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#### Details

Product Status	Active
Core Processor	S08
Core Size	8-Bit
Speed	40MHz
Connectivity	I <sup>2</sup> C, SCI, SPI
Peripherals	LVD, POR, PWM, WDT
Number of I/O	38
Program Memory Size	32KB (32K x 8)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	2K x 8
Voltage - Supply (Vcc/Vdd)	2.7V ~ 5.5V
Data Converters	A/D 8x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	48-VFQFN Exposed Pad
Supplier Device Package	48-QFN-EP (7x7)
Purchase URL	https://www.e-xfl.com/product-detail/nxp-semiconductors/s9s08aw32e5cfde

Email: info@E-XFL.COM

Address: Room A, 16/F, Full Win Commercial Centre, 573 Nathan Road, Mongkok, Hong Kong



#### **Chapter 4 Memory**





Figure 4-2. MC9S08AW32 and MC9S08AW16 Memory Map



be programmed to logic 0 to enable block protection. Therefore the value \$DE must be programmed into NVPROT to protect addresses \$E000 through \$FFFF.



#### Figure 4-5. Block Protection Mechanism

One use for block protection is to block protect an area of FLASH memory for a bootloader program. This bootloader program then can be used to erase the rest of the FLASH memory and reprogram it. Because the bootloader is protected, it remains intact even if MCU power is lost in the middle of an erase and reprogram operation.

### 4.4.7 Vector Redirection

Whenever any block protection is enabled, the reset and interrupt vectors will be protected. Vector redirection allows users to modify interrupt vector information without unprotecting bootloader and reset vector space. Vector redirection is enabled by programming the FNORED bit in the NVOPT register located at address \$FFBF to zero. For redirection to occur, at least some portion but not all of the FLASH memory must be block protected by programming the NVPROT register located at address \$FFBD. All of the interrupt vectors (memory locations \$FFC0-\$FFFD) are redirected, though the reset vector (\$FFFE:FFFF) is not.

For example, if 512 bytes of FLASH are protected, the protected address region is from \$FE00 through \$FFFF. The interrupt vectors (\$FFC0-\$FFFD) are redirected to the locations \$FDC0-\$FDFD. Now, if an SPI interrupt is taken for instance, the values in the locations \$FDE0:FDE1 are used for the vector instead of the values in the locations \$FFE0:FFE1. This allows the user to reprogram the unprotected portion of the FLASH with new program code including new interrupt vector values while leaving the protected area, which includes the default vector locations, unchanged.

### 4.5 Security

The MC9S08AW60 Series includes circuitry to prevent unauthorized access to the contents of FLASH and RAM memory. When security is engaged, FLASH and RAM are considered secure resources. Direct-page registers, high-page registers, and the background debug controller are considered unsecured resources. Programs executing within secure memory have normal access to any MCU memory locations and resources. Attempts to access a secure memory location with a program executing from an unsecured memory space or through the background debug interface are blocked (writes are ignored and reads return all 0s).

Security is engaged or disengaged based on the state of two nonvolatile register bits (SEC01:SEC00) in the FOPT register. During reset, the contents of the nonvolatile location NVOPT are copied from FLASH into the working FOPT register in high-page register space. A user engages security by programming the NVOPT location which can be done at the same time the FLASH memory is programmed. The 1:0 state disengages security and the other three combinations engage security. Notice the erased state (1:1) makes



I bit in the CCR is 0 to allow interrupts. The global interrupt mask (I bit) in the CCR is initially set after reset which masks (prevents) all maskable interrupt sources. The user program initializes the stack pointer and performs other system setup before clearing the I bit to allow the CPU to respond to interrupts.

