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"Embedded - Microcontrollers" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

#### Applications of "<u>Embedded -</u> <u>Microcontrollers</u>"

#### Details

Product Status	Active
Core Processor	8051
Core Size	8-Bit
Speed	48 MIPS
Connectivity	I <sup>2</sup> C, SPI, UART/USART, USB
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	25
Program Memory Size	64KB (64K x 8)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	4.25K x 8
Voltage - Supply (Vcc/Vdd)	2.7V ~ 5.25V
Data Converters	-
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	32-VFQFN Exposed Pad
Supplier Device Package	32-QFN (5x5)
Purchase URL	https://www.e-xfl.com/product-detail/silicon-labs/c8051f385-gm

Email: info@E-XFL.COM

Address: Room A, 16/F, Full Win Commercial Centre, 573 Nathan Road, Mongkok, Hong Kong

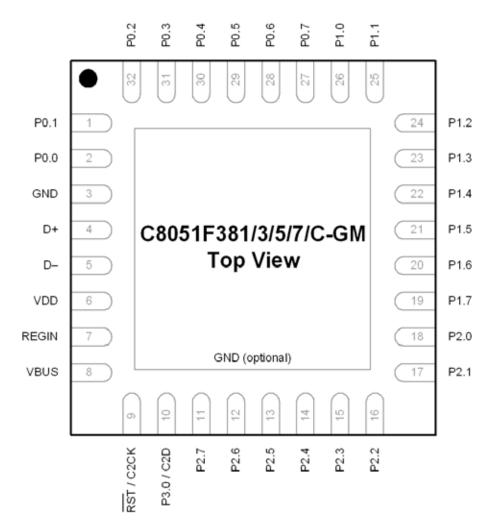


Figure 3.7. QFN-32 Pinout Diagram (Top View)



# SFR Definition 6.2. ADC0H: ADC0 Data Word MSB

Bit	7	6	5	4	3	2	1	0	
Name		ADC0H[7:0]							
Туре				R/	W				
Reset	0	0 0 0 0 0 0 0 0							

SFR Address = 0xBE; SFR Page = All Pages

Bit	Name	Function
7:0	ADC0H[7:0]	ADC0 Data Word High-Order Bits.
		For AD0LJST = 0: Bits 7–2 will read 000000b. Bits 1–0 are the upper 2 bits of the 10- bit ADC0 Data Word.
		For AD0LJST = 1: Bits 7–0 are the most-significant bits of the 10-bit ADC0 Data Word.

#### SFR Definition 6.3. ADC0L: ADC0 Data Word LSB

Bit	7	6	5	4	3	2	1	0	
Name		ADC0L[7:0]							
Туре				R/	W				
Reset	0	0 0 0 0 0 0 0 0							

#### SFR Address = 0xBD; SFR Page = All Pages

Bit	Name	Function
7:0		ADC0 Data Word Low-Order Bits. For AD0LJST = 0: Bits 7–0 are the lower 8 bits of the 10-bit Data Word. For AD0LJST = 1: Bits 7–6 are the lower 2 bits of the 10-bit Data Word. Bits 5–0 will read 000000b.



With the CIP-51's maximum system clock at 48 MHz, it has a peak throughput of 48 MIPS. The CIP-51 has a total of 109 instructions. The table below shows the total number of instructions that require each execution time.

Clocks to Execute	1	2	2/4	3	3/5	4	5	4/6	6	8
Number of Instructions	26	50	5	10	6	5	2	2	2	1

#### **Programming and Debugging Support**

In-system programming of the Flash program memory and communication with on-chip debug support logic is accomplished via the Silicon Labs 2-Wire Development Interface (C2).

The on-chip debug support logic facilitates full speed in-circuit debugging, allowing the setting of hardware breakpoints, starting, stopping and single stepping through program execution (including interrupt service routines), examination of the program's call stack, and reading/writing the contents of registers and memory. This method of on-chip debugging is completely non-intrusive, requiring no RAM, Stack, timers, or other on-chip resources. C2 details can be found in Section "28. C2 Interface" on page 316.

The CIP-51 is supported by development tools from Silicon Labs and third party vendors. Silicon Labs provides an integrated development environment (IDE) including editor, debugger and programmer. The IDE's debugger and programmer interface to the CIP-51 via the C2 interface to provide fast and efficient in-system device programming and debugging. Third party macro assemblers and C compilers are also available.

#### 11.1. Instruction Set

The instruction set of the CIP-51 System Controller is fully compatible with the standard MCS-51<sup>™</sup> instruction set. Standard 8051 development tools can be used to develop software for the CIP-51. All CIP-51 instructions are the binary and functional equivalent of their MCS-51<sup>™</sup> counterparts, including opcodes, addressing modes and effect on PSW flags. However, instruction timing is different than that of the standard 8051.

