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### Understanding [Embedded - FPGAs \(Field Programmable Gate Array\)](#)

Embedded - FPGAs, or Field Programmable Gate Arrays, are advanced integrated circuits that offer unparalleled flexibility and performance for digital systems. Unlike traditional fixed-function logic devices, FPGAs can be programmed and reprogrammed to execute a wide array of logical operations, enabling customized functionality tailored to specific applications. This reprogrammability allows developers to iterate designs quickly and implement complex functions without the need for custom hardware.

### Applications of Embedded - FPGAs

The versatility of Embedded - FPGAs makes them indispensable in numerous fields. In telecommunications.

#### Details

Product Status	Obsolete
Number of LABs/CLBs	576
Number of Logic Elements/Cells	1368
Total RAM Bits	18432
Number of I/O	113
Number of Gates	13000
Voltage - Supply	3V ~ 3.6V
Mounting Type	Surface Mount
Operating Temperature	-40°C ~ 100°C (TJ)
Package / Case	144-LQFP Exposed Pad
Supplier Device Package	144-TQFP (20x20)
Purchase URL	<a href="https://www.e-xfl.com/product-detail/xilinx/xc4013xl-1ht144i">https://www.e-xfl.com/product-detail/xilinx/xc4013xl-1ht144i</a>

## Detailed Functional Description

XC4000 Series devices achieve high speed through advanced semiconductor technology and improved architecture. The XC4000E and XC4000X support system clock rates of up to 80 MHz and internal performance in excess of 150 MHz. Compared to older Xilinx FPGA families, XC4000 Series devices are more powerful. They offer on-chip edge-triggered and dual-port RAM, clock enables on I/O flip-flops, and wide-input decoders. They are more versatile in many applications, especially those involving RAM. Design cycles are faster due to a combination of increased routing resources and more sophisticated software.

### Basic Building Blocks

Xilinx user-programmable gate arrays include two major configurable elements: configurable logic blocks (CLBs) and input/output blocks (IOBs).

- CLBs provide the functional elements for constructing the user's logic.
- IOBs provide the interface between the package pins and internal signal lines.

Three other types of circuits are also available:

- 3-State buffers (TBUFs) driving horizontal longlines are associated with each CLB.
- Wide edge decoders are available around the periphery of each device.
- An on-chip oscillator is provided.

Programmable interconnect resources provide routing paths to connect the inputs and outputs of these configurable elements to the appropriate networks.

The functionality of each circuit block is customized during configuration by programming internal static memory cells. The values stored in these memory cells determine the logic functions and interconnections implemented in the FPGA. Each of these available circuits is described in this section.

### Configurable Logic Blocks (CLBs)

Configurable Logic Blocks implement most of the logic in an FPGA. The principal CLB elements are shown in **Figure 1**. Two 4-input function generators (F and G) offer unrestricted versatility. Most combinatorial logic functions need four or fewer inputs. However, a third function generator (H) is provided. The H function generator has three inputs. Either zero, one, or two of these inputs can be the outputs of F and G; the other input(s) are from outside the CLB. The CLB can, therefore, implement certain functions of up to nine variables, like parity check or expandable-identity comparison of two sets of four inputs.

Each CLB contains two storage elements that can be used to store the function generator outputs. However, the storage elements and function generators can also be used independently. These storage elements can be configured as flip-flops in both XC4000E and XC4000X devices; in the XC4000X they can optionally be configured as latches. DIN can be used as a direct input to either of the two storage elements. H1 can drive the other through the H function generator. Function generator outputs can also drive two outputs independent of the storage element outputs. This versatility increases logic capacity and simplifies routing.

Thirteen CLB inputs and four CLB outputs provide access to the function generators and storage elements. These inputs and outputs connect to the programmable interconnect resources outside the block.

### Function Generators

Four independent inputs are provided to each of two function generators (F1 - F4 and G1 - G4). These function generators, with outputs labeled F' and G', are each capable of implementing any arbitrarily defined Boolean function of four inputs. The function generators are implemented as memory look-up tables. The propagation delay is therefore independent of the function implemented.

A third function generator, labeled H', can implement any Boolean function of its three inputs. Two of these inputs can optionally be the F' and G' functional generator outputs. Alternatively, one or both of these inputs can come from outside the CLB (H2, H0). The third input must come from outside the block (H1).

Signals from the function generators can exit the CLB on two outputs. F' or H' can be connected to the X output. G' or H' can be connected to the Y output.

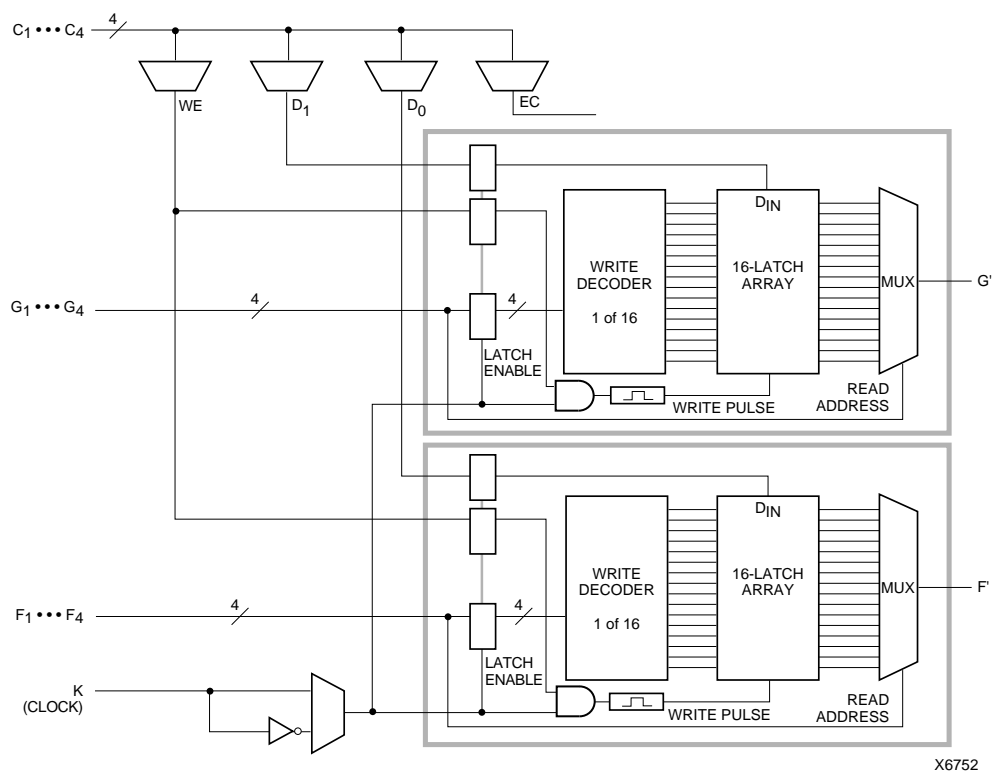
A CLB can be used to implement any of the following functions:

- any function of up to four variables, plus any second function of up to four unrelated variables, plus any third function of up to three unrelated variables<sup>1</sup>
- any single function of five variables
- any function of four variables together with some functions of six variables
- some functions of up to nine variables.

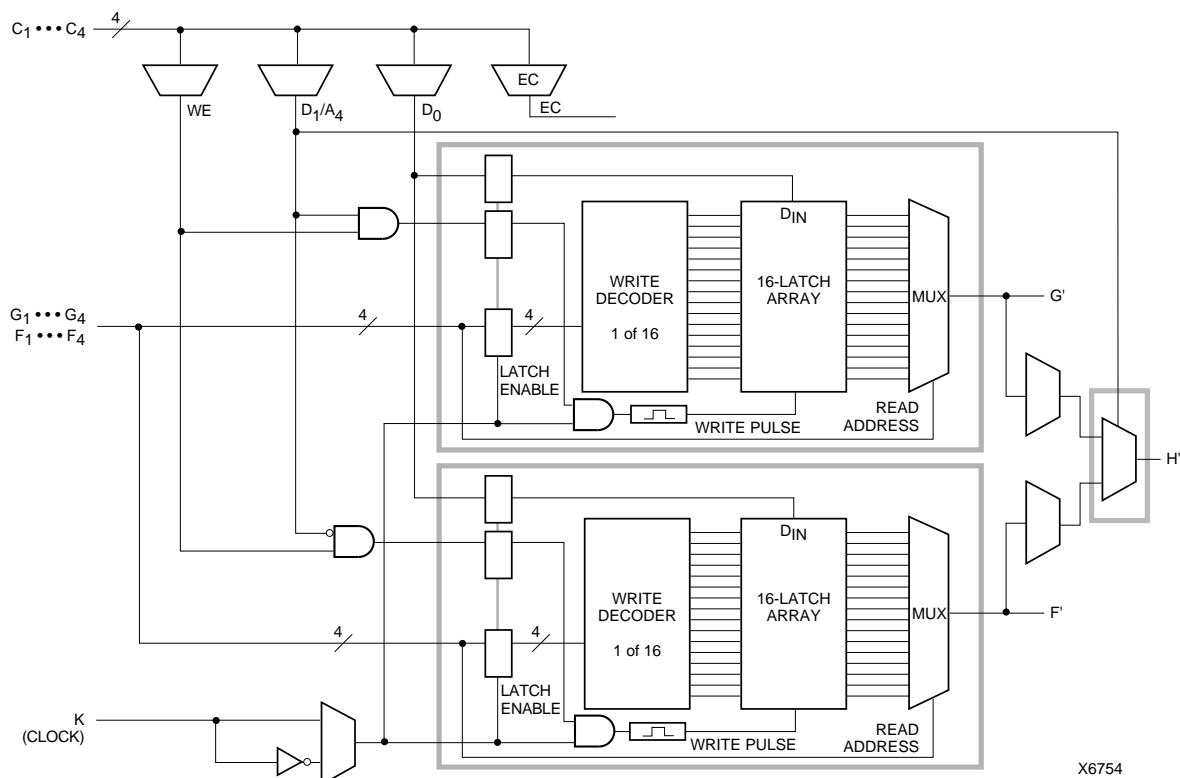
Implementing wide functions in a single block reduces both the number of blocks required and the delay in the signal path, achieving both increased capacity and speed.