When the CPU receives a qualified interrupt request, it completes the current instruction before responding to the interrupt. The interrupt sequence obeys the same cycle-by-cycle sequence as the SWI instruction and consists of:

- Saving the CPU registers on the stack
- Setting the I bit in the CCR to mask further interrupts
- Fetching the interrupt vector for the highest-priority interrupt that is currently pending
- Filling the instruction queue with the first three bytes of program information starting from the address fetched from the interrupt vector locations

While the CPU is responding to the interrupt, the I bit is automatically set to avoid the possibility of another interrupt interrupting the ISR itself (this is called nesting of interrupts). Normally, the I bit is restored to 0 when the CCR is restored from the value stacked on entry to the ISR. In rare cases, the I bit may be cleared inside an ISR (after clearing the status flag that generated the interrupt) so that other interrupts can be serviced without waiting for the first service routine to finish. This practice is not recommended for anyone other than the most experienced programmers because it can lead to subtle program errors that are difficult to debug.

The interrupt service routine ends with a return-from-interrupt (RTI) instruction which restores the CCR, A, X, and PC registers to their pre-interrupt values by reading the previously saved information off the stack.

#### NOTE

For compatibility with the M68HC08, the H register is not automatically saved and restored. It is good programming practice to push H onto the stack at the start of the interrupt service routine (ISR) and restore it immediately before the RTI that is used to return from the ISR.

When two or more interrupts are pending when the I bit is cleared, the highest priority source is serviced first (see Table 5-1).

#### 5.5.1 Interrupt Stack Frame

Figure 5-1 shows the contents and organization of a stack frame. Before the interrupt, the stack pointer (SP) points at the next available byte location on the stack. The current values of CPU registers are stored on the stack starting with the low-order byte of the program counter (PCL) and ending with the CCR. After stacking, the SP points at the next available location on the stack which is the address that is one less than the address where the CCR was saved. The PC value that is stacked is the address of the instruction in the main program that would have executed next if the interrupt had not occurred.



#### Chapter 6 Parallel Input/Output

_	7	6	5	4	3	2	1	0
R W	PTFDS7	PTFDS6	PTFDS5	PTFDS4	PTFDS3	PTFDS2	PTFDS1	PTFDS0
Reset	0	0	0	0	0	0	0	0

#### Figure 6-38. Output Drive Strength Selection for Port F (PTFDS)

#### Table 6-31. PTFDS Register Field Descriptions

Field	Description
7:0 PTFDS[7:0]	<ul> <li>Output Drive Strength Selection for Port F Bits — Each of these control bits selects between low and high output drive for the associated PTF pin.</li> <li>0 Low output drive enabled for port F bit n.</li> <li>1 High output drive enabled for port F bit n.</li> </ul>



#### Chapter 7 Central Processor Unit (S08CPUV2)

interrupt service routine, this would allow nesting of interrupts (which is not recommended because it leads to programs that are difficult to debug and maintain).

For compatibility with the earlier M68HC05 MCUs, the high-order half of the H:X index register pair (H) is not saved on the stack as part of the interrupt sequence. The user must use a PSHH instruction at the beginning of the service routine to save H and then use a PULH instruction just before the RTI that ends the interrupt service routine. It is not necessary to save H if you are certain that the interrupt service routine does not use any instructions or auto-increment addressing modes that might change the value of H.

The software interrupt (SWI) instruction is like a hardware interrupt except that it is not masked by the global I bit in the CCR and it is associated with an instruction opcode within the program so it is not asynchronous to program execution.

#### 7.4.3 Wait Mode Operation

The WAIT instruction enables interrupts by clearing the I bit in the CCR. It then halts the clocks to the CPU to reduce overall power consumption while the CPU is waiting for the interrupt or reset event that will wake the CPU from wait mode. When an interrupt or reset event occurs, the CPU clocks will resume and the interrupt or reset event will be processed normally.

If a serial BACKGROUND command is issued to the MCU through the background debug interface while the CPU is in wait mode, CPU clocks will resume and the CPU will enter active background mode where other serial background commands can be processed. This ensures that a host development system can still gain access to a target MCU even if it is in wait mode.

#### 7.4.4 Stop Mode Operation

Usually, all system clocks, including the crystal oscillator (when used), are halted during stop mode to minimize power consumption. In such systems, external circuitry is needed to control the time spent in stop mode and to issue a signal to wake up the target MCU when it is time to resume processing. Unlike the earlier M68HC05 and M68HC08 MCUs, the HCS08 can be configured to keep a minimum set of clocks running in stop mode. This optionally allows an internal periodic signal to wake the target MCU from stop mode.

