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"[Embedded - Microcontrollers](#)" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

Applications of "[Embedded - Microcontrollers](#)"

Details

Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	10MHz
Connectivity	-
Peripherals	POR, WDT
Number of I/O	5
Program Memory Size	1.75KB (1K x 14)
Program Memory Type	OTP
EEPROM Size	16 x 8
RAM Size	128 x 8
Voltage - Supply (Vcc/Vdd)	3V ~ 5.5V
Data Converters	A/D 4x8b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 125°C (TA)
Mounting Type	Through Hole
Package / Case	8-DIP (0.300", 7.62mm)
Supplier Device Package	8-PDIP
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic12ce673-10e-p

4.0 MEMORY ORGANIZATION

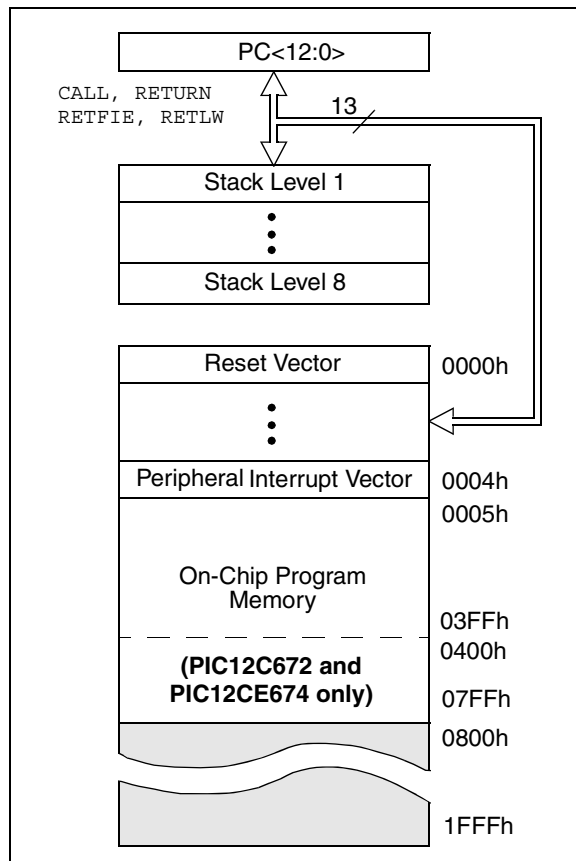
4.1 Program Memory Organization

The PIC12C67X has a 13-bit program counter capable of addressing an 8K x 14 program memory space.

For the PIC12C671 and the PIC12CE673, the first 1K x 14 (0000h-03FFh) is implemented.

For the PIC12C672 and the PIC12CE674, the first 2K x 14 (0000h-07FFh) is implemented. Accessing a location above the physically implemented address will cause a wraparound. The reset vector is at 0000h and the interrupt vector is at 0004h.

FIGURE 4-1: PIC12C67X PROGRAM MEMORY MAP AND STACK



4.2 Data Memory Organization

The data memory is partitioned into two banks, which contain the General Purpose Registers and the Special Function Registers. Bit RP0 is the bank select bit.

RP0 (STATUS<5>) = 1 → Bank 1

RP0 (STATUS<5>) = 0 → Bank 0

Each Bank extends up to 7Fh (128 bytes). The lower locations of each Bank are reserved for the Special Function Registers. Above the Special Function Registers are General Purpose Registers implemented as static RAM. Both Bank 0 and Bank 1 contain Special Function Registers. Some "high use" Special Function Registers from Bank 0 are mirrored in Bank 1 for code reduction and quicker access.

Also note that F0h through FFh on the PIC12C67X is mapped into Bank 0 registers 70h-7Fh as common RAM.

4.2.1 GENERAL PURPOSE REGISTER FILE

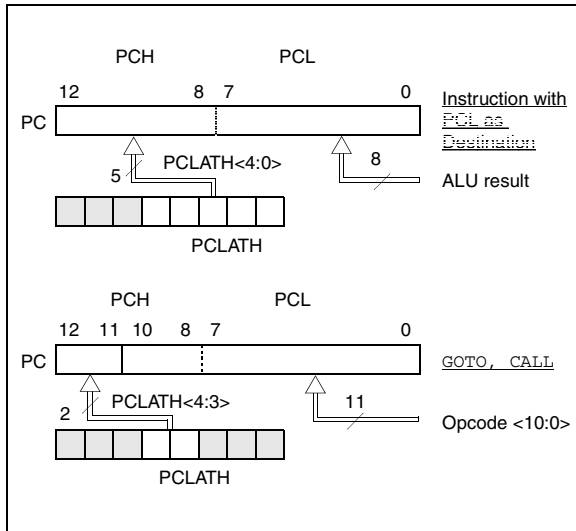
The register file can be accessed either directly or indirectly through the File Select Register FSR (Section 4.5).

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4.3 PCL and PCLATH

The Program Counter (PC) is 13-bits wide. The low byte comes from the PCL Register, which is a readable and writable register. The high byte (PC<12:8>) is not directly readable or writable and comes from PCLATH. On any reset, the PC is cleared. Figure 4-3 shows the two situations for the loading of the PC. The upper example in the figure shows how the PC is loaded on a write to PCL (PCLATH<4:0> → PCH). The lower example in the figure shows how the PC is loaded during a CALL or GOTO instruction (PCLATH<4:3> → PCH).

FIGURE 4-3: LOADING OF PC IN DIFFERENT SITUATIONS



4.3.1 COMPUTED GOTO

A Computed GOTO is accomplished by adding an offset to the program counter (`ADDWF PCL`). When doing a table read using a computed GOTO method, care should be exercised if the table location crosses a PCL memory boundary (each 256 byte block). Refer to the application note *“Implementing a Table Read”* (AN556).

4.3.2 STACK

The PIC12C67X family has an 8-level deep x 13-bit wide hardware stack. The stack space is not part of either program or data space and the stack pointer is not readable or writable. The PC is PUSHed onto the stack when a `CALL` instruction is executed or an interrupt causes a branch. The stack is POPed in the event of a `RETURN`, `RETLW` or a `RETFIE` instruction execution. PCLATH is not affected by a PUSH or POP operation.

