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"[Embedded - Microcontrollers](#)" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

Applications of "[Embedded - Microcontrollers](#)"

Details

Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	4MHz
Connectivity	-
Peripherals	Brown-out Detect/Reset, POR, WDT
Number of I/O	13
Program Memory Size	896B (512 x 14)
Program Memory Type	OTP
EEPROM Size	128 x 8
RAM Size	96 x 8
Voltage - Supply (Vcc/Vdd)	2.5V ~ 5.5V
Data Converters	-
Oscillator Type	External
Operating Temperature	0°C ~ 70°C (TA)
Mounting Type	Surface Mount
Package / Case	20-SSOP (0.209", 5.30mm Width)
Supplier Device Package	20-SSOP
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic16lce623-04-ss

PIC16CE62X

NOTES:

3.1 Clocking Scheme/Instruction Cycle

The clock input (OSC1/CLKIN pin) is internally divided by four to generate four non-overlapping quadrature clocks namely Q1, Q2, Q3 and Q4. Internally, the program counter (PC) is incremented every Q1, the instruction is fetched from the program memory and latched into the instruction register in Q4. The instruction is decoded and executed during the following Q1 through Q4. The clocks and instruction execution flow is shown in Figure 3-2.

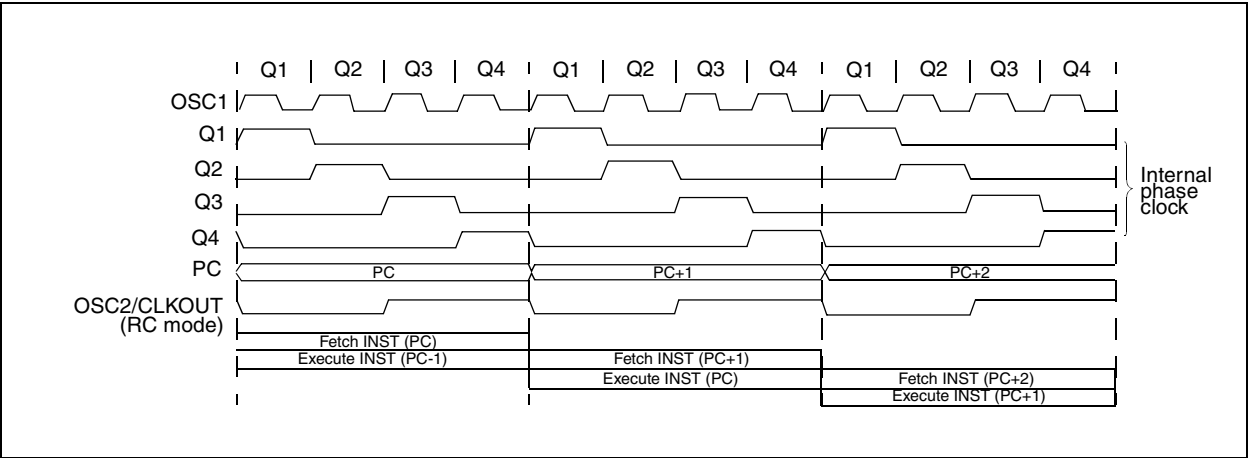
3.2 Instruction Flow/Pipelining

An “Instruction Cycle” consists of four Q cycles (Q1, Q2, Q3 and Q4). The instruction fetch and execute are pipelined such that fetch takes one instruction cycle, while decode and execute takes another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the program counter to change (i.e., GOTO) then two cycles are required to complete the instruction (Example 3-1).

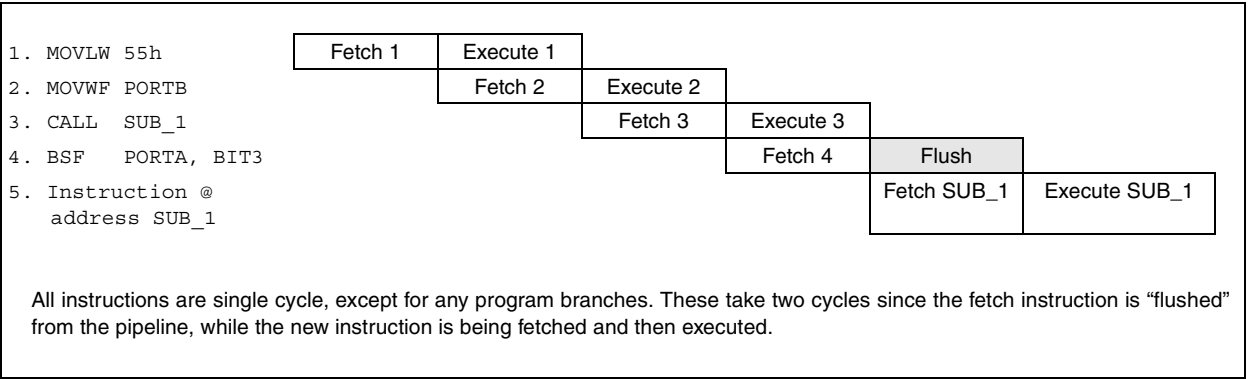
A fetch cycle begins with the program counter (PC) incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the “Instruction Register (IR)” in cycle Q1. This instruction is then decoded and executed during the Q2, Q3, and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).

FIGURE 3-2: CLOCK/INSTRUCTION CYCLE



EXAMPLE 3-1: INSTRUCTION PIPELINE FLOW



4.0 MEMORY ORGANIZATION

4.1 Program Memory Organization

The PIC16CE62X has a 13-bit program counter capable of addressing an 8K x 14 program memory space. Only the first 512 x 14 (0000h - 01FFh) for the PIC16CE623, 1K x 14 (0000h - 03FFh) for the PIC16CE624 and 2K x 14 (0000h - 07FFh) for the PIC16CE625 are physically implemented. Accessing a location above these boundaries will cause a wrap-around within the first 512 x 14 space (PIC16CE623) or 1K x 14 space (PIC16CE624) or 2K x 14 space (PIC16CE625). The reset vector is at 0000h and the interrupt vector is at 0004h (Figure 4-1, Figure 4-2, Figure 4-3).

