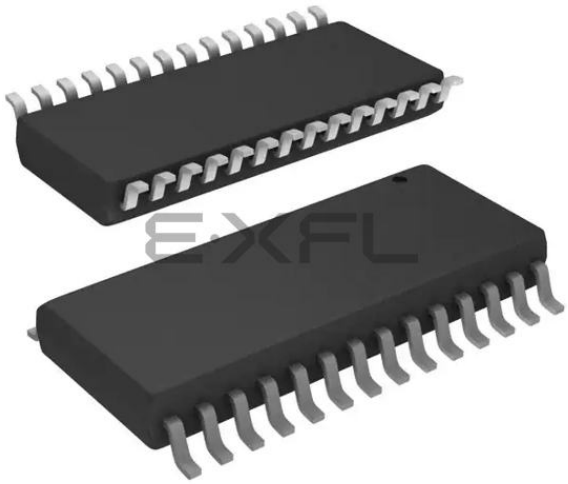

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Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	PIC
Core Size	16-Bit
Speed	32MHz
Connectivity	I²C, IrDA, LINbus, PMP, SPI, UART/USART, USB OTG
Peripherals	Brown-out Detect/Reset, DMA, LVD, POR, PWM, WDT
Number of I/O	20
Program Memory Size	128KB (43K x 24)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	8K x 8
Voltage - Supply (Vcc/Vdd)	2V ~ 3.6V
Data Converters	A/D 9x10b/12b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 125°C (TA)
Mounting Type	Surface Mount
Package / Case	28-SOIC (0.295", 7.50mm Width)
Supplier Device Package	28-SOIC
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic24fj128gb202-e-so

1.0 DEVICE OVERVIEW

This document contains device-specific information for the following devices:

- PIC24FJ64GB202
- PIC24FJ128GB202
- PIC24FJ64GB204
- PIC24FJ128GB204

The PIC24FJ128GB204 family expands the capabilities of the PIC24F family by adding a complete selection of Cryptographic Engines, ISO 7816 support and I^2S support to its existing features. This combination, along with its ultra low-power features, Direct Memory Access (DMA) for peripherals and USB On-The-Go, make this family the new standard for mixed-signal PIC® microcontrollers in one economical and power-saving package.

1.1 Core Features

1.1.1 16-BIT ARCHITECTURE

Central to all PIC24F devices is the 16-bit modified Harvard architecture, first introduced with Microchip's dsPIC® Digital Signal Controllers (DSCs). The PIC24F CPU core offers a wide range of enhancements, such as:

- 16-bit data and 24-bit address paths with the ability to move information between data and memory spaces
- Linear addressing of up to 12 Mbytes (program space) and 32 Kbytes (data)
- A 16-element Working register array with built-in software stack support
- A 17 x 17 hardware multiplier with support for integer math
- Hardware support for 32 by 16-bit division
- An instruction set that supports multiple addressing modes and is optimized for high-level languages, such as 'C'
- Operational performance up to 16 MIPS

1.1.2 XLP POWER-SAVING TECHNOLOGY

The PIC24FJ128GB204 family of devices introduces a greatly expanded range of power-saving operating modes for the ultimate in power conservation. The new modes include:

- Retention Sleep with essential circuits being powered from a separate low-voltage regulator
- Deep Sleep without RTCC for the lowest possible power consumption under software control
- V_{BAT} mode (with or without RTCC) to continue limited operation from a backup battery when V_{DD} is removed

Many of these new low-power modes also support the continuous operation of the low-power, on-chip Real-Time Clock/Calendar (RTCC), making it possible for an application to keep time while the device is otherwise asleep.

Aside from these new features, PIC24FJ128GB204 family devices also include all of the legacy power-saving features of previous PIC24F microcontrollers, such as:

- On-the-Fly Clock Switching, allowing the selection of a lower power clock during run time
- Doze Mode Operation, for maintaining peripheral clock speed while slowing the CPU clock
- Instruction-Based Power-Saving Modes, for quick invocation of Idle and the many Sleep modes

1.1.3 OSCILLATOR OPTIONS AND FEATURES

All of the devices in the PIC24FJ128GB204 family offer five different oscillator options, allowing users a range of choices in developing application hardware. These include:

- Two Crystal modes
- Two External Clock (EC) modes
- A Phase-Locked Loop (PLL) frequency multiplier, which allows clock speeds of up to 32 MHz
- A Fast Internal Oscillator (FRC) – Nominal 8 MHz output with multiple frequency divider options and automatic frequency self-calibration during run time
- A separate Low-Power Internal RC Oscillator (LPRC) – 31 kHz nominal for low-power, timing-insensitive applications.

The internal oscillator block also provides a stable reference source for the Fail-Safe Clock Monitor (FSCM). This option constantly monitors the main clock source against a reference signal provided by the internal oscillator and enables the controller to switch to the internal oscillator, allowing for continued low-speed operation or a safe application shutdown.

1.1.4 EASY MIGRATION

Regardless of the memory size, all devices share the same rich set of peripherals, allowing for a smooth migration path as applications grow and evolve. This extends the ability of applications to grow from the relatively simple, to the powerful and complex, yet still selecting a Microchip device.

4.3.2 DATA ACCESS FROM PROGRAM MEMORY USING TABLE INSTRUCTIONS

The TBLRDL and TBLWTL instructions offer a direct method of reading or writing the lower word of any address within the program space without going through Data Space. The TBLRDH and TBLWTH instructions are the only method to read or write the upper 8 bits of a program space word as data.

The PC is incremented by two for each successive 24-bit program word. This allows program memory addresses to directly map to Data Space addresses. Program memory can thus be regarded as two, 16-bit word-wide address spaces, residing side by side, each with the same address range. TBLRDL and TBLWTL access the space which contains the least significant data word, and TBLRDH and TBLWTH access the space which contains the upper data byte.

Two table instructions are provided to move byte or word-sized (16-bit) data to and from program space. Both function as either byte or word operations.

1. TBLRDL (Table Read Low): In Word mode, it maps the lower word of the program space location (P<15:0>) to a data address (D<15:0>).

   In Byte mode, either the upper or lower byte of the lower program word is mapped to the lower byte of a data address. The upper byte is selected when byte select is '1'; the lower byte is selected when it is '0'.

2. TBLRDH (Table Read High): In Word mode, it maps the entire upper word of a program address (P<23:16>) to a data address. Note that D<15:8>, the 'phantom' byte, will always be '0'.

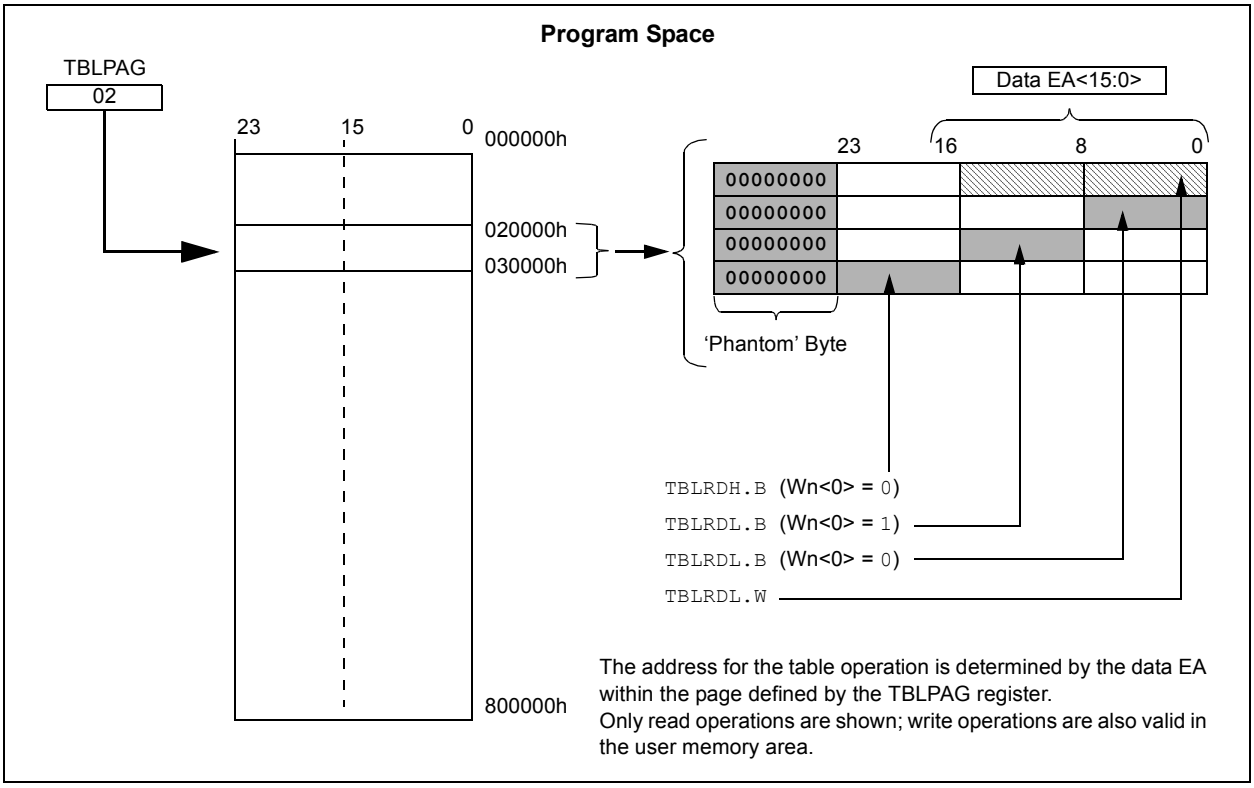
   In Byte mode, it maps the upper or lower byte of the program word to D<7:0> of the data address, as above. Note that the data will always be '0' when the upper 'phantom' byte is selected (Byte Select = 1).

In a similar fashion, two table instructions, TBLWTH and TBLWTL, are used to write individual bytes or words to a program space address. The details of their operation are described in **Section 6.0 "Flash Program Memory"**.

For all table operations, the area of program memory space to be accessed is determined by the Table Memory Page Address register (TBLPAG). TBLPAG covers the entire program memory space of the device, including user and configuration spaces. When TBLPAG<7> = 0, the table page is located in the user memory space. When TBLPAG<7> = 1, the page is located in configuration space.

Note:	Only Table Read operations will execute in the configuration memory space where Device IDs are located; Table Write operations are not allowed.

FIGURE 4-9: ACCESSING PROGRAM MEMORY WITH TABLE INSTRUCTIONS



The address for the table operation is determined by the data EA within the page defined by the TBLPAG register.
Only read operations are shown; write operations are also valid in the user memory area.

4.3.3 READING DATA FROM PROGRAM MEMORY USING EDS

The upper 32 Kbytes of Data Space may optionally be mapped into any 16K word page of the program space. This provides transparent access of stored constant data from the Data Space without the need to use special instructions (i.e., TBLRDL/H).

Program space access through the Data Space occurs when the MSb of EA is '1' and the DSRPAG<9> bit is also '1'. The lower 8 bits of DSRPAG are concatenated to the Wn<14:0> bits to form a 23-bit EA to access program memory. The DSRPAG<8> bit decides which word should be addressed; when the bit is '0', the lower word and when '1', the upper word of the program memory is accessed.

The entire program memory is divided into 512 EDS pages, from 200h to 3FFh, each consisting of 16K words of data. Pages, 200h to 2FFh, correspond to the lower words of the program memory, while 300h to 3FFh correspond to the upper words of the program memory.

Using this EDS technique, the entire program memory can be accessed. Previously, the access to the upper word of the program memory was not supported.

Table 4-37 provides the corresponding 23-bit EDS address for program memory with EDS page and source addresses.

For operations that use PSV and are executed outside a REPEAT loop, the MOV and MOV.D instructions will require one instruction cycle in addition to the specified execution time. All other instructions will require two instruction cycles in addition to the specified execution time.

For operations that use PSV, which are executed inside a REPEAT loop, there will be some instances that require two instruction cycles in addition to the specified execution time of the instruction:

- Execution in the first iteration
- Execution in the last iteration
- Execution prior to exiting the loop due to an interrupt
- Execution upon re-entering the loop after an interrupt is serviced

Any other iteration of the REPEAT loop will allow the instruction accessing data, using PSV, to execute in a single cycle.

TABLE 4-37: EDS PROGRAM ADDRESS WITH DIFFERENT PAGES AND ADDRESSES

DSRPAG (Data Space Read Register)	Source Address while Indirect Addressing	23-Bit EA Pointing to EDS	Comment
200h • • • 2FFh	8000h to FFFFh	000000h to 007FFEh • • • 7F8000h to 7FFFFEh	Lower words of 4M program instructions (8 Mbytes); for read operations only
300h • • • 3FFh		000001h to 007FFFh • • • 7F8001h to 7FFFFFh	Upper words of 4M program instructions (4 Mbytes remaining, 4 Mbytes are phantom bytes); for read operations only
000h		Invalid Address	Address error trap(1)

Note 1: When the source/destination address is above 8000h and DSRPAG/DSWPAG is '0', an address error trap will occur.

EXAMPLE 4-3: EDS READ CODE FROM PROGRAM MEMORY IN ASSEMBLY

```
; Set the EDS page from where the data to be read
    mov     #0x0202, w0
    mov     w0, DSRPAG                ;page 0x202, consisting lower words, is selected for read
    mov     #0x000A, w1               ;select the location (0x0A) to be read
    bset    w1, #15                   ;set the MSB of the base address, enable EDS mode
;Read a byte from the selected location
    mov.b   [w1++], w2                ;read Low byte
    mov.b   [w1++], w3                ;read High byte
;Read a word from the selected location
    mov     [w1], w2                  ;
;Read Double - word from the selected location
    mov.d   [w1], w2                  ;two word read, stored in w2 and w3
```

FIGURE 4-10: PROGRAM SPACE VISIBILITY OPERATION TO ACCESS LOWER WORD

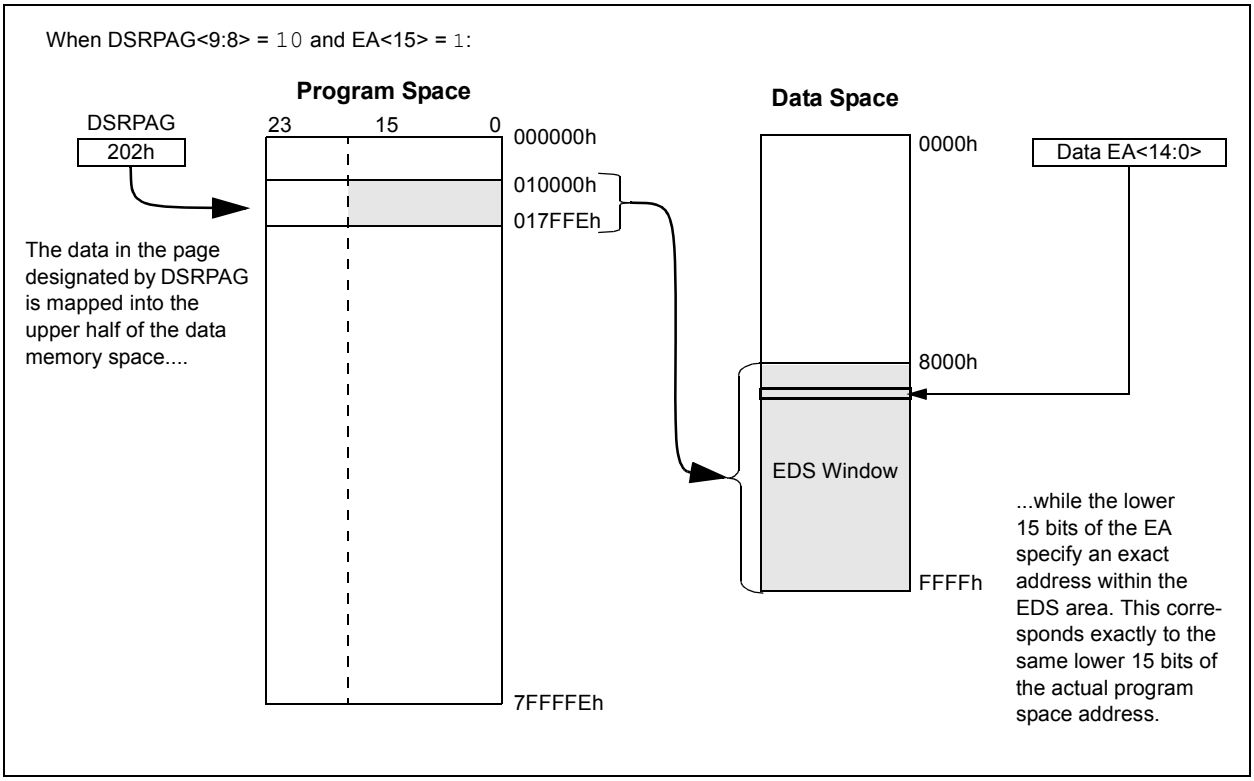
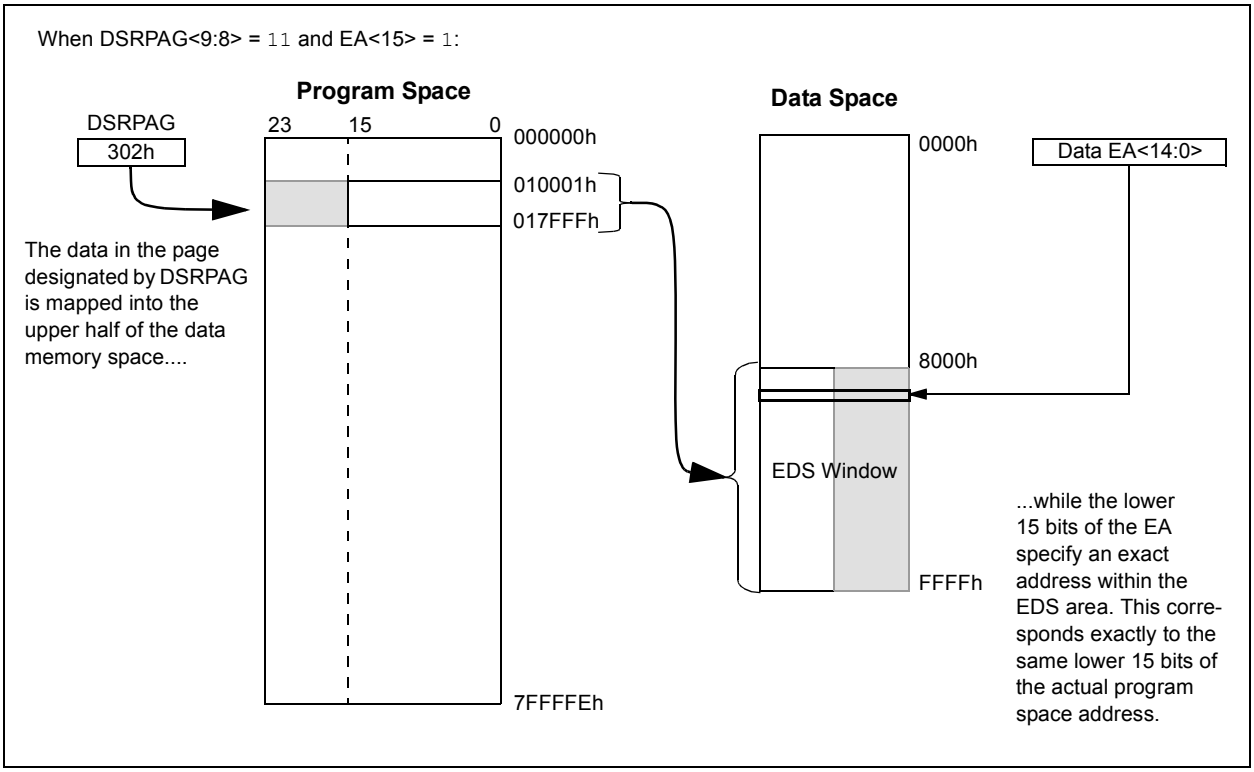


FIGURE 4-11: PROGRAM SPACE VISIBILITY OPERATION TO ACCESS UPPER WORD



REGISTER 5-2: DMACHn: DMA CHANNEL n CONTROL REGISTER

U-0	U-0	U-0	r-0	U-0	R/W-0	R/W-0	R/W-0
—	—	—	—	—	NULLW	RELOAD(1)	CHREQ(3)
bit 15							bit 8

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
SAMODE1	SAMODE0	DAMODE1	DAMODE0	TRMODE1	TRMODE0	SIZE	CHEN
bit 7							bit 0

Legend:	r = Reserved bit		
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 15-13 **Unimplemented:** Read as '0'

bit 12 **Reserved:** Maintain as '0'

bit 11 **Unimplemented:** Read as '0'

bit 10 **NULLW:** Null Write Mode bit
1 = A dummy write is initiated to DMASRCn for every write to DMADSTn
0 = No dummy write is initiated

bit 9 **RELOAD:** Address and Count Reload bit(1)
1 = DMASRCn, DMADSTn and DMACNTn registers are reloaded to their previous values upon the start of the next operation
0 = DMASRCn, DMADSTn and DMACNTn are not reloaded on the start of the next operation(2)

bit 8 **CHREQ:** DMA Channel Software Request bit(3)
1 = A DMA request is initiated by software; automatically cleared upon completion of a DMA transfer
0 = No DMA request is pending

bit 7-6 **SAMODE<1:0>:** Source Address Mode Selection bits
11 = DMASRCn is used in Peripheral Indirect Addressing and remains unchanged
10 = DMASRCn is decremented based on the SIZE bit after a transfer completion
01 = DMASRCn is incremented based on the SIZE bit after a transfer completion
00 = DMASRCn remains unchanged after a transfer completion

bit 5-4 **DAMODE<1:0>:** Destination Address Mode Selection bits
11 = DMADSTn is used in Peripheral Indirect Addressing and remains unchanged
10 = DMADSTn is decremented based on the SIZE bit after a transfer completion
01 = DMADSTn is incremented based on the SIZE bit after a transfer completion
00 = DMADSTn remains unchanged after a transfer completion

