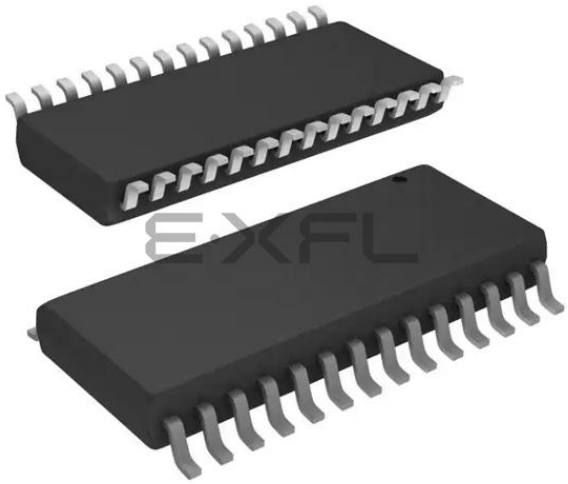
Microchip Technology - DSPIC30F2010T-20I/SOG Datasheet



Welcome to E-XFL.COM

What is "Embedded - Microcontrollers"?

"Embedded - Microcontrollers" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

Applications of "Embedded - Microcontrollers"

Details	
Product Status	Obsolete
Core Processor	dsPIC
Core Size	16-Bit
Speed	20 MIPS
Connectivity	I²C, SPI, UART/USART
Peripherals	Brown-out Detect/Reset, Motor Control PWM, QEI, POR, PWM, WDT
Number of I/O	20
Program Memory Size	12KB (4K x 24)
Program Memory Type	FLASH
EEPROM Size	1K x 8
RAM Size	512 x 8
Voltage - Supply (Vcc/Vdd)	2.5V ~ 5.5V
Data Converters	A/D 6x10b
Oscillator Type	Internal
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	28-SOIC (0.295", 7.50mm Width)
Supplier Device Package	28-SOIC
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/dspic30f2010t-20i-sog

5.5.3 PROGRAMMING VERIFICATION

Once code memory is programmed, the contents of memory can be verified to ensure that programming was successful. Verification requires code memory to be read back and compared against the copy held in the programmer's buffer.

The `READP` command can be used to read back all the programmed code memory.

Alternatively, you can have the programmer perform the verification once the entire device is programmed using a checksum computation, as described in **Section 6.8 "Checksum Computation"**.

5.6 Data EEPROM Programming

5.6.1 OVERVIEW

The panel architecture for the data EEPROM memory array consists of 128 rows of sixteen 16-bit data words. Each panel stores 2K words. All devices have either one or no memory panels. Devices with data EEPROM provide either 512 words, 1024 words or 2048 words of memory on the one panel (see Table 5-3).

TABLE 5-3: DATA EEPROM SIZE

Device	Data EEPROM Size (Words)	Number of Rows
dsPIC30F2010	512	32
dsPIC30F2011	0	0
dsPIC30F2012	0	0
dsPIC30F3010	512	32
dsPIC30F3011	512	32
dsPIC30F3012	512	32
dsPIC30F3013	512	32
dsPIC30F3014	512	32
dsPIC30F4011	512	32
dsPIC30F4012	512	32
dsPIC30F4013	512	32
dsPIC30F5011	512	32
dsPIC30F5013	512	32
dsPIC30F5015	512	32
dsPIC30F5016	512	32
dsPIC30F6010	2048	128
dsPIC30F6010A	2048	128
dsPIC30F6011	1024	64
dsPIC30F6011A	1024	64
dsPIC30F6012	2048	128
dsPIC30F6012A	2048	128
dsPIC30F6013	1024	64
dsPIC30F6013A	1024	64
dsPIC30F6014	2048	128
dsPIC30F6014A	2048	128
dsPIC30F6015	2048	128

5.6.2 PROGRAMMING METHODOLOGY

The programming executive uses the `PROGD` command to program the data EEPROM. Figure 5-4 illustrates the flowchart of the process. Firstly, the number of rows to program (RemainingRows) is based on the device size, and the destination address (DestAddress) is set to '0'. In this example, 128 rows (2048 words) of data EEPROM will be programmed.

The first `PROGD` command programs the first row of data EEPROM. Once the command completes successfully, 'RemainingRows' is decremented by 1 and compared with 0. Since there are 127 more rows to program, 'BaseAddress' is incremented by 0x20 to point to the next row of data EEPROM. This process is then repeated until all 128 rows of data EEPROM are programmed.

