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"Embedded - Microcontrollers" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

Applications of "<u>Embedded -</u> <u>Microcontrollers</u>"

#### Details

⊡XFI

Product Status	Obsolete
Core Processor	PIC
Core Size	16-Bit
Speed	20 MIPS
Connectivity	I <sup>2</sup> C, IrDA, LINbus, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	21
Program Memory Size	32KB (11K x 24)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	2K x 8
Voltage - Supply (Vcc/Vdd)	3V ~ 3.6V
Data Converters	A/D 10x10b/12b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 150°C (TA)
Mounting Type	Surface Mount
Package / Case	28-VQFN Exposed Pad
Supplier Device Package	28-QFN-S (6x6)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic24hj32gp202-h-mm

Email: info@E-XFL.COM

Address: Room A, 16/F, Full Win Commercial Centre, 573 Nathan Road, Mongkok, Hong Kong

# 1.0 DEVICE OVERVIEW

- Note 1: This data sheet summarizes the features of the PIC24HJ32GP202/204 and PIC24HJ16GP304 devices. It is not intended to be a comprehensive reference source. To complement the information in this data sheet, refer to the "dsPIC33F/PIC24H Family Reference Manual". Please see the Microchip web site (www.microchip.com) for the latest dsPIC33F/PIC24H Family Reference Manual sections.
  - 2: Some registers and associated bits described in this section may not be available on all devices. Refer to Section 4.0 "Memory Organization" in this data sheet for device-specific register and bit information.

This document contains device-specific information for the following devices:

- PIC24HJ32GP202
- PIC24HJ32GP204
- PIC24HJ16GP304

Figure 1-1 shows a general block diagram of the core and peripheral modules in the PIC24HJ32GP202/204 and PIC24HJ16GP304 family of devices. Table 1-1 lists the functions of the various pins shown in the pinout diagrams.

# 3.0 CPU

- Note 1: This data sheet summarizes the features of the PIC24HJ32GP202/204 and PIC24HJ16GP304 family of devices. However, it is not intended to be a comprehensive reference source. To complement the information in this data sheet, refer to "Section 2. CPU" (DS70204) of the "dsPIC33F/PIC24H Family Reference Manual", which is available from the Microchip website (www.microchip.com).
  - Some registers and associated bits described in this section may not be available on all devices. Refer to Section 4.0 "Memory Organization" in this data sheet for device-specific register and bit information.

The PIC24HJ32GP202/204 and PIC24HJ16GP304 CPU modules have a 16-bit (data) modified Harvard architecture with an enhanced instruction set and addressing modes. The CPU has a 24-bit instruction word with a variable length opcode field. The Program Counter (PC) is 23 bits wide and addresses up to 4M x 24 bits of user program memory space. The actual amount of program memory implemented varies by device. A single-cycle instruction prefetch mechanism is used to help maintain throughput and provides predictable execution. All instructions execute in a single cycle, with the exception of instructions that change the program flow, the double word move (MOV.D) instruction and the table instructions. Overhead-free, single-cycle program loop constructs are supported using the REPEAT instruction, which is interruptible at any time.

The PIC24HJ32GP202/204 and PIC24HJ16GP304 devices have sixteen, 16-bit working registers in the programmer's model. Each of the working registers can serve as a data, address or address offset register. The 16th working register (W15) operates as a software Stack Pointer (SP) for interrupts and calls.

The instruction set includes many addressing modes and is designed for optimum C compiler efficiency. For most instructions, the devices are capable of executing a data (or program data) memory read, a working register (data) read, a data memory write and a program (instruction) memory read per instruction cycle. As a result, three parameter instructions can be supported, allowing A + B = C operations to be executed in a single cycle.

A block diagram of the CPU is shown in Figure 3-1. The programmer's model for the PIC24HJ32GP202/204 and PIC24HJ16GP304 is shown in Figure 3-2.

# 3.1 Data Addressing Overview

The data space can be linearly addressed as 32K words or 64 Kbytes using an Address Generation Unit (AGU). The upper 32 Kbytes of the data space memory map can optionally be mapped into program space at any 16K program word boundary defined by the 8-bit Program Space Visibility Page register (PSVPAG). The program to data space mapping feature lets any instruction access program space as if it were data space.

The data space also includes 2 Kbytes of DMA RAM, which is primarily used for DMA data transfers, but this may be used as general purpose RAM.

# 3.2 Special MCU Features

The PIC24HJ32GP202/204 and PIC24HJ16GP304 devices feature a 17-bit by 17-bit, single-cycle multiplier. The multiplier can perform signed, unsigned and mixed-sign multiplication. Using a 17-bit by 17-bit multiplier for 16-bit by 16-bit multiplication makes mixed-sign multiplication possible.

The PIC24HJ32GP202/204 and PIC24HJ16GP304 devices support 16/16 and 32/16 integer divide operations. All divide instructions are iterative operations. They must be executed within a REPEAT loop, resulting in a total execution time of 19 instruction cycles. The divide operation can be interrupted during any of those 19 cycles without loss of data.

A multi-bit data shifter is used to perform up to a 16-bit, left or right shift in a single cycle.

# 4.2 Data Address Space

The CPU has a separate 16 bit wide data memory space. The data space is accessed using separate Address Generation Units (AGUs) for read and write operations. The data memory maps is shown in Figure 4-3.

All Effective Addresses (EAs) in the data memory space are 16 bits wide and point to the bytes within the data space. This arrangement gives a data space address range of 64 Kbytes or 32K words. The lower half of the data memory space (that is, when EA<15>=0) is used for implemented memory addresses, while the upper half (EA<15>=1) is reserved for the Program Space Visibility area (see Section 4.6.3 "Reading Data from Program Memory Using Program Space Visibility").

PIC24HJ32GP202/204 and PIC24HJ16GP304 devices implement up to 2 Kbytes of data memory. Should an EA point to a location outside of this area, an all-zero word or byte will be returned.

## 4.2.1 DATA SPACE WIDTH

The data memory space is organized in byte addressable, 16 bit wide blocks. Data is aligned in data memory and registers as 16-bit words, but all data space EAs resolve to bytes. The Least Significant Bytes (LSBs) of each word have even addresses, while the Most Significant Bytes (MSBs) have odd addresses.

### 4.2.2 DATA MEMORY ORGANIZATION AND ALIGNMENT

To maintain backward compatibility with PIC<sup>®</sup> devices and improve data space memory usage efficiency, the PIC24HJ32GP202/204 and PIC24HJ16GP304 instruction set supports both word and byte operations. As a consequence of byte accessibility, all effective address calculations are internally scaled to step through word-aligned memory. For example, the core recognizes that Post-Modified Register Indirect Addressing mode [Ws++] will result in a value of Ws + 1 for byte operations and Ws + 2 for word operations.

Data byte reads will read the complete word that contains the byte, using the LSB of any EA to determine which byte to select. The selected byte is placed onto the LSB of the data path. That is, data memory and registers are organized as two parallel byte-wide entities with shared (word) address decode, but separate write lines. Data byte writes only write to the corresponding side of the array or register that matches the byte address. All word accesses must be aligned to an even address. Misaligned word data fetches are not supported, so care must be taken when mixing byte and word operations, or when translating from 8-bit MCU code. If a misaligned read or write is attempted, an address error trap is generated. If the error occurred on a read, the instruction underway is completed. If the instruction occurred on a write, the instruction is executed but the write does not occur. In either case, a trap is then executed, allowing the system and/or user application to examine the machine state prior to execution of the address Fault.

All byte loads into any W register are loaded into the Least Significant Byte. The Most Significant Byte is not modified.

A sign-extend instruction (SE) is provided to allow users to translate 8-bit signed data to 16-bit signed values. Alternatively, for 16-bit unsigned data, user applications can clear the MSB of any W register by executing a zero-extend (ZE) instruction on the appropriate address.

### 4.2.3 SFR SPACE

The first 2 Kbytes of the Near Data Space, from 0x0000 to 0x07FF, is primarily occupied by Special Function Registers (SFRs). These are used by the PIC24HJ32GP202/204 and PIC24HJ16GP304 core and peripheral modules to control the operation of the device.

SFRs are distributed among the modules that they control, and are generally grouped together by module. Much of the SFR space contains unused addresses; these are read as '0'. A complete listing of implemented SFRs, including their addresses, is shown in Table 4-1 through Table 4-22.

Note:	The actual set of peripheral features and interrupts varies by the device. Refer to						
	the corresponding device tables and						
	pinout diagrams for device-specific						
	information.						

## 4.2.4 NEAR DATA SPACE

The 8 Kbyte area between 0x0000 and 0x1FFF is referred to as the Near Data Space. Locations in this space are directly addressable via 13-bit absolute address field within all memory direct instructions. Additionally, the whole data space is addressable using MOV instructions, which support Memory Direct Addressing mode with a 16-bit address field, or by using Indirect Addressing mode using a working register as an address pointer.

### FIGURE 4-3: DATA MEMORY MAP FOR PIC2 4HJ32GP202/204 AND PIC24HJ16GP304 DEVICES WITH 2 KB RAM



# 4.3 Program Memory Resources

Many useful resources are provided on the main product page of the Microchip web site for the devices listed in this data sheet. This product page, which can be accessed using this link, contains the latest updates and additional information.

Note:	In the event you are not able to access								
	the product page using the link above,								
	enter this URL in your browser:								
	http://www.microchip.com/wwwproducts/								
	Devices.aspx?dDocName=en530271								

# 4.3.1 KEY RESOURCES

- Section 4. "Program Memory" (DS70202)
- Code Samples
- Application Notes
- Software Libraries
- Webinars
- All related dsPIC33F/PIC24H Family Reference Manuals Sections
- Development Tools

### 4.4.1 SOFTWARE STACK

In addition to its use as a working register, the W15 register is also used as a software Stack Pointer. The Stack Pointer always points to the first available free word and grows from lower to higher addresses. It pre-decrements for stack pops and post-increments for stack pushes, as shown in Figure 4-4. For a PC push during any CALL instruction, the MSB of the PC is zero-extended before the push, ensuring that the MSB is always clear.

Note:	A PC push during exception processing
	concatenates the SRL register to the MSB
	of the PC prior to the push.

The Stack Pointer Limit register (SPLIM) associated with the Stack Pointer sets an upper address boundary for the stack. SPLIM is uninitialized at Reset. Similarly, the Stack Pointer, SPLIM<0> is forced to '0' because all stack operations must be word aligned.

When an EA is generated using W15 as a source or destination pointer, the resulting address is compared with the value in SPLIM. If the contents of the Stack Pointer (W15) and the SPLIM register are equal and a push operation is performed, a stack error trap will not occur. The stack error trap will occur on a subsequent push operation. For example, to cause a stack error trap when the stack grows beyond address 0x1000 in RAM, initialize the SPLIM with the value 0x0FFE.

Similarly, a Stack Pointer underflow (stack error) trap is generated when the Stack Pointer address is found to be lesser than 0x0800. This prevents the stack from interfering with the Special Function Register (SFR) space.

A write to the SPLIM register should not be immediately followed by an indirect read operation using W15.





## 4.4.2 DATA RAM PROTECTION FEATURE

The PIC24H product family supports Data RAM protection features that enable segments of RAM to be protected when used in conjunction with Boot and Secure Code Segment Security. BSRAM (Secure RAM segment for BS) is accessible only from the Boot Segment Flash code when enabled. SSRAM (Secure RAM segment for RAM) is accessible only from the Secure Segment Flash code when enabled. See Table 4-1 for an overview of the BSRAM and SSRAM SFRs.

## 4.5 Instruction Addressing Modes

The addressing modes shown in Table 4-23 form the basis of the addressing modes optimized to support the specific features of individual instructions. The addressing modes provided in the MAC class of instructions differ from those in the other instruction types.

## 4.5.1 FILE REGISTER INSTRUCTIONS

Most file register instructions use a 13-bit address field (f) to directly address data present in the first 8192 bytes of data memory (Near Data Space). Most file register instructions employ a working register, W0, which is denoted as WREG in these instructions. The destination is typically either the same file register or WREG (with the exception of the MUL instruction), which writes the result to a register or register pair. The MOV instruction allows additional flexibility and can access the entire data space.

## 4.5.2 MCU INSTRUCTIONS

The three-operand MCU instructions are of the form:

Operand 3 = Operand 1 <function> Operand 2 where:

Operand 1 is always a working register (that is, the addressing mode can only be register direct), which is referred to as Wb.

Operand 2 can be a W register, fetched from data memory, or a 5-bit literal. The result location can be either a W register or a data memory location. The following addressing modes are supported by MCU instructions:

- Register Direct
- Register Indirect
- Register Indirect Post-Modified
- Register Indirect Pre-Modified
- 5-bit or 10-bit Literal
- Note: Not all instructions support all the addressing modes given above. Individual instructions can support different subsets of these addressing modes.

### TABLE 4-23: FUNDAMENTAL ADDRESSING MODES SUPPORTED

Addressing Mode	Description
File Register Direct	The address of the file register is specified explicitly.
Register Direct	The contents of a register are accessed directly.
Register Indirect	The contents of Wn forms the Effective Address (EA.)
Register Indirect Post-Modified	The contents of Wn forms the EA. Wn is post-modified (incremented or decremented) by a constant value.
Register Indirect Pre-Modified	Wn is pre-modified (incremented or decremented) by a signed constant value to form the EA.
Register Indirect with Register Offset (Register Indexed)	The sum of Wn and Wb forms the EA.
Register Indirect with Literal Offset	The sum of Wn and a literal forms the EA.

## 4.5.3 MOVE (MOV) INSTRUCTION

Move instructions provide a greater degree of addressing flexibility than the other instructions. In addition to the Addressing modes supported by most MCU instructions, MOV instructions also support Register Indirect with Register Offset Addressing mode. This is also referred to as Register Indexed mode.

Note:	For the MOV instructions, the addressing mode specified in the instruction can differ						
	for the source and the destination EA.						
	However, the 4-bit Wb (Register Offset)						
	field is shared by both source and						
	destination (but typically only used by						
	one).						

In summary, move instructions support the following addressing modes:

- Register Direct
- Register Indirect
- Register Indirect Post-modified
- Register Indirect Pre-modified
- Register Indirect with Register Offset (Indexed)
- Register Indirect with Literal Offset
- 8-bit Literal
- 16-bit Literal

Note:	Not	all	instructions	support	all	the
	addr	essir	ng modes give	n above. I	ndivi	dual
	instr	uctio	ns may suppo	ort differen	t sub	sets
	of th	ese a	addressing mo	odes.		

# 4.5.4 OTHER INSTRUCTIONS

Besides the addressing modes outlined previously, some instructions use literal constants of various sizes. For example, BRA (branch) instructions use 16-bit signed literals to specify the branch destination directly, whereas the DISI instruction uses a 14-bit unsigned literal field. In some instructions, such as ADD Acc, the source of an operand or result is implied by the opcode itself. Certain operations, such as NOP, do not have any operands.

## 7.3 Interrupt Resources

Many useful resources are provided on the main product page of the Microchip web site for the devices listed in this data sheet. This product page, which can be accessed using this link, contains the latest updates and additional information.

Note:	In the event you are not able to access							
	the product page using the link above,							
	enter this URL in your browser:							
	http://www.microchip.com/wwwproducts/							
	Devices.aspx?dDocName=en530271							

#### 7.3.1 KEY RESOURCES

- Section 6. "Interrupts" (DS70184)
- Code Samples
- Application Notes
- Software Libraries
- Webinars
- All related dsPIC33F/PIC24H Family Reference Manuals Sections
- Development Tools

# 7.4 Interrupt Control and Status Registers

PIC24HJ32GP202/204 and PIC24HJ16GP304 devices implement a total of 17 registers for the interrupt controller:

- Interrupt Control Register 1 (INTCON1)
- Interrupt Control Register 2 (INTCON2)
- Interrupt Flag Status Registers (IFSx)
- Interrupt Enable Control Registers (IECx)
- Interrupt Priority Control Registers (IPCx)
- Interrupt Control and Status Register (INTTREG)

#### 7.4.1 INTCON1 AND INTCON2

Global interrupt control functions are controlled from INTCON1 and INTCON2. INTCON1 contains the Interrupt Nesting Disable bit (NSTDIS) as well as the control and status flags for the processor trap sources. The INTCON2 register controls the external interrupt request signal behavior and the use of the Alternate Interrupt Vector Table.

### 7.4.2 IFSx

The IFS registers maintain all the interrupt request flags. Each source of interrupt has a status bit, which is set by the respective peripherals or external signal and this is cleared via software.

#### 7.4.3 IECx

The IEC registers maintain all the interrupt enable bits. These control bits are used individually to enable interrupts from the peripherals or external signals.

### 7.4.4 IPCx

The IPC registers are used to set the interrupt priority level for each source of interrupt. Each user interrupt source can be assigned to one of the eight priority levels.

#### 7.4.5 INTTREG

The INTTREG register contains the associated interrupt vector number and the new CPU interrupt priority level, which are latched into vector number (VECNUM<6:0>) and Interrupt level (ILR<3:0>) bit fields in the INTTREG register. The new interrupt priority level is the priority of the pending interrupt.

The interrupt sources are assigned to the IFSx, IECx and IPCx registers in the same sequence that they are listed in Table 7-1. For example, the INT0 (External Interrupt 0) is shown as having vector number 8 and a natural order priority of 0. Thus, the INT0IF bit is found in IFS0<0>, the INT0IE bit in IEC0<0>, and the INT0IP bits in the first position of IPC0 (IPC0<2:0>).

## 7.4.6 STATUS REGISTERS

Although these are not specifically part of the interrupt control hardware, two of the CPU Control registers contain bits that control interrupt functionality:

- The CPU STATUS register, SR, contains the IPL<2:0> bits (SR<7:5>). These bits indicate the current CPU interrupt priority level. The user can change the current CPU priority level by writing to the IPL bits.
- The CORCON register contains the IPL3 bit which, together with IPL<2:0>, also indicates the current CPU priority level. IPL3 is a read-only bit, so that trap events cannot be masked by the user software.

All Interrupt registers are described in Register 7-1 through Register 7-19.

U-0	R/W-1	R/W-0	R/W-0	U-0	R/W-1	R/W-0	R/W-0		
_		T2IP<2:0>				OC2IP<2:0>			
bit 15							bit 8		
11-0	R/M-1	R/M-0	R/\/-0	11-0	11-0	11-0	11-0		
	10/00-1	IC2IP<2:0>	10/00-0		<u> </u>				
bit 7							bit 0		
Legend:									
R = Readab	le bit	W = Writable	bit	U = Unimple	mented bit, rea	id as '0'			
-n = Value a	t POR	'1' = Bit is set	t	'0' = Bit is cle	eared	x = Bit is unk	nown		
bit 15	Unimpleme	nted: Read as '	0'						
bit 14-12	T2IP<2:0>:	Timer2 Interrupt	Priority bits						
	111 = Inter	rupt is priority 7 (	highest prior	ity interrupt)					
	•								
	•								
	001 = Inter 000 = Inter	rupt is priority 1 rupt source is dis	abled						
bit 11	Unimpleme	nted: Read as '	0'						
bit 10-8	OC2IP<2:0	>: Output Compa	are Channel	2 Interrupt Prior	rity bits				
	111 = Inter	rupt is priority 7 (	highest prior	ity interrupt)					
	•								
	•								
	001 = Inter	rupt is priority 1							
	000 = Inter	rupt source is dis	abled						
bit 7	Unimpleme	nted: Read as '	0'						
bit 6-4	IC2IP<2:0>	IC2IP<2:0>: Input Capture Channel 2 Interrupt Priority bits							
	111 = Inter	rupt is priority 7 (	highest prior	ity interrupt)					
	•								
	•								
	001 = Inter	rupt is priority 1							
	000 = Inter	rupt source is dis	abled						
bit 3-0	Unimpleme	nted: Read as '	0'						

# REGISTER 7-12: IPC1: INTERRUPT PRIORITY CONTROL REGISTER 1

	5. I EEI D						
U-0	U-0	U-0	U-0	U-0	U-0	U-0	R/W-0
—		—	—	—	—	—	PLLDIV<8>
bit 15							bit 8
R/W-0	R/W-0	R/W-1	R/W-1	R/W-0	R/W-0	R/W-0	R/W-0
			PLLDI	V<7:0>			
bit 7							bit 0
Legend:							
R = Readable	bit	W = Writable	bit	U = Unimpler	mented bit, read	l as '0'	
-n = Value at F	POR	'1' = Bit is set		'0' = Bit is cleared x = Bit is unknown			
bit 15-9	Unimplement	ed: Read as '	0'				
bit 8-0	PLLDIV<8:0>	·: PLL Feedbac	ck Divisor bits	(also denoted	as 'M', PLL mul	ltiplier)	
	111111111 =	= 513					
	•						
	•						
	•						
	000110000 =	= 50 (default)					
	•						
	•						
	•	- 1					
	000000000000000000000000000000000000000	- <del>-</del> = 3					
	0000000000	= 2					

# REGISTER 8-3: PLLFBD: PLL FEEDBACK DIVISOR REGISTER <sup>(1)</sup>

Note 1: This register is reset only on a Power-on Reset (POR).

## 10.6 Peripheral Pin Select

A major challenge in general purpose devices is providing the largest possible set of peripheral features while minimizing the conflict of features on I/O pins. The challenge is even greater on low-pin count devices. In an application where more than one peripheral must be assigned to a single pin, inconvenient workarounds in application code or a complete redesign may be the only option.

Peripheral pin select configuration enables peripheral set selection and placement on a wide range of I/O pins. By increasing the pinout options available on a particular device, programmers can better tailor the microcontroller to their entire application, rather than trimming the application to fit the device.

The peripheral pin select configuration feature operates over a fixed subset of digital I/O pins. Programmers can independently map the input and/or output of most digital peripherals to any one of these I/O pins. Peripheral pin select is performed in software, and generally does not require the device to be reprogrammed. Hardware safeguards are included that prevent accidental or spurious changes to the peripheral mapping, once it has been established.

## 10.6.1 AVAILABLE PINS

The peripheral pin select feature is used with a range of up to 26 pins. The number of available pins depends on the particular device and its pin count. Pins that support the peripheral pin select feature include the designation "RPn" in their full pin designation, where "RP" designates a remappable peripheral and "n" is the remappable pin number.

### 10.6.2 CONTROLLING PERIPHERAL PIN SELECT

Peripheral pin select features are controlled through two sets of special function registers to map peripherals and to map outputs.

Since they are separately controlled, a particular peripheral's input and output (if the peripheral has both) can be placed on any selectable function pin without constraint.

The association of a peripheral to a peripheral selectable pin is handled in two different ways, depending on whether an input or output is being mapped.

## 10.6.2.1 Input Mapping

The inputs of the peripheral pin select options are mapped on the basis of the peripheral. A control register associated with a peripheral dictates the pin it will be mapped to. The RPINRx registers are used to configure peripheral input mapping (see Register 10-1 through Register 10-9). Each register contains sets of 5-bit fields, with each set associated with one of the remappable peripherals. Programming a given peripheral's bit field with an appropriate 5-bit value maps the RPn pin with that value to that peripheral. For any given device, the valid range of values for any bit field corresponds to the maximum number of peripheral pin selections supported by the device.

Figure 10-2 Illustrates remappable pin selection for U1RX input.

For input mapping only, the Peripheral Pin
Select (PPS) functionality does not have
priority over the TRISx settings. There-
fore, when configuring the RPn pin for
input, the corresponding bit in the TRISx
register must also be configured for input
(i.e., set to '1').

### FIGURE 10-2: REMAPPABLE MUX INPUT FOR U1RX



U-0	U-0	U-0	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1		
—		_			U1CTSR<4:	0>			
bit 15	·						bit		
U-0	U-0	U-0	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1		
—	—	—	U1RXR<4:0>						
bit 7							bit		
Legend:									
R = Readab	le bit	W = Writable I	bit	U = Unimpler	mented bit, rea	ad as '0'			
-n = Value a	t POR	'1' = Bit is set		'0' = Bit is cle	ared	x = Bit is unki	nown		
	11001 = Inpu • •	it fied to RP25							
	00001 = Inpu 00000 = Inpu	ut tied to RP1 ut tied to RP0							
bit 7-5	Unimplement	ed: Read as 'o	)'						
bit 4-0	U1RXR<4:0>	: Assign UART	1 Receive (U	11RX) to the co	orresponding F	RPn pin			
	11111 = Inpu 11001 = Inpu	11111 = Input tied to Vss 11001 = Input tied to RP25							
	•								
	•								
	•								
	00001 = Inpu 00000 = Inpu	ut tied to RP1 ut tied to RP0							

# REGISTER 10-7: RPINR18: PERIPHERAL PIN SELECT INPUT REGISTER 18

### REGISTER 10-21: RPOR11: PERIPHERAL PIN SELECT OUTPUT REGISTER 11

U-0	U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
	—	—			RP23R<4:0>		
bit 15							bit 8

U-0	U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
—	—	—			RP22R<4:0>		
bit 7							bit 0

Legend:				
R = Readable bit W = Writable bit		U = Unimplemented bit, read as '0'		
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown	

bit 15-13 Unimplemented: Read as '0'

bit 12-8 RP23R<4:0>: Peripheral Output Function is Assigned to RP23 Output Pin (see Table 10-2 for peripheral function numbers)

bit 7-5 Unimplemented: Read as '0'

bit 4-0 RP22R<4:0>: Peripheral Output Function is Assigned to RP22 Output Pin (see Table 10-2 for peripheral function numbers)

### REGISTER 10-22: RPOR12: PERIPHERAL PIN SELECT OUTPUT REGISTER 12

-n = Value at POR '1' = Bit is set		:	'0' = Bit is cleared x = Bit is unknown			nown	
R = Readable bit W = Writable bit		bit	U = Unimplemented bit, read as '0'				
Legend:							
bit 7	•						bit 0
_	—	_	RP24R<4:0>				
U-0	U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
511 15							
bit 15							bit 8
_	_	_			RP25R<4:0>	>	
U-0	U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0

bit 15-13 Unimplemented: Read as '0'

bit 12-8 RP25R<4:0>: Peripheral Output Function is Assigned to RP25 Output Pin (see Table 10-2 for peripheral function numbers)

bit 7-5 Unimplemented: Read as '0'

bit 4-0 RP24R<4:0>: Peripheral Output Function is Assigned to RP24 Output Pin (see Table 10-2 for peripheral function numbers)

# 11.0 TIMER1

- Note 1: This data sheet summarizes the features of the PIC24HJ32GP202/204 and PIC24HJ16GP304 family of devices. It is not intended to be a comprehensive reference source. To complement the information in this data sheet, refer to "Section 11. Timers" (DS70205) of the "dsPIC33F/PIC24H Family Reference Manual", which is available from the Microchip website (www.microchip.com).
  - 2: Some registers and associated bits described in this section may not be available on all devices. Refer to Section 4.0 "Memory Organization" in this data sheet for device-specific register and bit information.

The Timer1 module is a 16-bit timer, which can serve as the time counter for the real-time clock, or operate as a free-running interval timer/counter. Timer1 can operate in three modes:

- 16-bit Timer
- 16-bit Synchronous Counter
- 16-bit Asynchronous Counter

Timer1 also supports these features:

- Timer gate operation
- · Selectable prescaler settings
- Timer operation during CPU Idle and Sleep modes
- Interrupt on 16-bit Period register match or falling edge of external gate signal

Figure 11-1 shows a block diagram of the 16-bit timer module.

To configure Timer1 for operation:

- 1. Set the TON bit (= 1) in the T1CON register.
- 2. Select the timer prescaler ratio using the TCKPS<1:0> bits in the T1CON register.
- 3. Set the Clock and Gating modes using the TCS and TGATE bits in the T1CON register.
- 4. Set or clear the TSYNC bit in T1CON to select synchronous or asynchronous operation.
- 5. Load the timer period value into the PR1 register.
- 6. If interrupts are required, set the interrupt enable bit, T1IE. Use the priority bits, T1IP<2:0>, to set the interrupt priority.



## 12.2 Timer2/3 Resources

Many useful resources are provided on the main product page of the Microchip web site for the devices listed in this data sheet. This product page, which can be accessed using this link, contains the latest updates and additional information.

Note:	In the event you are not able to access
	the product page using the link above,
	enter this URL in your browser:
	http://www.microchip.com/wwwproducts/
	Devices.aspx?dDocName=en530271

## 12.2.1 KEY RESOURCES

- Section 11. "Timers" (DS70205)
- Code Samples
- Application Notes
- Software Libraries
- Webinars
- All related dsPIC33F/PIC24H Family Reference Manuals Sections
- Development Tools

## 13.1 Input Capture Resources

Many useful resources are provided on the main product page of the Microchip web site for the devices listed in this data sheet. This product page, which can be accessed using this link, contains the latest updates and additional information.

Note:	In the event you are not able to access
	the product page using the link above,
	enter this URL in your browser:
	http://www.microchip.com/wwwproducts/
	Devices.aspx?dDocName=en530271

#### 13.1.1 KEY RESOURCES

- Section 12. "Input Capture" (DS70198)
- Code Samples
- Application Notes
- Software Libraries
- Webinars
- All related dsPIC33F/PIC24H Family Reference Manuals Sections
- Development Tools

# 15.3 SPI Control Registers

# REGISTER 15-1: SPIx STAT: SPIx STATUS AND CONTROL REGISTER

R/W-0	U-0	R/W-0	U-0	U-0	U-0	U-0	U-0			
SPIEN	—	SPISIDL	—	—	—	—	—			
bit 15							bit 8			
U-0	R/C-0	U-0	U-0	U-0	U-0	R-0	R-0			
	SPIROV	—	—	—		SPITBF	SPIRBF			
bit 7							bit 0			
Г										
Legend:		C = Clearable	bit							
R = Readable	e bit	W = Writable	bit	U = Unimpler	mented bit, read	as '0'				
-n = Value at	POR	'1' = Bit is set		'0' = Bit is cle	ared	x = Bit is unkr	nown			
		En alta la it								
DIT 15	SPIEN: SPIX	Enable bit	figures CCK			ol port pipo				
	1 = Enables r 0 = Disables I	module and cor module	ingures SCKX	, SDOX, SDIX (	and SSX as sen	ai port pins				
bit 14	Unimplement	ed: Read as '	)'							
bit 13	SPISIDL: Sto	p in Idle Mode	bit							
	1 = Discontine	ue module ope	ration when d	evice enters Id	lle mode					
	0 = Continue	module operati	on in Idle mo	de						
bit 12-7	Unimplement	ed: Read as '	)'							
bit 6	SPIROV: Rec	eive Overflow	Flag bit							
	I = A new byte/word is completely received and discarded. The user software has not read the previous data in the SPIxBUF register									
	0 = No overflow has occurred.									
bit 5-2	Unimplement	ed: Read as '	)'							
bit 1	SPITBF: SPI	x Transmit Buff	er Full Status	bit						
	1 = Transmit	Transmit not yet started, SPIxTXB is full								
	0 = Transmit started, SPIxTXB is empty									
	Automatically set in hardware when CPU writes SPIXBUF location, loading SPIXTXB.									
hit 0		x Receive Buff	ar Full Status	hit						
bit 0	1 - Decoive complete SDIVDYB is full									
	0 = Receive is	s not complete,	SPIxRXB is (	empty						
	Automatically	set in hardwar	e when SPIx	transfers data f	from SPIxSR to	SPIxRXB.				
	Automatically	cleared in hard	dware when c	ore reads SPIx	BUF location, r	eading SPIxRX	(В.			

### REGISTER 15-2: SPIxCON1: SPIx CONTROL REGISTER 1 (CONTINUED)

- bit 4-2 SPRE<2:0>: Secondary Prescale bits (Master mode)<sup>(3)</sup> 111 = Secondary prescale 1:1 110 = Secondary prescale 2:1 • • • 000 = Secondary prescale 8:1 bit 1-0 PPRE<1:0>: Primary Prescale bits (Master mode)<sup>(3)</sup> 11 = Primary prescale 1:1 10 = Primary prescale 4:1
  - 01 = Primary prescale 16:1
  - 00 = Primary prescale 64:1
- Note 1: The CKE bit is not used in the Framed SPI modes. Program this bit to '0' for the Framed SPI modes (FRMEN = 1).
  - 2: This bit must be cleared when FRMEN = 1.
  - 3: Do not set both Primary and Secondary prescalers to a value of 1:1.

# 17.1 UART Helpful Tips

- 1. In multi-node direct-connect UART networks, receive inputs UART react to the complementary logic level defined by the URXINV bit (UxMODE<4>), which defines the idle state, the default of which is logic high, (i.e., URXINV = 0). Because remote devices do not initialize at the same time, it is likely that one of the devices, because the RX line is floating, will trigger a start bit detection and will cause the first byte received after the device has been initialized to be invalid. To avoid this situation, the user should use a pull-up or pull-down resistor on the RX pin depending on the value of the URXINV bit.
  - a) If URXINV = 0, use a pull-up resistor on the RX pin.
  - b) If URXINV = 1, use a pull-down resistor on the RX pin.
- 2. The first character received on a wake-up from Sleep mode caused by activity on the UxRX pin of the UART module will be invalid. In Sleep mode, peripheral clocks are disabled. By the time the oscillator system has restarted and stabilized from Sleep mode, the baud rate bit sampling clock relative to the incoming UxRX bit timing is no longer synchronized, resulting in the first character being invalid. This is to be expected.

## 17.2 UART Resources

Many useful resources are provided on the main product page of the Microchip web site for the devices listed in this data sheet. This product page, which can be accessed using this link, contains the latest updates and additional information.

Note: In the event you are not able to access the product page using the link above, enter this URL in your browser: http://www.microchip.com/wwwproducts/ Devices.aspx?dDocName=en530271

## 17.2.1 KEY RESOURCES

- Section 17. "UART" (DS70188)
- Code Samples
- Application Notes
- Software Libraries
- Webinars
- All related dsPIC33F/PIC24H Family Reference Manuals Sections
- Development Tools

## REGISTER 18-6: AD1CSSL: ADC1 IN PUT SCAN SELECT REGISTER LOW<sup>(1)</sup>

U-0	U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
—	—	—	CSS12	CSS11	CSS10	CSS9	CSS8
bit 15							bit 8
R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
CSS7	CSS6	CSS5	CSS4	CSS3	CSS2	CSS1	CSS0

bit 0

Legend:				
R = Readable bit W = Writable bit		U = Unimplemented bit, read as '0'		
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown	

bit 15-13 Unimplemented: Read as '0'

bit 7

bit 12-0 CSS<12:0>: ADC Input Scan Selection bits

1 =Select ANx for input scan

0 = Skip ANx for input scan

- Note 1: On devices without 13 analog inputs, all AD1CSSL bits can be selected by the user application. However, inputs selected for scan without a corresponding input on device converts VREFL.
  - 2: CSSx = ANx, where x = 0 through 12.

# REGISTER 18-7: AD1PCFGL: ADC1 PO RT CONFIGURATION REGISTER LOW<sup>(1,2,3)</sup>

U-0	U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
—	—	—	PCFG12	PCFG11	PCFG10	PCFG9	PCFG8
bit 15							bit 8
R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
PCFG7	PCFG6	PCFG5	PCFG4	PCFG3	PCFG2	PCFG1	PCFG0
bit 7							bit 0
Legend:							

R = Readable bit	W = Writable bit	U = Unimplemented bit, read	d as '0'
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 15-13 Unimplemented: Read as '0'

bit 12-0 PCFG<12:0>: ADC Port Configuration Control bits

1 = Port pin in Digital mode, port read input enabled, ADC input multiplexer connected to AVss

0 = Port pin in Analog mode, port read input disabled, ADC samples pin voltage

Note 1: On devices without 13 analog inputs, all PCFG bits are R/W by user software. However, the PCFG bits are ignored on ports without a corresponding input on device.

2: PCFGx = ANx, where x = 0 through 12.

3: The PCFGx bits have no effect if the ADC module is disabled by setting ADxMD bit in the PMDx register. In this case, all port pins multiplexed with ANx will be in Digital mode.

# FIGURE 22-13: SPIX SLAVE MODE (FULL-DUPLEX, CKE = 1, CKP = 0, SMP = 0) TIMING CHARACTERISTICS

