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Applications of "Embedded - Microcontrollers"

Details	
Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	25MHz
Connectivity	UART/USART
Peripherals	Brown-out Detect/Reset, LVD, POR, PWM, WDT
Number of I/O	16
Program Memory Size	4KB (2K x 16)
Program Memory Type	FLASH
EEPROM Size	128 x 8
RAM Size	256 x 8
Voltage - Supply (Vcc/Vdd)	4.2V ~ 5.5V
Data Converters	A/D 4x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 125°C (TA)
Mounting Type	Surface Mount
Package / Case	20-SSOP (0.209", 5.30mm Width)
Supplier Device Package	20-SSOP
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic18f1230-e-ss

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NOTES:

4.3 Sleep Mode

The power-managed Sleep mode in the PIC18F1230/1330 devices is identical to the legacy Sleep mode offered in all other PIC devices. It is entered by clearing the IDLEN bit (the default state on device Reset) and executing the SLEEP instruction. This shuts down the selected oscillator (Figure 4-5). All clock source status bits are cleared.

Entering the Sleep mode from any other mode does not require a clock switch. This is because no clocks are needed once the controller has entered Sleep. If the WDT is selected, the INTRC source will continue to operate. If the Timer1 oscillator is enabled, it will also continue to run.

When a wake event occurs in Sleep mode (by interrupt, Reset or WDT time-out), the device will not be clocked until the clock source selected by the SCS1:SCS0 bits becomes ready (see Figure 4-6), or it will be clocked from the internal oscillator block if either the Two-Speed Start-up or the Fail-Safe Clock Monitor are enabled (see **Section 20.0 "Special Features of the CPU"**). In either case, the OSTS bit is set when the primary clock is providing the device clocks. The IDLEN and SCS bits are not affected by the wake-up.

4.4 Idle Modes

The Idle modes allow the controller's CPU to be selectively shut down while the peripherals continue to operate. Selecting a particular Idle mode allows users to further manage power consumption.

If the IDLEN bit is set to a '1' when a SLEEP instruction is executed, the peripherals will be clocked from the clock source selected using the SCS1:SCS0 bits; however, the CPU will not be clocked. The clock source status bits are not affected. Setting IDLEN and executing a SLEEP instruction provides a quick method of switching from a given Run mode to its corresponding Idle mode.

If the WDT is selected, the INTRC source will continue to operate. If the Timer1 oscillator is enabled, it will also continue to run.

Since the CPU is not executing instructions, the only exits from any of the Idle modes are by interrupt, WDT time-out or a Reset. When a wake event occurs, CPU execution is delayed by an interval of T_{CSD} (parameter 38, Table 23-10) while it becomes ready to execute code. When the CPU begins executing code, it resumes with the same clock source for the current Idle mode. For example, when waking from RC_IDLE mode, the internal oscillator block will clock the CPU and peripherals (in other words, RC_RUN mode). The IDLEN and SCS bits are not affected by the wake-up.

While in any Idle mode or the Sleep mode, a WDT time-out will result in a WDT wake-up to the Run mode currently specified by the SCS1:SCS0 bits.

FIGURE 4-5: TRANSITION TIMING FOR ENTRY TO SLEEP MODE

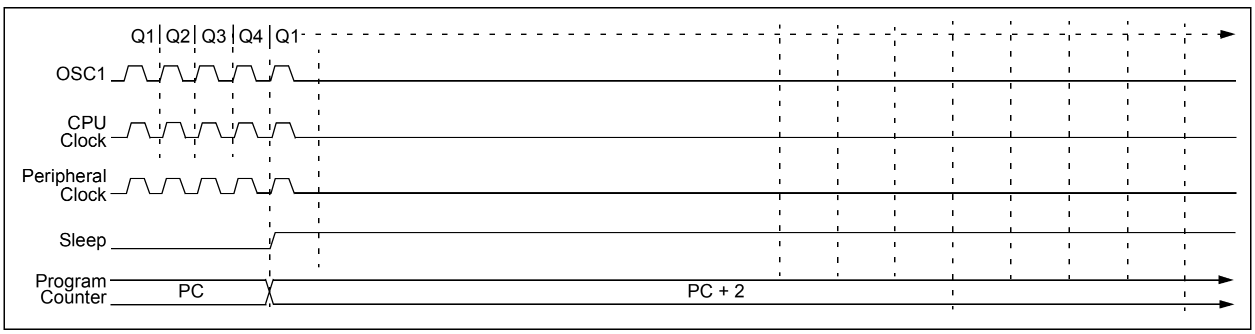
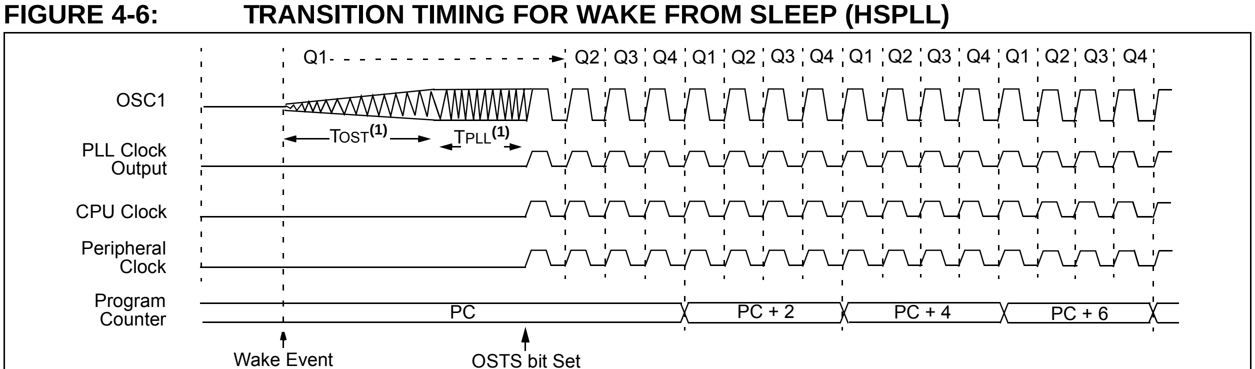


FIGURE 4-6: TRANSITION TIMING FOR WAKE FROM SLEEP (HSPLL)



Note 1: T_{OST} = 1024 T_{OSC}; T_{PLL} = 2 ms (approx). These intervals are not shown to scale.

4.4.1 PRI_IDLE MODE

This mode is unique among the three low-power Idle modes, in that it does not disable the primary device clock. For timing sensitive applications, this allows for the fastest resumption of device operation with its more accurate primary clock source, since the clock source does not have to "warm-up" or transition from another oscillator.

PRI_IDLE mode is entered from PRI_RUN mode by setting the IDLEN bit and executing a SLEEP instruction. If the device is in another Run mode, set IDLEN first, then clear the SCS bits and execute SLEEP. Although the CPU is disabled, the peripherals continue to be clocked from the primary clock source specified by the FOSC3:FOSC0 Configuration bits. The OSTS bit remains set (see Figure 4-7).

When a wake event occurs, the CPU is clocked from the primary clock source. A delay of interval TCSD is required between the wake event and when code execution starts. This is required to allow the CPU to become ready to execute instructions. After the wake-up, the OSTS bit remains set. The IDLEN and SCS bits are not affected by the wake-up (see Figure 4-8).

4.4.2 SEC_IDLE MODE

In SEC_IDLE mode, the CPU is disabled but the peripherals continue to be clocked from the Timer1 oscillator. This mode is entered from SEC_RUN by setting the IDLEN bit and executing a SLEEP instruction. If the device is in another Run mode, set the IDLEN bit first, then set the SCS1:SCS0 bits to '01' and execute SLEEP. When the clock source is switched to the Timer1 oscillator, the primary oscillator is shut down, the OSTS bit is cleared and the T1RUN bit is set.