#### 11.1.1. Instruction and CPU Timing

In many 8051 implementations, a distinction is made between machine cycles and clock cycles, with machine cycles varying from 2 to 12 clock cycles in length. However, the CIP-51 implementation is based solely on clock cycle timing. All instruction timings are specified in terms of clock cycles.

Due to the pipelined architecture of the CIP-51, most instructions execute in the same number of clock cycles as there are program bytes in the instruction. Conditional branch instructions take one less clock cycle to complete when the branch is not taken as opposed to when the branch is taken. Table 11.1 is the CIP-51 Instruction Set Summary, which includes the mnemonic, number of bytes, and number of clock cycles for each instruction.



# SFR Definition 11.6. PSW: Program Status Word

Bit	7	6	5	4	3	2	1	0	
Nam	e CY	AC	F0	RS	[1:0]	OV	F1	PARITY	
Туре	R/W	R/W	R/W	R	/W	R/W	R/W	R	
Rese	et 0	0	0	0	0	0	0	0	
SFR Address = 0xD0; SFR Page = All Pages; Bit-Addressable									
Bit	Name				Function				
7	CY	Carry Flag.							
		This bit is set row (subtraction						n) or a bor-	
6	AC	Auxiliary Car	ry Flag.						
		This bit is set							
		borrow from (s metic operatio	,	the high ord	er nibble. It	is cleared to I	logic 0 by al	I other arith-	
5	F0	User Flag 0.							
		This is a bit-ad	ddressable,	general purp	ose flag for	use under so	oftware cont	rol.	
4:3	RS[1:0]	Register Ban	k Select.						
		These bits sel		-	s used durir	ng register ac	cesses.		
		00: Bank 0, Ao 01: Bank 1, Ao							
		10: Bank 2, A							
		11: Bank 3, Ac	dresses 0x	18-0x1F					
2	OV	<b>Overflow Flag</b>	g.						
		This bit is set		-					
						n-change ove greater than 2			
				es a divide-by	•	•	00).		
		The OV bit is other cases.	cleared to 0	by the ADD,	ADDC, SU	BB, MUL, and	d DIV instru	ctions in all	
1	F1	User Flag 1.							
		This is a bit-ad	ddressable,	general purp	ose flag for	use under so	oftware cont	rol.	
0	PARITY	Parity Flag.							
		This bit is set t if the sum is e	-	ne sum of the	e eight bits ir	n the accumu	lator is odd	and cleared	



## SFR Definition 14.1. EMI0CN: External Memory Interface Control

	<b>F</b>									
Bit	7	6	5	4	3	2	1	0		
Name		PGSEL[7:0]								
Туре				R/	W					
Reset	0	0 0 0 0 0 0 0 0								

SFR Address = 0xAA; SFR Page = All Pages

Bit	Name	Function
7:0	PGSEL[7:0]	XRAM Page Select Bits.
		The XRAM Page Select Bits provide the high byte of the 16-bit external data mem- ory address when using an 8-bit MOVX command, effectively selecting a 256-byte page of RAM. 0x00: 0x0000 to 0x00FF 0x01: 0x0100 to 0x01FF  0xFE: 0xFE00 to 0xFEFF 0xFF: 0xFF00 to 0xFFFF



