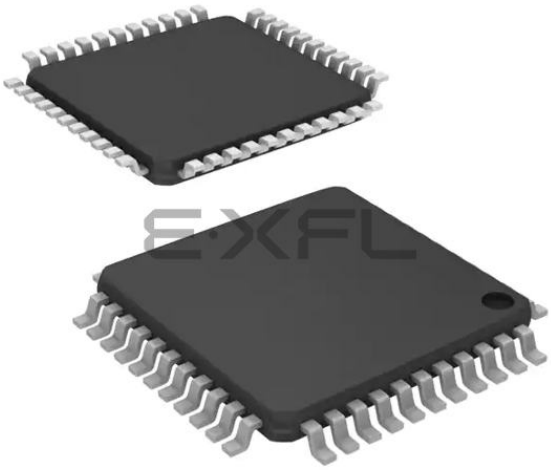

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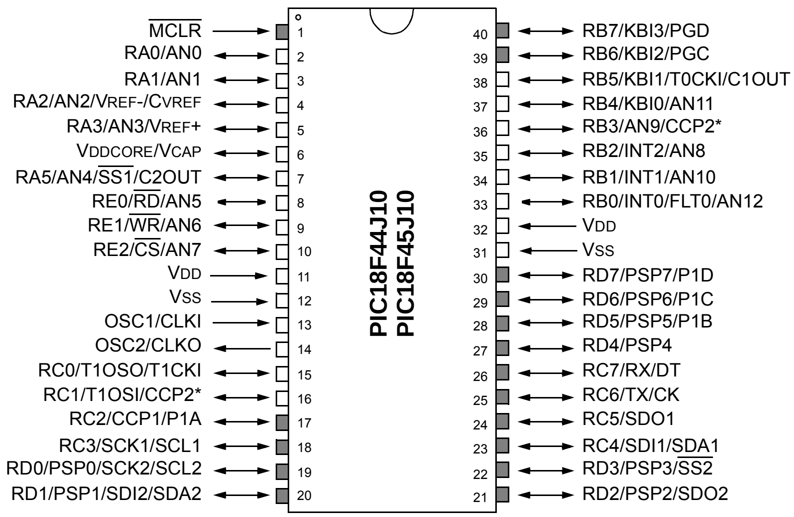
Applications of "Embedded - Microcontrollers"

Details	
Product Status	Obsolete
Core Processor	PIC
Core Size	8-Bit
Speed	40MHz
Connectivity	I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	32
Program Memory Size	32KB (16K x 16)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	1K x 8
Voltage - Supply (Vcc/Vdd)	2.7V ~ 3.6V
Data Converters	A/D 13x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 125°C (TA)
Mounting Type	Surface Mount
Package / Case	44-TQFP
Supplier Device Package	44-TQFP (10x10)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18f45j10t-e-pt

Pin Diagrams (Continued)

40-Pin PDIP (600 MIL)

■ = Pins are up to 5.5V tolerant

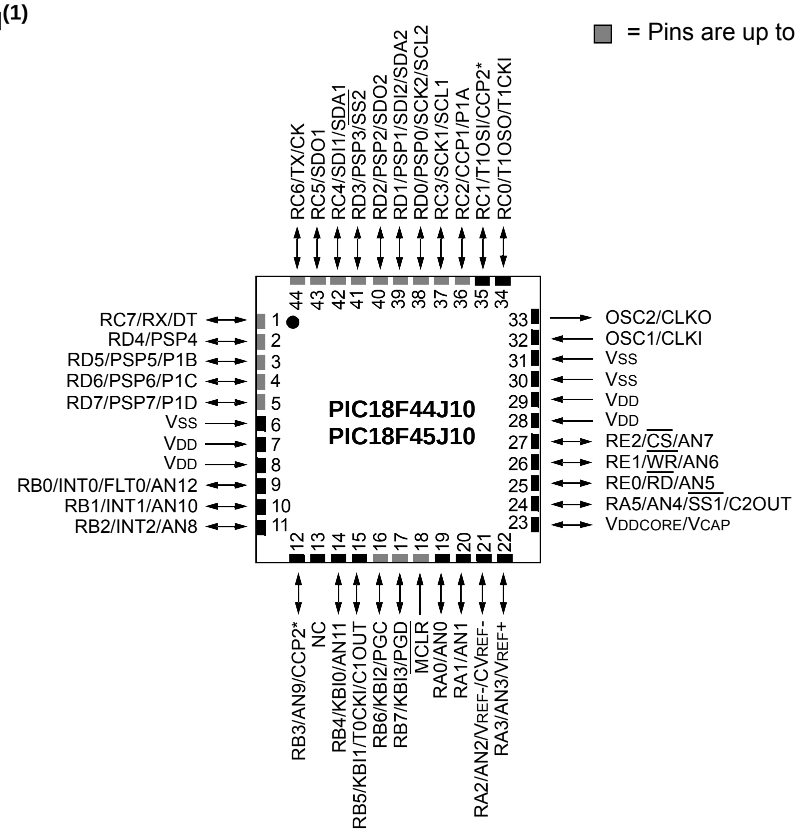


* Pin feature is dependent on device configuration.

.

44-Pin QFN(1)

■ = Pins are up to 5.5V tolerant



* Pin feature is dependent on device configuration.

Note 1: For the QFN package, it is recommended that the bottom pad be connected to VSS.

4.0 POWER-MANAGED MODES

The PIC18F45J10 family devices provide the ability to manage power consumption by simply managing clocking to the CPU and the peripherals. In general, a lower clock frequency and a reduction in the number of circuits being clocked constitutes lower consumed power. For the sake of managing power in an application, there are three primary modes of operation:

- Run mode
- Idle mode
- Sleep mode

These modes define which portions of the device are clocked and at what speed. The Run and Idle modes may use any of the three available clock sources (primary, secondary or internal oscillator block); the Sleep mode does not use a clock source.

The power-managed modes include several power-saving features offered on previous PIC® microcontrollers. One is the clock switching feature, offered in other PIC18 devices, allowing the controller to use the Timer1 oscillator in place of the primary oscillator. Also included is the Sleep mode, offered by all PIC microcontrollers, where all device clocks are stopped.

4.1 Selecting Power-Managed Modes

Selecting a power-managed mode requires two decisions: if the CPU is to be clocked or not and which clock source is to be used. The IDLEN bit (OSCCON<7>) controls CPU clocking, while the SCS<1:0> bits (OSCCON<1:0>) select the clock source. The individual modes, bit settings, clock sources and affected modules are summarized in Table 4-1.

4.1.1 CLOCK SOURCES

The SCS<1:0> bits allow the selection of one of three clock sources for power-managed modes. They are:

- the primary clock, as defined by the FOSC<1:0> Configuration bits
- the secondary clock (Timer1 oscillator)
- the internal oscillator

4.1.2 ENTERING POWER-MANAGED MODES

Switching from one power-managed mode to another begins by loading the OSCCON register. The SCS<1:0> bits select the clock source and determine which Run or Idle mode is to be used. Changing these bits causes an immediate switch to the new clock source, assuming that it is running. The switch may also be subject to clock transition delays. These are discussed in **Section 4.1.3 "Clock Transitions and Status Indicators"** and subsequent sections.

Entry to the power-managed Idle or Sleep modes is triggered by the execution of a SLEEP instruction. The actual mode that results depends on the status of the IDLEN bit.

Depending on the current mode and the mode being switched to, a change to a power-managed mode does not always require setting all of these bits. Many transitions may be done by changing the oscillator select bits, or changing the IDLEN bit, prior to issuing a SLEEP instruction. If the IDLEN bit is already configured correctly, it may only be necessary to perform a SLEEP instruction to switch to the desired mode.

