



Welcome to [E-XFL.COM](https://www.e-xfl.com)

What is "[Embedded - Microcontrollers](#)"?

"[Embedded - Microcontrollers](#)" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

Applications of "[Embedded - Microcontrollers](#)"

Details

Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	64MHz
Connectivity	CANbus, I ² C, LINbus, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, DMA, POR, PWM, WDT
Number of I/O	25
Program Memory Size	64KB (32K x 16)
Program Memory Type	FLASH
EEPROM Size	1K x 8
RAM Size	4K x 8
Voltage - Supply (Vcc/Vdd)	1.8V ~ 3.6V
Data Converters	A/D 24x12b; D/A 1x5b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	28-SSOP (0.209", 5.30mm Width)
Supplier Device Package	28-SSOP
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18lf26k83-i-ss

2.5 External Oscillator Pins

Many microcontrollers have options for at least two oscillators: a high-frequency primary oscillator and a low-frequency secondary oscillator (refer to **Section 7.0 “Oscillator Module (with Fail-Safe Clock Monitor)”** for details).

The oscillator circuit should be placed on the same side of the board as the device. Place the oscillator circuit close to the respective oscillator pins with no more than 0.5 inch (12 mm) between the circuit components and the pins. The load capacitors should be placed next to the oscillator itself, on the same side of the board.

Use a grounded copper pour around the oscillator circuit to isolate it from surrounding circuits. The grounded copper pour should be routed directly to the MCU ground. Do not run any signal traces or power traces inside the ground pour. Also, if using a two-sided board, avoid any traces on the other side of the board where the crystal is placed.

Layout suggestions are shown in Figure 2-3. In-line packages may be handled with a single-sided layout that completely encompasses the oscillator pins. With fine-pitch packages, it is not always possible to completely surround the pins and components. A suitable solution is to tie the broken guard sections to a mirrored ground layer. In all cases, the guard trace(s) must be returned to ground.

In planning the application's routing and I/O assignments, ensure that adjacent port pins, and other signals in close proximity to the oscillator, are benign (i.e., free of high frequencies, short rise and fall times, and other similar noise).

For additional information and design guidance on oscillator circuits, refer to these Microchip Application Notes, available at the corporate website (www.microchip.com):

- AN826, “Crystal Oscillator Basics and Crystal Selection for rPIC™ and PICmicro® Devices”
- AN849, “Basic PICmicro® Oscillator Design”
- AN943, “Practical PICmicro® Oscillator Analysis and Design”
- AN949, “Making Your Oscillator Work”

2.6 Unused I/Os

Unused I/O pins should be configured as outputs and driven to a logic low state. Alternatively, connect a 1 kΩ to 10 kΩ resistor to Vss on unused pins and drive the output to logic low.

FIGURE 2-3: SUGGESTED PLACEMENT OF THE OSCILLATOR CIRCUIT

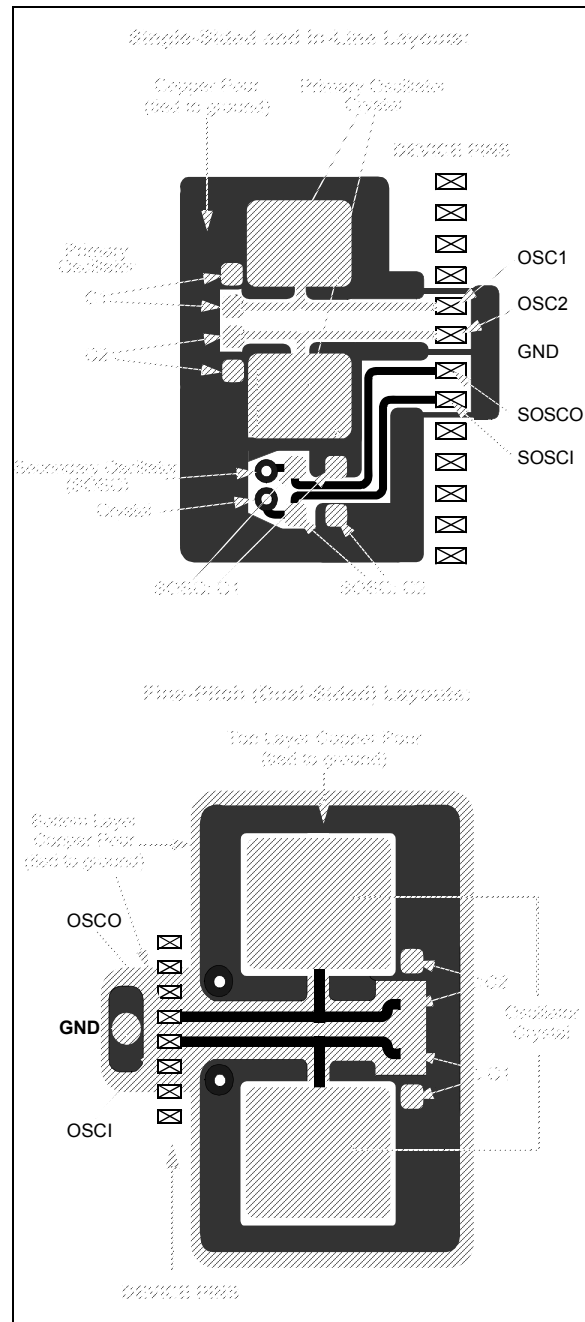


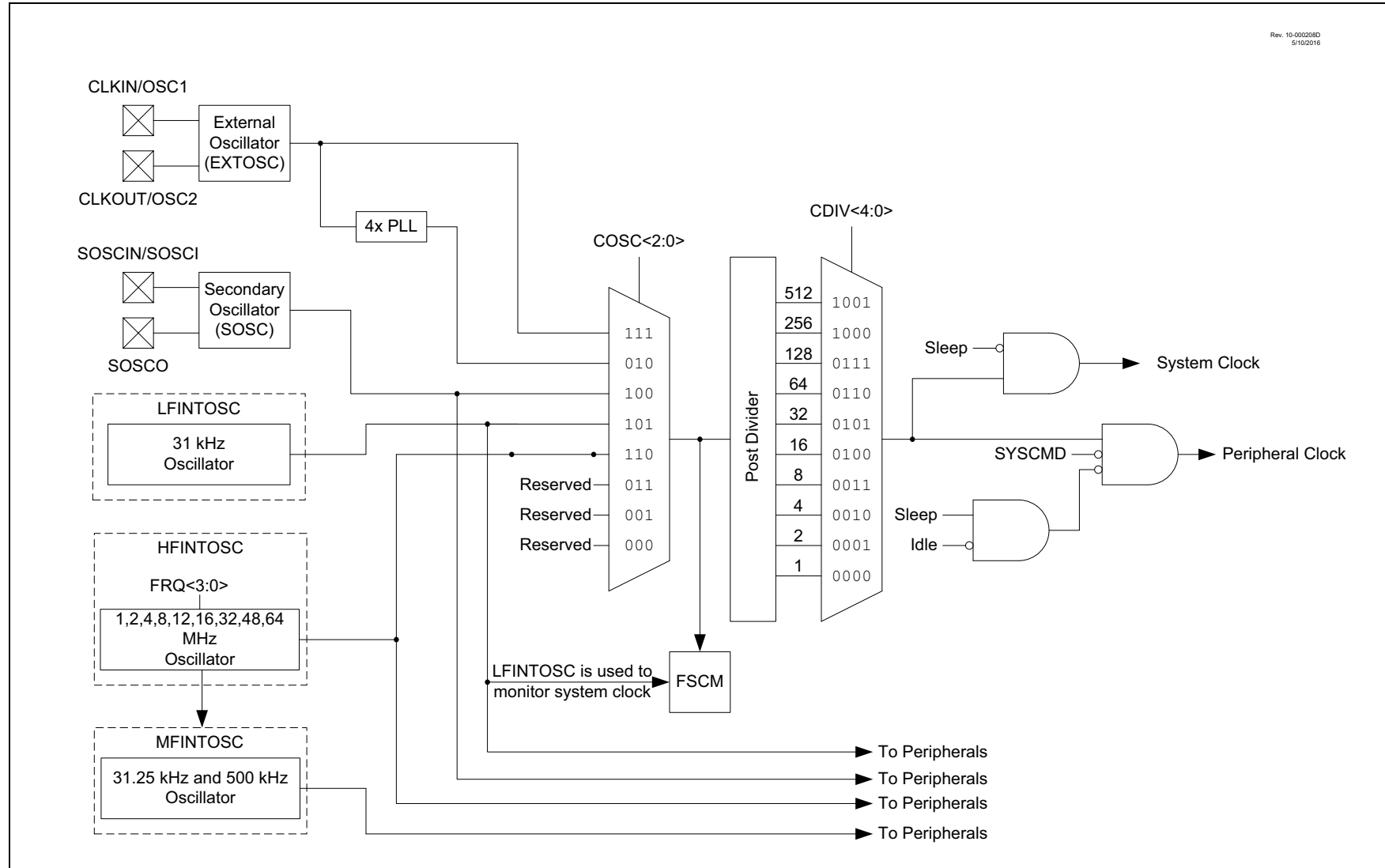
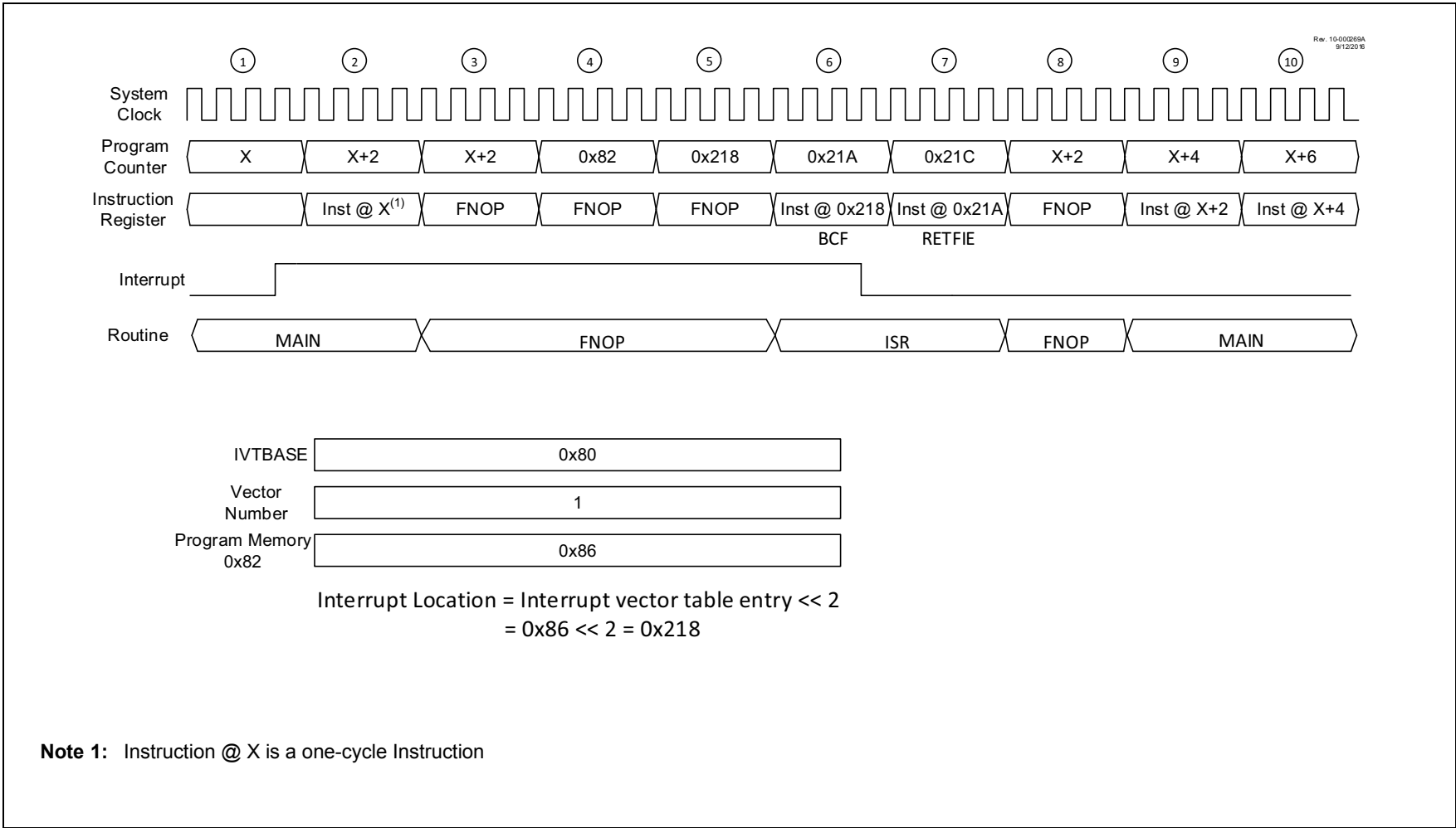
FIGURE 7-1: SIMPLIFIED PIC® MCU CLOCK SOURCE BLOCK DIAGRAM

FIGURE 9-7: INTERRUPT TIMING DIAGRAM – ONE CYCLE INSTRUCTION



REGISTER 9-39: IVTLOCK: INTERRUPT VECTOR TABLE LOCK REGISTER

U-0	U-0	U-0	U-0	U-0	U-0	U-0	R/W-0/0
—	—	—	—	—	—	—	IVTLOCKED ^(1,2)
bit 7							bit 0

Legend:

R = Readable bit

W = Writable bit

U = Unimplemented bit, read as '0'

u = Bit is unchanged

x = Bit is unknown

-n/n = Value at POR and BOR/Value at all other Resets

'1' = Bit is set

'0' = Bit is cleared

bit 7-1 **Unimplemented:** Read as '0'

bit 0 **IVTLOCKED:** IVT Registers Lock bits^(1,2)

1 = IVTBASE Registers are locked and cannot be written

0 = IVTBASE Registers can be modified by write operations

Note 1: The IVTLOCK bit can only be set or cleared after the unlock sequence in Example 9-1.

2: If IVT1WAY = 1, the IVTLOCK bit cannot be cleared after it has been set. See Register 5-3.

REGISTER 9-40: SHADCON: SHADOW CONTROL REGISTER

U-0	U-0	U-0	U-0	U-0	U-0	U-0	R/W-0/0
—	—	—	—	—	—	—	SHADLO
bit 7							bit 0

Legend:

R = Readable bit

W = Writable bit

U = Unimplemented bit, read as '0'

-n = Value at POR

'1' = Bit is set

'0' = Bit is cleared

x = Bit is unknown

bit 7-1 **Unimplemented:** Read as '0'

bit 0 **SHADLO:** Interrupt Shadow Register Access Switch bit

0 = Access Main Context for Interrupt Shadow Registers

1 = Access Low-Priority Interrupt Context for Interrupt Shadow Registers

13.1.4 NVM UNLOCK SEQUENCE

The unlock sequence is a mechanism that protects the NVM from unintended self-write programming or erasing. The sequence must be executed and completed without interruption to successfully complete any of the following operations:

- PFM Row Erase
- Write of PFM write latches to PFM memory
- Write of PFM write latches to User IDs
- Write to Data EEPROM Memory
- Write to Configuration Words

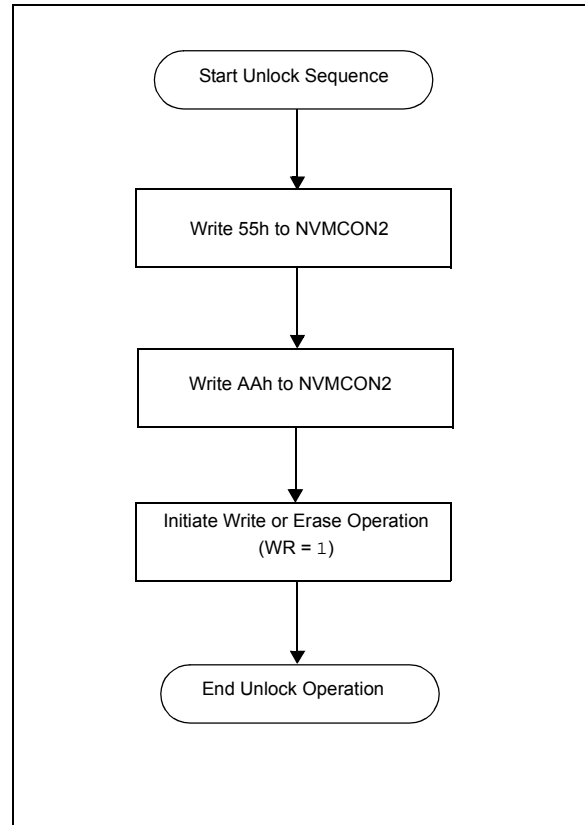
The unlock sequence consists of the following steps and must be completed in order:

- Write 55h to NVMCON2
- Write AAh to NVMCON2
- Set the WR bit of NVMCON1

Once the WR bit is set, the processor will stall internal operations until the operation is complete and then resume with the next instruction.

Since the unlock sequence must not be interrupted, global interrupts should be disabled prior to the unlock sequence and re-enabled after the unlock sequence is completed.

FIGURE 13-6: NVM UNLOCK SEQUENCE FLOWCHART



EXAMPLE 13-2: NVM UNLOCK SEQUENCE

```

BCF          INTCON0,GIE          ; Recommended so sequence is not interrupted
BANKSEL      NVMCON1
BSF          NVMCON1,WREN         ; Enable write/erase
MOVLW       55h                  ; Load 55h

MOVWF        NVMCON2              ; Step 1: Load 55h into NVMCON2
MOVLW       AAh                  ; Step 2: Load W with AAh
MOVWF        NVMCON2              ; Step 3: Load AAh into NVMCON2
BSF          INTCON1,WR           ; Step 4: Set WR bit to begin write/erase

BSF          INTCON0,GIE          ; Re-enable interrupts
  
```

Note 1: Sequence begins when NVMCON2 is written; steps 1-4 must occur in the cycle-accurate order shown. If the timing of the steps 1 to 4 is corrupted by an interrupt or a debugger Halt, the action will not take place.

2: Opcodes shown are illustrative; any instruction that has the indicated effect may be used.

15.0 DIRECT MEMORY ACCESS (DMA)

15.1 Introduction

The Direct Memory Access (DMA) module is designed to service data transfers between different memory regions directly without intervention from the CPU. By eliminating the need for CPU-intensive management of handling interrupts intended for data transfers, the CPU now can spend more time on other tasks.

PIC18(L)F25/26K83 family has two DMA modules which can be independently programmed to transfer data between different memory locations, move different data sizes, and use a wide range of hardware triggers to initiate transfers. The two DMA registers can even be programmed to work together, in order to carry out more complex data transfers without CPU overhead.

Key features of the DMA module include:

- Support access to the following memory regions:
 - GPR and SFR space (R/W)
 - Program Flash Memory (R only)
 - Data EEPROM Memory (R only)
- Programmable priority between the DMA and CPU Operations. Refer to **Section 3.1 “System Arbitration”** for details.
- Programmable Source and Destination address modes
 - Fixed address
 - Post-increment address
 - Post-decrement address
- Programmable Source and Destination sizes
- Source and destination pointer register, dynamically updated and reloadable
- Source and destination count register, dynamically updated and reloadable
- Programmable auto-stop based on Source or Destination counter
- Software triggered transfers
- Multiple user selectable sources for hardware triggered transfers
- Multiple user selectable sources for aborting DMA transfers

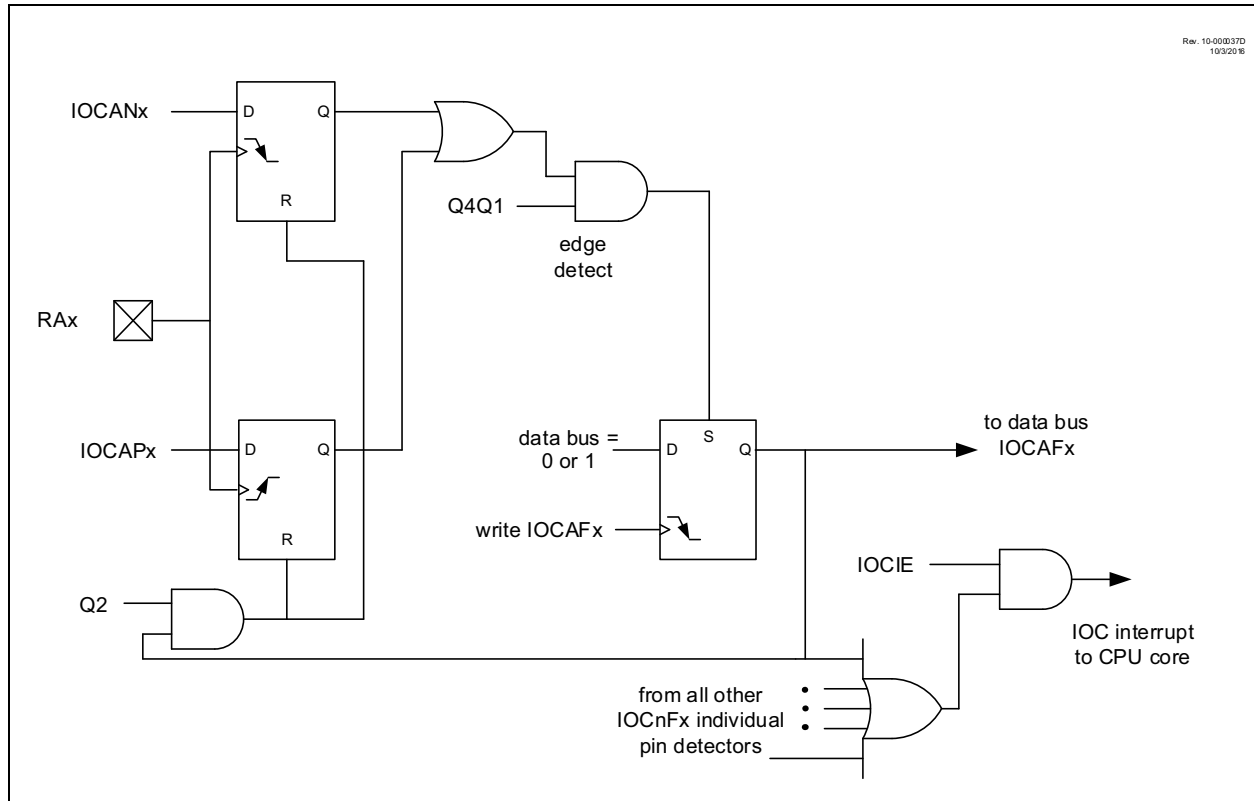
15.2 DMA Registers

The operation of the DMA module has the following registers:

- Control registers (DMAxCON0, DMAxCON1)
- Data buffer register (DMAxBUF)
- Source Start Address Register (DMAxSSAU:H:L)
- Source Pointer Register (DMAxSPTRU:H:L)
- Source Message Size Register (DMAxSSZH:L)
- Source Count Register (DMAxSCNTH:L)
- Destination Start Address Register (DMAxDSAH:L)
- Destination Pointer Register (DMAxDPTRH:L)
- Destination Message Size Register (DMAxDSZH:L)
- Destination Count Register (DMAxDCNTH:L)
- Start Interrupt Request Source Register (DMAxSIRQ)
- Abort Interrupt Request Source Register (DMAxAIRQ)

These registers are detailed in **Section 15.13 “Register definitions: DMA”**.

FIGURE 18-1: INTERRUPT-ON-CHANGE BLOCK DIAGRAM (PORTA EXAMPLE)



22.6 Timer2 Operation During Sleep

When PSYNC = 1, Timer2 cannot be operated while the processor is in Sleep mode. The contents of the T2TMR and T2PR registers will remain unchanged while processor is in Sleep mode.

When PSYNC = 0, Timer2 will operate in Sleep as long as the clock source selected is also still running. Selecting the LFINTOSC, MFINTOSC, or HFINTOSC oscillator as the timer clock source will keep the selected oscillator running during Sleep.

23.5 Register Definitions: CCP Control

Long bit name prefixes for the CCP peripherals are shown below. Refer to **Section 1.3.2.2 “Long Bit Names”** for more information.

Peripheral	Bit Name Prefix
CCP1	CCP1
CCP2	CCP2
CCP3	CCP3
CCP4	CCP4

REGISTER 23-1: CCPxCON: CCPx CONTROL REGISTER

R/W-0/0	U-0	R-x	R/W-0/0	R/W-0/0	R/W-0/0	R/W-0/0	R/W-0/0
EN	—	OUT	FMT	MODE<3:0>			
bit 7				bit 0			

Legend:

R = Readable bit W = Writable bit U = Unimplemented bit, read as '0'
 -n = Value at POR '1' = Bit is set '0' = Bit is cleared x = Bit is unknown

bit 7 **EN:** CCP Module Enable bit
 1 = CCP is enabled
 0 = CCP is disabled

bit 6 **Unimplemented:** Read as '0'

bit 5 **OUT:** CCPx Output Data bit (read-only)

bit 4 **FMT:** CCPW (pulse-width) Alignment bit
 MODE = Capture mode:
 Unused
 MODE = Compare mode:
 Unused
 MODE = PWM mode:
 1 = Left-aligned format
 0 = Right-aligned format

bit 3-0 **MODE<3:0>:** CCPx Mode Select bits

MODE	Operating Mode	Operation	Set CCPxIF
11xx	PWM	PWM operation	Yes
1011	Compare	Pulse output; clear TMR1 ⁽²⁾	Yes
1010		Pulse output	Yes
1001		Clear output ⁽¹⁾	Yes
1000		Set output ⁽¹⁾	Yes
0111	Capture	Every 16th rising edge of CCPx input	Yes
0110		Every 4th rising edge of CCPx input	Yes
0101		Every rising edge of CCPx input	Yes
0100		Every falling edge of CCPx input	Yes
0011		Every edge of CCPx input	Yes
0010	Compare	Toggle output	Yes
0001		Toggle output; clear TMR1 ⁽²⁾	Yes
0000	Disabled		—

Note 1: The set and clear operations of the Compare mode are reset by setting MODE = 4'b0000 or EN = 0.
Note 2: When MODE = 0001 or 1011, then the timer associated with the CCP module is cleared. TMR1 is the default selection for the CCP module, so it is used for indication purpose only.

FIGURE 25-7: PERIOD AND DUTY-CYCLE SINGLE ACQUISITION TIMING DIAGRAM

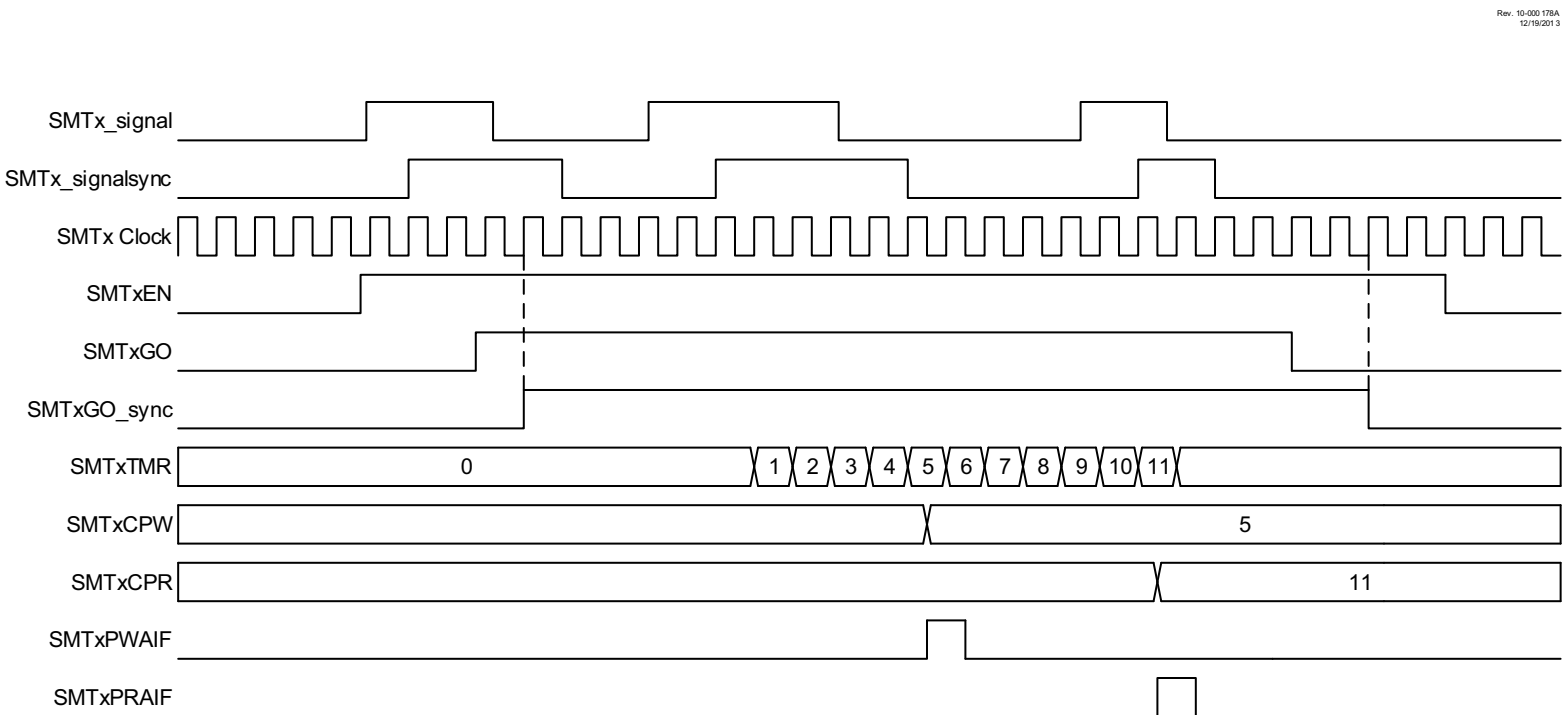
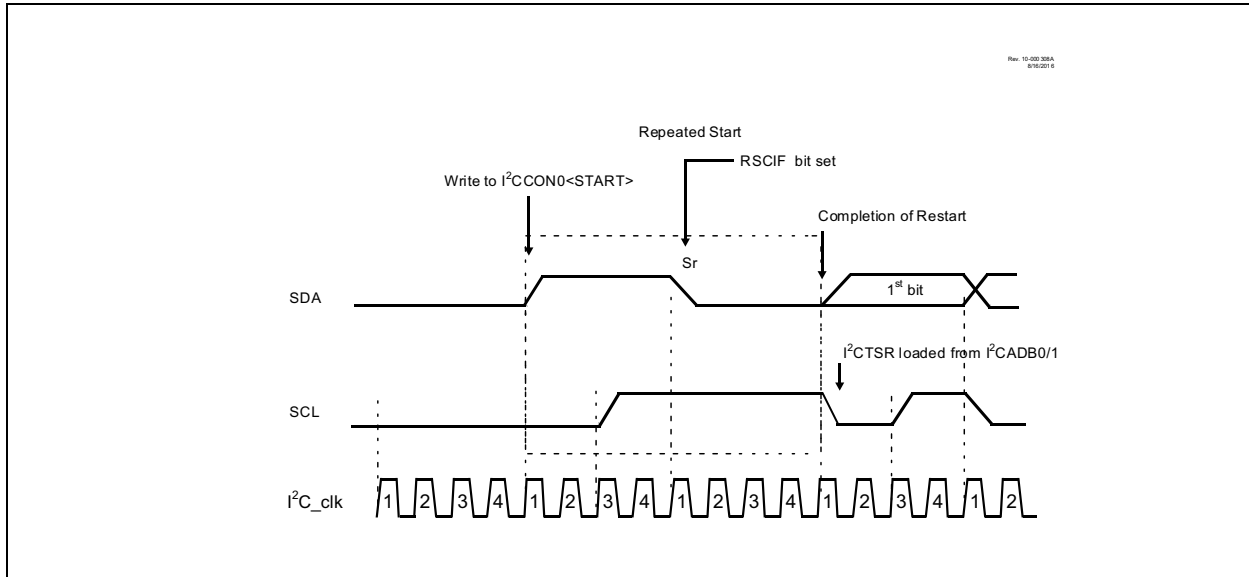


FIGURE 33-16: REPEATED START CONDITION TIMING

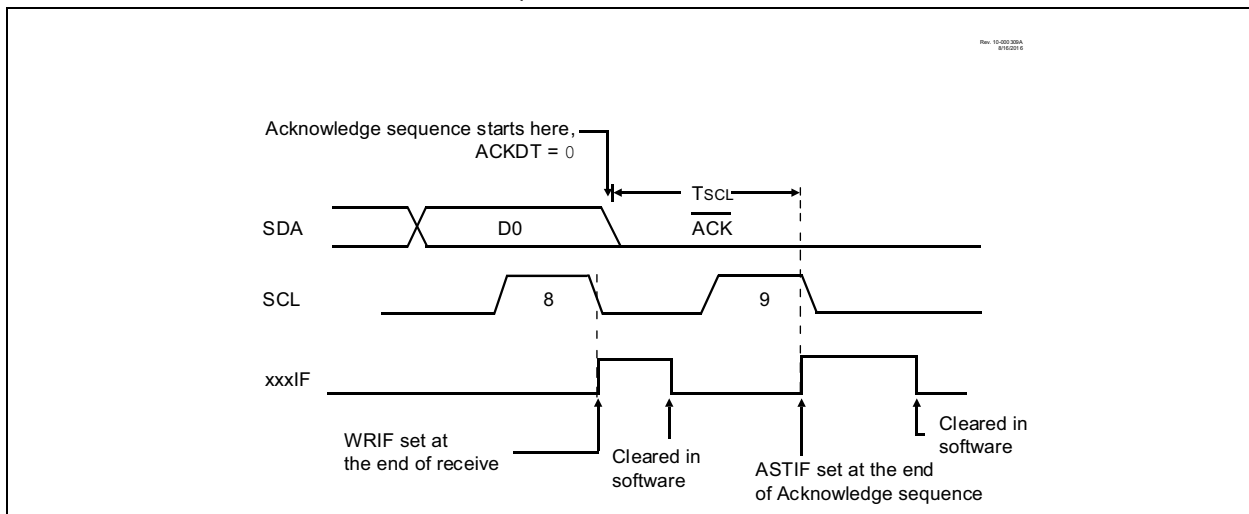


33.5.7 ACKNOWLEDGE SEQUENCE TIMING

An Acknowledge sequence is enabled automatically following an address/data byte transmission. The SCL pin is pulled low and the contents of the Acknowledge Data bits (ACKDT/ACKCNT) are presented on the SDA pin. If the user wishes to generate an Acknowledge, then the ACKDT bit should be cleared. If not, the user

should set the ACKDT bit before starting an Acknowledge sequence. The master then waits one clock period (T_{SCL}) and the SCL pin is released high. When the SCL pin is sampled high (clock arbitration), the master counts another T_{SCL} . The SCL pin is then pulled low. Figure 33-17 shows the timings for Acknowledge sequence.

FIGURE 33-17: ACKNOWLEDGE SEQUENCE TIMING



33.5.8 STOP CONDITION TIMING

A Stop bit is asserted on the SDA pin at the end of receive/transmit when $I2Cx\text{CNT} = 0$. After the last byte of a receive/transmit sequence, the SCL line is held low. The master asserts the SDA line low. The SCL pin is then released high $T_{SCL}/2$ later and is detected high. The SDA pin is then released. When the SDA pin tran-

sitions high while SCL is high, the PCIF bit of the $I2Cx\text{IF}$ register is set. Figure 33-18 shows the timings for a Stop condition.

REGISTER 33-7: I2CxSTAT1: I²C STATUS REGISTER 1

R/W/HS-0	U-0	R-1	U-0	R/W/HS-0	R/S-0/0	U-0	R-0
TXWE ⁽²⁾	—	TXBE ^(1, 3)	—	RXRE ⁽²⁾	CLRBF	—	RXBF ^(1,3)
bit 7							bit 0

Legend:

R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'
u = Bit is unchanged	x = Bit is unknown	-n/n = Value at POR and BOR/Value at all other Resets
'1' = Bit is set	'0' = Bit is cleared	HS = Hardware set HC = Hardware clear

- bit 7 **TXWE:** Transmit Write Error Status bit ⁽²⁾
1 = A new byte of data was written to I2CTXB when it was full (Must be cleared by software)
0 = No transmit write error
- bit 6 **Unimplemented:** Read as '0'
- bit 5 **TXBE:** Transmit Buffer Empty Status bit
1 = I2CTXB is empty (Cleared by writing the I2CTXB register)
0 = I2CTXB is full
- bit 4 **Unimplemented:** Read as '0'
- bit 3 **RXRE:** Receive Read Error Status bit
1 = A byte of data was read from I2CxRXB when it was empty. (Must be cleared by software)
0 = No receive overflow
- bit 2 **CLRBF:** Clear Buffer bit
Setting this bit clears/empties the receive and transmit buffers, causing reset of RXBF and TXBE.
Setting this bit clears the RXIF and TXIF interrupt flags.
This bit is set-only special function, and always reads '0'
- bit 1 **Unimplemented:** Read as '0'
- bit 0 **RXBF:** Receive Buffer Full Status bit
1 = I2CRXB has received new data (Cleared by reading the I2CRXB register)
0 = I2CRXB is empty

- Note 1:** The bits are held in Reset when I2CEN = 0.
- 2:** Will cause NACK to be sent for slave address and master/slave data read bytes.
- 3:** Used as triggers for DMA operation.

REGISTER 34-35: BnDLC: TX/RX BUFFER 'n' DATA LENGTH CODE REGISTERS IN TRANSMIT MODE [0 ≤ n ≤ 5, TXnEN (BSEL<n>) = 1]⁽¹⁾

U-0	R/W-x	U-0	U-0	R/W-x	R/W-x	R/W-x	R/W-x
—	TXRTR	—	—	DLC3	DLC2	DLC1	DLC0
bit 7				bit 0			

Legend:

R = Readable bit W = Writable bit U = Unimplemented bit, read as '0'
-n = Value at POR '1' = Bit is set '0' = Bit is cleared x = Bit is unknown

- bit 7 **Unimplemented:** Read as '0'
- bit 6 **TXRTR:** Transmitter Remote Transmission Request bit
1 = Transmitted message will have the RTR bit set
0 = Transmitted message will have the RTR bit cleared
- bit 5-4 **Unimplemented:** Read as '0'
- bit 3-0 **DLC<3:0>:** Data Length Code bits
1111-1001 = Reserved
1000 = Data length = 8 bytes
0111 = Data length = 7 bytes
0110 = Data length = 6 bytes
0101 = Data length = 5 bytes
0100 = Data length = 4 bytes
0011 = Data length = 3 bytes
0010 = Data length = 2 bytes
0001 = Data length = 1 byte
0000 = Data length = 0 bytes

Note 1: These registers are available in Mode 1 and 2 only.

REGISTER 34-36: BSEL0: BUFFER SELECT REGISTER 0⁽¹⁾

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	U-0	U-0
B5TXEN	B4TXEN	B3TXEN	B2TXEN	B1TXEN	B0TXEN	—	—
bit 7						bit 0	

Legend:

R = Readable bit W = Writable bit U = Unimplemented bit, read as '0'
-n = Value at POR '1' = Bit is set '0' = Bit is cleared x = Bit is unknown

- bit 7-2 **B<5:0>TXEN:** Buffer 5 to Buffer 0 Transmit Enable bits
1 = Buffer is configured in Transmit mode
0 = Buffer is configured in Receive mode
- bit 1-0 **Unimplemented:** Read as '0'

Note 1: These registers are available in Mode 1 and 2 only.

REGISTER 34-39: RXFnEIDH: RECEIVE ACCEPTANCE FILTER 'n' EXTENDED IDENTIFIER REGISTERS, HIGH BYTE [0 ≤ n ≤ 15]⁽¹⁾

R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x
EID15	EID14	EID13	EID12	EID11	EID10	EID9	EID8
bit 7							bit 0

Legend:

R = Readable bit W = Writable bit U = Unimplemented bit, read as '0'
 -n = Value at POR '1' = Bit is set '0' = Bit is cleared x = Bit is unknown

bit 7-0 **EID<15:8>**: Extended Identifier Filter bits

Note 1: Registers, RXF6EIDH:RXF15EIDH, are available in Mode 1 and 2 only.

REGISTER 34-40: RXFnEIDL: RECEIVE ACCEPTANCE FILTER 'n' EXTENDED IDENTIFIER REGISTERS, LOW BYTE [0 ≤ n ≤ 15]⁽¹⁾

R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x
EID7	EID6	EID5	EID4	EID3	EID2	EID1	EID0
bit 7							bit 0

Legend:

R = Readable bit W = Writable bit U = Unimplemented bit, read as '0'
 -n = Value at POR '1' = Bit is set '0' = Bit is cleared x = Bit is unknown

bit 7-0 **EID<7:0>**: Extended Identifier Filter bits

Note 1: Registers, RXF6EIDL:RXF15EIDL, are available in Mode 1 and 2 only.

REGISTER 34-41: RXMnSIDH: RECEIVE ACCEPTANCE MASK 'n' STANDARD IDENTIFIER MASK REGISTERS, HIGH BYTE [0 ≤ n ≤ 1]

R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x
SID10	SID9	SID8	SID7	SID6	SID5	SID4	SID3
bit 7							bit 0

Legend:

R = Readable bit W = Writable bit U = Unimplemented bit, read as '0'
 -n = Value at POR '1' = Bit is set '0' = Bit is cleared x = Bit is unknown

bit 7-0 **SID<10:3>**: Standard Identifier Mask bits or Extended Identifier Mask bits (EID<28:21>)

TABLE 42-2: INSTRUCTION SET (CONTINUED)

Mnemonic, Operands		Description	Cycles	16-Bit Instruction Word				Status Affected	Notes	
				MSb		LSb				
LITERAL INSTRUCTIONS										
ADDLW	k	Add literal and WREG	1	0000	1111	kkkk	kkkk	C, DC, Z, OV, N		
ANDLW	k	AND literal with WREG	1	0000	1011	kkkk	kkkk	Z, N		
IORLW	k	Inclusive OR literal with WREG	1	0000	1001	kkkk	kkkk	Z, N		
LFSR	f _n , k	Load FSR(f _n) with a 14-bit literal (k)	2	1110	1110	00ff	kkkk	None		
ADDFSR	f _n , k	Add FSR(f _n) with (k)	1	1110	1000	ffkk	kkkk	None		
SUBFSR	f _n , k	Subtract (k) from FSR(f _n)	1	1110	1001	ffkk	kkkk	None		
MOVLB	k	Move literal to BSR<5:0>	1	0000	0001	00kk	kkkk	None		
MOVLW	k	Move literal to WREG	1	0000	1110	kkkk	kkkk	None		
MULLW	k	Multiply literal with WREG	1	0000	1101	kkkk	kkkk	None		
RETLW	k	Return with literal in WREG	2	0000	1100	kkkk	kkkk	None		
SUBLW	k	Subtract WREG from literal	1	0000	1000	kkkk	kkkk	C, DC, Z, OV, N		
XORLW	k	Exclusive OR literal with WREG	1	0000	1010	kkkk	kkkk	Z, N		
DATA MEMORY – PROGRAM MEMORY INSTRUCTIONS										
TBLRD*		Table Read	2 - 5	0000	0000	0000	1000	None		
TBLRD*+		Table Read with post-increment		0000	0000	0000	1001	None		
TBLRD*-		Table Read with post-decrement		0000	0000	0000	1010	None		
TBLRD+*		Table Read with pre-increment	2 - 5	0000	0000	0000	1011	None		
TBLWT*		Table Write		0000	0000	0000	1100	None		
TBLWT*+		Table Write with post-increment		0000	0000	0000	1101	None		
TBLWT*-		Table Write with post-decrement		0000	0000	0000	1110	None		
TBLWT+*		Table Write with pre-increment		0000	0000	0000	1111	None		

- Note 1:** If Program Counter (PC) is modified or a conditional test is true, the instruction requires an additional cycle. The extra cycle is executed as a NOP.
- 2:** Some instructions are multi word instructions. The second/third words of these instructions will be decoded as a NOP, unless the first word of the instruction retrieves the information embedded in these 16-bits. This ensures that all program memory locations have a valid instruction.
- 3:** f_s and f_d do not cover the full memory range. 2 MSBs of bank selection are forced to 'b00 to limit the range of these instructions to lower 4k addressing space.

BRA Unconditional Branch

Syntax:	BRA n							
Operands:	$-1024 \leq n \leq 1023$							
Operation:	$(PC) + 2 + 2n \rightarrow PC$							
Status Affected:	None							
Encoding:	<table border="1"><tr><td>1101</td><td>0nnn</td><td>nnnn</td><td>nnnn</td></tr></table>				1101	0nnn	nnnn	nnnn
1101	0nnn	nnnn	nnnn					
Description:	Add the 2's complement number '2n' to the PC. Since the PC will have incremented to fetch the next instruction, the new address will be $PC + 2 + 2n$. This instruction is a 2-cycle instruction.							
Words:	1							
Cycles:	2							

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'n'	Process Data	Write to PC
No operation	No operation	No operation	No operation

Example: HERE BRA Jump

Before Instruction
PC = address (HERE)

After Instruction
PC = address (Jump)

BSF Bit Set f

Syntax:	BSF f, b {,a}				
Operands:	$0 \leq f \leq 255$ $0 \leq b \leq 7$ $a \in [0,1]$				
Operation:	$1 \rightarrow f \langle b \rangle$				
Status Affected:	None				
Encoding:	<table border="1"><tr><td>1000</td><td>bbba</td><td>ffff</td><td>ffff</td></tr></table>	1000	bbba	ffff	ffff
1000	bbba	ffff	ffff		
Description:	<p>Bit 'b' in register 'f' is set.</p> <p>If 'a' is '0', the Access Bank is selected.</p> <p>If 'a' is '1', the BSR is used to select the GPR bank.</p> <p>If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See Section 42.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode” for details.</p>				

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: BSF FLAG_REG, 7, 1

Before Instruction
FLAG_REG = 0Ah

After Instruction
FLAG_REG = 8Ah

CALLW Subroutine Call Using WREG

Syntax:	CALLW			
Operands:	None			
Operation:	(PC + 2) → TOS, (W) → PCL, (PCLATH) → PCH, (PCLATU) → PCU			
Status Affected:	None			
Encoding:	0000	0000	0001	0100
Description	<p>First, the return address (PC + 2) is pushed onto the return stack. Next, the contents of W are written to PCL; the existing value is discarded. Then, the contents of PCLATH and PCLATU are latched into PCH and PCU, respectively. The second cycle is executed as a NOP instruction while the new next instruction is fetched. Unlike CALL, there is no option to update W, Status or BSR.</p>			
Words:	1			
Cycles:	2			
Q Cycle Activity:				

Q1	Q2	Q3	Q4
Decode	Read WREG	PUSH PC to stack	No operation
No operation	No operation	No operation	No operation

Example: HERE CALLW

Before Instruction

PC = address (HERE)
PCLATH = 10h
PCLATU = 00h
W = 06h

After Instruction

PC = 001006h
TOS = address (HERE + 2)
PCLATH = 10h
PCLATU = 00h
W = 06h

CLRF Clear f

Syntax:	CLRF f{,a}				
Operands:	$0 \leq f \leq 255$ $a \in [0,1]$				
Operation:	$000h \rightarrow f$ $1 \rightarrow Z$				
Status Affected:	Z				
Encoding:	<table><tr><td>0110</td><td>101a</td><td>ffff</td><td>ffff</td></tr></table>	0110	101a	ffff	ffff
0110	101a	ffff	ffff		
Description:	<p>Clears the contents of the specified register.</p> <p>If 'a' is '0', the Access Bank is selected.</p> <p>If 'a' is '1', the BSR is used to select the GPR bank.</p> <p>If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See Section 42.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode” for details.</p>				
Words:	1				
Cycles:	1				
Q Cycle Activity:					