The stack operates as a circular buffer. This means that after the stack has been PUSHed eight times, the ninth push overwrites the value that was stored from the first push. The tenth push overwrites the second push (and so on).

Note 1: There are no status bits to indicate stack overflow or stack underflow conditions.

2: There are no instructions/mnemonics called PUSH or POP. These are actions that occur from the execution of the `CALL`, `RETURN`, `RETLW`, and `RETFIE` instructions, or the vectoring to an interrupt address.

4.4 Program Memory Paging

The PIC12C67X ignores both paging bits PCLATH<4:3>, which are used to access program memory when more than one page is available. The use of PCLATH<4:3> as general purpose read/write bits for the PIC12C67X is not recommended since this may affect upward compatibility with future products.

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6.3 Write Operations

6.3.1 BYTE WRITE

Following the start signal from the processor, the device code (4 bits), the don't care bits (3 bits), and the R/W bit (which is a logic low) are placed onto the bus by the processor. This indicates to the addressed EEPROM that a byte with a word address will follow after it has generated an acknowledge bit during the ninth clock cycle. Therefore, the next byte transmitted by the processor is the word address and will be written into the address pointer. Only the lower four address bits are used by the device, and the upper four bits are don't cares. If the address byte is acknowledged, the processor will then transmit the data word to be written into the addressed memory location. The memory acknowledges again and the processor generates a stop condition. This initiates the internal write cycle, and during this time will not generate acknowledge signals. After a byte write command, the internal address counter will not be incremented and will point to the same address location that was just written. If a stop bit sequence is transmitted to the device at any point in the write command sequence before the entire sequence is complete, then the command will abort and no data will be written. If more than 8 data bits are transmitted before the stop bit sequence is sent, then the device will clear the previously loaded byte and begin loading the data buffer again. If more than one data byte is transmitted to the device and a stop bit is sent before a full eight data bits have been transmitted, then the write command will abort and no data will be written. The EEPROM memory employs a VCC threshold detector circuit, which disables the internal erase/write logic if the VCC is below minimum VDD. Byte write operations must be preceded and immediately followed by a bus not busy bus cycle where both SDA and SCL are held high. (See Figure 6-7 for Byte Write operation.)

6.4 Acknowledge Polling

Since the EEPROM will not acknowledge during a write cycle, this can be used to determine when the cycle is complete (this feature can be used to maximize bus throughput). Once the stop condition for a write command has been issued from the processor, the device initiates the internally timed write cycle. ACK polling can be initiated immediately. This involves the processor sending a start condition followed by the control byte for a write command ($R/\bar{W} = 0$). If the device is still busy with the write cycle, then no ACK will be returned. If no ACK is returned, then the start bit and control byte must be re-sent. If the cycle is complete, then the device will return the ACK and the processor can then proceed with the next read or write command. (See Figure 6-6 for flow diagram.)

FIGURE 6-6: ACKNOWLEDGE POLLING FLOW



FIGURE 6-7: BYTE WRITE



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NOTES:

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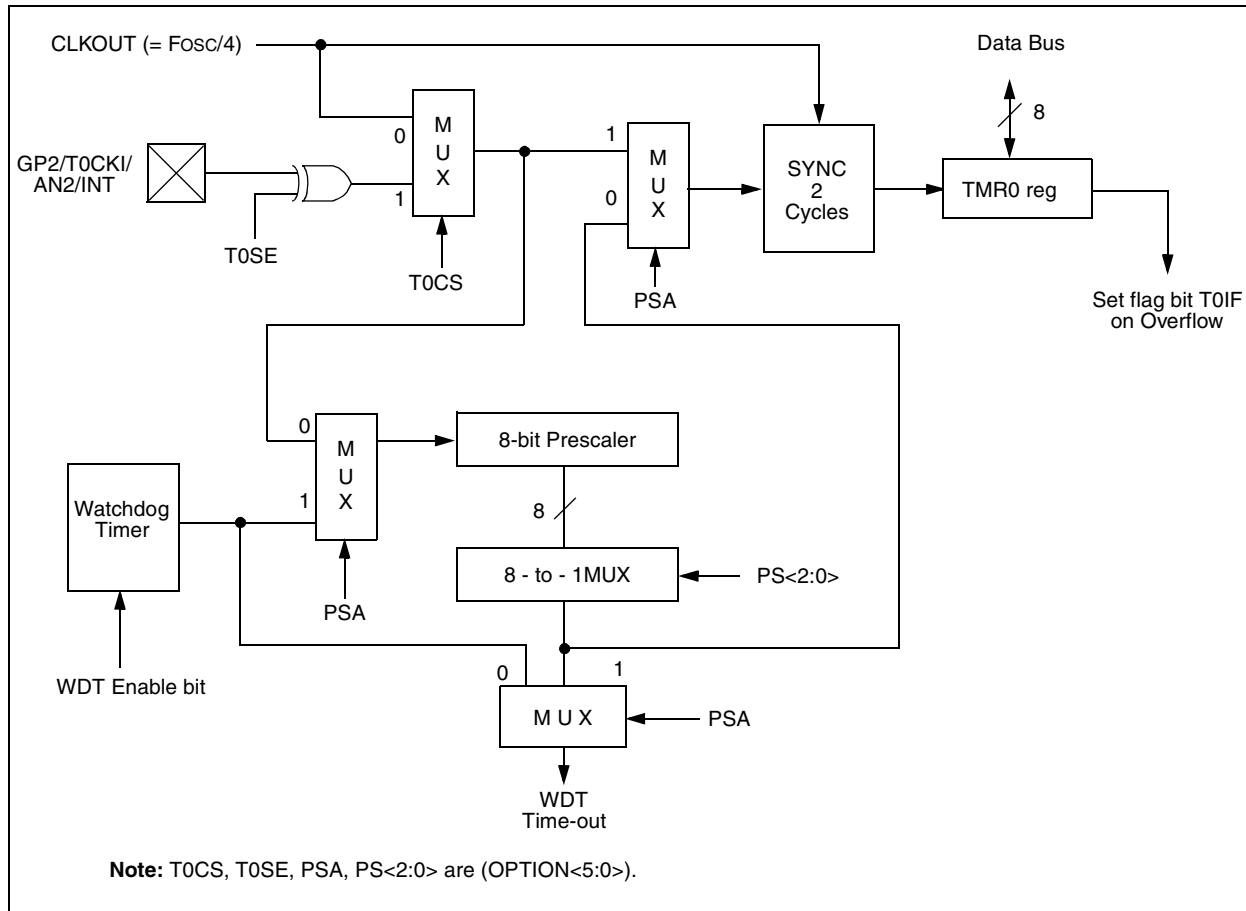
7.3 Prescaler

An 8-bit counter is available as a prescaler for the Timer0 module, or as a postscaler for the Watchdog Timer, respectively (Figure 7-6). For simplicity, this counter is being referred to as “prescaler” throughout this data sheet. Note that there is only one prescaler available which is mutually exclusively shared between the Timer0 module and the Watchdog Timer. Thus, a prescaler assignment for the Timer0 module means that there is no prescaler for the Watchdog Timer, and vice-versa.