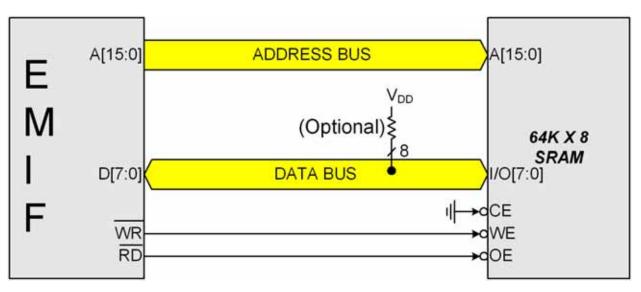


Figure 14.3. Non-multiplexed Configuration Example



Address	Page	0(8)	1(9)	2(A)	3(B)	4(C)	5(D)	6(E)	7(F)
F8		SPI0CN	PCA0L	PCA0H	PCA0CPL0	PCA0CPH0	PCA0CPL4	PCA0CPH4	VDM0CN
F0		В	P0MDIN	P1MDIN	P2MDIN	P3MDIN	P4MDIN	EIP1	EIP2
E8		ADC0CN	PCA0CPL1	PCA0CPH1	PCA0CPL2	PCA0CPH2	PCA0CPL3	PCA0CPH3	RSTSRC
E0	0 F	ACC	XBR0	XBR1	XBR2	IT01CF CKCON1	SMOD1	EIE1	EIE2
D8		PCA0CN	PCA0MD	PCA0CPM0	PCA0CPM1	PCA0CPM2	PCA0CPM3	PCA0CPM4	P3SKIP
D0		PSW	REF0CN	SCON1	SBUF1	P0SKIP	P1SKIP	P2SKIP	<b>USB0XCN</b>
C8	0	TMR2CN	REG01CN	TMR2RLL	TMR2RLH	TMR2L	TMR2H	SMB0ADM	SMB0ADR
	F	TMR5CN	REGUICIN	TMR5RLL	TMR5RLH	TMR5L	TMR5H	SMB1ADM	SMB1ADR
CO	0	SMB0CN	SMB0CF	SMB0DAT	ADC0GTL	ADC0GTH	ADC0LTL	ADC0LTH	P4
	F	SMB1CN	SMB1CF	SMB1DAT	ADCOGIL	ADCOGIN /	ADCULIL	ADCULIN	Γ4
B8	0	IP	CLKMUL	AMX0N	AMX0P	ADC0CF	ADC0L	ADC0H	SFRPAGE
Бо	F		SMBTC		AIVIAUE		ADCOL	ADCOIT	SERFAGE
B0		P3	OSCXCN	OSCICN	OSCICL	SBRLL1	SBRLH1	FLSCL	FLKEY
A8		IE	CLKSEL	EMIOCN		SBCON1		P4MDOUT	PFE0CN
A0		P2	SPI0CFG	SPI0CKR	SPI0DAT	POMDOUT	P1MDOUT	P2MDOUT	P3MDOUT
98		SCON0	SBUF0	CPT1CN	CPT0CN	CPT1MD	CPT0MD	CPT1MX	CPT0MX
90	0	P1	TMR3CN	TMR3RLL	TMR3RLH	TMR3L	TMR3H	USB0ADR	USB0DAT
90	F		TMR4CN	TMR4RLL	TMR4RLH	TMR4L	TMR4H	USBUADR	USBUDAI
88		TCON	TMOD	TL0	TL1	TH0	TH1	CKCON	PSCTL
80		P0	SP	DPL	DPH	EMI0TC	EMI0CF	OSCLCN	PCON
		0(8)	1(9)	2(A)	3(B)	4(C)	5(D)	6(E)	7(F)
Note					-	-			

#### Table 15.1. Special Function Register (SFR) Memory Map

Notes:

1. SFR Addresses ending in 0x0 or 0x8 are bit-addressable locations and can be used with bitwise instructions.

2. Unless indicated otherwise, SFRs are available on both page 0 and page F.



## 16.3. INTO and INT1 External Interrupt Sources

The INTO and INT1 external interrupt sources are configurable as active high or low, edge or level sensitive. The INOPL (INTO Polarity) and IN1PL (INT1 Polarity) bits in the IT01CF register select active high or active low; the IT0 and IT1 bits in TCON (Section "26.1. Timer 0 and Timer 1" on page 266) select level or edge sensitive. The table below lists the possible configurations.

IT0	IN0PL	INT0 Interrupt
1	0	Active low, edge sensitive
1	1	Active high, edge sensitive
0	0	Active low, level sensitive
0	1	Active high, level sensitive

IT1	IN1PL	INT1 Interrupt
1	0	Active low, edge sensitive
1	1	Active high, edge sensitive
0	0	Active low, level sensitive
0	1	Active high, level sensitive

INTO and INT1 are assigned to Port pins as defined in the IT01CF register (see SFR Definition 16.7). Note that INT0 and INT0 Port pin assignments are independent of any Crossbar assignments. INT0 and INT1 will monitor their assigned Port pins without disturbing the peripheral that was assigned the Port pin via the Crossbar. To assign a Port pin only to INT0 and/or INT1, configure the Crossbar to skip the selected pin(s). This is accomplished by setting the associated bit in register PnSKIP (see Section "20.1. Priority Crossbar Decoder" on page 154 for complete details on configuring the Crossbar).

IE0 (TCON.1) and IE1 (TCON.3) serve as the interrupt-pending flags for the INTO and INT1 external interrupts, respectively. If an INT0 or INT1 external interrupt is configured as edge-sensitive, the corresponding interrupt-pending flag is automatically cleared by the hardware when the CPU vectors to the ISR. When configured as level sensitive, the interrupt-pending flag remains logic 1 while the input is active as defined by the corresponding polarity bit (INOPL or IN1PL); the flag remains logic 0 while the input is inactive. The external interrupt source must hold the input active until the interrupt request is recognized. It must then deactivate the interrupt request before execution of the ISR completes or another interrupt request will be generated.



# SFR Definition 19.6. OSCXCN: External Oscillator Control

Bit	7	6	5	4	3	2	1	0
Name	XCLKVLD	XOSCMD[2:0]				XFCN[2:0]		
Туре	R	R/W			R	R/W		
Reset	0	0	0	0	0	0	0	0

