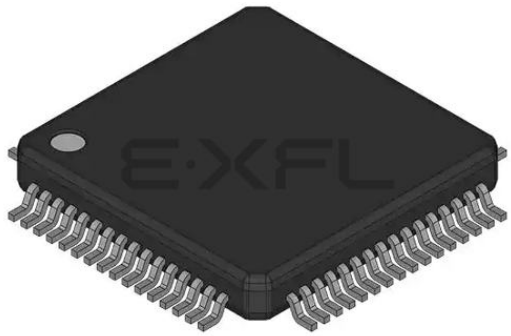


Welcome to E-XFL.COM

What is "Embedded - Microcontrollers"?

"Embedded - Microcontrollers" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	HC08
Core Size	8-Bit
Speed	8MHz
Connectivity	CANbus, SCI, SPI
Peripherals	LVD, POR, PWM
Number of I/O	53
Program Memory Size	32KB (32K x 8)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	1.5K x 8
Voltage - Supply (Vcc/Vdd)	3V ~ 5.5V
Data Converters	A/D 24x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 125°C (TA)
Mounting Type	Surface Mount
Package / Case	64-QFP
Supplier Device Package	64-QFP (14x14)
Purchase URL	https://www.e-xfl.com/pro/item?MUrl=&PartUrl=mc908gz32mfue

Chapter 15
System Integration Module (SIM)

Chapter 18
Timer Interface Module (TIM1)

Chapter 19
Timer Interface Module (TIM2)

Table 2-1. Vector Addresses (Continued)

Vector Priority	Vector	Address	Vector
	IF15	$FFDE	ADC Conversion Complete Vector (High)
		$FFDF	ADC Conversion Complete Vector (Low)
	IF14	$FFE0	Keyboard Vector (High)
		$FFE1	Keyboard Vector (Low)
	IF13	$FFE2	ESCI Transmit Vector (High)
		$FFE3	ESCI Transmit Vector (Low)
	IF12	$FFE4	ESCI Receive Vector (High)
		$FFE5	ESCI Receive Vector (Low)
	IF11	$FFE6	ESCI Error Vector (High)
		$FFE7	ESCI Error Vector (Low)
	IF10	$FFE8	SPI Transmit Vector (High)
		$FFE9	SPI Transmit Vector (Low)
	IF9	$FFEA	SPI Receive Vector (High)
		$FFEB	SPI Receive Vector (Low)
	IF8	$FFEC	TIM2 Overflow Vector (High)
		$FFED	TIM2 Overflow Vector (Low)
	IF7	$FFEE	TIM2 Channel 1 Vector (High)
		$FFEF	TIM2 Channel 1 Vector (Low)
	IF6	$FFF0	TIM2 Channel 0 Vector (High)
		$FFF1	TIM2 Channel 0 Vector (Low)
	IF5	$FFF2	TIM1 Overflow Vector (High)
		$FFF3	TIM1 Overflow Vector (Low)
	IF4	$FFF4	TIM1 Channel 1 Vector (High)
		$FFF5	TIM1 Channel 1 Vector (Low)
	IF3	$FFF6	TIM1 Channel 0 Vector (High)
		$FFF7	TIM1 Channel 0 Vector (Low)
	IF2	$FFF8	PLL Vector (High)
		$FFF9	PLL Vector (Low)
	IF1	$FFFA	$\overline{\text{IRQ}}$ Vector (High)
		$FFFB	$\overline{\text{IRQ}}$ Vector (Low)
	—	$FFFC	SWI Vector (High)
		$FFFD	SWI Vector (Low)
Highest	—	$FFFE	Reset Vector (High)
		$FFFF	Reset Vector (Low)

2.5 Random-Access Memory (RAM)

The RAM locations are broken into two non-continuous memory blocks. The RAM addresses locations are $0040–$043F and $0580–$097F. The location of the stack RAM is programmable. The 16-bit stack pointer allows the stack to be anywhere in the 64-Kbyte memory space.

NOTE

For correct operation, the stack pointer must point only to RAM locations.

Within page zero are 192 bytes of RAM. Because the location of the stack RAM is programmable, all page zero RAM locations can be used for I/O control and user data or code. When the stack pointer is moved from its reset location at $00FF out of page zero, direct addressing mode instructions can efficiently access all page zero RAM locations. Page zero RAM, therefore, provides ideal locations for frequently accessed global variables.

Before processing an interrupt, the CPU uses five bytes of the stack to save the contents of the CPU registers.

NOTE

For M6805 compatibility, the H register is not stacked.

During a subroutine call, the CPU uses two bytes of the stack to store the return address. The stack pointer decrements during pushes and increments during pulls.

NOTE

Be careful when using nested subroutines. The CPU may overwrite data in the RAM during a subroutine or during the interrupt stacking operation.

2.6 FLASH-1 Memory (FLASH-1)

This subsection describes the operation of the embedded FLASH-1 memory. This memory can be read, programmed, and erased from a single external supply. The program and erase operations are enabled through the use of an internal charge pump.

2.6.1 Functional Description

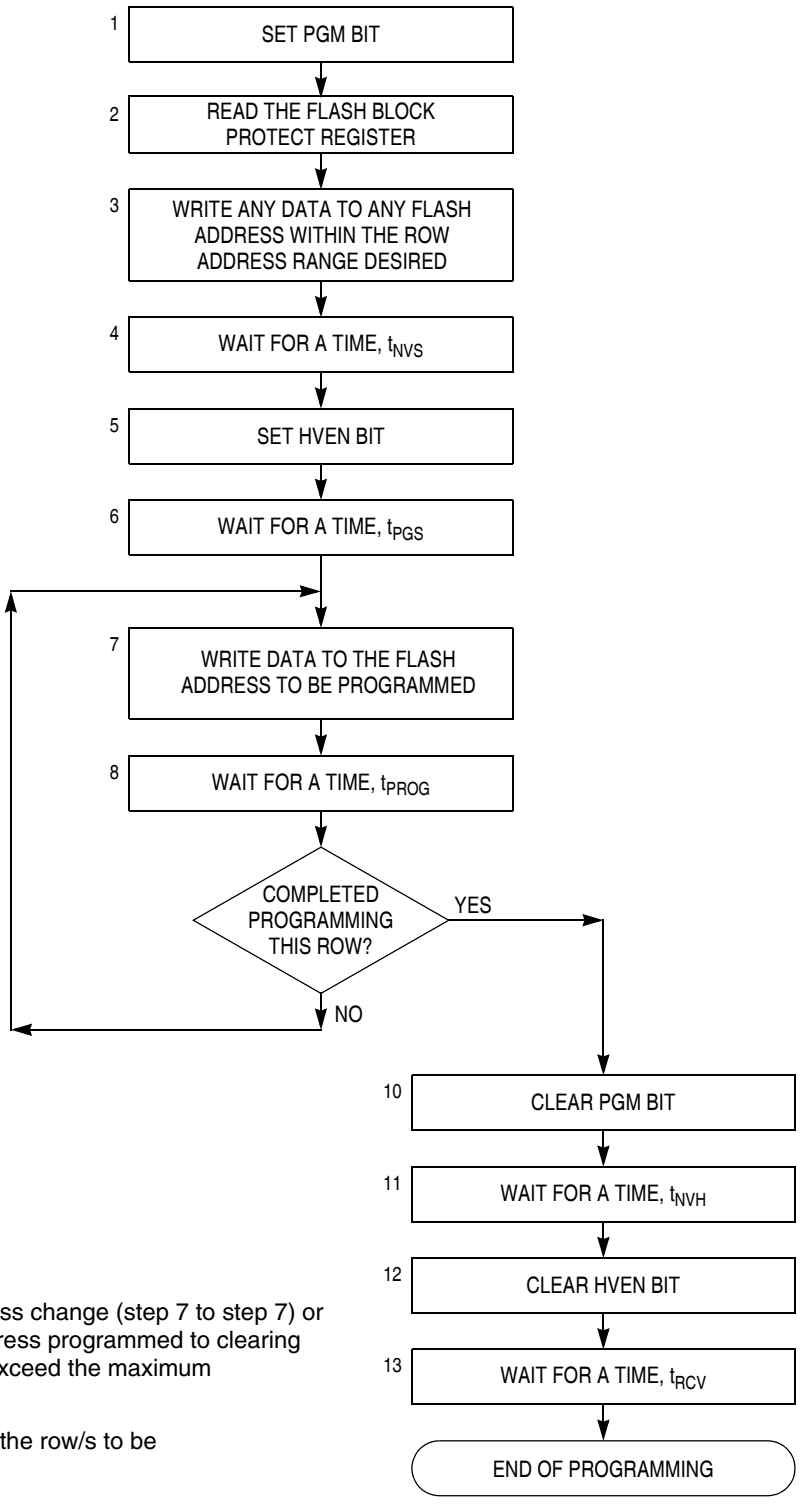
The FLASH-1 memory is an array of 32,256 bytes with two bytes of block protection (one byte for protecting areas within FLASH-1 array and one byte for protecting areas within FLASH-2 array) and an additional 52 bytes of user vectors. An erased bit reads as a 1 and a programmed bit reads as a 0.

Memory in the FLASH-1 array is organized into rows within pages. There are two rows of memory per page with 64 bytes per row. The minimum erase block size is a single page,128 bytes. Programming is performed on a per-row basis, 64 bytes at a time. Program and erase operations are facilitated through control bits in the FLASH-1 control register (FL1CR). Details for these operations appear later in this subsection.

The FLASH-1 memory map consists of:

- $8000–$FDFF: user memory (32,256 bytes)
- $FF80: FLASH-1 block protect register (FL1BPR)
- $FF81: FLASH-2 block protect register (FL2BPR)
- $FF88: FLASH-1 control register (FL1CR)
- $FFCC–$FFFF: these locations are reserved for user-defined interrupt and reset vectors (see Table 2-1 for details)

Algorithm for programming a row (64 bytes) of FLASH memory

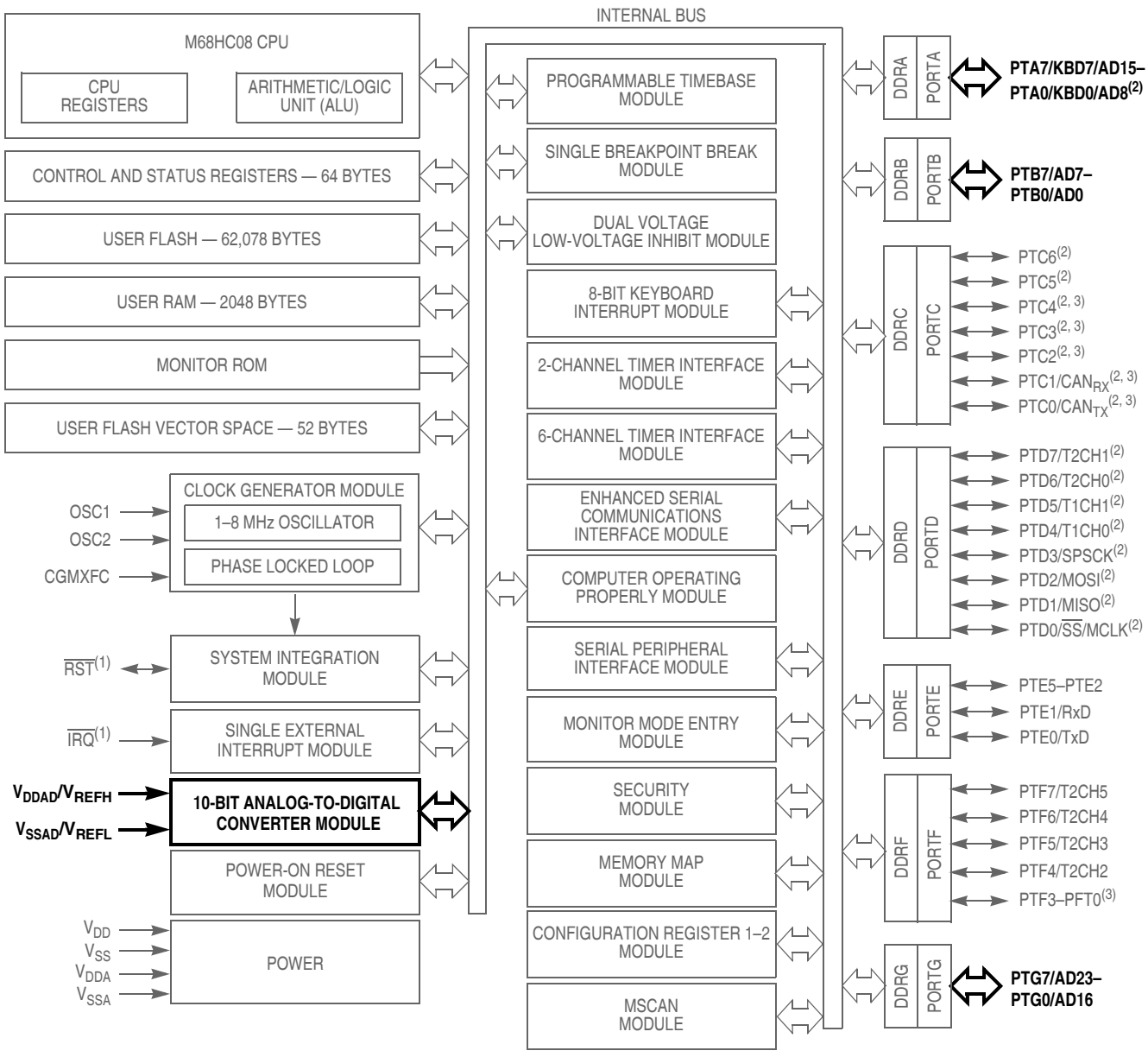


NOTES:

The time between each FLASH address change (step 7 to step 7) or the time between the last FLASH address programmed to clearing PGM bit (step 7 to step10) must not exceed the maximum programming time, t_{PROG}, maximum.

This row program algorithm assumes the row/s to be programmed are initially erased.

Figure 2-6. FLASH-1 Programming Algorithm Flowchart



1. Pin contains integrated pullup device.
2. Ports are software configurable with pullup device if input port or pullup/pulldown device for keyboard input.
3. Higher current drive port pins

Figure 3-1. Block Diagram Highlighting ADC Block and Pins

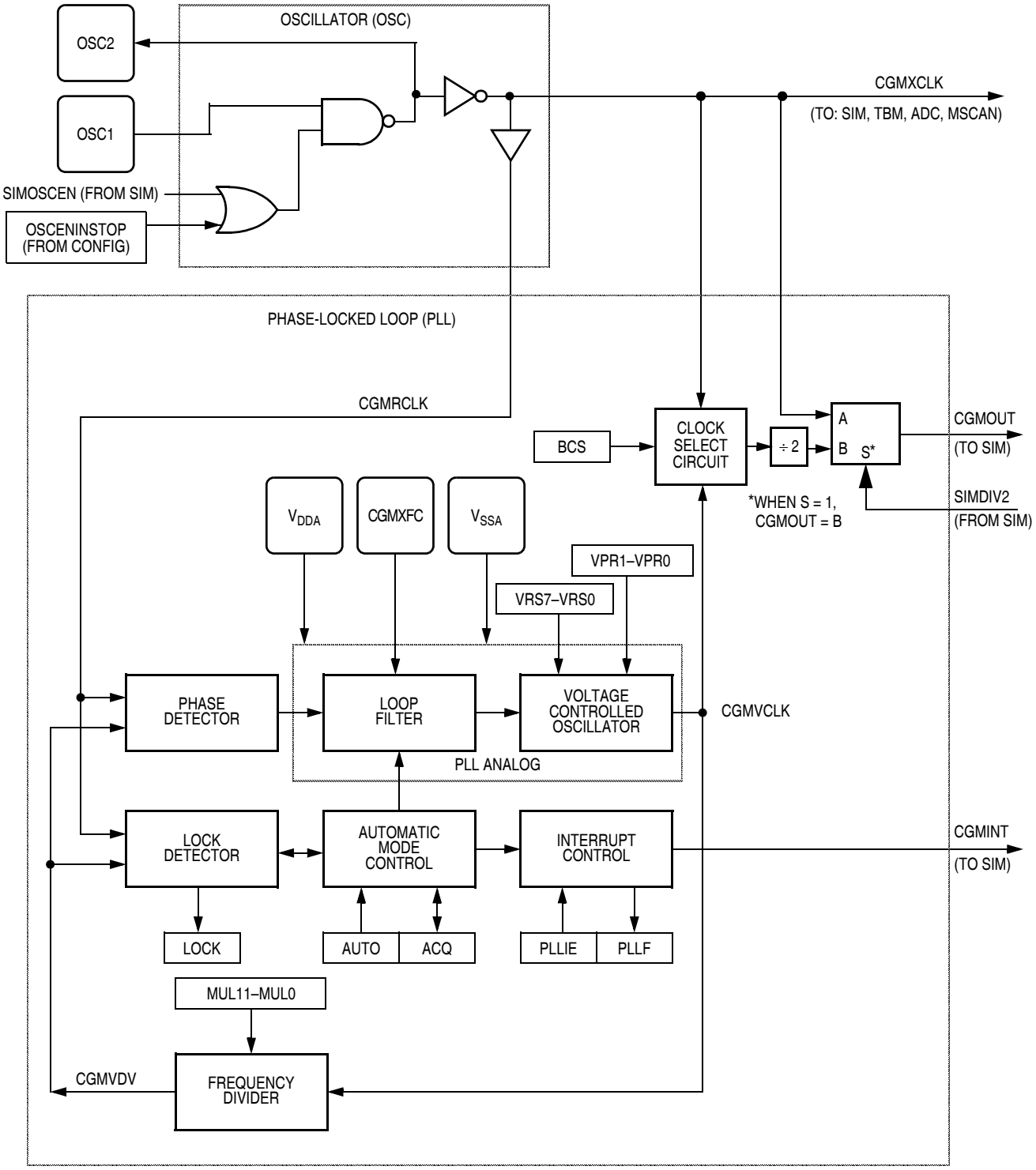


Figure 4-1. CGM Block Diagram

frequency, f_{RCLK}. The circuit determines the mode of the PLL and the lock condition based on this comparison.

4.3.4 Acquisition and Tracking Modes

The PLL filter is manually or automatically configurable into one of two operating modes:

- Acquisition mode — In acquisition mode, the filter can make large frequency corrections to the VCO. This mode is used at PLL start up or when the PLL has suffered a severe noise hit and the VCO frequency is far off the desired frequency. When in acquisition mode, the $\overline{ACQ}$ bit is clear in the PLL bandwidth control register. (See 4.5.2 PLL Bandwidth Control Register.)
- Tracking mode — In tracking mode, the filter makes only small corrections to the frequency of the VCO. PLL jitter is much lower in tracking mode, but the response to noise is also slower. The PLL enters tracking mode when the VCO frequency is nearly correct, such as when the PLL is selected as the base clock source. (See 4.3.8 Base Clock Selector Circuit.) The PLL is automatically in tracking mode when not in acquisition mode or when the $\overline{ACQ}$ bit is set.

4.3.5 Manual and Automatic PLL Bandwidth Modes

The PLL can change the bandwidth or operational mode of the loop filter manually or automatically. Automatic mode is recommended for most users.

In automatic bandwidth control mode (AUTO = 1), the lock detector automatically switches between acquisition and tracking modes. Automatic bandwidth control mode also is used to determine when the VCO clock, CGMVCLK, is safe to use as the source for the base clock, CGMOUT. (See 4.5.2 PLL Bandwidth Control Register.) If PLL interrupts are enabled, the software can wait for a PLL interrupt request and then check the LOCK bit. If interrupts are disabled, software can poll the LOCK bit continuously (for example, during PLL start up) or at periodic intervals. In either case, when the LOCK bit is set, the VCO clock is safe to use as the source for the base clock. (See 4.3.8 Base Clock Selector Circuit.) If the VCO is selected as the source for the base clock and the LOCK bit is clear, the PLL has suffered a severe noise hit and the software must take appropriate action, depending on the application. (See 4.6 Interrupts for information and precautions on using interrupts.)

The following conditions apply when the PLL is in automatic bandwidth control mode:

- The $\overline{ACQ}$ bit (See 4.5.2 PLL Bandwidth Control Register.) is a read-only indicator of the mode of the filter. (See 4.3.4 Acquisition and Tracking Modes.)
- The $\overline{ACQ}$ bit is set when the VCO frequency is within a certain tolerance and is cleared when the VCO frequency is out of a certain tolerance. (See 4.8 Acquisition/Lock Time Specifications for more information.)
- The LOCK bit is a read-only indicator of the locked state of the PLL.
- The LOCK bit is set when the VCO frequency is within a certain tolerance and is cleared when the VCO frequency is out of a certain tolerance. (See 4.8 Acquisition/Lock Time Specifications for more information.)
- CPU interrupts can occur if enabled (PLLIE = 1) when the PLL's lock condition changes, toggling the LOCK bit. (See 4.5.1 PLL Control Register.)

The PLL also may operate in manual mode (AUTO = 0). Manual mode is used by systems that do not require an indicator of the lock condition for proper operation. Such systems typically operate well below f_{BUSMAX}.

Address: $001E

	Bit 7	6	5	4	3	2	1	Bit 0
Read:	0	MCLKSEL	MCLK1	MCLK0	MSCANEN	TMBCLKSEL	OSCENINSTOP	SCIBDSRC
Write:								
Reset:	0	0	0	0	See note	0	0	1

Note: MSCANEN is only reset via POR (power-on reset).

= Unimplemented

Figure 5-1. Configuration Register 2 (CONFIG2)

MCLKSEL — MCLK Source Select Bit

1 = Crystal frequency
0 = Bus frequency

MCLK1 and MCLK0 — MCLK Output Select Bits

Setting the MCLK1 and MCLK0 bits enables the PTD0/$\overline{SS}$ pin to be used as a MCLK output clock. Once configured for MCLK, the PTD data direction register for PTD0 is used to enable and disable the MCLK output. See Table 5-1 for MCLK options.

Table 5-1. MCLK Output Select

MCLK1	MCLK0	MCLK Frequency
0	0	MCLK not enabled
0	1	Clock
1	0	Clock divided by 2
1	1	Clock divided by 4

MSCANEN— MSCAN08 Enable Bit

Setting the MSCANEN enables the MSCAN08 module and allows the MSCAN08 to use the PTC0/PTC1 pins. See Chapter 12 MSCAN08 Controller (MSCAN08) for a more detailed description of the MSCAN08 operation.

1 = Enables MSCAN08 module
0 = Disables the MSCAN08 module

NOTE

The MSCANEN bit is cleared by a power-on reset (POR) only. Other resets will leave this bit unaffected.

TMBCLKSEL— Timebase Clock Select Bit

TMBCLKSEL enables an extra divide-by-128 prescaler in the timebase module. Setting this bit enables the extra prescaler and clearing this bit disables it. See Chapter 17 Timebase Module (TBM) for a more detailed description of the external clock operation.

1 = Enables extra divide-by-128 prescaler in timebase module
0 = Disables extra divide-by-128 prescaler in timebase module

6.3.6 COPD (COP Disable)

The COPD signal reflects the state of the COP disable bit (COPD) in the configuration register. See Chapter 5 Configuration Register (CONFIG).

6.3.7 COPRS (COP Rate Select)

The COPRS signal reflects the state of the COP rate select bit (COPRS) in the configuration register. See Chapter 5 Configuration Register (CONFIG).

6.4 COP Control Register

The COP control register (COPCTL) is located at address $FFFF and overlaps the reset vector. Writing any value to $FFFF clears the COP counter and starts a new timeout period. Reading location $FFFF returns the low byte of the reset vector.

Address: $FFFF

	Bit 7	6	5	4	3	2	1	Bit 0
Read:	Low byte of reset vector							
Write:	Clear COP counter							
Reset:	Unaffected by reset							

Figure 6-2. COP Control Register (COPCTL)

6.5 Interrupts

The COP does not generate central processor unit (CPU) interrupt requests.

6.6 Monitor Mode

When monitor mode is entered with V_{TST} on the $\overline{IRQ}$ pin, the COP is disabled as long as V_{TST} remains on the $\overline{IRQ}$ pin or the $\overline{RST}$ pin. When monitor mode is entered by having blank reset vectors and not having V_{TST} on the $\overline{IRQ}$ pin, the COP is automatically disabled until a POR occurs.

6.7 Low-Power Modes

The WAIT and STOP instructions put the microcontroller unit (MCU) in low power-consumption standby modes.

6.7.1 Wait Mode

The COP remains active during wait mode. If COP is enabled, a reset will occur at COP timeout.

6.7.2 Stop Mode

Stop mode turns off the CGMXCLK input to the COP and clears the SIM counter. Service the COP immediately before entering or after exiting stop mode to ensure a full COP timeout period after entering or exiting stop mode.

7.3.5 Condition Code Register

The 8-bit condition code register contains the interrupt mask and five flags that indicate the results of the instruction just executed. Bits 6 and 5 are set permanently to 1. The following paragraphs describe the functions of the condition code register.

	Bit 7	6	5	4	3	2	1	Bit 0
Read: Write:	V	1	1	H	I	N	Z	C
Reset:	X	1	1	X	1	X	X	X

X = Indeterminate

Figure 7-6. Condition Code Register (CCR)

V — Overflow Flag

The CPU sets the overflow flag when a two's complement overflow occurs. The signed branch instructions BGT, BGE, BLE, and BLT use the overflow flag.

1 = Overflow
0 = No overflow

H — Half-Carry Flag

The CPU sets the half-carry flag when a carry occurs between accumulator bits 3 and 4 during an add-without-carry (ADD) or add-with-carry (ADC) operation. The half-carry flag is required for binary-coded decimal (BCD) arithmetic operations. The DAA instruction uses the states of the H and C flags to determine the appropriate correction factor.

1 = Carry between bits 3 and 4
0 = No carry between bits 3 and 4

I — Interrupt Mask

When the interrupt mask is set, all maskable CPU interrupts are disabled. CPU interrupts are enabled when the interrupt mask is cleared. When a CPU interrupt occurs, the interrupt mask is set automatically after the CPU registers are saved on the stack, but before the interrupt vector is fetched.

1 = Interrupts disabled
0 = Interrupts enabled

NOTE

To maintain M6805 Family compatibility, the upper byte of the index register (H) is not stacked automatically. If the interrupt service routine modifies H, then the user must stack and unstack H using the PSHH and PULH instructions.

After the I bit is cleared, the highest-priority interrupt request is serviced first.

A return-from-interrupt (RTI) instruction pulls the CPU registers from the stack and restores the interrupt mask from the stack. After any reset, the interrupt mask is set and can be cleared only by the clear interrupt mask software instruction (CLI).

N — Negative Flag

The CPU sets the negative flag when an arithmetic operation, logic operation, or data manipulation produces a negative result, setting bit 7 of the result.

1 = Negative result
0 = Non-negative result

Table 7-1. Instruction Set Summary (Sheet 2 of 6)

Source Form	Operation	Description	Effect on CCR						Address Mode	Opcode	Operand	Cycles
			V	H	I	N	Z	C				
BHS *rel*	Branch if Higher or Same (Same as BCC)	PC ← (PC) + 2 + *rel* ? (C) = 0	–	–	–	–	–	–	REL	24	rr	3
BIH *rel*	Branch if $\overline{\text{IRQ}}$ Pin High	PC ← (PC) + 2 + *rel* ? $\overline{\text{IRQ}}$ = 1	–	–	–	–	–	–	REL	2F	rr	3
BIL *rel*	Branch if $\overline{\text{IRQ}}$ Pin Low	PC ← (PC) + 2 + *rel* ? $\overline{\text{IRQ}}$ = 0	–	–	–	–	–	–	REL	2E	rr	3
BIT #*opr* BIT *opr* BIT *opr* BIT *opr*,X BIT *opr*,X BIT ,X BIT *opr*,SP BIT *opr*,SP	Bit Test	(A) & (M)	0	–	–	↕	↕	–	IMM DIR EXT IX2 IX1 IX SP1 SP2	A5 B5 C5 D5 E5 F5 9EE5 9ED5	ii dd hh ll ee ff ff ff ee ff	2 3 4 4 3 2 4 5
BLE *opr*	Branch if Less Than or Equal To (Signed Operands)	PC ← (PC) + 2 + *rel* ? (Z) ∣ (N ⊕ V) = 1	–	–	–	–	–	–	REL	93	rr	3
BLO *rel*	Branch if Lower (Same as BCS)	PC ← (PC) + 2 + *rel* ? (C) = 1	–	–	–	–	–	–	REL	25	rr	3
BLS *rel*	Branch if Lower or Same	PC ← (PC) + 2 + *rel* ? (C) ∣ (Z) = 1	–	–	–	–	–	–	REL	23	rr	3
BLT *opr*	Branch if Less Than (Signed Operands)	PC ← (PC) + 2 + *rel* ? (N ⊕ V) =1	–	–	–	–	–	–	REL	91	rr	3
BMC *rel*	Branch if Interrupt Mask Clear	PC ← (PC) + 2 + *rel* ? (I) = 0	–	–	–	–	–	–	REL	2C	rr	3
BMI *rel*	Branch if Minus	PC ← (PC) + 2 + *rel* ? (N) = 1	–	–	–	–	–	–	REL	2B	rr	3
BMS *rel*	Branch if Interrupt Mask Set	PC ← (PC) + 2 + *rel* ? (I) = 1	–	–	–	–	–	–	REL	2D	rr	3
BNE *rel*	Branch if Not Equal	PC ← (PC) + 2 + *rel* ? (Z) = 0	–	–	–	–	–	–	REL	26	rr	3
BPL *rel*	Branch if Plus	PC ← (PC) + 2 + *rel* ? (N) = 0	–	–	–	–	–	–	REL	2A	rr	3
BRA *rel*	Branch Always	PC ← (PC) + 2 + *rel*	–	–	–	–	–	–	REL	20	rr	3
BRCLR *n*,*opr*,*rel*	Branch if Bit *n* in M Clear	PC ← (PC) + 3 + *rel* ? (Mn) = 0	–	–	–	–	–	↕	DIR (b0) DIR (b1) DIR (b2) DIR (b3) DIR (b4) DIR (b5) DIR (b6) DIR (b7)	01 03 05 07 09 0B 0D 0F	dd rr dd rr dd rr dd rr dd rr dd rr dd rr dd rr	5 5 5 5 5 5 5 5
BRN *rel*	Branch Never	PC ← (PC) + 2	–	–	–	–	–	–	REL	21	rr	3
BRSET *n*,*opr*,*rel*	Branch if Bit *n* in M Set	PC ← (PC) + 3 + *rel* ? (Mn) = 1	–	–	–	–	–	↕	DIR (b0) DIR (b1) DIR (b2) DIR (b3) DIR (b4) DIR (b5) DIR (b6) DIR (b7)	00 02 04 06 08 0A 0C 0E	dd rr dd rr dd rr dd rr dd rr dd rr dd rr dd rr	5 5 5 5 5 5 5 5
BSET *n*,*opr*	Set Bit *n* in M	Mn ← 1	–	–	–	–	–	–	DIR (b0) DIR (b1) DIR (b2) DIR (b3) DIR (b4) DIR (b5) DIR (b6) DIR (b7)	10 12 14 16 18 1A 1C 1E	dd dd dd dd dd dd dd dd	4 4 4 4 4 4 4 4
BSR *rel*	Branch to Subroutine	PC ← (PC) + 2; push (PCL) SP ← (SP) – 1; push (PCH) SP ← (SP) – 1 PC ← (PC) + *rel*	–	–	–	–	–	–	REL	AD	rr	4
CBEQ *opr*,*rel* CBEQA #*opr*,*rel* CBEQX #*opr*,*rel* CBEQ *opr*,X+,*rel* CBEQ X+,*rel* CBEQ *opr*,SP,*rel*	Compare and Branch if Equal	PC ← (PC) + 3 + rel ? (A) – (M) = $00 PC ← (PC) + 3 + rel ? (A) – (M) = $00 PC ← (PC) + 3 + rel ? (X) – (M) = $00 PC ← (PC) + 3 + rel ? (A) – (M) = $00 PC ← (PC) + 2 + rel ? (A) – (M) = $00 PC ← (PC) + 4 + rel ? (A) – (M) = $00	–	–	–	–	–	–	DIR IMM IMM IX1+ IX+ SP1	31 41 51 61 71 9E61	dd rr ii rr ii rr ff rr rr ff rr	5 4 4 5 4 6
CLC	Clear Carry Bit	C ← 0	–	–	–	–	–	0	INH	98		1
CLI	Clear Interrupt Mask	I ← 0	–	–	0	–	–	–	INH	9A		2

