМх	OLx	Action
0	0	Timer disconnected from output pin logic
0	1	Toggle OCx output line
1	0	Clear OCx output line to zero
1	1	Set OCx output line to one

To operate the 16-bit pulse accumulator independently of input capture or output compare 7 and 0 respectively the user must set the corresponding bits IOSx = 1, OMx = 0 and OLx = 0. OC7M7 in the OC7M register must also be cleared.

To enable output action using the OM7 and OL7 bits on the timer port, the corresponding bit OC7M7 in the OC7M register must also be cleared. The settings for these bits can be seen in Table 15-11

Table 15-11. The OC7 and OCx event priority

OC7M7=0					OC7I	W7=1	
OC7Mx=1 OC7Mx=0		Vx=0	OC7	Mx=1	OC7	Mx=0	
TC7=TCx	TC7>TCx	TC7=TCx TC7>TCx		TC7=TCx	TC7>TCx	TC7=TCx	TC7>TCx
IOCx=OC7Dx IOC7=OM7/O L7	IOCx=OC7Dx +OMx/OLx IOC7=OM7/O L7	IOCx=0 IOC7=0	Mx/OLx M7/OL7	IOCx=OC7Dx IOC7=OC7D7	IOCx=OC7Dx +OMx/OLx IOC7=OC7D7	IOCx=C IOC7=0	Mx/OLx OC7D7

```
NP
```

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Chapter 17 16 Kbyte Flash Module (S12FTS16KV1)
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Figure 17-20. RESERVED6

All bits read 0 and are not writable.

17.4 Functional Description

17.4.1 Flash Command Operations

Write operations are used for the program, erase, and erase verify algorithms described in this section. The program and erase algorithms are controlled by a state machine whose timebase FCLK is derived from the oscillator clock via a programmable divider. The FCMD register as well as the associated FADDR and FDATA registers operate as a buffer and a register (2-stage FIFO) so that a new command along with the necessary data and address can be stored to the buffer while the previous command is still in progress. This pipelined operation allows a time optimization when programming more than one word on a specific row, as the high voltage generation can be kept active in between two programming commands. The pipelined operation allows a simplification of command launching. Buffer empty as well as command completion are signalled by flags in the FSTAT register with corresponding interrupts generated, if enabled.

The next sections describe:

- How to write the FCLKDIV register
- Command write sequence used to program, erase or erase verify the Flash array
- Valid Flash commands
- Errors resulting from illegal Flash operations

17.4.1.1 Writing the FCLKDIV Register

Prior to issuing any Flash command after a reset, it is first necessary to write the FCLKDIV register to divide the oscillator clock down to within the 150-kHz to 200-kHz range. Since the program and erase timings are also a function of the bus clock, the FCLKDIV determination must take this information into account.

If we define:

- FCLK as the clock of the Flash timing control block
- Tbus as the period of the bus clock
- INT(x) as taking the integer part of x (e.g., INT(4.323) = 4),



Table 19-14. FSTAT Field Descriptions

Field	Description
5 PVIOL	 Protection Violation — The PVIOL flag indicates an attempt was made to program or erase an address in a protected Flash array memory area. The PVIOL flag is cleared by writing a 1 to PVIOL. Writing a 0 to the PVIOL flag has no effect on PVIOL. While PVIOL is set, it is not possible to launch another command. 0 No protection violation detected 1 Protection violation has occurred
4 ACCERR	Access Error — The ACCERR flag indicates an illegal access to the Flash array caused by either a violation of the command write sequence, issuing an illegal command (illegal combination of the CMDBx bits in the FCMD register) or the execution of a CPU STOP instruction while a command is executing (CCIF=0). The ACCERR flag is cleared by writing a 1 to ACCERR. Writing a 0 to the ACCERR flag has no effect on ACCERR. While ACCERR is set, it is not possible to launch another command. 0 No access error detected 1 Access error has occurred
2 BLANK	 Flash Array Has Been Verified as Erased — The BLANK flag indicates that an erase verify command has checked the Flash array and found it to be erased. The BLANK flag is cleared by hardware when CBEIF is cleared as part of a new valid command write sequence. Writing to the BLANK flag has no effect on BLANK. 0 If an erase verify command has been requested, and the CCIF flag is set, then a 0 in BLANK indicates the array is not erased 1 Flash array verifies as erased
1 FAIL	Flag Indicating a Failed Flash Operation — In special modes, the FAIL flag will set if the erase verify operation fails (Flash array verified as not erased). Writing a 0 to the FAIL flag has no effect on FAIL. The FAIL flag is cleared by writing a 1 to FAIL. While FAIL is set, it is not possible to launch another command. 0 Flash operation completed without error 1 Flash operation failed
0 DONE	 Flag Indicating a Failed Operation is not Active — In special modes, the DONE flag will clear if a program, erase, or erase verify operation is active. 0 Flash operation is active 1 Flash operation is not active

19.3.2.7 Flash Command Register (FCMD)

The FCMD register defines the Flash commands.



