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### Understanding [Embedded - FPGAs \(Field Programmable Gate Array\)](#)

Embedded - FPGAs, or Field Programmable Gate Arrays, are advanced integrated circuits that offer unparalleled flexibility and performance for digital systems. Unlike traditional fixed-function logic devices, FPGAs can be programmed and reprogrammed to execute a wide array of logical operations, enabling customized functionality tailored to specific applications. This reprogrammability allows developers to iterate designs quickly and implement complex functions without the need for custom hardware.

### Applications of Embedded - FPGAs

The versatility of Embedded - FPGAs makes them indispensable in numerous fields. In telecommunications.

#### Details

Product Status	Obsolete
Number of LABs/CLBs	2304
Number of Logic Elements/Cells	5472
Total RAM Bits	73728
Number of I/O	352
Number of Gates	62000
Voltage - Supply	3V ~ 3.6V
Mounting Type	Surface Mount
Operating Temperature	-40°C ~ 100°C (TJ)
Package / Case	432-LBGA Exposed Pad, Metal
Supplier Device Package	432-MBGA (40x40)
Purchase URL	<a href="https://www.e-xfl.com/product-detail/xilinx/xc4062xl-3bg432i">https://www.e-xfl.com/product-detail/xilinx/xc4062xl-3bg432i</a>

## Detailed Functional Description

XC4000 Series devices achieve high speed through advanced semiconductor technology and improved architecture. The XC4000E and XC4000X support system clock rates of up to 80 MHz and internal performance in excess of 150 MHz. Compared to older Xilinx FPGA families, XC4000 Series devices are more powerful. They offer on-chip edge-triggered and dual-port RAM, clock enables on I/O flip-flops, and wide-input decoders. They are more versatile in many applications, especially those involving RAM. Design cycles are faster due to a combination of increased routing resources and more sophisticated software.

### Basic Building Blocks

Xilinx user-programmable gate arrays include two major configurable elements: configurable logic blocks (CLBs) and input/output blocks (IOBs).

- CLBs provide the functional elements for constructing the user's logic.
- IOBs provide the interface between the package pins and internal signal lines.

Three other types of circuits are also available:

- 3-State buffers (TBUFs) driving horizontal longlines are associated with each CLB.
- Wide edge decoders are available around the periphery of each device.
- An on-chip oscillator is provided.

Programmable interconnect resources provide routing paths to connect the inputs and outputs of these configurable elements to the appropriate networks.

The functionality of each circuit block is customized during configuration by programming internal static memory cells. The values stored in these memory cells determine the logic functions and interconnections implemented in the FPGA. Each of these available circuits is described in this section.

### Configurable Logic Blocks (CLBs)

Configurable Logic Blocks implement most of the logic in an FPGA. The principal CLB elements are shown in **Figure 1**. Two 4-input function generators (F and G) offer unrestricted versatility. Most combinatorial logic functions need four or fewer inputs. However, a third function generator (H) is provided. The H function generator has three inputs. Either zero, one, or two of these inputs can be the outputs of F and G; the other input(s) are from outside the CLB. The CLB can, therefore, implement certain functions of up to nine variables, like parity check or expandable-identity comparison of two sets of four inputs.

Each CLB contains two storage elements that can be used to store the function generator outputs. However, the storage elements and function generators can also be used independently. These storage elements can be configured as flip-flops in both XC4000E and XC4000X devices; in the XC4000X they can optionally be configured as latches. DIN can be used as a direct input to either of the two storage elements. H1 can drive the other through the H function generator. Function generator outputs can also drive two outputs independent of the storage element outputs. This versatility increases logic capacity and simplifies routing.

Thirteen CLB inputs and four CLB outputs provide access to the function generators and storage elements. These inputs and outputs connect to the programmable interconnect resources outside the block.

### Function Generators

Four independent inputs are provided to each of two function generators (F1 - F4 and G1 - G4). These function generators, with outputs labeled F' and G', are each capable of implementing any arbitrarily defined Boolean function of four inputs. The function generators are implemented as memory look-up tables. The propagation delay is therefore independent of the function implemented.

A third function generator, labeled H', can implement any Boolean function of its three inputs. Two of these inputs can optionally be the F' and G' functional generator outputs. Alternatively, one or both of these inputs can come from outside the CLB (H2, H0). The third input must come from outside the block (H1).

Signals from the function generators can exit the CLB on two outputs. F' or H' can be connected to the X output. G' or H' can be connected to the Y output.

A CLB can be used to implement any of the following functions:

- any function of up to four variables, plus any second function of up to four unrelated variables, plus any third function of up to three unrelated variables<sup>1</sup>
- any single function of five variables
- any function of four variables together with some functions of six variables
- some functions of up to nine variables.

Implementing wide functions in a single block reduces both the number of blocks required and the delay in the signal path, achieving both increased capacity and speed.

The versatility of the CLB function generators significantly improves system speed. In addition, the design-software tools can deal with each function generator independently. This flexibility improves cell usage.

1. When three separate functions are generated, one of the function outputs must be captured in a flip-flop internal to the CLB. Only two unregistered function generator outputs are available from the CLB.

### Set/Reset

An asynchronous storage element input (SR) can be configured as either set or reset. This configuration option determines the state in which each flip-flop becomes operational after configuration. It also determines the effect of a Global Set/Reset pulse during normal operation, and the effect of a pulse on the SR pin of the CLB. All three set/reset functions for any single flip-flop are controlled by the same configuration data bit.

The set/reset state can be independently specified for each flip-flop. This input can also be independently disabled for either flip-flop.

The set/reset state is specified by using the INIT attribute, or by placing the appropriate set or reset flip-flop library symbol.

SR is active High. It is not invertible within the CLB.

### Global Set/Reset

A separate Global Set/Reset line (not shown in Figure 1) sets or clears each storage element during power-up, re-configuration, or when a dedicated Reset net is driven active. This global net (GSR) does not compete with other routing resources; it uses a dedicated distribution network.

Each flip-flop is configured as either globally set or reset in the same way that the local set/reset (SR) is specified. Therefore, if a flip-flop is set by SR, it is also set by GSR. Similarly, a reset flip-flop is reset by both SR and GSR.



**Figure 2: Schematic Symbols for Global Set/Reset**

GSR can be driven from any user-programmable pin as a global reset input. To use this global net, place an input pad and input buffer in the schematic or HDL code, driving the GSR pin of the STARTUP symbol. (See Figure 2.) A specific pin location can be assigned to this input using a LOC attribute or property, just as with any other user-programmable pad. An inverter can optionally be inserted after the input buffer to invert the sense of the Global Set/Reset signal.

Alternatively, GSR can be driven from any internal node.

### Data Inputs and Outputs

The source of a storage element data input is programmable. It is driven by any of the functions F', G', and H', or by the Direct In (DIN) block input. The flip-flops or latches drive the XQ and YQ CLB outputs.

Two fast feed-through paths are available, as shown in Figure 1. A two-to-one multiplexer on each of the XQ and YQ outputs selects between a storage element output and any of the control inputs. This bypass is sometimes used by the automated router to repower internal signals.

### Control Signals

Multiplexers in the CLB map the four control inputs (C1 - C4 in Figure 1) into the four internal control signals (H1, DIN/H2, SR/H0, and EC). Any of these inputs can drive any of the four internal control signals.

When the logic function is enabled, the four inputs are:

- EC — Enable Clock
- SR/H0 — Asynchronous Set/Reset or H function generator Input 0
- DIN/H2 — Direct In or H function generator Input 2
- H1 — H function generator Input 1.

When the memory function is enabled, the four inputs are:

- EC — Enable Clock
- WE — Write Enable
- D0 — Data Input to F and/or G function generator
- D1 — Data input to G function generator (16x1 and 16x2 modes) or 5th Address bit (32x1 mode).

### Using FPGA Flip-Flops and Latches

The abundance of flip-flops in the XC4000 Series invites pipelined designs. This is a powerful way of increasing performance by breaking the function into smaller subfunctions and executing them in parallel, passing on the results through pipeline flip-flops. This method should be seriously considered wherever throughput is more important than latency.

To include a CLB flip-flop, place the appropriate library symbol. For example, FDCE is a D-type flip-flop with clock enable and asynchronous clear. The corresponding latch symbol (for the XC4000X only) is called LDCE.

In XC4000 Series devices, the flip flops can be used as registers or shift registers without blocking the function generators from performing a different, perhaps unrelated task. This ability increases the functional capacity of the devices.

The CLB setup time is specified between the function generator inputs and the clock input K. Therefore, the specified CLB flip-flop setup time includes the delay through the function generator.

### Using Function Generators as RAM

Optional modes for each CLB make the memory look-up tables in the F' and G' function generators usable as an array of Read/Write memory cells. Available modes are level-sensitive (similar to the XC4000/A/H families), edge-triggered, and dual-port edge-triggered. Depending on the selected mode, a single CLB can be configured as either a 16x2, 32x1, or 16x1 bit array.

### Fast Carry Logic

Each CLB F and G function generator contains dedicated arithmetic logic for the fast generation of carry and borrow signals. This extra output is passed on to the function generator in the adjacent CLB. The carry chain is independent of normal routing resources.

Dedicated fast carry logic greatly increases the efficiency and performance of adders, subtractors, accumulators, comparators and counters. It also opens the door to many new applications involving arithmetic operation, where the previous generations of FPGAs were not fast enough or too inefficient. High-speed address offset calculations in micro-processor or graphics systems, and high-speed addition in digital signal processing are two typical applications.

The two 4-input function generators can be configured as a 2-bit adder with built-in hidden carry that can be expanded to any length. This dedicated carry circuitry is so fast and efficient that conventional speed-up methods like carry generate/propagate are meaningless even at the 16-bit level, and of marginal benefit at the 32-bit level.

This fast carry logic is one of the more significant features of the XC4000 Series, speeding up arithmetic and counting into the 70 MHz range.

The carry chain in XC4000E devices can run either up or down. At the top and bottom of the columns where there are no CLBs above or below, the carry is propagated to the right. (See Figure 11.) In order to improve speed in the high-capacity XC4000X devices, which can potentially have very long carry chains, the carry chain travels upward only, as shown in Figure 12. Additionally, standard interconnect can be used to route a carry signal in the downward direction.

Figure 13 on page 19 shows an XC4000E CLB with dedicated fast carry logic. The carry logic in the XC4000X is similar, except that COUT exits at the top only, and the signal CINDOWN does not exist. As shown in Figure 13, the carry logic shares operand and control inputs with the function generators. The carry outputs connect to the function generators, where they are combined with the operands to form the sums.

