# E·XFL

## AMD Xilinx - XC4005-5PC84C Datasheet



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#### Understanding <u>Embedded - FPGAs (Field</u> <u>Programmable Gate Array)</u>

Embedded - FPGAs, or Field Programmable Gate Arrays, are advanced integrated circuits that offer unparalleled flexibility and performance for digital systems. Unlike traditional fixed-function logic devices, FPGAs can be programmed and reprogrammed to execute a wide array of logical operations, enabling customized functionality tailored to specific applications. This reprogrammability allows developers to iterate designs quickly and implement complex functions without the need for custom hardware.

#### **Applications of Embedded - FPGAs**

The versatility of Embedded - FPGAs makes them indispensable in numerous fields. In telecommunications.

#### Details

Details	
Product Status	Obsolete
Number of LABs/CLBs	196
Number of Logic Elements/Cells	466
Total RAM Bits	6272
Number of I/O	61
Number of Gates	5000
Voltage - Supply	4.75V ~ 5.25V
Mounting Type	Surface Mount
Operating Temperature	0°C ~ 85°C (TJ)
Package / Case	84-LCC (J-Lead)
Supplier Device Package	84-PLCC (29.31x29.31)
Purchase URL	https://www.e-xfl.com/product-detail/xilinx/xc4005-5pc84c

Email: info@E-XFL.COM

Address: Room A, 16/F, Full Win Commercial Centre, 573 Nathan Road, Mongkok, Hong Kong

## XC4000 Compared to XC3000A

For those readers already familiar with the XC3000A family of Xilinx Field Programmable Gate Arrays, here is a concise list of the major new features in the XC4000 family.

CLB has two **independent** 4-input function generators. A **third** function generator combines the outputs of the two other function generators with a ninth input. All function inputs are swappable, all have full access; none are mutually exclusive.

- CLB has very fast arithmetic carry capability.
- CLB function generator look-up table can also be used as high-speed **RAM**.
- CLB flip-flops have asynchronous set or reset.
- CLB has four outputs, two flip-flops, two combinatorial.
- CLB connections symmetrically located on all four edges.
- **IOB** has more versatile clocking polarity options.

**IOB** has programmable input set-up time: **long** to avoid potential hold time problems,

**short** to improve performance. **IOB** has Longline access through its own TBUF.

Outputs are **n-channel only**, lower V<sub>OH</sub> increases speed. XC4000 outputs can be paired to double sink current to **24 mA.** XC4000A and XC4000H outputs can each

sink 24 mA, can be paired for **48 mA** sink current.

IEEE 1149.1- type **boundary scan** is supported in the I/O.

Wide decoders on all four edges of the LCA device.

Increased number of interconnect resources.

All CLB inputs and outputs have access to most interconnect lines.

Switch Matrices are simplified to increase speed.

- **Eight global nets** can be used for clocking or distributing logic signals.
- **TBUF** output configuration is more versatile and 3-state control less confined.

**Program** is single-function input pin,overrides everything. **INIT** pin also acts as Configuration Error output.

Peripheral Synchronous Mode (8 bit) has been added. Peripheral Asynchronous Mode has improved handshake.

- **Start-up** can be **synchronized** to any user clock (this is a configuration option).
- No Powerdown, but instead a **Global 3-state input** that does not reset any flip-flops.

No on-chip crystal oscillator amplifier.

Configuration Bit Stream includes CRC error checking. Configuration Clock can be increased to >8 MHz.

- Configuration Clock is **fully static**, no constraint on the maximum Low time.
- **Readback** either ignores flip-flop content (avoids need for masking) or it takes a **snapshot** of all flip-flops at the start of Readback.
- Readback has same **polarity** as Configuration and can be **aborted.**

 Table 2. Three Generations of Xilinx Field-Programmable Gate Array Families

Parameter	XC4025	XC3195A	XC2018
Number of flip-flops	2,560	1,320	174
Max number of user I/O	256	176	74
Max number of RAM bits	32,768	0	0
Function generators per CLB	3	2	2
Number of logic inputs per CLB	9	5	4
Number of logic outputs per CLB	4	2	2
Number of low-skew global nets	8	2	2
Dedicated decoders	yes	no	no
Fast carry logic	yes	no	no
Internal 3-state drivers	yes	yes	no
Output slew-rate control	yes	yes	no
Power-down option	no	yes	yes
Crystal oscillator circuit	no	yes	yes

## **Architectural Overview**

The XC4000 families achieve high speed through advanced semiconductor technology and through improved architecture, and supports system clock rates of up to 50 MHz. Compared to older Xilinx FPGA families, the XC4000 families are more powerful, offering on-chip RAM and wide-input decoders. They are more versatile in their applications, and design cycles are faster due to a combination of increased routing resources and more sophisticated software. And last, but not least, they more than double the available complexity, up to the 20,000-gate level.

The XC4000 families have 16 members, ranging in complexity from 2,000 to 25,000 gates.

#### **Logic Cell Array Families**

Xilinx high-density user-programmable gate arrays include three major configurable elements: configurable logic blocks (CLBs), input/output blocks (IOBs), and interconnections. The CLBs provide the functional elements for constructing the user's logic. The IOBs provide the interface between the package pins and internal signal lines. The programmable interconnect resources provide routing paths to connect the inputs and outputs of the CLBs and IOBs onto the appropriate networks. Customized configuration is established by programming internal static memory cells that determine the logic functions and interconnections implemented in the LCA device.

The first generation of LCA devices, the XC2000 family, was introduced in 1985. It featured logic blocks consisting of a combinatorial function generator capable of implementing 4-input Boolean functions and a single storage element. The XC2000 family has two members ranging in complexity from 800 to 1500 gates.

In the second-generation XC3000A LCA devices, introduced in 1987, the logic block was expanded to implement wider Boolean functions and to incorporate a second flipflop in each logic block. Today, the XC3000 devices range in complexity from 1,300 to 10,000 usable gates. They have a maximum guaranteed toggle frequency ranging from 70 to 270 MHz, equivalent to maximum system clock frequencies of up to 80 MHz.

The third generation of LCA devices further extends this architecture with a yet more powerful and flexible logic block. I/O block functions and interconnection options have also been enhanced with each successive generation, further extending the range of applications that can be implemented with an LCA device.

This third-generation architecture forms the basis of the XC4000 families of devices that feature logic densities up to 25,000 usable gates and support system clock rates of

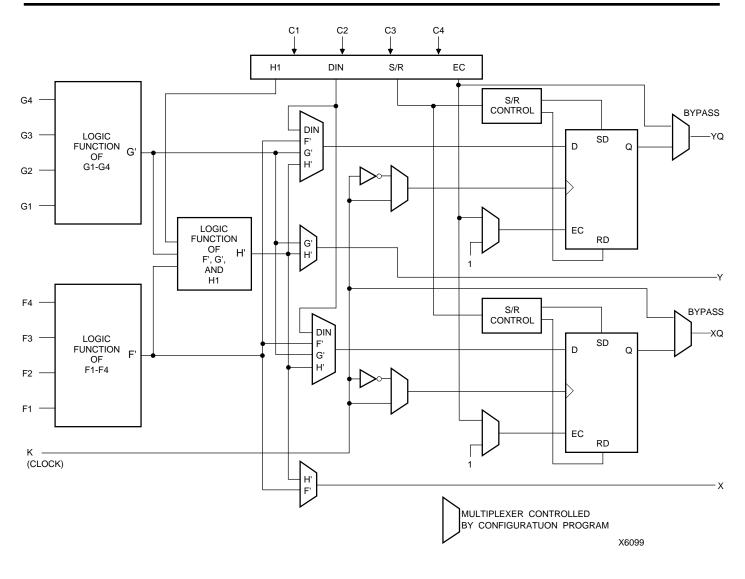
up to 50 MHz. The use of an advanced, sub-micron CMOS process technology as well as architectural improvements contribute to this increase in FPGA capabilities. However, achieving these high logic-density and performance levels also requires new and more powerful automated design tools. IC and software engineers collaborated during the definition of the third-generation LCA architecture to meet an important performance goal — an FPGA architecture and companion design tools for completely automatic placement and routing of 95% of all designs, plus a convenient way to complete the remaining few designs.

## **Configurable Logic Blocks**

A number of architectural improvements contribute to the increased logic density and performance levels of the XC4000 families. The most important one is a more powerful and flexible CLB surrounded by a versatile set of routing resources, resulting in more "effective gates per CLB." The principal CLB elements are shown in Figure 1. Each new CLB also packs a pair of flip-flops and two independent 4-input function generators. The two function generators offer designers plenty of flexibility because most combinatorial logic functions need less than four inputs. Consequently, the design-software tools can deal with each function generator independently, thus improving cell usage.

Thirteen CLB inputs and four CLB outputs provide access to the function generators and flip-flops. More than double the number available in the XC3000 families, these inputs and outputs connect to the programmable interconnect resources outside the block. Four independent inputs are provided to each of two function generators (F1 - F4 and G1 – G4). These function generators, whose outputs are labeled F' and G', are each capable of implementing any arbitrarily defined Boolean function of their four inputs. The function generators are implemented as memory look-up tables; therefore, the propagation delay is independent of the function being implemented. A third function generator, labeled H', can implement any Boolean function of its three inputs: F' and G' and a third input from outside the block (H1). Signals from the function generators can exit the CLB on two outputs; F' or H' can be connected to the X output, and G' or H' can be connected to the Y output. Thus, a CLB can be used to implement any two independent functions of up-to-four variables, or any single function of five variables, or any function of four variables together with some functions of five variables, or it can implement even some functions of up to nine variables. Implementing wide functions in a single block reduces both the number of blocks required and the delay in the signal path, achieving both increased density and speed.

The two storage elements in the CLB are edge-triggered D-type flip-flops with common clock (K) and clock enable (EC) inputs. A third common input (S/R) can be programmed as either an asynchronous set or reset signal





independently for each of the two registers; this input also can be disabled for either flip-flop. A separate global Set/ Reset line (not shown in Figure 1) sets or clears each register during power-up, reconfiguration, or when a dedicated Reset net is driven active. This Reset net does not compete with other routing resources; it can be connected to any package pin as a global reset input.

Each flip-flop can be triggered on either the rising or falling clock edge. The source of a flip-flop data input is programmable: it is driven either by the functions F', G', and H', or the Direct In (DIN) block input. The flip-flops drive the XQ and YQ CLB outputs.

In addition, each CLB F' and G' function generator contains dedicated arithmetic logic for the fast generation of carry and borrow signals, greatly increasing the efficiency and performance of adders, subtracters, accumulators, comparators and even counters.

Multiplexers in the CLB map the four control inputs, labeled C1 through C4 in Figure 1, into the four internal control signals (H1, DIN, S/R, and EC) in any arbitrary manner.

The flexibility and symmetry of the CLB architecture facilitates the placement and routing of a given application. Since the function generators and flip-flops have independent inputs and outputs, each can be treated as a separate entity during placement to achieve high packing density. Inputs, outputs, and the functions themselves can freely swap positions within a CLB to avoid routing congestion during the placement and routing operation.

#### Speed Is Enhanced Two Ways

Delays in LCA-based designs are layout dependent. While this makes it hard to predict a worst-case guaranteed performance, there is a rule of thumb designers can consider — the system clock rate should not exceed one third to one half of the specified toggle rate. Critical portions of a design, shift registers and simple counters, can run faster — approximately two thirds of the specified toggle rate.

The XC4000 family can run at synchronous system clock rates of up to 60 MHz. This increase in performance over the previous families stems from two basic improvements: improved architecture and more abundant routing resources.

#### Improved Architecture

More Inputs: The versatility of the CLB function generators improves system speed significantly. Table 3 shows how the XC4000 families implement many functions more efficiently and faster than is possible with XC3000 devices. A 9-bit parity checker, for example, can be implemented in one CLB with a propagation delay of 7 ns. Using a XC3000-family device, the same function requires two CLBs with a propagation delay of  $2 \times 5.5$  ns = 11 ns. One XC4000 CLB can determine whether two 4-bit words are identical, again with a 7-ns propagation delay. The ninth input can be used for simple ripple expansion of this identity comparator (25.5 ns over 16 bits, 51.5 ns over 32 bits), or a 2-layer identity comparator can generate the result of a 32-bit comparison in 15 ns, at the cost of a single extra CLB. Simpler functions like multiplexers also benefit from the greater flexibility of the XC4000-families CLB. A 16-input multiplexer uses 5 CLBs and has a delay of only 13.5 ns.

*More Outputs:* The CLB can pass the combinatorial output(s) to the interconnect network, but can also store the combinatorial result(s) or other incoming data in one or two flip-flops, and connect their outputs to the interconnect

network as well. With XC3000-families CLBs the designer has to make a choice, either output the combinatorial function or the stored value. In the XC4000 families, the flip flops can be used as registers or shift registers without blocking the function generators from performing a different, perhaps unrelated task. This increases the functional density of the devices.

When a function generator drives a flip-flop in a CLB, the combinatorial propagation delay *overlaps completely* with the set-up time of the flip-flop. The set-up time is specified between the function generator inputs and the clock input. This represents a performance advantage over competing technologies where combinatorial delays must be added to the flip-flop set-up time.

*Fast Carry:* As described earlier, each CLB includes highspeed carry logic that can be activated by configuration. The two 4-input function generators can be configured as a 2-bit adder with built-in hidden carry that can be expanded to any length. This dedicated carry circuitry is so fast and efficient that conventional speed-up methods like carry generate/propagate are meaningless even at the 16-bit level, and of marginal benefit at the 32-bit level.

A 16-bit adder requires nine CLBs and has a combinatorial carry delay of 20.5 ns. Compare that to the 30 CLBs and 50 ns, or 41 CLBs and 30 ns in the XC3000 family.

The fast-carry logic opens the door to many new applications involving arithmetic operation, where the previous generations of FPGAs were not fast and/or not efficient enough. High-speed address offset calculations in microprocessor or graphics systems, and high-speed addition in digital signal processing are two typical applications.

*Faster and More Efficient Counters:* The XC4000-families fast-carry logic puts two counter bits into each CLB and runs them at a clock rate of up to 42 MHz for 16 bits, whether the counters are loadable or not. For a 16-bit

		XC300	0 (-125)	XC4000 (-5)		
16-bit Decoder From Input P	ad	15 ns	4 CLBs	12 ns	0 CLBs	
24-bit Accumulator		17 MHz	46 CLBs	32 MHz	13 CLBs	
State Machine Benchmark*		18 MHz	34 CLBs	30 MHz	26 CLBs	
16:1 Multiplexer		16 ns	8 CLBs	16 ns	5 CLBs	
16-bit Unidirectional	Max Density	20 MHz	16 CLBs	40 MHz	8 CLBs	
Loadable Counter	Max Speed	34 MHz	23 CLBs	42 MHz	9 CLBs	
16-bit U/D Counter	Max Density	20 MHz	16 CLBs	40 MHz	8 CLBs	
	Max Speed	30 MHz	27 CLBs	40 MHz	8 CLBs	
16-bit Adder	Max Density	50 ns	30 CLBs	20.5 ns	9 CLBs	
	Max Speed	30 ns	41 CLBs	20.5 ns	9 CLBs	

#### Table 3. Density and Performance for Several Common Circuit Functions

\* 16 states, 40 transitions, 10 inputs, 8 outputs

pass through a global buffer before arriving at the IOB. This eliminates the possibility of a data hold-time requirement at the external pin. The I1 and I2 signals that exit the block can each carry either the direct or registered input signal.

Output signals can be inverted or not inverted, and can pass directly to the pad or be stored in an edge-triggered flip-flop. Optionally, an output enable signal can be used to place the output buffer in a high-impedance state, implementing 3-state outputs or bidirectional I/O. Under configuration control, the output (OUT) and output enable (OE) signals can be inverted, and the slew rate of the output buffer can be reduced to minimize power bus transients when switching non-critical signals. Each XC4000-families output buffer is capable of sinking 12 mA; two adjacent output buffers can be wire-ANDed externally to sink up to 24 mA. In the XC4000A and XC4000H families, each output buffer can sink 24 mA.

There are a number of other programmable options in the IOB. Programmable pull-up and pull-down resistors are useful for tying unused pins to  $V_{CC}$  or ground to minimize power consumption. Separate clock signals are provided for the input and output registers; these clocks can be inverted, generating either falling-edge or rising-edge triggered flip-flops. As is the case with the CLB registers, a global set/reset signal can be used to set or clear the input and output registers whenever the RESET net is active.

Embedded logic attached to the IOBs contains test structures compatible with IEEE Standard 1149.1 for boundaryscan testing, permitting easy chip and board-level testing.



All internal connections are composed of metal segments with programmable switching points to implement the desired routing. An abundance of different routing resources is provided to achieve efficient automated routing. The number of routing channels is scaled to the size of the array; i.e., it increases with array size.

In previous generations of LCAs, the logic-block inputs were located on the top, left, and bottom of the block; outputs exited the block on the right, favoring left-to-right data flow through the device. For the third-generation family, the CLB inputs and outputs are distributed on all four sides of the block, providing additional routing flexibility (Figure 6). In general, the entire architecture is more symmetrical and regular than that of earlier generations, and is more suited to well-established placement and routing algorithms developed for conventional mask- programmed gate-array design.

There are three main types of interconnect, distinguished by the relative length of their segments: single-length lines, double-length lines, and Longlines. Note: The number of routing channels shown in Figures 6 and 9 are for illustration purposes only; the actual number of routing channels varies with array size. The routing scheme was designed for minimum resistance and capacitance of the average routing path, resulting in significant performance improvements.

The single-length lines are a grid of horizontal and vertical lines that intersect at a Switch Matrix between each block. Figure 6 illustrates the single-length interconnect lines

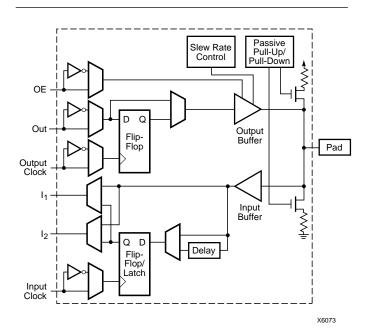


Figure 5. XC4000 and XC4000A Families Input/Output Block

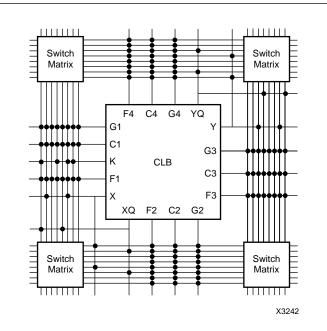


Figure 6. Typical CLB Connections to Adjacent Single-Length Lines

surrounding one CLB in the array. Each Switch Matrix consists of programmable n-channel pass transistors used to establish connections between the single-length lines (Figure 7). For example, a signal entering on the right side of the Switch Matrix can be routed to a single-length line on the top, left, or bottom sides, or any combination thereof, if multiple branches are required. Single-length lines are normally used to conduct signals within a localized area and to provide the branching for nets with fanout greater than one.

Compared to the previous generations of LCA architectures, the number of possible connections through the Switch Matrix has been reduced. This decreases capacitive loading and minimizes routing delays, thus increasing performance. However, a much more versatile set of connections between the single-length lines and the CLB inputs and outputs more than compensate for the reduction in Switch Matrix options, resulting in overall increased routability.

The function generator and control inputs to the CLB (F1-F4, G1-G4, and C1-C4) can be driven from any adjacent single-length line segment (Figure 6). The CLB clock (K) input can be driven from one-half of the adjacent singlelength lines. Each CLB output can drive several of the single-length lines, with connections to both the horizontal and vertical Longlines.

The double-length lines (Figure 8) consist of a grid of metal segments twice as long as the single-length lines; i.e, a double-length line runs past two CLBs before entering a Switch Matrix. Double-length lines are grouped in pairs with the Switch Matrices staggered so that each line goes through a Switch Matrix at every other CLB location in that row or column. As with single-length lines, all the CLB inputs except K can be driven from any adjacent double-length lines in both the vertical and horizontal planes. Double-length lines provide the most efficient implementation of intermediate length, point-to-point interconnections.

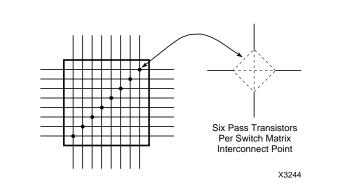
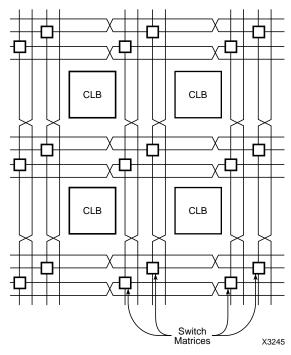


Figure 7. Switch Matrix



**Figure 8. Double-Length Lines** 

Longlines form a grid of metal interconnect segments that run the entire length or width of the array (Figure 9). Additional vertical longlines can be driven by special global buffers, designed to distribute clocks and other high fanout control signals throughout the array with minimal skew. Longlines are intended for high fan-out, time-critical signal nets. Each Longline has a programmable splitter switch at its center, that can separate the line into two independent routing channels, each running half the width or height of the array. CLB inputs can be driven from a subset of the adjacent Longlines; CLB outputs are routed to the Longlines via 3-state buffers or the single-length interconnected lines.

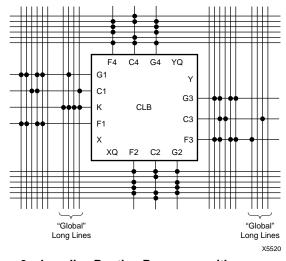


Figure 9. Longline Routing Resources with Typical CLB Connections

The XACT system also includes XDelay, a static timing analyzer. XDelay examines a design's logic and timing to calculate the performance along signal paths, identify possible race conditions, and detect set-up and hold-time violations. Timing analyzers do not require that the user generate input stimulus patterns or test vectors.

# Summary

The result of eight years of FPGA design experience and feedback from thousands of customers, the XC4000 families combine architectural versatility, on-chip RAM, increased speed and gate complexity with abundant routing resources and new, sophisticated software to achieve fully automated implementation of complex, high-performance designs.

7400 Equivalents		Barrel Shifters		Multiplexers	
	# of CLBs			m2-1e	1
'138	5	brlshft4	4	m4-1e	1
'139		brlshft8	13	m8-1e	
'147	2 5			m16-1e	3 5
'148	6	4-Bit Counters			C C
'150	5	cd4ce	3	Registers	
'151	5 3	cd4cle	5	-	
'152	3	cd4rle	6	rd4r	2
·153	2	cb4ce	3	rd8r	4
'154	16	cb4cle	6	rd16r	8
'157	2	cb4cle cb4re	5		
'158	2	CD4IE	5	Shift Registers	
'160	5	8- and 16-Bit Co	untore	sr8ce	4
'161	6	8- and 10-Bit CO	uniters	sroce sr16re	4 8
'162	8	cb8ce	6	SITOLE	0
·163	8	cb8re	10	RAMs	
'164	4	cc16ce	10	RANIS	
'165s	9	cc16cle	11	ram 16x4	2
'166	5	cc16cled	21		
'168	7			Explanation of co	ounter nomenclature
·174	3	Identity Compar	ators	-	
'194	5			cb = binary counter	
'195	3	comp4	1	cd = BCD counter	
'280	3	comp8	2	cc = cascadable b	pinary counter
<sup>283</sup>	3 8	comp16	5	d = bidirectional	
·298				I = loadable	
·352	2 2 3	Magnitude Com	parators	x = cascadable	
·390		compm4	4	e = clock enable	
·518	3	compm8	9	r = synchronous	reset
·521	3	compm16	20	c = asynchronou	s clear
521	5	compinite	20		
		Decoders			
		d2-4e	2		
		d3-8e	4		
		d4-16e	16		
		41.100			

Figure 10. CLB Count of Selected XC4000 Soft Macros

## **Detailed Functional Description**

## XC4000 and XC4000A Input/Output Blocks

(For XC4000H family, see page 2-82)

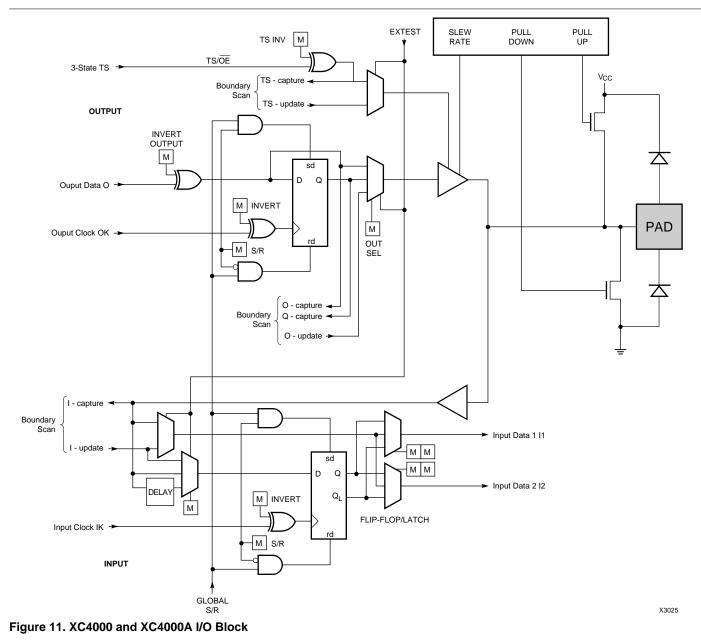
The IOB forms the interface between the internal logic and the I/O pads of the LCA device. Under configuration control, the output buffer receives either the logic signal (.out) routed from the internal logic to the IOB, or the complement of this signal, or this same data after it has been clocked into the output flip-flop.

As a configuration option, each flip-flop (CLB or IOB) is initialized as either set or reset, and is also forced into this programmable initialization state whenever the global Set/ Reset net is activated after configuration has been completed. The clock polarity of each IOB flip-flop can be configured individually, as can the polarity of the 3-state control for the output buffer. Each output buffer can be configured to be either fast or slew-rate limited, which reduces noise generation and ground bounce. Each I/O pin can be configured with either an internal pull-up or pull down resistor, or with no internal resistor. Independent of this choice, each IOB has a pullup resistor during the configuration process.

The 3-state output driver uses a totem pole n-channel output structure.  $V_{\rm OH}$  is one n-channel threshold lower than  $V_{\rm CC},$  which makes rise and fall delays more symmetrical.

Family	Per IOB Source	Per IOB Sink	Per IOB Pair Sink	# Slew Modes
XC4000	4	12	24	2
XC4000A	4	24	48	4
XC4000H	4	24*	48	2

\*XC4000H devices can sink only 4 mA configured for SoftEdge mode



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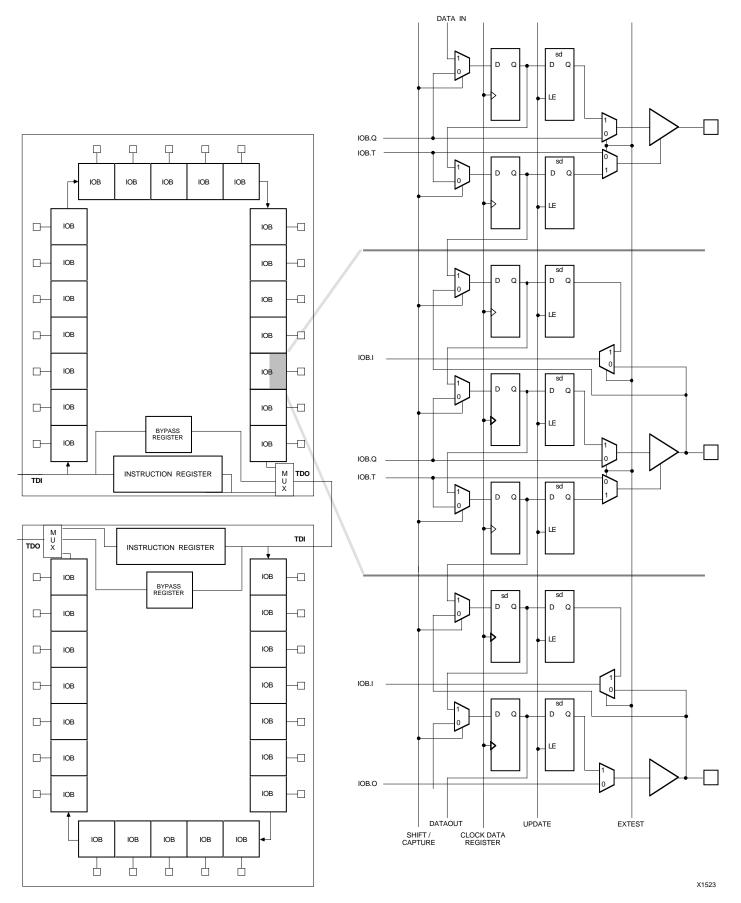


Figure 16. XC4000 Boundary Scan Logic. Includes three bits of Data Register per IOB, the IEEE 1149.1 Test Access Port controller, and the Instruction Register with decodes.

## Oscillator

An internal oscillator is used for clocking of the power-on time-out, configuration memory clearing, and as the source of CCLK in Master modes. This oscillator signal runs at a nominal 8 MHz and varies with process,  $V_{CC}$  and temperature between 10 MHz max and 4 MHz min. This signal is available on an output control net (OSCO) in the upper right corner of the chip, if the oscillator-run control bit is enabled in the configuration memory. Two of four resynchronized taps of the power-on time-out divider are also available on OSC1 and OSC2. These taps are at the fourth, ninth, fourteenth and nineteenth bits of the ripple divider. This can provide output signals of approximately 500 kHz, 16 kHz, 490 Hz and 15 Hz.

## **Special Purpose Pins**

The mode pins are sampled prior to configuration to determine the configuration mode and timing options. After configuration, these pins can be used as auxiliary connections: Mode 0 (MD0.I) and Mode 2 (MD2.I) as inputs and Mode 1 (MD1.O and MD1.T) as an output. The XACT development system will not use these resources unless they are explicitly specified in the design entry. These dedicated nets are located in the lower left chip corner and are near the readback nets. This allows convenient routing if compatibility with the XC2000 and XC3000 family conventions of M0/RT, M1/RD is desired.

#### **Table 6. Configuration Modes**

Mode	M2	M1	MO	CCLK	Data					
Master Serial Slave Serial	0 1	0 1	0 1	output input	Bit-Serial Bit-Serial					
Master Parallel up Master Parallel down				output output	Byte-Wide, 00000 ↑ Byte-Wide, 3FFFF↓					
Peripheral Synchr. Peripheral Asynchr.	0 1	1 0	1 1	input output	Byte-Wide Byte-Wide					
Reserved Reserved	0 0	1 0	0 1	_	=					
Peripheral Synchronous can be considered Slave Parallel										

# Configuration

Configuration is the process of loading design-specific programming data into one or more LCA devices to define the functional operation of the internal blocks and their interconnections. This is somewhat like loading the command registers of a programmable peripheral chip. The XC4000 families use about 350 bits of configuration data per CLB and its associated interconnects. Each configuration bit defines the state of a static memory cell that controls either a function look-up table bit, a multiplexer input, or an interconnect pass transistor. The XACT development system translates the design into a netlist file. It automatically partitions, places and routes the logic and generates the configuration data in PROM format.

## Modes

The XC4000 families have six configuration modes selected by a 3- bit input code applied to the M0, M1, and M2 inputs. There are three self-loading Master modes, two Peripheral modes and the Serial Slave mode used primarily for daisy-chained devices. During configuration, some of the I/O pins are used temporarily for the configuration process. See Table 6.

For a detailed description of these configuration modes, see pages 2-32 through 2-41.

#### Master

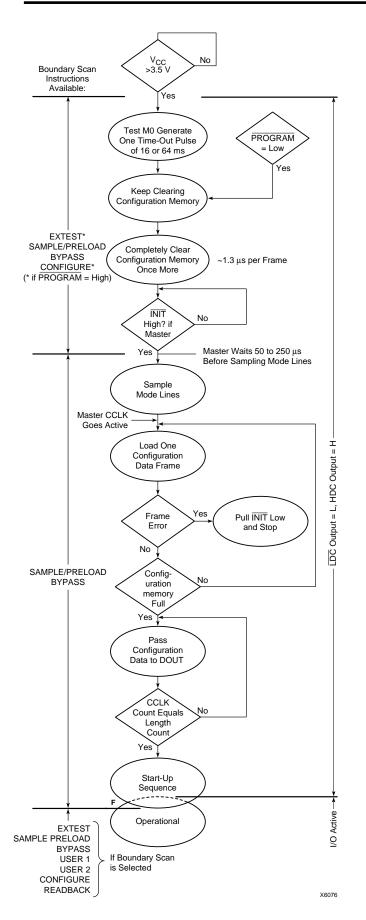
The Master modes use an internal oscillator to generate CCLK for driving potential slave devices, and to generate address and timing for external PROM(s) containing the configuration data. Master Parallel (up or down) modes generate the CCLK signal and PROM addresses and receive byte parallel data, which is internally serialized into the LCA data-frame format. The up and down selection generates starting addresses at either zero or 3FFFF, to be compatible with different microprocessor addressing conventions. The Master Serial mode generates CCLK and receives the configuration data in serial form from a Xilinx serial-configuration PROM.

## Peripheral

The two Peripheral modes accept byte-wide data from a bus. A READY/BUSY status is available as a handshake signal. In the asynchronous mode, the internal oscillator generates a CCLK burst signal that serializes the byte-wide data. In the synchronous mode, an externally supplied clock input to CCLK serializes the data.

#### Serial Slave

In the Serial Slave mode, the LCA device receives serialconfiguration data on the rising edge of CCLK and, after loading its configuration, passes additional data out, resynchronized on the next falling edge of CCLK. Multiple slave devices with identical configurations can be wired with parallel DIN inputs so that the devices can be configured simultaneously.



device when it recognizes the 0010 preamble. Following the length-count data, any LCA device outputs a High on DOUT until it has received its required number of data frames.

After an LCA device has received its configuration data, it passes on any additional frame start bits and configuration data on DOUT. When the total number of configuration clocks applied after memory initialization equals the value of the 24-bit length count, the LCA device(s) begin the start-up sequence and become operational together.

#### Configuration Sequence Configuration Memory Clear

When power is first applied or reapplied to an LCA device, an internal circuit forces initialization of the configuration logic. When  $V_{CC}$  reaches an operational level, and the circuit passes the write and read test of a sample pair of configuration bits, a nominal 16-ms time delay is started (four times longer when M0 is Low, i.e., in Master mode). During this time delay, or as long as the <u>PROGRAM</u> input is asserted, the configuration logic is held in a Configuration Memory Clear state. The configuration-memory frames are consecutively initialized, using the internal oscillator. At the end of each complete pass through the frame addressing, the power-on time-out delay circuitry and the level of the <u>PROGRAM</u> pin are tested. If neither is asserted, the logic initiates one additional clearing of the configuration frames and then tests the <u>INIT</u> input.

# Initialization

During initialization and configuration, user pins HDC, <u>LDC</u> and <u>INIT</u> provide status outputs for system interface. The outputs, <u>LDC</u>, <u>INIT</u> and DONE are held Low and HDC is held High starting at the initial application of power. The open drain <u>INIT</u> pin is released after the final initialization pass through the frame addresses. There is a deliberate delay of 50 to 250  $\mu$ s before a Master-mode device recognizes an inactive <u>INIT</u>. Two internal clocks after the <u>INIT</u> pin is recognized as High, the LCA device samples the three mode lines to determine the configuration mode. The appropriate interface lines become active and the configuration preamble and data can be loaded.

## Configuration

The 0010 preamble code indicates that the following 24 bits represent the length count, i.e., the total number of configuration clocks needed to load the total configuration data. After the preamble and the length count have been passed through to all devices in the daisy chain, DOUT is held High to prevent frame start bits from reaching any daisy-chained devices. A specific configuration bit, early in the first frame of a master device, controls the configuration-clock rate and can increase it by a factor of eight. Each frame has a Low start bit followed by the frame-configura-

Figure 20. Start-up Sequence

tion data bits and a 4-bit frame error field. If a frame data error is detected, the LCA device halts loading, and signals the error by pulling the open-drain <u>INIT</u> pin Low.

After all configuration frames have been loaded into an LCA device, DOUT again follows the input data so that the remaining data is passed on to the next device.

# Start-Up

Start-up is the transition from the configuration process to the intended user operation. This means a change from one clock source to another, and a change from interfacing parallel or serial configuration data where most outputs are 3-stated, to normal operation with I/O pins active in the user-system. Start-up must make sure that the user-logic "wakes up" gracefully, that the outputs become active without causing contention with the configuration signals, and that the internal flip-flops are released from the global Reset or Set at the right time.

Figure 21 describes Start-up timing for the three Xilinx families in detail.

The **XC2000** family goes through a fixed sequence:

DONE goes High and the internal global Reset is deactivated one CCLK period after the I/O become active.

The **XC3000A** family offers some flexibility: DONE can be programmed to go High one CCLK period before or after the I/O become active. Independent of DONE, the internal global Reset is de-activated one CCLK period before or after the I/O become active.

The **XC4000** family offers additional flexibility: The three events, DONE going High, the internal Reset/Set being de-activated, and the user I/O going active, can all occur in any arbitrary sequence, each of them one CCLK period before or after, or simultaneous with, any of the other.

The default option, and the most practical one, is for DONE to go High first, disconnecting the configuration data source and avoiding any contention when the I/Os become active one clock later. Reset/Set is then released another clock period later to make sure that user-operation starts from stable internal conditions. This is the most common sequence, shown with heavy lines in Figure 21, but the designer can modify it to meet particular requirements.

The XC4000 family offers another start-up clocking option: The three events described above don't have to be triggered by CCLK, they can, as a configuration option, be triggered by a user clock. This means that the device can wake up in synchronism with the user system. The XC4000 family introduces an additional option: When this option is enabled, the user can externally hold the open-drain DONE output Low, and thus stall all further progress in the Start-up sequence, until DONE is released and has gone High. This option can be used to force synchronization of several LCA devices to a common user clock, or to guarantee that all devices are successfully configured before any I/Os go active.

# Start-up Sequence

The Start-up sequence begins when the configuration memory is full, and the total number of configuration clocks received since INIT went High equals the loaded value of the length count. The next rising clock edge sets a flip-flop Q0 (see Figure 22), the leading bit of a 5-bit shift register.

The outputs of this register can be programmed to control three events.

- The release of the open-drain DONE output,
- The change of configuration-related pins to the user function, activating all IOBs.
- The termination of the global Set/Reset initialization of all CLB and IOB storage elements.

The DONE pin can also be wire-ANDed with DONE pins of other LCA devices or with other external signals, and can then be used as input to bit Q3 of the start-up register. This is called "Start-up Timing Synchronous to Done In" and labeled: CCLK\_SYNC or UCLK\_SYNC. When DONE is not used as an input, the operation is called Start-up Timing Not Synchronous to DONE In, and is labeled CCLK\_NOSYNC or UCLK\_NOSYNC. These labels are not intuitively obvious.

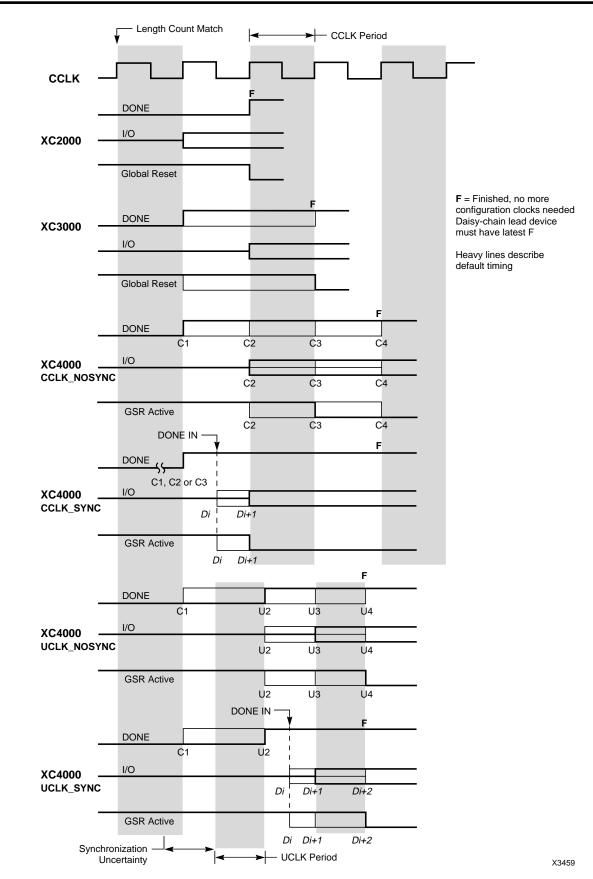
As a configuration option, the start-up control register beyond Q0 can be clocked either by subsequent CCLK pulses or from an on-chip user net called STARTUP.CLK.

# Start-up from CCLK

If CCLK is used to drive the start-up, Q0 through Q3 provide the timing. Heavy lines in Figure 21 show the default timing which is compatible with XC2000 and XC3000 devices using early DONE and late Reset. The thin lines indicate all other possible timing options.

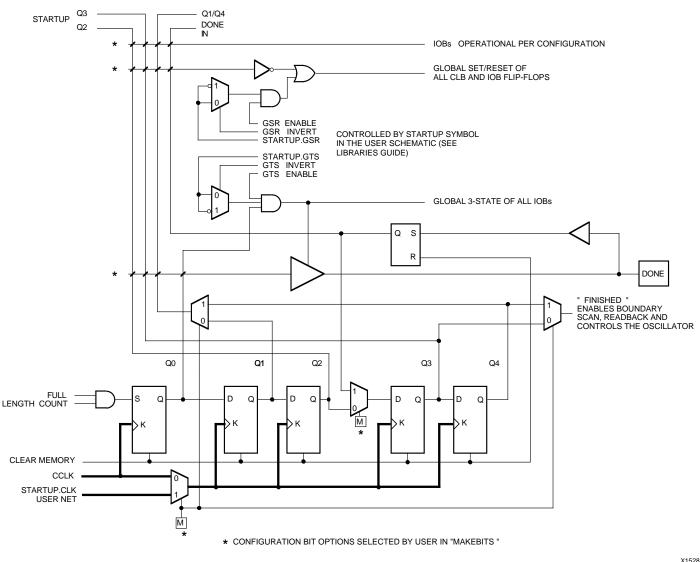
# Start-up from a User Clock (STARTUP.CLK)

When, instead of CCLK, a user-supplied start-up clock is selected, Q1 is used to bridge the unknown phase relationship between CCLK and the user clock. This arbitration causes an unavoidable one-cycle uncertainty in the timing of the rest of the start-up sequence.



Note: Thick lines are default option.





#### Figure 22. Start-up Logic

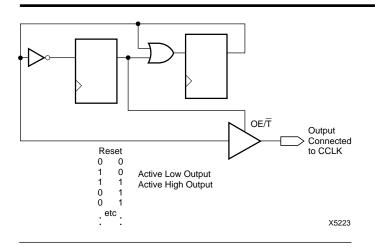
All Xilinx FPGAs of the XC2000, XC3000, XC4000 familiies use a compatible bitstream format and can, therefore, be connected in a daisy-chain in an arbitrary sequence. There is however one limitation. The lead device must belong to the highest family in the chain. If the chain contains XC4000 devices, the master cannot be an XC2000 or XC3000 device; if the daisy-chain contains XC3000 devices, the master cannot be an XC2000 device. The reason for this rule is shown in Figure 21 on the previous page. Since all devices in the chain store the same length count value and generate or receive one common sequence of CCLK pulses, they all recognize length-count match on the same CCLK edge, as indicated on the left edge of Figure 21. The master device will then drive additional CCLK pulses until it reaches its finish point F. The different families generate or require different numbers of additional CCLK pulses until they reach F.

Not reaching F means that the device does not really finish its configuration, although DONE may have gone High, the

outputs became active, and the internal RESET was released. The user has some control over the relative timing of these events and can, therefore, make sure that they occur early enough.

But, for XC4000, not reaching F means that READBACK cannot be initiated and most Boundary Scan instructions cannot be used. This limitation has been critized by designers who want to use an inexpensive lead device in peripheral mode and have the more precious I/O pins of the XC4000 devices all available for user I/O. Here is a solution for that case.

One CLB and one IOB in the lead XC3000 device are used to generate the additional CCLK pulse required by the XC4000 devices. When the lead device removes the internal RESET signal, the 2-bit shift register responds to its clock input and generates an active Low output signal for the duration of the subsequent clock period. An external connection between this output and CCLK thus creates



the extra CCLK pulse. This solution requires one CLB, one IOB and pin, and an internal oscillator with a frequency of up to 5 MHz as available clock source. Obviously, this XC3000 master device must be configured with late Internal Reset, which happens to be the default option.

#### Using Global Set/Reset and Global 3-State Nets

The global Set/Reset (STARTUP.GSR) net can be driven by the user at any time to re-initialize all CLBs and IOBs to the same state they had at the end of configuration. For CLBs that is the same state as the one driven by the individually programmable asynchronous Set/Reset inputs. The global 3-state net (STARTUP.GTS), whenever activated after configuration is completed, forces all LCA outputs to the high-impedance state, unless Boundary Scan is enabled and is executing an EXTEST instruction.

# Readback

The user can read back the content of configuration memory and the level of certain internal nodes without interfering with the normal operation of the device.

Readback reports not only the downloaded configuration bits, but can also include the present state of the device represented by the content of all used flip-flops and latches in CLBs and IOBs, as well as the content of function generators used as RAMs.

XC4000 Readback does not use any dedicated pins, but uses four internal nets (RDBK.TRIG, RDBK.DATA, RDBK.RIP and RDBK.CLK) that can be routed to any IOB.

After Readback has been initiated by a Low-to-High transition on RDBK.TRIG, the RDBK.RIP (Read In Progress) output goes High on the next rising edge of RDBK.CLK. Subsequent rising edges of this clock shift out Readback data on the RDBK.DATA net. Readback data does not include the preamble, but starts with five dummy bits (all High) followed by the Start bit (Low) of the first frame. The first two data bits of the first frame are always High.

Note that, in the XC4000 families, data is not inverted with respect to configuration the way it is in XC2000 and XC3000 families.

Each frame ends with four error check bits. They are read back as High. The last seven bits of the last frame are also read back as High. An additional Start bit (Low) and an 11-bit Cyclic Redundancy Check (CRC) signature follow, before RIP returns Low.

Readback options are: Read Capture, Read Abort, and Clock Select.

## **Read Capture**

When the Readback Capture option is selected, the readback data stream includes sampled values of CLB and IOB signals imbedded in the data stream. The rising edge of RDBK.TRIG located in the lower-left chip corner, captures, in latches, the inverted values of the four CLB outputs and the IOB output flip-flops and the input signals 11, 12. When the capture option is not selected, the values of the capture bits reflect the configuration data originally written to those memory locations. If the RAM capability of the CLBs is used, RAM data are available in readback, since they directly overwrite the F and G function-table configuration of the CLB.

## **Read Abort**

When the Readback Abort option is selected, a High-to-Low transition on RDBK.TRIG terminates the readback operation and prepares the logic to accept another trigger. After an aborted readback, additional clocks (up-to-one readback clock per configuration frame) may be required to re-initialize the control logic. The status of readback is indicated by the output control net (RDBK.RIP).

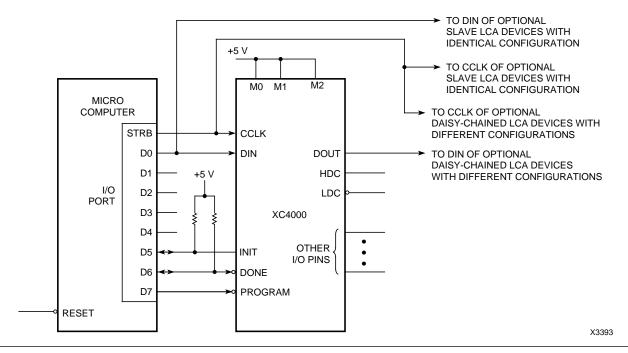
## **Clock Select**

Readback control and data are clocked on rising edges of RDBK.CLK located in the lower right chip corner. CCLK is an optional clock. If Readback must be inhibited for security reasons, the readback control nets are simply not connected.

## XChecker

The XChecker Universal Download/Readback Cable and Logic Probe uses the Readback feature for bitstream verification and for display of selected internal signals on the PC or workstation screen, effectively as a low-cost incircuit emulator.

# **Slave Serial Mode**



In Slave Serial mode, an external signal drives the CCLK input(s) of the LCA device(s). The serial configuration bitstream must be available at the DIN input of the lead LCA device a short set-up time before each rising CCLK edge. The lead LCA device then presents the preamble data (and all data that overflows the lead device) on its DOUT pin.

There is an internal delay of 0.5 CCLK periods, which means that DOUT changes on the falling CCLK edge, and the next LCA device in the daisy-chain accepts data on the subsequent rising CCLK edge.

#### How to Delay Configuration After Power-Up

There are two methods to delay configuration after powerup: Put a logic Low on the <u>PROGRAM</u> input, or pull the bidirectional <u>INIT</u> pin Low, using an open-collector (opendrain) driver. (See also Figure 20 on page 2-27.) A Low on the <u>PROGRAM</u> input is the more radical approach, and is recommended when the power-supply rise time is excessive or poorly defined. As long as <u>PROGRAM</u> is Low, the XC4000 device keeps clearing its configuration memory. When <u>PROGRAM</u> goes High, the configuration memory is cleared one more time, followed by the beginning of configuration, provided the INIT input is not externally held Low. Note that a Low on the <u>PROGRAM</u> input automatically forces a Low on the <u>INIT</u> output.

Using an open-collector or open-drain driver to hold <u>INIT</u> Low before the beginning of configuration, causes the LCA device to wait after having completed the configuration memory clear operation. When <u>INIT</u> is no longer held Low externally, the device determines its configuration mode by capturing its status inputs, and is ready to start the configuration process. A master device waits an additional max 250  $\mu$ s to make sure that all slaves in the potential daisy-chain have seen <u>INIT</u> being High.

#### CCLK INIT BYTE 0 BYTE BYTE 0 OUT BYTE 1 OUT 2 3 5 0 0 1 4 6 7 DOUT RDY/BUSY

# Synchronous Peripheral Mode Programming Switching Characteristics

X6096

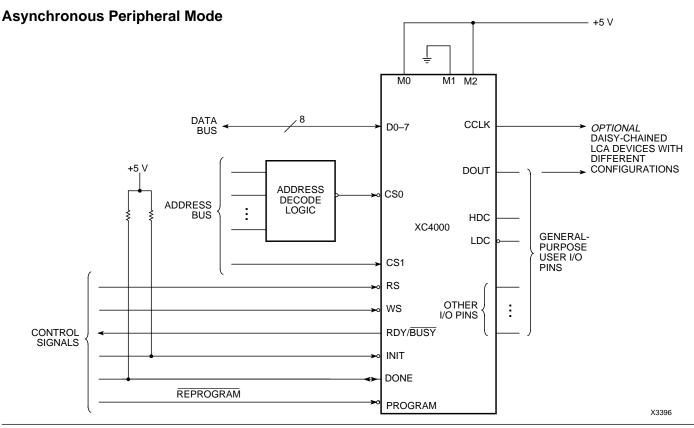
	Description	S	ymbol	Min	Мах	Units
CCLK	INIT (High) Setup time required	1	T <sub>IC</sub>	5		μs
	D0-D7 Setup time required	2	T <sub>DC</sub>	60		ns
	D0-D7 Hold time required	3	T <sub>CD</sub>	0		ns
	CCLK High time		Тссн	50		ns
	CCLK Low time		T <sub>CCL</sub>	60		ns
	CCLK Frequency		F <sub>cc</sub>		8	MHz

Notes: Peripheral Synchronous mode can be considered Slave Parallel mode. An external CCLK provides timing, clocking in the **first** data byte on the **second** rising edge of CCLK after <u>INIT</u> goes High. Subsequent data bytes are clocked in on every eighth consecutive rising edge of CCLK.

The RDY/<u>BUSY</u> line goes High for one CCLK period after data has been clocked in, although synchronous operation does not require such a response.

The pin name RDY/<u>BUSY</u> is a misnomer; in Synchronous Peripheral mode this is really an ACKNOWLEDGE signal.

Note that data starts to shift out serially on the DOUT pin 0.5 CLK periods after it was loaded in parallel. This obviously requires additional CCLK pulses after the last byte has been loaded.



# Write to LCA

Asynchronous Peripheral mode uses the trailing edge of the logic AND condition of the <u>CS0</u>, CS1 and <u>WS</u> inputs to accept byte-wide data from a microprocessor bus. In the lead LCA device, this data is loaded into a double-buffered UART-like parallel-to-serial converter and is serially shifted into the internal logic. The lead LCA device presents the preamble data (and all data that overflows the lead device) on the DOUT pin.

The RDY/<u>BUSY</u> output from the lead LCA device acts as a handshake signal to the microprocessor. RDY/<u>BUSY</u> goes Low when a byte has been received, and goes High again when the byte-wide input buffer has transferred its information into the shift register, and the buffer is ready to receive new data. The length of the <u>BUSY</u> signal depends on the activity in the UART. If the shift register had been empty when the new byte was received, the <u>BUSY</u> signal lasts for only two CCLK periods. If the shift register was still full when the new byte was received, the <u>BUSY</u> signal can be as long as nine CCLK periods.

Note that after the last byte has been entered, only seven of its bits are shifted out. CCLK remains High with DOUT equal to bit 6 (the next-to-last bit) of the last byte entered. The READY/<u>BUSY</u> handshake can be ignored if the delay from any one Write to the end of the next Write is guaranteed to be longer than 10 CCLK periods, i.e. longer than 20  $\mu$ s.

# Status Read

The logic AND condition of the <u>CS0</u>, CS1and <u>RS</u> inputs puts the device status on the Data bus.

- D7 = High indicates Ready
- D7 Low indicates Busy
- D0 through D6 go unconditionally High

It is mandatory that the whole start-up sequence be started and completed by one byte-wide input. Otherwise, the pins used as Write Strobe or Chip Enable might become active outputs and inteffere with the final byte transfer. If this transfer does not occur, the start-up sequence will not be completed all the way to the finish (point F in Figure 21 on page 2-29). At worst, the internal reset will not be released; at best, Readback and Boundary Scan will be inhibited. The length-count value, as generated by MAKEPROM, is supposed to ensure that these problems never occur.

Although RDY/<u>BUSY</u> is brought out as a separate signal, microprocessors can more easily read this information on one of the data lines. For this purpose, D7 represents the RDY/<u>BUSY</u> status when <u>RS</u> is Low, <u>WS</u> is High, and the two chip select lines are both active.

## How to Delay Configuration After Power-Up

There are two methods to delay configuration after powerup: Put a logic Low on the <u>PROGRAM</u> input, or pull the bidirectional <u>INIT</u> pin Low, using an open-collector (opendrain) driver. (See also Figure 20 on page 2-27). A Low on the <u>PROGRAM</u> input is the more radical approach, and is recommended when the power-supply rise time is excessive or poorly defined. As long as <u>PROGRAM</u> is Low, the XC4000 device keeps clearing its configuration memory. When <u>PROGRAM</u> goes High, the configuration memory is cleared one more time, followed by the beginning of configuration, provided the INIT input is not externally held Low. Note that a Low on the <u>PROGRAM</u> input automatically forces a Low on the <u>INIT</u> output.

Using an open-collector or open-drain driver to hold <u>INIT</u> Low before the beginning of configuration, causes the LCA device to wait after having completed the configuration memory clear operation. When <u>INIT</u> is no longer held Low externally, the device determines its configuration mode by capturing its status inputs, and is ready to start the configuration process. A master device waits an additional max 250  $\mu$ s to make sure that all slaves in the potential daisy-chain have seen <u>INIT</u> being High.

Asynch	ronous Peripheral Mode Programming Switching Characteristics Write to LCA Read Status
WS/CS0	
RS, CS1	$\begin{array}{c} & & \\ \hline \end{array} \\ \hline \\ \\ \hline \end{array} \\ \hline \end{array} \\ \hline \end{array} \\ \hline \end{array} \\ \hline \\ \\ \hline \end{array} \\ \hline \end{array} \\ \\ \\ \\$
D0-D7	$\begin{array}{c} & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & & \\ & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & \\ & & & & & & \\ & & & & & & \\ & & & & & & \\ & & &$
CCLK	
RDY/BUSY	$\xrightarrow{T_{WTRB}(4) \rightarrow  } \underbrace{\leftarrow}_{6} T_{BUSY}$
DOUT	Previous Byte D6         D7         D0         D1         D2

						X6097
	Description	S	ymbol	Min	Max	Units
Write	Effective Write time required $(\underline{CS0}, \underline{WS} = Low, \underline{RS}, CS1 = High)$	1	T <sub>CA</sub>	100		ns
	DIN Setup time required DIN Hold time required	2 3	T <sub>DC</sub> T <sub>CD</sub>	60 0		ns ns
	RDY/ <u>BUSY</u> delay after end of Write or Read RDY/BUSY active after begining of	4	T <sub>WTRB</sub>		60	ns
	Read				60	ns
RDY	Earliest next <u>WS</u> after end of <u>BUSY</u>	5	T <sub>RBWT</sub>	0		ns
	BUSY Low output (Note 4)	6	T <sub>BUSY</sub>	2	9	CCLK Periods

Notes: 1. Configuration must be delayed until the <u>INIT</u> of all LCA devices is High.

- 2. Time from end of <u>WS</u> to CCLK cycle for the new byte of data depends on completion of previous byte processing and the phase of the internal timing generator for CCLK.
  - 3. CCLK and DOUT timing is tested in slave mode.
  - 4. TBUSY indicates that the double-buffered parallel-to-serial converter is not yet ready to receive new data. The shortest TBUSY occurs when a byte is loaded into an empty parallel-to-serial converter. The longest TBUSY occurs when a new word is loaded into the input register before the second-level buffer has started shifting out data.

This timing diagram shows very relaxed requirements:

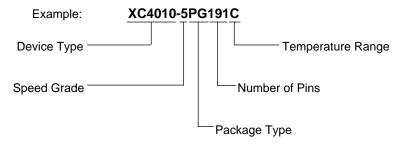
Data need not be held beyond the rising edge of <u>WS</u>. <u>BUSY</u> will go active within 60 ns after the end of <u>WS</u>. <u>WS</u> may be asserted immediately after the end of <u>BUSY</u>. For a detailed description of the device architecture, see page 2-9 through 2-31.

For a detailed description of the configuration modes and their timing, see pages 2-32 through 2-55.

For detailed lists of package pinouts, see pages 2-57 through 2-67, 2-70, 2-81 through 2-85, and 2-100 through 2-101.

For package physical dimensions and thermal data, see Section 4.

# **Ordering Information**



# **Component Availability**

PINS	84		100		120	144	156	160	164	191	196	20	)8	223	225	24	40	299	304
				TOP					TOP		TOP								
TYPE	PLAST. PLCC	PLAST. PQFP	PLAST. VQFP	BRAZED CQFP	CERAM. PGA	PLAST. TQFP	CERAM PGA	PLAST. PQFP	BRAZED CQFP	CERAM. PGA	BRAZED CQFP	PLAST. PQFP	METAL PQFP	CERAM. PGA	PLAST. BGA	PLAST. PQFP	METAL PQFP	METAL PQFP	н
	PLCC	PQFP	VQFP	CQFP	PGA	IQFP	PGA	PQFP	CQFP	PGA	CQFP	PQFP	PQFP	PGA	BGA	PQFP	PQFP	PQFP	QUAD
CODE	PC84		VQ100	CB100		TQ144	PG156	PQ160	CB164	PG191	CB196	PQ208	MQ208	PG223	BG225	PQ240	MQ240	PG299	HQ304
-6	CI	СІ			CI														
XC4003 -5 -4	C C	C C			C C														
-10		-			-		MВ		MВ										
XC4005 -6		CI					CIMB	CI	MВ			CI							
-5	CI C	C I C					C I C	C I C				CI C							
-6	СІ	-					CI	CI				CI							
XC4006 -5							CI	CI				CI							
-4	C CI						С	C CI		CI		C CI	CI						
XC4008 -5	CI							CI		CI		CI	CI						
-4	С							С		С		С	С						
-10 XC4010 -6								CI		M B CIM B	M B M B	CI	CI		CI				
-5	CI							CI		CI		CI	CI		CI				
-4	С							С		С		С	С		С				
-6 XC4010D -5	CI CI							C I C I				CI CI			C I C I				
-4								c				c			c				
-6								CI				CI	CI	CI (M B)	CI	CI	CI		
XC4013 -5 -4								C I C				CI C	C I C	C I C	C I C	C I CI	C I C		
-6								CI				CI	0	0	CI	CI	0		
XC4013D -5								CI				CI			CI	CI			
-4								С				C (C I)		(C I)	С	C (C I)		(C I)	
XC4020 -5												(CI) (CI)		(CI) (CI)		(C I) (C I)		(CI) (CI)	
-4												(C)		(C)		(C)		(C)	
-6 XC4025 -5														C C			CI	CI	CI
-4														C			01	CT	CT
-6	CI	CI	CI		CI														
XC4002A -5 -4		С	С		С														
-10				MВ	MВ														
XC4003A -6		CI	CI	MВ	СІМВ														
-5	C C	C C	C C		C C														
-6	CI	C	C		CI	CI		CI											
XC4004A -5	С				С	С		С											
-4	CI					CI	CI	CI				01							
-6 XC4005A -5	CI					CI	CI	CI				CI CI							
-4	С					С	С	С		<u><u> </u></u>		С	_						
XC4003H -6 -5										CI C		CI C							
XC4005H -6														CI		CI	CI		
-5		mercial	= 0° to +	85° C	I = Indus	strial = -4	0° to +10	0° C	M = Mil	Temp = -	55° to +1	25° C		С		С	С		
			3C Class				cate futu					_, ,							