FIGURE 4-1: PROGRAM MEMORY MAP AND STACK FOR THE PIC16CE623

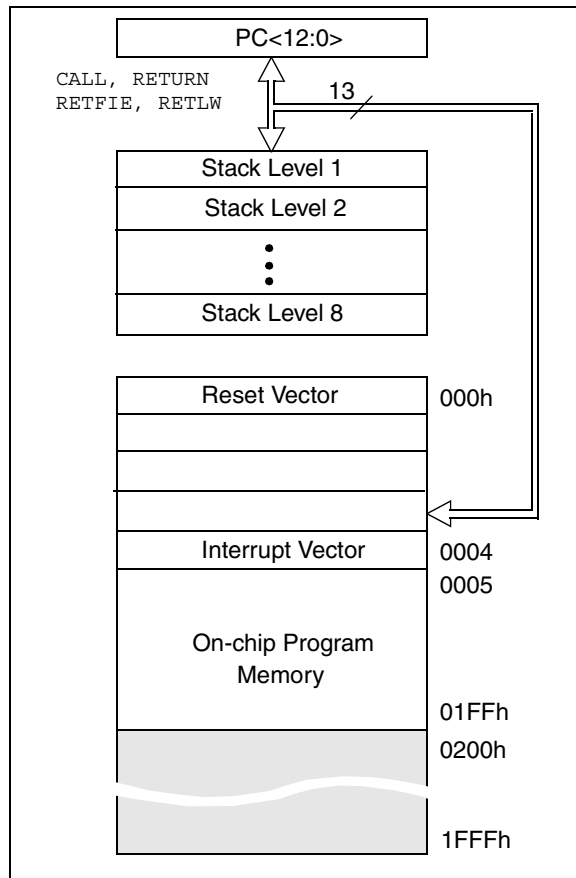


FIGURE 4-2: PROGRAM MEMORY MAP AND STACK FOR THE PIC16CE624

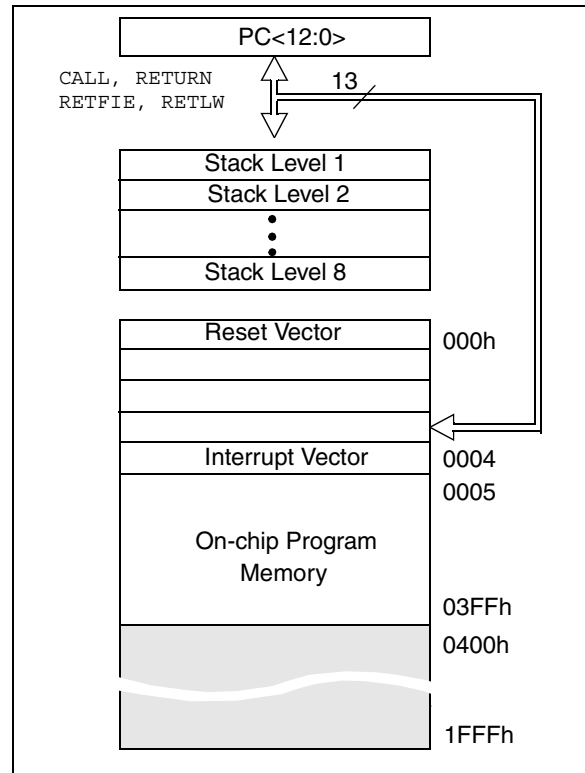
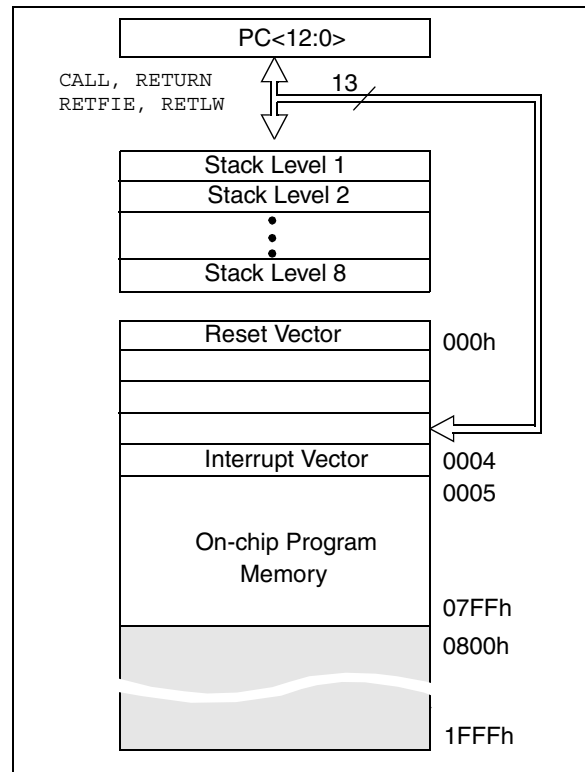


FIGURE 4-3: PROGRAM MEMORY MAP AND STACK FOR THE PIC16CE625



4.2.2.3 INTCON REGISTER

The INTCON register is a readable and writable register which contains the various enable and flag bits for all interrupt sources except the comparator module. See Section 4.2.2.4 and Section 4.2.2.5 for a description of the comparator enable and flag bits.

Note: Interrupt flag bits get set when an interrupt condition occurs, regardless of the state of its corresponding enable bit or the global enable bit, GIE (INTCON<7>).

REGISTER 4-3: INTCON REGISTER (ADDRESS 0BH OR 8BH)

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-x
GIE	PEIE	TOIE	INTE	RBIE	TOIF	INTF	RBIF
bit7							bit0
<p>bit 7: GIE: Global Interrupt Enable bit 1 = Enables all un-masked interrupts 0 = Disables all interrupts</p> <p>bit 6: PEIE: Peripheral Interrupt Enable bit 1 = Enables all un-masked peripheral interrupts 0 = Disables all peripheral interrupts</p> <p>bit 5: TOIE: TMR0 Overflow Interrupt Enable bit 1 = Enables the TMR0 interrupt 0 = Disables the TMR0 interrupt</p> <p>bit 4: INTE: RB0/INT External Interrupt Enable bit 1 = Enables the RB0/INT external interrupt 0 = Disables the RB0/INT external interrupt</p> <p>bit 3: RBIE: RB Port Change Interrupt Enable bit 1 = Enables the RB port change interrupt 0 = Disables the RB port change interrupt</p> <p>bit 2: TOIF: TMR0 Overflow Interrupt Flag bit 1 = TMR0 register has overflowed (must be cleared in software) 0 = TMR0 register did not overflow</p> <p>bit 1: INTF: RB0/INT External Interrupt Flag bit 1 = The RB0/INT external interrupt occurred (must be cleared in software) 0 = The RB0/INT external interrupt did not occur</p> <p>bit 0: RBIF: RB Port Change Interrupt Flag bit 1 = When at least one of the RB<7:4> pins changed state (must be cleared in software) 0 = None of the RB<7:4> pins have changed state</p>							
<p>R = Readable bit W = Writable bit U = Unimplemented bit, read as '0' -n = Value at POR reset -x = Unknown at POR reset</p>							

4.4 Indirect Addressing, INDF and FSR Registers

The INDF register is not a physical register. Addressing the INDF register will cause indirect addressing.

Indirect addressing is possible by using the INDF register. Any instruction using the INDF register actually accesses data pointed to by the File Select Register (FSR). Reading INDF itself indirectly will produce 00h. Writing to the INDF register indirectly results in a no-operation (although status bits may be affected). An effective 9-bit address is obtained by concatenating the 8-bit FSR register and the IRP bit (STATUS<7>), as shown in Figure 4-7. However, IRP is not used in the PIC16CE62X.

A simple program to clear RAM location 20h-2Fh using indirect addressing is shown in Example 4-1.

