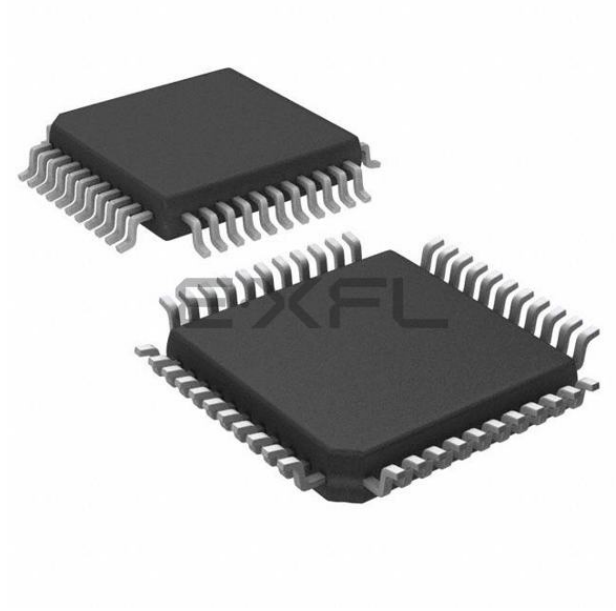


Welcome to E-XFL.COM

What is "Embedded - Microcontrollers"?

"Embedded - Microcontrollers" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	8MHz
Connectivity	UART/USART
Peripherals	POR, PWM, WDT
Number of I/O	33
Program Memory Size	16KB (8K x 16)
Program Memory Type	OTP
EEPROM Size	-
RAM Size	454 x 8
Voltage - Supply (Vcc/Vdd)	2.5V ~ 6V
Data Converters	-
Oscillator Type	External
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	44-QFP
Supplier Device Package	44-MQFP (10x10)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic17lc44t-08i-pq

5.2 Peripheral Interrupt Enable Register (PIE)

This register contains the individual flag bits for the Peripheral interrupts.

FIGURE 5-3: PIE REGISTER (ADDRESS: 17h, BANK 1)

R/W - 0	R/W - 0	R/W - 0	R/W - 0	R/W - 0	R/W - 0	R/W - 0	R/W - 0
RBIE	TMR3IE	TMR2IE	TMR1IE	CA2IE	CA1IE	TXIE	RCIE
bit7							bit0

R = Readable bit
W = Writable bit
-n = Value at POR reset

bit 7: **RBIE**: PORTB Interrupt on Change Enable bit
1 = Enable PORTB interrupt on change
0 = Disable PORTB interrupt on change

bit 6: **TMR3IE**: Timer3 Interrupt Enable bit
1 = Enable Timer3 interrupt
0 = Disable Timer3 interrupt

bit 5: **TMR2IE**: Timer2 Interrupt Enable bit
1 = Enable Timer2 interrupt
0 = Disable Timer2 interrupt

bit 4: **TMR1IE**: Timer1 Interrupt Enable bit
1 = Enable Timer1 interrupt
0 = Disable Timer1 interrupt

bit 3: **CA2IE**: Capture2 Interrupt Enable bit
1 = Enable Capture interrupt on RB1/CAP2 pin
0 = Disable Capture interrupt on RB1/CAP2 pin

bit 2: **CA1IE**: Capture1 Interrupt Enable bit
1 = Enable Capture interrupt on RB2/CAP1 pin
0 = Disable Capture interrupt on RB2/CAP1 pin

bit 1: **TXIE**: USART Transmit Interrupt Enable bit
1 = Enable Transmit buffer empty interrupt
0 = Disable Transmit buffer empty interrupt

bit 0: **RCIE**: USART Receive Interrupt Enable bit
1 = Enable Receive buffer full interrupt
0 = Disable Receive buffer full interrupt

5.4 Interrupt Operation

Global Interrupt Disable bit, GLINTD (CPUSTA<4>), enables all unmasked interrupts (if clear) or disables all interrupts (if set). Individual interrupts can be disabled through their corresponding enable bits in the INTSTA register. Peripheral interrupts need either the global peripheral enable PEIE bit disabled, or the specific peripheral enable bit disabled. Disabling the peripherals via the global peripheral enable bit, disables all peripheral interrupts. GLINTD is set on reset (interrupts disabled).

The RETFIE instruction allows returning from interrupt and re-enable interrupts at the same time.

When an interrupt is responded to, the GLINTD bit is automatically set to disable any further interrupt, the return address is pushed onto the stack and the PC is loaded with interrupt vector. There are four interrupt vectors to reduce interrupt latency.

The peripheral interrupt vector has multiple interrupt sources. Once in the peripheral interrupt service routine, the source(s) of the interrupt can be determined by polling the interrupt flag bits. The peripheral interrupt flag bit(s) must be cleared in software before re-enabling interrupts to avoid continuous interrupts.

The PIC17C4X devices have four interrupt vectors. These vectors and their hardware priority are shown in Table 5-1. If two enabled interrupts occur “at the same time”, the interrupt of the highest priority will be serviced first. This means that the vector address of that interrupt will be loaded into the program counter (PC).

TABLE 5-1: INTERRUPT VECTORS/ PRIORITIES

Address	Vector	Priority
0008h	External Interrupt on RA0/ INT pin (INTF)	1 (Highest)
0010h	TMR0 overflow interrupt (T0IF)	2
0018h	External Interrupt on T0CKI (T0CKIF)	3
0020h	Peripherals (PEIF)	4 (Lowest)

Note 1: Individual interrupt flag bits are set regardless of the status of their corresponding mask bit or the GLINTD bit.

Note 2: When disabling any of the INTSTA enable bits, the GLINTD bit should be set (disabled).

Note 3: For the PIC17C42 only:
If an interrupt occurs while the Global Interrupt Disable (GLINTD) bit is being set, the GLINTD bit may unintentionally be re-enabled by the user’s Interrupt Service Routine (the RETFIE instruction). The events that would cause this to occur are:

1. An interrupt occurs simultaneously with an instruction that sets the GLINTD bit.
2. The program branches to the Interrupt vector and executes the Interrupt Service Routine.
3. The Interrupt Service Routine completes with the execution of the RETFIE instruction. This causes the GLINTD bit to be cleared (enables interrupts), and the program returns to the instruction after the one which was meant to disable interrupts.

The method to ensure that interrupts are globally disabled is:

1. Ensure that the GLINTD bit was set by the instruction, as shown in the following code:

```
LOOP   BSF    CPUSTA, GLINTD ; Disable Global
                              ; Interrupt
       BTFSS  CPUSTA, GLINTD ; Global Interrupt
                              ; Disabled?
       GOTO   LOOP           ; NO, try again
                              ; YES, continue
                              ; with program
                              ; low
```

6.4.1 INDIRECT ADDRESSING REGISTERS

The PIC17C4X has four registers for indirect addressing. These registers are:

- INDF0 and FSR0
- INDF1 and FSR1

Registers INDF0 and INDF1 are not physically implemented. Reading or writing to these registers activates indirect addressing, with the value in the corresponding FSR register being the address of the data. The FSR is an 8-bit register and allows addressing anywhere in the 256-byte data memory address range. For banked memory, the bank of memory accessed is specified by the value in the BSR.

If file INDF0 (or INDF1) itself is read indirectly via an FSR, all '0's are read (Zero bit is set). Similarly, if INDF0 (or INDF1) is written to indirectly, the operation will be equivalent to a NOP, and the status bits are not affected.

6.4.2 INDIRECT ADDRESSING OPERATION

The indirect addressing capability has been enhanced over that of the PIC16CXX family. There are two control bits associated with each FSR register. These two bits configure the FSR register to:

- Auto-decrement the value (address) in the FSR after an indirect access
- Auto-increment the value (address) in the FSR after an indirect access
- No change to the value (address) in the FSR after an indirect access

These control bits are located in the ALUSTA register. The FSR1 register is controlled by the FS3:FS2 bits and FSR0 is controlled by the FS1:FS0 bits.

When using the auto-increment or auto-decrement features, the effect on the FSR is not reflected in the ALUSTA register. For example, if the indirect address causes the FSR to equal '0', the Z bit will not be set.

If the FSR register contains a value of 0h, an indirect read will read 0h (Zero bit is set) while an indirect write will be equivalent to a NOP (status bits are not affected).

Indirect addressing allows single cycle data transfers within the entire data space. This is possible with the use of the MOVPF and MOVFP instructions, where either 'p' or 'f' is specified as INDF0 (or INDF1).

If the source or destination of the indirect address is in banked memory, the location accessed will be determined by the value in the BSR.

A simple program to clear RAM from 20h - FFh is shown in Example 6-1.

EXAMPLE 6-1: INDIRECT ADDRESSING

```
    MOVLW    0x20          ;
    MOVWF    FSR0          ; FSR0 = 20h
    BCF      ALUSTA, FS1   ; Increment FSR
    BSF      ALUSTA, FS0   ; after access
    BCF      ALUSTA, C     ; C = 0
    MOVLW    END_RAM + 1   ;
LP  CLRF     INDF0         ; Addr(FSR) = 0
    CPFSEQ   FSR0          ; FSR0 = END_RAM+1?
    GOTO     LP            ; NO, clear next
    :                      ; YES, All RAM is
    :                      ; cleared
```

6.5 Table Pointer (TBLPTRL and TBLPTRH)

File registers TBLPTRL and TBLPTRH form a 16-bit pointer to address the 64K program memory space. The table pointer is used by instructions TABLWT and TABLRD.

The TABLRD and the TABLWT instructions allow transfer of data between program and data space. The table pointer serves as the 16-bit address of the data word within the program memory. For a more complete description of these registers and the operation of Table Reads and Table Writes, see Section 7.0.

6.6 Table Latch (TBLATH, TBLATL)

The table latch (TBLAT) is a 16-bit register, with TBLATH and TBLATL referring to the high and low bytes of the register. It is not mapped into data or program memory. The table latch is used as a temporary holding latch during data transfer between program and data memory (see descriptions of instructions TABLRD, TABLWT, TLRD and TLWT). For a more complete description of these registers and the operation of Table Reads and Table Writes, see Section 7.0.

6.7 Program Counter Module

The Program Counter (PC) is a 16-bit register. PCL, the low byte of the PC, is mapped in the data memory. PCL is readable and writable just as is any other register. PCH is the high byte of the PC and is not directly addressable. Since PCH is not mapped in data or program memory, an 8-bit register PCLATH (PC high latch) is used as a holding latch for the high byte of the PC. PCLATH is mapped into data memory. The user can read or write PCH through PCLATH.

The 16-bit wide PC is incremented after each instruction fetch during Q1 unless:

- Modified by GOTO, CALL, LCALL, RETURN, RETLW, or RETFIE instruction
- Modified by an interrupt response
- Due to destination write to PCL by an instruction

"Skips" are equivalent to a forced NOP cycle at the skipped address.

Figure 6-11 and Figure 6-12 show the operation of the program counter for various situations.

FIGURE 6-11: PROGRAM COUNTER OPERATION

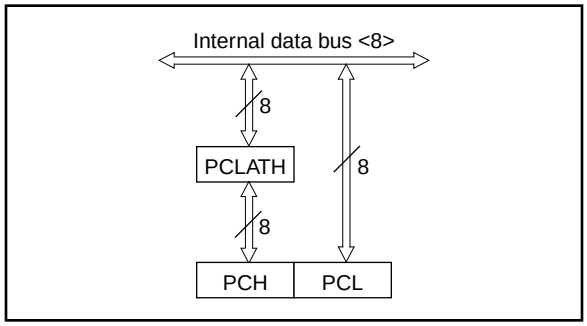
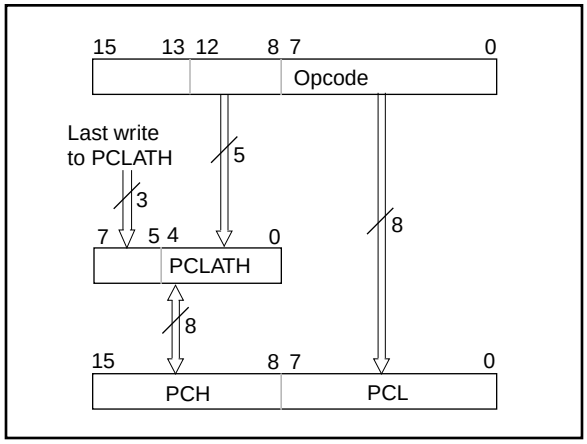


FIGURE 6-12: PROGRAM COUNTER USING THE CALL AND GOTO INSTRUCTIONS



Using Figure 6-11, the operations of the PC and PCLATH for different instructions are as follows:

a) LCALL instructions:
   An 8-bit destination address is provided in the instruction (opcode). PCLATH is unchanged.
   PCLATH → PCH
   Opcode<7:0> → PCL

b) Read instructions on PCL:
   Any instruction that reads PCL.
   PCL → data bus → ALU or destination
   PCH → PCLATH

c) Write instructions on PCL:
   Any instruction that writes to PCL.
   8-bit data → data bus → PCL
   PCLATH → PCH

d) Read-Modify-Write instructions on PCL:
   Any instruction that does a read-write-modify operation on PCL, such as ADDWF PCL.
   Read: PCL → data bus → ALU
   Write: 8-bit result → data bus → PCL
   PCLATH → PCH

e) RETURN instruction:
   PCH → PCLATH
   Stack<MRU> → PC<15:0>

Using Figure 6-12, the operation of the PC and PCLATH for GOTO and CALL instructions is a follows:

CALL, GOTO instructions:
A 13-bit destination address is provided in the instruction (opcode).
Opcode<12:0> → PC <12:0>
PC<15:13> → PCLATH<7:5>
Opcode<12:8> → PCLATH <4:0>

The read-modify-write only affects the PCL with the result. PCH is loaded with the value in the PCLATH. For example, ADDWF PCL will result in a jump within the current page. If PC = 03F0h, WREG = 30h and PCLATH = 03h before instruction, PC = 0320h after the instruction. To accomplish a true 16-bit computed jump, the user needs to compute the 16-bit destination address, write the high byte to PCLATH and then write the low value to PCL.

The following PC related operations do not change PCLATH:

a) LCALL, RETLW, and RETFIE instructions.
b) Interrupt vector is forced onto the PC.
c) Read-modify-write instructions on PCL (e.g. BSF PCL).

Example 8-3 shows the sequence to do a 16 x 16 unsigned multiply. Equation 8-1 shows the algorithm that is used. The 32-bit result is stored in 4 registers RES3:RES0.

EQUATION 8-1: 16 x 16 UNSIGNED MULTIPLICATION ALGORITHM

$$
\begin{aligned}
RES3:RES0 &= ARG1H:ARG1L * ARG2H:ARG2L \\
&= (ARG1H * ARG2H * 2^{16}) + \\
&\quad (ARG1H * ARG2L * 2^{8}) + \\
&\quad (ARG1L * ARG2H * 2^{8}) + \\
&\quad (ARG1L * ARG2L)
\end{aligned}
$$

EXAMPLE 8-3: 16 x 16 MULTIPLY ROUTINE

```
   MOVFP    ARG1L, WREG
   MULWF    ARG2L       ; ARG1L * ARG2L ->
                        ;   PRODH:PRODL
   MOVPF    PRODH, RES1 ;
   MOVPF    PRODL, RES0 ;
;
   MOVFP    ARG1H, WREG
   MULWF    ARG2H       ; ARG1H * ARG2H ->
                        ;   PRODH:PRODL
   MOVPF    PRODH, RES3 ;
   MOVPF    PRODL, RES2 ;
;
   MOVFP    ARG1L, WREG
   MULWF    ARG2H       ; ARG1L * ARG2H ->
                        ;   PRODH:PRODL
   MOVFP    PRODL, WREG ;
   ADDWF    RES1, F     ; Add cross
   MOVFP    PRODH, WREG ;    products
   ADDWFC   RES2, F     ;
   CLRF     WREG, F     ;
   ADDWFC   RES3, F     ;
;
   MOVFP    ARG1H, WREG ;
   MULWF    ARG2L       ; ARG1H * ARG2L ->
                        ;   PRODH:PRODL

   MOVFP    PRODL, WREG ;
   ADDWF    RES1, F     ; Add cross
   MOVFP    PRODH, WREG ;    products
   ADDWFC   RES2, F     ;
   CLRF     WREG, F     ;
   ADDWFC   RES3, F     ;
```

NOTES:

12.1.3 USING PULSE WIDTH MODULATION (PWM) OUTPUTS WITH TMR1 AND TMR2

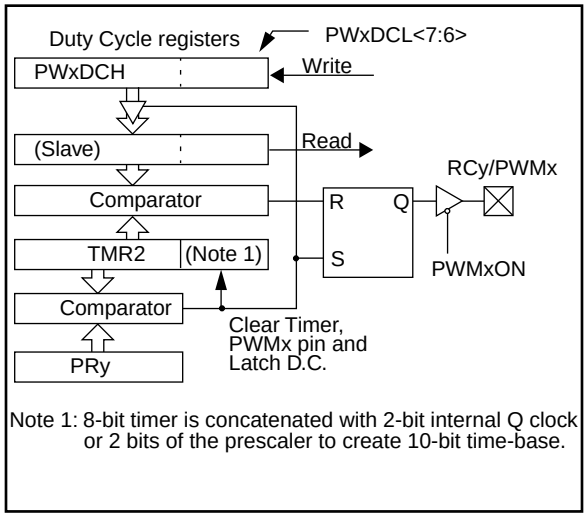
Two high speed pulse width modulation (PWM) outputs are provided. The PWM1 output uses Timer1 as its time-base, while PWM2 may be software configured to use either Timer1 or Timer2 as the time-base. The PWM outputs are on the RB2/PWM1 and RB3/PWM2 pins.

Each PWM output has a maximum resolution of 10-bits. At 10-bit resolution, the PWM output frequency is 24.4 kHz (@ 25 MHz clock) and at 8-bit resolution the PWM output frequency is 97.7 kHz. The duty cycle of the output can vary from 0% to 100%.

Figure 12-5 shows a simplified block diagram of the PWM module. The duty cycle register is double buffered for glitch free operation. Figure 12-6 shows how a glitch could occur if the duty cycle registers were not double buffered.

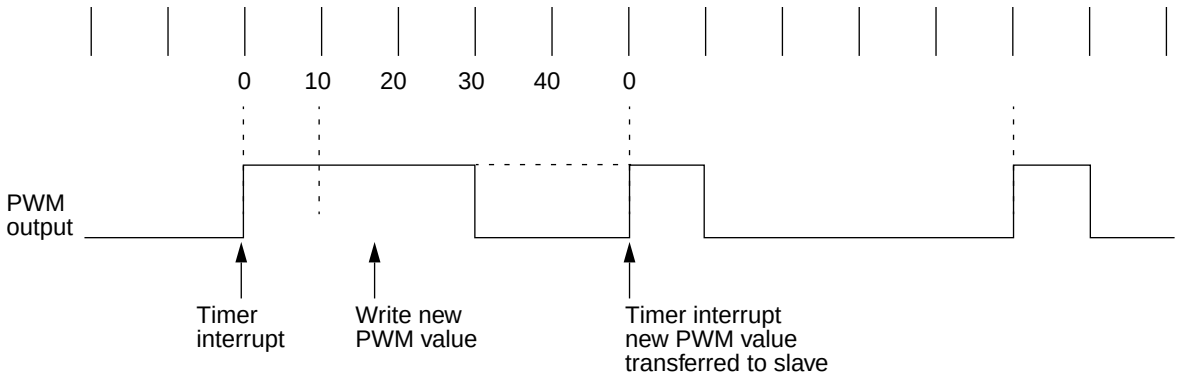
The user needs to set the PWM1ON bit (TCON2<4>) to enable the PWM1 output. When the PWM1ON bit is set, the RB2/PWM1 pin is configured as PWM1 output and forced as an output irrespective of the data direction bit (DDRB<2>). When the PWM1ON bit is clear, the pin behaves as a port pin and its direction is controlled by its data direction bit (DDRB<2>). Similarly, the PWM2ON (TCON2<5>) bit controls the configuration of the RB3/PWM2 pin.

FIGURE 12-5: SIMPLIFIED PWM BLOCK DIAGRAM



Note 1: 8-bit timer is concatenated with 2-bit internal Q clock or 2 bits of the prescaler to create 10-bit time-base.

FIGURE 12-6: PWM OUTPUT



Note The dotted line shows PWM output if duty cycle registers were not double buffered. If the new duty cycle is written after the timer has passed that value, then the PWM does not reset at all during the current cycle causing a “glitch”.

In this example, PWM period = 50. Old duty cycle is 30. New duty cycle value is 10.

12.2.3 EXTERNAL CLOCK INPUT FOR TIMER3

When TMR3CS is set, the 16-bit TMR3 increments on the falling edge of clock input TCLK3. The input on the RB5/TCLK3 pin is sampled and synchronized by the internal phase clocks twice every instruction cycle. This causes a delay from the time a falling edge appears on TCLK3 to the time TMR3 is actually incremented. For the external clock input timing requirements, see the Electrical Specification section. Figure 12-9 shows the timing diagram when operating from an external clock.

12.2.4 READING/WRITING TIMER3

Since Timer3 is a 16-bit timer and only 8-bits at a time can be read or written, care should be taken when reading or writing while the timer is running. The best method to read or write the timer is to stop the timer, perform any read or write operation, and then restart Timer3 (using the TMR3ON bit). However, if it is necessary to keep Timer3 free-running, care must be taken. For writing to the 16-bit TMR3, Example 12-2 may be used. For reading the 16-bit TMR3, Example 12-3 may be used. Interrupts must be disabled during this routine.

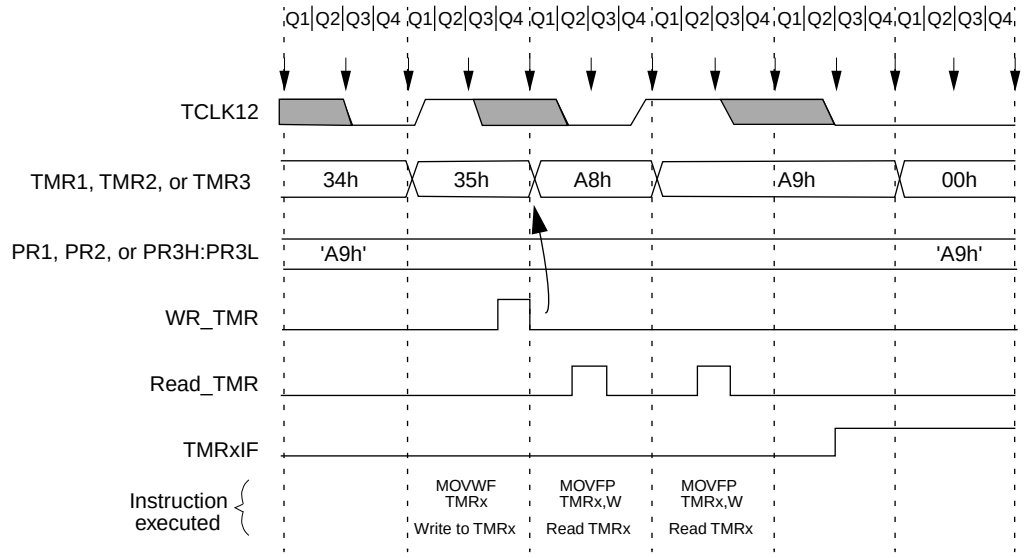
EXAMPLE 12-2: WRITING TO TMR3

```
BSF   CPUSTA, GLINTD ;Disable interrupt
MOVFP RAM_L,  TMR3L  ;
MOVFP RAM_H,  TMR3H  ;
BCF   CPUSTA, GLINTD ;Done,enable interrupt
```

EXAMPLE 12-3: READING FROM TMR3

```
MOVPF   TMR3L, TMPLO    ;read low tmr0
MOVPF   TMR3H, TMPHI    ;read high tmr0
MOVFP   TMPLO, WREG     ;tmplo -> wreg
CPFSLT  TMR3L, WREG     ;tmr0l < wreg?
RETURN                  ;no then return
MOVPF   TMR3L, TMPLO    ;read low tmr0
MOVPF   TMR3H, TMPHI    ;read high tmr0
RETURN                  ;return
```

FIGURE 12-9: TMR1, TMR2, AND TMR3 OPERATION IN EXTERNAL CLOCK MODE



Note 1: TCLK12 is sampled in Q2 and Q4.
2: ↓ indicates a sampling point.
3: The latency from TCLK12 ↓ to timer increment is between 2Tosc and 6Tosc.

NOTES:

TABLE 13-4: BAUD RATES FOR ASYNCHRONOUS MODE

BAUD RATE (K)	FOSC = 33 MHz			FOSC = 25 MHz			FOSC = 20 MHz			FOSC = 16 MHz		
	KBAUD	%ERROR	SPBRG value (decimal)	KBAUD	%ERROR	SPBRG value (decimal)	KBAUD	%ERROR	SPBRG value (decimal)	KBAUD	%ERROR	SPBRG value (decimal)
0.3	NA	—	—	NA	—	—	NA	—	—	NA	—	—
1.2	NA	—	—	NA	—	—	1.221	+1.73	255	1.202	+0.16	207
2.4	2.398	-0.07	214	2.396	0.14	162	2.404	+0.16	129	2.404	+0.16	103
9.6	9.548	-0.54	53	9.53	-0.76	40	9.469	-1.36	32	9.615	+0.16	25
19.2	19.09	-0.54	26	19.53	+1.73	19	19.53	+1.73	15	19.23	+0.16	12
76.8	73.66	-4.09	6	78.13	+1.73	4	78.13	+1.73	3	83.33	+8.51	2
96	103.12	+7.42	4	97.65	+1.73	3	104.2	+8.51	2	NA	—	—
300	257.81	-14.06	1	390.63	+30.21	0	312.5	+4.17	0	NA	—	—
500	515.62	+3.13	0	NA	—	—	NA	—	—	NA	—	—
HIGH	515.62	—	0	—	—	0	312.5	—	0	250	—	0
LOW	2.014	—	255	1.53	—	255	1.221	—	255	0.977	—	255

BAUD RATE (K)	FOSC = 10 MHz			FOSC = 7.159 MHz			FOSC = 5.068 MHz		
	KBAUD	%ERROR	SPBRG value (decimal)	KBAUD	%ERROR	SPBRG value (decimal)	KBAUD	%ERROR	SPBRG value (decimal)
0.3	NA	—	—	NA	—	—	0.31	+3.13	255
1.2	1.202	+0.16	129	1.203	_0.23	92	1.2	0	65
2.4	2.404	+0.16	64	2.380	-0.83	46	2.4	0	32
9.6	9.766	+1.73	15	9.322	-2.90	11	9.9	-3.13	7
19.2	19.53	+1.73	7	18.64	-2.90	5	19.8	+3.13	3
76.8	78.13	+1.73	1	NA	—	—	79.2	+3.13	0
96	NA	—	—	NA	—	—	NA	—	—
300	NA	—	—	NA	—	—	NA	—	—
500	NA	—	—	NA	—	—	NA	—	—
HIGH	156.3	—	0	111.9	—	0	79.2	—	0
LOW	0.610	—	255	0.437	—	255	0.309	—	255

BAUD RATE (K)	FOSC = 3.579 MHz			FOSC = 1 MHz			FOSC = 32.768 kHz		
	KBAUD	%ERROR	SPBRG value (decimal)	KBAUD	%ERROR	SPBRG value (decimal)	KBAUD	%ERROR	SPBRG value (decimal)
0.3	0.301	+0.23	185	0.300	+0.16	51	0.256	-14.67	1
1.2	1.190	-0.83	46	1.202	+0.16	12	NA	—	—
2.4	2.432	+1.32	22	2.232	-6.99	6	NA	—	—
9.6	9.322	-2.90	5	NA	—	—	NA	—	—
19.2	18.64	-2.90	2	NA	—	—	NA	—	—
76.8	NA	—	—	NA	—	—	NA	—	—
96	NA	—	—	NA	—	—	NA	—	—
300	NA	—	—	NA	—	—	NA	—	—
500	NA	—	—	NA	—	—	NA	—	—
HIGH	55.93	—	0	15.63	—	0	0.512	—	0
LOW	0.218	—	255	0.061	—	255	0.002	—	255

13.2.2 USART ASYNCHRONOUS RECEIVER

The receiver block diagram is shown in Figure 13-4. The data comes in the RA4/RX/DT pin and drives the data recovery block. The data recovery block is actually a high speed shifter operating at 16 times the baud rate, whereas the main receive serial shifter operates at the bit rate or at FOSC.

Once asynchronous mode is selected, reception is enabled by setting bit CREN (RCSTA<4>).

The heart of the receiver is the receive (serial) shift register (RSR). After sampling the stop bit, the received data in the RSR is transferred to the RCREG (if it is empty). If the transfer is complete, the interrupt bit RCIF (PIR<0>) is set. The actual interrupt can be enabled/disabled by setting/clearing the RCIE (PIE<0>) bit. RCIF is a read only bit which is cleared by the hardware. It is cleared when RCREG has been read and is empty. RCREG is a double buffered register; (i.e. it is a two deep FIFO). It is possible for two bytes of data to be received and transferred to the RCREG FIFO and a third byte begin shifting to the RSR. On detection of the stop bit of the third byte, if the RCREG is still full, then the overrun error bit, OERR (RCSTA<1>) will be set. The word in the RSR will be lost. RCREG can be read twice to retrieve the two bytes in the FIFO. The OERR bit has to be cleared in software which is done by resetting the receive logic (CREN is set). If the OERR bit is set, transfers from the RSR to RCREG are inhibited, so it is essential to clear the OERR bit if it is set. The framing error bit FERR (RCSTA<2>) is set if a stop bit is not detected.

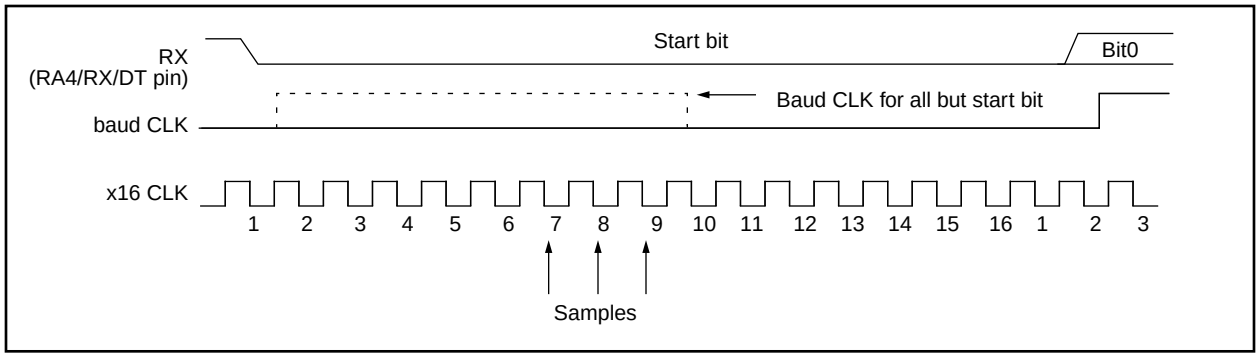
Note:	The FERR and the 9th receive bit are buffered the same way as the receive data. Reading the RCREG register will allow the RX9D and FERR bits to be loaded with values for the next received Received data; therefore, it is essential for the user to read the RCSTA register before reading RCREG in order not to lose the old FERR and RX9D information.

13.2.3 SAMPLING

The data on the RA4/RX/DT pin is sampled three times by a majority detect circuit to determine if a high or a low level is present at the RA4/RX/DT pin. The sampling is done on the seventh, eighth and ninth falling edges of a x16 clock (Figure 11-3).

The x16 clock is a free running clock, and the three sample points occur at a frequency of every 16 falling edges.

FIGURE 13-7: RX PIN SAMPLING SCHEME



DECF **Decrement f**

Syntax: [*label*] DECF f,d

Operands: $0 \le f \le 255$
$d \in [0,1]$

Operation: (f) – 1 → (dest)

Status Affected: OV, C, DC, Z

Encoding:

0000	011d	ffff	ffff

Description: Decrement register 'f'. If 'd' is 0 the result is stored in WREG. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Execute	Write to destination

Example:
```
DECF    CNT,   1
```
Before Instruction
CNT = 0x01
Z = 0

After Instruction
CNT = 0x00
Z = 1

DECFSZ **Decrement f, skip if 0**

Syntax: [*label*] DECFSZ f,d

Operands: $0 \le f \le 255$
$d \in [0,1]$

Operation: (f) – 1 → (dest);
skip if result = 0

Status Affected: None

Encoding:

0001	011d	ffff	ffff

Description: The contents of register 'f' are decremented. If 'd' is 0 the result is placed in WREG. If 'd' is 1 the result is placed back in register 'f'.
If the result is 0, the next instruction, which is already fetched, is discarded, and an NOP is executed instead making it a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Execute	Write to destination

Example:
```
HERE    DECFSZ  CNT, 1
        GOTO    LOOP
CONTINUE
```
Before Instruction
PC = Address (HERE)

After Instruction
CNT = CNT - 1
If CNT = 0;
PC = Address (CONTINUE)
If CNT ≠ 0;
PC = Address (HERE+1)

DCFSNZ **Decrement f, skip if not 0**

Syntax: [*label*] DCFSNZ f,d

Operands: $0 \le f \le 255$
$d \in [0,1]$

Operation: (f) – 1 → (dest);
skip if not 0

Status Affected: None

Encoding:

0010	011d	ffff	ffff

Description: The contents of register 'f' are decremented. If 'd' is 0 the result is placed in WREG. If 'd' is 1 the result is placed back in register 'f'.
If the result is not 0, the next instruction, which is already fetched, is discarded, and an NOP is executed instead making it a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Execute	Write to destination

If skip:

Q1	Q2	Q3	Q4
Forced NOP	NOP	Execute	NOP

Example:
```
HERE    DCFSNZ  TEMP, 1
ZERO    :
NZERO   :
```

Before Instruction
TEMP_VALUE = ?

After Instruction
TEMP_VALUE = TEMP_VALUE - 1,
If TEMP_VALUE = 0;
PC = Address (ZERO)
If TEMP_VALUE ≠ 0;
PC = Address (NZERO)

GOTO **Unconditional Branch**

Syntax: [*label*] GOTO k

Operands: $0 \le k \le 8191$

Operation: k → PC<12:0>;
k<12:8> → PCLATH<4:0>,
PC<15:13> → PCLATH<7:5>

Status Affected: None

Encoding:

110k	kkkk	kkkk	kkkk

Description: GOTO allows an unconditional branch anywhere within an 8K page boundary. The thirteen bit immediate value is loaded into PC bits <12:0>. Then the upper eight bits of PC are loaded into PCLATH. GOTO is always a two-cycle instruction.

Words: 1

Cycles: 2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'<7:0>	Execute	NOP
Forced NOP	NOP	Execute	NOP

Example:
```
GOTO THERE
```

After Instruction
PC = Address (THERE)

TABLRD **Table Read**

Example1: `TABLRD 1, 1, REG ;`

Before Instruction

REG	=	0x53
TBLATH	=	0xAA
TBLATL	=	0x55
TBLPTR	=	0xA356
MEMORY(TBLPTR)	=	0x1234

After Instruction (table write completion)

REG	=	0xAA
TBLATH	=	0x12
TBLATL	=	0x34
TBLPTR	=	0xA357
MEMORY(TBLPTR)	=	0x5678

Example2: `TABLRD 0, 0, REG ;`

Before Instruction

REG	=	0x53
TBLATH	=	0xAA
TBLATL	=	0x55
TBLPTR	=	0xA356
MEMORY(TBLPTR)	=	0x1234

After Instruction (table write completion)

REG	=	0x55
TBLATH	=	0x12
TBLATL	=	0x34
TBLPTR	=	0xA356
MEMORY(TBLPTR)	=	0x1234

TABLWT **Table Write**

Syntax: [*label*] TABLWT t,i,f

Operands: $0 \le f \le 255$
$i \in [0,1]$
$t \in [0,1]$

Operation: If t = 0,
f → TBLATL;
If t = 1,
f → TBLATH;
TBLAT → Prog Mem (TBLPTR);
If i = 1,
TBLPTR + 1 → TBLPTR

Status Affected: None

Encoding:

1010	11ti	ffff	ffff

Description:

1. Load value in 'f' into 16-bit table latch (TBLAT)
   If t = 0: load into low byte;
   If t = 1: load into high byte
2. The contents of TBLAT is written to the program memory location pointed to by TBLPTR
   If TBLPTR points to external program memory location, then the instruction takes two-cycle
   If TBLPTR points to an internal EPROM location, then the instruction is terminated when an interrupt is received.

Note: The $\overline{MCLR}/V_{PP}$ pin must be at the programming voltage for successful programming of internal memory.
If $\overline{MCLR}/V_{PP} = V_{DD}$
the programming sequence of internal memory will be executed, but will not be successful (although the internal memory location may be disturbed)

3. The TBLPTR can be automatically incremented
   If i = 0; TBLPTR is not incremented
   If i = 1; TBLPTR is incremented

Words: 1

Cycles: 2 (many if write is to on-chip EPROM program memory)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Execute	Write register TBLATH or TBLATL

Applicable Devices	42	R42	42A	43	R43	44

FIGURE 17-7: CAPTURE TIMINGS

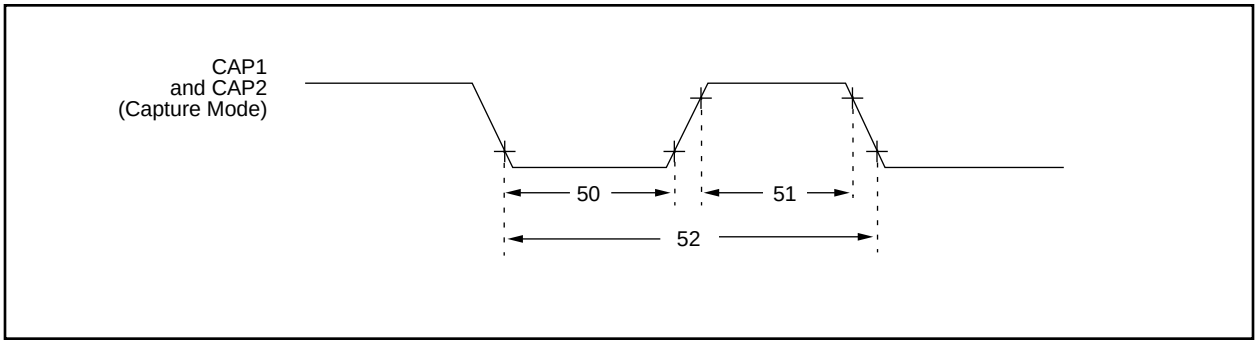


TABLE 17-7: CAPTURE REQUIREMENTS

Parameter No.	Sym	Characteristic	Min	Typ†	Max	Units	Conditions
50	TccL	Capture1 and Capture2 input low time	10 *	—	—	ns	
51	TccH	Capture1 and Capture2 input high time	10 *	—	—	ns	
52	TccP	Capture1 and Capture2 input period	$\frac{2\ T_{CY}}{N}$ §	—	—	ns	N = prescale value (4 or 16)

* These parameters are characterized but not tested.

† Data in "Typ" column is at 5V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

§ This specification ensured by design.

FIGURE 17-8: PWM TIMINGS

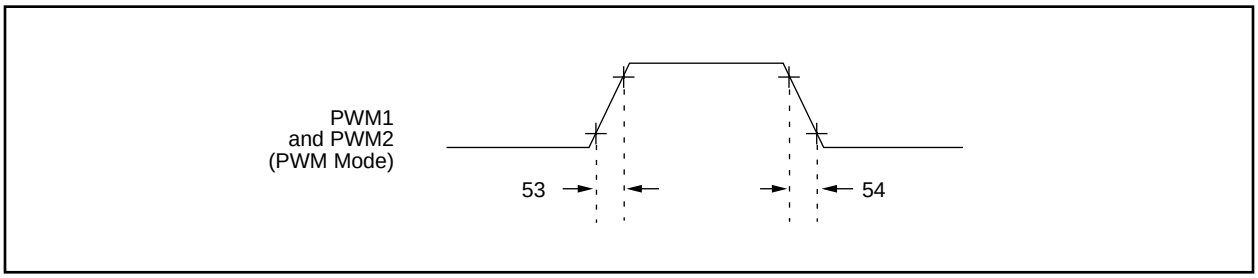


TABLE 17-8: PWM REQUIREMENTS

Parameter No.	Sym	Characteristic	Min	Typ†	Max	Units	Conditions
53	TccR	PWM1 and PWM2 output rise time	—	10 *	35 *§	ns	
54	TccF	PWM1 and PWM2 output fall time	—	10 *	35 *§	ns	

* These parameters are characterized but not tested.

† Data in "Typ" column is at 5V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

§ This specification ensured by design.

Applicable Devices	42	R42	42A	43	R43	44

FIGURE 17-11: MEMORY INTERFACE WRITE TIMING

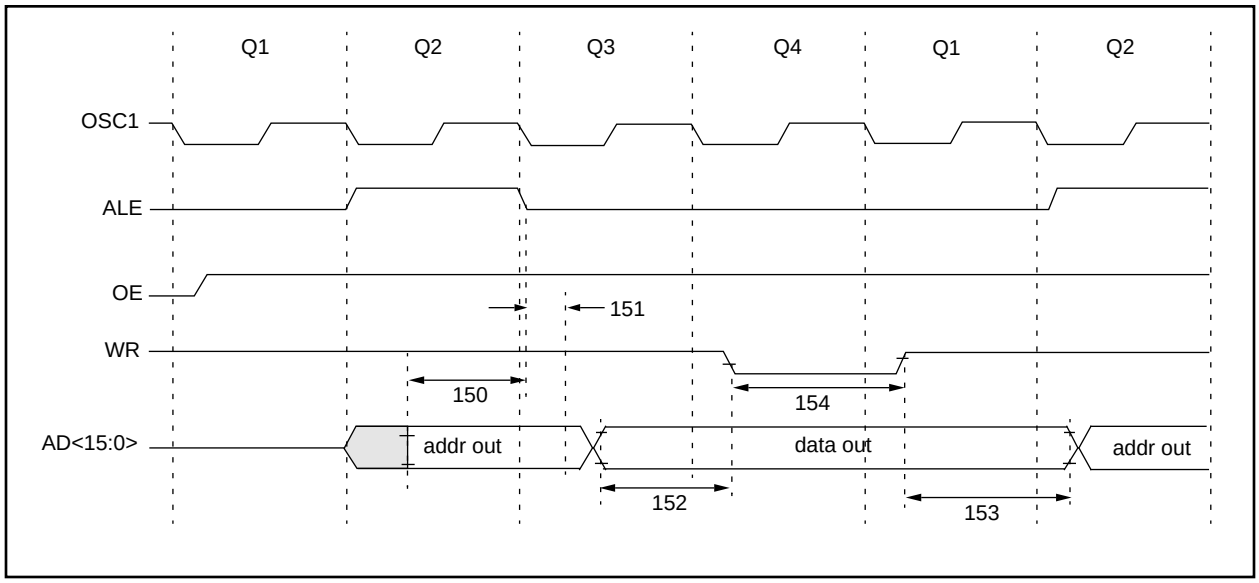


TABLE 17-11: MEMORY INTERFACE WRITE REQUIREMENTS

Parameter No.	Sym	Characteristic	Min	Typ†	Max	Units	Conditions
150	TadV2alL	AD<15:0> (address) valid to ALE↓ (address setup time)	0.25Tcy - 30	—	—	ns	
151	TalL2adI	ALE↓ to address out invalid (address hold time)	0	—	—	ns	
152	TadV2wrL	Data out valid to $\overline{WR}$↓ (data setup time)	0.25Tcy - 40	—	—	ns	
153	TwrH2adI	$\overline{WR}$↑ to data out invalid (data hold time)	—	0.25TCY §	—	ns	
154	TwrL	$\overline{WR}$ pulse width	—	0.25TCY §	—	ns	

* These parameters are characterized but not tested.

† Data in "Typ" column is at 5V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

§ This specification is guaranteed by design.

Applicable Devices	42	R42	42A	43	R43	44

FIGURE 18-7: TYPICAL IDD vs. FREQUENCY (EXTERNAL CLOCK 25°C)

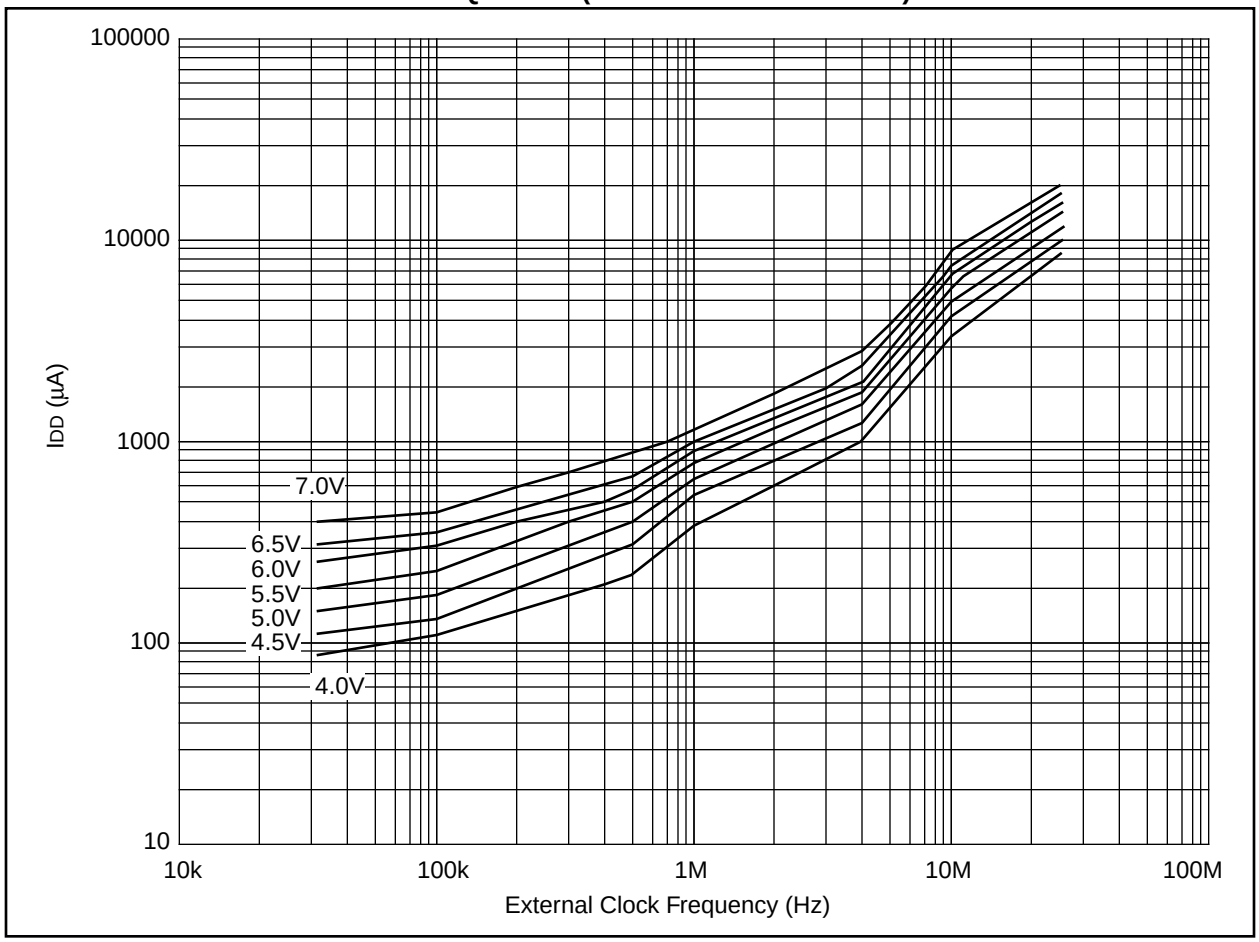
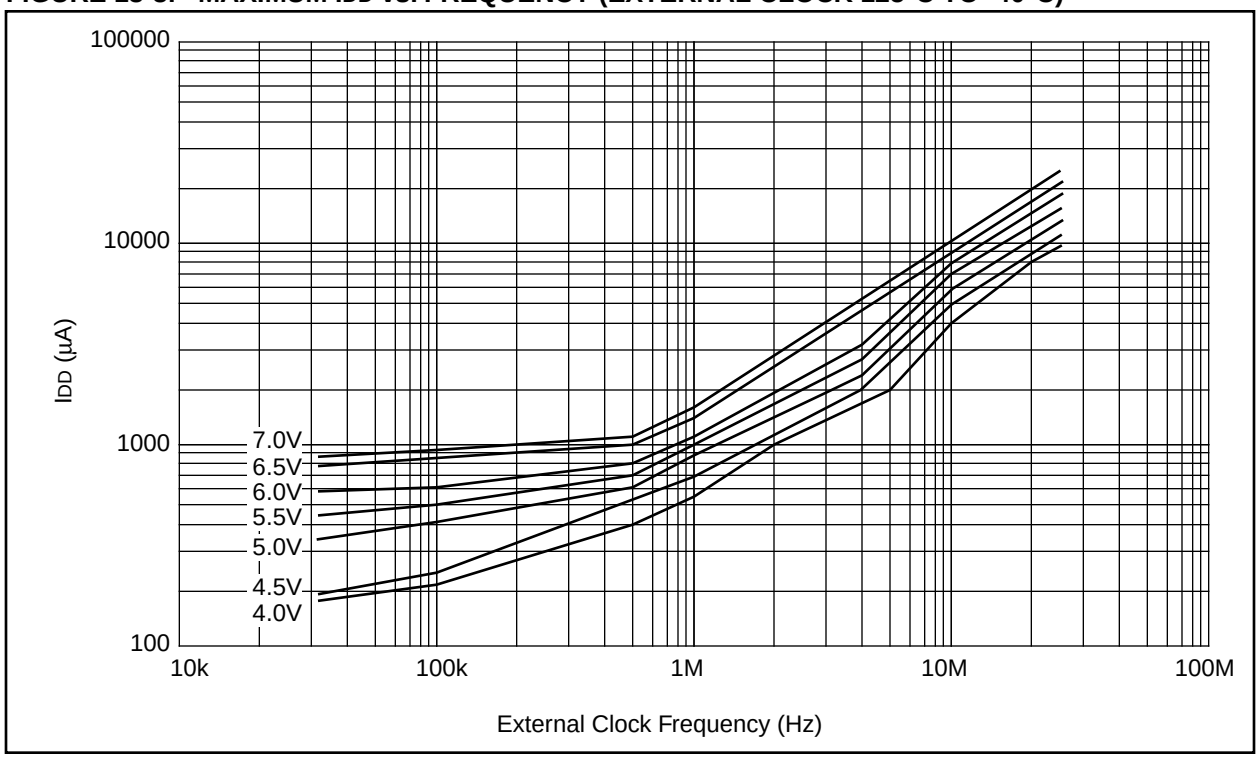


FIGURE 18-8: MAXIMUM IDD vs. FREQUENCY (EXTERNAL CLOCK 125°C TO -40°C)



Applicable Devices	42	R42	42A	43	R43	44

DC CHARACTERISTICS

Standard Operating Conditions (unless otherwise stated)
Operating temperature
-40°C ≤ TA ≤ +85°C for industrial and
0°C ≤ TA ≤ +70°C for commercial
Operating voltage VDD range as described in Section 19.1

Parameter No.	Sym	Characteristic	Min	Typ†	Max	Units	Conditions
		Output Low Voltage					
D080	VOL	I/O ports (except RA2 and RA3)					IOL = VDD/1.250 mA
			–	–	0.1VDD	V	4.5V ≤ VDD ≤ 6.0V
			–	–	0.1VDD *	V	VDD = 2.5V
D081		with TTL buffer	–	–	0.4	V	IOL = 6 mA, VDD = 4.5V Note 6
D082		RA2 and RA3	–	–	3.0	V	IOL = 60.0 mA, VDD = 6.0V
D083		OSC2/CLKOUT	–	–	0.4	V	IOL = 1 mA, VDD = 4.5V
D084		(RC and EC osc modes)	–	–	0.1VDD *	V	IOL = VDD/5 mA (PIC17LC43/LC44 only)
		Output High Voltage (Note 3)					
D090	VOH	I/O ports (except RA2 and RA3)					IOH = -VDD/2.500 mA
			0.9VDD	–	–	V	4.5V ≤ VDD ≤ 6.0V
			0.9VDD *	–	–	V	VDD = 2.5V
D091		with TTL buffer	2.4	–	–	V	IOH = -6.0 mA, VDD=4.5V Note 6
D092		RA2 and RA3	–	–	12	V	Pulled-up to externally applied voltage
D093		OSC2/CLKOUT	2.4	–	–	V	IOH = -5 mA, VDD = 4.5V
D094		(RC and EC osc modes)	0.9VDD *	–	–	V	IOH = -VDD/5 mA (PIC17LC43/LC44 only)
		Capacitive Loading Specs on Output Pins					
D100	COSC2	OSC2/CLKOUT pin	–	–	25	pF	In EC or RC osc modes when OSC2 pin is outputting CLKOUT. external clock is used to drive OSC1.
D101	CIO	All I/O pins and OSC2 (in RC mode)	–	–	50	pF	
D102	CAD	System Interface Bus (PORTC, PORTD and PORTE)	–	–	50	pF	In Microprocessor or Extended Microcontroller mode

* These parameters are characterized but not tested.

† Data in "Typ" column is at 5V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

‡ These parameters are for design guidance only and are not tested, nor characterized.

Note 1: In RC oscillator configuration, the OSC1/CLKIN pin is a Schmitt Trigger input. It is not recommended that the PIC17CXX devices be driven with external clock in RC mode.

2: The leakage current on the $\overline{MCLR}$ pin is strongly dependent on the applied voltage level. The specified levels represent normal operating conditions. Higher leakage current may be measured at different input voltages.

3: Negative current is defined as coming out of the pin.

4: These specifications are for the programming of the on-chip program memory EPROM through the use of the table write instructions. The complete programming specifications can be found in: PIC17CXX Programming Specifications (Literature number DS30139).

5: The $\overline{MCLR}$/VPP pin may be kept in this range at times other than programming, but is not recommended.

6: For TTL buffers, the better of the two specifications may be used.

Applicable Devices	42	R42	42A	43	R43	44

FIGURE 19-11: MEMORY INTERFACE WRITE TIMING (NOT SUPPORTED IN PIC17LC4X DEVICES)

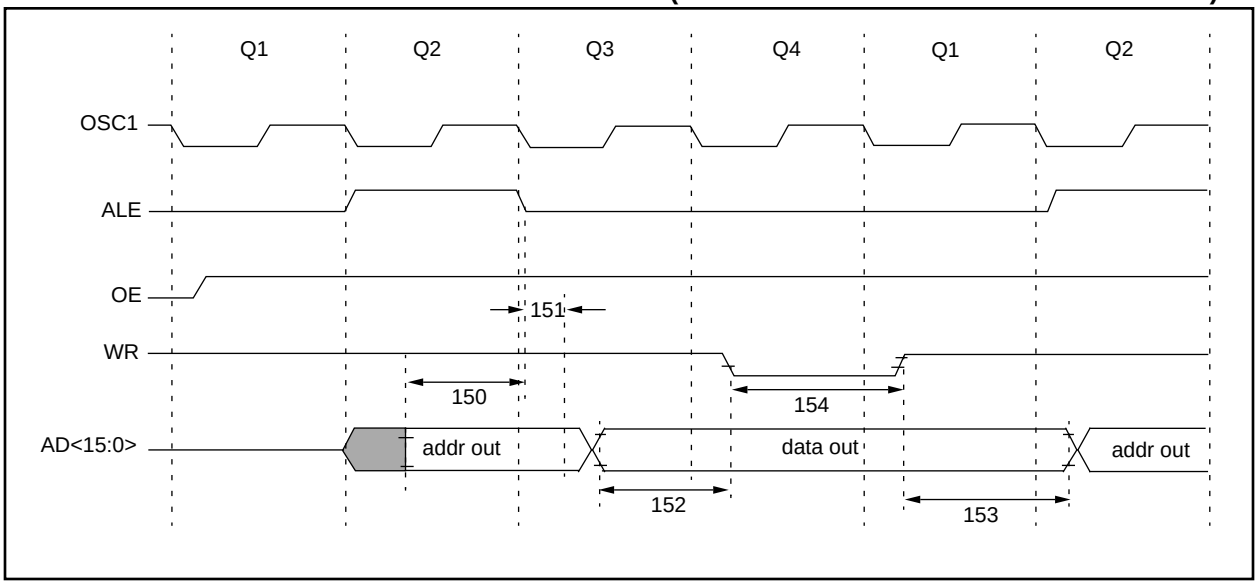


TABLE 19-11: MEMORY INTERFACE WRITE REQUIREMENTS (NOT SUPPORTED IN PIC17LC4X DEVICES)

Parameter No.	Sym	Characteristic	Min	Typ†	Max	Units	Conditions
150	TadV2alL	AD<15:0> (address) valid to ALE↓ (address setup time)	0.25Tcy - 10	—	—	ns	
151	TalL2adI	ALE↓ to address out invalid (address hold time)	0	—	—	ns	
152	TadV2wrL	Data out valid to $\overline{WR}$↓ (data setup time)	0.25Tcy - 40	—	—	ns	
153	TwrH2adI	$\overline{WR}$↑ to data out invalid (data hold time)	—	0.25TCY §	—	ns	
154	TwrL	$\overline{WR}$ pulse width	—	0.25TCY §	—	ns	

* These parameters are characterized but not tested.

† Data in "Typ" column is at 5V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

§ This specification ensured by design.

WORLDWIDE SALES AND SERVICE

AMERICAS

Corporate Office
2355 West Chandler Blvd.
Chandler, AZ 85224-6199
Tel: 480-792-7200 Fax: 480-792-7277
Technical Support: 480-792-7627
Web Address: http://www.microchip.com

Rocky Mountain
2355 West Chandler Blvd.
Chandler, AZ 85224-6199
Tel: 480-792-7966 Fax: 480-792-7456

Atlanta
500 Sugar Mill Road, Suite 200B
Atlanta, GA 30350
Tel: 770-640-0034 Fax: 770-640-0307

Boston
2 Lan Drive, Suite 120
Westford, MA 01886
Tel: 978-692-3848 Fax: 978-692-3821

Chicago
333 Pierce Road, Suite 180
Itasca, IL 60143
Tel: 630-285-0071 Fax: 630-285-0075

Dallas
4570 Westgrove Drive, Suite 160
Addison, TX 75001
Tel: 972-818-7423 Fax: 972-818-2924

Detroit
Tri-Atria Office Building
32255 Northwestern Highway, Suite 190
Farmington Hills, MI 48334
Tel: 248-538-2250 Fax: 248-538-2260

Kokomo
2767 S. Albright Road
Kokomo, Indiana 46902
Tel: 765-864-8360 Fax: 765-864-8387

Los Angeles
18201 Von Karman, Suite 1090
Irvine, CA 92612
Tel: 949-263-1888 Fax: 949-263-1338

New York
150 Motor Parkway, Suite 202
Hauppauge, NY 11788
Tel: 631-273-5305 Fax: 631-273-5335

San Jose
Microchip Technology Inc.
2107 North First Street, Suite 590
San Jose, CA 95131
Tel: 408-436-7950 Fax: 408-436-7955

Toronto
6285 Northam Drive, Suite 108
Mississauga, Ontario L4V 1X5, Canada
Tel: 905-673-0699 Fax: 905-673-6509

ASIA/PACIFIC

Australia
Microchip Technology Australia Pty Ltd
Suite 22, 41 Rawson Street
Epping 2121, NSW
Australia
Tel: 61-2-9868-6733 Fax: 61-2-9868-6755

China - Beijing
Microchip Technology Consulting (Shanghai) Co., Ltd., Beijing Liaison Office
Unit 915
Bei Hai Wan Tai Bldg.
No. 6 Chaoyangmen Beidajie
Beijing, 100027, No. China
Tel: 86-10-85282100 Fax: 86-10-85282104

China - Chengdu
Microchip Technology Consulting (Shanghai) Co., Ltd., Chengdu Liaison Office
Rm. 2401, 24th Floor,
Ming Xing Financial Tower
No. 88 TIDU Street
Chengdu 610016, China
Tel: 86-28-6766200 Fax: 86-28-6766599

China - Fuzhou
Microchip Technology Consulting (Shanghai) Co., Ltd., Fuzhou Liaison Office
Unit 28F, World Trade Plaza
No. 71 Wusi Road
Fuzhou 350001, China
Tel: 86-591-7503506 Fax: 86-591-7503521

China - Shanghai
Microchip Technology Consulting (Shanghai) Co., Ltd.
Room 701, Bldg. B
Far East International Plaza
No. 317 Xian Xia Road
Shanghai, 200051
Tel: 86-21-6275-5700 Fax: 86-21-6275-5060

China - Shenzhen
Microchip Technology Consulting (Shanghai) Co., Ltd., Shenzhen Liaison Office
Rm. 1315, 13/F, Shenzhen Kerry Centre,
Renminnan Lu
Shenzhen 518001, China
Tel: 86-755-2350361 Fax: 86-755-2366086

Hong Kong
Microchip Technology Hongkong Ltd.
Unit 901-6, Tower 2, Metroplaza
223 Hing Fong Road
Kwai Fong, N.T., Hong Kong
Tel: 852-2401-1200 Fax: 852-2401-3431

India
Microchip Technology Inc.
India Liaison Office
Divyasree Chambers
1 Floor, Wing A (A3/A4)
No. 11, O'Shaugnessey Road
Bangalore, 560 025, India
Tel: 91-80-2290061 Fax: 91-80-2290062

Japan
Microchip Technology Japan K.K.
Benex S-1 6F
3-18-20, Shinyokohama
Kohoku-Ku, Yokohama-shi
Kanagawa, 222-0033, Japan
Tel: 81-45-471- 6166 Fax: 81-45-471-6122

Korea
Microchip Technology Korea
168-1, Youngbo Bldg. 3 Floor
Samsung-Dong, Kangnam-Ku
Seoul, Korea 135-882
Tel: 82-2-554-7200 Fax: 82-2-558-5934

Singapore
Microchip Technology Singapore Pte Ltd.
200 Middle Road
#07-02 Prime Centre
Singapore, 188980
Tel: 65-334-8870 Fax: 65-334-8850

Taiwan
Microchip Technology Taiwan
11F-3, No. 207
Tung Hua North Road
Taipei, 105, Taiwan
Tel: 886-2-2717-7175 Fax: 886-2-2545-0139

EUROPE

Denmark
Microchip Technology Nordic ApS
Regus Business Centre
Lautrup hoj 1-3
Ballerup DK-2750 Denmark
Tel: 45 4420 9895 Fax: 45 4420 9910

France
Microchip Technology SARL
Parc d'Activite du Moulin de Massy
43 Rue du Saule Trapu
Batiment A - ler Etage
91300 Massy, France
Tel: 33-1-69-53-63-20 Fax: 33-1-69-30-90-79

Germany
Microchip Technology GmbH
Gustav-Heinemann Ring 125
D-81739 Munich, Germany
Tel: 49-89-627-144 0 Fax: 49-89-627-144-44

Italy
Microchip Technology SRL
Centro Direzionale Colleoni
Palazzo Taurus 1 V. Le Colleoni 1
20041 Agrate Brianza
Milan, Italy
Tel: 39-039-65791-1 Fax: 39-039-6899883

United Kingdom
Arizona Microchip Technology Ltd.
505 Eskdale Road
Winnersh Triangle
Wokingham
Berkshire, England RG41 5TU
Tel: 44 118 921 5869 Fax: 44-118 921-5820

01/18/02