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Details

Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	48MHz
Connectivity	I ² C, LINbus, SPI, UART/USART, USB
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	34
Program Memory Size	64KB (32K x 16)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	3.8K x 8
Voltage - Supply (Vcc/Vdd)	2.15V ~ 3.6V
Data Converters	A/D 13x10b/12b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	44-TQFP
Supplier Device Package	44-TQFP (10x10)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18f46j53t-i-pt

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TADLE 3-2.	INITIALIZATION CONDITIONS FOR ALL REGISTERS (CONTINUED)					
Register	Applicable Devices		Applicable Devices Power-on Reset, Brown-out Reset, Wake From Deep Sleep		Wake-up via WDT or Interrupt	
INDF2	PIC18F2XJ53	PIC18F4XJ53	N/A	N/A	N/A	
POSTINC2	PIC18F2XJ53	PIC18F4XJ53	N/A	N/A	N/A	
POSTDEC2	PIC18F2XJ53	PIC18F4XJ53	N/A	N/A	N/A	
PREINC2	PIC18F2XJ53	PIC18F4XJ53	N/A	N/A	N/A	
PLUSW2	PIC18F2XJ53	PIC18F4XJ53	N/A	N/A	N/A	
FSR2H	PIC18F2XJ53	PIC18F4XJ53	0000	0000	uuuu	
FSR2L	PIC18F2XJ53	PIC18F4XJ53	xxxx xxxx	uuuu uuuu	uuuu uuuu	
STATUS	PIC18F2XJ53	PIC18F4XJ53	x xxxx	u uuuu	u uuuu	
TMR0H	PIC18F2XJ53	PIC18F4XJ53	0000 0000	0000 0000	uuuu uuuu	
TMR0L	PIC18F2XJ53	PIC18F4XJ53	xxxx xxxx	uuuu uuuu	uuuu uuuu	
TOCON	PIC18F2XJ53	PIC18F4XJ53	1111 1111	1111 1111	uuuu uuuu	
OSCCON	PIC18F2XJ53	PIC18F4XJ53	0110 qq00	0110 qq00	0110 qq0u	
CM1CON	PIC18F2XJ53	PIC18F4XJ53	0001 1111	uuuu uuuu	uuuu uuuu	
CM2CON	PIC18F2XJ53	PIC18F4XJ53	0001 1111	uuuu uuuu	uuuu uuuu	
RCON ⁽⁴⁾	PIC18F2XJ53	PIC18F4XJ53	0-11 11qq	0-qq qquu	u-qq qquu	
TMR1H	PIC18F2XJ53	PIC18F4XJ53	xxxx xxxx	uuuu uuuu	uuuu uuuu	
TMR1L	PIC18F2XJ53	PIC18F4XJ53	xxxx xxxx	uuuu uuuu	uuuu uuuu	
T1CON	PIC18F2XJ53	PIC18F4XJ53	0000 0000	u0uu uuuu	uuuu uuuu	
TMR2	PIC18F2XJ53	PIC18F4XJ53	0000 0000	0000 0000	uuuu uuuu	
PR2	PIC18F2XJ53	PIC18F4XJ53	1111 1111	1111 1111	uuuu uuuu	
T2CON	PIC18F2XJ53	PIC18F4XJ53	-000 0000	-000 0000	-uuu uuuu	
SSP1BUF	PIC18F2XJ53	PIC18F4XJ53	XXXX XXXX	uuuu uuuu	uuuu uuuu	
SSP1ADD	PIC18F2XJ53	PIC18F4XJ53	0000 0000	0000 0000	uuuu uuuu	
SSP1MSK	PIC18F2XJ53	PIC18F4XJ53		uuuu uuuu	uuuu uuuu	
SSP1STAT	PIC18F2XJ53	PIC18F4XJ53	0000 0000	0000 0000	uuuu uuuu	
SSP1CON1	PIC18F2XJ53	PIC18F4XJ53	0000 0000	0000 0000	uuuu uuuu	
SSP1CON2	PIC18F2XJ53	PIC18F4XJ53	0000 0000	0000 0000	uuuu uuuu	
ADRESH	PIC18F2XJ53	PIC18F4XJ53	XXXX XXXX	uuuu uuuu	uuuu uuuu	
ADRESL	PIC18F2XJ53	PIC18F4XJ53	XXXX XXXX	uuuu uuuu	uuuu uuuu	
ADCON0	PIC18F2XJ53	PIC18F4XJ53	0000 0000	0000 0000	uuuu uuuu	
ADCON1	PIC18F2XJ53	PIC18F4XJ53	0000 0000	0000 0000	uuuu uuuu	
WDTCON	PIC18F2XJ53	PIC18F4XJ53	1qq0 0000	0qq0 0000	uqqu uuuu	

TABLE 5-2:	INITIALIZATION CONDITIONS FOR ALL REGISTERS (CONTINUED)

Legend: u = unchanged, x = unknown, - = unimplemented bit, read as '0', q = value depends on condition.

- 2: When the wake-up is due to an interrupt and the GIEL or GIEH bit is set, the PC is loaded with the interrupt vector (0008h or 0018h).
- 3: One or more bits in the INTCONx or PIRx registers will be affected (to cause wake-up).
- 4: See Table 5-1 for Reset value for specific condition.
- 5: Not implemented for PIC18F2XJ53 devices.
- 6: Not implemented for "LF" devices.

Note 1: When the wake-up is due to an interrupt and the GIEL or GIEH bit is set, the TOSU, TOSH and TOSL are updated with the current value of the PC. The STKPTR is modified to point to the next location in the hardware stack.

6.1.3 PROGRAM COUNTER

The Program Counter (PC) specifies the address of the instruction to fetch for execution. The PC is 21 bits wide and is contained in three separate 8-bit registers. The low byte, known as the PCL register, is both readable and writable. The high byte, or PCH register, contains the PC<15:8> bits; it is not directly readable or writable. Updates to the PCH register are performed through the PCLATH register. The upper byte is called PCU. This register contains the PC<20:16> bits; it is also not directly readable or writable. Updates to the PCH register. Updates to the PCU. This register contains the PC<20:16> bits; it is also not directly readable or writable. Updates to the PCU

The contents of PCLATH and PCLATU are transferred to the program counter by any operation that writes to PCL. Similarly, the upper 2 bytes of the program counter are transferred to PCLATH and PCLATU by an operation that reads PCL. This is useful for computed offsets to the PC (see **Section 6.1.6.1 "Computed GOTO"**).

The PC addresses bytes in the program memory. To prevent the PC from becoming misaligned with word instructions, the Least Significant bit (LSb) of PCL is fixed to a value of '0'. The PC increments by two to address sequential instructions in the program memory.

The CALL, RCALL, GOTO and program branch instructions write to the program counter directly. For these instructions, the contents of PCLATH and PCLATU are not transferred to the program counter.

6.1.4 RETURN ADDRESS STACK

The return address stack allows any combination of up to 31 program calls and interrupts to occur. The PC is pushed onto the stack when a CALL or RCALL instruction is executed, or an interrupt is Acknowledged. The PC value is pulled off the stack on a RETURN, RETLW or a RETFIE instruction (and on ADDULNK and SUBULNK instructions if the extended instruction set is enabled). PCLATU and PCLATH are not affected by any of the RETURN or CALL instructions. The stack operates as a 31-word by 21-bit RAM and a 5-bit Stack Pointer (SP), STKPTR. The stack space is not part of either program or data space. The Stack Pointer is readable and writable and the address on the top of the stack is readable and writable through the Top-of-Stack Special Function Registers (SFRs). Data can also be pushed to, or popped from, the stack using these registers.

A CALL type instruction causes a push onto the stack. The Stack Pointer is first incremented and the location pointed to by the Stack Pointer is written with the contents of the PC (already pointing to the instruction following the CALL). A RETURN type instruction causes a pop from the stack. The contents of the location pointed to by the STKPTR are transferred to the PC and then the Stack Pointer is decremented.

The Stack Pointer is initialized to '00000' after all Resets. There is no RAM associated with the location corresponding to a Stack Pointer value of '00000'; this is only a Reset value. Status bits indicate if the stack is full, has overflowed or has underflowed.

6.1.4.1 Top-of-Stack Access

Only the top of the return address stack (TOS) is readable and writable. A set of three registers, TOSU:TOSH:TOSL, holds the contents of the stack location pointed to by the STKPTR register (Figure 6-3). This allows users to implement a software stack if necessary. After a CALL, RCALL or interrupt (and ADDULNK and SUBULNK instructions if the extended instruction set is enabled), the software can read the pushed value by reading the TOSU:TOSH:TOSL registers. These values can be placed on a user-defined software stack. At return time, the software can return these values to TOSU:TOSH:TOSL and do a return.

The user must disable the global interrupt enable bits while accessing the stack to prevent inadvertent stack corruption.



FIGURE 6-3: RETURN ADDRESS STACK AND ASSOCIATED REGISTERS

6.1.4.4 Stack Full and Underflow Resets

Device Resets on stack overflow and stack underflow conditions are enabled by setting the STVREN bit in Configuration register 1L. When STVREN is set, a full or underflow condition sets the appropriate STKFUL or STKUNF bit and then causes a device Reset. When STVREN is cleared, a full or underflow condition sets the appropriate STKFUL or STKUNF bit, but does not cause a device Reset. The STKFUL or STKUNF bits are cleared by the user software or a POR.

6.1.5 FAST REGISTER STACK (FRS)

A Fast Register Stack (FRS) is provided for the STATUS, WREG and BSR registers to provide a "fast return" option for interrupts. This stack is only one level deep and is neither readable nor writable. It is loaded with the current value of the corresponding register when the processor vectors for an interrupt. All interrupt sources push values into the Stack registers. The values in the registers are then loaded back into the working registers if the RETFIE, FAST instruction is used to return from the interrupt.

If both low-priority and high-priority interrupts are enabled, the Stack registers cannot be used reliably to return from low-priority interrupts. If a high-priority interrupt occurs while servicing a low-priority interrupt, the Stack register values stored by the low-priority interrupt will be overwritten. In these cases, users must save the key registers in software during a low-priority interrupt.

If interrupt priority is not used, all interrupts may use the FRS for returns from interrupt. If no interrupts are used, the FRS can be used to restore the STATUS, WREG and BSR registers at the end of a subroutine call. To use the Fast Register Stack for a subroutine call, a CALL label, FAST instruction must be executed to save the STATUS, WREG and BSR registers to the Fast Register Stack. A RETURN, FAST instruction is then executed to restore these registers from the FRS.

Example 6-1 provides a source code example that uses the FRS during a subroutine call and return.

EXAMPLE 6-1: FAST REGISTER STACK CODE EXAMPLE

CALL SUB1, FAST	;STATUS, WREG, BSR ;SAVED IN FAST REGISTER ;STACK
SUB1 •	
RETURN FAST	;RESTORE VALUES SAVED ;IN FAST REGISTER STACK

6.1.6 LOOK-UP TABLES IN PROGRAM MEMORY

There may be programming situations that require the creation of data structures or look-up tables in program memory. For PIC18 devices, look-up tables can be implemented in two ways:

- Computed GOTO
- Table Reads

6.1.6.1 Computed GOTO

A computed GOTO is accomplished by adding an offset to the PC. An example is shown in Example 6-2.

A look-up table can be formed with an ADDWF PCL instruction and a group of RETLW nn instructions. The W register is loaded with an offset into the table before executing a call to that table. The first instruction of the called routine is the ADDWF PCL instruction. The next executed instruction will be one of the RETLW nn instructions that returns the value, 'nn', to the calling function.

The offset value (in WREG) specifies the number of bytes that the PC should advance and should be multiples of 2 (LSb = 0).

In this method, only one byte may be stored in each instruction location, but room on the return address stack is required.

EXAMPLE 6-2: COMPUTED GOTO USING AN OFFSET VALUE

	MOVF	OFFSET,	W
	CALL	TABLE	
ORG	nn00h		
TABLE	ADDWF	PCL	
	RETLW	nnh	
	RETLW	nnh	
	RETLW	nnh	
	•		
	•		

6.1.6.2 Table Reads

A better method of storing data in program memory allows two bytes to be stored in each instruction location.

Look-up table data may be stored two bytes per program word while programming. The Table Pointer (TBLPTR) specifies the byte address, and the Table Latch (TABLAT) contains the data that is read from the program memory. Data is transferred from program memory, one byte at a time.

Table read operation is discussed further in **Section 7.1 "Table Reads and Table Writes**".

6.2 PIC18 Instruction Cycle

6.2.1 CLOCKING SCHEME

The microcontroller clock input, whether from an internal or external source, is internally divided by '4' to generate four non-overlapping quadrature clocks (Q1, Q2, Q3 and Q4). Internally, the PC is incremented on every Q1; the instruction is fetched from the program memory and latched into the Instruction Register (IR) during Q4. The instruction is decoded and executed during the following Q1 through Q4. Figure 6-4 illustrates the clocks and instruction execution flow.

6.2.2 INSTRUCTION FLOW/PIPELINING

An "Instruction Cycle" consists of four Q cycles, Q1 through Q4. The instruction fetch and execute are pipelined in such a manner that a fetch takes one instruction cycle, while the decode and execute take another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the PC to change (e.g., GOTO), then two cycles are required to complete the instruction (Example 6-3).

A fetch cycle begins with the PC incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the IR in cycle, Q1. This instruction is then decoded and executed during the Q2, Q3 and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).



EXAMPLE 6-3: INSTRUCTION PIPELINE FLOW

		Toy 1	Toy2	Toy/2	Toy/	ToyE	
	ICYU	ICYT	ICYZ	1013	1614	1015	
1. MOVLW 55h	Fetch 1	Execute 1		_			
2. MOVWF PORTB		Fetch 2	Execute 2		_		
3. BRA SUB_1			Fetch 3	Execute 3		_	
4. BSF PORTA, BIT3 (Forced NOP) Fetch 4 Flush (NOP)							
5. Instruction @ add	5. Instruction @ address SUB_1 Fetch SUB_1 Execute SUB_1						
Note: All instructions are single-cycle, except for any program branches. These take two cycles since the fetch instruction is "flushed" from the pipeline while the new instruction is being fetched and then exe-							

fetch instru cuted.

7.3 Reading the Flash Program Memory

The TBLRD instruction is used to retrieve data from program memory and places it into data RAM. Table reads from program memory are performed one byte at a time.

The TBLPTR points to a byte address in program space. Executing TBLRD places the byte pointed to into TABLAT. In addition, the TBLPTR can be modified automatically for the next table read operation.

The internal program memory is typically organized by words. The LSb of the address selects between the high and low bytes of the word.

Figure 7-4 illustrates the interface between the internal program memory and the TABLAT.

FIGURE 7-4: READS FROM FLASH PROGRAM MEMORY



EXAMPLE 7-1: READING A FLASH PROGRAM MEMORY WORD

	MOVLW MOVWF MOVLW	CODE_ADDR_UPPER TBLPTRU CODE_ADDR_HIGH	; ;	Load TBLPTR with the base address of the word
	MOVWF	TBLPTRH		
	MOVLW	CODE_ADDR_LOW		
	MOVWF	TBLPTRL		
READ_WORD				
	TBLRD*-	+	;	read into TABLAT and increment
	MOVF	TABLAT, W	;	get data
	MOVWF	WORD_EVEN		
	TBLRD*-	÷	;	read into TABLAT and increment
	MOVF	TABLAT, W	;	get data
	MOVWF	WORD_ODD		

9.2 PIR Registers

The PIR registers contain the individual flag bits for the peripheral interrupts. Due to the number of peripheral interrupt sources, there are three Peripheral Interrupt Request (Flag) registers (PIR1, PIR2, PIR3).

- Note 1: Interrupt flag bits are set when an interrupt condition occurs, regardless of the state of its corresponding enable bit or the Global Interrupt Enable bit, GIE (INT-CON<7>).
 - 2: User software should ensure the appropriate interrupt flag bits are cleared prior to enabling an interrupt and after servicing that interrupt.

REGISTER 9-4: PIR1: PERIPHERAL INTERRUPT REQUEST (FLAG) REGISTER 1 (ACCESS F9Eh)

R/W-0	R/W-0	R-0	R-0	R/W-0	R/W-0	R/W-0	R/W-0
PMPIF ⁽¹⁾	ADIF	RC1IF	TX1IF	SSP1IF	CCP1IF	TMR2IF	TMR1IF
bit 7							bit 0

Legend:								
R = Readable	e bit	W = Writable bit	U = Unimplemented bit	, read as '0'				
-n = Value at	POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown				
bit 7	PMPIF: F	Parallel Master Port Read/Wri	ite Interrupt Flag bit ⁽¹⁾					
 1 = A read or a write operation has taken place (must be cleared in software) 0 = No read or write has occurred 								
bit 6	ADIF: A/	D Converter Interrupt Flag bit						
	1 = An A 0 = The	 A/D conversion completed (minimum) A/D conversion is not completed 	ust be cleared in software) ete					
bit 5	RC1IF: E	USART1 Receive Interrupt F	lag bit					
	1 = The 0 = The	EUSART1 receive buffer, RC EUSART1 receive buffer is e	REG1, is full (cleared when R mpty	CREG1 is read)				
bit 4	TX1IF: E	USART1 Transmit Interrupt F	lag bit					
	1 = The 0 = The	EUSART1 transmit buffer, TX EUSART1 transmit buffer is f	KREG1, is empty (cleared whe full	en TXREG1 is written)				
bit 3	SSP1IF: Master Synchronous Serial Port 1 Interrupt Flag bit							
	 1 = The transmission/reception is complete (must be cleared in software) 0 = Waiting to transmit/receive 							
bit 2	CCP1IF: ECCP1 Interrupt Flag bit							
<u>Capture mode:</u> 1 = A TMR1/TMR3 register capture occurred (must be cleared in software) 0 = No TMR1/TMR3 register capture occurred								
	<u>Compare mode:</u> 1 = A TMR1/TMR3 register compare match occurred (must be cleared in software) 0 = No TMR1/TMR3 register compare match occurred							
	<u>PWM mc</u> Unused i	<u>ode:</u> n this mode.						
bit 1	TMR2IF:	TMR2 to PR2 Match Interrup	ot Flag bit					
	1 = TMR 0 = No T	2 to PR2 match occurred (mi MR2 to PR2 match occurred	ust be cleared in software)					
bit 0	TMR1IF:	TMR1 Overflow Interrupt Fla	g bit					
	1 = TMR 0 = TMR	1 register overflowed (must k 1 register did not overflow	be cleared in software)					
Note 1: Th	ese bits ar	e unimplemented on 28-pin c	levices.					

10.0 I/O PORTS

Depending on the device selected and features enabled, there are up to five ports available. Some pins of the I/O ports are multiplexed with an alternate function from the peripheral features on the device. In general, when a peripheral is enabled, that pin may not be used as a general purpose I/O pin.

Each port has three registers for its operation. These registers are:

- TRIS register (Data Direction register)
- PORT register (reads the levels on the pins of the device)
- LAT register (Data Latch)

The Data Latch (LAT register) is useful for read-modifywrite operations on the value that the I/O pins are driving.

Figure 10-1 displays a simplified model of a generic I/O port, without the interfaces to other peripherals.

FIGURE 10-1: GENERIC I/O PORT OPERATION



10.1 I/O Port Pin Capabilities

When developing an application, the capabilities of the port pins must be considered. Outputs on some pins have higher output drive strength than others. Similarly, some pins can tolerate higher than VDD input levels.

10.1.1 PIN OUTPUT DRIVE

The output pin drive strengths vary for groups of pins intended to meet the needs for a variety of applications. PORTB and PORTC are designed to drive higher loads, such as LEDs. All other ports are designed for small loads, typically indication only. Table 10-1 summarizes the output capabilities. Refer to **Section 31.0 "Electrical Characteristics"** for more details.

TABLE 10-1:	OUTPUT DRIVE LEVELS

Port	Drive	Description		
PORTA				
PORTD	Minimum	Intended for indication.		
PORTE				
PORTB	High	Suitable for direct LED drive		
PORTC	підп	levels.		

10.1.2 INPUT PINS AND VOLTAGE CONSIDERATIONS

The voltage tolerance of pins used as device inputs is dependent on the pin's input function. Pins that are used as digital only inputs are able to handle DC voltages up to 5.5V; a level typical for digital logic circuits. In contrast, pins that also have analog input functions of any kind can only tolerate voltages up to VDD. Voltage excursions beyond VDD on these pins should be avoided. Table 10-2 summarizes the input capabilities. Refer to **Section 31.0 "Electrical Characteristics"** for more details.

TABLE 10-2: INPUT VOLTAGE LEVELS

Port or Pin Tolerated Input		Description
PORTA<7:0>		
PORTB<3:0>	Vpp	Only VDD input levels tolerated.
PORTC<2:0>	VDD	
PORTE<2:0>		
PORTB<7:4>		Tolerates input levels
PORTC<7:6>	5.5V	above VDD, useful for
PORTD<7:0>		most standard logic.
PORTC<5:4>	(USB)	Designed for USB specifications.

REGISTER 10-4: PADCFG1: PAD CONFIGURATION CONTROL REGISTER 1 (BANKED F3Ch)

U-0	U-0	U-0	U-0	U-0	R/W-0	R/W-0	R/W-0
—	_	—	_	_	RTSECSEL1 ⁽¹⁾	RTSECSEL0 ⁽¹⁾	PMPTTL
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read a	as 'O'
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 7-3	Unimplemented: Read as '0'
bit 2-1	RTSECSEL<1:0>: RTCC Seconds Clock Output Select bits ⁽¹⁾
	 11 = Reserved; do not use 10 = RTCC source clock is selected for the RTCC pin (can be INTRC, T1OSC or T1CKI depending upon the RTCOSC (CONFIG3L<1>) and T1OSCEN (T1CON<3>) bit settings) 01 = RTCC seconds clock is selected for the RTCC pin 00 = RTCC alarm pulse is selected for the RTCC pin
bit 0	PMPTTL: PMP Module TTL Input Buffer Select bit 1 = PMP module uses TTL input buffers 0 = PMP module uses Schmitt Trigger input buffers

Note 1: To enable the actual RTCC output, the RTCOE (RTCCFG<2>) bit needs to be set.

10.2 PORTA, TRISA and LATA Registers

PORTA is a 7-bit wide, bidirectional port. It may function as a 5-bit port, depending on the oscillator mode selected. Setting a TRISA bit (= 1) will make the corresponding PORTA pin an input (i.e., put the corresponding output driver in a High-impedance mode). Clearing a TRISA bit (= 0) will make the corresponding PORTA pin an output (i.e., put the contents of the output latch on the selected pin).

Reading the PORTA register reads the status of the pins, whereas writing to it, will write to the port latch.

The Data Latch (LATA) register is also memory mapped. Read-modify-write operations on the LATA register read and write the latched output value for PORTA.

The other PORTA pins are multiplexed with analog inputs, the analog VREF+ and VREF- inputs and the comparator voltage reference output. The operation of pins, RA<3:0> and RA5, as A/D Converter inputs is selected by clearing or setting the control bits in the ADCON0 register (A/D Port Configuration Register 0).

Pins, RA0, RA2, and RA3, may also be used as comparator inputs and by setting the appropriate bits in the CMCON register. To use RA<3:0> as digital inputs, it is also necessary to turn off the comparators.

Note: On a Power-on Reset (POR), RA5 and RA<3:0> are configured as analog inputs and read as '0'.

All PORTA pins have TTL input levels and full CMOS output drivers.

The TRISA register controls the direction of the PORTA pins, even when they are being used as analog inputs. The user must ensure the bits in the TRISA register are maintained set when using them as analog inputs.

EXAMPLE 10-2: INITIALIZING PORTA

CLRF	PORTA	; Initialize PORTA by
		; clearing output
		; data latches
CLRF	LATA	; Alternate method
		; to clear output
		; data latches
MOVLW	07h	; Configure A/D
MOVWF	ADCON0	; for digital inputs
MOVWF	07h	; Configure comparators
MOVWF	CMCON	; for digital input
MOVLW	0CFh	; Value used to
		; initialize data
		; direction
MOVWF	TRISA	; Set RA<3:0> as inputs
		; RA<5:4> as outputs

NOTES:

REGISTER 20-2: SSPxCON1: MSSPx CONTROL REGISTER 1 (SPI MODE) (1, ACCESS FC6h; 2, F72h)

R/W_0	R///_0	R/\\/_0	R/M/-0	R/\/_0	R/M/_0	R/M_0	R/\//_0
		SSDEN(2)		<u>SCDM3(3)</u>	SSDM2(3)	SSDM1(3)	SSDM0(3)
hit 7	33FUV.7	33FEN. /	UKF	33F1013*7	33FIVIZ	33FINITY	bit 0
							bit 0
Legend:							
R = Reada	able bit	W = Writable I	oit	U = Unimpler	nented bit, read	d as '0'	
-n = Value	at POR	'1' = Bit is set		'0' = Bit is clea	ared	x = Bit is unkr	nown
bit 7	WCOL: Write	Collision Detec	t bit				
	1 = The SSP	xBUF register i	s written while	e it is still transm	itting the previ	ous word (mus	t be cleared in
	software)						
		on					
bit 6	SSPOV: Rece	eive Overflow Ir	dicator bit()				
	<u>SPI Slave mo</u>	<u>de:</u> to in received w	hila tha SSDvl	21 IE register in a	till bolding the	araviava data. I	n anna af avar
	flow the	data in SSPxSF	R is lost Over	flow can only or	cur in Slave n	node The user	must read the
	SSPxBU	, even if only tr	ansmitting dat	a, to avoid settir	ng overflow (mu	ist be cleared in	n software).
	0 = No overfl	ow					
bit 5	SSPEN: Mast	er Synchronou	s Serial Port E	Enable bit ⁽²⁾			
	1 = Enables s 0 = Disables s	erial port and c serial port and c	onfigures SCI configures the	<pre><x, as="" i="" o="" p<="" pins="" pre="" sdix="" sdox,="" se=""></x,></pre>	and SSx as so ort pins	erial port pins	
bit 4	CKP: Clock P	olarity Select b	it		·		
	1 = Idle state	for clock is a hi	gh level				
	0 = Idle state	for clock is a lo	w level				
bit 3-0	SSPM<3:0>:	Master Synchro	onous Serial F	Port Mode Selec	t bits ⁽³⁾		
	0101 = SPI S	lave mode, clo	ck = SCKx pin	i; <u>SSx</u> pin contro	ol disabled, SS	x can be used	as I/O pin
	0100 = SPI S	lave mode, clo	ck = SCKx pin	i; SSx pin contro	ol enabled		
	0011 = SPIN	laster mode, cli laster mode, cli	DCK = IMRZ 0	utput/2			
	0001 = SPI M	laster mode, cl	bck = Fosc/16	5			
	1010 = SPI M	laster mode, cl	ock = Fosc/8				
	0000 = SPI M	laster mode, clo	ock = Fosc/4				
Note 1:	In Master mode, t	he overflow bit	is not set sind	e each new rec	eption (and tra	nsmission) is ir	nitiated by
	writing to the SSP	XBUF register.			. 、	,	2

- 2: When enabled, this pin must be properly configured as input or output.
- 3: Bit combinations not specifically listed here are either reserved or implemented in I²C mode only.

20.5 I²C Mode

The MSSP module in I²C mode fully implements all master and slave functions (including general call support), and provides interrupts on Start and Stop bits in hardware to determine a free bus (multi-master function). The MSSP module implements the standard mode specifications and 7-bit and 10-bit addressing.

Two pins are used for data transfer:

- Serial Clock (SCLx) RB4/CCP4/PMA1/KBI0/SCK1/SCL1/RP7 or RD0/PMD0/SCL2
- Serial Data (SDAx) RB5/CCP5/PMA0/KBI1/SDI1/SDA1/RP8 or RD1/PMD1/SDA2

The user must configure these pins as inputs by setting the associated TRIS bits. These pins are up to 5.5V tolerant, allowing direct use in I²C buses operating at voltages higher than VDD.

FIGURE 20-7: MSSPx BLOCK DIAGRAM (I²C MODE)



20.5.1 REGISTERS

The MSSP module has six registers for $\mathsf{I}^2\mathsf{C}$ operation. These are:

- MSSPx Control Register 1 (SSPxCON1)
- MSSPx Control Register 2 (SSPxCON2)
- MSSPx Status Register (SSPxSTAT)
- Serial Receive/Transmit Buffer Register (SSPxBUF)
- MSSPx Shift Register (SSPxSR) Not directly accessible
- MSSPx Address Register (SSPxADD)
- MSSPx 7-Bit Address Mask Register (SSPxMSK)

SSPxCON1, SSPxCON2 and SSPxSTAT are the control and status registers in I^2 C mode operation. The SSPxCON1 and SSPxCON2 registers are readable and writable. The lower six bits of the SSPxSTAT are read-only. The upper two bits of the SSPxSTAT are read/write.

SSPxSR is the shift register used for shifting data in or out. SSPxBUF is the buffer register to which data bytes are written to or read from.

SSPxADD contains the slave device address when the MSSP is configured in I²C Slave mode. When the MSSP is configured in Master mode, the lower seven bits of SSPxADD act as the Baud Rate Generator (BRG) reload value.

SSPxMSK holds the slave address mask value when the module is configured for 7-Bit Address Masking mode. While it is a separate register, it shares the same SFR address as SSPxADD; it is only accessible when the SSPM<3:0> bits are specifically set to permit access. Additional details are provided in Section 20.5.3.4 "7-Bit Address Masking Mode".

In receive operations, SSPxSR and SSPxBUF together, create a double-buffered receiver. When SSPxSR receives a complete byte, it is transferred to SSPxBUF and the SSPxIF interrupt is set.

During transmission, the SSPxBUF is not double-buffered. A write to SSPxBUF will write to both SSPxBUF and SSPxSR.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	
INTCON	GIE/GIEH	PEIE/GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF	
PIR1	PMPIF ⁽³⁾	ADIF	RC1IF	TX1IF	SSP1IF	CCP1IF	TMR2IF	TMR1IF	
PIE1	PMPIE ⁽³⁾	ADIE	RC1IE	TX1IE	SSP1IE	CCP1IE	TMR2IE	TMR1IE	
IPR1	PMPIP ⁽³⁾	ADIP	RC1IP	TX1IP	SSP1IP	CCP1IP	TMR2IP	TMR1IP	
PIR2	OSCFIF	CM2IF	CM1IF	USBIF	BCL1IF	HLVDIF	TMR3IF	CCP2IF	
PIE2	OSCFIE	CM2IE	CM1IE	USBIE	BCL1IE	HLVDIE	TMR3IE	CCP2IE	
IPR2	OSCFIP	CM2IP	CM1IP	USBIP	BCL1IP	HLVDIP	TMR3IP	CCP2IP	
PIR3	SSP2IF	BCL2IF	RC2IF	TX2IF	TMR4IF	CTMUIF	TMR3GIF	RTCCIF	
PIE3	SSP2IE	BCL2IE	RC2IE	TX2IE	TMR4IE	CTMUIE	TMR3GIE	RTCCIE	
IPR3	SSP2IP	BCL2IP	RC2IP	TX2IP	TMR4IP	CTMUIP	TMR3GIP	RTCCIP	
TRISC	TRISC7	TRISC6	—	—	—	TRISC2	TRISC1	TRISC0	
TRISB	TRISB7	TRISB6	TRISB5	TRISB4	TRISB3	TRISB2	TRISB1	TRISB0	
SSP1BUF	MSSP1 Rece	eive Buffer/Tra	nsmit Registe	r					
SSP1ADD	MSSP1 Addr	ess Register (I	² C Slave mod	e), MSSP1 Ba	ud Rate Reloa	ad Register (I ²	C Master mod	e)	
SSPxMSK ⁽¹⁾	MSK7	MSK6	MSK5	MSK4	MSK3	MSK2	MSK1	MSK0	
SSPxCON1	WCOL	SSPOV	SSPEN	CKP	SSPM3	SSPM2	SSPM1	SSPM0	
SSPxCON2	GCEN	ACKSTAT	ACKDT	ACKEN	RCEN	PEN	RSEN	SEN	
	GCEN	ACKSTAT	ADMSK5 ⁽²⁾	ADMSK4 ⁽²⁾	ADMSK3 ⁽²⁾	ADMSK2 ⁽²⁾	ADMSK1 ⁽²⁾	SEN	
SSPxSTAT	SMP	SMP CKE D/Ā P S R/W UA BF							
SSP2BUF	MSSP2 Receive Buffer/Transmit Register								
SSP2ADD	MSSP2 Addr	ess Register (I ² C Slave mo	de), MSSP2 E	Baud Rate Rel	oad Register	(I ² C Master m	node)	

TABLE 20-4: REGISTERS ASSOCIATED WITH I²C OPERATION

Legend: — = unimplemented, read as '0'. Shaded cells are not used by the MSSPx module in I²C mode.

Note 1: SSPxMSK shares the same address in SFR space as SSPxADD, but is only accessible in certain I²C Slave mode operations in 7-Bit Masking mode. See Section 20.5.3.4 "7-Bit Address Masking Mode" for more details.

2: Alternate bit definitions for use in I²C Slave mode operations only.

3: These bits are only available on 44-pin devices.

24.7 Comparator Operation During Sleep

When a comparator is active and the device is placed in Sleep mode, the comparator remains active and the interrupt is functional, if enabled. This interrupt will wake-up the device from Sleep mode when enabled. Each operational comparator will consume additional current. To minimize power consumption while in Sleep mode, turn off the comparators (CON = 0) before entering Sleep. If the device wakes up from Sleep, the contents of the CMxCON register are not affected.

24.8 Effects of a Reset

A device Reset forces the CMxCON registers to their Reset state. This forces both comparators and the voltage reference to the OFF state.

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0
INTCON	GIE/GIEH	PEIE/GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF
PIR2	OSCFIF	CM2IF	CM1IF	USBIF	BCL1IF	HLVDIF	TMR3IF	CCP2IF
PIR5	—	—	CM3IF	TMR8IF	TMR6IF	TMR5IF	TMR5GIF	TMR1GIF
PIE2	OSCFIE	CM2IE	CM1IE	USBIE	BCL1IE	HLVDIE	TMR3IE	CCP2IE
PIE5	_	—	CM3IE	TMR8IE	TMR6IE	TMR5IE	TMR5GIE	TMR1GIE
IPR2	OSCFIP	CM2IP	CM1IP	USBIP	BCL1IP	HLVDIP	TMR3IP	CCP2IP
IPR5	—	—	CM3IP	TMR8IP	TMR6IP	TMR5IP	TMR5GIP	TMR1GIP
CMxCON	CON	COE	CPOL	EVPOL1	EVPOL0	CREF	CCH1	CCH0
CVRCON	CVREN	CVROE	CVRR	CVRSS	CVR3	CVR2	CVR1	CVR0
CMSTAT	_	—	—	—	_	COUT3	COUT2	COUT1
ANCON0	PCFG7 ^{(Leg} end:)	PCFG6 ^{(Leg-} end:)	PCFG5 ^{(Le} gend:)	PCFG4	PCFG3	PCFG2	PCFG1	PCFG0
ANCON1	VBGEN		_	PCFG12	PCFG11	PCFG10	PCFG9	PCFG8
TRISA	TRISA7	TRISA6	TRISA5	_	TRISA3	TRISA2	TRISA1	TRISA0

TABLE 24-3: REGISTERS ASSOCIATED WITH COMPARATOR MODULE

Legend: — = unimplemented, read as '0'. Shaded cells are not related to comparator operation.

The CTMU current source may be trimmed with the trim bits in CTMUICON using an iterative process to get an exact desired current. Alternatively, the nominal value without adjustment may be used; it may be stored by the software for use in all subsequent capacitive or time measurements.

To calculate the value for *RCAL*, the nominal current must be chosen and then the resistance can be calculated. For example, if the A/D Converter reference voltage is 3.3V, use 70% of full scale or 2.31V as the desired approximate voltage to be read by the A/D Converter. If the range of the CTMU current source is selected to be 0.55 μ A, the resistor value needed is calculated as *RCAL* = 2.31V/0.55 μ A, for a value of 4.2 M Ω . Similarly, if the current source is chosen to be 5.5 μ A, *RCAL* would be 420,000 Ω , and 42,000 Ω if the current source is source is set to 55 μ A.

FIGURE 27-2: CTMU CURRENT SOURCE CALIBRATION CIRCUIT



A value of 70% of full-scale voltage is chosen to make sure that the A/D Converter was in a range that is well above the noise floor. Keep in mind that if an exact current is chosen that is to incorporate the trimming bits from CTMUICON, the resistor value of RCAL may need to be adjusted accordingly. RCAL may also be adjusted to allow for available resistor values. RCAL should be of the highest precision available, keeping in mind the amount of precision needed for the circuit that the CTMU will be used to measure. A recommended minimum would be 0.1% tolerance.

The following examples show one typical method for performing a CTMU current calibration. Example 27-1 demonstrates how to initialize the A/D Converter and the CTMU. This routine is typical for applications using both modules. Example 27-2 demonstrates one method for the actual calibration routine.

REGISTER 28-1: CONFIG1L: CONFIGURATION REGISTER 1 LOW (BYTE ADDRESS 300000h)

R/WO-1	R/WO-1	R/WO-1	R/WO-1	R/WO-1	R/WO-1	R/WO-1	R/WO-1
DEBUG	XINST	STVREN	CFGPLLEN	PLLDIV2	PLLDIV1	PLLDIV0	WDTEN
bit 7							bit 0

Legend:			
R = Readable	bit WO = Write-Once bit	U = Unimplemented bit,	read as '0'
-n = Value at F	OR '1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown
bit 7	DEBUG : Background Debugger Enabl	e bit	
	1 = Background debugger is disabled;0 = Background debugger is enabled;	RB6 and RB7 configured as RB6 and RB7 are dedicated	s general purpose I/O pins I to In-Circuit Debug
bit 6	XINST: Extended Instruction Set Enable	le bit	
	1 = Instruction set extension and Index0 = Instruction set extension and Index	xed Addressing mode are er xed Addressing mode are dia	nabled sabled
bit 5	STVREN: Stack Overflow/Underflow R	eset Enable bit	
	1 = Reset on stack overflow/underflow0 = Reset on stack overflow/underflow	ν is enabled ν is disabled	
bit 4	CFGPLLEN: PLL Enable bit 1 = 96 MHz PLL is disabled 0 = 96 MHz PLL is enabled		
bit 3-1	PLLDIV<2:0>: Oscillator Selection bits	6	
	Divider must be selected to provide a 4 111 = No divide – oscillator used direc 110 = Oscillator divided by 2 (8 MHz i 101 = Oscillator divided by 3 (12 MHz 100 = Oscillator divided by 4 (16 MHz 011 = Oscillator divided by 5 (20 MHz 010 = Oscillator divided by 6 (24 MHz 001 = Oscillator divided by 10 (40 MH 000 = Oscillator divided by 12 (48 MHz	4 MHz input into the 96 MHz ctly (4 MHz input) nput) : input) : input) : input) : input) lz input) lz input)	PLL.
bit 0	WDTEN: Watchdog Timer Enable bit 1 = WDT is enabled 0 = WDT is disabled (control is placed	on the SWDTEN bit)	

29.0 INSTRUCTION SET SUMMARY

The PIC18F47J53 family of devices incorporate the standard set of 75 PIC18 core instructions, as well as an extended set of 8 new instructions for the optimization of code that is recursive or that utilizes a software stack. The extended set is discussed later in this section.

29.1 Standard Instruction Set

The standard PIC18 MCU instruction set adds many enhancements to the previous PIC[®] MCU instruction sets, while maintaining an easy migration from these PIC MCU instruction sets. Most instructions are a single program memory word (16 bits), but there are four instructions that require two program memory locations.

Each single-word instruction is a 16-bit word divided into an opcode, which specifies the instruction type and one or more operands, which further specify the operation of the instruction.

The instruction set is highly orthogonal and is grouped into four basic categories:

- Byte-oriented operations
- Bit-oriented operations
- · Literal operations
- Control operations

The PIC18 instruction set summary in Table 29-2 lists **byte-oriented**, **bit-oriented**, **literal** and **control** operations. Table 29-1 shows the opcode field descriptions.

Most byte-oriented instructions have three operands:

- 1. The file register (specified by 'f')
- 2. The destination of the result (specified by 'd')
- 3. The accessed memory (specified by 'a')

The file register designator, 'f', specifies which file register is to be used by the instruction. The destination designator, 'd', specifies where the result of the operation is to be placed. If 'd' is zero, the result is placed in the WREG register. If 'd' is one, the result is placed in the file register specified in the instruction.

All **bit-oriented** instructions have three operands:

- 1. The file register (specified by 'f')
- 2. The bit in the file register (specified by 'b')
- 3. The accessed memory (specified by 'a')

The bit field designator, 'b', selects the number of the bit affected by the operation, while the file register designator, 'f', represents the number of the file in which the bit is located.

The **literal** instructions may use some of the following operands:

- A literal value to be loaded into a file register (specified by 'k')
- The desired FSR register to load the literal value into (specified by 'f')
- No operand required (specified by '—')

The **control** instructions may use some of the following operands:

- A program memory address (specified by 'n')
- The mode of the CALL or RETURN instructions (specified by 's')
- The mode of the table read and table write instructions (specified by 'm')
- No operand required (specified by '—')

All instructions are a single word, except for four double-word instructions. These instructions were made double-word to contain the required information in 32 bits. In the second word, the 4 MSbs are '1's. If this second word is executed as an instruction (by itself), it will execute as a NOP.

All single-word instructions are executed in a single instruction cycle, unless a conditional test is true or the program counter is changed as a result of the instruction. In these cases, the execution takes two instruction cycles with the additional instruction cycle(s) executed as a NOP.

The double-word instructions execute in two instruction cycles.

One instruction cycle consists of four oscillator periods. Thus, for an oscillator frequency of 4 MHz, the normal instruction execution time is 1 μ s. If a conditional test is true, or the program counter is changed as a result of an instruction, the instruction execution time is 2 μ s. Two-word branch instructions (if true) would take 3 μ s.

Figure 29-1 shows the general formats that the instructions can have. All examples use the convention 'nnh' to represent a hexadecimal number.

The Instruction Set Summary, shown in Table 29-2, lists the standard instructions recognized by the Microchip MPASM[™] Assembler.

Section 29.1.1 "Standard Instruction Set" provides a description of each instruction.

Mnem	onic,	Description	Cycles	16-	Bit Inst	ruction	Word	Status	Notos
Operands		Description	Cycles	MSb			LSb	Affected	Noles
LITERAL (
ADDLW	k	Add Literal and WREG	1	0000	1111	kkkk	kkkk	C, DC, Z, OV, N	
ANDLW	k	AND Literal with WREG	1	0000	1011	kkkk	kkkk	Z, N	
IORLW	k	Inclusive OR Literal with WREG	1	0000	1001	kkkk	kkkk	Z, N	
LFSR	f, k	Move literal (12-bit) 2nd word	2	1110	1110	00ff	kkkk	None	
		to FSR(f) 1st word		1111	0000	kkkk	kkkk		
MOVLB	k	Move Literal to BSR<3:0>	1	0000	0001	0000	kkkk	None	
MOVLW	k	Move Literal to WREG	1	0000	1110	kkkk	kkkk	None	
MULLW	k	Multiply Literal with WREG	1	0000	1101	kkkk	kkkk	None	
RETLW	k	Return with Literal in WREG	2	0000	1100	kkkk	kkkk	None	
SUBLW	k	Subtract WREG from Literal	1	0000	1000	kkkk	kkkk	C, DC, Z, OV, N	
XORLW	k	Exclusive OR Literal with WREG	1	0000	1010	kkkk	kkkk	Z, N	
DATA MEN	MORY ↔	PROGRAM MEMORY OPERATI	ONS					· · · · · · · · · · · · · · · · · · ·	
TBLRD*		Table Read	2	0000	0000	0000	1000	None	
TBLRD*+		Table Read with Post-Increment		0000	0000	0000	1001	None	
TBLRD*-		Table Read with Post-Decrement		0000	0000	0000	1010	None	
TBLRD+*		Table Read with Pre-Increment		0000	0000	0000	1011	None	
TBLWT*		Table Write	2	0000	0000	0000	1100	None	
TBLWT*+		Table Write with Post-Increment		0000	0000	0000	1101	None	
TBLWT*-		Table Write with Post-Decrement		0000	0000	0000	1110	None	
TBLWT+*		Table Write with Pre-Increment		0000	0000	0000	1111	None	

TABLE 29-2: PIC18F47J53 FAMILY INSTRUCTION SET (CONTINUED)

Note 1: When a PORT register is modified as a function of itself (e.g., MOVF PORTB, 1, 0), the value used will be that value present on the pins themselves. For example, if the data latch is '1' for a pin configured as input and is driven low by an external device, the data will be written back with a '0'.

2: If this instruction is executed on the TMR0 register (and, where applicable, d = 1), the prescaler will be cleared if assigned.

3: If Program Counter (PC) is modified or a conditional test is true, the instruction requires two cycles. The second cycle is executed as a NOP.

4: Some instructions are two-word instructions. The second word of these instructions will be executed as a NOP unless the first word of the instruction retrieves the information embedded in these 16 bits. This ensures that all program memory locations have a valid instruction.





TABLE 31-24: I²C BUS START/STOP BITS REQUIREMENTS (SLAVE MODE)

Param. No.	Symbol	Characteristic		Min.	Max.	Units	Conditions
90	TSU:STA	Start Condition	100 kHz mode	4700	_	ns	Only relevant for Repeated
		Setup Time	400 kHz mode	600	_		Start condition
91	THD:STA	Start Condition	100 kHz mode	4000		ns	After this period, the first
		Hold Time	400 kHz mode	600	_		clock pulse is generated
92	Tsu:sto	Stop Condition	100 kHz mode	4700	_	ns	
		Setup Time	400 kHz mode	600			
93	THD:STO	Stop Condition	100 kHz mode	4000	_	ns	
		Hold Time	400 kHz mode	600	_		

FIGURE 31-17: I²C BUS DATA TIMING



PRODUCT IDENTIFICATION SYSTEM

To order or obtain information, e.g., on pricing or delivery, refer to the factory or the listed sales office.

PART NO.	X /XX XXX Temperature Package Pattern Range	 Examples: a) PIC18F47J53-I/PT 301 = Industrial temp., TQFP package, QTP pattern #301. b) PIC18F47J53T-I/PT = Tape and reel, Industrial temp., TQFP package.
Device	PIC18F26J53, PIC18F27J53, PIC18F46J53 and PIC18F47J53 VDD range 2.15V to 3.6V PIC18LF26J53, PIC18LF27J53, PIC18LF46J53 and PIC18LF47J53 VDD range 2.0V to 3.6V	Note 1: F = Standard Voltage Range, with internal 2.5V core voltage regulator LF = Wide Voltage Range, with no internal core voltage regulator
Temperature Range	I = -40° C to $+85^{\circ}$ C (Industrial) E = -40° C to $+125^{\circ}$ C (Extended)	2: T = in tape and reel TQFP packages only.
Package	ML = QFN PT = TQFP (Thin Quad Flatpack) SO = SOIC SP = Skinny Plastic DIP SS = SSOP	
Pattern	QTP, SQTP, Code or Special Requirements (blank otherwise)	