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Understanding [Embedded - FPGAs \(Field Programmable Gate Array\)](#)

Embedded - FPGAs, or Field Programmable Gate Arrays, are advanced integrated circuits that offer unparalleled flexibility and performance for digital systems. Unlike traditional fixed-function logic devices, FPGAs can be programmed and reprogrammed to execute a wide array of logical operations, enabling customized functionality tailored to specific applications. This reprogrammability allows developers to iterate designs quickly and implement complex functions without the need for custom hardware.

Applications of Embedded - FPGAs

The versatility of Embedded - FPGAs makes them indispensable in numerous fields. In telecommunications.

Details

Product Status	Obsolete
Number of LABs/CLBs	784
Number of Logic Elements/Cells	1862
Total RAM Bits	25088
Number of I/O	145
Number of Gates	20000
Voltage - Supply	3V ~ 3.6V
Mounting Type	Surface Mount
Operating Temperature	-40°C ~ 100°C (TJ)
Package / Case	176-LQFP Exposed Pad
Supplier Device Package	176-TQFP (24x24)
Purchase URL	https://www.e-xfl.com/product-detail/xilinx/xc4020xl-3ht176i

Set/Reset

An asynchronous storage element input (SR) can be configured as either set or reset. This configuration option determines the state in which each flip-flop becomes operational after configuration. It also determines the effect of a Global Set/Reset pulse during normal operation, and the effect of a pulse on the SR pin of the CLB. All three set/reset functions for any single flip-flop are controlled by the same configuration data bit.

The set/reset state can be independently specified for each flip-flop. This input can also be independently disabled for either flip-flop.

The set/reset state is specified by using the INIT attribute, or by placing the appropriate set or reset flip-flop library symbol.

SR is active High. It is not invertible within the CLB.

Global Set/Reset

A separate Global Set/Reset line (not shown in Figure 1) sets or clears each storage element during power-up, re-configuration, or when a dedicated Reset net is driven active. This global net (GSR) does not compete with other routing resources; it uses a dedicated distribution network.

Each flip-flop is configured as either globally set or reset in the same way that the local set/reset (SR) is specified. Therefore, if a flip-flop is set by SR, it is also set by GSR. Similarly, a reset flip-flop is reset by both SR and GSR.

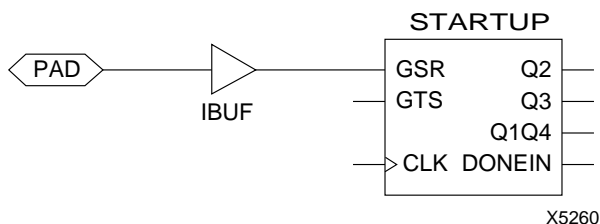


Figure 2: Schematic Symbols for Global Set/Reset

GSR can be driven from any user-programmable pin as a global reset input. To use this global net, place an input pad and input buffer in the schematic or HDL code, driving the GSR pin of the STARTUP symbol. (See Figure 2.) A specific pin location can be assigned to this input using a LOC attribute or property, just as with any other user-programmable pad. An inverter can optionally be inserted after the input buffer to invert the sense of the Global Set/Reset signal.

Alternatively, GSR can be driven from any internal node.

Data Inputs and Outputs

The source of a storage element data input is programmable. It is driven by any of the functions F', G', and H', or by the Direct In (DIN) block input. The flip-flops or latches drive the XQ and YQ CLB outputs.

Two fast feed-through paths are available, as shown in Figure 1. A two-to-one multiplexer on each of the XQ and YQ outputs selects between a storage element output and any of the control inputs. This bypass is sometimes used by the automated router to repower internal signals.

Control Signals

Multiplexers in the CLB map the four control inputs (C1 - C4 in Figure 1) into the four internal control signals (H1, DIN/H2, SR/H0, and EC). Any of these inputs can drive any of the four internal control signals.

When the logic function is enabled, the four inputs are:

- EC — Enable Clock
- SR/H0 — Asynchronous Set/Reset or H function generator Input 0
- DIN/H2 — Direct In or H function generator Input 2
- H1 — H function generator Input 1.

When the memory function is enabled, the four inputs are:

- EC — Enable Clock
- WE — Write Enable
- D0 — Data Input to F and/or G function generator
- D1 — Data input to G function generator (16x1 and 16x2 modes) or 5th Address bit (32x1 mode).

Using FPGA Flip-Flops and Latches

The abundance of flip-flops in the XC4000 Series invites pipelined designs. This is a powerful way of increasing performance by breaking the function into smaller subfunctions and executing them in parallel, passing on the results through pipeline flip-flops. This method should be seriously considered wherever throughput is more important than latency.

To include a CLB flip-flop, place the appropriate library symbol. For example, FDCE is a D-type flip-flop with clock enable and asynchronous clear. The corresponding latch symbol (for the XC4000X only) is called LDCE.

In XC4000 Series devices, the flip flops can be used as registers or shift registers without blocking the function generators from performing a different, perhaps unrelated task. This ability increases the functional capacity of the devices.

The CLB setup time is specified between the function generator inputs and the clock input K. Therefore, the specified CLB flip-flop setup time includes the delay through the function generator.

Using Function Generators as RAM

Optional modes for each CLB make the memory look-up tables in the F' and G' function generators usable as an array of Read/Write memory cells. Available modes are level-sensitive (similar to the XC4000/A/H families), edge-triggered, and dual-port edge-triggered. Depending on the selected mode, a single CLB can be configured as either a 16x2, 32x1, or 16x1 bit array.

tions of the CLB, with the exception of the redefinition of the control signals. In 16x2 and 16x1 modes, the H' function generator can be used to implement Boolean functions of F', G', and D1, and the D flip-flops can latch the F', G', H', or D0 signals.

Single-Port Edge-Triggered Mode

Edge-triggered (synchronous) RAM simplifies timing requirements. XC4000 Series edge-triggered RAM timing operates like writing to a data register. Data and address are presented. The register is enabled for writing by a logic High on the write enable input, WE. Then a rising or falling clock edge loads the data into the register, as shown in Figure 3.

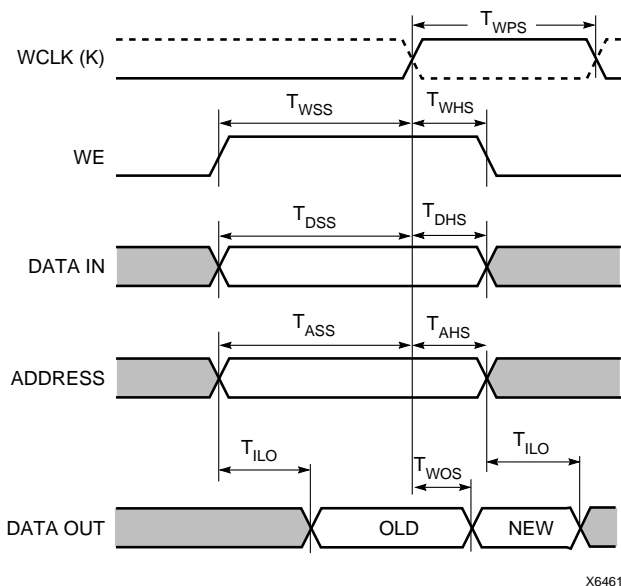


Figure 3: Edge-Triggered RAM Write Timing

Complex timing relationships between address, data, and write enable signals are not required, and the external write enable pulse becomes a simple clock enable. The active edge of WCLK latches the address, input data, and WE sig-

nals. An internal write pulse is generated that performs the write. See Figure 4 and Figure 5 for block diagrams of a CLB configured as 16x2 and 32x1 edge-triggered, single-port RAM.

The relationships between CLB pins and RAM inputs and outputs for single-port, edge-triggered mode are shown in Table 5.

The Write Clock input (WCLK) can be configured as active on either the rising edge (default) or the falling edge. It uses the same CLB pin (K) used to clock the CLB flip-flops, but it can be independently inverted. Consequently, the RAM output can optionally be registered within the same CLB either by the same clock edge as the RAM, or by the opposite edge of this clock. The sense of WCLK applies to both function generators in the CLB when both are configured as RAM.

The WE pin is active-High and is not invertible within the CLB.

Note: The pulse following the active edge of WCLK (T_{WPS} in Figure 3) must be less than one millisecond wide. For most applications, this requirement is not overly restrictive; however, it must not be forgotten. Stopping WCLK at this point in the write cycle could result in excessive current and even damage to the larger devices if many CLBs are configured as edge-triggered RAM.

Table 5: Single-Port Edge-Triggered RAM Signals

RAM Signal	CLB Pin	Function
D	D0 or D1 (16x2, 16x1), D0 (32x1)	Data In
A[3:0]	F1-F4 or G1-G4	Address
A[4]	D1 (32x1)	Address
WE	WE	Write Enable
WCLK	K	Clock
SPO (Data Out)	F' or G'	Single Port Out (Data Out)

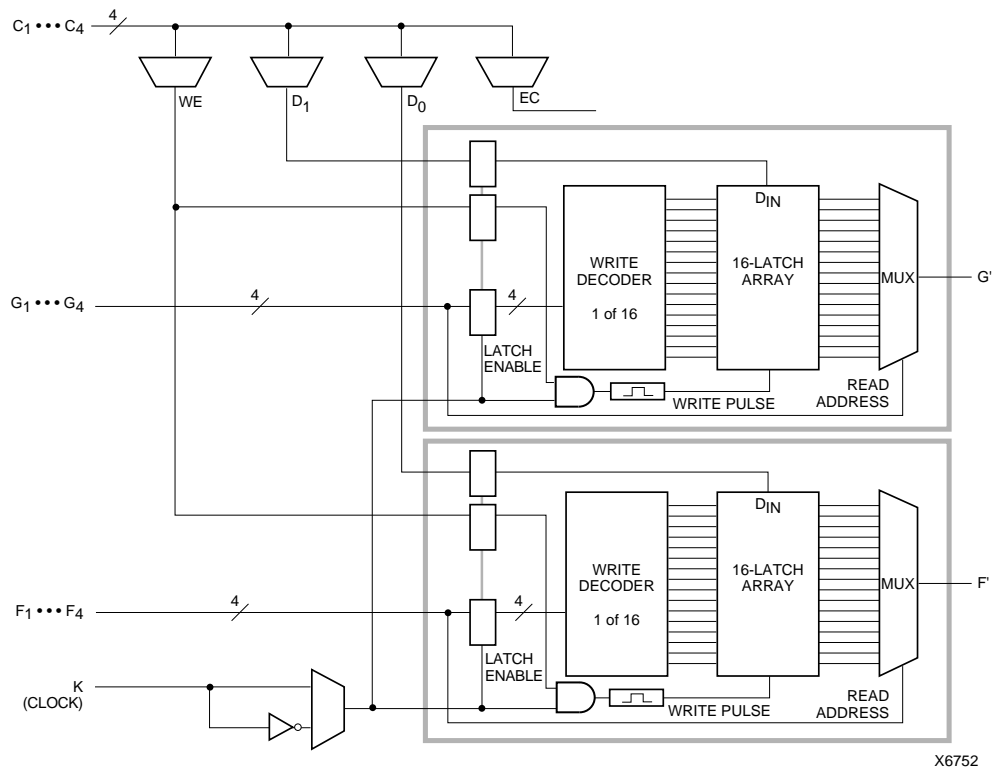


Figure 4: 16x2 (or 16x1) Edge-Triggered Single-Port RAM

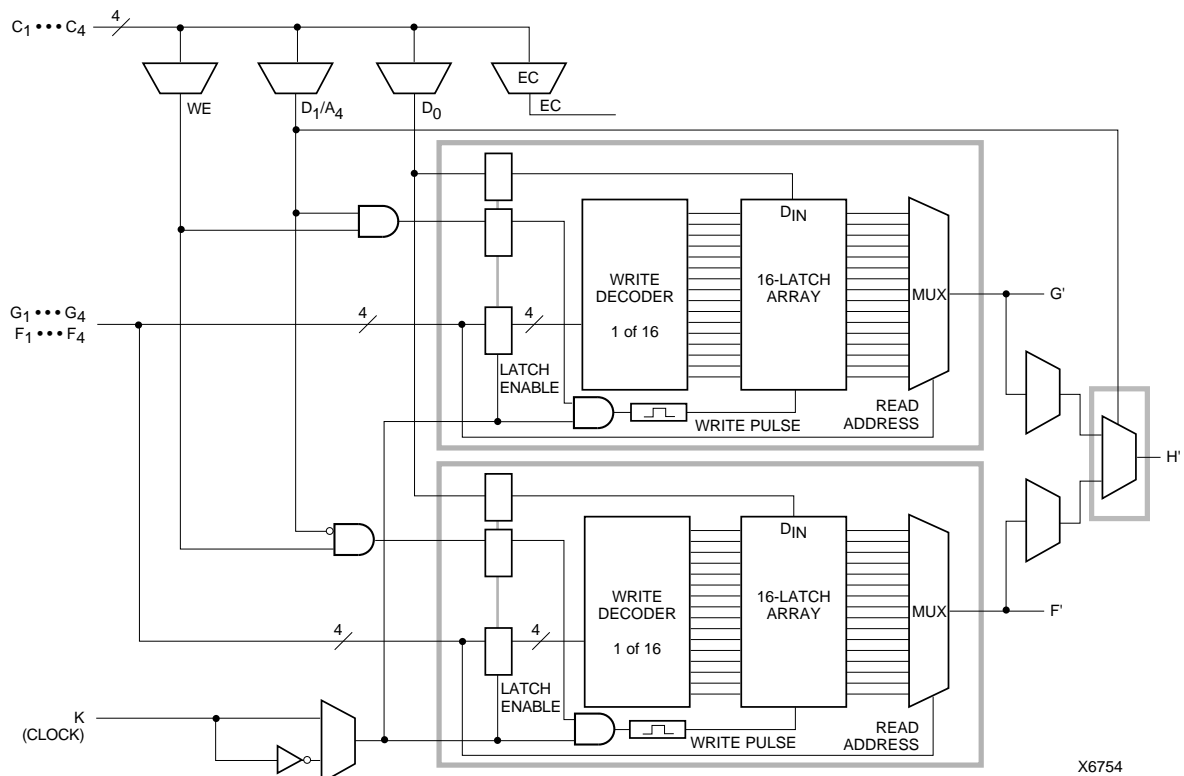


Figure 5: 32x1 Edge-Triggered Single-Port RAM (F and G addresses are identical)

Dual-Port Edge-Triggered Mode

In dual-port mode, both the F and G function generators are used to create a single 16x1 RAM array with one write port and two read ports. The resulting RAM array can be read and written simultaneously at two independent addresses. Simultaneous read and write operations at the same address are also supported.

Dual-port mode always has edge-triggered write timing, as shown in [Figure 3](#).

[Figure 6](#) shows a simple model of an XC4000 Series CLB configured as dual-port RAM. One address port, labeled A[3:0], supplies both the read and write address for the F function generator. This function generator behaves the same as a 16x1 single-port edge-triggered RAM array. The RAM output, Single Port Out (SPO), appears at the F function generator output. SPO, therefore, reflects the data at address A[3:0].

The other address port, labeled DPRA[3:0] for Dual Port Read Address, supplies the read address for the G function generator. The write address for the G function generator, however, comes from the address A[3:0]. The output from this 16x1 RAM array, Dual Port Out (DPO), appears at the G function generator output. DPO, therefore, reflects the data at address DPRA[3:0].

Therefore, by using A[3:0] for the write address and DPRA[3:0] for the read address, and reading only the DPO output, a FIFO that can read and write simultaneously is easily generated. Simultaneous access doubles the effective throughput of the FIFO.

The relationships between CLB pins and RAM inputs and outputs for dual-port, edge-triggered mode are shown in [Table 6](#). See [Figure 7 on page 16](#) for a block diagram of a CLB configured in this mode.

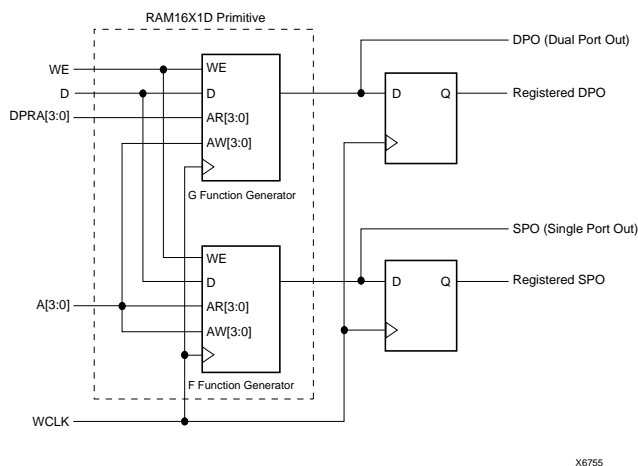


Figure 6: XC4000 Series Dual-Port RAM, Simple Model

Table 6: Dual-Port Edge-Triggered RAM Signals

RAM Signal	CLB Pin	Function
D	D0	Data In
A[3:0]	F1-F4	Read Address for F, Write Address for F and G
DPRA[3:0]	G1-G4	Read Address for G
WE	WE	Write Enable
WCLK	K	Clock
SPO	F'	Single Port Out (addressed by A[3:0])
DPO	G'	Dual Port Out (addressed by DPRA[3:0])

Note: The pulse following the active edge of WCLK (T_{WPS} in [Figure 3](#)) must be less than one millisecond wide. For most applications, this requirement is not overly restrictive; however, it must not be forgotten. Stopping WCLK at this point in the write cycle could result in excessive current and even damage to the larger devices if many CLBs are configured as edge-triggered RAM.

Single-Port Level-Sensitive Timing Mode

Note: Edge-triggered mode is recommended for all new designs. Level-sensitive mode, also called asynchronous mode, is still supported for XC4000 Series backward-compatibility with the XC4000 family.

Level-sensitive RAM timing is simple in concept but can be complicated in execution. Data and address signals are presented, then a positive pulse on the write enable pin (WE) performs a write into the RAM at the designated address. As indicated by the “level-sensitive” label, this RAM acts like a latch. During the WE High pulse, changing the data lines results in new data written to the old address. Changing the address lines while WE is High results in spurious data written to the new address—and possibly at other addresses as well, as the address lines inevitably do not all change simultaneously.

The user must generate a carefully timed WE signal. The delay on the WE signal and the address lines must be carefully verified to ensure that WE does not become active until after the address lines have settled, and that WE goes inactive before the address lines change again. The data must be stable before and after the falling edge of WE.

In practical terms, WE is usually generated by a 2X clock. If a 2X clock is not available, the falling edge of the system clock can be used. However, there are inherent risks in this approach, since the WE pulse must be guaranteed inactive before the next rising edge of the system clock. Several older application notes are available from Xilinx that discuss the design of level-sensitive RAMs.

However, the edge-triggered RAM available in the XC4000 Series is superior to level-sensitive RAM for almost every application.

Table 8: Supported Sources for XC4000 Series Device Inputs

Source	XC4000E/EX Series Inputs		XC4000XL Series Inputs
	5 V, TTL	5 V, CMOS	3.3 V CMOS
Any device, V _{CC} = 3.3 V, CMOS outputs	✓	Unreliable Data	✓
XC4000 Series, V _{CC} = 5 V, TTL outputs	✓		✓
Any device, V _{CC} = 5 V, TTL outputs (V _{oh} ≤ 3.7 V)	✓		✓
Any device, V _{CC} = 5 V, CMOS outputs	✓	✓	✓

XC4000XL 5-Volt Tolerant I/Os

The I/Os on the XC4000XL are fully 5-volt tolerant even though the V_{CC} is 3.3 volts. This allows 5 V signals to directly connect to the XC4000XL inputs without damage, as shown in [Table 8](#). In addition, the 3.3 volt V_{CC} can be applied before or after 5 volt signals are applied to the I/Os. This makes the XC4000XL immune to power supply sequencing problems.


Registered Inputs

The I1 and I2 signals that exit the block can each carry either the direct or registered input signal.

The input and output storage elements in each IOB have a common clock enable input, which, through configuration, can be activated individually for the input or output flip-flop, or both. This clock enable operates exactly like the EC pin on the XC4000 Series CLB. It cannot be inverted within the IOB.

The storage element behavior is shown in [Table 9](#).

Table 9: Input Register Functionality (active rising edge is shown)

Mode	Clock	Clock Enable	D	Q
Power-Up or GSR	X	X	X	SR
Flip-Flop		1*	D	D
	0	X	X	Q
Latch	1	1*	X	Q
	0	1*	D	D
Both	X	0	X	Q

Legend:

X 

SR

0*

1*

Don't care
Rising edge

Set or Reset value. Reset is default.

Input is Low or unconnected (default value)

Input is High or unconnected (default value)

Optional Delay Guarantees Zero Hold Time

The data input to the register can optionally be delayed by several nanoseconds. With the delay enabled, the setup time of the input flip-flop is increased so that normal clock routing does not result in a positive hold-time requirement. A positive hold time requirement can lead to unreliable, temperature- or processing-dependent operation.

The input flip-flop setup time is defined between the data measured at the device I/O pin and the clock input at the IOB (not at the clock pin). Any routing delay from the device clock pin to the clock input of the IOB must, therefore, be subtracted from this setup time to arrive at the real setup time requirement relative to the device pins. A short specified setup time might, therefore, result in a negative setup time at the device pins, i.e., a positive hold-time requirement.

When a delay is inserted on the data line, more clock delay can be tolerated without causing a positive hold-time requirement. Sufficient delay eliminates the possibility of a data hold-time requirement at the external pin. The maximum delay is therefore inserted as the default.

The XC4000E IOB has a one-tap delay element: either the delay is inserted (default), or it is not. The delay guarantees a zero hold time with respect to clocks routed through any of the XC4000E global clock buffers. (See [“Global Nets and Buffers \(XC4000E only\)” on page 35](#) for a description of the global clock buffers in the XC4000E.) For a shorter input register setup time, with non-zero hold, attach a NODELAY attribute or property to the flip-flop.

The XC4000X IOB has a two-tap delay element, with choices of a full delay, a partial delay, or no delay. The attributes or properties used to select the desired delay are shown in [Table 10](#). The choices are no added attribute, MEDDELAY, and NODELAY. The default setting, with no added attribute, ensures no hold time with respect to any of the XC4000X clock buffers, including the Global Low-Skew buffers. MEDDELAY ensures no hold time with respect to the Global Early buffers. Inputs with NODELAY may have a positive hold time with respect to all clock buffers. For a description of each of these buffers, see [“Global Nets and Buffers \(XC4000X only\)” on page 37](#).

Table 10: XC4000X IOB Input Delay Element

Value	When to Use
full delay (default, no attribute added)	Zero Hold with respect to Global Low-Skew Buffer, Global Early Buffer
MEDDELAY	Zero Hold with respect to Global Early Buffer
NODELAY	Short Setup, positive Hold time

Additional Input Latch for Fast Capture (XC4000X only)

The XC4000X IOB has an additional optional latch on the input. This latch, as shown in [Figure 16](#), is clocked by the output clock — the clock used for the output flip-flop — rather than the input clock. Therefore, two different clocks can be used to clock the two input storage elements. This additional latch allows the very fast capture of input data, which is then synchronized to the internal clock by the IOB flip-flop or latch.

To use this Fast Capture technique, drive the output clock pin (the Fast Capture latching signal) from the output of one of the Global Early buffers supplied in the XC4000X. The second storage element should be clocked by a Global Low-Skew buffer, to synchronize the incoming data to the internal logic. (See [Figure 17](#).) These special buffers are described in “Global Nets and Buffers (XC4000X only)” on [page 37](#).

The Fast Capture latch (FCL) is designed primarily for use with a Global Early buffer. For Fast Capture, a single clock signal is routed through both a Global Early buffer and a Global Low-Skew buffer. (The two buffers share an input pad.) The Fast Capture latch is clocked by the Global Early buffer, and the standard IOB flip-flop or latch is clocked by the Global Low-Skew buffer. This mode is the safest way to use the Fast Capture latch, because the clock buffers on both storage elements are driven by the same pad. There is no external skew between clock pads to create potential problems.

To place the Fast Capture latch in a design, use one of the special library symbols, ILFFX or ILFLX. ILFFX is a transparent-Low Fast Capture latch followed by an active-High input flip-flop. ILFLX is a transparent-Low Fast Capture latch followed by a transparent-High input latch. Any of the clock inputs can be inverted before driving the library element, and the inverter is absorbed into the IOB. If a single BUFG output is used to drive both clock inputs, the software automatically runs the clock through both a Global Low-Skew buffer and a Global Early buffer, and clocks the Fast Capture latch appropriately.

[Figure 16 on page 21](#) also shows a two-tap delay on the input. By default, if the Fast Capture latch is used, the Xilinx software assumes a Global Early buffer is driving the clock, and selects MEDDELAY to ensure a zero hold time. Select

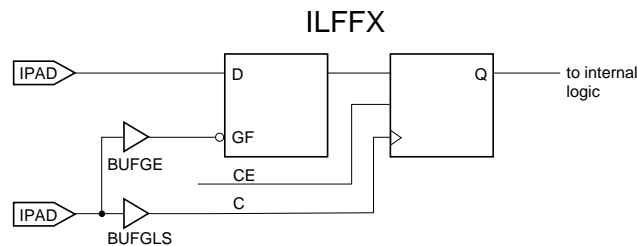


Figure 17: Examples Using XC4000X FCL

the desired delay based on the discussion in the previous subsection.

IOB Output Signals

Output signals can be optionally inverted within the IOB, and can pass directly to the pad or be stored in an edge-triggered flip-flop. The functionality of this flip-flop is shown in [Table 11](#).

An active-High 3-state signal can be used to place the output buffer in a high-impedance state, implementing 3-state outputs or bidirectional I/O. Under configuration control, the output (OUT) and output 3-state (T) signals can be inverted. The polarity of these signals is independently configured for each IOB.

The 4-mA maximum output current specification of many FPGAs often forces the user to add external buffers, which are especially cumbersome on bidirectional I/O lines. The XC4000E and XC4000EX/XL devices solve many of these problems by providing a guaranteed output sink current of 12 mA. Two adjacent outputs can be interconnected externally to sink up to 24 mA. The XC4000E and XC4000EX/XL FPGAs can thus directly drive buses on a printed circuit board.

By default, the output pull-up structure is configured as a TTL-like totem-pole. The High driver is an n-channel pull-up transistor, pulling to a voltage one transistor threshold below Vcc. Alternatively, the outputs can be globally configured as CMOS drivers, with p-channel pull-up transistors pulling to Vcc. This option, applied using the bitstream generation software, applies to all outputs on the device. It is not individually programmable. In the XC4000XL, all outputs are pulled to the positive supply rail.

Table 11: Output Flip-Flop Functionality (active rising edge is shown)

Mode	Clock	Clock Enable	T	D	Q
Power-Up or GSR	X	X	0*	X	SR
Flip-Flop	X	0	0*	X	Q
		1*	0*	D	D
	X	X	1	X	Z
	0	X	0*	X	Q

Legend:

X

Don't care

Rising edge

SR

Set or Reset value. Reset is default.

0*

Input is Low or unconnected (default value)

1*

Input is High or unconnected (default value)

Z

3-state

The oscillator output is optionally available after configuration. Any two of four resynchronized taps of a built-in divider are also available. These taps are at the fourth, ninth, fourteenth and nineteenth bits of the divider. Therefore, if the primary oscillator output is running at the nominal 8 MHz, the user has access to an 8 MHz clock, plus any two of 500 kHz, 16kHz, 490Hz and 15Hz (up to 10% lower for low-voltage devices). These frequencies can vary by as much as -50% or +25%.

These signals can be accessed by placing the OSC4 library element in a schematic or in HDL code (see [Figure 24](#)).

The oscillator is automatically disabled after configuration if the OSC4 symbol is not used in the design.

Programmable Interconnect

All internal connections are composed of metal segments with programmable switching points and switching matrices to implement the desired routing. A structured, hierarchical matrix of routing resources is provided to achieve efficient automated routing.

The XC4000E and XC4000X share a basic interconnect structure. XC4000X devices, however, have additional routing not available in the XC4000E. The extra routing resources allow high utilization in high-capacity devices. All XC4000X-specific routing resources are clearly identified throughout this section. Any resources not identified as XC4000X-specific are present in all XC4000 Series devices.

This section describes the varied routing resources available in XC4000 Series devices. The implementation software automatically assigns the appropriate resources based on the density and timing requirements of the design.

Interconnect Overview

There are several types of interconnect.

- CLB routing is associated with each row and column of the CLB array.
- IOB routing forms a ring (called a VersaRing) around the outside of the CLB array. It connects the I/O with the internal logic blocks.

- Global routing consists of dedicated networks primarily designed to distribute clocks throughout the device with minimum delay and skew. Global routing can also be used for other high-fanout signals.

Five interconnect types are distinguished by the relative length of their segments: single-length lines, double-length lines, quad and octal lines (XC4000X only), and longlines. In the XC4000X, direct connects allow fast data flow between adjacent CLBs, and between IOBs and CLBs.

Extra routing is included in the IOB pad ring. The XC4000X also includes a ring of octal interconnect lines near the IOBs to improve pin-swapping and routing to locked pins.

XC4000E/X devices include two types of global buffers. These global buffers have different properties, and are intended for different purposes. They are discussed in detail later in this section.

CLB Routing Connections

A high-level diagram of the routing resources associated with one CLB is shown in [Figure 25](#). The shaded arrows represent routing present only in XC4000X devices.

[Table 14](#) shows how much routing of each type is available in XC4000E and XC4000X CLB arrays. Clearly, very large designs, or designs with a great deal of interconnect, will route more easily in the XC4000X. Smaller XC4000E designs, typically requiring significantly less interconnect, do not require the additional routing.

[Figure 27 on page 30](#) is a detailed diagram of both the XC4000E and the XC4000X CLB, with associated routing. The shaded square is the programmable switch matrix, present in both the XC4000E and the XC4000X. The L-shaped shaded area is present only in XC4000X devices. As shown in the figure, the XC4000X block is essentially an XC4000E block with additional routing.

CLB inputs and outputs are distributed on all four sides, providing maximum routing flexibility. In general, the entire architecture is symmetrical and regular. It is well suited to established placement and routing algorithms. Inputs, outputs, and function generators can freely swap positions within a CLB to avoid routing congestion during the placement and routing operation.



Figure 28: Single- and Double-Length Lines, with Programmable Switch Matrices (PSMs)

Double-Length Lines

The double-length lines consist of a grid of metal segments, each twice as long as the single-length lines: they run past two CLBs before entering a switch matrix. Double-length lines are grouped in pairs with the switch matrices staggered, so that each line goes through a switch matrix at every other row or column of CLBs (see [Figure 28](#)).

There are four vertical and four horizontal double-length lines associated with each CLB. These lines provide faster signal routing over intermediate distances, while retaining routing flexibility. Double-length lines are connected by way of the programmable switch matrices. Routing connectivity is shown in [Figure 27](#).

Quad Lines (XC4000X only)

XC4000X devices also include twelve vertical and twelve horizontal quad lines per CLB row and column. Quad lines are four times as long as the single-length lines. They are interconnected via buffered switch matrices (shown as diamonds in [Figure 27 on page 30](#)). Quad lines run past four CLBs before entering a buffered switch matrix. They are grouped in fours, with the buffered switch matrices staggered, so that each line goes through a buffered switch matrix at every fourth CLB location in that row or column. (See [Figure 29](#).)

The buffered switch matrixes have four pins, one on each edge. All of the pins are bidirectional. Any pin can drive any or all of the other pins.

Each buffered switch matrix contains one buffer and six pass transistors. It resembles the programmable switch matrix shown in [Figure 26](#), with the addition of a programmable buffer. There can be up to two independent inputs

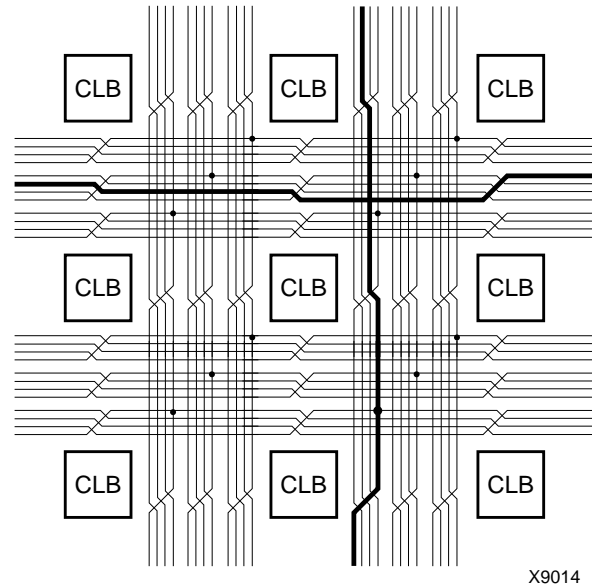


Figure 29: Quad Lines (XC4000X only)

and up to two independent outputs. Only one of the independent inputs can be buffered.

The place and route software automatically uses the timing requirements of the design to determine whether or not a quad line signal should be buffered. A heavily loaded signal is typically buffered, while a lightly loaded one is not. One scenario is to alternate buffers and pass transistors. This allows both vertical and horizontal quad lines to be buffered at alternating buffered switch matrices.

Due to the buffered switch matrices, quad lines are very fast. They provide the fastest available method of routing heavily loaded signals for long distances across the device.

Longlines

Longlines form a grid of metal interconnect segments that run the entire length or width of the array. Longlines are intended for high fan-out, time-critical signal nets, or nets that are distributed over long distances. In XC4000X devices, quad lines are preferred for critical nets, because the buffered switch matrices make them faster for high fan-out nets.

Two horizontal longlines per CLB can be driven by 3-state or open-drain drivers (TBUFs). They can therefore implement unidirectional or bidirectional buses, wide multiplexers, or wired-AND functions. (See ["Three-State Buffers" on page 26](#) for more details.)

Each horizontal longline driven by TBUFs has either two (XC4000E) or eight (XC4000X) pull-up resistors. To activate these resistors, attach a PULLUP symbol to the long-line net. The software automatically activates the appropriate number of pull-ups. There is also a weak keeper at each end of these two horizontal longlines. This

IOB inputs and outputs interface with the octal lines via the single-length interconnect lines. Single-length lines are also used for communication between the octals and double-length lines, quads, and longlines within the CLB array.

Segmentation into buffered octals was found to be optimal for distributing signals over long distances around the device.

Global Nets and Buffers

Both the XC4000E and the XC4000X have dedicated global networks. These networks are designed to distribute clocks and other high fanout control signals throughout the devices with minimal skew. The global buffers are described in detail in the following sections. The text descriptions and diagrams are summarized in [Table 15](#). The table shows which CLB and IOB clock pins can be sourced by which global buffers.

In both XC4000E and XC4000X devices, placement of a library symbol called BUFG results in the software choosing the appropriate clock buffer, based on the timing requirements of the design. The detailed information in these sections is included only for reference.

Global Nets and Buffers (XC4000E only)

Four vertical longlines in each CLB column are driven exclusively by special global buffers. These longlines are in addition to the vertical longlines used for standard interconnect. The four global lines can be driven by either of two types of global buffers. The clock pins of every CLB and IOB can also be sourced from local interconnect.

Two different types of clock buffers are available in the XC4000E:

- Primary Global Buffers (BUFGP)
- Secondary Global Buffers (BUFGS)

Four Primary Global buffers offer the shortest delay and negligible skew. Four Secondary Global buffers have slightly longer delay and slightly more skew due to potentially heavier loading, but offer greater flexibility when used to drive non-clock CLB inputs.

The Primary Global buffers must be driven by the semi-dedicated pads. The Secondary Global buffers can be sourced by either semi-dedicated pads or internal nets.

Each CLB column has four dedicated vertical Global lines. Each of these lines can be accessed by one particular Primary Global buffer, or by any of the Secondary Global buffers, as shown in [Figure 34](#). Each corner of the device has one Primary buffer and one Secondary buffer.

IOBs along the left and right edges have four vertical global longlines. Top and bottom IOBs can be clocked from the global lines in the adjacent CLB column.

A global buffer should be specified for all timing-sensitive global signal distribution. To use a global buffer, place a BUFGP (primary buffer), BUFGS (secondary buffer), or BUFG (either primary or secondary buffer) element in a schematic or in HDL code. If desired, attach a LOC attribute or property to direct placement to the designated location. For example, attach a LOC=L attribute or property to a BUFGS symbol to direct that a buffer be placed in one of the two Secondary Global buffers on the left edge of the device, or a LOC=BL to indicate the Secondary Global buffer on the bottom edge of the device, on the left.

Table 15: Clock Pin Access

	XC4000E		XC4000X			Local Inter-connect
	BUFGP	BUFGS	BUFGLS	L & R BUFGE	T & B BUFGE	
All CLBs in Quadrant	√	√	√	√	√	√
All CLBs in Device	√	√	√			√
IOBs on Adjacent Vertical Half Edge	√	√	√	√	√	√
IOBs on Adjacent Vertical Full Edge	√	√	√	√		√
IOBs on Adjacent Horizontal Half Edge (Direct)				√		√
IOBs on Adjacent Horizontal Half Edge (through CLB globals)	√	√	√	√	√	√
IOBs on Adjacent Horizontal Full Edge (through CLB globals)	√	√	√			√

L = Left, R = Right, T = Top, B = Bottom

Global Nets and Buffers (XC4000X only)

Eight vertical longlines in each CLB column are driven by special global buffers. These longlines are in addition to the vertical longlines used for standard interconnect. The global lines are broken in the center of the array, to allow faster distribution and to minimize skew across the whole array. Each half-column global line has its own buffered multiplexer, as shown in [Figure 35](#). The top and bottom global lines cannot be connected across the center of the device, as this connection might introduce unacceptable skew. The top and bottom halves of the global lines must be separately driven — although they can be driven by the same global buffer.

The eight global lines in each CLB column can be driven by either of two types of global buffers. They can also be driven by internal logic, because they can be accessed by single, double, and quad lines at the top, bottom, half, and quarter points. Consequently, the number of different clocks that can be used simultaneously in an XC4000X device is very large.

There are four global lines feeding the IOBs at the left edge of the device. IOBs along the right edge have eight global lines. There is a single global line along the top and bottom edges with access to the IOBs. All IOB global lines are broken at the center. They cannot be connected across the center of the device, as this connection might introduce unacceptable skew.

IOB global lines can be driven from two types of global buffers, or from local interconnect. Alternatively, top and bottom IOBs can be clocked from the global lines in the adjacent CLB column.

Two different types of clock buffers are available in the XC4000X:

- Global Low-Skew Buffers (BUFGSL)
- Global Early Buffers (BUFGE)

Global Low-Skew Buffers are the standard clock buffers. They should be used for most internal clocking, whenever a large portion of the device must be driven.

Global Early Buffers are designed to provide a faster clock access, but CLB access is limited to one-fourth of the device. They also facilitate a faster I/O interface.

[Figure 35](#) is a conceptual diagram of the global net structure in the XC4000X.

Global Early buffers and Global Low-Skew buffers share a single pad. Therefore, the same IPAD symbol can drive one buffer of each type, in parallel. This configuration is particularly useful when using the Fast Capture latches, as described in [“IOB Input Signals” on page 20](#). Paired Global

Early and Global Low-Skew buffers share a common input; they cannot be driven by two different signals.

Choosing an XC4000X Clock Buffer

The clocking structure of the XC4000X provides a large variety of features. However, it can be simple to use, without understanding all the details. The software automatically handles clocks, along with all other routing, when the appropriate clock buffer is placed in the design. In fact, if a buffer symbol called BUFG is placed, rather than a specific type of buffer, the software even chooses the buffer most appropriate for the design. The detailed information in this section is provided for those users who want a finer level of control over their designs.

If fine control is desired, use the following summary and [Table 15 on page 35](#) to choose an appropriate clock buffer.

- The simplest thing to do is to use a Global Low-Skew buffer.
- If a faster clock path is needed, try a BUFG. The software will first try to use a Global Low-Skew Buffer. If timing requirements are not met, a faster buffer will automatically be used.
- If a single quadrant of the chip is sufficient for the clocked logic, and the timing requires a faster clock than the Global Low-Skew buffer, use a Global Early buffer.

Global Low-Skew Buffers

Each corner of the XC4000X device has two Global Low-Skew buffers. Any of the eight Global Low-Skew buffers can drive any of the eight vertical Global lines in a column of CLBs. In addition, any of the buffers can drive any of the four vertical lines accessing the IOBs on the left edge of the device, and any of the eight vertical lines accessing the IOBs on the right edge of the device. (See [Figure 36 on page 38](#).)

IOBs at the top and bottom edges of the device are accessed through the vertical Global lines in the CLB array, as in the XC4000E. Any Global Low-Skew buffer can, therefore, access every IOB and CLB in the device.

The Global Low-Skew buffers can be driven by either semi-dedicated pads or internal logic.

To use a Global Low-Skew buffer, instantiate a BUFGSL element in a schematic or in HDL code. If desired, attach a LOC attribute or property to direct placement to the designated location. For example, attach a LOC=T attribute or property to direct that a BUFGSL be placed in one of the two Global Low-Skew buffers on the top edge of the device, or a LOC=TR to indicate the Global Low-Skew buffer on the top edge of the device, on the right.

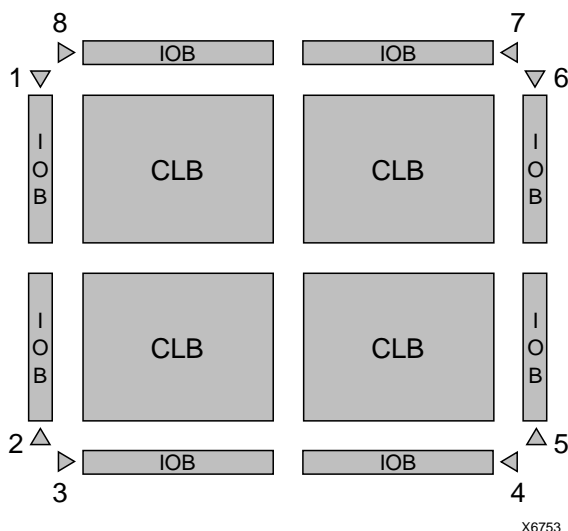


Figure 36: Any BUFGLS (GCK1 - GCK8) Can Drive Any or All Clock Inputs on the Device

Global Early Buffers

Each corner of the XC4000X device has two Global Early buffers. The primary purpose of the Global Early buffers is to provide an earlier clock access than the potentially heavily-loaded Global Low-Skew buffers. A clock source applied to both buffers will result in the Global Early clock edge occurring several nanoseconds earlier than the Global Low-Skew buffer clock edge, due to the lighter loading.

Global Early buffers also facilitate the fast capture of device inputs, using the Fast Capture latches described in **"IOB Input Signals"** on page 20. For Fast Capture, take a single clock signal, and route it through both a Global Early buffer and a Global Low-Skew buffer. (The two buffers share an input pad.) Use the Global Early buffer to clock the Fast Capture latch, and the Global Low-Skew buffer to clock the normal input flip-flop or latch, as shown in **Figure 17** on page 23.

The Global Early buffers can also be used to provide a fast Clock-to-Out on device output pins. However, an early clock in the output flip-flop IOB must be taken into consideration when calculating the internal clock speed for the design.

The Global Early buffers at the left and right edges of the chip have slightly different capabilities than the ones at the top and bottom. Refer to **Figure 37**, **Figure 38**, and **Figure 35** on page 36 while reading the following explanation.

Each Global Early buffer can access the eight vertical Global lines for all CLBs in the quadrant. Therefore, only one-fourth of the CLB clock pins can be accessed. This restriction is in large part responsible for the faster speed of the buffers, relative to the Global Low-Skew buffers.

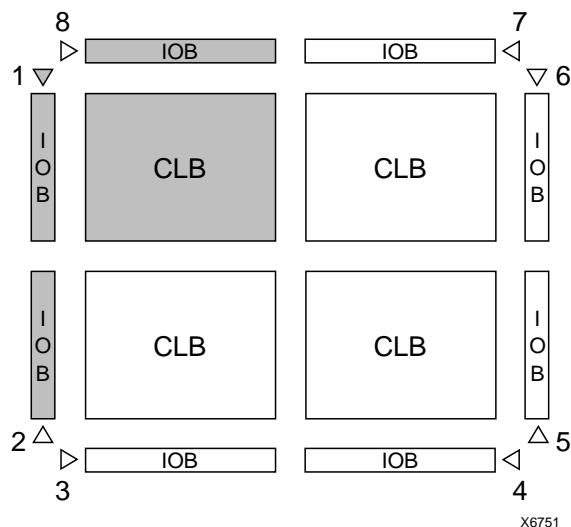


Figure 37: Left and Right BUFGEs Can Drive Any or All Clock Inputs in Same Quadrant or Edge (GCK1 is shown. GCK2, GCK5 and GCK6 are similar.)

The left-side Global Early buffers can each drive two of the four vertical lines accessing the IOBs on the entire left edge of the device. The right-side Global Early buffers can each drive two of the eight vertical lines accessing the IOBs on the entire right edge of the device. (See **Figure 37**.)

Each left and right Global Early buffer can also drive half of the IOBs along either the top or bottom edge of the device, using a dedicated line that can only be accessed through the Global Early buffers.

The top and bottom Global Early buffers can drive half of the IOBs along either the left or right edge of the device, as shown in **Figure 38**. They can only access the top and bottom IOBs via the CLB global lines.

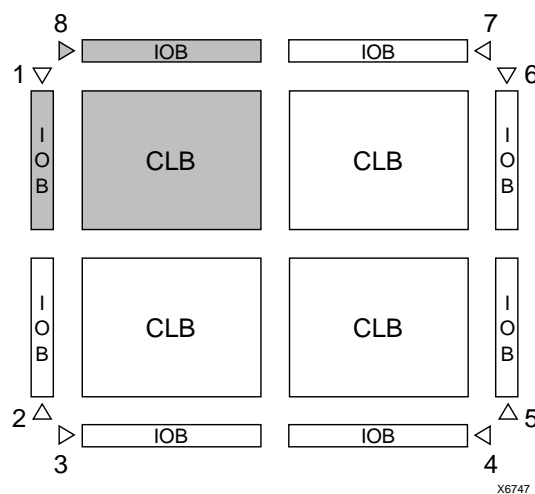


Figure 38: Top and Bottom BUFGEs Can Drive Any or All Clock Inputs in Same Quadrant (GCK8 is shown. GCK3, GCK4 and GCK7 are similar.)

is passed through and is captured by each FPGA when it recognizes the 0010 preamble. Following the length-count data, each FPGA outputs a High on DOUT until it has received its required number of data frames.

After an FPGA has received its configuration data, it passes on any additional frame start bits and configuration data on DOUT. When the total number of configuration clocks applied after memory initialization equals the value of the 24-bit length count, the FPGAs begin the start-up sequence and become operational together. FPGA I/O are normally released two CCLK cycles after the last configuration bit is received. **Figure 47 on page 53** shows the start-up timing for an XC4000 Series device.

The daisy-chained bitstream is not simply a concatenation of the individual bitstreams. The PROM file formatter must be used to combine the bitstreams for a daisy-chained configuration.

Multi-Family Daisy Chain

All Xilinx FPGAs of the XC2000, XC3000, and XC4000 Series use a compatible bitstream format and can, therefore, be connected in a daisy chain in an arbitrary sequence. There is, however, one limitation. The lead device must belong to the highest family in the chain. If the chain contains XC4000 Series devices, the master normally cannot be an XC2000 or XC3000 device.

The reason for this rule is shown in **Figure 47 on page 53**. Since all devices in the chain store the same length count value and generate or receive one common sequence of CCLK pulses, they all recognize length-count match on the same CCLK edge, as indicated on the left edge of **Figure 47**. The master device then generates additional CCLK pulses until it reaches its finish point F. The different families generate or require different numbers of additional CCLK pulses until they reach F. Not reaching F means that the device does not really finish its configuration, although DONE may have gone High, the outputs became active, and the internal reset was released. For the XC4000 Series device, not reaching F means that readback cannot be ini-

tiated and most boundary scan instructions cannot be used.

The user has some control over the relative timing of these events and can, therefore, make sure that they occur at the proper time and the finish point F is reached. Timing is controlled using options in the bitstream generation software.

XC3000 Master with an XC4000 Series Slave

Some designers want to use an inexpensive lead device in peripheral mode and have the more precious I/O pins of the XC4000 Series devices all available for user I/O. **Figure 44** provides a solution for that case.

This solution requires one CLB, one IOB and pin, and an internal oscillator with a frequency of up to 5 MHz as a clock source. The XC3000 master device must be configured with late Internal Reset, which is the default option.

One CLB and one IOB in the lead XC3000-family device are used to generate the additional CCLK pulse required by the XC4000 Series devices. When the lead device removes the internal RESET signal, the 2-bit shift register responds to its clock input and generates an active Low output signal for the duration of the subsequent clock period. An external connection between this output and CCLK thus creates the extra CCLK pulse.

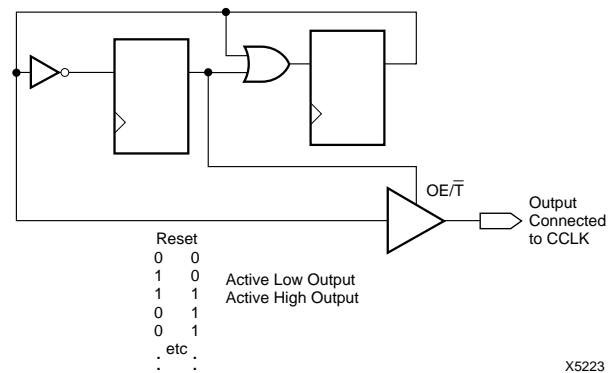


Figure 44: CCLK Generation for XC3000 Master Driving an XC4000 Series Slave

Table 20: XC4000E Program Data

Device	XC4003E	XC4005E	XC4006E	XC4008E	XC4010E	XC4013E	XC4020E	XC4025E
Max Logic Gates	3,000	5,000	6,000	8,000	10,000	13,000	20,000	25,000
CLBs (Row x Col.)	100 (10 x 10)	196 (14 x 14)	256 (16 x 16)	324 (18 x 18)	400 (20 x 20)	576 (24 x 24)	784 (28 x 28)	1,024 (32 x 32)
IOBs	80	112	128	144	160	192	224	256
Flip-Flops	360	616	768	936	1,120	1,536	2,016	2,560
Bits per Frame	126	166	186	206	226	266	306	346
Frames	428	572	644	716	788	932	1,076	1,220
Program Data	53,936	94,960	119,792	147,504	178,096	247,920	329,264	422,128
PROM Size (bits)	53,984	95,008	119,840	147,552	178,144	247,968	329,312	422,176

- Notes:
- Bits per Frame = (10 x number of rows) + 7 for the top + 13 for the bottom + 1 + 1 start bit + 4 error check bits
 Number of Frames = (36 x number of columns) + 26 for the left edge + 41 for the right edge + 1
 Program Data = (Bits per Frame x Number of Frames) + 8 postamble bits
 PROM Size = Program Data + 40 (header) + 8
 - The user can add more "one" bits as leading dummy bits in the header, or, if CRC = off, as trailing dummy bits at the end of any frame, following the four error check bits. However, the Length Count value **must** be adjusted for all such extra "one" bits, even for extra leading ones at the beginning of the header.

Table 21: XC4000EX/XL Program Data

Device	XC4002XL	XC4005	XC4010	XC4013	XC4020	XC4028	XC4036	XC4044	XC4052	XC4062	XC4085
Max Logic Gates	2,000	5,000	10,000	13,000	20,000	28,000	36,000	44,000	52,000	62,000	85,000
CLBs (Row x Column)	64 (8 x 8)	196 (14 x 14)	400 (20 x 20)	576 (24 x 24)	784 (28 x 28)	1,024 (32 x 32)	1,296 (36 x 36)	1,600 (40 x 40)	1,936 (44 x 44)	2,304 (48 x 48)	3,136 (56 x 56)
IOBs	64	112	160	192	224	256	288	320	352	384	448
Flip-Flops	256	616	1,120	1,536	2,016	2,560	3,168	3,840	4,576	5,376	7,168
Bits per Frame	133	205	277	325	373	421	469	517	565	613	709
Frames	459	741	1,023	1,211	1,399	1,587	1,775	1,963	2,151	2,339	2,715
Program Data	61,052	151,910	283,376	393,580	521,832	668,124	832,480	1,014,876	1,215,320	1,433,804	1,924,940
PROM Size (bits)	61,104	151,960	283,424	393,632	521,880	668,172	832,528	1,014,924	1,215,368	1,433,852	1,924,992

- Notes:
- Bits per frame = (13 x number of rows) + 9 for the top + 17 for the bottom + 8 + 1 start bit + 4 error check bits.
 Frames = (47 x number of columns) + 27 for the left edge + 52 for the right edge + 4.
 Program data = (bits per frame x number of frames) + 5 postamble bits.
 PROM size = (program data + 40 header bits + 8 start bits) rounded up to the nearest byte.
 - The user can add more "one" bits as leading dummy bits in the header, or, if CRC = off, as trailing dummy bits at the end of any frame, following the four error check bits. However, the Length Count value must be adjusted for all such extra "one" bits, even for extra leading "ones" at the beginning of the header.

Cyclic Redundancy Check (CRC) for Configuration and Readback

The Cyclic Redundancy Check is a method of error detection in data transmission applications. Generally, the transmitting system performs a calculation on the serial bitstream. The result of this calculation is tagged onto the data stream as additional check bits. The receiving system performs an identical calculation on the bitstream and compares the result with the received checksum.

Each data frame of the configuration bitstream has four error bits at the end, as shown in [Table 19](#). If a frame data error is detected during the loading of the FPGA, the con-

figuration process with a potentially corrupted bitstream is terminated. The FPGA pulls the $\overline{\text{INIT}}$ pin Low and goes into a Wait state.

During Readback, 11 bits of the 16-bit checksum are added to the end of the Readback data stream. The checksum is computed using the CRC-16 CCITT polynomial, as shown in [Figure 45](#). The checksum consists of the 11 most significant bits of the 16-bit code. A change in the checksum indicates a change in the Readback bitstream. A comparison to a previous checksum is meaningful only if the readback data is independent of the current device state. CLB outputs should not be included (Read Capture option not

used), and if RAM is present, the RAM content must be unchanged.

Statistically, one error out of 2048 might go undetected.

Configuration Sequence

There are four major steps in the XC4000 Series power-up configuration sequence.

- Configuration Memory Clear
- Initialization
- Configuration
- Start-Up

The full process is illustrated in Figure 46.

Configuration Memory Clear

When power is first applied or is reapplied to an FPGA, an internal circuit forces initialization of the configuration logic. When V_{CC} reaches an operational level, and the circuit passes the write and read test of a sample pair of configuration bits, a time delay is started. This time delay is nominally 16 ms, and up to 10% longer in the low-voltage devices. The delay is four times as long when in Master Modes (M0 Low), to allow ample time for all slaves to reach a stable V_{CC} . When all \overline{INIT} pins are tied together, as recommended, the longest delay takes precedence. Therefore, devices with different time delays can easily be mixed and matched in a daisy chain.

This delay is applied only on power-up. It is not applied when re-configuring an FPGA by pulsing the PROGRAM pin

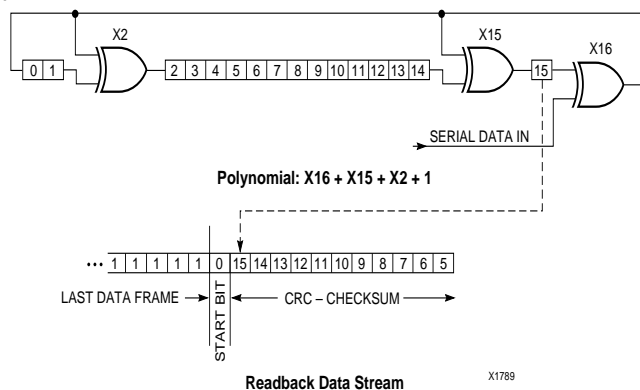


Figure 45: Circuit for Generating CRC-16

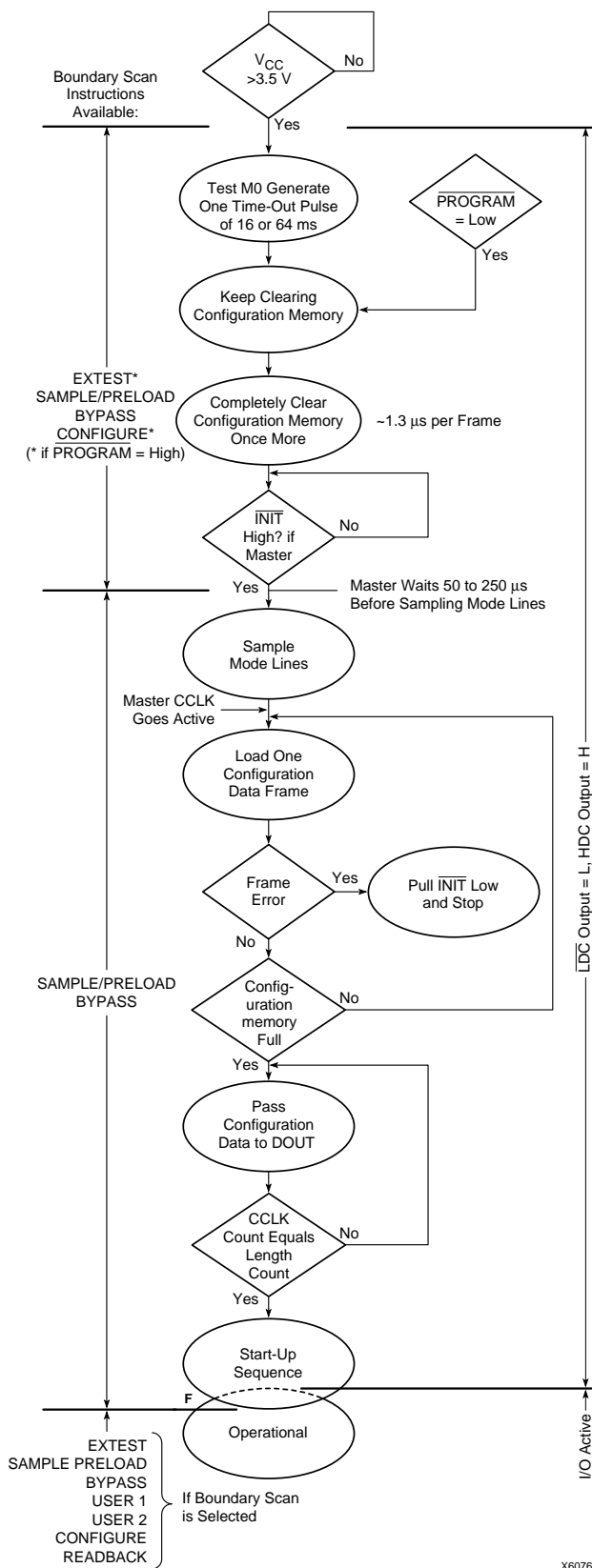


Figure 46: Power-up Configuration Sequence

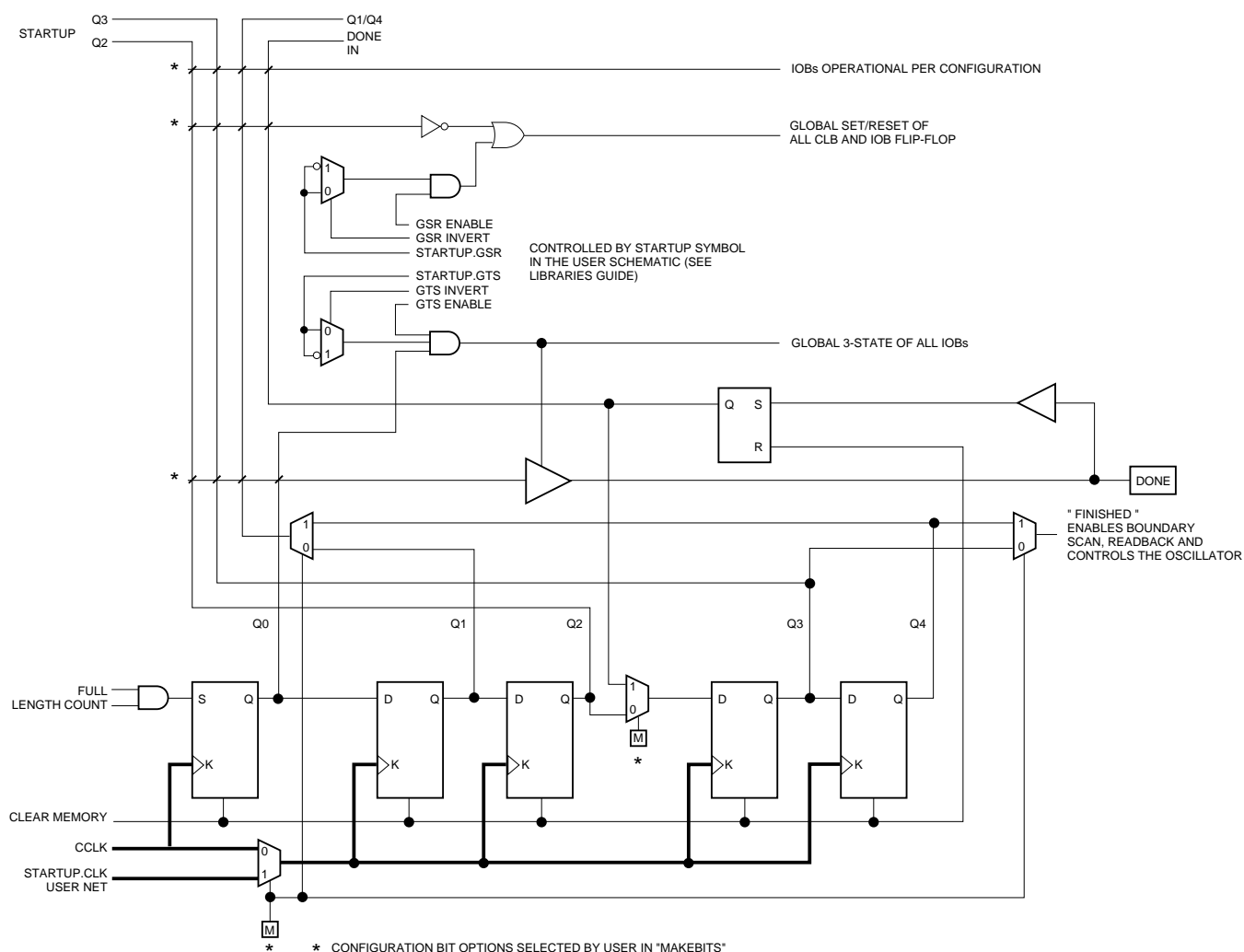


Figure 48: Start-up Logic

Readback

The user can read back the content of configuration memory and the level of certain internal nodes without interfering with the normal operation of the device.

Readback not only reports the downloaded configuration bits, but can also include the present state of the device, represented by the content of all flip-flops and latches in CLBs and IOBs, as well as the content of function generators used as RAMs.

Note that in XC4000 Series devices, configuration data is *not* inverted with respect to configuration as it is in XC2000 and XC3000 families.

XC4000 Series Readback does not use any dedicated pins, but uses four internal nets (RDBK.TRIG, RDBK.DATA, RDBK.RIP and RDBK.CLK) that can be routed to any IOB. To access the internal Readback signals, place the READ-

BACK library symbol and attach the appropriate pad symbols, as shown in **Figure 49**.

After Readback has been initiated by a High level on RDBK.TRIG after configuration, the RDBK.RIP (Read In Progress) output goes High on the next rising edge of RDBK.CLK. Subsequent rising edges of this clock shift out Readback data on the RDBK.DATA net.

Readback data does not include the preamble, but starts with five dummy bits (all High) followed by the Start bit (Low) of the first frame. The first two data bits of the first frame are always High.

Each frame ends with four error check bits. They are read back as High. The last seven bits of the last frame are also read back as High. An additional Start bit (Low) and an 11-bit Cyclic Redundancy Check (CRC) signature follow, before RDBK.RIP returns Low.

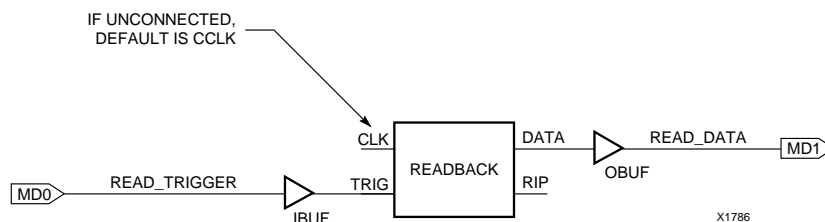


Figure 49: Readback Schematic Example

Readback Options

Readback options are: Read Capture, Read Abort, and Clock Select. They are set with the bitstream generation software.

Read Capture

When the Read Capture option is selected, the readback data stream includes sampled values of CLB and IOB signals. The rising edge of RDBK.TRIG latches the inverted values of the four CLB outputs, the IOB output flip-flops and the input signals I1 and I2. Note that while the bits describing configuration (interconnect, function generators, and RAM content) are *not* inverted, the CLB and IOB output signals *are* inverted.

When the Read Capture option is not selected, the values of the capture bits reflect the configuration data originally written to those memory locations.

If the RAM capability of the CLBs is used, RAM data are available in readback, since they directly overwrite the F and G function-table configuration of the CLB.

RDBK.TRIG is located in the lower-left corner of the device, as shown in [Figure 50](#).

Read Abort

When the Read Abort option is selected, a High-to-Low transition on RDBK.TRIG terminates the readback operation and prepares the logic to accept another trigger.

After an aborted readback, additional clocks (up to one readback clock per configuration frame) may be required to re-initialize the control logic. The status of readback is indicated by the output control net RDBK.RIP. RDBK.RIP is High whenever a readback is in progress.

Clock Select

CCLK is the default clock. However, the user can insert another clock on RDBK.CLK. Readback control and data are clocked on rising edges of RDBK.CLK. If readback must be inhibited for security reasons, the readback control nets are simply not connected.

RDBK.CLK is located in the lower right chip corner, as shown in [Figure 50](#).

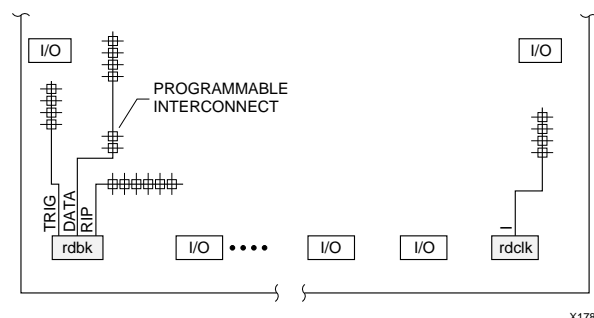


Figure 50: READBACK Symbol in Graphical Editor

Violating the Maximum High and Low Time Specification for the Readback Clock

The readback clock has a maximum High and Low time specification. In some cases, this specification cannot be met. For example, if a processor is controlling readback, an interrupt may force it to stop in the middle of a readback. This necessitates stopping the clock, and thus violating the specification.

The specification is mandatory only on clocking data at the end of a frame prior to the next start bit. The transfer mechanism will load the data to a shift register during the last six clock cycles of the frame, prior to the start bit of the following frame. This loading process is dynamic, and is the source of the maximum High and Low time requirements.

Therefore, the specification only applies to the six clock cycles prior to and including any start bit, including the clocks before the first start bit in the readback data stream. At other times, the frame data is already in the register and the register is not dynamic. Thus, it can be shifted out just like a regular shift register.

The user must precisely calculate the location of the readback data relative to the frame. The system must keep track of the position within a data frame, and disable interrupts before frame boundaries. Frame lengths and data formats are listed in [Table 19](#), [Table 20](#) and [Table 21](#).

Readback with the XChecker Cable

The XChecker Universal Download/Readback Cable and Logic Probe uses the readback feature for bitstream verification. It can also display selected internal signals on the PC or workstation screen, functioning as a low-cost in-circuit emulator.

Asynchronous Peripheral Mode

Write to FPGA

Asynchronous Peripheral mode uses the trailing edge of the logic AND condition of \overline{WS} and $\overline{CS0}$ being Low and \overline{RS} and $\overline{CS1}$ being High to accept byte-wide data from a microprocessor bus. In the lead FPGA, this data is loaded into a double-buffered UART-like parallel-to-serial converter and is serially shifted into the internal logic.

The lead FPGA presents the preamble data (and all data that overflows the lead device) on its DOUT pin. The RDY/BUSY output from the lead FPGA acts as a handshake signal to the microprocessor. RDY/BUSY goes Low when a byte has been received, and goes High again when the byte-wide input buffer has transferred its information into the shift register, and the buffer is ready to receive new data. A new write may be started immediately, as soon as the RDY/BUSY output has gone Low, acknowledging receipt of the previous data. Write may not be terminated until RDY/BUSY is High again for one CCLK period. Note that RDY/BUSY is pulled High with a high-impedance pull-up prior to \overline{INIT} going High.

The length of the \overline{BUSY} signal depends on the activity in the UART. If the shift register was empty when the new byte was received, the \overline{BUSY} signal lasts for only two CCLK periods. If the shift register was still full when the new byte was received, the \overline{BUSY} signal can be as long as nine CCLK periods.

Note that after the last byte has been entered, only seven of its bits are shifted out. CCLK remains High with DOUT equal to bit 6 (the next-to-last bit) of the last byte entered.

The RDY/ \overline{BUSY} handshake can be ignored if the delay from any one Write to the end of the next Write is guaranteed to be longer than 10 CCLK periods.

Status Read

The logic AND condition of the $\overline{CS0}$, $\overline{CS1}$ and \overline{RS} inputs puts the device status on the Data bus.

- D7 High indicates Ready
- D7 Low indicates Busy
- D0 through D6 go unconditionally High

It is mandatory that the whole start-up sequence be started and completed by one byte-wide input. Otherwise, the pins used as Write Strobe or Chip Enable might become active outputs and interfere with the final byte transfer. If this transfer does not occur, the start-up sequence is not completed all the way to the finish (point F in [Figure 47 on page 53](#)).

In this case, at worst, the internal reset is not released. At best, Readback and Boundary Scan are inhibited. The length-count value, as generated by the XACTstep software, ensures that these problems never occur.

Although RDY/ \overline{BUSY} is brought out as a separate signal, microprocessors can more easily read this information on one of the data lines. For this purpose, D7 represents the RDY/ \overline{BUSY} status when \overline{RS} is Low, \overline{WS} is High, and the two chip select lines are both active.

Asynchronous Peripheral mode is selected by a <101> on the mode pins (M2, M1, M0).

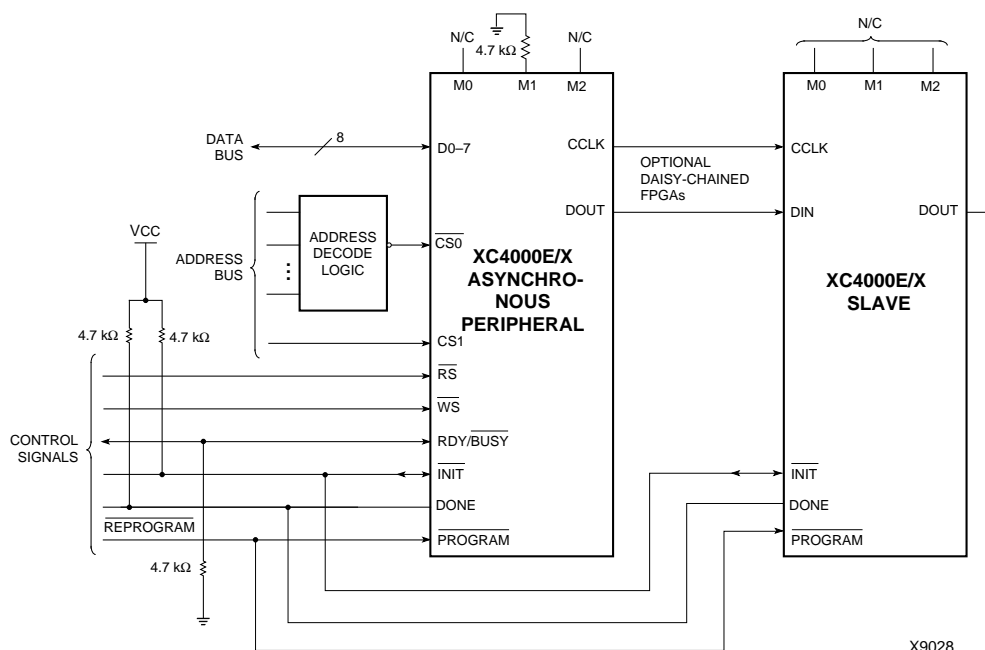
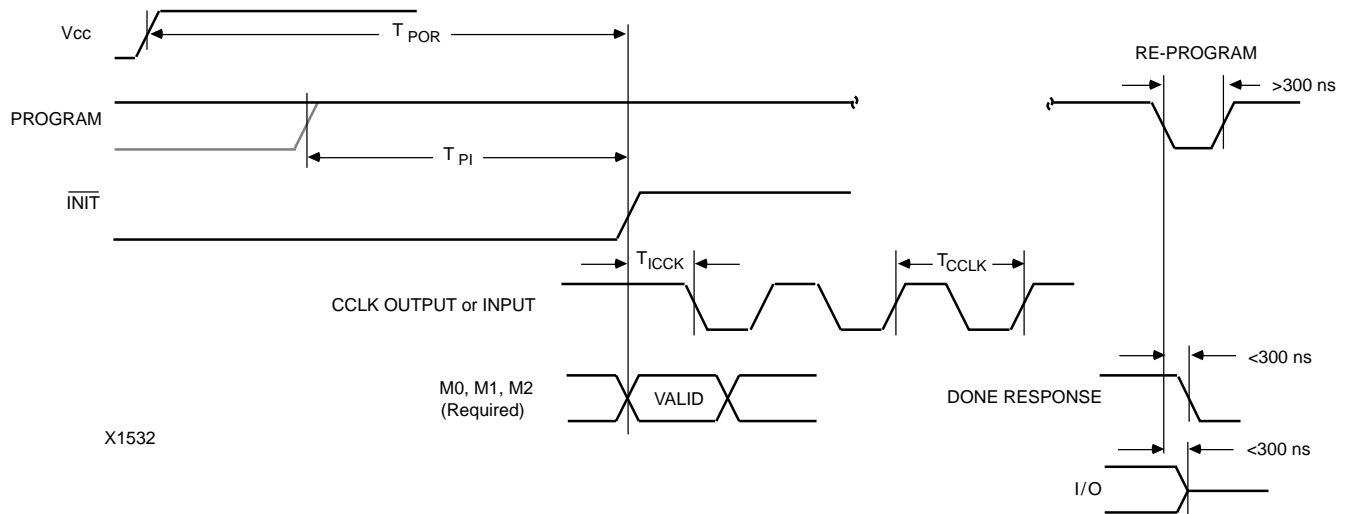


Figure 58: Asynchronous Peripheral Mode Circuit Diagram

Configuration Switching Characteristics



X1532

Master Modes (XC4000E/EX)

Description		Symbol	Min	Max	Units
Power-On Reset	M0 = High	T_{POR}	10	40	ms
	M0 = Low	T_{POR}	40	130	ms
Program Latency		T_{PI}	30	200	μ s per CLB column
CCLK (output) Delay		T_{ICCK}	40	250	μ s
CCLK (output) Period, slow		T_{CCLK}	640	2000	ns
CCLK (output) Period, fast		T_{CCLK}	80	250	ns

Master Modes (XC4000XL)

Description		Symbol	Min	Max	Units
Power-On Reset	M0 = High	T_{POR}	10	40	ms
	M0 = Low	T_{POR}	40	130	ms
Program Latency		T_{PI}	30	200	μ s per CLB column
CCLK (output) Delay		T_{ICCK}	40	250	μ s
CCLK (output) Period, slow		T_{CCLK}	540	1600	ns
CCLK (output) Period, fast		T_{CCLK}	67	200	ns

Slave and Peripheral Modes (All)

Description		Symbol	Min	Max	Units
Power-On Reset		T_{POR}	10	33	ms
Program Latency		T_{PI}	30	200	μ s per CLB column
CCLK (input) Delay (required)		T_{ICCK}	4		μ s
CCLK (input) Period (required)		T_{CCLK}	100		ns

User I/O Per Package

Table 27, Table 28, and Table 29 show the number of user I/Os available in each package for XC4000-Series devices. Call your local sales office for the latest availability information, or see the Xilinx website at <http://www.xilinx.com> for the latest revision of the specifications.

Table 27: User I/O Chart for XC4000XL FPGAs

Device	Max I/O	Maximum User Accessible I/O by Package Type																					
		PC84	PQ100	VQ100	TQ144	HT144	HQ160	PQ160	TQ176	HT176	HQ208	PQ208	HQ240	PQ240	BG256	PG299	HQ304	BG352	PG411	BG432	PG475	PG559	BG560
XC4002XL	64	61	64	64																			
XC4005XL	112	61	77	77	112			112				112											
XC4010XL	160	61	77		113			129	145			160			160								
XC4013XL	192					113		129		145		160		192	192								
XC4020XL	224					113		129		145		160		192	205								
XC4028XL	256						129				160		193		205	256	256	256					
XC4036XL	288						129				160		193				256	288	288	288			
XC4044XL	320						129				160		193				256	289	320	320			
XC4052XL	352												193				256		352	352			352
XC4062XL	384												193				256			352	384		384
XC4085XL	448																			352		448	448

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Table 28: User I/O Chart for XC4000E FPGAs

Device	Max I/O	Maximum User Accessible I/O by Package Type															
		PC84	PQ100	VQ100	PG120	TQ144	PG156	PQ160	PG191	HQ208	PQ208	PG223	BG225	HQ240	PQ240	PG299	HQ304
XC4003E	80	61	77	77	80												
XC4005E	112	61	77			112	112	112			112						
XC4006E	128	61				113	125	128			128						
XC4008E	144	61						129	144		144						
XC4010E	160	61						129	160	160	160		160				
XC4013E	192							129		160	160	192	192	192	192		
XC4020E	224									160		192		193			
XC4025E	256											192		193		256	256

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Table 29: User I/O Chart for XC4000EX FPGAs

Device	Max I/O	Maximum User Accessible I/O by Package Type						
		HQ208	HQ240	PG299	HQ304	BG352	PG411	BG432
XC4028EX	256	160	193	256	256	256		
XC4036EX	288		193		256	288	288	288

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XC4000 Series Electrical Characteristics and Device-Specific Pinout Table

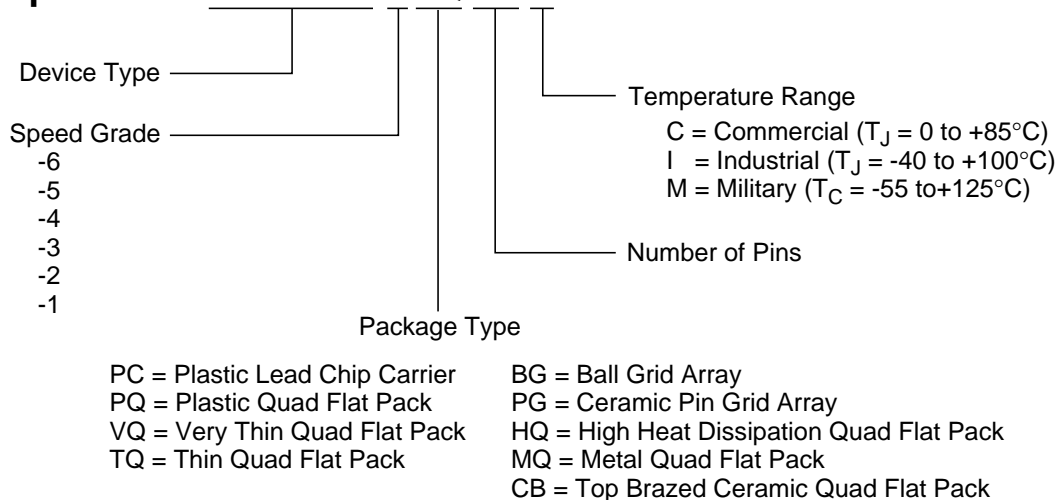
For the latest Electrical Characteristics and package/pinout information for each XC4000 Family, see the Xilinx web site at

http://www.xilinx.com/xlnx/xweb/xil_publications_index.jsp

Ordering Information

Example:

XC4013E-3HQ240C



X9020

Revision Control

Version	Description
3/30/98 (1.5)	Updated XC4000XL timing and added XC4002XL
1/29/99 (1.5)	Updated pin diagrams
5/14/99 (1.6)	Replaced Electrical Specification and pinout pages for E, EX, and XL families with separate updates and added URL link for electrical specifications/pinouts for Web users