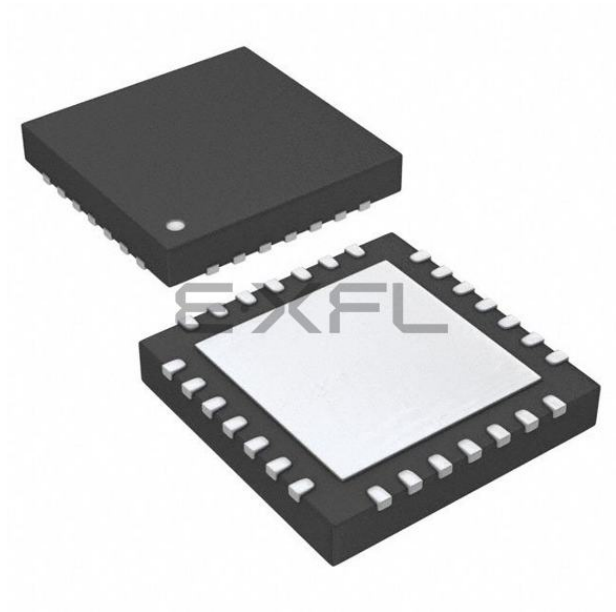

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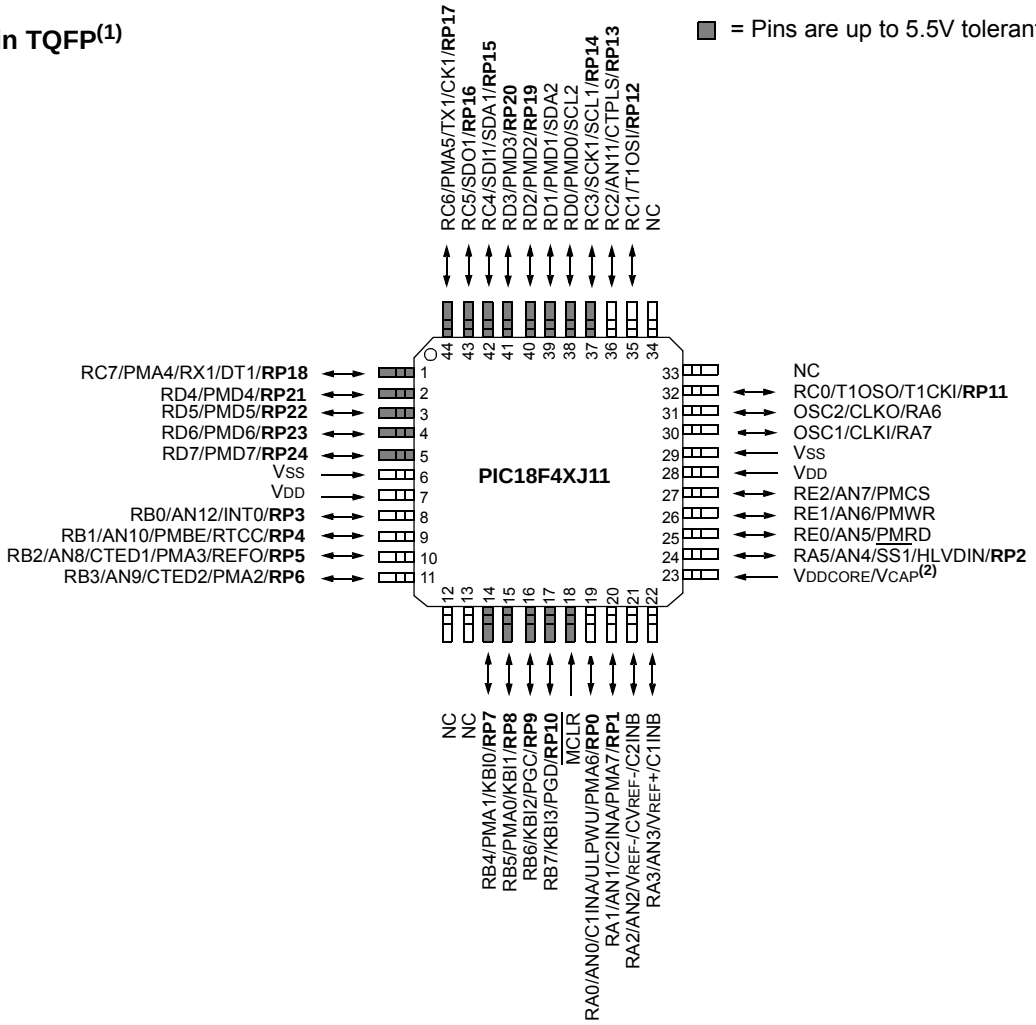
Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	48MHz
Connectivity	I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, POR, PWM, WDT
Number of I/O	16
Program Memory Size	16KB (8K x 16)
Program Memory Type	FLASH
EEPROM Size	-
RAM Size	3.8K x 8
Voltage - Supply (Vcc/Vdd)	2V ~ 3.6V
Data Converters	A/D 10x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	28-VQFN Exposed Pad
Supplier Device Package	28-QFN (6x6)
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18lf24j11-i-ml

Pin Diagrams (Continued)

44-Pin TQFP(1)

= Pins are up to 5.5V tolerant



Legend: **RPn** represents remappable pins.

Note 1: Some input and output functions are routed through the Peripheral Pin Select (PPS) module and can be dynamically assigned to any of the RPn pins. For a list of the input and output functions, see Table 10-13 and Table 10-14, respectively. For details on configuring the PPS module, see **Section 10.7 "Peripheral Pin Select (PPS)"**.

2: See **Section 26.3 "On-Chip Voltage Regulator"** for details on how to connect the VDDCORE/VCAP pin.

REGISTER 3-1: OSCTUNE: OSCILLATOR TUNING REGISTER (ACCESS F9Bh)

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
INTSRC	PLLEN	TUN5	TUN4	TUN3	TUN2	TUN1	TUN0
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 7 **INTSRC:** Internal Oscillator Low-Frequency Source Select bit
1 = 31.25 kHz device clock derived from 8 MHz INTOSC source (divide-by-256 enabled)
0 = 31 kHz device clock derived directly from INTRC internal oscillator

bit 6 **PLLEN:** Frequency Multiplier Enable bit
1 = PLL enabled
0 = PLL disabled

bit 5-0 **TUN<5:0>:** Frequency Tuning bits
011111 = Maximum frequency
011110
•
•
•
000001
000000 = Center frequency; oscillator module is running at the calibrated frequency
111111
•
•
•
100000 = Minimum frequency

3.3 Clock Sources and Oscillator Switching

Like previous PIC18 enhanced devices, the PIC18F46J11 family includes a feature that allows the device clock source to be switched from the main oscillator to an alternate, low-frequency clock source. PIC18F46J11 family devices offer two alternate clock sources. When an alternate clock source is enabled, the various power-managed operating modes are available.

Essentially, there are three clock sources for these devices:

- Primary Oscillators
- Secondary Oscillators
- Internal Oscillator Block

The **Primary Oscillators** include the External Crystal and Resonator modes, the External Clock modes and the internal oscillator block. The particular mode is defined by the FOSC<2:0> Configuration bits. The details of these modes are covered earlier in this chapter.

The **Secondary Oscillators** are external sources that are not connected to the OSC1 or OSC2 pins. These sources may continue to operate even after the controller is placed in a power-managed mode.

PIC18F46J11 family devices offer the Timer1 oscillator as a secondary oscillator. This oscillator, in all power-managed modes, is often the time base for functions such as a Real-Time Clock (RTC). Most often, a 32.768 kHz watch crystal is connected between the RC0/T1OSO/T1CKI/RP11 and RC1/T1OSI/RP12 pins. Like the HS Oscillator mode circuits, loading capacitors are also connected from each pin to ground. The Timer1 oscillator is discussed in more detail in **Section 13.5 "Timer1 Oscillator"**.

In addition to being a primary clock source, the **postscaled internal clock** is available as a power-managed mode clock source. The INTRC source is also used as the clock source for several special features, such as the WDT and Fail-Safe Clock Monitor (FSCM).

6.1.4.4 Stack Full and Underflow Resets

Device Resets on stack overflow and stack underflow conditions are enabled by setting the STVREN bit in Configuration register 1L. When STVREN is set, a full or underflow condition sets the appropriate STKFUL or STKUNF bit and then causes a device Reset. When STVREN is cleared, a full or underflow condition sets the appropriate STKFUL or STKUNF bit, but does not cause a device Reset. The STKFUL or STKUNF bits are cleared by the user software or a POR.

6.1.5 FAST REGISTER STACK (FRS)

A Fast Register Stack (FRS) is provided for the STATUS, WREG and BSR registers to provide a "fast return" option for interrupts. This stack is only one level deep and is neither readable nor writable. It is loaded with the current value of the corresponding register when the processor vectors for an interrupt. All interrupt sources push values into the Stack registers. The values in the registers are then loaded back into the working registers if the RETFIE, FAST instruction is used to return from the interrupt.

If both low-priority and high-priority interrupts are enabled, the Stack registers cannot be used reliably to return from low-priority interrupts. If a high-priority interrupt occurs while servicing a low-priority interrupt, the Stack register values stored by the low-priority interrupt will be overwritten. In these cases, users must save the key registers in software during a low-priority interrupt.

If interrupt priority is not used, all interrupts may use the FRS for returns from interrupt. If no interrupts are used, the FRS can be used to restore the STATUS, WREG and BSR registers at the end of a subroutine call. To use the Fast Register Stack for a subroutine call, a CALL label, FAST instruction must be executed to save the STATUS, WREG and BSR registers to the Fast Register Stack. A RETURN, FAST instruction is then executed to restore these registers from the FRS.

Example 6-1 provides a source code example that uses the FRS during a subroutine call and return.

EXAMPLE 6-1: FAST REGISTER STACK CODE EXAMPLE

```
CALL SUB1, FAST    ;STATUS, WREG, BSR
                   ;SAVED IN FAST REGISTER
                   ;STACK
       •
       •

SUB1   •
       •
   RETURN FAST     ;RESTORE VALUES SAVED
                   ;IN FAST REGISTER STACK
```

6.1.6 LOOK-UP TABLES IN PROGRAM MEMORY

There may be programming situations that require the creation of data structures or look-up tables in program memory. For PIC18 devices, look-up tables can be implemented in two ways:

- Computed GOTO
- Table Reads

6.1.6.1 Computed GOTO

A computed GOTO is accomplished by adding an offset to the PC. An example is shown in Example 6-2.

A look-up table can be formed with an ADDWF PCL instruction and a group of RETLW nn instructions. The W register is loaded with an offset into the table before executing a call to that table. The first instruction of the called routine is the ADDWF PCL instruction. The next executed instruction will be one of the RETLW nn instructions that returns the value 'nn' to the calling function.

The offset value (in WREG) specifies the number of bytes that the PC should advance and should be multiples of 2 (LSb = 0).

In this method, only one byte may be stored in each instruction location; room on the return address stack is required.

EXAMPLE 6-2: COMPUTED GOTO USING AN OFFSET VALUE

```
        MOVF    OFFSET, W
        CALL    TABLE
ORG     nn00h
TABLE   ADDWF   PCL
        RETLW   nnh
        RETLW   nnh
        RETLW   nnh
        .
        .
        .
```

6.1.6.2 Table Reads

A better method of storing data in program memory allows two bytes to be stored in each instruction location.

Look-up table data may be stored two bytes per program word while programming. The Table Pointer (TBLPTR) specifies the byte address and the Table Latch (TABLAT) contains the data that is read from the program memory. Data is transferred from program memory one byte at a time.

Table read operation is discussed further in **Section 7.1 "Table Reads and Table Writes"**.

6.2 PIC18 Instruction Cycle

6.2.1 CLOCKING SCHEME

The microcontroller clock input, whether from an internal or external source, is internally divided by '4' to generate four non-overlapping quadrature clocks (Q1, Q2, Q3 and Q4). Internally, the PC is incremented on every Q1; the instruction is fetched from the program memory and latched into the Instruction Register (IR) during Q4. The instruction is decoded and executed during the following Q1 through Q4. Figure 6-4 illustrates the clocks and instruction execution flow.

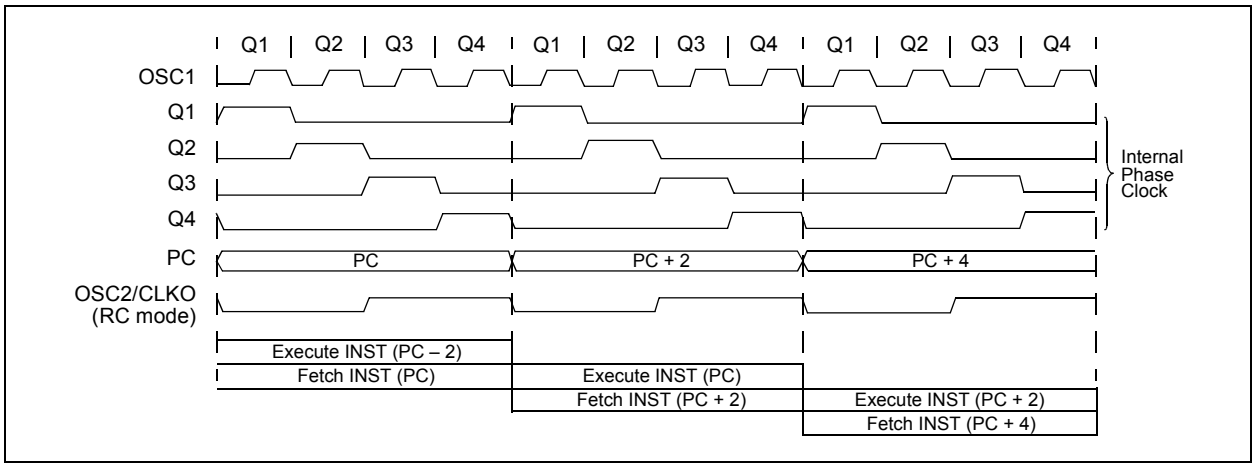
6.2.2 INSTRUCTION FLOW/PIPELINING

An "Instruction Cycle" consists of four Q cycles, Q1 through Q4. The instruction fetch and execute are pipelined in such a manner that a fetch takes one instruction cycle, while the decode and execute take another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the PC to change (e.g., GOTO), then two cycles are required to complete the instruction (Example 6-3).

A fetch cycle begins with the PC incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the IR in cycle Q1. This instruction is then decoded and executed during the Q2, Q3 and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).

FIGURE 6-4: CLOCK/INSTRUCTION CYCLE

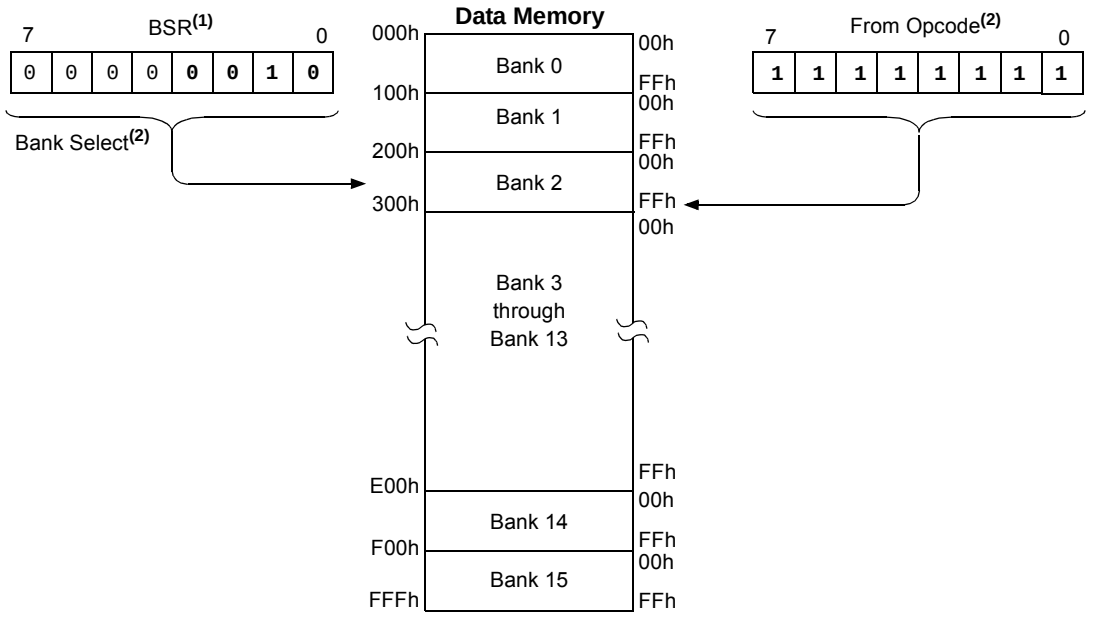


EXAMPLE 6-3: INSTRUCTION PIPELINE FLOW

	TCY0	TCY1	TCY2	TCY3	TCY4	TCY5
1. MOVLW 55h	Fetch 1	Execute 1				
2. MOVWF LATB		Fetch 2	Execute 2			
3. BRA SUB_1			Fetch 3	Execute 3		
4. BSF LATA, 3 (Forced NOP)				Fetch 4	Flush (NOP)	
5. Instruction @ address SUB_1					Fetch SUB_1	Execute SUB_1

Note: All instructions are single-cycle, except for any program branches. These take two cycles since the fetch instruction is "flushed" from the pipeline while the new instruction is being fetched and then executed.

FIGURE 6-7: USE OF THE BANK SELECT REGISTER (DIRECT ADDRESSING)



Note 1: The Access RAM bit of the instruction can be used to force an override of the selected bank (BSR<3:0>) to the registers of the Access Bank.

2: The MOVFF instruction embeds the entire 12-bit address in the instruction.

6.3.2 ACCESS BANK

While the use of the BSR with an embedded 8-bit address allows users to address the entire range of data memory, it also means that the user must always ensure that the correct bank is selected. Otherwise, data may be read from or written to the wrong location. This can be disastrous if a GPR is the intended target of an operation, but an SFR is written to instead. Verifying and/or changing the BSR for each read or write to data memory can become very inefficient.

To streamline access for the most commonly used data memory locations, the data memory is configured with an Access Bank, which allows users to access a mapped block of memory without specifying a BSR. The Access Bank consists of the first 96 bytes of memory (00h-5Fh) in Bank 0 and the last 160 bytes of memory (60h-FFh) in Bank 15. The lower half is known as the Access RAM and is composed of GPRs. The upper half is where the device's SFRs are mapped. These two areas are mapped contiguously in the Access Bank and can be addressed in a linear fashion by an 8-bit address (Figure 6-6).

The Access Bank is used by core PIC18 instructions that include the Access RAM bit (the 'a' parameter in the instruction). When 'a' is equal to '1', the instruction uses the BSR and the 8-bit address included in the opcode for the data memory address. When 'a' is '0', however, the instruction is forced to use the Access Bank address map; the current value of the BSR is ignored entirely.

Using this "forced" addressing allows the instruction to operate on a data address in a single cycle without updating the BSR first. For 8-bit addresses of 60h and above, this means that users can evaluate and operate on SFRs more efficiently. The Access RAM below 60h is a good place for data values that the user might need to access rapidly, such as immediate computational results or common program variables. Access RAM also allows for faster and more code efficient context saving and switching of variables.

The mapping of the Access Bank is slightly different when the extended instruction set is enabled (XINST Configuration bit = 1). This is discussed in more detail in **Section 6.6.3 "Mapping the Access Bank in Indexed Literal Offset Mode"**.

6.3.3 GENERAL PURPOSE REGISTER FILE

PIC18 devices may have banked memory in the GPR area. This is data RAM, which is available for use by all instructions. GPRs start at the bottom of Bank 0 (address 000h) and grow upward toward the bottom of the SFR area. GPRs are not initialized by a POR and are unchanged on all other Resets.

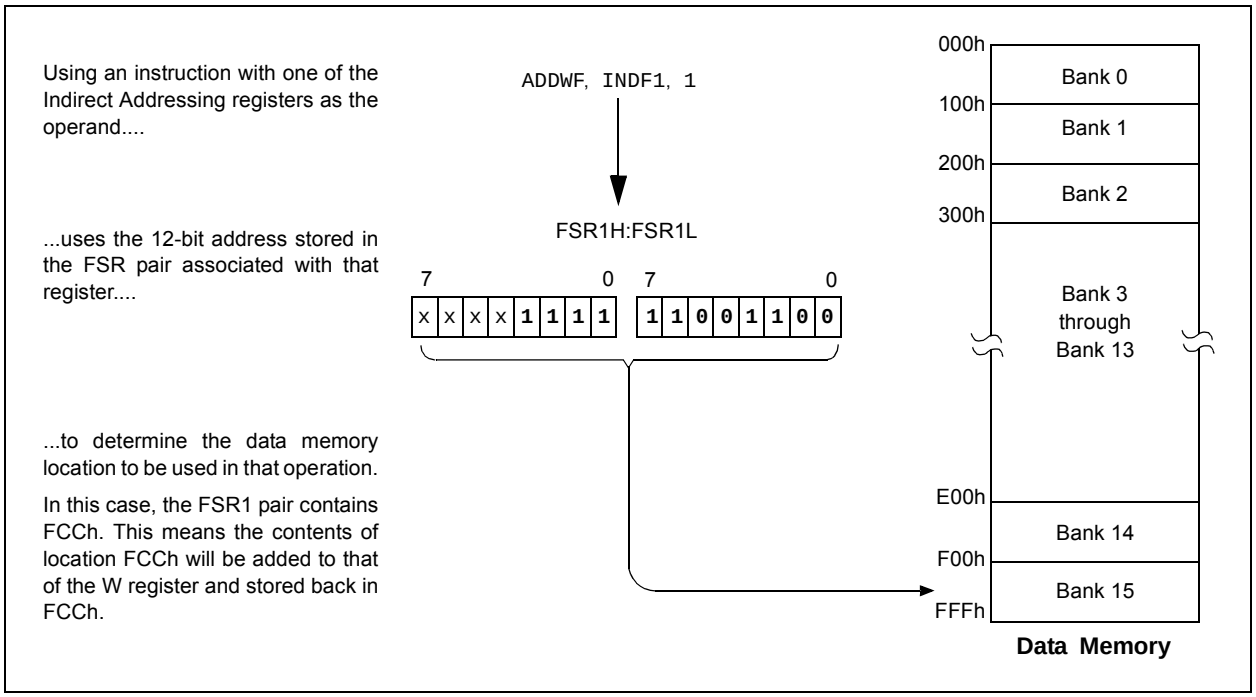
6.4.3.1 FSR Registers and the INDF Operand (INDF)

At the core of Indirect Addressing are three sets of registers: FSR0, FSR1 and FSR2. Each represents a pair of 8-bit registers, FSRnH and FSRnL. The four upper bits of the FSRnH register are not used, so each FSR pair holds a 12-bit value. This represents a value that can address the entire range of the data memory in a linear fashion. The FSR register pairs, then, serve as pointers to data memory locations.

Indirect Addressing is accomplished with a set of INDF operands, INDF0 through INDF2. These can be presumed to be "virtual" registers: they are mapped in the SFR space but are not physically implemented. Reading or writing to a particular INDF register actually accesses its corresponding FSR register pair. A read from INDF1, for example, reads the data at the address indicated by FSR1H:FSR1L. Instructions that use the INDF registers as operands actually use the contents of their corresponding FSR as a pointer to the instruction's target. The INDF operand is just a convenient way of using the pointer.

Because Indirect Addressing uses a full 12-bit address, data RAM banking is not necessary. Thus, the current contents of the BSR and the Access RAM bit have no effect on determining the target address.

FIGURE 6-8: INDIRECT ADDRESSING



10.2 PORTA, TRISA and LATA Registers

PORTA is a 7-bit wide, bidirectional port. It may function as a 5-bit port, depending on the oscillator mode selected. Setting a TRISA bit (= 1) will make the corresponding PORTA pin an input (i.e., put the corresponding output driver in a high-impedance mode). Clearing a TRISA bit (= 0) will make the corresponding PORTA pin an output (i.e., put the contents of the output latch on the selected pin).

Reading the PORTA register reads the status of the pins, whereas writing to it, will write to the port latch.

The Data Latch (LATA) register is also memory mapped. Read-modify-write operations on the LATA register read and write the latched output value for PORTA.

The other PORTA pins are multiplexed with analog inputs, the analog VREF+ and VREF- inputs and the comparator voltage reference output. The operation of pins, RA<3:0> and RA5, as A/D converter inputs is selected by clearing or setting the control bits in the ANCON0 register (A/D Port Configuration Register 0).

Pins, RA0 and RA3, may also be used as comparator inputs by setting the appropriate bits in the CMCON register. To use RA<3:0> as digital inputs, it is also necessary to turn off the comparators.

Note:	On a Power-on Reset (POR), RA5 and RA<3:0> are configured as analog inputs and read as '0'.

All PORTA pins have TTL input levels and full CMOS output drivers.

The TRISA register controls the direction of the PORTA pins, even when they are being used as analog inputs. The user must ensure the bits in the TRISA register are maintained set when using them as analog inputs.

EXAMPLE 10-2: INITIALIZING PORTA

```
CLRF   LATA     ; Initialize LATA
                ; to clear output
                ; data latches
MOVLB  0x0F     ; ANCONx register not in
                ; Access Bank
MOVLW  0x0F     ; Configure A/D
MOVWF  ANCON0   ; for digital inputs
MOVLW  0xCF     ; Value used to
                ; initialize data
                ; direction
MOVWF  TRISA    ; Set RA<3:0> as inputs
                ; RA<5:4> as outputs
```

TABLE 11-2: REGISTERS ASSOCIATED WITH PMP MODULE

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Reset Values on Page:
INTCON	GIE/GIEH	PEIE/GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF	69
PIR1	PMPIF(2)	ADIF	RC1IF	TX1IF	SSP1IF	CCP1IF	TMR2IF	TMR1IF	72
PIE1	PMPIE(2)	ADIE	RC1IE	TX1IE	SSP1IE	CCP1IE	TMR2IE	TMR1IE	72
IPR1	PMPIP(2)	ADIP	RC1IP	TX1IP	SSP1IP	CCP1IP	TMR2IP	TMR1IP	72
PMCONH(2)	PMPEN	—	—	ADRMUX1	ADRMUX0	PTBEEN	PTWREN	PTRDEN	73
PMCONL(2)	CSF1	CSF0	ALP	—	CS1P	BEP	WRSP	RDSP	73
PMADDRH(1,2)/	—	CS1	Parallel Master Port Address High Byte						73
PMDOUT1H(1,2)	Parallel Port Out Data High Byte (Buffer 1)								73
PMADDRL(1,2)/	Parallel Master Port Address Low Byte								73
PMDOUT1L(1,2)	Parallel Port Out Data Low Byte (Buffer 0)								73
PMDOUT2H(2)	Parallel Port Out Data High Byte (Buffer 3)								73
PMDOUT2L(2)	Parallel Port Out Data Low Byte (Buffer 2)								73
PMDIN1H(2)	Parallel Port In Data High Byte (Buffer 1)								73
PMDIN1L(2)	Parallel Port In Data Low Byte (Buffer 0)								73
PMDIN2H(2)	Parallel Port In Data High Byte (Buffer 3)								73
PMDIN2L(2)	Parallel Port In Data Low Byte (Buffer 2)								73
PMMODEH(2)	BUSY	IRQM1	IRQM0	INCM1	INCM0	MODE16	MODE1	MODE0	73
PMMODEL(2)	WAITB1	WAITB0	WAITM3	WAITM2	WAITM1	WAITM0	WAITE1	WAITE0	73
PMEH(2)	—	PTEN14	—	—	—	—	—	—	74
PMEL(2)	PTEN7	PTEN6	PTEN5	PTEN4	PTEN3	PTEN2	PTEN1	PTEN0	74
PMSTATH(2)	IBF	IBOV	—	—	IB3F	IB2F	IB1F	IB0F	74
PMSTATL(2)	OBE	OBUF	—	—	OB3E	OB2E	OB1E	OB0E	74
PADCFG1	—	—	—	—	—	RTSECSEL1	RTSECSEL0	PMPTTL	74

Legend: — = unimplemented, read as '0'. Shaded cells are not used during PMP operation.

Note 1: The PMADDRH/PMDOUT1H and PMADDRL/PMDOUT1L register pairs share the physical registers and addresses, but have different functions determined by the module's operating mode.

2: These bits and/or registers are only available in 44-pin devices.

13.2 Timer1 Operation

The Timer1 module is an 8-bit or 16-bit incrementing counter, which is accessed through the TMR1H:TMR1L register pair.

When used with an internal clock source, the module is a timer and increments on every instruction cycle. When used with an external clock source, the module can be used as either a timer or counter and increments on every selected edge of the external source.

Timer1 is enabled by configuring the TMR1ON and TMR1GE bits in the T1CON and T1GCON registers, respectively.

When Timer1 is enabled, the RC1/T1OSI/RP12 and RC0/T1OSO/T1CKI/RP11 pins become inputs. This means the values of TRISC<1:0> are ignored and the pins are read as '0'.

13.3 Clock Source Selection

The TMR1CS<1:0> and T1OSCEN bits of the T1CON register are used to select the clock source for Timer1. Register 13-1 displays the clock source selections.

When switching clock sources and using the clock prescaler, write to TMR1L afterwards to reset the internal prescaler count to 0.

13.3.1 INTERNAL CLOCK SOURCE

When the internal clock source is selected, the TMR1H:TMR1L register pair will increment on multiples of FOSC as determined by the Timer1 prescaler.

13.3.2 EXTERNAL CLOCK SOURCE

When the external clock source is selected, the Timer1 module may work as a timer or a counter.

When enabled to count, Timer1 is incremented on the rising edge of the external clock input, T1CKI, or the capacitive sensing oscillator signal. Either of these external clock sources can be synchronized to the microcontroller system clock or they can run asynchronously.

When used as a timer with a clock oscillator, an external 32.768 kHz crystal can be used in conjunction with the dedicated internal oscillator circuit.

Note: In Counter mode, a falling edge must be registered by the counter prior to the first incrementing rising edge after any one or more of the following conditions:

- Timer1 enabled after POR Reset
- Write to TMR1H or TMR1L
- Timer1 is disabled
- Timer1 is disabled (TMR1ON = 0)

when T1CKI is high, then Timer1 is enabled (TMR1ON = 1) when T1CKI is low.

TABLE 13-1: TIMER1 CLOCK SOURCE SELECTION

TMR1CS1	TMR1CS0	T1OSCEN	Clock Source
0	1	x	Clock Source (FOSC)
0	0	x	Instruction Clock (FOSC/4)
1	0	0	External Clock on T1CKI Pin
1	0	1	Oscillator Circuit on T1OSI/T1OSO Pin

In addition to the expanded range of modes available through the CCPxCON and ECCPxAS registers, the ECCP modules have two additional registers associated with Enhanced PWM operation and auto-shutdown features. They are:

- ECCPxDEL (Enhanced PWM Control)
- PSTRxCON (Pulse Steering Control)

18.1 ECCP Outputs and Configuration

The Enhanced CCP module may have up to four PWM outputs, depending on the selected operating mode. These outputs, designated PxA through PxD, are routed through the Peripheral Pin Select (PPS) module. Therefore, individual functions may be mapped to any of the remappable I/O pins, RPn. The outputs that are active depend on the ECCP operating mode selected. The pin assignments are summarized in Table 18-4.

To configure the I/O pins as PWM outputs, the proper PWM mode must be selected by setting the PxM<1:0> and CCPxM<3:0> bits. The appropriate TRIS direction bits for the port pins must also be set as outputs and the output functions need to be assigned to I/O pins in the PPS module. (For details on configuring the module, see **Section 10.7 "Peripheral Pin Select (PPS)"**.)

18.1.1 ECCP MODULE AND TIMER RESOURCES

The ECCP modules utilize Timers 1, 2, 3 or 4, depending on the mode selected. Timer1 and Timer3 are available to modules in Capture or Compare modes, while Timer2 and Timer4 are available for modules in PWM mode.

TABLE 18-1: ECCP MODE – TIMER RESOURCE

ECCP Mode	Timer Resource
Capture	Timer1 or Timer3
Compare	Timer1 or Timer3
PWM	Timer2 or Timer4

The assignment of a particular timer to a module is determined by the Timer-to-ECCP enable bits in the TCLKCON register (Register 13-3). The interactions between the two modules are depicted in Figure 18-1. Capture operations are designed to be used when the timer is configured for Synchronous Counter mode. Capture operations may not work as expected if the associated timer is configured for Asynchronous Counter mode.

18.2 Capture Mode

In Capture mode, the CCPRxH:CCPRxL register pair captures the 16-bit value of the TMR1 or TMR3 registers when an event occurs on the corresponding ECCPx pin. An event is defined as one of the following:

- Every falling edge
- Every rising edge
- Every 4th rising edge
- Every 16th rising edge

The event is selected by the mode select bits, CCPxM<3:0>, of the CCPxCON register. When a capture is made, the interrupt request flag bit, CCPxIF, is set; it must be cleared by software. If another capture occurs before the value in register CCPRx is read, the old captured value is overwritten by the new captured value.

18.2.1 ECCP PIN CONFIGURATION

In Capture mode, the appropriate ECCPx pin should be configured as an input by setting the corresponding TRIS direction bit.

Additionally, the ECCPx input function needs to be assigned to an I/O pin through the Peripheral Pin Select module. For details on setting up the remappable pins, see **Section 10.7 "Peripheral Pin Select (PPS)"**.

Note:	If the ECCPx pin is configured as an output, a write to the port can cause a capture condition.

18.2.2 TIMER1/TIMER3 MODE SELECTION

The timers that are to be used with the capture feature (Timer1 and/or Timer3) must be running in Timer mode or Synchronized Counter mode. In Asynchronous Counter mode, the capture operation may not work. The timer to be used with each ECCP module is selected in the TCLKCON register (Register 13-3).

18.2.3 SOFTWARE INTERRUPT

When the Capture mode is changed, a false capture interrupt may be generated. The user should keep the CCPxIE interrupt enable bit clear to avoid false interrupts. The interrupt flag bit, CCPxIF, should also be cleared following any such change in operating mode.

REGISTER 20-1: TXSTAx: TRANSMIT STATUS AND CONTROL REGISTER (ACCESS FADh/FA8h)

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R-1	R/W-0
CSRC	TX9	TXEN(1)	SYNC	SENDB	BRGH	TRMT	TX9D
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as '0'	
-n = Value at POR	'1' = Bit is set	'0' = Bit is cleared	x = Bit is unknown

bit 7 **CSRC:** Clock Source Select bit
Asynchronous mode:
Don't care.
Synchronous mode:
1 = Master mode (clock generated internally from BRG)
0 = Slave mode (clock from external source)

bit 6 **TX9:** 9-Bit Transmit Enable bit
1 = Selects 9-bit transmission
0 = Selects 8-bit transmission

bit 5 **TXEN:** Transmit Enable bit(1)
1 = Transmit is enabled and the TXx/CKx pin is configured as an output
0 = Transmit is disabled

bit 4 **SYNC:** EUSART Mode Select bit
1 = Synchronous mode
0 = Asynchronous mode

bit 3 **SENDB:** Send Break Character bit
Asynchronous mode:
1 = Send Sync Break on next transmission (cleared by hardware upon completion)
0 = Sync Break transmission completed
Synchronous mode:
Don't care.

bit 2 **BRGH:** High Baud Rate Select bit
Asynchronous mode:
1 = High speed
0 = Low speed
Synchronous mode:
Unused in this mode.

bit 1 **TRMT:** Transmit Shift Register Status bit
1 = TSR empty
0 = TSR full

bit 0 **TX9D:** 9th bit of Transmit Data
Can be address/data bit or a parity bit.

Note 1: SREN/CREN overrides TXEN in Sync mode.

TABLE 20-3: BAUD RATES FOR ASYNCHRONOUS MODES

BAUD RATE (K)	SYNC = 0, BRGH = 0, BRG16 = 0											
	FOSC = 40.000 MHz			FOSC = 20.000 MHz			FOSC = 10.000 MHz			FOSC = 8.000 MHz		
	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)
0.3	—	—	—	—	—	—	—	—	—	—	—	—
1.2	—	—	—	1.221	1.73	255	1.202	0.16	129	1.201	-0.16	103
2.4	2.441	1.73	255	2.404	0.16	129	2.404	0.16	64	2.403	-0.16	51
9.6	9.615	0.16	64	9.766	1.73	31	9.766	1.73	15	9.615	-0.16	12
19.2	19.531	1.73	31	19.531	1.73	15	19.531	1.73	7	—	—	—
57.6	56.818	-1.36	10	62.500	8.51	4	52.083	-9.58	2	—	—	—
115.2	125.000	8.51	4	104.167	-9.58	2	78.125	-32.18	1	—	—	—

BAUD RATE (K)	SYNC = 0, BRGH = 0, BRG16 = 0								
	FOSC = 4.000 MHz			FOSC = 2.000 MHz			FOSC = 1.000 MHz		
	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)
0.3	0.300	0.16	207	0.300	-0.16	103	0.300	-0.16	51
1.2	1.202	0.16	51	1.201	-0.16	25	1.201	-0.16	12
2.4	2.404	0.16	25	2.403	-0.16	12	—	—	—
9.6	8.929	-6.99	6	—	—	—	—	—	—
19.2	20.833	8.51	2	—	—	—	—	—	—
57.6	62.500	8.51	0	—	—	—	—	—	—
115.2	62.500	-45.75	0	—	—	—	—	—	—

BAUD RATE (K)	SYNC = 0, BRGH = 1, BRG16 = 0											
	FOSC = 40.000 MHz			FOSC = 20.000 MHz			FOSC = 10.000 MHz			FOSC = 8.000 MHz		
	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)
0.3	—	—	—	—	—	—	—	—	—	—	—	—
1.2	—	—	—	—	—	—	—	—	—	—	—	—
2.4	—	—	—	—	—	—	2.441	1.73	255	2.403	-0.16	207
9.6	9.766	1.73	255	9.615	0.16	129	9.615	0.16	64	9615.	-0.16	51
19.2	19.231	0.16	129	19.231	0.16	64	19.531	1.73	31	19.230	-0.16	25
57.6	58.140	0.94	42	56.818	-1.36	21	56.818	-1.36	10	55.555	3.55	8
115.2	113.636	-1.36	21	113.636	-1.36	10	125.000	8.51	4	—	—	—

BAUD RATE (K)	SYNC = 0, BRGH = 1, BRG16 = 0								
	FOSC = 4.000 MHz			FOSC = 2.000 MHz			FOSC = 1.000 MHz		
	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)	Actual Rate (K)	% Error	SPBRG value (decimal)
0.3	—	—	—	—	—	—	0.300	-0.16	207
1.2	1.202	0.16	207	1.201	-0.16	103	1.201	-0.16	51
2.4	2.404	0.16	103	2.403	-0.16	51	2.403	-0.16	25
9.6	9.615	0.16	25	9.615	-0.16	12	—	—	—
19.2	19.231	0.16	12	—	—	—	—	—	—
57.6	62.500	8.51	3	—	—	—	—	—	—
115.2	125.000	8.51	1	—	—	—	—	—	—

10. If any error occurred, clear the CREN bit.
11. If the device has been addressed, clear the ADDEN bit to allow all received data into the receive buffer and interrupt the CPU.

FIGURE 20-6: EUSARTx RECEIVE BLOCK DIAGRAM

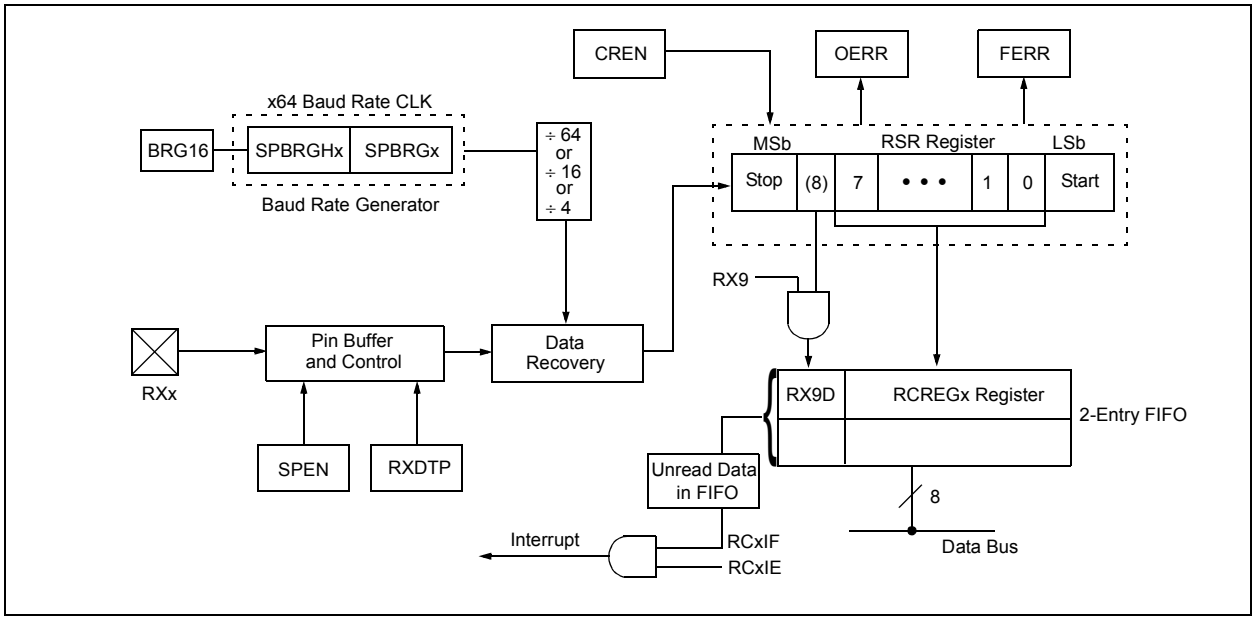
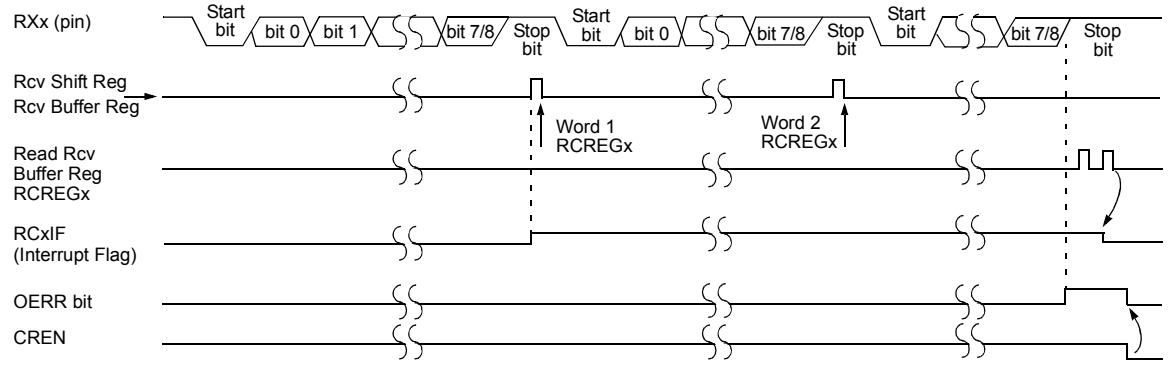


FIGURE 20-7: ASYNCHRONOUS RECEPTION



Note: This timing diagram shows three words appearing on the RXx input. The RCREGx (Receive Buffer) is read after the third word causing the OERR (Overrun) bit to be set.

22.0 COMPARATOR MODULE

The analog comparator module contains two comparators that can be independently configured in a variety of ways. The inputs can be selected from the analog inputs and two internal voltage references. The digital outputs are available at the pin level and can also be read through the control register. Multiple output and interrupt event generation is also available. Figure 22-1 provides a generic single comparator from the module.

Key features of the module are:

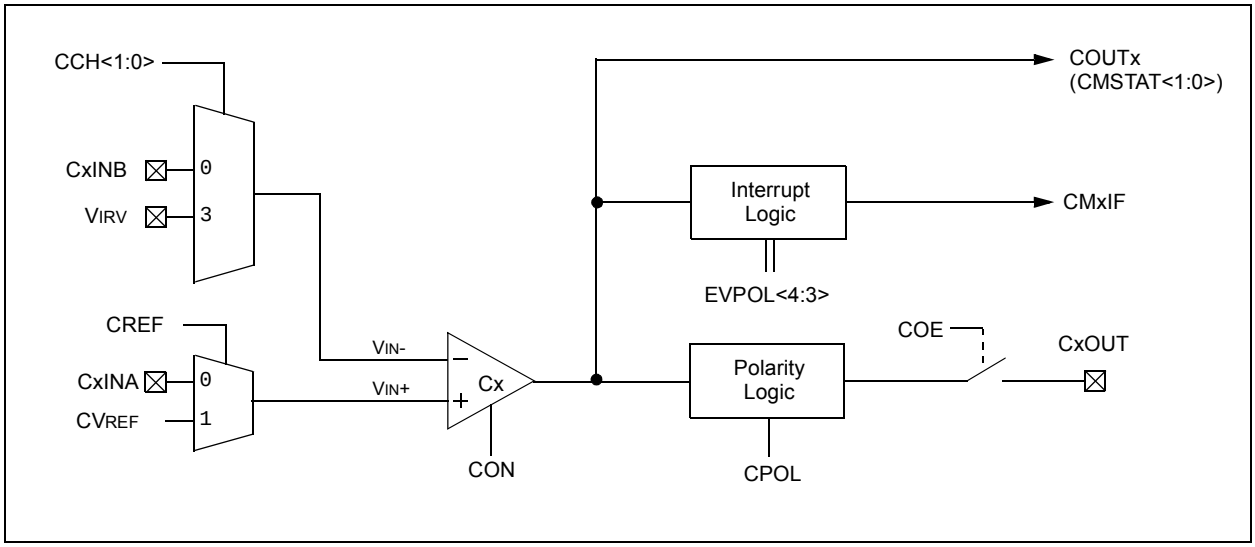
- Independent comparator control
- Programmable input configuration
- Output to both pin and register levels
- Programmable output polarity
- Independent interrupt generation for each comparator with configurable interrupt-on-change

22.1 Registers

The CMxCON registers (Register 22-1) select the input and output configuration for each comparator, as well as the settings for interrupt generation.

The CMSTAT register (Register 22-2) provides the output results of the comparators. The bits in this register are read-only.

FIGURE 22-1: COMPARATOR SIMPLIFIED BLOCK DIAGRAM



24.1 Operation

When the HLVD module is enabled, a comparator uses an internally generated reference voltage as the set point. The set point is compared with the trip point, where each node in the resistor divider represents a trip point voltage. The “trip point” voltage is the voltage level at which the device detects a high or low-voltage event, depending on the configuration of the module.

When the supply voltage is equal to the trip point, the voltage tapped off of the resistor array is equal to the internal reference voltage generated by the voltage reference module. The comparator then generates an interrupt signal by setting the LVDIF bit.

The trip point voltage is software programmable to any one of 8 values. The trip point is selected by programming the HLVDL<3:0> bits (HLVDCON<3:0>).

Additionally, the HLVD module allows the user to supply the trip voltage to the module from an external source. This mode is enabled when bits, HLVDL<3:0>, are set to ‘1111’. In this state, the comparator input is multiplexed from the external input pin, HLVDIN. This gives users flexibility because it allows them to configure the HLVD interrupt to occur at any voltage in the valid operating range.

FIGURE 24-1: HLVD MODULE BLOCK DIAGRAM (WITH EXTERNAL INPUT)

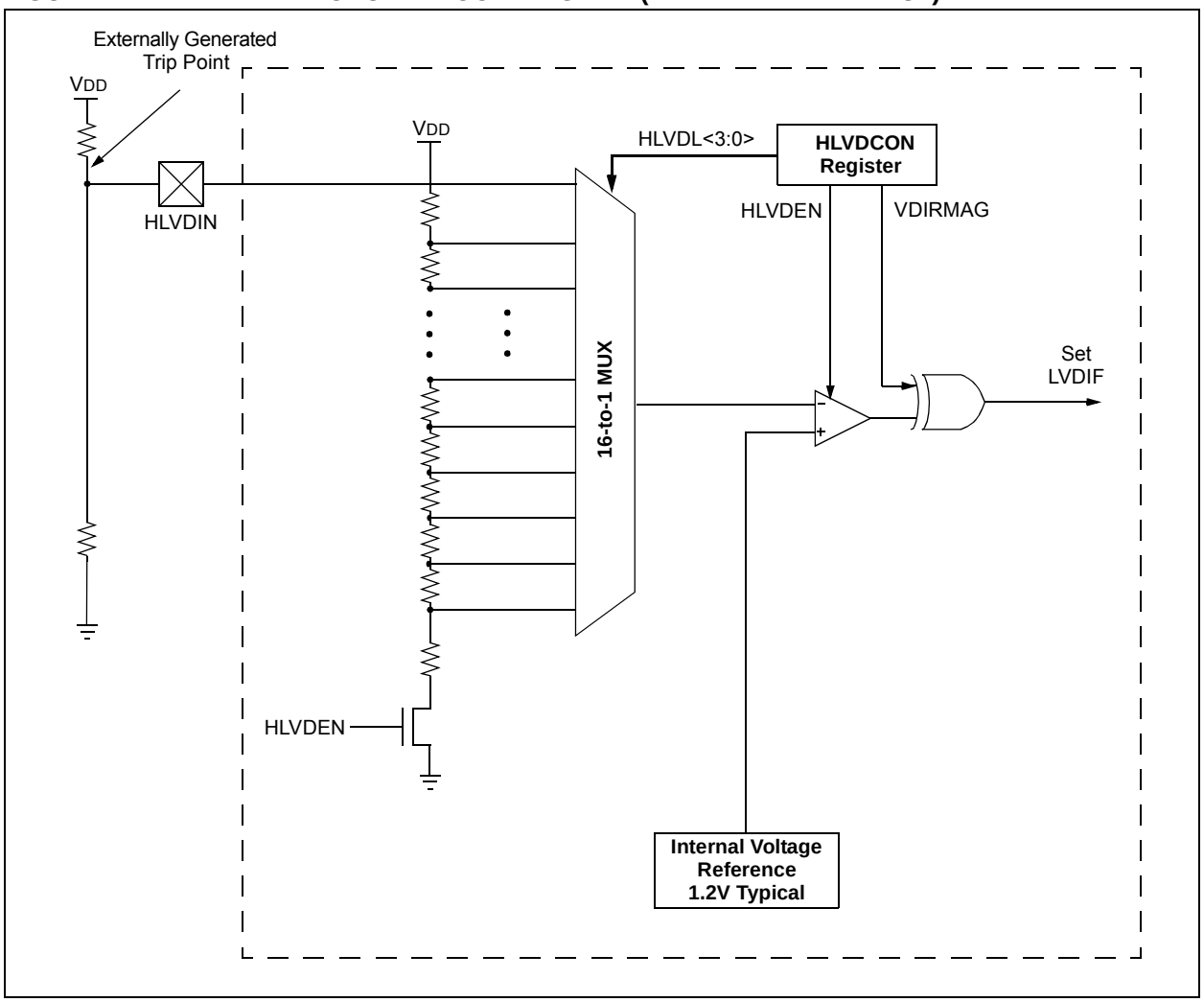
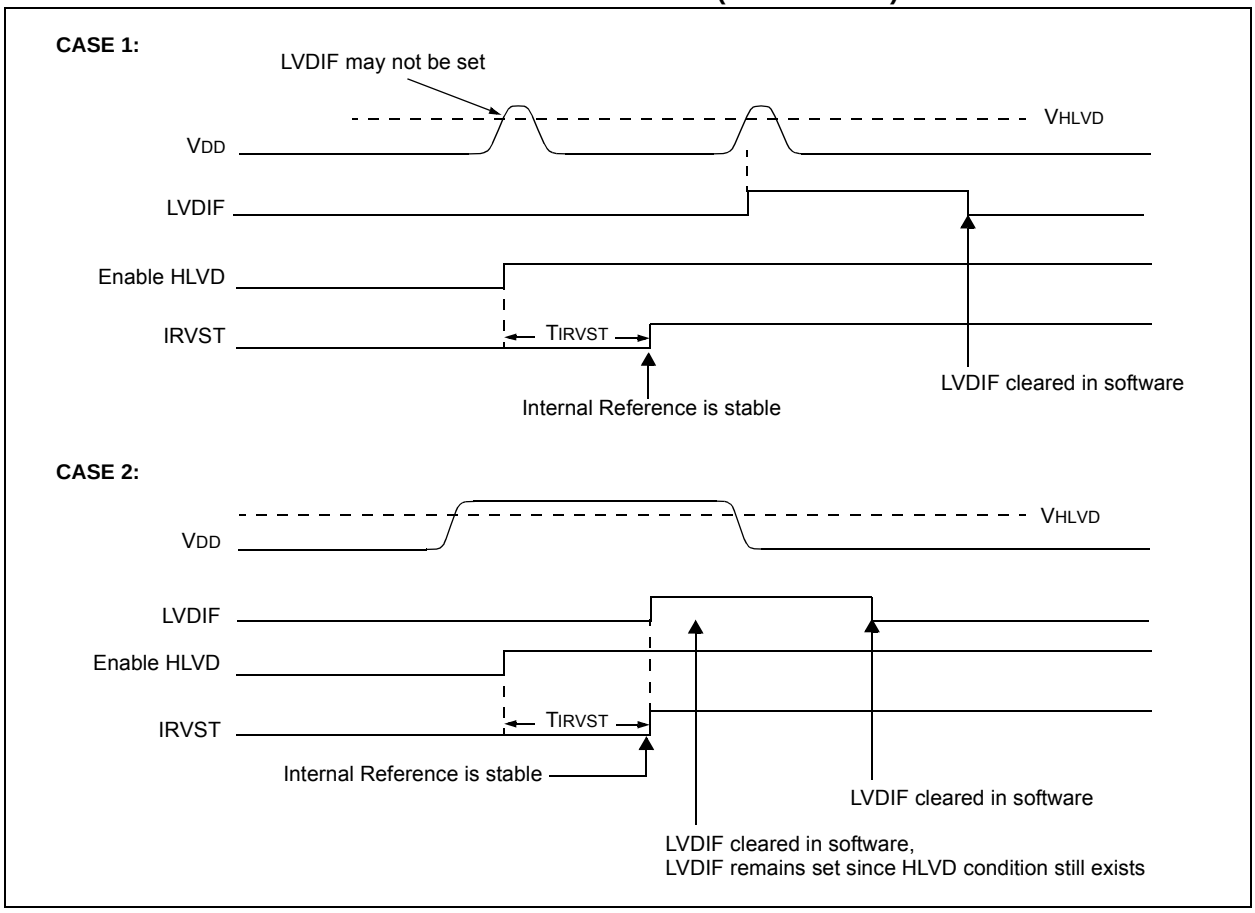


FIGURE 24-3: HIGH-VOLTAGE DETECT OPERATION (VDIRMAG = 1)



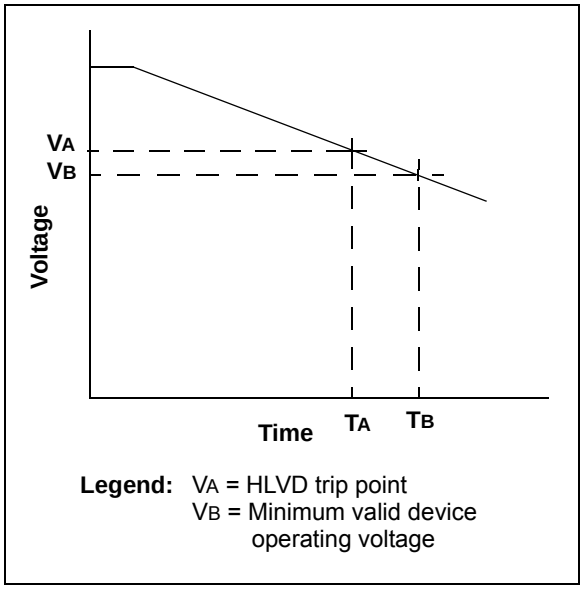
24.5 Applications

In many applications, it is desirable to have the ability to detect a drop below, or rise above, a particular threshold. For general battery applications, Figure 24-4 provides a possible voltage curve.

Over time, the device voltage decreases. When the device voltage reaches voltage, VA, the HLVD logic generates an interrupt at time, TA. The interrupt could cause the execution of an ISR, which would allow the application to perform "housekeeping tasks" and perform a controlled shutdown before the device voltage exits the valid operating range at TB.

The HLVD, thus, would give the application a time window, represented by the difference between TA and TB, to safely exit.

FIGURE 24-4: TYPICAL HIGH/ LOW-VOLTAGE DETECT APPLICATION



EXAMPLE 25-4: ROUTINE FOR CAPACITIVE TOUCH SWITCH

```
#include <p18cxxx.h>

#define COUNT 500                       //@ 8MHz = 125uS.
#define DELAY for(i=0;i<COUNT;i++)
#define OPENSW 1000                     //Un-pressed switch value
#define TRIP 300                        //Difference between pressed
                                        //and un-pressed switch
#define HYST 65                         //amount to change
                                        //from pressed to un-pressed
#define PRESSED 1
#define UNPRESSED 0

int main(void)
{
    unsigned int Vread;                 //storage for reading
    unsigned int switchState;
    int i;

    //assume CTMU and A/D have been setup correctly
    //see Example 25-1 for CTMU & A/D setup
    setup();

    CTMUCONHbits.CTMUEN = 1;            // Enable the CTMU
    CTMUCONLbits.EDG1STAT = 0;          // Set Edge status bits to zero
    CTMUCONLbits.EDG2STAT = 0;
    CTMUCONHbits.IDISSEN = 1;           //drain charge on the circuit
    DELAY;                              //wait 125us
    CTMUCONHbits.IDISSEN = 0;           //end drain of circuit

    CTMUCONLbits.EDG1STAT = 1;          //Begin charging the circuit
                                        //using CTMU current source
    DELAY;                              //wait for 125us
    CTMUCONLbits.EDG1STAT = 0;          //Stop charging circuit

    PIR1bits.ADIF = 0;                  //make sure A/D Int not set
    ADCON0bits.GO=1;                    //and begin A/D conv.
    while(!PIR1bits.ADIF);              //Wait for A/D convert complete

    Vread = ADRES;                      //Get the value from the A/D

    if(Vread < OPENSW - TRIP)
    {
        switchState = PRESSED;
    }
    else if(Vread > OPENSW - TRIP + HYST)
    {
        switchState = UNPRESSED;
    }
}
```

TBLWT **Table Write**

Syntax: TBLWT (*; *+; *-; +*)

Operands: None

Operation: if TBLWT*,
(TABLAT) → Holding Register,
TBLPTR – No Change;
if TBLWT*+,
(TABLAT) → Holding Register,
(TBLPTR) + 1 → TBLPTR;
if TBLWT*-,
(TABLAT) → Holding Register,
(TBLPTR) – 1 → TBLPTR;
if TBLWT+*,
(TBLPTR) + 1 → TBLPTR,
(TABLAT) → Holding Register

Status Affected: None

Encoding:

0000	0000	0000	11nn nn=0 * =1 *+ =2 *- =3 +*

Description: This instruction uses the 3 LSBs of TBLPTR to determine which of the 8 holding registers the TABLAT is written to. The holding registers are used to program the contents of Program Memory (P.M.). (Refer to **Section 6.0 "Memory Organization"** for additional details on programming Flash memory.)

The TBLPTR (a 21-bit pointer) points to each byte in the program memory. TBLPTR has a 2-Mbyte address range. The LSb of the TBLPTR selects which byte of the program memory location to access.

TBLPTR<0> = 0: Least Significant Byte of Program Memory Word

TBLPTR<0> = 1: Most Significant Byte of Program Memory Word

The TBLWT instruction can modify the value of TBLPTR as follows:

- no change
- post-increment
- post-decrement
- pre-increment

Words: 1

Cycles: 2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	No operation	No operation	No operation
No operation	No operation (Read TABLAT)	No operation	No operation (Write to Holding Register)

TBLWT **Table Write (Continued)**

Example 1: TBLWT *+

Before Instruction
TABLAT = 55h
TBLPTR = 00A356h
HOLDING REGISTER (00A356h) = FFh
After Instructions (table write completion)
TABLAT = 55h
TBLPTR = 00A357h
HOLDING REGISTER (00A356h) = 55h

Example 2: TBLWT +*

Before Instruction
TABLAT = 34h
TBLPTR = 01389Ah
HOLDING REGISTER (01389Ah) = FFh
HOLDING REGISTER (01389Bh) = FFh
After Instruction (table write completion)
TABLAT = 34h
TBLPTR = 01389Bh
HOLDING REGISTER (01389Ah) = FFh
HOLDING REGISTER (01389Bh) = 34h

27.2.2 EXTENDED INSTRUCTION SET

ADDFSR **Add Literal to FSR**

Syntax: ADDFSR f, k

Operands: $0 \le k \le 63$
$f \in [0, 1, 2]$

Operation: FSR(f) + k → FSR(f)

Status Affected: None

Encoding:

1110	1000	ffkk	kkkk

Description: The 6-bit literal 'k' is added to the contents of the FSR specified by 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to FSR

Example: `ADDFSR 2, 0x23`

Before Instruction
FSR2 = 03FFh
After Instruction
FSR2 = 0422h

ADDULNK **Add Literal to FSR2 and Return**

Syntax: ADDULNK k

Operands: $0 \le k \le 63$

Operation: FSR2 + k → FSR2,
(TOS) → PC

Status Affected: None

Encoding:

1110	1000	11kk	kkkk

Description: The 6-bit literal 'k' is added to the contents of FSR2. A RETURN is then executed by loading the PC with the TOS.

The instruction takes two cycles to execute; a NOP is performed during the second cycle.

This may be thought of as a special case of the ADDFSR instruction, where f = 3 (binary '11'); it operates only on FSR2.

Words: 1

Cycles: 2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process Data	Write to FSR
No Operation	No Operation	No Operation	No Operation

Example: `ADDULNK 0x23`

Before Instruction
FSR2 = 03FFh
PC = 0100h
After Instruction
FSR2 = 0422h
PC = (TOS)

Note: All PIC18 instructions may take an optional label argument preceding the instruction mnemonic for use in symbolic addressing. If a label is used, the instruction format then becomes: {label} instruction argument(s).

TABLE 29-30: A/D CONVERTER CHARACTERISTICS: PIC18F46J11 FAMILY (INDUSTRIAL)

Param No.	Symbol	Characteristic	Min	Typ	Max	Units	Conditions
A01	NR	Resolution	—	—	10	bit	$\Delta V_{REF} \geq 3.0V$
A03	EIL	Integral Linearity Error	—	—	<±1	LSb	$\Delta V_{REF} \geq 3.0V$
A04	EDL	Differential Linearity Error	—	—	<±1	LSb	$\Delta V_{REF} \geq 3.0V$
A06	EOFF	Offset Error	—	—	<±3	LSb	$\Delta V_{REF} \geq 3.0V$
A07	EGN	Gain Error	—	—	<±3.5	LSb	$\Delta V_{REF} \geq 3.0V$
A10		Monotonicity	Guaranteed(1)			—	$V_{SS} \leq V_{AIN} \leq V_{REF}$
A20	ΔVREF	Reference Voltage Range (VREFH – VREFL)	2.0 3	— —	— —	V V	$V_{DD} < 3.0V$ $V_{DD} \geq 3.0V$
A21	VREFH	Reference Voltage High	VREFL	—	VDD + 0.3V	V	
A22	VREFL	Reference Voltage Low	VSS – 0.3V	—	VREFH	V	
A25	VAIN	Analog Input Voltage	VREFL	—	VREFH	V	
A30	ZAIN	Recommended Impedance of Analog Voltage Source	—	—	2.5	kΩ	
A50	IREF	VREF Input Current(2)	— —	— —	5 150	μA μA	During VAIN acquisition. During A/D conversion cycle.

Note 1: The A/D conversion result never decreases with an increase in the input voltage and has no missing codes.

2: VREFH current is from RA3/AN3/VREF+/C1INB pin or VDD, whichever is selected as the VREFH source. VREFL current is from RA2/AN2/VREF-/CVREF/C2INB pin or VSS, whichever is selected as the VREFL source.