Zilog - Z8S18020FSG1960 Datasheet

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Applications of Embedded - Microprocessors

Embedded microprocessors are utilized across a broad spectrum of applications, making them indispensable in

Details	
Product Status	Active
Core Processor	Z8S180
Number of Cores/Bus Width	1 Core, 8-Bit
Speed	20MHz
Co-Processors/DSP	-
RAM Controllers	DRAM
Graphics Acceleration	No
Display & Interface Controllers	-
Ethernet	-
SATA	-
USB	-
Voltage - I/O	5.0V
Operating Temperature	0°C ~ 70°C (TA)
Security Features	-
Package / Case	80-BQFP
Supplier Device Package	80-QFP
Purchase URL	https://www.e-xfl.com/product-detail/zilog/z8s18020fsg1960

GENERAL DESCRIPTION (Continued)

Power connections follow the conventional descriptions below:

Connection	Circuit	Device
Power	V_{CC}	V_{DD}
Ground	GND	V_{SS}

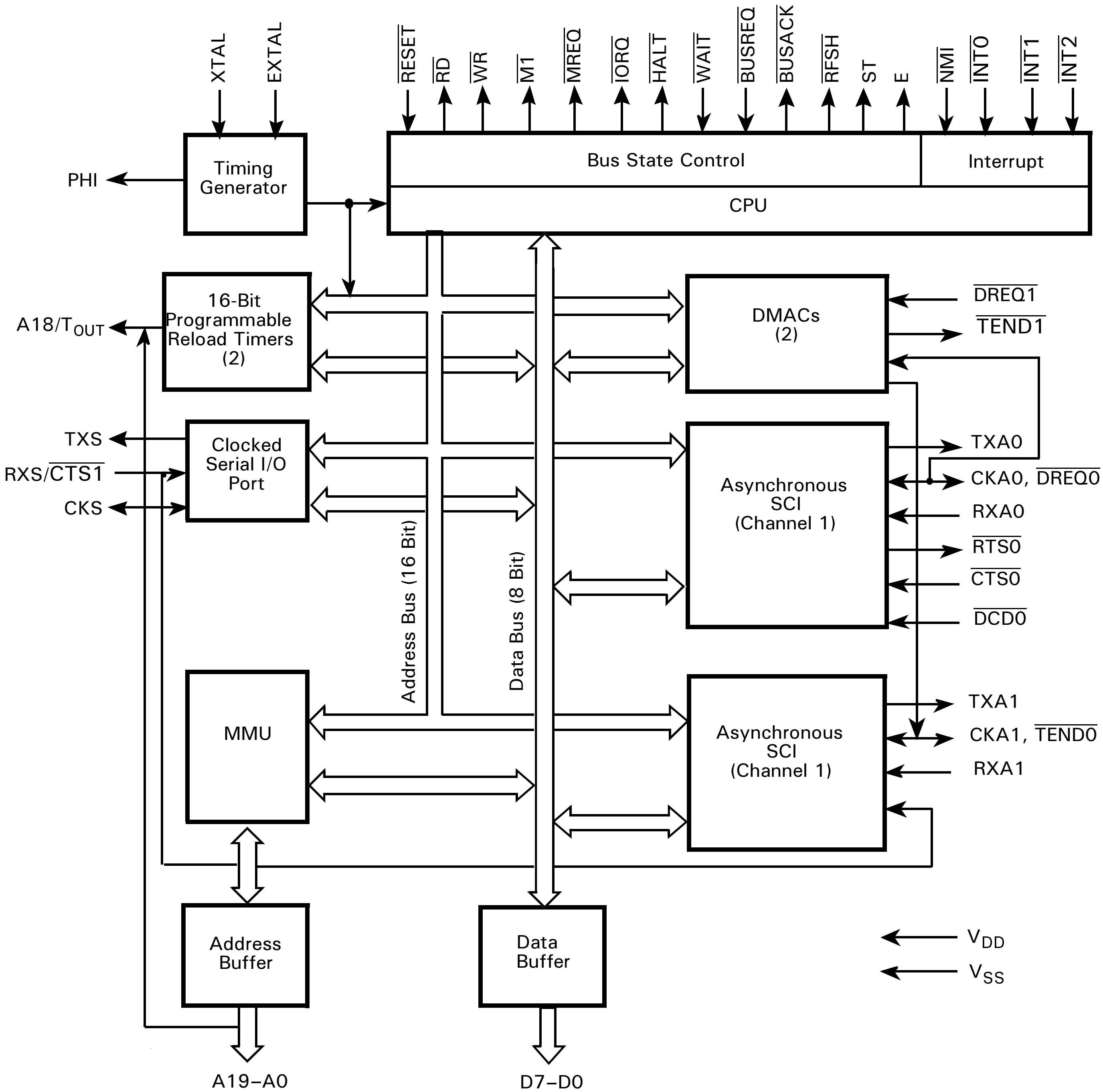

Figure 1. Z8S180/Z8L180 Functional Block Diagram

PIN IDENTIFICATION (Continued)

Table 2. Pin Status During RESET, BUSACK, and SLEEP Modes

Pin Number and Package Type					Pin Status		
QFP	PLCC	DIP	Default Function	Secondary Function	RESET	BUSACK	SLEEP
1	9	8	$\overline{NMI}$		IN	IN	IN
2			NC				
3			NC				
4	10	9	$\overline{INT0}$		IN	IN	IN
5	11	10	$\overline{INT1}$		IN	IN	IN
6	12	11	$\overline{INT2}$		IN	IN	IN
7	13	12	ST		High	High	High
8	14	13	A0		3T	3T	High
9	15	14	A1		3T	3T	High
10	16	15	A2		3T	3T	High
11	17	16	A3		3T	3T	High
12	18		V_{SS}		V_{SS}	V_{SS}	V_{SS}
13	19	17	A4		3T	3T	High
14			NC				
15	20	18	A5		3T	3T	High
16	21	19	A6		3T	3T	High
17	22	20	A7		3T	3T	High
18	23	21	A8		3T	3T	High
19	24	22	A9		3T	3T	High
20	25	23	A10		3T	3T	High
21	26	24	A11		3T	3T	High
22			NC				
23			NC				
24	27	25	A12		3T	3T	High
25	28	26	A13		3T	3T	High
26	29	27	A14		3T	3T	High
27	30	28	A15		3T	3T	High
28	31	29	A16		3T	3T	High
29	32	30	A17		3T	3T	High
30			NC				
31	33	31	A18		3T	3T	High
			T_{OUT}		N/A	OUT	OUT
32	34	32	V_{DD}		V_{DD}	V_{DD}	V_{DD}
33	35		A19		3T	3T	High
34	36	33	V_{SS}		V_{SS}	V_{SS}	V_{SS}
35	37	34	D0		3T	3T	3T
36	38	35	D1		3T	3T	3T
37	39	36	D2		3T	3T	3T
38	40	37	D3		3T	3T	3T

PIN IDENTIFICATION (Continued)

Table 2. Pin Status During RESET, BUSACK, and SLEEP Modes (Continued)

Pin Number and Package Type					Pin Status		
QFP	PLCC	DIP	Default Function	Secondary Function	RESET	BUSACK	SLEEP
76	4	3	EXTAL		IN	IN	IN
77	5	4	$\overline{\text{WAIT}}$		IN	IN	IN
78	6	5	$\overline{\text{BUSACK}}$		High	OUT	OUT
79	7	6	$\overline{\text{BUSREQ}}$		IN	IN	IN
80	8	7	$\overline{\text{RESET}}$		IN	IN	IN

PIN DESCRIPTIONS

A0–A19. Address Bus (Output, 3-state). A0–A19 form a 20-bit address bus. The Address Bus provides the address for memory data bus exchanges (up to 1 MB) and I/O data bus exchanges (up to 64 KB). The address bus enters a high–impedance state during reset and external bus acknowledge cycles. Address line A18 is multiplexed with the output of PRT channel 1 (T_{OUT}, selected as address output on reset), and address line A19 is not available in DIP versions of the Z8S180.

$\overline{\text{BUSACK}}$. Bus Acknowledge (Output, active Low). $\overline{\text{BUSACK}}$ indicates that the requesting device, the MPU address and data bus, and some control signals enter their high-impedance state.

