Freescale Semiconductor - T4160NSE7TTB Datasheet

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Understanding Embedded - Microprocessors

Embedded microprocessors are specialized computing chips designed to perform specific tasks within an embedded system. Unlike general-purpose microprocessors found in personal computers, embedded microprocessors are tailored for dedicated functions within larger systems, offering optimized performance, efficiency, and reliability. These microprocessors are integral to the operation of countless electronic devices, providing the computational power necessary for controlling processes, handling data, and managing communications.

Applications of Embedded - Microprocessors

Embedded microprocessors are utilized across a broad spectrum of applications, making them indispensable in

Details	
Product Status	Active
Core Processor	-
Number of Cores/Bus Width	-
Speed	-
Co-Processors/DSP	-
RAM Controllers	-
Graphics Acceleration	-
Display & Interface Controllers	-
Ethernet	-
SATA	-
USB	-
Voltage - I/O	-
Operating Temperature	-
Security Features	-
Package / Case	-
Supplier Device Package	-
Purchase URL	https://www.e-xfl.com/pro/item?MUrl=&PartUrl=t4160nse7ttb

The Frame Manager (FMAN), a primary element of the DPAA, parses headers from incoming packets and classifies and selects data buffers with optional policing and congestion management. The FMAN passes its work to the Queue Manager (QMAN), which assigns it to cores or accelerators with a multilevel scheduling hierarchy. The T4240 processor's implementation of the DPAA offers accelerations for cryptography, enhanced regular expression pattern matching, and compression/decompression.

2.4 System peripherals and networking

For networking, there are dual FMANs with an aggregate of up to 16 any-speed MAC controllers that connect to PHYs, switches, and backplanes over RGMII, SGMII, QSGMII, HiGig2, XAUI, XFI, and 10Gbase-KR. The FMAN also supports new quality of service features through egress traffic shaping and priority flow control for data center bridging in converged data center networking applications. High-speed system expansion is supported through four PCI Express controllers that support varieties of lane lengths for PCIe specification 3.0, including endpoint SR-IOV with 128 virtual functions. Other peripherals include:

- SRIO
- Interlaken-LA
- SATA
- SD/MMC
- I^2C
- UART
- SPI
- NOR/NAND controller
- GPIO
- 1866 MT/s DDR3/L controller

3 Application examples

This chip is well-suited for applications that are highly compute-intensive, I/O-intensive, or both.

3.1 1U security appliance

This figure shows a 1U security appliance built around a single SoC. The QorIQ DPAA accelerates basic packet classification, filtering, and packet queuing, while the crypto accelerator (SEC 5.0), regex accelerator (PME 2.1), and compression/decompression accelerator (DCE 1.0) perform high throughput content processing. The high single threaded and aggregate DMIPS of the core CPUs provide the processing horsepower for complex classification and flow state tracking required for proxying applications as well as heuristic traffic analysis and policy enforcement.

The SoC's massive integration significantly reduces system BOM cost. SATA hard drives connect directly to the SoC's integrated controllers, and an Ethernet switch is only required if more than 16 1 GE ports or 4 10 GE ports are required. The SoC supports PCIe and Serial RapidIO for expansion.

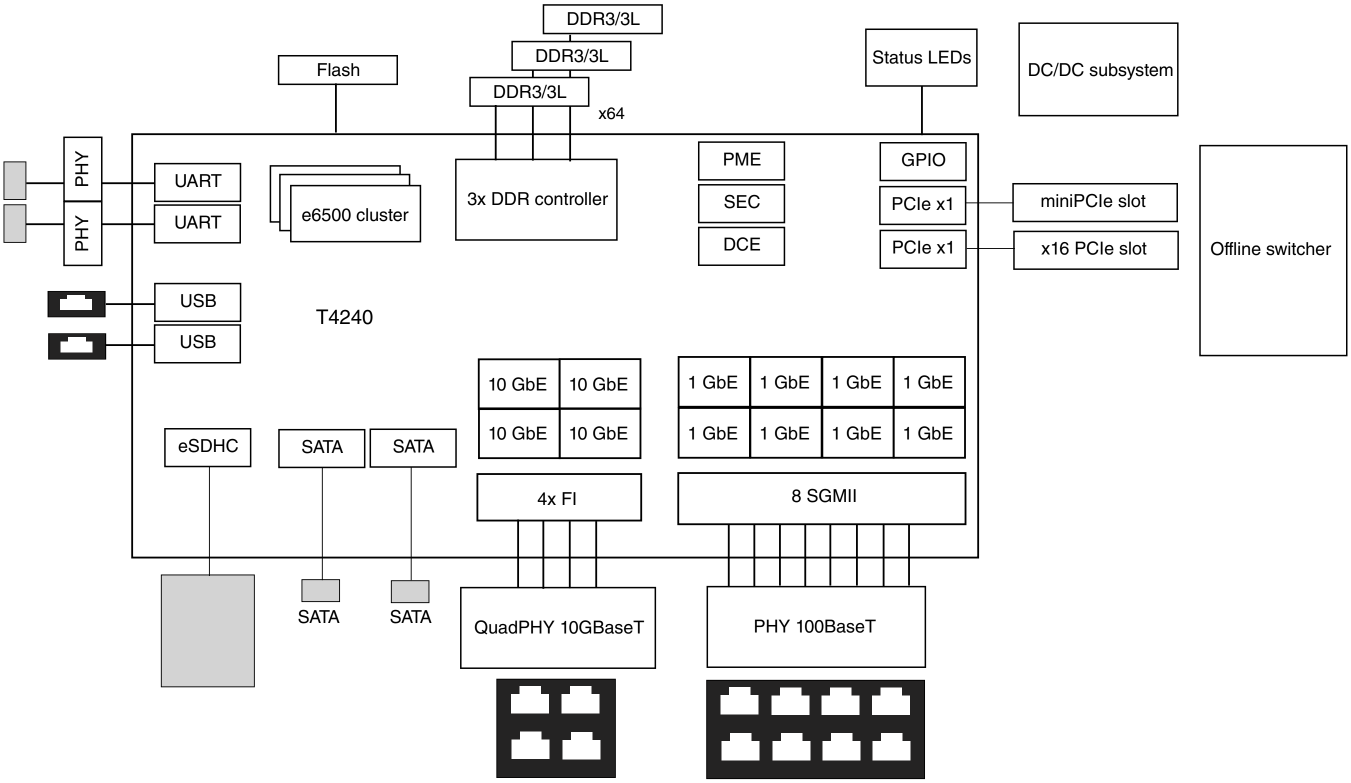

Figure 1. SoC 1U security appliance

3.2 Rack-mounted services blade

Networking and telecom systems are frequently modular in design, built from multiple standard dimension blades, which can be progressively added to a chassis to increase interface bandwidth or processing power. ATCA is a common standard form factor for chassis-based systems.

This figure shows a potential configuration for an ATCA blade with four chips and an Ethernet switch, which provides connectivity to the front panel and backplane, as well as between the chips. Potential systems enabled by chips in ATCA style modular architectures are described below.