12.13.8 MSCAN08 Transmitter Control Register

Address: $0507

	Bit 7	6	5	4	3	2	1	Bit 0
Read:	0	ABTRQ2	ABTRQ1	ABTRQ0	0	TXEIE2	TXEIE1	TXEIE0
Write:								
Reset:	0	0	0	0	0	0	0	0

= Unimplemented

Figure 12-23. Transmitter Control Register (CTCR)

ABTRQ2–ABTRQ0 — Abort Request

The CPU sets an ABTRQx bit to request that an already scheduled message buffer (TXE = 0) be aborted. The MSCAN08 will grant the request if the message has not already started transmission, or if the transmission is not successful (lost arbitration or error). When a message is aborted the associated TXE and the abort acknowledge flag (ABTAK) (see 12.13.7 MSCAN08 Transmitter Flag Register) will be set and an TXE interrupt is generated if enabled. The CPU cannot reset ABTRQx. ABTRQx is cleared implicitly whenever the associated TXE flag is set.

1 = Abort request pending
0 = No abort request

NOTE

The software must not clear one or more of the TXE flags in CTFLG and simultaneously set the respective ABTRQ bit(s).

TXEIE2–TXEIE0 — Transmitter Empty Interrupt Enable

1 = A transmitter empty (transmit buffer available for transmission) event results in a transmitter empty interrupt.
0 = No interrupt is generated from this event.

NOTE

The CTCR register is held in the reset state when the SFTRES bit in CMCR0 is set.

12.13.9 MSCAN08 Identifier Acceptance Control Register

Address: $0508

	Bit 7	6	5	4	3	2	1	Bit 0
Read:	0	IDAM2	IDAM1	IDAM0	0	IDHIT2	IDHIT1	IDHIT0
Write:								
Reset:	0	0	0	0	0	0	0	0

= Unimplemented

Figure 12-24. Identifier Acceptance Control Register (CIDAC)

14.4.2 Transmitter

Figure 14-5 shows the structure of the SCI transmitter and the registers are summarized in Figure 14-4. The baud rate clock source for the ESCI can be selected via the configuration bit, SCIBDSRC.