#### SFR Address = 0xB1; SFR Page = All Pages

Bit	Name	Function					
7	XCLKVLD	<ul> <li>External Oscillator Valid Flag.</li> <li>Provides External Oscillator status and is valid at all times for all modes of operation except External CMOS Clock Mode and External CMOS Clock Mode with divide by 2. In these modes, XCLKVLD always returns 0.</li> <li>0: External Oscillator is unused or not yet stable.</li> <li>1: External Oscillator is running and stable.</li> </ul>					
6:4	XOSCMD[2:0]	External Oscillator Mode Select.00x: External Oscillator circuit off.010: External CMOS Clock Mode.011: External CMOS Clock Mode with divide-by-2 stage.100: RC Oscillator Mode with divide-by-2 stage.101: Capacitor Oscillator Mode with divide-by-2 stage.110: Crystal Oscillator Mode.111: Crystal Oscillator Mode with divide-by-2 stage.					
3	Unused	Read = 0; Write = don't care					
2:0	XFCN[2:0]	External Oscillator Frequency Control Bits. Set according to the desired frequency for RC mode. Set according to the desired K Factor for C mode.					
		XFCN Crystal Mode RC Mode C Mo					
		000	f ≤ 20 kHz	f ≤ 25 kHz	K Factor = 0.87		
		001	20 kHz < f $\leq$ 58 kHz	$25 \text{ kHz} < f \le 50 \text{ kHz}$	K Factor = 2.6		
		010	58 kHz < f ≤ 155 kHz	50 kHz < f $\leq$ 100 kHz	K Factor = 7.7		
		011	155 kHz $<$ f $\leq$ 415 kHz	100 kHz $<$ f $\leq$ 200 kHz	K Factor = 22		
		100	415 kHz $<$ f $\leq$ 1.1 MHz	$200 \text{ kHz} < f \leq 400 \text{ kHz}$	K Factor = 65		
		101	$1.1 \text{ MHz} < f \le 3.1 \text{ MHz}$	400 kHz $<$ f $\le$ 800 kHz	K Factor = 180		
		110	$3.1 \text{ MHz} < f \le 8.2 \text{ MHz}$	800 kHz $< f \le 1.6$ MHz	K Factor = 664		
		111	111         8.2 MHz < f $\leq$ 25 MHz         1.6 MHz < f $\leq$ 3.2 MHz         K Factor = 1590				



### 21.9. The Serial Interface Engine

The Serial Interface Engine (SIE) performs all low level USB protocol tasks, interrupting the processor when data has successfully been transmitted or received. When receiving data, the SIE will interrupt the processor when a complete data packet has been received; appropriate handshaking signals are automatically generated by the SIE. When transmitting data, the SIE will interrupt the processor when a complete data packet has been received; appropriate handshaking signals are automatically generated by the SIE. When transmitting data, the SIE will interrupt the processor when a complete data packet has been transmitted and the appropriate handshake signal has been received.

The SIE will not interrupt the processor when corrupted/erroneous packets are received.

### 21.10. Endpoint0

Endpoint0 is managed through the USB register E0CSR (USB Register Definition 21.18). The INDEX register must be loaded with 0x00 to access the E0CSR register.

An Endpoint0 interrupt is generated when:

- 1. A data packet (OUT or SETUP) has been received and loaded into the Endpoint0 FIFO. The OPRDY bit (E0CSR.0) is set to 1 by hardware.
- 2. An IN data packet has successfully been unloaded from the Endpoint0 FIFO and transmitted to the host; INPRDY is reset to 0 by hardware.
- 3. An IN transaction is completed (this interrupt generated during the status stage of the transaction).
- 4. Hardware sets the STSTL bit (E0CSR.2) after a control transaction ended due to a protocol violation.
- 5. Hardware sets the SUEND bit (E0CSR.4) because a control transfer ended before firmware sets the DATAEND bit (E0CSR.3).

The E0CNT register (USB Register Definition 21.11) holds the number of received data bytes in the Endpoint0 FIFO.

Hardware will automatically detect protocol errors and send a STALL condition in response. Firmware may force a STALL condition to abort the current transfer. When a STALL condition is generated, the STSTL bit will be set to 1 and an interrupt generated. The following conditions will cause hardware to generate a STALL condition:

- 1. The host sends an OUT token during a OUT data phase after the DATAEND bit has been set to 1.
- 2. The host sends an IN token during an IN data phase after the DATAEND bit has been set to 1.
- 3. The host sends a packet that exceeds the maximum packet size for Endpoint0.
- 4. The host sends a non-zero length DATA1 packet during the status phase of an IN transaction.

Firmware sets the SDSTL bit (E0CSR.5) to 1.

#### 21.10.1. Endpoint0 SETUP Transactions

All control transfers must begin with a SETUP packet. SETUP packets are similar to OUT packets, containing an 8-byte data field sent by the host. Any SETUP packet containing a command field of anything other than 8 bytes will be automatically rejected by USB0. An Endpoint0 interrupt is generated when the data from a SETUP packet is loaded into the Endpoint0 FIFO. Software should unload the command from the Endpoint0 FIFO, decode the command, perform any necessary tasks, and set the SOPRDY bit to indicate that it has serviced the OUT packet.