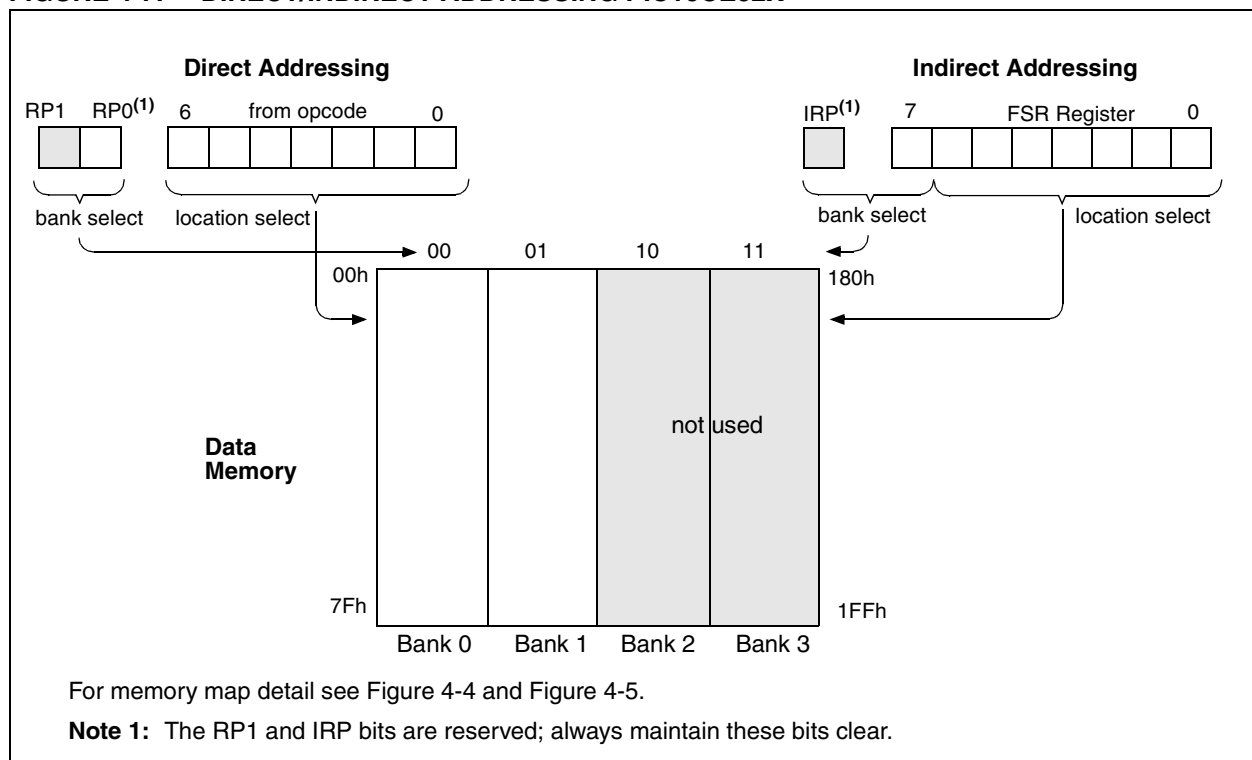
EXAMPLE 4-1: INDIRECT ADDRESSING

```

movlw 0x20    ;initialize pointer
movwf FSR     ;to RAM
NEXT        clrfs INDF    ;clear INDF register
            incf FSR      ;inc pointer
            btfss FSR,4    ;all done?
            goto NEXT     ;no clear next
                        ;yes continue
CONTINUE:

```

FIGURE 4-7: DIRECT/INDIRECT ADDRESSING PIC16CE62X



5.0 I/O PORTS

The PIC16CE62X parts have two ports, PORTA and PORTB. Some pins for these I/O ports are multiplexed with an alternate function for the peripheral features on the device. In general, when a peripheral is enabled, that pin may not be used as a general purpose I/O pin.

5.1 PORTA and TRISA Registers

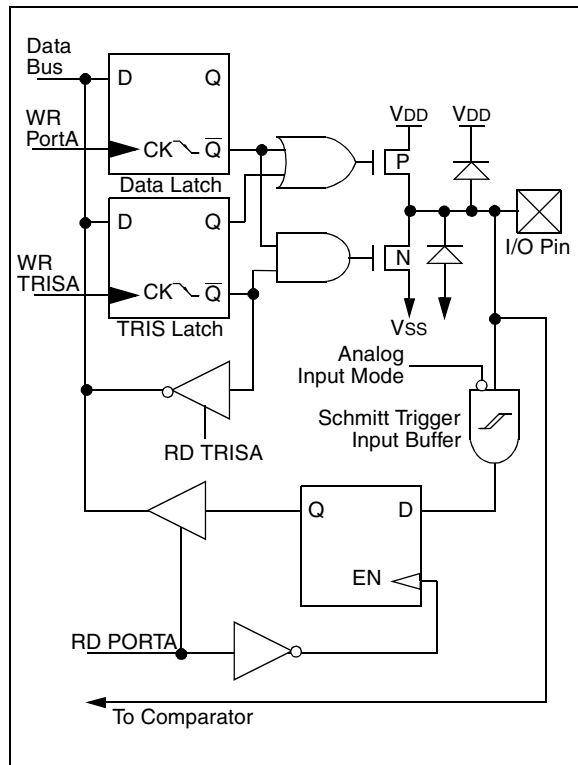
PORTA is a 5-bit wide latch. RA4 is a Schmitt Trigger input and an open drain output. Port RA4 is multiplexed with the T0CKI clock input. All other RA port pins have Schmitt Trigger input levels and full CMOS output drivers. All pins have data direction bits (TRIS registers), which can configure these pins as input or output.

A '1' in the TRISA register puts the corresponding output driver in a hi-impedance mode. A '0' in the TRISA register puts the contents of the output latch on the selected pin(s).

Reading the PORTA register reads the status of the pins, whereas writing to it will write to the port latch. All write operations are read-modify-write operations. So a write to a port implies that the port pins are first read, then this value is modified and written to the port data latch.

The PORTA pins are multiplexed with comparator and voltage reference functions. The operation of these pins are selected by control bits in the CMCON (Comparator Control Register) register and the VRCON (Voltage Reference Control Register) register. When selected as a comparator input, these pins will read as '0's.

FIGURE 5-1: BLOCK DIAGRAM OF RA<1:0> PINS



Note: On reset, the TRISA register is set to all inputs. The digital inputs are disabled and the comparator inputs are forced to ground to reduce excess current consumption.

TRISA controls the direction of the RA pins, even when they are being used as comparator inputs. The user must make sure to keep the pins configured as inputs when using them as comparator inputs.

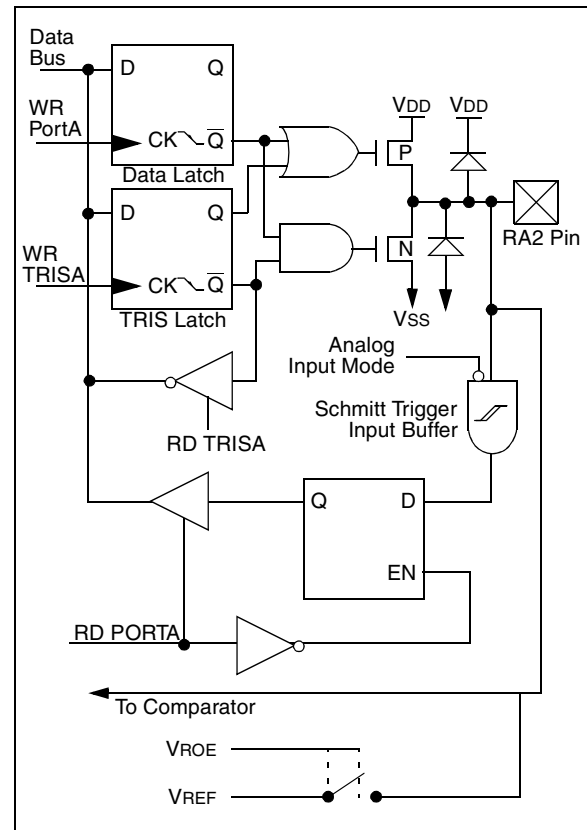
The RA2 pin will also function as the output for the voltage reference. When in this mode, the VREF pin is a very high impedance output. The user must configure TRISA<2> bit as an input and use high impedance loads.

In one of the comparator modes defined by the CMCON register, pins RA3 and RA4 become outputs of the comparators. The TRISA<4:3> bits must be cleared to enable outputs to use this function.

EXAMPLE 5-1: INITIALIZING PORTA

```
CLRF PORTA      ;Initialize PORTA by setting
                  ;output data latches
MOVWLW 0X07     ;Turn comparators off and
MOVWF CMCON     ;enable pins for I/O
                  ;functions
BSF STATUS, RP0 ;Select Bank1
MOVWLW 0x1F     ;Value used to initialize
                  ;data direction
MOVWF TRISA     ;Set RA<4:0> as inputs
                  ;TRISA<7:5> are always
                  ;read as '0'.
```

FIGURE 5-2: BLOCK DIAGRAM OF RA2 PIN



5.2 PORTB and TRISB Registers

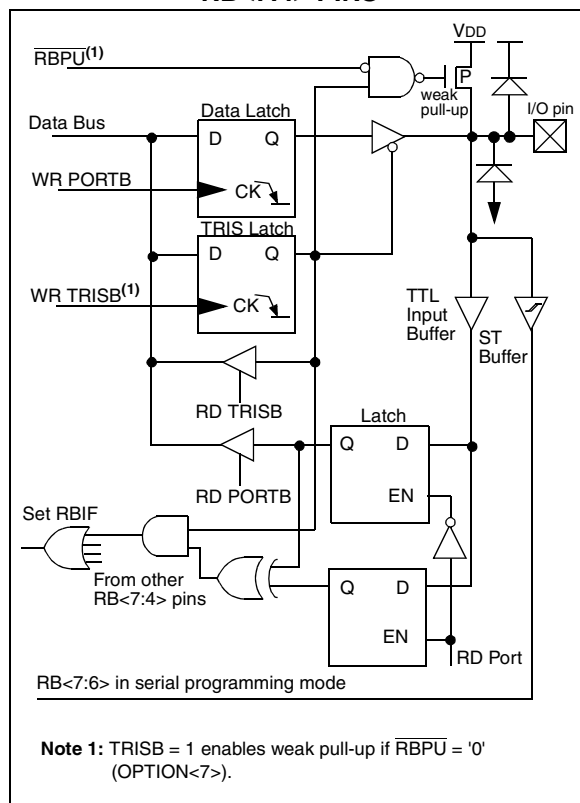
PORTB is an 8-bit wide, bi-directional port. The corresponding data direction register is TRISB. A '1' in the TRISB register puts the corresponding output driver in a high impedance mode. A '0' in the TRISB register puts the contents of the output latch on the selected pin(s).

Reading PORTB register reads the status of the pins, whereas writing to it will write to the port latch. All write operations are read-modify-write operations. So a write to a port implies that the port pins are first read, then this value is modified and written to the port data latch.

Each of the PORTB pins has a weak internal pull-up ($\approx 200 \mu\text{A}$ typical). A single control bit can turn on all the pull-ups. This is done by clearing the $\overline{\text{RBPU}}$ (OPTION<7>) bit. The weak pull-up is automatically turned off when the port pin is configured as an output. The pull-ups are disabled on Power-on Reset.

Four of PORTB's pins, RB<7:4>, have an interrupt on change feature. Only pins configured as inputs can cause this interrupt to occur (i.e., any RB<7:4> pin configured as an output is excluded from the interrupt on change comparison). The input pins of RB<7:4> are compared with the old value latched on the last read of PORTB. The "mismatch" outputs of RB<7:4> are OR'ed together to generate the RBIF interrupt (flag latched in INTCON<0>).

FIGURE 5-5: BLOCK DIAGRAM OF RB<7:4> PINS



This interrupt can wake the device from SLEEP. The user, in the interrupt service routine, can clear the interrupt in the following manner:

- Any read or write of PORTB. This will end the mismatch condition.
- Clear flag bit RBIF.

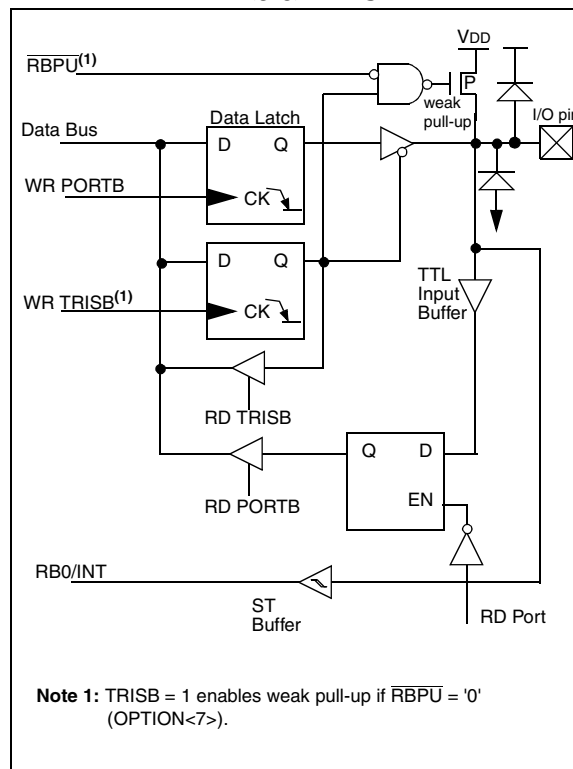
A mismatch condition will continue to set flag bit RBIF. Reading PORTB will end the mismatch condition and allow flag bit RBIF to be cleared.

This interrupt on mismatch feature, together with software configurable pull-ups on these four pins allow easy interface to a key pad and make it possible for wake-up on key-depression. (See AN552, "Implementing Wake-Up on Key Strokes".)

Note: If a change on the I/O pin should occur when the read operation is being executed (start of the Q2 cycle), then the RBIF interrupt flag may not get set.

The interrupt on change feature is recommended for wake-up on key depression operation and operations where PORTB is only used for the interrupt on change feature. Polling of PORTB is not recommended while using the interrupt on change feature.

FIGURE 5-6: BLOCK DIAGRAM OF RB<3:0> PINS



6.1 Bus Characteristics

In this section, the term “processor” refers to the portion of the PIC16CE62X that interfaces to the EEPROM through software manipulating the EEINTF register. The following **bus protocol** is to be used with the EEPROM data memory.

- Data transfer may be initiated only when the bus is not busy.
- During data transfer, the data line must remain stable whenever the clock line is HIGH. Changes in the data line while the clock line is HIGH will be interpreted by the EEPROM as a START or STOP condition.

Accordingly, the following bus conditions have been defined (Figure 6-1).

6.1.1 BUS NOT BUSY (A)

Both data and clock lines remain HIGH.

6.1.2 START DATA TRANSFER (B)

A HIGH to LOW transition of the SDA line while the clock (SCL) is HIGH determines a START condition. All commands must be preceded by a START condition.

6.1.3 STOP DATA TRANSFER (C)

A LOW to HIGH transition of the SDA line while the clock (SCL) is HIGH determines a STOP condition. All operations must be ended with a STOP condition.

6.1.4 DATA VALID (D)

The state of the data line represents valid data when, after a START condition, the data line is stable for the duration of the HIGH period of the clock signal.

The data on the line must be changed during the LOW period of the clock signal. There is one bit of data per clock pulse.

Each data transfer is initiated with a START condition and terminated with a STOP condition. The number of the data bytes transferred between the START and STOP conditions is determined by the processor and is theoretically unlimited, although only the last sixteen will be stored when doing a write operation. When an overwrite does occur, it will replace data in a first-in, first-out fashion.

6.1.5 ACKNOWLEDGE

The EEPROM will generate an acknowledge after the reception of each byte. The processor must generate an extra clock pulse which is associated with this acknowledge bit.

Note: Acknowledge bits are not generated if an internal programming cycle is in progress.
--

When the EEPROM acknowledges, it pulls down the SDA line during the acknowledge clock pulse in such a way that the SDA line is stable LOW during the HIGH period of the acknowledge related clock pulse. Of course, setup and hold times must be taken into account. The processor must signal an end of data to the EEPROM by not generating an acknowledge bit on the last byte that has been clocked out of the EEPROM. In this case, the EEPROM must leave the data line HIGH to enable the processor to generate the STOP condition (Figure 6-2).

7.2 Using Timer0 with External Clock

When an external clock input is used for Timer0, it must meet certain requirements. The external clock requirement is due to internal phase clock (TOSC) synchronization. Also, there is a delay in the actual incrementing of Timer0 after synchronization.

7.2.1 EXTERNAL CLOCK SYNCHRONIZATION

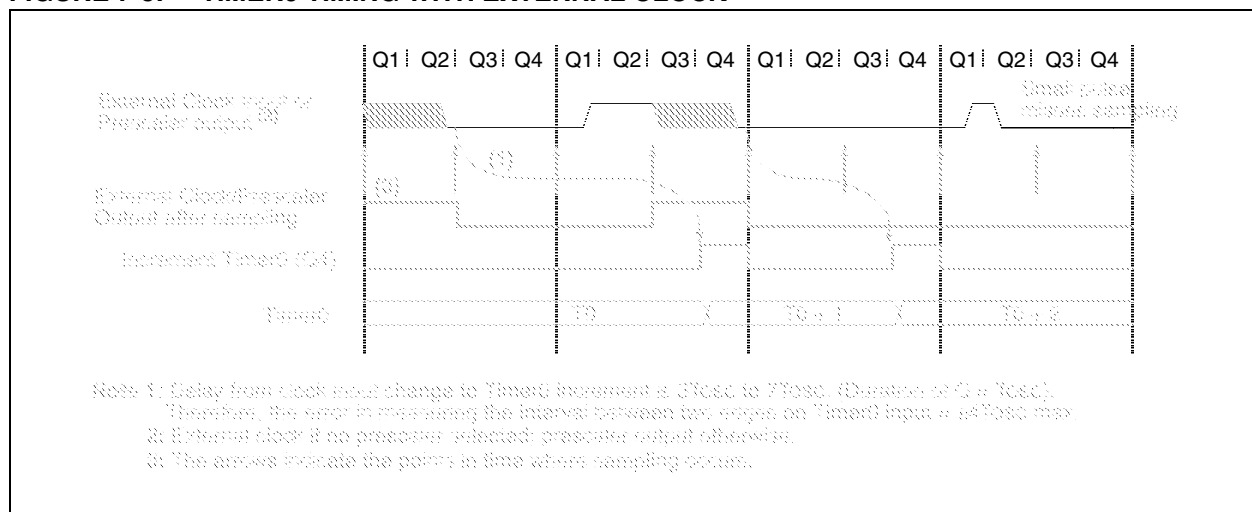
When no prescaler is used, the external clock input is the same as the prescaler output. The synchronization of T0CKI with the internal phase clocks is accomplished by sampling the prescaler output on the Q2 and Q4 cycles of the internal phase clocks (Figure 7-5). Therefore, it is necessary for T0CKI to be high for at least 2TOSC (and a small RC delay of 20 ns) and low for at least 2TOSC (and a small RC delay of 20 ns). Refer to the electrical specification of the desired device.

When a prescaler is used, the external clock input is divided by the asynchronous ripple-counter type prescaler so that the prescaler output is symmetrical. For the external clock to meet the sampling requirement, the ripple-counter must be taken into account. Therefore, it is necessary for T0CKI to have a period of at least 4TOSC (and a small RC delay of 40 ns) divided by the prescaler value. The only requirement on T0CKI high and low time is that they do not violate the minimum pulse width requirement of 10 ns. Refer to parameters 40, 41 and 42 in the electrical specification of the desired device.

7.2.2 TIMER0 INCREMENT DELAY

Since the prescaler output is synchronized with the internal clocks, there is a small delay from the time the external clock edge occurs to the time the TMR0 is actually incremented. Figure 7-5 shows the delay from the external clock edge to the timer incrementing.

FIGURE 7-5: TIMER0 TIMING WITH EXTERNAL CLOCK



10.0 SPECIAL FEATURES OF THE CPU

Special circuits to deal with the needs of real time applications are what sets a microcontroller apart from other processors. The PIC16CE62X family has a host of such features intended to maximize system reliability, minimize cost through elimination of external components, provide power saving operating modes and offer code protection.

These are:

1. OSC selection
2. Reset
 - Power-on Reset (POR)
 - Power-up Timer (PWRT)
 - Oscillator Start-Up Timer (OST)
 - Brown-out Reset (BOD)
3. Interrupts
4. Watchdog Timer (WDT)
5. SLEEP
6. Code protection
7. ID Locations
8. In-circuit serial programming

The PIC16CE62X has a Watchdog Timer which is controlled by configuration bits. It runs off its own RC oscillator for added reliability. There are two timers that offer necessary delays on power-up. One is the Oscillator Start-up Timer (OST), intended to keep the chip in reset until the crystal oscillator is stable. The other is the Power-up Timer (PWRT), which provides a fixed delay of 72 ms (nominal) on power-up only, and is designed to keep the part in reset while the power supply stabilizes. There is also circuitry to reset the device if a brown-out occurs, which provides at least a 72 ms reset. With these three functions on-chip, most applications need no external reset circuitry.

The SLEEP mode is designed to offer a very low current power-down mode. The user can wake-up from SLEEP through external reset, Watchdog Timer wake-up or through an interrupt. Several oscillator options are also made available to allow the part to fit the application. The RC oscillator option saves system cost, while the LP crystal option saves power. A set of configuration bits are used to select various options.

GOTO		Unconditional Branch							
Syntax:	[<i>label</i>] GOTO k								
Operands:	$0 \leq k \leq 2047$								
Operation:	$k \rightarrow PC<10:0>$ $PCLATH<4:3> \rightarrow PC<12:11>$								
Status Affected:	None								
Encoding:	<table><tr><td>10</td><td>1kkk</td><td>kkkk</td><td>kkkk</td></tr></table>					10	1kkk	kkkk	kkkk
10	1kkk	kkkk	kkkk						
Description:	<p>GOTO is an unconditional branch. The eleven bit immediate value is loaded into PC bits <10:0>. The upper bits of PC are loaded from PCLATH<4:3>.</p> <p>GOTO is a two-cycle instruction.</p>								
Words:	1								
Cycles:	2								
Example	<pre> GOTO THERE After Instruction PC = Address THERE</pre>								

INCFSZ		Increment f, Skip if 0	
Syntax:	[<i>label</i>] INCFSZ f,d		
Operands:	$0 \leq f \leq 127$ $d \in [0,1]$		
Operation:	$(f) + 1 \rightarrow (\text{dest})$, skip if result = 0		
Status Affected:	None		
Encoding:	00	1111	dfff ffff
Description:	<p>The contents of register 'f' are incremented. If 'd' is 0, the result is placed in the W register. If 'd' is 1, the result is placed back in register 'f'. If the result is 0, the next instruction, which is already fetched, is discarded. A NOP is executed instead making it a two-cycle instruction.</p>		
Words:	1		
Cycles:	1(2)		
Example	HERE	INCFSZ	CNT, 1

Before Instruction
 PC = address HERE
 After Instruction
 CNT = CNT + 1
 if CNT= 0,
 PC = address CONTINUE
 if CNT≠ 0,
 PC = address HERE +1