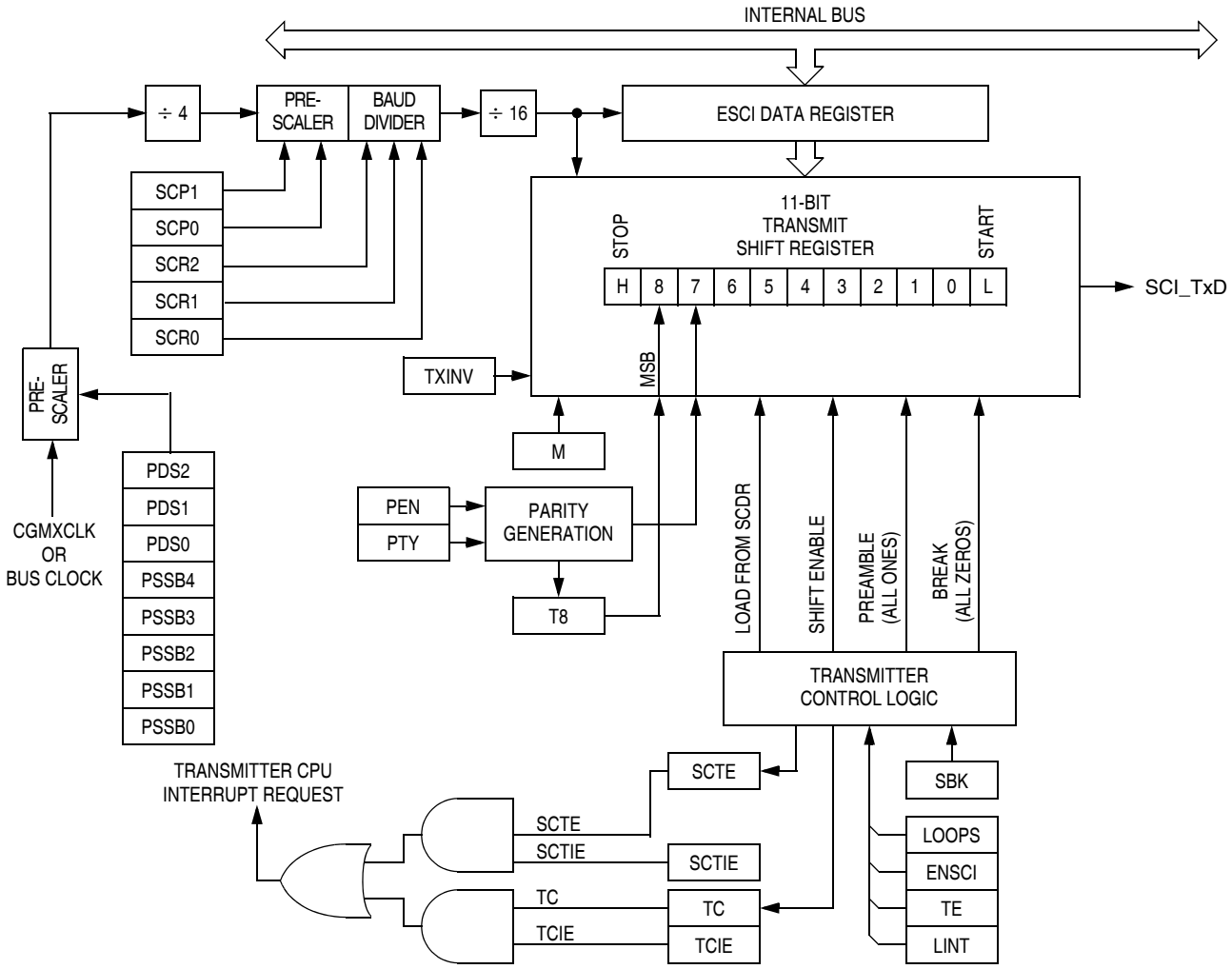


Figure 14-5. ESCI Transmitter

14.4.2.1 Character Length

The transmitter can accommodate either 8-bit or 9-bit data. The state of the M bit in ESCI control register 1 (SCC1) determines character length. When transmitting 9-bit data, bit T8 in ESCI control register 3 (SCC3) is the ninth bit (bit 8).

14.4.2.2 Character Transmission

During an ESCI transmission, the transmit shift register shifts a character out to the TxD pin. The ESCI data register (SCDR) is the write-only buffer between the internal data bus and the transmit shift register.

IDLE — Receiver Idle Bit

This clearable, read-only bit is set when 10 or 11 consecutive 1s appear on the receiver input. IDLE generates an ESCI receiver CPU interrupt request if the ILIE bit in SCC2 is also set. Clear the IDLE bit by reading SCS1 with IDLE set and then reading the SCDR. After the receiver is enabled, it must receive a valid character that sets the SCRF bit before an idle condition can set the IDLE bit. Also, after the IDLE bit has been cleared, a valid character must again set the SCRF bit before an idle condition can set the IDLE bit. Reset clears the IDLE bit.

1 = Receiver input idle
0 = Receiver input active (or idle since the IDLE bit was cleared)

OR — Receiver Overrun Bit

This clearable, read-only bit is set when software fails to read the SCDR before the receive shift register receives the next character. The OR bit generates an ESCI error CPU interrupt request if the ORIE bit in SCC3 is also set. The data in the shift register is lost, but the data already in the SCDR is not affected. Clear the OR bit by reading SCS1 with OR set and then reading the SCDR. Reset clears the OR bit.

1 = Receive shift register full and SCRF = 1
0 = No receiver overrun

Software latency may allow an overrun to occur between reads of SCS1 and SCDR in the flag-clearing sequence. Figure 14-14 shows the normal flag-clearing sequence and an example of an overrun caused by a delayed flag-clearing sequence. The delayed read of SCDR does not clear the OR bit because OR was not set when SCS1 was read. Byte 2 caused the overrun and is lost. The next flag-clearing sequence reads byte 3 in the SCDR instead of byte 2.

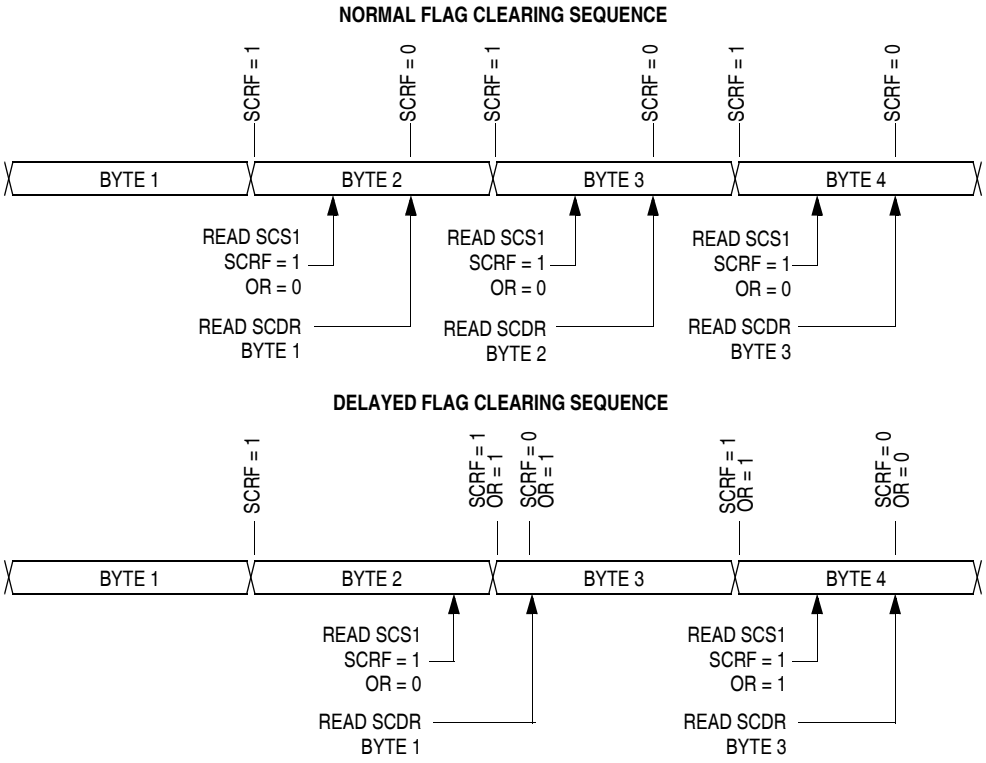
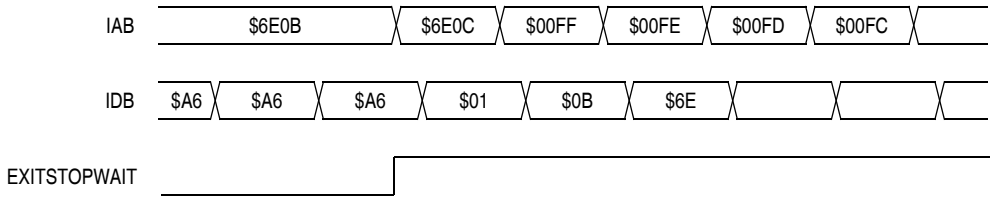


Figure 14-14. Flag Clearing Sequence

A module that is active during wait mode can wakeup the CPU with an interrupt if the interrupt is enabled. Stacking for the interrupt begins one cycle after the WAIT instruction during which the interrupt occurred. In wait mode, the CPU clocks are inactive. Refer to the wait mode subsection of each module to see if the module is active or inactive in wait mode. Some modules can be programmed to be active in wait mode.

Wait mode also can be exited by a reset or break. A break interrupt during wait mode sets the SIM break stop/wait bit, SBSW, in the SIM break status register (BSR). If the COP disable bit, COPD, in the CONFIG1 register is 0, then the computer operating properly module (COP) is enabled and remains active in wait mode.

Figure 15-17 and Figure 15-18 show the timing for WAIT recovery.