When a wake event occurs, the peripherals continue to be clocked from the Timer1 oscillator. After an interval of TCSD following the wake event, the CPU begins executing code being clocked by the Timer1 oscillator. The IDLEN and SCS bits are not affected by the wake-up; the Timer1 oscillator continues to run (see Figure 4-8).

Note:	The Timer1 oscillator should already be running prior to entering SEC_IDLE mode. If the T1OSCEN bit is not set when the SLEEP instruction is executed, the SLEEP instruction will be ignored and entry to SEC_IDLE mode will not occur. If the Timer1 oscillator is enabled but not yet running, peripheral clocks will be delayed until the oscillator has started. In such situations, initial oscillator operation is far from stable and unpredictable operation may result.

FIGURE 4-7: TRANSITION TIMING FOR ENTRY TO IDLE MODE

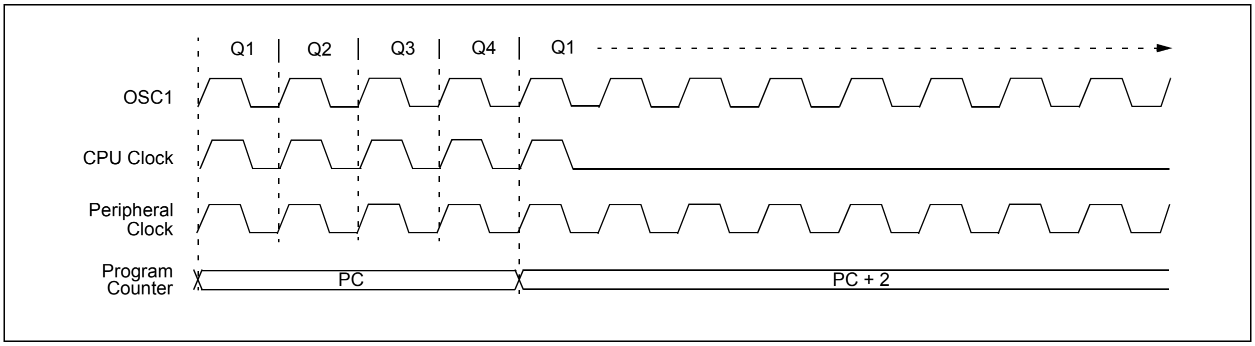
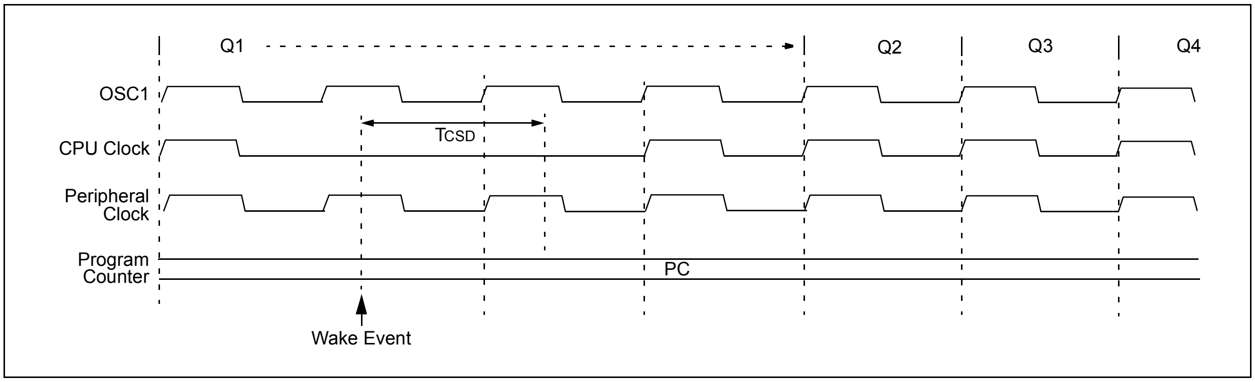


FIGURE 4-8: TRANSITION TIMING FOR WAKE FROM IDLE TO RUN MODE



6.0 MEMORY ORGANIZATION

There are three types of memory in PIC18 Enhanced microcontroller devices:

- Program Memory
- Data RAM
- Data EEPROM

As Harvard architecture devices, the data and program memories use separate busses; this allows for concurrent access of the two memory spaces. The data EEPROM, for practical purposes, can be regarded as a peripheral device, since it is addressed and accessed through a set of control registers.

Additional detailed information on the operation of the Flash program memory is provided in **Section 7.0 "Flash Program Memory"**. Data EEPROM is discussed separately in **Section 8.0 "Data EEPROM Memory"**.

6.1 Program Memory Organization

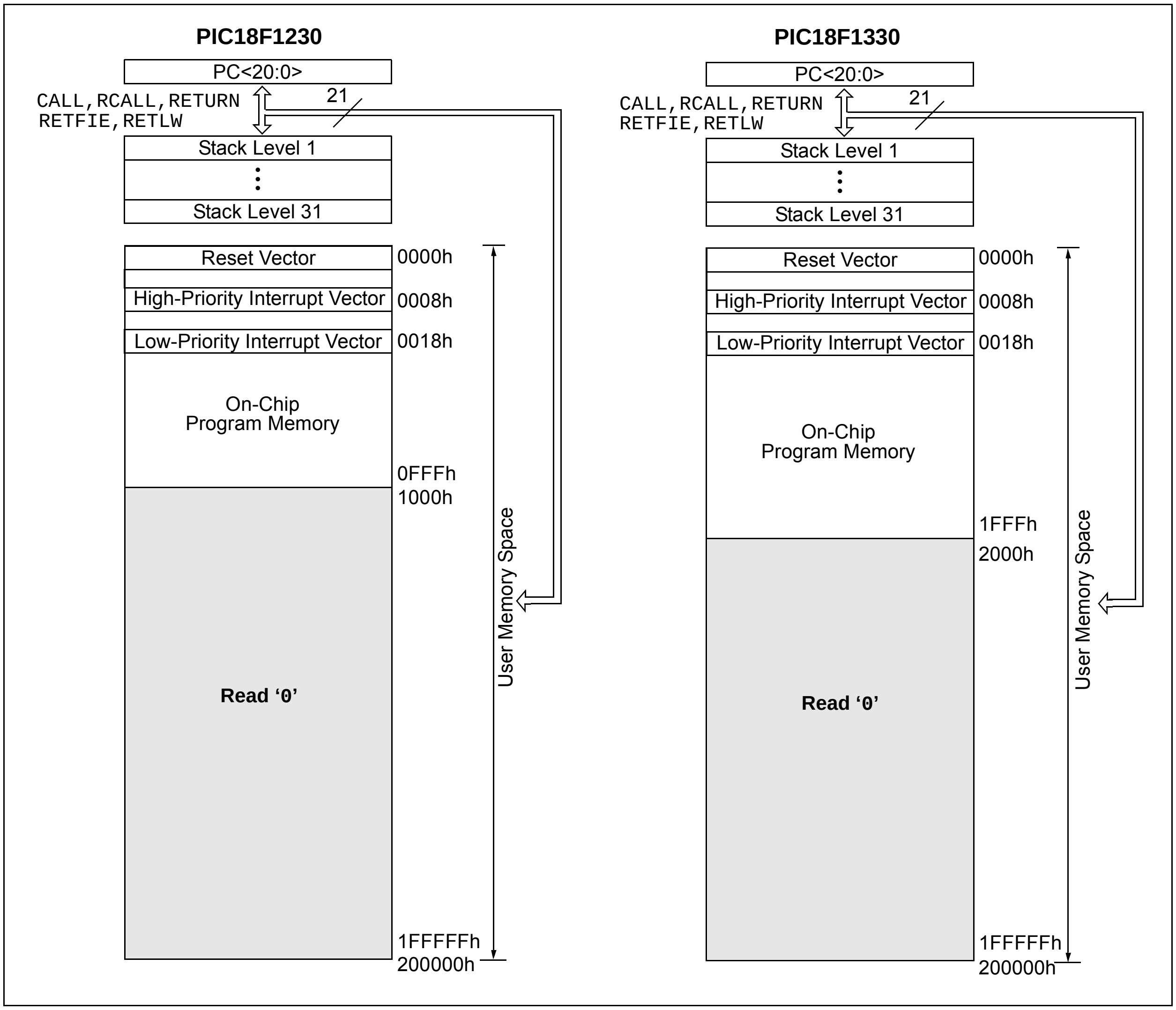
PIC18 microcontrollers implement a 21-bit program counter, which is capable of addressing a 2-Mbyte program memory space. Accessing a location between the upper boundary of the physically implemented memory and the 2-Mbyte address will return all '0's (a NOP instruction).