#### 21.10.2. Endpoint0 IN Transactions

When a SETUP request is received that requires USB0 to transmit data to the host, one or more IN requests will be sent by the host. For the first IN transaction, firmware should load an IN packet into the Endpoint0 FIFO, and set the INPRDY bit (E0CSR.1). An interrupt will be generated when an IN packet is transmitted successfully. Note that no interrupt will be generated if an IN request is received before firm-



All transactions are initiated by a master, with one or more addressed slave devices as the target. The master generates the START condition and then transmits the slave address and direction bit. If the transaction is a WRITE operation from the master to the slave, the master transmits the data a byte at a time waiting for an ACK from the slave at the end of each byte. For READ operations, the slave transmits the data waiting for an ACK from the master at the end of each byte. At the end of the data transfer, the master generates a STOP condition to terminate the transaction and free the bus. Figure 22.3 illustrates a typical SMBus transaction.

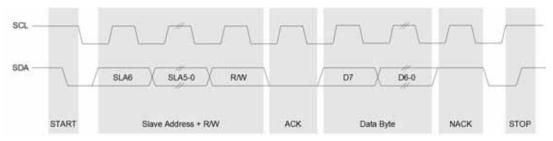


Figure 22.3. SMBus Transaction

#### 22.3.1. Transmitter Vs. Receiver

On the SMBus communications interface, a device is the "transmitter" when it is sending an address or data byte to another device on the bus. A device is a "receiver" when an address or data byte is being sent to it from another device on the bus. The transmitter controls the SDA line during the address or data byte. After each byte of address or data information is sent by the transmitter, the receiver sends an ACK or NACK bit during the ACK phase of the transfer, during which time the receiver controls the SDA line.

#### 22.3.2. Arbitration

A master may start a transfer only if the bus is free. The bus is free after a STOP condition or after the SCL and SDA lines remain high for a specified time (see Section "22.3.5. SCL High (SMBus Free) Timeout" on page 208). In the event that two or more devices attempt to begin a transfer at the same time, an arbitration scheme is employed to force one master to give up the bus. The master devices continue transmitting until one attempts a HIGH while the other transmits a LOW. Since the bus is open-drain, the bus will be pulled LOW. The master attempting the HIGH will detect a LOW SDA and lose the arbitration. The winning master continues its transmission without interruption; the losing master becomes a slave and receives the rest of the transfer if addressed. This arbitration scheme is non-destructive: one device always wins, and no data is lost.

#### 22.3.3. Clock Low Extension

SMBus provides a clock synchronization mechanism, similar to I2C, which allows devices with different speed capabilities to coexist on the bus. A clock-low extension is used during a transfer in order to allow slower slave devices to communicate with faster masters. The slave may temporarily hold the SCL line LOW to extend the clock low period, effectively decreasing the serial clock frequency.

#### 22.3.4. SCL Low Timeout

If the SCL line is held low by a slave device on the bus, no further communication is possible. Furthermore, the master cannot force the SCL line high to correct the error condition. To solve this problem, the SMBus protocol specifies that devices participating in a transfer must detect any clock cycle held low longer than 25 ms as a "timeout" condition. Devices that have detected the timeout condition must reset the communication no later than 10 ms after detecting the timeout condition.

For the SMBus0 interface, Timer 3 is used to implement SCL low timeouts. Timer 4 is used on the SMBus1 interface for SCL low timeouts. The SCL low timeout feature is enabled by setting the SMBnTOE bit in SMBnCF. The associated timer is forced to reload when SCL is high, and allowed to count when SCL is



#### 22.5.3. Write Sequence (Slave)

During a write sequence, an SMBus master writes data to a slave device. The slave in this transfer will be a receiver during the address byte, and a receiver during all data bytes. When slave events are enabled (INH = 0), the interface enters Slave Receiver Mode when a START followed by a slave address and direction bit (WRITE in this case) is received. If hardware ACK generation is disabled, upon entering Slave Receiver Mode, an interrupt is generated and the ACKRQ bit is set. The software must respond to the received slave address with an ACK, or ignore the received slave address with a NACK. If hardware ACK generation is enabled, the hardware will apply the ACK for a slave address which matches the criteria set up by SMB0ADR and SMB0ADM. The interrupt will occur after the ACK cycle.

If the received slave address is ignored (by software or hardware), slave interrupts will be inhibited until the next START is detected. If the received slave address is acknowledged, zero or more data bytes are received.

If hardware ACK generation is disabled, the ACKRQ is set to 1 and an interrupt is generated after each received byte. Software must write the ACK bit at that time to ACK or NACK the received byte.

With hardware ACK generation enabled, the SMBus hardware will automatically generate the ACK/NACK, and then post the interrupt. It is important to note that the appropriate ACK or NACK value should be set up by the software prior to receiving the byte when hardware ACK generation is enabled.

The interface exits Slave Receiver Mode after receiving a STOP. The interface will switch to Slave Transmitter Mode if SMB0DAT is written while an active Slave Receiver. Figure 22.7 shows a typical slave write sequence. Two received data bytes are shown, though any number of bytes may be received. Notice that the 'data byte transferred' interrupts occur at different places in the sequence, depending on whether hardware ACK generation is enabled. The interrupt occurs **before** the ACK with hardware ACK generation disabled, and **after** the ACK when hardware ACK generation is enabled.