Figure 14 on page 20 shows the details of the carry logic for the XC4000E. This diagram shows the contents of the box labeled "CARRY LOGIC" in Figure 13. The XC4000X carry logic is very similar, but a multiplexer on the pass-through carry chain has been eliminated to reduce delay. Additionally, in the XC4000X the multiplexer on the G4 path has a memory-programmable 0 input, which permits G4 to directly connect to COUT. G4 thus becomes an additional high-speed initialization path for carry-in.

The dedicated carry logic is discussed in detail in Xilinx document XAPP 013: "Using the Dedicated Carry Logic in

XC4000." This discussion also applies to XC4000E devices, and to XC4000X devices when the minor logic changes are taken into account.

The fast carry logic can be accessed by placing special library symbols, or by using Xilinx Relationally Placed Macros (RPMs) that already include these symbols.



X6687

**Figure 11: Available XC4000E Carry Propagation Paths**



X6610

**Figure 12: Available XC4000X Carry Propagation Paths (dotted lines use general interconnect)**

**Table 8: Supported Sources for XC4000 Series Device Inputs**

Source	XC4000E/EX Series Inputs		XC4000XL Series Inputs
	5 V, TTL	5 V, CMOS	3.3 V CMOS
Any device, V <sub>CC</sub> = 3.3 V, CMOS outputs	✓	Unreliable Data	✓
XC4000 Series, V <sub>CC</sub> = 5 V, TTL outputs	✓		✓
Any device, V <sub>CC</sub> = 5 V, TTL outputs (V <sub>oh</sub> ≤ 3.7 V)	✓		✓
Any device, V <sub>CC</sub> = 5 V, CMOS outputs	✓	✓	✓

#### XC4000XL 5-Volt Tolerant I/Os

The I/Os on the XC4000XL are fully 5-volt tolerant even though the V<sub>CC</sub> is 3.3 volts. This allows 5 V signals to directly connect to the XC4000XL inputs without damage, as shown in [Table 8](#). In addition, the 3.3 volt V<sub>CC</sub> can be applied before or after 5 volt signals are applied to the I/Os. This makes the XC4000XL immune to power supply sequencing problems.


#### Registered Inputs

The I1 and I2 signals that exit the block can each carry either the direct or registered input signal.

The input and output storage elements in each IOB have a common clock enable input, which, through configuration, can be activated individually for the input or output flip-flop, or both. This clock enable operates exactly like the EC pin on the XC4000 Series CLB. It cannot be inverted within the IOB.

The storage element behavior is shown in [Table 9](#).

**Table 9: Input Register Functionality (active rising edge is shown)**

Mode	Clock	Clock Enable	D	Q
Power-Up or GSR	X	X	X	SR
Flip-Flop		1*	D	D
	0	X	X	Q
Latch	1	1*	X	Q
	0	1*	D	D
Both	X	0	X	Q

Legend:

X 

SR

0\*

1\*

Don't care  
Rising edge

Set or Reset value. Reset is default.

Input is Low or unconnected (default value)

Input is High or unconnected (default value)

#### Optional Delay Guarantees Zero Hold Time

The data input to the register can optionally be delayed by several nanoseconds. With the delay enabled, the setup time of the input flip-flop is increased so that normal clock routing does not result in a positive hold-time requirement. A positive hold time requirement can lead to unreliable, temperature- or processing-dependent operation.

The input flip-flop setup time is defined between the data measured at the device I/O pin and the clock input at the IOB (not at the clock pin). Any routing delay from the device clock pin to the clock input of the IOB must, therefore, be subtracted from this setup time to arrive at the real setup time requirement relative to the device pins. A short specified setup time might, therefore, result in a negative setup time at the device pins, i.e., a positive hold-time requirement.

When a delay is inserted on the data line, more clock delay can be tolerated without causing a positive hold-time requirement. Sufficient delay eliminates the possibility of a data hold-time requirement at the external pin. The maximum delay is therefore inserted as the default.

The XC4000E IOB has a one-tap delay element: either the delay is inserted (default), or it is not. The delay guarantees a zero hold time with respect to clocks routed through any of the XC4000E global clock buffers. (See [“Global Nets and Buffers \(XC4000E only\)” on page 35](#) for a description of the global clock buffers in the XC4000E.) For a shorter input register setup time, with non-zero hold, attach a NODELAY attribute or property to the flip-flop.

The XC4000X IOB has a two-tap delay element, with choices of a full delay, a partial delay, or no delay. The attributes or properties used to select the desired delay are shown in [Table 10](#). The choices are no added attribute, MEDDELAY, and NODELAY. The default setting, with no added attribute, ensures no hold time with respect to any of the XC4000X clock buffers, including the Global Low-Skew buffers. MEDDELAY ensures no hold time with respect to the Global Early buffers. Inputs with NODELAY may have a positive hold time with respect to all clock buffers. For a description of each of these buffers, see [“Global Nets and Buffers \(XC4000X only\)” on page 37](#).

**Table 10: XC4000X IOB Input Delay Element**

Value	When to Use
full delay (default, no attribute added)	Zero Hold with respect to Global Low-Skew Buffer, Global Early Buffer
MEDDELAY	Zero Hold with respect to Global Early Buffer
NODELAY	Short Setup, positive Hold time



Any XC4000 Series 5-Volt device with its outputs configured in TTL mode can drive the inputs of any typical 3.3-Volt device. (For a detailed discussion of how to interface between 5 V and 3.3 V devices, see the 3V Products section of *The Programmable Logic Data Book*.)

Supported destinations for XC4000 Series device outputs are shown in [Table 12](#).

An output can be configured as open-drain (open-collector) by placing an OBUFT symbol in a schematic or HDL code, then tying the 3-state pin (T) to the output signal, and the input pin (I) to Ground. (See [Figure 18](#).)

**Table 12: Supported Destinations for XC4000 Series Outputs**

Destination	XC4000 Series Outputs		
	3.3 V, CMOS	5 V, TTL	5 V, CMOS
Any typical device, Vcc = 3.3 V, CMOS-threshold inputs	✓	✓	some <sup>1</sup>
Any device, Vcc = 5 V, TTL-threshold inputs	✓	✓	✓
Any device, Vcc = 5 V, CMOS-threshold inputs	Unreliable Data		✓

1. Only if destination device has 5-V tolerant inputs



**Figure 18: Open-Drain Output**

### Output Slew Rate

The slew rate of each output buffer is, by default, reduced, to minimize power bus transients when switching non-critical signals. For critical signals, attach a FAST attribute or property to the output buffer or flip-flop.

For XC4000E devices, maximum total capacitive load for simultaneous fast mode switching in the same direction is 200 pF for all package pins between each Power/Ground pin pair. For XC4000X devices, additional internal

Power/Ground pin pairs are connected to special Power and Ground planes within the packages, to reduce ground bounce. Therefore, the maximum total capacitive load is 300 pF between each external Power/Ground pin pair. Maximum loading may vary for the low-voltage devices.

For slew-rate limited outputs this total is two times larger for each device type: 400 pF for XC4000E devices and 600 pF for XC4000X devices. This maximum capacitive load should not be exceeded, as it can result in ground bounce of greater than 1.5 V amplitude and more than 5 ns duration. This level of ground bounce may cause undesired transient behavior on an output, or in the internal logic. This restriction is common to all high-speed digital ICs, and is not particular to Xilinx or the XC4000 Series.

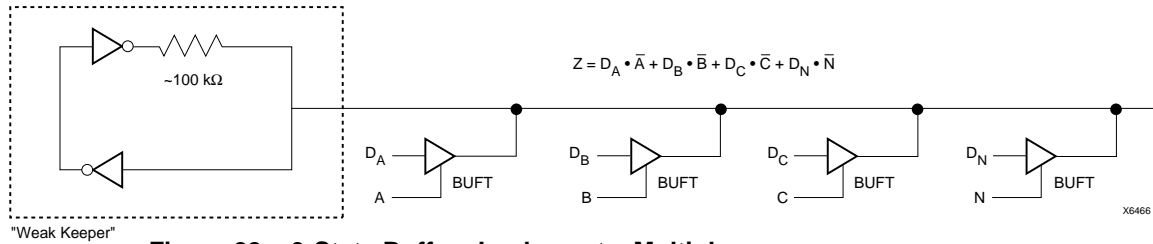
XC4000 Series devices have a feature called “Soft Start-up,” designed to reduce ground bounce when all outputs are turned on simultaneously at the end of configuration. When the configuration process is finished and the device starts up, the first activation of the outputs is automatically slew-rate limited. Immediately following the initial activation of the I/O, the slew rate of the individual outputs is determined by the individual configuration option for each IOB.

### Global Three-State

A separate Global 3-State line (not shown in [Figure 15](#) or [Figure 16](#)) forces all FPGA outputs to the high-impedance state, unless boundary scan is enabled and is executing an EXTEST instruction. This global net (GTS) does not compete with other routing resources; it uses a dedicated distribution network.

GTS can be driven from any user-programmable pin as a global 3-state input. To use this global net, place an input pad and input buffer in the schematic or HDL code, driving the GTS pin of the STARTUP symbol. A specific pin location can be assigned to this input using a LOC attribute or property, just as with any other user-programmable pad. An inverter can optionally be inserted after the input buffer to invert the sense of the Global 3-State signal. Using GTS is similar to GSR. See [Figure 2 on page 11](#) for details.

Alternatively, GTS can be driven from any internal node.



**Figure 22: 3-State Buffers Implement a Multiplexer**

## Wide Edge Decoders

Dedicated decoder circuitry boosts the performance of wide decoding functions. When the address or data field is wider than the function generator inputs, FPGAs need multi-level decoding and are thus slower than PALs. XC4000 Series CLBs have nine inputs. Any decoder of up to nine inputs is, therefore, compact and fast. However, there is also a need for much wider decoders, especially for address decoding in large microprocessor systems.

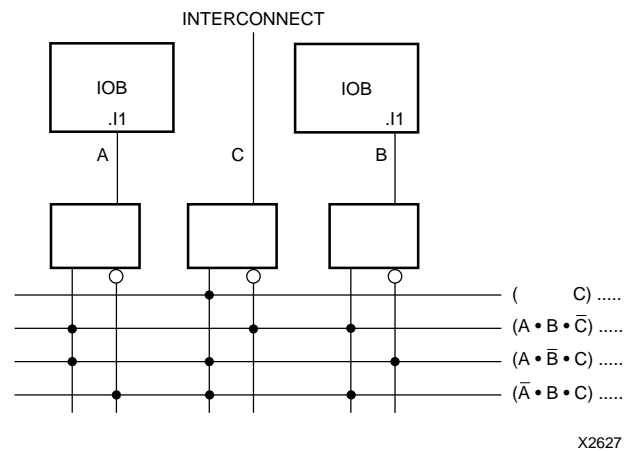
An XC4000 Series FPGA has four programmable decoders located on each edge of the device. The inputs to each decoder are any of the IOB I1 signals on that edge plus one local interconnect per CLB row or column. Each row or column of CLBs provides up to three variables or their complements., as shown in Figure 23. Each decoder generates a High output (resistor pull-up) when the AND condition of the selected inputs, or their complements, is true. This is analogous to a product term in typical PAL devices.

Each of these wired-AND gates is capable of accepting up to 42 inputs on the XC4005E and 72 on the XC4013E. There are up to 96 inputs for each decoder on the XC4028X and 132 on the XC4052X. The decoders may also be split in two when a larger number of narrower decoders are required, for a maximum of 32 decoders per device.

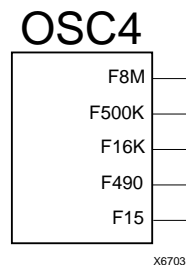
The decoder outputs can drive CLB inputs, so they can be combined with other logic to form a PAL-like AND/OR structure. The decoder outputs can also be routed directly to the chip outputs. For fastest speed, the output should be on the same chip edge as the decoder. Very large PALs can be emulated by ORing the decoder outputs in a CLB. This decoding feature covers what has long been considered a weakness of older FPGAs. Users often resorted to external PALs for simple but fast decoding functions. Now, the dedicated decoders in the XC4000 Series device can implement these functions fast and efficiently.

To use the wide edge decoders, place one or more of the WAND library symbols (WAND1, WAND4, WAND8, WAND16). Attach a DECODE attribute or property to each WAND symbol. Tie the outputs together and attach a PUL-

LUP symbol. Location attributes or properties such as L (left edge) or TR (right half of top edge) should also be used to ensure the correct placement of the decoder inputs.



**Figure 23: XC4000 Series Edge Decoding Example**



**Figure 24: XC4000 Series Oscillator Symbol**

## On-Chip Oscillator

XC4000 Series devices include an internal oscillator. This oscillator is used to clock the power-on time-out, for configuration memory clearing, and as the source of CCLK in Master configuration modes. The oscillator runs at a nominal 8 MHz frequency that varies with process, Vcc, and temperature. The output frequency falls between 4 and 10 MHz.

The oscillator output is optionally available after configuration. Any two of four resynchronized taps of a built-in divider are also available. These taps are at the fourth, ninth, fourteenth and nineteenth bits of the divider. Therefore, if the primary oscillator output is running at the nominal 8 MHz, the user has access to an 8 MHz clock, plus any two of 500 kHz, 16kHz, 490Hz and 15Hz (up to 10% lower for low-voltage devices). These frequencies can vary by as much as -50% or +25%.

These signals can be accessed by placing the OSC4 library element in a schematic or in HDL code (see [Figure 24](#)).

The oscillator is automatically disabled after configuration if the OSC4 symbol is not used in the design.

## Programmable Interconnect

All internal connections are composed of metal segments with programmable switching points and switching matrices to implement the desired routing. A structured, hierarchical matrix of routing resources is provided to achieve efficient automated routing.

The XC4000E and XC4000X share a basic interconnect structure. XC4000X devices, however, have additional routing not available in the XC4000E. The extra routing resources allow high utilization in high-capacity devices. All XC4000X-specific routing resources are clearly identified throughout this section. Any resources not identified as XC4000X-specific are present in all XC4000 Series devices.

This section describes the varied routing resources available in XC4000 Series devices. The implementation software automatically assigns the appropriate resources based on the density and timing requirements of the design.

## Interconnect Overview

There are several types of interconnect.

- CLB routing is associated with each row and column of the CLB array.
- IOB routing forms a ring (called a VersaRing) around the outside of the CLB array. It connects the I/O with the internal logic blocks.

- Global routing consists of dedicated networks primarily designed to distribute clocks throughout the device with minimum delay and skew. Global routing can also be used for other high-fanout signals.

Five interconnect types are distinguished by the relative length of their segments: single-length lines, double-length lines, quad and octal lines (XC4000X only), and longlines. In the XC4000X, direct connects allow fast data flow between adjacent CLBs, and between IOBs and CLBs.

Extra routing is included in the IOB pad ring. The XC4000X also includes a ring of octal interconnect lines near the IOBs to improve pin-swapping and routing to locked pins.

XC4000E/X devices include two types of global buffers. These global buffers have different properties, and are intended for different purposes. They are discussed in detail later in this section.

## CLB Routing Connections

A high-level diagram of the routing resources associated with one CLB is shown in [Figure 25](#). The shaded arrows represent routing present only in XC4000X devices.

[Table 14](#) shows how much routing of each type is available in XC4000E and XC4000X CLB arrays. Clearly, very large designs, or designs with a great deal of interconnect, will route more easily in the XC4000X. Smaller XC4000E designs, typically requiring significantly less interconnect, do not require the additional routing.

[Figure 27 on page 30](#) is a detailed diagram of both the XC4000E and the XC4000X CLB, with associated routing. The shaded square is the programmable switch matrix, present in both the XC4000E and the XC4000X. The L-shaped shaded area is present only in XC4000X devices. As shown in the figure, the XC4000X block is essentially an XC4000E block with additional routing.

CLB inputs and outputs are distributed on all four sides, providing maximum routing flexibility. In general, the entire architecture is symmetrical and regular. It is well suited to established placement and routing algorithms. Inputs, outputs, and function generators can freely swap positions within a CLB to avoid routing congestion during the placement and routing operation.





**Figure 28: Single- and Double-Length Lines, with Programmable Switch Matrices (PSMs)**

### Double-Length Lines

The double-length lines consist of a grid of metal segments, each twice as long as the single-length lines: they run past two CLBs before entering a switch matrix. Double-length lines are grouped in pairs with the switch matrices staggered, so that each line goes through a switch matrix at every other row or column of CLBs (see [Figure 28](#)).

There are four vertical and four horizontal double-length lines associated with each CLB. These lines provide faster signal routing over intermediate distances, while retaining routing flexibility. Double-length lines are connected by way of the programmable switch matrices. Routing connectivity is shown in [Figure 27](#).

### Quad Lines (XC4000X only)

XC4000X devices also include twelve vertical and twelve horizontal quad lines per CLB row and column. Quad lines are four times as long as the single-length lines. They are interconnected via buffered switch matrices (shown as diamonds in [Figure 27 on page 30](#)). Quad lines run past four CLBs before entering a buffered switch matrix. They are grouped in fours, with the buffered switch matrices staggered, so that each line goes through a buffered switch matrix at every fourth CLB location in that row or column. (See [Figure 29](#).)

The buffered switch matrixes have four pins, one on each edge. All of the pins are bidirectional. Any pin can drive any or all of the other pins.

Each buffered switch matrix contains one buffer and six pass transistors. It resembles the programmable switch matrix shown in [Figure 26](#), with the addition of a programmable buffer. There can be up to two independent inputs



**Figure 29: Quad Lines (XC4000X only)**

and up to two independent outputs. Only one of the independent inputs can be buffered.

The place and route software automatically uses the timing requirements of the design to determine whether or not a quad line signal should be buffered. A heavily loaded signal is typically buffered, while a lightly loaded one is not. One scenario is to alternate buffers and pass transistors. This allows both vertical and horizontal quad lines to be buffered at alternating buffered switch matrices.

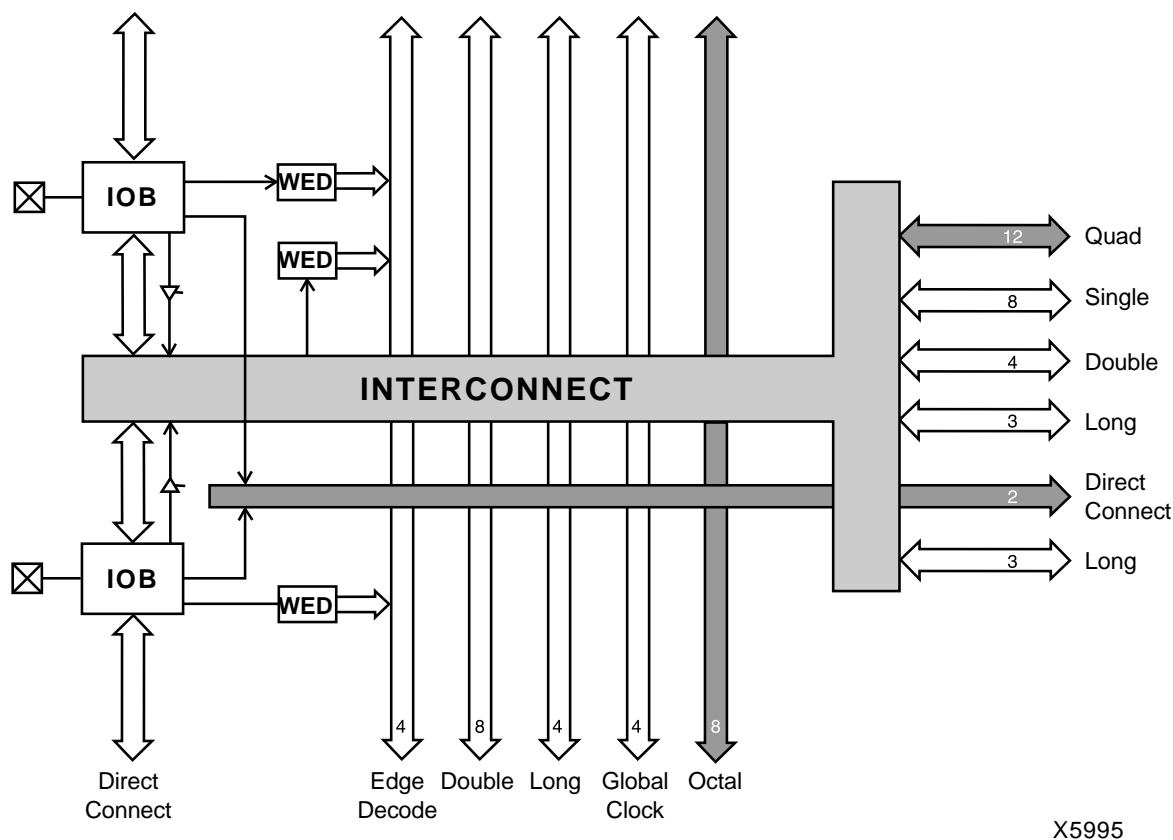
Due to the buffered switch matrices, quad lines are very fast. They provide the fastest available method of routing heavily loaded signals for long distances across the device.

### Longlines

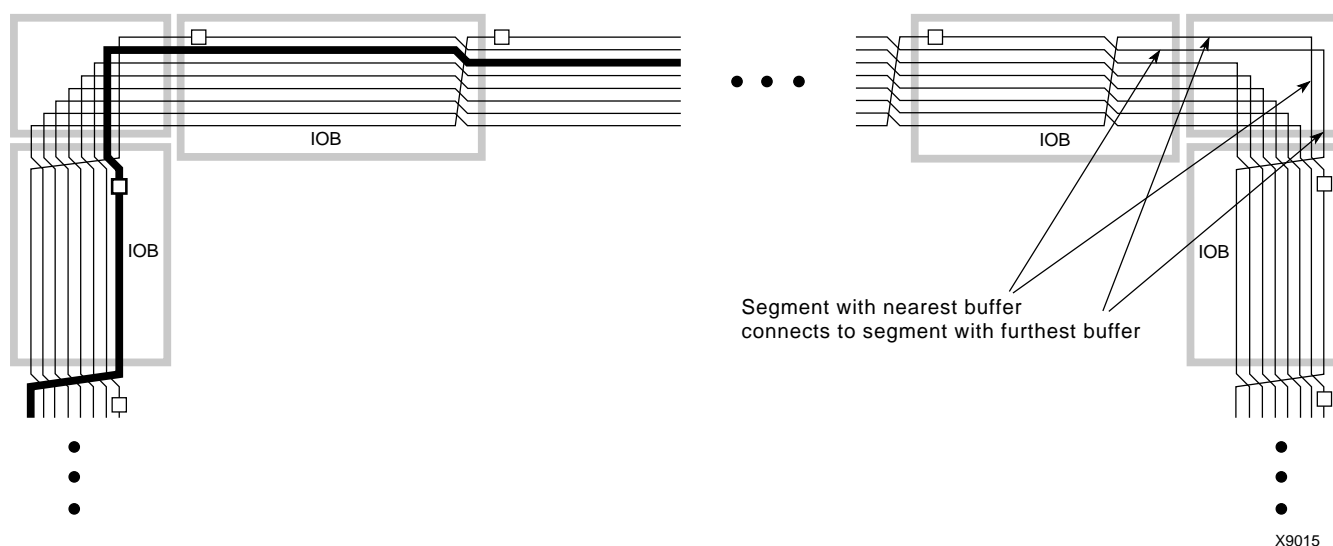
Longlines form a grid of metal interconnect segments that run the entire length or width of the array. Longlines are intended for high fan-out, time-critical signal nets, or nets that are distributed over long distances. In XC4000X devices, quad lines are preferred for critical nets, because the buffered switch matrices make them faster for high fan-out nets.

Two horizontal longlines per CLB can be driven by 3-state or open-drain drivers (TBUFs). They can therefore implement unidirectional or bidirectional buses, wide multiplexers, or wired-AND functions. (See [“Three-State Buffers” on page 26](#) for more details.)

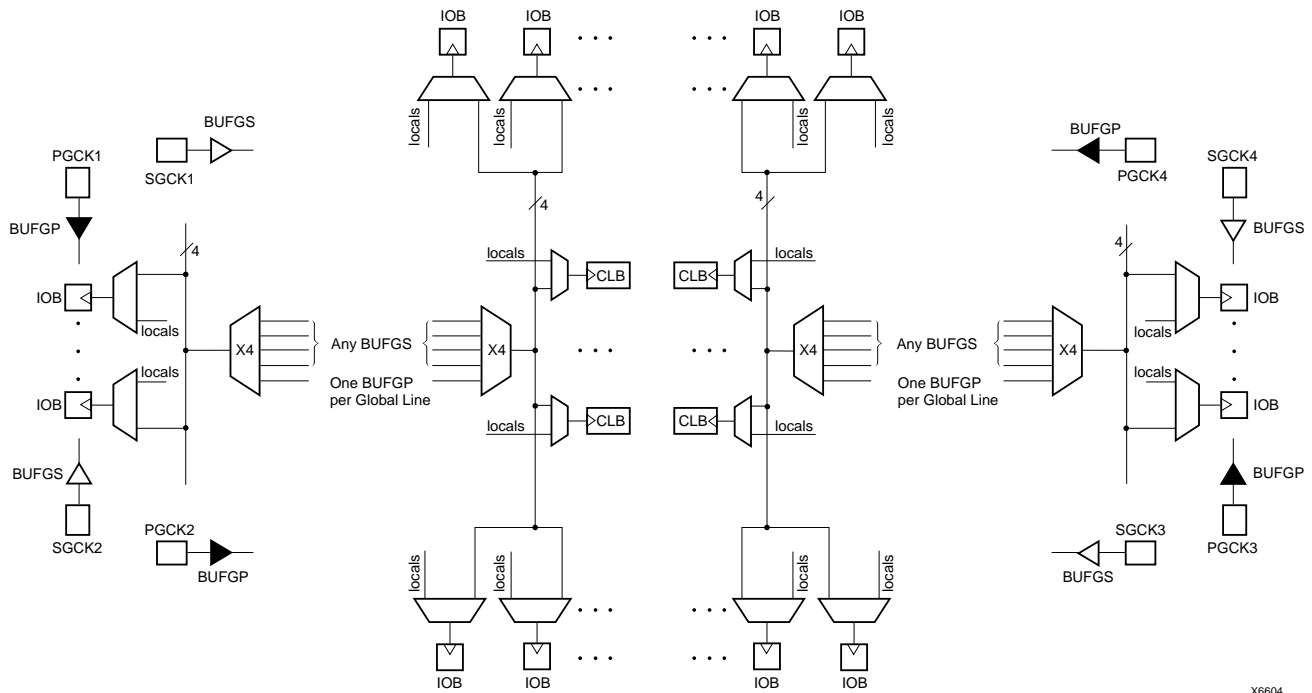
Each horizontal longline driven by TBUFs has either two (XC4000E) or eight (XC4000X) pull-up resistors. To activate these resistors, attach a PULLUP symbol to the long-line net. The software automatically activates the appropriate number of pull-ups. There is also a weak keeper at each end of these two horizontal longlines. This



**Figure 31: High-Level Routing Diagram of XC4000 Series VersaRing (Left Edge)**  
WED = Wide Edge Decoder, IOB = I/O Block (shaded arrows indicate XC4000X only)

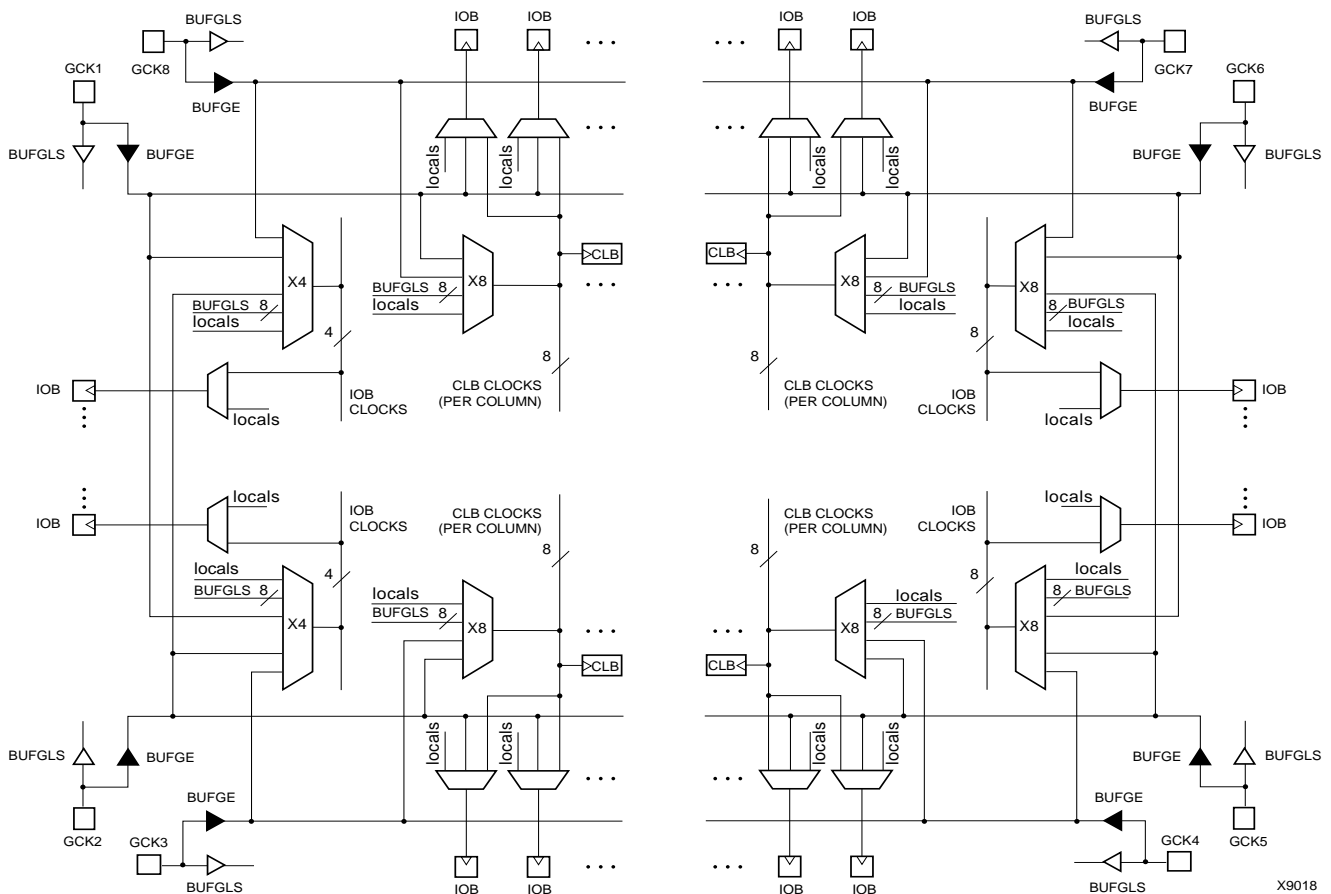


**Figure 32: XC4000X Octal I/O Routing**



X6604

**Figure 34: XC4000E Global Net Distribution**



X9018

**Figure 35: XC4000X Global Net Distribution**

### Global Nets and Buffers (XC4000X only)

Eight vertical longlines in each CLB column are driven by special global buffers. These longlines are in addition to the vertical longlines used for standard interconnect. The global lines are broken in the center of the array, to allow faster distribution and to minimize skew across the whole array. Each half-column global line has its own buffered multiplexer, as shown in [Figure 35](#). The top and bottom global lines cannot be connected across the center of the device, as this connection might introduce unacceptable skew. The top and bottom halves of the global lines must be separately driven — although they can be driven by the same global buffer.

The eight global lines in each CLB column can be driven by either of two types of global buffers. They can also be driven by internal logic, because they can be accessed by single, double, and quad lines at the top, bottom, half, and quarter points. Consequently, the number of different clocks that can be used simultaneously in an XC4000X device is very large.

There are four global lines feeding the IOBs at the left edge of the device. IOBs along the right edge have eight global lines. There is a single global line along the top and bottom edges with access to the IOBs. All IOB global lines are broken at the center. They cannot be connected across the center of the device, as this connection might introduce unacceptable skew.

IOB global lines can be driven from two types of global buffers, or from local interconnect. Alternatively, top and bottom IOBs can be clocked from the global lines in the adjacent CLB column.

Two different types of clock buffers are available in the XC4000X:

- Global Low-Skew Buffers (BUFGSL)
- Global Early Buffers (BUFGE)

Global Low-Skew Buffers are the standard clock buffers. They should be used for most internal clocking, whenever a large portion of the device must be driven.

Global Early Buffers are designed to provide a faster clock access, but CLB access is limited to one-fourth of the device. They also facilitate a faster I/O interface.

[Figure 35](#) is a conceptual diagram of the global net structure in the XC4000X.

Global Early buffers and Global Low-Skew buffers share a single pad. Therefore, the same IPAD symbol can drive one buffer of each type, in parallel. This configuration is particularly useful when using the Fast Capture latches, as described in ["IOB Input Signals" on page 20](#). Paired Global

Early and Global Low-Skew buffers share a common input; they cannot be driven by two different signals.

### Choosing an XC4000X Clock Buffer

The clocking structure of the XC4000X provides a large variety of features. However, it can be simple to use, without understanding all the details. The software automatically handles clocks, along with all other routing, when the appropriate clock buffer is placed in the design. In fact, if a buffer symbol called BUFG is placed, rather than a specific type of buffer, the software even chooses the buffer most appropriate for the design. The detailed information in this section is provided for those users who want a finer level of control over their designs.

If fine control is desired, use the following summary and [Table 15 on page 35](#) to choose an appropriate clock buffer.

- The simplest thing to do is to use a Global Low-Skew buffer.
- If a faster clock path is needed, try a BUFG. The software will first try to use a Global Low-Skew Buffer. If timing requirements are not met, a faster buffer will automatically be used.
- If a single quadrant of the chip is sufficient for the clocked logic, and the timing requires a faster clock than the Global Low-Skew buffer, use a Global Early buffer.

### Global Low-Skew Buffers

Each corner of the XC4000X device has two Global Low-Skew buffers. Any of the eight Global Low-Skew buffers can drive any of the eight vertical Global lines in a column of CLBs. In addition, any of the buffers can drive any of the four vertical lines accessing the IOBs on the left edge of the device, and any of the eight vertical lines accessing the IOBs on the right edge of the device. (See [Figure 36 on page 38](#).)

IOBs at the top and bottom edges of the device are accessed through the vertical Global lines in the CLB array, as in the XC4000E. Any Global Low-Skew buffer can, therefore, access every IOB and CLB in the device.

The Global Low-Skew buffers can be driven by either semi-dedicated pads or internal logic.

To use a Global Low-Skew buffer, instantiate a BUFGSL element in a schematic or in HDL code. If desired, attach a LOC attribute or property to direct placement to the designated location. For example, attach a LOC=T attribute or property to direct that a BUFGSL be placed in one of the two Global Low-Skew buffers on the top edge of the device, or a LOC=TR to indicate the Global Low-Skew buffer on the top edge of the device, on the right.

**Table 16: Pin Descriptions (Continued)**

Pin Name	I/O During Config.	I/O After Config.	Pin Description
$\overline{CS0}$ , CS1, $\overline{WS}$ , $\overline{RS}$	I	I/O	These four inputs are used in Asynchronous Peripheral mode. The chip is selected when $\overline{CS0}$ is Low and CS1 is High. While the chip is selected, a Low on Write Strobe ( $\overline{WS}$ ) loads the data present on the D0 - D7 inputs into the internal data buffer. A Low on Read Strobe ( $\overline{RS}$ ) changes D7 into a status output — High if Ready, Low if Busy — and drives D0 - D6 High. In Express mode, CS1 is used as a serial-enable signal for daisy-chaining. $\overline{WS}$ and $\overline{RS}$ should be mutually exclusive, but if both are Low simultaneously, the Write Strobe overrides. After configuration, these are user-programmable I/O pins.
A0 - A17	O	I/O	During Master Parallel configuration, these 18 output pins address the configuration EPROM. After configuration, they are user-programmable I/O pins.
A18 - A21 (XC4003XL to XC4085XL)	O	I/O	During Master Parallel configuration with an XC4000X master, these 4 output pins add 4 more bits to address the configuration EPROM. After configuration, they are user-programmable I/O pins. (See Master Parallel Configuration section for additional details.)
D0 - D7	I	I/O	During Master Parallel and Peripheral configuration, these eight input pins receive configuration data. After configuration, they are user-programmable I/O pins.
DIN	I	I/O	During Slave Serial or Master Serial configuration, DIN is the serial configuration data input receiving data on the rising edge of CCLK. During Parallel configuration, DIN is the D0 input. After configuration, DIN is a user-programmable I/O pin.
DOUT	O	I/O	During configuration in any mode but Express mode, DOUT is the serial configuration data output that can drive the DIN of daisy-chained slave FPGAs. DOUT data changes on the falling edge of CCLK, one-and-a-half CCLK periods after it was received at the DIN input. In Express mode for XC4000E and XC4000X only, DOUT is the status output that can drive the CS1 of daisy-chained FPGAs, to enable and disable downstream devices. After configuration, DOUT is a user-programmable I/O pin.
<b>Unrestricted User-Programmable I/O Pins</b>			
I/O	Weak Pull-up	I/O	These pins can be configured to be input and/or output after configuration is completed. Before configuration is completed, these pins have an internal high-value pull-up resistor (25 k $\Omega$ - 100 k $\Omega$ ) that defines the logic level as High.

## Boundary Scan

The 'bed of nails' has been the traditional method of testing electronic assemblies. This approach has become less appropriate, due to closer pin spacing and more sophisticated assembly methods like surface-mount technology and multi-layer boards. The IEEE Boundary Scan Standard 1149.1 was developed to facilitate board-level testing of electronic assemblies. Design and test engineers can imbed a standard test logic structure in their device to achieve high fault coverage for I/O and internal logic. This structure is easily implemented with a four-pin interface on any boundary scan-compatible IC. IEEE 1149.1-compatible devices may be serial daisy-chained together, connected in parallel, or a combination of the two.

The XC4000 Series implements IEEE 1149.1-compatible BYPASS, PRELOAD/SAMPLE and EXTEST boundary scan instructions. When the boundary scan configuration option is selected, three normal user I/O pins become dedicated inputs for these functions. Another user output pin becomes the dedicated boundary scan output. The details

of how to enable this circuitry are covered later in this section.

By exercising these input signals, the user can serially load commands and data into these devices to control the driving of their outputs and to examine their inputs. This method is an improvement over bed-of-nails testing. It avoids the need to over-drive device outputs, and it reduces the user interface to four pins. An optional fifth pin, a reset for the control logic, is described in the standard but is not implemented in Xilinx devices.

The dedicated on-chip logic implementing the IEEE 1149.1 functions includes a 16-state machine, an instruction register and a number of data registers. The functional details can be found in the IEEE 1149.1 specification and are also discussed in the Xilinx application note XAPP 017: "*Boundary Scan in XC4000 Devices*."

Figure 40 on page 43 shows a simplified block diagram of the XC4000E Input/Output Block with boundary scan implemented. XC4000X boundary scan logic is identical.



is passed through and is captured by each FPGA when it recognizes the 0010 preamble. Following the length-count data, each FPGA outputs a High on DOUT until it has received its required number of data frames.

After an FPGA has received its configuration data, it passes on any additional frame start bits and configuration data on DOUT. When the total number of configuration clocks applied after memory initialization equals the value of the 24-bit length count, the FPGAs begin the start-up sequence and become operational together. FPGA I/O are normally released two CCLK cycles after the last configuration bit is received. **Figure 47 on page 53** shows the start-up timing for an XC4000 Series device.

The daisy-chained bitstream is not simply a concatenation of the individual bitstreams. The PROM file formatter must be used to combine the bitstreams for a daisy-chained configuration.

### Multi-Family Daisy Chain

All Xilinx FPGAs of the XC2000, XC3000, and XC4000 Series use a compatible bitstream format and can, therefore, be connected in a daisy chain in an arbitrary sequence. There is, however, one limitation. The lead device must belong to the highest family in the chain. If the chain contains XC4000 Series devices, the master normally cannot be an XC2000 or XC3000 device.

The reason for this rule is shown in **Figure 47 on page 53**. Since all devices in the chain store the same length count value and generate or receive one common sequence of CCLK pulses, they all recognize length-count match on the same CCLK edge, as indicated on the left edge of **Figure 47**. The master device then generates additional CCLK pulses until it reaches its finish point F. The different families generate or require different numbers of additional CCLK pulses until they reach F. Not reaching F means that the device does not really finish its configuration, although DONE may have gone High, the outputs became active, and the internal reset was released. For the XC4000 Series device, not reaching F means that readback cannot be ini-

tiated and most boundary scan instructions cannot be used.

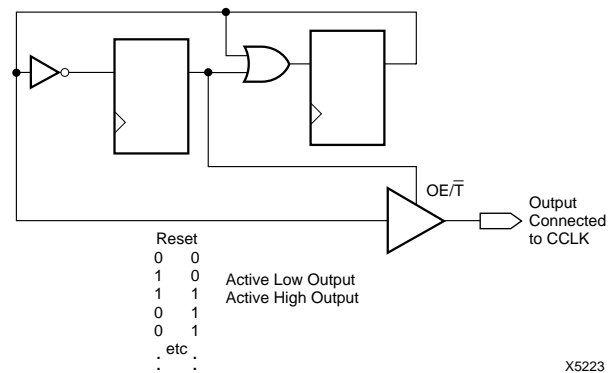
The user has some control over the relative timing of these events and can, therefore, make sure that they occur at the proper time and the finish point F is reached. Timing is controlled using options in the bitstream generation software.

### XC3000 Master with an XC4000 Series Slave

Some designers want to use an inexpensive lead device in peripheral mode and have the more precious I/O pins of the XC4000 Series devices all available for user I/O. **Figure 44** provides a solution for that case.

This solution requires one CLB, one IOB and pin, and an internal oscillator with a frequency of up to 5 MHz as a clock source. The XC3000 master device must be configured with late Internal Reset, which is the default option.

One CLB and one IOB in the lead XC3000-family device are used to generate the additional CCLK pulse required by the XC4000 Series devices. When the lead device removes the internal RESET signal, the 2-bit shift register responds to its clock input and generates an active Low output signal for the duration of the subsequent clock period. An external connection between this output and CCLK thus creates the extra CCLK pulse.



**Figure 44: CCLK Generation for XC3000 Master Driving an XC4000 Series Slave**

## Setting CCLK Frequency

For Master modes, CCLK can be generated in either of two frequencies. In the default slow mode, the frequency ranges from 0.5 MHz to 1.25 MHz for XC4000E and XC4000EX devices and from 0.6 MHz to 1.8 MHz for XC4000XL devices. In fast CCLK mode, the frequency ranges from 4 MHz to 10 MHz for XC4000E/EX devices and from 5 MHz to 15 MHz for XC4000XL devices. The frequency is selected by an option when running the bitstream generation software. If an XC4000 Series Master is driving an XC3000- or XC2000-family slave, slow CCLK mode must be used. In addition, an XC4000XL device driving a XC4000E or XC4000EX should use slow mode. Slow mode is the default.

**Table 19: XC4000 Series Data Stream Formats**

Data Type	All Other Modes (D0...)
Fill Byte	11111111b
Preamble Code	0010b
Length Count	COUNT(23:0)
Fill Bits	1111b
Start Field	0b
Data Frame	DATA(n-1:0)
CRC or Constant Field Check	xxxx (CRC) or 0110b
Extend Write Cycle	—
Postamble	01111111b
Start-Up Bytes	xxh
Legend:	
Not shaded	Once per bitstream
Light	Once per data frame
Dark	Once per device

## Data Stream Format

The data stream (“bitstream”) format is identical for all configuration modes.

The data stream formats are shown in [Table 19](#). Bit-serial data is read from left to right, and byte-parallel data is effectively assembled from this serial bitstream, with the first bit in each byte assigned to D0.

The configuration data stream begins with a string of eight ones, a preamble code, followed by a 24-bit length count and a separator field of ones. This header is followed by the actual configuration data in frames. The length and number of frames depends on the device type (see [Table 20](#) and [Table 21](#)). Each frame begins with a start field and ends with an error check. A postamble code is required to signal the end of data for a single device. In all cases, additional start-up bytes of data are required to provide four clocks for the startup sequence at the end of configuration. Long daisy chains require additional startup bytes to shift the last data through the chain. All startup bytes are don't-cares; these bytes are not included in bitstreams created by the Xilinx software.

A selection of CRC or non-CRC error checking is allowed by the bitstream generation software. The non-CRC error checking tests for a designated end-of-frame field for each frame. For CRC error checking, the software calculates a running CRC and inserts a unique four-bit partial check at the end of each frame. The 11-bit CRC check of the last frame of an FPGA includes the last seven data bits.

Detection of an error results in the suspension of data loading and the pulling down of the  $\overline{\text{INIT}}$  pin. In Master modes, CCLK and address signals continue to operate externally. The user must detect  $\overline{\text{INIT}}$  and initialize a new configuration by pulsing the  $\overline{\text{PROGRAM}}$  pin Low or cycling Vcc.

Table 20: XC4000E Program Data

Device	XC4003E	XC4005E	XC4006E	XC4008E	XC4010E	XC4013E	XC4020E	XC4025E
Max Logic Gates	3,000	5,000	6,000	8,000	10,000	13,000	20,000	25,000
CLBs (Row x Col.)	100 (10 x 10)	196 (14 x 14)	256 (16 x 16)	324 (18 x 18)	400 (20 x 20)	576 (24 x 24)	784 (28 x 28)	1,024 (32 x 32)
IOBs	80	112	128	144	160	192	224	256
Flip-Flops	360	616	768	936	1,120	1,536	2,016	2,560
Bits per Frame	126	166	186	206	226	266	306	346
Frames	428	572	644	716	788	932	1,076	1,220
Program Data	53,936	94,960	119,792	147,504	178,096	247,920	329,264	422,128
PROM Size (bits)	53,984	95,008	119,840	147,552	178,144	247,968	329,312	422,176

- Notes:
- Bits per Frame = (10 x number of rows) + 7 for the top + 13 for the bottom + 1 + 1 start bit + 4 error check bits  
 Number of Frames = (36 x number of columns) + 26 for the left edge + 41 for the right edge + 1  
 Program Data = (Bits per Frame x Number of Frames) + 8 postamble bits  
 PROM Size = Program Data + 40 (header) + 8
  - The user can add more "one" bits as leading dummy bits in the header, or, if CRC = off, as trailing dummy bits at the end of any frame, following the four error check bits. However, the Length Count value **must** be adjusted for all such extra "one" bits, even for extra leading ones at the beginning of the header.

Table 21: XC4000EX/XL Program Data

Device	XC4002XL	XC4005	XC4010	XC4013	XC4020	XC4028	XC4036	XC4044	XC4052	XC4062	XC4085
Max Logic Gates	2,000	5,000	10,000	13,000	20,000	28,000	36,000	44,000	52,000	62,000	85,000
CLBs (Row x Column)	64 (8 x 8)	196 (14 x 14)	400 (20 x 20)	576 (24 x 24)	784 (28 x 28)	1,024 (32 x 32)	1,296 (36 x 36)	1,600 (40 x 40)	1,936 (44 x 44)	2,304 (48 x 48)	3,136 (56 x 56)
IOBs	64	112	160	192	224	256	288	320	352	384	448
Flip-Flops	256	616	1,120	1,536	2,016	2,560	3,168	3,840	4,576	5,376	7,168
Bits per Frame	133	205	277	325	373	421	469	517	565	613	709
Frames	459	741	1,023	1,211	1,399	1,587	1,775	1,963	2,151	2,339	2,715
Program Data	61,052	151,910	283,376	393,580	521,832	668,124	832,480	1,014,876	1,215,320	1,433,804	1,924,940
PROM Size (bits)	61,104	151,960	283,424	393,632	521,880	668,172	832,528	1,014,924	1,215,368	1,433,852	1,924,992

- Notes:
- Bits per frame = (13 x number of rows) + 9 for the top + 17 for the bottom + 8 + 1 start bit + 4 error check bits.  
 Frames = (47 x number of columns) + 27 for the left edge + 52 for the right edge + 4.  
 Program data = (bits per frame x number of frames) + 5 postamble bits.  
 PROM size = (program data + 40 header bits + 8 start bits) rounded up to the nearest byte.
  - The user can add more "one" bits as leading dummy bits in the header, or, if CRC = off, as trailing dummy bits at the end of any frame, following the four error check bits. However, the Length Count value must be adjusted for all such extra "one" bits, even for extra leading "ones" at the beginning of the header.

## Cyclic Redundancy Check (CRC) for Configuration and Readback

The Cyclic Redundancy Check is a method of error detection in data transmission applications. Generally, the transmitting system performs a calculation on the serial bitstream. The result of this calculation is tagged onto the data stream as additional check bits. The receiving system performs an identical calculation on the bitstream and compares the result with the received checksum.

Each data frame of the configuration bitstream has four error bits at the end, as shown in [Table 19](#). If a frame data error is detected during the loading of the FPGA, the con-

figuration process with a potentially corrupted bitstream is terminated. The FPGA pulls the  $\overline{\text{INIT}}$  pin Low and goes into a Wait state.

During Readback, 11 bits of the 16-bit checksum are added to the end of the Readback data stream. The checksum is computed using the CRC-16 CCITT polynomial, as shown in [Figure 45](#). The checksum consists of the 11 most significant bits of the 16-bit code. A change in the checksum indicates a change in the Readback bitstream. A comparison to a previous checksum is meaningful only if the readback data is independent of the current device state. CLB outputs should not be included (Read Capture option not

used), and if RAM is present, the RAM content must be unchanged.

Statistically, one error out of 2048 might go undetected.

## Configuration Sequence

There are four major steps in the XC4000 Series power-up configuration sequence.

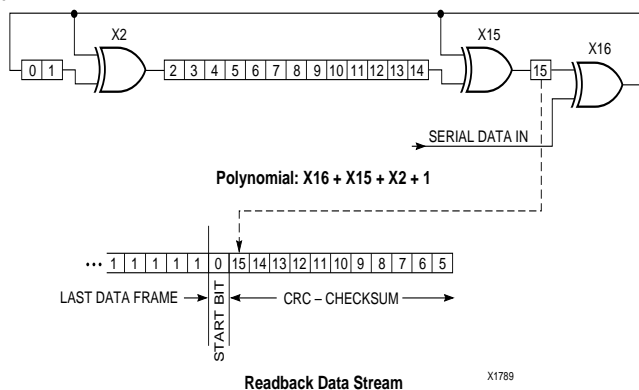
- Configuration Memory Clear
- Initialization
- Configuration
- Start-Up

The full process is illustrated in Figure 46.

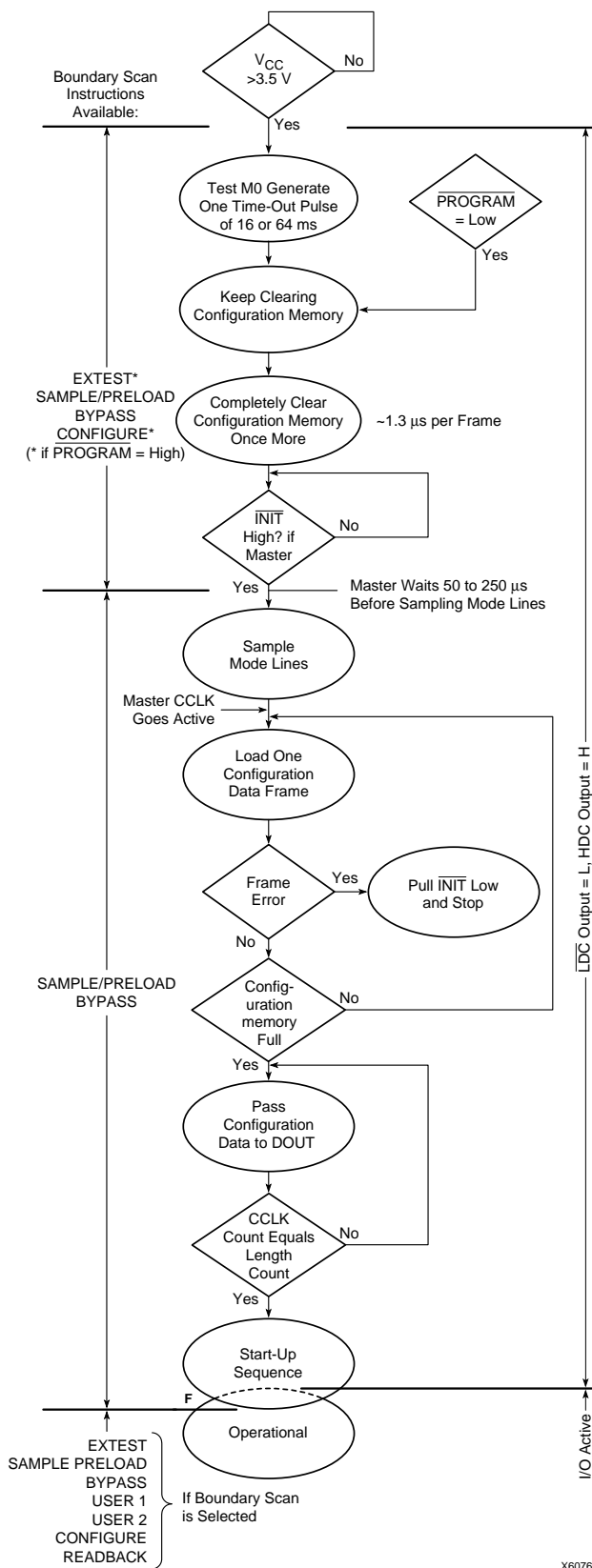
### Configuration Memory Clear

When power is first applied or is reapplied to an FPGA, an internal circuit forces initialization of the configuration logic. When  $V_{CC}$  reaches an operational level, and the circuit passes the write and read test of a sample pair of configuration bits, a time delay is started. This time delay is nominally 16 ms, and up to 10% longer in the low-voltage devices. The delay is four times as long when in Master Modes (M0 Low), to allow ample time for all slaves to reach a stable  $V_{CC}$ . When all  $\overline{INIT}$  pins are tied together, as recommended, the longest delay takes precedence. Therefore, devices with different time delays can easily be mixed and matched in a daisy chain.

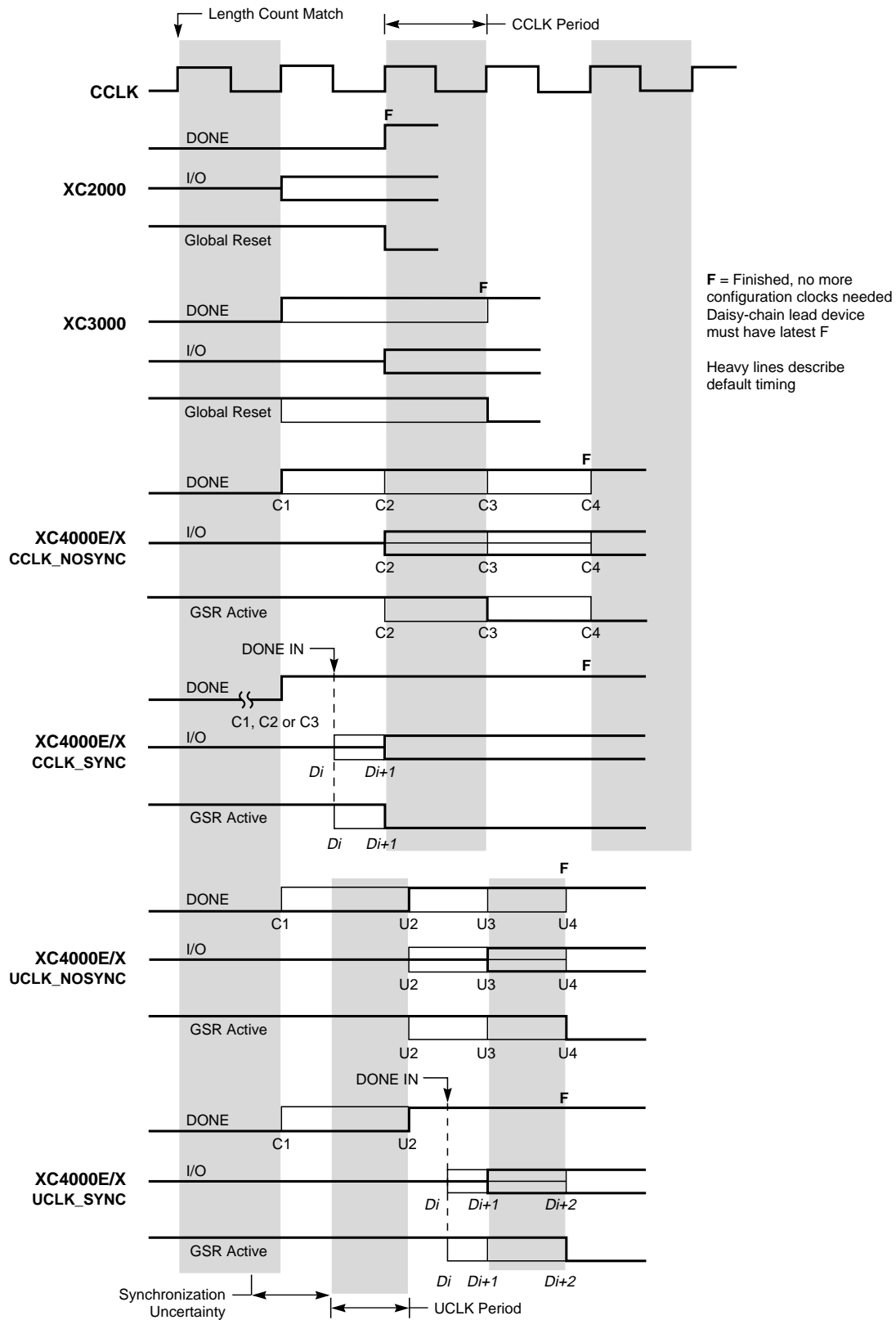
This delay is applied only on power-up. It is not applied when re-configuring an FPGA by pulsing the  $\overline{PROGRAM}$  pin



**Figure 45: Circuit for Generating CRC-16**



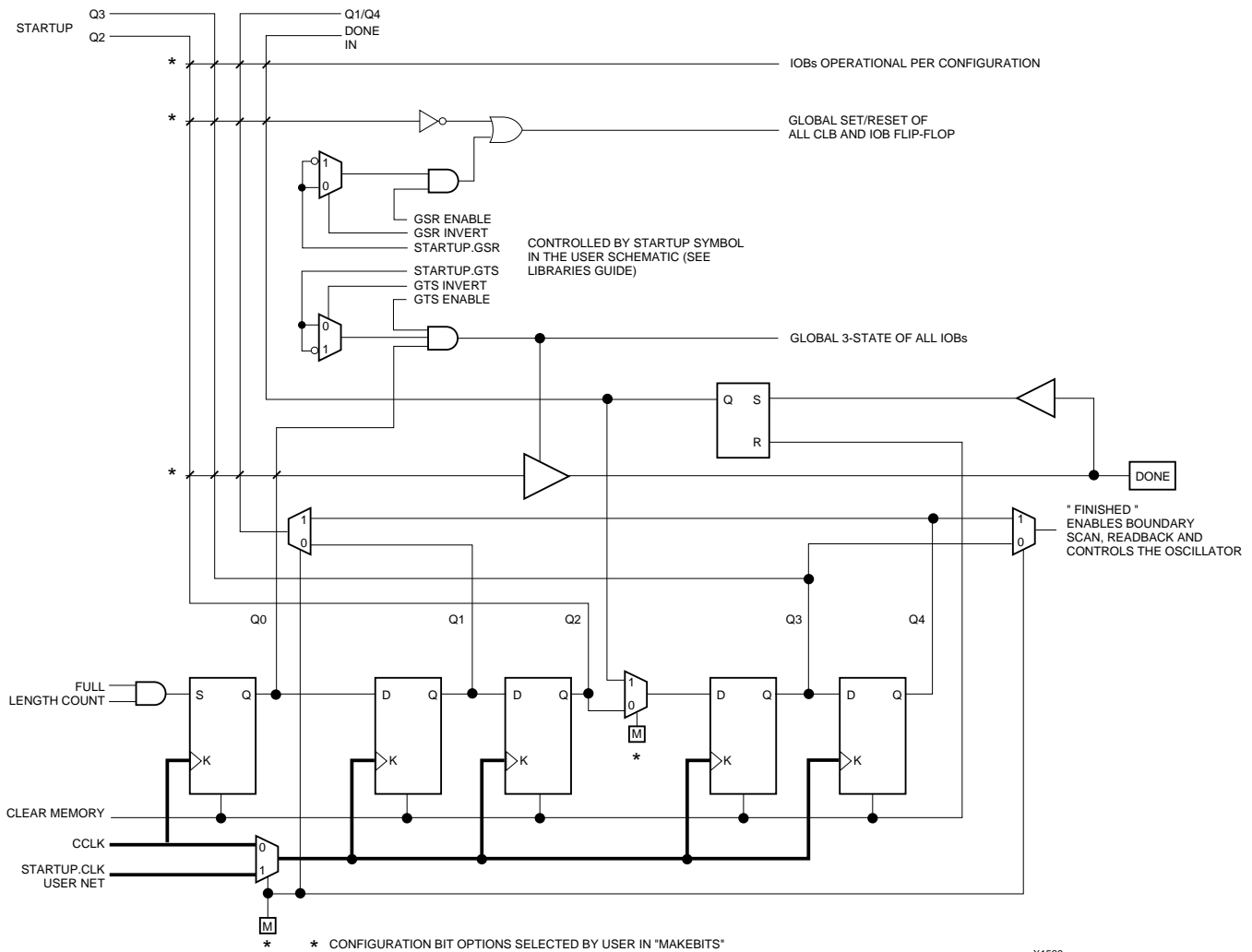
**Figure 46: Power-up Configuration Sequence**



X9024

**Figure 47: Start-up Timing**





**Figure 48: Start-up Logic**

## Readback

The user can read back the content of configuration memory and the level of certain internal nodes without interfering with the normal operation of the device.

Readback not only reports the downloaded configuration bits, but can also include the present state of the device, represented by the content of all flip-flops and latches in CLBs and IOBs, as well as the content of function generators used as RAMs.

Note that in XC4000 Series devices, configuration data is *not* inverted with respect to configuration as it is in XC2000 and XC3000 families.

XC4000 Series Readback does not use any dedicated pins, but uses four internal nets (RDBK.TRIG, RDBK.DATA, RDBK.RIP and RDBK.CLK) that can be routed to any IOB. To access the internal Readback signals, place the READ-

BACK library symbol and attach the appropriate pad symbols, as shown in [Figure 49](#).

After Readback has been initiated by a High level on RDBK.TRIG after configuration, the RDBK.RIP (Read In Progress) output goes High on the next rising edge of RDBK.CLK. Subsequent rising edges of this clock shift out Readback data on the RDBK.DATA net.

Readback data does not include the preamble, but starts with five dummy bits (all High) followed by the Start bit (Low) of the first frame. The first two data bits of the first frame are always High.

Each frame ends with four error check bits. They are read back as High. The last seven bits of the last frame are also read back as High. An additional Start bit (Low) and an 11-bit Cyclic Redundancy Check (CRC) signature follow, before RDBK.RIP returns Low.

**Table 22: Pin Functions During Configuration**

CONFIGURATION MODE <M2:M1:M0>						USER OPERATION
SLAVE SERIAL <1:1:1>	MASTER SERIAL <0:0:0>	SYNCH. PERIPHERAL <0:1:1>	ASYNCH. PERIPHERAL <1:0:1>	MASTER PARALLEL DOWN <1:1:0>	MASTER PARALLEL UP <1:0:0>	
M2(HIGH) (I)	M2(LOW) (I)	M2(LOW) (I)	M2(HIGH) (I)	M2(HIGH) (I)	M2(HIGH) (I)	(I)
M1(HIGH) (I)	M1(LOW) (I)	M1(HIGH) (I)	M1(LOW) (I)	M1(HIGH) (I)	M1(LOW) (I)	(O)
M0(HIGH) (I)	M0(LOW) (I)	M0(HIGH) (I)	M0(HIGH) (I)	M0(LOW) (I)	M0(LOW) (I)	(I)
HDC (HIGH)	HDC (HIGH)	HDC (HIGH)	HDC (HIGH)	HDC (HIGH)	HDC (HIGH)	I/O
LDC (LOW)	LDC (LOW)	LDC (LOW)	LDC (LOW)	LDC (LOW)	LDC (LOW)	I/O
INIT	INIT	INIT	INIT	INIT	INIT	I/O
DONE	DONE	DONE	DONE	DONE	DONE	DONE
PROGRAM (I)	PROGRAM (I)	PROGRAM (I)	PROGRAM (I)	PROGRAM (I)	PROGRAM (I)	PROGRAM
CCLK (I)	CCLK (O)	CCLK (I)	CCLK (O)	CCLK (O)	CCLK (O)	CCLK (I)
		RDY/BUSY (O)	RDY/BUSY (O)	RCLK (O)	RCLK (O)	I/O
			RS (I)			I/O
			CS0 (I)			I/O
		DATA 7 (I)	DATA 7 (I)	DATA 7 (I)	DATA 7 (I)	I/O
		DATA 6 (I)	DATA 6 (I)	DATA 6 (I)	DATA 6 (I)	I/O
		DATA 5 (I)	DATA 5 (I)	DATA 5 (I)	DATA 5 (I)	I/O
		DATA 4 (I)	DATA 4 (I)	DATA 4 (I)	DATA 4 (I)	I/O
		DATA 3 (I)	DATA 3 (I)	DATA 3 (I)	DATA 3 (I)	I/O
		DATA 2 (I)	DATA 2 (I)	DATA 2 (I)	DATA 2 (I)	I/O
		DATA 1 (I)	DATA 1 (I)	DATA 1 (I)	DATA 1 (I)	I/O
DIN (I)	DIN (I)	DATA 0 (I)	DATA 0 (I)	DATA 0 (I)	DATA 0 (I)	I/O
DOUT	DOUT	DOUT	DOUT	DOUT	DOUT	SGCK4-GCK6-I/O
TDI	TDI	TDI	TDI	TDI	TDI	TDI-I/O
TCK	TCK	TCK	TCK	TCK	TCK	TCK-I/O
TMS	TMS	TMS	TMS	TMS	TMS	TMS-I/O
TDO	TDO	TDO	TDO	TDO	TDO	TDO-(O)
			WS (I)	A0	A0	I/O
				A1	A1	PGCK4-GCK7-I/O
			CS1	A2	A2	I/O
				A3	A3	I/O
				A4	A4	I/O
				A5	A5	I/O
				A6	A6	I/O
				A7	A7	I/O
				A8	A8	I/O
				A9	A9	I/O
				A10	A10	I/O
				A11	A11	I/O
				A12	A12	I/O
				A13	A13	I/O
				A14	A14	I/O
				A15	A15	SGCK1-GCK8-I/O
				A16	A16	PGCK1-GCK1-I/O
				A17	A17	I/O
				A18*	A18*	I/O
				A19*	A19*	I/O
				A20*	A20*	I/O
				A21*	A21*	I/O
						ALL OTHERS

## Asynchronous Peripheral Mode

## Write to FPGA

Asynchronous Peripheral mode uses the trailing edge of the logic AND condition of  $\overline{WS}$  and  $\overline{CS0}$  being Low and RS and CS1 being High to accept byte-wide data from a micro-processor bus. In the lead FPGA, this data is loaded into a double-buffered UART-like parallel-to-serial converter and is serially shifted into the internal logic.

The lead FPGA presents the preamble data (and all data that overflows the lead device) on its DOUT pin. The RDY/BUSY output from the lead FPGA acts as a handshake signal to the microprocessor. RDY/BUSY goes Low when a byte has been received, and goes High again when the byte-wide input buffer has transferred its information into the shift register, and the buffer is ready to receive new data. A new write may be started immediately, as soon as the RDY/BUSY output has gone Low, acknowledging receipt of the previous data. Write may not be terminated until RDY/BUSY is High again for one CCLK period. Note that RDY/BUSY is pulled High with a high-impedance pull-up prior to INIT going High.

The length of the  $\overline{\text{BUSY}}$  signal depends on the activity in the UART. If the shift register was empty when the new byte was received, the  $\overline{\text{BUSY}}$  signal lasts for only two CCLK periods. If the shift register was still full when the new byte was received, the  $\overline{\text{BUSY}}$  signal can be as long as nine CCLK periods.

Note that after the last byte has been entered, only seven of its bits are shifted out. CCLK remains High with DOUT equal to bit 6 (the next-to-last bit) of the last byte entered.

The READY/BUSY handshake can be ignored if the delay from any one Write to the end of the next Write is guaranteed to be longer than 10 CCLK periods.

**Status Read**

The logic AND condition of the  $\overline{\text{CS0}}$ , CS1 and  $\overline{\text{RS}}$  inputs puts the device status on the Data bus.

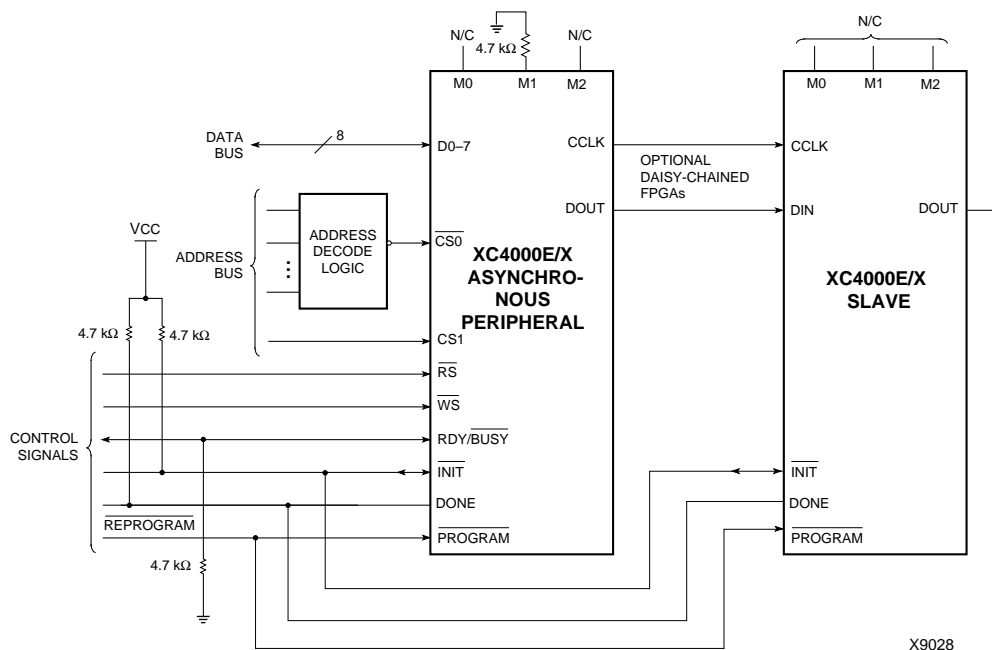
- D7 High indicates Ready
- D7 Low indicates Busy
- D0 through D6 go unconditionally High

It is mandatory that the whole start-up sequence be started and completed by one byte-wide input. Otherwise, the pins used as Write Strobe or Chip Enable might become active outputs and interfere with the final byte transfer. If this transfer does not occur, the start-up sequence is not completed all the way to the finish (point F in [Figure 47 on page 53](#)).

In this case, at worst, the internal reset is not released. At best, Readback and Boundary Scan are inhibited. The length-count value, as generated by the XACT<sup>step</sup> software, ensures that these problems never occur.

Although RDY/ $\overline{\text{BUSY}}$  is brought out as a separate signal, microprocessors can more easily read this information on one of the data lines. For this purpose, D7 represents the RDY/ $\overline{\text{BUSY}}$  status when  $\overline{\text{RS}}$  is Low,  $\overline{\text{WS}}$  is High, and the two chip select lines are both active.

Asynchronous Peripheral mode is selected by a <101> on the mode pins (M2, M1, M0).



**Figure 58: Asynchronous Peripheral Mode Circuit Diagram**