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Applications of "<u>Embedded -</u> <u>Microcontrollers</u>"

Details

Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	40MHz
Connectivity	CANbus, I ² C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, LVD, POR, PWM, WDT
Number of I/O	33
Program Memory Size	32KB (16K x 16)
Program Memory Type	FLASH
EEPROM Size	256 x 8
RAM Size	1.5K x 8
Voltage - Supply (Vcc/Vdd)	2V ~ 5.5V
Data Converters	A/D 8x10b
Oscillator Type	External
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	44-TQFP
Supplier Device Package	44-TQFP (10x10)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18lf458-i-pt

Email: info@E-XFL.COM

Address: Room A, 16/F, Full Win Commercial Centre, 573 Nathan Road, Mongkok, Hong Kong

4.0 MEMORY ORGANIZATION

There are three memory blocks in Enhanced MCU devices. These memory blocks are:

- Enhanced Flash Program Memory
- Data Memory
- EEPROM Data Memory

Data and program memory use separate busses, which allows concurrent access of these blocks. Additional detailed information on data EEPROM and Flash program memory is provided in Section 5.0 "Data EEPROM Memory" and Section 6.0 "Flash Program Memory", respectively.

4.1 Program Memory Organization

The PIC18F258/458 devices have a 21-bit program counter that is capable of addressing a 2-Mbyte program memory space.

The Reset vector address is at 0000h and the interrupt vector addresses are at 0008h and 0018h.

FIGURE 4-1: PROGRAM MEMORY MAP AND STACK FOR PIC18F248/448



Figure 4-1 shows the diagram for program memory map and stack for the PIC18F248 and PIC18F448. Figure 4-2 shows the diagram for the program memory map and stack for the PIC18F258 and PIC18F458.

4.1.1 INTERNAL PROGRAM MEMORY OPERATION

The PIC18F258 and the PIC18F458 have 32 Kbytes of internal Enhanced Flash program memory. This means that the PIC18F258 and the PIC18F458 can store up to 16K of single-word instructions. The PIC18F248 and PIC18F448 have 16 Kbytes of Enhanced Flash program memory. This translates into 8192 single-word instructions, which can be stored in the program memory. Accessing a location between the physically implemented memory and the 2-Mbyte address will cause a read of all '0's (a NOP instruction).



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4.2 Return Address Stack

The return address stack allows any combination of up to 31 program calls and interrupts to occur. The PC (Program Counter) is pushed onto the stack when a PUSH, CALL or RCALL instruction is executed, or an interrupt is Acknowledged. The PC value is pulled off the stack on a RETURN, RETLW or a RETFIE instruction. PCLATU and PCLATH are not affected by any of the RETURN instructions.

The stack operates as a 31-word by 21-bit stack memory and a 5-bit Stack Pointer register, with the Stack Pointer initialized to 00000b after all Resets. There is no RAM associated with Stack Pointer 00000b. This is only a Reset value. During a CALL type instruction, causing a push onto the stack, the Stack Pointer is first incremented and the RAM location pointed to by the Stack Pointer is written with the contents of the PC. During a RETURN type instruction, causing a pop from the stack, the contents of the RAM location indicated by the STKPTR are transferred to the PC and then the Stack Pointer is decremented.

The stack space is not part of either program or data space. The Stack Pointer is readable and writable and the data on the top of the stack is readable and writable through SFR registers. Status bits indicate if the stack pointer is at or beyond the 31 levels provided.

4.2.1 TOP-OF-STACK ACCESS

The top of the stack is readable and writable. Three register locations, TOSU, TOSH and TOSL allow access to the contents of the stack location indicated by the STKPTR register. This allows users to implement a software stack, if necessary. After a CALL, RCALL or interrupt, the software can read the pushed value by reading the TOSU, TOSH and TOSL registers. These values can be placed on a user defined software stack. At return time, the software can replace the TOSU, TOSH and TOSL and do a return.

The user should disable the global interrupt enable bits during this time to prevent inadvertent stack operations.

4.2.2 RETURN STACK POINTER (STKPTR)

The STKPTR register contains the Stack Pointer value, the STKFUL (Stack Full) status bit and the STKUNF (Stack Underflow) status bits. Register 4-1 shows the STKPTR register. The value of the Stack Pointer can be 0 through 31. The Stack Pointer increments when values are pushed onto the stack and decrements when values are popped off the stack. At Reset, the Stack Pointer value will be '0'. The user may read and write the Stack Pointer value. This feature can be used by a Real-Time Operating System for return stack maintenance.

After the PC is pushed onto the stack 31 times (without popping any values off the stack), the STKFUL bit is set. The STKFUL bit can only be cleared in software or by a POR.

The action that takes place when the stack becomes full depends on the state of the STVREN (Stack Overflow Reset Enable) configuration bit. Refer to **Section 21.0 "Comparator Module**" for a description of the device configuration bits. If STVREN is set (default), the 31st push will push the (PC + 2) value onto the stack, set the STKFUL bit and reset the device. The STKFUL bit will remain set and the Stack Pointer will be set to '0'.

If STVREN is cleared, the STKFUL bit will be set on the 31st push and the Stack Pointer will increment to 31. The 32nd push will overwrite the 31st push (and so on), while STKPTR remains at 31.

When the stack has been popped enough times to unload the stack, the next pop will return a value of zero to the PC and sets the STKUNF bit, while the stack pointer remains at '0'. The STKUNF bit will remain set until cleared in software or a POR occurs.

Note: Returning a value of zero to the PC on an underflow has the effect of vectoring the program to the Reset vector, where the stack conditions can be verified and appropriate actions can be taken.

