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Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	4MHz
Connectivity	-
Peripherals	Brown-out Detect/Reset, POR, WDT
Number of I/O	13
Program Memory Size	896B (512 x 14)
Program Memory Type	OTP
EEPROM Size	128 x 8
RAM Size	96 x 8
Voltage - Supply (Vcc/Vdd)	3V ~ 5.5V
Data Converters	-
Oscillator Type	External
Operating Temperature	-40°C ~ 125°C (TA)
Mounting Type	Surface Mount
Package / Case	20-SSOP (0.209", 5.30mm Width)
Supplier Device Package	20-SSOP
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic16ce623-04e-ss

NOTES:

3.1 Clocking Scheme/Instruction Cycle

The clock input (OSC1/CLKIN pin) is internally divided by four to generate four non-overlapping quadrature clocks namely Q1, Q2, Q3 and Q4. Internally, the program counter (PC) is incremented every Q1, the instruction is fetched from the program memory and latched into the instruction register in Q4. The instruction is decoded and executed during the following Q1 through Q4. The clocks and instruction execution flow is shown in Figure 3-2.

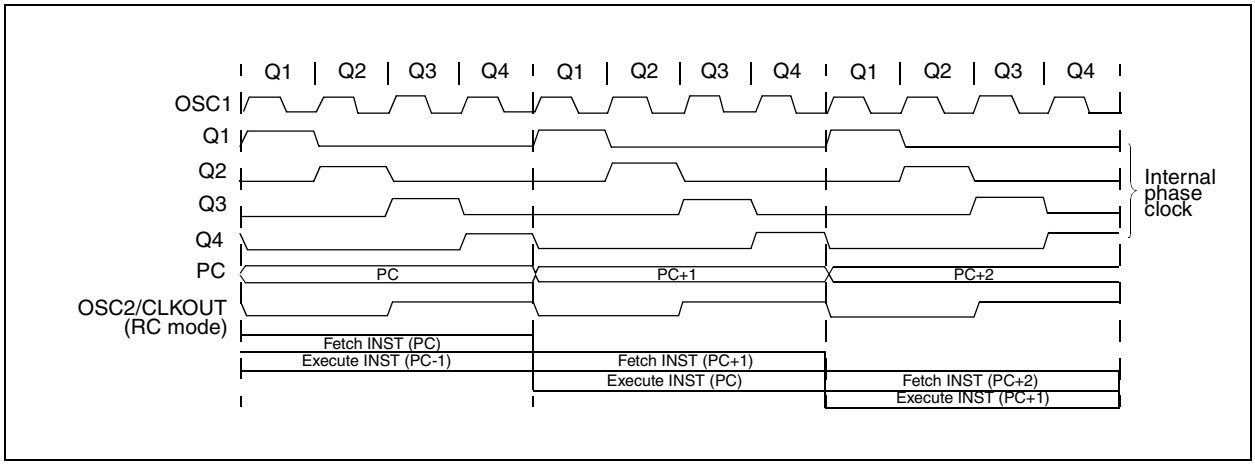
3.2 Instruction Flow/Pipelining

An "Instruction Cycle" consists of four Q cycles (Q1, Q2, Q3 and Q4). The instruction fetch and execute are pipelined such that fetch takes one instruction cycle, while decode and execute takes another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the program counter to change (i.e., GOTO) then two cycles are required to complete the instruction (Example 3-1).

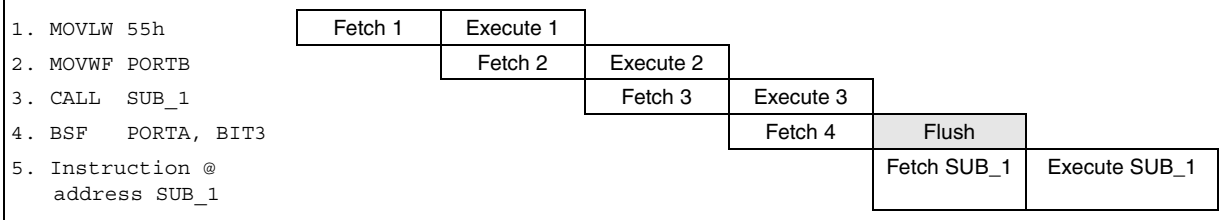
A fetch cycle begins with the program counter (PC) incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the "Instruction Register (IR)" in cycle Q1. This instruction is then decoded and executed during the Q2, Q3, and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).

FIGURE 3-2: CLOCK/INSTRUCTION CYCLE



EXAMPLE 3-1: INSTRUCTION PIPELINE FLOW



All instructions are single cycle, except for any program branches. These take two cycles since the fetch instruction is "flushed" from the pipeline, while the new instruction is being fetched and then executed.

FIGURE 4-4: DATA MEMORY MAP FOR THE PIC16CE623/624

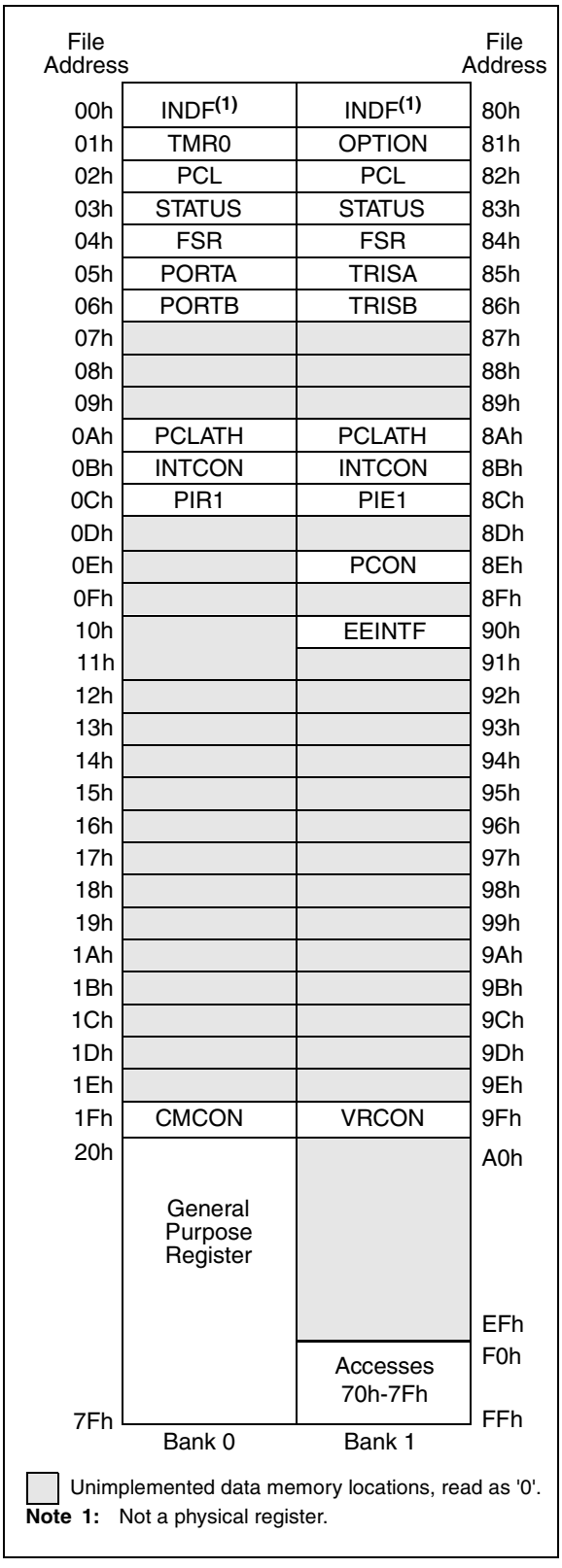
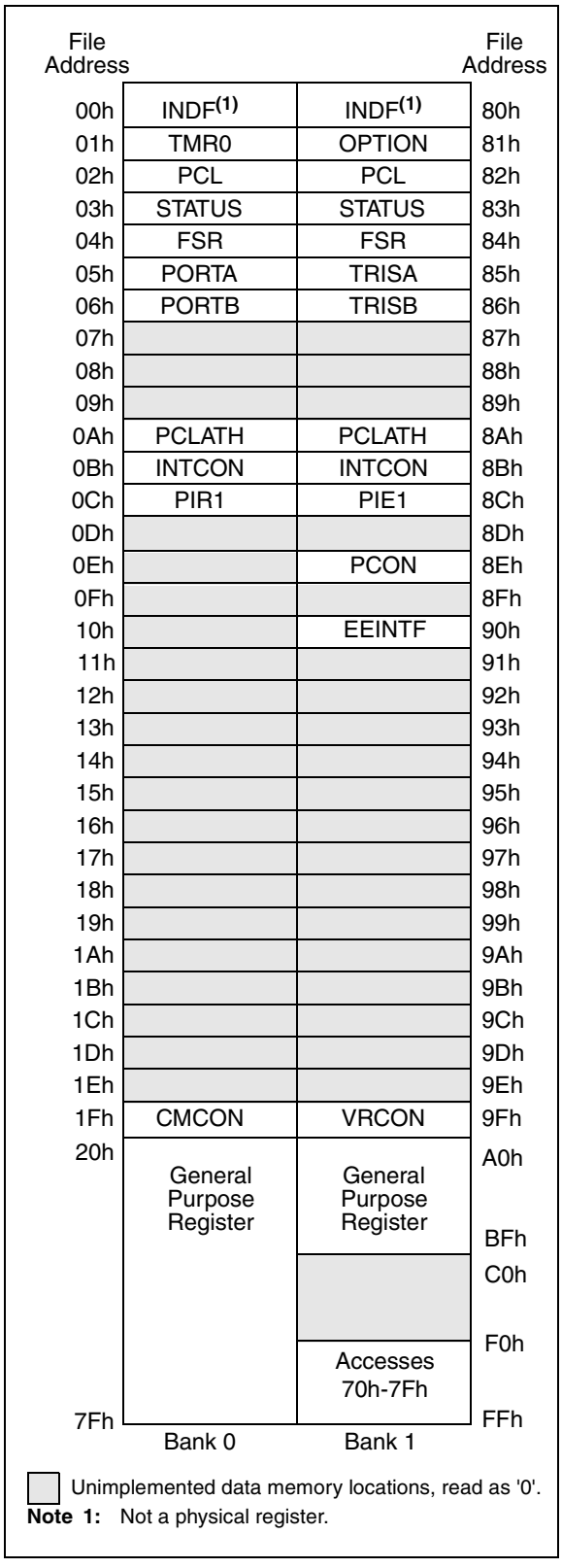


FIGURE 4-5: DATA MEMORY MAP FOR THE PIC16CE625



4.2.2.3 INTCON REGISTER

The INTCON register is a readable and writable register which contains the various enable and flag bits for all interrupt sources except the comparator module. See Section 4.2.2.4 and Section 4.2.2.5 for a description of the comparator enable and flag bits.

Note: Interrupt flag bits get set when an interrupt condition occurs, regardless of the state of its corresponding enable bit or the global enable bit, GIE (INTCON<7>).

REGISTER 4-3: INTCON REGISTER (ADDRESS 0BH OR 8BH)

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-x
GIE	PEIE	T0IE	INTE	RBIE	T0IF	INTF	RBIF
bit7							bit0

R = Readable bit
W = Writable bit
U = Unimplemented bit, read as '0'
-n = Value at POR reset
-x = Unknown at POR reset

bit 7: **GIE:** Global Interrupt Enable bit
1 = Enables all un-masked interrupts
0 = Disables all interrupts

bit 6: **PEIE**: Peripheral Interrupt Enable bit
1 = Enables all un-masked peripheral interrupts
0 = Disables all peripheral interrupts

bit 5: **T0IE**: TMR0 Overflow Interrupt Enable bit
1 = Enables the TMR0 interrupt
0 = Disables the TMR0 interrupt

bit 4: **INTE**: RB0/INT External Interrupt Enable bit
1 = Enables the RB0/INT external interrupt
0 = Disables the RB0/INT external interrupt

bit 3: **RBIE**: RB Port Change Interrupt Enable bit
1 = Enables the RB port change interrupt
0 = Disables the RB port change interrupt

bit 2: **T0IF**: TMR0 Overflow Interrupt Flag bit
1 = TMR0 register has overflowed (must be cleared in software)
0 = TMR0 register did not overflow

bit 1: **INTF**: RB0/INT External Interrupt Flag bit
1 = The RB0/INT external interrupt occurred (must be cleared in software)
0 = The RB0/INT external interrupt did not occur

bit 0: **RBIF**: RB Port Change Interrupt Flag bit
1 = When at least one of the RB<7:4> pins changed state (must be cleared in software)
0 = None of the RB<7:4> pins have changed state

5.0 I/O PORTS

The PIC16CE62X parts have two ports, PORTA and PORTB. Some pins for these I/O ports are multiplexed with an alternate function for the peripheral features on the device. In general, when a peripheral is enabled, that pin may not be used as a general purpose I/O pin.

5.1 PORTA and TRISA Registers

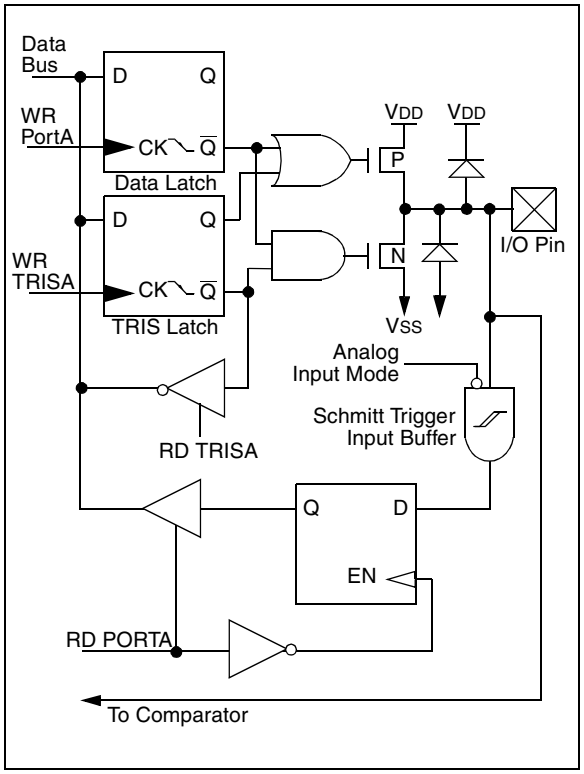
PORTA is a 5-bit wide latch. RA4 is a Schmitt Trigger input and an open drain output. Port RA4 is multiplexed with the T0CKI clock input. All other RA port pins have Schmitt Trigger input levels and full CMOS output drivers. All pins have data direction bits (TRIS registers), which can configure these pins as input or output.

A '1' in the TRISA register puts the corresponding output driver in a hi- impedance mode. A '0' in the TRISA register puts the contents of the output latch on the selected pin(s).

Reading the PORTA register reads the status of the pins, whereas writing to it will write to the port latch. All write operations are read-modify-write operations. So a write to a port implies that the port pins are first read, then this value is modified and written to the port data latch.

The PORTA pins are multiplexed with comparator and voltage reference functions. The operation of these pins are selected by control bits in the CMCON (Comparator Control Register) register and the VRCON (Voltage Reference Control Register) register. When selected as a comparator input, these pins will read as '0's.

FIGURE 5-1: BLOCK DIAGRAM OF RA<1:0> PINS



Note:	On reset, the TRISA register is set to all inputs. The digital inputs are disabled and the comparator inputs are forced to ground to reduce excess current consumption.

TRISA controls the direction of the RA pins, even when they are being used as comparator inputs. The user must make sure to keep the pins configured as inputs when using them as comparator inputs.

The RA2 pin will also function as the output for the voltage reference. When in this mode, the VREF pin is a very high impedance output. The user must configure TRISA<2> bit as an input and use high impedance loads.

In one of the comparator modes defined by the CMCON register, pins RA3 and RA4 become outputs of the comparators. The TRISA<4:3> bits must be cleared to enable outputs to use this function.

EXAMPLE 5-1: INITIALIZING PORTA

```
CLRF   PORTA       ;Initialize PORTA by setting
                   ;output data latches
MOVLW  0X07        ;Turn comparators off and
MOVWF  CMCON       ;enable pins for I/O
                   ;functions
BSF    STATUS, RP0 ;Select Bank1
MOVLW  0x1F        ;Value used to initialize
                   ;data direction
MOVWF  TRISA       ;Set RA<4:0> as inputs
                   ;TRISA<7:5> are always
                   ;read as '0'.
```

FIGURE 5-2: BLOCK DIAGRAM OF RA2 PIN

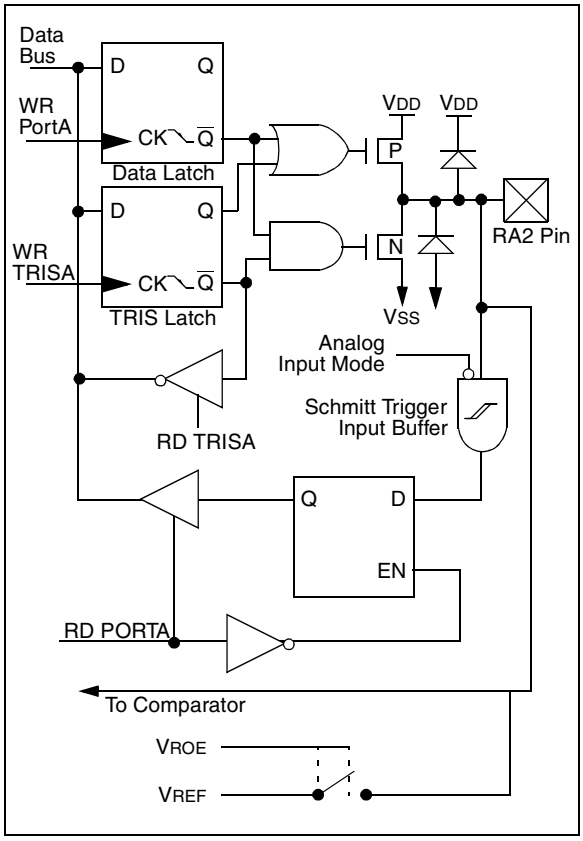


FIGURE 5-3: BLOCK DIAGRAM OF RA3 PIN

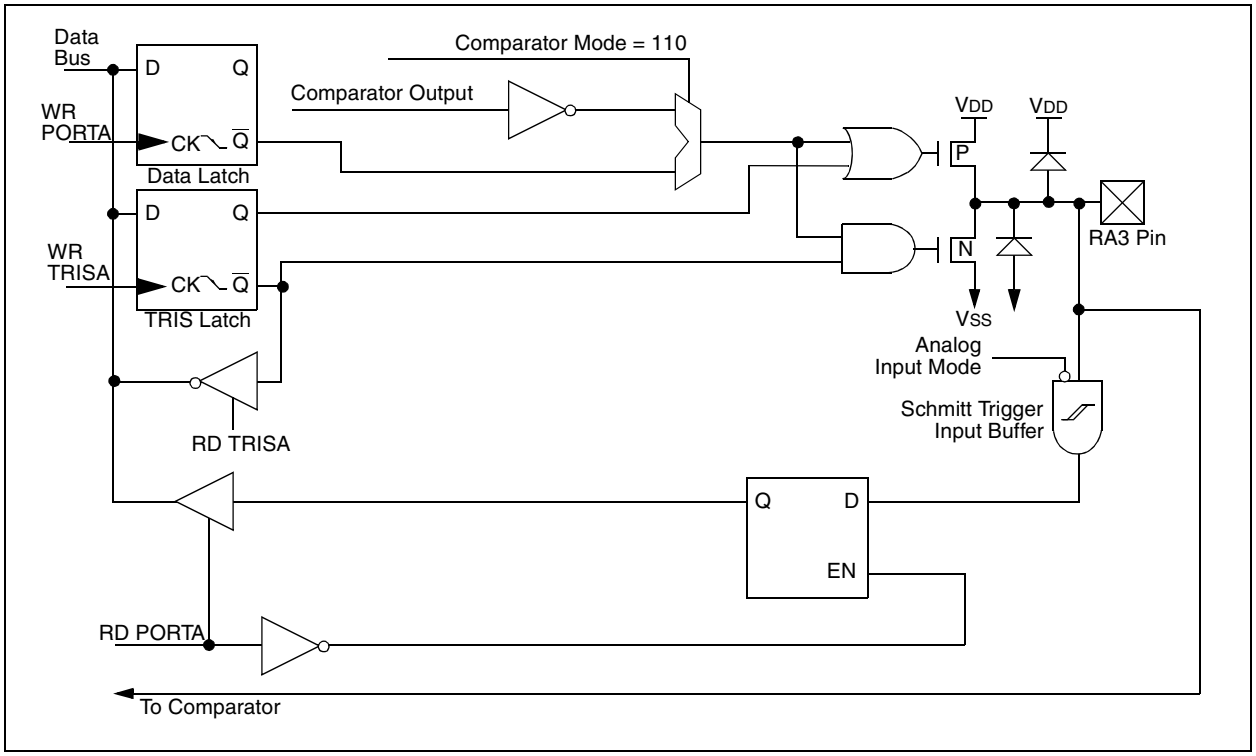
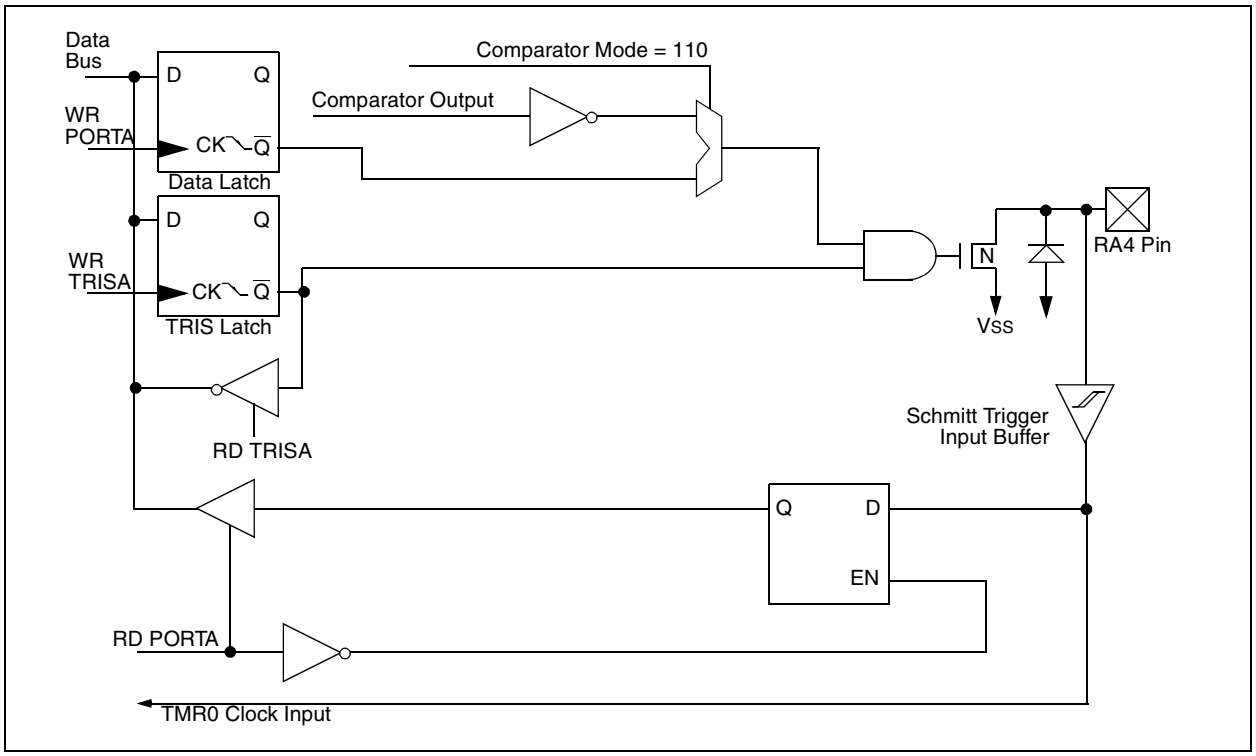


FIGURE 5-4: BLOCK DIAGRAM OF RA4 PIN



5.3 I/O Programming Considerations

5.3.1 BI-DIRECTIONAL I/O PORTS

Any instruction which writes, operates internally as a read followed by a write operation. The BCF and BSF instructions, for example, read the register into the CPU, execute the bit operation and write the result back to the register. Caution must be used when these instructions are applied to a port with both inputs and outputs defined. For example, a BSF operation on bit5 of PORTB will cause all eight bits of PORTB to be read into the CPU. Then the BSF operation takes place on bit5 and PORTB is written to the output latches. If another bit of PORTB is used as a bidirectional I/O pin (i.e., bit0) and it is defined as an input at this time, the input signal present on the pin itself would be read into the CPU and re-written to the data latch of this particular pin, overwriting the previous content. As long as the pin stays in the input mode, no problem occurs. However, if bit0 is switched into output mode later on, the content of the data latch may now be unknown.

Reading the port register, reads the values of the port pins. Writing to the port register writes the value to the port latch. When using read modify write instructions (i.e., BCF, BSF, etc.) on a port, the value of the port pins is read, the desired operation is done to this value, and this value is then written to the port latch.

Example 5-2 shows the effect of two sequential read-modify-write instructions (i.e., BCF, BSF, etc.) on an I/O port.

A pin actively outputting a Low or High should not be driven from external devices at the same time in order to change the level on this pin ("wired-or", "wired-and"). The resulting high output currents may damage the chip.

EXAMPLE 5-2: READ-MODIFY-WRITE INSTRUCTIONS ON AN I/O PORT

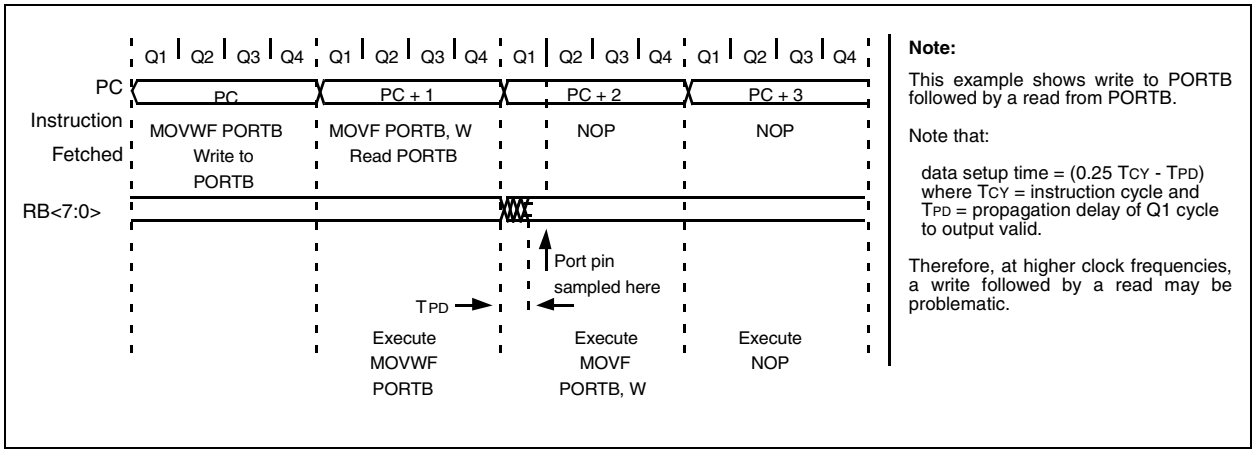
```
; Initial PORT settings:  PORTB<7:4> Inputs
;                         PORTB<3:0> Outputs
; PORTB<7:6> have external pull-up and are not
; connected to other circuitry
;
;                         PORT latch  PORT pins
;                         ----------  ----------

    BCF PORTB, 7        ; 01pp pppp   11pp pppp
    BCF PORTB, 6        ; 10pp pppp   11pp pppp
    BSF STATUS,RP0      ;
    BCF TRISB, 7        ; 10pp pppp   11pp pppp
    BCF TRISB, 6        ; 10pp pppp   10pp pppp
;
; Note that the user may have expected the pin
; values to be 00pp pppp. The 2nd BCF caused
; RB7 to be latched as the pin value (High).
```

5.3.2 SUCCESSIVE OPERATIONS ON I/O PORTS

The actual write to an I/O port happens at the end of an instruction cycle, whereas for reading, the data must be valid at the beginning of the instruction cycle (Figure 5-7). Therefore, care must be exercised if a write followed by a read operation is carried out on the same I/O port. The sequence of instructions should allow the pin voltage to stabilize (load dependent) before the next instruction causes that file to be read into the CPU. Otherwise, the previous state of that pin may be read into the CPU rather than the new state. When in doubt, it is better to separate these instructions with an NOP or another instruction not accessing this I/O port.

FIGURE 5-7: SUCCESSIVE I/O OPERATION



Note:

This example shows write to PORTB followed by a read from PORTB.

Note that:

data setup time = (0.25 TCY - TPD) where TCY = instruction cycle and TPD = propagation delay of Q1 cycle to output valid.

Therefore, at higher clock frequencies, a write followed by a read may be problematic.

6.0 EEPROM PERIPHERAL OPERATION

The PIC16CE623/624/625 each have 128 bytes of EEPROM data memory. The EEPROM data memory supports a bi-directional, 2-wire bus and data transmission protocol. These two-wires are serial data (SDA) and serial clock (SCL), and are mapped to bit1 and bit2, respectively, of the EEINTF register (SFR 90h). In addition, the power to the EEPROM can be controlled using bit0 (EEVDD) of the EEINTF register. For most applications, all that is required is calls to the following functions:

```
; Byte_Write: Byte write routine
;       Inputs: EEPROM Address     EEADDR
;               EEPROM Data        EEDATA
;       Outputs:    Return 01 in W if OK, else
;                   return 00 in W
;
; Read_Current: Read EEPROM at address
currently held by EE device.
;       Inputs: NONE
;       Outputs:    EEPROM Data    EEDATA
;                   Return 01 in W if OK, else
;                   return 00 in W
;
; Read_Random: Read EEPROM byte at supplied
; address
;       Inputs: EEPROM Address     EEADDR
;       Outputs:    EEPROM Data    EEDATA
;                   Return 01 in W if OK,
;                   else return 00 in W
```

The code for these functions is available on our web site (www.microchip.com). The code will be accessed by either including the source code FL62XINC.ASM or by linking FLASH62X.ASM. FLASH62.IMC provides external definition to the calling program.

6.0.1 SERIAL DATA

SDA is a bi-directional pin used to transfer addresses and data into and data out of the memory.

For normal data transfer, SDA is allowed to change only during SCL low. Changes during SCL high are reserved for indicating the START and STOP conditions.

6.0.2 SERIAL CLOCK

This SCL input is used to synchronize the data transfer to and from the memory.

6.0.3 EEINTF REGISTER

The EEINTF register (SFR 90h) controls the access to the EEPROM. Register 6-1 details the function of each bit. User code must generate the clock and data signals.

REGISTER 6-1: EEINTF REGISTER (ADDRESS 90h)

U-0	U-0	U-0	U-0	U-0	R/W-1	R/W-1	R/W-1
—	—	—	—	—	EESCL	EESDA	EEVDD
bit7							bit0

R = Readable bit
W = Writable bit
U = Unimplemented bit, read as '0'
- n = Value at POR reset

bit 7-3: **Unimplemented:** Read as '0'

bit 2: **EESCL**: Clock line to the EEPROM
1 = Clock high
0 = Clock low

bit 1: **EESDA**: Data line to EEPROM
1 = Data line is high (pin is tri-stated, line is pulled high by a pull-up resistor)
0 = Data line is low

bit 0: **EEVDD**: VDD control bit for EEPROM
1 = VDD is turned on to EEPROM
0 = VDD is turned off to EEPROM (all pins are tri-stated and the EEPROM is powered down)

Note: EESDA, EESCL and EEVDD will read '0' if EEVDD is turned off.

7.2 Using Timer0 with External Clock

When an external clock input is used for Timer0, it must meet certain requirements. The external clock requirement is due to internal phase clock (TOSC) synchronization. Also, there is a delay in the actual incrementing of Timer0 after synchronization.

7.2.1 EXTERNAL CLOCK SYNCHRONIZATION

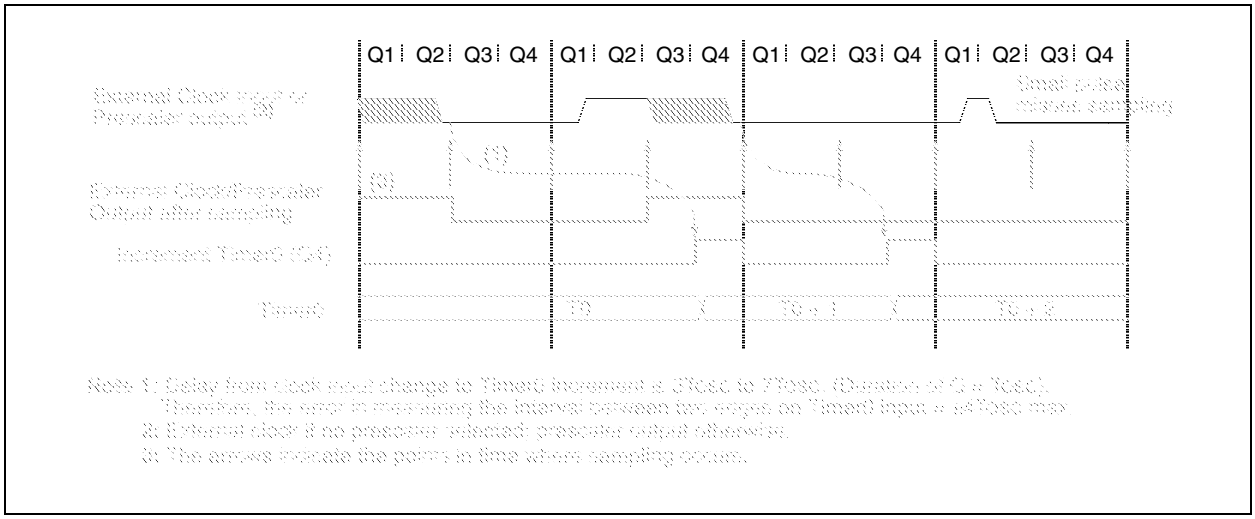
When no prescaler is used, the external clock input is the same as the prescaler output. The synchronization of T0CKI with the internal phase clocks is accomplished by sampling the prescaler output on the Q2 and Q4 cycles of the internal phase clocks (Figure 7-5). Therefore, it is necessary for T0CKI to be high for at least 2TOSC (and a small RC delay of 20 ns) and low for at least 2TOSC (and a small RC delay of 20 ns). Refer to the electrical specification of the desired device.

When a prescaler is used, the external clock input is divided by the asynchronous ripple-counter type prescaler so that the prescaler output is symmetrical. For the external clock to meet the sampling requirement, the ripple-counter must be taken into account. Therefore, it is necessary for T0CKI to have a period of at least 4TOSC (and a small RC delay of 40 ns) divided by the prescaler value. The only requirement on T0CKI high and low time is that they do not violate the minimum pulse width requirement of 10 ns. Refer to parameters 40, 41 and 42 in the electrical specification of the desired device.

7.2.2 TIMER0 INCREMENT DELAY

Since the prescaler output is synchronized with the internal clocks, there is a small delay from the time the external clock edge occurs to the time the TMR0 is actually incremented. Figure 7-5 shows the delay from the external clock edge to the timer incrementing.

FIGURE 7-5: TIMER0 TIMING WITH EXTERNAL CLOCK



Note 1: Delay from clock input change to Timer0 increment is 3Tosc to 7Tosc. (Duration of Q = Tosc).
Therefore, the error in measuring the interval between two edges on Timer0 input = ±4Tosc max.
2: External clock if no prescaler selected; prescaler output otherwise.
3: The arrows indicate the points in time where sampling occurs.

8.6 Comparator Interrupts

The comparator interrupt flag is set whenever there is a change in the output value of either comparator. Software will need to maintain information about the status of the output bits, as read from CMCON<7:6>, to determine the actual change that has occurred. The CMIF bit, PIR1<6>, is the comparator interrupt flag. The CMIF bit must be reset by clearing '0'. Since it is also possible to write a '1' to this register, a simulated interrupt may be initiated.

The CMIE bit (PIE1<6>) and the PEIE bit (INTCON<6>) must be set to enable the interrupt. In addition, the GIE bit must also be set. If any of these bits are clear, the interrupt is not enabled, though the CMIF bit will still be set if an interrupt condition occurs.

Note:	If a change in the CMCON register (C1OUT or C2OUT) should occur when a read operation is being executed (start of the Q2 cycle), then the CMIF (PIR1<6>) interrupt flag may not get set.

The user, in the interrupt service routine, can clear the interrupt in the following manner:

a) Any read or write of CMCON. This will end the mismatch condition.

b) Clear flag bit CMIF.

A mismatch condition will continue to set flag bit CMIF. Reading CMCON will end the mismatch condition, and allow flag bit CMIF to be cleared.

8.7 Comparator Operation During SLEEP

When a comparator is active and the device is placed in SLEEP mode, the comparator remains active and the interrupt is functional if enabled. This interrupt will wake-up the device from SLEEP mode when enabled. While the comparator is powered-up, higher sleep currents than shown in the power down current specification will occur. Each comparator that is operational will consume additional current as shown in the comparator specifications. To minimize power consumption while in SLEEP mode, turn off the comparators, CM<2:0> = 111, before entering sleep. If the device wakes-up from sleep, the contents of the CMCON register are not affected.

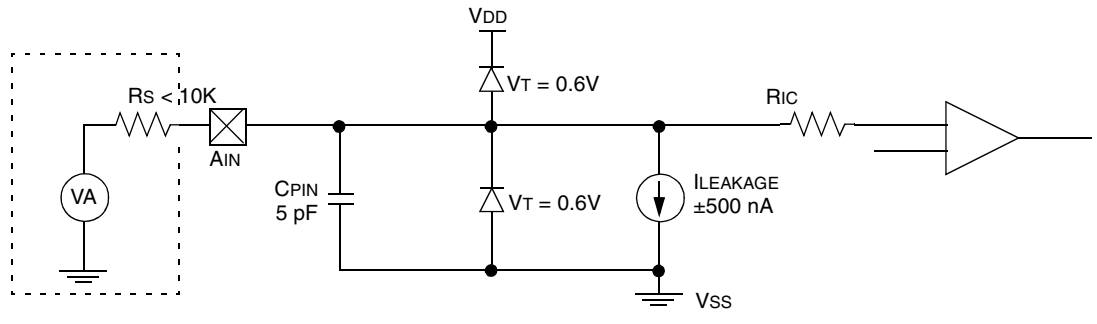
8.8 Effects of a RESET

A device reset forces the CMCON register to its reset state. This forces the comparator module to be in the comparator reset mode, CM<2:0> = 000. This ensures that all potential inputs are analog inputs. Device current is minimized when analog inputs are present at reset time. The comparators will be powered-down during the reset interval.

8.9 Analog Input Connection Considerations

A simplified circuit for an analog input is shown in Figure 8-4. Since the analog pins are connected to a digital output, they have reverse biased diodes to VDD and VSS. The analog input therefore, must be between VSS and VDD. If the input voltage deviates from this range by more than 0.6V in either direction, one of the diodes is forward biased and a latch-up may occur. A maximum source impedance of 10 kΩ is recommended for the analog sources. Any external component connected to an analog input pin, such as a capacitor or a Zener diode, should have very little leakage current.

FIGURE 8-4: ANALOG INPUT MODEL



Legend
CPIN = Input capacitance
VT = Threshold voltage
ILEAKAGE = Leakage current at the pin due to various junctions
RIC = Interconnect resistance
RS = Source impedance
VA = Analog voltage

10.2 Oscillator Configurations

10.2.1 OSCILLATOR TYPES

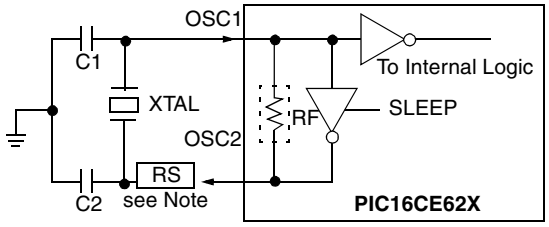
The PIC16CE62X can be operated in four different oscillator options. The user can program two configuration bits (FOSC1 and FOSC0) to select one of these four modes:

- LP Low Power Crystal
- XT Crystal/Resonator
- HS High Speed Crystal/Resonator
- RC Resistor/Capacitor

10.2.2 CRYSTAL OSCILLATOR / CERAMIC RESONATORS

In XT, LP or HS modes, a crystal or ceramic resonator is connected to the OSC1 and OSC2 pins to establish oscillation (Figure 10-1). The PIC16CE62X oscillator design requires the use of a parallel cut crystal. Use of a series cut crystal may give a frequency out of the crystal manufacturers specifications. When in XT, LP or HS modes, the device can have an external clock source to drive the OSC1 pin (Figure 10-2).

FIGURE 10-1: CRYSTAL OPERATION (OR CERAMIC RESONATOR) (HS, XT OR LP OSC CONFIGURATION)



See Table 10-1 and Table 10-2 for recommended values of C1 and C2.

Note: A series resistor may be required for AT strip cut crystals.

FIGURE 10-2: EXTERNAL CLOCK INPUT OPERATION (HS, XT OR LP OSC CONFIGURATION)

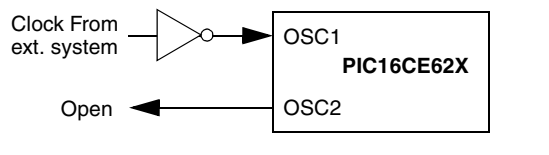


TABLE 10-1: CERAMIC RESONATORS, PIC16CE62X

Ranges Tested:			
Mode	**Freq**	**OSC1**	**OSC2**
XT	455 kHz 2.0 MHz 4.0 MHz	68 - 100 pF 15 - 68 pF 15 - 68 pF	68 - 100 pF 15 - 68 pF 15 - 68 pF
HS	8.0 MHz 16.0 MHz	10 - 68 pF 10 - 22 pF	10 - 68 pF 10 - 22 pF
These values are for design guidance only. See notes at bottom of page.			

TABLE 10-2: CAPACITOR SELECTION FOR CRYSTAL OSCILLATOR, PIC16CE62X

Osc Type	Crystal Freq	Cap. Range C1	Cap. Range C2
LP	32 kHz	33 pF	33 pF
	200 kHz	15 pF	15 pF
XT	200 kHz	47-68 pF	47-68 pF
	1 MHz	15 pF	15 pF
	4 MHz	15 pF	15 pF
HS	4 MHz	15 pF	15 pF
	8 MHz	15-33 pF	15-33 pF
	20 MHz	15-33 pF	15-33 pF
These values are for design guidance only. See notes at bottom of page.			

1. Recommended values of C1 and C2 are identical to the ranges tested table.
2. Higher capacitance increases the stability of oscillator, but also increases the start-up time.
3. Since each resonator/crystal has its own characteristics, the user should consult the resonator/crystal manufacturer for appropriate values of external components.
4. Rs may be required in HS mode, as well as XT mode, to avoid overdriving crystals with low drive level specification.

10.4.5 TIME-OUT SEQUENCE

On power-up, the time-out sequence is as follows: First PWRT time-out is invoked after POR has expired, then OST is activated. The total time-out will vary based on oscillator configuration and $\overline{\text{PWRTE}}$ bit status. For example, in RC mode with $\overline{\text{PWRTE}}$ bit erased (PWRT disabled), there will be no time-out at all. Figure 10-8, Figure 10-9 and Figure 10-10 depict time-out sequences.

Since the time-outs occur from the POR pulse, if $\overline{\text{MCLR}}$ is kept low long enough, the time-outs will expire. Then bringing $\overline{\text{MCLR}}$ high will begin execution immediately (see Figure 10-9). This is useful for testing purposes or to synchronize more than one PIC® device operating in parallel.

Table 10-5 shows the reset conditions for some special registers, while Table 10-6 shows the reset conditions for all the registers.

10.4.6 POWER CONTROL (PCON)/STATUS REGISTER

The power control/status register, PCON (address 8Eh) has two bits.

Bit0 is $\overline{\text{BOR}}$ (Brown-out). $\overline{\text{BOR}}$ is unknown on power-on-reset. It must then be set by the user and checked on subsequent resets to see if $\overline{\text{BOR}}$ = 0 indicating that a brown-out has occurred. The $\overline{\text{BOR}}$ status bit is a don't care and is not necessarily predictable if the brown-out circuit is disabled (by setting BODEN bit = 0 in the Configuration word).

Bit1 is $\overline{\text{POR}}$ (Power-on-reset). It is a '0' on power-on-reset and unaffected otherwise. The user must write a '1' to this bit following a power-on-reset. On a subsequent reset, if $\overline{\text{POR}}$ is '0', it will indicate that a power-on-reset must have occurred (V_{DD} may have gone too low).

TABLE 10-3: TIME-OUT IN VARIOUS SITUATIONS

Oscillator Configuration	Power-up		Brown-out Reset	Wake-up from SLEEP
	$\overline{\text{PWRTE}}$ = 0	$\overline{\text{PWRTE}}$ = 1		
XT, HS, LP	72 ms + 1024 TOSC	1024 TOSC	72 ms + 1024 TOSC	1024 TOSC
RC	72 ms	—	72 ms	—

TABLE 10-4: STATUS/PCON BITS AND THEIR SIGNIFICANCE

$\overline{\text{POR}}$	$\overline{\text{BOR}}$	$\overline{\text{TO}}$	$\overline{\text{PD}}$	
0	x	1	1	Power-on-reset
0	x	0	x	Illegal, $\overline{\text{TO}}$ is set on $\overline{\text{POR}}$
0	x	x	0	Illegal, $\overline{\text{PD}}$ is set on $\overline{\text{POR}}$
1	0	x	x	Brown-out Reset
1	1	0	u	WDT Reset
1	1	0	0	WDT Wake-up
1	1	u	u	$\overline{\text{MCLR}}$ reset during normal operation
1	1	1	0	$\overline{\text{MCLR}}$ reset during SLEEP

Legend: x = unknown, u = unchanged

TABLE 10-5: INITIALIZATION CONDITION FOR SPECIAL REGISTERS

Condition	Program Counter	STATUS Register	PCON Register
Power-on Reset	000h	0001 1xxx	---- --0x
$\overline{MCLR}$ reset during normal operation	000h	000u uuuu	---- --uu
$\overline{MCLR}$ reset during SLEEP	000h	0001 0uuu	---- --uu
WDT reset	000h	0000 uuuu	---- --uu
WDT Wake-up	PC + 1	uuu0 0uuu	---- --uu
Brown-out Reset	000h	000x xuuu	---- --u0
Interrupt Wake-up from SLEEP	PC + 1(1)	uuu1 0uuu	---- --uu

Legend: u = unchanged, x = unknown, - = unimplemented bit, reads as '0'.

Note 1: When the wake-up is due to an interrupt and global enable bit, GIE is set and the PC is loaded with the interrupt vector (0004h) after execution of PC+1.

TABLE 10-6: INITIALIZATION CONDITION FOR REGISTERS

Register	Address	Power-on Reset	• $\overline{MCLR}$ Reset during normal operation • $\overline{MCLR}$ Reset during SLEEP • WDT Reset • Brown-out Reset (1)	• Wake-up from SLEEP through interrupt • Wake-up from SLEEP through WDT time-out
W	-	xxxx xxxx	uuuu uuuu	uuuu uuuu
INDF	00h	-	-	-
TMR0	01h	xxxx xxxx	uuuu uuuu	uuuu uuuu
PCL	02h	0000 0000	0000 0000	PC + 1(3)
STATUS	03h	0001 1xxx	000q quuu(4)	uuuq quuu(4)
FSR	04h	xxxx xxxx	uuuu uuuu	uuuu uuuu
PORTA	05h	---x xxxx	---u uuuu	---u uuuu
PORTB	06h	xxxx xxxx	uuuu uuuu	uuuu uuuu
CMCON	1Fh	00-- 0000	00-- 0000	uu-- uuuu
PCLATH	0Ah	---0 0000	---0 0000	---u uuuu
INTCON	0Bh	0000 000x	0000 000u	uuuu uqqq(2)
PIR1	0Ch	-0-- ----	-0-- ----	-q-- ----(2,5)
OPTION	81h	1111 1111	1111 1111	uuuu uuuu
TRISA	85h	---1 1111	---1 1111	---u uuuu
TRISB	86h	1111 1111	1111 1111	uuuu uuuu
PIE1	8Ch	-0-- ----	-0-- ----	-u-- ----
PCON	8Eh	---- --0x	---- --uq(1,6)	---- --uu
EEINTF	90h	---- -111	---- -111	---- -111
VRCON	9Fh	000- 0000	000- 0000	uuu- uuuu

Legend: u = unchanged, x = unknown, - = unimplemented bit, reads as '0',q = value depends on condition.

Note 1: If VDD goes too low, Power-on Reset will be activated and registers will be affected differently.

2: One or more bits in INTCON, PIR1 and/or PIR2 will be affected (to cause wake-up).

3: When the wake-up is due to an interrupt and the GIE bit is set, the PC is loaded with the interrupt vector (0004h).

4: See Table 10-5 for reset value for specific condition.

5: If wake-up was due to comparator input changing , then bit 6 = 1. All other interrupts generating a wake-up will cause bit 6 = u.

6: If reset was due to brown-out, then PCON bit 0 = 0. All other resets will cause bit 0 = u.

10.9 Code Protection

If the code protection bit(s) have not been programmed, the on-chip program memory can be read out for verification purposes.

Note:	Microchip does not recommend code protecting windowed devices.

10.10 ID Locations

Four memory locations (2000h-2003h) are designated as ID locations where the user can store checksum or other code-identification numbers. These locations are not accessible during normal execution but are readable and writable during program/verify. Only the least significant 4 bits of the ID locations are used.

10.11 In-Circuit Serial Programming

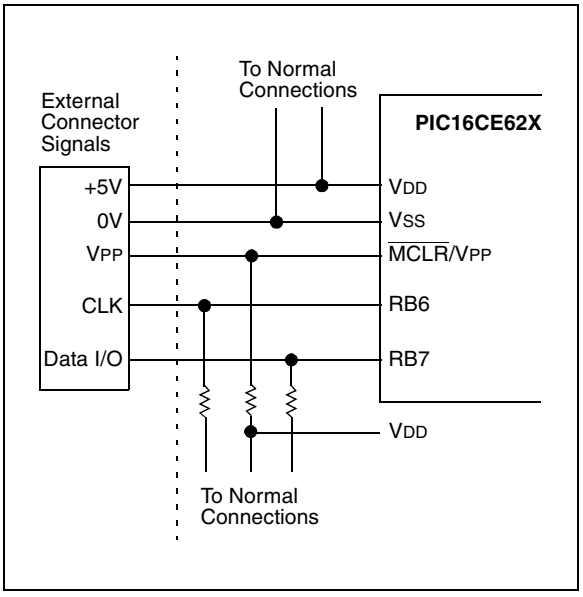
The PIC16CE62X microcontrollers can be serially programmed while in the end application circuit. This is simply done with two lines for clock and data, and three other lines for power, ground, and the programming voltage. This allows customers to manufacture boards with unprogrammed devices, and then program the microcontroller just before shipping the product. This also allows the most recent firmware or a custom firmware to be programmed.

The device is placed into a program/verify mode by holding the RB6 and RB7 pins low, while raising the $\overline{MCLR}$ (V_{PP}) pin from V_{IL} to V_{IHH} (see programming specification). RB6 becomes the programming clock and RB7 becomes the programming data. Both RB6 and RB7 are Schmitt Trigger inputs in this mode.

After reset, to place the device into programming/verify mode, the program counter (PC) is at location 00h. A 6-bit command is then supplied to the device. Depending on the command, 14-bits of program data are then supplied to or from the device, depending if the command was a load or a read. For complete details of serial programming, please refer to the PIC16C6X/7X/9XX Programming Specifications (Literature #DS30228).

A typical in-circuit serial programming connection is shown in Figure 10-20.

FIGURE 10-20: TYPICAL IN-CIRCUIT SERIAL PROGRAMMING CONNECTION



BCF **Bit Clear f**

Syntax: [*label*] BCF f,b

Operands: $0 \le f \le 127$
$0 \le b \le 7$

Operation: $0 \rightarrow (f<b>)$

Status Affected: None

Encoding:

01	00bb	bfff	ffff

Description: Bit 'b' in register 'f' is cleared.

Words: 1

Cycles: 1

Example

```
BCF     FLAG_REG, 7
```

Before Instruction
FLAG_REG = 0xC7
After Instruction
FLAG_REG = 0x47

BSF **Bit Set f**

Syntax: [*label*] BSF f,b

Operands: $0 \le f \le 127$
$0 \le b \le 7$

Operation: $1 \rightarrow (f<b>)$

Status Affected: None

Encoding:

01	01bb	bfff	ffff

Description: Bit 'b' in register 'f' is set.

Words: 1

Cycles: 1

Example

```
BSF     FLAG_REG,   7
```

Before Instruction
FLAG_REG = 0x0A
After Instruction
FLAG_REG = 0x8A

BTFSC **Bit Test, Skip if Clear**

Syntax: [*label*] BTFSC f,b

Operands: $0 \le f \le 127$
$0 \le b \le 7$

Operation: skip if (f<b>) = 0

Status Affected: None

Encoding:

01	10bb	bfff	ffff

Description: If bit 'b' in register 'f' is '0', then the next instruction is skipped.
If bit 'b' is '0', then the next instruction fetched during the current instruction execution is discarded, and a `NOP` is executed instead, making this a two-cycle instruction.

Words: 1

Cycles: 1(2)

Example

```
HERE    BTFSC   FLAG,1
FALSE   GOTO    PROCESS_CODE
TRUE    •
        •
        •
```

Before Instruction
PC = address `HERE`
After Instruction
if FLAG<1> = 0,
PC = address `TRUE`
if FLAG<1>=1,
PC = address `FALSE`

CLRWDT **Clear Watchdog Timer**

Syntax: [*label*] CLRWDT

Operands: None

Operation: 00h → WDT
0 → WDT prescaler,
1 → $\overline{TO}$
1 → $\overline{PD}$

Status Affected: $\overline{TO}$, $\overline{PD}$

Encoding:

00	0000	0110	0100

Description: CLRWDT instruction resets the Watchdog Timer. It also resets the prescaler of the WDT. Status bits $\overline{TO}$ and $\overline{PD}$ are set.

Words: 1

Cycles: 1

Example
```
CLRWDT
```
Before Instruction
WDT counter = ?
After Instruction
WDT counter = 0x00
WDT prescaler = 0
$\overline{TO}$ = 1
$\overline{PD}$ = 1

COMF **Complement f**

Syntax: [*label*] COMF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: ($\bar{f}$) → (dest)

Status Affected: Z

Encoding:

00	1001	dfff	ffff

Description: The contents of register 'f' are complemented. If 'd' is 0, the result is stored in W. If 'd' is 1, the result is stored back in register 'f'.

Words: 1

Cycles: 1

Example
```
COMF    REG1,0
```
Before Instruction
REG1 = 0x13
After Instruction
REG1 = 0x13
W = 0xEC

DECF **Decrement f**

Syntax: [*label*] DECF f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (f) - 1 → (dest)

Status Affected: Z

Encoding:

00	0011	dfff	ffff

Description: Decrement register 'f'. If 'd' is 0, the result is stored in the W register. If 'd' is 1, the result is stored back in register 'f'.

Words: 1

Cycles: 1

Example
```
DECF    CNT, 1
```
Before Instruction
CNT = 0x01
Z = 0
After Instruction
CNT = 0x00
Z = 1

DECFSZ **Decrement f, Skip if 0**

Syntax: [*label*] DECFSZ f,d

Operands: 0 ≤ f ≤ 127
d ∈ [0,1]

Operation: (f) - 1 → (dest); skip if result = 0

Status Affected: None

Encoding:

00	1011	dfff	ffff

Description: The contents of register 'f' are decremented. If 'd' is 0, the result is placed in the W register. If 'd' is 1, the result is placed back in register 'f'.
If the result is 0, the next instruction, which is already fetched, is discarded. A NOP is executed instead making it a two-cycle instruction.

Words: 1

Cycles: 1(2)

Example
```
HERE     DECFSZ   CNT, 1
         GOTO     LOOP
CONTINUE •
         •
         •
```
Before Instruction
PC = address HERE
After Instruction
CNT = CNT - 1
if CNT = 0,
PC = address CONTINUE
if CNT ≠ 0,
PC = address HERE+1

SUBLW **Subtract W from Literal**

Syntax: [*label*] SUBLW k

Operands: $0 \le k \le 255$

Operation: k - (W) → (W)

Status Affected: C, DC, Z

Encoding:

11	110x	kkkk	kkkk

Description: The W register is subtracted (2's complement method) from the eight bit literal 'k'. The result is placed in the W register.

Words: 1

Cycles: 1

Example 1:
```
SUBLW   0x02
```
Before Instruction
W = 1
C = ?
After Instruction
W = 1
C = 1; result is positive

Example 2: Before Instruction
W = 2
C = ?
After Instruction
W = 0
C = 1; result is zero

Example 3: Before Instruction
W = 3
C = ?
After Instruction
W = 0xFF
C = 0; result is negative

SUBWF **Subtract W from f**

Syntax: [*label*] SUBWF f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: (f) - (W) → (dest)

Status Affected: C, DC, Z

Encoding:

00	0010	dfff	ffff

Description: Subtract (2's complement method) W register from register 'f'. If 'd' is 0, the result is stored in the W register. If 'd' is 1, the result is stored back in register 'f'.

Words: 1

Cycles: 1

Example 1:
```
SUBWF   REG1,1
```
Before Instruction
REG1 = 3
W = 2
C = ?
After Instruction
REG1 = 1
W = 2
C = 1; result is positive

Example 2: Before Instruction
REG1 = 2
W = 2
C = ?
After Instruction
REG1 = 0
W = 2
C = 1; result is zero

Example 3: Before Instruction
REG1 = 1
W = 2
C = ?
After Instruction
REG1 = 0xFF
W = 2
C = 0; result is negative

NOTES:

APPENDIX A: CODE FOR ACCESSING EEPROM DATA MEMORY

Please check our web site at www.microchip.com for code availability.

APPENDIX B:REVISION HISTORY

Revision D (January 2013)

Added a note to each package outline drawing.

NOTES: