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### What is "[Embedded - Microcontrollers](#)"?

"[Embedded - Microcontrollers](#)" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

### Applications of "[Embedded - Microcontrollers](#)"

#### Details

Product Status	Not For New Designs
Core Processor	CPU32
Core Size	32-Bit Single-Core
Speed	20MHz
Connectivity	EBI/EMI, SCI, SPI, UART/USART
Peripherals	POR, PWM, WDT
Number of I/O	15
Program Memory Size	-
Program Memory Type	ROMless
EEPROM Size	-
RAM Size	2K x 8
Voltage - Supply (Vcc/Vdd)	4.5V ~ 5.5V
Data Converters	-
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	144-LQFP
Supplier Device Package	144-LQFP (20x20)
Purchase URL	<a href="https://www.e-xfl.com/product-detail/nxp-semiconductors/mc68332acag20">https://www.e-xfl.com/product-detail/nxp-semiconductors/mc68332acag20</a>

**Table 1 Ordering Information**

Package Type	TPU Type	Temperature	Frequency (MHz)	Package Order Quantity	Order Number
132-Pin PQFP	Motion Control	-40 to +85 °C	16 MHz	2 pc tray	SPAKMC332GCFC16
				36 pc tray	MC68332GCFC16
			20 MHz	2 pc tray	SPAKMC332GCFC20
				36 pc tray	MC68332GCFC20
		-40 to +105 °C	16 MHz	2 pc tray	SPAKMC332GVFC16
				36 pc tray	MC68332GVFC16
			20 MHz	2 pc tray	SPAKMC332GVFC20
				36 pc tray	MC68332GVFC20
		-40 to +125 °C	16 MHz	2 pc tray	SPAKMC332GMFC16
				36 pc tray	MC68332GMFC16
			20 MHz	2 pc tray	SPAKMC332GMFC20
				36 pc tray	MC68332GMFC20
	Standard	-40 to +85 °C	16 MHz	2 pc tray	SPAKMC332CFC16
				36 pc tray	MC68332CFC16
			20 MHz	2 pc tray	SPAKMC332CFC20
				36 pc tray	MC68332CFC20
		-40 to +105 °C	16 MHz	2 pc tray	SPAKMC332VFC16
				36 pc tray	MC68332VFC16
			20 MHz	2 pc tray	SPAKMC332VFC20
				36 pc tray	MC68332VFC20
		-40 to +125 °C	16 MHz	2 pc tray	SPAKMC332MFC16
				36 pc tray	MC68332MFC16
			20 MHz	2 pc tray	SPAKMC332MFC20
				36 pc tray	MC68332MFC20
	Std w/enhanced PPWA	-40 to +85 °C	16 MHz	2 pc tray	SPAKMC332ACFC16
				36 pc tray	MC68332ACFC16
			20 MHz	2 pc tray	SPAKMC332ACFC20
				36 pc tray	MC68332ACFC20
		-40 to +105 °C	16 MHz	2 pc tray	SPAKMC332AVFC16
				36 pc tray	MC68332AVFC16
			20 MHz	2 pc tray	SPAKMC332AVFC20
				36 pc tray	MC68332AVFC20
		-40 to +125 °C	16 MHz	2 pc tray	SPAKMC332AMFC16
				36 pc tray	MC68332AMFC16
			20 MHz	2 pc tray	SPAKMC332AMFC20
				36 pc tray	MC68332AMFC20

## 1.2 Block Diagram

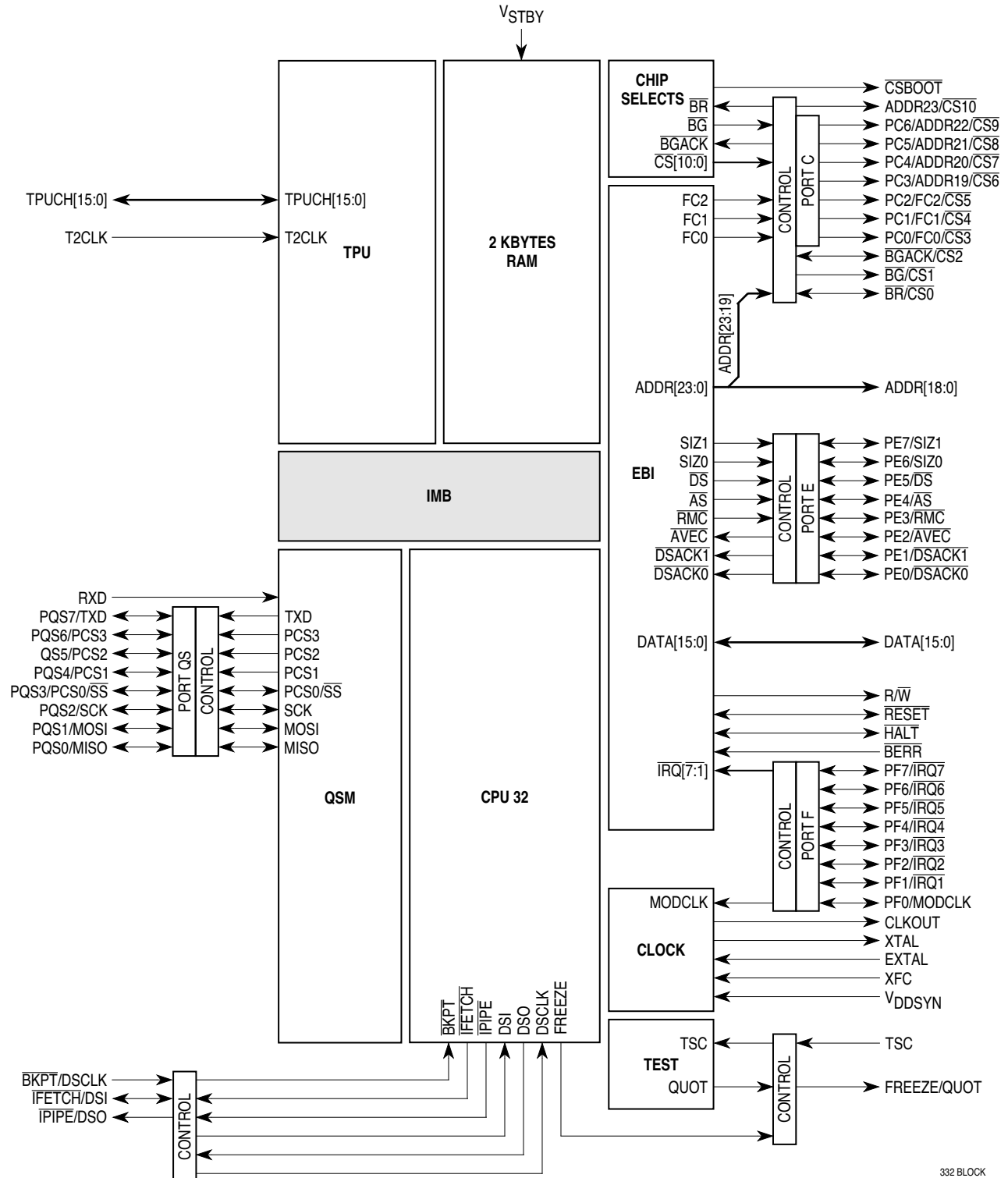


Figure 1 MCU Block Diagram

## 2 Signal Descriptions

### 2.1 Pin Characteristics

The following table shows MCU pins and their characteristics. All inputs detect CMOS logic levels. All inputs can be put in a high-impedance state, but the method of doing this differs depending upon pin function. Refer to the table, MCU Driver Types, for a description of output drivers. An entry in the discrete I/O column of the MCU Pin Characteristics table indicates that a pin has an alternate I/O function. The port designation is given when it applies. Refer to the MCU Block Diagram for information about port organization.

**Table 2 MCU Pin Characteristic**

Pin Mnemonic	Output Driver	Input Synchronized	Input Hysteresis	Discrete I/O	Port Designation
ADDR23/CS10/ECLK	A	Y	N	O	—
ADDR[22:19]/CS[9:6]	A	Y	N	O	PC[6:3]
ADDR[18:0]	A	Y	N	—	—
$\overline{AS}$	B	Y	N	I/O	PE5
$\overline{AVEC}$	B	Y	N	I/O	PE2
$\overline{BERR}$	B	Y	N	—	—
$\overline{BG/CS1}$	B	—	—	—	—
$\overline{BGACK/CS2}$	B	Y	N	—	—
$\overline{BKPT/DSCLK}$	—	Y	Y	—	—
$\overline{BR/CS0}$	B	Y	N	—	—
CLKOUT	A	—	—	—	—
$\overline{CSBOOT}$	B	—	—	—	—
DATA[15:0] <sup>1</sup>	Aw	Y	N	—	—
$\overline{DS}$	B	Y	N	I/O	PE4
$\overline{DSACK1}$	B	Y	N	I/O	PE1
$\overline{DSACK0}$	B	Y	N	I/O	PE0
DSI/IFETCH	A	Y	Y	—	—
DSO/IPIPE	A	—	—	—	—
EXTAL <sup>2</sup>	—	—	Special	—	—
FC[2:0]/CS[5:3]	A	Y	N	O	PC[2:0]
FREEZE/QUOT	A	—	—	—	—
$\overline{HALT}$	Bo	Y	N	—	—
$\overline{IRQ[7:1]}$	B	Y	Y	I/O	PF[7:1]
MISO	Bo	Y	Y	I/O	PQS0
MODCLK <sup>1</sup>	B	Y	N	I/O	PF0
MOSI	Bo	Y	Y	I/O	PQS1
PCS0/ $\overline{SS}$	Bo	Y	Y	I/O	PQS3
PCS[3:1]	Bo	Y	Y	I/O	PQS[6:4]
R/W	A	Y	N	—	—
$\overline{RESET}$	Bo	Y	Y	—	—
$\overline{RMC}$	B	Y	N	I/O	PE3
RXD	—	N	N	—	—
SCK	Bo	Y	Y	I/O	PQS2
SIZ[1:0]	B	Y	N	I/O	PE[7:6]

Table 7 SIM Address Map (Continued)

Access	Address	15	8 7	0
S	\$YFFA56	CHIP-SELECT OPTION 2 (CSOR2)		
S	\$YFFA58	CHIP-SELECT BASE 3 (CSBAR3)		
S	\$YFFA5A	CHIP-SELECT OPTION 3 (CSOR3)		
S	\$YFFA5C	CHIP-SELECT BASE 4 (CSBAR4)		
S	\$YFFA5E	CHIP-SELECT OPTION 4 (CSOR4)		
S	\$YFFA60	CHIP-SELECT BASE 5 (CSBAR5)		
S	\$YFFA62	CHIP-SELECT OPTION 5 (CSOR5)		
S	\$YFFA64	CHIP-SELECT BASE 6 (CSBAR6)		
S	\$YFFA66	CHIP-SELECT OPTION 6 (CSOR6)		
S	\$YFFA68	CHIP-SELECT BASE 7 (CSBAR7)		
S	\$YFFA6A	CHIP-SELECT OPTION 7 (CSOR7)		
S	\$YFFA6C	CHIP-SELECT BASE 8 (CSBAR8)		
S	\$YFFA6E	CHIP-SELECT OPTION 8 (CSOR8)		
S	\$YFFA70	CHIP-SELECT BASE 9 (CSBAR9)		
S	\$YFFA72	CHIP-SELECT OPTION 9 (CSOR9)		
S	\$YFFA74	CHIP-SELECT BASE 10 (CSBAR10)		
S	\$YFFA76	CHIP-SELECT OPTION 10 (CSOR10)		
	\$YFFA78	NOT USED		NOT USED
	\$YFFA7A	NOT USED		NOT USED
	\$YFFA7C	NOT USED		NOT USED
	\$YFFA7E	NOT USED		NOT USED

Y = M111, where M is the logic state of the module mapping (MM) bit in the SIMCR.

### 3.2 System Configuration and Protection

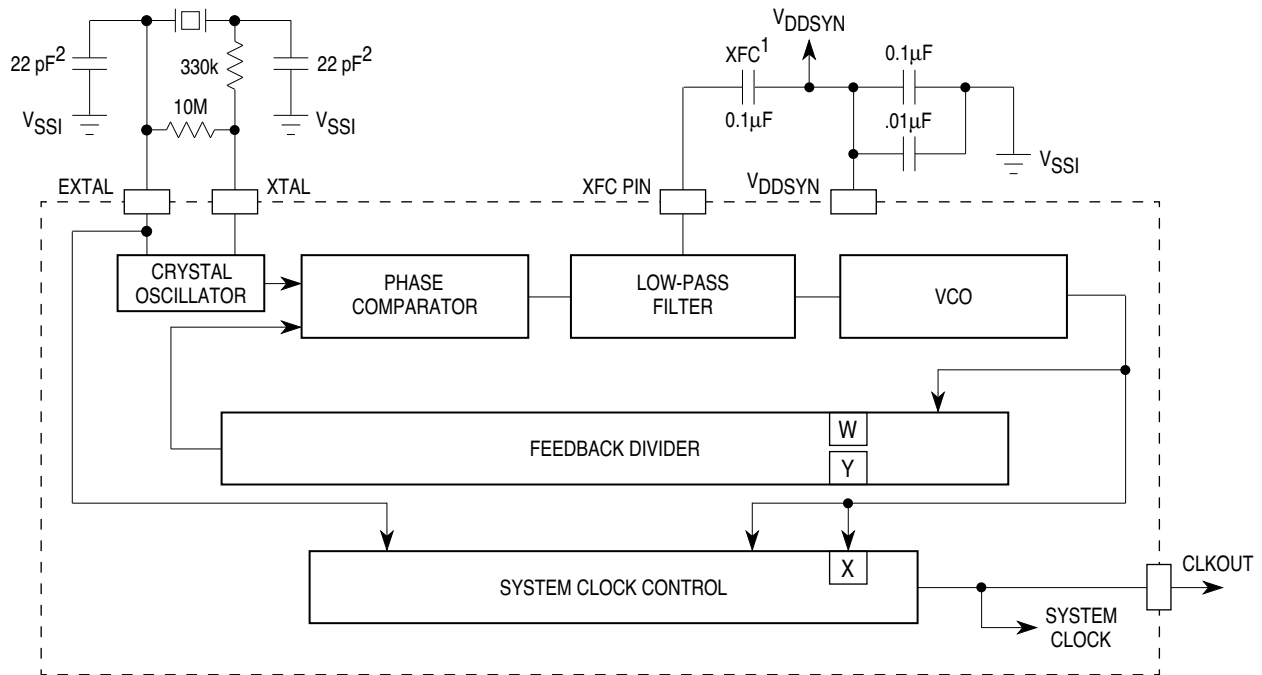
This functional block provides configuration control for the entire MCU. It also performs interrupt arbitration, bus monitoring, and system test functions. MCU system protection includes a bus monitor, a HALT monitor, a spurious interrupt monitor, and a software watchdog timer. These functions have been made integral to the microcontroller to reduce the number of external components in a complete control system.

## 3.3 System Clock

The system clock in the SIM provides timing signals for the IMB modules and for an external peripheral bus. Because MCU operation is fully static, register and memory contents are not affected when the clock rate changes. System hardware and software support changes in the clock rate during operation.

The system clock signal can be generated in three ways. An internal phase-locked loop can synthesize the clock from an internal or external frequency source, or the clock signal can be input from an external source.

Following is a block diagram of the clock submodule.



1. MUST BE LOW-LEAKAGE CAPACITOR (INSULATION RESISTANCE 30,000 MΩ OR GREATER).
2. RESISTANCE AND CAPACITANCE BASED ON A TEST CIRCUIT CONSTRUCTED WITH A DAISHINKU DMX-38 32.768-kHz CRYSTAL. SPECIFIC COMPONENTS MUST BE BASED ON CRYSTAL TYPE. CONTACT CRYSTAL VENDOR FOR EXACT CIRCUIT.

SYS CLOCK  
BLOCK 32KHZ

Figure 7 System Clock Block Diagram

### 3.3.1 Clock Sources

The state of the clock mode (MODCLK) pin during reset determines the clock source. When MODCLK is held high during reset, the clock synthesizer generates a clock signal from either a crystal oscillator or an external reference input. Clock synthesizer control register SYNCR determines operating frequency and various modes of operation. When MODCLK is held low during reset, the clock synthesizer is disabled, and an external system clock signal must be applied. When the synthesizer is disabled, SYNCR control bits have no effect.

A reference crystal must be connected between the EXTAL and XTAL pins to use the internal oscillator. Use of a 32.768-kHz crystal is recommended. These crystals are inexpensive and readily available. If an external reference signal or an external system clock signal is applied through the EXTAL pin, the XTAL pin must be left floating. External reference signal frequency must be less than or equal to maximum specified reference frequency. External system clock signal frequency must be less than or equal to maximum specified system clock frequency.

When an external system clock signal is applied (i.e., the PLL is not used), duty cycle of the input is critical, especially at near maximum operating frequencies. The relationship between clock signal duty cycle and clock signal period is expressed:

$$\text{Minimum external clock period} = \frac{\text{minimum external clock high/low time}}{50\% - \text{percentage variation of external clock input duty cycle}}$$

### 3.3.2 Clock Synthesizer Operation

A voltage controlled oscillator (VCO) generates the system clock signal. A portion of the clock signal is fed back to a divider/counter. The divider controls the frequency of one input to a phase comparator. The other phase comparator input is a reference signal, either from the internal oscillator or from an external source. The comparator generates a control signal proportional to the difference in phase between its two inputs. The signal is low-pass filtered and used to correct VCO output frequency.

The synthesizer locks when VCO frequency is identical to reference frequency. Lock time is affected by the filter time constant and by the amount of difference between the two comparator inputs. Whenever comparator input changes, the synthesizer must re-lock. Lock status is shown by the SLOCK bit in SYNCR.

The MCU does not come out of reset state until the synthesizer locks. Crystal type, characteristic frequency, and layout of external oscillator circuitry affect lock time.

The low-pass filter requires an external low-leakage capacitor, typically 0.1  $\mu\text{F}$ , connected between the XFC and  $V_{\text{DDSYN}}$  pins.

$V_{\text{DDSYN}}$  is used to power the clock circuits. A separate power source increases MCU noise immunity and can be used to run the clock when the MCU is powered down. Use a quiet power supply as the  $V_{\text{DDSYN}}$  source, since PLL stability depends on the VCO, which uses this supply. Place adequate external bypass capacitors as close as possible to the  $V_{\text{DDSYN}}$  pin to ensure stable operating frequency.

When the clock synthesizer is used, control register SYNCR determines operating frequency and various modes of operation. SYNCR can be read only when the processor is operating at the supervisor privilege level.

The SYNCR X bit controls a divide by two prescaler that is not in the synthesizer feedback loop. Setting X doubles clock speed without changing VCO speed. There is no VCO relock delay. The SYNCR W bit controls a 3-bit prescaler in the feedback divider. Setting W increases VCO speed by a factor of four. The SYNCR Y field determines the count modulus for a modulo 64 down counter, causing it to divide by a value of  $Y + 1$ . When either W or Y value changes, there is a VCO relock delay.

Clock frequency is determined by SYNCR bit settings as follows:

$$F_{\text{SYSTEM}} = F_{\text{REFERENCE}} [4(Y + 1)(2^{2W} + X)]$$

In order for the device to perform correctly, the clock frequency selected by the W, X, and Y bits must be within the limits specified for the MCU.

The VCO frequency is twice the system clock frequency if  $X = 1$  or four times the system clock frequency if  $X = 0$ .

The reset state of SYNCR (\$3F00) produces a modulus-64 count.

## DSACK — Data and Size Acknowledge

This field specifies the source of  $\overline{\text{DSACK}}$  in asynchronous mode. It also allows the user to adjust bus timing with internal  $\overline{\text{DSACK}}$  generation by controlling the number of wait states that are inserted to optimize bus speed in a particular application. The following table shows the  $\overline{\text{DSACK}}$  field encoding. The fast termination encoding (1110) is used for two-cycle access to external memory.

DSACK	Description
0000	No Wait States
0001	1 Wait State
0010	2 Wait States
0011	3 Wait States
0100	4 Wait States
0101	5 Wait States
0110	6 Wait States
0111	7 Wait States
1000	8 Wait States
1001	9 Wait States
1010	10 Wait States
1011	11 Wait States
1100	12 Wait States
1101	13 Wait States
1110	Fast Termination
1111	External $\overline{\text{DSACK}}$

## SPACE — Address Space

Use this option field to select an address space for the chip-select logic. The CPU32 normally operates in supervisor or user space, but interrupt acknowledge cycles must take place in CPU space.

Space Field	Address Space
00	CPU Space
01	User Space
10	Supervisor Space
11	Supervisor/User Space

## IPL — Interrupt Priority Level

If the space field is set for CPU space (00), chip-select logic can be used for interrupt acknowledge. During an interrupt acknowledge cycle, the priority level on address lines ADDR[3:1] is compared to the value in the IPL field. If the values are the same, a chip select is asserted, provided that other option register conditions are met. The following table shows IPL field encoding.

IPL	Description
000	Any Level
001	IPL1
010	IPL2
011	IPL3
100	IPL4
101	IPL5
110	IPL6
111	IPL7

This field only affects the response of chip selects and does not affect interrupt recognition by the CPU. Any level means that chip select is asserted regardless of the level of the interrupt acknowledge cycle.



## PPFAR — Port F Pin Assignment Register

\$YFFA1F

15	8	7	6	5	4	3	2	1	0
NOT USED		PFPA7	PFPA6	PFPA5	PFPA4	PFPA3	PFPA2	PFPA1	PFPA0

RESET:

DATA9 DATA9 DATA9 DATA9 DATA9 DATA9 DATA9 DATA9

The bits in this register control the function of each port F pin. Any bit cleared to zero defines the corresponding pin to be an I/O pin. Any bit set to one defines the corresponding pin to be an interrupt request signal or MODCLK. The MODCLK signal has no function after reset.

**Table 17 Port F Pin Assignments**

PFFAR Field	Port F Signal	Alternate Signal
PFFA7	PF7	$\overline{\text{IRQ7}}$
PFFA6	PF6	$\overline{\text{IRQ6}}$
PFFA5	PF5	$\overline{\text{IRQ5}}$
PFFA4	PF4	$\overline{\text{IRQ4}}$
PFFA3	PF3	$\overline{\text{IRQ3}}$
PFFA2	PF2	$\overline{\text{IRQ2}}$
PFFA1	PF1	$\overline{\text{IRQ1}}$
PFFA0	PF0	MODCLK

Data bus pin 9 controls the state of this register following reset. If DATA9 is set to one during reset, the register is set to \$FF, which defines all port F pins as interrupt request inputs. If DATA9 is cleared to zero during reset, this register is set to \$00, defining all port F pins as I/O pins.

### 3.7 Resets

Reset procedures handle system initialization and recovery from catastrophic failure. The MCU performs resets with a combination of hardware and software. The system integration module determines whether a reset is valid, asserts control signals, performs basic system configuration based on hardware mode-select inputs, then passes control to the CPU.

Reset occurs when an active low logic level on the  $\overline{\text{RESET}}$  pin is clocked into the SIM. Resets are gated by the CLKOUT signal. Asynchronous resets are assumed to be catastrophic. An asynchronous reset can occur on any clock edge. Synchronous resets are timed to occur at the end of bus cycles. If there is no clock when  $\overline{\text{RESET}}$  is asserted, reset does not occur until the clock starts. Resets are clocked in order to allow completion of write cycles in progress at the time  $\overline{\text{RESET}}$  is asserted.

Reset is the highest-priority CPU32 exception. Any processing in progress is aborted by the reset exception, and cannot be restarted. Only essential tasks are performed during reset exception processing. Other initialization tasks must be accomplished by the exception handler routine.

#### 3.7.1 SIM Reset Mode Selection

The logic states of certain data bus pins during reset determine SIM operating configuration. In addition, the state of the MODCLK pin determines system clock source and the state of the BKPT pin determines what happens during subsequent breakpoint assertions. The following table is a summary of reset mode selection options.

**Table 18 Reset Mode Selection**

Mode Select Pin	Default Function (Pin Left High)	Alternate Function (Pin Pulled Low)
-----------------	-------------------------------------	--

mask lower-priority interrupts during exception processing, and it is decoded by modules that have requested interrupt service to determine whether the current interrupt acknowledge cycle pertains to them.

Modules that have requested interrupt service decode the IP value placed on the address bus at the beginning of the interrupt acknowledge cycle, and if their requests are at the specified IP level, respond to the cycle. Arbitration between simultaneous requests of the same priority is performed by means of serial contention between module interrupt arbitration (IARB) field bit values.

Each module that can make an interrupt service request, including the SIM, has an IARB field in its configuration register. An IARB field can be assigned a value from %0001 (lowest priority) to %1111 (highest priority). A value of %0000 in an IARB field causes the CPU to process a spurious interrupt exception when an interrupt from that module is recognized.

Because the EBI manages external interrupt requests, the SIM IARB value is used for arbitration between internal and external interrupt requests. The reset value of IARB for the SIM is %1111, and the reset IARB value for all other modules is %0000. Initialization software must assign different IARB values in order to implement an arbitration scheme.

Each module must have a unique IARB value. When two or more IARB fields have the same nonzero value, the CPU interprets multiple vector numbers simultaneously, with unpredictable consequences.

Arbitration must always take place, even when a single source requests service. This point is important for two reasons: the CPU interrupt acknowledge cycle is not driven on the external bus unless the SIM wins contention, and failure to contend causes an interrupt acknowledge bus cycle to be terminated by a bus error, which causes a spurious interrupt exception to be taken.

When arbitration is complete, the dominant module must place an interrupt vector number on the data bus and terminate the bus cycle. In the case of an external interrupt request, because the interrupt acknowledge cycle is transferred to the external bus, an external device must decode the mask value and respond with a vector number, then generate bus cycle termination signals. If the device does not respond in time, a spurious interrupt exception is taken.

The periodic interrupt timer (PIT) in the SIM can generate internal interrupt requests of specific priority at predetermined intervals. By hardware convention, PIT interrupts are serviced before external interrupt service requests of the same priority. Refer to 3.2.7 Periodic Interrupt Timer for more information.

## 3.8.2 Interrupt Processing Summary

A summary of the interrupt processing sequence follows. When the sequence begins, a valid interrupt service request has been detected and is pending.

- A. The CPU finishes higher priority exception processing or reaches an instruction boundary.
- B. Processor state is stacked. The contents of the status register and program counter are saved.
- C. The interrupt acknowledge cycle begins:
  1. FC[2:0] are driven to %111 (CPU space) encoding.
  2. The address bus is driven as follows. ADDR[23:20] = %1111; ADDR[19:16] = %1111, which indicates that the cycle is an interrupt acknowledge CPU space cycle; ADDR[15:4] = %111111111111; ADDR[3:1] = the level of the interrupt request being acknowledged; and ADDR0 = %1.
  3. Request priority level is latched into the IP field in the status register from the address bus.
- D. Modules or external peripherals that have requested interrupt service decode the request level in ADDR[3:1]. If the request level of at least one interrupting module or device is the same as the value in ADDR[3:1], interrupt arbitration contention takes place. When there is no contention, the spurious interrupt monitor asserts BERR, and a spurious interrupt exception is processed.
- E. After arbitration, the interrupt acknowledge cycle can be completed in one of three ways:

## 4.6 Instruction Set Summary

Table 20 Instruction Set Summary

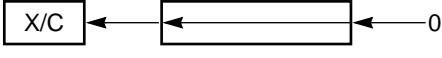
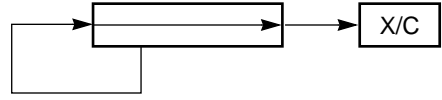
Instruction	Syntax	Operand Size	Operation
ABCD	Dn, Dn – (An), – (An)	8 8	Source <sub>10</sub> + Destination <sub>10</sub> + X ⇒ Destination
ADD	Dn, <ea> <ea>, Dn	8, 16, 32 8, 16, 32	Source + Destination ⇒ Destination
ADDA	<ea>, An	16, 32	Source + Destination ⇒ Destination
ADDI	#<data>, <ea>	8, 16, 32	Immediate data + Destination ⇒ Destination
ADDQ	# <data>, <ea>	8, 16, 32	Immediate data + Destination ⇒ Destination
ADDX	Dn, Dn – (An), – (An)	8, 16, 32 8, 16, 32	Source + Destination + X ⇒ Destination
AND	<ea>, Dn Dn, <ea>	8, 16, 32 8, 16, 32	Source • Destination ⇒ Destination
ANDI	# <data>, <ea>	8, 16, 32	Data • Destination ⇒ Destination
ANDI to CCR	# <data>, CCR	8	Source • CCR ⇒ CCR
ANDI to SR1 <sup>1</sup>	# <data>, SR	16	Source • SR ⇒ SR
ASL	Dn, Dn # <data>, Dn i	8, 16, 32 8, 16, 32 16	
ASR	Dn, Dn # <data>, Dn i	8, 16, 32 8, 16, 32 16	
Bcc	label	8, 16, 32	If condition true, then PC + d ⇒ PC
BCHG	Dn, <ea> # <data>, <ea>	8, 32 8, 32	(bit number) of destination ⇒ Z ⇒ bit of destination
BCLR	Dn, <ea> # <data>, <ea>	8, 32 8, 32	(bit number) of destination 0 ⇒ bit of destination
BGND	none	none	If background mode enabled, then enter background mode, else format/vector ⇒ – (SSP); PC ⇒ – (SSP); SR ⇒ – (SSP); (vector) ⇒ PC
BKPT	# <data>	none	If breakpoint cycle acknowledged, then execute returned operation word, else trap as illegal instruction
BRA	label	8, 16, 32	PC + d ⇒ PC
BSET	Dn, <ea> # <data>, <ea>	8, 32 8, 32	(bit number) of destination ⇒ Z; 1 ⇒ bit of destination
BSR	label	8, 16, 32	SP – 4 ⇒ SP; PC ⇒ (SP); PC + d ⇒ PC
BTST	Dn, <ea> # <data>, <ea>	8, 32 8, 32	(bit number) of destination ⇒ Z
CHK	<ea>, Dn	16, 32	If Dn < 0 or Dn > (ea), then CHK exception
CHK2	<ea>, Rn	8, 16, 32	If Rn < lower bound or Rn > upper bound, then CHK exception
CLR	i	8, 16, 32	0 ⇒ Destination
CMP	<ea>, Dn	8, 16, 32	(Destination – Source), CCR shows results
CPMA	<ea>, An	16, 32	(Destination – Source), CCR shows results
CMPI	# <data>, <ea>	8, 16, 32	(Destination – Data), CCR shows results
CMPM	(An) +, (An) +	8, 16, 32	(Destination – Source), CCR shows results
CMP2	<ea>, Rn	8, 16, 32	Lower bound ≤ Rn ≤ Upper bound, CCR shows result

Table 20 Instruction Set Summary(Continued)

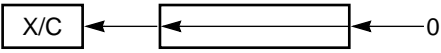
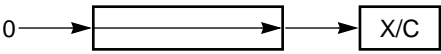
Instruction	Syntax	Operand Size	Operation
DBcc	Dn, label	16	If condition false, then Dn - 1 ⇒ PC; if Dn ≠ (-1), then PC + d ⇒ PC
DIVS/DIVU	<ea>, Dn	32/16 ⇒ 16 : 16	Destination / Source ⇒ Destination (signed or unsigned)
DIVSL/DIVUL	<ea>, Dr : Dq <ea>, Dq <ea>, Dr : Dq	64/32 ⇒ 32 : 32 32/32 ⇒ 32 32/32 ⇒ 32 : 32	Destination / Source ⇒ Destination (signed or unsigned)
EOR	Dn, <ea>	8, 16, 32	Source ⊕ Destination ⇒ Destination
EORI	# <data>, <ea>	8, 16, 32	Data ⊕ Destination ⇒ Destination
EORI to CCR	# <data>, CCR	8	Source ⊕ CCR ⇒ CCR
EORI to SR <sup>1</sup>	# <data>, SR	16	Source ⊕ SR ⇒ SR
EXG	Rn, Rn	32	Rn ⇒ Rn
EXT	Dn Dn	8 ⇒ 16 16 ⇒ 32	Sign extended Destination ⇒ Destination
EXTB	Dn	8 ⇒ 32	Sign extended Destination ⇒ Destination
ILLEGAL	none	none	SSP - 2 ⇒ SSP; vector offset ⇒ (SSP); SSP - 4 ⇒ SSP; PC ⇒ (SSP); SSP - 2 ⇒ SSP; SR ⇒ (SSP); Illegal instruction vector address ⇒ PC
JMP	í	none	Destination ⇒ PC
JSR	í	none	SP - 4 ⇒ SP; PC ⇒ (SP); destination ⇒ PC
LEA	<ea>, An	32	<ea> ⇒ An
LINK	An, # d	16, 32	SP - 4 ⇒ SP, An ⇒ (SP); SP ⇒ An, SP + d ⇒ SP
LPSTOP <sup>1</sup>	# <data>	16	Data ⇒ SR; interrupt mask ⇒ EBI; STOP
LSL	Dn, Dn # <data>, Dn í	8, 16, 32 8, 16, 32 16	
LSR	Dn, Dn # <data>, Dn í	8, 16, 32 8, 16, 32 16	
MOVE	<ea>, <ea>	8, 16, 32	Source ⇒ Destination
MOVEA	<ea>, An	16, 32 ⇒ 32	Source ⇒ Destination
MOVEA <sup>1</sup>	USP, An An, USP	32 32	USP ⇒ An An ⇒ USP
MOVE from CCR	CCR, <ea>	16	CCR ⇒ Destination
MOVE to CCR	<ea>, CCR	16	Source ⇒ CCR
MOVE from SR <sup>1</sup>	SR, <ea>	16	SR ⇒ Destination
MOVE to SR <sup>1</sup>	<ea>, SR	16	Source ⇒ SR
MOVE USP <sup>1</sup>	USP, An An, USP	32 32	USP ⇒ An An ⇒ USP
MOVEC <sup>1</sup>	Rc, Rn Rn, Rc	32 32	Rc ⇒ Rn Rn ⇒ Rc
MOVEM	list, <ea> <ea>, list	16, 32 16, 32 ⇒ 32	Listed registers ⇒ Destination Source ⇒ Listed registers
MOVEP	Dn, (d16, An)  (d16, An), Dn	16, 32	Dn [31 : 24] ⇒ (An + d); Dn [23 : 16] ⇒ (An + d + 2); Dn [15 : 8] ⇒ (An + d + 4); Dn [7 : 0] ⇒ (An + d + 6)  (An + d) ⇒ Dn [31 : 24]; (An + d + 2) ⇒ Dn [23 : 16]; (An + d + 4) ⇒ Dn [15 : 8]; (An + d + 6) ⇒ Dn [7 : 0]
MOVEQ	# <data>, Dn	8 ⇒ 32	Immediate data ⇒ Destination

Table 20 Instruction Set Summary(Continued)

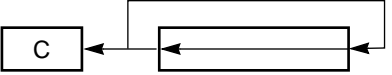

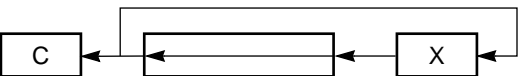
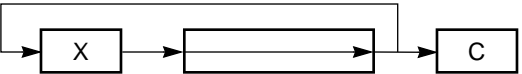
Instruction	Syntax	Operand Size	Operation
MOVES <sup>1</sup>	Rn, <ea> <ea>, Rn	8, 16, 32	Rn $\Rightarrow$ Destination using DFC Source using SFC $\Rightarrow$ Rn
MULS/MULU	<ea>, Dn <ea>, DI <ea>, Dh : DI	16 * 16 $\Rightarrow$ 32 32 * 32 $\Rightarrow$ 32 32 * 32 $\Rightarrow$ 64	Source * Destination $\Rightarrow$ Destination (signed or unsigned)
NBCD	$\dot{I}$	8 8	0 – Destination <sub>10</sub> – X $\Rightarrow$ Destination
NEG	$\dot{I}$	8, 16, 32	0 – Destination $\Rightarrow$ Destination
NEGX	$\dot{I}$	8, 16, 32	0 – Destination – X $\Rightarrow$ Destination
NOP	none	none	PC + 2 $\Rightarrow$ PC
NOT	$\dot{I}$	8, 16, 32	Destination $\Rightarrow$ Destination
OR	<ea>, Dn Dn, <ea>	8, 16, 32 8, 16, 32	Source + Destination $\Rightarrow$ Destination
ORI	#<data>, <ea>	8, 16, 32	Data + Destination $\Rightarrow$ Destination
ORI to CCR	#<data>, CCR	16	Source + CCR $\Rightarrow$ SR
ORI to SR <sup>1</sup>	#<data>, SR	16	Source ; SR $\Rightarrow$ SR
PEA	$\dot{I}$	32	SP – 4 $\Rightarrow$ SP; <ea> $\Rightarrow$ SP
RESET <sup>1</sup>	none	none	Assert RESET line
ROL	Dn, Dn #<data>, Dn $\dot{I}$	8, 16, 32 8, 16, 32 16	
ROR	Dn, Dn #<data>, Dn $\dot{I}$	8, 16, 32 8, 16, 32 16	
ROXL	Dn, Dn #<data>, Dn $\dot{I}$	8, 16, 32 8, 16, 32 16	
ROXR	Dn, Dn #<data>, Dn $\dot{I}$	8, 16, 32 8, 16, 32 16	
RTD	#d	16	(SP) $\Rightarrow$ PC; SP + 4 + d $\Rightarrow$ SP
RTE <sup>1</sup>	none	none	(SP) $\Rightarrow$ SR; SP + 2 $\Rightarrow$ SP; (SP) $\Rightarrow$ PC; SP + 4 $\Rightarrow$ SP; Restore stack according to format
RTR	none	none	(SP) $\Rightarrow$ CCR; SP + 2 $\Rightarrow$ SP; (SP) $\Rightarrow$ PC; SP + 4 $\Rightarrow$ SP
RTS	none	none	(SP) $\Rightarrow$ PC; SP + 4 $\Rightarrow$ SP
SBCD	Dn, Dn – (An), – (An)	8 8	Destination <sub>10</sub> – Source <sub>10</sub> – X $\Rightarrow$ Destination
Scc	$\dot{I}$	8	If condition true, then destination bits are set to 1; else, destination bits are cleared to 0
STOP <sup>1</sup>	#<data>	16	Data $\Rightarrow$ SR; STOP
SUB	<ea>, Dn Dn, <ea>	8, 16, 32	Destination – Source $\Rightarrow$ Destination
SUBA	<ea>, An	16, 32	Destination – Source $\Rightarrow$ Destination
SUBI	#<data>, <ea>	8, 16, 32	Destination – Data $\Rightarrow$ Destination
SUBQ	#<data>, <ea>	8, 16, 32	Destination – Data $\Rightarrow$ Destination
SUBX	Dn, Dn – (An), – (An)	8, 16, 32 8, 16, 32	Destination – Source – X $\Rightarrow$ Destination

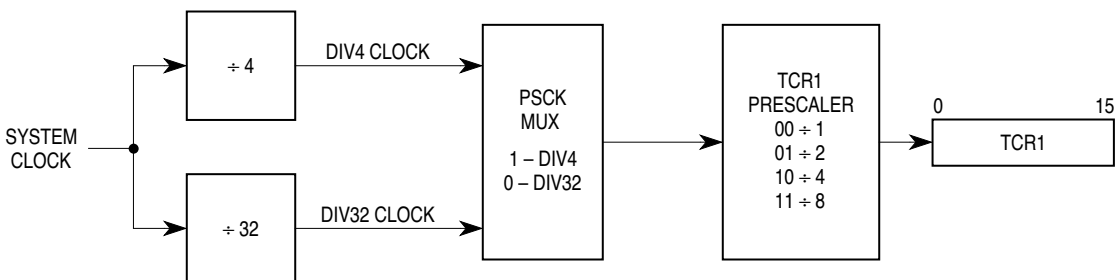
Table 20 Instruction Set Summary(Continued)

Instruction	Syntax	Operand Size	Operation
SWAP	Dn	16	
TAS	í	8	Destination Tested Condition Codes bit 7 of Destination
TBLS/TBLU	<ea>, Dn Dym : Dyn, Dn	8, 16, 32	$\text{Dyn} - \text{Dym} \Rightarrow \text{Temp}$ $(\text{Temp} * \text{Dn} [7 : 0]) \Rightarrow \text{Temp}$ $(\text{Dym} * 256) + \text{Temp} \Rightarrow \text{Dn}$
TBLSN/TBLUN	<ea>, Dn Dym : Dyn, Dn	8, 16, 32	$\text{Dyn} - \text{Dym} \Rightarrow \text{Temp}$ $(\text{Temp} * \text{Dn} [7 : 0]) / 256 \Rightarrow \text{Temp}$ $\text{Dym} + \text{Temp} \Rightarrow \text{Dn}$
TRAP	#<data>	none	$\text{SSP} - 2 \Rightarrow \text{SSP}$ ; format/vector offset $\Rightarrow (\text{SSP})$ ; $\text{SSP} - 4 \Rightarrow \text{SSP}$ ; PC $\Rightarrow (\text{SSP})$ ; SR $\Rightarrow (\text{SSP})$ ; vector address $\Rightarrow \text{PC}$
TRAPcc	none #<data>	none 16, 32	If cc true, then TRAP exception
TRAPV	none	none	If V set, then overflow TRAP exception
TST	í	8, 16, 32	Source – 0, to set condition codes
UNLK	An	32	$\text{An} \Rightarrow \text{SP}$ ; $(\text{SP}) \Rightarrow \text{An}$ , $\text{SP} + 4 \Rightarrow \text{SP}$

1. Privileged instruction.

## TCR1P — Timer Count Register 1 Prescaler Control

TCR1 is clocked from the output of a prescaler. The prescaler's input is the internal TPU system clock divided by either 4 or 32, depending on the value of the PSCK bit. The prescaler divides this input by 1, 2, 4, or 8. Channels using TCR1 have the capability to resolve down to the TPU system clock divided by 4.

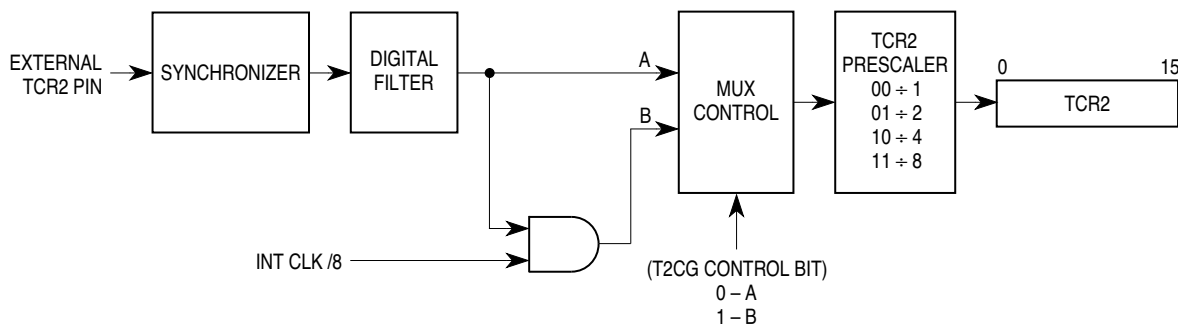


PRESCALER CTL BLOCK 1

TCR1 Prescaler	Divide By	PSCK = 0		PSCK = 1	
		Number of Clocks	Rate at 16 MHz	Number of Clocks	Rate at 16 MHz
00	1	32	2 ms	4	250 ns
01	2	64	4 ms	8	500 ns
10	4	128	8 ms	16	1 ms
11	8	256	16 ms	32	2 ms

## TCR2P — Timer Count Register 2 Prescaler Control

TCR2 is clocked from the output of a prescaler. If T2CG = 0, the input to the TCR2 prescaler is the external TCR2 clock source. If T2CG = 1, the input is the TPU system clock divided by eight. The TCR2P field specifies the value of the prescaler: 1, 2, 4, or 8. Channels using TCR2 have the capability to resolve down to the TPU system clock divided by 8. The following table is a summary of prescaler output.



PRESCALER CTL BLOCK 2

TCR2 Prescaler	Divide By	Internal Clock Divided By	External Clock Divided By
00	1	8	1
01	2	16	2
10	4	32	4
11	8	64	8

## HSQR0 — Host Sequence Register 0

\$YFFE14

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 15	CH 14	CH 13	CH 12	CH 11	CH 10	CH 9	CH 8	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

RESET:

## HSQR1 — Host Sequence Register 1

\$YFFE16

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

RESET:

## CH[15:0] — Encoded Host Sequence

The host sequence field selects the mode of operation for the time function selected on a given channel. The meaning of the host sequence bits depends on the time function specified.

## HSRR0 — Host Service Request Register 0

\$YFFE18

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 15	CH 14	CH 13	CH 12	CH 11	CH 10	CH 9	CH 8	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

RESET:

## HSRR1 — Host Service Request Register 1

\$YFFE1A

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

RESET:

## CH[15:0] — Encoded Type of Host Service

The host service request field selects the type of host service request for the time function selected on a given channel. The meaning of the host service request bits depends on the time function specified. A host service request field cleared to %00 signals the host that service is completed by the microengine on that channel. The host can request service on a channel by writing the corresponding host service request field to one of three nonzero states. The CPU should monitor the host service request register until the TPU clears the service request to %00 before the CPU changes any parameters or issues a new service request to the channel.

## CPR0 — Channel Priority Register 0

\$YFFE1C

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 15	CH 14	CH 13	CH 12	CH 11	CH 10	CH 9	CH 8	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

RESET:

## CPR1 — Channel Priority Register 1

\$YFFE1E

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

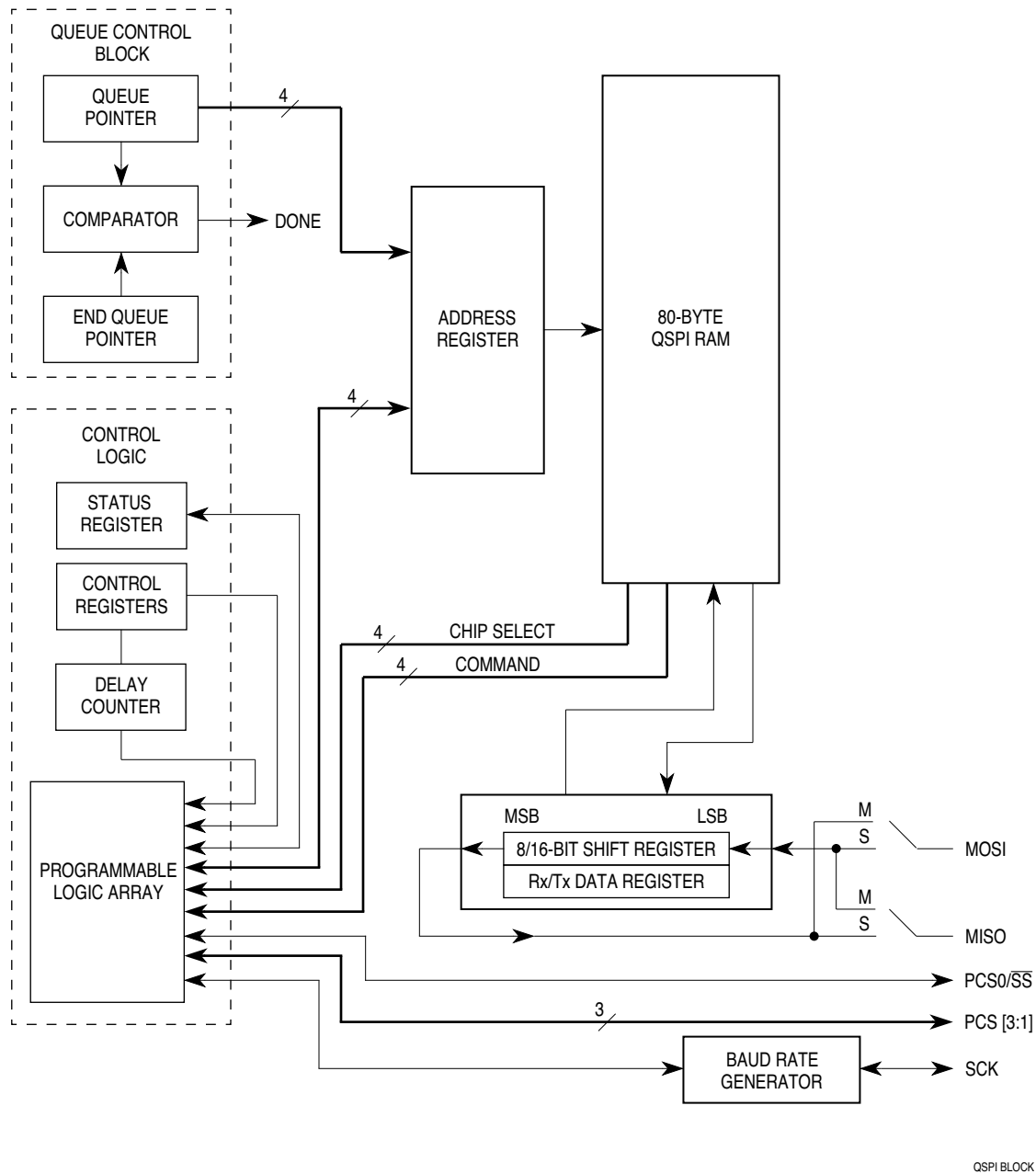
RESET:

## CH[15:0] — Encoded One of Three Channel Priority Levels



## 6.5 QSPI Submodule

The QSPI submodule communicates with external devices through a synchronous serial bus. The QSPI is fully compatible with the serial peripheral interface (SPI) systems found on other Motorola products. A block diagram of the QSPI is shown below.



**Figure 14 QSPI Block Diagram**

### 6.5.1 QSPI Pins

Seven pins are associated with the QSPI. When not needed for a QSPI application, they can be configured as general-purpose I/O pins. The **PCS0/SS** pin can function as a peripheral chip select output, slave select input, or general-purpose I/O. Refer to the following table for QSPI input and output pins and their functions.

Pin Names	Mnemonics	Mode	Function
Master In Slave Out	MISO	Master Slave	Serial Data Input to QSPI Serial Data Output from QSPI
Master Out Slave In	MOSI	Master Slave	Serial Data Output from QSPI Serial Data Input to QSPI
Serial Clock	SCK	Master Slave	Clock Output from QSPI Clock Input to QSPI
Peripheral Chip Selects	PCS[3:1]	Master	Select Peripherals
Peripheral Chip Select Slave Select	PCS0 SS	Master Master Slave	Selects Peripheral Causes Mode Fault Initiates Serial Transfer

## 6.5.2 QSPI Registers

The programmer's model for the QSPI submodule consists of the QSM global and pin control registers, four QSPI control registers, one status register, and the 80-byte QSPI RAM.

The CPU can read and write to registers and RAM. The four control registers must be initialized before the QSPI is enabled to ensure defined operation. SPCR1 should be written last because it contains QSPI enable bit SPE. Asserting this bit starts the QSPI. The QSPI control registers are reset to a defined state and can then be changed by the CPU. Reset values are shown below each register.

Refer to the following memory map of the QSPI.

Address	Name	Usage
\$YFFC18	SPCR0	QSPI Control Register 0
\$YFFC1A	SPCR1	QSPI Control Register 1
\$YFFC1C	SPCR2	QSPI Control Register 2
\$YFFC1E	SPCR3	QSPI Control Register 3
\$YFFC1F	SPSR	QSPI Status Register
\$YFFD00	RAM	QSPI Receive Data (16 Words)
\$YFFD20	RAM	QSPI Transmit Data (16 Words)
\$YFFD40	RAM	QSPI Command Control (8 Words)

Writing a different value into any control register except SPCR2 while the QSPI is enabled disrupts operation. SPCR2 is buffered to prevent disruption of the current serial transfer. After completion of the current serial transfer, the new SPCR2 values become effective.

Writing the same value into any control register except SPCR2 while the QSPI is enabled has no effect on QSPI operation. Rewriting NEWQP in SPCR2 causes execution to restart at the designated location.

### SPCR0 — QSPI Control Register 0

**\$YFFC18**

15	14	13		10	9	8	7								0
MSTR	WOMQ			BITS		CPOL	CPHA								SPBR

RESET:

0 0 0 0 0 0 0 1 0 0 0 0 0 0 1 0 0

SPCR0 contains parameters for configuring the QSPI before it is enabled. The CPU can read and write this register. The QSM has read-only access.

## MSTR — Master/Slave Mode Select

0 = QSPI is a slave device and only responds to externally generated serial data.

1 = QSPI is system master and can initiate transmission to external SPI devices.

MSTR configures the QSPI for either master or slave mode operation. This bit is cleared on reset and may only be written by the CPU.

## WOMQ — Wired-OR Mode for QSPI Pins

0 = Outputs have normal MOS drivers.

1 = Pins designated for output by DDRQS have open-drain drivers.

WOMQ allows the wired-OR function to be used on QSPI pins, regardless of whether they are used as general-purpose outputs or as QSPI outputs. WOMQ affects the QSPI pins regardless of whether the QSPI is enabled or disabled.

## BITS — Bits Per Transfer

In master mode, when BITSE in a command is set, the BITS field determines the number of data bits transferred. When BITSE is cleared, eight bits are transferred. Reserved values default to eight bits. BITSE is not used in slave mode.

The following table shows the number of bits per transfer.

BITS	Bits per Transfer
0000	16
0001	Reserved
0010	Reserved
0011	Reserved
0100	Reserved
0101	Reserved
0110	Reserved
0111	Reserved
1000	8
1001	9
1010	10
1011	11
1100	12
1101	13
1110	14
1111	15

## CPOL — Clock Polarity

0 = The inactive state value of SCK is logic level zero.

1 = The inactive state value of SCK is logic level one.

CPOL is used to determine the inactive state value of the serial clock (SCK). It is used with CPHA to produce a desired clock/data relationship between master and slave devices.

## CPHA — Clock Phase

0 = Data is captured on the leading edge of SCK and changed on the following edge of SCK.

1 = Data is changed on the leading edge of SCK and captured on the following edge of SCK.

CPHA determines which edge of SCK causes data to change and which edge causes data to be captured. CPHA is used with CPOL to produce a desired clock/data relationship between master and slave devices. CPHA is set at reset.

## SPBR — Serial Clock Baud Rate

The QSPI uses a modulus counter to derive SCK baud rate from the MCU system clock. Baud rate is selected by writing a value from 2 to 255 into the SPBR field. The following equation determines the

Command RAM is used by the QSPI when in master mode. The CPU writes one byte of control information to this segment for each QSPI command to be executed. The QSPI cannot modify information in command RAM.

Command RAM consists of 16 bytes. Each byte is divided into two fields. The peripheral chip-select field enables peripherals for transfer. The command control field provides transfer options.

A maximum of 16 commands can be in the queue. Queue execution by the QSPI proceeds from the address in NEWQP through the address in ENDQP. (Both of these fields are in SPCR2.)

CONT — Continue

- 0 = Control of chip selects returned to PORTQS after transfer is complete.
- 1 = Peripheral chip selects remain asserted after transfer is complete.

BITSE — Bits per Transfer Enable

- 0 = 8 bits
- 1 = Number of bits set in BITS field of SPCR0

DT — Delay after Transfer

The QSPI provides a variable delay at the end of serial transfer to facilitate the interface with peripherals that have a latency requirement. The delay between transfers is determined by the SPCR1 DTL field.

DSCK — PCS to SCK Delay

- 0 = PCS valid to SCK transition is one-half SCK.
- 1 = SPCR1 DSCKL field specifies delay from PCS valid to SCK.

PCS[3:0] — Peripheral Chip Select

Use peripheral chip-select bits to select an external device for serial data transfer. More than one peripheral chip select can be activated at a time, and more than one peripheral chip can be connected to each PCS pin, provided that proper fanout is observed.

$\overline{SS}$  — Slave Mode Select

Initiates slave mode serial transfer. If  $\overline{SS}$  is taken low when the QSPI is in master mode, a mode fault will be generated.

## 6.5.4 Operating Modes

The QSPI operates in either master or slave mode. Master mode is used when the MCU originates data transfers. Slave mode is used when an external device initiates serial transfers to the MCU through the QSPI. Switching between the modes is controlled by MSTR in SPCR0. Before entering either mode, appropriate QSM and QSPI registers must be properly initialized.

In master mode, the QSPI executes a queue of commands defined by control bits in each command RAM queue entry. Chip-select pins are activated, data is transmitted from transmit RAM and received into receive RAM.

In slave mode, operation proceeds in response to  $\overline{SS}$  pin activation by an external bus master. Operation is similar to master mode, but no peripheral chip selects are generated, and the number of bits transferred is controlled in a different manner. When the QSPI is selected, it automatically executes the next queue transfer to exchange data with the external device correctly.

Although the QSPI inherently supports multimaster operation, no special arbitration mechanism is provided. A mode fault flag (MODF) indicates a request for SPI master arbitration. System software must provide arbitration. Note that unlike previous SPI systems, MSTR is not cleared by a mode fault being set, nor are the QSPI pin output drivers disabled. The QSPI and associated output drivers must be disabled by clearing SPE in SPCR1.

## SBK — Send Break

0 = Normal operation

1 = Break frame(s) transmitted after completion of current frame

SBK provides the ability to transmit a break code from the SCI. If the SCI is transmitting when SBK is set, it will transmit continuous frames of zeros after it completes the current frame, until SBK is cleared. If SBK is toggled (one to zero in less than one frame interval), the transmitter sends only one or two break frames before reverting to idle line or beginning to send data.

## SCSR — SCI Status Register

**\$YFFC0C**

15	9	8	7	6	5	4	3	2	1	0
NOT USED									TDRE	TC
									RDRF	RAF
									IDLE	OR
									NF	FE
									PF	

RESET:

1 1 0 0 0 0 0 0 0 0

SCSR contains flags that show SCI operational conditions. These flags can be cleared either by hardware or by a special acknowledgment sequence. The sequence consists of SCSR read with flags set, followed by SCDR read (write in the case of TDRE and TC). A long-word read can consecutively access both SCSR and SCDR. This action clears receive status flag bits that were set at the time of the read, but does not clear TDRE or TC flags.

If an internal SCI signal for setting a status bit comes after the CPU has read the asserted status bits, but before the CPU has written or read register SCDR, the newly set status bit is not cleared. SCSR must be read again with the bit set. Also, SCDR must be written or read before the status bit is cleared.

Reading either byte of SCSR causes all 16 bits to be accessed. Any status bit already set in either byte will be cleared on a subsequent read or write of register SCDR.

## TDRE — Transmit Data Register Empty Flag

0 = Register TDR still contains data to be sent to the transmit serial shifter.

1 = A new character can now be written to register TDR.

TDRE is set when the byte in register TDR is transferred to the transmit serial shifter. If TDRE is zero, transfer has not occurred and a write to TDR will overwrite the previous value. New data is not transmitted if TDR is written without first clearing TDRE.

## TC — Transmit Complete Flag

0 = SCI transmitter is busy

1 = SCI transmitter is idle

TC is set when the transmitter finishes shifting out all data, queued preambles (mark/idle line), or queued breaks (logic zero). The interrupt can be cleared by reading SCSR when TC is set and then by writing the transmit data register (TDR) of SCDR.

## RDRF — Receive Data Register Full Flag

0 = Register RDR is empty or contains previously read data.

1 = Register RDR contains new data.

RDRF is set when the content of the receive serial shifter is transferred to the RDR. If one or more errors are detected in the received word, flag(s) NF, FE, and/or PF are set within the same clock cycle.

## RAF — Receiver Active Flag

0 = SCI receiver is idle

1 = SCI receiver is busy

RAF indicates whether the SCI receiver is busy. It is set when the receiver detects a possible start bit and is cleared when the chosen type of idle line is detected. RAF can be used to reduce collisions in systems with multiple masters.