Note: EXITSTOPWAIT = $\overline{\text{RST}}$ pin, CPU interrupt, or break interrupt

Figure 15-17. Wait Recovery from Interrupt or Break

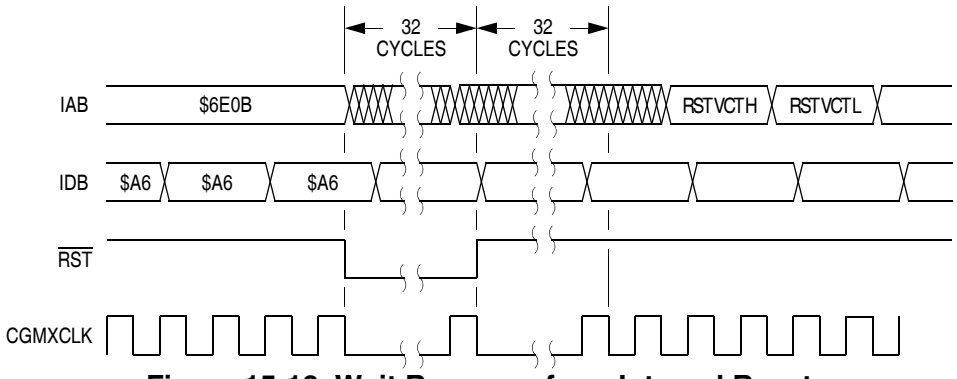


Figure 15-18. Wait Recovery from Internal Reset

15.6.2 Stop Mode

In stop mode, the SIM counter is reset and the system clocks are disabled. An interrupt request from a module can cause an exit from stop mode. Stacking for interrupts begins after the selected stop recovery time has elapsed. Reset also causes an exit from stop mode.

The SIM disables the clock generator module outputs (CGMOUT and CGMXCLK) in stop mode, stopping the CPU and peripherals. Stop recovery time is selectable using the SSREC bit in CONFIG1. If SSREC is set, stop recovery is reduced from the normal delay of 4096 CGMXCLK cycles down to 32. This is ideal for applications using canned oscillators that do not require long startup times from stop mode.

NOTE

External crystal applications should use the full stop recovery time by clearing the SSREC bit unless OSCENINSTOP bit is set in CONFIG2.

ELSxB and ELSxA — Edge/Level Select Bits

When channel x is an input capture channel, these read/write bits control the active edge-sensing logic on channel x.

When channel x is an output compare channel, ELSxB and ELSxA control the channel x output behavior when an output compare occurs.

When ELSxB and ELSxA are both clear, channel x is not connected to port D or port F, and pin PTDx/T2CHx or pin PTFx/T2CHx is available as a general- purpose I/O pin. Table 19-2 shows how ELSxB and ELSxA work. Reset clears the ELSxB and ELSxA bits.

NOTE

After initially enabling a TIM2 channel register for input capture operation and selecting the edge sensitivity, clear CHxF to ignore any erroneous edge detection flags.

TOVx — Toggle-On-Overflow Bit

When channel x is an output compare channel, this read/write bit controls the behavior of the channel x output when the TIM2 counter overflows. When channel x is an input capture channel, TOVx has no effect. Reset clears the TOVx bit.

1 = Channel x pin toggles on TIM2 counter overflow.
0 = Channel x pin does not toggle on TIM2 counter overflow.

NOTE

When TOVx is set, a TIM2 counter overflow takes precedence over a channel x output compare if both occur at the same time.

CHxMAX — Channel x Maximum Duty Cycle Bit

When the TOVx bit is at a 1 and clear output on compare is selected, setting the CHxMAX bit forces the duty cycle of buffered and unbuffered PWM signals to 100%. As Figure 19-9 shows, the CHxMAX bit takes effect in the cycle after it is set or cleared. The output stays at 100% duty cycle level until the cycle after CHxMAX is cleared.

NOTE

The 100% PWM duty cycle is defined as a continuous high level if the PWM polarity is 1 and a continuous low level if the PWM polarity is 0. Conversely, a 0% PWM duty cycle is defined as a continuous low level if the PWM polarity is 1 and a continuous high level if the PWM polarity is 0.

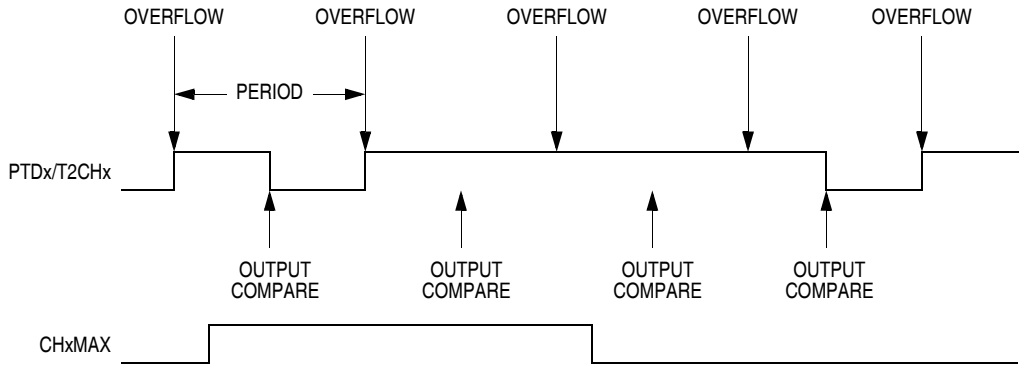
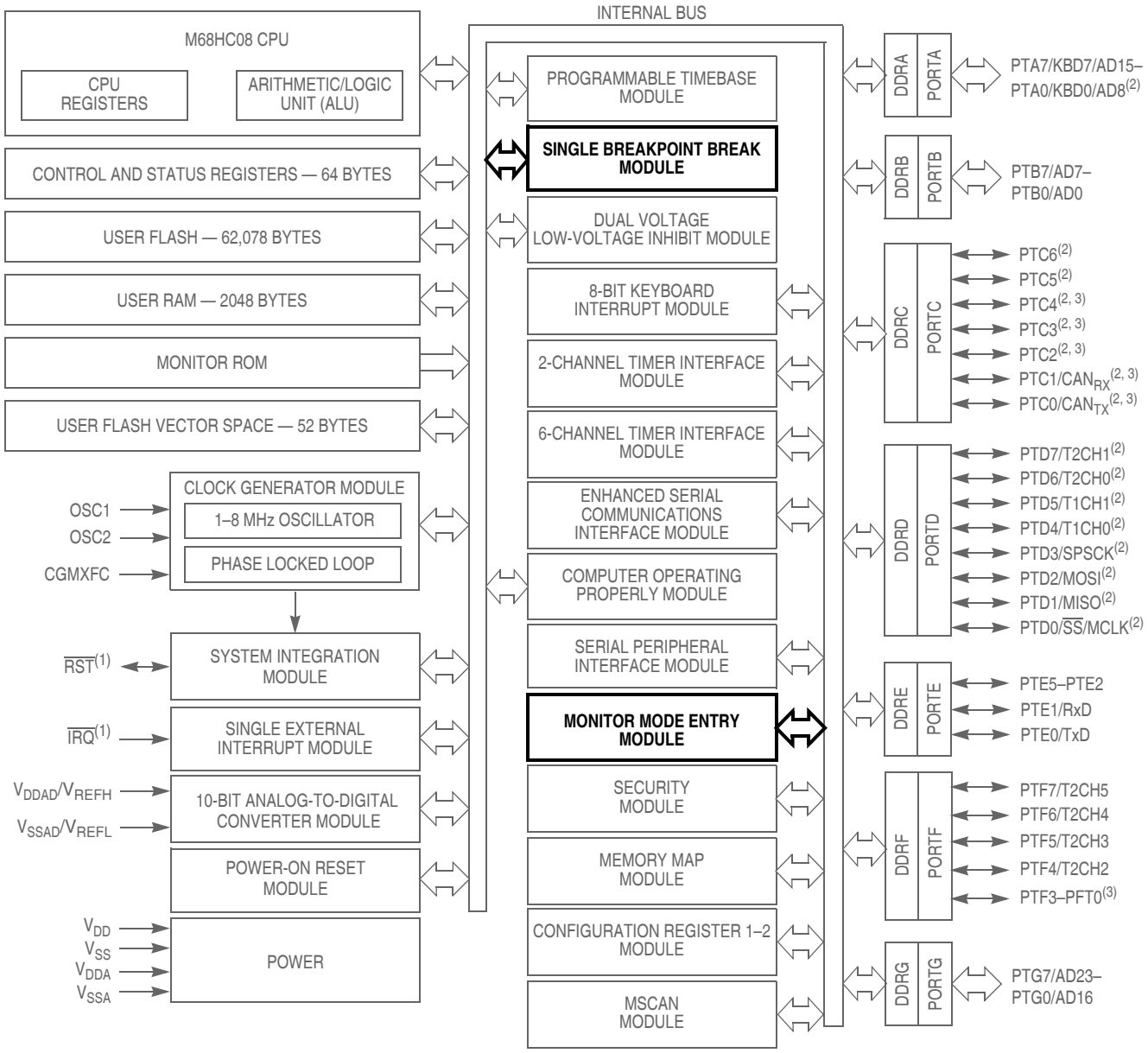


Figure 19-9. CHxMAX Latency



1. Pin contains integrated pullup device.
2. Ports are software configurable with pullup device if input port or pullup/pulldown device for keyboard input.
3. Higher current drive port pins

Figure 20-1. Block Diagram Highlighting BRK and MON Blocks

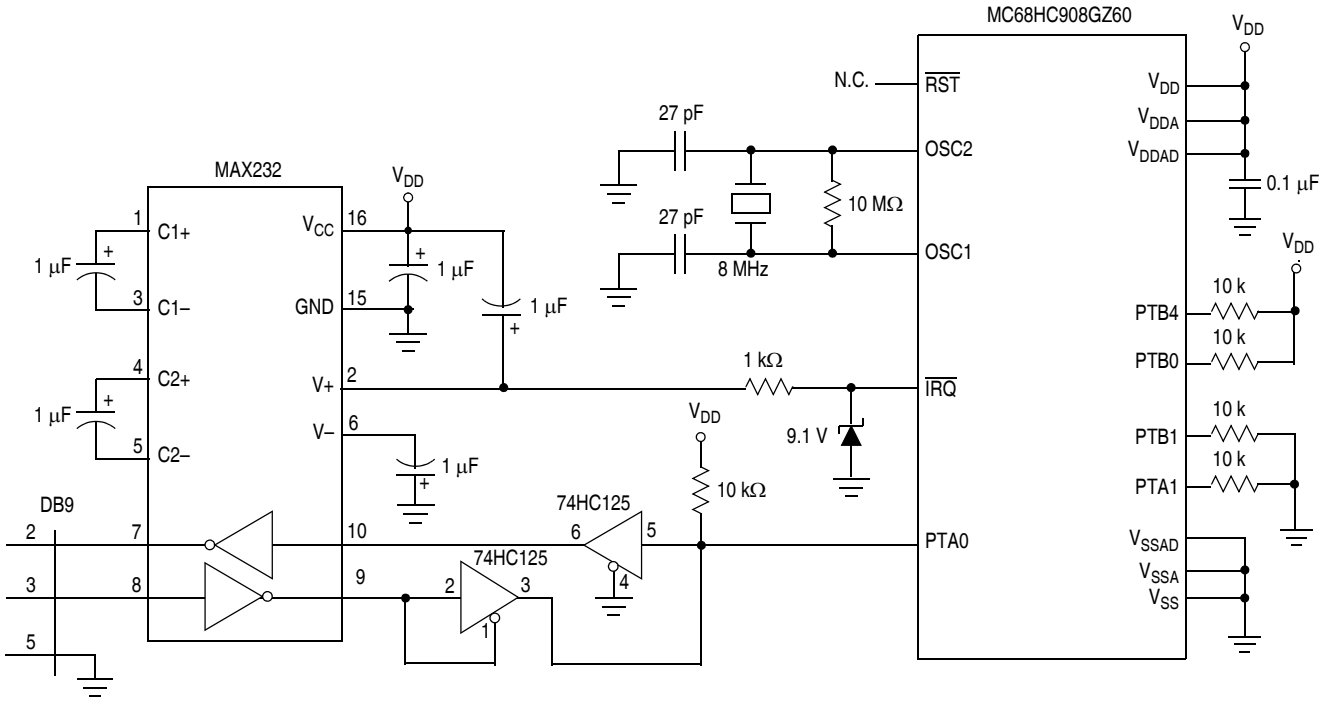


Figure 20-10. Normal Monitor Mode Circuit

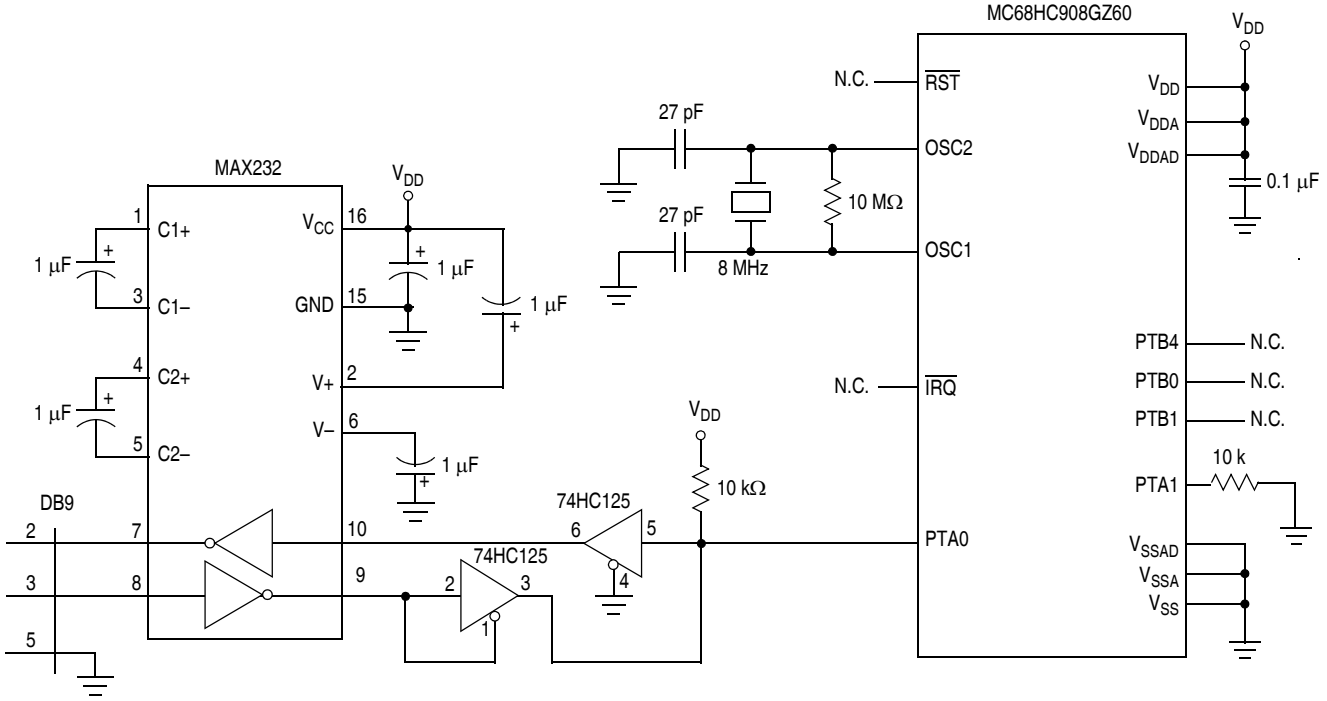


Figure 20-11. Forced Monitor Mode