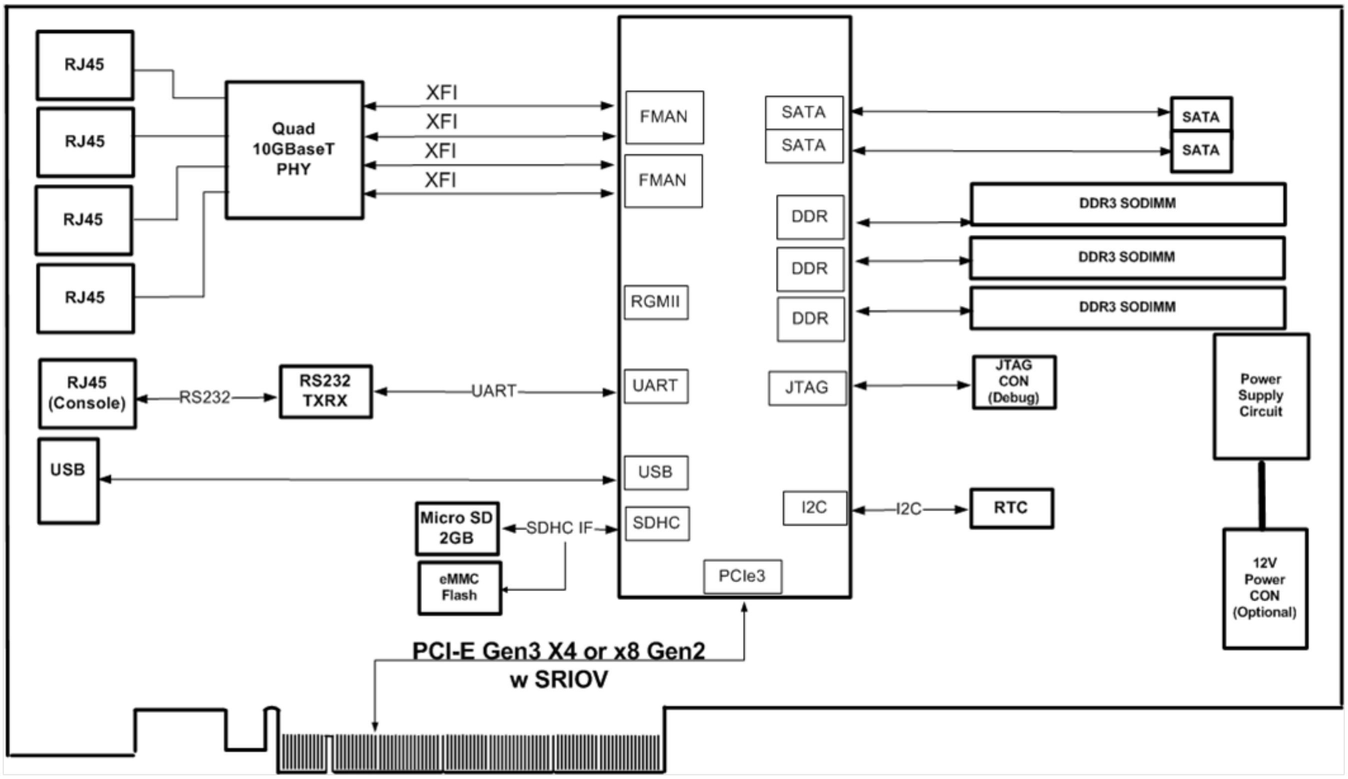

Figure 4. Intelligent network adapter

4 Multicore processing options

This flexible chip can be configured to meet many system application needs. The chip's CPUs (and hardware threads as virtual CPUs) can be combined as a fully-symmetric, multiprocessing, system-on-a-chip, or they can be operated with varying degrees of independence to perform asymmetric multiprocessing. High levels of processor independence, including the ability to independently boot and reset each core, is characteristic of the chip. The ability of the cores to run different operating systems, or run OS-less, provides the user with significant flexibility in partitioning between control, datapath, and applications processing. It also simplifies consolidation of functions previously spread across multiple discrete processors onto a single device.

While up to 24 Power Architecture threads (henceforth referred to as 'virtual CPUs', or 'vCPUs') offer a large amount of total, available computing performance, raw processing power is not enough to achieve multi-Gbps data rates in high-touch networking and telecom applications. To address this, this chip enhances the Freescale Data Path Acceleration Architecture (DPAA), further reducing data plane instructions per packet, and enabling more CPU cycles to work on value-added services as opposed to repetitive, low-level tasks. Combined with specialized accelerators for cryptography, pattern matching, and compression, the chip allows the user's software to perform complex packet processing at high data rates. There are many ways to map operating systems and I/O up to 24 chip vCPUs.

4.1 Asymmetric multiprocessing

As shown in this figure, the chip's vCPUs can be used in an asymmetric multi-processing model, with *n* copies of the same uni-processor OS, or *n* copies of OS 1, *n* copies of OS 2, and so on, up to 24 OS instances. The DPAA distributes work to the specific vCPUs based on basic classification or it puts work onto a common queue from which any vCPU can dequeue work.

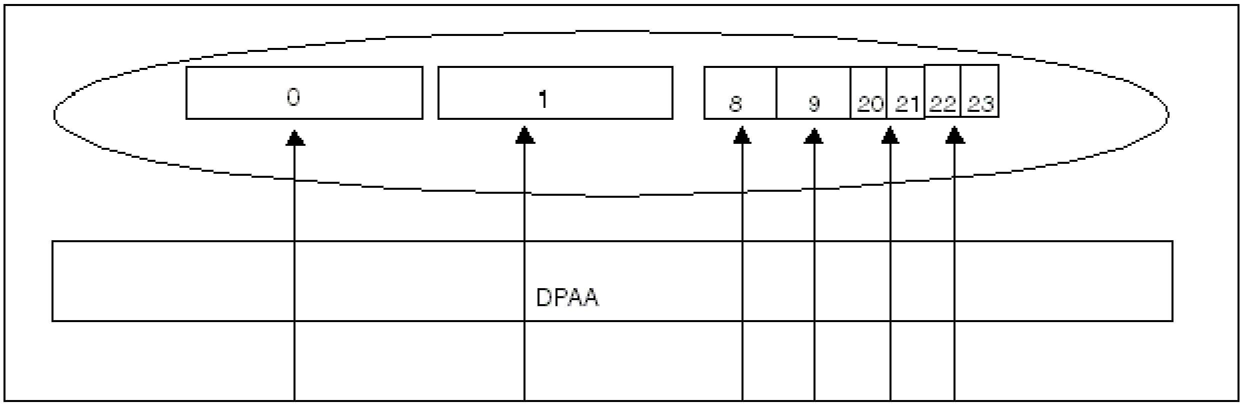

Figure 7. Mixed SMP and AMP option 2

5 Chip features

This section describes the key features and functionalities of the T4240 chip. See the T4160 and T4080 appendices for those device's specific block diagrams.

5.1 Block diagram

This figure shows the major functional units within the chip.

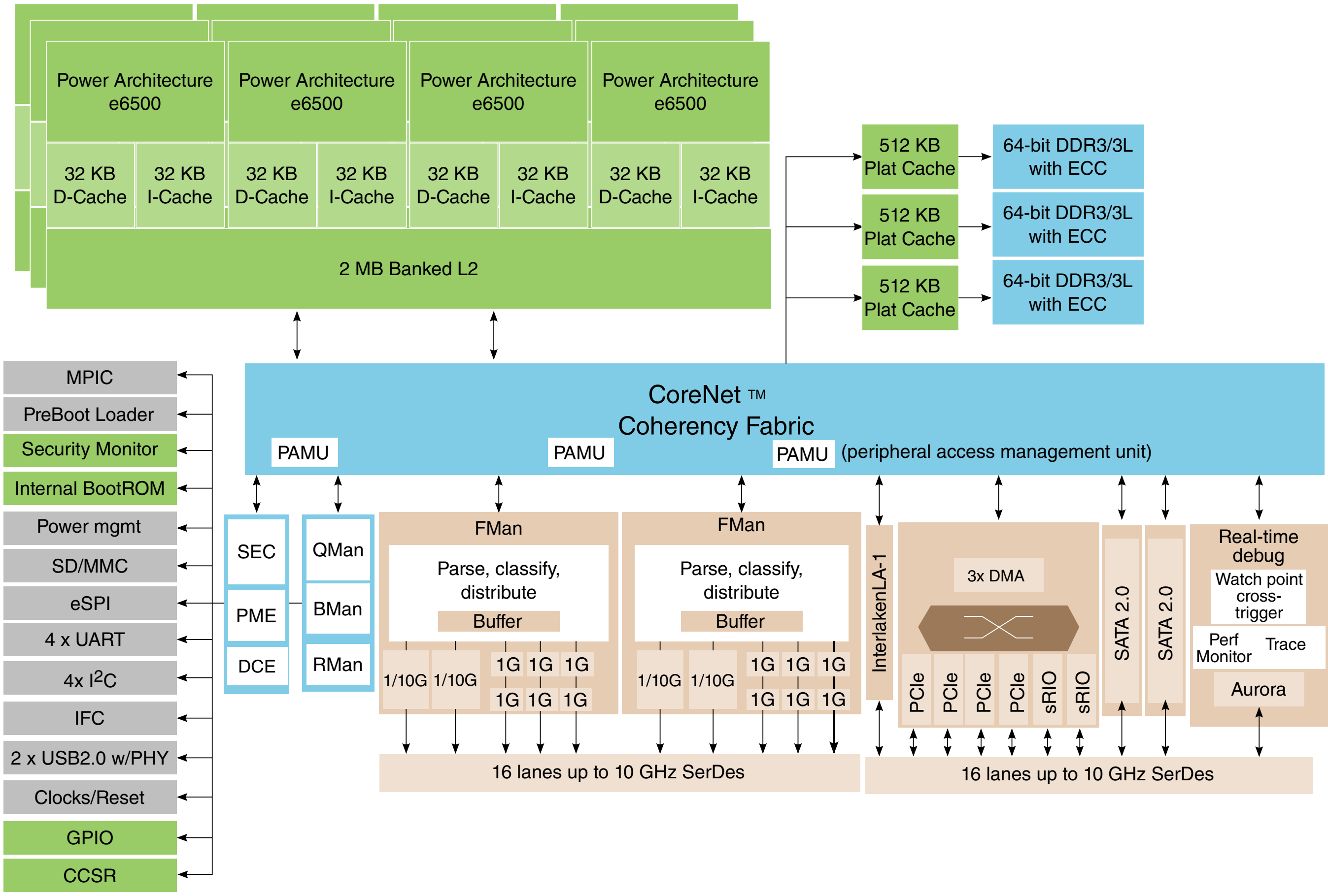

Figure 8. T4240 block diagram

5.2 Features summary

This chip includes the following functions and features:

- 12, dual-threaded e6500 cores for a total of 24/16/8 threads (T4240/T4160/T4080) built on Power Architecture® technology
    - Arranged as three clusters of four cores sharing a 2 MB L2 cache, 6 MB L2 cache total.
    - Up to 1.8 GHz with 64-bit ISA support (Power Architecture v2.06-compliant)
    - Three privilege levels of instruction: user, supervisor, and hypervisor
- Up to 1.5 MB CoreNet Platform Cache (CPC)
- Hierarchical interconnect fabric
    - CoreNet fabric supporting coherent and non-coherent transactions with prioritization and bandwidth allocation amongst CoreNet end-points
    - 1.46 Tbps coherent read bandwidth
- Up to three 64-bit DDR3/3L SDRAM memory controllers with ECC and interleaving support
    - Up to 1.867 GT/s data transfer rate
    - 64 GB per DDR controller
- Data Path Acceleration Architecture (DPAA) incorporating acceleration for the following functions:
    - Packet parsing, classification, and distribution (Frame Manager 1.1) up to 50 Gbps
    - Queue management for scheduling, packet sequencing, and congestion management (Queue Manager 1.1)
    - Queue Manager (QMan) fabric supporting packet-level queue management and quality of service scheduling
    - Hardware buffer management for buffer allocation and de-allocation (BMan 1.1)
    - Cryptography acceleration (SEC 5.0) at up to 40 Gbps

- RegEx Pattern Matching Acceleration (PME 2.1) at up to 10 Gbps
- Decompression/Compression Acceleration (DCE 1.0) at up to 20 Gbps
- DPAA chip-to-chip interconnect via RapidIO Message Manager (RMAN 1.0)
- Up to 32 SerDes lanes at up to 10.3125 GHz
- Ethernet interfaces
  - Up to four 10 Gbps Ethernet XAUI or 10GBase-KR XFI MACs
  - Up to sixteen 1 Gbps Ethernet MACs
  - Up to two 1Gbps Ethernet RGMII MACs
  - Maximum configuration of 4 x 10 GE (XFI) + 10 x 1 GE (SGMII) + 2 x 1 GE (RGMII)
- High-speed peripheral interfaces
  - Up to four PCI Express 2.0 controllers, two supporting 3.0
  - Two Serial RapidIO 2.0 controllers/ports running at up to 5 GHz with Type 11 messaging and Type 9 data streaming support
  - Interlaken look-aside interface for serial TCAM connection at 6.25 and 10.3125 Gbps per-lane rates.
- Additional peripheral interfaces
  - Two serial ATA (SATA 2.0) controllers
  - Two high-speed USB 2.0 controllers with integrated PHY
  - Enhanced secure digital host controller (SD/MMC/eMMC)
  - Enhanced serial peripheral interface (eSPI)
  - Four I2C controllers
  - Four 2-pin or two 4-pin UARTs
  - Integrated Flash controller supporting NAND and NOR flash
- Three eight-channel DMA engines.
- Support for hardware virtualization and partitioning enforcement
- QorIQ Platform's Trust Architecture 2.0

5.3 Critical performance parameters

This table lists key performance indicators that define a set of values used to measure SoC operation.

Table 1. Critical performance parameters

Indicator	Values(s)
Top speed bin core frequency	1.8 GHz
Maximum memory data rate	1867 MHz (DDR3)[1], 1600 MHz for DDR3L • 1.5 V for DDR3 • 1.35 V for DDR3L
Integrated flash controller (IFC)	1.8 V
Operating junction temperature range	0-105 C
Package	1932-pin, flip-chip plastic ball grid array (FC-PBGA), 45 x 45mm

1. Conforms to JEDEC standard

5.4 Core and CPU clusters

This chip offers 12, high-performance, 64-bit Power Architecture, Book E-compliant cores. Each CPU core supports two hardware threads, which software views as a virtual CPU. The core CPUs are arranged in clusters of four with a shared 2 MB L2 cache.

This table shows the computing metrics the core supports.

Table 2. Power architecture metrics

Metric	Per core	Per cluster	Full device
DMIPS	10,800	43,200	129,600
Single-precision GFLOPs	18	72	Up to 216
Double-precision GFLOPs	3.6	14.4	Up to 42.4

The core subsystem includes the following features:

- Up to 1.8 GHz
- Dual-thread with simultaneous multi-threading (SMT)
    - Threading can be disabled on a per CPU basis
- 40-bit physical addressing
- L2 MMU
    - Supporting 4 KB pages
    - TLB0; 8-way set-associative, 1024-entries (4 KB pages)
    - TLB1; fully associative, 64-entry, supporting variable size pages and indirect page table entries
- Hardware page table walk
- 64-byte cache line size
- L1 caches, running at core frequency
    - 32 KB instruction, 8-way set-associative
    - 32 KB data, 8-way set-associative
    - Each with data and tag parity protection
- Hardware support for memory coherency
- Five integer units: 4 simple (2 per thread), 1 complex (integer multiply and divide)
- Two load-store units: one per thread
- Classic double-precision floating-point unit
    - Uses 32 64-bit floating-point registers (FPRs) for scalar single- and double-precision floating-point arithmetic
    - Designed to comply with IEEE Std. 754™-1985 FPU for both single and double-precision operations
- AltiVec unit
    - 128-bit Vector SIMD engine
    - 32 128-bit VR registers
    - Operates on a vector of
        - Four 32-bit integers
        - Four 32-bit single precision floating-point units
        - Eight 16-bit integers
        - Sixteen 8-bit integers
    - Powerful permute unit
    - Enhancements include: Move from GPRs to VR, sum of absolute differences operation, extended support for misaligned vectors, handling head and tails of vectors
- Supports Data Path Acceleration Architecture (DPAA) data and context "stashing" into L1 and L2 caches
- User, supervisor, and hypervisor instruction level privileges
- Addition of Elemental Barriers and "wait on reservation" instructions
- New power-saving modes including "drowsy core" with state retention and nap
    - State retention power-saving mode allows core to quickly wake up and respond to service requests
- Processor facilities
    - Hypervisor APU
    - "Decorated Storage" APU for improved statistics support
        - Provides additional atomic operations, including a "fire-and-forget" atomic update of up to two 64-bit quantities by a single access
    - Addition of Logical to Real Address translation mechanism (LRAT) to accelerate hypervisor performance
    - Expanded interrupt model