When a host debug system is connected to the background debug pin (BKGD) and the ENBDM control bit has been set by a serial command through the background interface (or because the MCU was reset into active background mode), the oscillator is forced to remain active when the MCU enters stop mode. In this case, if a serial BACKGROUND command is issued to the MCU through the background debug interface while the CPU is in stop mode, CPU clocks will resume and the CPU will enter active background mode where other serial background commands can be processed. This ensures that a host development system can still gain access to a target MCU even if it is in stop mode.

Recovery from stop mode depends on the particular HCS08 and whether the oscillator was stopped in stop mode. Refer to the Modes of Operation chapter for more details.



Source				c	Eff on (	ec CC	t R		ess le	epe	and	cles <sup>1</sup>
Form	Form Operation Description				I	N	z	с	Addre Moc	Opco	Opera	Bus Cy
BRCLR n,opr8a,rel	Branch if Bit <i>n</i> in Memory Clear	Branch if (Mn) = 0	_	_	_	_	_	¢	DIR (b0) DIR (b1) DIR (b2) DIR (b3) DIR (b3) DIR (b4) DIR (b5) DIR (b6) DIR (b7)	01 03 05 07 09 0B 0D 0F	dd rr dd rr dd rr dd rr dd rr dd rr dd rr dd rr dd rr	5555555 555555555555555555555555555555
BRN <i>rel</i>	Branch Never	Uses 3 Bus Cycles	-	-	-	-	-	-	REL	21	rr	3
BRSET n,opr8a,rel	Branch if Bit <i>n</i> in Memory Set	Branch if (Mn) = 1	_	_	_	_	_	\$	DIR (b0) DIR (b1) DIR (b2) DIR (b3) DIR (b4) DIR (b5) DIR (b6) DIR (b7)	00 02 04 06 08 0A 0C 0E	dd rr dd rr dd rr dd rr dd rr dd rr dd rr dd rr dd rr	5555555555
BSET n,opr8a	Set Bit <i>n</i> in Memory	Mn ← 1	_	_	_	_	_	_	DIR (b0) DIR (b1) DIR (b2) DIR (b3) DIR (b3) DIR (b4) DIR (b5) DIR (b6) DIR (b7)	10 12 14 16 18 1A 1C 1E	dd dd dd dd dd dd dd dd dd	5 5 5 5 5 5 5 5 5 5 5 5 5
BSR rel	Branch to Subroutine	$\begin{array}{l} PC \leftarrow (PC) + 0 \texttt{x0002} \\ \texttt{push} \ (PCL) \texttt{; } SP \leftarrow (SP) - 0 \texttt{x0001} \\ \texttt{push} \ (PCH) \texttt{; } SP \leftarrow (SP) - 0 \texttt{x0001} \\ PC \leftarrow (PC) + \mathit{rel} \end{array}$	_	_	_	_	_	_	REL	AD	rr	5
CBEQ opr8a,rel CBEQA #opr8i,rel CBEQX #opr8i,rel CBEQ oprx8,X+,rel CBEQ ,X+,rel CBEQ oprx8,SP,rel	Compare and Branch if Equal	Branch if $(A) = (M)$ Branch if $(A) = (M)$ Branch if $(X) = (M)$ Branch if $(A) = (M)$ Branch if $(A) = (M)$ Branch if $(A) = (M)$	_	_	_	_	_	_	DIR IMM IX1+ IX+ SP1	31 41 51 61 71 9E61	dd rr ii rr ii rr ff rr rr ff rr	544556
CLC	Clear Carry Bit	$C \leftarrow 0$	-	-	-	-	-	0	INH	98		1
CLI	Clear Interrupt Mask Bit	$I \leftarrow 0$	-	-	0	-	-	-	INH	9A		1
CLR opr8a CLRA CLRX CLRH CLR oprx8,X CLR ,X CLR oprx8,SP	Clear	$\begin{array}{c} M \leftarrow 0x00 \\ A \leftarrow 0x00 \\ X \leftarrow 0x00 \\ H \leftarrow 0x00 \\ M \leftarrow 0x00 \\ M \leftarrow 0x00 \\ M \leftarrow 0x00 \\ M \leftarrow 0x00 \end{array}$	0	_	_	0	1	_	DIR INH INH INH IX1 IX SP1	3F 4F 5F 8C 6F 7F 9E6F	dd ff ff	5 1 1 5 4 6
CMP #opr8i CMP opr8a CMP opr16a CMP oprx16,X CMP oprx8,X CMP ,X CMP oprx16,SP CMP oprx8,SP	Compare Accumulator with Memory	(A) – (M) (CCR Updated But Operands Not Changed)	\$	_	_	\$	\$	\$	IMM DIR EXT IX2 IX1 IX SP2 SP1	A1 B1 C1 E1 F1 9ED1 9EE1	ii dd hh II ee ff ff ee ff ff	2 3 4 3 3 5 4
COM opr8a COMA COMX COM oprx8,X COM ,X COM oprx8,SP	Complement (One's Complement)	$\begin{array}{c} M \leftarrow (\overline{M}) = 0xFF - (M) \\ A \leftarrow (\overline{A}) = 0xFF - (A) \\ X \leftarrow (\overline{X}) = 0xFF - (X) \\ M \leftarrow (\overline{M}) = 0xFF - (M) \\ M \leftarrow (\overline{M}) = 0xFF - (M) \\ M \leftarrow (\overline{M}) = 0xFF - (M) \end{array}$	0	_	_	\$	\$	1	DIR INH INH IX1 IX SP1	33 43 53 63 73 9E63	dd ff ff	5 1 5 4 6
CPHX opr16a CPHX #opr16i CPHX opr8a CPHX oprx8,SP	Compare Index Register (H:X) with Memory	(H:X) – (M:M + 0x0001) (CCR Updated But Operands Not Changed)	\$	_	_	¢	\$	\$	EXT IMM DIR SP1	3E 65 75 9EF3	hh ll jj kk dd ff	6 3 5 6

