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Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	CPU32
Core Size	32-Bit Single-Core
Speed	25MHz
Connectivity	EBI/EMI, SCI, SPI, UART/USART
Peripherals	POR, PWM, WDT
Number of I/O	15
Program Memory Size	-
Program Memory Type	ROMless
EEPROM Size	-
RAM Size	2K x 8
Voltage - Supply (Vcc/Vdd)	4.5V ~ 5.5V
Data Converters	-
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	144-LQFP
Supplier Device Package	144-LQFP (20x20)
Purchase URL	https://www.e-xfl.com/pro/item?MUrl=&PartUrl=mc68332gcag25

2 Signal Descriptions

2.1 Pin Characteristics

The following table shows MCU pins and their characteristics. All inputs detect CMOS logic levels. All inputs can be put in a high-impedance state, but the method of doing this differs depending upon pin function. Refer to the table, MCU Driver Types, for a description of output drivers. An entry in the discrete I/O column of the MCU Pin Characteristics table indicates that a pin has an alternate I/O function. The port designation is given when it applies. Refer to the MCU Block Diagram for information about port organization.

Table 2 MCU Pin Characteristic

Pin Mnemonic	Output Driver	Input Synchronized	Input Hysteresis	Discrete I/O	Port Designation
ADDR23/$\overline{CS10}$/ECLK	A	Y	N	O	—
ADDR[22:19]/$\overline{CS[9:6]}$	A	Y	N	O	PC[6:3]
ADDR[18:0]	A	Y	N	—	—
$\overline{AS}$	B	Y	N	I/O	PE5
$\overline{AVEC}$	B	Y	N	I/O	PE2
$\overline{BERR}$	B	Y	N	—	—
$\overline{BG}$/$\overline{CS1}$	B	—	—	—	—
$\overline{BGACK}$/$\overline{CS2}$	B	Y	N	—	—
$\overline{BKPT}$/DSCLK	—	Y	Y	—	—
$\overline{BR}$/$\overline{CS0}$	B	Y	N	—	—
CLKOUT	A	—	—	—	—
$\overline{CSBOOT}$	B	—	—	—	—
DATA[15:0][1]	Aw	Y	N	—	—
$\overline{DS}$	B	Y	N	I/O	PE4
$\overline{DSACK1}$	B	Y	N	I/O	PE1
$\overline{DSACK0}$	B	Y	N	I/O	PE0
DSI/$\overline{IFETCH}$	A	Y	Y	—	—
DSO/$\overline{IPIPE}$	A	—	—	—	—
EXTAL[2]	—	—	Special	—	—
FC[2:0]/$\overline{CS[5:3]}$	A	Y	N	O	PC[2:0]
FREEZE/QUOT	A	—	—	—	—
$\overline{HALT}$	Bo	Y	N	—	—
$\overline{IRQ[7:1]}$	B	Y	Y	I/O	PF[7:1]
MISO	Bo	Y	Y	I/O	PQS0
MODCLK[1]	B	Y	N	I/O	PF0
MOSI	Bo	Y	Y	I/O	PQS1
PCS0/$\overline{SS}$	Bo	Y	Y	I/O	PQS3
PCS[3:1]	Bo	Y	Y	I/O	PQS[6:4]
R/$\overline{W}$	A	Y	N	—	—
$\overline{RESET}$	Bo	Y	Y	—	—
$\overline{RMC}$	B	Y	N	I/O	PE3
RXD	—	N	N	—	—
SCK	Bo	Y	Y	I/O	PQS2
SIZ[1:0]	B	Y	N	I/O	PE[7:6]

Table 5 MCU Signal Characteristics (Continued)

Signal Name	MCU Module	Signal Type	Active State
TSC	SIM	Input	—
TXD	QSM	Output	—
XFC	SIM	Input	—
XTAL	SIM	Output	—

2.5 Signal Function

Table 6 MCU Signal Function

Signal Name	Mnemonic	Function
Address Bus	ADDR[23:0]	24-bit address bus
Address Strobe	$\overline{AS}$	Indicates that a valid address is on the address bus
Autovector	$\overline{AVEC}$	Requests an automatic vector during interrupt acknowledge
Bus Error	$\overline{BERR}$	Indicates that a bus error has occurred
Bus Grant	$\overline{BG}$	Indicates that the MCU has relinquished the bus
Bus Grant Acknowledge	$\overline{BGACK}$	Indicates that an external device has assumed bus mastership
Breakpoint	$\overline{BKPT}$	Signals a hardware breakpoint to the CPU
Bus Request	$\overline{BR}$	Indicates that an external device requires bus mastership
System Clockout	CLKOUT	System clock output
Chip Selects	$\overline{CS[10:0]}$	Select external devices at programmed addresses
Boot Chip Select	$\overline{CSBOOT}$	Chip select for external boot start-up ROM
Data Bus	DATA[15:0]	16-bit data bus
Data Strobe	$\overline{DS}$	During a read cycle, indicates when it is possible for an external device to place data on the data bus. During a write cycle, indicates that valid data is on the data bus.
Data and Size Acknowledge	$\overline{DSACK[1:0]}$	Provide asynchronous data transfers and dynamic bus sizing
Development Serial In, Out, Clock	DSI, DSO, DSCLK	Serial I/O and clock for background debugging mode
Crystal Oscillator	EXTAL, XTAL	Connections for clock synthesizer circuit reference; a crystal or an external oscillator can be used
Function Codes	FC[2:0]	Identify processor state and current address space
Freeze	FREEZE	Indicates that the CPU has entered background mode
Halt	$\overline{HALT}$	Suspend external bus activity
Instruction Pipeline	$\overline{IFETCH}$ $\overline{IPIPE}$	Indicate instruction pipeline activity
Interrupt Request Level	$\overline{IRQ[7:1]}$	Provides an interrupt priority level to the CPU
Master In Slave Out	MISO	Serial input to QSPI in master mode; serial output from QSPI in slave mode
Clock Mode Select	MODCLK	Selects the source and type of system clock
Master Out Slave In	MOSI	Serial output from QSPI in master mode; serial input to QSPI in slave mode
Port C	PC[6:0]	SIM digital output port signals
Peripheral Chip Select	PCS[3:0]	QSPI peripheral chip selects
Port E	PE[7:0]	SIM digital I/O port signals
Port F	PF[7:0]	SIM digital I/O port signals
Port QS	PQS[7:0]	QSM digital I/O port signals

Table 7 SIM Address Map (Continued)

Access	Address	15 8	7 0
S	$YFFA56	CHIP-SELECT OPTION 2 (CSOR2)	
S	$YFFA58	CHIP-SELECT BASE 3 (CSBAR3)	
S	$YFFA5A	CHIP-SELECT OPTION 3 (CSOR3)	
S	$YFFA5C	CHIP-SELECT BASE 4 (CSBAR4)	
S	$YFFA5E	CHIP-SELECT OPTION 4 (CSOR4)	
S	$YFFA60	CHIP-SELECT BASE 5 (CSBAR5)	
S	$YFFA62	CHIP-SELECT OPTION 5 (CSOR5)	
S	$YFFA64	CHIP-SELECT BASE 6 (CSBAR6)	
S	$YFFA66	CHIP-SELECT OPTION 6 (CSOR6)	
S	$YFFA68	CHIP-SELECT BASE 7 (CSBAR7)	
S	$YFFA6A	CHIP-SELECT OPTION 7 (CSOR7)	
S	$YFFA6C	CHIP-SELECT BASE 8 (CSBAR8)	
S	$YFFA6E	CHIP-SELECT OPTION 8 (CSOR8)	
S	$YFFA70	CHIP-SELECT BASE 9 (CSBAR9)	
S	$YFFA72	CHIP-SELECT OPTION 9 (CSOR9)	
S	$YFFA74	CHIP-SELECT BASE 10 (CSBAR10)	
S	$YFFA76	CHIP-SELECT OPTION 10 (CSOR10)	
	$YFFA78	NOT USED	NOT USED
	$YFFA7A	NOT USED	NOT USED
	$YFFA7C	NOT USED	NOT USED
	$YFFA7E	NOT USED	NOT USED

Y = M111, where M is the logic state of the module mapping (MM) bit in the SIMCR.

3.2 System Configuration and Protection

This functional block provides configuration control for the entire MCU. It also performs interrupt arbitration, bus monitoring, and system test functions. MCU system protection includes a bus monitor, a HALT monitor, a spurious interrupt monitor, and a software watchdog timer. These functions have been made integral to the microcontroller to reduce the number of external components in a complete control system.

3.4.8 Data Transfer Mechanism

The MCU architecture supports byte, word, and long-word operands, allowing access to 8- and 16-bit data ports through the use of asynchronous cycles controlled by the data transfer and size acknowledge inputs ($\overline{\text{DSACK1}}$ and $\overline{\text{DSACK0}}$).

3.4.9 Dynamic Bus Sizing

The MCU dynamically interprets the port size of the addressed device during each bus cycle, allowing operand transfers to or from 8- and 16-bit ports. During an operand transfer cycle, the slave device signals its port size and indicates completion of the bus cycle to the MCU through the use of the $\overline{\text{DSACK0}}$ and $\overline{\text{DSACK1}}$ inputs, as shown in the following table.

Table 10 Effect of $\overline{\text{DSACK}}$ Signals

$\overline{\text{DSACK1}}$	$\overline{\text{DSACK0}}$	Result
1	1	Insert Wait States in Current Bus Cycle
1	0	Complete Cycle —Data Bus Port Size is 8 Bits
0	1	Complete Cycle —Data Bus Port Size is 16 Bits
0	0	Reserved

For example, if the MCU is executing an instruction that reads a long-word operand from a 16-bit port, the MCU latches the 16 bits of valid data and then runs another bus cycle to obtain the other 16 bits. The operation for an 8-bit port is similar, but requires four read cycles. The addressed device uses the $\overline{\text{DSACK0}}$ and $\overline{\text{DSACK1}}$ signals to indicate the port width. For instance, a 16-bit device always returns $\overline{\text{DSACK0}}$ = 1 and $\overline{\text{DSACK1}}$ = 0 for a 16-bit port, regardless of whether the bus cycle is a byte or word operation.

Dynamic bus sizing requires that the portion of the data bus used for a transfer to or from a particular port size be fixed. A 16-bit port must reside on data bus bits [15:0] and an 8-bit port must reside on data bus bits [15:8]. This minimizes the number of bus cycles needed to transfer data and ensures that the MCU transfers valid data.

The MCU always attempts to transfer the maximum amount of data on all bus cycles. For a word operation, it is assumed that the port is 16 bits wide when the bus cycle begins. Operand bytes are designated as shown in the following figure. OP0 is the most significant byte of a long-word operand, and OP3 is the least significant byte. The two bytes of a word-length operand are OP0 (most significant) and OP1. The single byte of a byte-length operand is OP0.

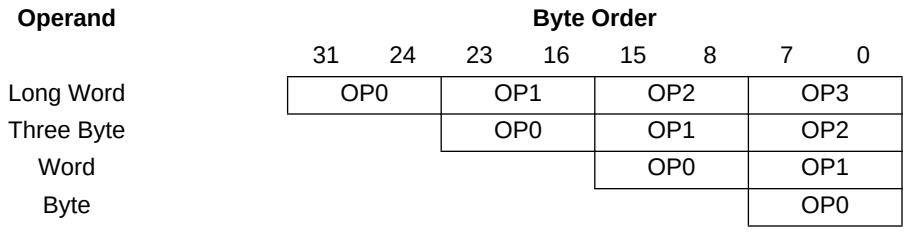


Figure 8 Operand Byte Order

3.4.10 Operand Alignment

The data multiplexer establishes the necessary connections for different combinations of address and data sizes. The multiplexer takes the two bytes of the 16-bit bus and routes them to their required positions. Positioning of bytes is determined by the size and address outputs. SIZ1 and SIZ0 indicate the remaining number of bytes to be transferred during the current bus cycle. The number of bytes transferred is equal to or less than the size indicated by SIZ1 and SIZ0, depending on port width.

Chip-select assertion can be synchronized with bus control signals to provide output enable, read/write strobes, or interrupt acknowledge signals. Logic can also generate $\overline{\text{DSACK}}$ signals internally. A single $\overline{\text{DSACK}}$ generator is shared by all circuits. Multiple chip selects assigned to the same address and control must have the same number of wait states.

Chip selects can also be synchronized with the ECLK signal available on ADDR23.

When a memory access occurs, chip-select logic compares address space type, address, type of access, transfer size, and interrupt priority (in the case of interrupt acknowledge) to parameters stored in chip-select registers. If all parameters match, the appropriate chip-select signal is asserted. Select signals are active low. Refer to the following block diagram of a single chip-select circuit.

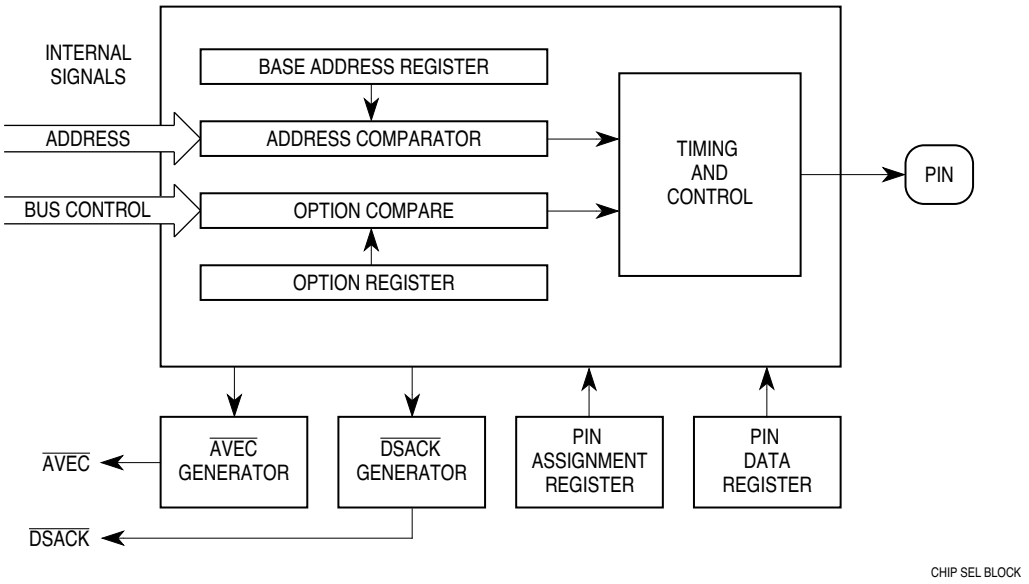


Figure 9 Chip-Select Circuit Block Diagram

The following table lists allocation of chip-selects and discrete outputs on the pins of the MCU.

Pin	Chip Select	Discrete Outputs
$\overline{\text{CSBOOT}}$	$\overline{\text{CSBOOT}}$	—
$\overline{\text{BR}}$	$\overline{\text{CS0}}$	—
$\overline{\text{BG}}$	$\overline{\text{CS1}}$	—
$\overline{\text{BGACK}}$	$\overline{\text{CS2}}$	—
FC0	$\overline{\text{CS3}}$	PC0
FC1	$\overline{\text{CS4}}$	PC1
FC2	$\overline{\text{CS5}}$	PC2
ADDR19	$\overline{\text{CS6}}$	PC3
ADDR20	$\overline{\text{CS7}}$	PC4
ADDR21	$\overline{\text{CS8}}$	PC5
ADDR22	$\overline{\text{CS9}}$	PC6
ADDR23	$\overline{\text{CS10}}$	ECLK

CSORBT —Chip-Select Option Register Boot ROM **$YFFA4A**

15	14	13	12	11	10	9	6	5	4	3	1	0
MODE	BYTE		R/$\overline{W}$		STRB	DSACK		SPACE		IPL		$\overline{AVEC}$
RESET:												
0	1	1	1	1	0	1	1 0	1	1	0	0	0 0

CSOR[10:0] —Chip-Select Option Registers **$YFFA4E–$YFFA76**

15	14	13	12	11	10	9	6	5	4	3	1	0
MODE	BYTE		R/$\overline{W}$		STRB	DSACK		SPACE		IPL		AVEC
RESET:												
0	0	0	0	0	0	0	0 0	0	0	0	0	0 0

CSORBT, the option register for $\overline{CSBOOT}$, contains special reset values that support bootstrap operations from peripheral memory devices.

The following bit descriptions apply to both CSORBT and CSOR[10:0] option registers.

MODE — Asynchronous/Synchronous Mode

0 = Asynchronous mode selected (chip-select assertion determined by internal or external bus control signals)

1 = Synchronous mode selected (chip-select assertion synchronized with ECLK signal)

In asynchronous mode, the chip select is asserted synchronized with $\overline{AS}$ or $\overline{DS}$.

The $\overline{DSACK}$ field is not used in synchronous mode because a bus cycle is only performed as a synchronous operation. When a match condition occurs on a chip select programmed for synchronous operation, the chip select signals the EBI that an ECLK cycle is pending.

BYTE — Upper/Lower Byte Option

This field is used only when the chip-select 16-bit port option is selected in the pin assignment register. The following table lists upper/lower byte options.

Byte	Description
00	Disable
01	Lower Byte
10	Upper Byte
11	Both Bytes

R/$\overline{W}$ — Read/Write

This field causes a chip select to be asserted only for a read, only for a write, or for both read and write. Refer to the following table for options available.

R/$\overline{W}$	Description
00	Reserved
01	Read Only
10	Write Only
11	Read/Write

STRB — Address Strobe/Data Strobe

0 = Address strobe

1 = Data strobe

This bit controls the timing for assertion of a chip select in asynchronous mode. Selecting address strobe causes chip select to be asserted synchronized with address strobe. Selecting data strobe causes chip select to be asserted synchronized with data strobe.

PEPAR — Port E Pin Assignment Register **$YFFA17**

15 … 8	7	6	5	4	3	2	1	0
NOT USED	PEPA7	PEPA6	PEPA5	PEPA4	PEPA3	PEPA2	PEPA1	PEPA0
RESET:	DATA8	DATA8	DATA8	DATA8	DATA8	DATA8	DATA8	DATA8

The bits in this register control the function of each port E pin. Any bit set to one configures the corresponding pin as a bus control signal, with the function shown in the following table. Any bit cleared to zero defines the corresponding pin to be an I/O pin, controlled by PORTE and DDRE.

Data bus bit 8 controls the state of this register following reset. If DATA8 is set to one during reset, the register is set to $FF, which defines all port E pins as bus control signals. If DATA8 is cleared to zero during reset, this register is set to $00, configuring all port E pins as I/O pins.

Any bit cleared to zero defines the corresponding pin to be an I/O pin. Any bit set to one defines the corresponding pin to be a bus control signal.

Table 16 Port E Pin Assignments

PEPAR Bit	Port E Signal	Bus Control Signal
PEPA7	PE7	SIZ1
PEPA6	PE6	SIZ0
PEPA5	PE5	$\overline{AS}$
PEPA4	PE4	$\overline{DS}$
PEPA3	PE3	$\overline{RMC}$
PEPA2	PE2	$\overline{AVEC}$
PEPA1	PE1	$\overline{DSACK1}$
PEPA0	PE0	$\overline{DSACK0}$

PORTF0, PORTF1 — Port F Data Register **$YFFA19, $YFFA1B**

15 … 8	7	6	5	4	3	2	1	0
NOT USED	PF7	PF6	PF5	PF4	PF3	PF2	PF1	PF0
RESET:	U	U	U	U	U	U	U	U

The write to the port F data register is stored in the internal data latch, and if any port F pin is configured as an output, the value stored for that bit is driven onto the pin. A read of the port F data register returns the value at the pin only if the pin is configured as a discrete input. Otherwise, the value read is the value stored in the register.

The port F data register is a single register that can be accessed in two locations. When accessed at $YFFA19, the register is referred to as PORTF0; when accessed at $YFFA1B, the register is referred to as PORTF1. The register can be read or written at any time. It is unaffected by reset.

DDRF — Port F Data Direction Register **$YFFA1D**

15 … 8	7	6	5	4	3	2	1	0
NOT USED	DDF7	DDF6	DDF5	DDF4	DDF3	DDF2	DDF1	DDF0
RESET:	0	0	0	0	0	0	0	0

The bits in this register control the direction of the pin drivers when the pins are configured for I/O. Any bit in this register set to one configures the corresponding pin as an output. Any bit in this register cleared to zero configures the corresponding pin as an input.

PFPAR — Port F Pin Assignment Register **$YFFA1F**

15 … 8	7	6	5	4	3	2	1	0
NOT USED	PFPA7	PFPA6	PFPA5	PFPA4	PFPA3	PFPA2	PFPA1	PFPA0
RESET:	DATA9	DATA9	DATA9	DATA9	DATA9	DATA9	DATA9	DATA9

The bits in this register control the function of each port F pin. Any bit cleared to zero defines the corresponding pin to be an I/O pin. Any bit set to one defines the corresponding pin to be an interrupt request signal or MODCLK. The MODCLK signal has no function after reset.

Table 17 Port F Pin Assignments

PFPAR Field	Port F Signal	Alternate Signal
PFPA7	PF7	$\overline{\text{IRQ7}}$
PFPA6	PF6	$\overline{\text{IRQ6}}$
PFPA5	PF5	$\overline{\text{IRQ5}}$
PFPA4	PF4	$\overline{\text{IRQ4}}$
PFPA3	PF3	$\overline{\text{IRQ3}}$
PFPA2	PF2	$\overline{\text{IRQ2}}$
PFPA1	PF1	$\overline{\text{IRQ1}}$
PFPA0	PF0	MODCLK

Data bus pin 9 controls the state of this register following reset. If DATA9 is set to one during reset, the register is set to $FF, which defines all port F pins as interrupt request inputs. If DATA9 is cleared to zero during reset, this register is set to $00, defining all port F pins as I/O pins.

3.7 Resets

Reset procedures handle system initialization and recovery from catastrophic failure. The MCU performs resets with a combination of hardware and software. The system integration module determines whether a reset is valid, asserts control signals, performs basic system configuration based on hardware mode-select inputs, then passes control to the CPU.

Reset occurs when an active low logic level on the $\overline{\text{RESET}}$ pin is clocked into the SIM. Resets are gated by the CLKOUT signal. Asynchronous resets are assumed to be catastrophic. An asynchronous reset can occur on any clock edge. Synchronous resets are timed to occur at the end of bus cycles. If there is no clock when $\overline{\text{RESET}}$ is asserted, reset does not occur until the clock starts. Resets are clocked in order to allow completion of write cycles in progress at the time $\overline{\text{RESET}}$ is asserted.

Reset is the highest-priority CPU32 exception. Any processing in progress is aborted by the reset exception, and cannot be restarted. Only essential tasks are performed during reset exception processing. Other initialization tasks must be accomplished by the exception handler routine.

3.7.1 SIM Reset Mode Selection

The logic states of certain data bus pins during reset determine SIM operating configuration. In addition, the state of the MODCLK pin determines system clock source and the state of the $\overline{\text{BKPT}}$ pin determines what happens during subsequent breakpoint assertions. The following table is a summary of reset mode selection options.

Table 18 Reset Mode Selection

Mode Select Pin	Default Function (Pin Left High)	Alternate Function (Pin Pulled Low)

3.7.3 Reset Timing

The $\overline{\text{RESET}}$ input must be asserted for a specified minimum period in order for reset to occur. External $\overline{\text{RESET}}$ assertion can be delayed internally for a period equal to the longest bus cycle time (or the bus monitor time-out period) in order to protect write cycles from being aborted by reset. While $\overline{\text{RESET}}$ is asserted, SIM pins are either in a disabled high-impedance state or are driven to their inactive states.

When an external device asserts $\overline{\text{RESET}}$ for the proper period, reset control logic clocks the signal into an internal latch. The control logic drives the $\overline{\text{RESET}}$ pin low for an additional 512 CLKOUT cycles after it detects that the $\overline{\text{RESET}}$ signal is no longer being externally driven, to guarantee this length of reset to the entire system.

If an internal source asserts a reset signal, the reset control logic asserts $\overline{\text{RESET}}$ for a minimum of 512 cycles. If the reset signal is still asserted at the end of 512 cycles, the control logic continues to assert $\overline{\text{RESET}}$ until the internal reset signal is negated.

After 512 cycles have elapsed, the reset input pin goes to an inactive, high-impedance state for ten cycles. At the end of this 10-cycle period, the reset input is tested. When the input is at logic level one, reset exception processing begins. If, however, the reset input is at logic level zero, the reset control logic drives the pin low for another 512 cycles. At the end of this period, the pin again goes to high-impedance state for ten cycles, then it is tested again. The process repeats until $\overline{\text{RESET}}$ is released.

3.7.4 Power-On Reset

When the SIM clock synthesizer is used to generate the system clock, power-on reset involves special circumstances related to application of system and clock synthesizer power. Regardless of clock source, voltage must be applied to clock synthesizer power input pin V_{DDSYN} in order for the MCU to operate. The following discussion assumes that V_{DDSYN} is applied before and during reset. This minimizes crystal start-up time. When V_{DDSYN} is applied at power-on, start-up time is affected by specific crystal parameters and by oscillator circuit design. V_{DD} ramp-up time also affects pin state during reset.

During power-on reset, an internal circuit in the SIM drives the internal (IMB) and external reset lines. The circuit releases the internal reset line as V_{DD} ramps up to the minimum specified value, and SIM pins are initialized. When V_{DD} reaches the specified minimum value, the clock synthesizer VCO begins operation. Clock frequency ramps up to the specified limp mode frequency. The external $\overline{\text{RESET}}$ line remains asserted until the clock synthesizer PLL locks and 512 CLKOUT cycles elapse.

The SIM clock synthesizer provides clock signals to the other MCU modules. After the clock is running and the internal reset signal is asserted for four clock cycles, these modules reset. V_{DD} ramp time and VCO frequency ramp time determine how long these four cycles take. Worst case is approximately 15 milliseconds. During this period, module port pins may be in an indeterminate state. While input-only pins can be put in a known state by means of external pull-up resistors, external logic on input/output or output-only pins must condition the lines during this time. Active drivers require high-impedance buffers or isolation resistors to prevent conflict.

3.7.5 Use of Three State Control Pin

Asserting the three-state control (TSC) input causes the MCU to put all output drivers in an inactive, high-impedance state. The signal must remain asserted for ten clock cycles in order for drivers to change state. There are certain constraints on use of TSC during power-on reset:

When the internal clock synthesizer is used (MODCLK held high during reset), synthesizer ramp-up time affects how long the ten cycles take. Worst case is approximately 20 milliseconds from TSC assertion.

When an external clock signal is applied (MODCLK held low during reset), pins go to high-impedance state as soon after TSC assertion as ten clock pulses have been applied to the EXTAL pin.

mask lower-priority interrupts during exception processing, and it is decoded by modules that have requested interrupt service to determine whether the current interrupt acknowledge cycle pertains to them.

Modules that have requested interrupt service decode the IP value placed on the address bus at the beginning of the interrupt acknowledge cycle, and if their requests are at the specified IP level, respond to the cycle. Arbitration between simultaneous requests of the same priority is performed by means of serial contention between module interrupt arbitration (IARB) field bit values.

Each module that can make an interrupt service request, including the SIM, has an IARB field in its configuration register. An IARB field can be assigned a value from %0001 (lowest priority) to %1111 (highest priority). A value of %0000 in an IARB field causes the CPU to process a spurious interrupt exception when an interrupt from that module is recognized.

Because the EBI manages external interrupt requests, the SIM IARB value is used for arbitration between internal and external interrupt requests. The reset value of IARB for the SIM is %1111, and the reset IARB value for all other modules is %0000. Initialization software must assign different IARB values in order to implement an arbitration scheme.

Each module must have a unique IARB value. When two or more IARB fields have the same nonzero value, the CPU interprets multiple vector numbers simultaneously, with unpredictable consequences.

Arbitration must always take place, even when a single source requests service. This point is important for two reasons: the CPU interrupt acknowledge cycle is not driven on the external bus unless the SIM wins contention, and failure to contend causes an interrupt acknowledge bus cycle to be terminated by a bus error, which causes a spurious interrupt exception to be taken.

When arbitration is complete, the dominant module must place an interrupt vector number on the data bus and terminate the bus cycle. In the case of an external interrupt request, because the interrupt acknowledge cycle is transferred to the external bus, an external device must decode the mask value and respond with a vector number, then generate bus cycle termination signals. If the device does not respond in time, a spurious interrupt exception is taken.

The periodic interrupt timer (PIT) in the SIM can generate internal interrupt requests of specific priority at predetermined intervals. By hardware convention, PIT interrupts are serviced before external interrupt service requests of the same priority. Refer to 3.2.7 Periodic Interrupt Timer for more information.

3.8.2 Interrupt Processing Summary

A summary of the interrupt processing sequence follows. When the sequence begins, a valid interrupt service request has been detected and is pending.

A. The CPU finishes higher priority exception processing or reaches an instruction boundary.
B. Processor state is stacked. The contents of the status register and program counter are saved.
C. The interrupt acknowledge cycle begins:
    1. FC[2:0] are driven to %111 (CPU space) encoding.
    2. The address bus is driven as follows. ADDR[23:20] = %1111; ADDR[19:16] = %1111, which indicates that the cycle is an interrupt acknowledge CPU space cycle; ADDR[15:4] = %111111111111; ADDR[3:1] = the level of the interrupt request being acknowledged; and ADDR0 = %1.
    3. Request priority level is latched into the IP field in the status register from the address bus.
D. Modules or external peripherals that have requested interrupt service decode the request level in ADDR[3:1]. If the request level of at least one interrupting module or device is the same as the value in ADDR[3:1], interrupt arbitration contention takes place. When there is no contention, the spurious interrupt monitor asserts $\overline{\text{BERR}}$, and a spurious interrupt exception is processed.
E. After arbitration, the interrupt acknowledge cycle can be completed in one of three ways:

1. The dominant interrupt source supplies a vector number and $\overline{\text{DSACK}}$ signals appropriate to the access. The CPU32 acquires the vector number.
2. The $\overline{\text{AVEC}}$ signal is asserted (the signal can be asserted by the dominant interrupt source or the pin can be tied low), and the CPU32 generates an autovector number corresponding to interrupt priority.
3. The bus monitor asserts $\overline{\text{BERR}}$ and the CPU32 generates the spurious interrupt vector number.

F. The vector number is converted to a vector address.
G. The content of the vector address is loaded into the PC, and the processor transfers control to the exception handler routine.

3.9 Factory Test Block

The test submodule supports scan-based testing of the various MCU modules. It is integrated into the SIM to support production testing.

Test submodule registers are intended for Motorola use. Register names and addresses are provided to indicate that these addresses are occupied.

SIMTR —System Integration Test Register **$YFFA02**

SIMTRE —System Integration Test Register (E Clock) **$YFFA08**

TSTMSRA —Master Shift Register A **$YFFA30**

TSTMSRB —Master Shift Register B **$YFFA32**

TSTSC —Test Module Shift Count **$YFFA34**

TSTRC —Test Module Repetition Count **$YFFA36**

CREG —Test Module Control Register **$YFFA38**

DREG —Test Module Distributed Register **$YFFA3A**

5.4 Parameter RAM

Parameter RAM occupies 256 bytes at the top of the TPU module address map. Channel parameters are organized as 128 16-bit words. However, only 100 words are actually implemented. The parameter RAM address map shows how parameter words are organized in memory.

Table 23 TPU Parameter RAM Address Map

Channel Number	Base Address	Parameter Address							
		0	1	2	3	4	5	6	7
0	$YFFFF##	00	02	04	06	08	0A	—	—
1	$YFFFF##	10	12	14	16	18	1A	—	—
2	$YFFFF##	20	22	24	26	28	2A	—	—
3	$YFFFF##	30	32	34	36	38	3A	—	—
4	$YFFFF##	40	42	44	46	48	4A	—	—
5	$YFFFF##	50	52	54	56	58	5A	—	—
6	$YFFFF##	60	62	64	66	68	6A	—	—
7	$YFFFF##	70	72	74	76	78	7A	—	—
8	$YFFFF##	80	82	84	86	88	8A	—	—
9	$YFFFF##	90	92	94	96	98	9A	—	—
10	$YFFFF##	A0	A2	A4	A6	A8	AA	—	—
11	$YFFFF##	B0	B2	B4	B6	B8	BA	—	—
12	$YFFFF##	C0	C2	C4	C6	C8	CA	—	—
13	$YFFFF##	D0	D2	D4	D6	D8	DA	—	—
14	$YFFFF##	E0	E2	E4	E6	E8	EA	EC	EE
15	$YFFFF##	F0	F2	F4	F6	F8	FA	FC	FE

—= Not Implemented
Y = M111, where M represents the logic state of the MM bit in the SIMCR.

5.5 TPU Registers

The TPU memory map contains three groups of registers:

System Configuration Registers
Channel Control and Status Registers
Development Support and Test Verification Registers

5.5.1 System Configuration Registers

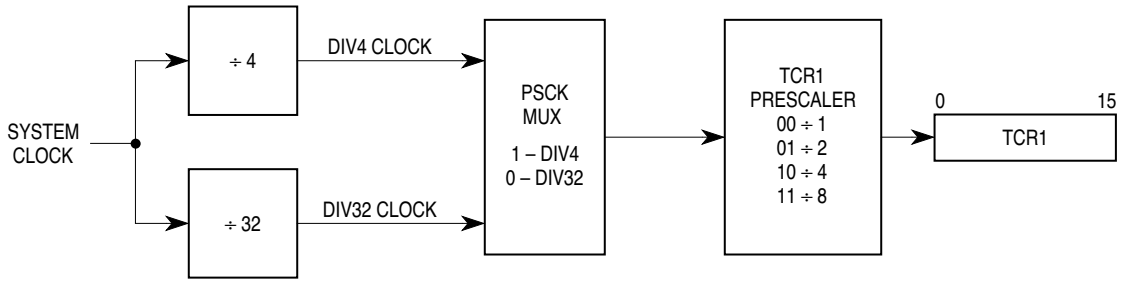
TPUMCR — TPU Module Configuration Register **$YFFE00**

15	14	13	12	11	10	9	8	7	6	5	4	3				0
STOP	TCR1P		TCR2P		EMU	T2CG	STF	SUPV	PSCK	0	0	IARB				
RESET:																
0	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	0

STOP — Stop Bit
0 = TPU operating normally
1 = Internal clocks shut down

TCR1P — Timer Count Register 1 Prescaler Control

TCR1 is clocked from the output of a prescaler. The prescaler's input is the internal TPU system clock divided by either 4 or 32, depending on the value of the PSCK bit. The prescaler divides this input by 1, 2, 4, or 8. Channels using TCR1 have the capability to resolve down to the TPU system clock divided by 4.

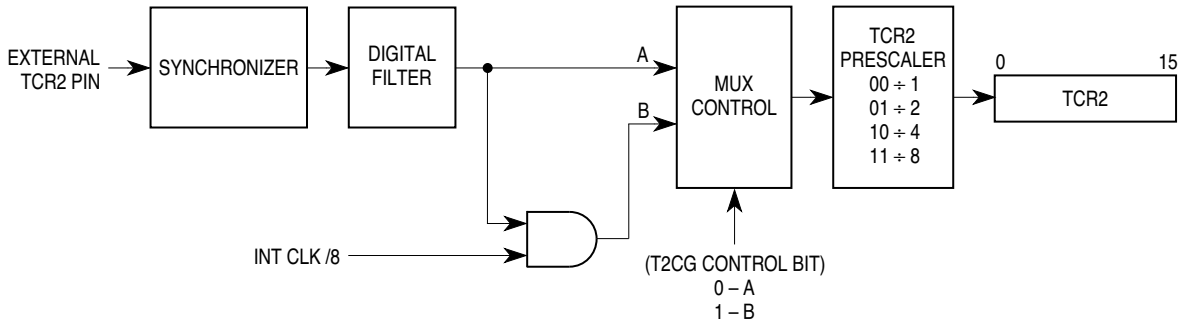


PRESCALER CTL BLOCK 1

TCR1 Prescaler	Divide By	PSCK = 0		PSCK = 1	
		Number of Clocks	Rate at 16 MHz	Number of Clocks	Rate at 16 MHz
00	1	32	2 ms	4	250 ns
01	2	64	4 ms	8	500 ns
10	4	128	8 ms	16	1 ms
11	8	256	16 ms	32	2 ms

TCR2P — Timer Count Register 2 Prescaler Control

TCR2 is clocked from the output of a prescaler. If T2CG = 0, the input to the TCR2 prescaler is the external TCR2 clock source. If T2CG = 1, the input is the TPU system clock divided by eight. The TCR2P field specifies the value of the prescaler: 1, 2, 4, or 8. Channels using TCR2 have the capability to resolve down to the TPU system clock divided by 8. The following table is a summary of prescaler output.



PRESCALER CTL BLOCK 2

TCR2 Prescaler	Divide By	Internal Clock Divided By	External Clock Divided By
00	1	8	1
01	2	16	2
10	4	32	4
11	8	64	8

5.5.2 Channel Control Registers

CIER — Channel Interrupt Enable Register **$YFFE0A**

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 15	CH 14	CH 13	CH 12	CH 11	CH 10	CH 9	CH 8	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
RESET:															
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

CH[15:0] — Channel Interrupt Enable/Disable
0 = Channel interrupts disabled
1 = Channel interrupts enabled

CISR — Channel Interrupt Status Register **$YFFE20**

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
CH 15	CH 14	CH 13	CH 12	CH 11	CH 10	CH 9	CH 8	CH 7	CH 6	CH 5	CH 4	CH 3	CH 2	CH 1	CH 0
RESET:															
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0

CH[15:0] — Channel Interrupt Status Bit
0 = Channel interrupt not asserted
1 = Channel interrupt asserted

CFSR0 — Channel Function Select Register 0 **$YFFE0C**

15	12	11	8	7	4	3	0
CHANNEL15		CHANNEL14		CHANNEL13		CHANNEL12	
RESET:							
0	0	0	0	0	0	0	0

CFSR1 — Channel Function Select Register 1 **$YFFE0E**

15	12	11	8	7	4	3	0
CHANNEL11		CHANNEL10		CHANNEL9		CHANNEL8	
RESET:							
0	0	0	0	0	0	0	0

CFSR2 — Channel Function Select Register 2 **$YFFE10**

15	12	11	8	7	4	3	0
CHANNEL7		CHANNEL6		CHANNEL5		CHANNEL4	
RESET:							
0	0	0	0	0	0	0	0

CFSR3 — Channel Function Select Register 3 **$YFFE12**

15	12	11	8	7	4	3	0
CHANNEL3		CHANNEL2		CHANNEL1		CHANNEL0	
RESET:							
0	0	0	0	0	0	0	0

CHANNEL[15:0] — Encoded Time Function for each Channel
Encoded 4-bit fields in the channel function select registers specify one of 16 time functions to be executed on the corresponding channel.

6.3 Pin Function

The following table is a summary of the functions of the QSM pins when they are not configured for general-purpose I/O. The QSM data direction register (DDRQS) designates each pin except RXD as an input or output.

	Pin	Mode	Pin Function
QSPI Pins	MISO	Master	Serial Data Input to QSPI
		Slave	Serial Data Output from QSPI
	MOSI	Master	Serial Data Output from QSPI
		Slave	Serial Data Input to QSPI
	SCK	Master	Clock Output from QSPI
		Slave	Clock Input to QSPI
	PCS0/$\overline{SS}$	Master	Input: Assertion Causes Mode Fault Output: Selects Peripherals
		Slave	Input: Selects the QSPI
	PCS[3:1]	Master	Output: Selects Peripherals
		Slave	None
SCI Pins	TXD	Transmit	Serial Data Output from SCI
	RXD	Receive	Serial Data Input to SCI

6.4 QSM Registers

QSM registers are divided into four categories: QSM global registers, QSM pin control registers, QSPI submodule registers, and SCI submodule registers. The QSPI and SCI registers are defined in separate sections below. Writes to unimplemented register bits have no meaning or effect, and reads from unimplemented bits always return a logic zero value.

The module mapping bit of the SIM configuration register (SIMCR) defines the most significant bit (ADDR23) of the address, shown in each register figure as Y (Y = $7 or $F). This bit, concatenated with the rest of the address given, forms the absolute address of each register. Refer to the SIM section of this technical summary for more information about how the state of MM affects the system.

6.4.1 Global Registers

The QSM global registers contain system parameters used by both the QSPI and the SCI submodules. These registers contain the bits and fields used to configure the QSM.

QSMCR — QSM Configuration Register **$YFFC00**

15	14	13	12	11	10	9	8	7	6	5	4	3				0
STOP	FRZ1	FRZ0	0	0	0	0	0	SUPV	0	0	0	IARB				
RESET:																
0	0	0	0	0	0	0	0	1	0	0	0	0	0	0	0	

The QSMCR contains parameters for the QSM/CPU/intermodule bus (IMB) interface.

STOP — Stop Enable

0 = Normal QSM clock operation

1 = QSM clock operation stopped

STOP places the QSM in a low-power state by disabling the system clock in most parts of the module. The QSMCR is the only register guaranteed to be readable while STOP is asserted. The QSPI RAM is not readable. However, writes to RAM or any register are guaranteed to be valid while STOP is asserted. STOP can be negated by the CPU and by reset.

SCK baud rate:

SCK Baud Rate = System Clock/(2SPBR)

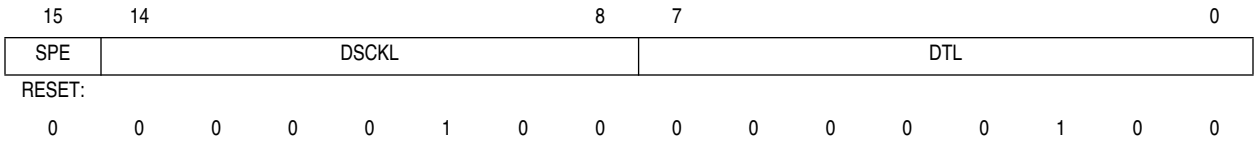
or

SPBR = System Clock/(2SCK)(Baud Rate Desired)

where SPBR equals {2, 3, 4,..., 255}

Giving SPBR a value of zero or one disables the baud rate generator. SCK is disabled and assumes its inactive state value. No serial transfers occur. At reset, baud rate is initialized to one eighth of the system clock frequency.

SPCR1 — QSPI Control Register 1 **$YFFC1A**



SPCR1 contains parameters for configuring the QSPI before it is enabled. The CPU can read and write this register, but the QSM has read access only, except for SPE, which is automatically cleared by the QSPI after completing all serial transfers, or when a mode fault occurs.

SPE — QSPI Enable

0 = QSPI is disabled. QSPI pins can be used for general-purpose I/O.
1 = QSPI is enabled. Pins allocated by PQSPAR are controlled by the QSPI.

DSCKL — Delay before SCK

When the DSCK bit in command RAM is set, this field determines the length of delay from PCS valid to SCK transition. PCS can be any of the four peripheral chip-select pins. The following equation determines the actual delay before SCK:

PCS to SCK Delay = [DSCKL/System Clock]

where DSCKL equals {1, 2, 3,..., 127}.

When the DSCK value of a queue entry equals zero, then DSCKL is not used. Instead, the PCS valid-to-SCK transition is one-half SCK period.

DTL — Length of Delay after Transfer

When the DT bit in command RAM is set, this field determines the length of delay after serial transfer. The following equation is used to calculate the delay:

Delay after Transfer = [(32DTL)/System Clock]

where DTL equals {1, 2, 3,..., 255}.

A zero value for DTL causes a delay-after-transfer value of 8192/System Clock.

If DT equals zero, a standard delay is inserted.

Standard Delay after Transfer = [17/System Clock]

Delay after transfer can be used to provide a peripheral deselect interval. A delay can also be inserted between consecutive transfers to allow serial A/D converters to complete conversion.

SBK — Send Break

0 = Normal operation

1 = Break frame(s) transmitted after completion of current frame

SBK provides the ability to transmit a break code from the SCI. If the SCI is transmitting when SBK is set, it will transmit continuous frames of zeros after it completes the current frame, until SBK is cleared. If SBK is toggled (one to zero in less than one frame interval), the transmitter sends only one or two break frames before reverting to idle line or beginning to send data.

SCSR — SCI Status Register **$YFFC0C**

15 … 9	8	7	6	5	4	3	2	1	0
NOT USED	TDRE	TC	RDRF	RAF	IDLE	OR	NF	FE	PF
RESET:	1	1	0	0	0	0	0	0	0

SCSR contains flags that show SCI operational conditions. These flags can be cleared either by hardware or by a special acknowledgment sequence. The sequence consists of SCSR read with flags set, followed by SCDR read (write in the case of TDRE and TC). A long-word read can consecutively access both SCSR and SCDR. This action clears receive status flag bits that were set at the time of the read, but does not clear TDRE or TC flags.

If an internal SCI signal for setting a status bit comes after the CPU has read the asserted status bits, but before the CPU has written or read register SCDR, the newly set status bit is not cleared. SCSR must be read again with the bit set. Also, SCDR must be written or read before the status bit is cleared.

Reading either byte of SCSR causes all 16 bits to be accessed. Any status bit already set in either byte will be cleared on a subsequent read or write of register SCDR.

TDRE — Transmit Data Register Empty Flag

0 = Register TDR still contains data to be sent to the transmit serial shifter.

1 = A new character can now be written to register TDR.

TDRE is set when the byte in register TDR is transferred to the transmit serial shifter. If TDRE is zero, transfer has not occurred and a write to TDR will overwrite the previous value. New data is not transmitted if TDR is written without first clearing TDRE.

TC — Transmit Complete Flag

0 = SCI transmitter is busy

1 = SCI transmitter is idle

TC is set when the transmitter finishes shifting out all data, queued preambles (mark/idle line), or queued breaks (logic zero). The interrupt can be cleared by reading SCSR when TC is set and then by writing the transmit data register (TDR) of SCDR.

RDRF — Receive Data Register Full Flag

0 = Register RDR is empty or contains previously read data.

1 = Register RDR contains new data.

RDRF is set when the content of the receive serial shifter is transferred to the RDR. If one or more errors are detected in the received word, flag(s) NF, FE, and/or PF are set within the same clock cycle.

RAF — Receiver Active Flag

0 = SCI receiver is idle

1 = SCI receiver is busy

RAF indicates whether the SCI receiver is busy. It is set when the receiver detects a possible start bit and is cleared when the chosen type of idle line is detected. RAF can be used to reduce collisions in systems with multiple masters.

7 Standby RAM with TPU Emulation RAM

The TPURAM module contains a 2-Kbyte array of fast (two bus cycle) static RAM, which is especially useful for system stacks and variable storage. Alternately, it can be used by the TPU as emulation RAM for new timer algorithms.

7.1 Overview

The TPURAM can be mapped to any 4-Kbyte boundary in the address map, but must not overlap the module control registers. (Overlap makes the registers inaccessible.) Data can be read or written in bytes, word, or long words. TPURAM responds to both program and data space accesses. Data can be read or written in bytes, words, or long words. The TPURAM is powered by V_{DD} in normal operation. During power-down, the TPURAM contents are maintained by power on standby voltage pin V_{STBY}. Power switching between sources is automatic.

Access to the TPURAM array is controlled by the RASP field in TRAMMCR. This field can be encoded so that TPURAM responds to both program and data space accesses. This allows code to be executed from TPURAM, and permits the use of program counter relative addressing mode for operand fetches from the array.

An address map of the TPURAM control registers follows. All TPURAM control registers are located in supervisor data space.

Table 28 TPURAM Control Register Address Map

Access	Address	15 8	7 0
S	$YFFB00	TPURAM MODULE CONFIGURATION REGISTER (TRAMMCR)	
S	$YFFB02	TPURAM TEST REGISTER (TRAMTST)	
S	$YFFB04	TPURAM BASE ADDRESS REGISTER (TRAMBAR)	
	$YFFB06–$YFFB3F	NOT USED	

Y = M111, where M is the logic state of the MM bit in the SIMCR.

7.2 TPURAM Register Block

There are three TPURAM control registers: the RAM module configuration register (TRAMMCR), the RAM test register (TRAMTST), and the RAM array base address registers (TRAMBAR).

There is an 8-byte minimum register block size for the module. Unimplemented register addresses are read as zeros, and writes have no effect.

7.3 TPURAM Registers

TRAMMCR —TPURAM Module Configuration Register **$YFFB00**

15	14	13	12	11	10	9	8	7 0
STOP	0	0	0	0	0	0	RASP	NOT USED
RESET:								
0	0	0	0	0	0	0	1	

TSTOP —Stop Control

0 = RAM array operates normally.

1 = RAM array enters low-power stop mode.

This bit controls whether the RAM array is in stop mode or normal operation. Reset state is zero, for normal operation. In stop mode, the array retains its contents, but cannot be read or written by the CPU.

8 Summary of Changes

This is a partial revision. Most of the publication remains the same, but the following changes were made to improve it. Typographical errors that do not affect content are not annotated. This document has also been reformatted for use on the web.

Pages 2-3	New Ordering Information included.
Page 6	New block diagram drawn.
Page 7	New 132-pin assignment diagram drawn.
Page 8	New 144-pin assignment diagram drawn.
Page 9	New address map drawn.
Pages 10-14	Added Signal Description section.
Pages 15-47	Expanded and revised SIM section. Made all register diagrams and bit mnemonics consistent. Incorporated new information concerning the system clock, resets, interrupts, and chip-selects circuits.
Page 48-56	Expanded and revised CPU section. Made all register diagrams and bit mnemonics consistent. Revised instruction set summary information.
Page 57-70	Expanded and revised TPU section. Made all register diagrams and bit mnemonics consistent. Revised time functions information to include both MC68332A and MC68332G microcode ROM applications.
Page 71-92	Expanded and revised QSM section. Made all register diagrams and bit mnemonics consistent. Added information concerning SPI and SCI operation.
Page 93-95	Revised Standby RAM with TPU Emulation RAM section. Made all register diagrams and bit mnemonics consistent.

Home Page:
www.freescale.com
email:
support@freescale.com
USA/Europe or Locations Not Listed:
Freescale Semiconductor
Technical Information Center, CH370
1300 N. Alma School Road
Chandler, Arizona 85224
(800) 521-6274
480-768-2130
support@freescale.com
Europe, Middle East, and Africa:
Freescale Halbleiter Deutschland GmbH
Technical Information Center
Schatzbogen 7
81829 Muenchen, Germany
+44 1296 380 456 (English)
+46 8 52200080 (English)
+49 89 92103 559 (German)
+33 1 69 35 48 48 (French)
support@freescale.com
Japan:
Freescale Semiconductor Japan Ltd.
Headquarters
ARCO Tower 15F
1-8-1, Shimo-Meguro, Meguro-ku
Tokyo 153-0064, Japan
0120 191014
+81 2666 8080
support.japan@freescale.com
Asia/Pacific:
Freescale Semiconductor Hong Kong Ltd.
Technical Information Center
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Tai Po Industrial Estate,
Tai Po, N.T., Hong Kong
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