5.6 CoreNet fabric and address map

The CoreNet fabric provides the following:

- A highly concurrent, fully cache coherent, multi-ported fabric
- Point-to-point connectivity with flexible protocol architecture allows for pipelined interconnection between CPUs, platform caches, memory controllers, and I/O and accelerators at up to 733 MHz
- The CoreNet fabric has been designed to overcome bottlenecks associated with shared bus architectures, particularly address issue and data bandwidth limitations. The chip's multiple, parallel address paths allow for high address bandwidth, which is a key performance indicator for large coherent multicore processors.
- Eliminates address retries, triggered by CPUs being unable to snoop within the narrow snooping window of a shared bus. This results in the chip having lower average memory latency.

This chip's 40-bit, physical address map consists of local space and external address space. For the local address map, 32 local access windows (LAWs) define mapping within the local 40-bit (1 TB) address space. Inbound and outbound translation windows can map the chip into a larger system address space such as the RapidIO or PCIe 64-bit address environment. This functionality is included in the address translation and mapping units (ATMUs).

5.7 Memory complex

The SoC's memory complex consists of up to three DDR controllers for main memory, and the memory controllers associated with the Integrated Flash Controller (IFC).

5.7.1 DDR memory controllers

The chip offers up to three 64-bit DDR controllers supporting ECC protected memories. These DDR controllers operate at up to 1.867 GT/s for DDR3, and, in more power sensitive applications, up to 1.6 GHz for DDR3L. Some key DDR controller features are as follows:

- Interleaving options
  - None, three fully independent controllers
  - Two interleaved, one independent
  - Three interleaved
  - Interleaving can be configured on 1 KB, 4 KB, and 8 KB granules
- Support x4, x8, and x16 memory widths
  - Programmable support for single, dual, and quad ranked devices and modules
  - Support for both unbuffered and registered DIMMs
  - 4 chip-selects per controller
  - 64 GB per controller, 192 GB per chip
- The SoC can be configured to retain the currently active SDRAM page for pipelined burst accesses. Page mode support of up to 64 simultaneously open pages can dramatically reduce access latencies for page hits. Depending on the memory system design and timing parameters, page mode can save up to ten memory clock cycles for subsequent burst accesses that hit in an active page.
- Using ECC, the SoC detects and corrects all single-bit errors and detects all double-bit errors and all errors within a nibble.
- Upon detection of a loss of power signal from external logic, the DDR controllers can put compliant DDR SDRAM DIMMs into self-refresh mode, allowing systems to implement battery-backed main memory protection.
- In addition, the DDR controllers offer an initialization bypass feature for use by system designers to prevent re-initialization of main memory during system power-on after an abnormal shutdown.
- Support active zeroization of system memory upon detection of a user-defined security violation.

5.7.1.1 DDR bandwidth optimizations

Multicore SoCs are able to increase CPU and network interface bandwidths faster than commodity DRAM technologies are improving. As a result, it becomes increasingly important to maximize utilization of main memory interfaces to avoid a memory bottleneck. The T4 family's DDR controllers are Freescale-developed IP, optimized for the QorIQ SoC architecture, with the goal of improving DDR bandwidth utilization by fifty percent when compared to first generation QorIQ SoCs.

Most of the WRITE bandwidth improvement and approximately half of the READ bandwidth improvement is met through target queue enhancements; in specific, changes to the scheduling algorithm, improvements in the bank hashing scheme, support for more transaction re-ordering, and additional proprietary techniques.

The remainder of the READ bandwidth improvement is due to the addition of an intelligent data prefetcher in the memory subsystem.

5.7.1.2 Prefetch Manager (PMan)

NOTE

All transactions to DDR pass through the CPC; this means the CPC can miss (and trigger prefetching) even on data that is not intended for allocation into the CPC.

The PMAN monitors CPC misses for opportunities to prefetch, using a "confidence"-based algorithm to determine its degree of aggressiveness. It can be configured to monitor multiple memory regions (each of different size) for prefetch opportunities. Multiple CPC misses on accesses to a tracked region for consecutive cache blocks increases confidence to start prefetching, and a CPC miss of a tracked region with same stride will instantly cause prefetching.

The PMan uses feedback to increase or decrease its aggressiveness. When the data it prefetches is being used, it prefetches further ahead. If the request stride length changes or previously prefetched data isn't consumed, prefetching slows or stops (at least for that region/requesting device/transaction type).

5.7.2 PreBoot Loader and nonvolatile memory interfaces

The PreBoot Loader (PBL) operates similarly to an I^2C boot sequencer but on behalf of a large number of interfaces.

It supports IFC, I^2C, eSPI, eSDHC.

The PBL's functions include the following:

- Simplifies boot operations, replacing pin strapping resistors with configuration data loaded from nonvolatile memory
- Uses the configuration data to initialize other system logic and to copy data from low speed memory interfaces (I^2C, IFC, eSPI, and SD/MMC) into fully initialized DDR or the 2 MB front-side cache

5.7.2.1 Integrated Flash Controller

The SoC incorporates an Integrated Flash Controller similar to the one used in some previous generation QorIQ SoCs. The IFC supports both NAND and NOR flash, as well as a general purpose memory mapped interface for connecting low speed ASICs and FPGAs.

5.7.2.1.1 NAND Flash features

- x8/x16 NAND Flash interface
- Optional ECC generation/checking
- Flexible timing control to allow interfacing with proprietary NAND devices
- SLC and MLC Flash devices support with configurable page sizes of up to 4 KB
- Support advance NAND commands like cache, copy-back, and multiplane programming

- Boot chip-select (CS0) available after system reset, with boot block size of 8 KB, for execute-in-place boot loading from NAND Flash
- Up to terabyte Flash devices supported

5.7.2.1.2 NOR Flash features

- Data bus width of 8/16/32
- Compatible with asynchronous NOR Flash
- Directly memory mapped
- Supports address data multiplexed (ADM) NOR device
- Flexible timing control allows interfacing with proprietary NOR devices
- Boot chip-select (CS0) available at system reset

5.7.2.1.3 General-purpose chip-select machine (GPCM)

The IFC's GPCM supports the following features:

- Normal GPCM
  - Support for x8/16/32-bit device
  - Compatible with general purpose addressable device, for example, SRAM and ROM
  - External clock is supported with programmable division ratio (2, 3, 4, and so on, up to 16)
- Generic ASIC Interface
  - Support for x8/16/32-bit device
  - Address and Data are shared on I/O bus
  - Following address and data sequences are supported on I/O bus:
    - 32-bit I/O: AD
    - 16-bit I/O: AADD
    - 8-bit I/O: AAAADDDD

5.7.2.2 Serial memory controllers

In addition to the parallel NAND and NOR flash supported by the IFC, the SoC supports serial flash using eSPI, I^2C and SD/MMC/eMMC card and device interfaces. The SD/MMC/eMMC controller includes a DMA engine, allowing it to move data from serial flash to external or internal memory following straightforward initiation by software.

Detailed features of the eSDHC include the following:

- Conforms to the SD Host Controller Standard Specification version 2.0, including Test event register support
- Compatible with the MMC System Specification version 4.2
- Compatible with the SD Memory Card Specification version 2.0, and supports the high capacity SD memory card
- Designed to work with SD memory, SD combo, MMC, and their variants like mini and micro.
- Card bus clock frequency up to 52 MHz
- Supports 1-/4-bit SD, 1-/4-/8-bit MMC modes
- Supports single-block and multi-block read, and write data transfer
- Supports block sizes of 1-2048 bytes
- Supports the mechanical write protect detection. In the case where write protect is enabled, the host will not initiate any write data command to the card
- Supports both synchronous and asynchronous abort
- Supports pause during the data transfer at block gap
- Supports Auto CMD12 for multi-block transfer
- Host can initiate command that do not use data lines, while data transfer is in progress
- Embodies a configurable 128x32-bit FIFO for read/write data
- Supports SDMA, ADMA1, and ADMA2 capabilities

5.9.2 Serial RapidIO

The Serial RapidIO interface is based on the *RapidIO Interconnect Specification, Revision 2.1* . RapidIO is a high-performance, point-to-point, low-pin-count, packet-switched system-level interconnect that can be used in a variety of applications as an open standard. The rich feature set includes high data bandwidth, low-latency capability, and support for high-performance I/O devices as well as message-passing and software-managed programming models. Receive and transmit ports operate independently, and with 2 x 4 Serial RapidIO controllers, the aggregate theoretical bandwidth is 32 Gbps.

The chip offers two Serial RapidIO controllers, muxed onto the SerDes blocks. The Serial RapidIO interface is based on the *RapidIO Interconnect Specification, Revision 2.1* . Receive and transmit ports operate independently and with 2 x 4 Serial RapidIO controllers; the aggregate theoretical bandwidth is 32 Gbps. The Serial RapidIO controllers can be used in conjunction with "Rapid IO Message Manager (RMAN), as described in RapidIO Message Manager (RMan)."

Key features of the Serial RapidIO interface unit include the following:

- Support for *RapidIO Interconnect Specification, Revision 2.1* (All transaction flows and priorities.)
- 2x, and 4x LP-serial link interfaces, with transmission rates of 2.5, 3.125, or 5.0 Gbaud (data rates of 1.0, 2.0, 2.5, or 4.0 Gbps) per lane
- Auto-detection of 1x, 2x, or 4x mode operation during port initialization
- 34-bit addressing and up to 256-byte data payload
- Support for SWRITE, NWRITE, NWRITE_R and Atomic transactions
- Receiver-controlled flow control
- RapidIO error injection
- Internal LP-serial and application interface-level loopback modes

The Serial RapidIO controller also supports the following capabilities, many of which are leveraged by the RMan to efficient chip-to-chip communication through the DPAA:

- Support for RapidIO Interconnect Specification 2.1, "Part 2: Message Passing Logical Specification"
- Supports RapidIO Interconnect Specification 2.1, "Part 10: Data Streaming Logical Specification"
- Supports RapidIO Interconnect Specification 2.1, "Annex 2: Session Management Protocol"
    - Supports basic stream management flow control (XON/XOFF) using extended header message format
- Up to 16 concurrent inbound reassembly operations
    - One additional reassembly context is reservable to a specific transaction type
- Support for outbound Type 11 messaging
- Support for outbound Type 5 NWRITE and Type 6 SWRITE transactions
- Support for inbound Type 11 messaging
- Support for inbound Type 9 data streaming transactions
- Support for outbound Type 9 data streaming transactions
    - Up to 64 KB total payload
- Support for inbound Type 10 doorbell transactions
    - Transaction steering through doorbell header classification
- Support for outbound Type 10 doorbell transactions
    - Ordering can be maintained with respect to other types of traffic.
- Support for inbound and outbound port-write transactions
    - Data payloads of 4 to 64 bytes

5.9.3 SATA

Each of the SoC's two SATA controllers is compliant with the *Serial ATA 2.6 Specification*. Each of the SATA controllers has the following features:

- Supports speeds: 1.5 Gbps (first-generation SATA), and 3Gbps (second-generation SATA)
- Supports advanced technology attachment packet interface (ATAPI) devices
- Contains high-speed descriptor-based DMA controller
- Supports native command queuing (NCQ) commands

5.10.1 Packet distribution and queue/congestion management

This table lists some packet distribution and queue/congestion management offload functions.

Table 3. Offload functions

Function type	Definition
Data buffer management	Supports allocation and deallocation of buffers belonging to pools originally created by software with configurable depletion thresholds. Implemented in a module called the Buffer Manager (BMan).
Queue management	Supports queuing and quality-of-service scheduling of frames to CPUs, network interfaces and DPAA logic blocks, maintains packet ordering within flows. Implemented in a module called the Queue Manager (QMan). The QMan, besides providing flow-level queuing, is also responsible for congestion management functions such as RED/WRED, congestion notifications and tail discards.
Packet distribution	Supports in-line packet parsing and general classification to enable policing and QoS-based packet distribution to the CPUs for further processing of the packets. This function is implemented in the block called the Frame Manager (FMan).
Policing	Supports in-line rate-limiting by means of two-rate, three-color marking (RFC 2698). Up to 256 policing profiles are supported. This function is also implemented in the FMan.
Egress Scheduling	Supports hierarchical scheduling and shaping, with committed and excess rates. This function is supported in the QMan, although the FMan performs the actual transmissions.

5.10.2 Accelerating content processing

Properly implemented acceleration logic can provide significant performance advantages over most optimized software with acceleration factors on the order of 10-100x. Accelerators in this category typically touch most of the bytes of a packet (not just headers). To avoid consuming CPU cycles in order to move data to the accelerators, these engines include well-pipelined DMAs. This table lists some specific content-processing accelerators on the chip.

Table 4. Content-processing accelerators

Interface	Definition
SEC	Crypto-acceleration for protocols such as IPsec, SSL, and 3GPP RLC
PME	Regex style pattern matching for unanchored searches, including cross-packet stateful patterns
DCE	Compression/Decompression acceleration for ZLib and deflate

5.10.3 Enhancements of T4240 compared to first generation DPAA

A short summary of T4240 enhancements over the first generation DPAA (as implemented in the P4080) is provided below:

- Frame Manager
    - 2x performance increase (up to 25 Gbps per FMan)
    - Storage profiles.
    - HiGig (3.125 GHz) and HiGig2 (3.125 GHz and 3.75 GHz)
    - Energy Efficient Ethernet
- SEC 5.0
    - 2x performance increase for symmetric encryption and protocol processing

This figure is a logical view of the DPAA.

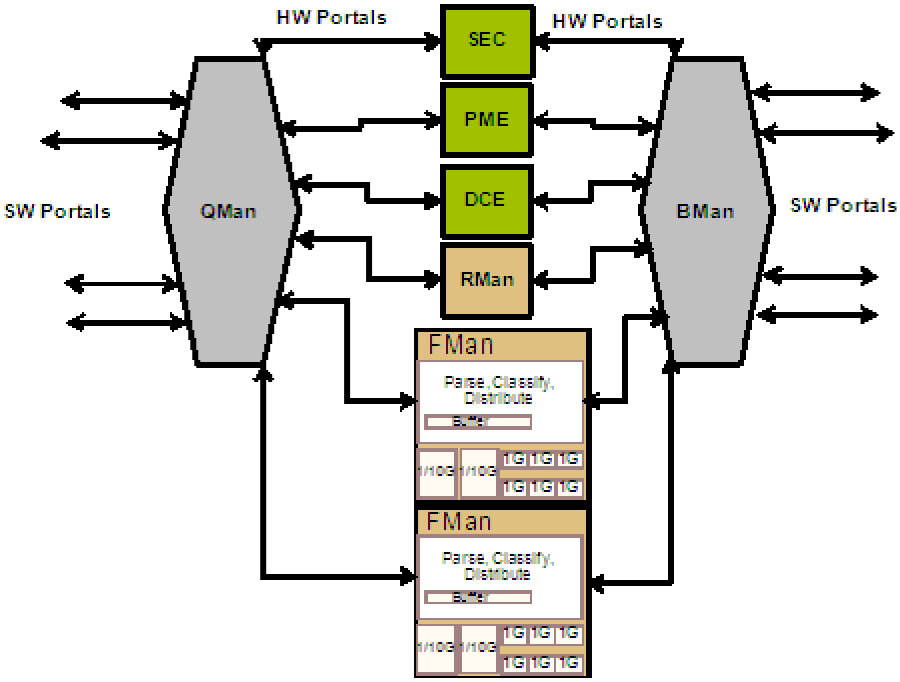

Figure 10. Logical representation of DPAA

5.10.5.1 Frame Manager and network interfaces

The chip incorporates two enhanced Frame Managers. The Frame Manager improves on the bandwidth and functionality offered in the P4080.