Table 7-2. HCS08 Instruction Set Summary (Sheet 3 of 7)



Chapter 8 Internal Clock Generator (S08ICGV4)



Figure 8-15. ICG Initialization and Stop Recovery for Example #2



Chapter 8 Internal Clock Generator (S08ICGV4)



CPWMS	MSnB:MSnA	ELSnB:ELSnA	Mode	Configuration				
Х	ХХ	00	Pin not used for TPM channel; use as an external clock for the TPM or revert to general-purpose I/O					
0	00	01	Input capture	Capture on rising edge only				
		10		Capture on falling edge only				
		11		Capture on rising or falling edge				
	01	00	Output	Software compare only				
		01	compare	Toggle output on compare				
		10		Clear output on compare				
		11		Set output on compare				
	1X	10	Edge-aligned	High-true pulses (clear output on compare)				
		X1	PWM	Low-true pulses (set output on compare)				
1	XX	10	Center-aligned	High-true pulses (clear output on compare-up)				
		X1	PWM	Low-true pulses (set output on compare-up)				

Tahle	10-5	Mode	Edge	and I	l evel	Selection
lane	10-5.	woue,	Luye,	anu	LEVEI	Selection

If the associated port pin is not stable for at least two bus clock cycles before changing to input capture mode, it is possible to get an unexpected indication of an edge trigger. Typically, a program would clear status flags after changing channel configuration bits and before enabling channel interrupts or using the status flags to avoid any unexpected behavior.

### 10.4.5 Timer x Channel Value Registers (TPMxCnVH:TPMxCnVL)

These read/write registers contain the captured TPM counter value of the input capture function or the output compare value for the output compare or PWM functions. The channel value registers are cleared by reset.

_	7	6	5	4	3	2	1	0
R W	Bit 15	14	13	12	11	10	9	Bit 8
Reset	0	0	0	0	0	0	0	0



	7	6	5	4	3	2	1	0
R W	Bit 7	6	5	4	3	2	1	Bit 0
Reset	0	0	0	0	0	0	0	0

Figure 10-10. Timer Channel Value Register Low (TPMxCnVL)

In input capture mode, reading either byte (TPMxCnVH or TPMxCnVL) latches the contents of both bytes into a buffer where they remain latched until the other byte is read. This latching mechanism also resets (becomes unlatched) when the TPMxCnSC register is written.



Chapter 12 Serial Peripheral Interface (S08SPIV3)

The most common uses of the SPI system include connecting simple shift registers for adding input or output ports or connecting small peripheral devices such as serial A/D or D/A converters. Although Figure 12-2 shows a system where data is exchanged between two MCUs, many practical systems involve simpler connections where data is unidirectionally transferred from the master MCU to a slave or from a slave to the master MCU.

#### 12.0.2.2 SPI Module Block Diagram

Figure 12-3 is a block diagram of the SPI module. The central element of the SPI is the SPI shift register. Data is written to the double-buffered transmitter (write to SPI1D) and gets transferred to the SPI shift register at the start of a data transfer. After shifting in a byte of data, the data is transferred into the double-buffered receiver where it can be read (read from SPI1D). Pin multiplexing logic controls connections between MCU pins and the SPI module.

When the SPI is configured as a master, the clock output is routed to the SPSCK pin, the shifter output is routed to MOSI, and the shifter input is routed from the MISO pin.

When the SPI is configured as a slave, the SPSCK pin is routed to the clock input of the SPI, the shifter output is routed to MISO, and the shifter input is routed from the MOSI pin.

In the external SPI system, simply connect all SPSCK pins to each other, all MISO pins together, and all MOSI pins together. Peripheral devices often use slightly different names for these pins.

ADCH	Channel	Input	Pin Control	ADCH	Channel	Input	Pin Contro
01110	AD14	PTD6/ADC1P14/ TPM1CLK	ADPC14	11110	V <sub>REFL</sub>	V <sub>REFL</sub>	N/A
01111	AD15	PTD7ADC1P15/ KBI1P7	ADPC15	11111	Module disabled	None	N/A

<sup>1</sup> For more information, see Section 14.2.3, "Temperature Sensor."

#### NOTE

Selecting the internal bandgap channel requires BGBE =1 in SPMSC1 see Section 5.9.8, "System Power Management Status and Control 1 Register (SPMSC1)." For value of bandgap voltage reference see Section Table A-7., "DC Characteristics."

#### 14.2.1 Alternate Clock

The ADC module is capable of performing conversions using the MCU bus clock, the bus clock divided by two, the local asynchronous clock (ADACK) within the module, or the alternate clock, ALTCLK. The alternate clock for the MC9S08AW60 Series MCU devices is the external reference clock (ICGERCLK) from the internal clock generator (ICG) module.

Because ICGERCLK is active only while an external clock source is enabled, the ICG must be configured for either FBE or FEE mode (CLKS1 = 1). ICGERCLK must run at a frequency such that the ADC conversion clock (ADCK) runs at a frequency within its specified range ( $f_{ADCK}$ ) after being divided down from the ALTCLK input as determined by the ADIV bits. For example, if the ADIV bits are set up to divide by four, then the minimum frequency for ALTCLK (ICGERCLK) is four times the minimum value for  $f_{ADCK}$  and the maximum frequency is four times the maximum value for  $f_{ADCK}$ . Because of the minimum frequency requirement, when an oscillator circuit is used it must be configured for high range operation (RANGE = 1).

ALTCLK is active while the MCU is in wait mode provided the conditions described above are met. This allows ALTCLK to be used as the conversion clock source for the ADC while the MCU is in wait mode.

ALTCLK cannot be used as the ADC conversion clock source while the MCU is in stop3.

#### 14.2.2 Hardware Trigger

The ADC hardware trigger, ADHWT, is output from the real time interrupt (RTI) counter. The RTI counter can be clocked by either ICGERCLK or a nominal 1 kHz clock source within the RTI block. The 1-kHz clock source can be used with the MCU in run, wait, or stop3. With the ICG configured for either FBE or FEE mode, ICGERCLK can be used with the MCU in run or wait.

The period of the RTI is determined by the input clock frequency and the RTIS bits. When the ADC hardware trigger is enabled, a conversion is initiated upon an RTI counter overflow. The RTI counter is a free running counter that generates an overflow at the RTI rate determined by the RTIS bits.



Field	Description
5 ACFE	<ul> <li>Compare Function Enable — ACFE is used to enable the compare function.</li> <li>0 Compare function disabled</li> <li>1 Compare function enabled</li> </ul>
4 ACFGT	<ul> <li>Compare Function Greater Than Enable — ACFGT is used to configure the compare function to trigger when the result of the conversion of the input being monitored is greater than or equal to the compare value. The compare function defaults to triggering when the result of the compare of the input being monitored is less than the compare value.</li> <li>0 Compare triggers when input is less than compare level</li> <li>1 Compare triggers when input is greater than or equal to compare level</li> </ul>

Table 14-4. ADC1SC2 Register Field Descriptions (continued)

#### 14.4.3 Data Result High Register (ADC1RH)

ADC1RH contains the upper two bits of the result of a 10-bit conversion. When configured for 8-bit conversions both ADR8 and ADR9 are equal to zero. ADC1RH is updated each time a conversion completes except when automatic compare is enabled and the compare condition is not met. In 10-bit MODE, reading ADC1RH prevents the ADC from transferring subsequent conversion results into the result registers until ADC1RL is read. If ADC1RL is not read until after the next conversion is completed, then the intermediate conversion result will be lost. In 8-bit mode there is no interlocking with ADC1RL. In the case that the MODE bits are changed, any data in ADC1RH becomes invalid.



Figure 14-6. Data Result High Register (ADC1RH)

#### 14.4.4 Data Result Low Register (ADC1RL)

ADC1RL contains the lower eight bits of the result of a 10-bit conversion, and all eight bits of an 8-bit conversion. This register is updated each time a conversion completes except when automatic compare is enabled and the compare condition is not met. In 10-bit mode, reading ADC1RH prevents the ADC from transferring subsequent conversion results into the result registers until ADC1RL is read. If ADC1RL is not read until the after next conversion is completed, then the intermediate conversion results will be lost. In 8-bit mode, there is no interlocking with ADC1RH. In the case that the MODE bits are changed, any data in ADC1RL becomes invalid.





Figure 14-7. Data Result Low Register (ADC1RL)

# 14.4.5 Compare Value High Register (ADC1CVH)

This register holds the upper two bits of the 10-bit compare value. These bits are compared to the upper two bits of the result following a conversion in 10-bit mode when the compare function is enabled. In 8-bit operation, ADC1CVH is not used during compare.



Figure 14-8. Compare Value High Register (ADC1CVH)

### 14.4.6 Compare Value Low Register (ADC1CVL)

This register holds the lower 8 bits of the 10-bit compare value, or all 8 bits of the 8-bit compare value. Bits ADCV7:ADCV0 are compared to the lower 8 bits of the result following a conversion in either 10-bit or 8-bit mode.



Figure 14-9. Compare Value Low Register(ADC1CVL)

# 14.4.7 Configuration Register (ADC1CFG)

ADC1CFG is used to select the mode of operation, clock source, clock divide, and configure for low power or long sample time.



# 15.3 On-Chip Debug System (DBG)

Because HCS08 devices do not have external address and data buses, the most important functions of an in-circuit emulator have been built onto the chip with the MCU. The debug system consists of an 8-stage FIFO that can store address or data bus information, and a flexible trigger system to decide when to capture bus information and what information to capture. The system relies on the single-wire background debug system to access debug control registers and to read results out of the eight stage FIFO.

The debug module includes control and status registers that are accessible in the user's memory map. These registers are located in the high register space to avoid using valuable direct page memory space.

Most of the debug module's functions are used during development, and user programs rarely access any of the control and status registers for the debug module. The one exception is that the debug system can provide the means to implement a form of ROM patching. This topic is discussed in greater detail in Section 15.3.6, "Hardware Breakpoints."

#### 15.3.1 Comparators A and B

Two 16-bit comparators (A and B) can optionally be qualified with the R/W signal and an opcode tracking circuit. Separate control bits allow you to ignore R/W for each comparator. The opcode tracking circuitry optionally allows you to specify that a trigger will occur only if the opcode at the specified address is actually executed as opposed to only being read from memory into the instruction queue. The comparators are also capable of magnitude comparisons to support the inside range and outside range trigger modes. Comparators are disabled temporarily during all BDC accesses.

The A comparator is always associated with the 16-bit CPU address. The B comparator compares to the CPU address or the 8-bit CPU data bus, depending on the trigger mode selected. Because the CPU data bus is separated into a read data bus and a write data bus, the RWAEN and RWA control bits have an additional purpose, in full address plus data comparisons they are used to decide which of these buses to use in the comparator B data bus comparisons. If RWAEN = 1 (enabled) and RWA = 0 (write), the CPU's write data bus is used. Otherwise, the CPU's read data bus is used.

The currently selected trigger mode determines what the debugger logic does when a comparator detects a qualified match condition. A match can cause:

- Generation of a breakpoint to the CPU
- Storage of data bus values into the FIFO
- Starting to store change-of-flow addresses into the FIFO (begin type trace)
- Stopping the storage of change-of-flow addresses into the FIFO (end type trace)

#### 15.3.2 Bus Capture Information and FIFO Operation

The usual way to use the FIFO is to setup the trigger mode and other control options, then arm the debugger. When the FIFO has filled or the debugger has stopped storing data into the FIFO, you would read the information out of it in the order it was stored into the FIFO. Status bits indicate the number of words of valid information that are in the FIFO as data is stored into it. If a trace run is manually halted by writing 0 to ARM before the FIFO is full (CNT = 1:0:0:0), the information is shifted by one position and



Chapter 15 Development Support

### 15.4.3.9 Debug Status Register (DBGS)

This is a read-only status register.



#### Figure 15-9. Debug Status Register (DBGS)

#### Table 15-6. DBGS Register Field Descriptions

Field	Description
7 AF	<ul> <li>Trigger Match A Flag — AF is cleared at the start of a debug run and indicates whether a trigger match A condition was met since arming.</li> <li>0 Comparator A has not matched</li> <li>1 Comparator A match</li> </ul>
6 BF	<ul> <li>Trigger Match B Flag — BF is cleared at the start of a debug run and indicates whether a trigger match B condition was met since arming.</li> <li>0 Comparator B has not matched</li> <li>1 Comparator B match</li> </ul>
5 ARMF	<ul> <li>Arm Flag — While DBGEN = 1, this status bit is a read-only image of ARM in DBGC. This bit is set by writing 1 to the ARM control bit in DBGC (while DBGEN = 1) and is automatically cleared at the end of a debug run. A debug run is completed when the FIFO is full (begin trace) or when a trigger event is detected (end trace). A debug run can also be ended manually by writing 0 to ARM or DBGEN in DBGC.</li> <li>0 Debugger not armed</li> <li>1 Debugger armed</li> </ul>
3:0 CNT[3:0]	FIFO Valid Count — These bits are cleared at the start of a debug run and indicate the number of words of valid data in the FIFO at the end of a debug run. The value in CNT does not decrement as data is read out of the FIFO. The external debug host is responsible for keeping track of the count as information is read out of the FIFO. 0000 Number of valid words in FIFO = No valid data 0001 Number of valid words in FIFO = 1 0010 Number of valid words in FIFO = 2 0011 Number of valid words in FIFO = 3 0100 Number of valid words in FIFO = 4 0101 Number of valid words in FIFO = 5 0110 Number of valid words in FIFO = 5 0110 Number of valid words in FIFO = 7 1000 Number of valid words in FIFO = 7 1000 Number of valid words in FIFO = 8



# A.10 AC Characteristics

This section describes ac timing characteristics for each peripheral system. For detailed information about how clocks for the bus are generated, see Chapter 8, "Internal Clock Generator (S08ICGV4)."

### A.10.1 Control Timing

Num	С	Parameter	Symbol	Min	Typ <sup>1</sup>	Max	Unit
1		Bus frequency (t <sub>cyc</sub> = 1/f <sub>Bus</sub> )	f <sub>Bus</sub>	dc	_	20	MHz
2	Р	Real-time interrupt internal oscillator period	t <sub>RTI</sub>	700		1300	μs
3		External reset pulse width <sup>2</sup> ( $t_{cyc} = 1/f_{Self\_reset}$ )	t <sub>extrst</sub>	1.5 x t <sub>Self_reset</sub>		_	ns
4		Reset low drive <sup>3</sup>	t <sub>rstdrv</sub>	34 x t <sub>cyc</sub>		_	ns
5		Active background debug mode latch setup time	t <sub>MSSU</sub>	25		—	ns
6		Active background debug mode latch hold time	t <sub>MSH</sub>	25		_	ns
7		IRQ pulse width Asynchronous path <sup>2</sup> Synchronous path <sup>4</sup>	t <sub>ILIH,</sub> t <sub>IHIL</sub>	100 1.5 x t <sub>cyc</sub>	_	_	ns
8		KBIPx pulse width Asynchronous path <sup>2</sup> Synchronous path <sup>3</sup>	t <sub>ILIH</sub> , tIHIL	100 1.5 x t <sub>cyc</sub>	_	_	ns
9	т	Port rise and fall time — Low output drive (PTxDS = 0) (load = 50 pF) <sup>5</sup> Slew rate control disabled (PTxSE = 0) Slew rate control enabled (PTxSE = 1)	t <sub>Rise</sub> , t <sub>Fall</sub>		40 75		ns
		Port rise and fall time — High output drive (PTxDS = 1) (load = 50 pF) Slew rate control disabled (PTxSE = 0) Slew rate control enabled (PTxSE = 1)	t <sub>Rise</sub> , t <sub>Fall</sub>		11 35		ns

Table A-13. Control Tir	ning
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<sup>1</sup> Typical values are based on characterization data at  $V_{DD} = 5.0V$ , 25°C unless otherwise stated.

<sup>2</sup> This is the shortest pulse that is guaranteed to be recognized as a reset pin request. Shorter pulses are not guaranteed to override reset requests from internal sources.

<sup>3</sup> When any reset is initiated, internal circuitry drives the reset pin low for about 34 bus cycles and then samples the level on the reset pin about 38 bus cycles later to distinguish external reset requests from internal requests.

<sup>4</sup> This is the minimum pulse width that is guaranteed to pass through the pin synchronization circuitry. Shorter pulses may or may not be recognized. In stop mode, the synchronizer is bypassed so shorter pulses can be recognized in that case.

 $^5~$  Timing is shown with respect to 20%  $V_{DD}$  and 80%  $V_{DD}$  levels. Temperature range –40°C to 125°C.



Appendix A Electrical Characteristics and Timing Specifications



Figure A-10. Typical RTI Clock Period vs. Temperature



Figure A-11. Reset Timing







Figure A-13. IRQ/KBIPx Timing

MC9S08AW60 Data Sheet, Rev 2



Appendix B Ordering Information and Mechanical Drawings

# B.2 Orderable Part Numbering System

### B.2.1 Consumer and Industrial Orderable Part Numbering System



# B.2.2 Automotive Orderable Part Numbering System



# B.3 Mechanical Drawings

This following pages contain mechanical specifications for MC9S08AW60 Series package options. See Table B-3 for the document numbers that correspond to each package type.

Pin Count	Туре	Designator	Document No.
44	LQFP	FG	98ASS23225W
48	QFN	FD	98ARH99048A
64	LQFP	PU	98ASS23234W
64	QFP	FU	98ASB42844B

Table B-3. Package Information





DETAIL "A"

SECTION B-B



DETAIL "C"

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TITLE:		DOCUMENT NE	]: 98ASB42844B	REV∶B	
64LD QFP (14 X	CASE NUMBER	2: 840B-01	20 MAY 2005		
		STANDARD: NE	IN-JEDEC		