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example: CLRF FLAG_REG, 1

Before Instruction

FLAG_REG = 5Ah

After Instruction

FLAG_REG = 00h

NEGF

Negate f

Syntax:	NEGF f{,a}			
Operands:	$0 \leq f \leq 255$ $a \in [0,1]$			
Operation:	$(\bar{f}) + 1 \rightarrow f$			
Status Affected:	N, OV, C, DC, Z			
Encoding:	0110	110a	ffff	ffff
Description:	<p>Location 'f' is negated using two's complement. The result is placed in the data memory location 'f'.</p> <p>If 'a' is '0', the Access Bank is selected.</p> <p>If 'a' is '1', the BSR is used to select the GPR bank.</p> <p>If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \leq 95$ (5Fh). See Section 42.2.3 “Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode” for details.</p>			
Words:	1			
Cycles:	1			

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	Write register 'f'

Example:

NEGF REG, 1

Before Instruction

REG = 0011 1010 [3Ah]

After Instruction

REG = 1100 0110 [C6h]

NOP

No Operation

Syntax:	NOP											
Operands:	None											
Operation:	No operation											
Status Affected:	None											
Encoding:	<table><tr><td>0000</td><td>0000</td><td>0000</td><td>0000</td></tr><tr><td>1111</td><td>xxxx</td><td>xxxx</td><td>xxxx</td></tr></table>	0000	0000	0000	0000	1111	xxxx	xxxx	xxxx			
0000	0000	0000	0000									
1111	xxxx	xxxx	xxxx									
Description:	No operation.											
Words:	1											
Cycles:	1											

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	No operation	No operation

Example:

None.

44.2 MPLAB XC Compilers

The MPLAB XC Compilers are complete ANSI C compilers for all of Microchip's 8, 16, and 32-bit MCU and DSC devices. These compilers provide powerful integration capabilities, superior code optimization and ease of use. MPLAB XC Compilers run on Windows, Linux or MAC OS X.

For easy source level debugging, the compilers provide debug information that is optimized to the MPLAB X IDE.

The free MPLAB XC Compiler editions support all devices and commands, with no time or memory restrictions, and offer sufficient code optimization for most applications.

MPLAB XC Compilers include an assembler, linker and utilities. The assembler generates relocatable object files that can then be archived or linked with other relocatable object files and archives to create an executable file. MPLAB XC Compiler uses the assembler to produce its object file. Notable features of the assembler include:

- Support for the entire device instruction set
- Support for fixed-point and floating-point data
- Command-line interface
- Rich directive set
- Flexible macro language
- MPLAB X IDE compatibility

44.3 MPASM Assembler

The MPASM Assembler is a full-featured, universal macro assembler for PIC10/12/16/18 MCUs.

The MPASM Assembler generates relocatable object files for the MPLINK Object Linker, Intel® standard HEX files, MAP files to detail memory usage and symbol reference, absolute LST files that contain source lines and generated machine code, and COFF files for debugging.

The MPASM Assembler features include:

- Integration into MPLAB X IDE projects
- User-defined macros to streamline assembly code
- Conditional assembly for multipurpose source files
- Directives that allow complete control over the assembly process

44.4 MPLINK Object Linker/ MPLIB Object Librarian

The MPLINK Object Linker combines relocatable objects created by the MPASM Assembler. It can link relocatable objects from precompiled libraries, using directives from a linker script.

The MPLIB Object Librarian manages the creation and modification of library files of precompiled code. When a routine from a library is called from a source file, only the modules that contain that routine will be linked in with the application. This allows large libraries to be used efficiently in many different applications.

The object linker/librarian features include:

- Efficient linking of single libraries instead of many smaller files
- Enhanced code maintainability by grouping related modules together
- Flexible creation of libraries with easy module listing, replacement, deletion and extraction

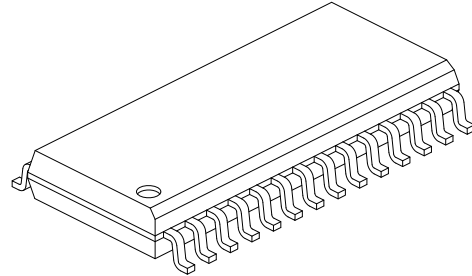
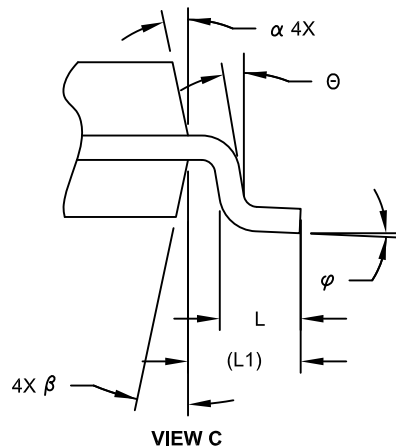
44.5 MPLAB Assembler, Linker and Librarian for Various Device Families

MPLAB Assembler produces relocatable machine code from symbolic assembly language for PIC24, PIC32 and dsPIC DSC devices. MPLAB XC Compiler uses the assembler to produce its object file. The assembler generates relocatable object files that can then be archived or linked with other relocatable object files and archives to create an executable file. Notable features of the assembler include:

- Support for the entire device instruction set
- Support for fixed-point and floating-point data
- Command-line interface
- Rich directive set
- Flexible macro language
- MPLAB X IDE compatibility

28-Lead Plastic Small Outline (SO) - Wide, 7.50 mm Body [SOIC]

Note: For the most current package drawings, please see the Microchip Packaging Specification located at <http://www.microchip.com/packaging>



Dimension	Units	MILLIMETERS		
		MIN	NOM	MAX
Number of Pins	N	28		
Pitch	e	1.27 BSC		
Overall Height	A	-	-	2.65
Molded Package Thickness	A2	2.05	-	-
Standoff §	A1	0.10	-	0.30
Overall Width	E	10.30 BSC		
Molded Package Width	E1	7.50 BSC		
Overall Length	D	17.90 BSC		
Chamfer (Optional)	h	0.25	-	0.75
Foot Length	L	0.40	-	1.27
Footprint	L1	1.40 REF		
Lead Angle	θ	0°	-	-
Foot Angle	φ	0°	-	8°
Lead Thickness	c	0.18	-	0.33
Lead Width	b	0.31	-	0.51
Mold Draft Angle Top	α	5°	-	15°
Mold Draft Angle Bottom	β	5°	-	15°

Notes:

- Pin 1 visual index feature may vary, but must be located within the hatched area.
- § Significant Characteristic
- Dimension D does not include mold flash, protrusions or gate burrs, which shall not exceed 0.15 mm per end. Dimension E1 does not include interlead flash or protrusion, which shall not exceed 0.25 mm per side.
- Dimensioning and tolerancing per ASME Y14.5M
 - BSC: Basic Dimension. Theoretically exact value shown without tolerances.
 - REF: Reference Dimension, usually without tolerance, for information purposes only.
- Datums A & B to be determined at Datum H.

Microchip Technology Drawing C04-052C Sheet 2 of 2