Each Frame Manager, or FMan, combines Ethernet MACs with packet parsing and classification logic to provide intelligent distribution and queuing decisions for incoming traffic. Each FMan supports PCD at 37.2 Mpps, supporting line rate 2x10G + 2x2.5G at minimum frame size.

These Ethernet combinations are supported:

- 10 Gbps Ethernet MACs are supported with XAUI (four lanes at 3.125 GHz) or XFI (one lane at 10.3125 GHz SerDes).
- 1 Gbps Ethernet MACs are supported with SGMII (one lane at 1.25 GHz with 3.125 GHz option for 2.5 Gbps Ethernet).
    - SGMIIs can be run at 3.125 GHz so long as the total Ethernet bandwidth does not exceed 25 Gbps on the associated FMan.
    - If not already assigned to SGMII, two MACs can be used with RGMII.
- Four x1Gbps Ethernet MACs can be supported using a single lane at 5 GHz (QSGMII).
- HiGig is supported using four lanes at 3.125 GHz or 3.75 GHz (HiGig2).

The Frame Manager's Ethernet functionality also supports the following:

- 1588v2 hardware timestamping mechanism in conjunction with IEEE Std. 802.3bf (Ethernet support for time synchronization protocol)
- Energy Efficient Ethernet (IEEE Std. 802.3az)
- IEEE Std. 802.3bd (MAC control frame support for priority based flow control)
- IEEE Std. 802.1Qbb (Priority-based flow control) for up to eight queues/priorities
- IEEE Std. 802.1Qaz (Enhanced transmission selection) for three or more traffic classes

5.10.5.1.1 Receiver functionality: parsing, classification, and distribution

Each Frame Manager matches its 25 Gbps Ethernet connectivity with 25 Gbps (37.2 Mpps) of Parsing, Classification, and Distribution (PCD) performance. PCD is the process by which the Frame Manager identifies the frame queue on which received packets should be enqueued. The consumer of the data on the frame queues is determined by Queue Manager configuration; however, these activities are closely linked and managed by the FMan Driver and FMan Configuration Tool, as in previous QorIQ SoCs.

This figure provides a logical view of the FMan's processing flow, illustrating the PCD features.

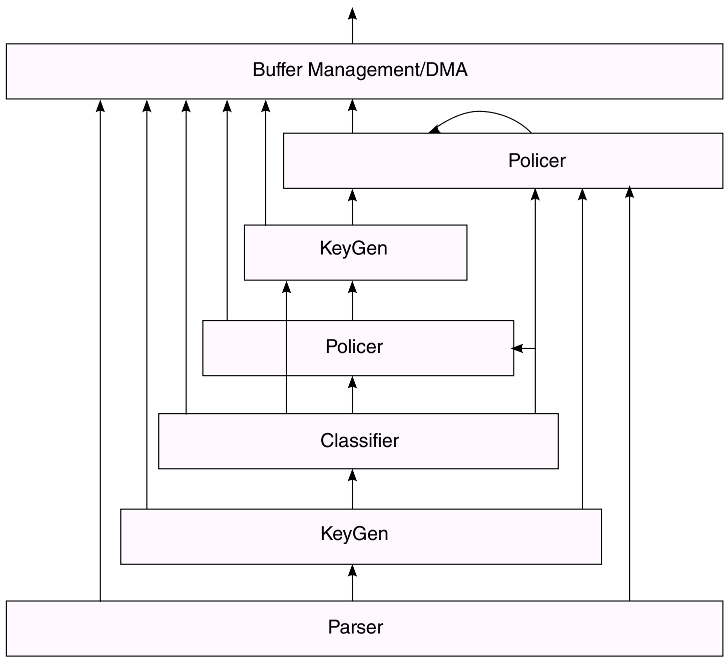

Figure 11. Logical view of FMan processing

Each frame received by the FMan is buffered internally while the Parser, KeyGen, and Classification functions operate.

The parse function can parse many standard protocols, including options and tunnels, and it supports a generic configurable capability to allow proprietary or future protocols to be parsed. Hard parsing of the standard protocol headers can be augmented with user-defined soft parsing rules to handle proprietary header fields. Hard and soft parsing occurs at wire speed.

This table defines several types of parser headers.

Table 6. Parser header types

Header type	Definition
Self-describing	Announced by proprietary values of Ethertype, protocol identifier, next header, and other standard fields. They are self-describing in that the frame contains information that describes the presence of the proprietary header.
Non-self-describing	Does not contain any information that indicates the presence of the header.

Table continues on the next page...

Table 6. Parser header types (continued)

Header type	Definition
	For example, a frame that always contains a proprietary header before the Ethernet header would be non-self-describing. Both self-describing and non-self-describing headers are supported by means of parsing rules in the FMan.
Proprietary	Can be defined as being self-describing or non-self-describing

The underlying notion is that different frames may require different treatment, and only through detailed parsing of the frame can proper treatment be determined.

Parse results can (optionally) be passed to software.

5.10.5.1.2 FMan distribution and policing

After parsing is complete, there are two options for treatment, as shown in this table.

Table 7. Post-parsing treatment options

Treatment	Function	Benefits
Hash	• Hashes select fields in the frame as part of a spreading mechanism. • The result is a specific frame queue identifier. • To support added control, this FQID can be indexed by values found in the frame, such as TOS or p-bits, or any other desired field(s).	Useful when spreading traffic while obeying QoS constraints is required
Classification look-up	• Looks up certain fields in the frame to determine subsequent action to take, including policing. • The FMan contains internal memory that holds small tables for this purpose. • The user configures the sets of lookups to perform, and the parse results dictate which one of those sets to use. • Lookups can be chained together such that a successful look-up can provide key information for a subsequent look-up. After all the look-ups are complete, the final classification result provides either a hash key to use for spreading, or a FQ ID directly.	• Useful when hash distribution is insufficient and a more detailed examination of the frame is required • Can determine whether policing is required and the policing context to use

Key benefits of the FMan policing function are as follows:

- Because the FMan has up to 256 policing profiles, any frame queue or group of frame queues can be policed to either drop or mark packets if the flow exceeds a preconfigured rate.
- Policing and classification can be used in conjunction to mitigate Distributed Denial of Service Attack (DDOS).
- The policing is based on the two-rate-three-color marking algorithm (RFC2698). The sustained and peak rates, as well as the burst sizes, are user-configurable. Therefore, the policing function can rate-limit traffic to conform to the rate that the flow is mapped to at flow set-up time. By prioritizing and policing traffic prior to software processing, CPU cycles can focus on important and urgent traffic ahead of other traffic.

Each FMan also supports PCD on traffic arriving from within the chip. This is referred to as off-line parsing, and it is useful for reclassification following decapsulation of encrypted or compressed packets.

FMan PCD supports virtualization and strong partitioning by delaying buffer pool selection until after classification. In addition to determining the FQ ID for the classified packet, the FMan also determines the 'storage profile.' Configuration of storage profiles (up to 32 per physical port) allows the FMan to store received packets using buffer pools owned by a single software partition, and enqueue the associated Frame Descriptor to a frame queue serviced by only that software partition.