BUSREQ. Bus Request (Input, active Low). This input is used by external devices (such as DMA controllers) to request access to the system bus. This request demands a higher priority than $\overline{\text{NMI}}$ and is always recognized at the end of the current machine cycle. This signal stops the CPU from executing further instructions, places addresses, data buses, and other control signals into the high-impedance state.

CKA0, CKA1. Asynchronous Clock 0 and 1 (bidirectional). When in output mode, these pins are the transmit and receive clock outputs from the ASCI baud rate generators. When in input mode, these pins serve as the external clock inputs for the ASCI baud rate generators. CKA0 is multiplexed with $\overline{\text{DREQ0}}$, and CKA1 is multiplexed with $\overline{\text{TEND0}}$.

CKS. Serial Clock (bidirectional). This line is the clock for the CSI/O channel.

$\overline{\text{CTS0}}$–$\overline{\text{CTS1}}$. Clear to send 0 and 1 (Inputs, active Low). These lines are modem control signals for the ASCI channels. $\overline{\text{CTS1}}$ is multiplexed with RXS.

D0–D7. Data Bus = (bidirectional, 3-state). D0–D7 constitute an 8-bit bidirectional data bus, used for the transfer of information to and from I/O and memory devices. The data bus enters the high-impedance state during reset and external bus acknowledge cycles.

$\overline{\text{DCD0}}$. Data Carrier Detect 0 (Input, active Low); a programmable modem control signal for ASCI channel 0.

$\overline{\text{DREQ0}}$, $\overline{\text{DREQ1}}$. DMA Request 0 and 1 (Input, active Low). $\overline{\text{DREQ}}$ is used to request a DMA transfer from one of the on-chip DMA channels. The DMA channels monitor these inputs to determine when an external device is ready for a READ or WRITE operation. These inputs can be programmed to be either level or edge sensed. $\overline{\text{DREQ0}}$ is multiplexed with CKA0.

E. Enable Clock (Output). This pin functions as a synchronous, machine-cycle clock output during bus transactions.

EXTAL. External Clock Crystal (Input). Crystal oscillator connections. An external clock can be input to the Z8S180/Z8L180 on this pin when a crystal is not used. This input is Schmitt triggered.

$\overline{\text{HALT}}$. HALT/SLEEP (Output, active Low). This output is asserted after the CPU executes either the HALT or SLEEP instruction and is waiting for either a nonmaskable or a maskable interrupt before operation can resume. It is also used with the $\overline{\text{M1}}$ and ST signals to decode the status of the CPU machine cycle.

$\overline{\text{INT0}}$. Maskable Interrupt Request 0 (Input, active Low). This signal is generated by external I/O devices. The CPU honors these requests at the end of the current instruction cycle as long as the $\overline{\text{NMI}}$ and $\overline{\text{BUSREQ}}$ signals are inactive. The CPU acknowledges this interrupt request with an interrupt acknowledge cycle. During this cycle, both the $\overline{\text{M1}}$ and $\overline{\text{IORQ}}$ signals become active.

$\overline{\text{INT1}}$, $\overline{\text{INT2}}$. Maskable Interrupt Request 1 and 2 (Inputs, active Low). This signal is generated by external I/O devices. The CPU honors these requests at the end of the current instruction cycle as long as the $\overline{\text{NMI}}$, $\overline{\text{BUSREQ}}$, and $\overline{\text{INT0}}$ signals are inactive. The CPU acknowledges these requests with an interrupt acknowledge cycle. Unlike the acknowledgment for $\overline{\text{INT0}}$, neither the $\overline{\text{M1}}$ or $\overline{\text{IORQ}}$ signals become active during this cycle.

$\overline{\text{IORQ}}$. I/O Request (Output, active Low, 3-state). $\overline{\text{IORQ}}$ indicates that the address bus contains a valid I/O address for an I/O READ or I/O WRITE operation. $\overline{\text{IORQ}}$ is also generated, along with $\overline{\text{M1}}$, during the acknowledgment of the $\overline{\text{INT0}}$ input signal to indicate that an interrupt response vector can be place onto the data bus. This signal is analogous to the $\overline{\text{IOE}}$ signal of the Z64180.

$\overline{\text{M1}}$. Machine Cycle 1 (Output, active Low). Together with $\overline{\text{MREQ}}$, $\overline{\text{M1}}$ indicates that the current cycle is the opcode-fetch cycle of instruction execution. Together with $\overline{\text{IORQ}}$, $\overline{\text{M1}}$ indicates that the current cycle is for interrupt acknowledgment. It is also used with the $\overline{\text{HALT}}$ and ST signal to decode the status of the CPU machine cycle. This signal is analogous to the $\overline{\text{LIR}}$ signal of the Z64180.

$\overline{\text{MREQ}}$. Memory Request (Output, active Low, 3-state). $\overline{\text{MREQ}}$ indicates that the address bus holds a valid address for a memory READ or memory WRITE operation. This signal is analogous to the $\overline{\text{ME}}$ signal of Z64180.

$\overline{\text{NMI}}$. Nonmaskable Interrupt (Input, negative edge triggered). $\overline{\text{NMI}}$ demands a higher priority than $\overline{\text{INT}}$ and is al-

PIN DESCRIPTIONS (Continued)

ways recognized at the end of an instruction, regardless of the state of the interrupt-enable flip-flops. This signal forces CPU execution to continue at location 0066H.

PHI. System Clock (Output). The output is used as a reference clock for the MPU and the external system. The frequency of this output may be one-half, equal to, or twice the crystal or input clock frequency.

$\overline{\text{RD}}$. Read (Output, active Low, 3-state). $\overline{\text{RD}}$ indicates that the CPU wants to read data from either memory or an I/O device. The addressed I/O or memory device should use this signal to gate data onto the CPU data bus.

$\overline{\text{RFSH}}$. Refresh (Output, active Low). Together with $\overline{\text{MREQ}}$, $\overline{\text{RFSH}}$ indicates that the current CPU machine cycle and the contents of the address bus should be used for refresh of dynamic memories. The low-order 8 bits of the address bus (A7–A0) contain the refresh address. *This signal is analogous to the $\overline{REF}$ signal of the Z64180.*

$\overline{\text{RTS0}}$. Request to Send 0 (Output, active Low); a programmable MODEM control signal for ASCI channel 0.

RXA0, RXA1. Receive Data 0 and 1 (Input). These signals are the receive data for the ASCI channels.

RXS. Clocked Serial Receive Data (Input). This line is the receive data for the CSI/O channel. RXS is multiplexed with the $\overline{\text{CTS1}}$ signal for ASCI channel 1.

ST. Status (Output). This signal is used with the $\overline{\text{M1}}$ and $\overline{\text{HALT}}$ output to decode the status of the CPU machine cycle. See Table 3.

Table 3. Status Summary

ST	$\overline{\text{HALT}}$	$\overline{\text{M1}}$	Operation
0	1	0	CPU Operation (1st Opcode Fetch)
1	1	0	CPU Operation (2nd Opcode and 3rd Opcode Fetch)
1	1	1	CPU Operation (MC Except Opcode Fetch)
0	X	1	DMA Operation
0	0	0	HALT Mode
1	0	1	SLEEP Mode (Including SYSTEM STOP Mode)

Notes:
X = Do not care.
MC = Machine Cycle.

$\overline{\text{TEND0}}$, $\overline{\text{TEND1}}$. Transfer End 0 and 1 (Outputs, active Low). This output is asserted active during the most recent WRITE cycle of a DMA operation. It is used to indicate the end of the block transfer. $\overline{\text{TEND0}}$ is multiplexed with CKA1.

TEST. Test (Output, not in DIP version). This pin is for test and should be left open.

T_{OUT}. Timer Out (Output). T_{OUT} is the output from PRT channel 1. This line is multiplexed with A18 of the address bus.

TXA0, TXA1. Transmit Data 0 and 1 (Outputs). These signals are the transmitted data from the ASCI channels. Transmitted data changes are with respect to the falling edge of the transmit clock.

TXS. Clocked Serial Transmit Data (Output). This line is the transmitted data from the CSI/O channel.

$\overline{\text{WAIT}}$. Wait (Input, active Low). $\overline{\text{WAIT}}$ indicates to the MPU that the addressed memory or I/O devices are not ready for data transfer. This input is sampled on the falling edge of T2 (and subsequent WAIT states). If the input is sampled Low, then the additional WAIT states are inserted until the $\overline{\text{WAIT}}$ input is sampled High, at which time execution continues.

$\overline{\text{WR}}$. WRITE (Output, active Low, 3-state). $\overline{\text{WR}}$ indicates that the CPU data bus holds valid data to be stored at the addressed I/O or memory location.

XTAL. Crystal Oscillator Connection (Input). This pin should be left open if an external clock is used instead of a crystal. The oscillator input is not a TTL level (see DC Characteristics).

Several pins are used for different conditions, depending on the circumstance.

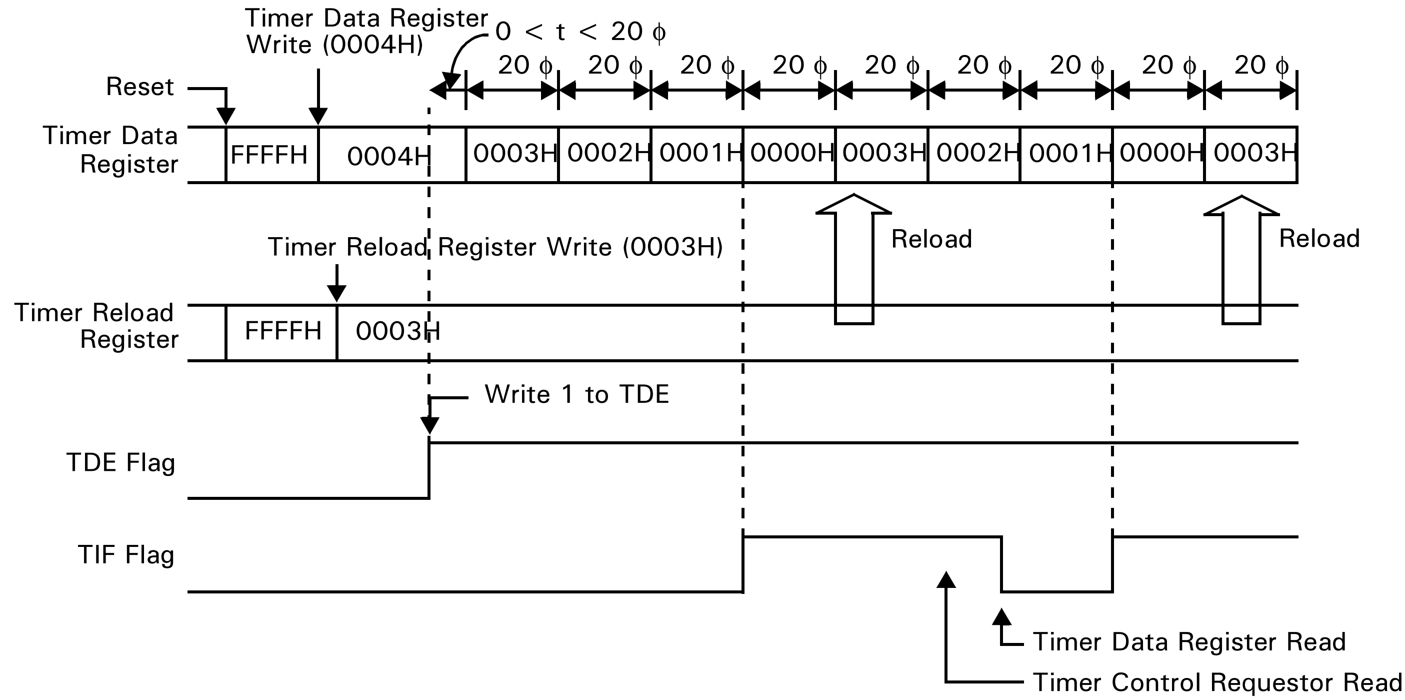

Figure 5. Timer Initialization, Count Down, and Reload Timing

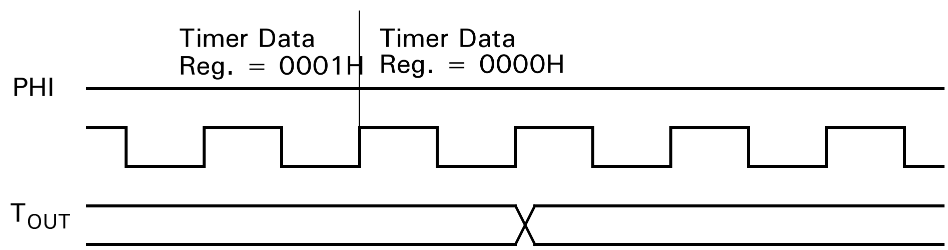

Figure 6. Timer Output Timing

Clocked Serial I/O (CSI/O). The CSI/O channel provides a half-duplex serial transmitter and receiver. This channel can be used for simple high-speed data connection to another microprocessor or microcomputer. TRDR is used for both CSI/O transmission and reception. Thus, the system design must ensure that the constraints of half-duplex operation are met (Transmit and Receive operation cannot occur simultaneously). For example, if a CSI/O transmission is attempted while the CSI/O is receiving data, a CSI/O does not work.

Note: TRDR is not buffered. Performing a CSI/O transmit while the previous transmission is still in progress causes the data to be immediately updated and corrupts the transmit operation. Similarly, reading TRDR while a transmit or receive is in progress should be avoided.

OPERATION MODES (Continued)

Table 5. RETI Control Signal States

Machine Cycle	States	Address	Data	$\overline{RD}$	$\overline{WR}$	$\overline{MREQ}$	$\overline{IORQ}$	$\overline{M1}$ M1E = 1	$\overline{M1}$ M1E = 0	$\overline{HALT}$	ST
1	T1–T3	1st Opcode	EDH	0	1	0	1	0	1	1	0
2	T1–T3	2nd Opcode	4DH	0	1	0	1	0	1	1	0
	Ti	NA	3-state	1	1	1	1	1	1	1	1
	Ti	NA	3-state	1	1	1	1	1	1	1	1
	Ti	NA	3-state	1	1	1	1	1	1	1	1
3	T1–T3	1st Opcode	EDH	0	1	0	1	0	0	1	1
	Ti	NA	3-state	1	1	1	1	1	1	1	1
4	T1–T3	2nd Opcode	4DH	0	1	0	1	0	1	1	1
5	T1–T3	SP	Data	0	1	0	1	1	1	1	1
6	T1–T3	SP + 1	Data	0	1	0	1	1	1	1	1

$\overline{M1TE}$ ($\overline{M1}$ Temporary Enable). This bit controls the temporary assertion of the $\overline{M1}$ signal. It is always read back as a 1 and is set to 1 during RESET.

When M1E is set to 0 to accommodate certain external Z80 peripheral(s), those same device(s) may require a pulse on $\overline{M1}$ after programming certain of their registers to complete the function being programmed.

For example, when a control word is written to the Z80 PIO to enable interrupts, no enable actually takes place until the PIO sees an active $\overline{M1}$ signal. When $\overline{M1TE}$ = 1, there is no change in the operation of the $\overline{M1}$ signal, and M1E controls its function. When $\overline{M1TE}$ = 0, the $\overline{M1}$ output is asserted during the next opcode fetch cycle regardless of the state programmed into the M1E bit. This condition is only momentary (one time) and it is not necessary to preprogram a 1 to disable the function (see Figure 10).

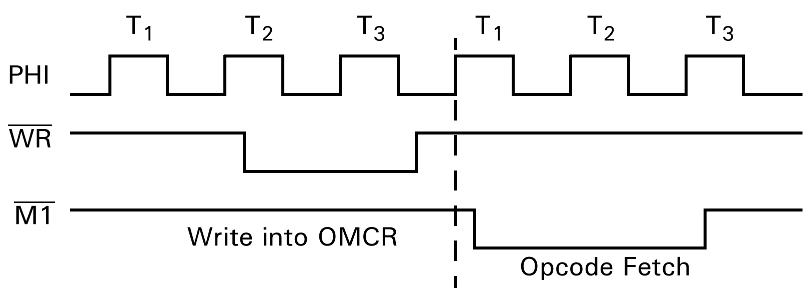

Figure 10. M1 Temporary Enable Timing

IOC (I/O Compatibility). This bit controls the timing of the $\overline{IORQ}$ and $\overline{RD}$ signals. The bit is set to 1 by RESET.