FIGURE 5-4: FLOWCHART FOR PROGRAMMING dsPIC30F6014A DATA EEPROM

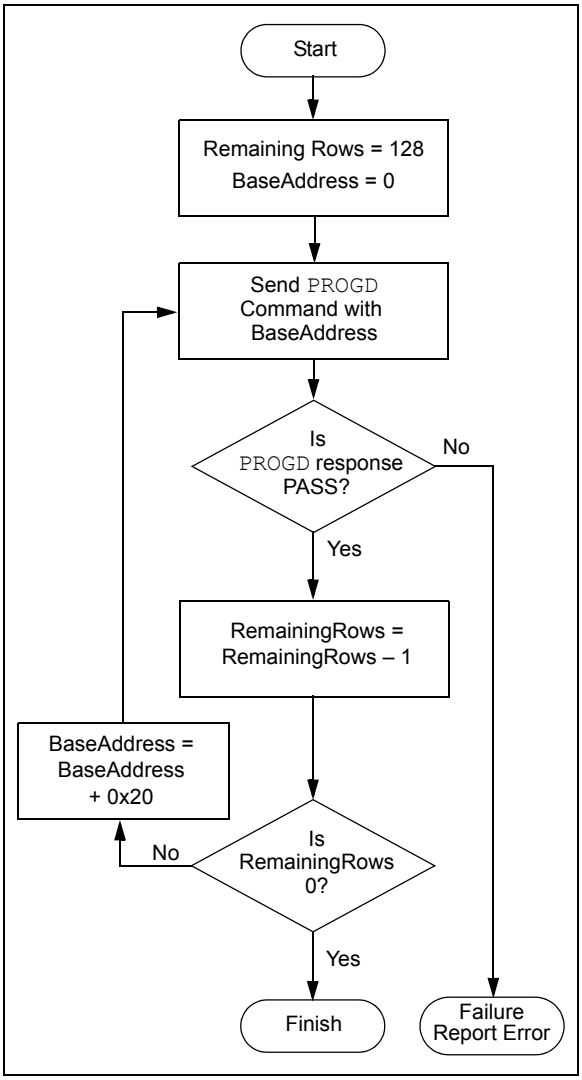


TABLE 5-7: CONFIGURATION BITS DESCRIPTION

Bit Field	Register	Description
FWPSA<1:0>	FWDT	**Watchdog Timer Prescaler A** 11 = 1:512 10 = 1:64 01 = 1:8 00 = 1:1
FWPSB<3:0>	FWDT	**Watchdog Timer Prescaler B** 1111 = 1:16 1110 = 1:15 • • • 0001 = 1:2 0000 = 1:1
FWDTEN	FWDT	**Watchdog Enable** 1 = Watchdog enabled (LPRC oscillator cannot be disabled. Clearing the SWDTEN bit in the RCON register will have no effect) 0 = Watchdog disabled (LPRC oscillator can be disabled by clearing the SWDTEN bit in the RCON register)
MCLREN	FBORPOR	**Master Clear Enable** $\overline{\text{MCLR}}$ 1 = Master Clear pin ($\overline{\text{MCLR}}$) is enabled 0 = $\overline{\text{MCLR}}$ pin is disabled
PWMPIN	FBORPOR	**Motor Control PWM Module Pin Mode** 1 = PWM module pins controlled by PORT register at device Reset (tri-stated) 0 = PWM module pins controlled by PWM module at device Reset (configured as output pins)
HPOL	FBORPOR	**Motor Control PWM Module High-Side Polarity** 1 = PWM module high-side output pins have active-high output polarity 0 = PWM module high-side output pins have active-low output polarity
LPOL	FBORPOR	**Motor Control PWM Module Low-Side Polarity** 1 = PWM module low-side output pins have active-high output polarity 0 = PWM module low-side output pins have active-low output polarity
BOREN	FBORPOR	**PBOR Enable** 1 = PBOR enabled 0 = PBOR disabled
BORV<1:0>	FBORPOR	**Brown-out Voltage Select** 11 = 2.0V (not a valid operating selection) 10 = 2.7V 01 = 4.2V 00 = 4.5V
FPWRT<1:0>	FBORPOR	**Power-on Reset Timer Value Select** 11 = PWRT = 64 ms 10 = PWRT = 16 ms 01 = PWRT = 4 ms 00 = Power-up Timer disabled
RBS<1:0>	FBS	**Boot Segment Data RAM Code Protection (only present in dsPIC30F5011/5013/6010A/6011A/6012A/6013A/6014A/6015)** 11 = No Data RAM is reserved for Boot Segment 10 = Small-sized Boot RAM [128 bytes of RAM are reserved for Boot Segment] 01 = Medium-sized Boot RAM [256 bytes of RAM are reserved for Boot Segment] 00 = Large-sized Boot RAM [512 bytes of RAM are reserved for Boot Segment in dsPIC30F5011/5013, and 1024 bytes in dsPIC30F6010A/6011A/6012A/6013A/6014A/6015]

5.7.2 PROGRAMMING METHODOLOGY

System operation Configuration bits are inherently different than all other memory cells. Unlike code memory, data EEPROM and code-protect Configuration bits, the system operation bits cannot be erased. If the chip is erased with the ERASEB command, the system-operation bits retain their previous value. Consequently, you should make no assumption about the value of the system operation bits. They should always be programmed to their desired setting.

Configuration bits are programmed as a single word at a time using the PROGC command. The PROGC command specifies the configuration data and Configuration register address. When Configuration bits are programmed, any unimplemented bits must be programmed with a '0', and any reserved bits must be programmed with a '1'.

Four PROGC commands are required to program all the Configuration bits. Figure 5-5 illustrates the flowchart of Configuration bit programming.

Note: If the General Code Segment Code Protect (GCP) bit is programmed to '0', code memory is code-protected and cannot be read. Code memory must be verified before enabling read protection. See **Section 5.7.4 "Code-Protect Configuration Bits"** for more information about code-protect Configuration bits.

5.7.3 PROGRAMMING VERIFICATION

Once the Configuration bits are programmed, the contents of memory should be verified to ensure that the programming was successful. Verification requires the Configuration bits to be read back and compared against the copy held in the programmer's buffer. The READD command reads back the programmed Configuration bits and verifies whether the programming was successful.

Any unimplemented Configuration bits are read-only and read as '0'.

5.7.4 CODE-PROTECT CONFIGURATION BITS

The FBS, FSS and FGS Configuration registers are special Configuration registers that control the size and level of code protection for the Boot Segment, Secure Segment and General Segment, respectively. For each segment, two main forms of code protection are provided. One form prevents code memory from being written (write protection), while the other prevents code memory from being read (read protection).

The BWRP, SWRP and GWRP bits control write protection; and BSS<2:0>, SSS<2:0> and GSS<1:0> bits control read protection. The Chip Erase ERASEB command sets all the code protection bits to '1', which allows the device to be programmed.

When write protection is enabled, any programming operation to code memory will fail. When read protection is enabled, any read from code memory will cause a '0x0' to be read, regardless of the actual contents of code memory. Since the programming executive always verifies what it programs, attempting to program code memory with read protection enabled will also result in failure.

It is imperative that all code protection bits are '1' while the device is being programmed and verified. Only after the device is programmed and verified should any of the above bits be programmed to '0' (see **Section 5.7 "Configuration Bits Programming"**).

In addition to code memory protection, parts of data EEPROM and/or data RAM can be configured to be accessible only by code resident in the Boot Segment and/or Secure Segment. The sizes of these "reserved" sections are user-configurable, using the EBS, RBS<1:0>, ESS<1:0> and RSS<1:0> bits.

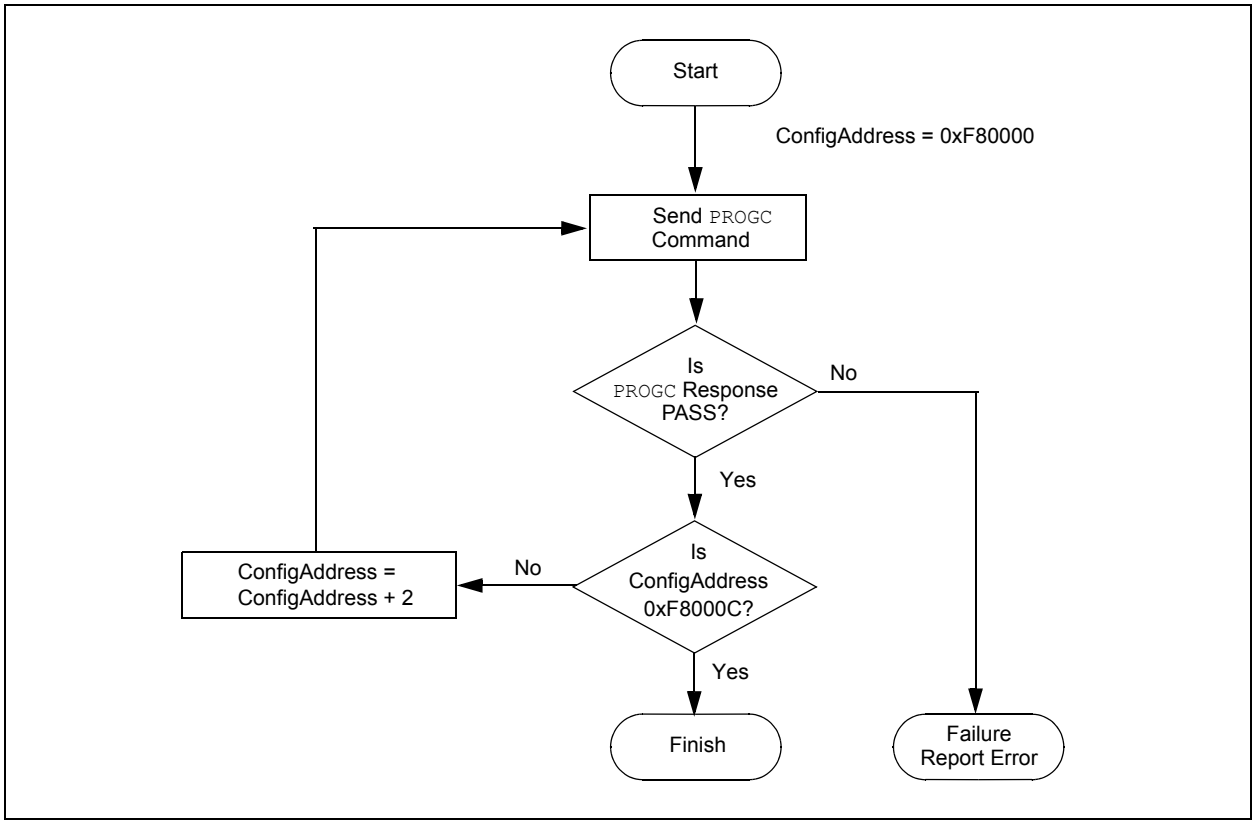
Note 1: All bits in the FBS, FSS and FGS Configuration registers can only be programmed to a value of '0'. ERASEB is the only way to reprogram code-protect bits from ON ('0') to OFF ('1').

2: If any of the code-protect bits in FBS, FSS, or FGS are clear, the entire device must be erased before it can be reprogrammed.

5.8 Exiting Enhanced ICSP Mode

The Enhanced ICSP mode is exited by removing power from the device or bringing MCLR to VIL. When normal user mode is next entered, the program that was stored using Enhanced ICSP will execute.

FIGURE 5-5: CONFIGURATION BIT PROGRAMMING FLOW



8.0 PROGRAMMING EXECUTIVE COMMANDS

8.1 Command Set

The programming executive command set is shown in Table 8-1. This table contains the opcode, mnemonic, length, time out and description for each command. Functional details on each command are provided in the command descriptions (see **Section 8.5 "Command Descriptions"**).

8.2 Command Format

All programming executive commands have a general format consisting of a 16-bit header and any required data for the command (see Figure 8-1). The 16-bit header consists of a 4-bit opcode field, which is used to identify the command, followed by a 12-bit command length field.

FIGURE 8-1: COMMAND FORMAT

15 12	11 0
Opcode	Length
Command Data First Word (if required)	
•	
•	
Command Data Last Word (if required)	

The command opcode must match one of those in the command set. Any command that is received which does not match the list in Table 8-1 will return a "NACK" response (see **Section 9.2.1 "Opcode Field"**).

The command length is represented in 16-bit words since the SPI operates in 16-bit mode. The programming executive uses the Command Length field to determine the number of words to read from the SPI port. If the value of this field is incorrect, the command will not be properly received by the programming executive.

8.3 Packed Data Format

When 24-bit instruction words are transferred across the 16-bit SPI interface, they are packed to conserve space using the format shown in Figure 8-2. This format minimizes traffic over the SPI and provides the programming executive with data that is properly aligned for performing table write operations.

FIGURE 8-2: PACKED INSTRUCTION WORD FORMAT

15 8	7 0
lsw1	
MSB2	MSB1
lsw2	

lswx: Least significant 16 bits of instruction word
MSBx: Most Significant Byte of instruction word

Note:	When the number of instruction words transferred is odd, MSB2 is zero and lsw2 cannot be transmitted.

8.4 Programming Executive Error Handling

The programming executive will "NACK" all unsupported commands. Additionally, due to the memory constraints of the programming executive, no checking is performed on the data contained in the Programmer command. It is the responsibility of the programmer to command the programming executive with valid command arguments, or the programming operation may fail. Additional information on error handling is provided in **Section 9.2.3 "QE_Code Field"**.

8.5.3 READP COMMAND

15 12	11 8	7 0
Opcode	Length	
N		
Reserved	Addr_MSB	
Addr_LS		

Field	Description
Opcode	0x2
Length	0x4
N	Number of 24-bit instructions to read (max of 32768)
Reserved	0x0
Addr_MSB	MSB of 24-bit source address
Addr_LS	LS 16 bits of 24-bit source address

The READP command instructs the programming executive to read N 24-bit words of code memory starting from the 24-bit address specified by Addr_MSB and Addr_LS. This command can only be used to read 24-bit data. All data returned in response to this command uses the packed data format described in **Section 8.3 "Packed Data Format"**.

Expected Response (2 + 3 * N/2 words for N even):
0x1200
2 + 3 * N/2
Least significant program memory word 1
...
Least significant data word N

Expected Response (4 + 3 * (N – 1)/2 words for N odd):
0x1200
4 + 3 * (N – 1)/2
Least significant program memory word 1
...
MSB of program memory word N (zero padded)

Note:	Reading unimplemented memory will cause the programming executive to reset.

8.5.4 PROGD COMMAND

15 12	11 8	7 0
Opcode	Length	
Reserved	Addr_MSB	
Addr_LS		
D_1		
D_2		
...		
D_16		

Field	Description
Opcode	0x4
Length	0x13
Reserved	0x0
Addr_MSB	MSB of 24-bit destination address
Addr_LS	LS 16 bits of 24-bit destination address
D_1	16-bit data word 1
D_2	16-bit data word 2
...	16-bit data words 3 through 15
D_16	16-bit data word 16

The PROGD command instructs the programming executive to program one row of data EEPROM. The data to be programmed is specified by the 16 data words (D_1, D_2,..., D_16) and is programmed to the destination address specified by Addr_MSB and Addr_LSB. The destination address should be a multiple of 0x20.

Once the row of data EEPROM has been programmed, the programming executive verifies the programmed data against the data in the command.

Expected Response (2 words):
0x1400
0x0002

Note:	Refer to Table 5-3 for data EEPROM size information.

8.5.7 ERASEB COMMAND

15 12	11 2	0
Opcode	Length	
Reserved		MS

Field	Description
Opcode	0x7
Length	0x2
Reserved	0x0
MS	Select memory to erase: 0x0 = All Code in General Segment 0x1 = All Data EEPROM in General Segment 0x2 = All Code and Data EEPROM in General Segment, interrupt vectors and FGS Configuration register 0x3 = Full Chip Erase 0x4 = All Code and Data EEPROM in Boot, Secure and General Segments, and FBS, FSS and FGS Configuration registers 0x5 = All Code and Data EEPROM in Secure and General Segments, and FSS and FGS Configuration registers 0x6 = All Data EEPROM in Boot Segment 0x7 = All Data EEPROM in Secure Segment

The ERASEB command performs a Bulk Erase. The MS field selects the memory to be bulk erased, with options for erasing Code and/or Data EEPROM in individual memory segments.

When Full Chip Erase is selected, the following memory regions are erased:

- All code memory (even if code-protected)
- All data EEPROM
- All code-protect Configuration registers

Only the executive code memory, Unit ID, device ID and Configuration registers that are not code-protected remain intact after a Chip Erase.

Expected Response (2 words):
0x1700
0x0002

Note:	A Full Chip Erase cannot be performed in low-voltage programming systems (V_{DD} less than 4.5 volts). ERASED and ERASEP must be used to erase code memory, executive memory and data memory. Alternatively, individual Segment Erase operations may be performed.

8.5.8 ERASED COMMAND

15 12	11 8	7 0
Opcode	Length	
Num_Rows		Addr_MSB
Addr_LS		

Field	Description
Opcode	0x8
Length	0x3
Num_Rows	Number of rows to erase (max of 128)
Addr_MSB	MSB of 24-bit base address
Addr_LS	LS 16 bits of 24-bit base address

The ERASED command erases the specified number of rows of data EEPROM from the specified base address. The specified base address must be a multiple of 0x20. Since the data EEPROM is mapped to program space, a 24-bit base address must be specified.

After the erase is performed, all targeted bytes of data EEPROM will contain 0xFF.

Expected Response (2 words):
0x1800
0x0002

Note:	The ERASED command cannot be used to erase the Configuration registers or device ID. Code-protect Configuration registers can only be erased with the ERASEB command, while the device ID is read-only.

8.5.9 ERASEP COMMAND

15 12	11 8	7 0
Opcode	Length	
Num_Rows		Addr_MSB
Addr_LS		

Field	Description
Opcode	0x9
Length	0x3
Num_Rows	Number of rows to erase
Addr_MSB	MSB of 24-bit base address
Addr_LS	LS 16 bits of 24-bit base address

The ERASEP command erases the specified number of rows of code memory from the specified base address. The specified base address must be a multiple of 0x40.

Once the erase is performed, all targeted words of code memory contain 0xFFFFFF.

Expected Response (2 words):
0x1900
0x0002

Note: The ERASEP command cannot be used to erase the Configuration registers or device ID. Code-protect Configuration registers can only be erased with the ERASEB command, while the device ID is read-only.

8.5.10 QBLANK COMMAND

15 12	11 0
Opcode	Length
PSize	
Reserved	DSize

Field	Description
Opcode	0xA
Length	0x3
PSize	Length of program memory to check (in 24-bit words), max of 49152
Reserved	0x0
DSize	Length of data memory to check (in 16-bit words), max of 2048

The QBLANK command queries the programming executive to determine if the contents of code memory and data EEPROM are blank (contains all '1's). The size of code memory and data EEPROM to check must be specified in the command.

The Blank Check for code memory begins at 0x0 and advances toward larger addresses for the specified number of instruction words. The Blank Check for data EEPROM begins at 0x7FFFFE and advances toward smaller addresses for the specified number of data words.

QBLANK returns a QE_Code of 0xF0 if the specified code memory and data EEPROM are blank. Otherwise, QBLANK returns a QE_Code of 0x0F.

Expected Response (2 words for blank device):
0x1AF0
0x0002

Expected Response (2 words for non-blank device):
0x1A0F
0x0002

Note: The QBLANK command does not check the system Configuration registers. The READD command must be used to determine the state of the Configuration registers.

8.5.11 QVER COMMAND

15 12	11 0
Opcode	Length

Field	Description
Opcode	0xB
Length	0x1

The QVER command queries the version of the programming executive software stored in test memory. The “version.revision” information is returned in the response’s QE_Code using a single byte with the following format: main version in upper nibble and revision in the lower nibble (i.e., 0x23 is version 2.3 of programming executive software).

Expected Response (2 words):
0x1BMN (where “MN” stands for version M.N)
0x0002

9.0 PROGRAMMING EXECUTIVE RESPONSES

9.1 Overview

The programming executive sends a response to the programmer for each command that it receives. The response indicates if the command was processed correctly, and includes any required response or error data.

The programming executive response set is shown in Table 9-1. This table contains the opcode, mnemonic and description for each response. The response format is described in **Section 9.2 “Response Format”**.

TABLE 9-1: PROGRAMMING EXECUTIVE RESPONSE SET

Opcode	Mnemonic	Description
0x1	PASS	Command successfully processed.
0x2	FAIL	Command unsuccessfully processed.
0x3	NACK	Command not known.

9.2 Response Format

As shown in Example 9-1, all programming executive responses have a general format consisting of a two word header and any required data for the command. Table 9-2 lists the fields and their descriptions.

EXAMPLE 9-1: FORMAT

15 12	11 8	7 0
Opcode	Last_Cmd	QE_Code
Length		
D_1 (if applicable)		
...		
D_N (if applicable)		

TABLE 9-2: FIELDS AND DESCRIPTIONS

Field	Description
Opcode	Response opcode.
Last_Cmd	Programmer command that generated the response.
QE_Code	Query code or Error code.
Length	Response length in 16-bit words (includes 2 header words.)
D_1	First 16-bit data word (if applicable).
D_N	Last 16-bit data word (if applicable).

9.2.1 Opcode FIELD

The Opcode is a 4-bit field in the first word of the response. The Opcode indicates how the command was processed (see Table 9-1). If the command is processed successfully, the response opcode is PASS. If there is an error in processing the command, the response opcode is FAIL, and the QE_Code indicates the reason for the failure. If the command sent to the programming executive is not identified, the programming executive returns a NACK response.

9.2.2 Last_Cmd FIELD

The Last_Cmd is a 4-bit field in the first word of the response and indicates the command that the programming executive processed. Since the programming executive can only process one command at a time, this field is technically not required. However, it can be used to verify whether the programming executive correctly received the command that the programmer transmitted.

9.2.3 QE_Code FIELD

The QE_Code is a byte in the first word of the response. This byte is used to return data for query commands, and error codes for all other commands.

When the programming executive processes one of the two query commands (QBLANK or QVER), the returned opcode is always PASS and the QE_Code holds the query response data. The format of the QE_Code for both queries is shown in Table 9-3.

TABLE 9-3: QE_Code FOR QUERIES

Query	QE_Code
QBLANK	0x0F = Code memory and data EEPROM are NOT blank 0xF0 = Code memory and data EEPROM are blank
QVER	0xMN, where programming executive software version = M.N (i.e., 0x32 means software version 3.2)

When the programming executive processes any command other than a Query, the QE_Code represents an error code. Supported error codes are shown in Table 9-4. If a command is successfully processed, the returned QE_Code is set to 0x0, which indicates that there was no error in the command processing. If the verify of the programming for the PROGD, PROGP or PROGC command fails, the QE_Code is set to 0x1. For all other programming executive errors, the QE_Code is 0x2.

TABLE 9-4: QE_Code FOR NON-QUERY COMMANDS

QE_Code	Description
0x0	No error
0x1	Verify failed
0x2	Other error

9.2.4 RESPONSE LENGTH

The response length indicates the length of the programming executive's response in 16-bit words. This field includes the 2 words of the response header.

With the exception of the response for the READD and READP commands, the length of each response is only 2 words.

The response to the READD command is N + 2 words, where N is the number of words specified in the READD command.

The response to the READP command uses the packed instruction word format described in **Section 8.3 "Packed Data Format"**. When reading an odd number of program memory words (N odd), the response to the READP command is (3 • (N + 1)/2 + 2) words. When reading an even number of program memory words (N even), the response to the READP command is (3 • N/2 + 2) words.

10.0 DEVICE ID

The device ID region is 2 x 16 bits and can be read using the `READD` command. This region of memory is read-only and can also be read when code protection is enabled.

Table 10-1 shows the device ID for each device, Table 10-2 shows the device ID registers and Table 10-3 describes the bit field of each register.

TABLE 10-1: DEVICE IDS

Device	DEVID	Silicon Revision							
		A0	A1	A2	A3	A4	B0	B1	B2
dsPIC30F2010	0x0040	0x1000	0x1001	0x1002	0x1003	0x1004	—	—	—
dsPIC30F2011	0x0240	—	0x1001	—	—	—	—	—	—
dsPIC30F2012	0x0241	—	0x1001	—	—	—	—	—	—
dsPIC30F3010	0x01C0	0x1000	0x1001	0x1002	—	—	—	—	—
dsPIC30F3011	0x01C1	0x1000	0x1001	0x1002	—	—	—	—	—
dsPIC30F3012	0x00C1	—	—	—	—	—	0x1040	0x1041	—
dsPIC30F3013	0x00C3	—	—	—	—	—	0x1040	0x1041	—
dsPIC30F3014	0x0160	—	0x1001	0x1002	—	—	—	—	—
dsPIC30F4011	0x0101	—	0x1001	0x1002	0x1003	0x1003	—	—	—
dsPIC30F4012	0x0100	—	0x1001	0x1002	0x1003	0x1003	—	—	—
dsPIC30F4013	0x0141	—	0x1001	0x1002	—	—	—	—	—
dsPIC30F5011	0x0080	—	0x1001	0x1002	0x1003	0x1003	—	—	—
dsPIC30F5013	0x0081	—	0x1001	0x1002	0x1003	0x1003	—	—	—
dsPIC30F5015	0x0200	0x1000	—	—	—	—	—	—	—
dsPIC30F5016	0x0201	0x1000	—	—	—	—	—	—	—
dsPIC30F6010	0x0188	—	—	—	—	—	—	0x1040	0x1042
dsPIC30F6010A	0x0281	—	—	0x1002	0x1003	0x1004	—	—	—
dsPIC30F6011	0x0192	—	—	—	0x1003	—	—	0x1040	0x1042
dsPIC30F6011A	0x02C0	—	—	0x1002	—	—	0x1040	0x1041	—
dsPIC30F6012	0x0193	—	—	—	0x1003	—	—	0x1040	0x1042
dsPIC30F6012A	0x02C2	—	—	0x1002	—	—	0x1040	0x1041	—
dsPIC30F6013	0x0197	—	—	—	0x1003	—	—	0x1040	0x1042
dsPIC30F6013A	0x02C1	—	—	0x1002	—	—	0x1040	0x1041	—
dsPIC30F6014	0x0198	—	—	—	0x1003	—	—	0x1040	0x1042
dsPIC30F6014A	0x02C3	—	—	0x1002	—	—	0x1040	0x1041	—
dsPIC30F6015	0x0280	—	—	0x1002	0x1003	0x1004	—	—	—

TABLE 10-2: dsPIC30F DEVICE ID REGISTERS

Address	Name	Bit															
		15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
0xFF0000	DEVID	DEVID<15:0>															
0xFF0002	DEVREV	PROC<3:0>				REV<5:0>						DOT<5:0>					

11.6 Erasing Program Memory in Low-Voltage Systems

The procedure for erasing program memory (all code memory and data memory) in low-voltage systems (with VDD between 2.5 volts and 4.5 volts) is quite different than the procedure for erasing program memory in normal-voltage systems. Instead of using a Bulk Erase operation, each region of memory must be individually erased by row. Namely, all of the code memory, executive memory and data memory must be erased one row at a time. This procedure is detailed in Table 11-5.

Due to security restrictions, the FBS, FSS and FGS register cannot be erased in low-voltage systems. Once any bits in the FGS register are programmed to '0', they can only be set back to '1' by performing a Bulk Erase in a normal-voltage system. Alternatively, a Segment Erase operation can be performed instead of a Bulk Erase.

Normal-voltage systems can also be used to erase program memory as shown in Table 11-5. However, since this method is more time-consuming and does not clear the code-protect bits, it is not recommended.

Note: Program memory must be erased before writing any data to program memory.

TABLE 11-5: SERIAL INSTRUCTION EXECUTION FOR ERASING PROGRAM MEMORY (EITHER IN LOW-VOLTAGE OR NORMAL-VOLTAGE SYSTEMS)

Command (Binary)	Data (Hexadecimal)	Description
Step 1: Exit the Reset vector.		
0000	040100	GOTO 0x100
0000	040100	GOTO 0x100
0000	000000	NOP
Step 2: Initialize NVMADR and NVMADRU to erase code memory and initialize W7 for row address updates.		
0000	EB0300	CLR W6
0000	883B16	MOV W6, NVMADR
0000	883B26	MOV W6, NVMADRU
0000	200407	MOV #0x40, W7
Step 3: Set NVMCON to erase 1 row of code memory.		
0000	24071A	MOV #0x4071, W10
0000	883B0A	MOV W10, NVMCON
Step 4: Unlock the NVMCON to erase 1 row of code memory.		
0000	200558	MOV #0x55, W8
0000	883B38	MOV W8, NVMKEY
0000	200AA9	MOV #0xAA, W9
0000	883B39	MOV W9, NVMKEY
Step 5: Initiate the erase cycle.		
0000	A8E761	BSET NVMCON, #WR
0000	000000	NOP
0000	000000	NOP
–	–	Externally time 'P13a' ms (see **Section 13.0 "AC/DC Characteristics and Timing Requirements"**)
0000	000000	NOP
0000	000000	NOP
0000	A9E761	BCLR NVMCON, #WR
0000	000000	NOP
0000	000000	NOP

11.7 Writing Configuration Memory

The FOSC, FWDT, FBORPOR and FICD registers are not erasable. It is recommended that all Configuration registers be set to a default value after erasing program memory. The FWDT, FBORPOR and FICD registers can be set to a default all '1's value by programming 0xFFFF to each register. Since these registers contain unimplemented bits that read as '0' the default values shown in Table 11-6 will be read instead of 0xFFFF. The recommended default FOSC value is 0xC100, which selects the FRC clock oscillator setting.

The FGS, FBS and FSS Configuration registers are special since they enable code protection for the device. For security purposes, once any bit in these registers is programmed to '0' (to enable some code protection feature), it can only be set back to '1' by performing a Bulk Erase or Segment Erase as described in **Section 11.5 "Erasing Program Memory in Normal-Voltage Systems"**. Programming these bits from a '0' to '1' is not possible, but they may be programmed from a '1' to a '0' to enable code protection.

Table 11-7 shows the ICSP programming details for clearing the Configuration registers. In Step 1, the Reset vector is exited. In Step 2, the write pointer (W7) is loaded with 0x0000, which is the original destination address (in TBLPAG 0xF8 of program memory). In Step 3, the NVMCON is set to program one Configuration register. In Step 4, the TBLPAG register is initialized, to 0xF8, for writing to the Configuration registers. In Step 5, the value to write to the each Configuration register (0xFFFF) is loaded to W6. In Step 6, the Configuration register data is written to the write latch using the TBLWTL instruction. In Steps 7 and 8, the NVMCON is unlocked for programming and the programming cycle is initiated, as described in **Section 11.4 "Flash Memory Programming in ICSP Mode"**. In Step 9, the internal PC is set to 0x100 as a safety measure to prevent the PC from incrementing into unimplemented memory. Lastly, Steps 3-9 are repeated six times until all seven Configuration registers are cleared.

TABLE 11-6: DEFAULT CONFIGURATION REGISTER VALUES

Address	Register	Default Value
0xF80000	FOSC	0xC100
0xF80002	FWDT	0x803F
0xF80004	FBORPOR	0x87B3
0xF80006	FBS	0x310F
0xF80008	FSS	0x330F
0xF8000A	FGS	0x0007
0xF8000C	FICD	0xC003

TABLE 11-7: SERIAL INSTRUCTION EXECUTION FOR WRITING CONFIGURATION REGISTERS

Command (Binary)	Data (Hexadecimal)	Description
Step 1: Exit the Reset vector.		
0000 0000 0000	040100 040100 000000	GOTO 0x100 GOTO 0x100 NOP
Step 2: Initialize the write pointer (W7) for the TBLWT instruction.		
0000	200007	MOV #0x0000, W7
Step 3: Set the NVMCON to program 1 Configuration register.		
0000 0000	24008A 883B0A	MOV #0x4008, W10 MOV W10, NVMCON
Step 4: Initialize the TBLPAG register.		
0000 0000	200F80 880190	MOV #0xF8, W0 MOV W0, TBLPAG
Step 5: Load the Configuration register data to W6.		
0000 0000	2xxxx0 000000	MOV #<CONFIG_VALUE>, W0 NOP

11.8 Writing Code Memory

The procedure for writing code memory is similar to the procedure for clearing the Configuration registers, except that 32 instruction words are programmed at a time. To facilitate this operation, working registers W0:W5 are used as temporary holding registers for the data to be programmed.

Table 11-8 shows the ICSP programming details, including the serial pattern with the ICSP command code, which must be transmitted Least Significant bit first using the PGC and PGD pins (see Figure 11-2). In Step 1, the Reset vector is exited. In Step 2, the NVMCON register is initialized for single-panel programming of code memory. In Step 3, the 24-bit starting destination address for programming is loaded into the TBLPAG register and W7 register. The upper byte of the starting destination address is stored to TBLPAG, while the lower 16 bits of the destination address are stored to W7.

To minimize the programming time, the same packed instruction format that the programming executive uses is utilized (Figure 8-2). In Step 4, four packed instruction words are stored to working registers W0:W5 using the `MOV` instruction and the read pointer W6 is initialized. The contents of W0:W5 holding the packed instruction word data is shown in Figure 11-4.

In Step 5, eight `TBLWT` instructions are used to copy the data from W0:W5 to the write latches of code memory. Since code memory is programmed 32 instruction words at a time, Steps 4 and 5 are repeated eight times to load all the write latches (Step 6).

After the write latches are loaded, programming is initiated by writing to the NVMKEY and NVMCON registers in Steps 7 and 8. In Step 9, the internal PC is reset to 0x100. This is a precautionary measure to prevent the PC from incrementing into unimplemented memory when large devices are being programmed. Lastly, in Step 10, Steps 2-9 are repeated until all of code memory is programmed.

FIGURE 11-5: PACKED INSTRUCTION WORDS IN W0:W5

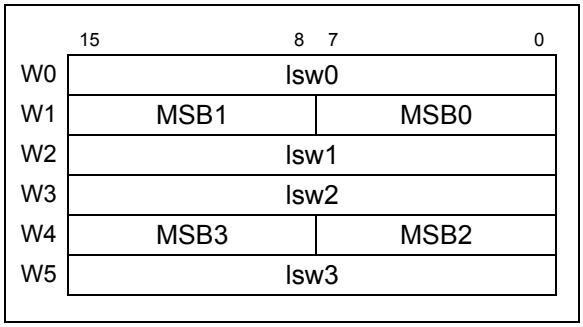


TABLE 11-8: SERIAL INSTRUCTION EXECUTION FOR WRITING CODE MEMORY

Command (Binary)	Data (Hexadecimal)	Description
Step 1: Exit the Reset vector.		
0000 0000 0000	040100 040100 000000	GOTO 0x100 GOTO 0x100 NOP
Step 2: Set the NVMCON to program 32 instruction words.		
0000 0000	24001A 883B0A	MOV #0x4001, W10 MOV W10, NVMCON
Step 3: Initialize the write pointer (W7) for `TBLWT` instruction.		

0000
0000
0000
	200xx0
	880190
	2xxxx7
		MOV #<DestinationAddress23:16>, W0
		MOV W0, TBLPAG
		MOV #<DestinationAddress15:0>, W7

Step 4: Initialize the read pointer (W6) and load W0:W5 with the next 4 instruction words to program.		

0000
0000
0000
0000
0000
0000
	2xxxx0
	2xxxx1
	2xxxx2
	2xxxx3
	2xxxx4
	2xxxx5
		MOV #<LSW0>, W0
		MOV #<MSB1:MSB0>, W1
		MOV #<LSW1>, W2
		MOV #<LSW2>, W3
		MOV #<MSB3:MSB2>, W4
		MOV #<LSW3>, W5

TABLE 11-8: SERIAL INSTRUCTION EXECUTION FOR WRITING CODE MEMORY (CONTINUED)

Command (Binary)	Data (Hexadecimal)	Description
Step 5: Set the read pointer (W6) and load the (next set of) write latches.		
0000	EB0300	CLR W6
0000	000000	NOP
0000	BB0BB6	TBLWTL [W6++], [W7]
0000	000000	NOP
0000	000000	NOP
0000	BBDBB6	TBLWTH.B [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BBEBB6	TBLWTH.B [W6++], [++W7]
0000	000000	NOP
0000	000000	NOP
0000	BB1BB6	TBLWTL [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BB0BB6	TBLWTL [W6++], [W7]
0000	000000	NOP
0000	000000	NOP
0000	BBDBB6	TBLWTH.B [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BBEBB6	TBLWTH.B [W6++], [++W7]
0000	000000	NOP
0000	000000	NOP
0000	BB1BB6	TBLWTL [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
Step 6: Repeat steps 4-5 eight times to load the write latches for 32 instructions.		
Step 7: Unlock the NVMCON for writing.		
0000	200558	MOV #0x55, W8
0000	883B38	MOV W8, NVMKEY
0000	200AA9	MOV #0xAA, W9
0000	883B39	MOV W9, NVMKEY
Step 8: Initiate the write cycle.		
0000	A8E761	BSET NVMCON, #WR
0000	000000	NOP
0000	000000	NOP
–	–	Externally time 'P12a' ms (see **Section 13.0 "AC/DC Characteristics and Timing Requirements"**)
0000	000000	NOP
0000	000000	NOP
0000	A9E761	BCLR NVMCON, #WR
0000	000000	NOP
0000	000000	NOP
Step 9: Reset device internal PC.		
0000	040100	GOTO 0x100
0000	000000	NOP
Step 10: Repeat steps 2-9 until all code memory is programmed.		

11.9 Writing Data EEPROM

The procedure for writing data EEPROM is very similar to the procedure for writing code memory, except that fewer words are programmed in each operation. When writing data EEPROM, one row of data EEPROM is programmed at a time. Each row consists of sixteen 16-bit data words. Since fewer words are programmed during each operation, only working registers W0:W3 are used as temporary holding registers for the data to be programmed.

Table 11-9 shows the ICSP programming details for writing data EEPROM. Note that a different NVMCON value is required to write to data EEPROM, and that the TBLPAG register is hard-coded to 0x7F (the upper byte address of all locations of data EEPROM).

TABLE 11-9: SERIAL INSTRUCTION EXECUTION FOR WRITING DATA EEPROM

Command (Binary)	Data (Hexadecimal)	Description
Step 1: Exit the Reset vector.		
0000	040100	GOTO 0x100
0000	040100	GOTO 0x100
0000	000000	NOP
Step 2: Set the NVMCON to write 16 data words.		
0000	24005A	MOV #0x4005, W10
0000	883B0A	MOV W10, NVMCON
Step 3: Initialize the write pointer (W7) for TBLWT instruction.		
0000	2007F0	MOV #0x7F, W0
0000	880190	MOV W0, TBLPAG
0000	2xxxx7	MOV #<DestinationAddress15:0>, W7
Step 4: Load W0:W3 with the next 4 data words to program.		
0000	2xxxx0	MOV #<WORD0>, W0
0000	2xxxx1	MOV #<WORD1>, W1
0000	2xxxx2	MOV #<WORD2>, W2
0000	2xxxx3	MOV #<WORD3>, W3
Step 5: Set the read pointer (W6) and load the (next set of) write latches.		
0000	EB0300	CLR W6
0000	000000	NOP
0000	BB1BB6	TBLWTL [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BB1BB6	TBLWTL [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BB1BB6	TBLWTL [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
0000	BB1BB6	TBLWTL [W6++], [W7++]
0000	000000	NOP
0000	000000	NOP
Step 6: Repeat steps 4-5 four times to load the write latches for 16 data words.		

TABLE 11-9: SERIAL INSTRUCTION EXECUTION FOR WRITING DATA EEPROM (CONTINUED)

Command (Binary)	Data (Hexadecimal)	Description
Step 7: Unlock the NVMCON for writing.		
0000	200558	MOV #0x55, W8
0000	883B38	MOV W8, NVMKEY
0000	200AA9	MOV #0xAA, W9
0000	883B39	MOV W9, NVMKEY
Step 8: Initiate the write cycle.		
0000	A8E761	BSET NVMCON, #WR
0000	000000	NOP
0000	000000	NOP
—	—	Externally time 'P12a' ms (see **Section 13.0 "AC/DC Characteristics and Timing Requirements"**)
0000	000000	NOP
0000	000000	NOP
0000	A9E761	BCLR NVMCON, #WR
0000	000000	NOP
0000	000000	NOP
Step 9: Reset device internal PC.		
0000	040100	GOTO 0x100
0000	000000	NOP
Step 10: Repeat steps 2-9 until all data memory is programmed.		

APPENDIX A: DEVICE-SPECIFIC INFORMATION

A.1 Checksum Computation

The checksum computation is described in **Section 6.8 "Checksum Computation"**. Table A-1 shows how this 16-bit computation can be made for each dsPIC30F device. Computations for read code protection are shown both enabled and disabled. The checksum values assume that the Configuration registers are also erased. However, when code protection is enabled, the value of the FGS register is assumed to be 0x5.

A.2 dsPIC30F5011 and dsPIC30F5013

A.2.1 ICSP PROGRAMMING

The dsPIC30F5011 and dsPIC30F5013 processors require that the FBS and FSS registers be programmed with 0x0000 before the device is chip erased. The steps to perform this action are shown in Table 11-4.

A.2.2 ENHANCED ICSP PROGRAMMING

The dsPIC30F5011 and dsPIC30F5013 processors require that the FBS and FSS registers be programmed with 0x0000 using the `PROGC` command before the `ERASEB` command is used to erase the chip.

TABLE A-1: CHECKSUM COMPUTATION

Device	Read Code Protection	Checksum Computation	Erased Value	Value with 0xAAAAAA at 0x0 and Last Code Address
dsPIC30F2010	Disabled	CFGB+SUM(0:001FFF)	0xD406	0xD208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F2011	Disabled	CFGB+SUM(0:001FFF)	0xD406	0xD208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F2012	Disabled	CFGB+SUM(0:001FFF)	0xD406	0xD208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F3010	Disabled	CFGB+SUM(0:003FFF)	0xA406	0xA208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F3011	Disabled	CFGB+SUM(0:003FFF)	0xA406	0xA208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F3012	Disabled	CFGB+SUM(0:003FFF)	0xA406	0xA208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F3013	Disabled	CFGB+SUM(0:003FFF)	0xA406	0xA208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F3014	Disabled	CFGB+SUM(0:003FFF)	0xA406	0xA208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F4011	Disabled	CFGB+SUM(0:007FFF)	0x4406	0x4208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F4012	Disabled	CFGB+SUM(0:007FFF)	0x4406	0x4208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F4013	Disabled	CFGB+SUM(0:007FFF)	0x4406	0x4208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F5011	Disabled	CFGB+SUM(0:00AFFF)	0xFC06	0xFA08
	Enabled	CFGB	0x0404	0x0404
dsPIC30F5013	Disabled	CFGB+SUM(0:00AFFF)	0xFC06	0xFA08
	Enabled	CFGB	0x0404	0x0404
dsPIC30F5015	Disabled	CFGB+SUM(0:00AFFF)	0xFC06	0xFA08
	Enabled	CFGB	0x0404	0x0404

Item Description:

SUM(a:b) = Byte sum of locations a to b inclusive (all 3 bytes of code memory)

CFGB = **Configuration Block (masked)** = Byte sum of ((FOSC&0xC10F) + (FWDT&0x803F) + (FBORPOR&0x87B3) + (FBS&0x310F) + (FSS&0x330F) + (FGS&0x0007) + (FICD&0xC003))

TABLE A-1: CHECKSUM COMPUTATION (CONTINUED)

Device	Read Code Protection	Checksum Computation	Erased Value	Value with 0xAAAAAA at 0x0 and Last Code Address
dsPIC30F5016	Disabled	CFGB+SUM(0:00AFFF)	0xFC06	0xFA08
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6010	Disabled	CFGB+SUM(0:017FFF)	0xC406	0xC208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6010A	Disabled	CFGB+SUM(0:017FFF)	0xC406	0xC208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6011	Disabled	CFGB+SUM(0:015FFF)	0xF406	0xF208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6011A	Disabled	CFGB+SUM(0:015FFF)	0xF406	0xF208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6012	Disabled	CFGB+SUM(0:017FFF)	0xC406	0xC208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6012A	Disabled	CFGB+SUM(0:017FFF)	0xC406	0xC208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6013	Disabled	CFGB+SUM(0:015FFF)	0xF406	0xF208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6013A	Disabled	CFGB+SUM(0:015FFF)	0xF406	0xF208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6014	Disabled	CFGB+SUM(0:017FFF)	0xC406	0xC208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6014A	Disabled	CFGB+SUM(0:017FFF)	0xC406	0xC208
	Enabled	CFGB	0x0404	0x0404
dsPIC30F6015	Disabled	CFGB+SUM(0:017FFF)	0xC406	0xC208
	Enabled	CFGB	0x0404	0x0404

Item Description:

SUM(a:b) = Byte sum of locations a to b inclusive (all 3 bytes of code memory)

CFGB = **Configuration Block (masked)** = Byte sum of ((FOSC&0xC10F) + (FWDT&0x803F) + (FBORPOR&0x87B3) + (FBS&0x310F) + (FSS&0x330F) + (FGS&0x0007) + (FICD&0xC003))

APPENDIX B: HEX FILE FORMAT

Flash programmers process the standard HEX format used by the Microchip development tools. The format supported is the Intel® HEX 32 Format (INHX32). Please refer to Appendix A in the "*MPASM User's Guide"* (DS33014) for more information about hex file formats.

The basic format of the hex file is:

```
:BBAAAATTHHHH...HHHHCC
```

Each data record begins with a 9-character prefix and always ends with a 2-character checksum. All records begin with '`:`' regardless of the format. The individual elements are described below.

- `BB` - is a two-digit hexadecimal byte count representing the number of data bytes that appear on the line. Divide this number by two to get the number of words per line.
- `AAAA` - is a four-digit hexadecimal address representing the starting address of the data record. Format is high byte first followed by low byte. The address is doubled because this format only supports 8-bits. Divide the value by two to find the real device address.
- `TT` - is a two-digit record type that will be '00' for data records, '01' for end-of-file records and '04' for extended-address record.
- `HHHH` - is a four-digit hexadecimal data word. Format is low byte followed by high byte. There will be `BB`/2 data words following `TT`.
- `CC` - is a two-digit hexadecimal checksum that is the two's complement of the sum of all the preceding bytes in the line record.

Because the Intel hex file format is byte-oriented, and the 16-bit program counter is not, program memory sections require special treatment. Each 24-bit program word is extended to 32 bits by inserting a so-called "phantom byte". Each program memory address is multiplied by 2 to yield a byte address.

As an example, a section that is located at 0x100 in program memory will be represented in the hex file as 0x200.

The hex file will be produced with the following contents:

```
:020000040000fa
:040200003322110096
:00000001FF
```

Notice that the data record (line 2) has a load address of 0200, while the source code specified address 0x100. Note also that the data is represented in "little-endian" format, meaning the Least Significant Byte (LSB) appears first. The phantom byte appears last, just before the checksum.