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Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	40MHz
Connectivity	EBI/EMI, I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, LVD, POR, PWM, WDT
Number of I/O	68
Program Memory Size	32KB (16K x 16)
Program Memory Type	FLASH
EEPROM Size	1K x 8
RAM Size	2K x 8
Voltage - Supply (Vcc/Vdd)	2V ~ 5.5V
Data Converters	A/D 16x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	80-TQFP
Supplier Device Package	80-TQFP (12x12)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18lf8520-i-pt

4.12 Indirect Addressing, INDF and FSR Registers

Indirect addressing is a mode of addressing data memory, where the data memory address in the instruction is not fixed. An FSR register is used as a pointer to the data memory location that is to be read or written. Since this pointer is in RAM, the contents can be modified by the program. This can be useful for data tables in the data memory and for software stacks. Figure 4-9 shows the operation of indirect addressing. This shows the moving of the value to the data memory address, specified by the value of the FSR register.

Indirect addressing is possible by using one of the INDF registers. Any instruction using the INDF register actually accesses the register pointed to by the File Select Register, FSR. Reading the INDF register itself, indirectly (FSR = 0), will read 00h. Writing to the INDF register indirectly, results in a no operation. The FSR register contains a 12-bit address, which is shown in Figure 4-10.

The INDFn register is not a physical register. Addressing INDFn actually addresses the register whose address is contained in the FSRn register (FSRn is a pointer). This is indirect addressing.

Example 4-4 shows a simple use of indirect addressing to clear the RAM in Bank 1 (locations 100h-1FFh) in a minimum number of instructions.

EXAMPLE 4-4: HOW TO CLEAR RAM (BANK 1) USING INDIRECT ADDRESSING

```
        LFSR  FSR0 ,0x100  ;
NEXT    CLRF  POSTINC0     ; Clear INDF
                           ; register and
                           ; inc pointer
        BTFSS FSR0H, 1     ; All done with
                           ; Bank 1?
        GOTO  NEXT         ; NO, clear next
CONTINUE                   ; YES, continue
```

There are three indirect addressing registers. To address the entire data memory space (4096 bytes), these registers are 12 bits wide. To store the 12 bits of addressing information, two 8-bit registers are required. These indirect addressing registers are:

1. FSR0: composed of FSR0H:FSR0L
2. FSR1: composed of FSR1H:FSR1L
3. FSR2: composed of FSR2H:FSR2L

In addition, there are registers INDF0, INDF1 and INDF2, which are not physically implemented. Reading or writing to these registers activates indirect addressing, with the value in the corresponding FSR register being the address of the data. If an instruction writes a value to INDF0, the value will be written to the address pointed to by FSR0H:FSR0L. A read from INDF1 reads the data from the address pointed to by FSR1H:FSR1L. INDFn can be used in code anywhere an operand can be used.

If INDF0, INDF1 or INDF2 are read indirectly via an FSR, all '0's are read (zero bit is set). Similarly, if INDF0, INDF1 or INDF2 are written to indirectly, the operation will be equivalent to a NOP instruction and the Status bits are not affected.

4.12.1 INDIRECT ADDRESSING OPERATION

Each FSR register has an INDF register associated with it, plus four additional register addresses. Performing an operation on one of these five registers determines how the FSR will be modified during indirect addressing.

When data access is done to one of the five INDFn locations, the address selected will configure the FSRn register to:

- Do nothing to FSRn after an indirect access (no change) – INDFn.
- Auto-decrement FSRn after an indirect access (post-decrement) – POSTDECn.
- Auto-increment FSRn after an indirect access (post-increment) – POSTINCn.
- Auto-increment FSRn before an indirect access (pre-increment) – PREINCn.
- Use the value in the WREG register as an offset to FSRn. Do not modify the value of the WREG or the FSRn register after an indirect access (no change) – PLUSWn.

When using the auto-increment or auto-decrement features, the effect on the FSR is not reflected in the Status register. For example, if the indirect address causes the FSR to equal '0', the Z bit will not be set.

Incrementing or decrementing an FSR affects all 12 bits. That is, when FSRnL overflows from an increment, FSRnH will be incremented automatically.

Adding these features allows the FSRn to be used as a stack pointer, in addition to its uses for table operations in data memory.

Each FSR has an address associated with it that performs an indexed indirect access. When a data access to this INDFn location (PLUSWn) occurs, the FSRn is configured to add the signed value in the WREG register and the value in FSR to form the address before an indirect access. The FSR value is not changed.

If an FSR register contains a value that points to one of the INDFn, an indirect read will read 00h (zero bit is set), while an indirect write will be equivalent to a NOP (Status bits are not affected).

If an indirect addressing operation is done where the target address is an FSRnH or FSRnL register, the write operation will dominate over the pre- or post-increment/decrement functions.

4.13 Status Register

The Status register, shown in Register 4-3, contains the arithmetic status of the ALU. The Status register can be the destination for any instruction, as with any other register. If the Status register is the destination for an instruction that affects the Z, DC, C, OV or N bits, then the write to these five bits is disabled. These bits are set or cleared according to the device logic. Therefore, the result of an instruction with the Status register as destination may be different than intended.

For example, `CLRF STATUS` will clear the upper three bits and set the Z bit. This leaves the Status register as `000u uluu` (where `u` = unchanged).

It is recommended, therefore, that only `BCF`, `BSF`, `SWAPF`, `MOVFF` and `MOVWF` instructions are used to alter the Status register, because these instructions do not affect the Z, C, DC, OV or N bits from the Status register. For other instructions not affecting any status bits, see Table 24-1.

Note: The C and DC bits operate as a $\overline{\text{borrow}}$ and $\overline{\text{digit borrow}}$ bit respectively, in subtraction.

REGISTER 4-3: STATUS REGISTER

U-0	U-0	U-0	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x
—	—	—	N	OV	Z	DC	C
bit 7							bit 0

bit 7-5 **Unimplemented:** Read as '0'

bit 4 **N:** Negative bit
This bit is used for signed arithmetic (2's complement). It indicates whether the result was negative (ALU MSB = 1).
1 = Result was negative
0 = Result was positive

bit 3 **OV:** Overflow bit
This bit is used for signed arithmetic (2's complement). It indicates an overflow of the 7-bit magnitude, which causes the sign bit (bit 7) to change state.
1 = Overflow occurred for signed arithmetic (in this arithmetic operation)
0 = No overflow occurred

bit 2 **Z:** Zero bit
1 = The result of an arithmetic or logic operation is zero
0 = The result of an arithmetic or logic operation is not zero

bit 1 **DC:** Digit carry/$\overline{\text{borrow}}$ bit
For `ADDWF`, `ADDLW`, `SUBLW` and `SUBWF` instructions:
1 = A carry-out from the 4th low-order bit of the result occurred
0 = No carry-out from the 4th low-order bit of the result

Note: For $\overline{\text{borrow}}$, the polarity is reversed. A subtraction is executed by adding the 2's complement of the second operand. For rotate (`RRF`, `RLF`) instructions, this bit is loaded with either bit 4 or bit 3 of the source register.

bit 0 **C:** Carry/$\overline{\text{borrow}}$ bit
For `ADDWF`, `ADDLW`, `SUBLW` and `SUBWF` instructions:
1 = A carry-out from the Most Significant bit of the result occurred
0 = No carry-out from the Most Significant bit of the result occurred

Note: For $\overline{\text{borrow}}$, the polarity is reversed. A subtraction is executed by adding the 2's complement of the second operand. For rotate (`RRF`, `RLF`) instructions, this bit is loaded with either the high or low-order bit of the source register.

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
- n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

6.2.2 16-BIT WORD WRITE MODE

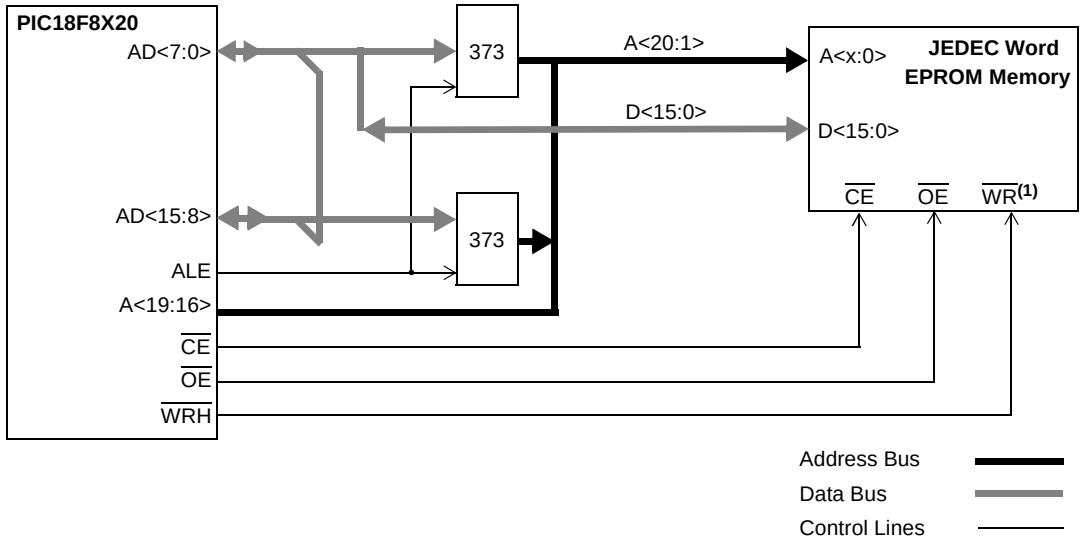
Figure 6-2 shows an example of 16-bit Word Write mode for PIC18F8X20 devices. This mode is used for word-wide memories, which includes some of the EPROM and Flash type memories. This mode allows opcode fetches and table reads from all forms of 16-bit memory and table writes to any type of word-wide external memories. This method makes a distinction between TBLWT cycles to even or odd addresses.

During a TBLWT cycle to an even address (TBLPTR<0> = 0), the TABLAT data is transferred to a holding latch and the external address data bus is tri-stated for the data portion of the bus cycle. No write signals are activated.

During a TBLWT cycle to an odd address (TBLPTR<0> = 1), the TABLAT data is presented on the upper byte of the AD15:AD0 bus. The contents of the holding latch are presented on the lower byte of the AD15:AD0 bus.

The $\overline{WRH}$ signal is strobed for each write cycle; the $\overline{WRL}$ pin is unused. The signal on the BA0 pin indicates the LSb of TBLPTR, but it is left unconnected. Instead, the $\overline{UB}$ and $\overline{LB}$ signals are active to select both bytes. The obvious limitation to this method is that the table write must be done in pairs on a specific word boundary to correctly write a word location.

FIGURE 6-2: 16-BIT WORD WRITE MODE EXAMPLE



Note 1: This signal only applies to table writes. See **Section 5.1 "Table Reads and Table Writes"**.

REGISTER 9-8: PIE2: PERIPHERAL INTERRUPT ENABLE REGISTER 2

U-0	R/W-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
—	CMIE	—	EEIE	BCLIE	LVDIE	TMR3IE	CCP2IE
bit 7							bit 0

bit 7 **Unimplemented:** Read as '0'

bit 6 **CMIE:** Comparator Interrupt Enable bit
1 = Enables the comparator interrupt
0 = Disables the comparator interrupt

bit 5 **Unimplemented:** Read as '0'

bit 4 **EEIE:** Data EEPROM/Flash Write Operation Interrupt Enable bit
1 = Enables the write operation interrupt
0 = Disables the write operation interrupt

bit 3 **BCLIE:** Bus Collision Interrupt Enable bit
1 = Enables the bus collision interrupt
0 = Disables the bus collision interrupt

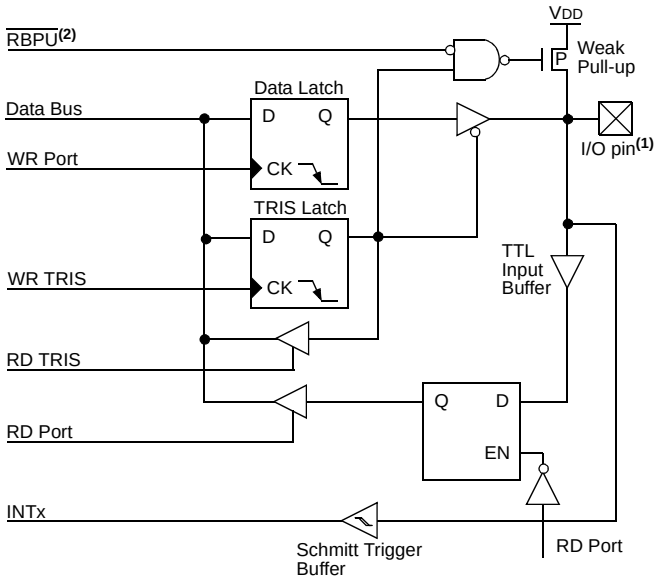
bit 2 **LVDIE:** Low-Voltage Detect Interrupt Enable bit
1 = Enables the Low-Voltage Detect interrupt
0 = Disables the Low-Voltage Detect interrupt

bit 1 **TMR3IE:** TMR3 Overflow Interrupt Enable bit
1 = Enables the TMR3 overflow interrupt
0 = Disables the TMR3 overflow interrupt

bit 0 **CCP2IE:** CCP2 Interrupt Enable bit
1 = Enables the CCP2 interrupt
0 = Disables the CCP2 interrupt

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
- n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

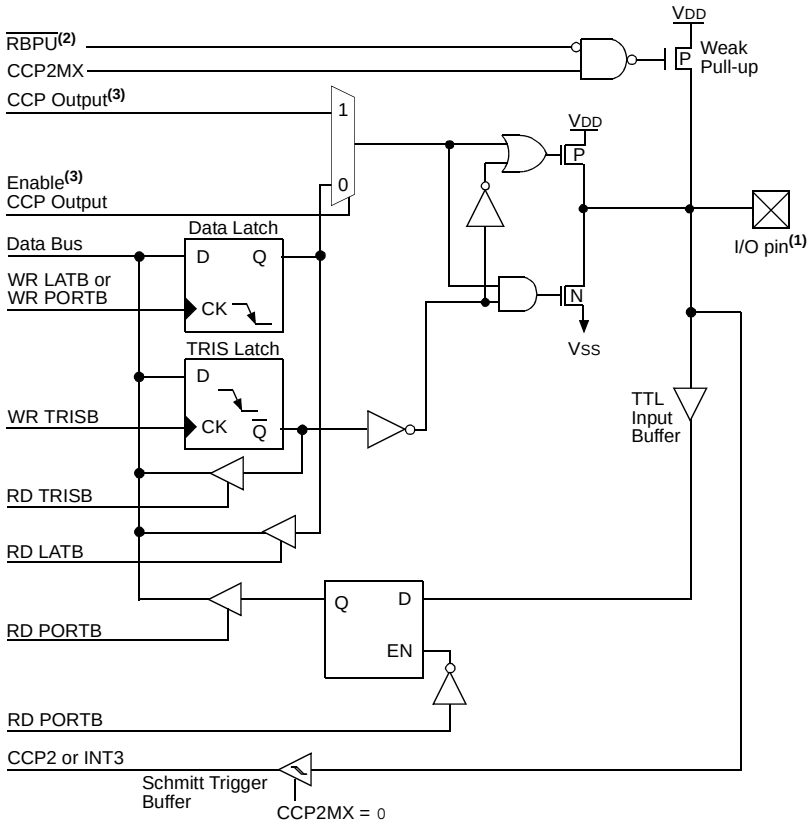
FIGURE 10-6: BLOCK DIAGRAM OF RB2:RB0 PINS



Note 1: I/O pins have diode protection to VDD and VSS.

2: To enable weak pull-ups, set the appropriate TRIS bit(s) and clear the $\overline{\text{RBPU}}$ bit (OPTION_REG<7>).

FIGURE 10-7: BLOCK DIAGRAM OF RB3 PIN



Note 1: I/O pin has diode protection to VDD and VSS.

2: To enable weak pull-ups, set the appropriate TRIS bit(s) and clear the $\overline{\text{RBPU}}$ bit (INTCON2<7>).

3: The CCP2 input/output is multiplexed with RB3 if the CCP2MX bit is enabled (= 0) in the Configuration register and the device is operating in Microprocessor, Microprocessor with Boot Block or Extended Microcontroller mode.

TABLE 10-5: PORTC FUNCTIONS

Name	Bit#	Buffer Type	Function
RC0/T1OSO/T13CKI	bit 0	ST	Input/output port pin, Timer1 oscillator output or Timer1/Timer3 clock input.
RC1/T1OSI/CCP2(1)	bit 1	ST	Input/output port pin, Timer1 oscillator input or Capture2 input/ Compare2 output/PWM output (when CCP2MX configuration bit is disabled).
RC2/CCP1	bit 2	ST	Input/output port pin or Capture1 input/Compare1 output/ PWM1 output.
RC3/SCK/SCL	bit 3	ST	RC3 can also be the synchronous serial clock for both SPI and I^2C modes.
RC4/SDI/SDA	bit 4	ST	RC4 can also be the SPI data in (SPI mode) or data I/O (I^2C mode).
RC5/SDO	bit 5	ST	Input/output port pin or synchronous serial port data output.
RC6/TX1/CK1	bit 6	ST	Input/output port pin, addressable USART1 asynchronous transmit or addressable USART1 synchronous clock.
RC7/RX1/DT1	bit 7	ST	Input/output port pin, addressable USART1 asynchronous receive or addressable USART1 synchronous data.

Legend: ST = Schmitt Trigger input

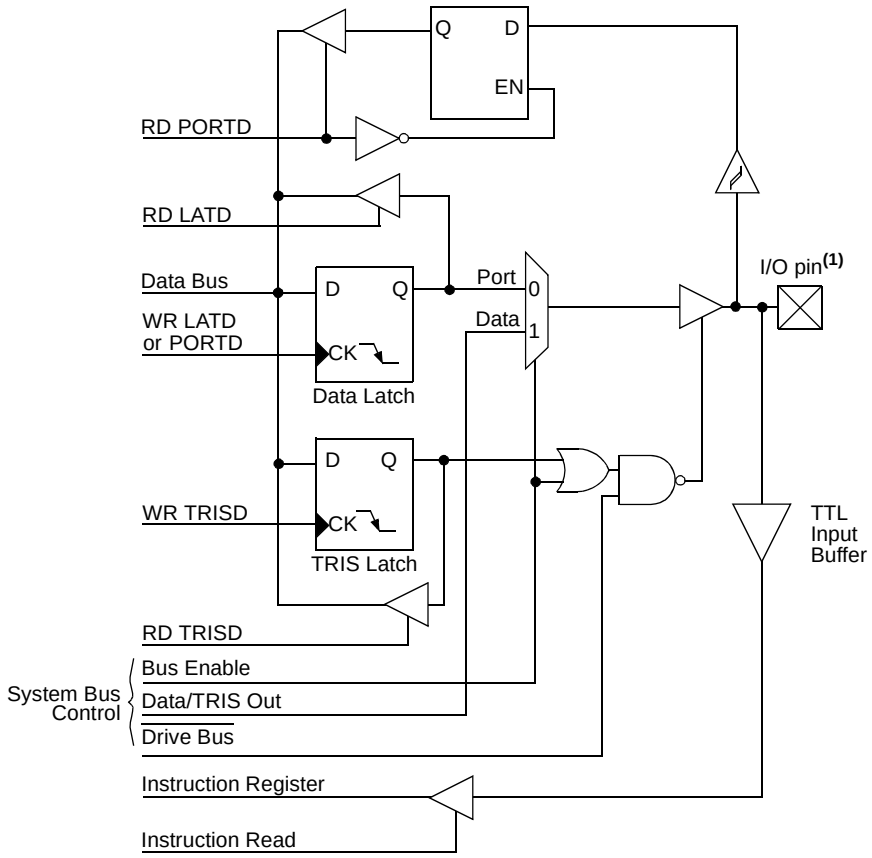
Note 1: RB3 is the alternate assignment for CCP2 when CCP2MX is set.

TABLE 10-6: SUMMARY OF REGISTERS ASSOCIATED WITH PORTC

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	Value on all other Resets
PORTC	RC7	RC6	RC5	RC4	RC3	RC2	RC1	RC0	xxxx xxxx	uuuu uuuu
LATC	LATC Data Output Register								xxxx xxxx	uuuu uuuu
TRISC	PORTC Data Direction Register								1111 1111	1111 1111

Legend: x = unknown, u = unchanged

FIGURE 10-10: PORTD BLOCK DIAGRAM IN SYSTEM BUS MODE



Note 1: I/O pins have protection diodes to VDD and VSS.

TABLE 10-13: PORTG FUNCTIONS

Name	Bit#	Buffer Type	Function
RG0/CCP3	bit 0	ST	Input/output port pin or Capture3 input/Compare3 output/PWM3 output.
RG1/TX2/CK2	bit 1	ST	Input/output port pin, addressable USART2 asynchronous transmit or addressable USART2 synchronous clock.
RG2/RX2/DT2	bit 2	ST	Input/output port pin, addressable USART2 asynchronous receive or addressable USART2 synchronous data.
RG3/CCP4	bit 3	ST	Input/output port pin or Capture4 input/Compare4 output/PWM4 output.
RG4/CCP5	bit 4	ST	Input/output port pin or Capture5 input/Compare5 output/PWM5 output.

Legend: ST = Schmitt Trigger input

TABLE 10-14: SUMMARY OF REGISTERS ASSOCIATED WITH PORTG

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	Value on all other Resets
PORTG	—	—	—	Read PORTF pin/Write PORTF Data Latch					---x xxxx	---u uuuu
LATG	—	—	—	LATG Data Output Register					---x xxxx	---u uuuu
TRISG	—	—	—	Data Direction Control Register for PORTG					---1 1111	---1 1111

Legend: x = unknown, u = unchanged

TABLE 10-15: PORTH FUNCTIONS

Name	Bit#	Buffer Type	Function
RH0/A16	bit 0	ST/TTL(1)	Input/output port pin or address bit 16 for external memory interface.
RH1/A17	bit 1	ST/TTL(1)	Input/output port pin or address bit 17 for external memory interface.
RH2/A18	bit 2	ST/TTL(1)	Input/output port pin or address bit 18 for external memory interface.
RH3/A19	bit 3	ST/TTL(1)	Input/output port pin or address bit 19 for external memory interface.
RH4/AN12	bit 4	ST	Input/output port pin or analog input channel 12.
RH5/AN13	bit 5	ST	Input/output port pin or analog input channel 13.
RH6/AN14	bit 6	ST	Input/output port pin or analog input channel 14.
RH7/AN15	bit 7	ST	Input/output port pin or analog input channel 15.

Legend: ST = Schmitt Trigger input, TTL = TTL input

Note 1: Input buffers are Schmitt Triggers when in I/O mode and TTL buffers when in System Bus or Parallel Slave Port mode.

TABLE 10-16: SUMMARY OF REGISTERS ASSOCIATED WITH PORTH

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	Value on all other Resets
TRISH	PORTH Data Direction Control Register								1111 1111	1111 1111
PORTH	Read PORTH pin/Write PORTH Data Latch								xxxx xxxx	uuuu uuuu
LATH	Read PORTH Data Latch/Write PORTH Data Latch								xxxx xxxx	uuuu uuuu
ADCON1	—	—	VCFG1	VCFG0	PCFG3	PCFG2	PCFG1	PCFG0	--00 0000	--00 0000
MEMCON	EBDIS	—	WAIT1	WAIT0	—	—	WM1	WM0	0-00 --00	0-00 --00

Legend: x = unknown, u = unchanged, – = unimplemented. Shaded cells are not used by PORTH.

FIGURE 17-5: SPI MODE WAVEFORM (SLAVE MODE WITH CKE = 0)

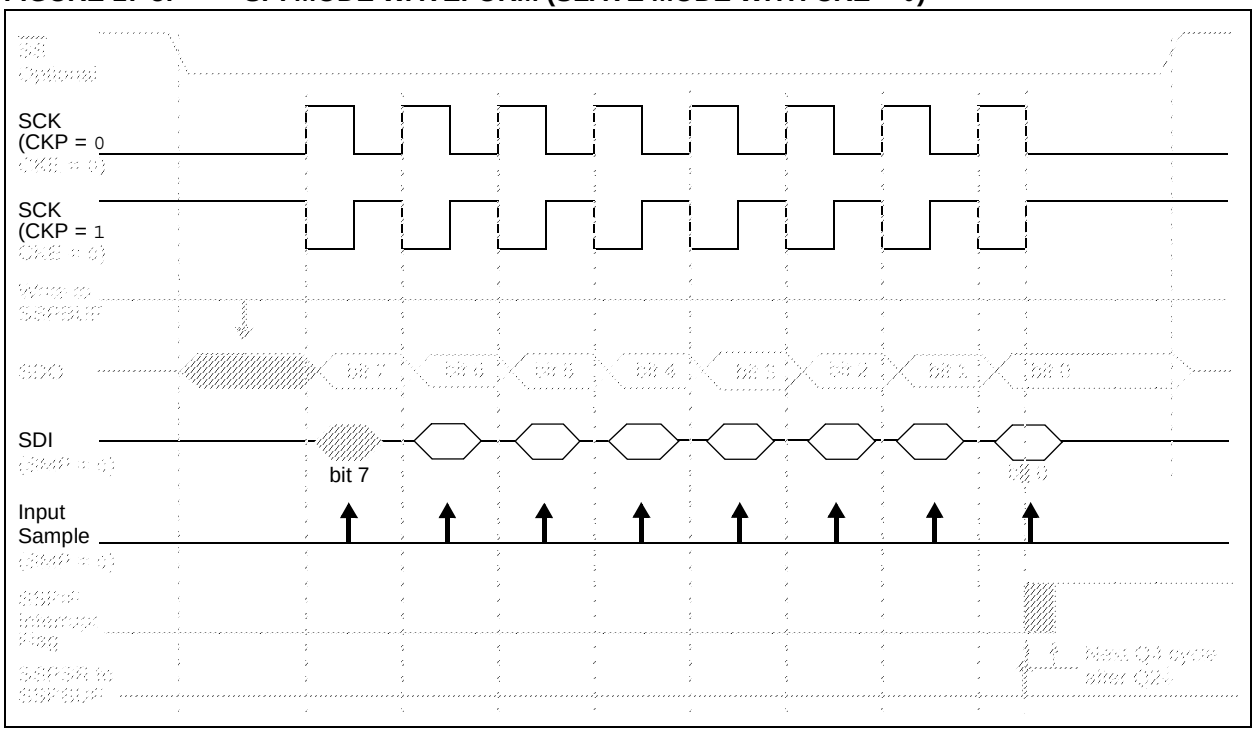
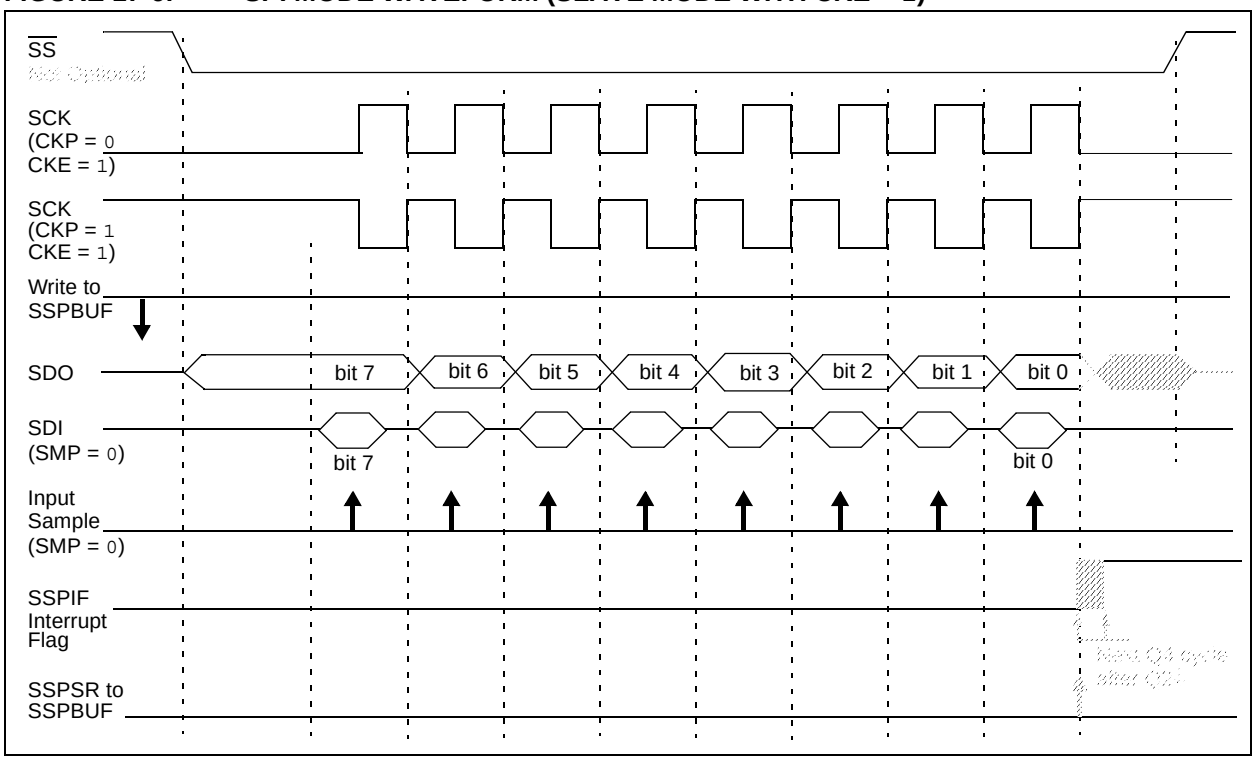


FIGURE 17-6: SPI MODE WAVEFORM (SLAVE MODE WITH CKE = 1)



18.3 USART Synchronous Master Mode

In Synchronous Master mode, the data is transmitted in a half-duplex manner (i.e., transmission and reception do not occur at the same time). When transmitting data, the reception is inhibited and vice versa. Synchronous mode is entered by setting bit SYNC (TXSTAx<4>). In addition, enable bit SPEN (RCSTAx<7>) is set in order to configure the appropriate I/O pins to CK (clock) and DT (data) lines, respectively. The Master mode indicates that the processor transmits the master clock on the CK line. The Master mode is entered by setting bit CSRC (TXSTAx<7>).

18.3.1 USART SYNCHRONOUS MASTER TRANSMISSION

The USART transmitter block diagram is shown in Figure 18-1. The heart of the transmitter is the Transmit (Serial) Shift Register (TSR). The shift register obtains its data from the Read/Write Transmit Buffer register, TXREG. The TXREGx register is loaded with data in software. The TSR register is not loaded until the last bit has been transmitted from the previous load. As soon as the last bit is transmitted, the TSR is loaded with new data from the TXREGx (if available). Once the TXREGx register transfers the data to the TSR register (occurs in one TCYCLE), the TXREGx is empty and interrupt bit TXxIF (PIR1<4> for USART1, PIR3<4> for USART2) is set. The interrupt can be enabled/disabled by setting/clearing enable bit TXxIE (PIE1<4> for USART1, PIE3<4> for USART2). Flag bit TXxIF will be set, regardless of the state of enable bit TXxIE and cannot be cleared in software. It will reset only when new data is loaded into the TXREGx register. While flag bit TXxIF indicates the status of the TXREGx register, another bit TRMT (TXSTAx<1>) shows the status of the TSR register. TRMT is a read-only bit, which is set when the TSR is empty. No interrupt logic is tied to this bit, so the user has to poll this bit in order to determine if the TSR register is empty. The TSR is not mapped in data memory, so it is not available to the user.

To set up a Synchronous Master Transmission:

1. Initialize the SPBRG register for the appropriate baud rate (**Section 18.1 “USART Baud Rate Generator (BRG)”**).
2. Enable the synchronous master serial port by setting bits SYNC, SPEN and CSRC.
3. If interrupts are desired, set enable bit TXxIE in the appropriate PIE register.
4. If 9-bit transmission is desired, set bit TX9.
5. Enable the transmission by setting bit TXEN.
6. If 9-bit transmission is selected, the ninth bit should be loaded in bit TX9D.
7. Start transmission by loading data to the TXREGx register.

Note:	TXIF is not cleared immediately upon loading data into the transmit buffer TXREG. The flag bit becomes valid in the second instruction cycle following the load instruction.

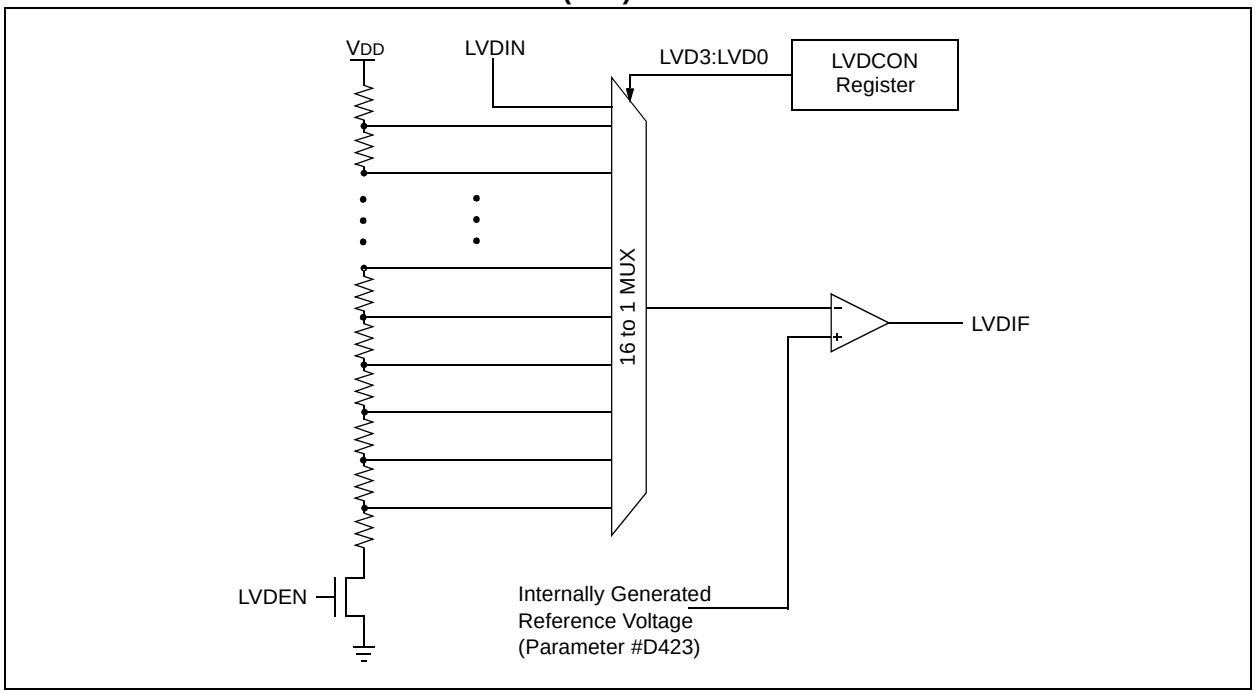
TABLE 18-8: REGISTERS ASSOCIATED WITH SYNCHRONOUS MASTER TRANSMISSION

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR, BOR	Value on all other Resets
INTCON	GIE/ GIEH	PEIE/ GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF	0000 0000	0000 0000
PIR1	PSPIF	ADIF	RCIF	TXIF	SSPIF	CCP1IF	TMR2IF	TMR1IF	0000 0000	0000 0000
PIE1	PSPIE	ADIE	RCIE	TXIE	SSPIE	CCP1IE	TMR2IE	TMR1IE	0000 0000	0000 0000
IPR1	PSPIP	ADIP	RCIP	TXIP	SSPIP	CCP1IP	TMR2IP	TMR1IP	0111 1111	0111 1111
PIR3	—	—	RC2IF	TX2IF	TMR4IF	CCP5IF	CCP4IF	CCP3IF	--00 0000	--00 0000
PIE3	—	—	RC2IE	TX2IE	TMR4IE	CCP5IE	CCP4IE	CCP3IE	--00 0000	--00 0000
IPR3	—	—	RC2IP	TX2IP	TMR4IP	CCP5IP	CCP4IP	CCP3IP	--11 1111	--11 1111
RCSTAx(1)	SPEN	RX9	SREN	CREN	ADDEN	FERR	OERR	RX9D	0000 000x	0000 000x
TXREGx(1)	USART Transmit Register								0000 0000	0000 0000
TXSTAx(1)	CSRC	TX9	TXEN	SYNC	—	BRGH	TRMT	TX9D	0000 -010	0000 -010
SPBRGx(1)	Baud Rate Generator Register								0000 0000	0000 0000

Legend: x = unknown, – = unimplemented, read as ‘0’. Shaded cells are not used for synchronous master transmission.

Note 1: Register names generically refer to both of the identically named registers for the two USART modules, where ‘x’ indicates the particular module. Bit names and Reset values are identical between modules.

FIGURE 22-2: LOW-VOLTAGE DETECT (LVD) BLOCK DIAGRAM



The LVD module has an additional feature that allows the user to supply the trip voltage to the module from an external source. This mode is enabled when bits LVDL3:LVDL0 are set to '1111'. In this state, the comparator input is multiplexed from the external input pin, LVDIN (Figure 22-3). This gives users flexibility because it allows them to configure the Low-Voltage Detect interrupt to occur at any voltage in the valid operating range.

FIGURE 22-3: LOW-VOLTAGE DETECT (LVD) WITH EXTERNAL INPUT BLOCK DIAGRAM

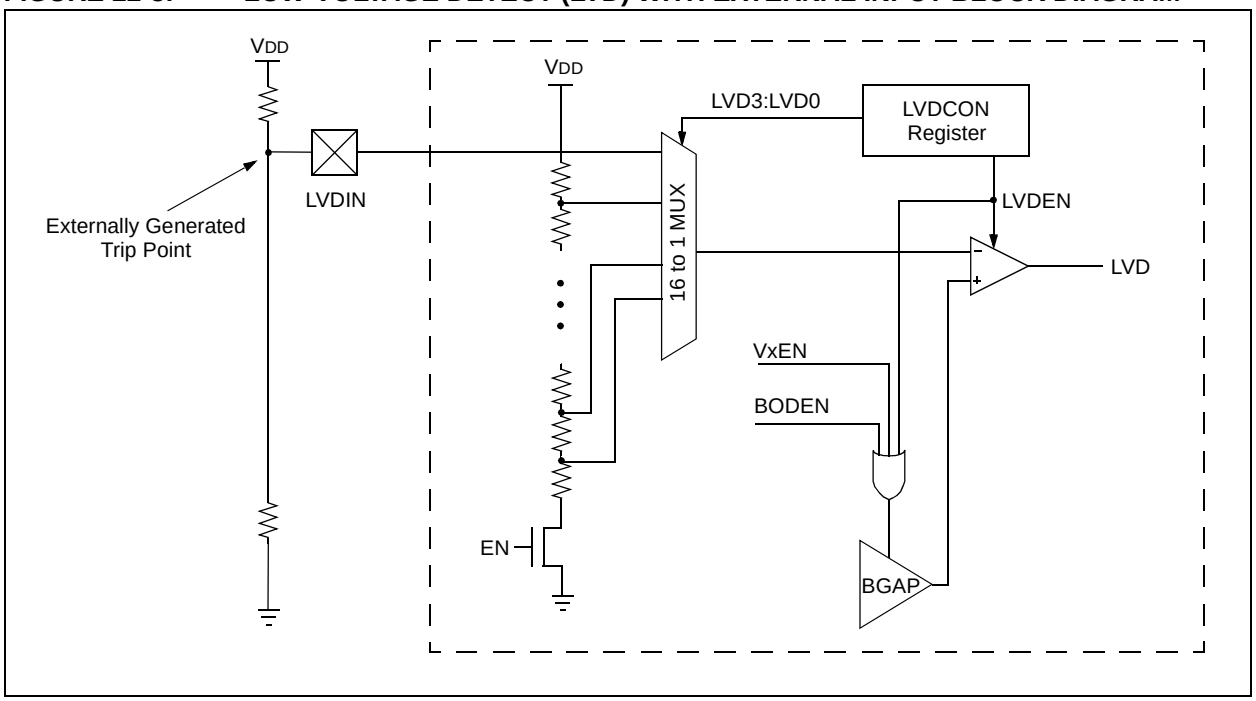


TABLE 23-1: CONFIGURATION BITS AND DEVICE IDS

File Name		Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Default/ Unprogrammed Value
300001h	CONFIG1H	—	—	OSCSEN	—	—	FOSC2	FOSC1	FOSC0	--1- -111
300002h	CONFIG2L	—	—	—	—	BORV1	BORV0	BODEN	PWRTEN	---- 1111
300003h	CONFIG2H	—	—	—	—	WDTPS2	WDTPS1	WDTPS0	WDTEN	---- 1111
300004h(1)	CONFIG3L	WAIT	—	—	—	—	—	PM1	PM0	1--- --11
300005h	CONFIG3H	—	—	—	—	—	—	r(3)	CCP2MX	---- --11
300006h	CONFIG4L	DEBUG	—	—	—	—	LVP	—	STVREN	1--- -1-1
300008h	CONFIG5L	CP7(2)	CP6(2)	CP5(2)	CP4(2)	CP3	CP2	CP1	CP0	1111 1111
300009h	CONFIG5H	CPD	CPB	—	—	—	—	—	—	11-- ----
30000Ah	CONFIG6L	WRT7(2)	WRT6(2)	WRT5(2)	WRT4(2)	WRT3	WRT2	WRT1	WRT0	1111 1111
30000Bh	CONFIG6H	WRTD	WRTB	WRTC	—	—	—	—	—	111- ----
30000Ch	CONFIG7L	EBTR7(2)	EBTR6(2)	EBTR5(2)	EBTR4(2)	EBTR3	EBTR2	EBTR1	EBTR0	1111 1111
30000Dh	CONFIG7H	—	EBTRB	—	—	—	—	—	—	-1-- ----
3FFFFEh	DEVID1	DEV2	DEV1	DEV0	REV4	REV3	REV2	REV1	REV0	(4)
3FFFFFh	DEVID2	DEV10	DEV9	DEV8	DEV7	DEV6	DEV5	DEV4	DEV3	0000 0110

Legend: x = unknown, u = unchanged, – = unimplemented, q = value depends on condition, r = reserved.
Shaded cells are unimplemented, read as '0'.

Note 1: Unimplemented in PIC18F6X20 devices; maintain this bit set.
2: Unimplemented in PIC18FX520 and PIC18FX620 devices; maintain this bit set.
3: Unimplemented in PIC18FX620 and PIC18FX720 devices; maintain this bit set.
4: See Register 23-13 for DEVID1 values.

REGISTER 23-1: CONFIG1H: CONFIGURATION REGISTER 1 HIGH (BYTE ADDRESS 300001h)

U-0	U-0	R/P-1	U-0	U-0	R/P-1	R/P-1	R/P-1
—	—	$\overline{\text{OSCSEN}}$	—	—	FOSC2	FOSC1	FOSC0
bit 7							bit 0

bit 7-6 **Unimplemented:** Read as '0'

bit 5 **$\overline{\text{OSCSEN}}$**: Oscillator System Clock Switch Enable bit
1 = Oscillator system clock switch option is disabled (main oscillator is source)
0 = Timer1 Oscillator system clock switch option is enabled (oscillator switching is enabled)

bit 4-3 **Unimplemented:** Read as '0'

bit 2-0 **FOSC2:FOSC0**: Oscillator Selection bits
111 = RC oscillator w/ OSC2 configured as RA6
110 = HS oscillator with PLL enabled; clock frequency = (4 x FOSC)
101 = EC oscillator w/ OSC2 configured as RA6
100 = EC oscillator w/ OSC2 configured as divide-by-4 clock output
011 = RC oscillator w/ OSC2 configured as divide-by-4 clock output
010 = HS oscillator
001 = XT oscillator
000 = LP oscillator

Legend:		
R = Readable bit	P = Programmable bit	U = Unimplemented bit, read as '0'
- n = Value when device is unprogrammed		u = Unchanged from programmed state

TABLE 24-1: OPCODE FIELD DESCRIPTIONS

Field	Description
a	RAM access bit a = 0: RAM location in Access RAM (BSR register is ignored) a = 1: RAM bank is specified by BSR register
bbb	Bit address within an 8-bit file register (0 to 7).
BSR	Bank Select Register. Used to select the current RAM bank.
d	Destination select bit d = 0: store result in WREG d = 1: store result in file register f
dest	Destination either the WREG register or the specified register file location.
f	8-bit Register file address (0x00 to 0xFF).
fs	12-bit Register file address (0x000 to 0xFFF). This is the source address.
fd	12-bit Register file address (0x000 to 0xFFF). This is the destination address.
k	Literal field, constant data or label (may be either an 8-bit, 12-bit or a 20-bit value).
label	Label name.
mm * *+ *- +*	The mode of the TBLPTR register for the table read and table write instructions. Only used with table read and table write instructions: No Change to register (such as TBLPTR with table reads and writes) Post-Increment register (such as TBLPTR with table reads and writes) Post-Decrement register (such as TBLPTR with table reads and writes) Pre-Increment register (such as TBLPTR with table reads and writes)
n	The relative address (2's complement number) for relative branch instructions, or the direct address for Call/Branch and Return instructions.
PRODH	Product of Multiply High Byte.
PRODL	Product of Multiply Low Byte.
s	Fast Call/Return mode select bit s = 0: do not update into/from shadow registers s = 1: certain registers loaded into/from shadow registers (Fast mode)
u	Unused or Unchanged.
WREG	Working register (accumulator).
x	Don't care ('0' or '1'). The assembler will generate code with x = 0. It is the recommended form of use for compatibility with all Microchip software tools.
TBLPTR	21-bit Table Pointer (points to a Program Memory location).
TABLAT	8-bit Table Latch.
TOS	Top-of-Stack.
PC	Program Counter.
PCL	Program Counter Low Byte.
PCH	Program Counter High Byte.
PCLATH	Program Counter High Byte Latch.
PCLATU	Program Counter Upper Byte Latch.
GIE	Global Interrupt Enable bit.
WDT	Watchdog Timer.
$\overline{TO}$	Time-out bit.
$\overline{PD}$	Power-down bit.
C, DC, Z, OV, N	ALU Status bits: Carry, Digit Carry, Zero, Overflow, Negative.
[ ]	Optional.
( )	Contents.
→	Assigned to.
< >	Register bit field.
∈	In the set of.
italics	User defined term (font is courier).

ANDWF **AND W with f**

Syntax: [*label*] ANDWF f [,d [,a]

Operands: $0 \le f \le 255$
$d \in [0,1]$
$a \in [0,1]$

Operation: (W) .AND. (f) → dest

Status Affected: N, Z

Encoding:

0001	01da	ffff	ffff

Description: The contents of W are AND'ed with register 'f'. If 'd' is '0', the result is stored in W. If 'd' is '1', the result is stored back in register 'f' (default). If 'a' is '0', the Access Bank will be selected. If 'a' is '1', the BSR will not be overridden (default).

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write to destination

Example: `ANDWF REG, 0, 0`

Before Instruction
W = 0x17
REG = 0xC2

After Instruction
W = 0x02
REG = 0xC2

BC **Branch if Carry**

Syntax: [*label*] BC n

Operands: $-128 \le n \le 127$

Operation: if Carry bit is '1'
(PC) + 2 + 2n → PC

Status Affected: None

Encoding:

1110	0010	nnnn	nnnn

Description: If the Carry bit is '1', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC+2+2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example: `HERE BC 5`

Before Instruction
PC = address (HERE)

After Instruction
If Carry = 1;
PC = address (HERE+12)
If Carry = 0;
PC = address (HERE+2)

BTG **Bit Toggle f**

Syntax: [*label*] BTG f,b[,a]

Operands: $0 \le f \le 255$
$0 \le b < 7$
$a \in [0,1]$

Operation: $(\overline{f<b>}) \rightarrow f<b>$

Status Affected: None

Encoding:

0111	bbba	ffff	ffff

Description: Bit 'b' in data memory location 'f' is inverted. If 'a' is '0', the Access Bank will be selected, overriding the BSR value. If 'a' = 1, then the bank will be selected as per the BSR value (default).

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: `BTG PORTC, 4, 0`

Before Instruction:
PORTC = 0111 0101 [0x75]
After Instruction:
PORTC = 0110 0101 [0x65]

BOV **Branch if Overflow**

Syntax: [*label*] BOV n

Operands: $-128 \le n \le 127$

Operation: if Overflow bit is '1'
$(PC) + 2 + 2n \rightarrow PC$

Status Affected: None

Encoding:

1110	0100	nnnn	nnnn

Description: If the Overflow bit is '1', then the program will branch.
The 2's complement number '2n' is added to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be PC+2+2n. This instruction is then a two-cycle instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:
If Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

If No Jump:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	No operation

Example:
```
HERE    BOV  Jump
```
Before Instruction
PC = address (HERE)
After Instruction
If Overflow = 1;
PC = address (Jump)
If Overflow = 0;
PC = address (HERE+2)

25.2 MPLAB C Compilers for Various Device Families

The MPLAB C Compiler code development systems are complete ANSI C compilers for Microchip's PIC18, PIC24 and PIC32 families of microcontrollers and the dsPIC30 and dsPIC33 families of digital signal controllers. These compilers provide powerful integration capabilities, superior code optimization and ease of use.

For easy source level debugging, the compilers provide symbol information that is optimized to the MPLAB IDE debugger.

25.3 HI-TECH C for Various Device Families

The HI-TECH C Compiler code development systems are complete ANSI C compilers for Microchip's PIC family of microcontrollers and the dsPIC family of digital signal controllers. These compilers provide powerful integration capabilities, omniscient code generation and ease of use.

For easy source level debugging, the compilers provide symbol information that is optimized to the MPLAB IDE debugger.

The compilers include a macro assembler, linker, preprocessor, and one-step driver, and can run on multiple platforms.

25.4 MPASM Assembler

The MPASM Assembler is a full-featured, universal macro assembler for PIC10/12/16/18 MCUs.

The MPASM Assembler generates relocatable object files for the MPLINK Object Linker, Intel® standard HEX files, MAP files to detail memory usage and symbol reference, absolute LST files that contain source lines and generated machine code and COFF files for debugging.

The MPASM Assembler features include:

- Integration into MPLAB IDE projects
- User-defined macros to streamline assembly code
- Conditional assembly for multi-purpose source files
- Directives that allow complete control over the assembly process

25.5 MPLINK Object Linker/ MPLIB Object Librarian

The MPLINK Object Linker combines relocatable objects created by the MPASM Assembler and the MPLAB C18 C Compiler. It can link relocatable objects from precompiled libraries, using directives from a linker script.

The MPLIB Object Librarian manages the creation and modification of library files of precompiled code. When a routine from a library is called from a source file, only the modules that contain that routine will be linked in with the application. This allows large libraries to be used efficiently in many different applications.

The object linker/library features include:

- Efficient linking of single libraries instead of many smaller files
- Enhanced code maintainability by grouping related modules together
- Flexible creation of libraries with easy module listing, replacement, deletion and extraction

25.6 MPLAB Assembler, Linker and Librarian for Various Device Families

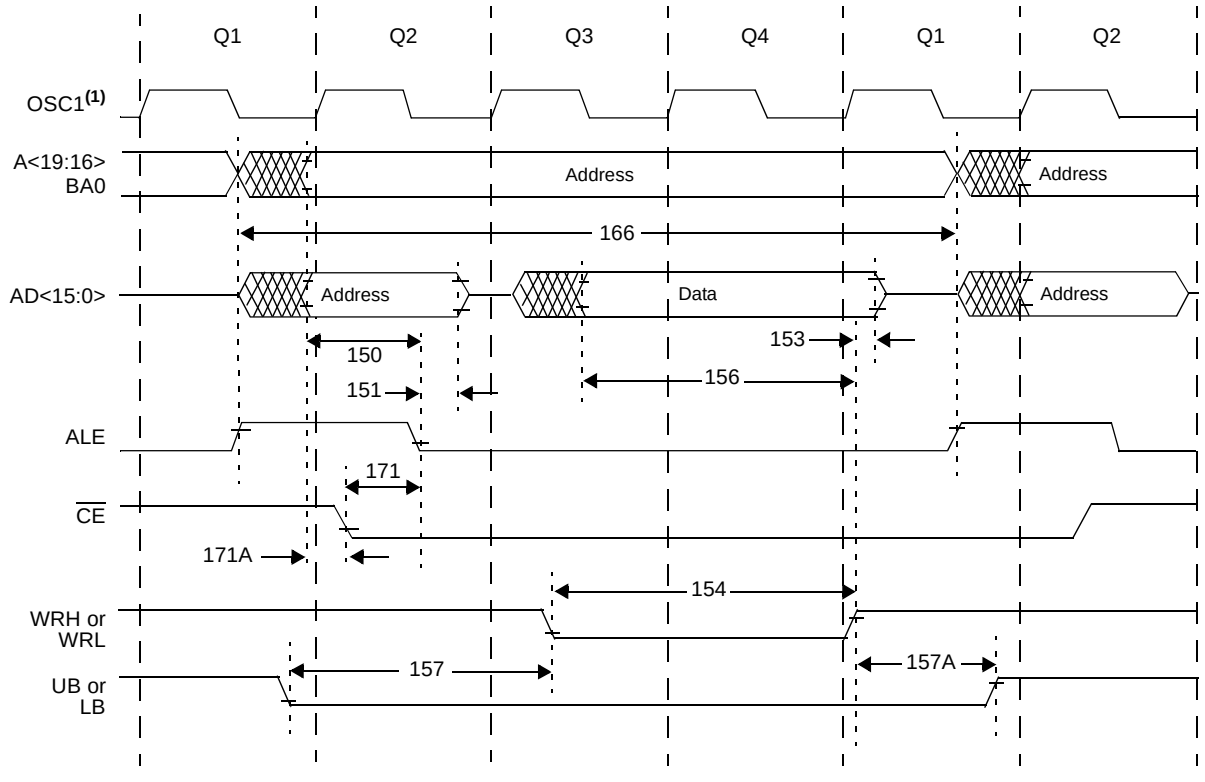
MPLAB Assembler produces relocatable machine code from symbolic assembly language for PIC24, PIC32 and dsPIC devices. MPLAB C Compiler uses the assembler to produce its object file. The assembler generates relocatable object files that can then be archived or linked with other relocatable object files and archives to create an executable file. Notable features of the assembler include:

- Support for the entire device instruction set
- Support for fixed-point and floating-point data
- Command line interface
- Rich directive set
- Flexible macro language
- MPLAB IDE compatibility

TABLE 26-9: CLKO AND I/O TIMING REQUIREMENTS

Param No.	Symbol	Characteristics	Min	Typ	Max	Units
150	TADV2ALL	Address Out Valid to ALE ↓ (address setup time)	0.25 TCY – 10	—	—	ns
151	TALL2ADL	ALE ↓ to Address Out Invalid (address hold time)	5	—	—	ns
155	TALL2OEL	ALE ↓ to $\overline{OE}$ ↓	10	0.125 TCY	—	ns
160	TADZ2OEL	AD high-Z to $\overline{OE}$ ↓ (bus release to $\overline{OE}$)	0	—	—	ns
161	TOEH2ADD	$\overline{OE}$ ↑ to AD Driven	0.125 TCY – 5	—	—	ns
162	TADV2OEH	LS Data Valid before $\overline{OE}$ ↑ (data setup time)	20	—	—	ns
163	TOEH2ADL	$\overline{OE}$ ↑ to Data In Invalid (data hold time)	0	—	—	ns
164	TALH2ALL	ALE Pulse Width	—	0.25 TCY	—	ns
165	TOEL2OEH	$\overline{OE}$ Pulse Width	0.5 TCY – 5	0.5 TCY	—	ns
166	TALH2ALH	ALE ↑ to ALE ↑ (cycle time)	—	TCY	—	ns
167	TACC	Address Valid to Data Valid	0.75 TCY – 25	—	—	ns
168	TOE	$\overline{OE}$ ↓ to Data Valid		—	0.5 TCY – 25	ns
169	TALL2OEH	ALE ↓ to $\overline{OE}$ ↑	0.625 TCY – 10	—	0.625 TCY + 10	ns
171	TALH2CSL	Chip Enable Active to ALE ↓	—	—	10	ns
171A	TUBL2OEH	AD Valid to Chip Enable Active	0.25 TCY – 20	—	—	ns

FIGURE 26-10: PROGRAM MEMORY WRITE TIMING DIAGRAM



Operating Conditions: 2.0V < VCC < 5.5V, -40°C < TA < +125°C unless otherwise stated.

Note 1: Maximum speed of FOSC is 25 MHz for external program memory read.

FIGURE 27-25: A/D NONLINEARITY vs. VREFH (VDD = VREFH, -40°C TO +125°C)

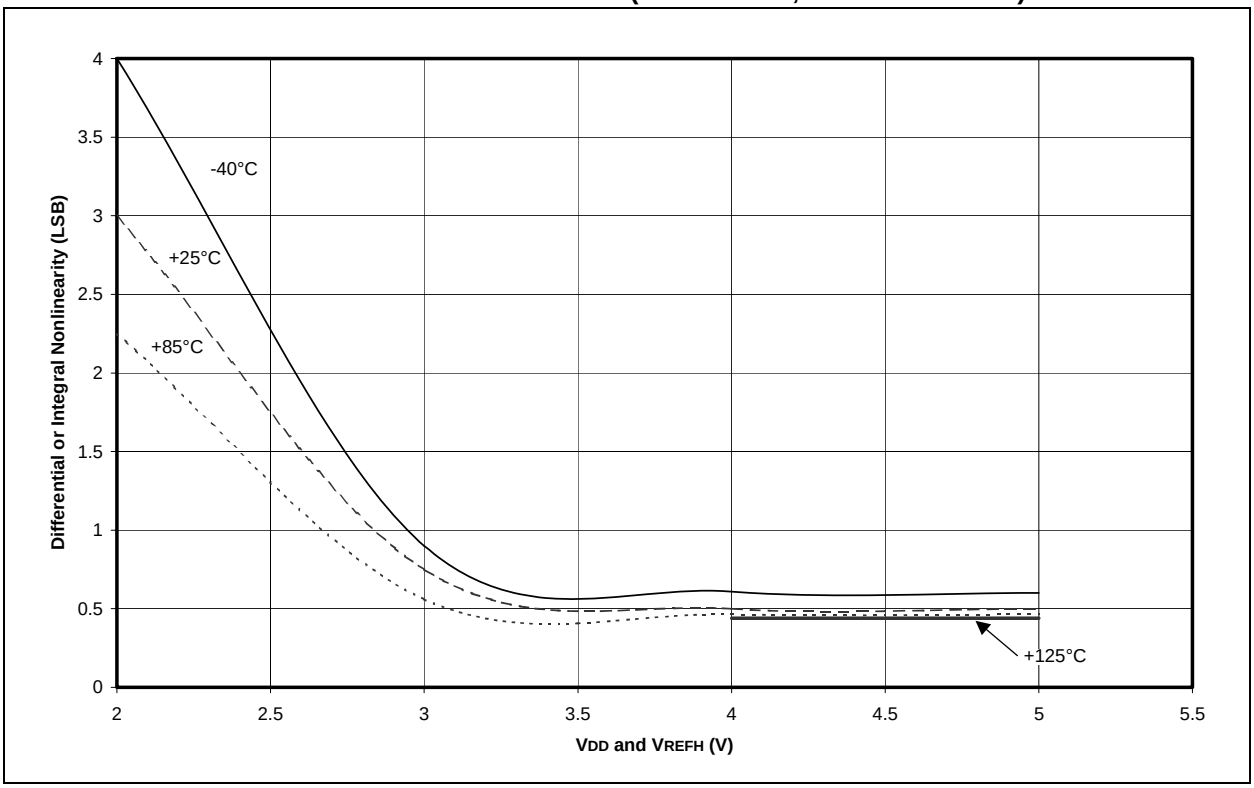
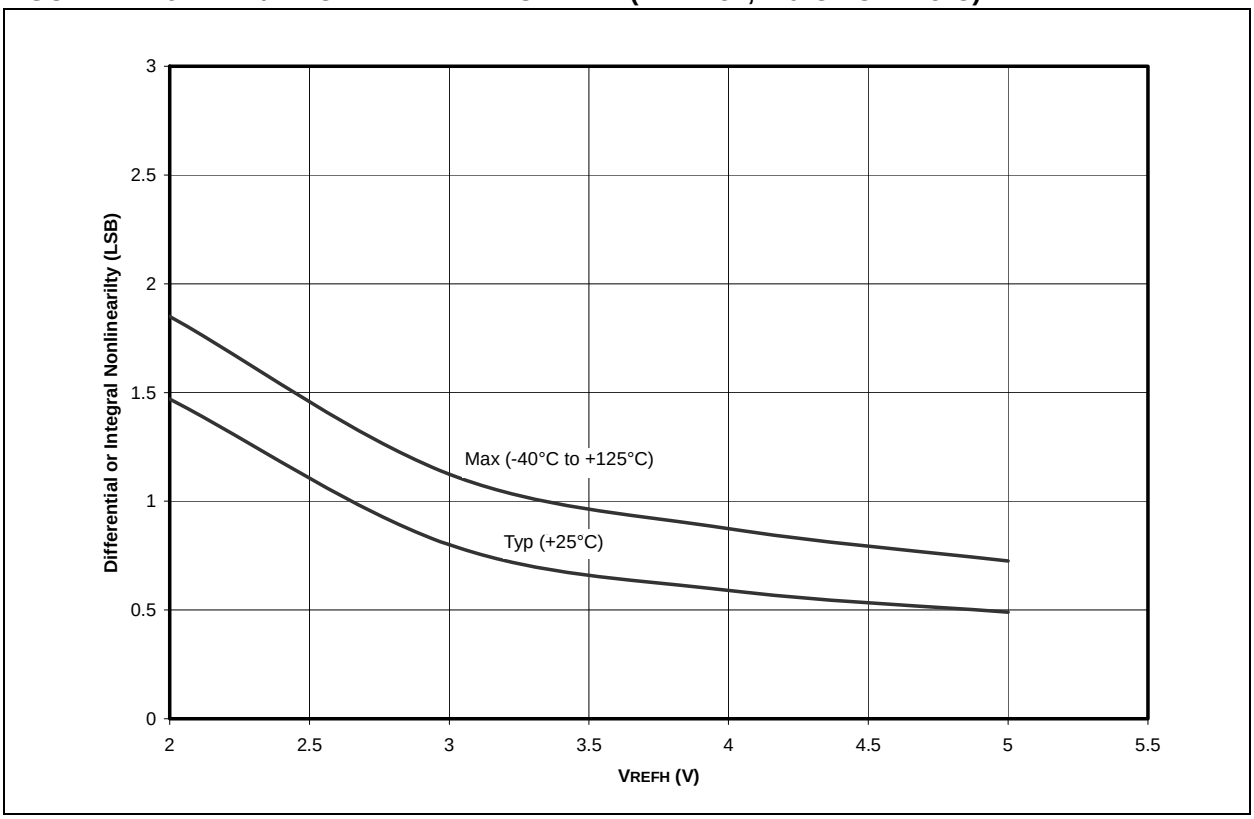


FIGURE 27-26: A/D NONLINEARITY vs. VREFH (VDD = 5V, -40°C TO +125°C)



Note the following details of the code protection feature on Microchip devices:

- Microchip products meet the specification contained in their particular Microchip Data Sheet.
- Microchip believes that its family of products is one of the most secure families of its kind on the market today, when used in the intended manner and under normal conditions.
- There are dishonest and possibly illegal methods used to breach the code protection feature. All of these methods, to our knowledge, require using the Microchip products in a manner outside the operating specifications contained in Microchip's Data Sheets. Most likely, the person doing so is engaged in theft of intellectual property.
- Microchip is willing to work with the customer who is concerned about the integrity of their code.
- Neither Microchip nor any other semiconductor manufacturer can guarantee the security of their code. Code protection does not mean that we are guaranteeing the product as "unbreakable."

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