The PIC18F1230 has 4 Kbytes of Flash memory and can store up to 2,048 single-word instructions. The PIC18F1330 has 8 Kbytes of Flash memory and can store up to 4,096 single-word instructions.

PIC18 devices have two interrupt vectors. The Reset vector address is at 0000h and the interrupt vector addresses are at 0008h and 0018h.

The program memory maps for PIC18F1230 and PIC18F1330 devices are shown in Figure 6-1.

FIGURE 6-1: PROGRAM MEMORY MAP AND STACK FOR PIC18F1230/1330 DEVICES



7.0 FLASH PROGRAM MEMORY

The Flash program memory is readable, writable and erasable during normal operation over the entire VDD range.

A read from program memory is executed on one byte at a time. A write to program memory is executed on blocks of 8 bytes at a time. Program memory is erased in blocks of 64 bytes at a time. A bulk erase operation may not be issued from user code.

Writing or erasing program memory will cease instruction fetches until the operation is complete. The program memory cannot be accessed during the write or erase, therefore, code cannot execute. An internal programming timer terminates program memory writes and erases.

A value written to program memory does not need to be a valid instruction. Executing a program memory location that forms an invalid instruction results in a NOP.

7.1 Table Reads and Table Writes

In order to read and write program memory, there are two operations that allow the processor to move bytes between the program memory space and the data RAM:

- Table Read (TBLRD)
- Table Write (TBLWT)

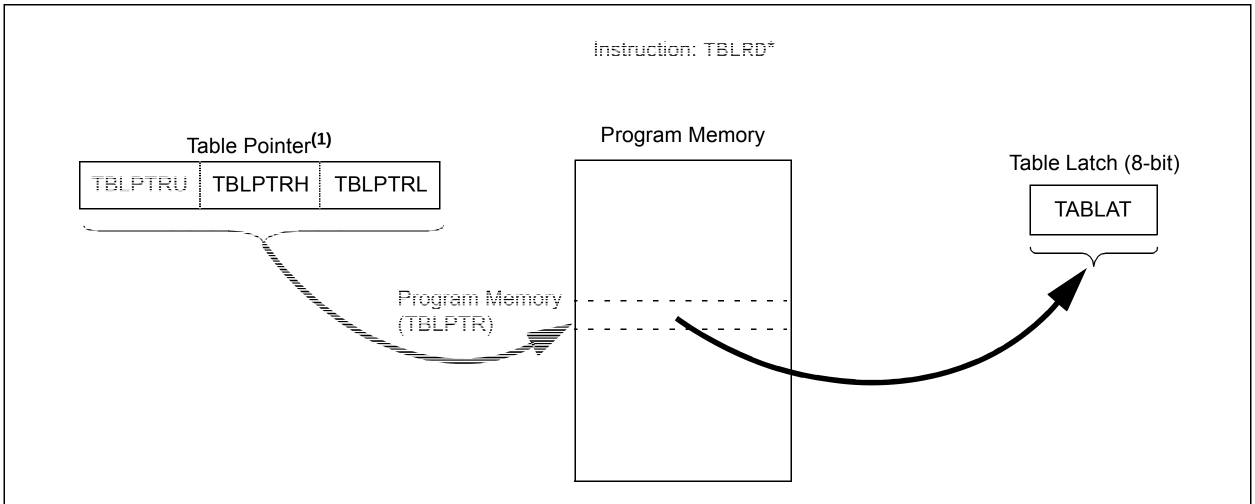
The program memory space is 16 bits wide, while the data RAM space is 8 bits wide. Table reads and table writes move data between these two memory spaces through an 8-bit register (TABLAT).

Table read operations retrieve data from program memory and place it into the data RAM space. Figure 7-1 shows the operation of a table read with program memory and data RAM.

Table write operations store data from the data memory space into holding registers in program memory. The procedure to write the contents of the holding registers into program memory is detailed in **Section 7.5 "Writing to Flash Program Memory"**. Figure 7-2 shows the operation of a table write with program memory and data RAM.

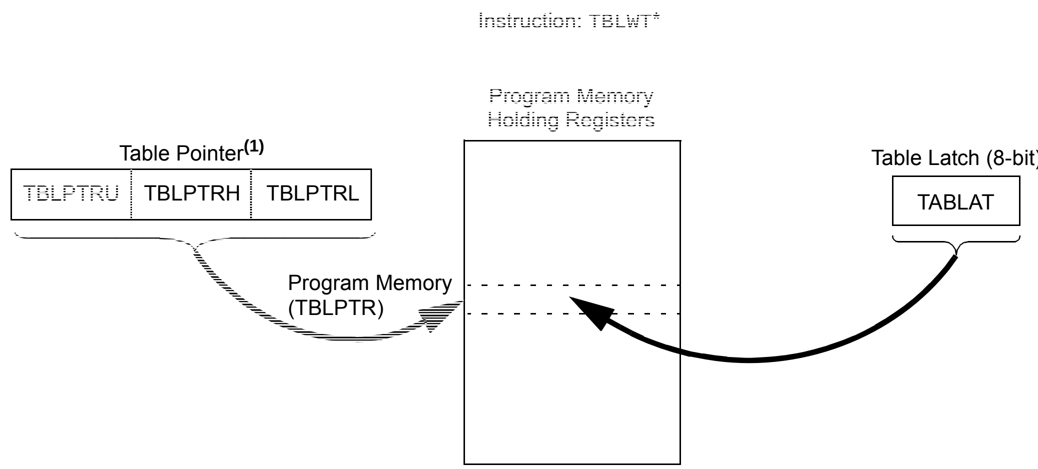
Table operations work with byte entities. A table block containing data, rather than program instructions, is not required to be word-aligned. Therefore, a table block can start and end at any byte address. If a table write is being used to write executable code into program memory, program instructions will need to be word-aligned.

FIGURE 7-1: TABLE READ OPERATION



Note 1: Table Pointer register points to a byte in program memory.

FIGURE 7-2: TABLE WRITE OPERATION



Note 1: Table Pointer actually points to one of 8 holding registers, the address of which is determined by TBLPTRL<5:0>. The process for physically writing data to the program memory array is discussed in **Section 7.5 "Writing to Flash Program Memory"**.

7.2 Control Registers

Several control registers are used in conjunction with the TBLRD and TBLWT instructions. These include the:

- EECON1 register
- EECON2 register
- TABLAT register
- TBLPTR registers

7.2.1 EECON1 AND EECON2 REGISTERS

The EECON1 register (Register 7-1) is the control register for memory accesses. The EECON2 register is not a physical register; it is used exclusively in the memory write and erase sequences. Reading EECON2 will read all '0's.

The EEPGD control bit determines if the access will be a program or data EEPROM memory access. When clear, any subsequent operations will operate on the data EEPROM memory. When set, any subsequent operations will operate on the program memory.