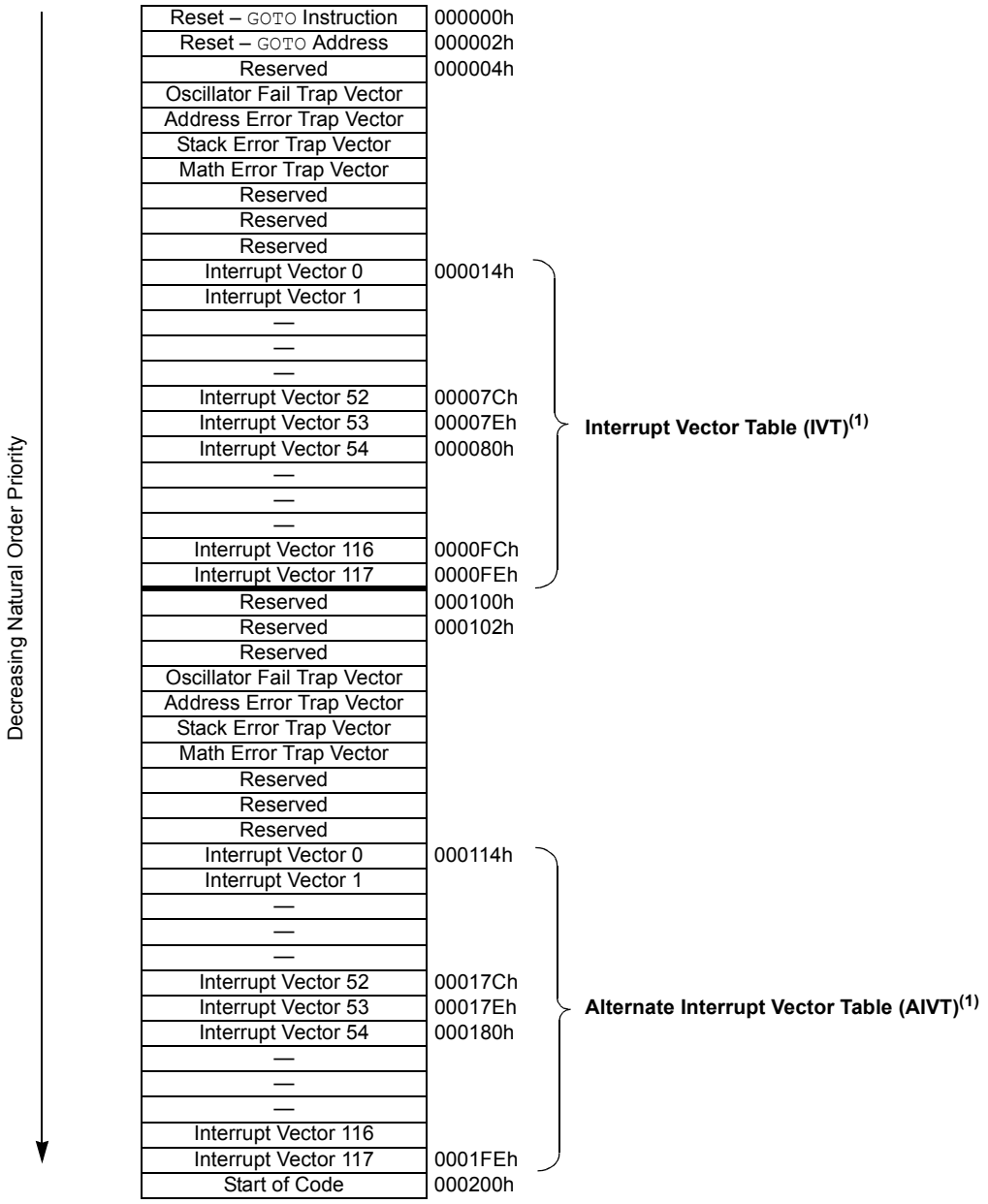
bit 3-2 **TRMODE<1:0>:** Transfer Mode Selection bits
11 = Repeated Continuous mode
10 = Continuous mode
01 = Repeated One-Shot mode
00 = One-Shot mode

bit 1 **SIZE:** Data Size Selection bit
1 = Byte (8-bit)
0 = Word (16-bit)

bit 0 **CHEN:** DMA Channel Enable bit
1 = The corresponding channel is enabled
0 = The corresponding channel is disabled

Note 1: Only the original DMACNTn is required to be stored to recover the original DMASRCn and DMADSTn values.
2: DMACNTn will always be reloaded in Repeated mode transfers, regardless of the state of the RELOAD bit.
3: The number of transfers executed while CHREQ is set depends on the configuration of TRMODE<1:0>.

FIGURE 8-1: PIC24F INTERRUPT VECTOR TABLE



Note 1: See Table 8-2 for the interrupt vector list.

TABLE 8-1: TRAP VECTOR DETAILS

Vector Number	IVT Address	AIVT Address	Trap Source
0	000004h	000104h	Reserved
1	000006h	000106h	Oscillator Failure
2	000008h	000108h	Address Error
3	00000Ah	00010Ah	Stack Error
4	00000Ch	00010Ch	Math Error
5	00000Eh	00010Eh	Reserved
6	000010h	000110h	Reserved
7	000012h	000112h	Reserved

REGISTER 8-8: IFS3: INTERRUPT FLAG STATUS REGISTER 3 (CONTINUED)

bit 5 **INT3IF:** External Interrupt 3 Flag Status bit
1 = Interrupt request has occurred
0 = Interrupt request has not occurred

bit 4-3 **Unimplemented:** Read as '0'

bit 2 **MI2C2IF:** Master I2C2 Event Interrupt Flag Status bit
1 = Interrupt request has occurred
0 = Interrupt request has not occurred

bit 1 **SI2C2IF:** Slave I2C2 Event Interrupt Flag Status bit
1 = Interrupt request has occurred
0 = Interrupt request has not occurred

bit 0 **Unimplemented:** Read as '0'

REGISTER 11-22: RPINR31: PERIPHERAL PIN SELECT INPUT REGISTER 31

U-0	U-0	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1
—	—	MDC2R5	MDC2R4	MDC2R3	MDC2R2	MDC2R1	MDC2R0
bit 15							bit 8

U-0	U-0	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1	R/W-1
—	—	MDC1R5	MDC1R4	MDC1R3	MDC1R2	MDC21R1	MDC1R0
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as ‘0’	
-n = Value at POR	‘1’ = Bit is set	‘0’ = Bit is cleared	x = Bit is unknown

bit 15-14 **Unimplemented:** Read as ‘0’

bit 13-8 **MDC2R<5:0>:** Assign TX Carrier 2 Input (MDCIN2) to Corresponding RPn or RPIn Pin bits

bit 7-6 **Unimplemented:** Read as ‘0’

bit 5-0 **MDC1R<5:0>:** Assign TX Carrier 1 Input (MDCIN1) to Corresponding RPn or RPIn Pin bits

REGISTER 11-35: RPOR12: PERIPHERAL PIN SELECT OUTPUT REGISTER 12(1)

U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
—	—	RP25R5	RP25R4	RP25R3	RP25R2	RP25R1	RP25R0
bit 15							bit 8

U-0	U-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
—	—	RP24R5	RP24R4	RP24R3	RP24R2	RP24R1	RP24R0
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 15-14 **Unimplemented:** Read as '0'

bit 13-8 **RP25R<5:0>:** RP25 Output Pin Mapping bits
Peripheral Output Number n is assigned to pin, RP25 (see Table 11-4 for peripheral function numbers).

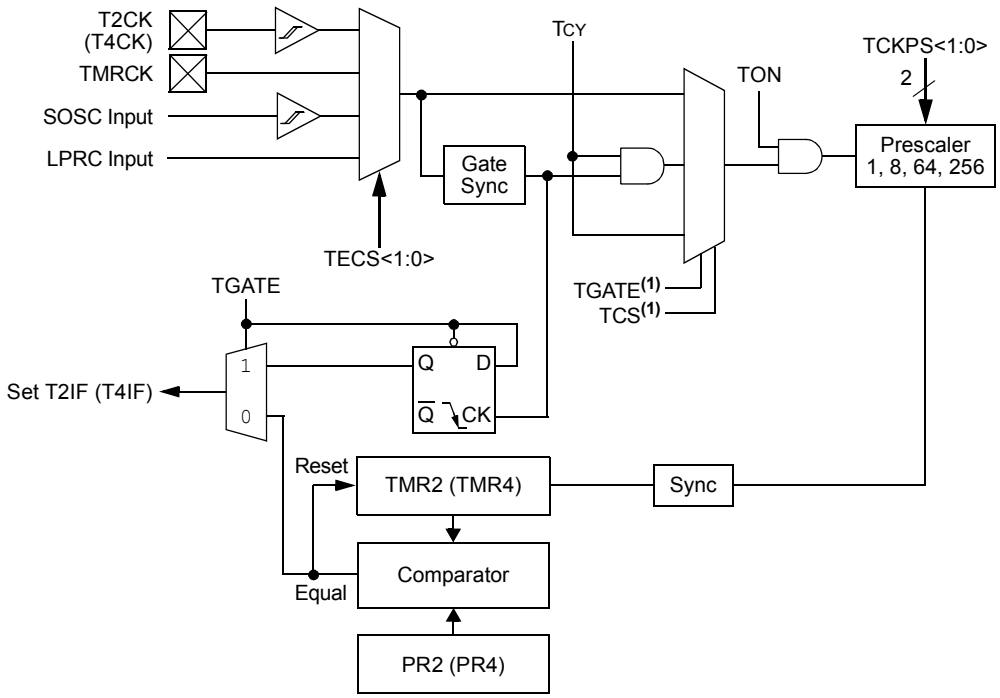
bit 7-6 **Unimplemented:** Read as '0'

bit 5-0 **RP24R<5:0>:** RP24 Output Pin Mapping bits
Peripheral Output Number n is assigned to pin, RP24 (see Table 11-4 for peripheral function numbers).

Note 1: These pins are not available in 28-pin devices.

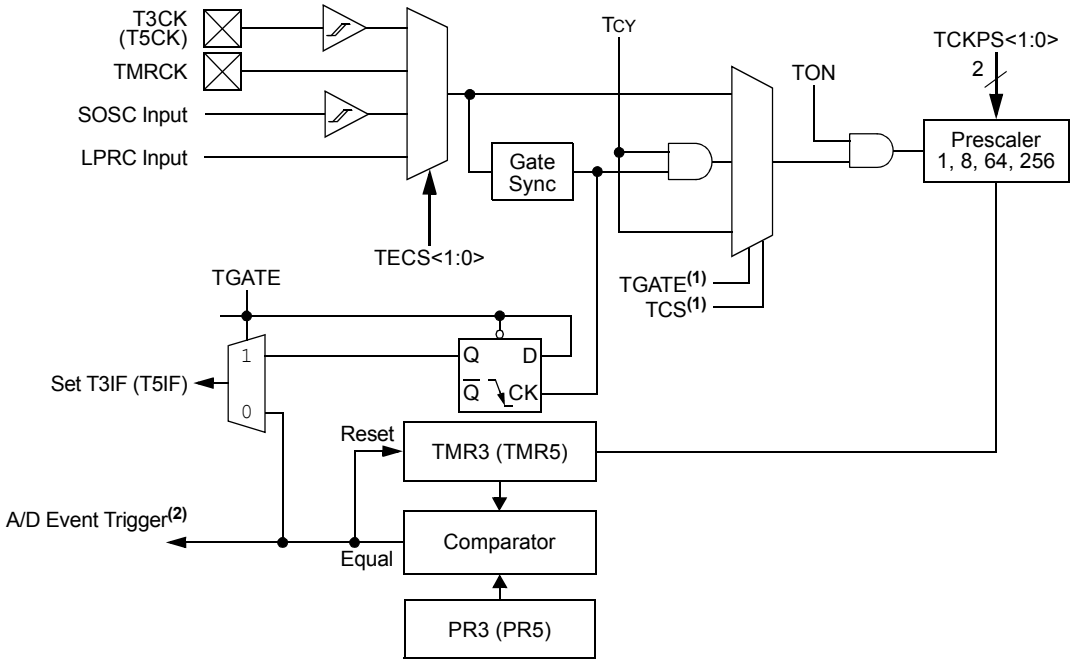
NOTES:

FIGURE 13-2: TIMER2 AND TIMER4 (16-BIT SYNCHRONOUS) BLOCK DIAGRAM



Note 1: The timer clock input must be assigned to an available RPn/RPIn pin before use. See **Section 11.4 "Peripheral Pin Select (PPS)"** for more information.

FIGURE 13-3: TIMER3 AND TIMER5 (16-BIT ASYNCHRONOUS) BLOCK DIAGRAM



Note 1: The timer clock input must be assigned to an available RPn/RPIn pin before use. See **Section 11.4 "Peripheral Pin Select (PPS)"** for more information.

2: The A/D Event Trigger is available only on Timer3.

REGISTER 14-2: ICxCON2: INPUT CAPTURE x CONTROL REGISTER 2

U-0	U-0	U-0	U-0	U-0	U-0	U-0	R/W-0
—	—	—	—	—	—	—	IC32
bit 15							bit 8

R/W-0	R/W-0, HS	U-0	R/W-0	R/W-1	R/W-1	R/W-0	R/W-1
ICTRIG	TRIGSTAT	—	SYNCSEL4	SYNCSEL3	SYNCSEL2	SYNCSEL1	SYNCSEL0
bit 7							bit 0

Legend:	HS = Hardware Settable bit		
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 15-9 **Unimplemented:** Read as '0'

bit 8 **IC32:** Cascade Two IC Modules Enable bit (32-bit operation)
1 = ICx and ICy operate in cascade as a 32-bit module (this bit must be set in both modules)
0 = ICx functions independently as a 16-bit module

bit 7 **ICTRIG:** Input Capture x Sync/Trigger Select bit
1 = Triggers ICx from the source designated by the SYNCSELx bits
0 = Synchronizes ICx with the source designated by the SYNCSELx bits

bit 6 **TRIGSTAT:** Timer Trigger Status bit
1 = Timer source has been triggered and is running (set in hardware, can be set in software)
0 = Timer source has not been triggered and is being held clear

bit 5 **Unimplemented:** Read as '0'

Note 1: Use these inputs as trigger sources only and never as sync sources.
2: Never use an IC module as its own trigger source by selecting this mode.

19.4.2 RECEIVING AN IN TOKEN IN DEVICE MODE

1. Attach to a USB host and enumerate as described in Chapter 9 of the *"USB 2.0 Specification"*.
2. Create a data buffer and populate it with the data to send to the host.
3. In the appropriate (Even or Odd) TX BD for the desired endpoint:
   a) Set up the status register (BDnSTAT) with the correct data toggle (DATA0/1) value and the byte count of the data buffer.
   b) Set up the address register (BDnADR) with the starting address of the data buffer.
   c) Set the UOWN bit of the status register to '1'.
4. When the USB module receives an IN token, it automatically transmits the data in the buffer. Upon completion, the module updates the status register (BDnSTAT) and sets the Transfer Done Interrupt Flag, TRNIF (U1IR<3>).

19.4.3 RECEIVING AN OUT TOKEN IN DEVICE MODE

1. Attach to a USB host and enumerate as described in Chapter 9 of the *"USB 2.0 Specification"*.
2. Create a data buffer with the amount of data you are expecting from the host.
3. In the appropriate (Even or Odd) TX BD for the desired endpoint:
   a) Set up the status register (BDnSTAT) with the correct data toggle (DATA0/1) value and the byte count of the data buffer.
   b) Set up the address register (BDnADR) with the starting address of the data buffer.
   c) Set the UOWN bit of the status register to '1'.
4. When the USB module receives an OUT token, it automatically receives the data sent by the host to the buffer. Upon completion, the module updates the status register (BDnSTAT) and sets the Transfer Done Interrupt Flag, TRNIF (U1IR<3>).

19.5 Host Mode Operation

The following sections describe how to perform common Host mode tasks. In Host mode, USB transfers are invoked explicitly by the host software. The host software is responsible for the Acknowledge portion of the transfer. Also, all transfers are performed using the USB Endpoint 0 Control register (U1EP0) and Buffer Descriptors.

19.5.1 ENABLE HOST MODE AND DISCOVER A CONNECTED DEVICE

1. Enable Host mode by setting the HOSTEN bit (U1CON<3>). This causes the Host mode control bits in other USB OTG registers to become available.
2. Enable the D+ and D- pull-down resistors by setting the DPPULDWN and DMPULDWN bits (U1OTGCON<5:4>). Disable the D+ and D- pull-up resistors by clearing the DPPULUP and DMPULUP bits (U1OTGCON<7:6>).
3. At this point, Start-of-Frame (SOF) generation begins with the SOF counter loaded with 12,000. Eliminate noise on the USB by clearing the SOFEN bit (U1CON<0>) to disable Start-of-Frame packet generation.
4. Enable the device attached interrupt by setting the ATTACHIE bit (U1IE<6>).
5. Wait for the device attached interrupt (U1IR<6> = 1). This is signaled by the USB device changing the state of D+ or D- from '0' to '1' (SE0 to J-state). After it occurs, wait 100 ms for the device power to stabilize.
6. Check the state of the JSTATE and SE0 bits in U1CON. If the JSTATE bit (U1CON<7>) is '0', the connecting device is low speed. If the connecting device is low speed, set the LSPDEN and LSPD bits (U1ADDR<7> and U1EP0<7>, respectively) to enable low-speed operation.
7. Reset the USB device by setting the USBRST bit (U1CON<4>) for at least 50 ms, sending Reset signaling on the bus. After 50 ms, terminate the Reset by clearing USBRST.
8. In order to keep the connected device from going into suspend, enable the SOF packet generation by setting the SOFEN bit.
9. Wait 10 ms for the device to recover from Reset.
10. Perform enumeration as described by Chapter 9 of the *"USB 2.0 Specification"*.

REGISTER 19-13: U1CNFG2: USB CONFIGURATION REGISTER 2

U-0	U-0	U-0	U-0	U-0	U-0	U-0	U-0
—	—	—	—	—	—	—	—
bit 15							bit 8

U-0	U-0	U-0	R/W-0	R/W-0	R/W-0	U-0	R/W-0
—	—	—	PUVBUS	EXTI2CEN	UVBUSDIS(1)	—	UTRDIS(1)
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 15-5 **Unimplemented:** Read as '0'

bit 4 **PUVBUS:** VBUS Pull-up Enable bit
1 = Pull-up on VBUS pin is enabled
0 = Pull-up on VBUS pin is disabled

bit 3 **EXTI2CEN:** I2C™ Interface for External Module Control Enable bit
1 = External module(s) is controlled via the I2C interface
0 = External module(s) is controlled via the dedicated pins

bit 2 **UVBUSDIS:** USB On-Chip 5V Boost Regulator Builder Disable bit(1)
1 = On-chip boost regulator builder is disabled; digital output control interface is enabled
0 = On-chip boost regulator builder is active

bit 1 **Unimplemented:** Read as '0'

bit 0 **UTRDIS:** USB On-Chip Transceiver Disable bit(1)
1 = On-chip transceiver is disabled; digital transceiver interface is enabled
0 = On-chip transceiver is active

Note 1: Never change these bits while the USBPWR bit is set (U1PWRC<0> = 1).

20.0 DATA SIGNAL MODULATOR (DSM)

Note: This data sheet summarizes the features of this group of PIC24F devices. It is not intended to be a comprehensive reference source. For more information, refer to the *"dsPIC33/PIC24 Family Reference Manual"*, **"Data Signal Modulator (DSM)"** (DS39744). The information in this data sheet supersedes the information in the FRM.

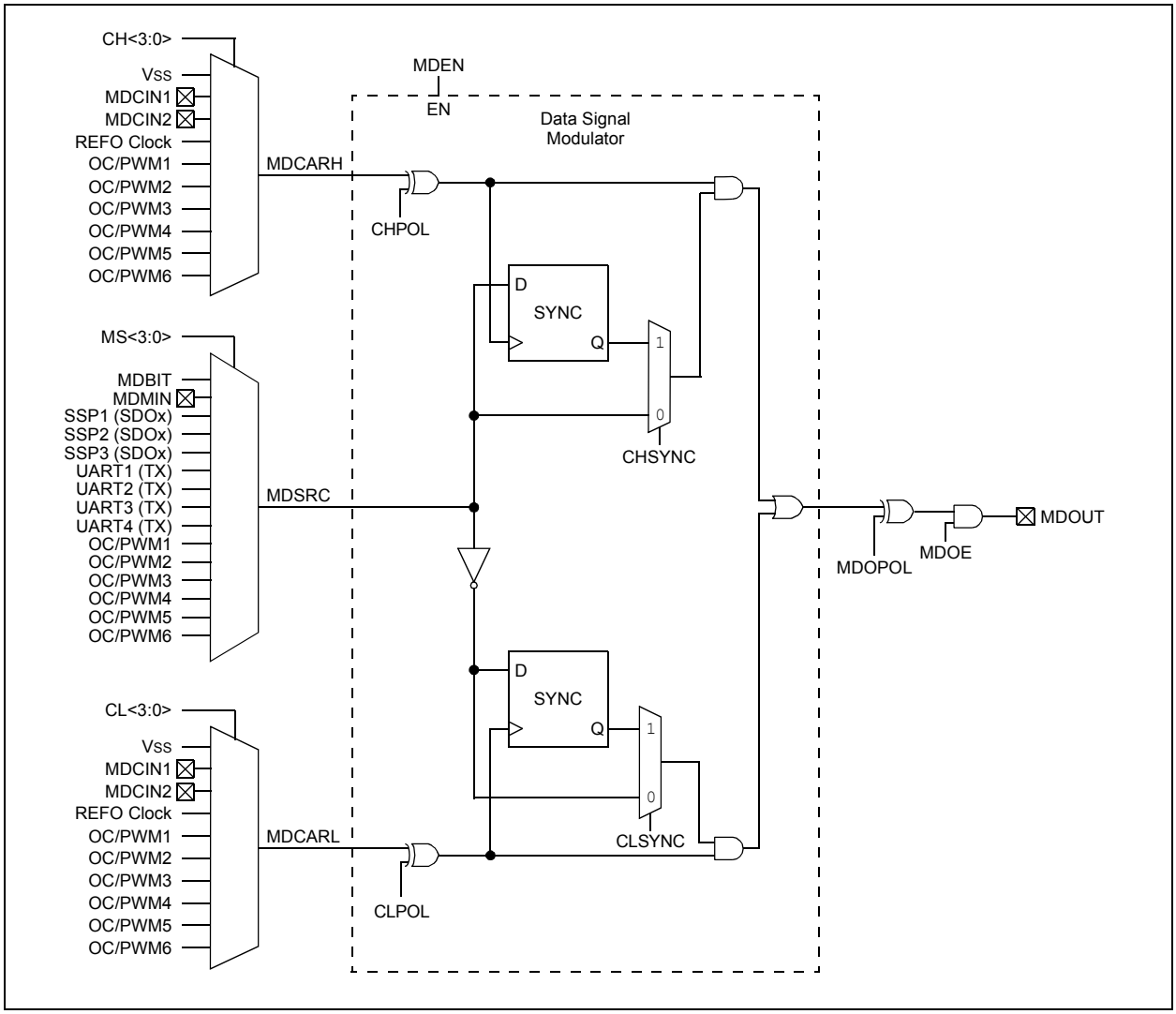
The Data Signal Modulator (DSM) allows the user to mix a digital data stream (the "modulator signal") with a carrier signal to produce a modulated output. Both the carrier and the modulator signals are supplied to the DSM module, either internally from the output of a peripheral, or externally through an input pin.

The modulated output signal is generated by performing a logical AND operation of both the carrier and modulator signals, and then it is provided to the MDOUT pin. Using this method, the DSM can generate the following types of key modulation schemes:

- Frequency Shift Keying (FSK)
- Phase-Shift Keying (PSK)
- On-Off Keying (OOK)

Figure 20-1 shows a simplified block diagram of the Data Signal Modulator peripheral.

FIGURE 20-1: SIMPLIFIED BLOCK DIAGRAM OF THE DATA SIGNAL MODULATOR



24.1 User Interface

24.1.1 POLYNOMIAL INTERFACE

The CRC module can be programmed for CRC polynomials of up to the 32nd order, using up to 32 bits.

Polynomial length, which reflects the highest exponent in the equation, is selected by the PLEN<4:0> bits (CRCCON2<4:0>).

The CRCXORL and CRCXORH registers control which exponent terms are included in the equation. Setting a particular bit includes that exponent term in the equation. Functionally, this includes an XOR operation on the corresponding bit in the CRC engine. Clearing the bit disables the XOR.

For example, consider two CRC polynomials, one is a 16-bit and the other is a 32-bit equation.

EQUATION 24-1: 16-BIT, 32-BIT CRC POLYNOMIALS

$$X16 + X12 + X5 + 1$$

and

$$X32+X26 + X23 + X22 + X16 + X12 + X11 + X10 + X8 + X7 + X5 + X4 + X2 + X + 1$$

To program these polynomials into the CRC generator, set the register bits, as shown in Table 24-1.

Note that the appropriate positions are set to '1' to indicate that they are used in the equation (for example, X26 and X23). The '0' bit required by the equation is always XORed; thus, X0 is a don't care. For a polynomial of length 32, it is assumed that the 32nd bit will be used. Therefore, the X<31:1> bits do not have the 32nd bit.

24.1.2 DATA INTERFACE

The module incorporates a FIFO that works with a variable data width. Input data width can be configured to any value between 1 and 32 bits using the DWIDTH<4:0> bits (CRCCON2<12:8>). When the data width is greater than 15, the FIFO is 4 words deep. When the DWIDTHx bits are between 15 and 8, the FIFO is 8 words deep. When the DWIDTHx bits are less than 8, the FIFO is 16 words deep.

The data for which the CRC is to be calculated must first be written into the FIFO. Even if the data width is less than 8, the smallest data element that can be written into the FIFO is 1 byte. For example, if the DWIDTHx bits are 5, then the size of the data is DWIDTH<4:0> + 1 or 6. The data is written as a whole byte; the two unused upper bits are ignored by the module.

Once data is written into the MSb of the CRCDAT registers (that is, the MSb as defined by the data width), the value of the VWORD<4:0> bits (CRCCON1<12:8>) increments by one. For example, if the DWIDTHx bits are 24, the VWORDx bits will increment when bit 7 of CRCDATH is written. Therefore, CRCDATL must always be written to before CRCDATH.

The CRC engine starts shifting data when the CRCGO bit is set and the value of the VWORDx bits is greater than zero.

Each word is copied out of the FIFO into a buffer register, which decrements the VWORDx bits. The data is then shifted out of the buffer. The CRC engine continues shifting at a rate of two bits per instruction cycle, until the VWORDx bits reach zero. This means that for a given data width, it takes half that number of instructions for each word to complete the calculation. For example, it takes 16 cycles to calculate the CRC for a single word of 32-bit data.

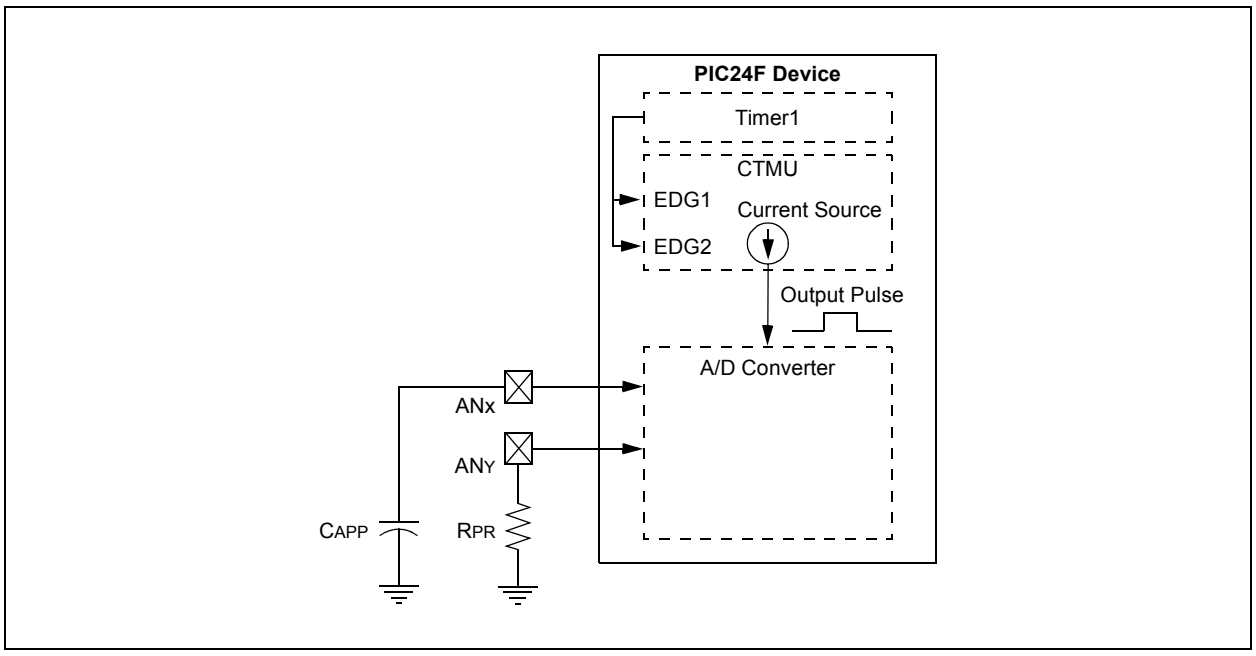
When the VWORDx bits reach the maximum value for the configured value of the DWIDTHx bits (4, 8 or 16), the CRCFUL bit becomes set. When the VWORDx bits reach zero, the CRCMPT bit becomes set. The FIFO is emptied and the VWORD<4:0> bits are set to '00000' whenever CRCEN is '0'.

At least one instruction cycle must pass after a write to CRCWDAT before a read of the VWORDx bits is done.

TABLE 24-1: CRC SETUP EXAMPLES FOR 16 AND 32-BIT POLYNOMIALS

CRC Control Bits	Bit Values	
	16-Bit Polynomial	32-Bit Polynomial
PLEN<4:0>	01111	11111
X<31:16>	0000 0000 0000 0001	0000 0100 1100 0001
X<15:0>	0001 0000 0010 00x	0001 1101 1011 011x

FIGURE 28-1: TYPICAL CONNECTIONS AND INTERNAL CONFIGURATION FOR CAPACITANCE MEASUREMENT



REGISTER 30-3: CW3: FLASH CONFIGURATION WORD 3 (CONTINUED)

bit 7 **Reserved:** Always maintain as '1'

bit 6-0 **WPFP<6:0>:** Write-Protected Code Segment Boundary Page bits(3)

Designates the 512 instruction words page boundary of the protected Code Segment.

If WPEND = 1:

Specifies the lower page boundary of the protected Code Segment; the last page being the last implemented page in the device.

If WPEND = 0:

Specifies the upper page boundary of the protected Code Segment; Page 0 being the lower boundary.

Note 1: Regardless of WPCFG status, if WPEND = 1 or if the WPFP<6:0> bits correspond to the Configuration Word page, the Configuration Word page is protected.

2: Ensure that the SCLKI pin is made a digital input while using this configuration (see Table 11-1).

3: For the 64K devices (PIC24FJ64GB2XX), maintain WPFP6 as '0'.

4: This Configuration bit only takes effect when PLL is not being used.

REGISTER 30-5: DEVID: DEVICE ID REGISTER

U-1	U-1	U-1	U-1	U-1	U-1	U-1	U-1
—	—	—	—	—	—	—	—
bit 23							bit 16

R	R	R	R	R	R	R	R
FAMID7	FAMID6	FAMID5	FAMID4	FAMID3	FAMID2	FAMID1	FAMID0
bit 15							bit 8

R	R	R	R	R	R	R	R
DEV7	DEV6	DEV5	DEV4	DEV3	DEV2	DEV1	DEV0
bit 7							bit 0

Legend:	R = Readable bit	U = Unimplemented bit, read as '1'

bit 23-16 **Unimplemented:** Read as '1'

bit 15-8 **FAMID<7:0>:** Device Family Identifier bits

0100 1100 = PIC24FJ128GB204 family

bit 7-0 **DEV<7:0>:** Individual Device Identifier bits

0101 1000 = PIC24FJ64GB202
0101 1010 = PIC24FJ128GB202
0101 1001 = PIC24FJ64GB204
0101 1011 = PIC24FJ128GB204

REGISTER 30-6: DEVREV: DEVICE REVISION REGISTER

U-0	U-0	U-0	U-0	U-0	U-0	U-0	U-0
—	—	—	—	—	—	—	—
bit 23							bit 16

U-0	U-0	U-0	U-0	U-0	U-0	U-0	U-0
—	—	—	—	—	—	—	—
bit 15							bit 8

U-0	U-0	U-0	U-0	R	R	R	R
—	—	—	—	REV3	REV2	REV1	REV0
bit 7							bit 0

Legend:	R = Readable bit	U = Unimplemented bit, read as '1'

bit 23-4 **Unimplemented:** Read as '1'

bit 3-0 **REV<3:0>:** Device Revision Identifier bits

30.4 Program Verification and Code Protection

PIC24FJ128GB204 family devices provide two complimentary methods to protect application code from overwrites and erasures. These also help to protect the device from inadvertent configuration changes during run time.

30.4.1 GENERAL SEGMENT PROTECTION

For all devices in the PIC24FJ128GB204 family, the on-chip program memory space is treated as a single block, known as the General Segment (GS). Code protection for this block is controlled by one Configuration bit, GCP. This bit inhibits external reads and writes to the program memory space. It has no direct effect in normal execution mode.

Write protection is controlled by the GWRP bit in Configuration Word 1. When GWRP is programmed to '0', internal write and erase operations to program memory are blocked.

30.4.2 CODE SEGMENT PROTECTION

In addition to global General Segment protection, a separate subrange of the program memory space can be individually protected against writes and erases. This area can be used for many purposes where a separate block of write and erase-protected code is needed, such as bootloader applications. Unlike common boot block implementations, the specially protected segment in the PIC24FJ128GB204 family devices can be located by the user anywhere in the program space and configured in a wide range of sizes.

Code Segment (CS) protection provides an added level of protection to a designated area of program memory by disabling the NVM safety interlock whenever a write or erase address falls within a specified range. It does not override General Segment protection controlled by the GCP or GWRP bits. For example, if GCP and GWRP are enabled, enabling segmented code protection for the bottom half of program memory does not undo General Segment protection for the top half.

The size and type of protection for the segmented code range are configured by the WPFPx, WPEND, WPCFG and WPDIS bits in Configuration Word 3. Code Segment protection is enabled by programming the WPDIS bit (= 0). The WPFPx bits specify the size of the segment to be protected, by specifying the 512-word code page that is the start or end of the protected segment. The specified region is inclusive, therefore, this page will also be protected.

The WPEND bit determines if the protected segment uses the top or bottom of the program space as a boundary. Programming WPEND (= 0) sets the bottom of program memory (000000h) as the lower boundary of the protected segment. Leaving WPEND unprogrammed (= 1) protects the specified page through the last page of implemented program memory, including the Configuration Word locations.

A separate bit, WPCFG, is used to protect the last page of program space, including the Flash Configuration Words. Programming WPCFG (= 0) protects the last page in addition to the pages selected by the WPEND and WPFP<6:0> bits setting. This is useful in circumstances where write protection is needed for both the Code Segment in the bottom of the memory and the Flash Configuration Words.

The various options for segment code protection are shown in Table 30-2.

TABLE 30-2: CODE SEGMENT PROTECTION CONFIGURATION OPTIONS

Segment Configuration Bits			Write/Erase Protection of Code Segment
WPDIS	WPEND	WPCFG	
1	x	x	No additional protection is enabled; all program memory protection is configured by GCP and GWRP.
0	1	x	Addresses from the first address of the code page are defined by WPFP<6:0> through the end of implemented program memory (inclusive); erase/write-protected, including Flash Configuration Words.
0	0	1	Address, 000000h through the last address of the code page, is defined by WPFP<6:0> (inclusive); erase/write-protected.
0	0	0	Address, 000000h through the last address of the code page, is defined by WPFP<6:0> (inclusive); erase/write-protected and the last page, including the Flash Configuration Words, are erase/write-protected.