4.10 Access Bank

The Access Bank is an architectural enhancement that is very useful for C compiler code optimization. The techniques used by the C compiler are also useful for programs written in assembly.

This data memory region can be used for:

- · Intermediate computational values
- · Local variables of subroutines
- · Faster context saving/switching of variables
- · Common variables
- Faster evaluation/control of SFRs (no banking)

The Access Bank is comprised of the upper 160 bytes in Bank 15 (SFRs) and the lower 96 bytes in Bank 0. These two sections will be referred to as Access Bank High and Access Bank Low, respectively. Figure 4-6 indicates the Access Bank areas.

A bit in the instruction word specifies if the operation is to occur in the bank specified by the BSR register or in the Access Bank.

When forced in the Access Bank (a = 0), the last address in Access Bank Low is followed by the first address in Access Bank High. Access Bank High maps most of the Special Function Registers so that these registers can be accessed without any software overhead.

4.11 Bank Select Register (BSR)

The need for a large general purpose memory space dictates a RAM banking scheme. The data memory is partitioned into sixteen banks. When using direct addressing, the BSR should be configured for the desired bank.

BSR<3:0> holds the upper 4 bits of the 12-bit RAM address. The BSR<7:4> bits will always read '0's and writes will have no effect.

A MOVLB instruction has been provided in the instruction set to assist in selecting banks.

If the currently selected bank is not implemented, any read will return all '0's and all writes are ignored. The Status register bits will be set/cleared as appropriate for the instruction performed.

Each Bank extends up to FFh (256 bytes). All data memory is implemented as static RAM.

A MOVFF instruction ignores the BSR since the 12-bit addresses are embedded into the instruction word.

Section 4.12 "Indirect Addressing, INDF and FSR Registers" provides a description of indirect addressing, which allows linear addressing of the entire RAM space.



FIGURE 4-7: DIRECT ADDRESSING

WRITE_WORD_TO_	HREGS		
	MOVFW	POSTINCO, W	; get low byte of buffer data
	MOVWF	TABLAT	; present data to table latch
	TBLWT+*	*	; write data, perform a short write
			; to internal TBLWT holding register.
	DECFSZ	COUNTER	; loop until buffers are full
	BRA	WRITE_WORD_TO_HREGS	
PROGRAM_MEMORY			
_	BSF	EECON1, EEPGD	; point to FLASH program memory
	BCF	EECON1, CFGS	; access FLASH program memory
	BSF	EECON1, WREN	; enable write to memory
	BCF	INTCON, GIE	; disable interrupts
	MOVLW	55h	; write 55h
Required	MOVWF	EECON2	
Sequence	MOVLW	0AAh	; write OAAh
	MOVWF	EECON2	; start program (CPU stall)
	BSF	EECON1, WR	
	NOP		
	BSF	INTCON, GIE	; re-enable interrupts
	DECFSZ	COUNTER_HI	; loop until done
	BRA	PROGRAM LOOP	
	BCF	EECON1, WREN	; disable write to memory

EXAMPLE 6-3: WRITING TO FLASH PROGRAM MEMORY (CONTINUED)

6.5.2 WRITE VERIFY

Depending on the application, good programming practice may dictate that the value written to the memory should be verified against the original value. This should be used in applications where excessive writes can stress bits near the specification limit.

6.5.3 UNEXPECTED TERMINATION OF WRITE OPERATION

If a write is terminated by an unplanned event, such as loss of power or an unexpected Reset, the memory location just programmed should be verified and reprogrammed if needed. The WRERR bit is set when a write operation is interrupted by a MCLR Reset or a WDT Time-out Reset during normal operation. In these situations, users can check the WRERR bit and rewrite the location.

6.5.4 PROTECTION AGAINST SPURIOUS WRITES

To reduce the probability against spurious writes to Flash program memory, the write initiate sequence must also be followed. See **Section 24.0** "**Special Features of the CPU**" for more detail.

6.6 Flash Program Operation During Code Protection

See **Section 24.0** "**Special Features of the CPU**" for details on code protection of Flash program memory.

TABLE 6-2:	REGISTERS ASSOCIATED WITH PROGRAM FLASH MEMORY

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on: POR, BOR	Value on all other Resets
TBLPTRU	—		bit 21	Program M (TBLPTR<	lemory Tabl 20:16>)	00 0000	00 0000			
TBPLTRH	Program M	lemory Tab	le Pointer I	High Byte (1	FBLPTR<15	:8>)			0000 0000	0000 0000
TBLPTRL	Program M	lemory Tab	le Pointer L	ow Byte (T	BLPTR<7:0))			0000 0000	0000 0000
TABLAT	Program M	lemory Tab	le Latch						0000 0000	0000 0000
INTCON	GIE/GIEH	PEIE/ GIEL	TMR0IE	INTE	RBIE	TMR0IF	INTF	RBIF	0000 000x	0000 000u
EECON2	EEPROM	Control Reg	gister 2 (no	t a physical	register)				_	_
EECON1	EEPGD	CFGS		FREE	WRERR	WREN	WR	RD	xx-0 x000	uu-0 u000
IPR2	—	CMIP	_	EEIP	BCLIP	LVDIP	TMR3IP	ECCP1IP ⁽¹⁾	-1-1 1111	-1-1 1111
PIR2	_	CMIF	_	EEIF	BCLIF	LVDIF	TMR3IF	ECCP1IF ⁽¹⁾	-0-0 0000	-0-0 0000
PIE2	_	CMIE	_	EEIE	BCLIE	LVDIE	TMR3IE	ECCP1IE ⁽¹⁾	-0-0 0000	-0-0 0000

Legend: x = unknown, u = unchanged, - = unimplemented, read as '0'.

Shaded cells are not used during Flash/EEPROM access.

Note 1: These registers or register bits are not implemented on the PIC18F248 and PIC18F258 and read as '0's.

8.0 INTERRUPTS

The PIC18FXX8 devices have multiple interrupt sources and an interrupt priority feature that allows each interrupt source to be assigned a high priority level or a low priority level. The high priority interrupt vector is at 000008h and the low priority interrupt vector is at 000018h. High priority interrupt events will override any low priority interrupts that may be in progress.

There are 13 registers that are used to control interrupt operation. These registers are:

- RCON
- INTCON
- INTCON2
- INTCON3
- PIR1, PIR2, PIR3
- PIE1, PIE2, PIE3
- IPR1, IPR2, IPR3

It is recommended that the Microchip header files, supplied with MPLAB[®] IDE, be used for the symbolic bit names in these registers. This allows the assembler/ compiler to automatically take care of the placement of these bits within the specified register.

Each interrupt source has three bits to control its operation. The functions of these bits are:

- Flag bit to indicate that an interrupt event occurred
- Enable bit that allows program execution to branch to the interrupt vector address when the flag bit is set
- Priority bit to select high priority or low priority

The interrupt priority feature is enabled by setting the IPEN bit (RCON register). When interrupt priority is enabled, there are two bits that enable interrupts globally. Setting the GIEH bit (INTCON<7>) enables all interrupts. Setting the GIEL bit (INTCON register) enables all interrupts that have the priority bit cleared. When the interrupt flag, enable bit and appropriate global interrupt enable bit are set, the interrupt will vector immediately to address 000008h or 000018h, depending on the priority level. Individual interrupts can be disabled through their corresponding enable bits.

When the IPEN bit is cleared (default state), the interrupt priority feature is disabled and interrupts are compatible with PICmicro[®] mid-range devices. In Compatibility mode, the interrupt priority bits for each source have no effect. The PEIE bit (INTCON register) enables/disables all peripheral interrupt sources. The GIE bit (INTCON register) enables/disables all interrupt sources. All interrupts branch to address 00008h in Compatibility mode.

When an interrupt is responded to, the global interrupt enable bit is cleared to disable further interrupts. If the IPEN bit is cleared, this is the GIE bit. If interrupt priority levels are used, this will be either the GIEH or GIEL bit. High priority interrupt sources can interrupt a low priority interrupt.

The return address is pushed onto the stack and the PC is loaded with the interrupt vector address (000008h or 000018h). Once in the Interrupt Service Routine, the source(s) of the interrupt can be determined by polling the interrupt flag bits. The interrupt flag bits must be cleared in software before re-enabling interrupts to avoid recursive interrupts.

The "return from interrupt" instruction, RETFIE, exits the interrupt routine and sets the GIE bit (GIEH or GIEL if priority levels are used), which re-enables interrupts.

For external interrupt events, such as the INT pins or the PORTB input change interrupt, the interrupt latency will be three to four instruction cycles. The exact latency is the same for one or two-cycle instructions. Individual interrupt flag bits are set regardless of the status of their corresponding enable bit or the GIE bit.

Note: Do not use the MOVFF instruction to modify any of the interrupt control registers while **any** interrupt is enabled. Doing so may cause erratic microcontroller behavior.

REGISTER 8-5:	PIR2: PE	RIPHERAL	INTERRU	IPT REQU	EST (FLA	G) REGIS	TER 2				
	U-0	R/W-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0			
	—	CMIF ⁽¹⁾	_	EEIF	BCLIF	LVDIF	TMR3IF	ECCP1IF ⁽¹⁾			
	bit 7				•			bit 0			
bit 7	Unimplen	nented: Rea	d as '0'								
bit 6	CMIF: Co	mparator Inte	errupt Flag	oit(1)							
	1 = Comp 0 = Comp	arator input l arator input l	nas change nas not cha	d nged							
bit 5	Unimplen	nented: Rea	d as '0'								
bit 4	EEIF: EEF	PROM Write	Operation I	nterrupt Fla	g bit						
	1 = Write 0 = Write	operation is o operation is r	complete (m not complet	nust be clea e	red in softw	are)					
bit 3	BCLIF: Bus Collision Interrupt Flag bit										
	 1 = A bus collision occurred (must be cleared in software) 0 = No bus collision occurred 										
bit 2	LVDIF: Low-Voltage Detect Interrupt Flag bit										
	 1 = A low-voltage condition occurred (must be cleared in software) 0 = The device voltage is above the Low-Voltage Detect trip point 										
bit 1	TMR3IF:	TMR3 Overfl	ow Interrup	t Flag bit							
	1 = TMR3 0 = TMR3	register ove register did	rflowed (mu not overflov	ist be cleare v	ed in softwa	re)					
bit 0	ECCP1IF:	ECCP1 Inte	errupt Flag b	_{oit} (1)							
	<u>Capture mode:</u> 1 = A TMR1 (TMR3) register capture occurred (must be cleared in software) 0 = No TMR1 (TMR3) register capture occurred										
	<u>Compare</u> 1 = A TMF 0 = No TM	<u>mode:</u> R1 register co /IR1 register (ompare mat	ch occurred atch occurre	d (must be c ed	leared in so	ftware)				
	<u>PWM moo</u> Unused in	<u>de:</u> 1 this mode.									
	Note 1	This hit is	only availat	le on PIC1	8F4X8 devi	ces For PIC	18F2X8 de	wices this hit			

Note 1: This bit is only available on PIC18F4X8 devices. For PIC18F2X8 devices, this bit is unimplemented and reads as '0'.

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented b	oit, read as '0'
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

12.1 Timer1 Operation

Timer1 can operate in one of these modes:

- As a timer
- As a synchronous counter
- As an asynchronous counter

The operating mode is determined by the clock select bit, TMR1CS (T1CON register).

When TMR1CS is clear, Timer1 increments every instruction cycle. When TMR1CS is set, Timer1 increments on every rising edge of the external clock input or the Timer1 oscillator, if enabled.

When the Timer1 oscillator is enabled (T1OSCEN is set), the RC1/T1OSI and RC0/T1OSO/T1CKI pins become inputs. That is, the TRISC<1:0> value is ignored.

Timer1 also has an internal "Reset input". This Reset can be generated by the CCP module (**Section 15.1** "**CCP1 Module**").



FIGURE 12-2: TIMER1 BLOCK DIAGRAM: 16-BIT READ/WRITE MODE



FIGURE 12-1: TIMER1 BLOCK DIAGRAM

TABLE 12-2: REGISTERS ASSOCIATED WITH TIMER1 AS A TIMER/COUNTER

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Valu POR,	e on BOR	Valu all c Res	e on other sets
INTCON	GIE/GIEH	PEIE/GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF	0000	000x	0000	000u
PIR1	PSPIF ⁽¹⁾	ADIF	RCIF	TXIF	SSPIF	CCP1IF	TMR2IF	TMR1IF	0000	0000	0000	0000
PIE1	PSPIE ⁽¹⁾	ADIE	RCIE	TXIE	SSPIE	CCP1IE	TMR2IE	TMR1IE	0000	0000	0000	0000
IPR1	PSPIP ⁽¹⁾	ADIP	RCIP	TXIP	SSPIP	CCP1IP	TMR2IP	TMR1IP	1111	1111	1111	1111
TMR1L	Holding Register for the Least Significant Byte of the 16-bit TMR1 Register								xxxx	xxxx	uuuu	uuuu
TMR1H	Holding Register for the Most Significant Byte of the 16-bit TMR1 Register							xxxx	xxxx	uuuu	uuuu	
T1CON	RD16	_	T1CKPS1	T1CKPS0	T1OSCEN	T1SYNC	TMR1CS	TMR1ON	0-00	0000	u-uu	uuuu

 $\label{eq:logend: Legend: Legend: x = unknown, u = unchanged, - = unimplemented, read as `0`. Shaded cells are not used by the Timer1 module.$

Note 1: These registers or register bits are not implemented on the PIC18F248 and PIC18F258 and read as '0's.

15.2.3 SOFTWARE INTERRUPT

When the Capture mode is changed, a false capture interrupt may be generated. The user should keep bit CCP1IE (PIE registers) clear to avoid false interrupts and should clear the flag bit CCP1IF, following any such change in operating mode.

15.2.4 CCP1 PRESCALER

There are four prescaler settings specified by bits CCP1M3:CCP1M0. Whenever the CCP1 module is turned off, or the CCP1 module is not in Capture mode, the prescaler counter is cleared. This means that any Reset will clear the prescaler counter.

Switching from one capture prescaler to another may generate an interrupt. Also, the prescaler counter will not be cleared; therefore, the first capture may be from a non-zero prescaler. Example 15-1 shows the recommended method for switching between capture prescalers. This example also clears the prescaler counter and will not generate the "false" interrupt.

15.2.5 CAN MESSAGE TIME-STAMP

The CAN capture event occurs when a message is received in either of the receive buffers. The CAN module provides a rising edge to the CCP1 module to cause a capture event. This feature is provided to time-stamp the received CAN messages.

This feature is enabled by setting the CANCAP bit of the CAN I/O control register (CIOCON<4>). The message receive signal from the CAN module then takes the place of the events on RC2/CCP1.

EXAMPLE 15-1: CHANGING BETWEEN CAPTURE PRESCALERS

CLRF	CCP1CON, F	; Turn CCP module off
MOVLW	NEW_CAPT_PS	; Load WREG with the
		; new prescaler mode
		; value and CCP ON
MOVWF	CCP1CON	; Load CCP1CON with
		; this value





17.4.2 OPERATION

The MSSP module functions are enabled by setting MSSP Enable bit, SSPEN (SSPCON1<5>).

The SSPCON1 register allows control of the I^2C operation. Four mode selection bits (SSPCON1<3:0>) allow one of the following I^2C modes to be selected:

- I²C Master mode, clock = OSC/4 (SSPADD +1)
- I²C Slave mode (7-bit address)
- I²C Slave mode (10-bit address)
- I²C Slave mode (7-bit address) with Start and Stop bit interrupts enabled
- I²C Slave mode (10-bit address) with Start and Stop bit interrupts enabled
- I²C Firmware Controlled Master mode, slave is Idle

Selection of any I²C mode with the SSPEN bit set forces the SCL and SDA pins to be open-drain, provided these pins are programmed to inputs by setting the appropriate TRISC bits. To ensure proper operation of the module, pull-up resistors must be provided externally to the SCL and SDA pins.

17.4.3 SLAVE MODE

In Slave mode, the SCL and SDA pins must be configured as inputs (TRISC<4:3> set). The MSSP module will override the input state with the output data when required (slave-transmitter).

The I²C Slave mode hardware will always generate an interrupt on an address match. Through the mode select bits, the user can also choose to interrupt on Start and Stop bits.

When an address is matched, or the data transfer after an address match is received, the hardware automatically will generate the Acknowledge (\overline{ACK}) pulse and load the SSPBUF register with the received value currently in the SSPSR register.

Any combination of the following conditions will cause the MSSP module not to give this ACK pulse:

- The Buffer Full bit, BF (SSPSTAT<0>), was set before the transfer was received.
- The overflow bit, SSPOV (SSPCON1<6>), was set before the transfer was received.

In this case, the SSPSR register value is not loaded into the SSPBUF, but bit SSPIF (PIR1<3>) is set. The BF bit is cleared by reading the SSPBUF register, while bit SSPOV is cleared through software.

The SCL clock input must have a minimum high and low for proper operation. The high and low times of the I^2C specification, as well as the requirement of the MSSP module, are shown in timing parameter #100 and parameter #101.

17.4.3.1 Addressing

Once the MSSP module has been enabled, it waits for a Start condition to occur. Following the Start condition, the 8 bits are shifted into the SSPSR register. All incoming bits are sampled with the rising edge of the clock (SCL) line. The value of register SSPSR<7:1> is compared to the value of the SSPADD register. The address is compared on the falling edge of the eighth clock (SCL) pulse. If the addresses match and the BF and SSPOV bits are clear, the following events occur:

- 1. The SSPSR register value is loaded into the SSPBUF register.
- 2. The Buffer Full bit BF is set.
- 3. An ACK pulse is generated.
- MSSP Interrupt Flag bit, SSPIF (PIR1<3>), is set (interrupt is generated if enabled) on the falling edge of the ninth SCL pulse.

In 10-bit Address mode, two address bytes need to be received by the slave. The five Most Significant bits (MSbs) of the first address byte specify if this is a 10-bit address. Bit R/W (SSPSTAT<2>) must specify a write so the slave device will receive the second address byte. For a 10-bit address, the first byte would equal '11110 A9 A8 0', where 'A9' and 'A8' are the two MSbs of the address. The sequence of events for 10-bit address is as follows, with steps 7 through 9 for the slave-transmitter:

- 1. Receive first (high) byte of address (bits SSPIF, BF and bit UA (SSPSTAT<1>) are set).
- 2. Update the SSPADD register with second (low) byte of address (clears bit UA and releases the SCL line).
- 3. Read the SSPBUF register (clears bit BF) and clear flag bit SSPIF.
- 4. Receive second (low) byte of address (bits SSPIF, BF and UA are set).
- 5. Update the SSPADD register with the first (high) byte of address. If match releases SCL line, this will clear bit UA.
- 6. Read the SSPBUF register (clears bit BF) and clear flag bit SSPIF.
- 7. Receive Repeated Start condition.
- 8. Receive first (high) byte of address (bits SSPIF and BF are set).
- 9. Read the SSPBUF register (clears bit BF) and clear flag bit SSPIF.



17.4.14 SLEEP OPERATION

While in Sleep mode, the I^2C module can receive addresses or data and when an address match or complete byte transfer occurs, wake the processor from Sleep (if the MSSP interrupt is enabled).

17.4.15 EFFECT OF A RESET

A Reset disables the MSSP module and terminates the current transfer.

17.4.16 MULTI-MASTER MODE

In Multi-Master mode, the interrupt generation on the detection of the Start and Stop conditions allows the determination of when the bus is free. The Stop (P) and Start (S) bits are cleared from a Reset or when the MSSP module is disabled. Control of the I^2C bus may be taken when the P bit (SSPSTAT<4>) is set, or the bus is ldle, with both the S and P bits clear. When the bus is busy, enabling the SSP interrupt will generate the interrupt when the Stop condition occurs.

In multi-master operation, the SDA line must be monitored for arbitration to see if the signal level is the expected output level. This check is performed in hardware with the result placed in the BCLIF bit.

The states where arbitration can be lost are:

- Address Transfer
- Data Transfer
- A Start Condition
- A Repeated Start Condition
- An Acknowledge Condition

17.4.17 MULTI -MASTER COMMUNICATION, BUS COLLISION AND BUS ARBITRATION

Multi-Master mode support is achieved by bus arbitration. When the master outputs address/data bits onto the SDA pin, arbitration takes place when the master outputs a '1' on SDA by letting SDA float high and another master asserts a '0'. When the SCL pin floats high, data should be stable. If the expected data on SDA is a '1' and the data sampled on the SDA pin = 0, then a bus collision has taken place. The master will set the Bus Collision Interrupt Flag BCLIF and reset the I^2C port to its Idle state (Figure 17-25).

If a transmit was in progress when the bus collision occurred, the transmission is halted, the BF flag is cleared, the SDA and SCL lines are deasserted and the SSPBUF can be written to. When the user services the bus collision Interrupt Service Routine and if the l^2C bus is free, the user can resume communication by asserting a Start condition.

If a Start, Repeated Start, Stop or Acknowledge condition was in progress when the bus collision occurred, the condition is aborted, the SDA and SCL lines are deasserted and the respective control bits in the SSPCON2 register are cleared. When the user services the bus collision Interrupt Service Routine and if the l^2C bus is free, the user can resume communication by asserting a Start condition.

The master will continue to monitor the SDA and SCL pins. If a Stop condition occurs, the SSPIF bit will be set.

A write to the SSPBUF will start the transmission of data at the first data bit regardless of where the transmitter left off when the bus collision occurred.

In Multi-Master mode, the interrupt generation on the detection of Start and Stop conditions allows the determination of when the bus is free. Control of the I^2C bus can be taken when the P bit is set in the SSPSTAT register or the bus is Idle and the S and P bits are cleared.