INCF		Increment f								
Syntax:	[<i>label</i>] INCF f,d									
Operands:	$0 \leq f \leq 127$ $d \in [0,1]$									
Operation:	$(f) + 1 \rightarrow (\text{dest})$									
Status Affected:	Z									
Encoding:	<table><tr><td>00</td><td>1010</td><td>dfff</td><td>ffff</td></tr></table>						00	1010	dfff	ffff
00	1010	dfff	ffff							
Description:	The contents of register 'f' are incremented. If 'd' is 0, the result is placed in the W register. If 'd' is 1, the result is placed back in register 'f'.									
Words:	1									
Cycles:	1									
Example	INCF CNT, 1									
	Before Instruction									
	CNT	=	0xFF							
	Z	=	0							
	After Instruction									
	CNT	=	0x00							
	Z	=	1							

IORLW		Inclusive OR Literal with W							
Syntax:	[<i>label</i>] IORLW k								
Operands:	$0 \leq k \leq 255$								
Operation:	(W) .OR. k \rightarrow (W)								
Status Affected:	Z								
Encoding:	<table border="1"><tr><td>11</td><td>1000</td><td>kkkk</td><td>kkkk</td></tr></table>				11	1000	kkkk	kkkk	
11	1000	kkkk	kkkk						
Description:	The contents of the W register are OR'ed with the eight bit literal 'k'. The result is placed in the W register.								
Words:	1								
Cycles:	1								
Example	IORLW 0x35								
	Before Instruction								
	W = 0x9A								
	After Instruction								
	W = 0xBF								
	Z = 1								

RETURN Return from Subroutine

Syntax: [*label*] RETURN

Operands: None

Operation: TOS → PC

Status Affected: None

Encoding:

00	0000	0000	1000
----	------	------	------

Description: Return from subroutine. The stack is POPed and the top of the stack (TOS) is loaded into the program counter. This is a two cycle instruction.

Words: 1

Cycles: 2

Example
 RETURN
 After Interrupt
 PC = TOS

RRF Rotate Right f through Carry

Syntax: [*label*] RRF f,d

Operands: $0 \leq f \leq 127$
 $d \in [0,1]$

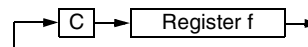
Operation: See description below

Status Affected: C

Encoding:

00	1100	dfff	ffff
----	------	------	------

Description: The contents of register 'f' are rotated one bit to the right through the Carry Flag. If 'd' is 0, the result is placed in the W register. If 'd' is 1, the result is placed back in register 'f'.



Words: 1

Cycles: 1

Example RRF REG1, 0

Before Instruction

REG1 = 1110 0110
 C = 0

After Instruction

REG1 = 1110 0110
 W = 0111 0011
 C = 0

RLF Rotate Left f through Carry

Syntax: [*label*] RLF f,d

Operands: $0 \leq f \leq 127$
 $d \in [0,1]$

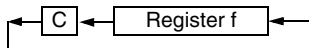
Operation: See description below

Status Affected: C

Encoding:

00	1101	dfff	ffff
----	------	------	------

Description: The contents of register 'f' are rotated one bit to the left through the Carry Flag. If 'd' is 0, the result is placed in the W register. If 'd' is 1, the result is stored back in register 'f'.



Words: 1

Cycles: 1

Example RLF REG1, 0
 Before Instruction
 REG1 = 1110 0110
 C = 0
 After Instruction
 REG1 = 1110 0110
 W = 1100 1100
 C = 1

SLEEP

Syntax: [*label*] SLEEP

Operands: None

Operation: 00h → WDT,
 0 → WDT prescaler,
 1 → $\overline{\text{TO}}$,
 0 → $\overline{\text{PD}}$

Status Affected: $\overline{\text{TO}}$, $\overline{\text{PD}}$

Encoding:

00	0000	0110	0011
----	------	------	------

Description: The power-down status bit, $\overline{\text{PD}}$ is cleared. Time-out status bit, $\overline{\text{TO}}$ is set. Watchdog Timer and its prescaler are cleared. The processor is put into SLEEP mode with the oscillator stopped. See Section 10.8 for more details.

Words: 1

Cycles: 1

Example: SLEEP

MPLIB is a librarian for pre-compiled code to be used with MPLINK. When a routine from a library is called from another source file, only the modules that contains that routine will be linked in with the application. This allows large libraries to be used efficiently in many different applications. MPLIB manages the creation and modification of library files.

MPLINK features include:

- MPLINK works with MPASM and MPLAB-C17 and MPLAB-C18.
- MPLINK allows all memory areas to be defined as sections to provide link-time flexibility.

MPLIB features include:

- MPLIB makes linking easier because single libraries can be included instead of many smaller files.
- MPLIB helps keep code maintainable by grouping related modules together.
- MPLIB commands allow libraries to be created and modules to be added, listed, replaced, deleted, or extracted.

12.5 MPLAB-SIM Software Simulator

The MPLAB-SIM Software Simulator allows code development in a PC host environment by simulating the PIC series microcontrollers on an instruction level. On any given instruction, the data areas can be examined or modified and stimuli can be applied from a file or user-defined key press to any of the pins. The execution can be performed in single step, execute until break, or trace mode.

MPLAB-SIM fully supports symbolic debugging using MPLAB-C17 and MPLAB-C18 and MPASM. The Software Simulator offers the flexibility to develop and debug code outside of the laboratory environment making it an excellent multi-project software development tool.

12.6 MPLAB-ICE High Performance Universal In-Circuit Emulator with MPLAB IDE

The MPLAB-ICE Universal In-Circuit Emulator is intended to provide the product development engineer with a complete microcontroller design tool set for PIC microcontrollers (MCUs). Software control of MPLAB-ICE is provided by the MPLAB Integrated Development Environment (IDE), which allows editing, “make” and download, and source debugging from a single environment.

Interchangeable processor modules allow the system to be easily reconfigured for emulation of different processors. The universal architecture of the MPLAB-ICE allows expansion to support new PIC microcontrollers.

The MPLAB-ICE Emulator System has been designed as a real-time emulation system with advanced features that are generally found on more expensive devel-

opment tools. The PC platform and Microsoft® Windows 3.x/95/98 environment were chosen to best make these features available to you, the end user.

MPLAB-ICE 2000 is a full-featured emulator system with enhanced trace, trigger, and data monitoring features. Both systems use the same processor modules and will operate across the full operating speed range of the PIC MCU.

12.7 PICMASTER/PICMASTER CE

The PICMASTER system from Microchip Technology is a full-featured, professional quality emulator system. This flexible in-circuit emulator provides a high-quality, universal platform for emulating Microchip 8-bit PIC microcontrollers (MCUs). PICMASTER systems are sold worldwide, with a CE compliant model available for European Union (EU) countries.

12.8 ICEPIC

ICEPIC is a low-cost in-circuit emulation solution for the Microchip Technology PIC16C5X, PIC16C6X, PIC16C7X, and PIC16CXXX families of 8-bit one-time-programmable (OTP) microcontrollers. The modular system can support different subsets of PIC16C5X or PIC16CXXX products through the use of interchangeable personality modules or daughter boards. The emulator is capable of emulating without target application circuitry being present.

12.9 MPLAB-ICD In-Circuit Debugger

Microchip's In-Circuit Debugger, MPLAB-ICD, is a powerful, low-cost run-time development tool. This tool is based on the flash PIC16F877 and can be used to develop for this and other PIC microcontrollers from the PIC16CXXX family. MPLAB-ICD utilizes the In-Circuit Debugging capability built into the PIC16F87X. This feature, along with Microchip's In-Circuit Serial Programming protocol, offers cost-effective in-circuit flash programming and debugging from the graphical user interface of the MPLAB Integrated Development Environment. This enables a designer to develop and debug source code by watching variables, single-stepping and setting break points. Running at full speed enables testing hardware in real-time. The MPLAB-ICD is also a programmer for the flash PIC16F87X family.