When $\overline{IOC}$ = 1, the $\overline{IORQ}$ and $\overline{RD}$ signals function the same as the Z64180 (Figure 11).

OPERATION MODES (Continued)

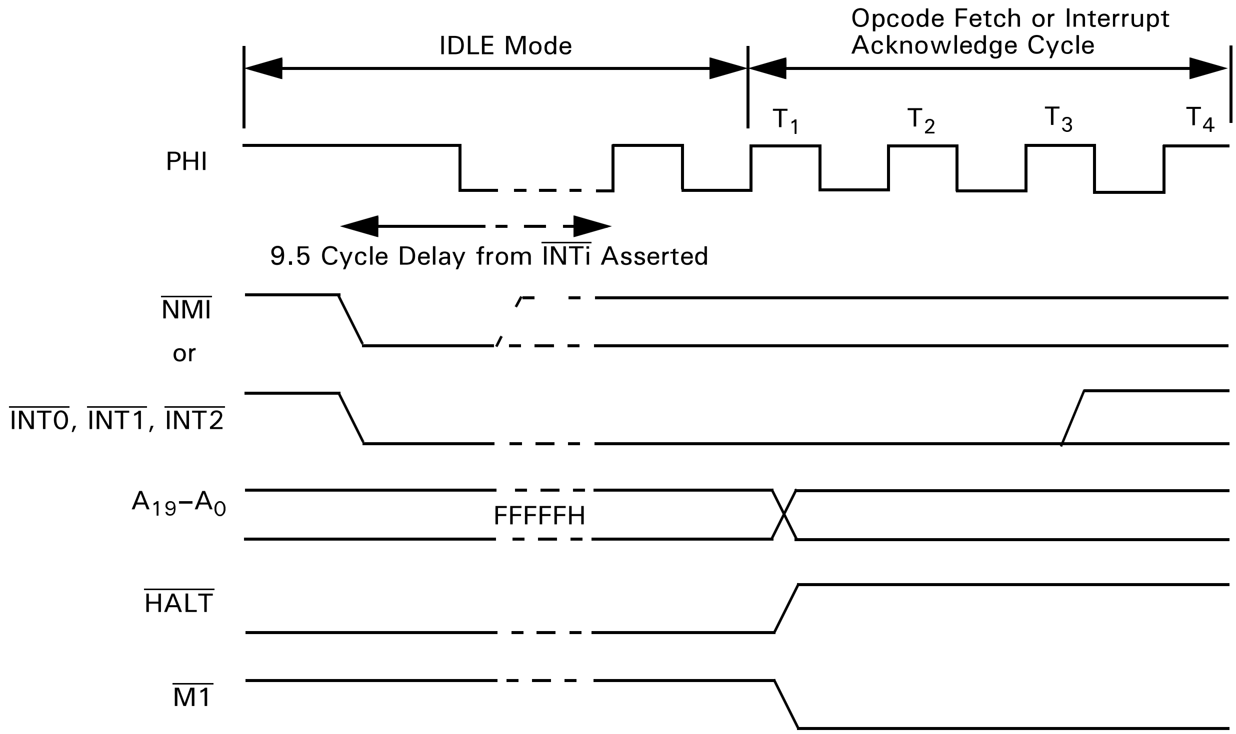

Figure 15. Z8S180/Z8L180 IDLE Mode Exit Due To External Interrupt

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While the Z8S180/Z8L180 is in IDLE mode, it grants the bus to an external Master if the BREXT bit (CCR5) is 1. Figure 16 depicts the timing for this sequence.

Note: A response to a bus request takes 8 clock cycles longer than in normal operation.

After the external Master negates the Bus Request, the Z8S180/Z8L180 disables the PHI clock and remains in IDLE mode.

ing the bus to an external Master during STANDBY mode, when the BREXT bit in the CPU Control Register (CCR5) is 1.

As described previously for SLEEP and IDLE modes, when the MPU leaves STANDBY mode due to $\overline{\text{NMI}}$ Low or an enabled $\overline{\text{INT0}}$–$\overline{\text{INT2}}$ Low when the IEF, flag is 1 due to an IE instruction, it starts by performing the interrupt with the return address being that of the instruction following the SLP instruction. If the Z8S180/Z8L180 leaves STANDBY mode due to an external interrupt request that's enabled in the INT/TRAP Control Register, but the IEF, bit is 0 due to a DI instruction, the processor restarts by executing the instruction(s) following the SLP instruction. If $\overline{\text{INT0}}$, or $\overline{\text{INT1}}$ or $\overline{\text{INT2}}$ goes inactive before the end of the clock stabilization delay, the Z8S180/Z8L180 stays in STANDBY mode.

Figure 17 indicates the timing for leaving STANDBY mode due to an interrupt request.

Note: The Z8S180/Z8L180 takes either 64 or 2^{17} (131,072) clocks to restart, depending on the CCR3 bit.

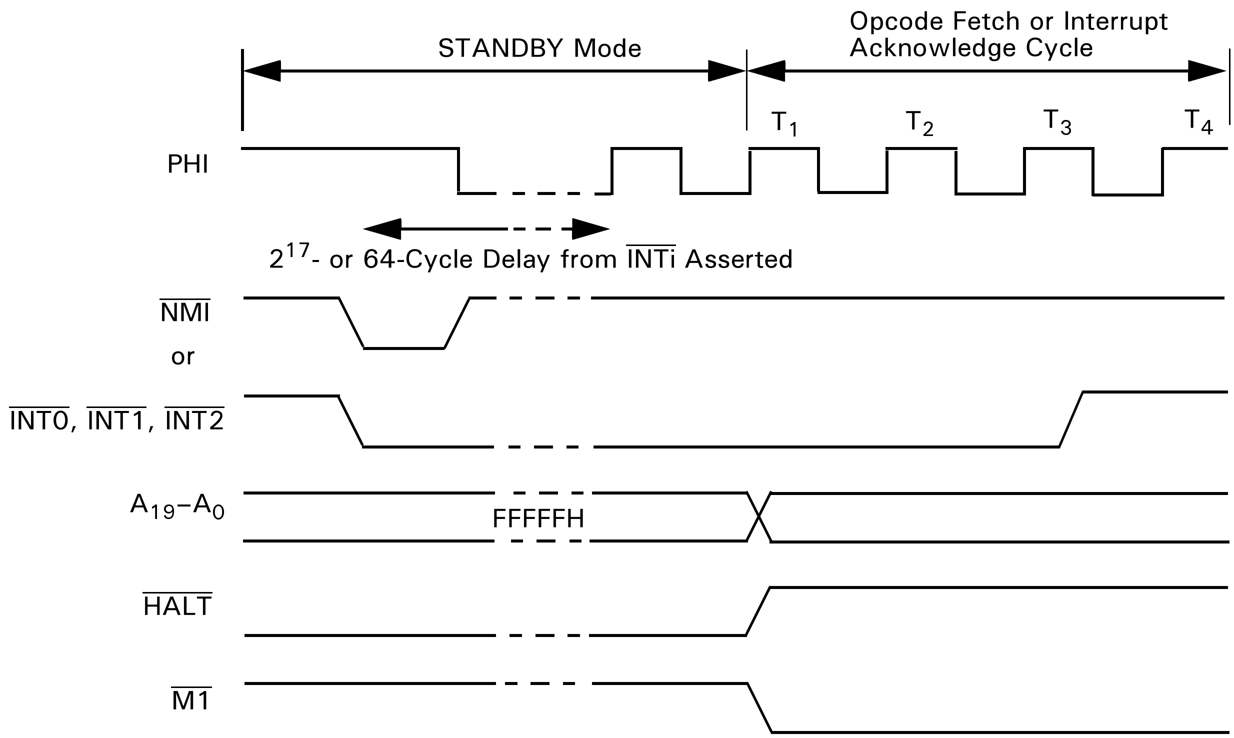

Figure 17. Z8S180/Z8L180 STANDBY Mode Exit Due to External Interrupt

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While the Z8S180/Z8L180 is in STANDBY mode, it grants the bus to an external Master if the BREXT bit (CCR5) is 1. Figure 18 indicates the timing of this sequence. The device takes 64 or 2^{17} (131,072) clock cycles to grant the bus depending on the CCR3 bit. The latter (not the QUICK RECOVERY) case may be prohibitive for many demand-driven external Masters. If so, QUICK RECOVERY or IDLE mode can be used.

ASCI REGISTER DESCRIPTION

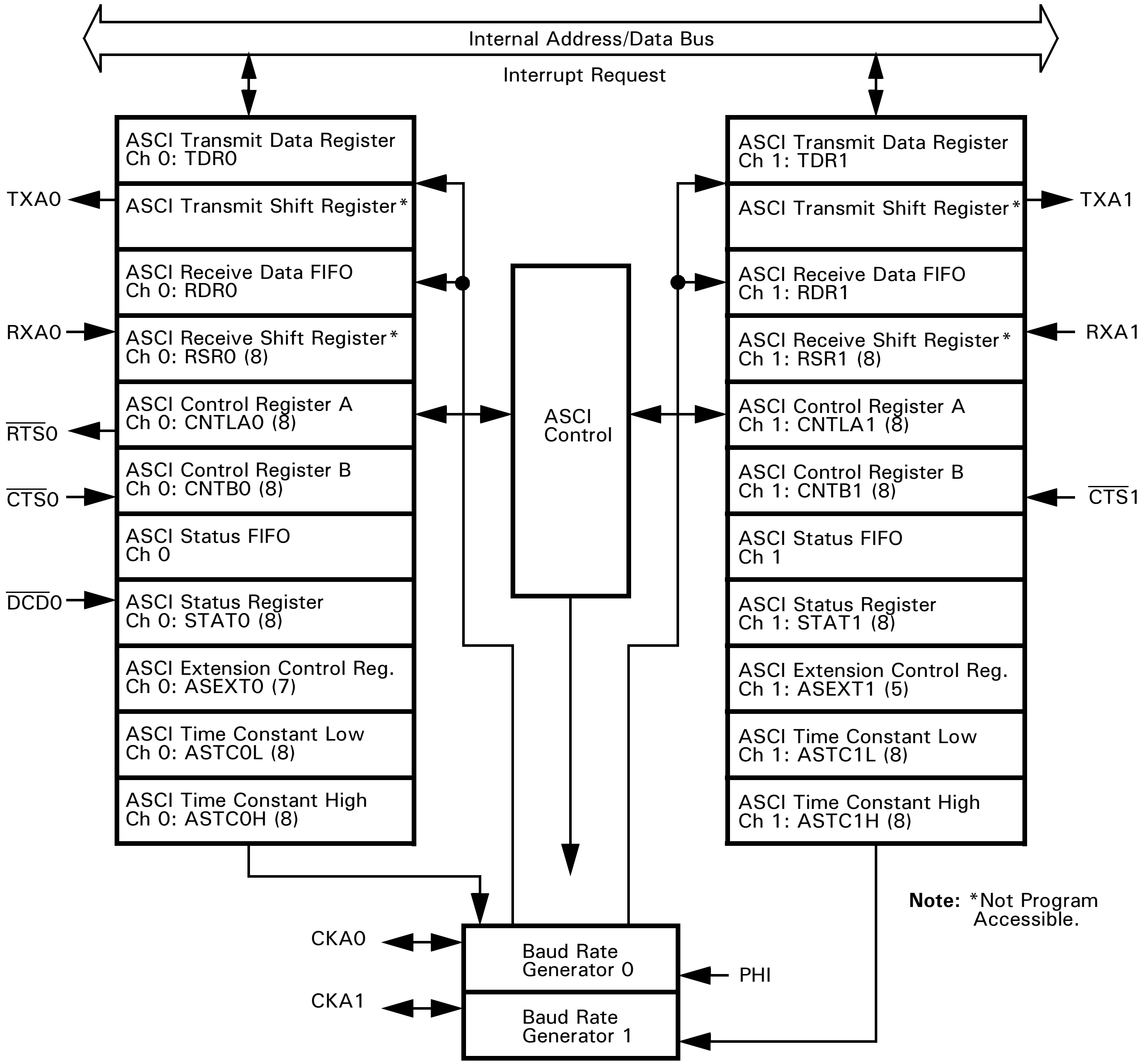

Figure 32. ASCI Block Diagram

ASCI Transmit Shift Register 0,1. When the ASCI Transmit Shift Register (TSR) receives data from the ASCI Transmit Data Register (TDR), the data is shifted out to the TXA pin. When transmission is completed, the next byte (if available) is automatically loaded from TDR into TSR and the next transmission starts. If no data is available for transmission, TSR idles by outputting a continuous High level. This register is not program-accessible

ASCI Transmit Data Register 0,1 (TDR0, 1: I/O address = 06H, 07H). Data written to the ASCI Transmit Data Register is transferred to the TSR as soon as TSR is empty. Data can be written while TSR is shifting out the previous byte of data. Thus, the ASCI transmitter is double buffered.

ASCI0 requests an interrupt when $\overline{\text{DCD0}}$ goes High. RIE is cleared to 0 by RESET.

$\overline{\text{DCD0}}$: Data Carrier Detect (Bit 2 STAT0). This bit is set to 1 when the pin is High. It is cleared to 0 on the first READ of STAT0 following the pin's transition from High to Low and during RESET. When bit 6 of the ASEXT0 register is 0 to select auto-enabling, and the pin is negated (High), the receiver is reset and its operation is inhibited.

$\overline{\text{CTS1E}}$: Clear To Send (Bit 2 STAT1). Channel 1 features an external $\overline{\text{CTS1}}$ input, which is multiplexed with the receive data pin RSX for the CSI/O. Setting this bit to 1 selects the CTS1 function; clearing the bit to 0 selects the RXS function.

TDRE: Transmit Data Register Empty (Bit 1). TDRE = 1 indicates that the TDR is empty and the next transmit data byte is written to TDR. After the byte is written to TDR, TDRE is cleared to 0 until the ASCI transfers the byte from TDR to the TSR and then TDRE is again set to 1. TDRE is set to 1 in IOSTOP mode and during RESET. On ASCI0, if the $\overline{\text{CTS0}}$ pin is auto-enabled in the ASEXT0 register and the pin is High, TDRE is reset to 0.

TIE: Transmit Interrupt Enable (Bit 0). TIE should be set to 1 to enable ASCI transmit interrupt requests. If TIE = 1, an interrupt is requested when TDRE = 1. TIE is cleared to 0 during RESET.

ASCI TRANSMIT DATA REGISTERS

Register addresses 06H and 07H hold the ASCI transmit data for channel 0 and channel 1, respectively.

ASCI Transmit Data Registers Channel 0

Mnemonic TDR0
Address 06H

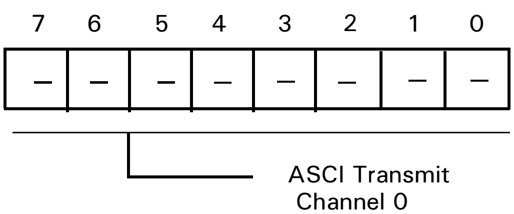

Figure 36. ASCI Register

ASCI Transmit Data Registers Channel 1

Mnemonic TDR1
Address 07H

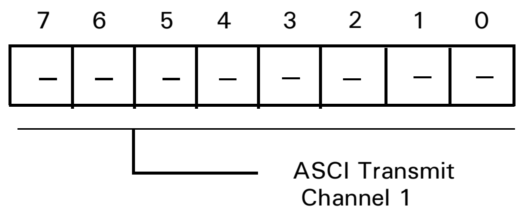

Figure 37. ASCI Register

ASCI RECEIVE REGISTER

Register addresses 08H and 09H hold the ASCI receive data for channel 0 and channel 1, respectively.