The PSA and PS<2:0> bits (OPTION<3:0>) determine the prescaler assignment and prescale ratio.

When assigned to the Timer0 module, all instructions writing to the TMR0 register (i.e., `CLRF 1`, `MOVWF 1`, `BSF 1,x`, etc.) will clear the prescaler. When assigned to WDT, a `CLRWDT` instruction will clear the prescaler along with the Watchdog Timer. The prescaler is not readable or writable.

FIGURE 7-6: BLOCK DIAGRAM OF THE TIMER0/WDT PRESCALER



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8.4 A/D Conversions

Example 8-2 shows how to perform an A/D conversion. The GPIO pins are configured as analog inputs. The analog reference (VREF) is the device VDD. The A/D interrupt is enabled and the A/D conversion clock is FRC. The conversion is performed on the GP0 channel.

Note: The GO/DONE bit should **NOT** be set in the same instruction that turns on the A/D.

Clearing the GO/DONE bit during a conversion will abort the current conversion. The ADRES register will NOT be updated with the partially completed A/D conversion sample. That is, the ADRES register will continue to contain the value of the last completed conversion (or the last value written to the ADRES register). After the A/D conversion is aborted, a 2TAD wait is required before the next acquisition is started. After this 2TAD wait, an acquisition is automatically started on the selected channel.

EXAMPLE 8-2: DOING AN A/D CONVERSION

```
BSF    STATUS, RP0        ; Select Page 1
CLRFB  ADCON1             ; Configure A/D inputs
BSF    PIE1, ADIE        ; Enable A/D interrupts
BCF    STATUS, RP0        ; Select Page 0
MOVLW  0xC1               ; RC Clock, A/D is on, Channel 0 is selected
MOVWF  ADCON0             ;
BCF    PIR1, ADIF        ; Clear A/D interrupt flag bit
BSF    INTCON, PEIE      ; Enable peripheral interrupts
BSF    INTCON, GIE       ; Enable all interrupts
;
; Ensure that the required sampling time for the selected input channel has elapsed.
; Then the conversion may be started.
;
BSF    ADCON0, GO        ; Start A/D Conversion
:      ; The ADIF bit will be set and the GO/DONE bit
:      ; is cleared upon completion of the A/D Conversion.
```

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FIGURE 8-4: FLOWCHART OF A/D OPERATION

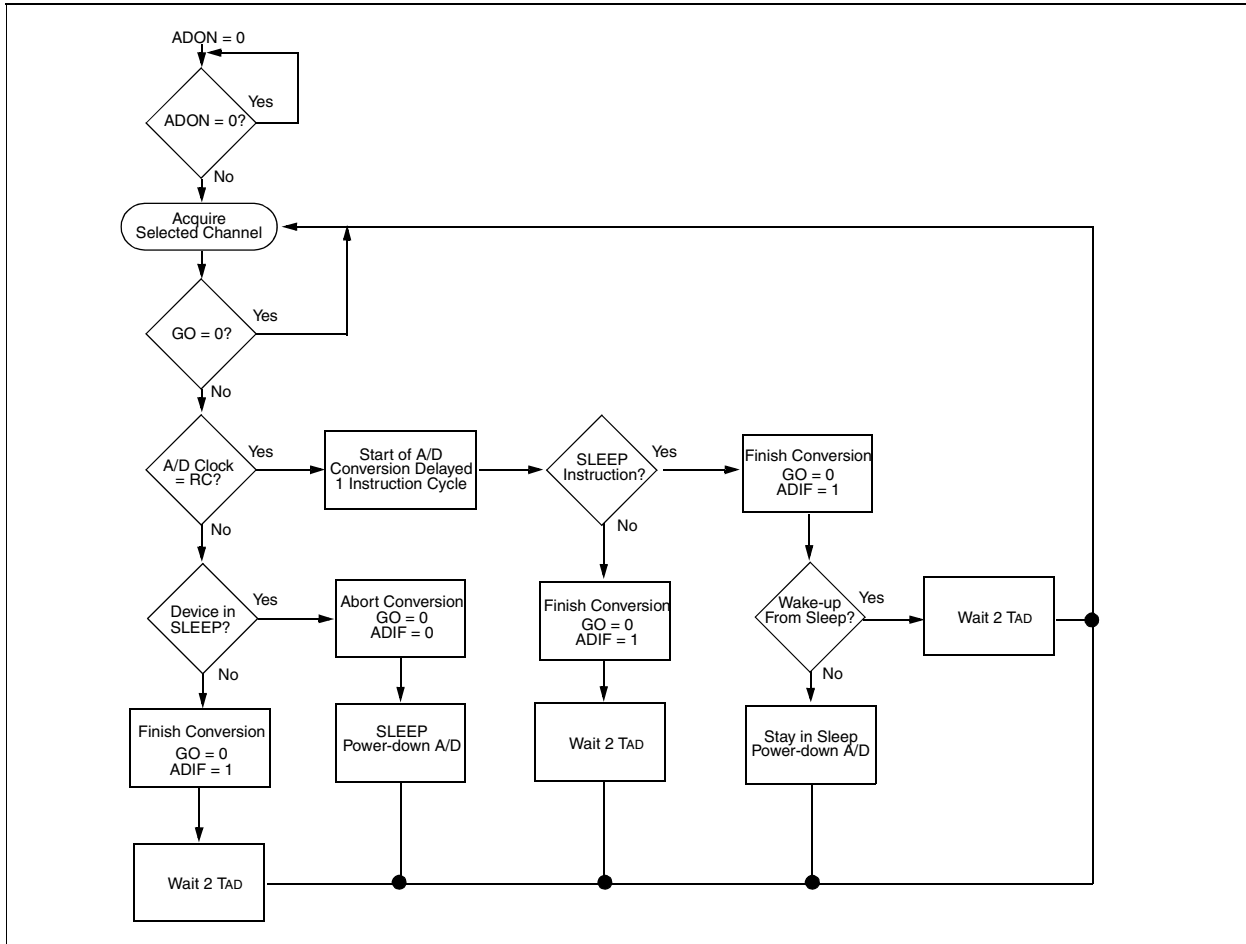


TABLE 8-2: SUMMARY OF A/D REGISTERS

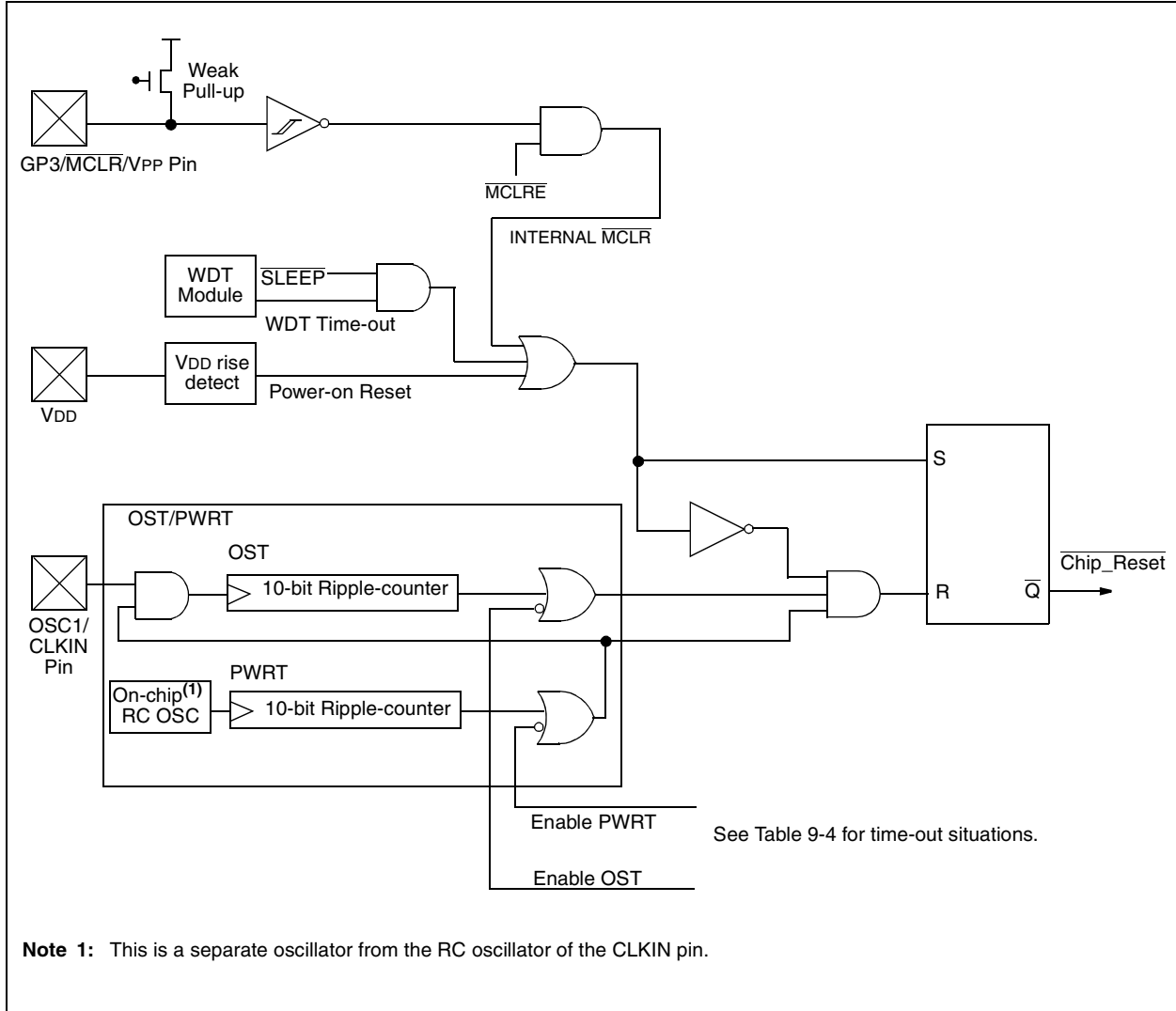
Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on Power-on Reset	Value on all other Resets
0Bh/8Bh	INTCON ⁽¹⁾	GIE	PEIE	TOIE	INTE	GPIE	TOIF	INTF	GPIF	0000 000x	0000 000u
0Ch	PIR1	—	ADIF	—	—	—	—	—	—	-0-- ----	-0-- ----
8Ch	PIE1	—	ADIE	—	—	—	—	—	—	-0-- ----	-0-- ----
1Eh	ADRES	A/D Result Register								xxxx xxxx	uuuu uuuu
1Fh	ADCON0	ADCS1	ADCS0	reserved	CHS1	CHS0	GO/DONE	reserved	ADON	0000 0000	0000 0000
9Fh	ADCON1	—	—	—	—	—	PCFG2	PCFG1	PCFG0	---- -000	---- -000
05h	GPIO	SCL ⁽²⁾	SDA ⁽²⁾	GP5	GP4	GP3	GP2	GP1	GP0	11xx xxxx	11uu uuuu
85h	TRIS	—	—	TRIS5	TRIS4	TRIS3	TRIS2	TRIS1	TRIS0	--11 1111	--11 1111

Legend: x = unknown, u = unchanged, - = unimplemented read as '0'. Shaded cells are not used for A/D conversion.

Note 1: These registers can be addressed from either bank.

2: The SCL (GP7) and SDA (GP6) bits are unimplemented on the PIC12C671/672 and read as '0'.

