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Understanding Embedded - FPGAs (Field Programmable Gate Array)

Embedded - FPGAs, or Field Programmable Gate Arrays, are advanced integrated circuits that offer unparalleled flexibility and performance for digital systems. Unlike traditional fixed-function logic devices, FPGAs can be programmed and reprogrammed to execute a wide array of logical operations, enabling customized functionality tailored to specific applications. This reprogrammability allows developers to iterate designs quickly and implement complex functions without the need for custom hardware.

Applications of Embedded - FPGAs

The versatility of Embedded - FPGAs makes them indispensable in numerous fields. In telecommunications,

Details	
Product Status	Obsolete
Number of LABs/CLBs	196
Number of Logic Elements/Cells	466
Total RAM Bits	6272
Number of I/O	112
Number of Gates	5000
Voltage - Supply	4.75V ~ 5.25V
Mounting Type	Through Hole
Operating Temperature	0°C ~ 85°C (TJ)
Package / Case	156-BCPGA
Supplier Device Package	156-CPGA (42.16x42.16)
Purchase URL	https://www.e-xfl.com/product-detail/xilinx/xc4005e-1pg156c

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Address: Room A, 16/F, Full Win Commercial Centre, 573 Nathan Road, Mongkok, Hong Kong

Supported CLB memory configurations and timing modes for single- and dual-port modes are shown in Table 3.

XC4000 Series devices are the first programmable logic devices with edge-triggered (synchronous) and dual-port RAM accessible to the user. Edge-triggered RAM simplifies system timing. Dual-port RAM doubles the effective throughput of FIFO applications. These features can be individually programmed in any XC4000 Series CLB.

Advantages of On-Chip and Edge-Triggered RAM

The on-chip RAM is extremely fast. The read access time is the same as the logic delay. The write access time is slightly slower. Both access times are much faster than any off-chip solution, because they avoid I/O delays.

Edge-triggered RAM, also called synchronous RAM, is a feature never before available in a Field Programmable Gate Array. The simplicity of designing with edge-triggered RAM, and the markedly higher achievable performance, add up to a significant improvement over existing devices with on-chip RAM.

Three application notes are available from Xilinx that discuss edge-triggered RAM: "*XC4000E Edge-Triggered and Dual-Port RAM Capability*," "*Implementing FIFOs in XC4000E RAM*," and "*Synchronous and Asynchronous FIFO Designs*." All three application notes apply to both XC4000E and XC4000X RAM.

Table 3: Supported RAM Modes

	16 x 1	16 x 2	32 x 1	Edge-Triggered Timing	Level-Sensitive Timing
Single-Port	√	√	√	√	√
Dual-Port	√			√	

RAM Configuration Options

The function generators in any CLB can be configured as RAM arrays in the following sizes:

- Two 16x1 RAMs: two data inputs and two data outputs with identical or, if preferred, different addressing for each RAM
- One 32x1 RAM: one data input and one data output.

One F or G function generator can be configured as a 16x1 RAM while the other function generators are used to implement any function of up to 5 inputs.

Additionally, the XC4000 Series RAM may have either of two timing modes:

- Edge-Triggered (Synchronous): data written by the designated edge of the CLB clock. WE acts as a true clock enable.
- Level-Sensitive (Asynchronous): an external WE signal acts as the write strobe.

The selected timing mode applies to both function generators within a CLB when both are configured as RAM.

The number of read ports is also programmable:

- Single Port: each function generator has a common read and write port
- Dual Port: both function generators are configured together as a single 16x1 dual-port RAM with one write port and two read ports. Simultaneous read and write operations to the same or different addresses are supported.

RAM configuration options are selected by placing the appropriate library symbol.

Choosing a RAM Configuration Mode

The appropriate choice of RAM mode for a given design should be based on timing and resource requirements, desired functionality, and the simplicity of the design process. Recommended usage is shown in Table 4.

The difference between level-sensitive, edge-triggered, and dual-port RAM is only in the write operation. Read operation and timing is identical for all modes of operation.

Table 4: RAM Mode Selection

	Level-Sensitive	Edge-Triggered	Dual-Port Edge-Triggered
Use for New Designs?	No	Yes	Yes
Size (16x1, Registered)	1/2 CLB	1/2 CLB	1 CLB
Simultaneous Read/Write	No	No	Yes
Relative Performance	X	2X	2X (4X effective)

RAM Inputs and Outputs

The F1-F4 and G1-G4 inputs to the function generators act as address lines, selecting a particular memory cell in each look-up table.

The functionality of the CLB control signals changes when the function generators are configured as RAM. The DIN/H2, H1, and SR/H0 lines become the two data inputs (D0, D1) and the Write Enable (WE) input for the 16x2 memory. When the 32x1 configuration is selected, D1 acts as the fifth address bit and D0 is the data input.

The contents of the memory cell(s) being addressed are available at the F' and G' function-generator outputs. They can exit the CLB through its X and Y outputs, or can be captured in the CLB flip-flop(s).

Configuring the CLB function generators as Read/Write memory does not affect the functionality of the other por-

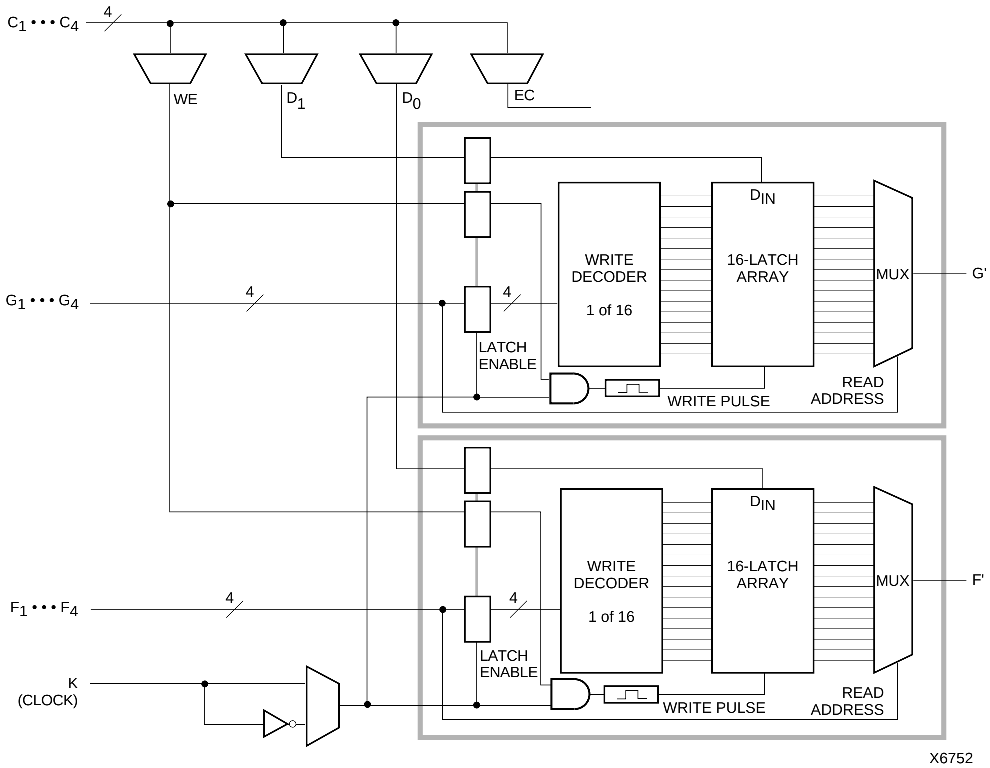


Figure 4: 16x2 (or 16x1) Edge-Triggered Single-Port RAM

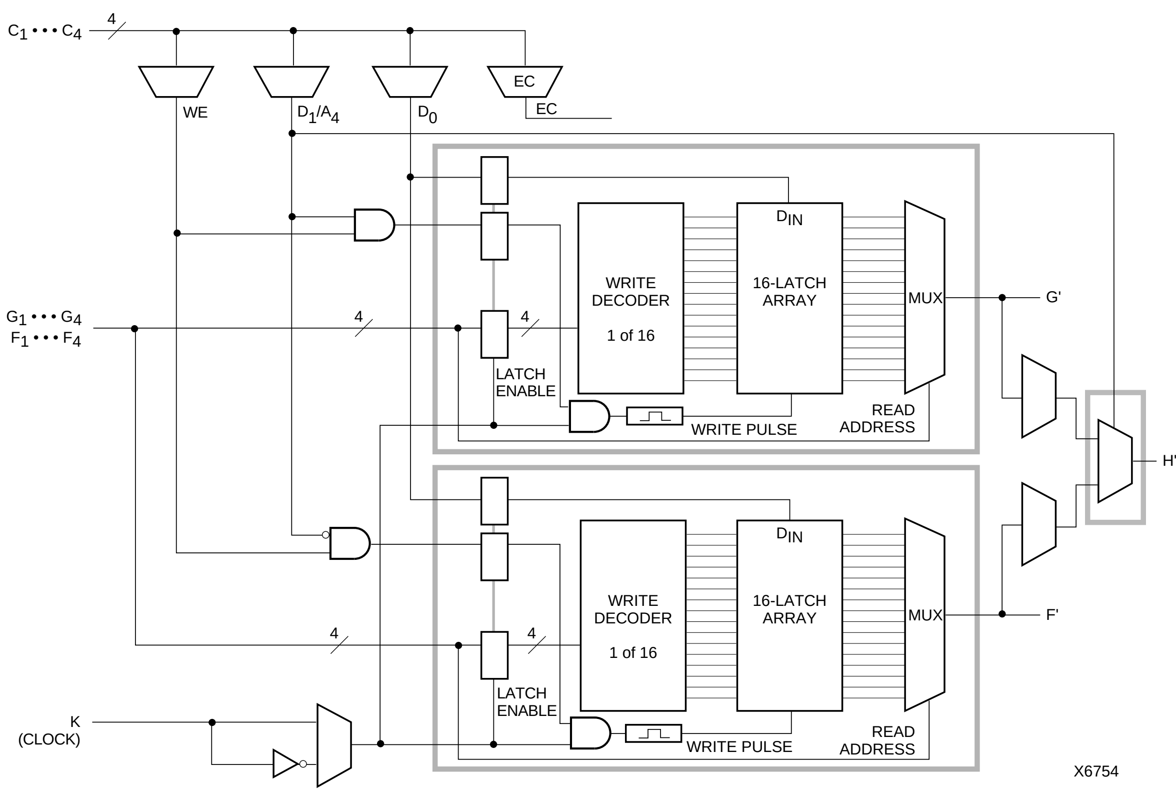


Figure 5: 32x1 Edge-Triggered Single-Port RAM (F and G addresses are identical)