ASCI Receive Register Channel 0

Mnemonic RDR0
Address 08H

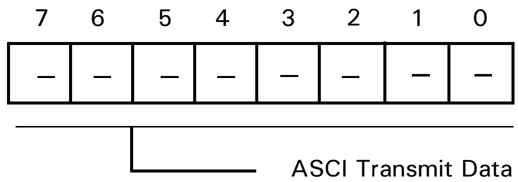

Figure 38. ASCI Receive Register Channel 0

ASCI Receive Register Channel 1

Mnemonic RDR1
Address 09H

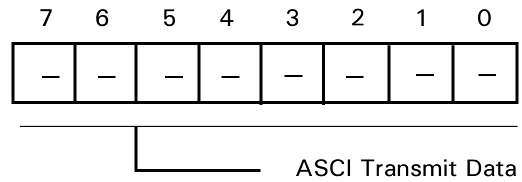

Figure 39. ASCI Receive Register Channel 1

CSI/O CONTROL/STATUS REGISTER

The CSI/O Control/Status Register (CNTR) is used to monitor CSI/O status, enable and disable the CSI/O, enable and disable interrupt generation, and select the data clock speed and source.

Bit	7	6	5	4	3	2	1	0
	EF	EIE	RE	TE	—	SS2	SS1	SS0
	R	R/W	R/W	R/W		R/W	R/W	R/W

Figure 40. CSI/O Control Register (CNTR: I/O Address = 000AH)

EF: End Flag (Bit 7). EF is set to 1 by the CSI/O to indicate completion of an 8-bit data transmit or receive operation. If End Interrupt Enable (EIE) bit = 1 when EF is set to 1, a CPU interrupt request is generated. Program access of TRDR only occurs if EF = 1. The CSI/O clears EF to 0 when TRDR is read or written. EF is cleared to 0 during RESET and IOSTOP mode.

EIE: End Interrupt Enable (Bit 6). EIE is set to 1 to generate a CPU interrupt request. The interrupt request is inhibited if EIE is reset to 0. EIE is cleared to 0 during RESET.

RE: Receive Enable (Bit 5). A CSI/O receive operation is started by setting RE to 1. When RE is set to 1, the data clock is enabled. In internal clock mode, the data clock is output from the CKS pin. In external clock mode, the clock is input on the CKS pin. In either case, data is shifted in on the RXS pin in synchronization with the (internal or external) data clock. After receiving 8 bits of data, the CSI/O automatically clears RE to 0, EF is set to 1, and an interrupt (if enabled by EIE = 1) is generated. RE and TE are never both set to 1 at the same time. RE is cleared to 0 during RESET and IOSTOP mode.

TE: Transmit Enable (Bit 4). A CSI/O transmit operation is started by setting TE to 1. When TE is set to 1, the data clock is enabled. When in internal clock mode, the data clock is output from the CKS pin. In external clock mode, the clock is input on the CKS pin. In either case, data is shifted out on the TXS pin synchronous with the (internal or external) data clock. After transmitting 8 bits of data, the CSI/O automatically clears TE to 0, sets EF to 1, and requests an interrupt if enabled by EIE = 1. TE and RE are

DMA BYTE COUNT REGISTER CHANNEL 0

The DMA Byte Count Register Channel 0 specifies the number of bytes to be transferred. This register contains 16 bits and may specify up to 64-KB transfers. When one byte is transferred, the register is decremented by one. If n bytes should be transferred, n must be stored before the DMA operation.

Note: All DMA Count Register channels are undefined during RESET.

DMA Byte Count Register Channel 0 Low

Mnemonic BCR0L
Address 26H

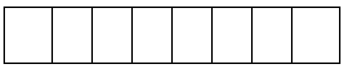

Figure 61. DMA Byte Count Register 0 Low

DMA Byte Count Register Channel 0 High

Mnemonic BCR0H
Address 27H

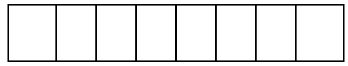

Figure 62. DMA Byte Count Register 0 High

DMA Byte Count Register Channel 1 Low

Mnemonic BCR1L
Address 2EH

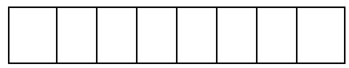

Figure 63. DMA Byte Count Register 1 Low

DMA Byte Count Register Channel 1 High

Mnemonic BCR1H
Address 2FH

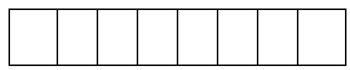

Figure 64. DMA Byte Count Register 1 High

DMA I/O ADDRESS REGISTER

The DMA I/O Address Register specifies the I/O device for channel 1 transfers. This address may be a destination or source I/O device. IAR1L and IAR1H each contain 8 address bits. The most significant byte identifies the Request Handshake signal and controls the Alternating Channel feature.

DMA I/O Address Register Channel 1 Low

Mnemonic IAR1L
Address 2BH

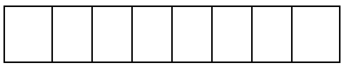

Figure 68. DMA I/O Address Register Channel 1 Low

DMA I/O Address Register Channel 1 High

Mnemonic IAR1H
Address 2CH

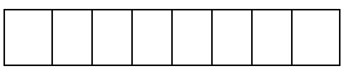

Figure 69. DMA I/O Address Register Channel 1 High

DMA I/O Address Register Channel 1 B

Mnemonic IAR1B
Address 2DH

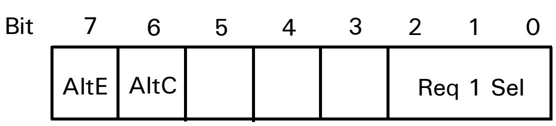

Figure 70. DMA I/O Address Register Channel 1 B

AltE. The AltE bit should be set only when both DMA channels are programmed for the same I/O source or I/O destination. In this case, a *channel end* condition (byte count = 0) on channel 0 sets bit 6 (AltC), which subsequently enables the channel 1 request and blocks the channel 0 request. Similarly, a channel end condition on channel 1 clears bit 6 (AltC), which then enables the channel 0 request and blocks the channel 1 request. For external requests, the request from the device must be routed or connected to both the DREQ0 and DREQ1 pins.

AltC. If bit (AltE) is 0, the AltC bit has no effect. When bit 7 (AltE) is 1 and the AltC bit is 0, the request signal selected by bits 2–0 is not presented to channel 1; however, the channel 0 request operates normally. When AltE is 1 and AltC is 1, the request selected by SAR18–16 or DAR18–16 is not presented to channel 0; however, the channel 1 request operates normally. The AltC bit can be written by software to select which channel should operate first; however, this operation should be executed only when both channels are stopped (both DE1 and DE0 are 0).

Req1Sel. If bit DIM1 in the DCNTL register is 1, indicating an I/O source, the following bits select which source handshake signal should control the transfer:

000	DREQ1 pin
001	ASCI0 RDRF
010	ASCI1 RDRF
Other	Reserved, do not program

If DIM1 is 0, indicating an I/O destination, the following bits select which destination handshake signal should control the transfer:

000	DREQ1 pin
001	ASCI0 TDRE
010	ASCI1 TDRE
Other	Reserved, do not program

Table 16 indicates all DMA transfer mode combinations of DM0, DM1, SM0, and SM1. Because I/O to/from I/O transfers are not implemented, 12 combinations are available.

Table 16. Transfer Mode Combinations

DM1	DM0	SM1	SM0	Transfer Mode	Address Increment/Decrement
0	0	0	0	Memory→Memory	SAR0 + 1, DAR0 + 1
0	0	0	1	Memory→Memory	SAR0–1, DAR0 + 1
0	0	1	0	Memory*→Memory	SAR0 fixed, DAR0 + 1
0	0	1	1	I/O→Memory	SAR0 fixed, DAR0 + 1
0	1	0	0	Memory→Memory	SAR0 + 1, DAR0–1
0	1	0	1	Memory→Memory	SAR0–1, DAR0–1
0	1	1	0	Memory*→Memory	SAR0 fixed, DAR0–1
0	1	1	1	I/O→Memory	SAR0 fixed, DAR0–1
1	0	0	0	Memory→Memory*	SAR0 + 1, DAR0 fixed
1	0	0	1	Memory→Memory*	SAR0–1, DAR0 fixed
1	0	1	0	Reserved	
1	0	1	1	Reserved	
1	1	0	0	Memory→I/O	SAR0 + 1, DAR0 fixed
1	1	0	1	Memory→I/O	SAR0–1, DAR0 fixed
1	1	1	0	Reserved	
1	1	1	1	Reserved	

Note: * Includes memory mapped I/O.