FIGURE 17-25: BUS COLLISION TIMING FOR TRANSMIT AND ACKNOWLEDGE





bit 7 bit 6

bit 5-3 bit 2-0

REGISTER 19-31: BRGCON3: BAUD RATE CONTROL REGISTER 3

••••	Diracon	IO. DAOD						
	U-0	R/W-0	U-0	U-0	U-0	R/W-0	R/W-0	R/W-0
		WAKFIL	—	—	_	SEG2PH2 ⁽¹⁾	SEG2PH1 ⁽¹⁾	SEG2PH0 ⁽¹⁾
	bit 7							bit 0
	Unimpler	nented: Re	ad as '0'					
	WAKFIL:	Selects CA	N bus Line	e Filter for V	Wake-up bi	it		
	1 = Use C	AN bus lin	e filter for v	vake-up				
	0 = CAN t	ous line filte	er is not us	ed for wake	e-up			
	Unimpler	nented: Re	ead as '0'					
	SEG2PH2	2:SEG2PH	0: Phase S	egment 2	Time Selec	t bits ⁽¹⁾		
	111 = Pha	ase Segme	nt 2 Time =	= 8 x Tq				

		•••g•=		-	~	• ~
110 =	Phase	Segment 2	Time =	7	х	ΤQ

- 101 = Phase Segment 2 Time = 6 x TQ
- 100 = Phase Segment 2 Time = 5 x TQ
- 011 = Phase Segment 2 Time = 4 x TQ
- 010 = Phase Segment 2 Time = 3 x TQ
- 001 = Phase Segment 2 Time = 2 x TQ
- 000 = Phase Segment 2 Time = 1 x TQ

Note 1: Ignored if SEG2PHTS bit (BRGCON2<7>) is clear.

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented	bit, read as '0'
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

GENERAL FORMAT FOR INSTRUCTIONS FIGURE 25-1: Byte-oriented file register operations **Example Instruction** 10 9 8 7 15 0 OPCODE f (FILE #) ADDWF MYREG, W, B d а d = 0 for result destination to be WREG register d = 1 for result destination to be file register (f) a = 0 to force Access Bank a = 1 for BSR to select bank f = 8-bit file register address Byte to Byte move operations (2-word) 15 12 11 0 OPCODE f (Source FILE #) MOVFF MYREG1, MYREG2 15 12 11 0 f (Destination FILE #) 1111 f = 12-bit file register address Bit-oriented file register operations 15 12 11 987 0 OPCODE b (BIT #) f (FILE #) а BSF MYREG, bit, B b = 3-bit position of bit in file register (f) a = 0 to force Access Bank a = 1 for BSR to select bank f = 8-bit file register address Literal operations 15 8 7 0 OPCODE k (literal) MOVLW 0x7F k = 8-bit immediate value **Control** operations CALL, GOTO and Branch operations 15 8 7 0 OPCODE n<7:0> (literal) GOTO Label 15 12 11 0 1111 n<19:8> (literal) n = 20-bit immediate value 15 8 7 0 CALL MYFUNC OPCODE S n<7:0> (literal) 15 12 11 0 n<19:8> (literal) S = Fast bit 15 11 10 0 OPCODE BRA MYFUNC n<10:0> (literal) 15 8 7 0 OPCODE n<7:0> (literal) BC MYFUNC

INCF	SZ	Increment	f, Skip if 0						
Synta	ax:	[label] IN	ICFSZ f[,d [,a]]					
Oper	ands:	$0 \le f \le 255$ $d \in [0,1]$ $a \in [0,1]$	$0 \le f \le 255$ $d \in [0,1]$ $a \in [0,1]$						
Oper	ation:	(f) + 1 \rightarrow de skip if result	(f) + 1 → dest, skip if result = 0						
Statu	s Affected:	None							
Enco	ding:	0011	11da :	fff	ffff				
Desc	nption:	I he content incremented placed in W placed back If the result which is alre and a NOP is it a two-cycl Access Ban overriding th the bank will BSR value of	s of registe d. If 'd' is '0 . If 'd' is '1' a in register is '0', the n eady fetche s executed e instructio k will be se ne BSR vali I be selecto (default).	r T are , the res 'f' (defa ext instead instead n. If 'a' elected, ue. If 'a' ed as pe	sult is sult is ault). ruction carded I, making is '0', the = 1, then er the				
Word	ls:	1	. ,						
Cycle	es:	1(2) Note: 3 cy by a	cles if skip 2-word ins	and foll	owed				
QC	vcle Activity:								
	Q1	Q2	Q3		Q4				
	Decode	Read	Process	, N	/rite to				
16 - 1-		register t	Data	des	stination				
IT SK	ID: 01	02	03		04				
	No	No	No		No				
	operation	operation	operation	n op	eration				
lf sk	ip and followe	d by 2-word in	struction:						
	Q1	Q2	Q3		Q4				
	No	No	No		No				
	operation	operation	operatior	п ор	eration				
	No operation	No operation	No operatior	n op	No eration				
Exan	nple:	HERE I NZERO ZERO	INCFSZ	CNT					
	Before Instruc PC	tion = Address	6 (HERE)						
	After Instructio CNT If CNT	on = CNT + 1 = 0;	1						
	If CNT	$=$ Address \neq 0;	s (ZERO)						
	PC	= Address	s (NZERO)						