12.10 PRO MATE II Universal Programmer

The PRO MATE II Universal Programmer is a full-featured programmer capable of operating in stand-alone mode as well as PC-hosted mode. PRO MATE II is CE compliant.

The PRO MATE II has programmable VDD and VPP supplies which allows it to verify programmed memory at VDD min and VDD max for maximum reliability. It has an LCD display for instructions and error messages, keys to enter commands and a modular detachable socket assembly to support various package types. In

TABLE 12-1: DEVELOPMENT TOOLS FROM MICROCHIP

	PIC12CXX	PIC14000	PIC16C5X	PIC16C6X	PIC16CXX	PIC16F62X	PIC16C7X	PIC16C7XX	PIC16C8X	PIC16F8XX	PIC16C9XX	PIC17C4X	PIC17C7XX	PIC18CXX2	24CXX/ 25CXX/ 93CXX	HCSXX	MCRFXX	MCP2510
Demo Boards and Eval Kits	MPLAB® Integrated Development Environment	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓				✓
	MPLAB® C17 Compiler											✓	✓					
	MPLAB® C18 Compiler													✓				
	MPASM/MPLINK	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓		
Emulators	MPLAB®-ICE	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓			
	PICMASTER/PICMASTER-CE	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓			
Debuggers	ICEPIC™ Low-Cost In-Circuit Emulator	✓	✓	✓	✓	✓	✓	✓	✓		✓							
	MPLAB®-ICD In-Circuit Debugger				✓		✓			✓								
Programmers	PICSTART® Plus Low-Cost Universal Dev. Kit	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓			
	PRO MATE® II Universal Programmer	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓		
Demo Boards and Eval Kits	SIMICE	✓	✓	✓														
	PICDEM-1		✓	✓			†		✓			✓						
	PICDEM-2				†		†							✓				
	PICDEM-3										✓							
	PICDEM-14A		✓										✓					
	PICDEM-17												✓					
	KEELOO® Evaluation Kit														✓			
	KEELOO Transponder Kit														✓			
	microID™ Programmer's Kit															✓		
	125 kHz microID Developer's Kit																✓	
	125 kHz Anticollision microID Developer's Kit																✓	
	13.56 MHz Anticollision microID Developer's Kit																✓	
	MCP2510 CAN Developer's Kit																✓	✓

* Contact the Microchip Technology Inc. web site at www.microchip.com for information on how to use the MPLAB®-ICD In-Circuit Debugger (DV164001) with PIC16C62, 63, 64, 65, 72, 73, 74, 76, 77

** Contact Microchip Technology Inc. for availability date.

† Development tool is available on select devices.

FIGURE 13-3: PIC16LCE62X VOLTAGE-FREQUENCY GRAPH, $-40^{\circ}\text{C} \leq T_A < +125^{\circ}\text{C}$

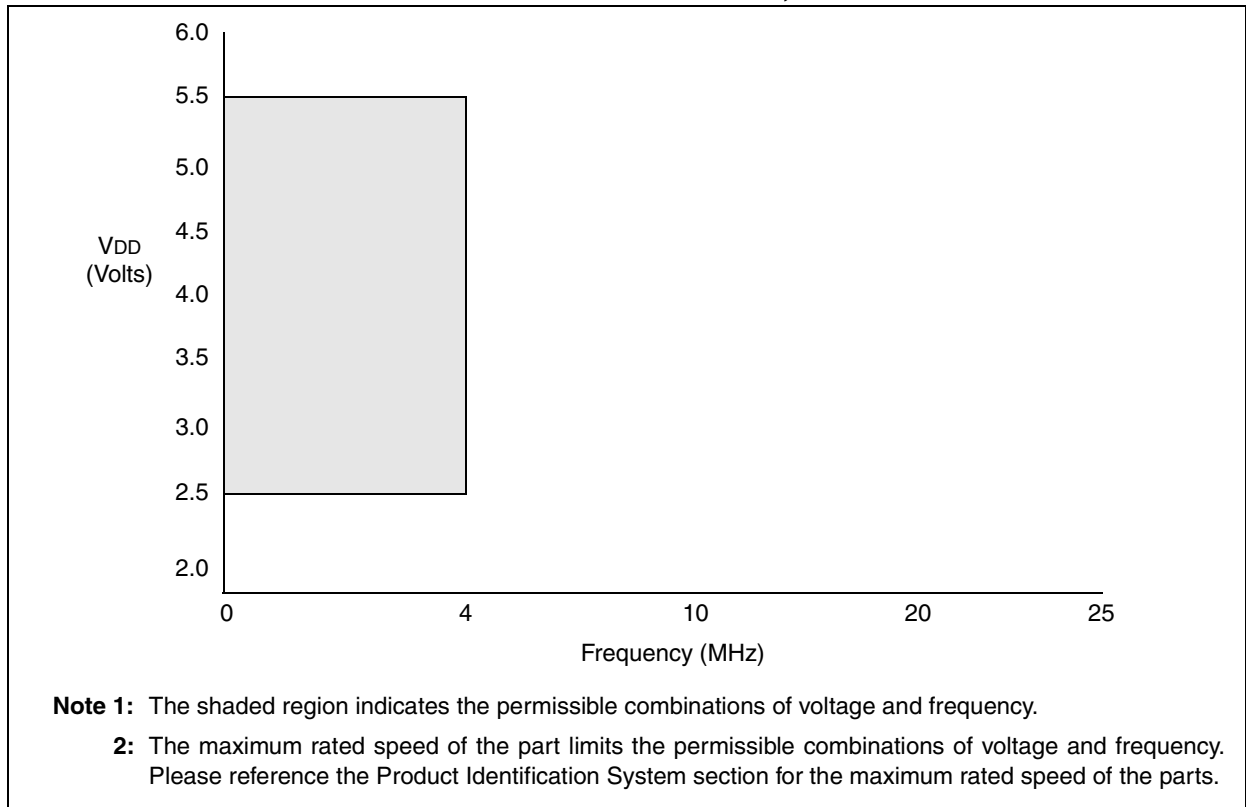


TABLE 13-1: COMPARATOR SPECIFICATIONS

Operating Conditions: VDD range as described in Table 12-1, -40°C<TA<+125°C. .

Param No.	Characteristics	Sym	Min	Typ	Max	Units	Comments
D300	Input offset voltage	VIOFF		± 5.0	± 10	mV	
D301	Input common mode voltage	VICM	0		VDD - 1.5	V	
D302	CMRR	CMRR	+55*			db	
300	Response Time ⁽¹⁾	TRESP		150*	400*	ns	PIC16CE62X
301	Comparator Mode Change to Output Valid	TMC2OV			10*	µs	

* These parameters are characterized but not tested.

Note 1: Response time measured with one comparator input at (VDD - 1.5)/2 while the other input transitions from VSS to VDD.

TABLE 13-2: VOLTAGE REFERENCE SPECIFICATIONS

Operating Conditions: VDD range as described in Table 12-1, -40°C<TA<+125°C.