FIGURE 9-6: SIMPLIFIED BLOCK DIAGRAM OF ON-CHIP RESET CIRCUIT



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9.5 Interrupts

There are four sources of interrupt:

Interrupt Sources
TMR0 Overflow Interrupt
External Interrupt GP2/INT pin
GPIO Port Change Interrupts (pins GP0, GP1, GP3)
A/D Interrupt

The Interrupt Control Register (INTCON) records individual interrupt requests in flag bits. It also has individual and global interrupt enable bits.

Note: Individual interrupt flag bits are set, regardless of the status of their corresponding mask bit or the GIE bit.

A global interrupt enable bit, GIE (INTCON<7>), enables (if set) all un-masked interrupts or disables (if cleared) all interrupts. When bit GIE is enabled and an interrupt's flag bit and mask bit are set, the interrupt will vector immediately. Individual interrupts can be disabled through their corresponding enable bits in various registers. Individual interrupt flag bits are set, regardless of the status of their corresponding mask bit or the GIE bit. The GIE bit is cleared on reset.

The "return-from-interrupt" instruction, `RETFIE`, exits the interrupt routine, as well as sets the GIE bit, which re-enables interrupts.

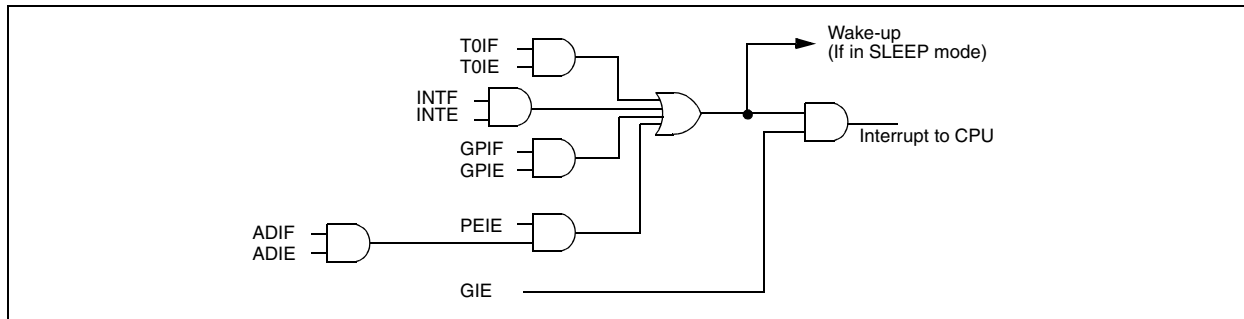
The GP2/INT, GPIO port change interrupt and the TMR0 overflow interrupt flags are contained in the INTCON register.

The peripheral interrupt flag ADIF, is contained in the Special Function Register PIR1. The corresponding interrupt enable bit is contained in Special Function Register PIE1, and the peripheral interrupt enable bit is contained in Special Function Register INTCON.

When an interrupt is responded to, the GIE bit is cleared to disable any further interrupt, the return address is pushed onto the stack and the PC is loaded with 0004h. Once in the interrupt service routine, the source(s) of the interrupt can be determined by polling the interrupt flag bits. The interrupt flag bit(s) must be cleared in software before re-enabling interrupts to avoid repeated interrupts.

For external interrupt events, such as GPIO change interrupt, the interrupt latency will be three or four instruction cycles. The exact latency depends on when the interrupt event occurs (Figure 9-14). The latency is the same for one or two cycle instructions. Individual interrupt flag bits are set, regardless of the status of their corresponding mask bit or the GIE bit.

FIGURE 9-13: INTERRUPT LOGIC



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9.5.1 TMR0 INTERRUPT

An overflow (FFh → 00h) in the TMR0 register will set flag bit T0IF (INTCON<2>). The interrupt can be enabled/disabled by setting/clearing enable bit T0IE (INTCON<5>) (Section 7.0). The flag bit T0IF (INTCON<2>) will be set, regardless of the state of the enable bits. If used, this flag must be cleared in software.

9.5.2 INT INTERRUPT

External interrupt on GP2/INT pin is edge triggered; either rising if bit INTEDG (OPTION<6>) is set, or falling, if the INTEDG bit is clear. When a valid edge appears on the GP2/INT pin, flag bit INTF (INTCON<1>) is set. This interrupt can be disabled by clearing enable bit INTE (INTCON<4>). Flag bit INTF must be cleared in software in the interrupt service routine before re-enabling this interrupt. The INT interrupt can wake-up the processor from SLEEP, if bit INTE was set prior to going into SLEEP. The status of global interrupt enable bit GIE decides whether or not the processor branches to the interrupt vector following wake-up. See Section 9.8 for details on SLEEP mode.

9.5.3 GPIO INTCON CHANGE

An input change on GP3, GP1 or GP0 sets flag bit GPIF (INTCON<0>). The interrupt can be enabled/disabled by setting/clearing enable bit GPIE (INTCON<3>) (Section 5.1). This flag bit GPIF (INTCON<0>) will be set, regardless of the state of the enable bits. If used, this flag must be cleared in software.

9.6 Context Saving During Interrupts

During an interrupt, only the return PC value is saved on the stack. Typically, users may wish to save key registers during an interrupt (i.e., W register and STATUS register). This will have to be implemented in software.

Example 9-1 shows the storing and restoring of the STATUS and W registers. The register, W_TEMP, must be defined in both banks and must be defined at the same offset from the bank base address (i.e., if W_TEMP is defined at 0x20 in bank 0, it must also be defined at 0xA0 in bank 1).

Example 9-2 shows the saving and restoring of STATUS and W using RAM locations 0x70 - 0x7F. W_TEMP is defined at 0x70 and STATUS_TEMP is defined at 0x71.

The example:

- Stores the W register.
- Stores the STATUS register in bank 0.
- Executes the ISR code.
- Restores the STATUS register (and bank select bit).
- Restores the W register.
- Returns from interrupt.