The versatility of the CLB function generators significantly improves system speed. In addition, the design-software tools can deal with each function generator independently. This flexibility improves cell usage.

1. When three separate functions are generated, one of the function outputs must be captured in a flip-flop internal to the CLB. Only two unregistered function generator outputs are available from the CLB.



**Figure 4: 16x2 (or 16x1) Edge-Triggered Single-Port RAM**



**Figure 5: 32x1 Edge-Triggered Single-Port RAM (F and G addresses are identical)**

### Fast Carry Logic

Each CLB F and G function generator contains dedicated arithmetic logic for the fast generation of carry and borrow signals. This extra output is passed on to the function generator in the adjacent CLB. The carry chain is independent of normal routing resources.

Dedicated fast carry logic greatly increases the efficiency and performance of adders, subtractors, accumulators, comparators and counters. It also opens the door to many new applications involving arithmetic operation, where the previous generations of FPGAs were not fast enough or too inefficient. High-speed address offset calculations in micro-processor or graphics systems, and high-speed addition in digital signal processing are two typical applications.

The two 4-input function generators can be configured as a 2-bit adder with built-in hidden carry that can be expanded to any length. This dedicated carry circuitry is so fast and efficient that conventional speed-up methods like carry generate/propagate are meaningless even at the 16-bit level, and of marginal benefit at the 32-bit level.

This fast carry logic is one of the more significant features of the XC4000 Series, speeding up arithmetic and counting into the 70 MHz range.

The carry chain in XC4000E devices can run either up or down. At the top and bottom of the columns where there are no CLBs above or below, the carry is propagated to the right. (See Figure 11.) In order to improve speed in the high-capacity XC4000X devices, which can potentially have very long carry chains, the carry chain travels upward only, as shown in Figure 12. Additionally, standard interconnect can be used to route a carry signal in the downward direction.

Figure 13 on page 19 shows an XC4000E CLB with dedicated fast carry logic. The carry logic in the XC4000X is similar, except that COUT exits at the top only, and the signal CINDOWN does not exist. As shown in Figure 13, the carry logic shares operand and control inputs with the function generators. The carry outputs connect to the function generators, where they are combined with the operands to form the sums.

Figure 14 on page 20 shows the details of the carry logic for the XC4000E. This diagram shows the contents of the box labeled "CARRY LOGIC" in Figure 13. The XC4000X carry logic is very similar, but a multiplexer on the pass-through carry chain has been eliminated to reduce delay. Additionally, in the XC4000X the multiplexer on the G4 path has a memory-programmable 0 input, which permits G4 to directly connect to COUT. G4 thus becomes an additional high-speed initialization path for carry-in.

The dedicated carry logic is discussed in detail in Xilinx document XAPP 013: "Using the Dedicated Carry Logic in

XC4000." This discussion also applies to XC4000E devices, and to XC4000X devices when the minor logic changes are taken into account.

The fast carry logic can be accessed by placing special library symbols, or by using Xilinx Relationally Placed Macros (RPMs) that already include these symbols.



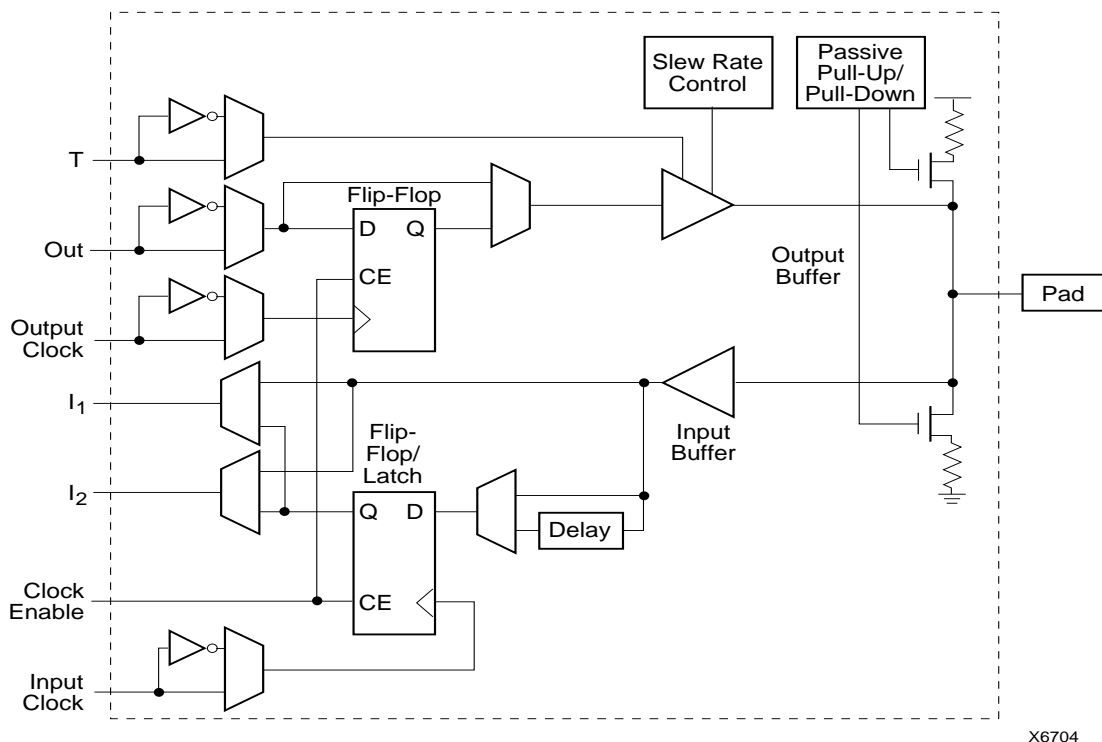
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**Figure 11: Available XC4000E Carry Propagation Paths**

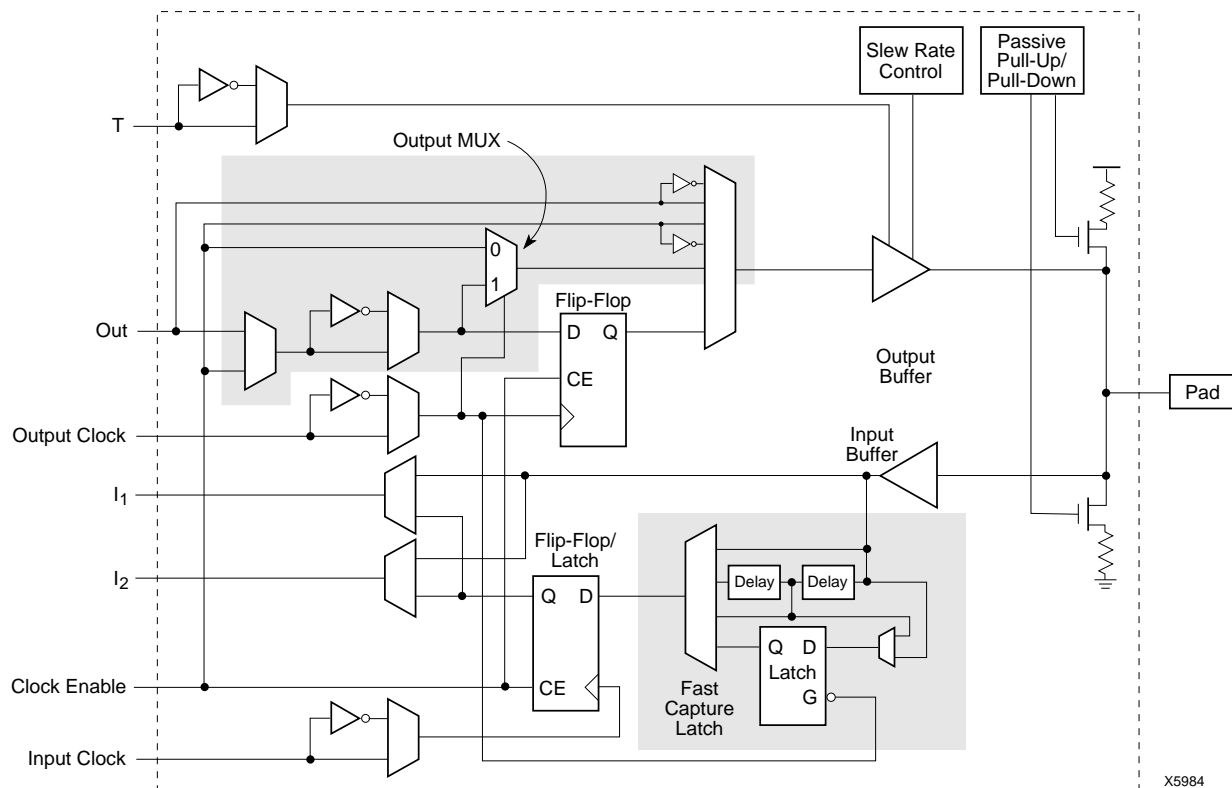


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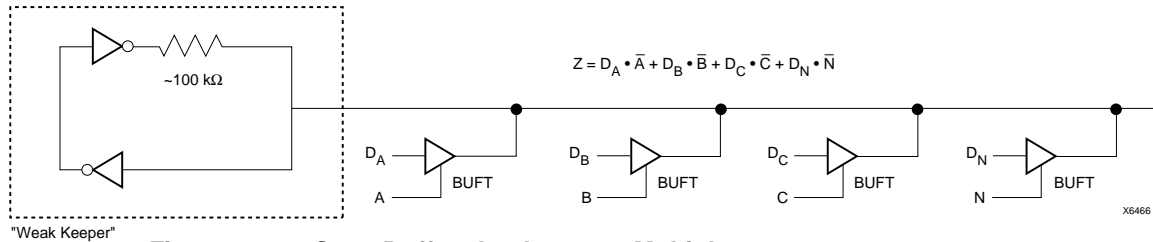
**Figure 12: Available XC4000X Carry Propagation Paths (dotted lines use general interconnect)**



**Figure 15: Simplified Block Diagram of XC4000E IOB**



**Figure 16: Simplified Block Diagram of XC4000X IOB (shaded areas indicate differences from XC4000E)**



**Figure 22: 3-State Buffers Implement a Multiplexer**

## Wide Edge Decoders

Dedicated decoder circuitry boosts the performance of wide decoding functions. When the address or data field is wider than the function generator inputs, FPGAs need multi-level decoding and are thus slower than PALs. XC4000 Series CLBs have nine inputs. Any decoder of up to nine inputs is, therefore, compact and fast. However, there is also a need for much wider decoders, especially for address decoding in large microprocessor systems.

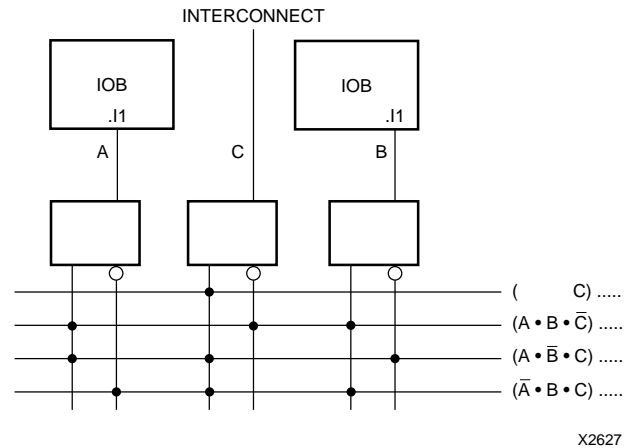
An XC4000 Series FPGA has four programmable decoders located on each edge of the device. The inputs to each decoder are any of the IOB I1 signals on that edge plus one local interconnect per CLB row or column. Each row or column of CLBs provides up to three variables or their complements., as shown in Figure 23. Each decoder generates a High output (resistor pull-up) when the AND condition of the selected inputs, or their complements, is true. This is analogous to a product term in typical PAL devices.

Each of these wired-AND gates is capable of accepting up to 42 inputs on the XC4005E and 72 on the XC4013E. There are up to 96 inputs for each decoder on the XC4028X and 132 on the XC4052X. The decoders may also be split in two when a larger number of narrower decoders are required, for a maximum of 32 decoders per device.

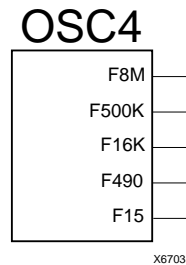
The decoder outputs can drive CLB inputs, so they can be combined with other logic to form a PAL-like AND/OR structure. The decoder outputs can also be routed directly to the chip outputs. For fastest speed, the output should be on the same chip edge as the decoder. Very large PALs can be emulated by ORing the decoder outputs in a CLB. This decoding feature covers what has long been considered a weakness of older FPGAs. Users often resorted to external PALs for simple but fast decoding functions. Now, the dedicated decoders in the XC4000 Series device can implement these functions fast and efficiently.

To use the wide edge decoders, place one or more of the WAND library symbols (WAND1, WAND4, WAND8, WAND16). Attach a DECODE attribute or property to each WAND symbol. Tie the outputs together and attach a PUL-

LUP symbol. Location attributes or properties such as L (left edge) or TR (right half of top edge) should also be used to ensure the correct placement of the decoder inputs.



**Figure 23: XC4000 Series Edge Decoding Example**



**Figure 24: XC4000 Series Oscillator Symbol**

## On-Chip Oscillator

XC4000 Series devices include an internal oscillator. This oscillator is used to clock the power-on time-out, for configuration memory clearing, and as the source of CCLK in Master configuration modes. The oscillator runs at a nominal 8 MHz frequency that varies with process, Vcc, and temperature. The output frequency falls between 4 and 10 MHz.

The oscillator output is optionally available after configuration. Any two of four resynchronized taps of a built-in divider are also available. These taps are at the fourth, ninth, fourteenth and nineteenth bits of the divider. Therefore, if the primary oscillator output is running at the nominal 8 MHz, the user has access to an 8 MHz clock, plus any two of 500 kHz, 16kHz, 490Hz and 15Hz (up to 10% lower for low-voltage devices). These frequencies can vary by as much as -50% or +25%.

These signals can be accessed by placing the OSC4 library element in a schematic or in HDL code (see [Figure 24](#)).

The oscillator is automatically disabled after configuration if the OSC4 symbol is not used in the design.

## Programmable Interconnect

All internal connections are composed of metal segments with programmable switching points and switching matrices to implement the desired routing. A structured, hierarchical matrix of routing resources is provided to achieve efficient automated routing.

The XC4000E and XC4000X share a basic interconnect structure. XC4000X devices, however, have additional routing not available in the XC4000E. The extra routing resources allow high utilization in high-capacity devices. All XC4000X-specific routing resources are clearly identified throughout this section. Any resources not identified as XC4000X-specific are present in all XC4000 Series devices.

This section describes the varied routing resources available in XC4000 Series devices. The implementation software automatically assigns the appropriate resources based on the density and timing requirements of the design.

## Interconnect Overview

There are several types of interconnect.

- CLB routing is associated with each row and column of the CLB array.
- IOB routing forms a ring (called a VersaRing) around the outside of the CLB array. It connects the I/O with the internal logic blocks.

- Global routing consists of dedicated networks primarily designed to distribute clocks throughout the device with minimum delay and skew. Global routing can also be used for other high-fanout signals.

Five interconnect types are distinguished by the relative length of their segments: single-length lines, double-length lines, quad and octal lines (XC4000X only), and longlines. In the XC4000X, direct connects allow fast data flow between adjacent CLBs, and between IOBs and CLBs.

Extra routing is included in the IOB pad ring. The XC4000X also includes a ring of octal interconnect lines near the IOBs to improve pin-swapping and routing to locked pins.

XC4000E/X devices include two types of global buffers. These global buffers have different properties, and are intended for different purposes. They are discussed in detail later in this section.

## CLB Routing Connections

A high-level diagram of the routing resources associated with one CLB is shown in [Figure 25](#). The shaded arrows represent routing present only in XC4000X devices.

[Table 14](#) shows how much routing of each type is available in XC4000E and XC4000X CLB arrays. Clearly, very large designs, or designs with a great deal of interconnect, will route more easily in the XC4000X. Smaller XC4000E designs, typically requiring significantly less interconnect, do not require the additional routing.

[Figure 27 on page 30](#) is a detailed diagram of both the XC4000E and the XC4000X CLB, with associated routing. The shaded square is the programmable switch matrix, present in both the XC4000E and the XC4000X. The L-shaped shaded area is present only in XC4000X devices. As shown in the figure, the XC4000X block is essentially an XC4000E block with additional routing.

CLB inputs and outputs are distributed on all four sides, providing maximum routing flexibility. In general, the entire architecture is symmetrical and regular. It is well suited to established placement and routing algorithms. Inputs, outputs, and function generators can freely swap positions within a CLB to avoid routing congestion during the placement and routing operation.





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**Figure 25: High-Level Routing Diagram of XC4000 Series CLB (shaded arrows indicate XC4000X only)**

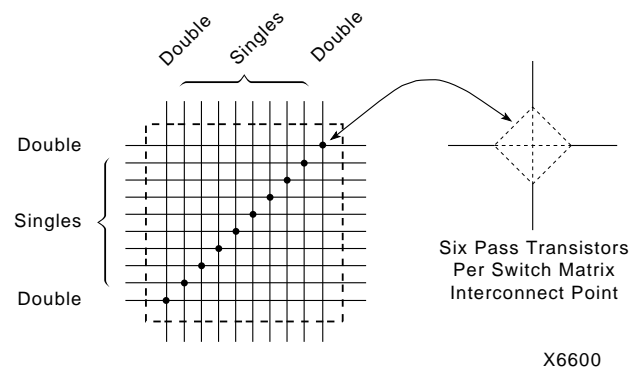
**Table 14: Routing per CLB in XC4000 Series Devices**

	XC4000E		XC4000X	
	Vertical	Horizontal	Vertical	Horizontal
Singles	8	8	8	8
Doubles	4	4	4	4
Quads	0	0	12	12
Longlines	6	6	10	6
Direct Connects	0	0	2	2
Globals	4	0	8	0
Carry Logic	2	0	1	0
Total	24	18	45	32

### Programmable Switch Matrices

The horizontal and vertical single- and double-length lines intersect at a box called a programmable switch matrix (PSM). Each switch matrix consists of programmable pass transistors used to establish connections between the lines (see Figure 26).

For example, a single-length signal entering on the right side of the switch matrix can be routed to a single-length line on the top, left, or bottom sides, or any combination thereof, if multiple branches are required. Similarly, a double-length signal can be routed to a double-length line on any or all of the other three edges of the programmable switch matrix.



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**Figure 26: Programmable Switch Matrix (PSM)**

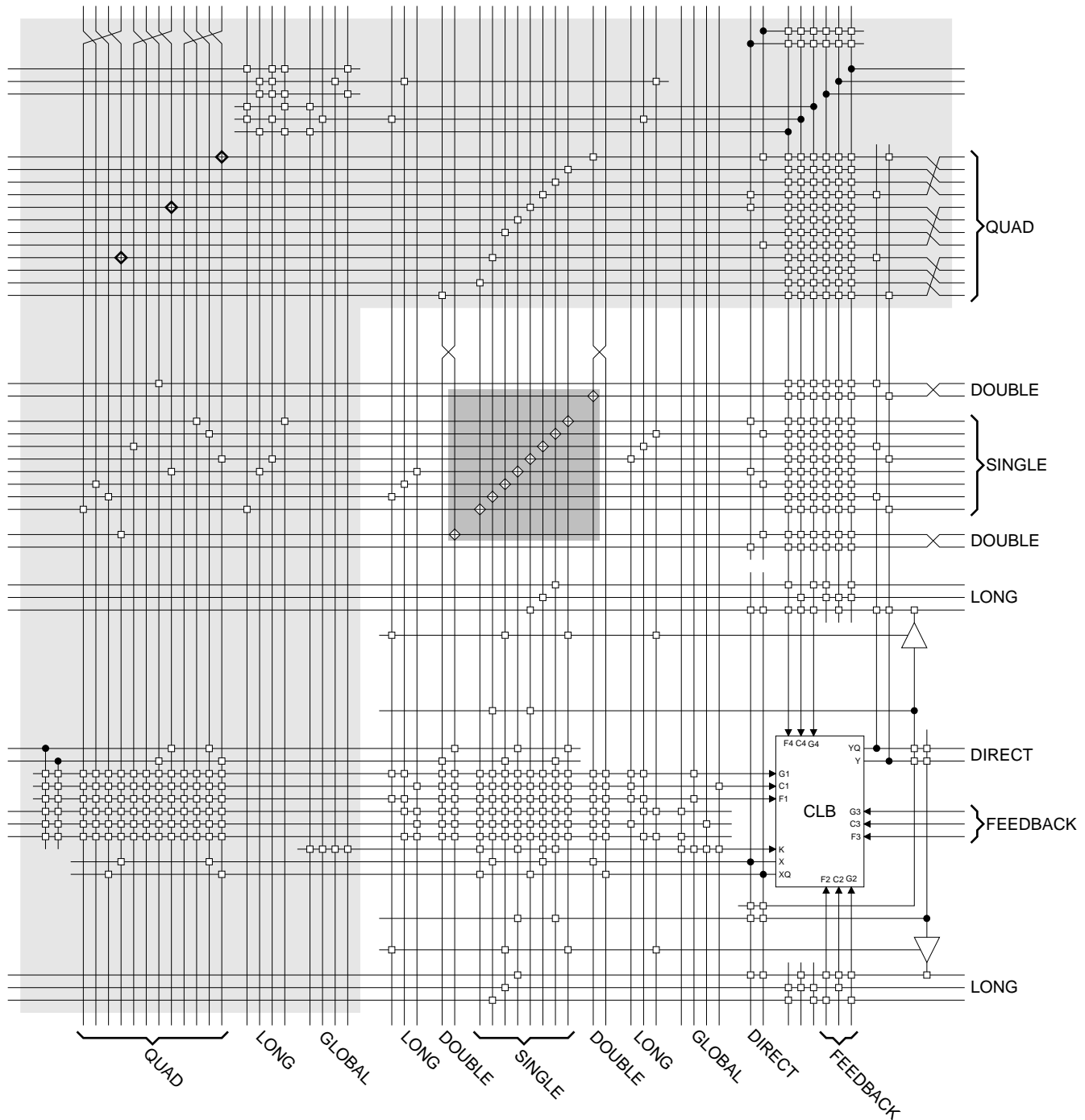
### Single-Length Lines

Single-length lines provide the greatest interconnect flexibility and offer fast routing between adjacent blocks. There are eight vertical and eight horizontal single-length lines associated with each CLB. These lines connect the switching matrices that are located in every row and a column of CLBs.

Single-length lines are connected by way of the programmable switch matrices, as shown in Figure 28. Routing connectivity is shown in Figure 27.

Single-length lines incur a delay whenever they go through a switching matrix. Therefore, they are not suitable for routing signals for long distances. They are normally used to conduct signals within a localized area and to provide the branching for nets with fanout greater than one.





**Figure 27: Detail of Programmable Interconnect Associated with XC4000 Series CLB**

circuit prevents undefined floating levels. However, it is overridden by any driver, even a pull-up resistor.

Each XC4000E longline has a programmable splitter switch at its center, as does each XC4000X longline driven by TBUFs. This switch can separate the line into two independent routing channels, each running half the width or height of the array.

Each XC4000X longline not driven by TBUFs has a buffered programmable splitter switch at the 1/4, 1/2, and 3/4 points of the array. Due to the buffering, XC4000X longline performance does not deteriorate with the larger array sizes. If the longline is split, the resulting partial longlines are independent.

Routing connectivity of the longlines is shown in [Figure 27 on page 30](#).

### **Direct Interconnect (XC4000X only)**

The XC4000X offers two direct, efficient and fast connections between adjacent CLBs. These nets facilitate a data flow from the left to the right side of the device, or from the top to the bottom, as shown in [Figure 30](#). Signals routed on the direct interconnect exhibit minimum interconnect propagation delay and use no general routing resources.

The direct interconnect is also present between CLBs and adjacent IOBs. Each IOB on the left and top device edges has a direct path to the nearest CLB. Each CLB on the right and bottom edges of the array has a direct path to the nearest two IOBs, since there are two IOBs for each row or column of CLBs.

The place and route software uses direct interconnect whenever possible, to maximize routing resources and minimize interconnect delays.



**Figure 30: XC4000X Direct Interconnect**

### **I/O Routing**

XC4000 Series devices have additional routing around the IOB ring. This routing is called a VersaRing. The VersaRing facilitates pin-swapping and redesign without affecting board layout. Included are eight double-length lines spanning two CLBs (four IOBs), and four longlines. Global lines and Wide Edge Decoder lines are provided. XC4000X devices also include eight octal lines.

A high-level diagram of the VersaRing is shown in [Figure 31](#). The shaded arrows represent routing present only in XC4000X devices.

[Figure 33 on page 34](#) is a detailed diagram of the XC4000E and XC4000X VersaRing. The area shown includes two IOBs. There are two IOBs per CLB row or column, therefore this diagram corresponds to the CLB routing diagram shown in [Figure 27 on page 30](#). The shaded areas represent routing and routing connections present only in XC4000X devices.

### **Octal I/O Routing (XC4000X only)**

Between the XC4000X CLB array and the pad ring, eight interconnect tracks provide for versatility in pin assignment and fixed pinout flexibility. (See [Figure 32 on page 33](#).)

These routing tracks are called octals, because they can be broken every eight CLBs (sixteen IOBs) by a programmable buffer that also functions as a splitter switch. The buffers are staggered, so each line goes through a buffer at every eighth CLB location around the device edge.

The octal lines bend around the corners of the device. The lines cross at the corners in such a way that the segment most recently buffered before the turn has the farthest distance to travel before the next buffer, as shown in [Figure 32](#).

IOB inputs and outputs interface with the octal lines via the single-length interconnect lines. Single-length lines are also used for communication between the octals and double-length lines, quads, and longlines within the CLB array.

Segmentation into buffered octals was found to be optimal for distributing signals over long distances around the device.

### Global Nets and Buffers

Both the XC4000E and the XC4000X have dedicated global networks. These networks are designed to distribute clocks and other high fanout control signals throughout the devices with minimal skew. The global buffers are described in detail in the following sections. The text descriptions and diagrams are summarized in [Table 15](#). The table shows which CLB and IOB clock pins can be sourced by which global buffers.

In both XC4000E and XC4000X devices, placement of a library symbol called BUFG results in the software choosing the appropriate clock buffer, based on the timing requirements of the design. The detailed information in these sections is included only for reference.

#### Global Nets and Buffers (XC4000E only)

Four vertical longlines in each CLB column are driven exclusively by special global buffers. These longlines are in addition to the vertical longlines used for standard interconnect. The four global lines can be driven by either of two types of global buffers. The clock pins of every CLB and IOB can also be sourced from local interconnect.

Two different types of clock buffers are available in the XC4000E:

- Primary Global Buffers (BUFGP)
- Secondary Global Buffers (BUFGS)

Four Primary Global buffers offer the shortest delay and negligible skew. Four Secondary Global buffers have slightly longer delay and slightly more skew due to potentially heavier loading, but offer greater flexibility when used to drive non-clock CLB inputs.

The Primary Global buffers must be driven by the semi-dedicated pads. The Secondary Global buffers can be sourced by either semi-dedicated pads or internal nets.

Each CLB column has four dedicated vertical Global lines. Each of these lines can be accessed by one particular Primary Global buffer, or by any of the Secondary Global buffers, as shown in [Figure 34](#). Each corner of the device has one Primary buffer and one Secondary buffer.

IOBs along the left and right edges have four vertical global longlines. Top and bottom IOBs can be clocked from the global lines in the adjacent CLB column.

A global buffer should be specified for all timing-sensitive global signal distribution. To use a global buffer, place a BUFGP (primary buffer), BUFGS (secondary buffer), or BUFG (either primary or secondary buffer) element in a schematic or in HDL code. If desired, attach a LOC attribute or property to direct placement to the designated location. For example, attach a LOC=L attribute or property to a BUFGS symbol to direct that a buffer be placed in one of the two Secondary Global buffers on the left edge of the device, or a LOC=BL to indicate the Secondary Global buffer on the bottom edge of the device, on the left.

**Table 15: Clock Pin Access**

	XC4000E		XC4000X			Local Inter-connect
	BUFGP	BUFGS	BUFGLS	L & R BUFGE	T & B BUFGE	
All CLBs in Quadrant	√	√	√	√	√	√
All CLBs in Device	√	√	√			√
IOBs on Adjacent Vertical Half Edge	√	√	√	√	√	√
IOBs on Adjacent Vertical Full Edge	√	√	√	√		√
IOBs on Adjacent Horizontal Half Edge (Direct)				√		√
IOBs on Adjacent Horizontal Half Edge (through CLB globals)	√	√	√	√	√	√
IOBs on Adjacent Horizontal Full Edge (through CLB globals)	√	√	√			√

L = Left, R = Right, T = Top, B = Bottom

The top and bottom Global Early buffers are about 1 ns slower clock to out than the left and right Global Early buffers.

The Global Early buffers can be driven by either semi-dedicated pads or internal logic. They share pads with the Global Low-Skew buffers, so a single net can drive both global buffers, as described above.

To use a Global Early buffer, place a BUFGE element in a schematic or in HDL code. If desired, attach a LOC attribute or property to direct placement to the designated location. For example, attach a LOC=T attribute or property to direct that a BUFGE be placed in one of the two Global Early buffers on the top edge of the device, or a LOC=TR to indicate the Global Early buffer on the top edge of the device, on the right.

## Power Distribution

Power for the FPGA is distributed through a grid to achieve high noise immunity and isolation between logic and I/O. Inside the FPGA, a dedicated Vcc and Ground ring surrounding the logic array provides power to the I/O drivers, as shown in [Figure 39](#). An independent matrix of Vcc and Ground lines supplies the interior logic of the device.

This power distribution grid provides a stable supply and ground for all internal logic, providing the external package power pins are all connected and appropriately de-coupled. Typically, a 0.1  $\mu$ F capacitor connected between each Vcc pin and the board's Ground plane will provide adequate de-coupling.

Output buffers capable of driving/sinking the specified 12 mA loads under specified worst-case conditions may be capable of driving/sinking up to 10 times as much current under best case conditions.

Noise can be reduced by minimizing external load capacitance and reducing simultaneous output transitions in the same direction. It may also be beneficial to locate heavily loaded output buffers near the Ground pads. The I/O Block output buffers have a slew-rate limited mode (default) which should be used where output rise and fall times are not speed-critical.



Figure 39: XC4000 Series Power Distribution

## Pin Descriptions

There are three types of pins in the XC4000 Series devices:

- Permanently dedicated pins
- User I/O pins that can have special functions
- Unrestricted user-programmable I/O pins.

Before and during configuration, all outputs not used for the configuration process are 3-stated with a 50 k $\Omega$  - 100 k $\Omega$  pull-up resistor.

After configuration, if an IOB is unused it is configured as an input with a 50 k $\Omega$  - 100 k $\Omega$  pull-up resistor.

XC4000 Series devices have no dedicated Reset input. Any user I/O can be configured to drive the Global Set/Reset net, GSR. See ["Global Set/Reset" on page 11](#) for more information on GSR.

XC4000 Series devices have no Powerdown control input, as the XC3000 and XC2000 families do. The XC3000/XC2000 Powerdown control also 3-stated all of the device

I/O pins. For XC4000 Series devices, use the global 3-state net, GTS, instead. This net 3-states all outputs, but does not place the device in low-power mode. See ["IOB Output Signals" on page 23](#) for more information on GTS.

Device pins for XC4000 Series devices are described in [Table 16](#). Pin functions during configuration for each of the seven configuration modes are summarized in [Table 22 on page 58](#), in the "Configuration Timing" section.

is passed through and is captured by each FPGA when it recognizes the 0010 preamble. Following the length-count data, each FPGA outputs a High on DOUT until it has received its required number of data frames.

After an FPGA has received its configuration data, it passes on any additional frame start bits and configuration data on DOUT. When the total number of configuration clocks applied after memory initialization equals the value of the 24-bit length count, the FPGAs begin the start-up sequence and become operational together. FPGA I/O are normally released two CCLK cycles after the last configuration bit is received. **Figure 47 on page 53** shows the start-up timing for an XC4000 Series device.

The daisy-chained bitstream is not simply a concatenation of the individual bitstreams. The PROM file formatter must be used to combine the bitstreams for a daisy-chained configuration.

### Multi-Family Daisy Chain

All Xilinx FPGAs of the XC2000, XC3000, and XC4000 Series use a compatible bitstream format and can, therefore, be connected in a daisy chain in an arbitrary sequence. There is, however, one limitation. The lead device must belong to the highest family in the chain. If the chain contains XC4000 Series devices, the master normally cannot be an XC2000 or XC3000 device.

The reason for this rule is shown in **Figure 47 on page 53**. Since all devices in the chain store the same length count value and generate or receive one common sequence of CCLK pulses, they all recognize length-count match on the same CCLK edge, as indicated on the left edge of **Figure 47**. The master device then generates additional CCLK pulses until it reaches its finish point F. The different families generate or require different numbers of additional CCLK pulses until they reach F. Not reaching F means that the device does not really finish its configuration, although DONE may have gone High, the outputs became active, and the internal reset was released. For the XC4000 Series device, not reaching F means that readback cannot be ini-

tiated and most boundary scan instructions cannot be used.

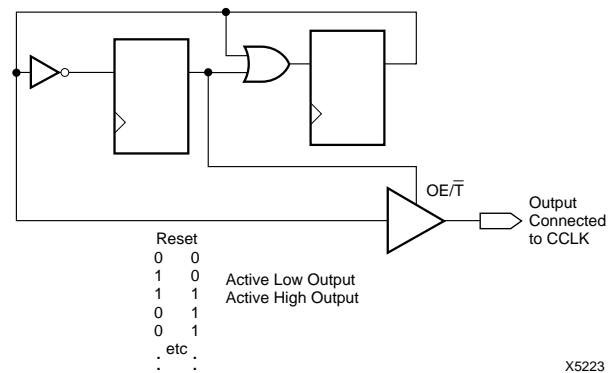
The user has some control over the relative timing of these events and can, therefore, make sure that they occur at the proper time and the finish point F is reached. Timing is controlled using options in the bitstream generation software.

### XC3000 Master with an XC4000 Series Slave

Some designers want to use an inexpensive lead device in peripheral mode and have the more precious I/O pins of the XC4000 Series devices all available for user I/O. **Figure 44** provides a solution for that case.

This solution requires one CLB, one IOB and pin, and an internal oscillator with a frequency of up to 5 MHz as a clock source. The XC3000 master device must be configured with late Internal Reset, which is the default option.

One CLB and one IOB in the lead XC3000-family device are used to generate the additional CCLK pulse required by the XC4000 Series devices. When the lead device removes the internal RESET signal, the 2-bit shift register responds to its clock input and generates an active Low output signal for the duration of the subsequent clock period. An external connection between this output and CCLK thus creates the extra CCLK pulse.



**Figure 44: CCLK Generation for XC3000 Master Driving an XC4000 Series Slave**

Table 20: XC4000E Program Data

Device	XC4003E	XC4005E	XC4006E	XC4008E	XC4010E	XC4013E	XC4020E	XC4025E
Max Logic Gates	3,000	5,000	6,000	8,000	10,000	13,000	20,000	25,000
CLBs (Row x Col.)	100 (10 x 10)	196 (14 x 14)	256 (16 x 16)	324 (18 x 18)	400 (20 x 20)	576 (24 x 24)	784 (28 x 28)	1,024 (32 x 32)
IOBs	80	112	128	144	160	192	224	256
Flip-Flops	360	616	768	936	1,120	1,536	2,016	2,560
Bits per Frame	126	166	186	206	226	266	306	346
Frames	428	572	644	716	788	932	1,076	1,220
Program Data	53,936	94,960	119,792	147,504	178,096	247,920	329,264	422,128
PROM Size (bits)	53,984	95,008	119,840	147,552	178,144	247,968	329,312	422,176

- Notes:
- Bits per Frame = (10 x number of rows) + 7 for the top + 13 for the bottom + 1 + 1 start bit + 4 error check bits  
 Number of Frames = (36 x number of columns) + 26 for the left edge + 41 for the right edge + 1  
 Program Data = (Bits per Frame x Number of Frames) + 8 postamble bits  
 PROM Size = Program Data + 40 (header) + 8
  - The user can add more "one" bits as leading dummy bits in the header, or, if CRC = off, as trailing dummy bits at the end of any frame, following the four error check bits. However, the Length Count value **must** be adjusted for all such extra "one" bits, even for extra leading ones at the beginning of the header.

Table 21: XC4000EX/XL Program Data

Device	XC4002XL	XC4005	XC4010	XC4013	XC4020	XC4028	XC4036	XC4044	XC4052	XC4062	XC4085
Max Logic Gates	2,000	5,000	10,000	13,000	20,000	28,000	36,000	44,000	52,000	62,000	85,000
CLBs (Row x Column)	64 (8 x 8)	196 (14 x 14)	400 (20 x 20)	576 (24 x 24)	784 (28 x 28)	1,024 (32 x 32)	1,296 (36 x 36)	1,600 (40 x 40)	1,936 (44 x 44)	2,304 (48 x 48)	3,136 (56 x 56)
IOBs	64	112	160	192	224	256	288	320	352	384	448
Flip-Flops	256	616	1,120	1,536	2,016	2,560	3,168	3,840	4,576	5,376	7,168
Bits per Frame	133	205	277	325	373	421	469	517	565	613	709
Frames	459	741	1,023	1,211	1,399	1,587	1,775	1,963	2,151	2,339	2,715
Program Data	61,052	151,910	283,376	393,580	521,832	668,124	832,480	1,014,876	1,215,320	1,433,804	1,924,940
PROM Size (bits)	61,104	151,960	283,424	393,632	521,880	668,172	832,528	1,014,924	1,215,368	1,433,852	1,924,992

- Notes:
- Bits per frame = (13 x number of rows) + 9 for the top + 17 for the bottom + 8 + 1 start bit + 4 error check bits.  
 Frames = (47 x number of columns) + 27 for the left edge + 52 for the right edge + 4.  
 Program data = (bits per frame x number of frames) + 5 postamble bits.  
 PROM size = (program data + 40 header bits + 8 start bits) rounded up to the nearest byte.
  - The user can add more "one" bits as leading dummy bits in the header, or, if CRC = off, as trailing dummy bits at the end of any frame, following the four error check bits. However, the Length Count value must be adjusted for all such extra "one" bits, even for extra leading "ones" at the beginning of the header.

## Cyclic Redundancy Check (CRC) for Configuration and Readback

The Cyclic Redundancy Check is a method of error detection in data transmission applications. Generally, the transmitting system performs a calculation on the serial bitstream. The result of this calculation is tagged onto the data stream as additional check bits. The receiving system performs an identical calculation on the bitstream and compares the result with the received checksum.

Each data frame of the configuration bitstream has four error bits at the end, as shown in [Table 19](#). If a frame data error is detected during the loading of the FPGA, the con-

figuration process with a potentially corrupted bitstream is terminated. The FPGA pulls the  $\overline{\text{INIT}}$  pin Low and goes into a Wait state.

During Readback, 11 bits of the 16-bit checksum are added to the end of the Readback data stream. The checksum is computed using the CRC-16 CCITT polynomial, as shown in [Figure 45](#). The checksum consists of the 11 most significant bits of the 16-bit code. A change in the checksum indicates a change in the Readback bitstream. A comparison to a previous checksum is meaningful only if the readback data is independent of the current device state. CLB outputs should not be included (Read Capture option not



Low. During this time delay, or as long as the  $\overline{\text{PROGRAM}}$  input is asserted, the configuration logic is held in a Configuration Memory Clear state. The configuration-memory frames are consecutively initialized, using the internal oscillator.

At the end of each complete pass through the frame addressing, the power-on time-out delay circuitry and the level of the  $\overline{\text{PROGRAM}}$  pin are tested. If neither is asserted, the logic initiates one additional clearing of the configuration frames and then tests the  $\overline{\text{INIT}}$  input.

### Initialization

During initialization and configuration, user pins  $\text{HDC}$ ,  $\overline{\text{LDC}}$ ,  $\overline{\text{INIT}}$  and  $\text{DONE}$  provide status outputs for the system interface. The outputs  $\overline{\text{LDC}}$ ,  $\overline{\text{INIT}}$  and  $\text{DONE}$  are held Low and  $\text{HDC}$  is held High starting at the initial application of power.

The open drain  $\overline{\text{INIT}}$  pin is released after the final initialization pass through the frame addresses. There is a deliberate delay of 50 to 250  $\mu\text{s}$  (up to 10% longer for low-voltage devices) before a Master-mode device recognizes an inactive  $\overline{\text{INIT}}$ . Two internal clocks after the  $\overline{\text{INIT}}$  pin is recognized as High, the FPGA samples the three mode lines to determine the configuration mode. The appropriate interface lines become active and the configuration preamble and data can be loaded. Configuration

The 0010 preamble code indicates that the following 24 bits represent the length count. The length count is the total number of configuration clocks needed to load the complete configuration data. (Four additional configuration clocks are required to complete the configuration process, as discussed below.) After the preamble and the length count have been passed through to all devices in the daisy chain,  $\text{DOUT}$  is held High to prevent frame start bits from reaching any daisy-chained devices.

A specific configuration bit, early in the first frame of a master device, controls the configuration-clock rate and can increase it by a factor of eight. Therefore, if a fast configuration clock is selected by the bitstream, the slower clock rate is used until this configuration bit is detected.

Each frame has a start field followed by the frame-configuration data bits and a frame error field. If a frame data error is detected, the FPGA halts loading, and signals the error by pulling the open-drain  $\overline{\text{INIT}}$  pin Low. After all configuration frames have been loaded into an FPGA,  $\text{DOUT}$  again follows the input data so that the remaining data is passed on to the next device.

### Delaying Configuration After Power-Up

There are two methods of delaying configuration after power-up: put a logic Low on the  $\overline{\text{PROGRAM}}$  input, or pull the bidirectional  $\overline{\text{INIT}}$  pin Low, using an open-collector (open-drain) driver. (See [Figure 46 on page 50](#).)

A Low on the  $\overline{\text{PROGRAM}}$  input is the more radical approach, and is recommended when the power-supply

rise time is excessive or poorly defined. As long as  $\overline{\text{PROGRAM}}$  is Low, the FPGA keeps clearing its configuration memory. When  $\overline{\text{PROGRAM}}$  goes High, the configuration memory is cleared one more time, followed by the beginning of configuration, provided the  $\overline{\text{INIT}}$  input is not externally held Low. Note that a Low on the  $\overline{\text{PROGRAM}}$  input automatically forces a Low on the  $\overline{\text{INIT}}$  output. The XC4000 Series  $\overline{\text{PROGRAM}}$  pin has a permanent weak pull-up.

Using an open-collector or open-drain driver to hold  $\overline{\text{INIT}}$  Low before the beginning of configuration causes the FPGA to wait after completing the configuration memory clear operation. When  $\overline{\text{INIT}}$  is no longer held Low externally, the device determines its configuration mode by capturing its mode pins, and is ready to start the configuration process. A master device waits up to an additional 250  $\mu\text{s}$  to make sure that any slaves in the optional daisy chain have seen that  $\overline{\text{INIT}}$  is High.

### Start-Up

Start-up is the transition from the configuration process to the intended user operation. This transition involves a change from one clock source to another, and a change from interfacing parallel or serial configuration data where most outputs are 3-stated, to normal operation with I/O pins active in the user-system. Start-up must make sure that the user-logic 'wakes up' gracefully, that the outputs become active without causing contention with the configuration signals, and that the internal flip-flops are released from the global Reset or Set at the right time.

[Figure 47](#) describes start-up timing for the three Xilinx families in detail. The configuration modes can use any of the four timing sequences.

To access the internal start-up signals, place the  $\text{STARTUP}$  library symbol.

### Start-up Timing

Different FPGA families have different start-up sequences.

The XC2000 family goes through a fixed sequence.  $\text{DONE}$  goes High and the internal global Reset is de-activated one CCLK period after the I/O become active.

The XC3000A family offers some flexibility.  $\text{DONE}$  can be programmed to go High one CCLK period before or after the I/O become active. Independent of  $\text{DONE}$ , the internal global Reset is de-activated one CCLK period before or after the I/O become active.

The XC4000 Series offers additional flexibility. The three events —  $\text{DONE}$  going High, the internal Set/Reset being de-activated, and the user I/O going active — can all occur in any arbitrary sequence. Each of them can occur one CCLK period before or after, or simultaneous with, any of the others. This relative timing is selected by means of software options in the bitstream generation software.



### **Start-up from a User Clock (STARTUP.CLK)**

When, instead of CCLK, a user-supplied start-up clock is selected, Q1 is used to bridge the unknown phase relationship between CCLK and the user clock. This arbitration causes an unavoidable one-cycle uncertainty in the timing of the rest of the start-up sequence.

### **DONE Goes High to Signal End of Configuration**

XC4000 Series devices read the expected length count from the bitstream and store it in an internal register. The length count varies according to the number of devices and the composition of the daisy chain. Each device also counts the number of CCLKs during configuration.

Two conditions have to be met in order for the DONE pin to go high:

- the chip's internal memory must be full, and
- the configuration length count must be met, *exactly*.

This is important because the counter that determines when the length count is met begins with the very first CCLK, not the first one after the preamble.

Therefore, if a stray bit is inserted before the preamble, or the data source is not ready at the time of the first CCLK, the internal counter that holds the number of CCLKs will be one ahead of the actual number of data bits read. At the end of configuration, the configuration memory will be full, but the number of bits in the internal counter will not match the expected length count.

As a consequence, a Master mode device will continue to send out CCLKs until the internal counter turns over to zero, and then reaches the correct length count a second time. This will take several seconds [ $2^{24} * \text{CCLK period}$ ] — which is sometimes interpreted as the device not configuring at all.

If it is not possible to have the data ready at the time of the first CCLK, the problem can be avoided by increasing the number in the length count by the appropriate value. The *XACT User Guide* includes detailed information about manually altering the length count.

Note that DONE is an open-drain output and does not go High unless an internal pull-up is activated or an external pull-up is attached. The internal pull-up is activated as the default by the bitstream generation software.

### **Release of User I/O After DONE Goes High**

By default, the user I/O are released one CCLK cycle after the DONE pin goes High. If CCLK is not clocked after DONE goes High, the outputs remain in their initial state — 3-stated, with a 50 k $\Omega$  - 100 k $\Omega$  pull-up. The delay from DONE High to active user I/O is controlled by an option to the bitstream generation software.

### **Release of Global Set/Reset After DONE Goes High**

By default, Global Set/Reset (GSR) is released two CCLK cycles after the DONE pin goes High. If CCLK is not clocked twice after DONE goes High, all flip-flops are held in their initial set or reset state. The delay from DONE High to GSR inactive is controlled by an option to the bitstream generation software.

### **Configuration Complete After DONE Goes High**

Three full CCLK cycles are required after the DONE pin goes High, as shown in [Figure 47 on page 53](#). If CCLK is not clocked three times after DONE goes High, readback cannot be initiated and most boundary scan instructions cannot be used.

### **Configuration Through the Boundary Scan Pins**

XC4000 Series devices can be configured through the boundary scan pins. The basic procedure is as follows:

- Power up the FPGA with  $\overline{\text{INIT}}$  held Low (or drive the  $\overline{\text{PROGRAM}}$  pin Low for more than 300 ns followed by a High while holding  $\overline{\text{INIT}}$  Low). Holding  $\overline{\text{INIT}}$  Low allows enough time to issue the CONFIG command to the FPGA. The pin can be used as I/O after configuration if a resistor is used to hold  $\overline{\text{INIT}}$  Low.
- Issue the CONFIG command to the TMS input
- Wait for  $\overline{\text{INIT}}$  to go High
- Sequence the boundary scan Test Access Port to the SHIFT-DR state
- Toggle TCK to clock data into TDI pin.

The user must account for all TCK clock cycles after INIT goes High, as all of these cycles affect the Length Count compare.

For more detailed information, refer to the Xilinx application note XAPP017, “*Boundary Scan in XC4000 Devices*.” This application note also applies to XC4000E and XC4000X devices.

**Table 22: Pin Functions During Configuration**

CONFIGURATION MODE <M2:M1:M0>						USER OPERATION
SLAVE SERIAL <1:1:1>	MASTER SERIAL <0:0:0>	SYNCH. PERIPHERAL <0:1:1>	ASYNCH. PERIPHERAL <1:0:1>	MASTER PARALLEL DOWN <1:1:0>	MASTER PARALLEL UP <1:0:0>	
M2(HIGH) (I)	M2(LOW) (I)	M2(LOW) (I)	M2(HIGH) (I)	M2(HIGH) (I)	M2(HIGH) (I)	(I)
M1(HIGH) (I)	M1(LOW) (I)	M1(HIGH) (I)	M1(LOW) (I)	M1(HIGH) (I)	M1(LOW) (I)	(O)
M0(HIGH) (I)	M0(LOW) (I)	M0(HIGH) (I)	M0(HIGH) (I)	M0(LOW) (I)	M0(LOW) (I)	(I)
HDC (HIGH)	HDC (HIGH)	HDC (HIGH)	HDC (HIGH)	HDC (HIGH)	HDC (HIGH)	I/O
LDC (LOW)	LDC (LOW)	LDC (LOW)	LDC (LOW)	LDC (LOW)	LDC (LOW)	I/O
INIT	INIT	INIT	INIT	INIT	INIT	I/O
DONE	DONE	DONE	DONE	DONE	DONE	DONE
PROGRAM (I)	PROGRAM (I)	PROGRAM (I)	PROGRAM (I)	PROGRAM (I)	PROGRAM (I)	PROGRAM
CCLK (I)	CCLK (O)	CCLK (I)	CCLK (O)	CCLK (O)	CCLK (O)	CCLK (I)
		RDY/BUSY (O)	RDY/BUSY (O)	RCLK (O)	RCLK (O)	I/O
			RS (I)			I/O
			CS0 (I)			I/O
		DATA 7 (I)	DATA 7 (I)	DATA 7 (I)	DATA 7 (I)	I/O
		DATA 6 (I)	DATA 6 (I)	DATA 6 (I)	DATA 6 (I)	I/O
		DATA 5 (I)	DATA 5 (I)	DATA 5 (I)	DATA 5 (I)	I/O
		DATA 4 (I)	DATA 4 (I)	DATA 4 (I)	DATA 4 (I)	I/O
		DATA 3 (I)	DATA 3 (I)	DATA 3 (I)	DATA 3 (I)	I/O
		DATA 2 (I)	DATA 2 (I)	DATA 2 (I)	DATA 2 (I)	I/O
		DATA 1 (I)	DATA 1 (I)	DATA 1 (I)	DATA 1 (I)	I/O
DIN (I)	DIN (I)	DATA 0 (I)	DATA 0 (I)	DATA 0 (I)	DATA 0 (I)	I/O
DOUT	DOUT	DOUT	DOUT	DOUT	DOUT	SGCK4-GCK6-I/O
TDI	TDI	TDI	TDI	TDI	TDI	TDI-I/O
TCK	TCK	TCK	TCK	TCK	TCK	TCK-I/O
TMS	TMS	TMS	TMS	TMS	TMS	TMS-I/O
TDO	TDO	TDO	TDO	TDO	TDO	TDO-(O)
			WS (I)	A0	A0	I/O
				A1	A1	PGCK4-GCK7-I/O
			CS1	A2	A2	I/O
				A3	A3	I/O
				A4	A4	I/O
				A5	A5	I/O
				A6	A6	I/O
				A7	A7	I/O
				A8	A8	I/O
				A9	A9	I/O
				A10	A10	I/O
				A11	A11	I/O
				A12	A12	I/O
				A13	A13	I/O
				A14	A14	I/O
				A15	A15	SGCK1-GCK8-I/O
				A16	A16	PGCK1-GCK1-I/O
				A17	A17	I/O
				A18*	A18*	I/O
				A19*	A19*	I/O
				A20*	A20*	I/O
				A21*	A21*	I/O
						ALL OTHERS

## Master Parallel Modes

In the two Master Parallel modes, the lead FPGA directly addresses an industry-standard byte-wide EPROM, and accepts eight data bits just before incrementing or decrementing the address outputs.

The eight data bits are serialized in the lead FPGA, which then presents the preamble data—and all data that overflows the lead device—on its DOUT pin. There is an internal delay of 1.5 CCLK periods, after the rising CCLK edge that accepts a byte of data (and also changes the EPROM address) until the falling CCLK edge that makes the LSB (D0) of this byte appear at DOUT. This means that DOUT changes on the falling CCLK edge, and the next FPGA in the daisy chain accepts data on the subsequent rising CCLK edge.

The PROM address pins can be incremented or decremented, depending on the mode pin settings. This option allows the FPGA to share the PROM with a wide variety of microprocessors and micro controllers. Some processors must boot from the bottom of memory (all zeros) while others must boot from the top. The FPGA is flexible and can load its configuration bitstream from either end of the memory.

Master Parallel Up mode is selected by a <100> on the mode pins (M2, M1, M0). The EPROM addresses start at 00000 and increment.

Master Parallel Down mode is selected by a <110> on the mode pins. The EPROM addresses start at 3FFFF and decrement.

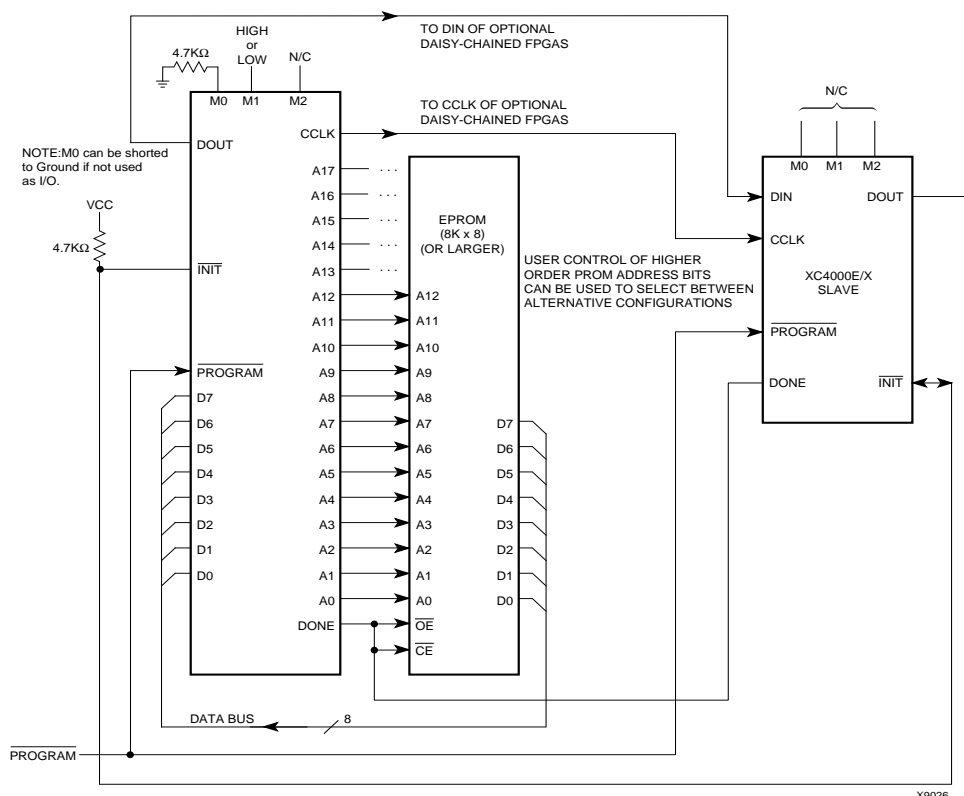
### Additional Address lines in XC4000 devices

The XC4000X devices have additional address lines (A18-A21) allowing the additional address space required to daisy-chain several large devices.

The extra address lines are programmable in XC4000EX devices. By default these address lines are not activated. In the default mode, the devices are compatible with existing XC4000 and XC4000E products. If desired, the extra address lines can be used by specifying the address lines option in bitgen as 22 (bitgen -g AddressLines:22). The lines (A18-A21) are driven when a master device detects, via the bitstream, that it should be using all 22 address lines. Because these pins will initially be pulled high by internal pull-ups, designers using Master Parallel Up mode should use external pull down resistors on pins A18-A21. If Master Parallel Down mode is used external resistors are not necessary.

All 22 address lines are always active in Master Parallel modes with XC4000XL devices. The additional address lines behave identically to the lower order address lines. If the Address Lines option in bitgen is set to 18, it will be ignored by the XC4000XL device.

The additional address lines (A18-A21) are not available in the PC84 package.



**Figure 54: Master Parallel Mode Circuit Diagram**

## Synchronous Peripheral Mode

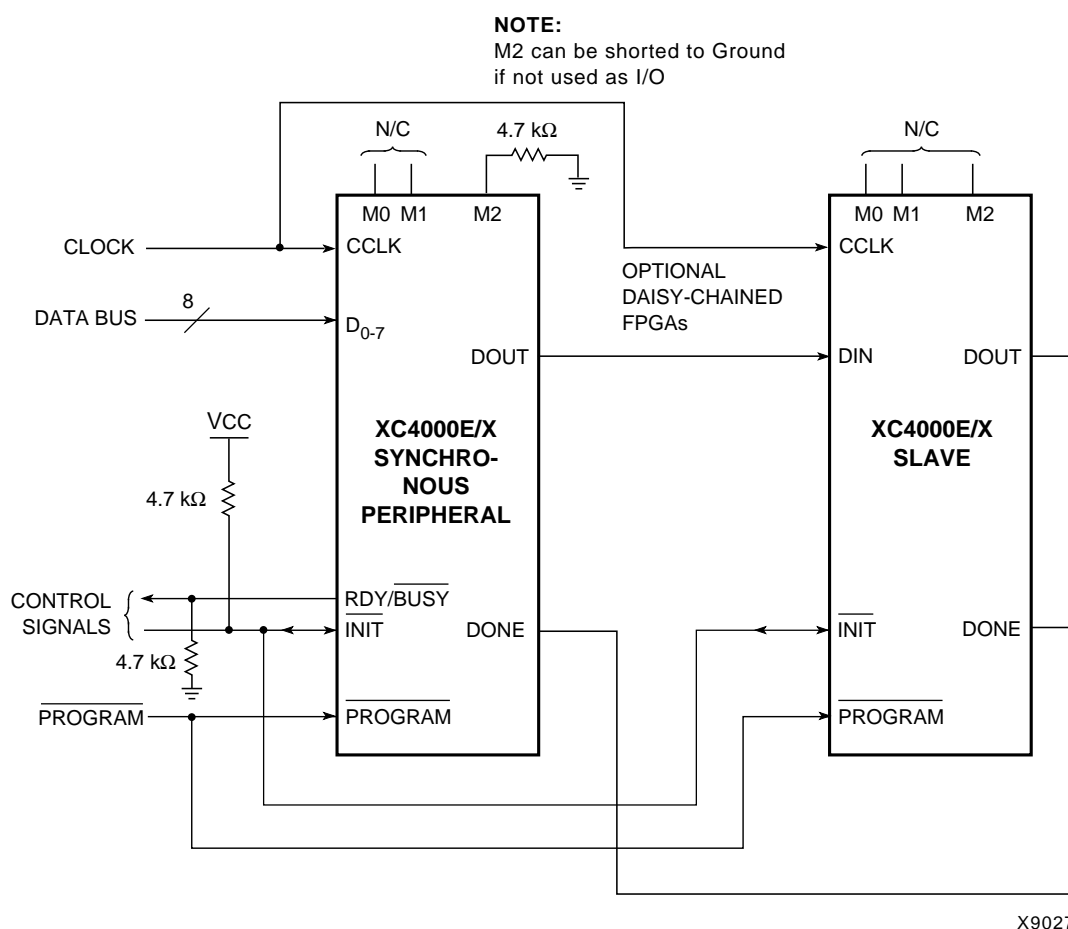
Synchronous Peripheral mode can also be considered Slave Parallel mode. An external signal drives the CCLK input(s) of the FPGA(s). The first byte of parallel configuration data must be available at the Data inputs of the lead FPGA a short setup time before the rising CCLK edge. Subsequent data bytes are clocked in on every eighth consecutive rising CCLK edge.

The same CCLK edge that accepts data, also causes the RDY/ $\overline{\text{BUSY}}$  output to go High for one CCLK period. The pin name is a misnomer. In Synchronous Peripheral mode it is really an ACKNOWLEDGE signal. Synchronous operation does not require this response, but it is a meaningful signal for test purposes. Note that RDY/ $\overline{\text{BUSY}}$  is pulled High with a high-impedance pullup prior to  $\overline{\text{INIT}}$  going High.

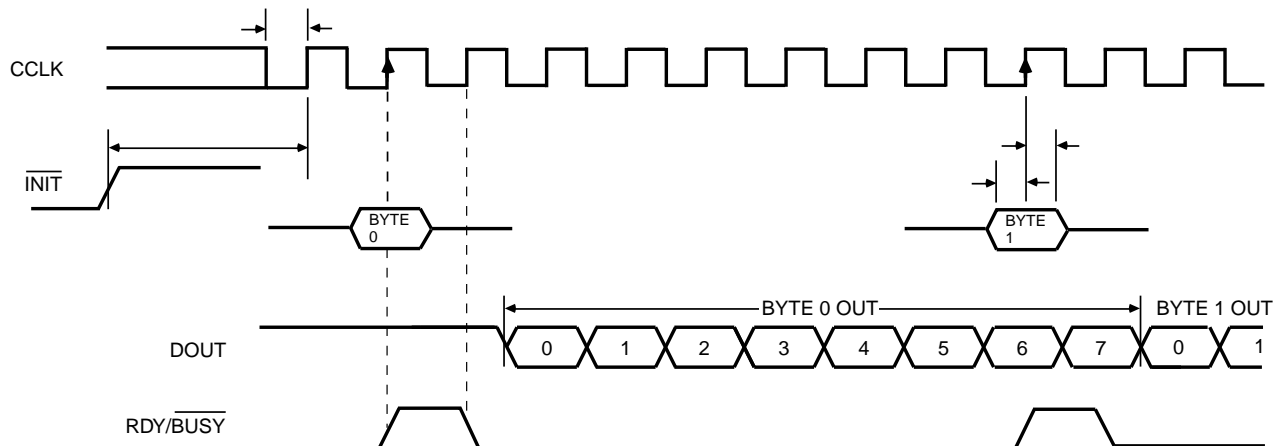
The lead FPGA serializes the data and presents the preamble data (and all data that overflows the lead device) on its DOUT pin. There is an internal delay of 1.5 CCLK periods, which means that DOUT changes on the falling CCLK edge, and the next FPGA in the daisy chain accepts data on the subsequent rising CCLK edge.

In order to complete the serial shift operation, 10 additional CCLK rising edges are required after the last data byte has been loaded, plus one more CCLK cycle for each daisy-chained device.

Synchronous Peripheral mode is selected by a <011> on the mode pins (M2, M1, M0).



**Figure 56: Synchronous Peripheral Mode Circuit Diagram**



X6096

	Description	Symbol	Min	Max	Units
CCLK	INIT (High) setup time	$T_{IC}$	5		$\mu s$
	D0 - D7 setup time	$T_{DC}$	60		ns
	D0 - D7 hold time	$T_{CD}$	0		ns
	CCLK High time	$T_{CCH}$	50		ns
	CCLK Low time	$T_{CCL}$	60		ns
	CCLK Frequency	$F_{CC}$		8	MHz

- Notes:
1. Peripheral Synchronous mode can be considered Slave Parallel mode. An external CCLK provides timing, clocking in the **first** data byte on the **second** rising edge of CCLK after INIT goes High. Subsequent data bytes are clocked in on every eighth consecutive rising edge of CCLK.
  2. The RDY/BUSY line goes High for one CCLK period after data has been clocked in, although synchronous operation does not require such a response.
  3. The pin name RDY/BUSY is a misnomer. In Synchronous Peripheral mode this is really an ACKNOWLEDGE signal.
  4. Note that data starts to shift out serially on the DOUT pin 0.5 CCLK periods after it was loaded in parallel. Therefore, additional CCLK pulses are clearly required after the last byte has been loaded.

Figure 57: Synchronous Peripheral Mode Programming Switching Characteristics

## XC4000 Series Electrical Characteristics and Device-Specific Pinout Table

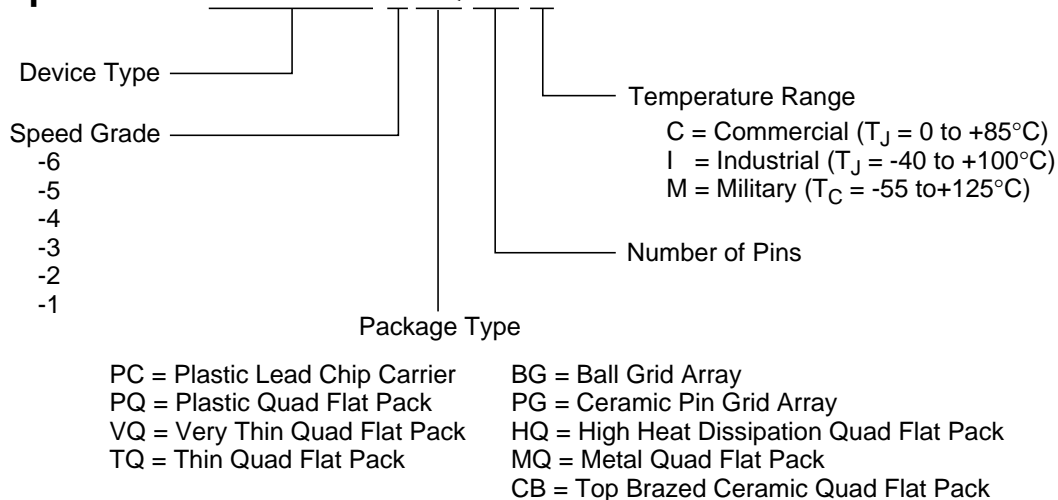
For the latest Electrical Characteristics and package/pinout information for each XC4000 Family, see the Xilinx web site at

[http://www.xilinx.com/xlnx/xweb/xil\\_publications\\_index.jsp](http://www.xilinx.com/xlnx/xweb/xil_publications_index.jsp)

## Ordering Information

### Example:

# XC4013E-3HQ240C



X9020

## Revision Control

Version	Description
3/30/98 (1.5)	Updated XC4000XL timing and added XC4002XL
1/29/99 (1.5)	Updated pin diagrams
5/14/99 (1.6)	Replaced Electrical Specification and pinout pages for E, EX, and XL families with separate updates and added URL link for electrical specifications/pinouts for Web users