TABLE 4-1: POWER-MANAGED MODES

Mode	OSCCON bits		Module Clocking		Available Clock and Oscillator Source
	IDLEN<7>(1)	SCS<1:0>	CPU	Peripherals	
Sleep	0	N/A	Off	Off	None – All clocks are disabled
PRI_RUN	N/A	10	Clocked	Clocked	Primary – HS, EC; this is the normal full-power execution mode
SEC_RUN	N/A	01	Clocked	Clocked	Secondary – Timer1 Oscillator
RC_RUN	N/A	11	Clocked	Clocked	Internal Oscillator
PRI_IDLE	1	10	Off	Clocked	Primary – HS, EC
SEC_IDLE	1	01	Off	Clocked	Secondary – Timer1 Oscillator
RC_IDLE	1	11	Off	Clocked	Internal Oscillator

Note 1: IDLEN reflects its value when the SLEEP instruction is executed.

FIGURE 5-4: TIME-OUT SEQUENCE ON POWER-UP ($\overline{MCLR}$ NOT TIED TO VDD): CASE 1

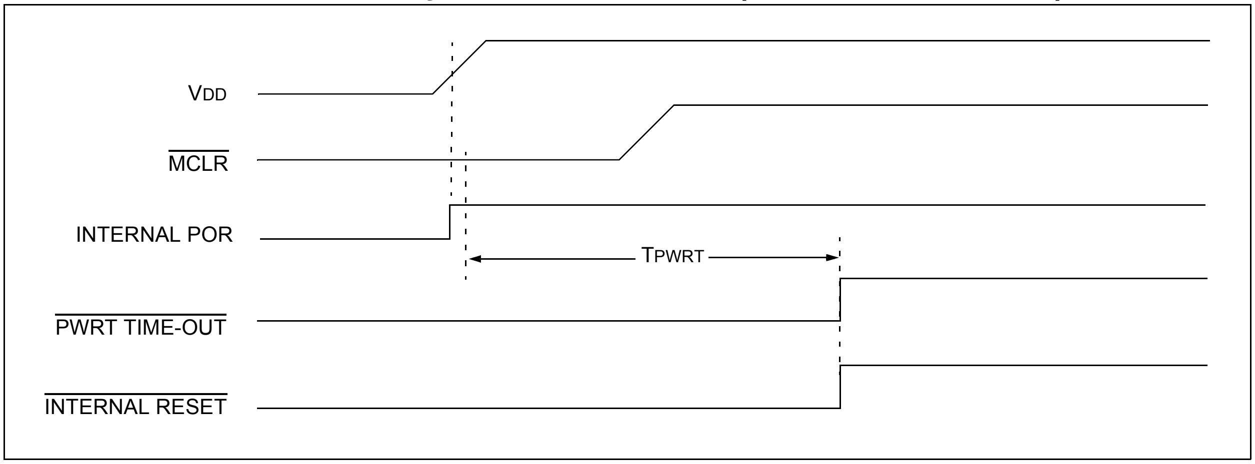


FIGURE 5-5: TIME-OUT SEQUENCE ON POWER-UP ($\overline{MCLR}$ NOT TIED TO VDD): CASE 2

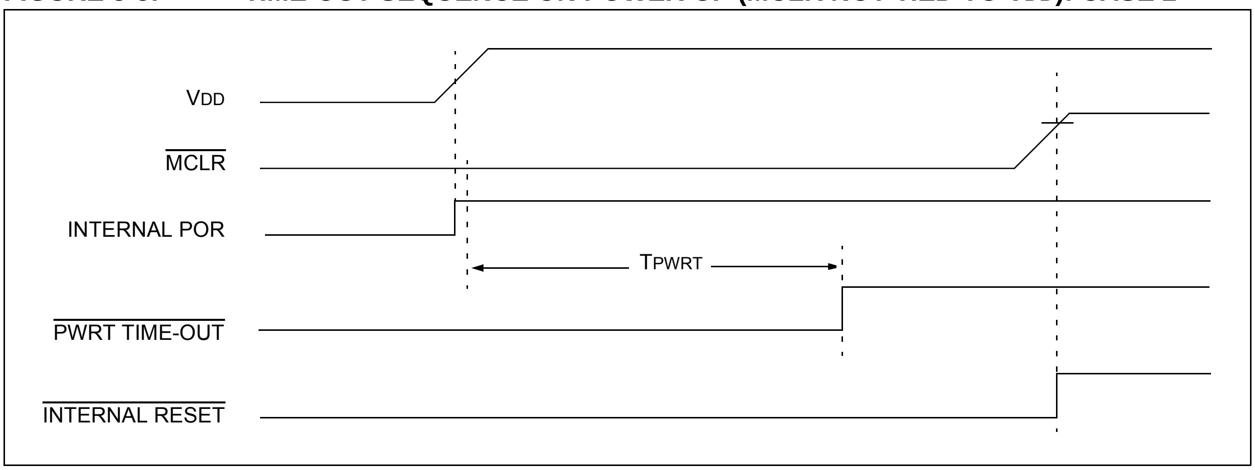
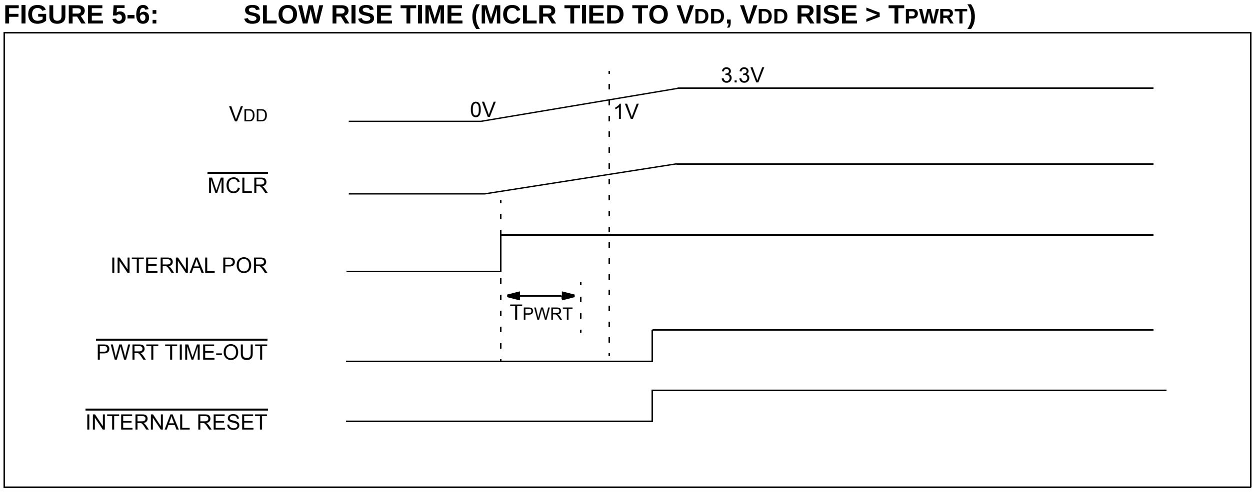


FIGURE 5-6: SLOW RISE TIME ($\overline{MCLR}$ TIED TO VDD, VDD RISE > TPWRT)



6.1.4.2 Return Stack Pointer (STKPTR)

The STKPTR register (Register 6-1) contains the Stack Pointer value, the STKFUL (Stack Overflow) status bit and the STKUNF (Stack Underflow) status bits. The value of the Stack Pointer can be 0 through 31. The Stack Pointer increments before values are pushed onto the stack and decrements after values are popped off the stack. On Reset, the Stack Pointer value will be zero. The user may read and write the Stack Pointer value. This feature can be used by a Real-Time Operating System (RTOS) for return stack maintenance.

After the PC is pushed onto the stack 31 times (without popping any values off the stack), the STKFUL bit is set. The STKFUL bit is cleared by software or by a POR.

The action that takes place when the stack becomes full depends on the state of the STVREN (Stack Overflow Reset Enable) Configuration bit. (Refer to **Section 21.1 "Configuration Bits"** for a description of the device Configuration bits.) If STVREN is set (default), the 31st push will push the (PC + 2) value onto the stack, set the STKFUL bit and reset the device. The STKFUL bit will remain set and the Stack Pointer will be set to zero.

If STVREN is cleared, the STKFUL bit will be set on the 31st push and the Stack Pointer will increment to 31. Any additional pushes will not overwrite the 31st push and the STKPTR will remain at 31.

When the stack has been popped enough times to unload the stack, the next pop will return a value of zero to the PC and sets the STKUNF bit, while the Stack Pointer remains at zero. The STKUNF bit will remain set until cleared by software or until a POR occurs.

Note:	Returning a value of zero to the PC on an underflow has the effect of vectoring the program to the Reset vector, where the stack conditions can be verified and appropriate actions can be taken. This is not the same as a Reset, as the contents of the SFRs are not affected.

6.1.4.3 `PUSH` and `POP` Instructions

Since the Top-of-Stack is readable and writable, the ability to push values onto the stack and pull values off the stack without disturbing normal program execution is a desirable feature. The PIC18 instruction set includes two instructions, `PUSH` and `POP`, that permit the TOS to be manipulated under software control. TOSU, TOSH and TOSL can be modified to place data or a return address on the stack.

The `PUSH` instruction places the current PC value onto the stack. This increments the Stack Pointer and loads the current PC value onto the stack.

The `POP` instruction discards the current TOS by decrementing the Stack Pointer. The previous value pushed onto the stack then becomes the TOS value.

REGISTER 6-1: STKPTR: STACK POINTER REGISTER

R/C-0	R/C-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
STKFUL[(1)]	STKUNF[(1)]	—	SP4	SP3	SP2	SP1	SP0
bit 7							bit 0

Legend:	C = Clearable bit		
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 7 **STKFUL:** Stack Full Flag bit[(1)]
`1` = Stack became full or overflowed
`0` = Stack has not become full or overflowed

bit 6 **STKUNF:** Stack Underflow Flag bit[(1)]
`1` = Stack underflow occurred
`0` = Stack underflow did not occur

bit 5 **Unimplemented:** Read as '`0`'

bit 4-0 **SP<4:0>:** Stack Pointer Location bits

Note 1: Bit 7 and bit 6 are cleared by user software or by a POR.

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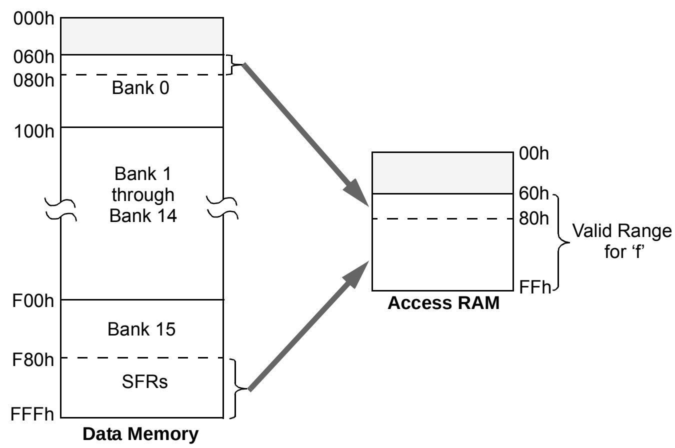
FIGURE 6-9: COMPARING ADDRESSING OPTIONS FOR BIT-ORIENTED AND BYTE-ORIENTED INSTRUCTIONS (EXTENDED INSTRUCTION SET ENABLED)

Example Instruction: `ADDWF, f, d, a` (Opcode: `0010 01da ffff ffff`)

When 'a' = `0` and f ≥ 60h:

The instruction executes in Direct Forced mode. 'f' is interpreted as a location in the Access RAM between 060h and 0FFh. This is the same as locations 060h to 07Fh (Bank 0) and F80h to FFFh (Bank 15) of data memory.

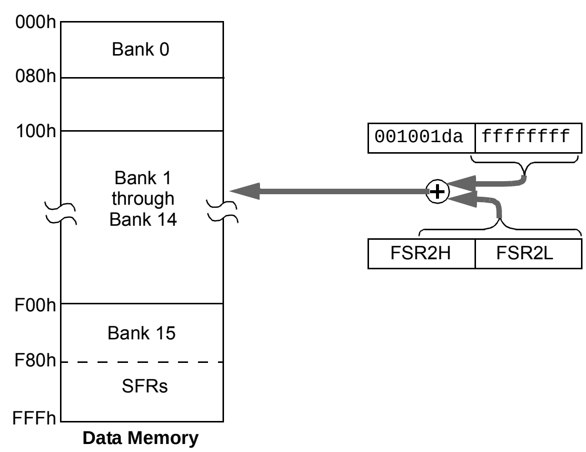
Locations below 60h are not available in this addressing mode.



When 'a' = `0` and f ≤ 5Fh:

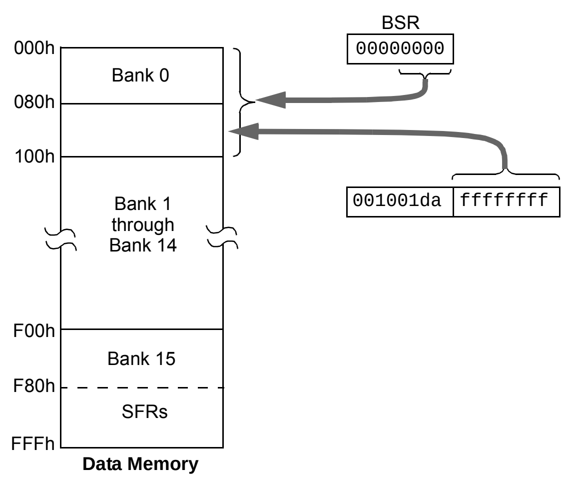
The instruction executes in Indexed Literal Offset mode. 'f' is interpreted as an offset to the address value in FSR2. The two are added together to obtain the address of the target register for the instruction. The address can be anywhere in the data memory space.

Note that in this mode, the correct syntax is now:
`ADDWF [k], d`
where 'k' is the same as 'f'.



When 'a' = `1` (all values of f):

The instruction executes in Direct mode (also known as Direct Long mode). 'f' is interpreted as a location in one of the 16 banks of the data memory space. The bank is designated by the Bank Select Register (BSR). The address can be in any implemented bank in the data memory space.



9.2 PIR Registers

The PIR registers contain the individual flag bits for the peripheral interrupts. Due to the number of peripheral interrupt sources, there are three Peripheral Interrupt Request (Flag) registers (PIR1, PIR2, PIR3).

> **Note 1:** Interrupt flag bits are set when an interrupt condition occurs regardless of the state of its corresponding enable bit or the Global Interrupt Enable bit, GIE (INTCON<7>).
>
> **2:** User software should ensure the appropriate interrupt flag bits are cleared prior to enabling an interrupt and after servicing that interrupt.

REGISTER 9-4: PIR1: PERIPHERAL INTERRUPT REQUEST (FLAG) REGISTER 1

R/W-0	R/W-0	R-0	R-0	R/W-0	R/W-0	R/W-0	R/W-0
PSPIF(1)	ADIF	RCIF	TXIF	SSP1IF	CCP1IF	TMR2IF	TMR1IF
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as ‘0’	
-n = Value at POR	‘1’ = Bit is set	‘0’ = Bit is cleared	x = Bit is unknown

bit 7 **PSPIF:** Parallel Slave Port Read/Write Interrupt Flag bit(1)
1 = A read or a write operation has taken place (must be cleared in software)
0 = No read or write has occurred

bit 6 **ADIF:** A/D Converter Interrupt Flag bit
1 = An A/D conversion completed (must be cleared in software)
0 = The A/D conversion is not complete

bit 5 **RCIF:** EUSART Receive Interrupt Flag bit
1 = The EUSART receive buffer, RCREG, is full (cleared when RCREG is read)
0 = The EUSART receive buffer is empty

bit 4 **TXIF:** EUSART Transmit Interrupt Flag bit
1 = The EUSART transmit buffer, TXREG, is empty (cleared when TXREG is written)
0 = The EUSART transmit buffer is full

bit 3 **SSP1IF:** Master Synchronous Serial Port 1 Interrupt Flag bit
1 = The transmission/reception is complete (must be cleared in software)
0 = Waiting to transmit/receive

bit 2 **CCP1IF:** ECCP1/CCP1 Interrupt Flag bit
<u>Capture mode:</u>
1 = A TMR1 register capture occurred (must be cleared in software)
0 = No TMR1 register capture occurred
<u>Compare mode:</u>
1 = A TMR1 register compare match occurred (must be cleared in software)
0 = No TMR1 register compare match occurred
<u>PWM mode:</u>
Unused in this mode.

bit 1 **TMR2IF:** TMR2 to PR2 Match Interrupt Flag bit
1 = TMR2 to PR2 match occurred (must be cleared in software)
0 = No TMR2 to PR2 match occurred

bit 0 **TMR1IF:** TMR1 Overflow Interrupt Flag bit
1 = TMR1 register overflowed (must be cleared in software)
0 = TMR1 register did not overflow

Note 1: This bit is not implemented on 28-pin devices and should be read as ‘0’.

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REGISTER 9-5: PIR2: PERIPHERAL INTERRUPT REQUEST (FLAG) REGISTER 2

R/W-0	R/W-0	U-0	U-0	R/W-0	U-0	U-0	R/W-0
OSCFIF	CMIF	—	—	BCLIF	—	—	CCP2IF
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as ‘0’	
-n = Value at POR	‘1’ = Bit is set	‘0’ = Bit is cleared	x = Bit is unknown

bit 7 **OSCFIF:** Oscillator Fail Interrupt Flag bit
1 = Device oscillator failed, clock input has changed to INTOSC (must be cleared in software)
0 = Device clock operating

bit 6 **CMIF:** Comparator Interrupt Flag bit
1 = Comparator input has changed (must be cleared in software)
0 = Comparator input has not changed

bit 5-4 **Unimplemented:** Read as ‘0’

bit 3 **BCLIF:** Bus Collision Interrupt Flag bit (MSSP1 module)
1 = A bus collision occurred (must be cleared in software)
0 = No bus collision occurred

bit 2-1 **Unimplemented:** Read as ‘0’

bit 0 **CCP2IF:** CCP2 Interrupt Flag bit
Capture mode:
1 = A TMR1 register capture occurred (must be cleared in software)
0 = No TMR1 register capture occurred
Compare mode:
1 = A TMR1 register compare match occurred (must be cleared in software)
0 = No TMR1 register compare match occurred
PWM mode:
Unused in this mode.

REGISTER 9-6: PIR3: PERIPHERAL INTERRUPT REQUEST (FLAG) REGISTER 3

R/W-0	R/W-0	U-0	U-0	U-0	U-0	U-0	U-0
SSP2IF	BCL2IF	—	—	—	—	—	—
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as ‘0’	
-n = Value at POR	‘1’ = Bit is set	‘0’ = Bit is cleared	x = Bit is unknown

bit 7 **SSP2IF:** Master Synchronous Serial Port 2 Interrupt Flag bit
1 = The transmission/reception is complete (must be cleared in software)
0 = Waiting to transmit/receive

bit 6 **BCL2IF:** Bus Collision Interrupt Flag bit (MSSP2 module)
1 = A bus collision occurred (must be cleared in software)
0 = No bus collision occurred

bit 5-0 **Unimplemented:** Read as ‘0’

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EXAMPLE 12-1: IMPLEMENTING A REAL-TIME CLOCK USING A TIMER1 INTERRUPT SERVICE

```
RTCinit
            MOVLW     80h               ; Preload TMR1 register pair
            MOVWF     TMR1H             ; for 1 second overflow
            CLRF      TMR1L
            MOVLW     b'00001111'       ; Configure for external clock,
            MOVWF     T1CON             ; Asynchronous operation, external oscillator
            CLRF      secs              ; Initialize timekeeping registers
            CLRF      mins              ;
            MOVLW     .12
            MOVWF     hours
            BSF       PIE1, TMR1IE      ; Enable Timer1 interrupt
            RETURN
RTCisr
            BSF       TMR1H, 7          ; Preload for 1 sec overflow
            BCF       PIR1, TMR1IF      ; Clear interrupt flag
            INCF      secs, F           ; Increment seconds
            MOVLW     .59               ; 60 seconds elapsed?
            CPFSGT    secs
            RETURN                      ; No, done
            CLRF      secs              ; Clear seconds
            INCF      mins, F           ; Increment minutes
            MOVLW     .59               ; 60 minutes elapsed?
            CPFSGT    mins
            RETURN                      ; No, done
            CLRF      mins              ; clear minutes
            INCF      hours, F          ; Increment hours
            MOVLW     .23               ; 24 hours elapsed?
            CPFSGT    hours
            RETURN                      ; No, done
            CLRF      hours             ; Reset hours
            RETURN                      ; Done
```

TABLE 12-2: REGISTERS ASSOCIATED WITH TIMER1 AS A TIMER/COUNTER

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Reset Values on page
INTCON	GIE/GIEH	PEIE/GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF	47
PIR1	PSPIF(1)	ADIF	RCIF	TXIF	SSP1IF	CCP1IF	TMR2IF	TMR1IF	49
PIE1	PSPIE(1)	ADIE	RCIE	TXIE	SSP1IE	CCP1IE	TMR2IE	TMR1IE	49
IPR1	PSPIP(1)	ADIP	RCIP	TXIP	SSP1IP	CCP1IP	TMR2IP	TMR1IP	49
TMR1L	Timer1 Register Low Byte								48
TMR1H	Timer1 Register High Byte								48
T1CON	RD16	T1RUN	T1CKPS1	T1CKPS0	T1OSCEN	$\overline{\text{T1SYNC}}$	TMR1CS	TMR1ON	48

Legend: Shaded cells are not used by the Timer1 module.

Note 1: These bits are not implemented on 28-pin devices and should be read as '0'.

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The CCPRxH register and a 2-bit internal latch are used to double-buffer the PWM duty cycle. This double-buffering is essential for glitchless PWM operation.

When the CCPRxH and 2-bit latch match TMR2, concatenated with an internal 2-bit Q clock or 2 bits of the TMR2 prescaler, the CCPx pin is cleared.

The maximum PWM resolution (bits) for a given PWM frequency is given by the equation:

EQUATION 14-3:

$$\text{PWM Resolution (max)} = \frac{\log\left(\frac{F_{OSC}}{F_{PWM}}\right)}{\log(2)} \text{ bits}$$

> **Note:** If the PWM duty cycle value is longer than the PWM period, the CCP2 pin will not be cleared.

TABLE 14-4: EXAMPLE PWM FREQUENCIES AND RESOLUTIONS AT 40 MHz

PWM Frequency	2.44 kHz	9.77 kHz	39.06 kHz	156.25 kHz	312.50 kHz	416.67 kHz
Timer Prescaler (1, 4, 16)	16	4	1	1	1	1
PR2 Value	FFh	FFh	FFh	3Fh	1Fh	17h
Maximum Resolution (bits)	10	10	10	8	7	6.58

14.4.3 PWM AUTO-SHUTDOWN (CCP1 ONLY)

The PWM auto-shutdown features of the Enhanced CCP module are also available to CCP1 in 28-pin devices. The operation of this feature is discussed in detail in **Section 15.4.7 "Enhanced PWM Auto-Shutdown"**.

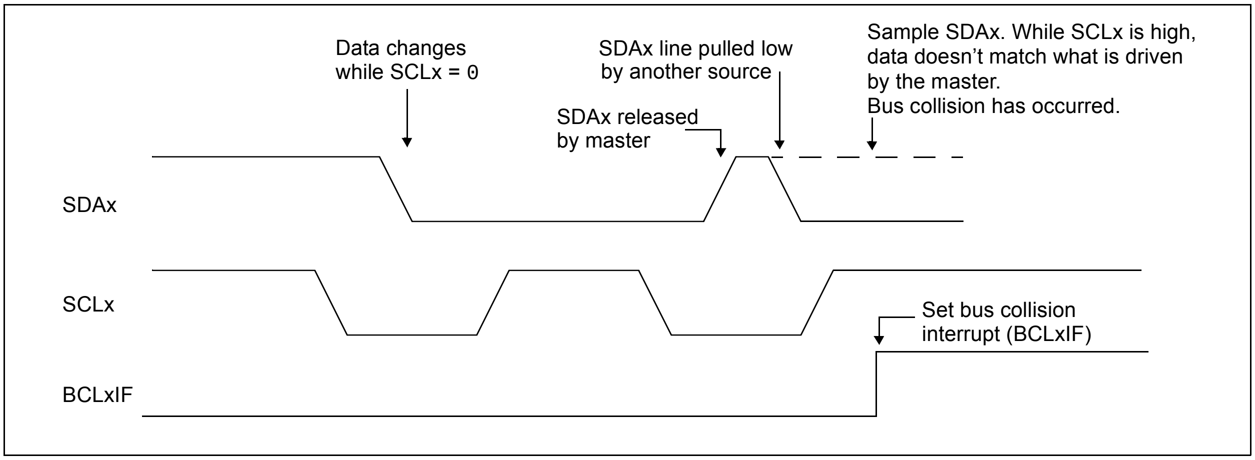
Auto-shutdown features are not available for CCP2.

14.4.4 SETUP FOR PWM OPERATION

The following steps should be taken when configuring the CCP module for PWM operation:

1. Set the PWM period by writing to the PR2 register.
2. Set the PWM duty cycle by writing to the CCPRxL register and CCPxCON<5:4> bits.
3. Make the CCPx pin an output by clearing the appropriate TRIS bit.
4. Set the TMR2 prescale value, then enable Timer2 by writing to T2CON.
5. Configure the CCPx module for PWM operation.

FIGURE 16-25: BUS COLLISION TIMING FOR TRANSMIT AND ACKNOWLEDGE



16.4.17.3 Bus Collision During a Stop Condition

Bus collision occurs during a Stop condition if:

a) After the SDAx pin has been deasserted and allowed to float high, SDAx is sampled low after the BRG has timed out.
b) After the SCLx pin is deasserted, SCLx is sampled low before SDAx goes high.

The Stop condition begins with SDAx asserted low. When SDAx is sampled low, the SCLx pin is allowed to float. When the pin is sampled high (clock arbitration), the Baud Rate Generator is loaded with SSPxADD<6:0> and counts down to 0. After the BRG times out, SDAx is sampled. If SDAx is sampled low, a bus collision has occurred. This is due to another master attempting to drive a data '0' (Figure 16-31). If the SCLx pin is sampled low before SDAx is allowed to float high, a bus collision occurs. This is another case of another master attempting to drive a data '0' (Figure 16-32).

FIGURE 16-31: BUS COLLISION DURING A STOP CONDITION (CASE 1)

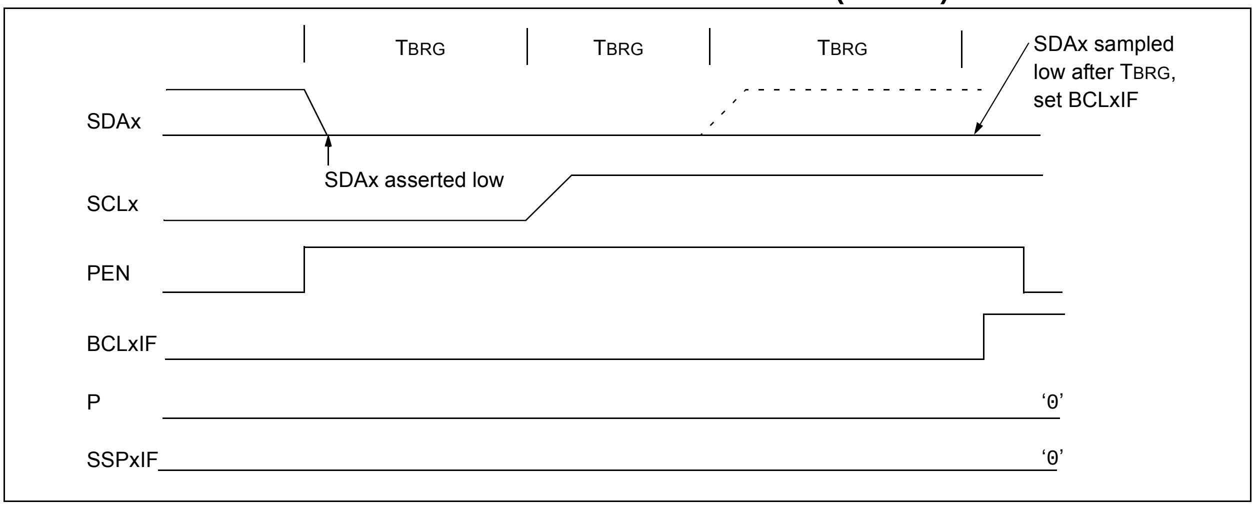
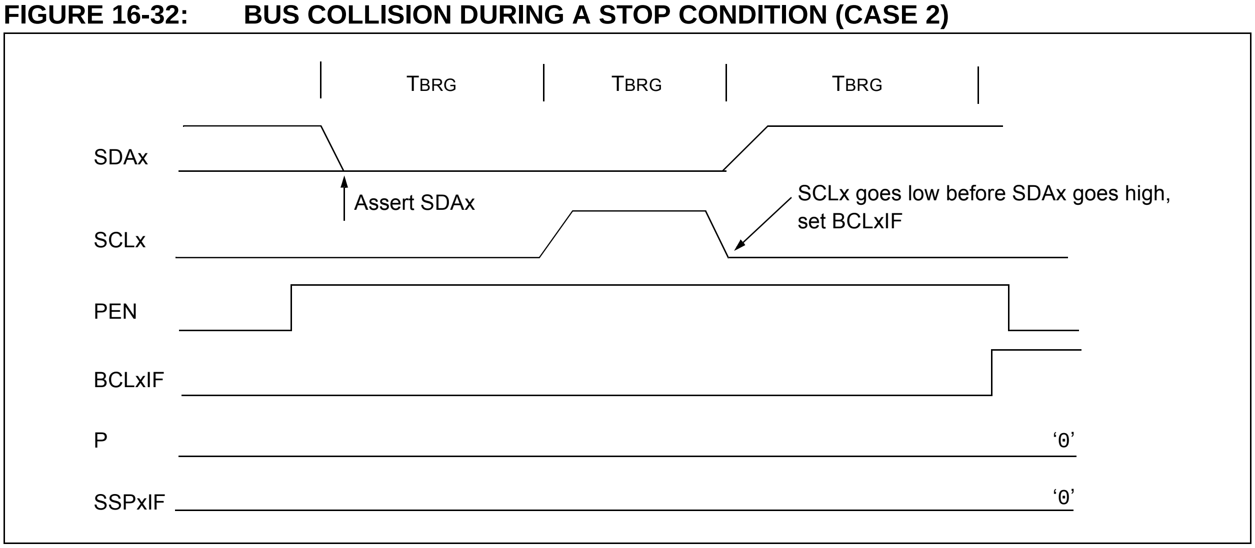


FIGURE 16-32: BUS COLLISION DURING A STOP CONDITION (CASE 2)



NOTES:

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20.0 COMPARATOR VOLTAGE REFERENCE MODULE

The comparator voltage reference is a 16-tap resistor ladder network that provides a selectable reference voltage. Although its primary purpose is to provide a reference for the analog comparators, it may also be used independently of them.

A block diagram of the module is shown in Figure 20-1. The resistor ladder is segmented to provide two ranges of CVREF values and has a power-down function to conserve power when the reference is not being used. The module's supply reference can be provided from either device VDD/VSS or an external voltage reference.

20.1 Configuring the Comparator Voltage Reference

The voltage reference module is controlled through the CVRCON register (Register 20-1). The comparator voltage reference provides two ranges of output voltage, each with 16 distinct levels. The range to be used is selected by the CVRR bit (CVRCON<5>). The primary difference between the ranges is the size of the steps selected by the CVREF Selection bits (CVR<3:0>), with one range offering finer resolution. The equations used to calculate the output of the comparator voltage reference are as follows:

If CVRR = 1:
CVREF = ((CVR<3:0>)/24) x CVRSRC

If CVRR = 0:
CVREF = (CVRSRC x 1/4) + (((CVR<3:0>)/32) x CVRSRC)

The comparator reference supply voltage can come from either VDD and VSS, or the external VREF+ and VREF- that are multiplexed with RA2 and RA3. The voltage source is selected by the CVRSS bit (CVRCON<4>).

The settling time of the comparator voltage reference must be considered when changing the CVREF output (see Table 24-3 in **Section 24.0 "Electrical Characteristics"**).

REGISTER 20-1: CVRCON: COMPARATOR VOLTAGE REFERENCE CONTROL REGISTER

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
CVREN	CVROE(1)	CVRR	CVRSS	CVR3	CVR2	CVR1	CVR0
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 7 **CVREN**: Comparator Voltage Reference Enable bit
1 = CVREF circuit powered on
0 = CVREF circuit powered down

bit 6 **CVROE**: Comparator VREF Output Enable bit(1)
1 = CVREF voltage level is also output on the RA2/AN2/VREF-/CVREF pin
0 = CVREF voltage is disconnected from the RA2/AN2/VREF-/CVREF pin

bit 5 **CVRR**: Comparator VREF Range Selection bit
1 = 0 to 0.667 CVRSRC, with CVRSRC/24 step size (low range)
0 = 0.25 CVRSRC to 0.75 CVRSRC, with CVRSRC/32 step size (high range)

bit 4 **CVRSS**: Comparator VREF Source Selection bit
1 = Comparator reference source, CVRSRC = (VREF+) – (VREF-)
0 = Comparator reference source, CVRSRC = VDD – VSS

bit 3-0 **CVR<3:0>:** Comparator VREF Value Selection bits (0 ≤ (CVR<3:0>) ≤ 15)
When CVRR = 1:
CVREF = ((CVR<3:0>)/24) • (CVRSRC)
When CVRR = 0:
CVREF = (CVRSRC/4) + ((CVR<3:0>)/32) • (CVRSRC)

Note 1: CVROE overrides the TRISA<2> bit setting.

TABLE 22-2: PIC18FXXXX INSTRUCTION SET

Mnemonic, Operands		Description	Cycles	16-Bit Instruction Word				Status Affected	Notes
				MSb			LSb		
BYTE-ORIENTED OPERATIONS									
ADDWF	f, d, a	Add WREG and f	1	`0010`	`01da`	`ffff`	`ffff`	C, DC, Z, OV, N	1, 2
ADDWFC	f, d, a	Add WREG and Carry bit to f	1	`0010`	`00da`	`ffff`	`ffff`	C, DC, Z, OV, N	1, 2
ANDWF	f, d, a	AND WREG with f	1	`0001`	`01da`	`ffff`	`ffff`	Z, N	1,2
CLRF	f, a	Clear f	1	`0110`	`101a`	`ffff`	`ffff`	Z	2
COMF	f, d, a	Complement f	1	`0001`	`11da`	`ffff`	`ffff`	Z, N	1, 2
CPFSEQ	f, a	Compare f with WREG, Skip =	1 (2 or 3)	`0110`	`001a`	`ffff`	`ffff`	None	4
CPFSGT	f, a	Compare f with WREG, Skip >	1 (2 or 3)	`0110`	`010a`	`ffff`	`ffff`	None	4
CPFSLT	f, a	Compare f with WREG, Skip <	1 (2 or 3)	`0110`	`000a`	`ffff`	`ffff`	None	1, 2
DECF	f, d, a	Decrement f	1	`0000`	`01da`	`ffff`	`ffff`	C, DC, Z, OV, N	1, 2, 3, 4
DECFSZ	f, d, a	Decrement f, Skip if 0	1 (2 or 3)	`0010`	`11da`	`ffff`	`ffff`	None	1, 2, 3, 4
DCFSNZ	f, d, a	Decrement f, Skip if Not 0	1 (2 or 3)	`0100`	`11da`	`ffff`	`ffff`	None	1, 2
INCF	f, d, a	Increment f	1	`0010`	`10da`	`ffff`	`ffff`	C, DC, Z, OV, N	1, 2, 3, 4
INCFSZ	f, d, a	Increment f, Skip if 0	1 (2 or 3)	`0011`	`11da`	`ffff`	`ffff`	None	4
INFSNZ	f, d, a	Increment f, Skip if Not 0	1 (2 or 3)	`0100`	`10da`	`ffff`	`ffff`	None	1, 2
IORWF	f, d, a	Inclusive OR WREG with f	1	`0001`	`00da`	`ffff`	`ffff`	Z, N	1, 2
MOVF	f, d, a	Move f	1	`0101`	`00da`	`ffff`	`ffff`	Z, N	1
MOVFF	f_s, f_d	Move f_s (source) to 1st Word f_d (destination) 2nd Word	2	`1100` `1111`	`ffff` `ffff`	`ffff` `ffff`	`ffff` `ffff`	None	
MOVWF	f, a	Move WREG to f	1	`0110`	`111a`	`ffff`	`ffff`	None	
MULWF	f, a	Multiply WREG with f	1	`0000`	`001a`	`ffff`	`ffff`	None	1, 2
NEGF	f, a	Negate f	1	`0110`	`110a`	`ffff`	`ffff`	C, DC, Z, OV, N	
RLCF	f, d, a	Rotate Left f through Carry	1	`0011`	`01da`	`ffff`	`ffff`	C, Z, N	1, 2
RLNCF	f, d, a	Rotate Left f (No Carry)	1	`0100`	`01da`	`ffff`	`ffff`	Z, N	
RRCF	f, d, a	Rotate Right f through Carry	1	`0011`	`00da`	`ffff`	`ffff`	C, Z, N	
RRNCF	f, d, a	Rotate Right f (No Carry)	1	`0100`	`00da`	`ffff`	`ffff`	Z, N	
SETF	f, a	Set f	1	`0110`	`100a`	`ffff`	`ffff`	None	1, 2
SUBFWB	f, d, a	Subtract f from WREG with Borrow	1	`0101`	`01da`	`ffff`	`ffff`	C, DC, Z, OV, N	
SUBWF	f, d, a	Subtract WREG from f	1	`0101`	`11da`	`ffff`	`ffff`	C, DC, Z, OV, N	1, 2
SUBWFB	f, d, a	Subtract WREG from f with Borrow	1	`0101`	`10da`	`ffff`	`ffff`	C, DC, Z, OV, N	
SWAPF	f, d, a	Swap Nibbles in f	1	`0011`	`10da`	`ffff`	`ffff`	None	4
TSTFSZ	f, a	Test f, Skip if 0	1 (2 or 3)	`0110`	`011a`	`ffff`	`ffff`	None	1, 2
XORWF	f, d, a	Exclusive OR WREG with f	1	`0001`	`10da`	`ffff`	`ffff`	Z, N	

Note 1: When a PORT register is modified as a function of itself (e.g., `MOVF PORTB, 1, 0`), the value used will be that value present on the pins themselves. For example, if the data latch is ‘`1`’ for a pin configured as input and is driven low by an external device, the data will be written back with a ‘`0`’.

2: If this instruction is executed on the TMR0 register (and where applicable, ‘d’ = `1`), the prescaler will be cleared if assigned.

3: If the Program Counter (PC) is modified or a conditional test is true, the instruction requires two cycles. The second cycle is executed as a `NOP`.

4: Some instructions are two-word instructions. The second word of these instructions will be executed as a `NOP` unless the first word of the instruction retrieves the information embedded in these 16 bits. This ensures that all program memory locations have a valid instruction.

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22.1.1 STANDARD INSTRUCTION SET

ADDLW **ADD Literal to W**

Syntax:	ADDLW k
Operands:	$0 \leq k \leq 255$
Operation:	(W) + k → W
Status Affected:	N, OV, C, DC, Z
Encoding:	`0000` `1111` `kkkk` `kkkk`
Description:	The contents of W are added to the 8-bit literal 'k' and the result is placed in W.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to W

Example: `ADDLW 15h`

Before Instruction
W = 10h
After Instruction
W = 25h

ADDWF **ADD W to f**

Syntax:	ADDWF f {,d {,a}}
Operands:	$0 \leq f \leq 255$ d ∈ [0,1] a ∈ [0,1]
Operation:	(W) + (f) → dest
Status Affected:	N, OV, C, DC, Z
Encoding:	`0010` `01da` `ffff` `ffff`
Description:	Add W to register 'f'. If 'd' is '0', the result is stored in W. If 'd' is '1', the result is stored back in register 'f' (default). If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank (default). If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See **Section 22.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `ADDWF REG, 0, 0`

Before Instruction
W = 17h
REG = 0C2h
After Instruction
W = 0D9h
REG = 0C2h

Note:	All PIC18 instructions may take an optional label argument preceding the instruction mnemonic for use in symbolic addressing. If a label is used, the instruction format then becomes: {label} instruction argument(s).

POP — Pop Top of Return Stack

Syntax:	POP
Operands:	None
Operation:	(TOS) $\rightarrow$ bit bucket
Status Affected:	None
Encoding:	`0000` `0000` `0000` `0110`
Description:	The TOS value is pulled off the return stack and is discarded. The TOS value then becomes the previous value that was pushed onto the return stack. This instruction is provided to enable the user to properly manage the return stack to incorporate a software stack.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	POP TOS value	No operation

Example:
```
POP
GOTO    NEW
```

Before Instruction
TOS = 0031A2h
Stack (1 level down) = 014332h

After Instruction
TOS = 014332h
PC = NEW

PUSH — Push Top of Return Stack

Syntax:	PUSH
Operands:	None
Operation:	(PC + 2) $\rightarrow$ TOS
Status Affected:	None
Encoding:	`0000` `0000` `0000` `0101`
Description:	The PC + 2 is pushed onto the top of the return stack. The previous TOS value is pushed down on the stack. This instruction allows implementing a software stack by modifying TOS and then pushing it onto the return stack.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	PUSH PC + 2 onto return stack	No operation	No operation

Example:
```
PUSH
```

Before Instruction
TOS = 345Ah
PC = 0124h

After Instruction
PC = 0126h
TOS = 0126h
Stack (1 level down) = 345Ah

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RRNCF — Rotate Right f (No Carry)

Syntax: RRNCF f {,d {,a}}

Operands: $0 \le f \le 255$
d ∈ [0,1]
a ∈ [0,1]

Operation: (f<n>) → dest<n – 1>,
(f<0>) → dest<7>

Status Affected: N, Z

Encoding:

0100	00da	ffff	ffff

Description: The contents of register ‘f’ are rotated one bit to the right. If ‘d’ is ‘0’, the result is placed in W. If ‘d’ is ‘1’, the result is placed back in register ‘f’ (default).
If ‘a’ is ‘0’, the Access Bank will be selected, overriding the BSR value. If ‘a’ is ‘1’, then the bank will be selected as per the BSR value (default).
If ‘a’ is ‘0’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 22.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write to destination

Example 1: `RRNCF REG, 1, 0`

Before Instruction
REG = 1101 0111
After Instruction
REG = 1110 1011

Example 2: `RRNCF REG, 0, 0`

Before Instruction
W = ?
REG = 1101 0111
After Instruction
W = 1110 1011
REG = 1101 0111

SETF — Set f

Syntax: SETF f {,a}

Operands: $0 \le f \le 255$
a ∈ [0,1]

Operation: FFh → f

Status Affected: None

Encoding:

0110	100a	ffff	ffff

Description: The contents of the specified register are set to FFh.
If ‘a’ is ‘0’, the Access Bank is selected. If ‘a’ is ‘1’, the BSR is used to select the GPR bank (default).
If ‘a’ is ‘0’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh). See **Section 22.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	Write register ‘f’

Example: `SETF REG, 1`

Before Instruction
REG = 5Ah
After Instruction
REG = FFh

TSTFSZ	Test f, Skip if 0
Syntax:	TSTFSZ f {,a}
Operands:	0 ≤ f ≤ 255 a ∈ [0,1]
Operation:	skip if f = 0
Status Affected:	None
Encoding:	0110 011a ffff ffff
Description:	If ‘f’ = 0, the next instruction fetched during the current instruction execution is discarded and a NOP is executed, making this a two-cycle instruction. If ‘a’ is ‘0’, the Access Bank is selected. If ‘a’ is ‘1’, the BSR is used to select the GPR bank (default). If ‘a’ is ‘0’ and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever f ≤ 95 (5Fh). See **Section 22.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode”** for details.
Words:	1
Cycles:	1(2) **Note:** 3 cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register ‘f’	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE     TSTFSZ  CNT, 1
NZERO    :
ZERO     :
```

Before Instruction
PC = Address (HERE)
After Instruction
If CNT = 00h,
PC = Address (ZERO)
If CNT ≠ 00h,
PC = Address (NZERO)

XORLW	Exclusive OR Literal with W
Syntax:	XORLW k
Operands:	0 ≤ k ≤ 255
Operation:	(W) .XOR. k → W
Status Affected:	N, Z
Encoding:	0000 1010 kkkk kkkk
Description:	The contents of W are XORed with the 8-bit literal ‘k’. The result is placed in W.
Words:	1
Cycles:	1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal ‘k’	Process Data	Write to W

Example: XORLW 0AFh

Before Instruction
W = B5h
After Instruction
W = 1Ah

23.7 MPLAB ICE 2000 High-Performance In-Circuit Emulator

The MPLAB ICE 2000 In-Circuit Emulator is intended to provide the product development engineer with a complete microcontroller design tool set for PIC microcontrollers. Software control of the MPLAB ICE 2000 In-Circuit Emulator is advanced by the MPLAB Integrated Development Environment, which allows editing, building, downloading and source debugging from a single environment.

The MPLAB ICE 2000 is a full-featured emulator system with enhanced trace, trigger and data monitoring features. Interchangeable processor modules allow the system to be easily reconfigured for emulation of different processors. The architecture of the MPLAB ICE 2000 In-Circuit Emulator allows expansion to support new PIC microcontrollers.

The MPLAB ICE 2000 In-Circuit Emulator system has been designed as a real-time emulation system with advanced features that are typically found on more expensive development tools. The PC platform and Microsoft® Windows® 32-bit operating system were chosen to best make these features available in a simple, unified application.

23.8 MPLAB REAL ICE In-Circuit Emulator System

MPLAB REAL ICE In-Circuit Emulator System is Microchip's next generation high-speed emulator for Microchip Flash DSC and MCU devices. It debugs and programs PIC® Flash MCUs and dsPIC® Flash DSCs with the easy-to-use, powerful graphical user interface of the MPLAB Integrated Development Environment (IDE), included with each kit.

The MPLAB REAL ICE probe is connected to the design engineer's PC using a high-speed USB 2.0 interface and is connected to the target with either a connector compatible with the popular MPLAB ICD 2 system (RJ11) or with the new high-speed, noise tolerant, Low-Voltage Differential Signal (LVDS) interconnection (CAT5).

MPLAB REAL ICE is field upgradeable through future firmware downloads in MPLAB IDE. In upcoming releases of MPLAB IDE, new devices will be supported, and new features will be added, such as software breakpoints and assembly code trace. MPLAB REAL ICE offers significant advantages over competitive emulators including low-cost, full-speed emulation, real-time variable watches, trace analysis, complex breakpoints, a ruggedized probe interface and long (up to three meters) interconnection cables.

23.9 MPLAB ICD 2 In-Circuit Debugger

Microchip's In-Circuit Debugger, MPLAB ICD 2, is a powerful, low-cost, run-time development tool, connecting to the host PC via an RS-232 or high-speed USB interface. This tool is based on the Flash PIC MCUs and can be used to develop for these and other PIC MCUs and dsPIC DSCs. The MPLAB ICD 2 utilizes the in-circuit debugging capability built into the Flash devices. This feature, along with Microchip's In-Circuit Serial Programming™ (ICSP™) protocol, offers cost-effective, in-circuit Flash debugging from the graphical user interface of the MPLAB Integrated Development Environment. This enables a designer to develop and debug source code by setting breakpoints, single stepping and watching variables, and CPU status and peripheral registers. Running at full speed enables testing hardware and applications in real time. MPLAB ICD 2 also serves as a development programmer for selected PIC devices.

23.10 MPLAB PM3 Device Programmer

The MPLAB PM3 Device Programmer is a universal, CE compliant device programmer with programmable voltage verification at V_{DDMIN} and V_{DDMAX} for maximum reliability. It features a large LCD display (128 x 64) for menus and error messages and a modular, detachable socket assembly to support various package types. The ICSP™ cable assembly is included as a standard item. In Stand-Alone mode, the MPLAB PM3 Device Programmer can read, verify and program PIC devices without a PC connection. It can also set code protection in this mode. The MPLAB PM3 connects to the host PC via an RS-232 or USB cable. The MPLAB PM3 has high-speed communications and optimized algorithms for quick programming of large memory devices and incorporates an SD/MMC card for file storage and secure data applications.

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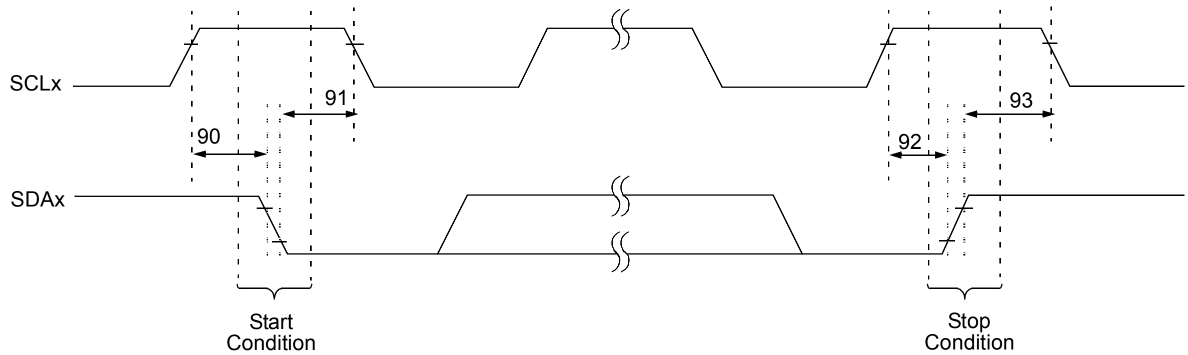
FIGURE 24-14: I²C™ BUS START/STOP BITS TIMING

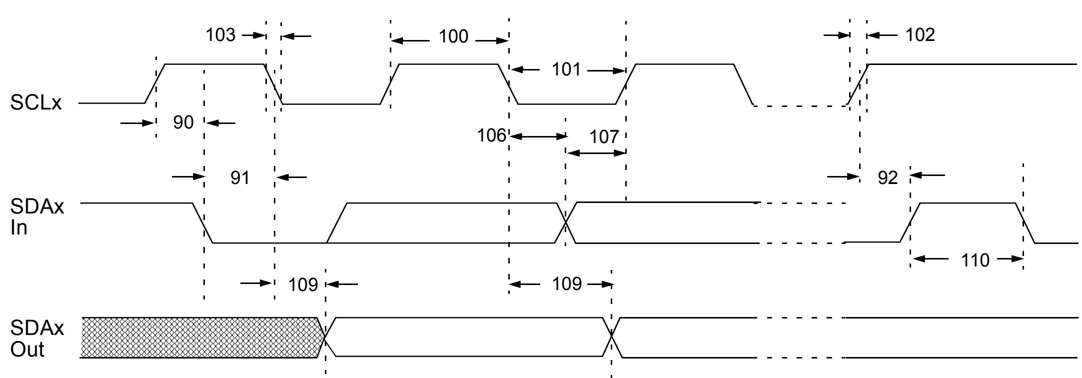


Note: Refer to Figure 24-3 for load conditions.

TABLE 24-18: I²C™ BUS START/STOP BITS REQUIREMENTS (SLAVE MODE)

Param. No.	Symbol	Characteristic		Min	Max	Units	Conditions
90	TSU:STA	Start Condition Setup Time	100 kHz mode	4700	—	ns	Only relevant for Repeated Start condition
			400 kHz mode	600	—		
91	THD:STA	Start Condition Hold Time	100 kHz mode	4000	—	ns	After this period, the first clock pulse is generated
			400 kHz mode	600	—		
92	TSU:STO	Stop Condition Setup Time	100 kHz mode	4700	—	ns	
			400 kHz mode	600	—		
93	THD:STO	Stop Condition Hold Time	100 kHz mode	4000	—	ns	
			400 kHz mode	600	—		

FIGURE 24-15: I²C™ BUS DATA TIMING



Note: Refer to Figure 24-3 for load conditions.