MMOD: Memory Mode Channel 0 (Bit 1). When channel 0 is configured for memory to/from memory transfers there is no Request Handshake signal to control the transfer timing. Instead, two automatic transfer timing modes are selectable: burst (MMOD = 1) and cycle steal (MMOD = 0). For burst memory to/from memory transfers, the DMAC takes control of the bus continuously until the DMA transfer completes (as indicated by the byte count register = 0). In cycle steal mode, the CPU is provided a cycle for each DMA byte transfer cycle until the transfer is completed.

For channel 0 DMA with I/O source or destination, the selected Request signal times the transfer ignoring MMOD. MMOD is cleared to 0 during RESET.

All TRAPs occur after fetching an undefined second opcode byte following one of the prefix opcodes (CBH, DDH, EDH, or FDH) or after fetching an undefined third opcode byte following one of the double-prefix opcodes (DDCBH or FDCBH).

The state of the Undefined Fetch Object (UFO) bit in ITC allows TRAP software to correctly *adjust* the stacked PC, depending on whether the second or third byte of the opcode generated the TRAP. If UFO = 0, the starting address of the invalid instruction is the stacked PC−1. If UFO = 1, the starting address of the invalid instruction is equal to the stacked PC−2.

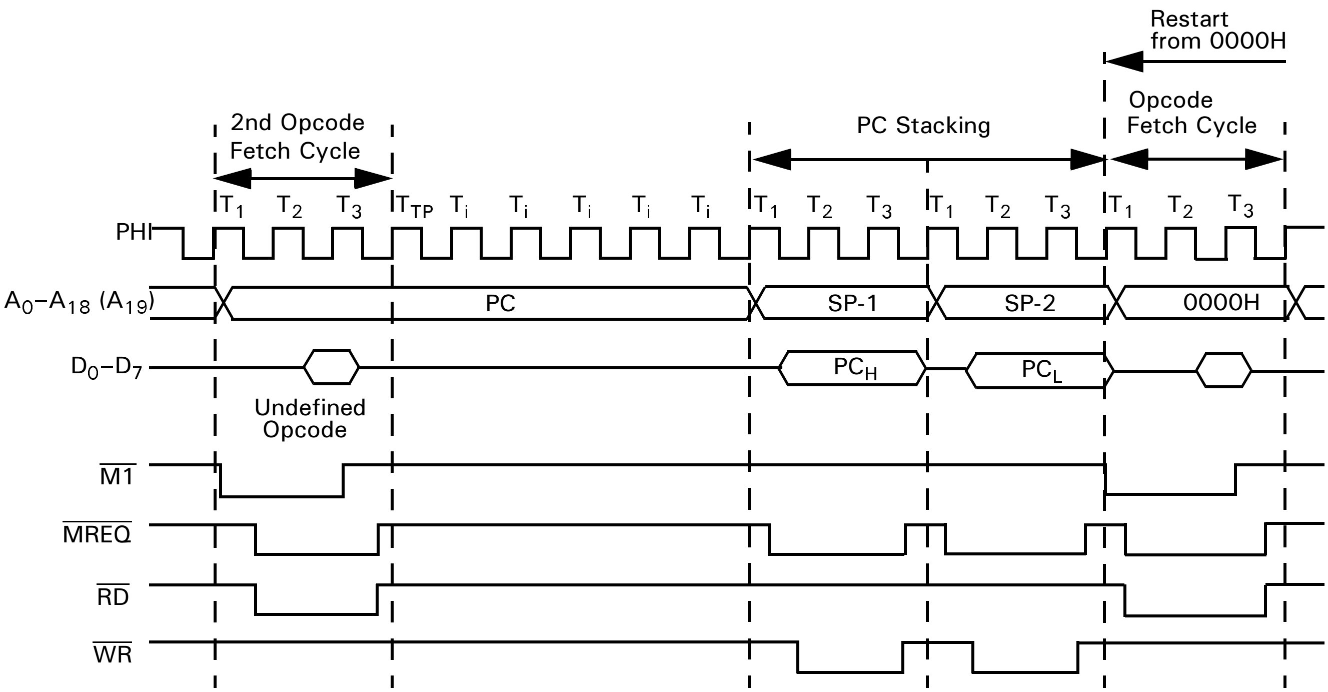

Figure 75. TRAP Timing—2nd Opcode Undefined

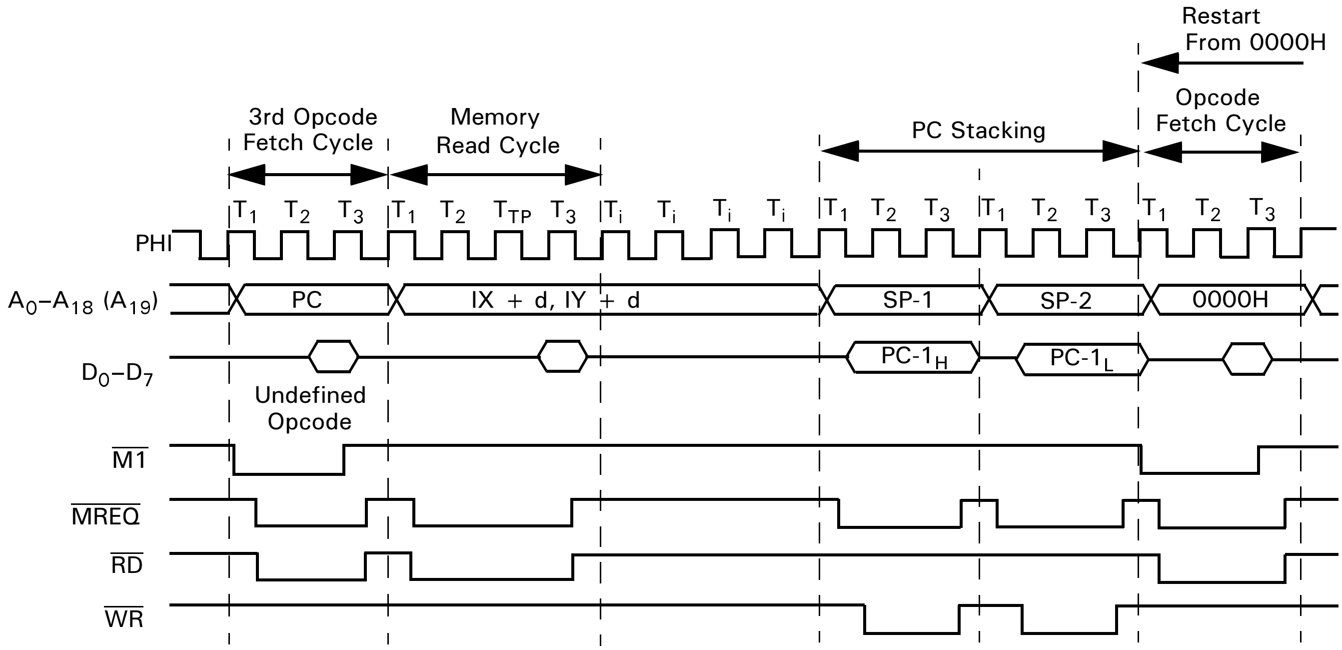

Figure 76. TRAP Timing—3rd Opcode Undefined

REFRESH CONTROL REGISTER

Mnemonic RCR
Address 36H

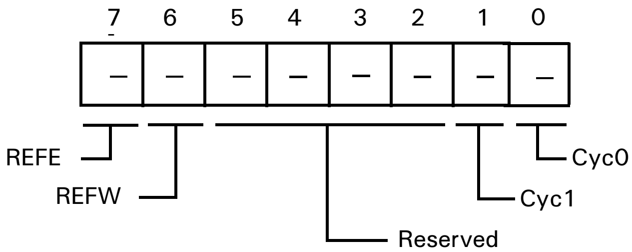

Figure 77. Refresh Control Register (RCR: I/O Address = 36H)

The Refresh Control Register (RCR) specifies the interval and length of refresh cycles, while enabling or disabling the refresh function.

REFE: Refresh Enable (Bit 7). REFE = 0 disables the refresh controller, while REFE = 1 enables refresh cycle insertion. REFE is set to 1 during RESET.

REFW: Refresh Wait (Bit 6). REFW = 0 causes the refresh cycle to be two clocks in duration. REFW = 1 causes the refresh cycle to be three clocks in duration by adding a refresh wait cycle (TRW). REFW is set to 1 during RESET.

