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Details

Product Status	Obsolete
Core Processor	HCS08
Core Size	8-Bit
Speed	40MHz
Connectivity	I ² C, SCI, SPI
Peripherals	LVD, POR, PWM, WDT
Number of I/O	38
Program Memory Size	60KB (60K x 8)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	2K x 8
Voltage - Supply (Vcc/Vdd)	2.7V ~ 5.5V
Data Converters	A/D 8x10b SAR
Oscillator Type	External, Internal
Operating Temperature	-40°C ~ 125°C (TA)
Mounting Type	Surface Mount
Package / Case	48-VFQFN Exposed Pad
Supplier Device Package	48-QFN (7x7)
Purchase URL	https://www.e-xfl.com/product-detail/nxp-semiconductors/s9s08aw60e5mfde

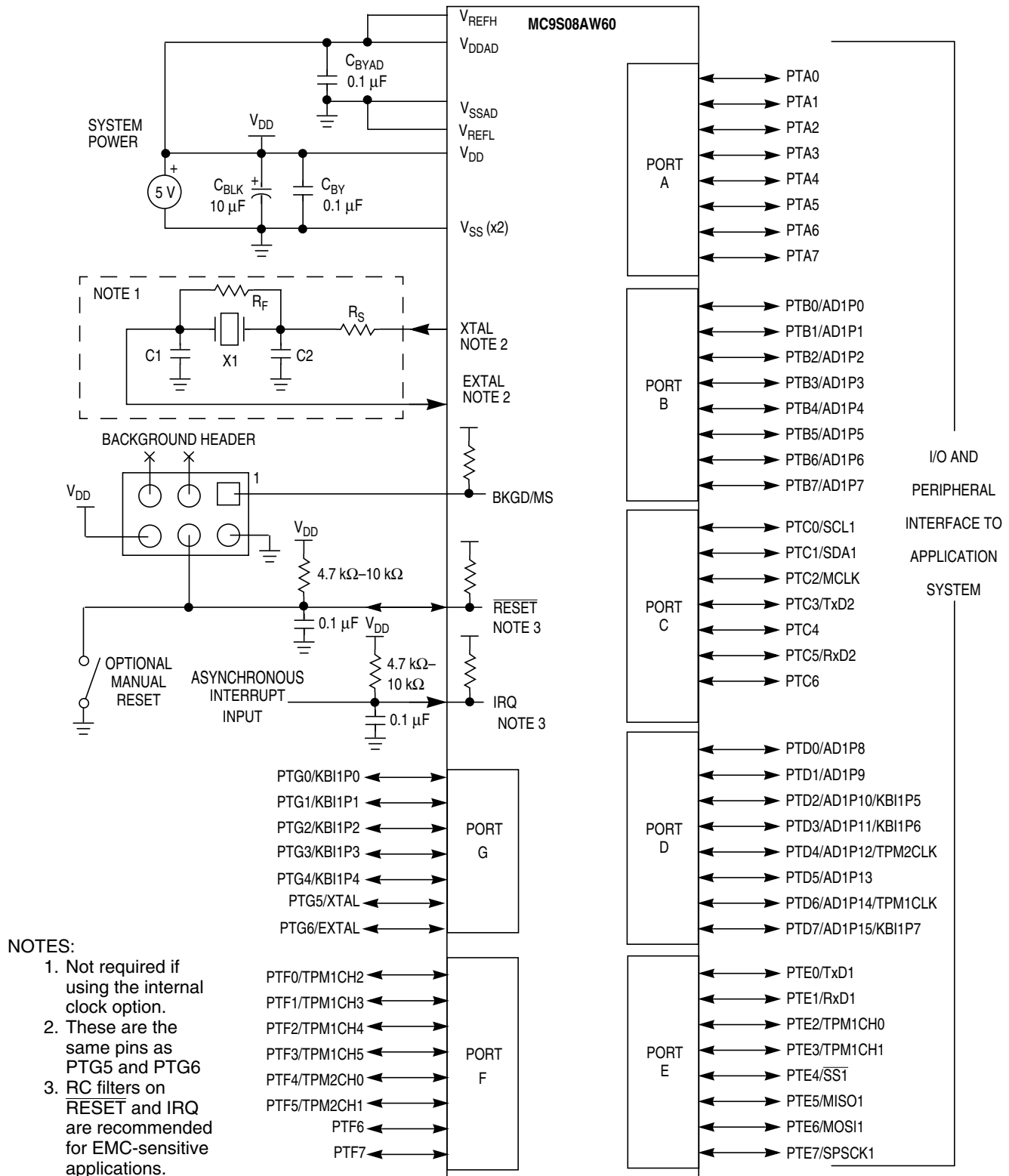


Figure 2-4. Basic System Connections



Table 4-2. Direct-Page Register Summary (Sheet 1 of 3)

Address	Register Name	Bit 7	6	5	4	3	2	1	Bit 0
\$0000	PTAD	PTAD7	PTAD6	PTAD5	PTAD4	PTAD3	PTAD2	PTAD1	PTAD0
\$0001	PTADD	PTADD7	PTADD6	PTADD5	PTADD4	PTADD3	PTADD2	PTADD1	PTADD0
\$0002	PTBD	PTBD7	PTBD6	PTBD5	PTBD4	PTBD3	PTBD2	PTBD1	PTBD0
\$0003	PTBDD	PTBDD7	PTBDD6	PTBDD5	PTBDD4	PTBDD3	PTBDD2	PTBDD1	PTBDD0
\$0004	PTCD	0	PTCD6	PTCD5	PTCD4	PTCD3	PTCD2	PTCD1	PTCD0
\$0005	PTCDD	0	PTCDD6	PTCDD5	PTCDD4	PTCDD3	PTCDD2	PTCDD1	PTCDD0
\$0006	PTDD	PTDD7	PTDD6	PTDD5	PTDD4	PTDD3	PTDD2	PTDD1	PTDD0
\$0007	PTDDD	PTDDD7	PTDDD6	PTDDD5	PTDDD4	PTDDD3	PTDDD2	PTDDD1	PTDDD0
\$0008	PTED	PTED7	PTED6	PTED5	PTED4	PTED3	PTED2	PTED1	PTED0
\$0009	PTEDD	PTEDD7	PTEDD6	PTEDD5	PTEDD4	PTEDD3	PTEDD2	PTEDD1	PTEDD0
\$000A	PTFD	PTFD7	PTFD6	PTFD5	PTFD4	PTFD3	PTFD2	PTFD1	PTFD0
\$000B	PTFDD	PTFDD7	PTFDD6	PTFDD5	PTFDD4	PTFDD3	PTFDD2	PTFDD1	PTFDD0
\$000C	PTGD	0	PTGD6	PTGD5	PTGD4	PTGD3	PTGD2	PTGD1	PTGD0
\$000D	PTGDD	0	PTGDD6	PTGDD5	PTGDD4	PTGDD3	PTGDD2	PTGDD1	PTGDD0
\$000E– \$000F	Reserved	—	—	—	—	—	—	—	—
\$0010	ADC1SC1	COCO	AIEN	ADCO	ADCH				
\$0011	ADC1SC2	ADACT	ADTRG	ACFE	ACFGT	0	0	R	R
\$0012	ADC1RH	0	0	0	0	0	0	ADR9	ADR8
\$0013	ADC1RL	ADR7	ADR6	ADR5	ADR4	ADR3	ADR2	ADR1	ADR0
\$0014	ADC1CVH	0	0	0	0	0	0	ADCV9	ADCV8
\$0015	ADC1CVL	ADCV7	ADCV6	ADCV5	ADCV4	ADCV3	ADCV2	ADCV1	ADCV0
\$0016	ADC1CFG	ADLPC	ADIV		ADLSMP	MODE		ADICLK	
\$0017	APCTL1	ADPC7	ADPC6	ADPC5	ADPC4	ADPC3	ADPC2	ADPC1	ADPC0
\$0018	APCTL2	ADPC15	ADPC14	ADPC13	ADPC12	ADPC11	ADPC10	ADPC9	ADPC8
\$0019	APCTL3	ADPC23	ADPC22	ADPC21	ADPC20	ADPC19	ADPC18	ADPC17	ADPC16
\$001A– \$001B	Reserved	—	—	—	—	—	—	—	—
\$001C	IRQSC	0	0	IRQEDG	IRQPE	IRQF	IRQACK	IRQIE	IRQMOD
\$001D	Reserved	—	—	—	—	—	—	—	—
\$001E	KBI1SC	KBEDG7	KBEDG6	KBEDG5	KBEDG4	KBF	KBACK	KBIE	KBIMOD
\$001F	KBI1PE	KBIPE7	KBIPE6	KBIPE5	KBIPE4	KBIPE3	KBIPE2	KBIPE1	KBIPE0
\$0020	TPM1SC	TOF	TOIE	CPWMS	CLKSB	CLKSA	PS2	PS1	PS0
\$0021	TPM1CNTH	Bit 15	14	13	12	11	10	9	Bit 8
\$0022	TPM1CNTL	Bit 7	6	5	4	3	2	1	Bit 0
\$0023	TPM1MODH	Bit 15	14	13	12	11	10	9	Bit 8
\$0024	TPM1MODL	Bit 7	6	5	4	3	2	1	Bit 0
\$0025	TPM1C0SC	CH0F	CH0IE	MS0B	MS0A	ELS0B	ELS0A	0	0
\$0026	TPM1C0VH	Bit 15	14	13	12	11	10	9	Bit 8
\$0027	TPM1C0VL	Bit 7	6	5	4	3	2	1	Bit 0

High-page registers, shown in Table 4-3, are accessed much less often than other I/O and control registers so they have been located outside the direct addressable memory space, starting at \$1800.

Table 4-3. High-Page Register Summary (Sheet 1 of 2)

Address	Register Name	Bit 7	6	5	4	3	2	1	Bit 0
\$1800	SRS	POR	PIN	COP	ILOP	0	ICG	LVD	0
\$1801	SBD FR	0	0	0	0	0	0	0	BDFR
\$1802	SOPT	COPE	COPT	STOPE	—	0	0	—	—
\$1803	SMCLK	0	0	0	MPE	0	MCSEL		
\$1804 — \$1805	Reserved	—	—	—	—	—	—	—	—
\$1806	SDIDH	REV3	REV2	REV1	REV0	ID11	ID10	ID9	ID8
\$1807	SDIDL	ID7	ID6	ID5	ID4	ID3	ID2	ID1	ID0
\$1808	SRTISC	RTIF	RTIACK	RTICLKS	RTIE	0	RTIS2	RTIS1	RTIS0
\$1809	SPMSC1	LVDF	LVDACK	LVDIE	LVDRE	LVDSE	LVDE	0 ¹	BGBE
\$180A	SPMSC2	LVWF	LVWACK	LVDV	LVWV	PPDF	PPDACK	—	PPDC
\$180B — \$180F	Reserved	—	—	—	—	—	—	—	—
\$1810	DBGCAH	Bit 15	14	13	12	11	10	9	Bit 8
\$1811	DBGCAL	Bit 7	6	5	4	3	2	1	Bit 0
\$1812	DBGCBH	Bit 15	14	13	12	11	10	9	Bit 8
\$1813	DBGCBL	Bit 7	6	5	4	3	2	1	Bit 0
\$1814	DBGFHH	Bit 15	14	13	12	11	10	9	Bit 8
\$1815	DBGFL	Bit 7	6	5	4	3	2	1	Bit 0
\$1816	DBG C	DBGEN	ARM	TAG	BRKEN	RWA	RWAEN	RWB	RWBEN
\$1817	DBG T	TRGSEL	BEGIN	0	0	TRG3	TRG2	TRG1	TRG0
\$1818	DBG S	AF	BF	ARMF	0	CNT3	CNT2	CNT1	CNT0
\$1819 — \$181F	Reserved	—	—	—	—	—	—	—	—
\$1820	FCDIV	DIVLD	PRDIV8	DIV5	DIV4	DIV3	DIV2	DIV1	DIV0
\$1821	FOPT	KEYEN	FNORED	0	0	0	0	SEC01	SEC00
\$1822	Reserved	—	—	—	—	—	—	—	—
\$1823	FCNFG	0	0	KEYACC	0	0	0	0	0
\$1824	FPROT	FPS7	FPS6	FPS5	FPS4	FPS3	FPS2	FPS1	FPDIS
\$1825	FSTAT	FCBEF	FCCF	FPVIOL	FACCERR	0	FBLANK	0	0
\$1826	FCMD	FCMD7	FCMD6	FCMD5	FCMD4	FCMD3	FCMD2	FCMD1	FCMD0
\$1827 — \$183F	Reserved	—	—	—	—	—	—	—	—
\$1840	PTAPE	PTAPE7	PTAPE6	PTAPE5	PTAPE4	PTAPE3	PTAPE2	PTAPE1	PTAPE0
\$1841	PTASE	PTASE7	PTASE6	PTASE5	PTASE4	PTASE3	PTASE2	PTASE1	PTASE0
\$1842	PTADS	PTADS7	PTADS6	PTADS5	PTADS4	PTADS3	PTADS2	PTADS1	PTADS0
\$1843	Reserved	—	—	—	—	—	—	—	—
\$1844	PTBPE	PTBPE7	PTBPE6	PTBPE5	PTBPE4	PTBPE3	PTBPE2	PTBPE1	PTBPE0
\$1845	PTBSE	PTBSE7	PTBSE6	PTBSE5	PTBSE4	PTBSE3	PTBSE2	PTBSE1	PTBSE0

Table 4-12. FSTAT Register Field Descriptions

Field	Description
7 FCBEF	FLASH Command Buffer Empty Flag — The FCBEF bit is used to launch commands. It also indicates that the command buffer is empty so that a new command sequence can be executed when performing burst programming. The FCBEF bit is cleared by writing a one to it or when a burst program command is transferred to the array for programming. Only burst program commands can be buffered. 0 Command buffer is full (not ready for additional commands). 1 A new burst program command may be written to the command buffer.
6 FCCF	FLASH Command Complete Flag — FCCF is set automatically when the command buffer is empty and no command is being processed. FCCF is cleared automatically when a new command is started (by writing 1 to FCBEF to register a command). Writing to FCCF has no meaning or effect. 0 Command in progress 1 All commands complete
5 FPVIOL	Protection Violation Flag — FPVIOL is set automatically when FCBEF is cleared to register a command that attempts to erase or program a location in a protected block (the erroneous command is ignored). FPVIOL is cleared by writing a 1 to FPVIOL. 0 No protection violation. 1 An attempt was made to erase or program a protected location.
4 FACCERR	Access Error Flag — FACCERR is set automatically when the proper command sequence is not obeyed exactly (the erroneous command is ignored), if a program or erase operation is attempted before the FCDIV register has been initialized, or if the MCU enters stop while a command was in progress. For a more detailed discussion of the exact actions that are considered access errors, see Section 4.4.5, “Access Errors.” FACCERR is cleared by writing a 1 to FACCERR. Writing a 0 to FACCERR has no meaning or effect. 0 No access error. 1 An access error has occurred.
2 FBLANK	FLASH Verified as All Blank (erased) Flag — FBLANK is set automatically at the conclusion of a blank check command if the entire FLASH array was verified to be erased. FBLANK is cleared by clearing FCBEF to write a new valid command. Writing to FBLANK has no meaning or effect. 0 After a blank check command is completed and FCCF = 1, FBLANK = 0 indicates the FLASH array is not completely erased. 1 After a blank check command is completed and FCCF = 1, FBLANK = 1 indicates the FLASH array is completely erased (all \$FF).

6.4 Parallel I/O Control

Reading and writing of parallel I/O is done through the port data registers. The direction, input or output, is controlled through the port data direction registers. The parallel I/O port function for an individual pin is illustrated in the block diagram below.

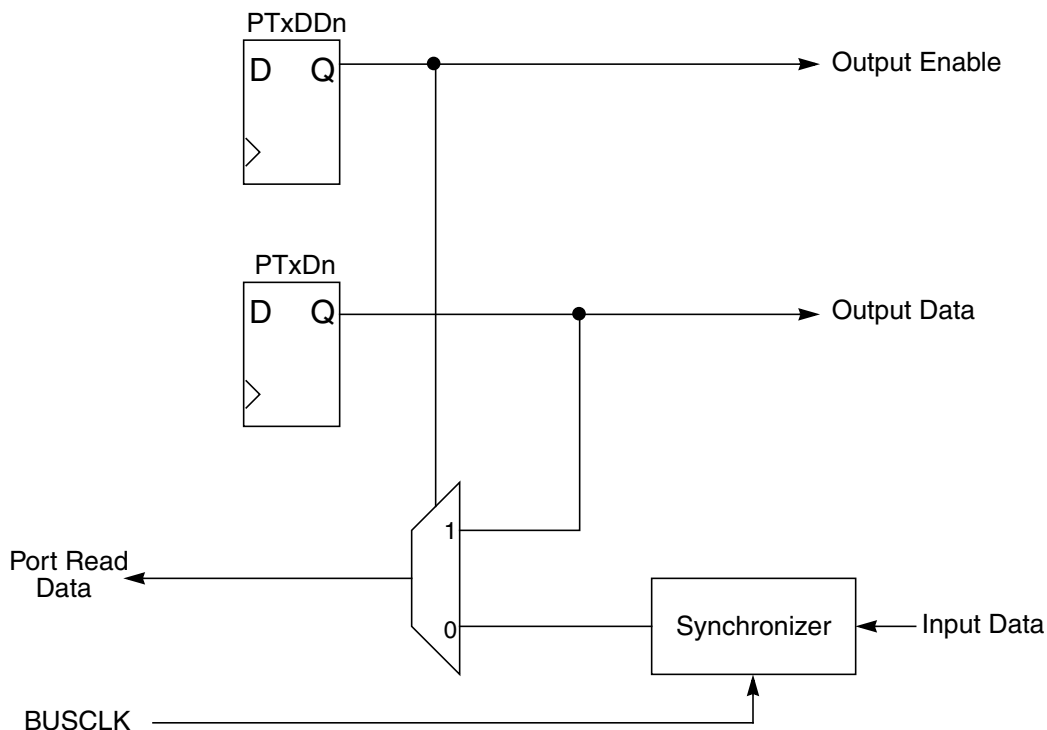


Figure 6-8. Parallel I/O Block Diagram

The data direction control bits determine whether the pin output driver is enabled, and they control what is read for port data register reads. Each port pin has a data direction register bit. When $PTxDDn = 0$, the corresponding pin is an input and reads of $PTxD$ return the pin value. When $PTxDDn = 1$, the corresponding pin is an output and reads of $PTxD$ return the last value written to the port data register. When a peripheral module or system function is in control of a port pin, the data direction register bit still controls what is returned for reads of the port data register, even though the peripheral system has overriding control of the actual pin direction.

When a shared analog function is enabled for a pin, all digital pin functions are disabled. A read of the port data register returns a value of 0 for any bits which have shared analog functions enabled. In general, whenever a pin is shared with both an alternate digital function and an analog function, the analog function has priority such that if both the digital and analog functions are enabled, the analog function controls the pin.

It is a good programming practice to write to the port data register before changing the direction of a port pin to become an output. This ensures that the pin will not be driven momentarily with an old data value that happened to be in the port data register.

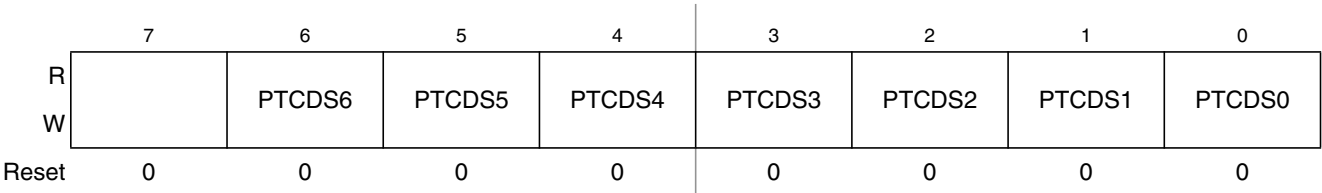


Figure 6-23. Output Drive Strength Selection for Port C (PTCDS)

Table 6-16. PTCDS Register Field Descriptions

Field	Description
6:0 PTCDS[6:0]	Output Drive Strength Selection for Port C Bits — Each of these control bits selects between low and high output drive for the associated PTC pin. 0 Low output drive enabled for port C bit n. 1 High output drive enabled for port C bit n.

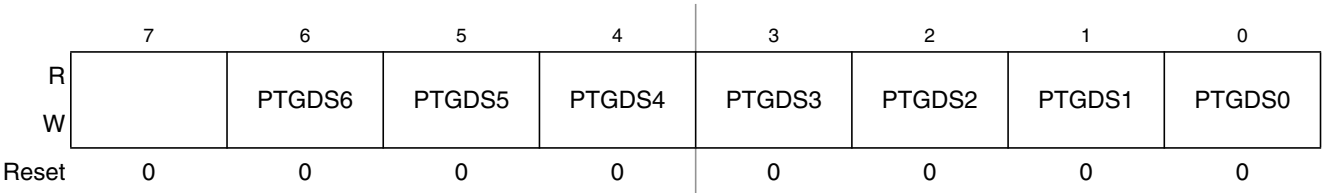


Figure 6-43. Output Drive Strength Selection for Port G (PTGDS)

Table 6-36. PTGDS Register Field Descriptions

Field	Description
6:0 PTGDS[6:0]	Output Drive Strength Selection for Port G Bits — Each of these control bits selects between low and high output drive for the associated PTG pin. 0 Low output drive enabled for port G bit n. 1 High output drive enabled for port G bit n.

Chapter 7

Central Processor Unit (S08CPUV2)

7.1 Introduction

This section provides summary information about the registers, addressing modes, and instruction set of the CPU of the HCS08 family. For a more detailed discussion, refer to the *HCS08 Family Reference Manual, volume 1*, Freescale Semiconductor document order number HCS08RMV1/D.

The HCS08 CPU is fully source- and object-code-compatible with the M68HC08 CPU. Several instructions and enhanced addressing modes were added to improve C compiler efficiency and to support a new background debug system which replaces the monitor mode of earlier M68HC08 microcontrollers (MCU).

7.1.1 Features

Features of the HCS08 CPU include:

- Object code fully upward-compatible with M68HC05 and M68HC08 Families
- All registers and memory are mapped to a single 64-Kbyte address space
- 16-bit stack pointer (any size stack anywhere in 64-Kbyte address space)
- 16-bit index register (H:X) with powerful indexed addressing modes
- 8-bit accumulator (A)
- Many instructions treat X as a second general-purpose 8-bit register
- Seven addressing modes:
 - Inherent — Operands in internal registers
 - Relative — 8-bit signed offset to branch destination
 - Immediate — Operand in next object code byte(s)
 - Direct — Operand in memory at 0x0000–0x00FF
 - Extended — Operand anywhere in 64-Kbyte address space
 - Indexed relative to H:X — Five submodes including auto increment
 - Indexed relative to SP — Improves C efficiency dramatically
- Memory-to-memory data move instructions with four address mode combinations
- Overflow, half-carry, negative, zero, and carry condition codes support conditional branching on the results of signed, unsigned, and binary-coded decimal (BCD) operations
- Efficient bit manipulation instructions
- Fast 8-bit by 8-bit multiply and 16-bit by 8-bit divide instructions
- STOP and WAIT instructions to invoke low-power operating modes

7.3.6.7 SP-Relative, 16-Bit Offset (SP2)

This variation of indexed addressing uses the 16-bit value in the stack pointer (SP) plus a 16-bit offset included in the instruction as the address of the operand needed to complete the instruction.

7.4 Special Operations

The CPU performs a few special operations that are similar to instructions but do not have opcodes like other CPU instructions. In addition, a few instructions such as STOP and WAIT directly affect other MCU circuitry. This section provides additional information about these operations.

7.4.1 Reset Sequence

Reset can be caused by a power-on-reset (POR) event, internal conditions such as the COP (computer operating properly) watchdog, or by assertion of an external active-low reset pin. When a reset event occurs, the CPU immediately stops whatever it is doing (the MCU does not wait for an instruction boundary before responding to a reset event). For a more detailed discussion about how the MCU recognizes resets and determines the source, refer to the Resets, Interrupts, and System Configuration chapter.

The reset event is considered concluded when the sequence to determine whether the reset came from an internal source is done and when the reset pin is no longer asserted. At the conclusion of a reset event, the CPU performs a 6-cycle sequence to fetch the reset vector from 0xFFFFE and 0xFFFF and to fill the instruction queue in preparation for execution of the first program instruction.

7.4.2 Interrupt Sequence

When an interrupt is requested, the CPU completes the current instruction before responding to the interrupt. At this point, the program counter is pointing at the start of the next instruction, which is where the CPU should return after servicing the interrupt. The CPU responds to an interrupt by performing the same sequence of operations as for a software interrupt (SWI) instruction, except the address used for the vector fetch is determined by the highest priority interrupt that is pending when the interrupt sequence started.

The CPU sequence for an interrupt is:

1. Store the contents of PCL, PCH, X, A, and CCR on the stack, in that order.
2. Set the I bit in the CCR.
3. Fetch the high-order half of the interrupt vector.
4. Fetch the low-order half of the interrupt vector.
5. Delay for one free bus cycle.
6. Fetch three bytes of program information starting at the address indicated by the interrupt vector to fill the instruction queue in preparation for execution of the first instruction in the interrupt service routine.

After the CCR contents are pushed onto the stack, the I bit in the CCR is set to prevent other interrupts while in the interrupt service routine. Although it is possible to clear the I bit with an instruction in the

8.3.2 ICG Control Register 2 (ICGC2)

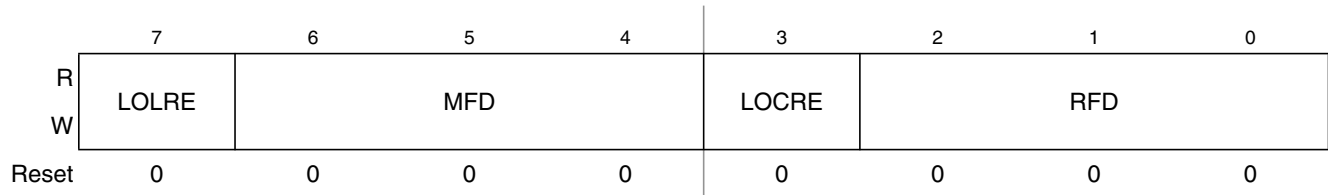


Figure 8-7. ICG Control Register 2 (ICGC2)

Table 8-2. ICGC2 Register Field Descriptions

Field	Description
7 LOLRE	Loss of Lock Reset Enable — The LOLRE bit determines what type of request is made by the ICG following a loss of lock indication. The LOLRE bit only has an effect when LOLS is set. 0 Generate an interrupt request on loss of lock. 1 Generate a reset request on loss of lock.
6:4 MFD	Multiplication Factor — The MFD bits control the programmable multiplication factor in the FLL loop. The value specified by the MFD bits establishes the multiplication factor (N) applied to the reference frequency. Writes to the MFD bits will not take effect if a previous write is not complete. Select a low enough value for N such that $f_{ICGDCLK}$ does not exceed its maximum specified value. 000 Multiplication factor = 4 001 Multiplication factor = 6 010 Multiplication factor = 8 011 Multiplication factor = 10 100 Multiplication factor = 12 101 Multiplication factor = 14 110 Multiplication factor = 16 111 Multiplication factor = 18
3 LOCRE	Loss of Clock Reset Enable — The LOCRE bit determines how the system manages a loss of clock condition. 0 Generate an interrupt request on loss of clock. 1 Generate a reset request on loss of clock.
2:0 RFD	Reduced Frequency Divider — The RFD bits control the value of the divider following the clock select circuitry. The value specified by the RFD bits establishes the division factor (R) applied to the selected output clock source. Writes to the RFD bits will not take effect if a previous write is not complete. 000 Division factor = 1 001 Division factor = 2 010 Division factor = 4 011 Division factor = 8 100 Division factor = 16 101 Division factor = 32 110 Division factor = 64 111 Division factor = 128

entering off mode. If CLKS bits are set to 01 or 11 coming out of the Off state, the ICG enters this mode until ICGDCLK is stable as determined by the DCOS bit. After ICGDCLK is considered stable, the ICG automatically closes the loop by switching to FLL engaged (internal or external) as selected by the CLKS bits.

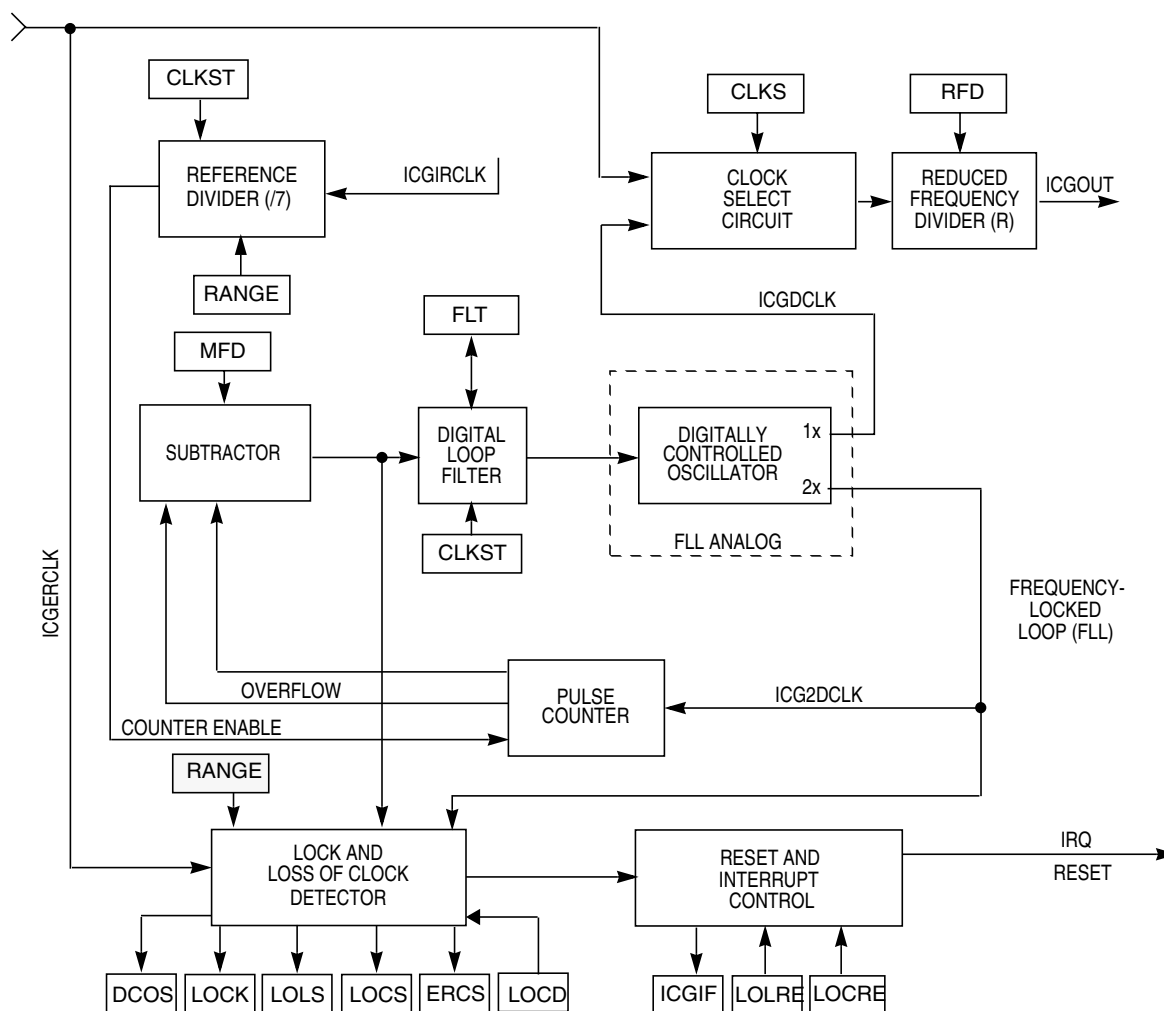


Figure 8-13. Detailed Frequency-Locked Loop Block Diagram

8.4.3 FLL Engaged, Internal Clock (FEI) Mode

FLL engaged internal (FEI) is entered when any of the following conditions occur:

- CLKS bits are written to 01
- The DCO clock stabilizes ($DCOS = 1$) while in SCM upon exiting the off state with $CLKS = 01$

In FLL engaged internal mode, the reference clock is derived from the internal reference clock ICGIRCLK, and the FLL loop will attempt to lock the ICGDCLK frequency to the desired value, as selected by the MFD bits.

ICGTRM = \$xx

Bit 7:0 TRIM

Only need to write when trimming internal oscillator; done in separate operation (see example #4)

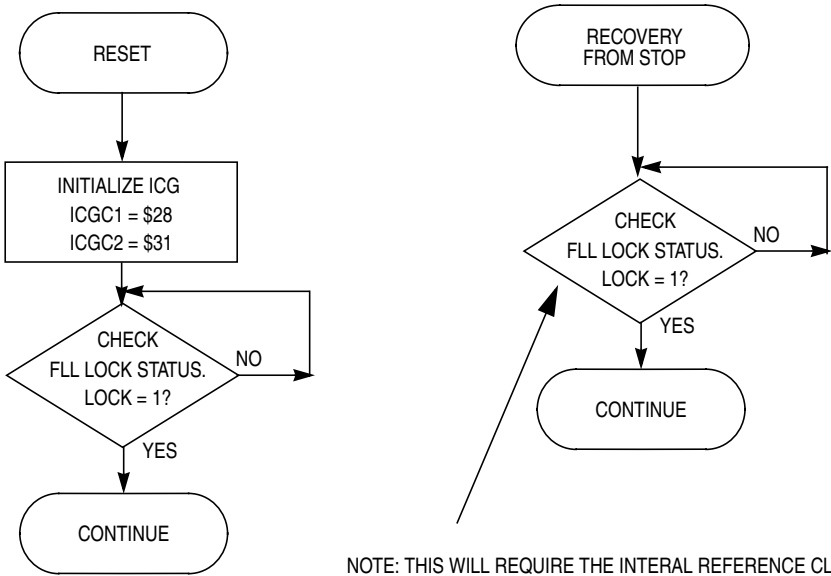


Figure 8-16. ICG Initialization and Stop Recovery for Example #3



9.4.1 KBI Status and Control Register (KBI1SC)

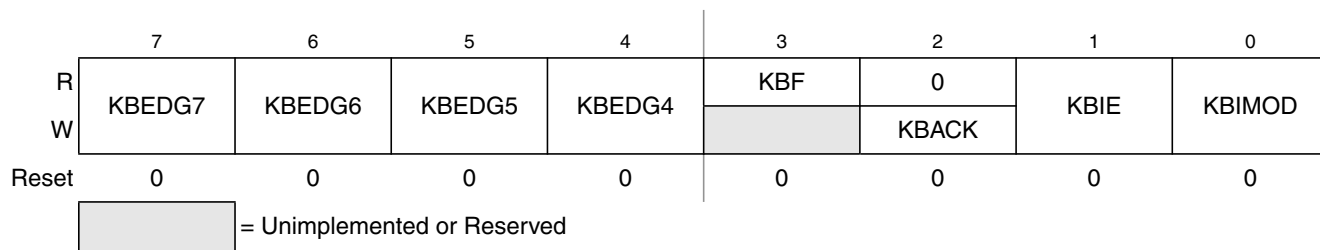


Figure 9-3. KBI Status and Control Register (KBI1SC)

Table 9-2. KBI1SC Register Field Descriptions

Field	Description
7:4 KBEDG[7:4]	Keyboard Edge Select for KBI Port Bits — Each of these read/write bits selects the polarity of the edges and/or levels that are recognized as trigger events on the corresponding KBI port pin when it is configured as a keyboard interrupt input (KBIPEn = 1). Also see the KBIMOD control bit, which determines whether the pin is sensitive to edges-only or edges and levels. 0 Falling edges/low levels 1 Rising edges/high levels
3 KBF	Keyboard Interrupt Flag — This read-only status flag is set whenever the selected edge event has been detected on any of the enabled KBI port pins. This flag is cleared by writing a 1 to the KBACK control bit. The flag will remain set if KBIMOD = 1 to select edge-and-level operation and any enabled KBI port pin remains at the asserted level. KBF can be used as a software pollable flag (KBIE = 0) or it can generate a hardware interrupt request to the CPU (KBIE = 1). 0 No KBI interrupt pending 1 KBI interrupt pending
2 KBACK	Keyboard Interrupt Acknowledge — This write-only bit (reads always return 0) is used to clear the KBF status flag by writing a 1 to KBACK. When KBIMOD = 1 to select edge-and-level operation and any enabled KBI port pin remains at the asserted level, KBF is being continuously set so writing 1 to KBACK does not clear the KBF flag.
1 KBIE	Keyboard Interrupt Enable — This read/write control bit determines whether hardware interrupts are generated when the KBF status flag equals 1. When KBIE = 0, no hardware interrupts are generated, but KBF can still be used for software polling. 0 KBF does not generate hardware interrupts (use polling) 1 KBI hardware interrupt requested when KBF = 1
KBIMOD	Keyboard Detection Mode — This read/write control bit selects either edge-only detection or edge-and-level detection. KBI port bits 3 through 0 can detect falling edges-only or falling edges and low levels. KBI port bits 7 through 4 can be configured to detect either: - Rising edges-only or rising edges and high levels (KBEDGn = 1) - Falling edges-only or falling edges and low levels (KBEDGn = 0) 0 Edge-only detection 1 Edge-and-level detection

transferred to the corresponding timer channel registers only after both 8-bit bytes of a 16-bit register have been written and the timer counter overflows (reverses direction from up-counting to down-counting at the end of the terminal count in the modulus register). This TPMxCNT overflow requirement only applies to PWM channels, not output compares.

Optionally, when $\text{TPMxCNTH}:\text{TPMxCNTL} = \text{TPMxMODH}:\text{TPMxMODL}$, the TPM can generate a TOF interrupt at the end of this count. The user can choose to reload any number of the PWM buffers, and they will all update simultaneously at the start of a new period.

Writing to TPMxSC cancels any values written to TPMxMODH and/or TPMxMODL and resets the coherency mechanism for the modulo registers. Writing to TPMxCnSC cancels any values written to the channel value registers and resets the coherency mechanism for TPMxCnVH:TPMxCnVL.

10.6 TPM Interrupts

The TPM generates an optional interrupt for the main counter overflow and an interrupt for each channel. The meaning of channel interrupts depends on the mode of operation for each channel. If the channel is configured for input capture, the interrupt flag is set each time the selected input capture edge is recognized. If the channel is configured for output compare or PWM modes, the interrupt flag is set each time the main timer counter matches the value in the 16-bit channel value register. See the Resets, Interrupts, and System Configuration chapter for absolute interrupt vector addresses, priority, and local interrupt mask control bits.

For each interrupt source in the TPM, a flag bit is set on recognition of the interrupt condition such as timer overflow, channel input capture, or output compare events. This flag may be read (polled) by software to verify that the action has occurred, or an associated enable bit (TOIE or CHnIE) can be set to enable hardware interrupt generation. While the interrupt enable bit is set, a static interrupt will be generated whenever the associated interrupt flag equals 1. It is the responsibility of user software to perform a sequence of steps to clear the interrupt flag before returning from the interrupt service routine.

10.6.1 Clearing Timer Interrupt Flags

TPM interrupt flags are cleared by a 2-step process that includes a read of the flag bit while it is set (1) followed by a write of 0 to the bit. If a new event is detected between these two steps, the sequence is reset and the interrupt flag remains set after the second step to avoid the possibility of missing the new event.

10.6.2 Timer Overflow Interrupt Description

The conditions that cause TOF to become set depend on the counting mode (up or up/down). In up-counting mode, the 16-bit timer counter counts from 0x0000 through 0xFFFF and overflows to 0x0000 on the next counting clock. TOF becomes set at the transition from 0xFFFF to 0x0000. When a modulus limit is set, TOF becomes set at the transition from the value set in the modulus register to 0x0000. When the counter is operating in up-/down-counting mode, the TOF flag gets set as the counter changes direction at the transition from the value set in the modulus register and the next lower count value. This corresponds to the end of a PWM period. (The 0x0000 count value corresponds to the center of a period.)

systems with modems to determine when it is safe to turn off the modem. If the transmit complete interrupt enable (TCIE) bit is set, a hardware interrupt will be requested whenever $TC = 1$. Instead of hardware interrupts, software polling may be used to monitor the TDRE and TC status flags if the corresponding TIE or TCIE local interrupt masks are 0s.

When a program detects that the receive data register is full ($RDRF = 1$), it gets the data from the receive data register by reading SCIxD. The RDRF flag is cleared by reading SCIxS1 while $RDRF = 1$ and then reading SCIxD. If the SCI is configured to operate in 9-bit mode, an additional read to the SCIxC3 register is required to clear RDRF.

When polling is used, this sequence is naturally satisfied in the normal course of the user program. If hardware interrupts are used, SCIxS1 must be read in the interrupt service routine (ISR). Normally, this is done in the ISR anyway to check for receive errors, so the sequence is automatically satisfied.

The IDLE status flag includes logic that prevents it from getting set repeatedly when the RxD line remains idle for an extended period of time. IDLE is cleared by reading SCIxS1 while $IDLE = 1$ and then reading SCIxD. After IDLE has been cleared, it cannot become set again until the receiver has received at least one new character and has set RDRF.

If the associated error was detected in the received character that caused RDRF to be set, the error flags — noise flag (NF), framing error (FE), and parity error flag (PF) — get set at the same time as RDRF. These flags are not set in overrun cases.

If RDRF was already set when a new character is ready to be transferred from the receive shifter to the receive data buffer, the overrun (OR) flag gets set instead and the data and any associated NF, FE, or PF condition is lost.

11.3.5 Additional SCI Functions

The following sections describe additional SCI functions.

11.3.5.1 8- and 9-Bit Data Modes

The SCI system (transmitter and receiver) can be configured to operate in 9-bit data mode by setting the M control bit in SCIxC1. In 9-bit mode, there is a ninth data bit to the left of the MSB of the SCI data register. For the transmit data buffer, this bit is stored in T8 in SCIxC3. For the receiver, the ninth bit is held in R8 in SCIxC3.

When transmitting 9-bit data, write to the T8 bit before writing to SCIxD for coherent writes to the transmit data buffer. If the bit value to be transmitted as the ninth bit of a new character is the same as for the previous character, it is not necessary to write to T8 again. When data is transferred from the transmit data buffer to the transmit shifter, the value in T8 is copied at the same time data is transferred from SCIxD to the shifter.

When receiving 9-bit data, clear the RDRF bit by reading both R8 and SCIxD. R8 and SCIxD can be read in either order.

12.4 Functional Description

An SPI transfer is initiated by checking for the SPI transmit buffer empty flag (SPTEF = 1) and then writing a byte of data to the SPI data register (SPI1D) in the master SPI device. When the SPI shift register is available, this byte of data is moved from the transmit data buffer to the shifter, SPTEF is set to indicate there is room in the buffer to queue another transmit character if desired, and the SPI serial transfer starts.

During the SPI transfer, data is sampled (read) on the MISO pin at one SPSCCK edge and shifted, changing the bit value on the MOSI pin, one-half SPSCCK cycle later. After eight SPSCCK cycles, the data that was in the shift register of the master has been shifted out the MOSI pin to the slave while eight bits of data were shifted in the MISO pin into the master's shift register. At the end of this transfer, the received data byte is moved from the shifter into the receive data buffer and SPRF is set to indicate the data can be read by reading SPI1D. If another byte of data is waiting in the transmit buffer at the end of a transfer, it is moved into the shifter, SPTEF is set, and a new transfer is started.

Normally, SPI data is transferred most significant bit (MSB) first. If the least significant bit first enable (LSBFE) bit is set, SPI data is shifted LSB first.

When the SPI is configured as a slave, its $\overline{SS}$ pin must be driven low before a transfer starts and $\overline{SS}$ must stay low throughout the transfer. If a clock format where CPHA = 0 is selected, $\overline{SS}$ must be driven to a logic 1 between successive transfers. If CPHA = 1, $\overline{SS}$ may remain low between successive transfers. See Section 12.4.1, "SPI Clock Formats" for more details.

Because the transmitter and receiver are double buffered, a second byte, in addition to the byte currently being shifted out, can be queued into the transmit data buffer, and a previously received character can be in the receive data buffer while a new character is being shifted in. The SPTEF flag indicates when the transmit buffer has room for a new character. The SPRF flag indicates when a received character is available in the receive data buffer. The received character must be read out of the receive buffer (read SPI1D) before the next transfer is finished or a receive overrun error results.

In the case of a receive overrun, the new data is lost because the receive buffer still held the previous character and was not ready to accept the new data. There is no indication for such an overrun condition so the application system designer must ensure that previous data has been read from the receive buffer before a new transfer is initiated.

12.4.1 SPI Clock Formats

To accommodate a wide variety of synchronous serial peripherals from different manufacturers, the SPI system has a clock polarity (CPOL) bit and a clock phase (CPHA) control bit to select one of four clock formats for data transfers. CPOL selectively inserts an inverter in series with the clock. CPHA chooses between two different clock phase relationships between the clock and data.

Figure 12-10 shows the clock formats when CPHA = 1. At the top of the figure, the eight bit times are shown for reference with bit 1 starting at the first SPSCCK edge and bit 8 ending one-half SPSCCK cycle after the sixteenth SPSCCK edge. The MSB first and LSB first lines show the order of SPI data bits depending on the setting in LSBFE. Both variations of SPSCCK polarity are shown, but only one of these waveforms applies for a specific transfer, depending on the value in CPOL. The SAMPLE IN waveform applies to the MOSI input of a slave or the MISO input of a master. The MOSI waveform applies to the MOSI output

13.1.1 Features

The IIC includes these distinctive features:

- Compatible with IIC bus standard
- Multi-master operation
- Software programmable for one of 64 different serial clock frequencies
- Software selectable acknowledge bit
- Interrupt driven byte-by-byte data transfer
- Arbitration lost interrupt with automatic mode switching from master to slave
- Calling address identification interrupt
- START and STOP signal generation/detection
- Repeated START signal generation
- Acknowledge bit generation/detection
- Bus busy detection

13.1.2 Modes of Operation

The IIC functions the same in normal and monitor modes. A brief description of the IIC in the various MCU modes is given here.

- Run mode — This is the basic mode of operation. To conserve power in this mode, disable the module.
- Wait mode — The module will continue to operate while the MCU is in wait mode and can provide a wake-up interrupt.
- Stop mode — The IIC is inactive in stop3 mode for reduced power consumption. The STOP instruction does not affect IIC register states. Stop2 will reset the register contents.

converter yields the lower code (and vice-versa). However, even very small amounts of system noise can cause the converter to be indeterminate (between two codes) for a range of input voltages around the transition voltage. This range is normally around $1/2\text{LSB}$ and will increase with noise. This error may be reduced by repeatedly sampling the input and averaging the result. Additionally the techniques discussed in Section 14.7.2.3 will reduce this error.

Non-monotonicity is defined as when, except for code jitter, the converter converts to a lower code for a higher input voltage. Missing codes are those values which are never converted for any input value.

In 8-bit or 10-bit mode, the ADC is guaranteed to be monotonic and to have no missing codes.