INFSNZ			Increment f, Skip if not 0							
Syntax:	[<i>la</i>	[<i>label</i>] INFSNZ f[,d[,a]]								
Operands:			$0 \le f \le 255$ $d \in [0,1]$ $a \in [0,1]$							
Operation:		(f) ski	(f) + 1 \rightarrow dest, skip if result \neq 0							
Status A	ffected:	No	None							
Encodin	ig:		0100 10da ffff ff:							
Description:		Th inc pla pla lf t ins dis ins wil val sel	The contents of register 'f' are incremented. If 'd' is '0', the result is placed in W. If 'd' is '1', the result is placed back in register 'f' (default). If the result is not '0', the next instruction which is already fetched is discarded and a NOP is executed instead, making it a two-cycle instruction. If 'a' is '0', the Access Bank will be selected, overriding the BSR value. If 'a' = 1, then the bank will be selected as per the BSR value (default).							
Words:		1								
Cycles:			1(2) Note: 3 cycles if skip and followed by a 2-word instruction.							
Q Cycle	e Activity:									
	Q1	1	Q2	Q	3		Q4			
	Decode	F reg	lead ister 'f'	Process Data		W des	rite to stination			
lf skip:										
	Q1		Q2	Q	3		Q4			
	No		No	No			No			
	peration	ope	eration	opera	tion	ор	eration			
If skip and followed		u by i				04				
	No	No.		No			No.			
0	peration	ope	eration	operation		ор	eration			
	No		No	No			No			
0	peration	оре	eration	operation		ор	eration			
Example:		HE ZE NZ	RE RO IERO	NFSNZ REG						
Before Instruction PC = Address (HERE)										
Aft	er Instructio REG If REG PC If REG PC	on ≠ = = =	REG + 0; Address 0; Address	1 s (nzem s (zero	RO) D)					

27.1 DC Characteristics (Continued)

PIC18LFXX8 (Industrial)			Standard Operating Conditions (unless otherwise stated)Operating temperature $-40^{\circ}C \le TA \le +85^{\circ}C$ for industrial							
PIC18FXX8 (Industrial, Extended)			$\begin{array}{llllllllllllllllllllllllllllllllllll$							
Param No.	Symbol	Characteristic/ Device	Min Typ Max Units Conditions			Conditions				
	Δ IWDT	Module Differential Current								
D022		Watchdog Timer PIC18LFXX8		0.75 0.8 7	1.5 8 25	μΑ μΑ μΑ	VDD = 2.5V, +25°C VDD = 2.0V, -40°C to +85°C VDD = 4.2V, -40°C to +85°C			
D022		Watchdog Timer PIC18FXX8		7 7 7	25 25 45	μΑ μΑ μΑ	VDD = 4.2V, +25°C VDD = 4.2V, -40°C to +85°C VDD = 4.2V, -40°C to +125°C			
D022A	ΔIBOR	Brown-out Reset ⁽⁵⁾ PIC18LFXX8		38 42 49	50 55 65	μΑ μΑ μΑ	VDD = 2.0V, +25°C VDD = 2.0V, -40°C to +85°C VDD = 4.2V, -40°C to +85°C			
D022A		Brown-out Reset ⁽⁵⁾ PIC18FXX8		46 49 50	65 65 75	μΑ μΑ μΑ	VDD = 4.2V, +25°C VDD = 4.2V, -40°C to +85°C VDD = 4.2V, -40°C to +125°C			
D022B	ΔILVD	Low-Voltage Detect ⁽⁵⁾ PIC18LFXX8	 	36 40 47	50 55 65	μΑ μΑ μΑ	VDD = 2.0V, +25°C VDD = 2.0V, -40°C to +85°C VDD = 4.2V, -40°C to +85°C			
D022B		Low-Voltage Detect ⁽⁵⁾ PIC18FXX8		44 47 47	65 65 75	μΑ μΑ μΑ	VDD = 4.2V, +25°C VDD = 4.2V, -40°C to +85°C VDD = 4.2V, -40°C to +125°C			
D025	ΔITMR1	Timer1 Oscillator PIC18LFXX8	- - -	6.2 6.2 7.5	40 45 55	μΑ μΑ μΑ	VDD = 2.0V, +25°C VDD = 2.0V, -40°C to +85°C VDD = 4.2V, -40°C to +85°C			
D025		Timer1 Oscillator PIC18FXX8		7.5 7.5 7.5	55 55 65	μΑ μΑ μΑ	VDD = 4.2V, +25°C VDD = 4.2V, -40°C to +85°C VDD = 4.2V, -40°C to +125°C			

Legend: Rows are shaded for improved readability.

Note 1: This is the limit to which VDD can be lowered in Sleep mode, or during a device Reset, without losing RAM data.

2: The supply current is mainly a function of the operating voltage and frequency. Other factors, such as I/O pin loading and switching rate, oscillator type, internal code execution pattern and temperature, also have an impact on the current consumption.

The test conditions for all IDD measurements in active operation mode are: $\frac{OSC1}{MCLR} = \text{external square wave, from rail-to-rail; all I/O pins tri-stated, pulled to VDD}$ $\frac{MCLR}{MCLR} = \text{VDD; WDT enabled/disabled as specified.}$

3: The power-down current in Sleep mode does not depend on the oscillator type. Power-down current is measured with the part in Sleep mode, with all I/O pins in high-impedance state and tied to VDD and VSS and all features that add delta current disabled (such as WDT, Timer1 Oscillator, BOR, ...).

- 4: For RC oscillator configuration, current through REXT is not included. The current through the resistor can be estimated by the formula Ir = VDD/2 REXT (mA) with REXT in kOhm.
- 5: The LVD and BOR modules share a large portion of circuitry. The △IBOR and △ILVD currents are not additive. Once one of these modules is enabled, the other may also be enabled without further penalty.