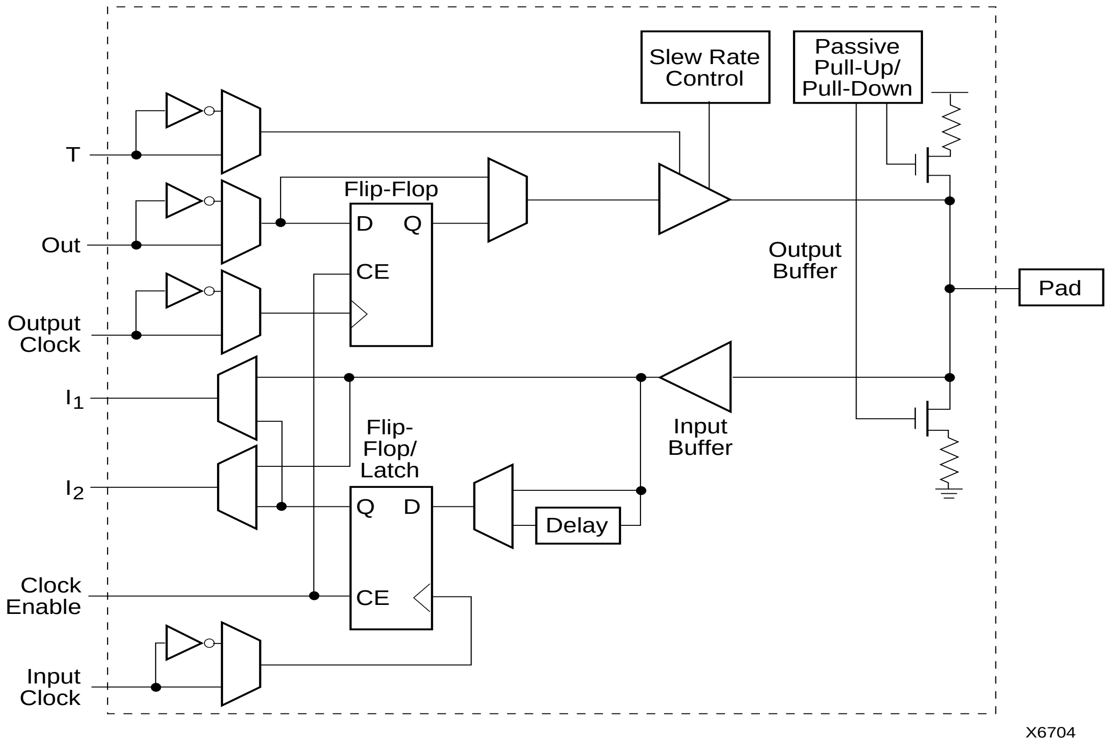


Figure 15: Simplified Block Diagram of XC4000E IOB

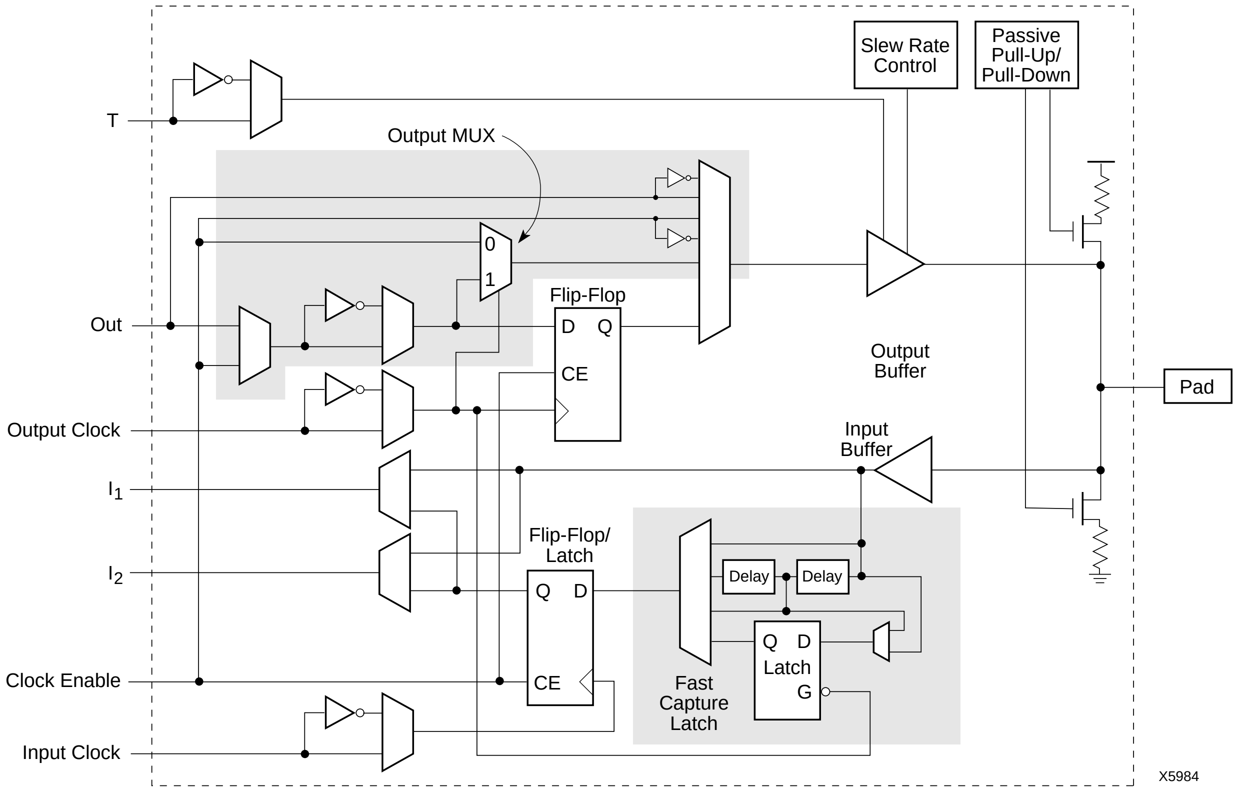


Figure 16: Simplified Block Diagram of XC4000X IOB (shaded areas indicate differences from XC4000E)

Table 8: Supported Sources for XC4000 Series Device Inputs

Source	XC4000E/EX Series Inputs		XC4000XL Series Inputs
	5 V, TTL	5 V, CMOS	3.3 V CMOS
Any device, Vcc = 3.3 V, CMOS outputs	√	Unreli-able Data	√
XC4000 Series, Vcc = 5 V, TTL outputs	√		√
Any device, Vcc = 5 V, TTL outputs (Voh ≤ 3.7 V)	√		√
Any device, Vcc = 5 V, CMOS outputs	√	√	√

XC4000XL 5-Volt Tolerant I/Os

The I/Os on the XC4000XL are fully 5-volt tolerant even though the V_{CC} is 3.3 volts. This allows 5 V signals to directly connect to the XC4000XL inputs without damage, as shown in Table 8. In addition, the 3.3 volt V_{CC} can be applied before or after 5 volt signals are applied to the I/Os. This makes the XC4000XL immune to power supply sequencing problems.

Registered Inputs

The I1 and I2 signals that exit the block can each carry either the direct or registered input signal.

The input and output storage elements in each IOB have a common clock enable input, which, through configuration, can be activated individually for the input or output flip-flop, or both. This clock enable operates exactly like the EC pin on the XC4000 Series CLB. It cannot be inverted within the IOB.

The storage element behavior is shown in Table 9.

Table 9: Input Register Functionality (active rising edge is shown)

Mode	Clock	Clock Enable	D	Q
Power-Up or GSR	X	X	X	SR
Flip-Flop	__/‾	1*	D	D
	0	X	X	Q
Latch	1	1*	X	Q
	0	1*	D	D
Both	X	0	X	Q

Legend:
X — Don't care
__/‾ — Rising edge
SR — Set or Reset value. Reset is default.
0* — Input is Low or unconnected (default value)
1* — Input is High or unconnected (default value)

Optional Delay Guarantees Zero Hold Time

The data input to the register can optionally be delayed by several nanoseconds. With the delay enabled, the setup time of the input flip-flop is increased so that normal clock routing does not result in a positive hold-time requirement. A positive hold time requirement can lead to unreliable, temperature- or processing-dependent operation.

The input flip-flop setup time is defined between the data measured at the device I/O pin and the clock input at the IOB (not at the clock pin). Any routing delay from the device clock pin to the clock input of the IOB must, therefore, be subtracted from this setup time to arrive at the real setup time requirement relative to the device pins. A short specified setup time might, therefore, result in a negative setup time at the device pins, i.e., a positive hold-time requirement.

When a delay is inserted on the data line, more clock delay can be tolerated without causing a positive hold-time requirement. Sufficient delay eliminates the possibility of a data hold-time requirement at the external pin. The maximum delay is therefore inserted as the default.

The XC4000E IOB has a one-tap delay element: either the delay is inserted (default), or it is not. The delay guarantees a zero hold time with respect to clocks routed through any of the XC4000E global clock buffers. (See "Global Nets and Buffers (XC4000E only)" on page 35 for a description of the global clock buffers in the XC4000E.) For a shorter input register setup time, with non-zero hold, attach a NODELAY attribute or property to the flip-flop.

The XC4000X IOB has a two-tap delay element, with choices of a full delay, a partial delay, or no delay. The attributes or properties used to select the desired delay are shown in Table 10. The choices are no added attribute, MEDDELAY, and NODELAY. The default setting, with no added attribute, ensures no hold time with respect to any of the XC4000X clock buffers, including the Global Low-Skew buffers. MEDDELAY ensures no hold time with respect to the Global Early buffers. Inputs with NODELAY may have a positive hold time with respect to all clock buffers. For a description of each of these buffers, see "Global Nets and Buffers (XC4000X only)" on page 37.

Table 10: XC4000X IOB Input Delay Element

Value	When to Use
full delay (default, no attribute added)	Zero Hold with respect to Global Low-Skew Buffer, Global Early Buffer
MEDDELAY	Zero Hold with respect to Global Early Buffer
NODELAY	Short Setup, positive Hold time

Output Multiplexer/2-Input Function Generator (XC4000X only)

As shown in Figure 16 on page 21, the output path in the XC4000X IOB contains an additional multiplexer not available in the XC4000E IOB. The multiplexer can also be configured as a 2-input function generator, implementing a pass-gate, AND-gate, OR-gate, or XOR-gate, with 0, 1, or 2 inverted inputs. The logic used to implement these functions is shown in the upper gray area of Figure 16.

When configured as a multiplexer, this feature allows two output signals to time-share the same output pad; effectively doubling the number of device outputs without requiring a larger, more expensive package.

When the MUX is configured as a 2-input function generator, logic can be implemented within the IOB itself. Combined with a Global Early buffer, this arrangement allows very high-speed gating of a single signal. For example, a wide decoder can be implemented in CLBs, and its output gated with a Read or Write Strobe Driven by a BUFGE buffer, as shown in Figure 19. The critical-path pin-to-pin delay of this circuit is less than 6 nanoseconds.

As shown in Figure 16, the IOB input pins Out, Output Clock, and Clock Enable have different delays and different flexibilities regarding polarity. Additionally, Output Clock sources are more limited than the other inputs. Therefore, the Xilinx software does not move logic into the IOB function generators unless explicitly directed to do so.

The user can specify that the IOB function generator be used, by placing special library symbols beginning with the letter "O." For example, a 2-input AND-gate in the IOB function generator is called OAND2. Use the symbol input pin labelled "F" for the signal on the critical path. This signal is placed on the OK pin — the IOB input with the shortest delay to the function generator. Two examples are shown in Figure 20.

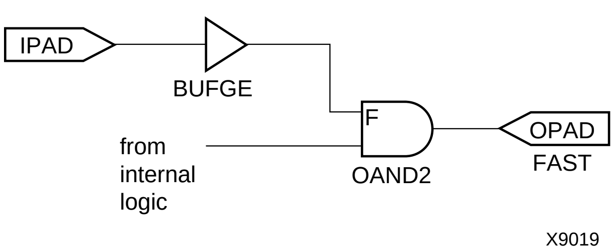


Figure 19: Fast Pin-to-Pin Path in XC4000X

Figure 20: AND & MUX Symbols in XC4000X IOB

Other IOB Options

There are a number of other programmable options in the XC4000 Series IOB.

Pull-up and Pull-down Resistors

Programmable pull-up and pull-down resistors are useful for tying unused pins to Vcc or Ground to minimize power consumption and reduce noise sensitivity. The configurable pull-up resistor is a p-channel transistor that pulls to Vcc. The configurable pull-down resistor is an n-channel transistor that pulls to Ground.

The value of these resistors is 50 kΩ – 100 kΩ. This high value makes them unsuitable as wired-AND pull-up resistors.

The pull-up resistors for most user-programmable IOBs are active during the configuration process. See Table 22 on page 58 for a list of pins with pull-ups active before and during configuration.

After configuration, voltage levels of unused pads, bonded or un-bonded, must be valid logic levels, to reduce noise sensitivity and avoid excess current. Therefore, by default, unused pads are configured with the internal pull-up resistor active. Alternatively, they can be individually configured with the pull-down resistor, or as a driven output, or to be driven by an external source. To activate the internal pull-up, attach the PULLUP library component to the net attached to the pad. To activate the internal pull-down, attach the PULLDOWN library component to the net attached to the pad.

Independent Clocks

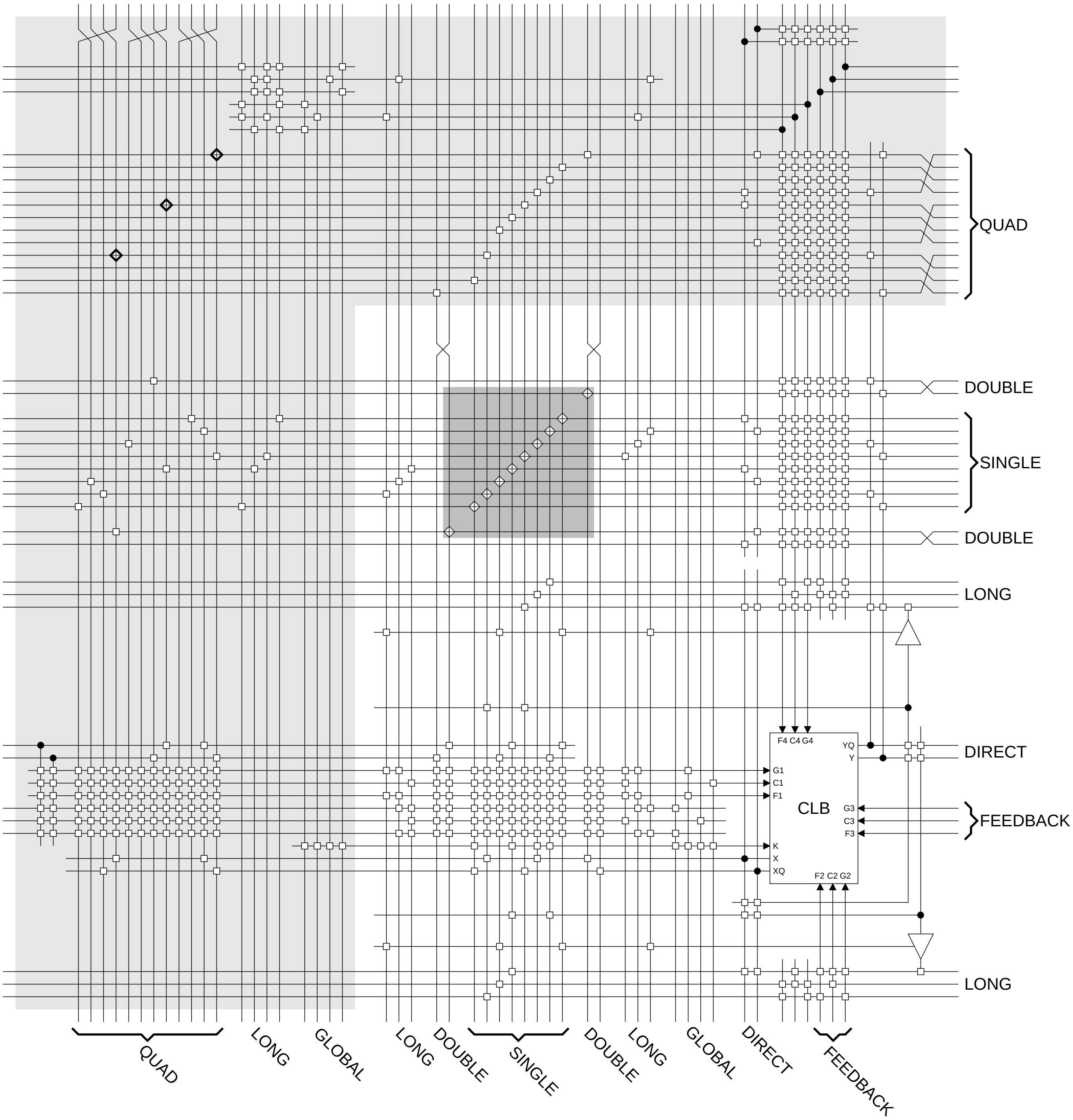
Separate clock signals are provided for the input and output flip-flops. The clock can be independently inverted for each flip-flop within the IOB, generating either falling-edge or rising-edge triggered flip-flops. The clock inputs for each IOB are independent, except that in the XC4000X, the Fast Capture latch shares an IOB input with the output clock pin.

Early Clock for IOBs (XC4000X only)

Special early clocks are available for IOBs. These clocks are sourced by the same sources as the Global Low-Skew buffers, but are separately buffered. They have fewer loads and therefore less delay. The early clock can drive either the IOB output clock or the IOB input clock, or both. The early clock allows fast capture of input data, and fast clock-to-output on output data. The Global Early buffers that drive these clocks are described in "Global Nets and Buffers (XC4000X only)" on page 37.

Global Set/Reset

As with the CLB registers, the Global Set/Reset signal (GSR) can be used to set or clear the input and output registers, depending on the value of the INIT attribute or property. The two flip-flops can be individually configured to set



Common to XC4000E and XC4000X

XC4000X only

Programmable Switch Matrix

Figure 27: Detail of Programmable Interconnect Associated with XC4000 Series CLB

circuit prevents undefined floating levels. However, it is overridden by any driver, even a pull-up resistor.

Each XC4000E longline has a programmable splitter switch at its center, as does each XC4000X longline driven by TBUFs. This switch can separate the line into two independent routing channels, each running half the width or height of the array.

Each XC4000X longline not driven by TBUFs has a buffered programmable splitter switch at the 1/4, 1/2, and 3/4 points of the array. Due to the buffering, XC4000X longline performance does not deteriorate with the larger array sizes. If the longline is split, the resulting partial longlines are independent.

Routing connectivity of the longlines is shown in Figure 27 on page 30.

Direct Interconnect (XC4000X only)

The XC4000X offers two direct, efficient and fast connections between adjacent CLBs. These nets facilitate a data flow from the left to the right side of the device, or from the top to the bottom, as shown in Figure 30. Signals routed on the direct interconnect exhibit minimum interconnect propagation delay and use no general routing resources.

The direct interconnect is also present between CLBs and adjacent IOBs. Each IOB on the left and top device edges has a direct path to the nearest CLB. Each CLB on the right and bottom edges of the array has a direct path to the nearest two IOBs, since there are two IOBs for each row or column of CLBs.

The place and route software uses direct interconnect whenever possible, to maximize routing resources and minimize interconnect delays.

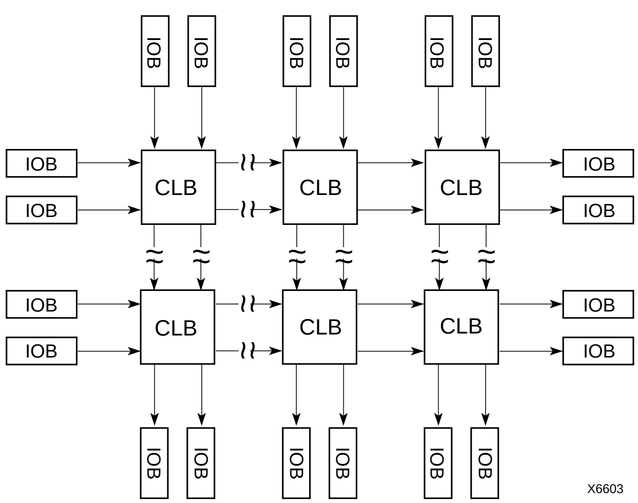


Figure 30: XC4000X Direct Interconnect

I/O Routing

XC4000 Series devices have additional routing around the IOB ring. This routing is called a VersaRing. The VersaRing facilitates pin-swapping and redesign without affecting board layout. Included are eight double-length lines spanning two CLBs (four IOBs), and four longlines. Global lines and Wide Edge Decoder lines are provided. XC4000X devices also include eight octal lines.

A high-level diagram of the VersaRing is shown in Figure 31. The shaded arrows represent routing present only in XC4000X devices.

Figure 33 on page 34 is a detailed diagram of the XC4000E and XC4000X VersaRing. The area shown includes two IOBs. There are two IOBs per CLB row or column, therefore this diagram corresponds to the CLB routing diagram shown in Figure 27 on page 30. The shaded areas represent routing and routing connections present only in XC4000X devices.

Octal I/O Routing (XC4000X only)

Between the XC4000X CLB array and the pad ring, eight interconnect tracks provide for versatility in pin assignment and fixed pinout flexibility. (See Figure 32 on page 33.)

These routing tracks are called octals, because they can be broken every eight CLBs (sixteen IOBs) by a programmable buffer that also functions as a splitter switch. The buffers are staggered, so each line goes through a buffer at every eighth CLB location around the device edge.

The octal lines bend around the corners of the device. The lines cross at the corners in such a way that the segment most recently buffered before the turn has the farthest distance to travel before the next buffer, as shown in Figure 32.

Global Nets and Buffers (XC4000X only)

Eight vertical longlines in each CLB column are driven by special global buffers. These longlines are in addition to the vertical longlines used for standard interconnect. The global lines are broken in the center of the array, to allow faster distribution and to minimize skew across the whole array. Each half-column global line has its own buffered multiplexer, as shown in Figure 35. The top and bottom global lines cannot be connected across the center of the device, as this connection might introduce unacceptable skew. The top and bottom halves of the global lines must be separately driven — although they can be driven by the same global buffer.

The eight global lines in each CLB column can be driven by either of two types of global buffers. They can also be driven by internal logic, because they can be accessed by single, double, and quad lines at the top, bottom, half, and quarter points. Consequently, the number of different clocks that can be used simultaneously in an XC4000X device is very large.

There are four global lines feeding the IOBs at the left edge of the device. IOBs along the right edge have eight global lines. There is a single global line along the top and bottom edges with access to the IOBs. All IOB global lines are broken at the center. They cannot be connected across the center of the device, as this connection might introduce unacceptable skew.

IOB global lines can be driven from two types of global buffers, or from local interconnect. Alternatively, top and bottom IOBs can be clocked from the global lines in the adjacent CLB column.

Two different types of clock buffers are available in the XC4000X:

- Global Low-Skew Buffers (BUFGLS)
- Global Early Buffers (BUFGE)

Global Low-Skew Buffers are the standard clock buffers. They should be used for most internal clocking, whenever a large portion of the device must be driven.

Global Early Buffers are designed to provide a faster clock access, but CLB access is limited to one-fourth of the device. They also facilitate a faster I/O interface.

Figure 35 is a conceptual diagram of the global net structure in the XC4000X.

Global Early buffers and Global Low-Skew buffers share a single pad. Therefore, the same IPAD symbol can drive one buffer of each type, in parallel. This configuration is particularly useful when using the Fast Capture latches, as described in "IOB Input Signals" on page 20. Paired Global Early and Global Low-Skew buffers share a common input; they cannot be driven by two different signals.

Choosing an XC4000X Clock Buffer

The clocking structure of the XC4000X provides a large variety of features. However, it can be simple to use, without understanding all the details. The software automatically handles clocks, along with all other routing, when the appropriate clock buffer is placed in the design. In fact, if a buffer symbol called BUFG is placed, rather than a specific type of buffer, the software even chooses the buffer most appropriate for the design. The detailed information in this section is provided for those users who want a finer level of control over their designs.

If fine control is desired, use the following summary and Table 15 on page 35 to choose an appropriate clock buffer.

- The simplest thing to do is to use a Global Low-Skew buffer.
- If a faster clock path is needed, try a BUFG. The software will first try to use a Global Low-Skew Buffer. If timing requirements are not met, a faster buffer will automatically be used.
- If a single quadrant of the chip is sufficient for the clocked logic, and the timing requires a faster clock than the Global Low-Skew buffer, use a Global Early buffer.

Global Low-Skew Buffers

Each corner of the XC4000X device has two Global Low-Skew buffers. Any of the eight Global Low-Skew buffers can drive any of the eight vertical Global lines in a column of CLBs. In addition, any of the buffers can drive any of the four vertical lines accessing the IOBs on the left edge of the device, and any of the eight vertical lines accessing the IOBs on the right edge of the device. (See Figure 36 on page 38.)

IOBs at the top and bottom edges of the device are accessed through the vertical Global lines in the CLB array, as in the XC4000E. Any Global Low-Skew buffer can, therefore, access every IOB and CLB in the device.

The Global Low-Skew buffers can be driven by either semi-dedicated pads or internal logic.

To use a Global Low-Skew buffer, instantiate a BUFGLS element in a schematic or in HDL code. If desired, attach a LOC attribute or property to direct placement to the designated location. For example, attach a LOC=T attribute or property to direct that a BUFGLS be placed in one of the two Global Low-Skew buffers on the top edge of the device, or a LOC=TR to indicate the Global Low-Skew buffer on the top edge of the device, on the right.

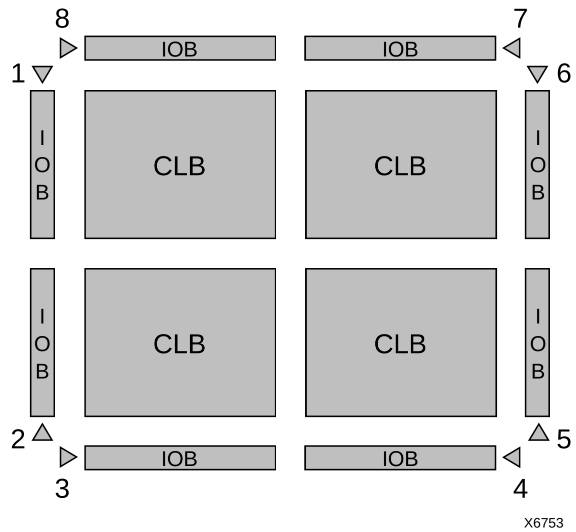


Figure 36: Any BUFGLS (GCK1 - GCK8) Can Drive Any or All Clock Inputs on the Device

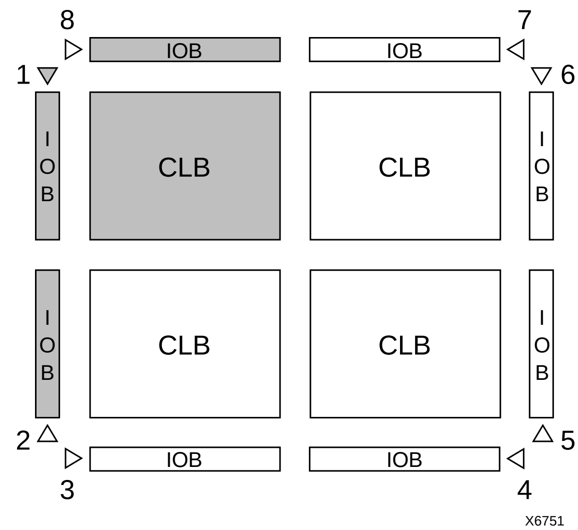


Figure 37: Left and Right BUFGEs Can Drive Any or All Clock Inputs in Same Quadrant or Edge (GCK1 is shown. GCK2, GCK5 and GCK6 are similar.)

Global Early Buffers

Each corner of the XC4000X device has two Global Early buffers. The primary purpose of the Global Early buffers is to provide an earlier clock access than the potentially heavily-loaded Global Low-Skew buffers. A clock source applied to both buffers will result in the Global Early clock edge occurring several nanoseconds earlier than the Global Low-Skew buffer clock edge, due to the lighter loading.

Global Early buffers also facilitate the fast capture of device inputs, using the Fast Capture latches described in "IOB Input Signals" on page 20. For Fast Capture, take a single clock signal, and route it through both a Global Early buffer and a Global Low-Skew buffer. (The two buffers share an input pad.) Use the Global Early buffer to clock the Fast Capture latch, and the Global Low-Skew buffer to clock the normal input flip-flop or latch, as shown in Figure 17 on page 23.

The Global Early buffers can also be used to provide a fast Clock-to-Out on device output pins. However, an early clock in the output flip-flop IOB must be taken into consideration when calculating the internal clock speed for the design.

The Global Early buffers at the left and right edges of the chip have slightly different capabilities than the ones at the top and bottom. Refer to Figure 37, Figure 38, and Figure 35 on page 36 while reading the following explanation.

Each Global Early buffer can access the eight vertical Global lines for all CLBs in the quadrant. Therefore, only one-fourth of the CLB clock pins can be accessed. This restriction is in large part responsible for the faster speed of the buffers, relative to the Global Low-Skew buffers.

The left-side Global Early buffers can each drive two of the four vertical lines accessing the IOBs on the entire left edge of the device. The right-side Global Early buffers can each drive two of the eight vertical lines accessing the IOBs on the entire right edge of the device. (See Figure 37.)

Each left and right Global Early buffer can also drive half of the IOBs along either the top or bottom edge of the device, using a dedicated line that can only be accessed through the Global Early buffers.

The top and bottom Global Early buffers can drive half of the IOBs along either the left or right edge of the device, as shown in Figure 38. They can only access the top and bottom IOBs via the CLB global lines.

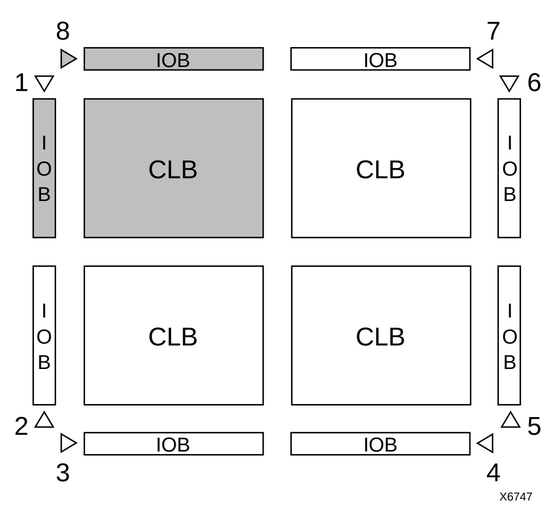


Figure 38: Top and Bottom BUFGEs Can Drive Any or All Clock Inputs in Same Quadrant (GCK8 is shown. GCK3, GCK4 and GCK7 are similar.)

Figure 41 on page 44 is a diagram of the XC4000 Series boundary scan logic. It includes three bits of Data Register per IOB, the IEEE 1149.1 Test Access Port controller, and the Instruction Register with decodes.

XC4000 Series devices can also be configured through the boundary scan logic. See "Readback" on page 55.

Data Registers

The primary data register is the boundary scan register. For each IOB pin in the FPGA, bonded or not, it includes three bits for In, Out and 3-State Control. Non-IOB pins have appropriate partial bit population for In or Out only. PROGRAM, CCLK and DONE are not included in the boundary scan register. Each EXTEST CAPTURE-DR state captures all In, Out, and 3-state pins.

The data register also includes the following non-pin bits: TDO.T, and TDO.O, which are always bits 0 and 1 of the data register, respectively, and BSCANT.UPD, which is always the last bit of the data register. These three boundary scan bits are special-purpose Xilinx test signals.

The other standard data register is the single flip-flop BYPASS register. It synchronizes data being passed through the FPGA to the next downstream boundary scan device.

The FPGA provides two additional data registers that can be specified using the BSCAN macro. The FPGA provides two user pins (BSCAN.SEL1 and BSCAN.SEL2) which are the decodes of two user instructions. For these instructions, two corresponding pins (BSCAN.TDO1 and BSCAN.TDO2) allow user scan data to be shifted out on TDO. The data register clock (BSCAN.DRCK) is available for control of test logic which the user may wish to implement with CLBs. The NAND of TCK and RUN-TEST-IDLE is also provided (BSCAN.IDLE).

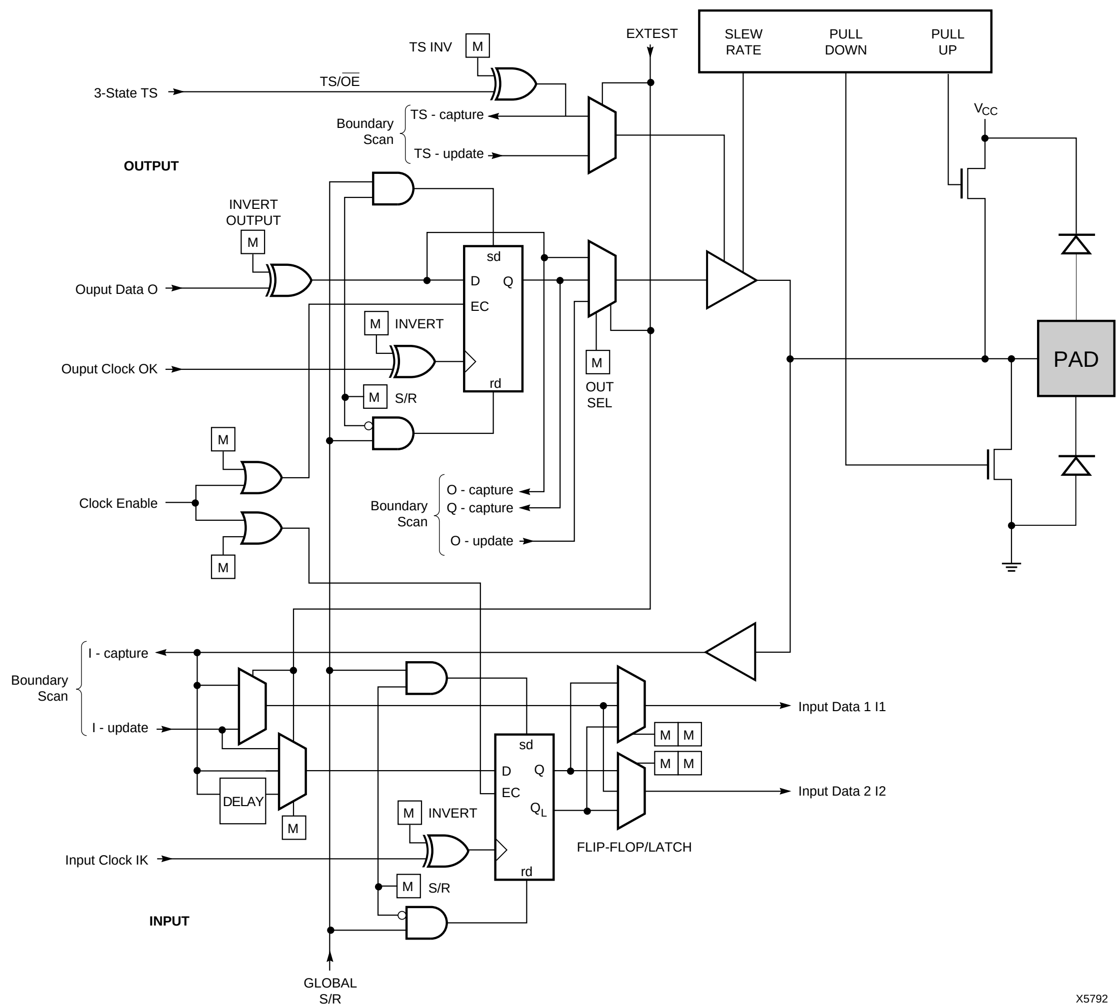


Figure 40: Block Diagram of XC4000E IOB with Boundary Scan (some details not shown). XC4000X Boundary Scan Logic is Identical.

Configuration Modes

XC4000E devices have six configuration modes. XC4000X devices have the same six modes, plus an additional configuration mode. These modes are selected by a 3-bit input code applied to the M2, M1, and M0 inputs. There are three self-loading Master modes, two Peripheral modes, and a Serial Slave mode, which is used primarily for daisy-chained devices. The coding for mode selection is shown in Table 18.

Table 18: Configuration Modes

Mode	M2	M1	M0	CCLK	Data
Master Serial	0	0	0	output	Bit-Serial
Slave Serial	1	1	1	input	Bit-Serial
Master Parallel Up	1	0	0	output	Byte-Wide, increment from 00000
Master Parallel Down	1	1	0	output	Byte-Wide, decrement from 3FFFF
Peripheral Synchronous*	0	1	1	input	Byte-Wide
Peripheral Asynchronous	1	0	1	output	Byte-Wide
Reserved	0	1	0	—	—
Reserved	0	0	1	—	—

* Can be considered byte-wide Slave Parallel

A detailed description of each configuration mode, with timing information, is included later in this data sheet. During configuration, some of the I/O pins are used temporarily for the configuration process. All pins used during configuration are shown in Table 22 on page 58.

Master Modes

The three Master modes use an internal oscillator to generate a Configuration Clock (CCLK) for driving potential slave devices. They also generate address and timing for external PROM(s) containing the configuration data.

Master Parallel (Up or Down) modes generate the CCLK signal and PROM addresses and receive byte parallel data. The data is internally serialized into the FPGA data-frame format. The up and down selection generates starting addresses at either zero or 3FFFF (3FFFFF when 22 address lines are used), for compatibility with different microprocessor addressing conventions. The Master Serial mode generates CCLK and receives the configuration data in serial form from a Xilinx serial-configuration PROM.

CCLK speed is selectable as either 1 MHz (default) or 8 MHz. Configuration always starts at the default slow frequency, then can switch to the higher frequency during the first frame. Frequency tolerance is -50% to +25%.

Additional Address lines in XC4000 devices

The XC4000X devices have additional address lines (A18-A21) allowing the additional address space required to daisy-chain several large devices.

The extra address lines are programmable in XC4000EX devices. By default these address lines are not activated. In the default mode, the devices are compatible with existing XC4000 and XC4000E products. If desired, the extra address lines can be used by specifying the address lines option in bitgen as 22 (bitgen -g AddressLines:22). The lines (A18-A21) are driven when a master device detects, via the bitstream, that it should be using all 22 address lines. Because these pins will initially be pulled high by internal pull-ups, designers using Master Parallel Up mode should use external pull down resistors on pins A18-A21. If Master Parallel Down mode is used external resistors are not necessary.

All 22 address lines are always active in Master Parallel modes with XC4000XL devices. The additional address lines behave identically to the lower order address lines. If the Address Lines option in bitgen is set to 18, it will be ignored by the XC4000XL device.

The additional address lines (A18-A21) are not available in the PC84 package.

Peripheral Modes

The two Peripheral modes accept byte-wide data from a bus. A RDY/$\overline{\text{BUSY}}$ status is available as a handshake signal. In Asynchronous Peripheral mode, the internal oscillator generates a CCLK burst signal that serializes the byte-wide data. CCLK can also drive slave devices. In the synchronous mode, an externally supplied clock input to CCLK serializes the data.

Slave Serial Mode

In Slave Serial mode, the FPGA receives serial configuration data on the rising edge of CCLK and, after loading its configuration, passes additional data out, resynchronized on the next falling edge of CCLK.

Multiple slave devices with identical configurations can be wired with parallel DIN inputs. In this way, multiple devices can be configured simultaneously.

Serial Daisy Chain

Multiple devices with different configurations can be connected together in a “daisy chain,” and a single combined bitstream used to configure the chain of slave devices.

To configure a daisy chain of devices, wire the CCLK pins of all devices in parallel, as shown in Figure 51 on page 60. Connect the DOUT of each device to the DIN of the next. The lead or master FPGA and following slaves each passes resynchronized configuration data coming from a single source. The header data, including the length count,

Table 20: XC4000E Program Data

Device	XC4003E	XC4005E	XC4006E	XC4008E	XC4010E	XC4013E	XC4020E	XC4025E
Max Logic Gates	3,000	5,000	6,000	8,000	10,000	13,000	20,000	25,000
CLBs (Row x Col.)	100 (10 x 10)	196 (14 x 14)	256 (16 x 16)	324 (18 x 18)	400 (20 x 20)	576 (24 x 24)	784 (28 x 28)	1,024 (32 x 32)
IOBs	80	112	128	144	160	192	224	256
Flip-Flops	360	616	768	936	1,120	1,536	2,016	2,560
Bits per Frame	126	166	186	206	226	266	306	346
Frames	428	572	644	716	788	932	1,076	1,220
Program Data	53,936	94,960	119,792	147,504	178,096	247,920	329,264	422,128
PROM Size (bits)	53,984	95,008	119,840	147,552	178,144	247,968	329,312	422,176

Notes: 1. Bits per Frame = (10 x number of rows) + 7 for the top + 13 for the bottom + 1 + 1 start bit + 4 error check bits
Number of Frames = (36 x number of columns) + 26 for the left edge + 41 for the right edge + 1
Program Data = (Bits per Frame x Number of Frames) + 8 postamble bits
PROM Size = Program Data + 40 (header) + 8
2. The user can add more “one” bits as leading dummy bits in the header, or, if CRC = off, as trailing dummy bits at the end of any frame, following the four error check bits. However, the Length Count value **must** be adjusted for all such extra “one” bits, even for extra leading ones at the beginning of the header.

Table 21: XC4000EX/XL Program Data

Device	XC4002XL	XC4005	XC4010	XC4013	XC4020	XC4028	XC4036	XC4044	XC4052	XC4062	XC4085
Max Logic Gates	2,000	5,000	10,000	13,000	20,000	28,000	36,000	44,000	52,000	62,000	85,000
CLBs (Row x Column)	64 (8 x 8)	196 (14 x 14)	400 (20 x 20)	576 (24 x 24)	784 (28 x 28)	1,024 (32 x 32)	1,296 (36 x 36)	1,600 (40 x 40)	1,936 (44 x 44)	2,304 (48 x 48)	3,136 (56 x 56)
IOBs	64	112	160	192	224	256	288	320	352	384	448
Flip-Flops	256	616	1,120	1,536	2,016	2,560	3,168	3,840	4,576	5,376	7,168
Bits per Frame	133	205	277	325	373	421	469	517	565	613	709
Frames	459	741	1,023	1,211	1,399	1,587	1,775	1,963	2,151	2,339	2,715
Program Data	61,052	151,910	283,376	393,580	521,832	668,124	832,480	1,014,876	1,215,320	1,433,804	1,924,940
PROM Size (bits)	61,104	151,960	283,424	393,632	521,880	668,172	832,528	1,014,924	1,215,368	1,433,852	1,924,992

Notes: 1. Bits per frame = (13 x number of rows) + 9 for the top + 17 for the bottom + 8 + 1 start bit + 4 error check bits.
Frames = (47 x number of columns) + 27 for the left edge + 52 for the right edge + 4.
Program data = (bits per frame x number of frames) + 5 postamble bits.
PROM size = (program data + 40 header bits + 8 start bits) rounded up to the nearest byte.
2. The user can add more “one” bits as leading dummy bits in the header, or, if CRC = off, as trailing dummy bits at the end of any frame, following the four error check bits. However, the Length Count value must be adjusted for all such extra “one” bits, even for extra leading “ones” at the beginning of the header.

Cyclic Redundancy Check (CRC) for Configuration and Readback

The Cyclic Redundancy Check is a method of error detection in data transmission applications. Generally, the transmitting system performs a calculation on the serial bitstream. The result of this calculation is tagged onto the data stream as additional check bits. The receiving system performs an identical calculation on the bitstream and compares the result with the received checksum.

Each data frame of the configuration bitstream has four error bits at the end, as shown in Table 19. If a frame data error is detected during the loading of the FPGA, the configuration process with a potentially corrupted bitstream is terminated. The FPGA pulls the $\overline{\text{INIT}}$ pin Low and goes into a Wait state.

During Readback, 11 bits of the 16-bit checksum are added to the end of the Readback data stream. The checksum is computed using the CRC-16 CCITT polynomial, as shown in Figure 45. The checksum consists of the 11 most significant bits of the 16-bit code. A change in the checksum indicates a change in the Readback bitstream. A comparison to a previous checksum is meaningful only if the readback data is independent of the current device state. CLB outputs should not be included (Read Capture option not

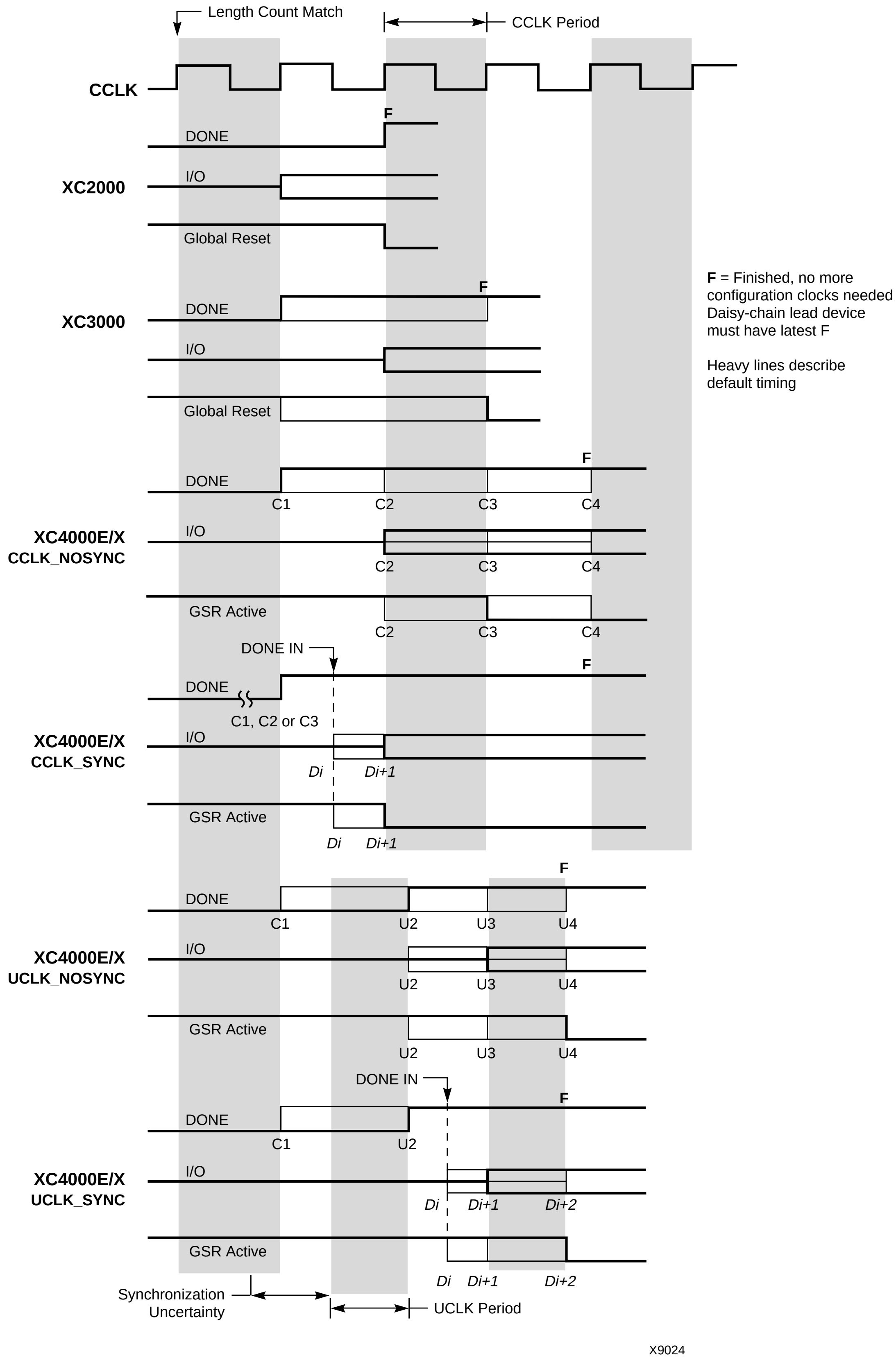


Figure 47: Start-up Timing

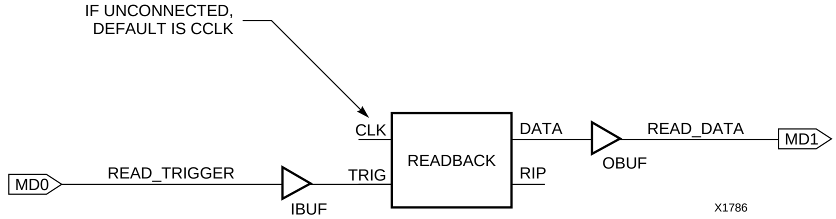


Figure 49: Readback Schematic Example

Readback Options

Readback options are: Read Capture, Read Abort, and Clock Select. They are set with the bitstream generation software.

Read Capture

When the Read Capture option is selected, the readback data stream includes sampled values of CLB and IOB signals. The rising edge of RDBK.TRIG latches the inverted values of the four CLB outputs, the IOB output flip-flops and the input signals I1 and I2. Note that while the bits describing configuration (interconnect, function generators, and RAM content) are *not* inverted, the CLB and IOB output signals *are* inverted.

When the Read Capture option is not selected, the values of the capture bits reflect the configuration data originally written to those memory locations.

If the RAM capability of the CLBs is used, RAM data are available in readback, since they directly overwrite the F and G function-table configuration of the CLB.

RDBK.TRIG is located in the lower-left corner of the device, as shown in Figure 50.

Read Abort

When the Read Abort option is selected, a High-to-Low transition on RDBK.TRIG terminates the readback operation and prepares the logic to accept another trigger.

After an aborted readback, additional clocks (up to one readback clock per configuration frame) may be required to re-initialize the control logic. The status of readback is indicated by the output control net RDBK.RIP. RDBK.RIP is High whenever a readback is in progress.

Clock Select

CCLK is the default clock. However, the user can insert another clock on RDBK.CLK. Readback control and data are clocked on rising edges of RDBK.CLK. If readback must be inhibited for security reasons, the readback control nets are simply not connected.

RDBK.CLK is located in the lower right chip corner, as shown in Figure 50.

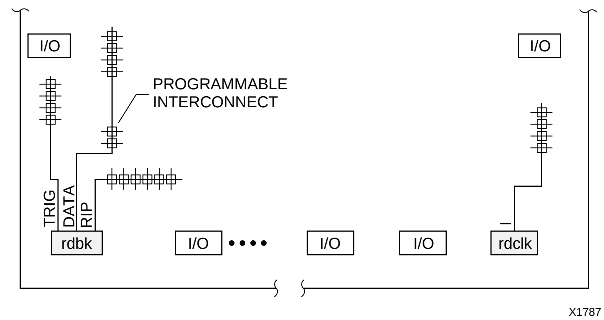


Figure 50: READBACK Symbol in Graphical Editor

Violating the Maximum High and Low Time Specification for the Readback Clock

The readback clock has a maximum High and Low time specification. In some cases, this specification cannot be met. For example, if a processor is controlling readback, an interrupt may force it to stop in the middle of a readback. This necessitates stopping the clock, and thus violating the specification.

The specification is mandatory only on clocking data at the end of a frame prior to the next start bit. The transfer mechanism will load the data to a shift register during the last six clock cycles of the frame, prior to the start bit of the following frame. This loading process is dynamic, and is the source of the maximum High and Low time requirements.

Therefore, the specification only applies to the six clock cycles prior to and including any start bit, including the clocks before the first start bit in the readback data stream. At other times, the frame data is already in the register and the register is not dynamic. Thus, it can be shifted out just like a regular shift register.

The user must precisely calculate the location of the readback data relative to the frame. The system must keep track of the position within a data frame, and disable interrupts before frame boundaries. Frame lengths and data formats are listed in Table 19, Table 20 and Table 21.

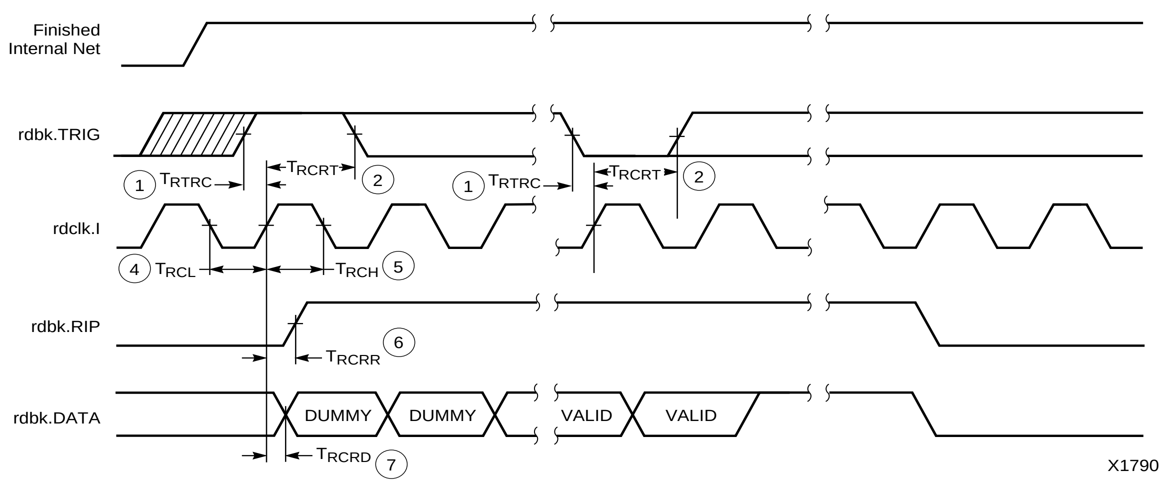
Readback with the XChecker Cable

The XChecker Universal Download/Readback Cable and Logic Probe uses the readback feature for bitstream verification. It can also display selected internal signals on the PC or workstation screen, functioning as a low-cost in-circuit emulator.

XC4000E/EX/XL Program Readback Switching Characteristic Guidelines

Testing of the switching parameters is modeled after testing methods specified by MIL-M-38510/605. All devices are 100% functionally tested. Internal timing parameters are not measured directly. They are derived from benchmark timing patterns that are taken at device introduction, prior to any process improvements.

The following guidelines reflect worst-case values over the recommended operating conditions.



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E/EX

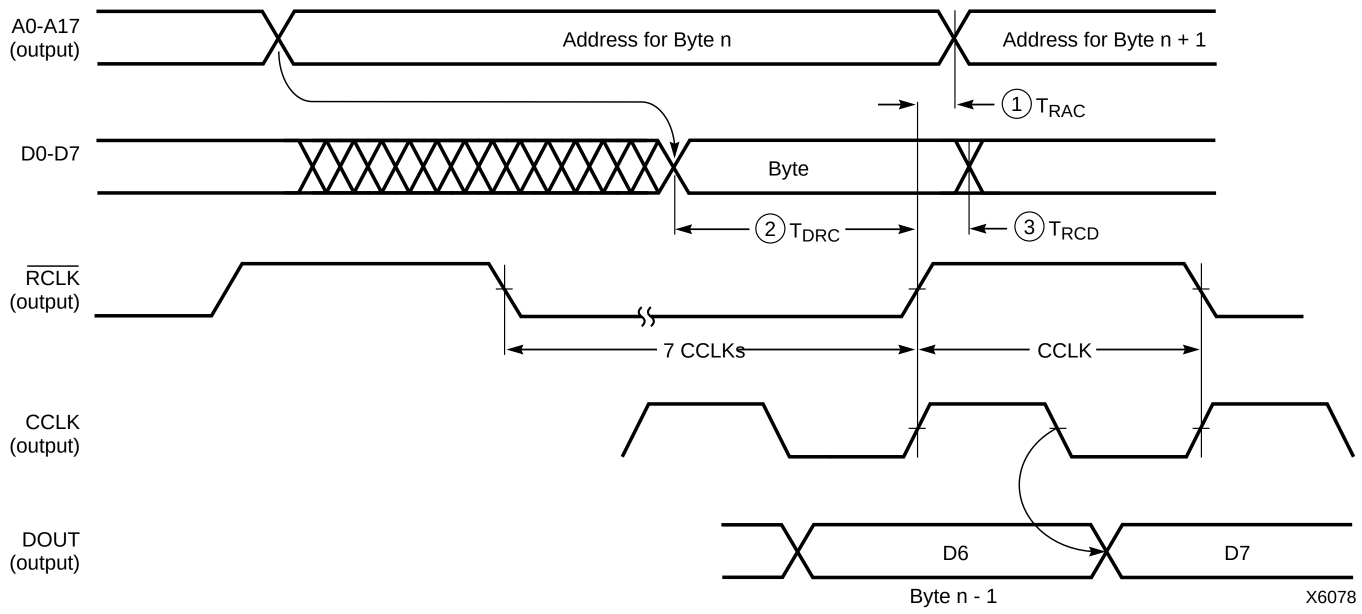
	Description	Symbol		Min	Max	Units
rdbk.TRIG	rdbk.TRIG setup to initiate and abort Readback	1	T_{RTRC}	200	-	ns
	rdbk.TRIG hold to initiate and abort Readback	2	T_{RCRT}	50	-	ns
rdclk.1	rdbk.DATA delay	7	T_{RCRD}	-	250	ns
	rdbk.RIP delay	6	T_{RCRR}	-	250	ns
	High time	5	T_{RCH}	250	500	ns
	Low time	4	T_{RCL}	250	500	ns

Note 1: Timing parameters apply to all speed grades.
Note 2: If rdbk.TRIG is High prior to Finished, Finished will trigger the first Readback.

XL

	Description	Symbol		Min	Max	Units
rdbk.TRIG	rdbk.TRIG setup to initiate and abort Readback	1	T_{RTRC}	200	-	ns
	rdbk.TRIG hold to initiate and abort Readback	2	T_{RCRT}	50	-	ns
rdclk.1	rdbk.DATA delay	7	T_{RCRD}	-	250	ns
	rdbk.RIP delay	6	T_{RCRR}	-	250	ns
	High time	5	T_{RCH}	250	500	ns
	Low time	4	T_{RCL}	250	500	ns

Note 1: Timing parameters apply to all speed grades.
Note 2: If rdbk.TRIG is High prior to Finished, Finished will trigger the first Readback.



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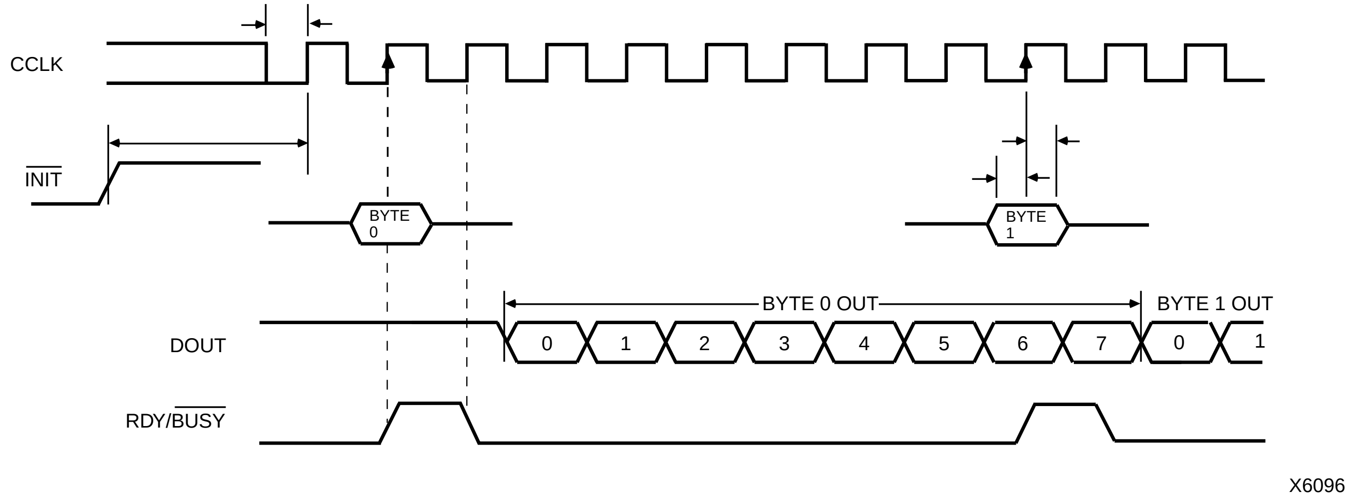
	Description	Symbol		Min	Max	Units
RCLK	Delay to Address valid	1	T_{RAC}	0	200	ns
	Data setup time	2	T_{DRC}	60		ns
	Data hold time	3	T_{RCD}	0		ns

Notes: 1. At power-up, Vcc must rise from 2.0 V to Vcc min in less than 25 ms, otherwise delay configuration by pulling $\overline{\text{PROGRAM}}$ Low until Vcc is valid.
2. The first Data byte is loaded and CCLK starts at the end of the first $\overline{\text{RCLK}}$ active cycle (rising edge).

This timing diagram shows that the EPROM requirements are extremely relaxed. EPROM access time can be longer than 500 ns. EPROM data output has no hold-time requirements.

Figure 55: Master Parallel Mode Programming Switching Characteristics



	Description	Symbol	Min	Max	Units
CCLK	INIT (High) setup time	T_{IC}	5		µs
	D0 - D7 setup time	T_{DC}	60		ns
	D0 - D7 hold time	T_{CD}	0		ns
	CCLK High time	T_{CCH}	50		ns
	CCLK Low time	T_{CCL}	60		ns
	CCLK Frequency	F_{CC}		8	MHz

Notes:
1. Peripheral Synchronous mode can be considered Slave Parallel mode. An external CCLK provides timing, clocking in the **first** data byte on the **second** rising edge of CCLK after INIT goes High. Subsequent data bytes are clocked in on every eighth consecutive rising edge of CCLK.
2. The RDY/BUSY line goes High for one CCLK period after data has been clocked in, although synchronous operation does not require such a response.
3. The pin name RDY/BUSY is a misnomer. In Synchronous Peripheral mode this is really an ACKNOWLEDGE signal.
4. Note that data starts to shift out serially on the DOUT pin 0.5 CCLK periods after it was loaded in parallel. Therefore, additional CCLK pulses are clearly required after the last byte has been loaded.

Figure 57: Synchronous Peripheral Mode Programming Switching Characteristics

Asynchronous Peripheral Mode

Write to FPGA

Asynchronous Peripheral mode uses the trailing edge of the logic AND condition of WS and CS0 being Low and RS and CS1 being High to accept byte-wide data from a microprocessor bus. In the lead FPGA, this data is loaded into a double-buffered UART-like parallel-to-serial converter and is serially shifted into the internal logic.

The lead FPGA presents the preamble data (and all data that overflows the lead device) on its DOUT pin. The RDY/BUSY output from the lead FPGA acts as a handshake signal to the microprocessor. RDY/BUSY goes Low when a byte has been received, and goes High again when the byte-wide input buffer has transferred its information into the shift register, and the buffer is ready to receive new data. A new write may be started immediately, as soon as the RDY/BUSY output has gone Low, acknowledging receipt of the previous data. Write may not be terminated until RDY/BUSY is High again for one CCLK period. Note that RDY/BUSY is pulled High with a high-impedance pull-up prior to INIT going High.

The length of the BUSY signal depends on the activity in the UART. If the shift register was empty when the new byte was received, the BUSY signal lasts for only two CCLK periods. If the shift register was still full when the new byte was received, the BUSY signal can be as long as nine CCLK periods.

Note that after the last byte has been entered, only seven of its bits are shifted out. CCLK remains High with DOUT equal to bit 6 (the next-to-last bit) of the last byte entered.

The READY/BUSY handshake can be ignored if the delay from any one Write to the end of the next Write is guaranteed to be longer than 10 CCLK periods.

Status Read

The logic AND condition of the CS0, CS1and RS inputs puts the device status on the Data bus.

- D7 High indicates Ready
- D7 Low indicates Busy
- D0 through D6 go unconditionally High

It is mandatory that the whole start-up sequence be started and completed by one byte-wide input. Otherwise, the pins used as Write Strobe or Chip Enable might become active outputs and interfere with the final byte transfer. If this transfer does not occur, the start-up sequence is not completed all the way to the finish (point F in Figure 47 on page 53).

In this case, at worst, the internal reset is not released. At best, Readback and Boundary Scan are inhibited. The length-count value, as generated by the XACT*step* software, ensures that these problems never occur.

Although RDY/BUSY is brought out as a separate signal, microprocessors can more easily read this information on one of the data lines. For this purpose, D7 represents the RDY/BUSY status when RS is Low, WS is High, and the two chip select lines are both active.

Asynchronous Peripheral mode is selected by a <101> on the mode pins (M2, M1, M0).

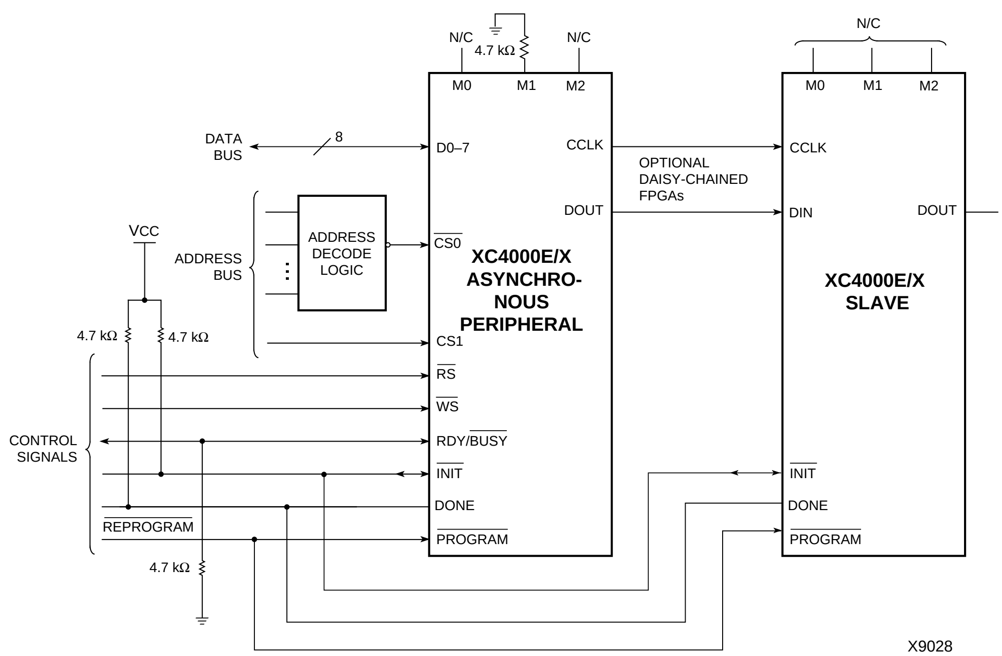
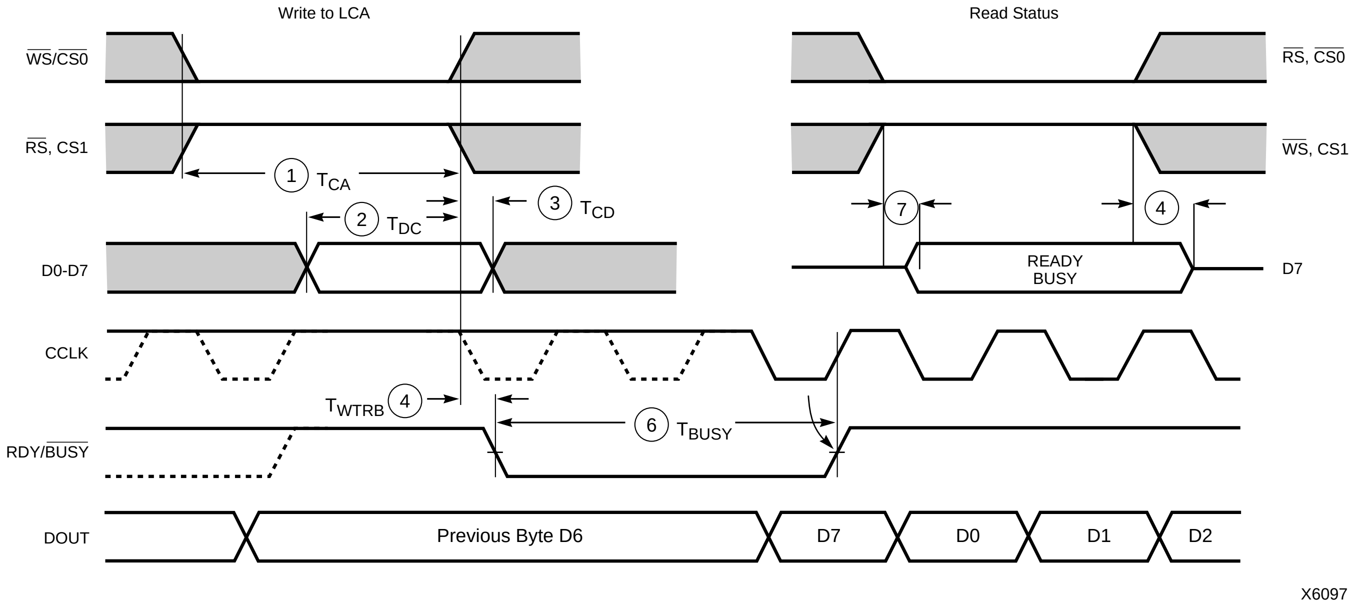


Figure 58: Asynchronous Peripheral Mode Circuit Diagram



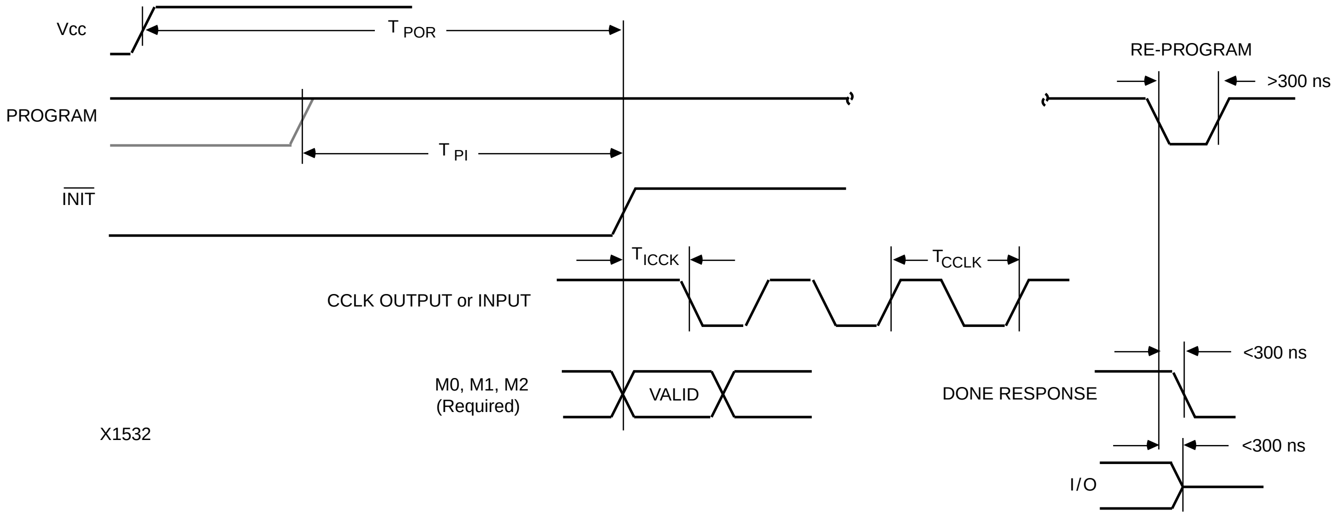
	Description	Symbol		Min	Max	Units
Write	Effective Write time ($\overline{\text{CS0}}$, $\overline{\text{WS}}$=Low; $\overline{\text{RS}}$, CS1=High)	1	T_{CA}	100		ns
	DIN setup time	2	T_{DC}	60		ns
	DIN hold time	3	T_{CD}	0		ns
RDY	RDY/$\overline{\text{BUSY}}$ delay after end of Write or Read	4	T_{WTRB}		60	ns
	RDY/$\overline{\text{BUSY}}$ active after beginning of Read	7			60	ns
	RDY/$\overline{\text{BUSY}}$ Low output (Note 4)	6	T_{BUSY}	2	9	CCLK periods

Notes:
1. Configuration must be delayed until the $\overline{\text{INIT}}$ pins of all daisy-chained FPGAs are High.
2. The time from the end of $\overline{\text{WS}}$ to CCLK cycle for the new byte of data depends on the completion of previous byte processing and the phase of the internal timing generator for CCLK.
3. CCLK and DOUT timing is tested in slave mode.
4. T_{BUSY} indicates that the double-buffered parallel-to-serial converter is not yet ready to receive new data. The shortest T_{BUSY} occurs when a byte is loaded into an empty parallel-to-serial converter. The longest T_{BUSY} occurs when a new word is loaded into the input register before the second-level buffer has started shifting out data

This timing diagram shows very relaxed requirements. Data need not be held beyond the rising edge of $\overline{\text{WS}}$. RDY/$\overline{\text{BUSY}}$ will go active within 60 ns after the end of $\overline{\text{WS}}$. A new write may be asserted immediately after RDY/$\overline{\text{BUSY}}$ goes Low, but write may not be terminated until RDY/$\overline{\text{BUSY}}$ has been High for one CCLK period.

Figure 59: Asynchronous Peripheral Mode Programming Switching Characteristics

Configuration Switching Characteristics



Master Modes (XC4000E/EX)

Description		Symbol	Min	Max	Units
Power-On Reset	M0 = High	T_{POR}	10	40	ms
	M0 = Low	T_{POR}	40	130	ms
Program Latency		T_{PI}	30	200	µs per CLB column
CCLK (output) Delay		T_{ICCK}	40	250	µs
CCLK (output) Period, slow		T_{CCLK}	640	2000	ns
CCLK (output) Period, fast		T_{CCLK}	80	250	ns

Master Modes (XC4000XL)

Description		Symbol	Min	Max	Units
Power-On Reset	M0 = High	T_{POR}	10	40	ms
	M0 = Low	T_{POR}	40	130	ms
Program Latency		T_{PI}	30	200	µs per CLB column
CCLK (output) Delay		T_{ICCK}	40	250	µs
CCLK (output) Period, slow		T_{CCLK}	540	1600	ns
CCLK (output) Period, fast		T_{CCLK}	67	200	ns

Slave and Peripheral Modes (All)

Description	Symbol	Min	Max	Units
Power-On Reset	T_{POR}	10	33	ms
Program Latency	T_{PI}	30	200	µs per CLB column
CCLK (input) Delay (required)	T_{ICCK}	4		µs
CCLK (input) Period (required)	T_{CCLK}	100		ns