EXAMPLE 9-1: SAVING STATUS AND W REGISTERS USING GENERAL PURPOSE RAM (0x20 - 0x6F)

```
MOVWF    W_TEMP          ;Copy W to TEMP register, could be bank one or zero
SWAPF    STATUS,W        ;Swap status to be saved into W
BCF      STATUS,RP0      ;Change to bank zero, regardless of current bank
MOVWF    STATUS_TEMP     ;Save status to bank zero STATUS_TEMP register
:
:(ISR)
:
SWAPF    STATUS_TEMP,W   ;Swap STATUS_TEMP register into W
                        ;(sets bank to original state)
MOVWF    STATUS          ;Move W into STATUS register
SWAPF    W_TEMP,F        ;Swap W_TEMP
SWAPF    W_TEMP,W        ;Swap W_TEMP into W
RETFIE   ;Return from interrupt
```

EXAMPLE 9-2: SAVING STATUS AND W REGISTERS USING SHARED RAM (0x70 - 0x7F)

```
MOVWF    W_TEMP          ;Copy W to TEMP register (bank independent)
MOVWF    STATUS,W        ;Move STATUS register into W
MOVWF    STATUS_TEMP     ;Save contents of STATUS register
:
:(ISR)
:
MOVWF    STATUS_TEMP,W   ;Retrieve copy of STATUS register
MOVWF    STATUS          ;Restore pre-isr STATUS register contents
SWAPF    W_TEMP,F        ;
SWAPF    W_TEMP,W        ;Restore pre-isr W register contents
RETFIE   ;Return from interrupt
```

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BTFS **Bit Test f, Skip if Set**

Syntax: [*label*] BTFS *f*,*b*

Operands: $0 \leq f \leq 127$
 $0 \leq b < 7$

Operation: skip if (*f*<*b*>) = 1

Status Affected: None

Encoding:

01	11bb	bfff	ffff
----	------	------	------

Description: If bit '*b*' in register '*f*' is '1', then the next instruction is skipped. If bit '*b*' is '1', then the next instruction fetched during the current instruction execution, is discarded and a NOP is executed instead, making this a 2 cycle instruction.

Words: 1

Cycles: 1(2)

Example

```

HERE   BTFS   FLAG,1
FALSE  GOTO  PROCESS_CO
TRUE   •      DE
        •
        •
    
```

Before Instruction
 PC = address HERE

After Instruction
 if FLAG<1> = 0,
 PC = address FALSE
 if FLAG<1> = 1,
 PC = address TRUE

CALL **Call Subroutine**

Syntax: [*label*] CALL *k*

Operands: $0 \leq k \leq 2047$

Operation: (PC)+ 1 → TOS,
k → PC<10:0>,
 (PCLATH<4:3>) → PC<12:11>

Status Affected: None

Encoding:

10	0kkk	kkkk	kkkk
----	------	------	------

Description: Call Subroutine. First, return address (PC+1) is pushed onto the stack. The eleven bit immediate address is loaded into PC bits <10:0>. The upper bits of the PC are loaded from PCLATH. CALL is a two cycle instruction.

Words: 1

Cycles: 2

Example

```

HERE   CALL
        THER
        E
    
```

Before Instruction
 PC = Address HERE

After Instruction
 PC = Address THERE
 TOS = Address HERE+1

CLRF **Clear f**

Syntax: [*label*] CLRF *f*

Operands: $0 \leq f \leq 127$

Operation: 00h → (*f*)
 1 → Z

Status Affected: Z

Encoding:

00	0001	1fff	ffff
----	------	------	------

Description: The contents of register '*f*' are cleared and the Z bit is set.

Words: 1

Cycles: 1

Example

```

CLRF   FLAG_REG
    
```

Before Instruction
 FLAG_REG = 0x5A

After Instruction
 FLAG_REG = 0x00
 Z = 1

CLRW **Clear W**

Syntax: [*label*] CLRW

Operands: None

Operation: 00h → (W)
 1 → Z

Status Affected: Z

Encoding:

00	0001	0000	0011
----	------	------	------

Description: W register is cleared. Zero bit (Z) is set.

Words: 1

Cycles: 1

Example

```

CLRW
    
```

Before Instruction
 W = 0x5A

After Instruction
 W = 0x00
 Z = 1

RETURN Return from Subroutine

Syntax: [label] RETURN

Operands: None

Operation: TOS → PC

Status Affected: None

Encoding:

00	0000	0000	1000
----	------	------	------

Description: Return from subroutine. The stack is POPed and the top of the stack (TOS) is loaded into the program counter. This is a two cycle instruction.

Words: 1

Cycles: 2

Example

```

RETURN
After Interrupt
PC = TOS
    
```

RRF Rotate Right f through Carry

Syntax: [label] RRF f,d

Operands: $0 \leq f \leq 127$
 $d \in [0,1]$

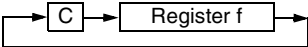
Operation: See description below

Status Affected: C

Encoding:

00	1100	dfff	ffff
----	------	------	------

Description: The contents of register 'f' are rotated one bit to the right through the Carry Flag. If 'd' is 0, the result is placed in the W register. If 'd' is 1, the result is placed back in register 'f'.



Words: 1

Cycles: 1

Example

```

RRF    REG1,
      0
    
```

Before Instruction

```

REG1 = 1110 0110
C = 0
    
```

After Instruction

```

REG1 = 1110 0110
W = 0111 0011
C = 0
    
```

RLF Rotate Left f through Carry

Syntax: [label] RLF f,d

Operands: $0 \leq f \leq 127$
 $d \in [0,1]$

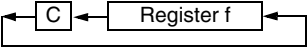
Operation: See description below

Status Affected: C

Encoding:

00	1101	dfff	ffff
----	------	------	------

Description: The contents of register 'f' are rotated one bit to the left through the Carry Flag. If 'd' is 0, the result is placed in the W register. If 'd' is 1, the result is stored back in register 'f'.



Words: 1

Cycles: 1

Example

```

RLF    REG1,0
    
```

Before Instruction

```

REG1 = 1110 0110
C = 0
    
```

After Instruction

```

REG1 = 1110 0110
W = 1100 1100
C = 1
    
```

SLEEP

Syntax: [label] SLEEP

Operands: None

Operation: 00h → WDT,
0 → WDT prescaler,
1 → \overline{TO} ,
0 → PD

Status Affected: \overline{TO} , PD

Encoding:

00	0000	0110	0011
----	------	------	------

Description: The power-down status bit, \overline{PD} is cleared. Time-out status bit, \overline{TO} is set. Watchdog Timer and its prescaler are cleared. The processor is put into SLEEP mode with the oscillator stopped.