The SEC 5.0 can perform full protocol processing for the following security protocols:

- IPsec
- SSL/TLS
- 3GPP RLC encryption/decryption
- LTE PDCP
- SRTP
- IEEE 802.1AE MACSec
- IEEE 802.16e WiMax MAC layer

The SEC 5.0 supports the following algorithms, modes, and key lengths as raw modes, or in combination with the security protocol processing described above.

- Public Key Hardware Accelerators (PKHA)
  - RSA and Diffie-Hellman (to 4096b)
  - Elliptic curve cryptography (1023b)
- Data Encryption Standard Accelerators (DESA)
  - DES, 3DES (2-key, 3-key)
  - ECB, CBC, OFB, and CFB modes
- Advanced Encryption Standard Accelerators (AESA)
  - Key lengths of 128-bit, 192-bit, and 256-bit
  - ECB, CBC, CTR, CCM, GCM, CMAC, OFB, CFB, xcbc-mac, and XTS
- ARC Four Hardware Accelerators (AFHA)
  - Compatible with RC4 algorithm
- Message Digest Hardware Accelerators (MDHA)
  - SHA-1, SHA-256, 384, 512-bit digests
  - MD5 128-bit digest
  - HMAC with all algorithms
- Kasumi/F8 Hardware Accelerators (KFHA)
  - F8, F9 as required for 3GPP
  - A5/3 for GSM and EDGE, GEA-3 for GPRS
- Snow 3G Hardware Accelerators (SNOWf8 and SNOWf9)
  - Implements Snow 3.0, F8 and F9 modes
- ZUC Hardware Accelerators (ZUCE and ZUCA)
  - Implements 128-EEA3 & 128-EIA3
- CRC Unit
  - Standard and user-defined polynomials
- Random Number Generator
  - Incorporates TRNG entropy generator for seeding and deterministic engine (SHA-256)
  - Supports random IV generation

The SEC 5.0 is designed to support bulk encryption at up to 40 Gbps, large packet/record IPsec/SSL at up to 30 Gbps, and 20 Gbps for IPsec ESP at Imix packet sizes. 3G and LTE algorithms are supported at 10 Gbps or more.

The SEC dequeues data from its QMan hardware portal and, based on FQ configuration, also dequeues associated instructions and operands in the Shared Descriptor. The SEC processes the data then enqueues it to the configured output FQ. The SEC uses the Status/CMD word in the output Frame Descriptor to inform the next consumer of any errors encountered during processing (for example, received packet outside the anti-replay window.)

    - All standard modes of decompression
    - No compression
    - Static Huffman codes
    - Dynamic Huffman codes
- Provides option to return original compressed Frame along with the uncompressed Frame or release the buffers to BMan
- Does not support use of ZLIB preset dictionaries (FDICT flag = 1 is treated as an error).
- Base 64 decoding (RFC4648) prior to decompression

The DCE 1.0 is designed to support up to 8.8 Gbps for either compression or decompression, or 17.5 Gbps aggregate at ~4 KB data sizes.

5.10.6 DPAA capabilities

Some DPAA features and capabilities have been described in the sections covering individual DPAA components. This section describes some capabilities enabled by DPAA components working together.

5.10.6.1 Ingress policing and congestion management

In addition to selecting FQ ID and storage profile, classification can determine whether policing is required for a received packet, along with the specific policing context to be used.

FMan policing capabilities include the following:

- RFC2698: two-rate, three-color marking algorithm
- RFC4115: Differentiated service two-rate, three-color marker with efficient handling of in-profile traffic
- Up to 256 internal profiles

The sustained and peak rates, and burst size for each policing profile are user-configurable.

5.10.6.2 Customer-edge egress-traffic management (CEETM)

Customer-edge egress-traffic management (CEETM) is a DPAA enhancement first appearing in the T4240. T4240 continues to support the work queue and frame queue scheduling functionality available in the P4080 and other first generation QorIQ chips, but introduces alternative functionary, CEETM, that can be mode selected on a network interface basis to support the shaping and scheduling requirements of carrier Ethernet connected systems.

5.10.6.2.1 CEETM features

Each instance of CEETM (one per FMan) provides the following features:

- Supports hierarchical multi-level scheduling and shaping, which:
    - is performed in an atomic manner; all context at all levels is examined and updated synchronously.
    - employs no intermediate buffering between class queues and the direct connect portal to the FMan.
- Supports dual-rate shaping (paired committed rate (CR) shaper and excess rate (ER) shaper) at all shaping points.
    - Shapers are token bucket based with configurable rate and burst limit.
    - Paired CR/ER shapers may be configured as independent or coupled on a per pair basis; coupled means that credits to the CR shaper in excess of its token bucket limit is credited to the ER bucket
- Supports eight logical network interfaces (LNI)
    - Each LNI:
        - aggregates frames from one or more channels.
        - priority schedules unshaped frames (aggregated from unshaped channels), CR frames, and ER frames (aggregated from shaped channels)

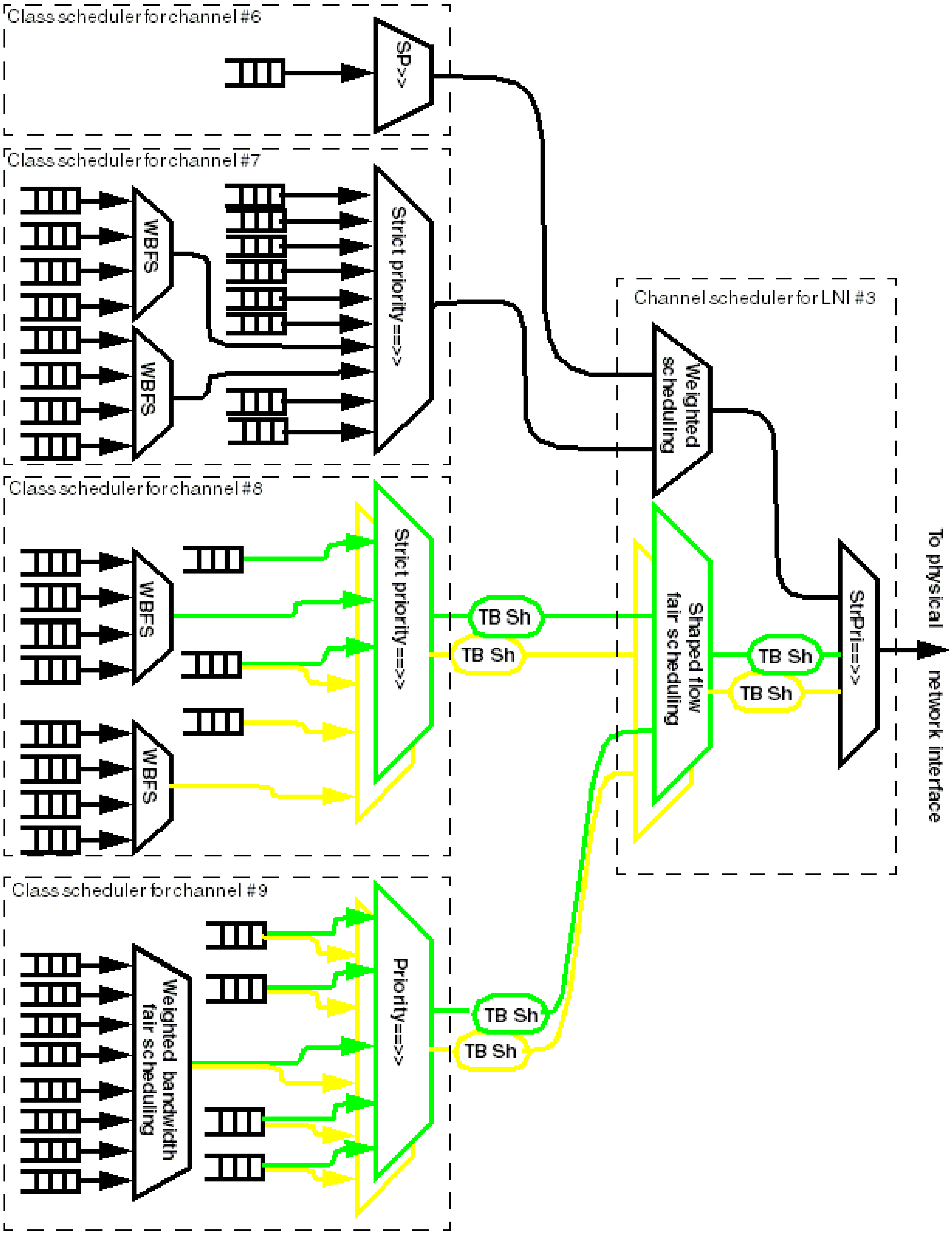

Figure 13. CEETM scheduler: illustrative configuration scenario

Figure 13 illustrates the following scenario:

5.13 Debug support

The reduced number of external buses enabled by the move to multicore chips greatly simplifies board level lay-out and eliminates many concerns over signal integrity. Even though the board designer may embrace multicore CPUs, software engineers have real concerns over the potential to lose debug visibility. Despite the problems external buses can cause for the hardware engineer, they provide software developers with the ultimate confirmation that the proper instructions and data are passing between processing elements.

Processing on a multicore chip with shared caches and peripherals also leads to greater concurrency and an increased potential for unintended interactions between device components. To ensure that software developers have the same or better visibility into the device as they would with multiple discrete communications processors, Freescale developed an Advanced Multicore Debug Architecture.

The debugging and performance monitoring capability enabled by the device hardware coexists within a debug ecosystem that offers a rich variety of tools at different levels of the hardware/software stack. Software development and debug tools from Freescale (CodeWarrior), as well as third-party vendors, provide a rich set of options for configuring, controlling, and analyzing debug and performance related events.

6 Conclusion

Featuring 24 virtual cores, and based on the dual-threaded e6500 Power Architecture core, the T4240 processor, along with its 16 (T4160) and 8 (T4080) virtual-core variants, offers frequencies up to 1.8 GHz, large caches, hardware acceleration, and advanced system peripherals. All three devices target applications that benefit from consolidation of control and data plane processing in a single chip. In addition, each e6500 core implements the Freescale AltiVec technology SIMD engine, dramatically boosting the performance of math-intensive algorithms without using additional DSP components on the board. A wide variety of applications can benefit from the processing, I/O integration, and power management offered for the T4 series processors. Similar to other QorIQ devices, the T4 family processors' high level of integration offers significant space, weight, and power benefits compared to multiple discrete devices. Freescale also offers fully featured development support, which includes the QorIQ T4240 QDS Development System, QorIQ T4240 Reference Design Board, Linux SDK for QorIQ Processors, as well as popular operating systems and development tools from a variety of vendors. See the Freescale website for the latest information on tools and SW availability.

For more information about the QorIQ T4 family, contact your Freescale sales representative.

Appendix A T4160

A.1 Introduction

The T4160 is a lower power version of the T4240. The T4160 combines eight dual threaded Power Architecture e6500 cores and two memory complexes (CoreNet platform cache and DDR3 memory controller) with the same high-performance datapath acceleration, networking, and peripheral bus interfaces.

This figure shows the major functional units within the chip.

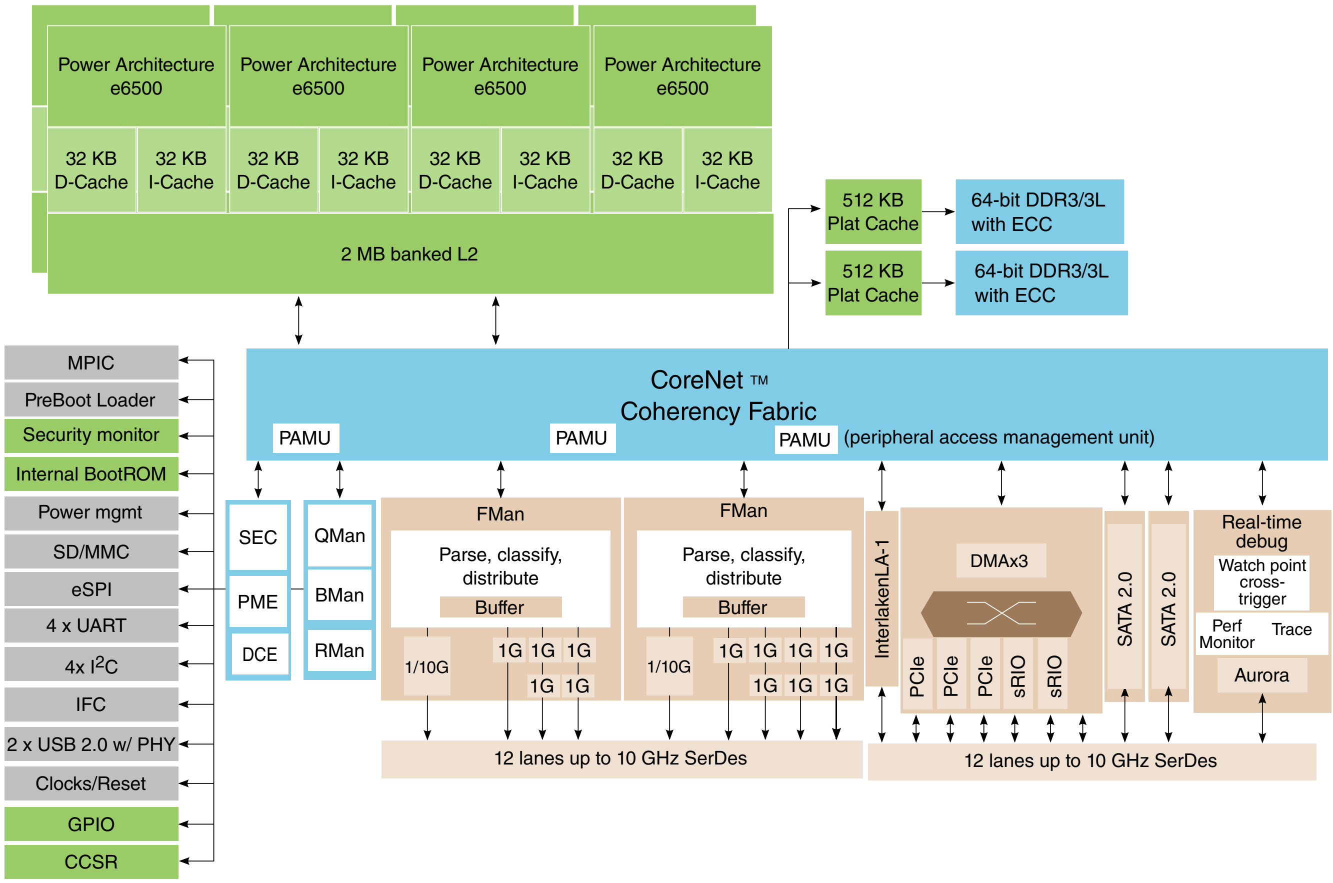

Figure A-1. T4160 block diagram

A.2 Overview of differences between T4240 and T4160

Table A-1. Differences between T4240 and T4160

Feature	T4240	T4160
Cores		
Number of physical cores	12	8
Number of threads	24	16
Number of clusters	3	2
Memory subsystem		
Total CPC memory	3 x 512 KB	2 x 512 KB
Number of DDR controllers	3	2
Peripherals		
Number of Frame Managers	2	2
Total number of Anyspeed MACs	8 per Frame Manager	6 (FMan1) and 8 (FMan2)

Table continues on the next page...