CYC1, 0: Cycle Interval (Bit 1,0). CYC1 and CYC0 specify the interval (in clock cycles) between refresh cycles. When dynamic RAM requires 128 refresh cycles every 2 ms (or 256 cycles in every 4 ms), the required refresh interval is less than or equal to 15.625 µs. Thus, the underlined values indicate the best refresh interval depending on CPU clock frequency. CYC0 and CYC1 are cleared to 0 during RESET (see Table 18).

Table 18. DRAM Refresh Intervals

			Time Interval				
CYC1	**CYC0**	**Insertion Interval**	**PHI: 10 MHz**	**8 MHz**	**6 MHz**	**4 MHz**	**2.5 MHz**
0	0	10 states	(1.0 µs)*	(1.25 µs)*	1.66 µs	2.5 µs	4.0 µs
0	1	20 states	(2.0 µs)*	(2.5 µs)*	3.3 µs	5.0 µs	8.0 µs
1	0	40 states	(4.0 µs)*	(5.0 µs)*	6.6 µs	10.0 µs	16.0 µs
1	1	80 states	(8.0 µs)*	(10.0 µs)*	13.3 µs	20.0 µs	32.0 µs

Note: *calculated interval.

Refresh Control and Reset. After RESET, based on the initialized value of RCR, refresh cycles occur with an interval of 10 clock cycles and be 3 clock cycles in duration.

Dynamic RAM Refresh Operation

1. Refresh Cycle insertion is stopped when the CPU is in the following states:
   a. During RESET
   b. When the bus is released in response to $\overline{\text{BUSREQ}}$
   c. During SLEEP mode
   d. During $\overline{\text{WAIT}}$ states
2. Refresh cycles are suppressed when the bus is released in response to $\overline{\text{BUSREQ}}$. However, the refresh timer continues to operate. The time at which the first refresh cycle occurs after the Z8S180/Z8L180 reacquires the bus depends on the refresh timer. This cycle offers no timing relationship with the bus exchange.
3. Refresh cycles are suppressed during SLEEP mode. If a refresh cycle is requested during SLEEP mode, the refresh cycle request is internally latched (until replaced with the next refresh request). The latched refresh cycle is inserted at the end of the first machine cycle after SLEEP mode is exited. After this initial cycle, the time at which the next refresh cycle occurs depends on the refresh time and offers no relationship with the exit from SLEEP mode.
4. The refresh address is incremented by one for each successful refresh cycle, not for each refresh. Thus, independent of the number of missed refresh requests, each refresh bus cycle uses a refresh address incremented by one from that of the previous refresh bus cycles.

CA3–CA0:CA (Bits 7–4). CA specifies the start (Low) address (on 4-KB boundaries) for Common Area 1. This condition also determines the most recent address of the Bank Area. All bits of CA are set to 1 during RESET.

BA3–BA0 (Bits 3–0). BA specifies the start (Low) address (on 4-KB boundaries) for the Bank Area. This condition also determines the most recent address of Common Area 0. All bits of BA are set to 1 during RESET.

OPERATION MODE CONTROL REGISTER

The Z8S180/Z8L180 is descended from two different ancestor processors, ZiLOG's original Z80 and the Hitachi 64180. The Operating Mode Control Register (OMCR) can be programmed to select between certain differences between the Z80 and the 64180.

Operation Mode Control Register

Mnemonic OMCR
Address 3EH

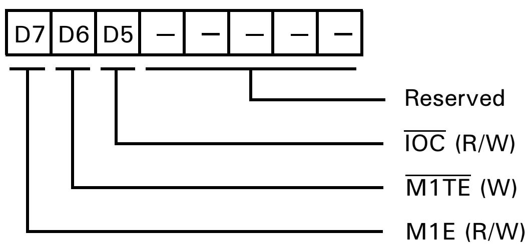

Figure 81. Operating Control Register (OMCR: I/O Address = 3EH)

M1E ($\overline{\text{M1}}$ Enable). This bit controls the $\overline{\text{M1}}$ output and is set to a 1 during reset.

When M1E = 1, the $\overline{\text{M1}}$ output is asserted Low during the opcode fetch cycle, the $\overline{\text{INT0}}$ acknowledge cycle, and the first machine cycle of the $\overline{\text{NMI}}$ acknowledge.

On the Z8S180/Z8L180, this choice makes the processor fetch one RETI instruction. When fetching a RETI from zero-wait-state memory, the processor uses three clock machine cycles that are not fully Z80-timing-compatible.

When M1E = 0, the processor does not drive $\overline{\text{M1}}$ Low during instruction fetch cycles. After fetching one RETI instruction with normal timing, the processor returns and refetches the instruction using Z80-compatible cycles that drive $\overline{\text{M1}}$ Low. This timing compatibility may be required by external Z80 peripherals to properly decode the RETI instruction.

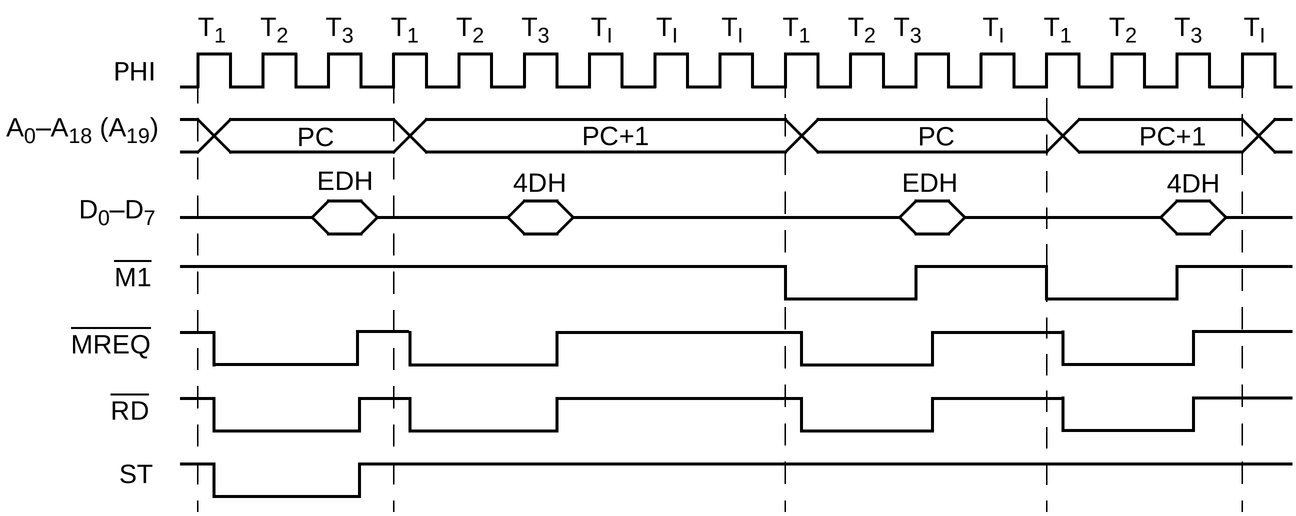

Figure 82. RETI Instruction Sequence with M1E = 0

ORDERING INFORMATION

Codes	
Speed	10 = 10 MHz 20 = 20 MHz 33 = 33 MHz
Package	P = 60-Pin Plastic DIP V = 68-Pin PLCC F = 80-Pin QFP
Temperature	S = 0°C to +70°C E = –40°C to +85°C
Environmental	C = Plastic Standard

For fast results, contact your local ZiLOG sales office for assistance in ordering the part(s) required.

Example:

Z 8S180 10 P S C   is a Z8S180 10-MHz 60-Pin DIP, 0° to +70°C, Plastic Standard Flow

- C — Environmental Flow
- S — Temperature
- P — Package
- 10 — Speed
- 8S180 — Product Number
- Z — ZiLOG Prefix

Pre-Characterization Product

The product represented by this document is newly introduced and ZiLOG has not completed the full characterization of the product. The document states what ZiLOG knows about this product at this time, but additional features or non-conformance with some aspects of the document may be found, either by ZiLOG or its customers in the course of further application and characterization work. In addition, ZiLOG cautions that delivery may be uncertain at times, due to start-up yield issues.

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