Words: 1

Cycles: 1

Example: SLEEP

and test the sample code. In addition, PICDEM-17 supports down-loading of programs to and executing out of external FLASH memory on board. The PICDEM-17 is also usable with the MPLAB-ICE or PICMASTER emulator, and all of the sample programs can be run and modified using either emulator. Additionally, a generous prototype area is available for user hardware.

11.17 SEEVAL Evaluation and Programming System

The SEEVAL SEEPROM Designer's Kit supports all Microchip 2-wire and 3-wire Serial EEPROMs. The kit includes everything necessary to read, write, erase or program special features of any Microchip SEEPROM product including Smart Serials™ and secure serials. The Total Endurance™ Disk is included to aid in trade-off analysis and reliability calculations. The total kit can significantly reduce time-to-market and result in an optimized system.

11.18 KEELOQ Evaluation and Programming Tools

KEELOQ evaluation and programming tools support Microchips HCS Secure Data Products. The HCS evaluation kit includes an LCD display to show changing codes, a decoder to decode transmissions, and a programming interface to program test transmitters.

PIC12C67X

NOTES:

12.0 ELECTRICAL SPECIFICATIONS FOR PIC12C67X

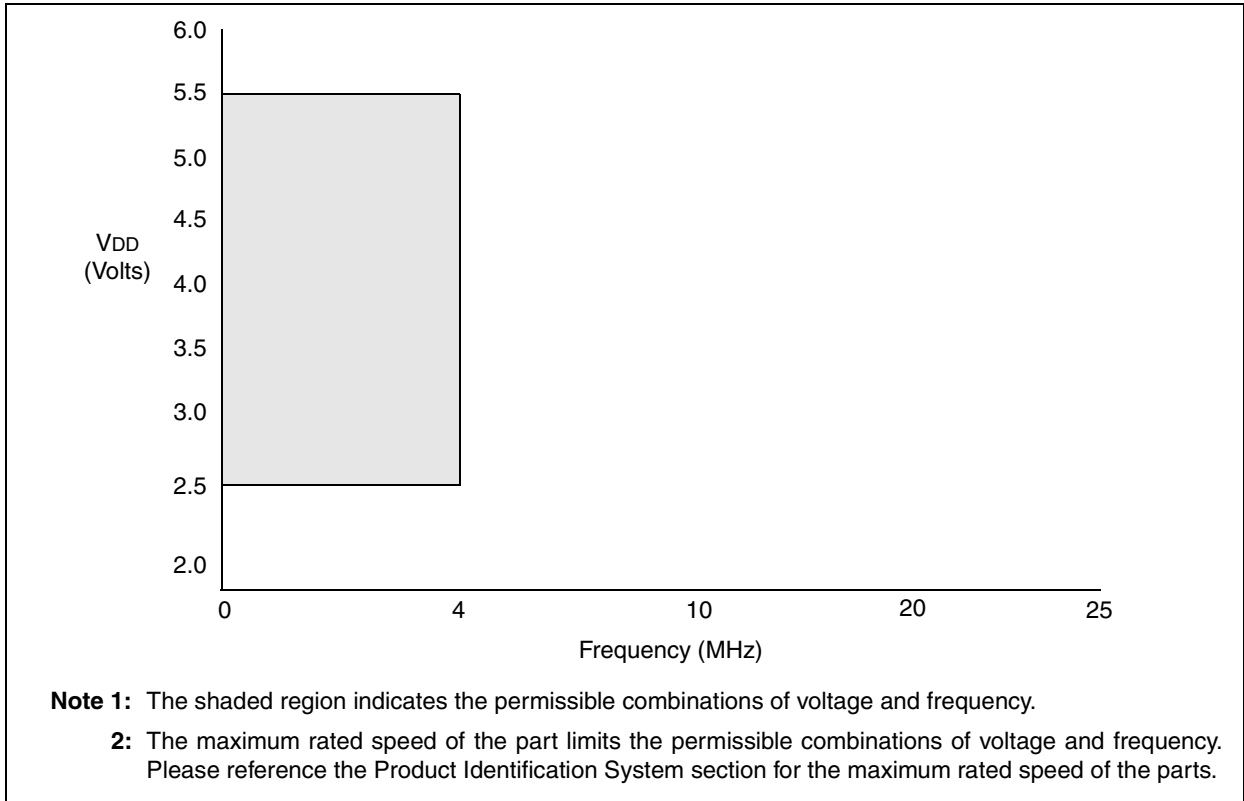
Absolute Maximum Ratings †

Ambient temperature under bias	-40° to +125°C
Storage temperature	-65°C to +150°C
Voltage on any pin with respect to VSS (except VDD and \overline{MCLR}).....	-0.3V to (VDD + 0.3V)
Voltage on VDD with respect to VSS	0 to +7.0V
Voltage on \overline{MCLR} with respect to VSS (Note 2).....	0 to +14V
Total power dissipation (Note 1).....	700 mW
Maximum current out of VSS pin	200 mA
Maximum current into VDD pin	150 mA
Input clamp current, I _{IK} (V _I < 0 or V _I > VDD).....	± 20 mA
Output clamp current, I _{OK} (V _O < 0 or V _O > VDD)	± 20 mA
Maximum output current sunk by any I/O pin.....	25 mA
Maximum output current sourced by any I/O pin	25 mA
Maximum current sunk by GPIO pins combined	100 mA
Maximum current sourced by GPIO pins combined.....	100 mA

Note 1: Power dissipation is calculated as follows: $P_{dis} = V_{DD} \times \{I_{DD} - \sum I_{OH}\} + \sum \{(V_{DD} - V_{OH}) \times I_{OH}\} + \sum (V_{OL} \times I_{OL})$.

† NOTICE: Stresses above those listed under “Absolute Maximum Ratings” may cause permanent damage to the device. This is a stress rating only and functional operation of the device at those or any other conditions above those indicated in the operation listings of this specification is not implied. Exposure to maximum rating conditions for extended periods may affect device reliability.

FIGURE 12-3: PIC12LC67X VOLTAGE-FREQUENCY GRAPH, $-40^{\circ}\text{C} \leq T_A \leq +85^{\circ}\text{C}$



PIC12C67X

Standard Operating Conditions (unless otherwise specified)							
DC CHARACTERISTICS		Operating temperature					
		0°C ≤ TA ≤ +70°C (commercial)					
		-40°C ≤ TA ≤ +85°C (industrial)					
						-40°C ≤ TA ≤ +125°C (extended)	
Operating voltage VDD range as described in DC spec Section 12.1 and Section 12.2.							
Param No.	Characteristic	Sym	Min	Typ†	Max	Units	Conditions
D090	Output High Voltage I/O ports (Note 3)	VOH	VDD - 0.7	—	—	V	IOH = -3.0 mA, VDD = 4.5V, -40°C to +85°C
D090A			VDD - 0.7	—	—	V	
D092	OSC2/CLKOUT		VDD - 0.7	—	—	V	IOH = 1.3 mA, VDD = 4.5V, -40°C to +85°C
D092A			VDD - 0.7	—	—	V	
Capacitive Loading Specs on Output Pins							
D100	OSC2 pin	Cosc2	—	—	15	pF	In XT and LP modes when external clock is used to drive OSC1.
D101	All I/O pins	CIO	—	—	50	pF	

† Data in "Typ" column is at 5V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

- Note 1:** In EXTRC oscillator configuration, the OSC1/CLKIN pin is a Schmitt Trigger input. It is not recommended that the PIC12C67X be driven with external clock in RC mode.
- 2:** The leakage current on the MCLR pin is strongly dependent on the applied voltage level. The specified levels represent normal operating conditions. Higher leakage current may be measured at different input voltages.
- 3:** Negative current is defined as coming out of the pin.
- 4:** Does not include GP3. For GP3 see parameters D061 and D061A.
- 5:** This spec. applies to GP3/MCLR configured as external MCLR and GP3/MCLR configured as input with internal pull-up enabled.
- 6:** This spec. applies when GP3/MCLR is configured as an input with pull-up disabled. The leakage current of the MCLR circuit is higher than the standard I/O logic.

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FIGURE 12-8: TIMER0 CLOCK TIMINGS

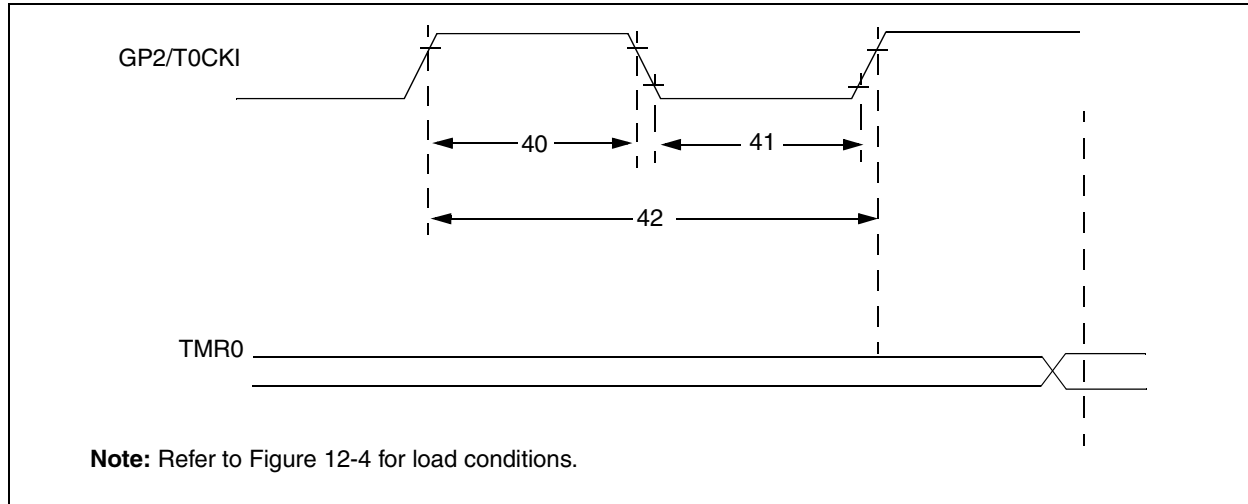


TABLE 12-5: TIMER0 AND TIMER1 EXTERNAL CLOCK REQUIREMENTS

Param No.	Sym	Characteristic	Min	Typ†	Max	Units	Conditions	
40*	Tt0H	T0CKI High Pulse Width	No Prescaler	$0.5T_{CY} + 20$	—	—	ns	Must also meet parameter 42
			With Prescaler	10	—	—	ns	
41*	Tt0L	T0CKI Low Pulse Width	No Prescaler	$0.5T_{CY} + 20$	—	—	ns	Must also meet parameter 42
			With Prescaler	10	—	—	ns	
42*	Tt0P	T0CKI Period	No Prescaler	$T_{CY} + 40$	—	—	ns	N = prescale value (2, 4, ..., 256)
			With Prescaler	Greater of: 20 or $\frac{T_{CY} + 40}{N}$	—	—	ns	
48	TCKE2tmr1	Delay from external clock edge to timer increment	$2T_{osc}$	—	$7T_{osc}$	—		

* These parameters are characterized but not tested.

† Data in "Typ" column is at 5V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

TABLE 12-6: GPIO PULL-UP RESISTOR RANGES

VDD (Volts)	Temperature (°C)	Min	Typ	Max	Units
GP0/GP1					
2.5	-40	38K	42K	63K	Ω
	25	42K	48K	63K	Ω
	85	42K	49K	63K	Ω
	125	50K	55K	63K	Ω
5.5	-40	15K	17K	20K	Ω
	25	18K	20K	23K	Ω
	85	19K	22K	25K	Ω
	125	22K	24K	28K	Ω
GP3					
2.5	-40	285K	346K	417K	Ω
	25	343K	414K	532K	Ω
	85	368K	457K	532K	Ω
	125	431K	504K	593K	Ω
5.5	-40	247K	292K	360K	Ω
	25	288K	341K	437K	Ω
	85	306K	371K	448K	Ω
	125	351K	407K	500K	Ω

* These parameters are characterized but not tested.

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