Param No.	Characteristics	Sym	Min	Typ	Max	Units	Comments
D310	Resolution	VRES	VDD/24		VDD/32	LSB	
D311	Absolute Accuracy	VRAA			±1/4 ±1/2	LSB LSB	Low Range (VRR=1) High Range (VRR=0)
D312	Unit Resistor Value (R)	VRUR		2K*		Ω	Figure 9-1
310	Settling Time ⁽¹⁾	TSET			10*	µs	

* These parameters are characterized but not tested.

Note 1: Settling time measured while VRR = 1 and VR<3:0> transitions from 0000 to 1111.

FIGURE 13-6: CLKOUT AND I/O TIMING

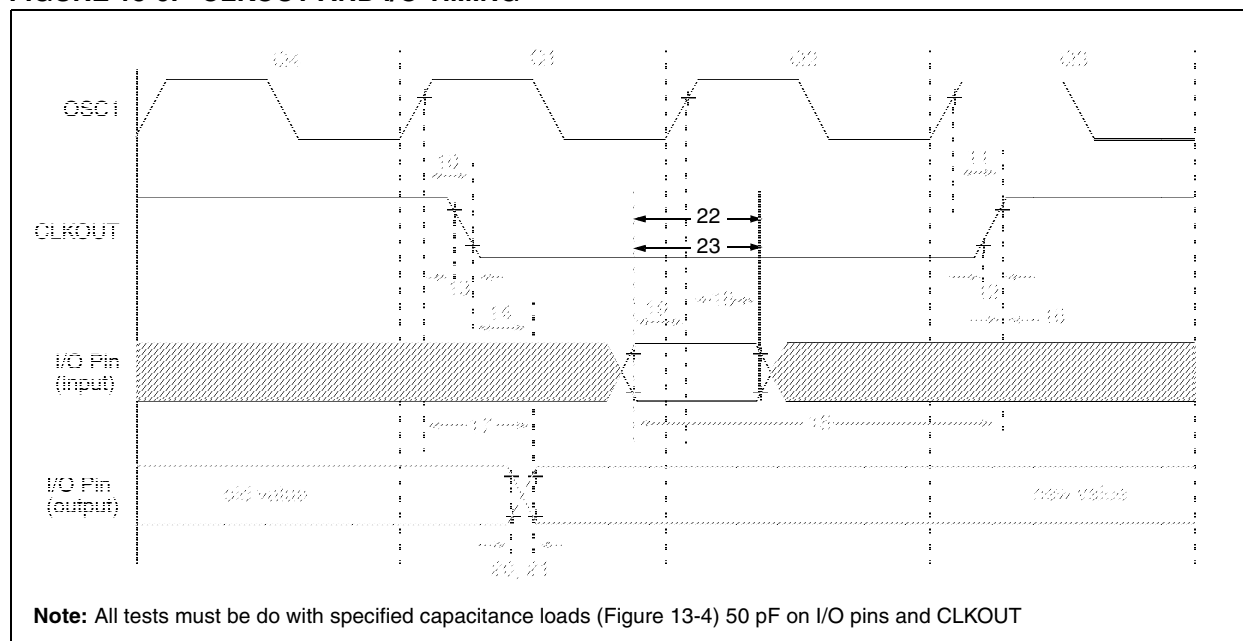


TABLE 13-4: CLKOUT AND I/O TIMING REQUIREMENTS

Parameter #	Sym	Characteristic	Min	Typ†	Max	Units
10*	TosH2ckL	OSC1↑ to CLKOUT↓ (1)	—	75	200	ns
11*	TosH2ckH	OSC1↑ to CLKOUT↑ (1)	—	75	200	ns
12*	TckR	CLKOUT rise time (1)	—	35	100	ns
13*	TckF	CLKOUT fall time (1)	—	35	100	ns
14*	TckL2ioV	CLKOUT ↓ to Port out valid (1)	—	—	20	ns
15*	TioV2ckH	Port in valid before CLKOUT ↑ (1)	Tosc +200 ns	—	—	ns
16*	TckH2ioI	Port in hold after CLKOUT ↑ (1)	0	—	—	ns
17*	TosH2ioV	OSC1↑ (Q1 cycle) to Port out valid	—	50	150	ns
18*	TosH2ioI	OSC1↑ (Q2 cycle) to Port input invalid (I/O in hold time)	100	—	—	ns
19*	TioV2osH	Port input valid to OSC1↑ (I/O in setup time)	0	—	—	ns
20*	TioR	Port output rise time	—	10	40	ns
21*	TioF	Port output fall time	—	10	40	ns
22*	Tinp	RB0/INT pin high or low time	25	—	—	ns
23	Trbp	RB<7:4> change interrupt high or low time	Tcy	—	—	ns

* These parameters are characterized but not tested

† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

Note 1: Measurements are taken in RC Mode where CLKOUT output is 4 x TOSC.

FIGURE 13-9: TIMER0 CLOCK TIMING

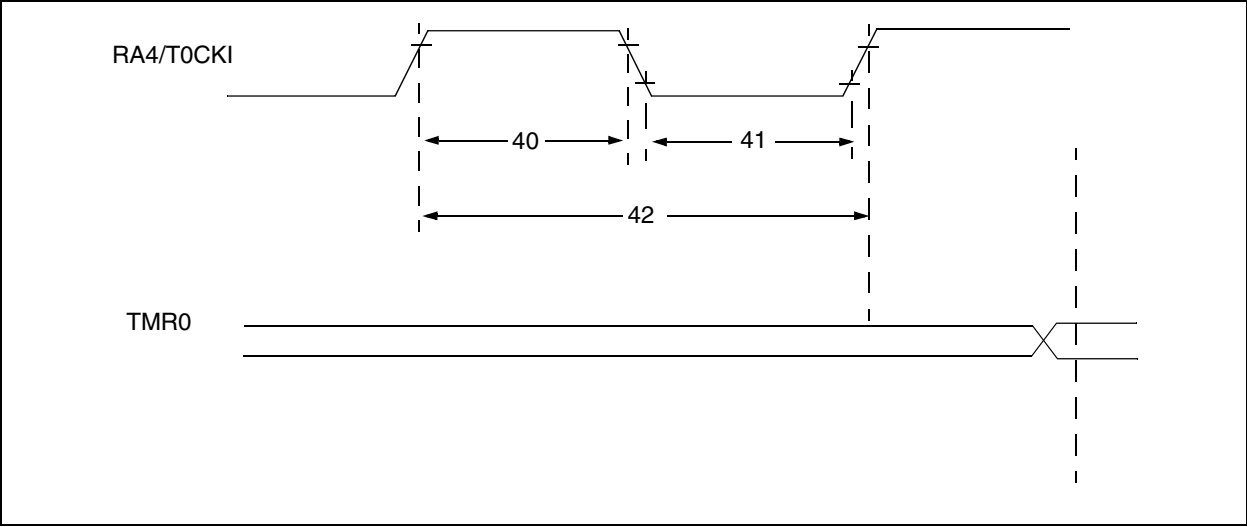


TABLE 13-6: TIMER0 CLOCK REQUIREMENTS

Parameter No.	Sym	Characteristic		Min	Typ†	Max	Units	Conditions
40	Tt0H	T0CKI High Pulse Width	No Prescaler	0.5 TCY + 20*	—	—	ns	
			With Prescaler	10*	—	—	ns	
41	Tt0L	T0CKI Low Pulse Width	No Prescaler	0.5 TCY + 20*	—	—	ns	
			With Prescaler	10*	—	—	ns	
42	Tt0P	T0CKI Period		$\frac{TCY + 40^*}{N}$	—	—	ns	N = prescale value (1, 2, 4, ..., 256)

* These parameters are characterized but not tested.
† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

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NOTES: