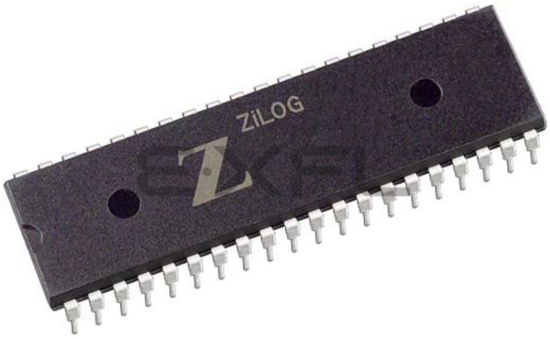

Understanding Embedded - Microprocessors

Embedded microprocessors are specialized computing chips designed to perform specific tasks within an embedded system. Unlike general-purpose microprocessors found in personal computers, embedded microprocessors are tailored for dedicated functions within larger systems, offering optimized performance, efficiency, and reliability. These microprocessors are integral to the operation of countless electronic devices, providing the computational power necessary for controlling processes, handling data, and managing communications.

Applications of Embedded - Microprocessors

Embedded microprocessors are utilized across a broad spectrum of applications, making them indispensable in

Details	
Product Status	Active
Core Processor	Z80
Number of Cores/Bus Width	1 Core, 8-Bit
Speed	20MHz
Co-Processors/DSP	-
RAM Controllers	-
Graphics Acceleration	No
Display & Interface Controllers	-
Ethernet	-
SATA	-
USB	-
Voltage - I/O	5.0V
Operating Temperature	-40°C ~ 100°C (TA)
Security Features	-
Package / Case	40-DIP (0.620", 15.75mm)
Supplier Device Package	40-PDIP
Purchase URL	https://www.e-xfl.com/product-detail/zilog/z84c0020peg

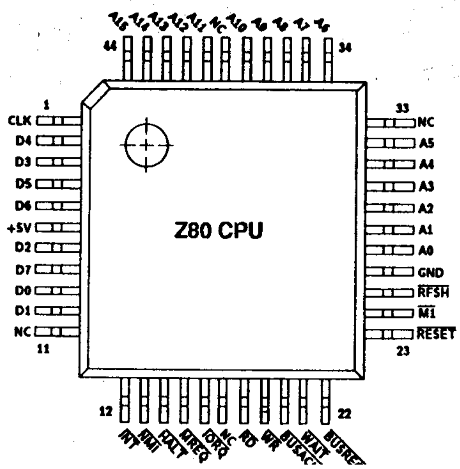

Figure 2a. 44-Pin LQFP, Pin Assignments (Only available for 84C00)

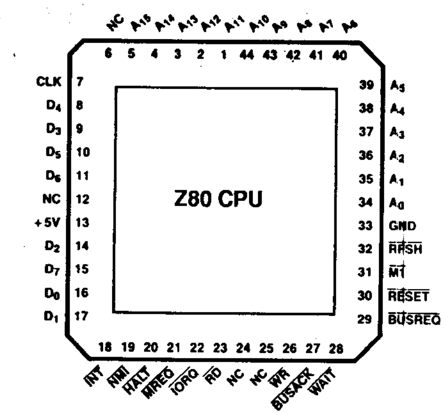

Figure 2b. 44-Pin Chip Carrier Pin Assignments

GENERAL DESCRIPTION

The CPUs are fourth-generation enhanced microprocessors with exceptional computational power. They offer higher system throughput and more efficient memory utilization than comparable second- and third-generation microprocessors. The internal registers contain 208 bits of read/write memory that are accessible to the programmer. These registers include two sets of six general-purpose registers which may be used individually as either 8-bit registers or as 16-bit register pairs. In addition, there are two sets of accumulator and flag registers. A group of "Exchange" instructions makes either set of main or alternate registers accessible to the programmer. The alternate set allows operation in foreground-background mode or it may be reserved for very fast interrupt response.

The CPU also contains a Stack Pointer, Program Counter, two index registers, a Refresh register (counter), and an Interrupt register. The CPU is easy to incorporate into a system since it requires only a single +5V power source. All output signals are fully decoded and timed to control standard memory or peripheral circuits; the CPU is supported by an extensive family of peripheral controllers. The internal block diagram (Figure 3) shows the primary functions of the processors. Subsequent text provides more detail on the I/O controller family, registers, instruction set, interrupts and daisy chaining, and CPU timing.

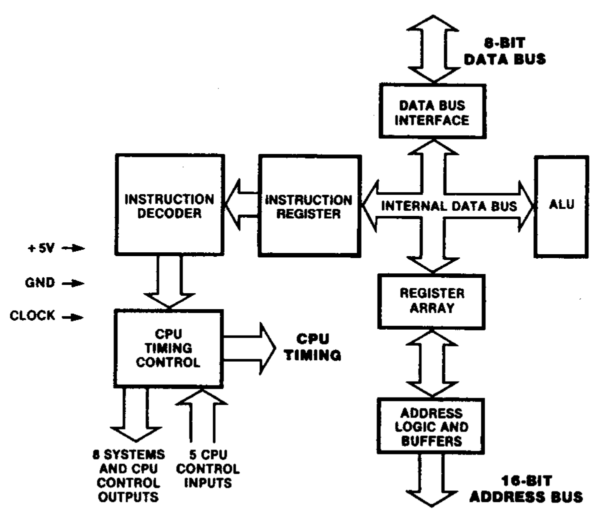

Figure 3. Z80C CPU Block Diagram

address. This flexibility in selecting the interrupt service routine address allows the peripheral device to use several different types of service routines. These routines may be located at any available location in memory. Since the interrupting device supplies the low-order byte of the 2-byte vector, bit 0 (A_0) must be a zero.

Interrupt Enable/Disable Operation. Two flip-flops, IFF_1 and IFF_2, referred to in the register description, are used to signal the CPU interrupt status. Operation of the two flip-flops is described in Table 2. For more details, refer to the *Z80 CPU Technical Manual* (03-0029-01) and *Z80 Assembly Language Programming Manual* (03-0002-01).

Table 2. State of Flip-Flops

Action	IFF_1	IFF_2	Comments
CPU Reset	0	0	Maskable interrupt $\overline{INT}$ disabled
DI instruction execution	0	0	Maskable interrupt $\overline{INT}$ disabled
EI instruction execution	1	1	Maskable interrupt $\overline{INT}$ enabled
LD A,I instruction execution	•	•	$IFF_2 \rightarrow$ Parity flag
LD A,R instruction execution	•	•	$IFF_2 \rightarrow$ Parity flag
Accept $\overline{NMI}$	0	•	Maskable interrupt $\overline{INT}$ disabled
RETN instruction execution	IFF_2	•	$IFF_2 \rightarrow IFF_1$ at completion of an $\overline{NMI}$ service routine.

INSTRUCTION SET

The microprocessor has one of the most powerful and versatile instruction sets available in any 8-bit microprocessor. It includes such unique operations as a block move for fast, efficient data transfers within memory, or between memory and I/O. It also allows operations on any bit in any location in memory.

The following is a summary of the instruction set which shows the assembly language mnemonic, the operation, the flag status, and gives comments on each instruction. For an explanation of flag notations and symbols for mnemonic tables, see the Symbolic Notations section which follows these tables. The *Z80 CPU Technical Manual* (03-0029-01), the *Programmer's Reference Guide* (03-0012-03), and *Assembly Language Programming Manual* (03-0002-01) contain significantly more details for programming use.

The instructions are divided into the following categories:

- 8-bit loads
- 16-bit loads
- Exchanges, block transfers, and searches
- 8-bit arithmetic and logic operations
- General-purpose arithmetic and CPU control
- 16-bit arithmetic operations
- Rotates and shifts
- Bit set, reset, and test operations
- Jumps
- Calls, returns, and restarts
- Input and output operations

A variety of addressing modes are implemented to permit efficient and fast data transfer between various registers, memory locations, and input/output devices. These addressing modes include:

- Immediate
- Immediate extended
- Modified page zero
- Relative
- Extended
- Indexed
- Register
- Register indirect
- Implied
- Bit

8-BIT LOAD GROUP

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments	
LD r, r′	r ← r′	•	•	X	•	X	•	•	•	01	r	r′		1	1	4	r, r′	Reg.
LD r, n	r ← n	•	•	X	•	X	•	•	•	00	r	110		2	2	7	000	B
											← n →						001	C
LD r, (HL)	r ← (HL)	•	•	X	•	X	•	•	•	01	r	110		1	2	7	010	D
LD r, (IX+d)	r ← (IX+d)	•	•	X	•	X	•	•	•	11	011	101	DD	3	5	19	011	E
										01	r	110					100	H
											← d →						101	L
LD r, (IY+d)	r ← (IY+d)	•	•	X	•	X	•	•	•	11	111	101	FD	3	5	19	111	A
										01	r	110						
											← d →							
LD (HL), r	(HL) ← r	•	•	X	•	X	•	•	•	01	110	r		1	2	7		
LD (IX+d), r	(IX+d) ← r	•	•	X	•	X	•	•	•	11	011	101	DD	3	5	19		
										01	110	r						
											← d →							
LD (IY+d), r	(IY+d) ← r	•	•	X	•	X	•	•	•	11	111	101	FD	3	5	19		
										01	110	r						
											← d →							
LD (HL), n	(HL) ← n	•	•	X	•	X	•	•	•	00	110	110	36	2	3	10		
											← n →							
LD (IX+d), n	(IX+d) ← n	•	•	X	•	X	•	•	•	11	011	101	DD	4	5	19		
										00	110	110	36					
											← d →							
											← n →							

8-BIT LOAD GROUP (Continued)

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
LD (IY+d), n	(IY+d) ← n	•	•	X	•	X	•	•	•	11	111	101	FD	4	5	19	
										00	110	110	36				
											← d →						
											← n →						
LD A, (BC)	A ← (BC)	•	•	X	•	X	•	•	•	00	001	010	0A	1	2	7	
LD A, (DE)	A ← (DE)	•	•	X	•	X	•	•	•	00	011	010	1A	1	2	7	
LD A, (nn)	A ← (nn)	•	•	X	•	X	•	•	•	00	111	010	3A	3	4	13	
											← n →						
											← n →						
LD (BC), A	(BC) ← A	•	•	X	•	X	•	•	•	00	000	010	02	1	2	7	
LD (DE), A	(DE) ← A	•	•	X	•	X	•	•	•	00	010	010	12	1	2	7	
LD (nn), A	(nn) ← A	•	•	X	•	X	•	•	•	00	110	010	32	3	4	13	
											← n →						
											← n →						
LD A, I	A ← I	↕	↕	X	0	X	IFF	0	•	11	101	101	ED	2	2	9	
										01	010	111	57				
LD A, R	A ← R	↕	↕	X	0	X	IFF	0	•	11	101	101	ED	2	2	9	
										01	011	111	5F				
LD I, A	I ← A	•	•	X	•	X	•	•	•	11	101	101	ED	2	2	9	
										01	000	111	47				
LD R, A	R ← A	•	•	X	•	X	•	•	•	11	101	101	ED	2	2	9	
										01	001	111	4F				

NOTE: IFF, the content of the interrupt enable flip-flop, (IFF_2), is copied into the P/V flag.

16-BIT LOAD GROUP

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments	
LD dd, nn	dd ← nn	•	•	X	•	X	•	•	•	00	dd0	001		3	3	10	dd	Pair
											← n →						00	BC
											← n →						01	DE
LD IX, nn	IX ← nn	•	•	X	•	X	•	•	•	11	011	101	DD	4	4	14	10	HL
										00	100	001	21				11	SP
											← n →							
											← n →							
LD IY, nn	IY ← nn	•	•	X	•	X	•	•	•	11	111	101	FD	4	4	14		
										00	100	001	21					
											← n →							
											← n →							
LD HL, (nn)	H ← (nn+1)	•	•	X	•	X	•	•	•	00	101	010	2A	3	5	16		
	L ← (nn)										← n →							
											← n →							
LD dd, (nn)	dd_H ← (nn+1)	•	•	X	•	X	•	•	•	11	101	101	ED	4	6	20		
	dd_L ← (nn)									01	dd1	011						
											← n →							
											← n →							

NOTE: $(PAIR)_H$, $(PAIR)_L$ refer to high order and low order eight bits of the register pair respectively. e.g., BC_L = C, AF_H = A.

16-BIT LOAD GROUP (Continued)

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
LD IX, (nn)	$IX_H \leftarrow (nn+1)$	•	•	X	•	X	•	•	•	11	011	101	DD	4	6	20	
	$IX_L \leftarrow (nn)$									00	101	010	2A				
											← n →						
											← n →						
LD IY, (nn)	$IY_H \leftarrow (nn+1)$	•	•	X	•	X	•	•	•	11	111	101	FD	4	6	20	
	$IY_L \leftarrow (nn)$									00	101	010	2A				
											← n →						
											← n →						
LD (nn), HL	$(nn+1) \leftarrow H$	•	•	X	•	X	•	•	•	00	100	010	22	3	5	16	
	$(nn) \leftarrow L$										← n →						
											← n →						
LD (nn), dd	$(nn+1) \leftarrow dd_H$	•	•	X	•	X	•	•	•	11	101	101	ED	4	6	20	
	$(nn) \leftarrow dd_L$									01	dd0	011					
											← n →						
											← n →						
LD (nn), IX	$(nn+1) \leftarrow IX_H$	•	•	X	•	X	•	•	•	11	011	101	DD	4	6	20	
	$(nn) \leftarrow IX_L$									00	100	010	22				
											← n →						
											← n →						
LD (nn), IY	$(nn+1) \leftarrow IY_H$	•	•	X	•	X	•	•	•	11	111	101	FD	4	6	20	
	$(nn) \leftarrow IY_L$									00	100	010	22				
											← n →						
											← n →						
LD SP, HL	$SP \leftarrow HL$	•	•	X	•	X	•	•	•	11	111	001	F9	1	1	6	
LD SP, IX	$SP \leftarrow IX$	•	•	X	•	X	•	•	•	11	011	101	DD	2	2	10	
										11	111	001	F9				
LD SP, IY	$SP \leftarrow IY$	•	•	X	•	X	•	•	•	11	111	101	FD	2	2	10	
										11	111	001	F9				qq Pair
PUSH qq	$(SP-2) \leftarrow qq_L$	•	•	X	•	X	•	•	•	11	qq0	101		1	3	11	00 BC
	$(SP-1) \leftarrow qq_H$																01 DE
	$SP \rightarrow SP-2$																10 HL
PUSH IX	$(SP-2) \leftarrow IX_L$	•	•	X	•	X	•	•	•	11	011	101	DD	2	4	15	11 AF
	$(SP-1) \leftarrow IX_H$									11	100	101	E5				
	$SP \rightarrow SP-2$																
PUSH IY	$(SP-2) \leftarrow IY_L$	•	•	X	•	X	•	•	•	11	111	101	FD	2	4	15	
	$(SP-1) \leftarrow IY_H$									11	100	101	E5				
	$SP \rightarrow SP-2$																
POP qq	$qq_H \leftarrow (SP+1)$	•	•	X	•	X	•	•	•	11	qq0	001		1	3	10	
	$qqL \leftarrow (SP)$																
	$SP \rightarrow SP+2$																
POP IX	$IX_H \leftarrow (SP+1)$	•	•	X	•	X	•	•	•	11	011	101	DD	2	4	14	
	$IX_L \leftarrow (SP)$									11	100	001	E1				
	$SP \rightarrow SP+2$																
POP IY	$IY_H \leftarrow (SP+1)$	•	•	X	•	X	•	•	•	11	111	101	FD	2	4	14	
	$IY_L \leftarrow (SP)$									11	100	001	E1				
	$SP \rightarrow SP+2$																

NOTE: $(PAIR)_H$, $(PAIR)_L$ refer to high order and low order eight bits of the register pair respectively, e.g., BC_L = C, AF_H = A.

EXCHANGE, BLOCK TRANSFER, BLOCK SEARCH GROUPS (Continued)

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
			③				①										
CPIR	A − (HL)	↕	↕	X	↕	X	↕	1	•	11	101	101	ED	2	5	21	If BC ≠ 0 and
																	A ≠ (HL)
	HL ← HL + 1									10	110	001	B1	2	4	16	If BC = 0 or
	BC ← BC − 1																A = (HL)
	Repeat until																
	A = (HL) or																
	BC = 0																
			③				①										
CPD	A − (HL)	↕	↕	X	↕	X	↕	1	•	11	101	101	ED	2	4	16	
	HL ← HL − 1									10	101	001	A9				
	BC ← BC − 1																
			③				①										
CPDR	A − (HL)	↕	↕	X	↕	X	↕	1	•	11	101	101	ED	2	5	21	If BC ≠ 0 and
																	A ≠ (HL)
	HL ← HL − 1									10	111	001	B9	2	4	16	If BC = 0 or
	BC ← BC − 1																A = (HL)
	Repeat until																
	A = (HL) or																
	BC = 0																

NOTE: ① P/V flag is 0 if the result of BC − 1 = 0, otherwise P/V = 1.
② P/V flag is 0 only at completion of instruction.
③ Z flag is 1 if A = (HL), otherwise Z = 0.

8-BIT ARITHMETIC AND LOGICAL GROUP

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
ADD A, r	A ← A + r	↕	↕	X	↕	X	V	0	↕	10	000	r		1	1	4	r Reg.
ADD A, n	A ← A + n	↕	↕	X	↕	X	V	0	↕	11	000	110		2	2	7	000 B
											← n →						001 C
																	010 D
ADD A, (HL)	A ← A + (HL)	↕	↕	X	↕	X	V	0	↕	10	000	110		1	2	7	011 E
ADD A, (IX + d)	A ← A + (IX + d)	↕	↕	X	↕	X	V	0	↕	11	011	101	DD	3	5	19	100 H
										10	000	110					101 L
											← d →						111 A
ADD A, (IY + d)	A ← A + (IY + d)	↕	↕	X	↕	X	V	0	↕	11	111	101	FD	3	5	19	
										10	000	110					
											← d →						
ADC A, s	A ← A + s + CY	↕	↕	X	↕	X	V	0	↕		001						s is any of r, n,
SUB s	A ← A − s	↕	↕	X	↕	X	V	1	↕		010						(HL), (IX + d),
SBC A, s	A ← A − s − CY	↕	↕	X	↕	X	V	1	↕		011						(IY + d) as
AND s	A ← A > s	↕	↕	X	1	X	P	0	0		100						shown for ADD
OR s	A ← A > s	↕	↕	X	0	X	P	0	0		110						instruction. The
XOR s	A ← A ⊕ s	↕	↕	X	0	X	P	0	0		101						indicated bits
CP s	A − s	↕	↕	X	↕	X	V	1	↕		111						replace the
																	000 in the
																	ADD set above.

8-BIT ARITHMETIC AND LOGICAL GROUP (Continued)

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
INC r	r ← r + 1	↕	↕	X	↕	X	V	0	•	00	r	[100]		1	1	4	
INC (HL)	(HL) ← (HL) + 1	↕	↕	X	↕	X	V	0	•	00	110	[100]		1	3	11	
INC (IX + d)	(IX + d) ← (IX + d) + 1	↕	↕	X	↕	X	V	0	•	11 00	011 110 ← d →	101 [100]	DD	3	6	23	
INC (IY + d)	(IY + d) ← (IY + d) + 1	↕	↕	X	↕	X	V	0	•	11 00	111 110 ← d →	101 [100]	FD	3	6	23	
DEC m	m ← m − 1	↕	↕	X	↕	X	V	1	•			[101]					

NOTE: m is any of r, (HL), (IX + d), (IY + d) as shown for INC. DEC same format and states as INC. Replace [100] with [101] in opcode.

GENERAL-PURPOSE ARITHMETIC AND CPU CONTROL GROUPS

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
DAA	@	↕	↕	X	↕	X	P	•	↕	00	100	111	27	1	1	4	Decimal adjust accumulator.
CPL	A ← $\overline{A}$	•	•	X	1	X	•	1	•	00	101	111	2F	1	1	4	Complement accumulator (one's complement).
NEG	A ← 0 − A	↕	↕	X	↕	X	V	1	↕	11 01	101 000	101 100	ED 44	2	2	8	Negate acc. (two's complement).
CCF	CY ← $\overline{CY}$	•	•	X	X	X	•	0	↕	00	111	111	3F	1	1	4	Complement carry flag.
SCF	CY ← 1	•	•	X	0	X	•	0	1	00	110	111	37	1	1	4	Set carry flag.
NOP	No operation	•	•	X	•	X	•	•	•	00	000	000	00	1	1	4	
HALT	CPU halted	•	•	X	•	X	•	•	•	01	110	110	76	1	1	4	
DI ★	IFF ← 0	•	•	X	•	X	•	•	•	11	110	011	F3	1	1	4	
EI ★	IFF ← 1	•	•	X	•	X	•	•	•	11	111	011	FB	1	1	4	
IM 0	Set interrupt mode 0	•	•	X	•	X	•	•	•	11 01	101 000	101 110	ED 46	2	2	8	
IM 1	Set interrupt mode 1	•	•	X	•	X	•	•	•	11 01	101 010	101 110	ED 56	2	2	8	
IM 2	Set interrupt mode 2	•	•	X	•	X	•	•	•	11 01	101 011	101 110	ED 5E	2	2	8	

NOTES: @ converts accumulator content into packed BCD following add or subtract with packed BCD operands.
IFF indicates the interrupt enable flip-flop.
CY indicates the carry flip-flop.
★ indicates interrupts are not sampled at the end of EI or DI.

BIT SET, RESET AND TEST GROUP

	Symbolic	Flags								Opcode				No. of	No. of M	No. of T		
Mnemonic	**Operation**	**S**	**Z**		**H**		**P/V**	**N**	**C**	**76**	**543**	**210**	**Hex**	**Bytes**	**Cycles**	**States**	**Comments**	
BIT b, r	$Z \leftarrow r_b$	X	‡	X	1	X	X	0	•	11	001	011	CB	2	2	8	r	Reg.
										01	b	r					000	B
BIT b, (HL)	$Z \leftarrow (HL)_b$	X	‡	X	1	X	X	0	•	11	001	011	CB	2	3	12	001	C
										01	b	110					010	D
BIT b,(IX+d)$_b$	$Z \leftarrow (IX+d)_b$	X	‡	X	1	X	X	0	•	11	011	101	DD	4	5	20	011	E
										11	001	011	CB				100	H
											←d→						101	L
										01	b	110					111	A
																	b	Bit Tested
BIT b, (IY+d)$_b$	$Z \leftarrow (IY+d)_b$	X	‡	X	1	X	X	0	•	11	111	101	FD	4	5	20	000	0
										11	001	011	CB				001	1
											←d→						010	2
										01	b	110					011	3
SET b, r	$r_b \leftarrow 1$	•	•	X	•	X	•	•	•	11	001	011	CB	2	2	8	100	4
										[11]	b	r					101	5
SET b, (HL)	$(HL)_b \leftarrow 1$	•	•	X	•	X	•	•	•	11	001	011	CB	2	4	15	110	6
										[11]	b	110					111	7
SET b, (IX+d)	$(IX+d)_b \leftarrow 1$	•	•	X	•	X	•	•	•	11	011	101	DD	4	6	23		
										11	001	011	CB					
											←d→							
										[11]	b	110						
SET b, (IY+d)	$(IY+d)_b \leftarrow 1$	•	•	X	•	X	•	•	•	11	111	101	FD	4	6	23		
										11	001	011	CB					
											←d→							
										[11]	b	110						
RES b, m	$m_b \leftarrow 0$	•	•	X	•	X	•	•	•	[10]							To form new	
	m≡r, (HL),																opcode replace	
	(IX+d), (IY+d)																[11] of SET b, s	
																	with [10]. Flags	
																	and time	
																	states for SET	
																	instruction.	

NOTE: The notation m_b indicates location m, bit b (0 to 7).

JUMP GROUP

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
JP nn	PC ← nn	•	•	X	•	X	•	•	•	11	000	011	C3	3	3	10	cc Condition
											← n →						000 NZ (non-zero)
											← n →						001 Z (zero)
JP cc, nn	If condition cc	•	•	X	•	X	•	•	•	11	cc	010		3	3	10	010 NC (non-carry)
	is true PC←nn,										← n →						011 C (carry)
	otherwise										← n →						100 PO (parity odd)
	continue																101 PE (parity even)
JR e	PC ← PC+e	•	•	X	•	X	•	•	•	00	011	000	18	2	3	12	110 P (sign positive)
											← e−2 →						111 M (sign negative)
JR C, e	If C=0,	•	•	X	•	X	•	•	•	00	111	000	38	2	2	7	If condition not met.
	continue										← e−2 →						
	If C=1,													2	3	12	If condition is met.
	PC ← PC+e																
JR NC, e	IF C=1,	•	•	X	•	X	•	•	•	00	110	000	30	2	2	7	If condition not met.
	continue										← e−2 →						
	If C=0,													2	3	12	If condition is met.
	PC ← PC+e																
JP Z, e	If Z=0	•	•	X	•	X	•	•	•	00	101	000	28	2	2	7	If condition not met.
	continue										← e−2 →						
	If Z=1,													2	3	12	If condition is met.
	PC ← PC+e																
JR NZ, e	If Z=1,	•	•	X	•	X	•	•	•	00	100	000	20	2	2	7	If condition not met.
	continue										← e−2 →						
	If Z=0,													2	3	12	If condition is met.
	PC ← PC+e																
JP (HL)	PC ← HL	•	•	X	•	X	•	•	•	11	101	001	E9	1	1	4	
JP (IX)	PC ← IX	•	•	X	•	X	•	•	•	11	011	101	DD	2	2	8	
										11	101	001	E9				
JP (IY)	PC ← IY	•	•	X	•	X	•	•	•	11	111	101	FD	2	2	8	
										11	101	001	E9				
DJNZ, e	B ← B−1	•	•	X	•	X	•	•	•	00	010	000	10	2	2	8	If B=0
	If B=0,										← e−2 →						
	continue																
	If B≠0,													2	3	13	If B≠0.
	PC ← PC+e																

NOTES: e represents the extension in the relative addressing mode.
e is a signal two's complement number in the range < − 126, 129 >.
e − 2 in the opcode provides an effective address of pc + e as PC is incremented by 2 prior to the addition of e.

CALL AND RETURN GROUP

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	Opcode 76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
CALL nn	(SP−1)← PC_H	•	•	X	•	X	•	•	•	11	001	101	CD	3	5	17	
	(SP−2)← PC_L										← n →						
	PC ← nn,										← n →						
CALL cc,nn	If condition	•	•	X	•	X	•	•	•	11	cc	100		3	3	10	If cc is false.
	cc is false										← n →						
	continue,										← n →			3	5	17	If cc is true.
	otherwise																
	same as																
	CALL nn																
RET	PC_L ← (SP)	•	•	X	•	X	•	•	•	11	001	001	C9	1	3	10	
	PC_H ← (SP+1)																
RET cc	If condition	•	•	X	•	X	•	•	•	11	cc	000		1	1	5	If cc is false.
	cc is false																
	continue,													1	3	11	If cc is true.
	otherwise																
	same as RET																
RETI	Return from	•	•	X	•	X	•	•	•	11	101	101	ED	2	4	14	
	interrupt									01	001	101	4D				
RETN[1]	Return from	•	•	X	•	X	•	•	•	11	101	101	ED	2	4	14	
	non-maskable									01	000	101	45				
	interrupt																
RST p	(SP−1)← PC_H	•	•	X	•	X	•	•	•	11	t	111		1	3	11	
	(SP−2)← PC_L																
	PC_H ← 0																
	PC_L ← p																

cc	Condition
000	NZ (non-zero)
001	Z (zero)
010	NC (non-carry)
011	C (carry)
100	PO (parity odd)
101	PE (parity even)
110	P (sign positive)
111	M (sign negative)

t	p
000	00H
001	08H
010	10H
011	18H
100	20H
101	28H
110	30H
111	38H

NOTE: [1]RETN loads $IFF_2 \rightarrow IFF_1$

INPUT AND OUTPUT GROUP

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	76	543	210	Hex	No. of Bytes	No. of M Cycles	No. of T States	Comments
IN A, (n)	A ← (n)	•	•	X	•	X	•	•	•	11	011 ← n →	01	DB	2	3	11	n to $A_0 \sim A_7$ Acc. to $A_8 \sim A_{15}$
IN r, (C)	r ← (C) if r = 110 only the flags will be affected	↕	↕	X	↕	X	P	0	•	11 01	101 r	101 000	ED	2	3	12	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
INI	(HL) ← (C) B ← B − 1 HL ← HL + 1	X	① ↕	X	X	X	X	1	X	11 10	101 100	101 010	ED A2	2	4	16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
INIR	(HL) ← (C) B ← B − 1 HL ← HL + 1 Repeat until B = 0	X	② 1	X	X	X	X	1	X	11 10	101 110	101 010	ED B2	2 2	5 (If B≠0) 4 (If B = 0)	21 16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
IND	(HL) ← (C) B ← B − 1 HL ← HL − 1	X	① ↕	X	X	X	X	1	X	11 10	101 101	101 010	ED AA	2	4	16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
INDR	(HL) ← (C) B ← B − 1 HL ← HL − 1 Repeat until B = 0	X	② 1	X	X	X	X	1	X	11 10	101 111	101 010	ED BA	2 2	5 (If B≠0) 4 (If B = 0)	21 16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
OUT (n), A	(n) ← A	•	•	X	•	X	•	•	•	11	010 ← n →	011	D3	2	3	11	n to $A_0 \sim A_7$ Acc. to $A_8 \sim A_{15}$
OUT (C), r	(C) ← r	•	•	X	•	X	•	•	•	11 01	101 r	101 001	ED	2	3	12	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
OUTI	(C) ← (HL) B ← B − 1 HL ← HL + 1	X	① ↕	X	X	X	X	1	X	11 10	101 100	101 011	ED A3	2	4	16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
OTIR	(C) ← (HL) B ← B − 1 HL ← HL + 1 Repeat until B = 0	X	② 1	X	X	X	X	1	X	11 10	101 110	101 011	ED B3	2 2	5 (If B≠0) 4 (If B = 0)	21 16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
OUTD	(C) ← (HL) B ← B − 1 HL ← HL − 1	X	① ↕	X	X	X	X	1	X	11 10	101 101	101 011	ED AB	2	4	16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
OTDR	(C) ← (HL) B ← B − 1 HL ← HL − 1 Repeat until B = 0	X	② 1	X	X	X	X	1	X	11 10	101 111	101 011	ED	2 2	5 (If B≠0) 4 (If B = 0)	21 16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$

NOTES: ① If the result of B − 1 is zero, the Z flag is set; otherwise it is reset.
② Z flag is set upon instruction completion only.

CPU REGISTERS

Figure 4 shows three groups of registers within the CPU. The first group consists of duplicate sets of 8-bit registers: a principal set and an alternate set [designated by ' (prime), e.g., A']. Both sets consist of the Accumulator register, the Flag register, and six general-purpose registers. Transfer of data between these duplicate sets of registers is accomplished by use of "Exchange" instructions. The result is faster response to interrupts and easy, efficient implementation of such versatile programming techniques as background-foreground data processing. The second set of registers consists of six registers with assigned functions. These are the I (Interrupt register), the R (Refresh register), the IX and IY (Index registers), the SP (Stack Pointer), and the PC (Program Counter). The third group consists of two interrupt status flip-flops, plus an additional pair of flip-flops which assists in identifying the interrupt mode at any particular time. Table 1 provides further information on these registers.

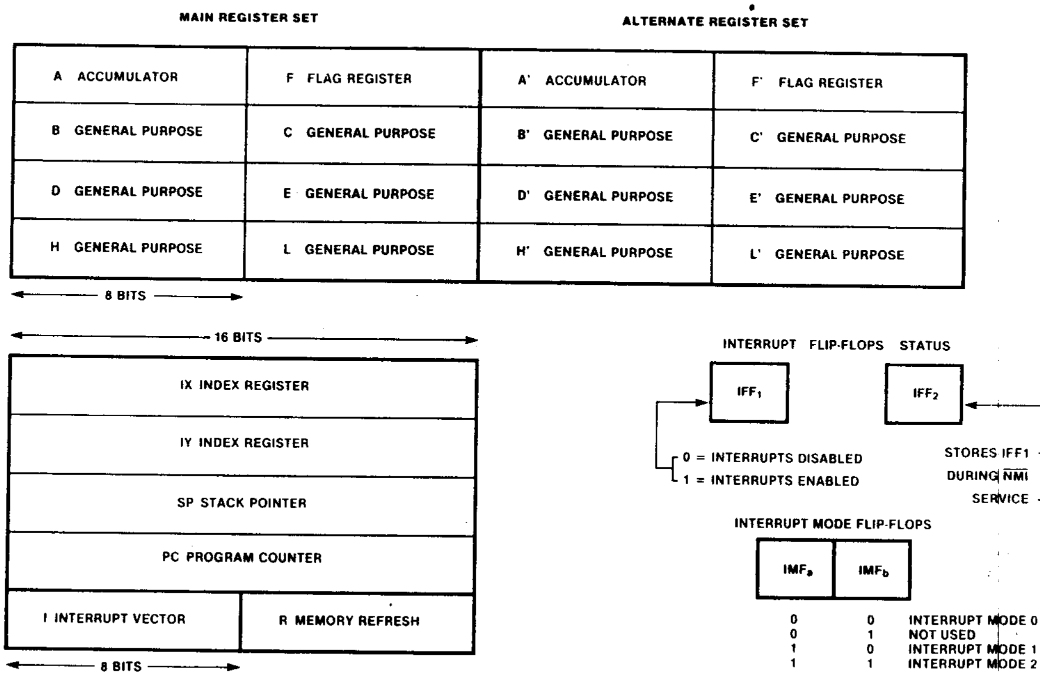

Figure 4. CPU Registers

INTERRUPTS: GENERAL OPERATION

The CPU accepts two interrupt input signals: $\overline{NMI}$ and $\overline{INT}$. The $\overline{NMI}$ is a non-maskable interrupt and has the highest priority. $\overline{INT}$ is a lower priority interrupt and it requires that interrupts be enabled in software in order to operate. $\overline{INT}$ can be connected to multiple peripheral devices in a wired-OR configuration.

The Z80 has a single response mode for interrupt service on the non-maskable interrupt. The maskable interrupt, $\overline{INT}$, has three programmable response modes available. These are:

- Mode 0 — similar to the 8080 microprocessor.
- Mode 1 — Peripheral Interrupt service, for use with non-8080/Z80 systems.
- **Mode 2 - a vectored interrupt scheme, usually daisy-chained, for use with the Z80 Family and compatible peripheral devices.**

The CPU services interrupts by sampling the $\overline{NMI}$ and $\overline{INT}$ signals at the rising edge of the last clock of an instruction. Further interrupt service processing depends upon the type of interrupt that was detected. Details on interrupt responses are shown in the CPU Timing Section.

Non-Maskable Interrupt ($\overline{NMI}$). The nonmaskable interrupt cannot be disabled by program control and therefore will be accepted at all times by the CPU. $\overline{NMI}$ is usually reserved for servicing only the highest priority type interrupts, such as that for orderly shutdown after power

Memory Read or Write Cycles. Figure 6 shows the timing of memory read or write cycles other than an opcode fetch ($\overline{M1}$) cycle. The $\overline{MREQ}$ and $\overline{RD}$ signals function exactly as in the fetch cycle. In a memory write cycle, $\overline{MREQ}$ also becomes active when the address bus is stable. The $\overline{WR}$ line is active when the data bus is stable, so that it can be used directly as an R/$\overline{W}$ pulse to most semiconductor memories.

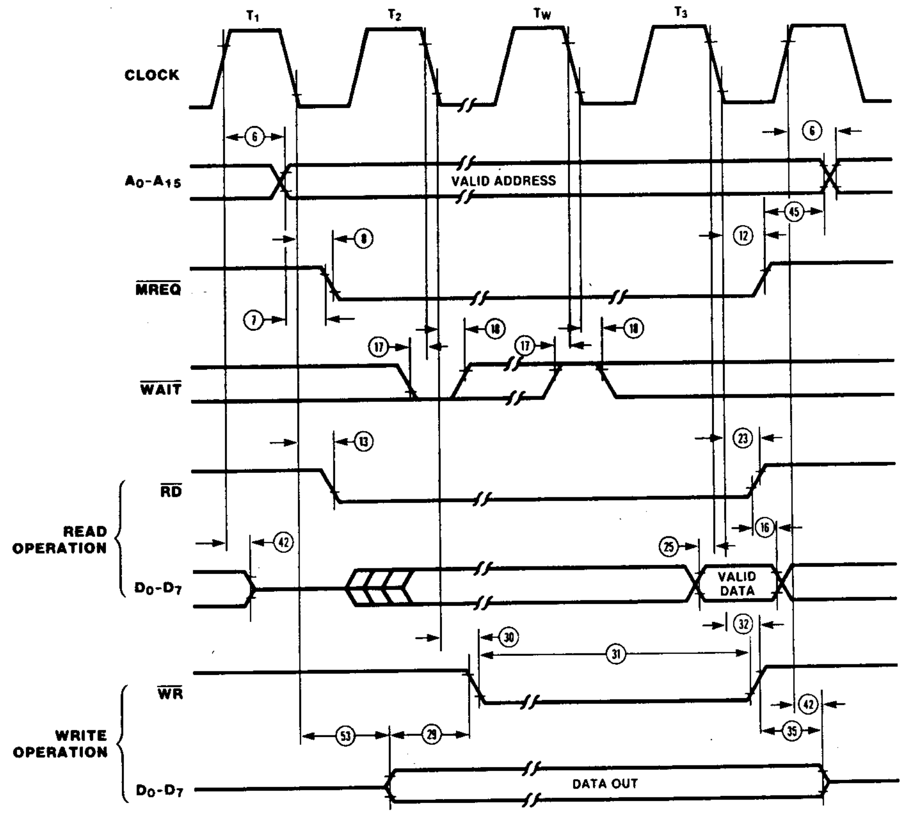

Figure 6. Memory Read or Write Cycles

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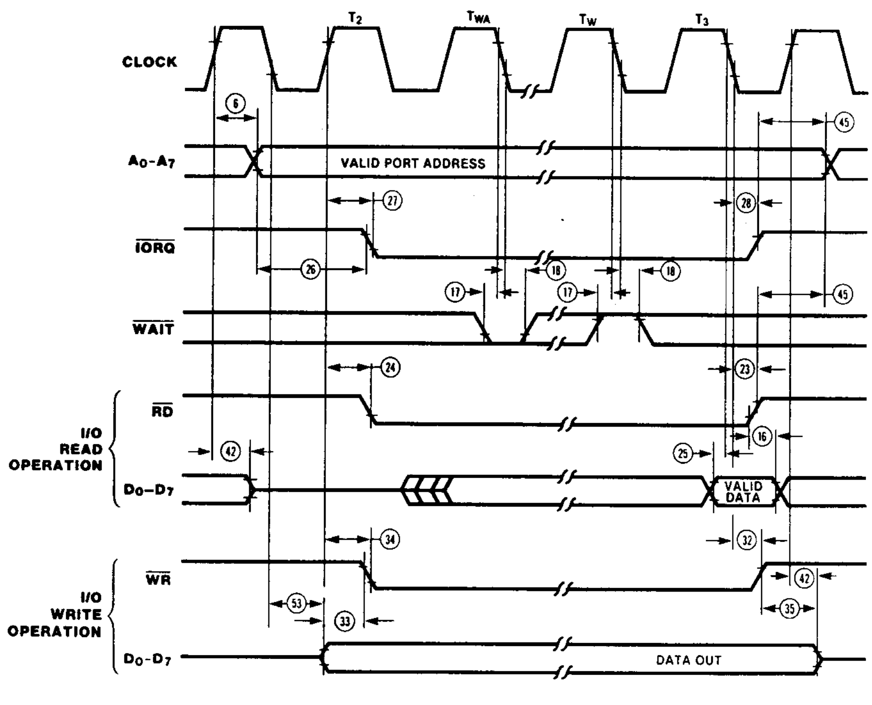
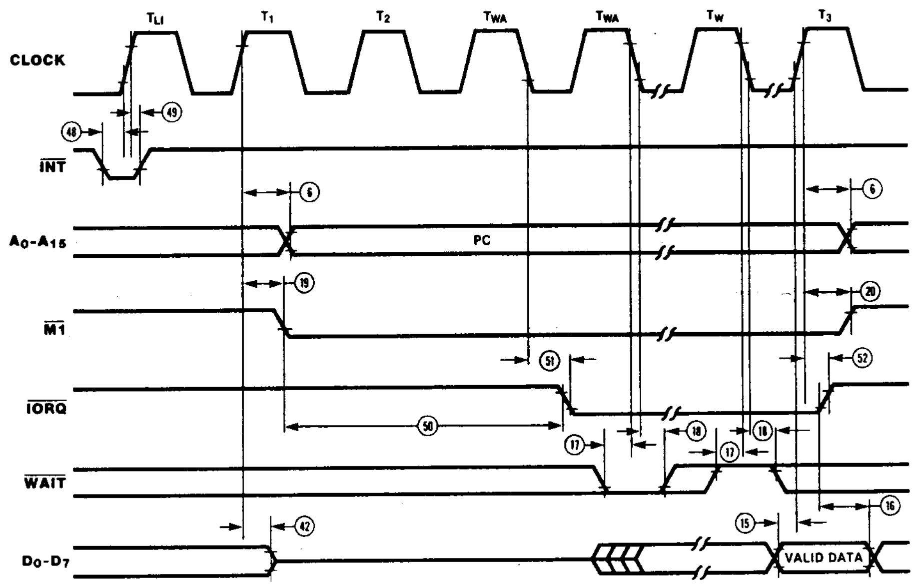
Input or Output Cycles. Figure 7 shows the timing for an I/O read or I/O write operation. During I/O operations, the CPU automatically inserts a single Wait state (T_{WA}). This extra Wait state allows sufficient time for an I/O port to decode the address from the port address lines.



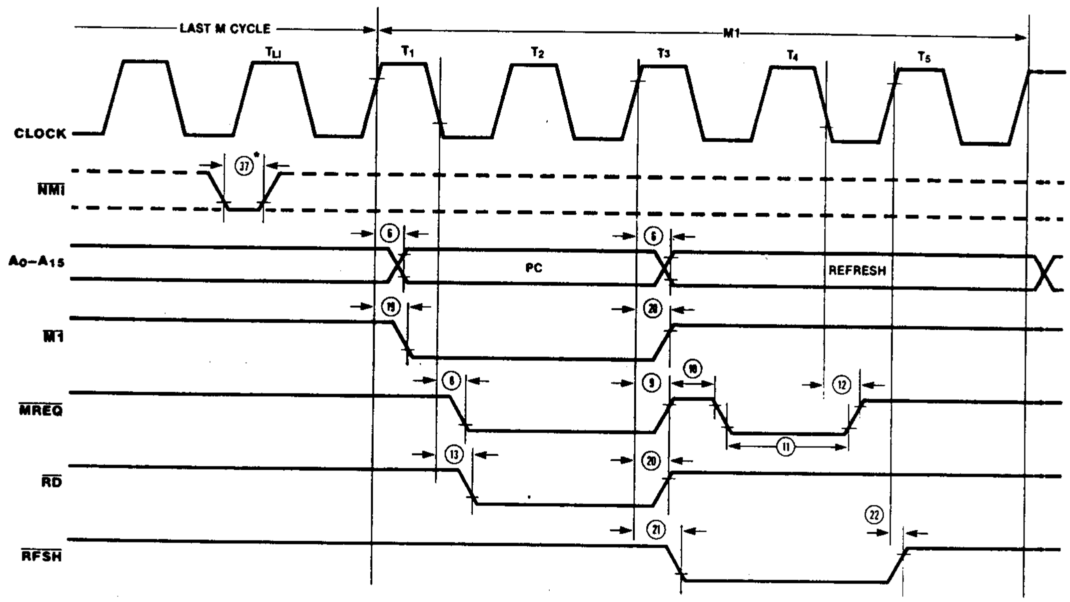
T_{WA} = One wait cycle automatically inserted by CPU.

Figure 7. Input or Output Cycles

Interrupt Request/Acknowledge Cycle. The CPU samples the interrupt signal with the rising edge of the last clock cycle at the end of any instruction (Figure 8). When an interrupt is accepted, a special $\overline{M1}$ cycle is generated. During this $\overline{M1}$ cycle, $\overline{IORQ}$ becomes active (instead of $\overline{MREQ}$) to indicate that the interrupting device can place an 8-bit vector on the data bus. The CPU automatically adds two Wait states to this cycle.



Non-Maskable Interrupt Request Cycle. $\overline{NMI}$ is sampled at the same time as the maskable interrupt input $\overline{INT}$ but has higher priority and cannot be disabled under software control. The subsequent timing is similar to that of a normal memory read operation except that data put on the bus by the memory is ignored. The CPU instead executes a restart (RST) operation and jumps to the $\overline{NMI}$ service routine located at address 0066H (Figure 9).



*Although $\overline{NMI}$ is an asynchronous input, to guarantee its being recognized on the following machine cycle, $\overline{NMI}$'s falling edge must occur no later than the rising edge of the clock cycle preceding the last state of any instruction cycle (T_{LI}).

Figure 9. Non-Maskable Interrupt Request Operation

ABSOLUTE MAXIMUM RATINGS

Voltage on V_{CC} with respect to V_{SS} . . . . – 0.3V to + 7V
Voltages on all inputs with respect
to V_{SS} . . . . – 0.3V to V_{CC} + 0.3V
Operating Ambient
Temperature . . . . .See Ordering Information
Storage Temperature . . . . – 65 °C to + 150 °C

Stresses greater than those listed under Absolute Maximum Ratings may cause permanent damage to the device. This is a stress rating only; operation of the device at any condition above those indicated in the operational sections of these specifications is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

STANDARD TEST CONDITIONS

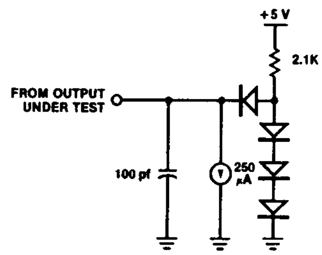
The DC Characteristics and capacitance sections below apply for the following standard test conditions, unless otherwise noted. All voltages are referenced to GND (0V). Positive current flows into the referenced pin.

Available operating temperature ranges are:

- **S = 0°C to +70°C**
  Voltage Supply Range:
  NMOS: +4.75V ≤ VCC ≤ +5.25V
  CMOS: +4.50V ≤ VCC ≤ +5.50V
- E= -40°C to 100°C, +4.50V ≤ VCC ≤ +5.50V

All ac parameters assume a load capacitance of 100 pf. Add 10 ns delay for each 50 pf increase in load up to a maximum of 200 pf for the data bus and 100 pf for address and control lines. AC timing measurements are referenced to 1.5 volts (except for clock, which is referenced to the 10% and 90% points).

The Ordering Information section lists temperature ranges and product numbers. Package drawings are in the Package Information section. Refer to the Literature List for additional documentation.



AC CHARACTERISTICS† (Z84C00/CMOS Z80 CPU)

V_{cc}=5.0V ± 10%, unless otherwise specified

No	Symbol	Parameter	Z84C0004** Min	Max	Z84C0006 Min	Max	Z84C0008 Min	Max	Z84C0010 Min	Max	Z84C0020[1] Min	Max	Unit	Note
1	TcC	Clock Cycle time	250*	DC	162*	DC	125*	DC	100*	DC	50*	DC	nS	
2	TwCh	Clock Pulse width (high)	110	DC	65	DC	55	DC	40	DC	20	DC	nS	
3	TwCl	Clock Pulse width (low)	110	DC	65	DC	55	DC	40	DC	20	DC	nS	
4	TfC	Clock Fall time		30		20		10		10		10	nS	
5	TrC	Clock Rise time		30		20		10		10		10	nS	
6	TdCr(A)	Address vaild from Clock Rise		110		90		80		65		57	nS	[2]
7	TdA(MREQf)	Address valid to /MREQ Fall	65*		35*		20*		5*		-15*		nS	
8	TdCf(MREQf)	Clock Fall to /MREQ Fall delay		85		70		60		55		40	nS	
9	TdCr(MREQr)	Clock Rise to /MREQ Rise delay		85		70		60		55		40	nS	
10	TwMREQh	/MREQ pulse width (High)	110*		65*		45**		30*		10*		nS	[3]
11	TwMREQl	/MREQ pulse width (low)	220*		132*		100*		75*		25*		nS	[3]
12	TdCf(MERQr)	Clock Fall to /MREQ Rise delay		85		70		60		55		40	nS	
13	TdCf(RDf)	Clock Fall to /RD Fall delay		95		80		70		65		40	nS	
14	TdCr(RDr)	Clock Rise to /RD Rise delay		85		70		60		55		40	nS	
15	TsD(Cr)	Data setup time to Clock Rise	35		30		30		25		12		nS	
16	ThD(RDr)	Data hold time after /RD Rise	0		0		0		0		0		nS	
17	TsWAIT(Cf)	/WAIT setup time to Clock Fall	70		60		50		20		7.5		nS	
18	ThWAIT(Cf)	/WAIT hold time after Clock Fall	10		10		10		10		10		nS	
19	TdCr(M1f)	Clock Rise to /M1 Fall delay		100		80		70		65		45	nS	
20	TdCr(M1r)	Clock Rise to /M1 Rise delay		100		80		70		65		45	nS	
21	TdCr(RFSHf)	Clock Rise to /RFSH Fall delay		130		110		95		80		60	nS	
22	TdCr(RFSHr)	Clock Rise to /RFSH Rise delay		120		100		85		80		60	nS	
23	TdCf(RDr)	Clock Fall to /RD Rise delay		85		70		60		55		40	nS	
24	TdCr(RDf)	Clock Rise to /RD Fall delay		85		70		60		55		40	nS	
25	TsD(Cf)	Data setup to Clock Fall during M2, M3, M4 or M5 cycles	50		40		30		25		12		nS	
26	TdA(IORQf)	Address stable prior to /IORQ Fall	180*		107*		75*		50*		0*		nS	
27	TdCr(IORQf)	Clock Rise to /IORQ Fall delay		75		65		55		50		40	nS	
28	TdCf(IORQr)	Clock Fall to /IORQ Rise delay		85		70		60		55		40	nS	
29	TdD(WRf)Mw	Data stable prior to /WR Fall	80*		22*		5*		40*		-10*		nS	
30	TdCf(WRf)	Clock Fall to /WR Fall delay		80		70		60		55		40	nS	
31	TwWR	/WR pulse width	220*		132*		100*		75*		25*		nS	
32	TdCf(WRr)	Clock Fall to /WR Rise delay		80		70		60		55		40	nS	
33	TdD(WRf)IO	Data stable prior to /WR Fall	-10*		-55*		-55*		-10*		-30*		nS	
34	TdCr(WRf)	Clock Rise to /WR Fall delay		65		60		60		50		40	nS	
35	TdWRr(D)	Data stable from /WR Rise	60*		30*		15*		10*		0*		nS	
36	TdCf(HALT)	Clock Fall to /HALT 'L' or 'H'		300		260		225		90		70	nS	
37	TwNMI	/NMI pulse width	80		60		60		60		60		nS	
38	TsBUSREQ (Cr)	/BUSREQ setup time to Clock Rise	50		50		40		30		15		nS	

*For clock periods other than the minimums shown, calculate parameters using the table on the following page. Calculated values above assumed TrC = TfC = 20 ns.

†Units in nanoseconds (ns).

†† **For loading ≥ 50 pf. Decrease width by 10 ns for each additional 50 pf.**

**4 MHz CMOS Z80 is obsoleted and replaced by 6 MHz

AC CHARACTERISTICS† (Z8400/NMOS Z80 CPU)

Number	Symbol	Parameter	Z0840004 Min	Z0840004 Max	Z0840006 Min	Z0840006 Max	Z0840008 Min	Z0840008 Max
1	TcC	Clock Cycle Time	250*		162*		125*	
2	TwCh	Clock Pulse Width (High)	110	2000	65	2000	55	2000
3	TwCl	Clock Pulse Width (Low)	110	2000	65	2000	55	2000
4	TfC	Clock Fall Time		30		20		10
5	TrC	Clock Rise Time		30		20		10
6	TdCr(A)	Clock ↑ to Address Valid Delay		110		90		80
7	TdA(MREQf)	Address Valid to $\overline{MREQ}$ ↓ Delay	65*		35*		20*	
8	TdCf(MREQf)	Clock ↓ to $\overline{MREQ}$ ↓ Delay		85		70		60
9	TdCr(MREQr)	Clock ↑ to $\overline{MREQ}$ ↑ Delay		85		70		60
10	TwMREQh	$\overline{MREQ}$ Pulse Width (High)	110*††		65*††		45*††	
11	TwMREQl	$\overline{MREQ}$ Pulse Width (Low)	220*††		135*††		100*††	
12	TdCf(MREQr)	Clock ↓ to $\overline{MREQ}$ ↑ Delay		85		70		60
13	TdCf(RDf)	Clock ↓ to $\overline{RD}$ ↓ Delay		95		80		70
14	TdCr(RDr)	Clock ↑ to $\overline{RD}$ ↑ Delay		85		70		60
15	TsD(Cr)	Data Setup Time to Clock ↑	35		30		30	
16	ThD(RDr)	Data Hold Time to $\overline{RD}$ ↑		0		0		0
17	TsWAIT(Cf)	$\overline{WAIT}$ Setup Time to Clock ↓	70		60		50	
18	ThWAIT(Cf)	$\overline{WAIT}$ Hold Time after Clock ↓		0		0		0
19	TdCr(M1f)	Clock ↑ to $\overline{M1}$ ↓ Delay		100		80		70
20	TdCr(M1r)	Clock ↑ to $\overline{M1}$ ↑ Delay		100		80		70
21	TdCr(RFSHf)	Clock ↑ to $\overline{RFSH}$ ↓ Delay		130		110		95
22	TdCr(RFSHr)	Clock ↑ to $\overline{RFSH}$ ↑ Delay		120		100		85
23	TdCf(RDr)	Clock ↓ to $\overline{RD}$ ↑ Delay		85		70		60
24	TdCr(RDf)	Clock ↑ to $\overline{RD}$ ↓ Delay		85		70		60
25	TsD(Cf)	Data Setup to Clock ↓ during M_2, M_3, M_4, or M_5 Cycles	50		40		30	
26	TdA(IORQf)	Address Stable prior to $\overline{IORQ}$ ↓	180*		110*		75*	
27	TdCr(IORQf)	Clock ↑ to $\overline{IORQ}$ ↓ Delay		75		65		55
28	TdCf(IORQr)	Clock ↓ to $\overline{IORQ}$ ↑ Delay		85		70		60
29	TdD(WRf)	Data Stable prior to $\overline{WR}$ ↓	80*		25*		5*	
30	TdCf(WRf)	Clock ↓ to $\overline{WR}$ ↓ Delay		80		70		60
31	TwWR	$\overline{WR}$ Pulse Width	220*		135*		100*	
32	TdCf(WRr)	Clock ↓ to $\overline{WR}$ ↑ Delay		80		70		60
33	TdD(WRf)	Data Stable prior to $\overline{WR}$ ↓	−10*		−55*		55*	
34	TdCr(WRf)	Clock ↑ to $\overline{WR}$ ↓ Delay		65		60		55
35	TdWRr(D)	Data Stable from $\overline{WR}$ ↑	60*		30*		15*	
36	TdCf(HALT)	Clock ↓ to $\overline{HALT}$ ↑ or ↓		300		260		225
37	TwNMI	$\overline{NMI}$ Pulse Width	80		70		60*	
38	TsBUSREQ(Cr)	$\overline{BUSREQ}$ Setup Time to Clock ↑	50		50		40	

*For clock periods other than the minimums shown, calculate parameters using the table on the following page. Calculated values above assumed TrC = TfC = 20 ns.

†Units in nanoseconds (ns).

†† For loading ≥ 50 pf., Decrease width by 10 ns for each additional 50 pf.

Customer Support

For answers to technical questions about the product, documentation, or any other issues with Zilog's offerings, please visit Zilog's Knowledge Base at http://www.zilog.com/kb.

For any comments, detail technical questions, or reporting problems, please visit Zilog's Technical Support at http://support.zilog.com.