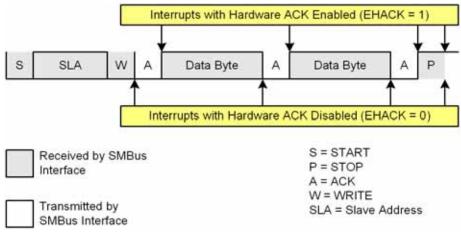


Figure 22.7. Typical Slave Write Sequence



# 23. UART0

UART0 is an asynchronous, full duplex serial port offering modes 1 and 3 of the standard 8051 UART. Enhanced baud rate support allows a wide range of clock sources to generate standard baud rates (details in Section "23.1. Enhanced Baud Rate Generation" on page 233). Received data buffering allows UART0 to start reception of a second incoming data byte before software has finished reading the previous data byte.

UART0 has two associated SFRs: Serial Control Register 0 (SCON0) and Serial Data Buffer 0 (SBUF0). The single SBUF0 location provides access to both transmit and receive registers. Writes to SBUF0 always access the Transmit register. Reads of SBUF0 always access the buffered Receive register; it is not possible to read data from the Transmit register.

With UART0 interrupts enabled, an interrupt is generated each time a transmit is completed (TI0 is set in SCON0), or a data byte has been received (RI0 is set in SCON0). The UART0 interrupt flags are not cleared by hardware when the CPU vectors to the interrupt service routine. They must be cleared manually by software, allowing software to determine the cause of the UART0 interrupt (transmit complete or receive complete).

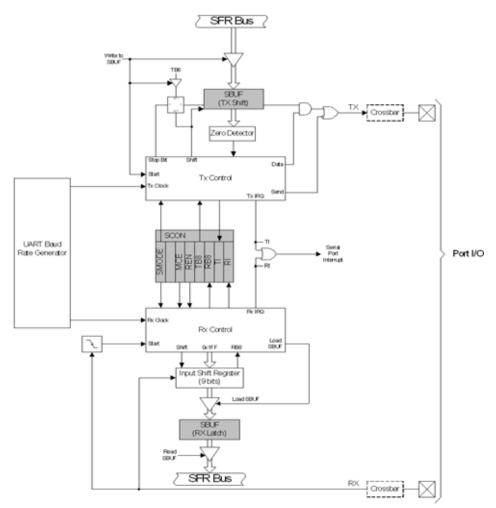


Figure 23.1. UART0 Block Diagram



# 24. UART1

UART1 is an asynchronous, full duplex serial port offering a variety of data formatting options. A dedicated baud rate generator with a 16-bit timer and selectable prescaler is included, which can generate a wide range of baud rates (details in Section "24.1. Baud Rate Generator" on page 241). A received data FIFO allows UART1 to receive up to three data bytes before data is lost and an overflow occurs.

UART1 has six associated SFRs. Three are used for the Baud Rate Generator (SBCON1, SBRLH1, and SBRLL1), two are used for data formatting, control, and status functions (SCON1, SMOD1), and one is used to send and receive data (SBUF1). The single SBUF1 location provides access to both the transmit holding register and the receive FIFO. Writes to SBUF1 always access the Transmit Holding Register. Reads of SBUF1 always access the first byte of the Receive FIFO; it is not possible to read data from the Transmit Holding Register.

With UART1 interrupts enabled, an interrupt is generated each time a transmit is completed (TI1 is set in SCON1), or a data byte has been received (RI1 is set in SCON1). The UART1 interrupt flags are not cleared by hardware when the CPU vectors to the interrupt service routine. They must be cleared manually by software, allowing software to determine the cause of the UART1 interrupt (transmit complete or receive complete). Note that if additional bytes are available in the Receive FIFO, the RI1 bit cannot be cleared by software.

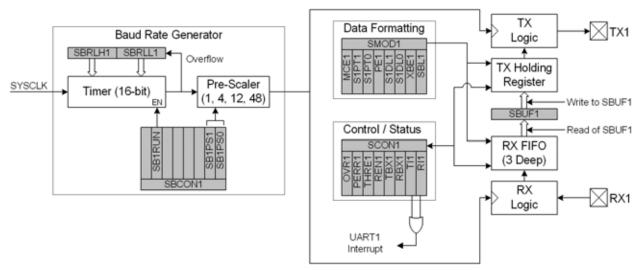


Figure 24.1. UART1 Block Diagram



# SFR Definition 24.2. SMOD1: UART1 Mode

Bit	7	6	5	4	3	2	1	0		
Nam	e MCE1	S1P	[1:0]	PE1	S1D	L[1:0]	XBE1	SBL1		
Туре	e R/W	R/	W	R/W	R	/W	R/W	R/W		
Rese	et 0	0	0	0	1	1	0	0		
SFR A	Address = 0	E5; SFR Page = All Pages								
Bit	Name	Function								
7	MCE1	Multiprocessor Communication Enable.								
		<ul> <li>0: RI will be activated if stop bit(s) are 1.</li> <li>1: RI will be activated if stop bit(s) and extra bit are 1 (extra bit must be enabled using XBE1).</li> <li>Note: This function is not available when hardware parity is enabled.</li> </ul>								
6:5	S1PT[1:0]	Parity Type Bits.								
		00: Odd								
		01: Even 10: Mark								
		11: Space								
4	PE1	Parity Enable.								
		<ul> <li>This bit activates hardware parity generation and checking. The parity type is selected by bits S1PT1-0 when parity is enabled.</li> <li>0: Hardware parity is disabled.</li> <li>1: Hardware parity is enabled.</li> </ul>								
3:2	S1DL[1:0]	Data Length.								
		00: 5-bit data								
		01: 6-bit data 10: 7-bit data								
		11: 8-bit data								
1	XBE1	Extra Bit Enable.								
		When enabled, the value of TBX1 will be appended to the data field.								
		0: Extra Bit Di 1: Extra Bit Er								
0	SBL1									
		Stop Bit Length. 0: Short—Stop bit is active for one bit time.								
		1: Long—Stop (data length =	bit is active			ngth = 6, 7, o	or 8 bits), or 1	.5 bit times		



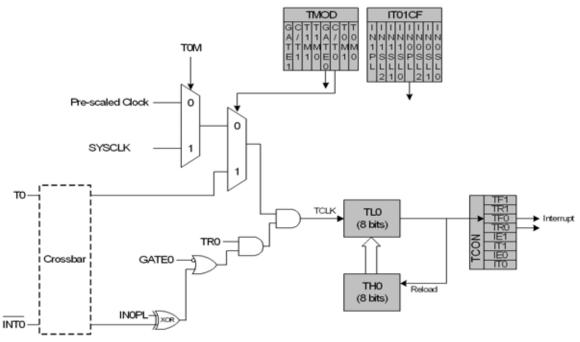


Figure 26.2. T0 Mode 2 Block Diagram

#### 26.1.4. Mode 3: Two 8-bit Counter/Timers (Timer 0 Only)

In Mode 3, Timer 0 is configured as two separate 8-bit counter/timers held in TL0 and TH0. The counter/timer in TL0 is controlled using the Timer 0 control/status bits in TCON and TMOD: TR0, C/T0, GATE0 and TF0. TL0 can use either the system clock or an external input signal as its timebase. The TH0 register is restricted to a timer function sourced by the system clock or prescaled clock. TH0 is enabled using the Timer 1 run control bit TR1. TH0 sets the Timer 1 overflow flag TF1 on overflow and thus controls the Timer 1 interrupt.

Timer 1 is inactive in Mode 3. When Timer 0 is operating in Mode 3, Timer 1 can be operated in Modes 0, 1 or 2, but cannot be clocked by external signals nor set the TF1 flag and generate an interrupt. However, the Timer 1 overflow can be used to generate baud rates or overflow conditions for other peripherals. While Timer 0 is operating in Mode 3, Timer 1 run control is handled through its mode settings. To run Timer 1 while Timer 0 is in Mode 3, set the Timer 1 Mode as 0, 1, or 2. To disable Timer 1, configure it for Mode 3.



### 26.3. Timer 3

Timer 3 is a 16-bit timer formed by two 8-bit SFRs: TMR3L (low byte) and TMR3H (high byte). Timer 3 may operate in 16-bit auto-reload mode, (split) 8-bit auto-reload mode, USB Start-of-Frame (SOF) capture mode, or Low-Frequency Oscillator (LFO) Rising Edge capture mode. The Timer 3 operation mode is defined by the T3SPLIT (TMR3CN.3), T3CE (TMR3CN.4) bits, and T3CSS (TMR3CN.1) bits.

Timer 3 may be clocked by the system clock, the system clock divided by 12, or the external oscillator source divided by 8. The external clock mode is ideal for real-time clock (RTC) functionality, where the internal oscillator drives the system clock while Timer 3 (and/or the PCA) is clocked by an external precision oscillator. Note that the external oscillator source divided by 8 is synchronized with the system clock.

#### 26.3.1. 16-bit Timer with Auto-Reload

When T3SPLIT (TMR3CN.3) is zero, Timer 3 operates as a 16-bit timer with auto-reload. Timer 3 can be clocked by SYSCLK, SYSCLK divided by 12, or the external oscillator clock source divided by 8. As the 16-bit timer register increments and overflows from 0xFFFF to 0x0000, the 16-bit value in the Timer 3 reload registers (TMR3RLH and TMR3RLL) is loaded into the Timer 3 register as shown in Figure 26.8, and the Timer 3 High Byte Overflow Flag (TMR3CN.7) is set. If Timer 3 interrupts are enabled (if EIE1.7 is set), an interrupt will be generated on each Timer 3 overflow. Additionally, if Timer 3 interrupts are enabled and the TF3LEN bit is set (TMR3CN.5), an interrupt will be generated each time the lower 8 bits (TMR3L) overflow from 0xFF to 0x00.

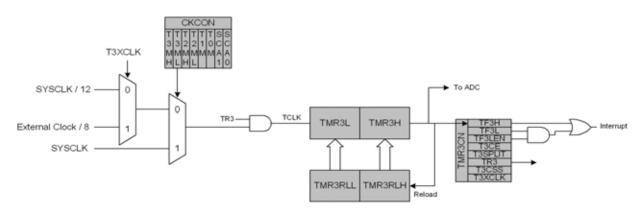


Figure 26.8. Timer 3 16-Bit Mode Block Diagram



## 27.2. PCA0 Interrupt Sources

Figure 27.3 shows a diagram of the PCA interrupt tree. There are six independent event flags that can be used to generate a PCA0 interrupt. They are: the main PCA counter overflow flag (CF), which is set upon a 16-bit overflow of the PCA0 counter and the individual flags for each PCA channel (CCF0, CCF1, CCF2, CCF3, and CCF4), which are set according to the operation mode of that module. These event flags are always set when the trigger condition occurs. Each of these flags can be individually selected to generate a PCA0 interrupt, using the corresponding interrupt enable flag (ECF for CF, and ECCFn for each CCFn). PCA0 interrupts must be globally enabled before any individual interrupt sources are recognized by the processor. PCA0 interrupts are globally enabled by setting the EA bit and the EPCA0 bit to logic 1.

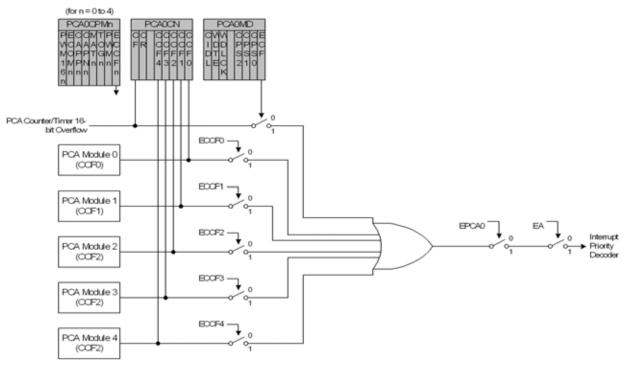


Figure 27.3. PCA Interrupt Block Diagram



### 27.4. Watchdog Timer Mode

A programmable watchdog timer (WDT) function is available through the PCA Module 4. The WDT is used to generate a reset if the time between writes to the WDT update register (PCA0CPH4) exceed a specified limit. The WDT can be configured and enabled/disabled as needed by software.

With the WDTE bit set in the PCA0MD register, Module 4 operates as a watchdog timer (WDT). The Module 4 high byte is compared to the PCA counter high byte; the Module 4 low byte holds the offset to be used when WDT updates are performed. The Watchdog Timer is enabled on reset. Writes to some PCA registers are restricted while the Watchdog Timer is enabled. The WDT will generate a reset shortly after code begins execution. To avoid this reset, the WDT should be explicitly disabled (and optionally re-configured and re-enabled if it is used in the system).

#### 27.4.1. Watchdog Timer Operation

While the WDT is enabled:

- PCA counter is forced on.
- Writes to PCA0L and PCA0H are not allowed.
- PCA clock source bits (CPS2–CPS0) are frozen.
- PCA Idle control bit (CIDL) is frozen.
- Module 4 is forced into software timer mode.
- Writes to the Module 4 mode register (PCA0CPM4) are disabled.

While the WDT is enabled, writes to the CR bit will not change the PCA counter state; the counter will run until the WDT is disabled. The PCA counter run control bit (CR) will read zero if the WDT is enabled but user software has not enabled the PCA counter. If a match occurs between PCA0CPH4 and PCA0H while the WDT is enabled, a reset will be generated. To prevent a WDT reset, the WDT may be updated with a write of any value to PCA0CPH4. Upon a PCA0CPH4 write, PCA0H plus the offset held in PCA0CPL4 is loaded into PCA0CPH4 (See Figure 27.10).

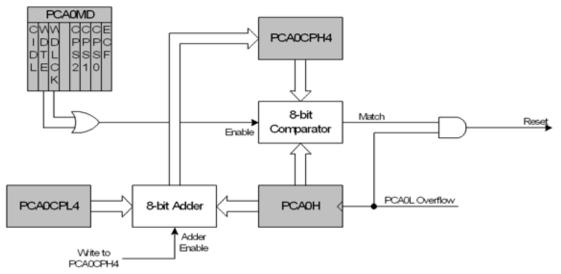


Figure 27.10. PCA Module 4 with Watchdog Timer Enabled





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