27.2 DC Characteristics: PIC18FXX8 (Industrial, Extended) PIC18LFXX8 (Industrial) (Continued)

DC CHARACTERISTICS			$\begin{array}{l} \mbox{Standard Operating Conditions (unless otherwise stated)} \\ \mbox{Operating temperature} & -40^{\circ}C \leq TA \leq +85^{\circ}C \mbox{ for industrial} \\ -40^{\circ}C \leq TA \leq +125^{\circ}C \mbox{ for extended} \end{array}$				
Param No.	Symbol	Characteristic/ Device	Min Max Unit		Units	Conditions	
	Vol	Output Low Voltage					
D080		I/O ports	—	0.6	V	IOL = 8.5 mA, VDD = 4.2V, -40°C to +85°C	
D080A			—	0.6	V	IOL = 7.0 mA, VDD = 4.2V, -40°C to +125°C	
D083		OSC2/CLKO (RC mode)	—	0.6	V	IOL = 1.6 mA, VDD = 4.2V, -40°C to +85°C	
D083A			—	0.6	V	IOL = 1.2 mA, VDD = 4.2V, -40°C to +125°C	
	Vон	Output High Voltage ⁽³⁾					
D090		I/O ports	Vdd - 0.7	—	V	IOH = -3.0 mA, VDD = 4.2V, -40°C to +85°C	
D090A			Vdd - 0.7	—	V	IOH = -2.5 mA, VDD = 4.2V, -40°С to +125°С	
D092		OSC2/CLKO (RC mode)	Vdd - 0.7	—	V	IOH = -1.3 mA, VDD = 4.2V, -40°С to +85°С	
D092A			Vdd - 0.7	—	V	IOH = -1.0 mA, VDD = 4.2V, -40°С to +125°С	
D150	Vod	Open-Drain High Voltage	—	7.5	V	RA4 pin	
		Capacitive Loading Specs on Output Pins					
D101	Сю	All I/O pins and OSC2 (in RC mode)	—	50	pF	To meet the AC Timing Specifications	
D102	Св	SCL, SDA	—	400	pF	In I ² C™ mode	

Note 1: In RC oscillator configuration, the OSC1/CLKI pin is a Schmitt Trigger input. It is not recommended that the PICmicro[®] device be driven with an external clock while in RC mode.

2: The leakage current on the MCLR pin is strongly dependent on the applied voltage level. The specified levels represent normal operating conditions. Higher leakage current may be measured at different input voltages.

3: Negative current is defined as current sourced by the pin.

BSF	3 4
BTFSS	4
BTG	5
BZ	6
C	
C Compilers	
MPLAB C17	4
MPLAB C18	4
MPLAB C30	4
CALL	6
CAN Module	-
Aborting Transmission	8
Acknowledge Error 23	7
Baud Rate Registers	8
Baud Bate Setting 23	3
Bit Frror 23	7
Bit Time Partitioning (diagram)	3
Bit Timing Configuration Begisters 23	6
BBGCON1 23	6
BRGCON2 23	6
BRGCON3 23	6
Calculating To, Nominal Bit Bate and	0
Nominal Bit Time	л
Configuration Mode	4 6
Control and Status Pagistars 20	1
Controller Register Man	5
	7
Diachla Mada	6
Disable Mode	0 7
Error Medee and Error Ocurators	7
Error Modes and Error Counters	/
Error Modes State (diagram)	8
Error Recognition Mode	7
Error States	1
Fliter/Mask Truth (table)	2
Form Error	7
Hard Synchronization	5
I/O Control Register	1
Information Processing Time	4
Initiating Transmission	8
Internal Message Reception	
Flowchart	1
Internal Transmit Message	_
Flowchart	9
Interrupt Acknowledge	9
Interrupt Registers	2
Interrupts23	8
Bus Activity Wake-up23	9
Bus-Off23	9
Code Bits23	8
Error	9
Message Error23	9
Receive	8
Receiver Bus Passive	9
Receiver Overflow23	9
Receiver Warning23	9
Transmit23	8
Transmitter Bus Passive23	9
Transmitter Warning23	9
Lengthening a Bit Period	
(diagram)23	5

Listen Only Mode	226
Loopback Mode	227
Message Acceptance Filters	
and Masks 215,	232
Message Acceptance Mask and	
Filter Operation (diagram)	232
Message Reception	230
Message Time-Stamping	230
Message Transmission	227
Modes of Operation	226
Normal Mode	226
Oscillator Tolerance	236
Overview	100
Phase Buffer Segments	234
Programming Time Segments	236
Propagation Segment	200
Propagation Segment	234
Receive Buller Registers	210
Receive bullers	230
Receive Message Buffering	230
Receive Priority	230
Registers	201
Resynchronization	235
Sample Point	234
Shortening a Bit Period (diagram)	236
Stuff Bit Error	237
Synchronization	235
Synchronization Rules	235
Synchronization Segment	234
Time Quanta	234
Transmit Buffer Registers	206
Transmit Buffers	227
Transmit Priority	227
Transmit/Receive Buffers	199
Values for ICODE (table)	239
Capture (CCP Module)	124
CAN Message Time-Stamp	125
CCP Pin Configuration	124
CCP1 Prescaler	125
CCPR1H:CCPR1L Registers	124
Software Interrupt	125
Timer1/Timer3 Mode Selection	124
Capture (ECCP Module)	133
CAN Message Time-Stamp	133
Capture/Compare/PWM (CCP)	123
Capture Mode See Capture	
(CCP Module)	
CCP1 Module	124
Timer Besources	124
CCPR1H Register	124
CCPR1L Begister	124
Compare Mode See Compare	124
(CCP Module)	
(UUF Module).	
	104
EUUF I WUUUIES	124
E vvivi ivioue. See E vvivi (CCE Module).	
Denance Testad	47
Clocking Schome	17
	41
	297
	297