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"Embedded - Microcontrollers" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

#### Applications of "<u>Embedded -</u> <u>Microcontrollers</u>"

#### Details

Product Status	Not For New Designs
Core Processor	CPU32
Core Size	32-Bit Single-Core
Speed	25MHz
Connectivity	EBI/EMI, SCI, SPI, UART/USART
Peripherals	POR, PWM, WDT
Number of I/O	15
Program Memory Size	-
Program Memory Type	ROMIess
EEPROM Size	-
RAM Size	2K x 8
Voltage - Supply (Vcc/Vdd)	4.5V ~ 5.5V
Data Converters	-
Oscillator Type	Internal
Operating Temperature	-40°C ~ 105°C (TA)
Mounting Type	Surface Mount
Package / Case	144-LQFP
Supplier Device Package	144-LQFP (20x20)
Purchase URL	https://www.e-xfl.com/product-detail/nxp-semiconductors/mc68332avag25

Email: info@E-XFL.COM

Address: Room A, 16/F, Full Win Commercial Centre, 573 Nathan Road, Mongkok, Hong Kong



Package Type	TPU Type	Temperature	Frequency	Package	Order Number
			(11112)	Quantity	
144-Pin QFP	Motion Control	-40 to +85 °C	16 MHz	2 pc tray	SPAKMC332GCFV16
				44 pc tray	MC68332GCFVV16
			20 MHz	2 pc tray	SPAKMC332GCFV20
				44 pc tray	MC68332GCFV20
		-40 to +105 °C	16 MHz	2 pc tray	SPAKMC332GVFV16
				44 pc tray	MC68332GVFV16
			20 MHz	2 pc tray	SPAKMC332GVFV20
				44 pc tray	MC68332GVFV20
		–40 to +125 °C	16 MHz	2 pc tray	SPAKMC332GMFV16
				44 pc tray	MC68332GMFV16
			20 MHz	2 pc tray	SPAKMC332GMFV20
				44 pc tray	MC68332GMFVV20
	Standard	–40 to +85 °C	16 MHz	2 pc tray	SPAKMC332CFV16
				44 pc tray	MC68332CFV16
			20 MHz	2 pc tray	SPAKMC332CFVV20
				44 pc tray	MC68332CFV20
		–40 to +105 °C	16 MHz	2 pc tray	SPAKMC332VFV16
				44 pc tray	MC68332VFV16
			20 MHz	2 pc tray	SPAKMC332VFV20
				44 pc tray	MC68332VFV20
		-40 to +125 °C	16 MHz	2 pc tray	SPAKMC332MFV16
				44 pc tray	MC68332MFV16
			20 MHz	2 pc tray	SPAKMC332MFV20
				44 pc tray	MC68332MFV20
	Std w/enhanced	–40 to +85 °C	16 MHz	2 pc tray	SPAKMC332ACFV16
	PPWA			44 pc tray	MC68332ACFV16
			20 MHz	2 pc tray	SPAKMC332ACFV20
				44 pc tray	MC68332ACFV20
		–40 to +105 °C	16 MHz	2 pc tray	SPAKMC332AVFV16
				44 pc tray	MC68332AVFV16
			20 MHz	2 pc tray	SPAKMC332AVFC20
				44 pc tray	MC68332AVFV20
		–40 to +125 °C	16 MHz	2 pc tray	SPAKMC332AMFV16
				44 pc tray	MC68332AMFV16
			20 MHz	2 pc tray	SPAKMC332AMFV20
				44 pc tray	MC68332AMFV20

### Table 1 Ordering Information (Continued)



1.2 Block Diagram



Figure 1 MCU Block Diagram



#### 3.2.3 Bus Monitor

The internal bus monitor checks for excessively long DSACK response times during normal bus cycles and for excessively long DSACK or AVEC response times during interrupt acknowledge cycles. The monitor asserts BERR if response time is excessive.

DSACK and AVEC response times are measured in clock cycles. The maximum allowable response time can be selected by setting the BMT field.

The monitor does not check DSACK response on the external bus unless the CPU initiates the bus cycle. The BME bit in the SYPCR enables the internal bus monitor for internal to external bus cycles. If a system contains external bus masters, an external bus monitor must be implemented and the internal to external bus monitor option must be disabled.

#### 3.2.4 Halt Monitor

The halt monitor responds to an assertion of  $\overline{HALT}$  on the internal bus. A flag in the reset status register (RSR) indicates that the last reset was caused by the halt monitor. The halt monitor reset can be inhibited by the HME bit in the SYPCR.

#### 3.2.5 Spurious Interrupt Monitor

The spurious interrupt monitor issues **BERR** if no interrupt arbitration occurs during an interrupt-acknowledge cycle.

#### 3.2.6 Software Watchdog

The software watchdog is controlled by SWE in the SYPCR. Once enabled, the watchdog requires that a service sequence be written to SWSR on a periodic basis. If servicing does not take place, the watchdog times out and issues a reset. This register can be written at any time, but returns zeros when read.

SWSR — Software Service Register								\$YF	FFA27
15	8	7	6	5	4	3	2	1	0
NOT USED		0	0	0	0	0	0	0	0
RESET:									
		0	0	0	0	0	0	0	0

Register shown with read value

Perform a software watchdog service sequence as follows:

- a. Write \$55 to SWSR.
- b. Write \$AA to SWSR.

Both writes must occur before time-out in the order listed, but any number of instructions can be executed between the two writes.

The watchdog clock rate is affected by SWP and SWT in SYPCR. When SWT[1:0] are modified, a watchdog service sequence must be performed before the new time-out period takes effect.

The reset value of SWP is affected by the state of the MODCLK pin on the rising edge of reset, as shown in the following table.

MODCLK	SWP
0	1
1	0



#### 3.3 System Clock

The system clock in the SIM provides timing signals for the IMB modules and for an external peripheral bus. Because MCU operation is fully static, register and memory contents are not affected when the clock rate changes. System hardware and software support changes in the clock rate during operation.

The system clock signal can be generated in three ways. An internal phase-locked loop can synthesize the clock from an internal or external frequency source, or the clock signal can be input from an external source.

Following is a block diagram of the clock submodule.



SYS CLOCK BLOCK 32KHZ

#### Figure 7 System Clock Block Diagram

#### 3.3.1 Clock Sources

The state of the clock mode (MODCLK) pin during reset determines the clock source. When MODCLK is held high during reset, the clock synthesizer generates a clock signal from either a crystal oscillator or an external reference input. Clock synthesizer control register SYNCR determines operating frequency and various modes of operation. When MODCLK is held low during reset, the clock synthesizer is disabled, and an external system clock signal must be applied. When the synthesizer is disabled, SYN-CR control bits have no effect.

A reference crystal must be connected between the EXTAL and XTAL pins to use the internal oscillator. Use of a 32.768-kHz crystal is recommended. These crystals are inexpensive and readily available. If an external reference signal or an external system clock signal is applied through the EXTAL pin, the XTAL pin must be left floating. External reference signal frequency must be less than or equal to maximum specified reference frequency. External system clock signal frequency must be less than or equal to maximum specified system clock frequency.



#### 3.4 External Bus Interface

The external bus interface (EBI) transfers information between the internal MCU bus and external devices. The external bus has 24 address lines and 16 data lines.

The EBI provides dynamic sizing between 8-bit and 16-bit data accesses. It supports byte, word, and long-word transfers. Ports are accessed through the use of asynchronous cycles controlled by the data transfer (SIZ1 and SIZ0) and data size acknowledge pins (DSACK1 and DSACK0). Multiple bus cycles may be required for a transfer to or from an 8-bit port.

Port width is the maximum number of bits accepted or provided during a bus transfer. External devices must follow the handshake protocol described below. Control signals indicate the beginning of the cycle, the address space, the size of the transfer, and the type of cycle. The selected device controls the length of the cycle. Strobe signals, one for the address bus and another for the data bus, indicate the validity of an address and provide timing information for data. The EBI operates in an asynchronous mode for any port width.

To add flexibility and minimize the necessity for external logic, MCU chip-select logic can be synchronized with EBI transfers. Chip-select logic can also provide internally-generated bus control signals for these accesses. Refer to **3.5 Chip Selects** for more information.

#### 3.4.1 Bus Control Signals

The CPU initiates a bus cycle by driving the address, size, function code, and read/write outputs. At the beginning of the cycle, size signals SIZ0 and SIZ1 are driven along with the function code signals. The size signals indicate the number of bytes remaining to be transferred during an operand cycle. They are valid while the address strobe ( $\overline{AS}$ ) is asserted. The following table shows SIZ0 and SIZ1 encoding. The read/write (R/W) signal determines the direction of the transfer during a bus cycle. This signal changes state, when required, at the beginning of a bus cycle, and is valid while  $\overline{AS}$  is asserted. R/W only changes state when a write cycle is preceded by a read cycle or vice versa. The signal can remain low for two consecutive write cycles.

SIZ1	SIZ0	Transfer Size
0	1	Byte
1	0	Word
1	1	Three Byte
0	0	Long Word

#### Table 8 Size Signal Encoding

#### 3.4.2 Function Codes

The CPU32 automatically generates function code signals FC[2:0]. The function codes can be considered address extensions that automatically select one of eight address spaces to which an address applies. These spaces are designated as either user or supervisor, and program or data spaces. Address space 7 is designated CPU space. CPU space is used for control information not normally associated with read or write bus cycles. Function codes are valid while  $\overline{AS}$  is asserted.



CSPAR	<b>1</b> —Ch	ip Sele	ect Pin	Assig	nment	Registe	er 1							\$YF	FA46
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	CSPA	.1[4]	CSPA	1[3]	CSPA	1[2]	CSPA	.1[1]	CSPA	1[0]
RESET:								•						•	
0	0	0	0	0	0	DATA7	1	DATA [7:6]	1	DATA [7:5]	1	DATA [7:4]	1	DATA [7:3]	1

CSPAR1 contains five 2-bit fields that determine the functions of corresponding chip-select pins. CSPAR1[15:10] are not used. These bits always read zero; writes have no effect.

CSPAR0 Field	Chip Select Signal	Alternate Signal	Discrete Output
CSPA1[4]	CS10	ADDR23	ECLK
CSPA1[3]	CS9	ADDR22	PC6
CSPA1[2]	CS8	ADDR21	PC5
CSPA1[1]	CS7	ADDR20	PC4
CSPA1[0]	CS6	ADDR19	PC3

#### Table 14 CSPAR1 Pin Assignments

At reset, either the alternate function (01) or chip-select function (11) can be encoded. DATA pins are driven to logic level one by a weak interval pull-up during reset. Encoding is for chip-select function unless a data line is held low during reset. Note that bus loading can overcome the weak pull-up and hold pins low during reset. The following table shows the hierarchical selection method that determines the reset functions of pins controlled by CSPAR1.

	Data B	Bus Pins at	Reset		Chip-Select/Address Bus Pin Function					
DATA7	DATA6	DATA5	DATA4	DATA3	CS10/ ADDR23	CS9/ ADDR22	CS8/ ADDR21	CS7/ ADDR20	CS6/ ADDR19	
1	1	1	1	1	CS10	CS9	CS8	CS7	CS6	
1	1	1	1	0	CS10	CS9	CS8	CS7	ADDR19	
1	1	1	0	X	CS10	CS9	CS8	ADDR20	ADDR19	
1	1	0	X	X	CS10	CS9	ADDR21	ADDR20	ADDR19	
1	0	Х	Х	Х	CS10	ADDR22	ADDR21	ADDR20	ADDR19	
0	Х	Х	Х	X	ADDR23	ADDR22	ADDR21	ADDR20	ADDR19	

#### Table 15 Reset Pin Function of CS[10:6]

A pin programmed as a discrete output drives an external signal to the value specified in the port C pin data register (PORTC), with the following exceptions:

- 1. No discrete output function is available on pins BR, BG, or BGACK.
- 2. ADDR23 provides E-clock output rather than a discrete output signal.

When a pin is programmed for discrete output or alternate function, internal chip-select logic still functions and can be used to generate DSACK or AVEC internally on an address match.

Port size is determined when a pin is assigned as a chip select. When a pin is assigned to an 8-bit port, the chip select is asserted at all addresses within the block range. If a pin is assigned to a 16-bit port, the upper/lower byte field of the option register selects the byte with which the chip select is associated.



mask lower-priority interrupts during exception processing, and it is decoded by modules that have requested interrupt service to determine whether the current interrupt acknowledge cycle pertains to them.

Modules that have requested interrupt service decode the IP value placed on the address bus at the beginning of the interrupt acknowledge cycle, and if their requests are at the specified IP level, respond to the cycle. Arbitration between simultaneous requests of the same priority is performed by means of serial contention between module interrupt arbitration (IARB) field bit values.

Each module that can make an interrupt service request, including the SIM, has an IARB field in its configuration register. An IARB field can be assigned a value from %0001 (lowest priority) to %1111 (highest priority). A value of %0000 in an IARB field causes the CPU to process a spurious interrupt exception when an interrupt from that module is recognized.

Because the EBI manages external interrupt requests, the SIM IARB value is used for arbitration between internal and external interrupt requests. The reset value of IARB for the SIM is %1111, and the reset IARB value for all other modules is %0000. Initialization software must assign different IARB values in order to implement an arbitration scheme.

Each module must have a unique IARB value. When two or more IARB fields have the same nonzero value, the CPU interprets multiple vector numbers simultaneously, with unpredictable consequences.

Arbitration must always take place, even when a single source requests service. This point is important for two reasons: the CPU interrupt acknowledge cycle is not driven on the external bus unless the SIM wins contention, and failure to contend causes an interrupt acknowledge bus cycle to be terminated by a bus error, which causes a spurious interrupt exception to be taken.

When arbitration is complete, the dominant module must place an interrupt vector number on the data bus and terminate the bus cycle. In the case of an external interrupt request, because the interrupt acknowledge cycle is transferred to the external bus, an external device must decode the mask value and respond with a vector number, then generate bus cycle termination signals. If the device does not respond in time, a spurious interrupt exception is taken.

The periodic interrupt timer (PIT) in the SIM can generate internal interrupt requests of specific priority at predetermined intervals. By hardware convention, PIT interrupts are serviced before external interrupt service requests of the same priority. Refer to 3.2.7 Periodic Interrupt Timer for more information.

### 3.8.2 Interrupt Processing Summary

A summary of the interrupt processing sequence follows. When the sequence begins, a valid interrupt service request has been detected and is pending.

- A. The CPU finishes higher priority exception processing or reaches an instruction boundary.
- B. Processor state is stacked. The contents of the status register and program counter are saved.
- C. The interrupt acknowledge cycle begins:
  - 1. FC[2:0] are driven to %111 (CPU space) encoding.
  - 2. The address bus is driven as follows. ADDR[23:20] = %1111; ADDR[19:16] = %1111, which indicates that the cycle is an interrupt acknowledge CPU space cycle; ADDR[15:4] = %111111111111; ADDR[3:1] = the level of the interrupt request being acknowledged; and ADDR0 = %1.
  - 3. Request priority level is latched into the IP field in the status register from the address bus.
- D. Modules or external peripherals that have requested interrupt service decode the request level in ADDR[3:1]. If the request level of at least one interrupting module or device is the same as the value in ADDR[3:1], interrupt arbitration contention takes place. When there is no contention, the spurious interrupt monitor asserts BERR, and a spurious interrupt exception is processed.
- E. After arbitration, the interrupt acknowledge cycle can be completed in one of three ways:



## **4 Central Processor Unit**

Based on the powerful MC68020, the CPU32 processing module provides enhanced system performance and also uses the extensive software base for the Motorola M68000 family.

#### 4.1 Overview

The CPU32 is fully object code compatible with the M68000 Family, which excels at processing calculation-intensive algorithms and supporting high-level languages. The CPU32 supports all of the MC68010 and most of the MC68020 enhancements, such as virtual memory support, loop mode operation, instruction pipeline, and 32-bit mathematical operations. Powerful addressing modes provide compatibility with existing software programs and increase the efficiency of high-level language compilers. Special instructions, such as table lookup and interpolate and low-power stop, support the specific requirements of controller applications. Also included is the background debugging mode, an alternate operating mode that suspends normal operation and allows the CPU to accept debugging commands from the development system.

Ease of programming is an important consideration in using a microcontroller. The CPU32 instruction set is optimized for high performance. The eight 32-bit general-purpose data registers readily support 8-bit (byte), 16-bit (word), and 32-bit (long word) operations. Ease of program checking and diagnosis is further enhanced by trace and trap capabilities at the instruction level.

Use of high-level languages is increasing as controller applications become more complex and control programs become larger. High-level languages aid rapid development of software, with less error, and are readily portable. The CPU32 instruction set supports high-level languages.

#### 4.2 Programming Model

The CPU32 has sixteen 32-bit general registers, a 32-bit program counter, one 32-bit supervisor stack pointer, a 16-bit status register, two alternate function code registers, and a 32-bit vector base register.

The programming model of the CPU32 consists of a user model and supervisor model, corresponding to the user and supervisor privilege levels. Some instructions available at the supervisor level are not available at the user level, allowing the supervisor to protect system resources from uncontrolled access. Bit S in the status register determines the privilege level.

The user programming model remains unchanged from previous M68000 Family microprocessors. Application software written to run at the non-privileged user level migrates without modification to the CPU32 from any M68000 platform. The move from SR instruction, however, is privileged in the CPU32. It is not privileged in the M68000.

For More Information On This Product, Go to: www.freescale.com





**Freescale Semiconductor, Inc.** 



		1	
Instruction	Syntax	Operand Size	Operation
MOVES <sup>1</sup>	Rn, <ea> <ea>, Rn</ea></ea>	8, 16, 32	$Rn \Rightarrow$ Destination using DFC Source using SFC $\Rightarrow$ Rn
MULS/MULU	<ea>, Dn <ea>, Dl <ea>, Dh : Dl</ea></ea></ea>	$16 * 16 \Rightarrow 32$ $32 * 32 \Rightarrow 32$ $32 * 32 \Rightarrow 64$	Source $*$ Destination $\Rightarrow$ Destination (signed or unsigned)
NBCD	Í	8 8	$0 - \text{Destination}_{10} - X \Rightarrow \text{Destination}$
NEG	Í	8, 16, 32	$0 - Destination \Rightarrow Destination$
NEGX	Í	8, 16, 32	$0 - Destination - X \Rightarrow Destination$
NOP	none	none	$PC + 2 \Rightarrow PC$
NOT	Í	8, 16, 32	$\overline{\text{Destination}} \Rightarrow \text{Destination}$
OR	<ea>, Dn Dn, <ea></ea></ea>	8, 16, 32 8, 16, 32	Source + Destination $\Rightarrow$ Destination
ORI	# <data>, <ea></ea></data>	8, 16, 32	Data + Destination $\Rightarrow$ Destination
ORI to CCR	# <data>, CCR</data>	16	Source + CCR $\Rightarrow$ SR
ORI to SR <sup>1</sup>	# <data>, SR</data>	16	Source ; SR $\Rightarrow$ SR
PEA	Í	32	$SP - 4 \Rightarrow SP; \langle ea \rangle \Rightarrow SP$
RESET <sup>1</sup>	none	none	Assert RESET line
ROL	Dn, Dn # <data>, Dn Í</data>	8, 16, 32 8, 16, 32 16	
ROR	Dn, Dn # <data>, Dn Í</data>	8, 16, 32 8, 16, 32 16	
ROXL	Dn, Dn # <data>, Dn Í</data>	8, 16, 32 8, 16, 32 16	
ROXR	Dn, Dn # <data>, Dn Í</data>	8, 16, 32 8, 16, 32 16	
RTD	#d	16	$(SP) \Rightarrow PC; SP + 4 + d \Rightarrow SP$
RTE <sup>1</sup>	none	none	$(SP) \Rightarrow SR; SP + 2 \Rightarrow SP; (SP) \Rightarrow PC;$ SP + 4 $\Rightarrow$ SP; Restore stack according to format
RTR	none	none	$(SP) \Rightarrow CCR; SP + 2 \Rightarrow SP; (SP) \Rightarrow PC; SP + 4 \Rightarrow SP$
RTS	none	none	$(SP) \Rightarrow PC; SP + 4 \Rightarrow SP$
SBCD	Dn, Dn – (An), – (An)	8 8	Destination10 – Source10 – $X \Rightarrow$ Destination
Scc	Í	8	If condition true, then destination bits are set to 1; else, destination bits are cleared to 0
STOP <sup>1</sup>	# <data></data>	16	Data $\Rightarrow$ SR; STOP
SUB	<ea>, Dn Dn, <ea></ea></ea>	8, 16, 32	Destination – Source $\Rightarrow$ Destination
SUBA	<ea>, An</ea>	16, 32	Destination – Source $\Rightarrow$ Destination
SUBI	# <data>, <ea></ea></data>	8, 16, 32	Destination – Data $\Rightarrow$ Destination
SUBQ	# <data>, <ea></ea></data>	8, 16, 32	Destination – Data $\Rightarrow$ Destination
SUBX	Dn, Dn – (An), – (An)	8, 16, 32 8, 16, 32	Destination – Source – $X \Rightarrow$ Destination

#### Table 20 Instruction Set Summary(Continued)



Instruction	Syntax	Operand Size	Operation
SWAP	Dn	16	MSW LSW
TAS	Í	8	Destination Tested Condition Codes bit 7 of Destination
TBLS/TBLU	<ea>, Dn Dym : Dyn, Dn</ea>	8, 16, 32	$Dyn - Dym \Rightarrow Temp$ (Temp * Dn [7 : 0]) $\Rightarrow$ Temp (Dym * 256) + Temp $\Rightarrow$ Dn
TBLSN/TBLUN	<ea>, Dn Dym : Dyn, Dn</ea>	8, 16, 32	$Dyn - Dym \Rightarrow Temp$ (Temp * Dn [7 : 0]) / 256 $\Rightarrow$ Temp Dym + Temp $\Rightarrow$ Dn
TRAP	# <data></data>	none	$\begin{array}{l} \text{SSP} - 2 \Rightarrow \text{SSP}; \text{ format/vector offset} \Rightarrow (\text{SSP});\\ \text{SSP} - 4 \Rightarrow \text{SSP}; \text{PC} \Rightarrow (\text{SSP}); \text{SR} \Rightarrow (\text{SSP});\\ \text{vector address} \Rightarrow \text{PC} \end{array}$
TRAPcc	none # <data></data>	none 16, 32	If cc true, then TRAP exception
TRAPV	none	none	If V set, then overflow TRAP exception
TST	Í	8, 16, 32	Source – 0, to set condition codes
UNLK	An	32	$An \Rightarrow SP; (SP) \Rightarrow An, SP + 4 \Rightarrow SP$

Table 20 Instruction Set Summary(Continued)

1. Privileged instruction.



### 5.1.7 Period Measurement with Missing Transition Detect (PMM)

Period measurement with missing transition detect allows a special-purpose 23-bit period measurement. It detects the occurrence of a missing transition (caused by a missing tooth on the sensed wheel), indicated by a period measurement that is greater than a programmable ratio of the previous period measurement. Once detected, this condition can be counted and compared to a programmable number of additional transitions detected before TCR2 is reset to \$FFFF. In addition, one byte at an address specified by a channel parameter can be read and used as a flag. A nonzero value of the flag indicates that TCR2 is to be reset to \$FFFF once the next missing transition is detected.

### 5.1.8 Position-Synchronized Pulse Generator (PSP)

Any channel of the TPU can generate an output transition or pulse, which is a projection in time based on a reference period previously calculated on another channel. Both TCRs are used in this algorithm: TCR1 is internally clocked, and TCR2 is clocked by a position indicator in the user's device. An example of a TCR2 clock source is a sensor that detects special teeth on the flywheel of an automobile using PMA or PMM. The teeth are placed at known degrees of engine rotation; hence, TCR2 is a coarse representation of engine degrees, i.e., each count represents some number of degrees.

Up to 15 position-synchronized pulse generator function channels can operate with a single input reference channel executing a PMA or PMM input function. The input channel measures and stores the time period between the flywheel teeth and resets TCR2 when the engine reaches a reference position. The output channel uses the period calculated by the input channel to project output transitions at specific engine degrees. Because the flywheel teeth might be 30 or more degrees apart, a fractional multiplication operation resolves down to the desired degrees. Two modes of operation allow pulse length to be determined either by angular position or by time.

### 5.1.9 Stepper Motor (SM)

The stepper motor control algorithm provides for linear acceleration and deceleration control of a stepper motor with a programmable number of step rates of up to 14. Any group of channels, up to eight, can be programmed to generate the control logic necessary to drive a stepper motor.

The time period between steps (P) is defined as:

$$P(r) = K1 - K2 * r$$

where r is the current step rate (1-14), and K1 and K2 are supplied as parameters.

After providing the desired step position in a 16-bit parameter, the CPU issues a step request. Next, the TPU steps the motor to the desired position through an acceleration/deceleration profile defined by parameters. The parameter indicating the desired position can be changed by the CPU while the TPU is stepping the motor. This algorithm changes the control state every time a new step command is received.

A 16-bit parameter initialized by the CPU for each channel defines the output state of the associated pin. The bit pattern written by the CPU defines the method of stepping, such as full stepping or half stepping. With each transition, the 16-bit parameter rotates one bit. The period of each transition is defined by the programmed step rate.

### 5.1.10 Period/Pulse-Width Accumulator (PPWA)

The period/pulse-width accumulator algorithm accumulates a 16-bit or 24-bit sum of either the period or the pulse width of an input signal over a programmable number of periods or pulses (from 1 to 255). After an accumulation period, the algorithm can generate a link to a sequential block of up to eight channels. The user specifies a starting channel of the block and number of channels within the block. Generation of links depends on the mode of operation. Any channel can be used to measure an accumulated number of periods of an input signal. A maximum of 24 bits can be used for the accumu-



EMU — Emulation Control

In emulation mode, the TPU executes microinstructions from MCU TPURAM exclusively. Access to the TPURAM module through the IMB by a host is blocked, and the TPURAM module is dedicated for use by the TPU. After reset, this bit can be written only once.

- 0 = TPU and TPURAM not in emulation mode
- 1 = TPU and TPURAM in emulation mode

### T2CG — TCR2 Clock/Gate Control

When the T2CG bit is set, the external TCR2 pin functions as a gate of the DIV8 clock (the TPU system clock divided by 8). In this case, when the external TCR2 pin is low, the DIV8 clock is blocked, preventing it from incrementing TCR2. When the external TCR2 pin is high, TCR2 is incremented at the frequency of the DIV8 clock. When T2CG is cleared, an external clock from the TCR2 pin, which has been synchronized and fed through a digital filter, increments TCR2.

- 0 = TCR2 pin used as clock source for TCR2
- 1 = TCR2 pin used as gate of DIV8 clock for TCR2
- STF Stop Flag
  - 0 = TPU operating
  - 1 = TPU stopped (STOP bit has been asserted)
- SUPV Supervisor Data Space
  - 0 = Assignable registers are unrestricted (FC2 is ignored)
  - 1 = Assignable registers are restricted (FC2 is decoded)

### PSCK — Prescaler Clock

- 0 = System clock/32 is input to TCR1 prescaler
- 1 = System clock/4 is input to TCR1 prescaler

#### IARB — Interrupt Arbitration Identification Number

The IARB field is used to arbitrate between simultaneous interrupt requests of the same priority. Each module that can generate interrupt requests must be assigned a unique, non-zero IARB field value. Refer to the **3.8 Interrupts** for more information.

TICR —	TPU	Interrupt	Configuration	Register

### \$YFFE08

15		11	10		8	7			4	3		0
	NOT USED			CIRL			CI	3V			NOT USED	
RESET:												
			0	0	0	0	0	0	0			

#### CIRL — Channel Interrupt Request Level

The interrupt request level for all channels is specified by this 3-bit encoded field. Level seven for this field indicates a nonmaskable interrupt; level zero indicates that all channel interrupts are disabled.

### CIBV — Channel Interrupt Base Vector

The TPU is assigned 16 unique interrupt vector numbers, one vector number for each channel. The CIBV field specifies the most significant nibble of all 16 TPU channel interrupt vector numbers. The lower nibble of the TPU interrupt vector number is determined by the channel number on which the interrupt occurs.



5.5.2 CI	5.5.2 Channel Control Registers														
CIER —	- Chan	nel Inte	errupt I	Enable	Regis	ter								\$YF	FE0A
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 15	CH 14	CH 13	CH 12	CH 11	CH 10	CH 9	CH 8	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
RESET:															
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
CH[15:0	CH[15:0] — Channel Interrupt Enable/Disable 0 = Channel interrupts disabled 1 = Channel interrupts enabled														
CISR — Channel Interrupt Status Register \$YFFE20													FE20		
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 15	CH 14	CH 13	CH 12	CH 11	CH 10	CH 9	CH 8	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
RESET:															
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
CH[15:0	CH[15:0] — Channel Interrupt Status Bit 0 = Channel interrupt not asserted 1 = Channel interrupt asserted														
CFSR0 — Channel Function Select Register 0 \$YFFE0C															
15			12	11			8	7			4	3			0
	CHANN	IEL15			CHAN	NEL14			CHAN	NEL13			CHAN	NEL12	
RESET:															
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
CFSR1	— Cha	annel F	unctio	n Sele	ct Regi	ster 1								\$YF	FE0E
15			12	11			8	7			4	3			0
	CHANN	IEL11			CHAN	NEL10			CHAN	INEL9			CHAN	INEL8	
RESET:															
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
CFSR2	— Cha	annel F	unctio	n Sele	ct Regi	ster 2								\$YI	FE10
15			12	11			8	7			4	3			0
	CHANN	NEL7			CHAN	INEL6			CHAN	INEL5			CHAN	INEL4	
RESET:												1			
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
CFSR3	— Cha	annel F	unctio	n Sele	ct Regi	ster 3								\$YI	FE12
15			12	11			8	7			4	3			0
CHANNEL3 CHANNEL2							CHAN	INEL1			CHAN	INEL0			
RESET:															
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

CHANNEL[15:0] — Encoded Time Function for each Channel

Encoded 4-bit fields in the channel function select registers specify one of 16 time functions to be executed on the corresponding channel.



#### 6.5 QSPI Submodule

The QSPI submodule communicates with external devices through a synchronous serial bus. The QSPI is fully compatible with the serial peripheral interface (SPI) systems found on other Motorola products. A block diagram of the QSPI is shown below.



#### Figure 14 QSPI Block Diagram

#### 6.5.1 QSPI Pins

Seven pins are associated with the QSPI. When not needed for a QSPI application, they can be configured as general-purpose I/O pins. The PCS0/SS pin can function as a peripheral chip select output, slave select input, or general-purpose I/O. Refer to the following table for QSPI input and output pins and their functions.



#### 6.6 SCI Submodule

The SCI submodule is used to communicate with external devices through an asynchronous serial bus. The SCI is fully compatible with the SCI systems found on other Motorola MCUs, such as the M68HC11 and M68HC05 Families.

#### 6.6.1 SCI Pins

There are two unidirectional pins associated with the SCI. The SCI controls the transmit data (TXD) pin when enabled, whereas the receive data (RXD) pin remains a dedicated input pin to the SCI. TXD is available as a general-purpose I/O pin when the SCI transmitter is disabled. When used for I/O, TXD can be configured either as input or output, as determined by QSM register DDRQS.

The following table shows SCI pins and their functions.

Pin Names	Mnemonics	Mode	Function		
Receive Data	RXD	Receiver Disabled Receiver Enabled	Not Used Serial Data Input to SCI		
Transmit Data	TXD	Transmitter Disabled Transmitter Enabled	General-Purpose I/O Serial Data Output from SCI		

#### 6.6.2 SCI Registers

The SCI programming model includes QSM global and pin control registers, and four SCI registers. There are two SCI control registers, one status register, and one data register. All registers can be read or written at any time by the CPU.

Changing the value of SCI control bits during a transfer operation may disrupt operation. Before changing register values, allow the transmitter to complete the current transfer, then disable the receiver and transmitter. Status flags in the SCSR may be cleared at any time.

SCCR0 — SCI Control Register 0 \$Y														\$YF	FC08
15	14	13	12												0
0	0	0							SCBR						
RESET:	·														
0	0	0	0	0	0	0	0	0	0	0	0	0	1	0	0

SCCR0 contains a baud rate selection parameter. Baud rate must be set before the SCI is enabled. The CPU can read and write this register at any time.

#### Bits [15:13] - Not Implemented

#### SCBR — Baud Rate

SCI baud rate is programmed by writing a 13-bit value to BR. The baud rate is derived from the MCU system clock by a modulus counter.

The SCI receiver operates asynchronously. An internal clock is necessary to synchronize with an incoming data stream. The SCI baud rate generator produces a receiver sampling clock with a frequency 16 times that of the expected baud rate of the incoming data. The SCI determines the position of bit boundaries from transitions within the received waveform, and adjusts sampling points to the proper positions within the bit period. Receiver sampling rate is always 16 times the frequency of the SCI baud rate, which is calculated as follows:

SCI Baud Rate = System Clock/(32SCBR)

or

SCBR = System Clock(32SCK)(Baud Rate desired)

where SCBR is in the range {1, 2, 3, ..., 8191}

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PE — Parity Enable

0 = SCI parity disabled

1 = SCI parity enabled

PE determines whether parity is enabled or disabled for both the receiver and the transmitter. If the received parity bit is not correct, the SCI sets the PF error flag in SCSR.

When PE is set, the most significant bit (MSB) of the data field is used for the parity function, which results in either seven or eight bits of user data, depending on the condition of M bit. The following table lists the available choices.

М	PE	Result				
0	0	8 Data Bits				
0	1	7 Data Bits, 1 Parity Bit				
1	0	9 Data Bits				
1	1	8 Data Bits, 1 Parity Bit				

### M — Mode Select

0 = SCI frame: 1 start bit, 8 data bits, 1 stop bit (10 bits total)

1 = SCI frame: 1 start bit, 9 data bits, 1 stop bit (11 bits total)

### WAKE — Wakeup by Address Mark

- 0 = SCI receiver awakened by idle-line detection
- 1 = SCI receiver awakened by address mark (last bit set)

### TIE — Transmit Interrupt Enable

- 0 = SCI TDRE interrupts inhibited
- 1 = SCI TDRE interrupts enabled
- TCIE Transmit Complete Interrupt Enable
  - 0 = SCI TC interrupts inhibited
  - 1 = SCI TC interrupts enabled

### RIE — Receiver Interrupt Enable

- 0 = SCI RDRF interrupt inhibited
- 1 = SCI RDRF interrupt enabled
- ILIE Idle-Line Interrupt Enable
  - 0 = SCI IDLE interrupts inhibited
  - 1 = SCI IDLE interrupts enabled

### TE — Transmitter Enable

0 = SCI transmitter disabled (TXD pin may be used as I/O)

1 = SCI transmitter enabled (TXD pin dedicated to SCI transmitter)

The transmitter retains control of the TXD pin until completion of any character transfer that was in progress when TE is cleared.

### RE — Receiver Enable

- 0 = SCI receiver disabled (status bits inhibited)
- 1 = SCI receiver enabled
- RWU Receiver Wakeup
  - 0 = Normal receiver operation (received data recognized)
  - 1 = Wakeup mode enabled (received data ignored until awakened)

Setting RWU enables the wakeup function, which allows the SCI to ignore received data until awakened by either an idle line or address mark (as determined by WAKE). When in wakeup mode, the receiver status flags are not set, and interrupts are inhibited. This bit is cleared automatically (returned to normal mode) when the receiver is awakened.



SBK — Send Break

0 = Normal operation

1 = Break frame(s) transmitted after completion of current frame

SBK provides the ability to transmit a break code from the SCI. If the SCI is transmitting when SBK is set, it will transmit continuous frames of zeros after it completes the current frame, until SBK is cleared. If SBK is toggled (one to zero in less than one frame interval), the transmitter sends only one or two break frames before reverting to idle line or beginning to send data.

SCSR — S	SCI Status	Register
----------	------------	----------

\$YFFC0C	
----------	--

15	9	8	7	6	5	4	3	2	1	0
NOT USED		TDRE	TC	RDRF	RAF	IDLE	OR	NF	FE	PF
RESET:										
		1	1	0	0	0	0	0	0	0

SCSR contains flags that show SCI operational conditions. These flags can be cleared either by hardware or by a special acknowledgment sequence. The sequence consists of SCSR read with flags set, followed by SCDR read (write in the case of TDRE and TC). A long-word read can consecutively access both SCSR and SCDR. This action clears receive status flag bits that were set at the time of the read, but does not clear TDRE or TC flags.

If an internal SCI signal for setting a status bit comes after the CPU has read the asserted status bits, but before the CPU has written or read register SCDR, the newly set status bit is not cleared. SCSR must be read again with the bit set. Also, SCDR must be written or read before the status bit is cleared.

Reading either byte of SCSR causes all 16 bits to be accessed. Any status bit already set in either byte will be cleared on a subsequent read or write of register SCDR.

### TDRE — Transmit Data Register Empty Flag

- 0 = Register TDR still contains data to be sent to the transmit serial shifter.
- 1 = A new character can now be written to register TDR.

TDRE is set when the byte in register TDR is transferred to the transmit serial shifter. If TDRE is zero, transfer has not occurred and a write to TDR will overwrite the previous value. New data is not transmitted if TDR is written without first clearing TDRE.

TC — Transmit Complete Flag

- 0 = SCI transmitter is busy
- 1 = SCI transmitter is idle

TC is set when the transmitter finishes shifting out all data, queued preambles (mark/idle line), or queued breaks (logic zero). The interrupt can be cleared by reading SCSR when TC is set and then by writing the transmit data register (TDR) of SCDR.

## RDRF — Receive Data Register Full Flag

0 = Register RDR is empty or contains previously read data.

1 = Register RDR contains new data.

RDRF is set when the content of the receive serial shifter is transferred to the RDR. If one or more errors are detected in the received word, flag(s) NF, FE, and/or PF are set within the same clock cycle.

### RAF — Receiver Active Flag

- 0 = SCI receiver is idle
- 1 = SCI receiver is busy

RAF indicates whether the SCI receiver is busy. It is set when the receiver detects a possible start bit and is cleared when the chosen type of idle line is detected. RAF can be used to reduce collisions in systems with multiple masters.



#### IDLE — Idle-Line Detected Flag

0 = SCI receiver did not detect an idle-line condition.

1 = SCI receiver detected an idle-line condition.

IDLE is disabled when RWU in SCCR1 is set. IDLE is set when the SCI receiver detects the idle-line condition specified by ILT in SCCR1. If cleared, IDLE will not set again until after RDRF is set. RDRF is set when a break is received, so that a subsequent idle line can be detected.

#### OR — Overrun Error Flag

0 = RDRF is cleared before new data arrives.

1 = RDRF is not cleared before new data arrives.

OR is set when a new byte is ready to be transferred from the receive serial shifter to the RDR, and RDRF is still set. Data transfer is inhibited until OR is cleared. Previous data in RDR remains valid, but data received during overrun condition (including the byte that set OR) is lost.

### NF — Noise Error Flag

- 0 = No noise detected on the received data
- 1 = Noise occurred on the received data

NF is set when the SCI receiver detects noise on a valid start bit, on any data bit, or on a stop bit. It is not set by noise on the idle line or on invalid start bits. Each bit is sampled three times. If none of the three samples are the same logic level, the majority value is used for the received data value, and NF is set. NF is not set until an entire frame is received and RDRF is set.

### FE — Framing Error Flag

0 = No framing error on the received data.

1 = Framing error or break occurred on the received data.

FE is set when the SCI receiver detects a zero where a stop bit was to have occurred. FE is not set until the entire frame is received and RDRF is set. A break can also cause FE to be set. It is possible to miss a framing error if RXD happens to be at logic level one at the time the stop bit is expected.

#### PF — Parity Error Flag

0 = No parity error on the received data

1 = Parity error occurred on the received data

PF is set when the SCI receiver detects a parity error. PF is not set until the entire frame is received and RDRF is set.

SCDR — SCI Data Register\$YF															FC0E
15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	0	0	0	R8/T8	R7/T7	R6/T6	R5/T5	R4/T4	R3/T3	R2/T2	R1/T1	R0/T0
RESET:															
0	0	0	0	0	0	0	U	U	U	U	U	U	U	U	U

SCDR contains two data registers at the same address. Receive data register (RDR) is a read-only register that contains data received by the SCI. The data comes into the receive serial shifter and is transferred to RDR. Transmit data register (TDR) is a write-only register that contains data to be transmitted. The data is first written to TDR, then transferred to the transmit serial shifter, where additional format bits are added before transmission. R[7:0]/T[7:0] contain either the first eight data bits received when SCDR is read, or the first eight data bits to be transmitted when SCDR is written. R8/T8 are used when the SCI is configured for 9-bit operation. When it is configured for 8-bit operation, they have no meaning or effect.



## 7 Standby RAM with TPU Emulation RAM

The TPURAM module contains a 2-Kbyte array of fast (two bus cycle) static RAM, which is especially useful for system stacks and variable storage. Alternately, it can be used by the TPU as emulation RAM for new timer algorithms.

#### 7.1 Overview

The TPURAM can be mapped to any 4-Kbyte boundary in the address map, but must not overlap the module control registers. (Overlap makes the registers inaccessible.) Data can be read or written in bytes, word, or long words. TPURAM responds to both program and data space accesses. Data can be read or written in bytes, words, or long words. The TPURAM is powered by  $V_{DD}$  in normal operation. During power-down, the TPURAM contents are maintained by power on standby voltage pin  $V_{STBY}$ . Power switching between sources is automatic.

Access to the TPURAM array is controlled by the RASP field in TRAMMCR. This field can be encoded so that TPURAM responds to both program and data space accesses. This allows code to be executed from TPURAM, and permits the use of program counter relative addressing mode for operand fetches from the array.

An address map of the TPURAM control registers follows. All TPURAM control registers are located in supervisor data space.

Access	Address	15	8	7	0					
S	\$YFFB00		TPURAM MODULE CONFIGURATION REGISTER (TRAMMCR)							
S	\$YFFB02		TPURAM TEST REGISTER (TRAMTST)							
S	\$YFFB04		TPURAM BASE ADDRESS REGISTER (TRAMBAR)							
	\$YFFB06– \$YFFB3F	0– NOT USED 3F								

#### Table 28 TPURAM Control Register Address Map

Y = M111, where M is the logic state of the MM bit in the SIMCR.

#### 7.2 TPURAM Register Block

There are three TPURAM control registers: the RAM module configuration register (TRAMMCR), the RAM test register (TRAMTST), and the RAM array base address registers (TRAMBAR).

There is an 8-byte minimum register block size for the module. Unimplemented register addresses are read as zeros, and writes have no effect.

### 7.3 TPURAM Registers

TRAMM	ICR —	TPUR	AM Mc		\$YFFB00					
15	14	13	12	11	10	9	8	7		0
STOP	0	0	0	0	0	0	RASP		NOT USED	
RESET:					•					
٥	٥	0	٥	0	0	0	1			

TSTOP —Stop Control

0 = RAM array operates normally.

1 = RAM array enters low-power stop mode.

This bit controls whether the RAM array is in stop mode or normal operation. Reset state is zero, for normal operation. In stop mode, the array retains its contents, but cannot be read or written by the CPU.