Figure 19-13. Flash Command Register (FCMD)

Bits CMDB6, CMDB5, CMDB2, and CMDB0 are readable and writable during a command write sequence while bits 7, 4, 3, and 1 read 0 and are not writable.

Module Base + 0x0006



19.3.2.14 RESERVED6

This register is reserved for factory testing and is not accessible to the user.

Module Base + 0x000F



Figure 19-23. RESERVED6

All bits read 0 and are not writable.

19.4 Functional Description

19.4.1 Flash Command Operations

Write operations are used for the program, erase, and erase verify algorithms described in this section. The program and erase algorithms are controlled by a state machine whose timebase FCLK is derived from the oscillator clock via a programmable divider. The FCMD register as well as the associated FADDR and FDATA registers operate as a buffer and a register (2-stage FIFO) so that a new command along with the necessary data and address can be stored to the buffer while the previous command is still in progress. This pipelined operation allows a time optimization when programming more than one word on a specific row, as the high voltage generation can be kept active in between two programming commands. The pipelined operation allows a simplification of command launching. Buffer empty as well as command completion are signalled by flags in the FSTAT register with corresponding interrupts generated, if enabled.

The next sections describe:

- How to write the FCLKDIV register
- Command write sequence used to program, erase or erase verify the Flash array
- Valid Flash commands
- Errors resulting from illegal Flash operations

19.4.1.1 Writing the FCLKDIV Register

Prior to issuing any Flash command after a reset, it is first necessary to write the FCLKDIV register to divide the oscillator clock down to within the 150-kHz to 200-kHz range. Since the program and erase timings are also a function of the bus clock, the FCLKDIV determination must take this information into account.

If we define:

• FCLK as the clock of the Flash timing control block



19.4.1.3 Valid Flash Commands

Table 19-17 summarizes the valid Flash commands along with the effects of the commands on the Flash array.

FCMD	Meaning	Function on Flash Array
0x05	Erase Verify	Verify all bytes in the Flash array are erased. If the Flash array is erased, the BLANK bit will set in the FSTAT register upon command completion.
0x20	Program	Program a word (2 bytes) in the Flash array.
0x40	Sector Erase	Erase all 1024 bytes in a sector of the Flash array.
0x41	Mass Erase	Erase all bytes in the Flash array. A mass erase of the full Flash array is only possible when FPLDIS, FPHDIS, and FPOPEN bits in the FPROT register are set prior to launching the command.

Table 19-17. Valid Flash Commands

CAUTION

A Flash word must be in the erased state before being programmed. Cumulative programming of bits within a Flash word is not allowed.



Figure 19-28. Example Mass Erase Command Flow



21.4.1.3.4 Mass Erase Command

The mass erase operation will erase all addresses in a Flash array using an embedded algorithm.

An example flow to execute the mass erase operation is shown in Figure 21-25. The mass erase command write sequence is as follows:

- 1. Write to a Flash array address to start the command write sequence for the mass erase command. The address and data written will be ignored.
- 2. Write the mass erase command, 0x41, to the FCMD register.
- 3. Clear the CBEIF flag in the FSTAT register by writing a 1 to CBEIF to launch the mass erase command.

If a Flash array to be erased contains any protected area, the PVIOL flag in the FSTAT register will set and the mass erase command will not launch. Once the mass erase command has successfully launched, the CCIF flag in the FSTAT register will set after the mass erase operation has completed unless a new command write sequence has been buffered.





----- Flash

A.6 SPI

This section provides electrical parametrics and ratings for the SPI.

In Table A-20 the measurement conditions are listed.

Table A-20. Measurement Conditions

Description	Value	Unit
Drive mode	Full drive mode	—
Load capacitance C _{LOAD,} on all outputs	50	pF
Thresholds for delay measurement points	(20% / 80%) V _{DDX}	V

A.6.1 Master Mode

In Figure A-6 the timing diagram for master mode with transmission format CPHA=0 is depicted.