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Operating Temperature	-
Mounting Type	-
Package / Case	-
Supplier Device Package	-
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C165

Revision History: **2000-12** V2.0

Previous Version:	1998-12	Update 0.5 μ technology
	01.96	3 Volt Addendum
	07.95	25 MHz Addendum
	09.94	Data Sheet

Page	Subjects (major changes since last revision)
All	Converted to Infineon layout
2	ROM derivatives removed, 25-MHz derivatives and 3 V derivatives included
6ff	Pin numbers for TQFP added
14	Address window arbitration and master/slave mode introduced
32	New standard layout for section “Absolute Maximum Ratings”
33	Section “Operating Conditions” added
34f	Parameter “ $\overline{\text{RSTIN}}$ pullup” replaced by “ $\overline{\text{RSTIN}}$ current”
36f	DC Characteristics for reduced supply voltage added
38f	Separate specification for power consumption with greatly improved values
40ff	Description of clock generation improved
45 , 55 , 65	Timing adapted to 25 MHz
48 , 58 , 66	Timing for reduced supply voltage added

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This document describes several derivatives of the C165 group. **Table 1** enumerates these derivatives and summarizes the differences. As this document refers to all of these derivatives, some descriptions may not apply to a specific product.

Table 1 C165 Derivative Synopsis

Derivative ¹⁾	Max. Operating Frequency	Operating Voltage	Package
SAF-C165-LM	20 MHz	4.5 to 5.5 V	MQFP-100
SAB-C165-LM	20 MHz	4.5 to 5.5 V	MQFP-100
SAF-C165-L25M	25 MHz	4.5 to 5.5 V	MQFP-100
SAB-C165-L25M	25 MHz	4.5 to 5.5 V	MQFP-100
SAF-C165-LF	20 MHz	4.5 to 5.5 V	TQFP-100
SAB-C165-LF	20 MHz	4.5 to 5.5 V	TQFP-100
SAF-C165-L25F	25 MHz	4.5 to 5.5 V	TQFP-100
SAB-C165-L25F	25 MHz	4.5 to 5.5 V	TQFP-100
SAF-C165-LM3V	20 MHz	3.0 to 3.6 V	MQFP-100
SAB-C165-LM3V	20 MHz	3.0 to 3.6 V	MQFP-100
SAF-C165-LF3V	20 MHz	3.0 to 3.6 V	TQFP-100
SAB-C165-LF3V	20 MHz	3.0 to 3.6 V	TQFP-100

¹⁾ This Data Sheet is valid for devices starting with and including design step HA.

For simplicity all versions are referred to by the term **C165** throughout this document.

Ordering Information

The ordering code for Infineon microcontrollers provides an exact reference to the required product. This ordering code identifies:

- the derivative itself, i.e. its function set, the temperature range, and the supply voltage
- the package and the type of delivery.

For the available ordering codes for the C165 please refer to the “**Product Catalog Microcontrollers**”, which summarizes all available microcontroller variants.

Note: The ordering codes for Mask-ROM versions are defined for each product after verification of the respective ROM code.

Pin Configuration TQFP Package
(top view)

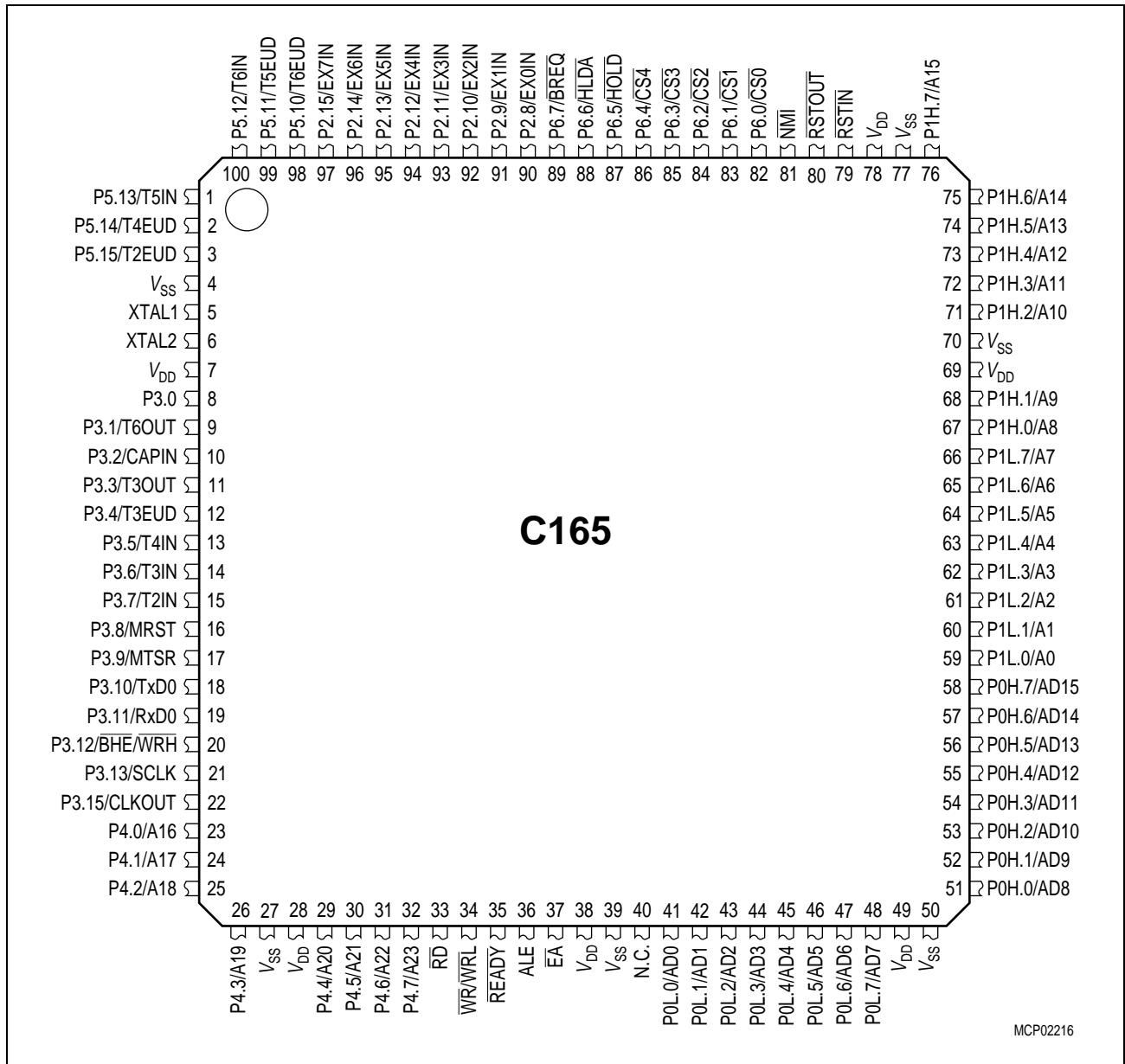


Figure 2

Table 2 Pin Definitions and Functions (cont'd)

Symbol	Pin Nr TQFP	Pin Nr MQFP	Input Outp.	Function
P4			IO	Port 4 is an 8-bit bidirectional I/O port. It is bit-wise programmable for input or output via direction bits. For a pin configured as input, the output driver is put into high-impedance state. Port 4 can be used to output the segment address lines:
P4.0	23	25	O	A16 Least Significant Segment Address Line
P4.1	24	26	O	A17 Segment Address Line
P4.2	25	27	O	A18 Segment Address Line
P4.3	26	28	O	A19 Segment Address Line
P4.4	29	31	O	A20 Segment Address Line
P4.5	30	32	O	A21 Segment Address Line
P4.6	31	33	O	A22 Segment Address Line
P4.7	32	34	O	A23 Most Significant Segment Address Line
\overline{RD}	33	35	O	External Memory Read Strobe. \overline{RD} is activated for every external instruction or data read access.
$\overline{WR}/$ \overline{WRL}	34	36	O	External Memory Write Strobe. In \overline{WR} -mode this pin is activated for every external data write access. In \overline{WRL} -mode this pin is activated for low byte data write accesses on a 16-bit bus, and for every data write access on an 8-bit bus. See WRCFG in register SYSCON for mode selection.
\overline{READY}	35	37	I	Ready Input. When the Ready function is enabled, a high level at this pin during an external memory access will force the insertion of memory cycle waitstates until the pin returns to a low level. An internal pullup device holds this pin high when nothing is driving it.
ALE	36	38	O	Address Latch Enable Output. Can be used for latching the address into external memory or an address latch in the multiplexed bus modes.
\overline{EA}	37	39	I	External Access Enable pin. A low level at this pin during and after Reset forces the C165 to begin instruction execution out of external memory. A high level forces execution out of the internal program memory. "ROMless" versions must have this pin tied to '0'.

Table 2 Pin Definitions and Functions (cont'd)

Symbol	Pin Nr TQFP	Pin Nr MQFP	Input Outp.	Function																		
NC	40	42	–	This pin is not connected in the C165. No connection to the PCB is required.																		
PORT0 P0L.0-7 P0H.0-7	41-48 51-58	43-50 53-60	IO	<p>PORT0 consists of the two 8-bit bidirectional I/O ports P0L and P0H. It is bit-wise programmable for input or output via direction bits. For a pin configured as input, the output driver is put into high-impedance state. In case of an external bus configuration, PORT0 serves as the address (A) and address/data (AD) bus in multiplexed bus modes and as the data (D) bus in demultiplexed bus modes.</p> <p>Demultiplexed bus modes:</p> <table> <tr> <td>Data Path Width:</td> <td>8-bit</td> <td>16-bit</td> </tr> <tr> <td>P0L.0 – P0L.7:</td> <td>D0 – D7</td> <td>D0 – D7</td> </tr> <tr> <td>P0H.0 – P0H.7:</td> <td>I/O</td> <td>D8 – D15</td> </tr> </table> <p>Multiplexed bus modes:</p> <table> <tr> <td>Data Path Width:</td> <td>8-bit</td> <td>16-bit</td> </tr> <tr> <td>P0L.0 – P0L.7:</td> <td>AD0 – AD7</td> <td>AD0 – AD7</td> </tr> <tr> <td>P0H.0 – P0H.7:</td> <td>A8 – A15</td> <td>AD8 – AD15</td> </tr> </table>	Data Path Width:	8-bit	16-bit	P0L.0 – P0L.7:	D0 – D7	D0 – D7	P0H.0 – P0H.7:	I/O	D8 – D15	Data Path Width:	8-bit	16-bit	P0L.0 – P0L.7:	AD0 – AD7	AD0 – AD7	P0H.0 – P0H.7:	A8 – A15	AD8 – AD15
Data Path Width:	8-bit	16-bit																				
P0L.0 – P0L.7:	D0 – D7	D0 – D7																				
P0H.0 – P0H.7:	I/O	D8 – D15																				
Data Path Width:	8-bit	16-bit																				
P0L.0 – P0L.7:	AD0 – AD7	AD0 – AD7																				
P0H.0 – P0H.7:	A8 – A15	AD8 – AD15																				
PORT1 P1L.0-7 P1H.0-7	59-66 67,68, 71-76	61-68 69-70, 73-78	IO	<p>PORT1 consists of the two 8-bit bidirectional I/O ports P1L and P1H. It is bit-wise programmable for input or output via direction bits. For a pin configured as input, the output driver is put into high-impedance state. PORT1 is used as the 16-bit address bus (A) in demultiplexed bus modes and also after switching from a demultiplexed bus mode to a multiplexed bus mode.</p>																		

Table 2 Pin Definitions and Functions (cont'd)

Symbol	Pin Nr TQFP	Pin Nr MQFP	Input Outp.	Function
$\overline{\text{RSTIN}}$	79	81	I/O	<p>Reset Input with Schmitt-Trigger characteristics. A low level at this pin while the oscillator is running resets the C165. An internal pullup resistor permits power-on reset using only a capacitor connected to V_{SS}. A spike filter suppresses input pulses < 10 ns. Input pulses >100 ns safely pass the filter. The minimum duration for a safe recognition should be 100 ns + 2 CPU clock cycles.</p> <p>In bidirectional reset mode (enabled by setting bit BDRSTEN in register SYSCON) the $\overline{\text{RSTIN}}$ line is internally pulled low for the duration of the internal reset sequence upon any reset (HW, SW, WDT). See note below this table.</p> <p><i>Note: To let the reset configuration of PORT0 settle a reset duration of ca. 1 ms is recommended.</i></p>
$\overline{\text{RSTOUT}}$	80	82	O	<p>Internal Reset Indication Output. This pin is set to a low level when the part is executing either a hardware-, a software- or a watchdog timer reset. $\overline{\text{RSTOUT}}$ remains low until the EINIT (end of initialization) instruction is executed.</p>
$\overline{\text{NMI}}$	81	83	I	<p>Non-Maskable Interrupt Input. A high to low transition at this pin causes the CPU to vector to the NMI trap routine. When the PWRDN (power down) instruction is executed, the $\overline{\text{NMI}}$ pin must be low in order to force the C165 to go into power down mode. If $\overline{\text{NMI}}$ is high, when PWRDN is executed, the part will continue to run in normal mode.</p> <p>If not used, pin $\overline{\text{NMI}}$ should be pulled high externally.</p>

The CPU has a register context consisting of up to 16 wordwide GPRs at its disposal. These 16 GPRs are physically allocated within the on-chip RAM area. A Context Pointer (CP) register determines the base address of the active register bank to be accessed by the CPU at any time. The number of register banks is only restricted by the available internal RAM space. For easy parameter passing, a register bank may overlap others.

A system stack of up to 1024 words is provided as a storage for temporary data. The system stack is allocated in the on-chip RAM area, and it is accessed by the CPU via the stack pointer (SP) register. Two separate SFRs, STKOV and STKUN, are implicitly compared against the stack pointer value upon each stack access for the detection of a stack overflow or underflow.

The high performance offered by the hardware implementation of the CPU can efficiently be utilized by a programmer via the highly efficient C165 instruction set which includes the following instruction classes:

- Arithmetic Instructions
- Logical Instructions
- Boolean Bit Manipulation Instructions
- Compare and Loop Control Instructions
- Shift and Rotate Instructions
- Prioritize Instruction
- Data Movement Instructions
- System Stack Instructions
- Jump and Call Instructions
- Return Instructions
- System Control Instructions
- Miscellaneous Instructions

The basic instruction length is either 2 or 4 bytes. Possible operand types are bits, bytes and words. A variety of direct, indirect or immediate addressing modes are provided to specify the required operands.

Interrupt System

With an interrupt response time within a range from just 5 to 12 CPU clocks (in case of internal program execution), the C165 is capable of reacting very fast to the occurrence of non-deterministic events.

The architecture of the C165 supports several mechanisms for fast and flexible response to service requests that can be generated from various sources internal or external to the microcontroller. Any of these interrupt requests can be programmed to being serviced by the Interrupt Controller or by the Peripheral Event Controller (PEC).

In contrast to a standard interrupt service where the current program execution is suspended and a branch to the interrupt vector table is performed, just one cycle is 'stolen' from the current CPU activity to perform a PEC service. A PEC service implies a single byte or word data transfer between any two memory locations with an additional increment of either the PEC source or the destination pointer. An individual PEC transfer counter is implicitly decremented for each PEC service except when performing in the continuous transfer mode. When this counter reaches zero, a standard interrupt is performed to the corresponding source related vector location. PEC services are very well suited, for example, for supporting the transmission or reception of blocks of data. The C165 has 8 PEC channels each of which offers such fast interrupt-driven data transfer capabilities.

A separate control register which contains an interrupt request flag, an interrupt enable flag and an interrupt priority bitfield exists for each of the possible interrupt sources. Via its related register, each source can be programmed to one of sixteen interrupt priority levels. Once having been accepted by the CPU, an interrupt service can only be interrupted by a higher prioritized service request. For the standard interrupt processing, each of the possible interrupt sources has a dedicated vector location.

Fast external interrupt inputs are provided to service external interrupts with high precision requirements. These fast interrupt inputs feature programmable edge detection (rising edge, falling edge or both edges).

Software interrupts are supported by means of the 'TRAP' instruction in combination with an individual trap (interrupt) number.

Table 3 shows all of the possible C165 interrupt sources and the corresponding hardware-related interrupt flags, vectors, vector locations and trap (interrupt) numbers.

Note: Interrupt nodes which are not used by associated peripherals, may be used to generate software controlled interrupt requests by setting the respective interrupt request bit (xIR).

The C165 also provides an excellent mechanism to identify and to process exceptions or error conditions that arise during run-time, so-called ‘Hardware Traps’. Hardware traps cause immediate non-maskable system reaction which is similar to a standard interrupt service (branching to a dedicated vector table location). The occurrence of a hardware trap is additionally signified by an individual bit in the trap flag register (TFR). Except when another higher prioritized trap service is in progress, a hardware trap will interrupt any actual program execution. In turn, hardware trap services can normally not be interrupted by standard or PEC interrupts.

Table 4 shows all of the possible exceptions or error conditions that can arise during run-time:

Table 4 Hardware Trap Summary

Exception Condition	Trap Flag	Trap Vector	Vector Location	Trap Number	Trap Priority
Reset Functions: – Hardware Reset – Software Reset – W-dog Timer Overflow	–	RESET RESET RESET	00'0000 _H 00'0000 _H 00'0000 _H	00 _H 00 _H 00 _H	III III III
Class A Hardware Traps: – Non-Maskable Interrupt – Stack Overflow – Stack Underflow	NMI STKOF STKUF	NMITRAP STOTRAP STUTRAP	00'0008 _H 00'0010 _H 00'0018 _H	02 _H 04 _H 06 _H	II II II
Class B Hardware Traps: – Undefined Opcode – Protected Instruction Fault – Illegal Word Operand Access – Illegal Instruction Access – Illegal External Bus Access	UNDOPC PRTFLT ILLOPA ILLINA ILLBUS	BTRAP BTRAP BTRAP BTRAP BTRAP	00'0028 _H 00'0028 _H 00'0028 _H 00'0028 _H 00'0028 _H	0A _H 0A _H 0A _H 0A _H 0A _H	I I I I I
Reserved	–	–	[2C _H – 3C _H]	[0B _H – 0F _H]	–
Software Traps – TRAP Instruction	–	–	Any [00'0000 _H – 00'01FC _H] in steps of 4 _H	Any [00 _H – 7F _H]	Current CPU Priority

Instruction Set Summary

Table 5 lists the instructions of the C165 in a condensed way.

The various addressing modes that can be used with a specific instruction, the operation of the instructions, parameters for conditional execution of instructions, and the opcodes for each instruction can be found in the “**C166 Family Instruction Set Manual**”.

This document also provides a detailed description of each instruction.

Table 5 Instruction Set Summary

Mnemonic	Description	Bytes
ADD(B)	Add word (byte) operands	2 / 4
ADDC(B)	Add word (byte) operands with Carry	2 / 4
SUB(B)	Subtract word (byte) operands	2 / 4
SUBC(B)	Subtract word (byte) operands with Carry	2 / 4
MUL(U)	(Un)Signed multiply direct GPR by direct GPR (16-16-bit)	2
DIV(U)	(Un)Signed divide register MDL by direct GPR (16-/16-bit)	2
DIVL(U)	(Un)Signed long divide reg. MD by direct GPR (32-/16-bit)	2
CPL(B)	Complement direct word (byte) GPR	2
NEG(B)	Negate direct word (byte) GPR	2
AND(B)	Bitwise AND, (word/byte operands)	2 / 4
OR(B)	Bitwise OR, (word/byte operands)	2 / 4
XOR(B)	Bitwise XOR, (word/byte operands)	2 / 4
BCLR	Clear direct bit	2
BSET	Set direct bit	2
BMOV(N)	Move (negated) direct bit to direct bit	4
BAND, BOR, BXOR	AND/OR/XOR direct bit with direct bit	4
BCMP	Compare direct bit to direct bit	4
BFLDH/L	Bitwise modify masked high/low byte of bit-addressable direct word memory with immediate data	4
CMP(B)	Compare word (byte) operands	2 / 4
CMPD1/2	Compare word data to GPR and decrement GPR by 1/2	2 / 4
CMPI1/2	Compare word data to GPR and increment GPR by 1/2	2 / 4
PRIOR	Determine number of shift cycles to normalize direct word GPR and store result in direct word GPR	2
SHL / SHR	Shift left/right direct word GPR	2
ROL / ROR	Rotate left/right direct word GPR	2
ASHR	Arithmetic (sign bit) shift right direct word GPR	2

Special Function Registers Overview

The following table lists all SFRs which are implemented in the C165 in alphabetical order.

Bit-addressable SFRs are marked with the letter “b” in column “Name”. SFRs within the **Extended SFR-Space** (ESFRs) are marked with the letter “E” in column “Physical Address”. Registers within on-chip X-peripherals are marked with the letter “X” in column “Physical Address”.

An SFR can be specified via its individual mnemonic name. Depending on the selected addressing mode, an SFR can be accessed via its physical address (using the Data Page Pointers), or via its short 8-bit address (without using the Data Page Pointers).

Table 6 C165 Registers, Ordered by Name

Name	Physical Address	8-Bit Addr.	Description	Reset Value
ADDRSEL1	FE18 _H	0C _H	Address Select Register 1	0000 _H
ADDRSEL2	FE1A _H	0D _H	Address Select Register 2	0000 _H
ADDRSEL3	FE1C _H	0E _H	Address Select Register 3	0000 _H
ADDRSEL4	FE1E _H	0F _H	Address Select Register 4	0000 _H
BUSCON0 b	FF0C _H	86 _H	Bus Configuration Register 0	0XX0 _H
BUSCON1 b	FF14 _H	8A _H	Bus Configuration Register 1	0000 _H
BUSCON2 b	FF16 _H	8B _H	Bus Configuration Register 2	0000 _H
BUSCON3 b	FF18 _H	8C _H	Bus Configuration Register 3	0000 _H
BUSCON4 b	FF1A _H	8D _H	Bus Configuration Register 4	0000 _H
CAPREL	FE4A _H	25 _H	GPT2 Capture/Reload Register	0000 _H
CC10IC b	FF8C _H	C6 _H	EX2IN Interrupt Control Register	0000 _H
CC11IC b	FF8E _H	C7 _H	EX3IN Interrupt Control Register	0000 _H
CC12IC b	FF90 _H	C8 _H	EX4IN Interrupt Control Register	0000 _H
CC13IC b	FF92 _H	C9 _H	EX5IN Interrupt Control Register	0000 _H
CC14IC b	FF94 _H	CA _H	EX6IN Interrupt Control Register	0000 _H
CC15IC b	FF96 _H	CB _H	EX7IN Interrupt Control Register	0000 _H
CC29IC b	F184 _H E	C2 _H	Software Interrupt Control Register	0000 _H
CC30IC b	F18C _H E	C6 _H	Software Interrupt Control Register	0000 _H
CC31IC b	F194 _H E	CA _H	Software Interrupt Control Register	0000 _H
CC8IC b	FF88 _H	C4 _H	EX0IN Interrupt Control Register	0000 _H
CC9IC b	FF8A _H	C5 _H	EX1IN Interrupt Control Register	0000 _H

Table 6 C165 Registers, Ordered by Name (cont'd)

Name		Physical Address	8-Bit Addr.	Description	Reset Value
P1H	b	FF06 _H	83 _H	Port 1 High Reg. (Upper half of PORT1)	00 _H
P1L	b	FF04 _H	82 _H	Port 1 Low Reg.(Lower half of PORT1)	00 _H
P2	b	FFC0 _H	E0 _H	Port 2 Register	0000 _H
P3	b	FFC4 _H	E2 _H	Port 3 Register	0000 _H
P4	b	FFC8 _H	E4 _H	Port 4 Register (8 bits)	00 _H
P5	b	FFA2 _H	D1 _H	Port 5 Register (read only)	XXXX _H
P6	b	FFCC _H	E6 _H	Port 6 Register (8 bits)	00 _H
PECC0		FEC0 _H	60 _H	PEC Channel 0 Control Register	0000 _H
PECC1		FEC2 _H	61 _H	PEC Channel 1 Control Register	0000 _H
PECC2		FEC4 _H	62 _H	PEC Channel 2 Control Register	0000 _H
PECC3		FEC6 _H	63 _H	PEC Channel 3 Control Register	0000 _H
PECC4		FEC8 _H	64 _H	PEC Channel 4 Control Register	0000 _H
PECC5		FECA _H	65 _H	PEC Channel 5 Control Register	0000 _H
PECC6		FECC _H	66 _H	PEC Channel 6 Control Register	0000 _H
PECC7		FECE _H	67 _H	PEC Channel 7 Control Register	0000 _H
PSW	b	FF10 _H	88 _H	CPU Program Status Word	0000 _H
RP0H	b	F108 _H	E 84 _H	System Startup Config. Reg. (Rd. only)	XX _H
S0BG		FEB4 _H	5A _H	Serial Channel 0 Baud Rate Generator Reload Register	0000 _H
S0CON	b	FFB0 _H	D8 _H	Serial Channel 0 Control Register	0000 _H
S0EIC	b	FF70 _H	B8 _H	Serial Channel 0 Error Interrupt Ctrl. Reg	0000 _H
S0RBUF		FEB2 _H	59 _H	Serial Channel 0 Receive Buffer Reg. (read only)	XX _H
S0RIC	b	FF6E _H	B7 _H	Serial Channel 0 Receive Interrupt Control Register	0000 _H
S0TBIC	b	F19C _H	E CE _H	Serial Channel 0 Transmit Buffer Interrupt Control Register	0000 _H
S0TBUF		FEB0 _H	58 _H	Serial Channel 0 Transmit Buffer Register (write only)	00 _H
S0TIC	b	FF6C _H	B6 _H	Serial Channel 0 Transmit Interrupt Control Register	0000 _H

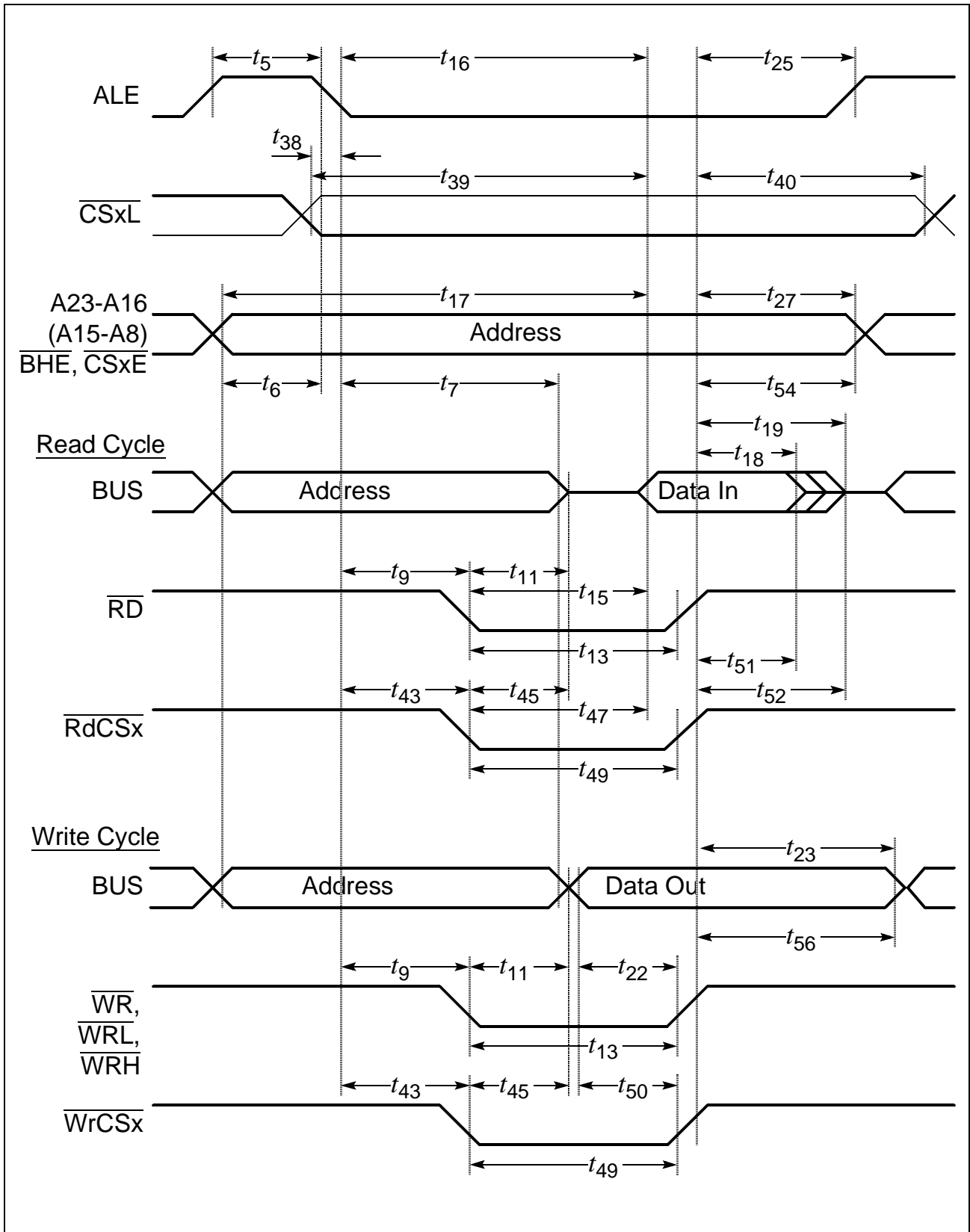


Figure 16 External Memory Cycle: Multiplexed Bus, No Read/Write Delay, Extended ALE

AC Characteristics

Demultiplexed Bus (Standard Supply Voltage Range)

(Operating Conditions apply)

ALE cycle time = $4 \text{ TCL} + 2t_A + t_C + t_F$ (80 ns at 25 MHz CPU clock without waitstates)

Parameter	Symbol	Max. CPU Clock = 25 MHz		Variable CPU Clock 1 / 2TCL = 1 to 25 MHz		Unit
		min.	max.	min.	max.	
ALE high time	t_5 CC	$10 + t_A$	–	$\text{TCL} - 10 + t_A$	–	ns
Address setup to ALE	t_6 CC	$4 + t_A$	–	$\text{TCL} - 16 + t_A$	–	ns
ALE falling edge to $\overline{\text{RD}}$, $\overline{\text{WR}}$ (with RW-delay)	t_8 CC	$10 + t_A$	–	$\text{TCL} - 10 + t_A$	–	ns
ALE falling edge to $\overline{\text{RD}}$, $\overline{\text{WR}}$ (no RW-delay)	t_9 CC	$-10 + t_A$	–	$-10 + t_A$	–	ns
$\overline{\text{RD}}$, $\overline{\text{WR}}$ low time (with RW-delay)	t_{12} CC	$30 + t_C$	–	$2\text{TCL} - 10 + t_C$	–	ns
$\overline{\text{RD}}$, $\overline{\text{WR}}$ low time (no RW-delay)	t_{13} CC	$50 + t_C$	–	$3\text{TCL} - 10 + t_C$	–	ns
$\overline{\text{RD}}$ to valid data in (with RW-delay)	t_{14} SR	–	$20 + t_C$	–	$2\text{TCL} - 20 + t_C$	ns
$\overline{\text{RD}}$ to valid data in (no RW-delay)	t_{15} SR	–	$40 + t_C$	–	$3\text{TCL} - 20 + t_C$	ns
ALE low to valid data in	t_{16} SR	–	$40 + t_A + t_C$	–	$3\text{TCL} - 20 + t_A + t_C$	ns
Address to valid data in	t_{17} SR	–	$50 + 2t_A + t_C$	–	$4\text{TCL} - 30 + 2t_A + t_C$	ns
Data hold after $\overline{\text{RD}}$ rising edge	t_{18} SR	0	–	0	–	ns
Data float after $\overline{\text{RD}}$ rising edge (with RW-delay ¹⁾)	t_{20} SR	–	$26 + 2t_A + t_F^{(1)}$	–	$2\text{TCL} - 14 + 22t_A + t_F^{(1)}$	ns
Data float after $\overline{\text{RD}}$ rising edge (no RW-delay ¹⁾)	t_{21} SR	–	$10 + 2t_A + t_F^{(1)}$	–	$\text{TCL} - 10 + 22t_A + t_F^{(1)}$	ns

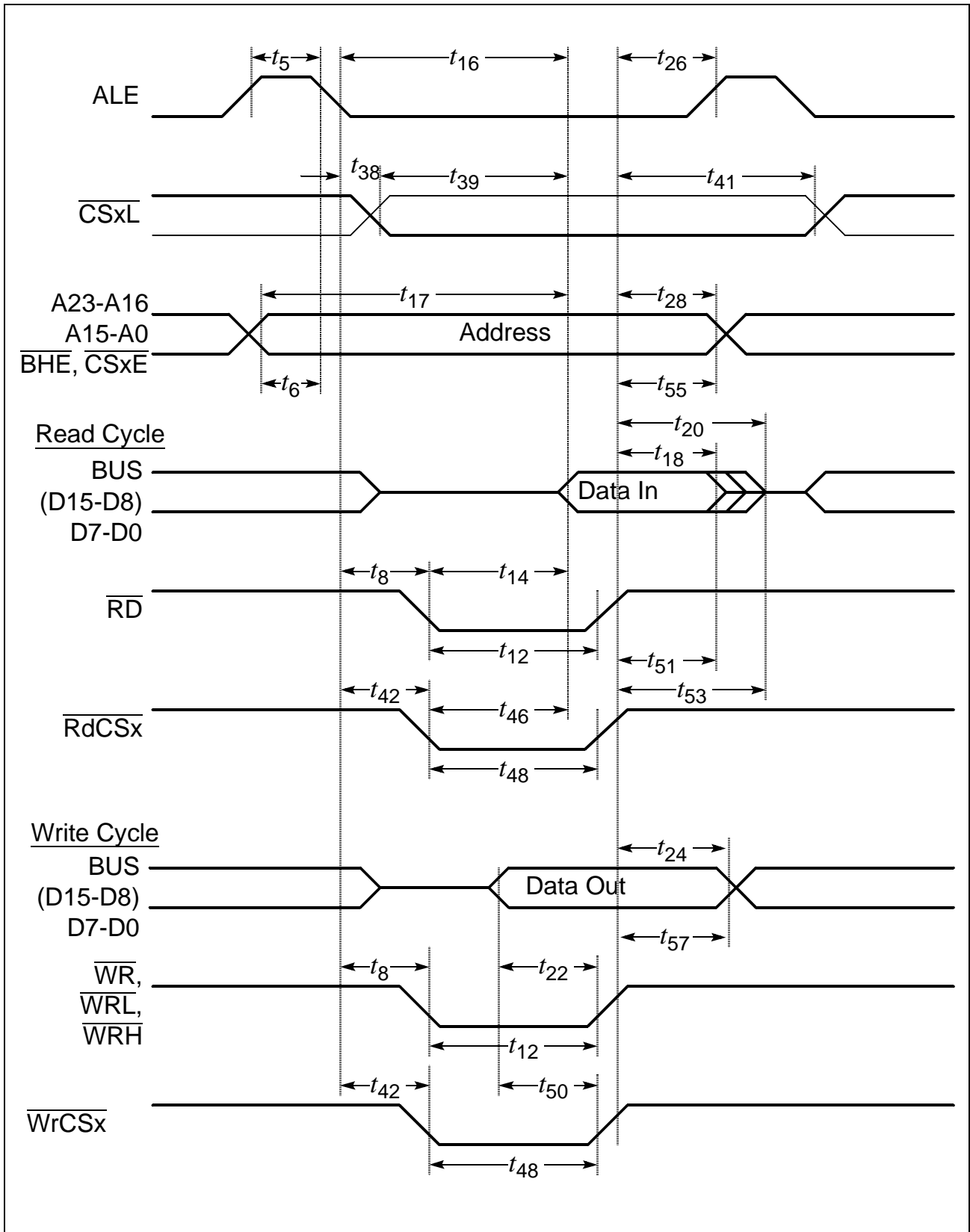


Figure 17 External Memory Cycle: Demultiplexed Bus, With Read/Write Delay, Normal ALE

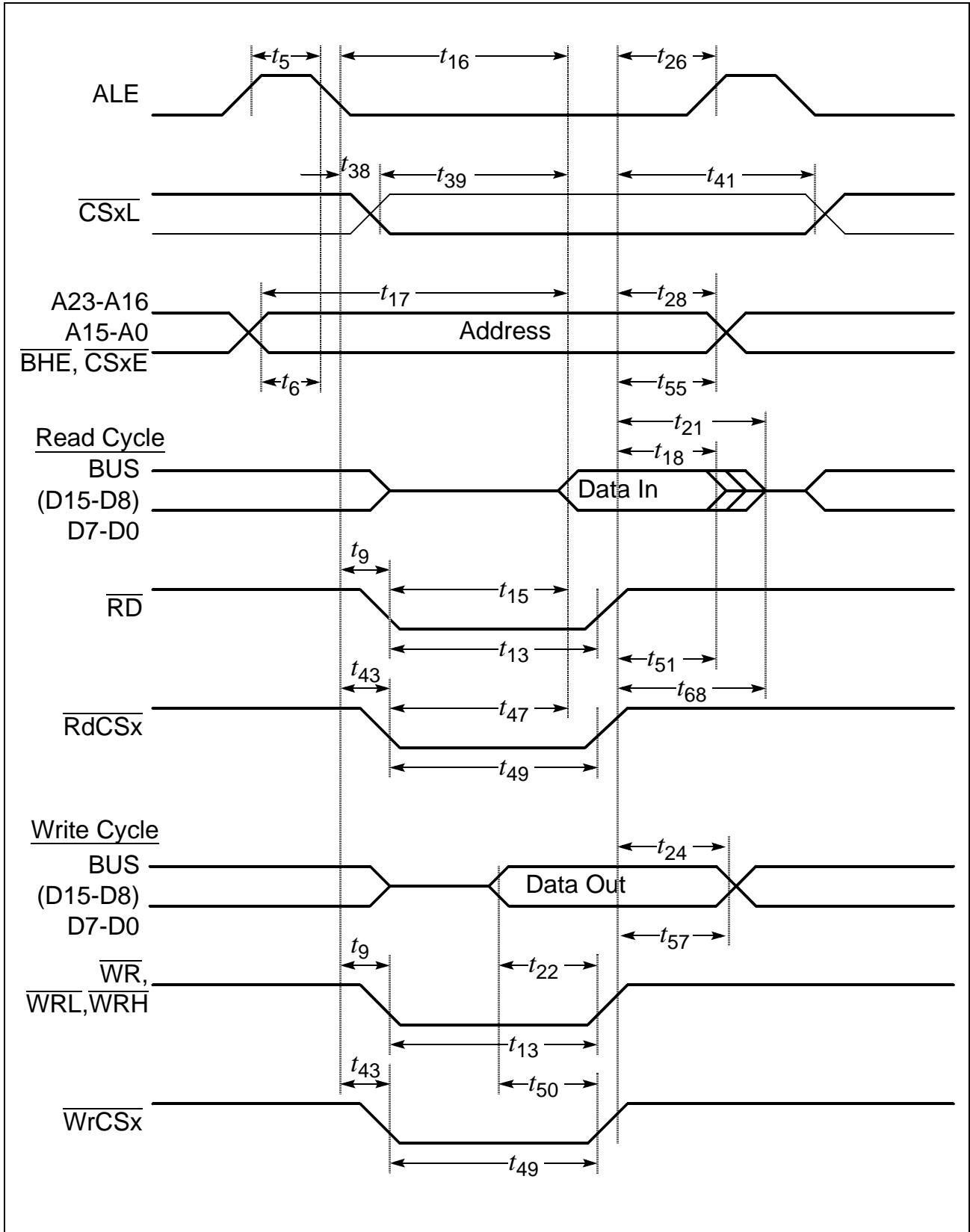


Figure 19 External Memory Cycle: Demultiplexed Bus, No Read/Write Delay, Normal ALE

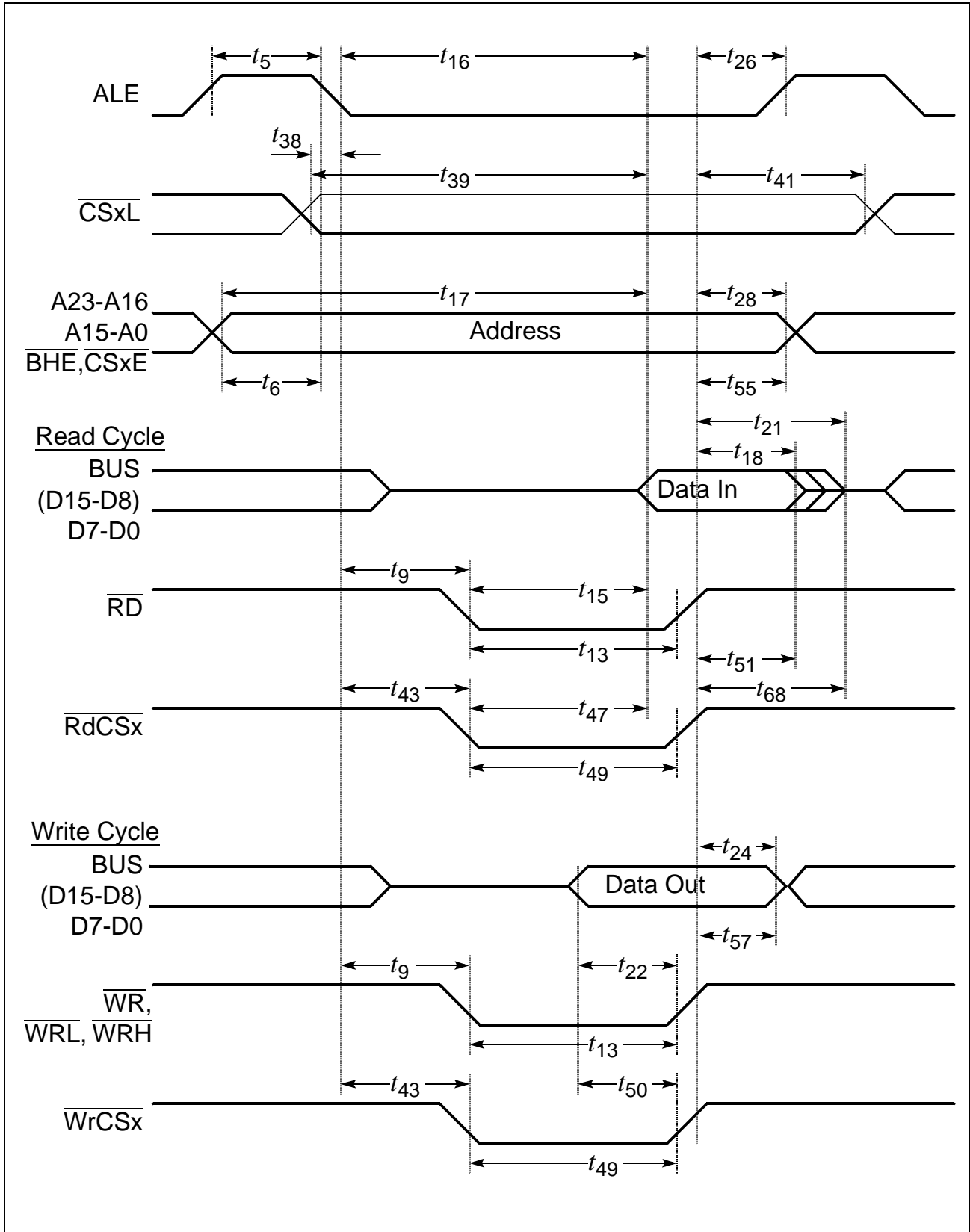


Figure 20 External Memory Cycle: Demultiplexed Bus, No Read/Write Delay, Extended ALE

AC Characteristics

CLKOUT and $\overline{\text{READY}}$ (Standard Supply Voltage)

(Operating Conditions apply)

Parameter	Symbol	Max. CPU Clock = 25 MHz		Variable CPU Clock 1 / 2TCL = 1 to 25 MHz		Unit
		min.	max.	min.	max.	
CLKOUT cycle time	t_{29} CC	40	40	2TCL	2TCL	ns
CLKOUT high time	t_{30} CC	14	–	TCL - 6	–	ns
CLKOUT low time	t_{31} CC	10	–	TCL - 10	–	ns
CLKOUT rise time	t_{32} CC	–	4	–	4	ns
CLKOUT fall time	t_{33} CC	–	4	–	4	ns
CLKOUT rising edge to ALE falling edge	t_{34} CC	$0 + t_A$	$10 + t_A$	$0 + t_A$	$10 + t_A$	ns
Synchronous $\overline{\text{READY}}$ setup time to CLKOUT	t_{35} SR	14	–	14	–	ns
Synchronous $\overline{\text{READY}}$ hold time after CLKOUT	t_{36} SR	4	–	4	–	ns
Asynchronous $\overline{\text{READY}}$ low time	t_{37} SR	54	–	$2\text{TCL} + t_{58}$	–	ns
Asynchronous $\overline{\text{READY}}$ setup time ¹⁾	t_{58} SR	14	–	14	–	ns
Asynchronous $\overline{\text{READY}}$ hold time ¹⁾	t_{59} SR	4	–	4	–	ns
Async. $\overline{\text{READY}}$ hold time after $\overline{\text{RD}}$, $\overline{\text{WR}}$ high (Demultiplexed Bus) ²⁾	t_{60} SR	0	$0 + 2t_A + t_C + t_F$ ²⁾	0	$\text{TCL} - 20 + 2t_A + t_C + t_F$ ²⁾	ns

1) These timings are given for test purposes only, in order to assure recognition at a specific clock edge.

2) Demultiplexed bus is the worst case. For multiplexed bus 2TCL are to be added to the maximum values. This adds even more time for deactivating $\overline{\text{READY}}$.

The $2t_A$ and t_C refer to the next following bus cycle, t_F refers to the current bus cycle.

The maximum limit for t_{60} must be fulfilled if the next following bus cycle is $\overline{\text{READY}}$ controlled.

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Dr. Ulrich Schumacher

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