The CFGS control bit determines if the access will be to the Configuration/Calibration registers or to program memory/data EEPROM memory. When set, subsequent operations will operate on Configuration registers regardless of EEPGD (see **Section 20.0 "Special Features of the CPU"**). When clear, memory selection access is determined by EEPGD.

The FREE bit, when set, will allow a program memory erase operation. When FREE is set, the erase operation is initiated on the next WR command. When FREE is clear, only writes are enabled.

The WREN bit, when set, will allow a write operation. On power-up, the WREN bit is clear. The WRERR bit is set in hardware when the WR bit is set and cleared when the internal programming timer expires and the write operation is complete.

Note:	During normal operation, the WRERR may read as '1'. This can indicate that a write operation was prematurely terminated by a Reset, or a write operation was attempted improperly.

The WR control bit initiates write operations. The bit cannot be cleared, only set, in software; it is cleared in hardware at the completion of the write operation.

Note:	The EEIF interrupt flag bit (PIR2<4>) is set when the write is complete. It must be cleared in software.

EXAMPLE 7-3: WRITING TO FLASH PROGRAM MEMORY (CONTINUED)

```
PROGRAM_MEMORY
               BSF     EECON1, EEPGD            ; point to Flash program memory
               BCF     EECON1, CFGS             ; access Flash program memory
               BSF     EECON1, WREN             ; enable write to memory
               BCF     INTCON, GIE              ; disable interrupts
               MOVLW   55h
    Required   MOVWF   EECON2                   ; write 55h
    Sequence   MOVLW   0AAh
               MOVWF   EECON2                   ; write 0AAh
               BSF     EECON1, WR               ; start program (CPU stall)
               BSF     INTCON, GIE              ; re-enable interrupts
               BCF     EECON1, WREN             ; disable write to memory
```

7.5.2 WRITE VERIFY

Depending on the application, good programming practice may dictate that the value written to the memory should be verified against the original value. This should be used in applications where excessive writes can stress bits near the specification limit.

7.5.3 UNEXPECTED TERMINATION OF WRITE OPERATION

If a write is terminated by an unplanned event, such as loss of power or an unexpected Reset, the memory location just programmed should be verified and reprogrammed, if needed. If the write operation is interrupted by a $\overline{MCLR}$ Reset or a WDT Time-out Reset during normal operation, the user can check the WRERR bit and rewrite the location(s) as needed.

7.5.4 PROTECTION AGAINST SPURIOUS WRITES

To protect against spurious writes to Flash program memory, the write initiate sequence must also be followed. See **Section 20.0 “Special Features of the CPU”** for more detail.

7.6 Flash Program Operation During Code Protection

See **Section 20.5 “Program Verification and Code Protection”** for details on code protection of Flash program memory.

TABLE 7-2: REGISTERS ASSOCIATED WITH PROGRAM FLASH MEMORY

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Reset Values on Page:
TBLPTRU	—	—	bit 21	Program Memory Table Pointer Upper Byte (TBLPTR<20:16>)					47
TBPLTRH	Program Memory Table Pointer High Byte (TBLPTR<15:8>)								47
TBLPTRL	Program Memory Table Pointer Low Byte (TBLPTR<7:0>)								47
TABLAT	Program Memory Table Latch								47
INTCON	GIE/GIEH	PEIE/GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF	47
EECON2	EEPROM Control Register 2 (not a physical register)								49
EECON1	EEPGD	CFGS	—	FREE	WRERR	WREN	WR	RD	49
IPR2	OSCFIP	—	—	EEIP	—	LVDIP	—	—	49
PIR2	OSCFIF	—	—	EEIF	—	LVDIF	—	—	49
PIE2	OSCFIE	—	—	EEIE	—	LVDIE	—	—	49

Legend: — = unimplemented, read as ‘0’. Shaded cells are not used during Flash/EEPROM access.

10.0 I/O PORTS

Depending on the device selected and features enabled, there are up to five ports available. Some pins of the I/O ports are multiplexed with an alternate function from the peripheral features on the device. In general, when a peripheral is enabled, that pin may not be used as a general purpose I/O pin.

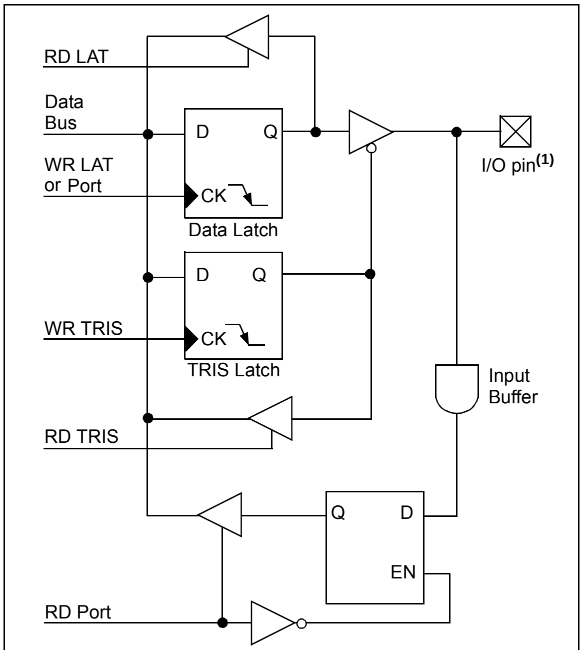
Each port has three registers for its operation. These registers are:

- TRIS register (Data Direction register)
- PORT register (reads the levels on the pins of the device)
- LAT register (Output Latch register)

The Output Latch (LAT register) is useful for read-modify-write operations on the value that the I/O pins are driving.

A simplified model of a generic I/O port, without the interfaces to other peripherals, is shown in Figure 10-1.

FIGURE 10-1: GENERIC I/O PORT OPERATION



Note 1: I/O pins have diode protection to VDD and VSS.

10.1 PORTA, TRISA and LATA Registers

PORTA is an 8-bit wide, bidirectional port. The corresponding Data Direction register is TRISA. Setting a TRISA bit (= `1`) will make the corresponding PORTA pin an input (i.e., put the corresponding output driver in a high-impedance mode). Clearing a TRISA bit (= `0`) will make the corresponding PORTA pin an output (i.e., put the contents of the output latch on the selected pin).

Reading the PORTA register reads the status of the pins, whereas writing to it, will write to the port latch.

The Output Latch (LATA) register is also memory mapped. Read-modify-write operations on the LATA register read and write the latched output value for PORTA.

Pins RA6 and RA7 are multiplexed with the main oscillator pins; they are enabled as oscillator or I/O pins by the selection of the main oscillator in the Configuration register (see **Section 20.1 “Configuration Bits”** for details). When they are not used as port pins, RA6 and RA7 and their associated TRIS and LAT bits are read as ‘`0`’.

The RA0 pin is multiplexed with one of the analog inputs, one of the external interrupt inputs, one of the interrupt-on-change inputs and one of the analog comparator inputs to become RA0/AN0/INT0/KBI0/CMP0 pin.

The RA1 pin is multiplexed with one of the analog inputs, one of the external interrupt inputs and one of the interrupt-on-change inputs to become RA1/AN1/INT1/KBI1 pin.

Pins RA2 and RA3 are multiplexed with the Enhanced USART transmission and reception input (see **Section 20.1 “Configuration Bits”** for details).

The RA4 pin is multiplexed with the Timer0 module clock input, one of the analog inputs and the analog VREF+ input to become the RA4/T0CKI/AN2/VREF+ pin.

The Fault detect input for PWM $\overline{\text{FLTA}}$ is multiplexed with pins RA5 and RA7. Its placement is decided by clearing or setting the FLTAMX bit of Configuration Register 3H.

> **Note:** On a Power-on Reset, RA0, RA1, RA4 and RA5 are configured as analog inputs and read as ‘`0`’. RA2 and RA3 are configured as digital inputs.

The TRISA register controls the direction of the PORTA pins, even when they are being used as analog inputs. The user must ensure the bits in the TRISA register are maintained set when using them as analog inputs.

EXAMPLE 10-1: INITIALIZING PORTA

```
CLRF    PORTA   ; Initialize PORTA by
                ; clearing output
                ; data latches
CLRF    LATA    ; Alternate method
                ; to clear output
                ; data latches
MOVLW   07h     ; Configure A/D
MOVWF   ADCON1  ; for digital inputs
MOVWF   07h     ; Configure comparators
MOVWF   CMCON   ; for digital input
MOVLW   0CFh    ; Value used to
                ; initialize data
                ; direction
MOVWF   TRISA   ; Set RA<7:6,3:0> as inputs
                ; RA<5:4> as outputs
```

11.5 RCON Register

The RCON register contains flag bits which are used to determine the cause of the last Reset or wake-up from Idle or Sleep modes. RCON also contains the IPEN bit which enables interrupt priorities.

The operation of the SBOREN bit and the Reset flag bits is discussed in more detail in **Section 5.1 “RCON Register”**.

REGISTER 11-13: RCON: RESET CONTROL REGISTER

R/W-0	R/W-1(1)	U-0	R/W-1	R-1	R-1	R/W-0(2)	R/W-0
IPEN	SBOREN	—	$\overline{\text{RI}}$	$\overline{\text{TO}}$	$\overline{\text{PD}}$	$\overline{\text{POR}}$	$\overline{\text{BOR}}$
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as ‘0’	
-n = Value at POR	‘1’ = Bit is set	‘0’ = Bit is cleared	x = Bit is unknown

bit 7 **IPEN:** Interrupt Priority Enable bit
1 = Enable priority levels on interrupts
0 = Disable priority levels on interrupts (PIC16CXXX Compatibility mode)

bit 6 **SBOREN:** BOR Software Enable bit(1)
For details of bit operation, see Register 5-1.

bit 5 **Unimplemented:** Read as ‘0’

bit 4 **$\overline{\text{RI}}$:** RESET Instruction Flag bit
For details of bit operation, see Register 5-1.

bit 3 **$\overline{\text{TO}}$:** Watchdog Time-out Flag bit
For details of bit operation, see Register 5-1.

bit 2 **$\overline{\text{PD}}$:** Power-Down Detection Flag bit
For details of bit operation, see Register 5-1.

bit 1 **$\overline{\text{POR}}$:** Power-on Reset Status bit(2)
For details of bit operation, see Register 5-1.

bit 0 **$\overline{\text{BOR}}$:** Brown-out Reset Status bit
For details of bit operation, see Register 5-1.

Note 1: If SBOREN is enabled, its Reset state is ‘1’; otherwise, it is ‘0’. See Register 5-1 for additional information.
2: The actual Reset value of $\overline{\text{POR}}$ is determined by the type of device Reset. See Register 5-1 for additional information.

EXAMPLE 13-1: IMPLEMENTING A REAL-TIME CLOCK USING A TIMER1 INTERRUPT SERVICE

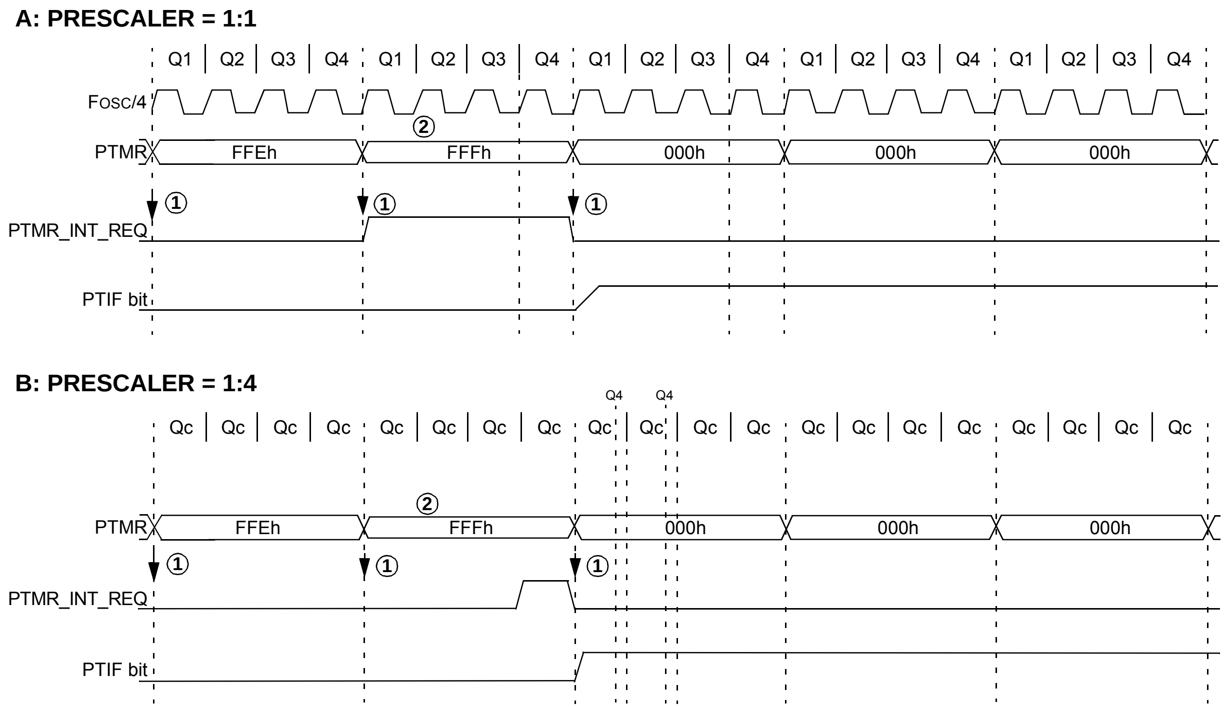
```
RTCinit
        MOVLW   0x80                ; Preload TMR1 register pair
        MOVWF   TMR1H               ; for 1 second overflow
        CLRF    TMR1L
        MOVLW   b'00001111'         ; Configure for external clock,
        MOVWF   T1CON               ; Asynchronous operation, external oscillator
        CLRF    secs                ; Initialize timekeeping registers
        CLRF    mins                ;
        MOVLW   .12
        MOVWF   hours
        BSF     PIE1, TMR1IE        ; Enable Timer1 interrupt
        RETURN
RTCisr
        BSF     TMR1H, 7            ; Preload for 1 sec overflow
        BCF     PIR1, TMR1IF        ; Clear interrupt flag
        INCF    secs, F             ; Increment seconds
        MOVLW   .59                 ; 60 seconds elapsed?
        CPFSGT  secs
        RETURN                      ; No, done
        CLRF    secs                ; Clear seconds
        INCF    mins, F             ; Increment minutes
        MOVLW   .59                 ; 60 minutes elapsed?
        CPFSGT  mins
        RETURN                      ; No, done
        CLRF    mins                ; clear minutes
        INCF    hours, F            ; Increment hours
        MOVLW   .23                 ; 24 hours elapsed?
        CPFSGT  hours
        RETURN                      ; No, done
        MOVLW   .01                 ; Reset hours to 1
        MOVWF   hours
        RETURN                      ; Done
```

TABLE 13-2: REGISTERS ASSOCIATED WITH TIMER1 AS A TIMER/COUNTER

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Reset Values on Page:
INTCON	GIE/GIEH	PEIE/GIEL	TMR0IE	INT0IE	RBIE	TMR0IF	INT0IF	RBIF	47
PIR1	—	ADIF	RCIF	TXIF	CMP2IF	CMP1IF	CMP0IF	TMR1IF	49
PIE1	—	ADIE	RCIE	TXIE	CMP2IE	CMP1IE	CMP0IE	TMR1IE	49
IPR1	—	ADIP	RCIP	TXIP	CMP2IP	CMP1IP	CMP0IP	TMR1IP	49
TMR1L	Timer1 Register Low Byte								48
TMR1H	Timer1 Register High Byte								48
T1CON	RD16	T1RUN	T1CKPS1	T1CKPS0	T1OSCEN	$\overline{\text{T1SYNC}}$	TMR1CS	TMR1ON	48

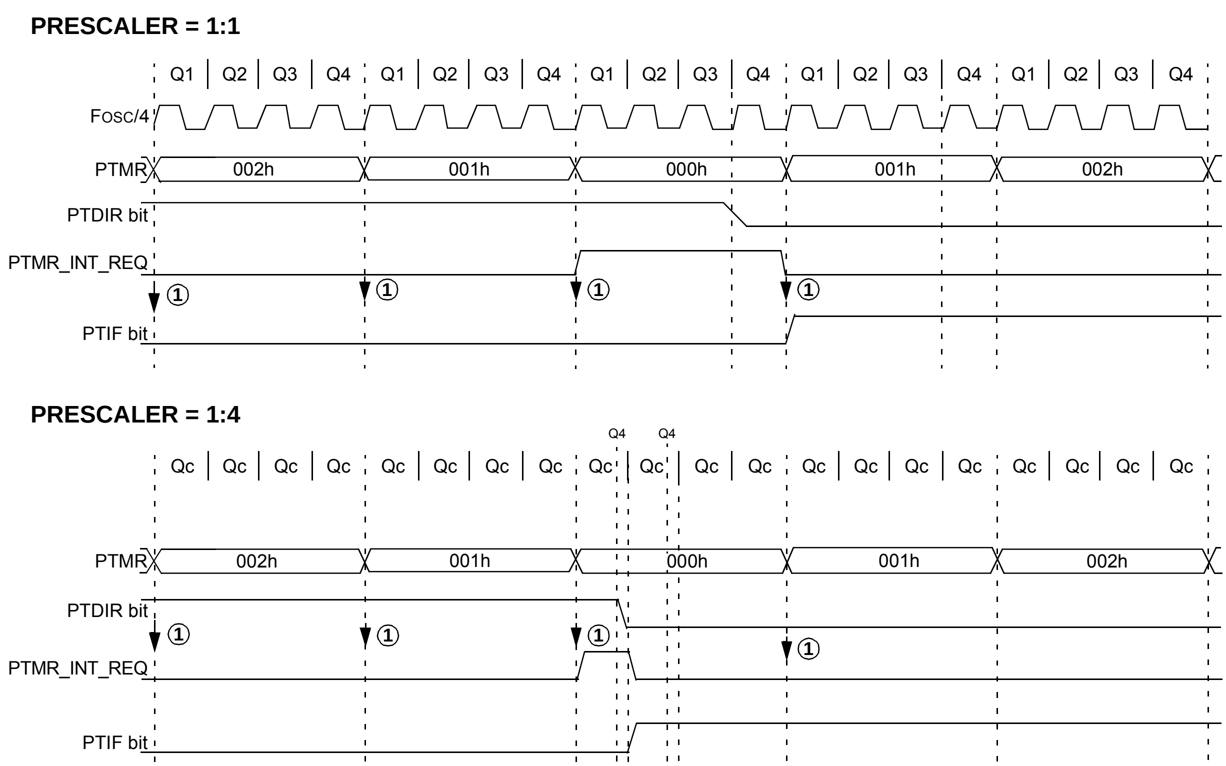
Legend: — = unimplemented, read as '0'. Shaded cells are not used by the Timer1 module.

FIGURE 14-6: PWM TIME BASE INTERRUPT TIMING, SINGLE-SHOT MODE



Note 1: Interrupt flag bit, PTIF, is sampled here (every Q1).
2: PWM Time Base Period register, PTPER, is loaded with the value FFFh for this example.

FIGURE 14-7: PWM TIME BASE INTERRUPTS, CONTINUOUS UP/DOWN COUNT MODE



Note 1: Interrupt flag bit, PTIF, is sampled here (every Q1).

15.2.5 BREAK CHARACTER SEQUENCE

The EUSART module has the capability of sending the special Break character sequences that are required by the LIN/J2602 bus standard. The Break character transmit consists of a Start bit, followed by twelve '0' bits and a Stop bit. The Frame Break character is sent whenever the SENDB and TXEN bits (TXSTA<3> and TXSTA<5>) are set while the Transmit Shift register is loaded with data. Note that the value of data written to TXREG will be ignored and all '0's will be transmitted.

The SENDB bit is automatically reset by hardware after the corresponding Stop bit is sent. This allows the user to preload the transmit FIFO with the next transmit byte following the Break character (typically, the Sync character in the LIN/J2602 specification).

Note that the data value written to the TXREG for the Break character is ignored. The write simply serves the purpose of initiating the proper sequence.

The TRMT bit indicates when the transmit operation is active or Idle, just as it does during normal transmission. See Figure 15-10 for the timing of the Break character sequence.

15.2.5.1 Break and Sync Transmit Sequence

The following sequence will send a message frame header made up of a Break, followed by an Auto-Baud Sync byte. This sequence is typical of a LIN/J2602 bus master.

1. Configure the EUSART for the desired mode.
2. Set the TXEN and SENDB bits to set up the Break character.
3. Load the TXREG with a dummy character to initiate transmission (the value is ignored).
4. Write '55h' to TXREG to load the Sync character into the transmit FIFO buffer.
5. After the Break has been sent, the SENDB bit is reset by hardware. The Sync character now transmits in the preconfigured mode.

When the TXREG becomes empty, as indicated by the TXIF, the next data byte can be written to TXREG.

15.2.6 RECEIVING A BREAK CHARACTER

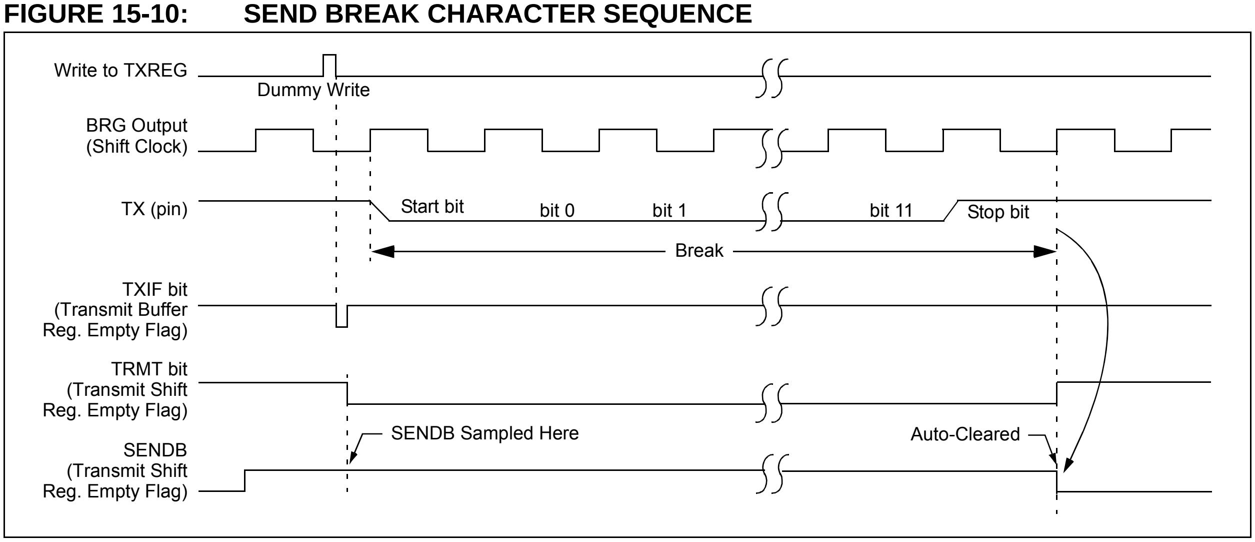
The Enhanced USART module can receive a Break character in two ways.

The first method forces configuration of the baud rate at a frequency of 9/13 the typical speed. This allows for the Stop bit transition to be at the correct sampling location (13 bits for Break versus Start bit and 8 data bits for typical data).

The second method uses the auto-wake-up feature described in **Section 15.2.4 "Auto-wake-up on Sync Break Character"**. By enabling this feature, the EUSART will sample the next two transitions on RX/DT, cause an RCIF interrupt and receive the next data byte followed by another interrupt.

Note that following a Break character, the user will typically want to enable the Auto-Baud Rate Detect feature. For both methods, the user can set the ABDEN bit once the TXIF interrupt is observed.

FIGURE 15-10: SEND BREAK CHARACTER SEQUENCE



15.3 EUSART Synchronous Master Mode

The Master mode indicates that the processor transmits the master clock on the CK line. The Synchronous Master mode is entered by setting the CSRC bit (TXSTA<7>). In this mode, the data is transmitted in a half-duplex manner (i.e., transmission and reception do not occur at the same time). When transmitting data, the reception is inhibited and vice versa. Synchronous mode is entered by setting bit SYNC (TXSTA<4>). In addition, enable bit, SPEN (RCSTA<7>), is set in order to configure the TX and RX pins to CK (clock) and DT (data) lines, respectively.

The Master mode indicates that the processor transmits the master clock on the CK line. Clock polarity is selected with the SCKP bit (BAUDCON<4>). Setting SCKP sets the Idle state on CK as high, while clearing the bit sets the Idle state as low.

15.3.1 EUSART SYNCHRONOUS MASTER TRANSMISSION

The EUSART transmitter block diagram is shown in Figure 15-3. The heart of the transmitter is the Transmit (Serial) Shift Register (TSR). The Shift register obtains its data from the Read/Write Transmit Buffer register, TXREG. The TXREG register is loaded with data in software. The TSR register is not loaded until the last bit has been transmitted from the previous load. As soon as the last bit is transmitted, the TSR is loaded with new data from the TXREG (if available).

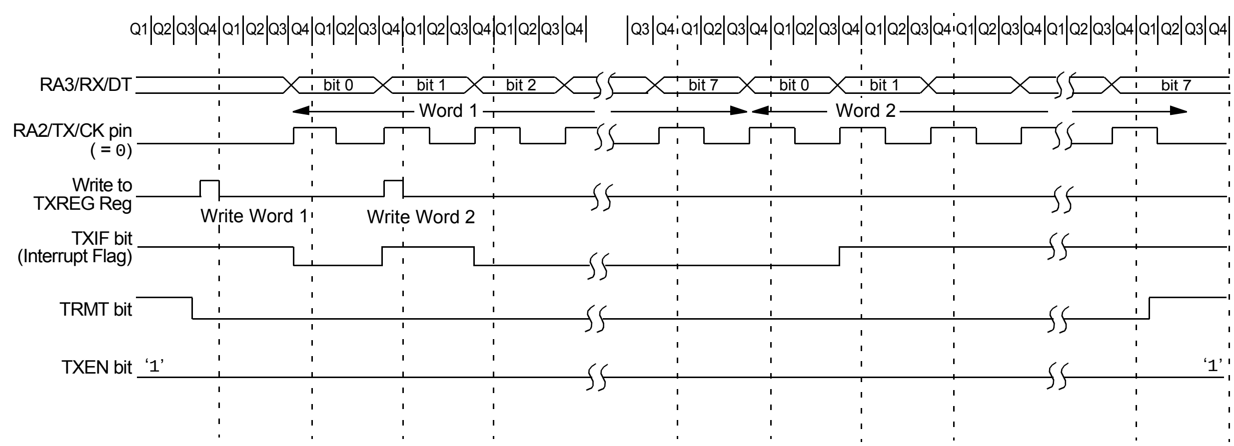
Once the TXREG register transfers the data to the TSR register (occurs in one TCY), the TXREG is empty and the TXIF flag bit (PIR1<4>) is set. The interrupt can be enabled or disabled by setting or clearing the interrupt enable bit, TXIE (PIE1<4>). TXIF is set regardless of the state of enable bit, TXIE; it cannot be cleared in software. It will reset only when new data is loaded into the TXREG register.

While flag bit TXIF indicates the status of the TXREG register, another bit, TRMT (TXSTA<1>), shows the status of the TSR register. TRMT is a read-only bit which is set when the TSR is empty. No interrupt logic is tied to this bit so the user has to poll this bit in order to determine if the TSR register is empty. The TSR is not mapped in data memory so it is not available to the user.

To set up a Synchronous Master Transmission:

1. Initialize the SPBRGH:SPBRG registers for the appropriate baud rate. Set or clear the BRG16 bit, as required, to achieve the desired baud rate.
2. Enable the synchronous master serial port by setting bits, SYNC, SPEN and CSRC.
3. If interrupts are desired, set enable bit, TXIE.
4. If 9-bit transmission is desired, set bit, TX9.
5. Enable the transmission by setting bit, TXEN.
6. If 9-bit transmission is selected, the ninth bit should be loaded in bit, TX9D.
7. Start transmission by loading data to the TXREG register.
8. If using interrupts, ensure that the GIE and PEIE bits in the INTCON register (INTCON<7:6>) are set.

FIGURE 15-11: SYNCHRONOUS TRANSMISSION



Note: Sync Master mode, SPBRG = 0, continuous transmission of two 8-bit words.

16.5 Operation in Power-Managed Modes

The selection of the automatic acquisition time and A/D conversion clock is determined in part by the clock source and frequency while in a power-managed mode.

If the A/D is expected to operate while the device is in a power-managed mode, the ACQT2:ACQT0 and ADCS2:ADCS0 bits in ADCON2 should be updated in accordance with the clock source to be used in that mode. After entering the mode, an A/D acquisition or conversion may be started. Once started, the device should continue to be clocked by the same clock source until the conversion has been completed.

If desired, the device may be placed into the corresponding Idle mode during the conversion. If the device clock frequency is less than 1 MHz, the A/D RC clock source should be selected.

Operation in Sleep mode requires the A/D FRC clock to be selected. If bits ACQT2:ACQT0 are set to '`000`' and a conversion is started, the conversion will be delayed one instruction cycle to allow execution of the `SLEEP` instruction and entry to Sleep mode. The IDLEN bit (OSCCON<7>) must have already been cleared prior to starting the conversion.

16.6 Configuring Analog Port Pins

The ADCON1 and TRISA registers configure the A/D port pins. The port pins needed as analog inputs must have their corresponding TRIS bits set (input). If the TRIS bit is cleared (output), the digital output level (V_{OH} or V_{OL}) will be converted.

The A/D operation is independent of the state of the CHS1:CHS0 bits and the TRIS bits.

Note 1:	When reading the PORT register, all pins configured as analog input channels will read as cleared (a low level). Pins configured as digital inputs will convert as analog inputs. Analog levels on a digitally configured input will be accurately converted.
2:	Analog levels on any pin defined as a digital input may cause the digital input buffer to consume current out of the device's specification limits.

REGISTER 20-11: CONFIG7L: CONFIGURATION REGISTER 7 LOW (BYTE ADDRESS 30000Ch)

U-0	U-0	U-0	U-0	U-0	U-0	R/C-1	R/C-1
—	—	—	—	—	—	EBTR1[(1)]	EBTR0[(1)]
bit 7							bit 0

Legend:		
R = Readable bit	C = Clearable bit	U = Unimplemented bit, read as ‘0’
-n = Value when device is unprogrammed		u = Unchanged from programmed state

bit 7-2 **Unimplemented:** Read as ‘0’

bit 1 **EBTR1:** Table Read Protection bit (Block 1 Code Memory Area)
1 = Block 1 is not protected from table reads executed in other blocks
0 = Block 1 is protected from table reads executed in other blocks

bit 0 **EBTR0:** Table Read Protection bit (Block 0 Code Memory Area)
1 = Block 0 is not protected from table reads executed in other blocks
0 = Block 0 is protected from table reads executed in other blocks

Note 1: It is recommended to enable the corresponding CPx bit to protect block from external read operations.

REGISTER 20-12: CONFIG7H: CONFIGURATION REGISTER 7 HIGH (BYTE ADDRESS 30000Dh)

U-0	R/C-1	U-0	U-0	U-0	U-0	U-0	U-0
—	EBTRB[(1)]	—	—	—	—	—	—
bit 7							bit 0

Legend:		
R = Readable bit	C = Clearable bit	U = Unimplemented bit, read as ‘0’
-n = Value when device is unprogrammed		u = Unchanged from programmed state

bit 7 **Unimplemented:** Read as ‘0’

bit 6 **EBTRB:** Table Read Protection bit (Boot Block Memory Area)
1 = Boot Block is not protected from table reads executed in other blocks
0 = Boot Block is protected from table reads executed in other blocks

bit 5-0 **Unimplemented:** Read as ‘0’

Note 1: It is recommended to enable the corresponding CPx bit to protect block from external read operations.

REGISTER 20-15: WDTCON: WATCHDOG TIMER CONTROL REGISTER

U-0	U-0	U-0	U-0	U-0	U-0	U-0	R/W-0
—	—	—	—	—	—	—	SWDTEN[(1)]
bit 7							bit 0

Legend:			
R = Readable bit	W = Writable bit	U = Unimplemented bit, read as ‘0’	
-n = Value at POR	‘1’ = Bit is set	‘0’ = Bit is cleared	x = Bit is unknown

bit 7-1 **Unimplemented:** Read as ‘0’

bit 0 **SWDTEN:** Software Controlled Watchdog Timer Enable bit[(1)]
1 = Watchdog Timer is on
0 = Watchdog Timer is off

Note 1: This bit has no effect if the Configuration bit, WDTEN, is enabled.

TABLE 20-2: SUMMARY OF WATCHDOG TIMER REGISTERS

Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Reset Values on Page:
RCON	IPEN	SBOREN[(1)]	—	$\overline{\text{RI}}$	$\overline{\text{TO}}$	$\overline{\text{PD}}$	$\overline{\text{POR}}$	$\overline{\text{BOR}}$	48
WDTCON	—	—	—	—	—	—	—	SWDTEN[(2)]	48

Legend: — = unimplemented, read as ‘0’. Shaded cells are not used by the Watchdog Timer.

Note 1: The SBOREN bit is only available when the BOREN1:BOREN0 Configuration bits = 01; otherwise, it is disabled and reads as ‘0’. See **Section 5.4 “Brown-out Reset (BOR)”**.

2: This bit has no effect if the Configuration bit, WDTEN, is enabled.

BTFSC — Bit Test File, Skip if Clear

Syntax: BTFSC f, b {,a}

Operands: $0 \le f \le 255$
$0 \le b \le 7$
$a \in [0,1]$

Operation: skip if (f<b>) = 0

Status Affected: None

Encoding:

1011	bbba	ffff	ffff

Description: If bit 'b' in register 'f' is '0', then the next instruction is skipped. If bit 'b' is '0', then the next instruction fetched during the current instruction execution is discarded and a NOP is executed instead, making this a two-cycle instruction.
If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank.
If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh).
See **Section 22.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.

Words: 1

Cycles: 1(2)
Note: 3 cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE    BTFSC   FLAG, 1, 0
FALSE   :
TRUE    :
```

Before Instruction
PC = address (HERE)
After Instruction
If FLAG<1> = 0;
PC = address (TRUE)
If FLAG<1> = 1;
PC = address (FALSE)

BTFSS — Bit Test File, Skip if Set

Syntax: BTFSS f, b {,a}

Operands: $0 \le f \le 255$
$0 \le b < 7$
$a \in [0,1]$

Operation: skip if (f<b>) = 1

Status Affected: None

Encoding:

1010	bbba	ffff	ffff

Description: If bit 'b' in register 'f' is '1', then the next instruction is skipped. If bit 'b' is '1', then the next instruction fetched during the current instruction execution is discarded and a NOP is executed instead, making this a two-cycle instruction.
If 'a' is '0', the Access Bank is selected. If 'a' is '1', the BSR is used to select the GPR bank.
If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever $f \le 95$ (5Fh).
See **Section 22.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.

Words: 1

Cycles: 1(2)
Note: 3 cycles if skip and followed by a 2-word instruction.

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process Data	No operation

If skip:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation

If skip and followed by 2-word instruction:

Q1	Q2	Q3	Q4
No operation	No operation	No operation	No operation
No operation	No operation	No operation	No operation

Example:

```
HERE    BTFSS   FLAG, 1, 0
FALSE   :
TRUE    :
```

Before Instruction
PC = address (HERE)
After Instruction
If FLAG<1> = 0;
PC = address (FALSE)
If FLAG<1> = 1;
PC = address (TRUE)

CALLW — Subroutine Call Using WREG

Syntax:	CALLW
Operands:	None
Operation:	(PC + 2) $\rightarrow$ TOS, (W) $\rightarrow$ PCL, (PCLATH) $\rightarrow$ PCH, (PCLATU) $\rightarrow$ PCU
Status Affected:	None
Encoding:	`0000` `0000` `0001` `0100`
Description	First, the return address (PC + 2) is pushed onto the return stack. Next, the contents of W are written to PCL; the existing value is discarded. Then, the contents of PCLATH and PCLATU are latched into PCH and PCU, respectively. The second cycle is executed as a `NOP` instruction while the new next instruction is fetched. Unlike `CALL`, there is no option to update W, STATUS or BSR.
Words:	1
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read WREG	PUSH PC to stack	No operation
No operation	No operation	No operation	No operation

Example: `HERE CALLW`

Before Instruction

PC	=	address (`HERE`)
PCLATH	=	10h
PCLATU	=	00h
W	=	06h

After Instruction

PC	=	001006h
TOS	=	address (`HERE + 2`)
PCLATH	=	10h
PCLATU	=	00h
W	=	06h

MOVSF — Move Indexed to f

Syntax:	MOVSF [z_s], f_d
Operands:	$0 \le z_s \le 127$ $0 \le f_d \le 4095$
Operation:	$((FSR2) + z_s) \rightarrow f_d$
Status Affected:	None
Encoding: 1st word (source) 2nd word (destin.)	 `1110` `1011` `0zzz` `zzzz`$_s$ `1111` `ffff` `ffff` `ffff`$_d$
Description:	The contents of the source register are moved to destination register 'f_d'. The actual address of the source register is determined by adding the 7-bit literal offset 'z_s' in the first word to the value of FSR2. The address of the destination register is specified by the 12-bit literal 'f_d' in the second word. Both addresses can be anywhere in the 4096-byte data space (000h to FFFh). The `MOVSF` instruction cannot use the PCL, TOSU, TOSH or TOSL as the destination register. If the resultant source address points to an indirect addressing register, the value returned will be 00h.
Words:	2
Cycles:	2

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Determine source addr	Determine source addr	Read source reg
Decode	No operation No dummy read	No operation	Write register 'f' (dest)

Example: `MOVSF [05h], REG2`

Before Instruction

FSR2	=	80h
Contents of 85h	=	33h
REG2	=	11h

After Instruction

FSR2	=	80h
Contents of 85h	=	33h
REG2	=	33h

NOTES: