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Understanding **Embedded - FPGAs (Field Programmable Gate Array)**

Embedded - FPGAs, or Field Programmable Gate Arrays, are advanced integrated circuits that offer unparalleled flexibility and performance for digital systems. Unlike traditional fixed-function logic devices, FPGAs can be programmed and reprogrammed to execute a wide array of logical operations, enabling customized functionality tailored to specific applications. This reprogrammability allows developers to iterate designs quickly and implement complex functions without the need for custom hardware.

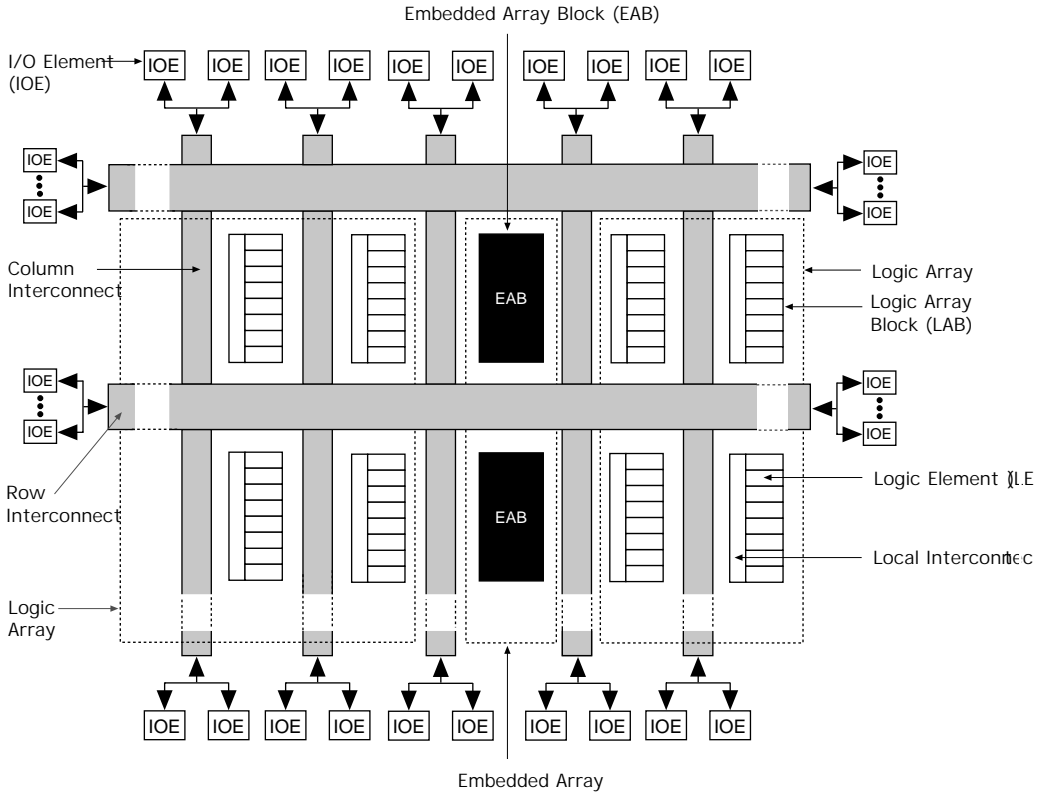
Applications of Embedded - FPGAs

The versatility of Embedded - FPGAs makes them indispensable in numerous fields. In telecommunications.

Details

Product Status	Obsolete
Number of LABs/CLBs	624
Number of Logic Elements/Cells	4992
Total RAM Bits	49152
Number of I/O	333
Number of Gates	257000
Voltage - Supply	2.375V ~ 2.625V
Mounting Type	Surface Mount
Operating Temperature	0°C ~ 70°C (TA)
Package / Case	484-BBGA
Supplier Device Package	484-FBGA (23x23)
Purchase URL	https://www.e-xfl.com/product-detail/intel/ep1k100fc484-3n

Figure 1. ACEX 1K Device Block Diagram



ACEX 1K devices provide six dedicated inputs that drive the flipflops' control inputs and ensure the efficient distribution of high-speed, low-skew (less than 1.0 ns) control signals. These signals use dedicated routing channels that provide shorter delays and lower skews than the FastTrack Interconnect routing structure. Four of the dedicated inputs drive four global signals. These four global signals can also be driven by internal logic, providing an ideal solution for a clock divider or an internally generated asynchronous clear signal that clears many registers in the device.

Embedded Array Block

The EAB is a flexible block of RAM, with registers on the input and output ports, that is used to implement common gate array megafunctions. Because it is large and flexible, the EAB is suitable for functions such as multipliers, vector scalars, and error correction circuits. These functions can be combined in applications such as digital filters and microcontrollers.

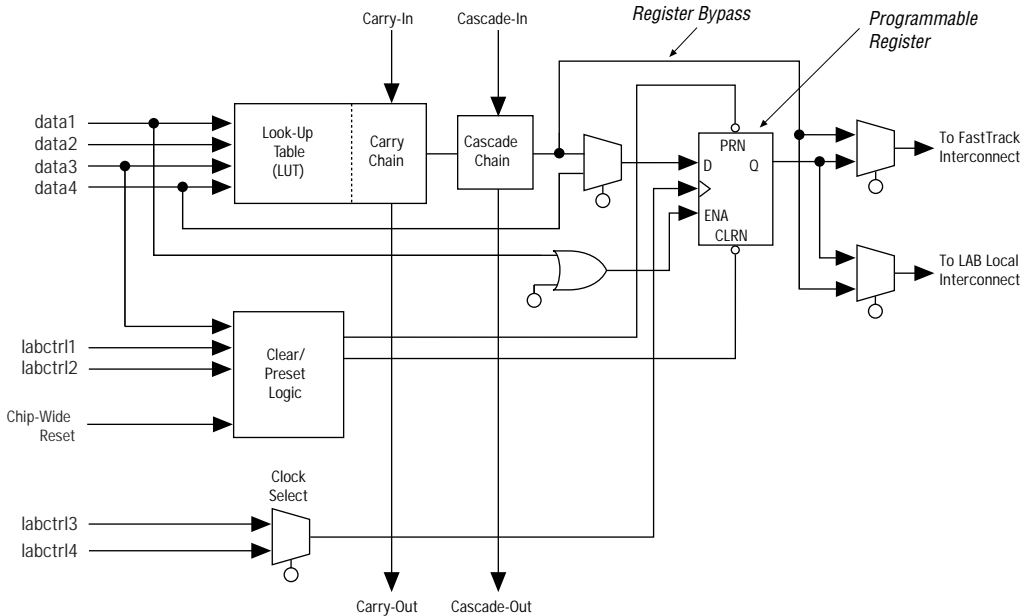
Logic functions are implemented by programming the EAB with a read-only pattern during configuration, thereby creating a large LUT. With LUTs, combinatorial functions are implemented by looking up the results rather than by computing them. This implementation of combinatorial functions can be faster than using algorithms implemented in general logic, a performance advantage that is further enhanced by the fast access times of EABs. The large capacity of EABs enables designers to implement complex functions in a single logic level without the routing delays associated with linked LEs or field-programmable gate array (FPGA) RAM blocks. For example, a single EAB can implement any function with 8 inputs and 16 outputs. Parameterized functions, such as LPM functions, can take advantage of the EAB automatically.

The ACEX 1K enhanced EAB supports dual-port RAM. The dual-port structure is ideal for FIFO buffers with one or two clocks. The ACEX 1K EAB can also support up to 16-bit-wide RAM blocks. The ACEX 1K EAB can act in dual-port or single-port mode. When in dual-port mode, separate clocks may be used for EAB read and write sections, allowing the EAB to be written and read at different rates. It also has separate synchronous clock enable signals for the EAB read and write sections, which allow independent control of these sections.

The EAB can also be used for bidirectional, dual-port memory applications where two ports read or write simultaneously. To implement this type of dual-port memory, two EABs are used to support two simultaneous reads or writes.

Alternatively, one clock and clock enable can be used to control the input registers of the EAB, while a different clock and clock enable control the output registers (see Figure 2).

Figure 8. ACEX 1K Logic Element

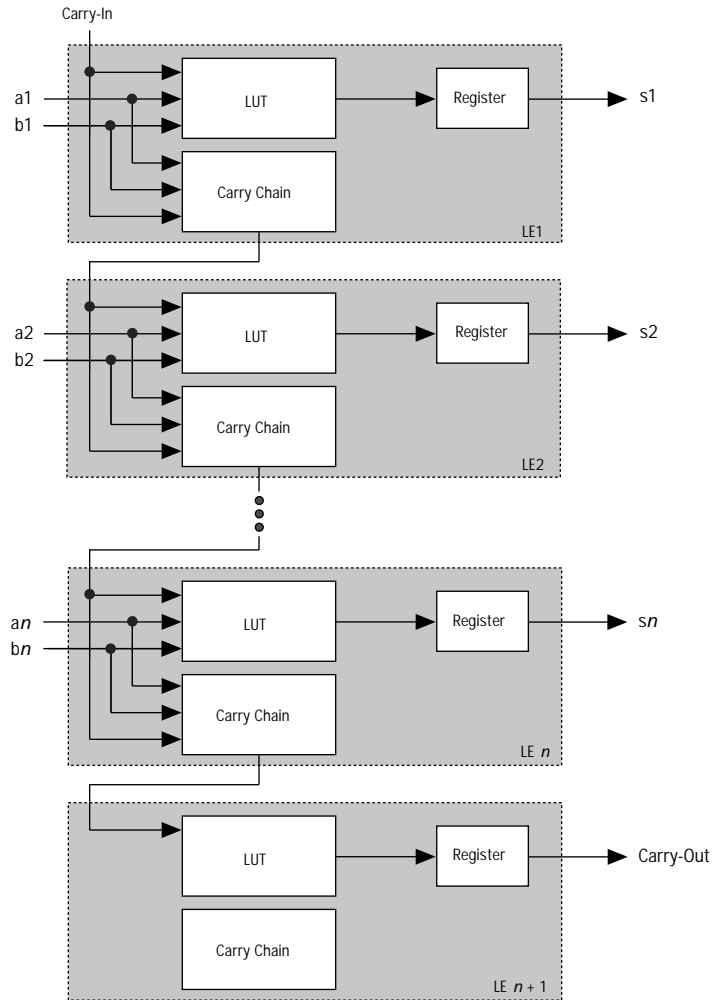


The programmable flipflop in the LE can be configured for D, T, JK, or SR operation. The clock, clear, and preset control signals on the flipflop can be driven by global signals, general-purpose I/O pins, or any internal logic. For combinational functions, the flipflop is bypassed and the LUT's output drives the LE's output.

The LE has two outputs that drive the interconnect: one drives the local interconnect, and the other drives either the row or column FastTrack Interconnect routing structure. The two outputs can be controlled independently. For example, the LUT can drive one output while the register drives the other output. This feature, called register packing, can improve LE utilization because the register and the LUT can be used for unrelated functions.

The ACEX 1K architecture provides two types of dedicated high-speed data paths that connect adjacent LEs without using local interconnect paths: carry chains and cascade chains. The carry chain supports high-speed counters and adders, and the cascade chain implements wide-input functions with minimum delay. Carry and cascade chains connect all LEs in a LAB and all LABs in the same row. Intensive use of carry and cascade chains can reduce routing flexibility. Therefore, the use of these chains should be limited to speed-critical portions of a design.

Figure 9. ACEX 1K Carry Chain Operation (n-Bit Full Adder)



Clearable Counter Mode

The clearable counter mode is similar to the up/down counter mode, but it supports a synchronous clear instead of the up/down control. The clear function is substituted for the cascade-in signal in the up/down counter mode. Two 3-input LUTs are used; one generates the counter data, and the other generates the fast carry bit. Synchronous loading is provided by a 2-to-1 multiplexer. The output of this multiplexer is ANDed with a synchronous clear signal.

Internal Tri-State Emulation

Internal tri-state emulation provides internal tri-states without the limitations of a physical tri-state bus. In a physical tri-state bus, the tri-state buffers' output enable (OE) signals select which signal drives the bus. However, if multiple OE signals are active, contending signals can be driven onto the bus. Conversely, if no OE signals are active, the bus will float. Internal tri-state emulation resolves contending tri-state buffers to a low value and floating buses to a high value, thereby eliminating these problems. The Altera software automatically implements tri-state bus functionality with a multiplexer.

Clear & Preset Logic Control

Logic for the programmable register's clear and preset functions is controlled by the DATA3 LABCTRL1 and LABCTRL2 inputs to the LE. The clear and preset control structure of the LE asynchronously loads signals into a register. Either LABCTRL1 or LABCTRL2 can control the asynchronous clear. Alternatively, the register can be set up so that LABCTRL1 implements an asynchronous load. The data to be loaded is driven to DATA3 when LABCTRL1 is asserted, DATA3 is loaded into the register.

During compilation, the compiler automatically selects the best control signal implementation. Because the clear and preset functions are active-low, the Compiler automatically assigns a logic high to an unused clear or preset.

The clear and preset logic is implemented in one of the following six modes chosen during design entry:

- Asynchronous clear
- Asynchronous preset
- Asynchronous clear and preset
- Asynchronous load with clear
- Asynchronous load with preset
- Asynchronous load without clear or preset

On all ACEX 1K devices, the input path from the I/O pad to the FastTrack Interconnect has a programmable delay element that can be used to guarantee a zero hold time. Depending on the placement of the IOE relative to what it is driving, the designer may choose to turn on the programmable delay to ensure a zero hold time or turn it off to minimize setup time. This feature is used to reduce setup time for complex pin-to-register paths (e.g., PCI designs).

Each IOE selects the clock, clear, clock enable, and output enable controls from a network of I/O control signals called the peripheral control bus. The peripheral control bus uses high-speed drivers to minimize signal skew across devices and provides up to 12 peripheral control signals that can be allocated as follows:

- Up to eight output enable signals
- Up to six clock enable signals
- Up to two clock signals
- Up to two clear signals

If more than six clock-enable or eight output-enable signals are required, each IOE on the device can be controlled by clock enable and output enable signals driven by specific LEs. In addition to the two clock signals available on the peripheral control bus, each IOE can use one of two dedicated clock pins. Each peripheral control signal can be driven by any of the dedicated input pins or the first LE of each LAB in a particular row. In addition, a LE in a different row can drive a column interconnect, which causes a row interconnect to drive the peripheral control signal. The chip-wide reset signal resets all IOE registers, overriding any other control signals.

When a dedicated clock pin drives IOE registers, it can be inverted for all IOEs in the device. All IOEs must use the same sense of the clock. For example, if any IOE uses the inverted clock, all IOEs must use the inverted clock, and no IOE can use the non-inverted clock. However, LEs can still use the true or complement of the clock on an LAB-by-LAB basis.

The incoming signal may be inverted at the dedicated clock pin and will drive all IOEs. For the true and complement of a clock to be used to drive IOEs, drive it into both global clock pins. One global clock pin will supply the true, and the other will supply the complement.

When the true and complement of a dedicated input drives IOE clocks, two signals on the peripheral control bus are consumed, one for each sense of the clock.

Table 12. ClockLock & ClockBoost Parameters for -2 Speed-Grade Devices

Symbol	Parameter	Condition	Min	Typ	Max	Unit
t_R	Input rise time				5	ns
t_F	Input fall time				5	ns
t_{INDUTY}	Input duty cycle		40		60	%
f_{CLK1}	Input clock frequency (ClockBoost clock multiplication factor equals 1)		25		80	MHz
f_{CLK2}	Input clock frequency (ClockBoost clock multiplication factor equals 2)		16		40	MHz
f_{CLKDEV}	Input deviation from user specification in the software (1)				25,000	PPM
$t_{INCLKSTB}$	Input clock stability (measured between adjacent clocks)				100	ps
t_{LOCK}	Time required for ClockLock or ClockBoost to acquire lock (3)				10	∞ s
t_{JITTER}	Jitter on ClockLock or ClockBoost-generated clock (4)	$t_{INCLKSTB} < 100$			250 (4)	ps
		$t_{INCLKSTB} < 50$			200 (4)	ps
$t_{OUTDUTY}$	Duty cycle for ClockLock or ClockBoost-generated clock		40	50	60	%

Notes to tables:

- (1) To implement the ClockLock and ClockBoost circuitry with the Altera software, designers must specify the input frequency. The Altera software tunes the PLL in the ClockLock and ClockBoost circuitry to this frequency. The f_{CLKDEV} parameter specifies how much the incoming clock can differ from the specified frequency during device operation. Simulation does not reflect this parameter.
- (2) Twenty-five thousand parts per million (PPM) equates to 2.5% of input clock period.
- (3) During device configuration, the ClockLock and ClockBoost circuitry is configured before the rest of the device. If the incoming clock is supplied during configuration, the ClockLock and ClockBoost circuitry locks during configuration because the t_{LOCK} value is less than the time required for configuration.
- (4) The t_{JITTER} specification is measured under long-term observation. The maximum value for t_{JITTER} is 200 ps if $t_{INCLKSTB}$ is lower than 50 ps.

I/O Configuration

This section discusses the PCI pull-up clamping diode option, slew-rate control, open-drain output option, and MultiVolt I/O interface for ACEX 1K devices. The PCI pull-up clamping diode, slew-rate control, and open-drain output options are controlled pin-by-pin via Altera software logic options. The MultiVolt I/O interface is controlled by connecting V_{CCIO} to a different voltage than V_{CCINT} . Its effect can be simulated in the Altera software via the Global Project Device Options dialog box (Assign menu).

Table 16. 32-Bit IDCODE for ACEX 1K Devices Note (1)

Device	IDCODE (32 Bits)			
	Version (4 Bits)	Part Number (16 Bits)	Manufacturer's Identity (11 Bits)	1 (1 Bit) (2)
EP1K10	0001	0001 0000 0001 0000	00001101110	1
EP1K30	0001	0001 0000 0011 0000	00001101110	1
EP1K50	0001	0001 0000 0101 0000	00001101110	1
EP1K100	0010	0000 0001 0000 0000	00001101110	1

Notes to tables:

- (1) The most significant bit (MSB) is on the left.
- (2) The least significant bit (LSB) for all JTAG IDCODEs is 1.

ACEX 1K devices include weak pull-up resistors on the JTAG pins.

f For more information, see the following documents:

- Application Note 39 (IEEE Std. 1149.1 (JTAG) Boundary-Scan Testing in Altera Devices)*
- ByteBlasterMV Parallel Port Download Cable Data Sheet*
- BitBlaster Serial Download Cable Data Sheet*
- Jam Programming & Test Language Specification*

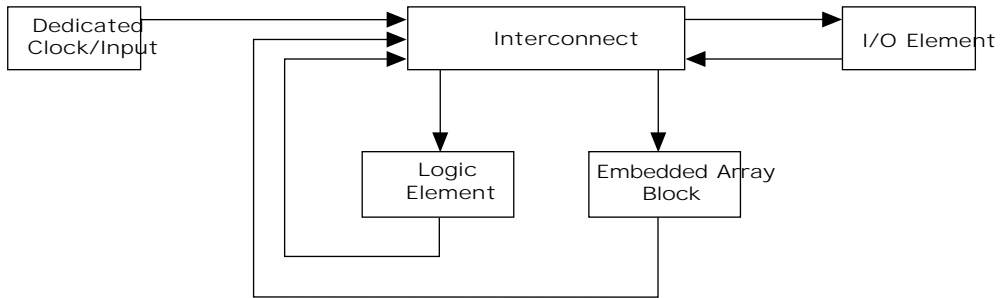
Figure 20 shows the timing requirements for the JTAG signals.

Notes to tables:

- (1) See the *Operating Requirements for Altera Devices Data Sheet*.
- (2) Minimum DC input voltage is -0.5 V. During transitions, the inputs may undershoot to -2.0 V for input currents less than 100 mA and periods shorter than 20 ns.
- (3) Numbers in parentheses are for industrial- and extended-temperature-range devices.
- (4) Maximum V_{CC} rise time is 100 ms, and V_{CC} must rise monotonically.
- (5) All pins, including dedicated inputs, clock, I/O, and JTAG pins, may be driven before V_{CCINT} and V_{CCIO} are powered.
- (6) Typical values are for $T_A = 25^\circ$ C, $V_{CCINT} = 2.5$ V, and $V_{CCIO} = 2.5$ V or 3.3 V.
- (7) These values are specified under the ACEX 1K Recommended Operating Conditions shown in Table 19 on page 46.

Figure 24 shows the overall timing model, which maps the possible paths to and from the various elements of the ACEX 1K device.

Figure 24. ACEX 1K Device Timing Model



Figures 25 through 28 show the delays that correspond to various paths and functions within the LE, IOE, EAB, and bidirectional timing models.

Figure 25. ACEX 1K Device LE Timing Model

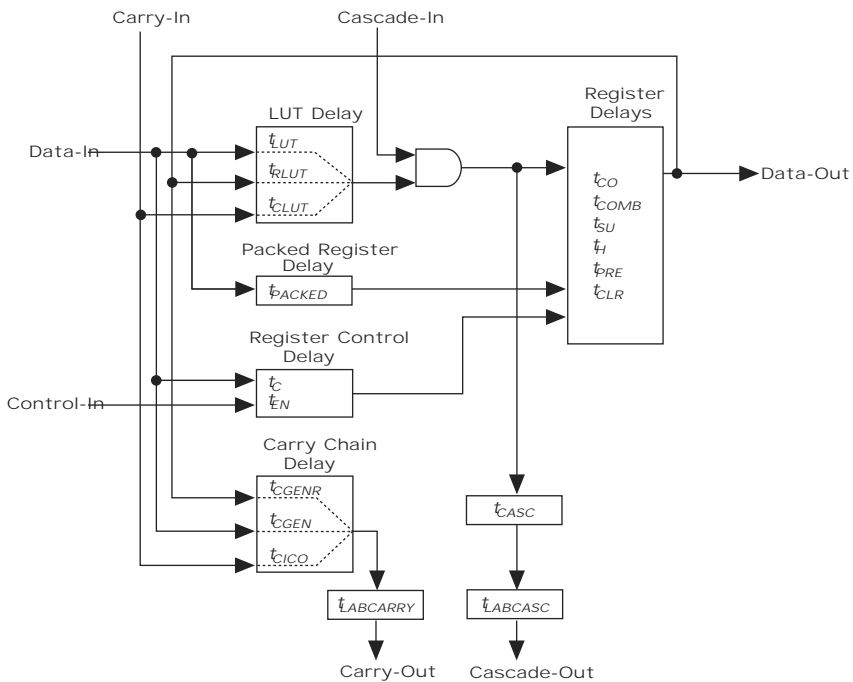


Figure 26. ACEX 1K Device IOE Timing Model

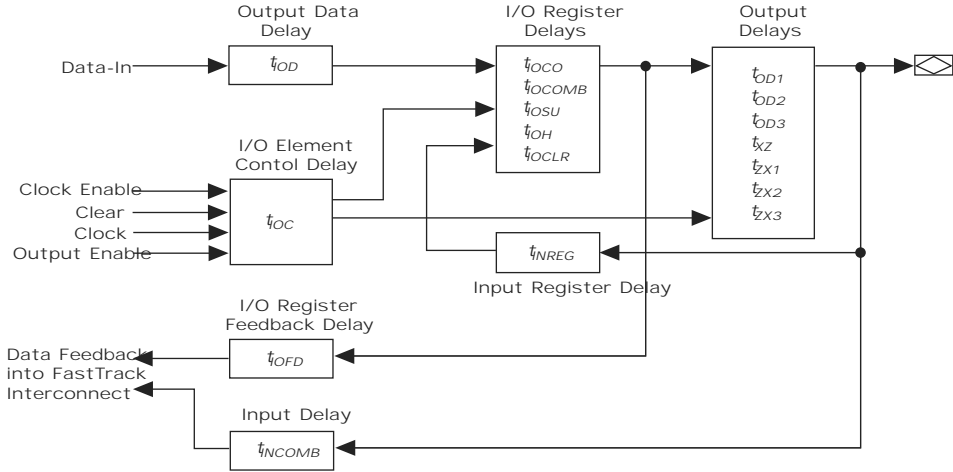


Figure 27. ACEX 1K Device EAB Timing Model

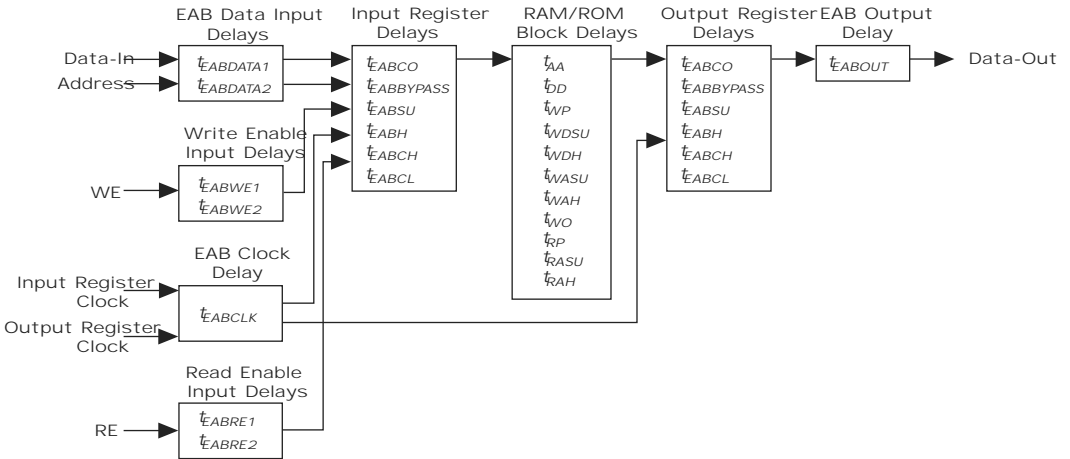


Table 25. EAB Timing Macroparameters Notes (1), (6)

Symbol	Parameter	Conditions
t_{EABAA}	EAB address access delay	
$t_{EABRCCOMB}$	EAB asynchronous read cycle time	
$t_{EABRCREG}$	EAB synchronous read cycle time	
t_{EABWP}	EAB write pulse width	
$t_{EABWCCOMB}$	EAB asynchronous write cycle time	
$t_{EABWCREG}$	EAB synchronous write cycle time	
t_{EABDD}	EAB data-in to data-out valid delay	
$t_{EABDATACO}$	EAB clock-to-output delay when using output registers	
$t_{EABDATASU}$	EAB data/address setup time before clock when using input register	
$t_{EABDATAH}$	EAB data/address hold time after clock when using input register	
$t_{EABWESU}$	EAB WEsetup time before clock when using input register	
t_{EABWEH}	EAB WEhold time after clock when using input register	
$t_{EABWDSU}$	EAB data setup time before falling edge of write pulse when not using input registers	
t_{EABWDH}	EAB data hold time after falling edge of write pulse when not using input registers	
$t_{EABWASU}$	EAB address setup time before rising edge of write pulse when not using input registers	
t_{EABWAH}	EAB address hold time after falling edge of write pulse when not using input registers	
t_{EABWO}	EAB write enable to data output valid delay	

