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Understanding Embedded - FPGAs (Field Programmable Gate Array)

Embedded - FPGAs, or Field Programmable Gate Arrays, are advanced integrated circuits that offer unparalleled flexibility and performance for digital systems. Unlike traditional fixed-function logic devices, FPGAs can be programmed and reprogrammed to execute a wide array of logical operations, enabling customized functionality tailored to specific applications. This reprogrammability allows developers to iterate designs quickly and implement complex functions without the need for custom hardware.

Applications of Embedded - FPGAs

The versatility of Embedded - FPGAs makes them indispensable in numerous fields. In telecommunications.

Details	
Product Status	Obsolete
Number of LABs/CLBs	216
Number of Logic Elements/Cells	1728
Total RAM Bits	24576
Number of I/O	102
Number of Gates	119000
Voltage - Supply	2.375V ~ 2.625V
Mounting Type	Surface Mount
Operating Temperature	-40°C ~ 85°C (TA)
Package / Case	144-LQFP
Supplier Device Package	144-TQFP (20x20)
Purchase URL	https://www.e-xfl.com/product-detail/intel/ep1k30ti144-2

General Description

Altera® ACEX 1K devices provide a die-efficient, low-cost architecture by combining look-up table (LUT) architecture with EABs. LUT-based logic provides optimized performance and efficiency for data-path, register intensive, mathematical, or digital signal processing (DSP) designs, while EABs implement RAM, ROM, dual-port RAM, or first-in first-out (FIFO) functions. These elements make ACEX 1K suitable for complex logic functions and memory functions such as digital signal processing, wide data-path manipulation, data transformation and microcontrollers, as required in high-performance communications applications. Based on reconfigurable CMOS SRAM elements, the ACEX 1K architecture incorporates all features necessary to implement common gate array megafunctions, along with a high pin count to enable an effective interface with system components. The advanced process and the low voltage requirement of the 2.5-V core allow ACEX 1K devices to meet the requirements of low-cost, high-volume applications ranging from DSL modems to low-cost switches.

The ability to reconfigure ACEX 1K devices enables complete testing prior to shipment and allows the designer to focus on simulation and design verification. ACEX 1K device reconfigurability eliminates inventory management for gate array designs and test vector generation for fault coverage.

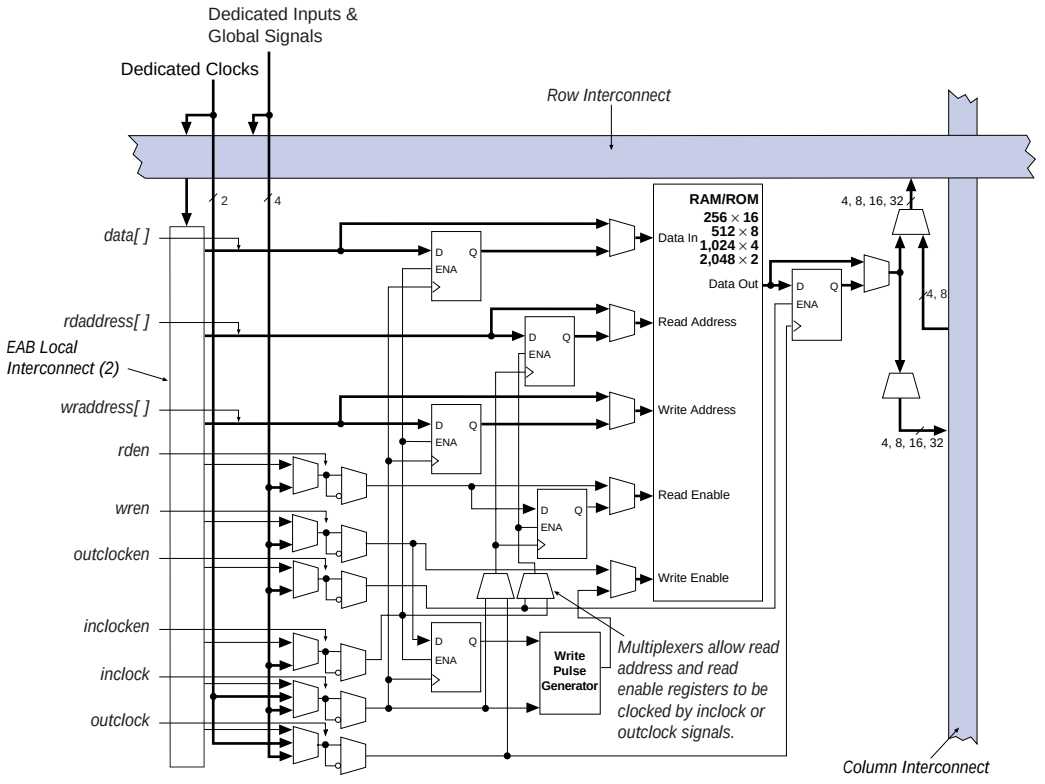
Table 4 shows ACEX 1K device performance for some common designs. All performance results were obtained with Synopsys DesignWare or LPM functions. Special design techniques are not required to implement the applications; the designer simply infers or instantiates a function in a Verilog HDL, VHDL, Altera Hardware Description Language (AHDL), or schematic design file.

Table 4. ACEX 1K Device Performance

Application	Resources Used		Performance			
	LEs	EABs	Speed Grade			Units
			-1	-2	-3	
16-bit loadable counter	16	0	285	232	185	MHz
16-bit accumulator	16	0	285	232	185	MHz
16-to-1 multiplexer *(1)*	10	0	3.5	4.5	6.6	ns
16-bit multiplier with 3-stage pipeline*(2)*	592	0	156	131	93	MHz
256 × 16 RAM read cycle speed *(2)*	0	1	278	196	143	MHz
256 × 16 RAM write cycle speed *(2)*	0	1	185	143	111	MHz

Notes:
(1) This application uses combinatorial inputs and outputs.
(2) This application uses registered inputs and outputs.

Figure 2. ACEX 1K Device in Dual-Port RAM Mode *Note (1)*



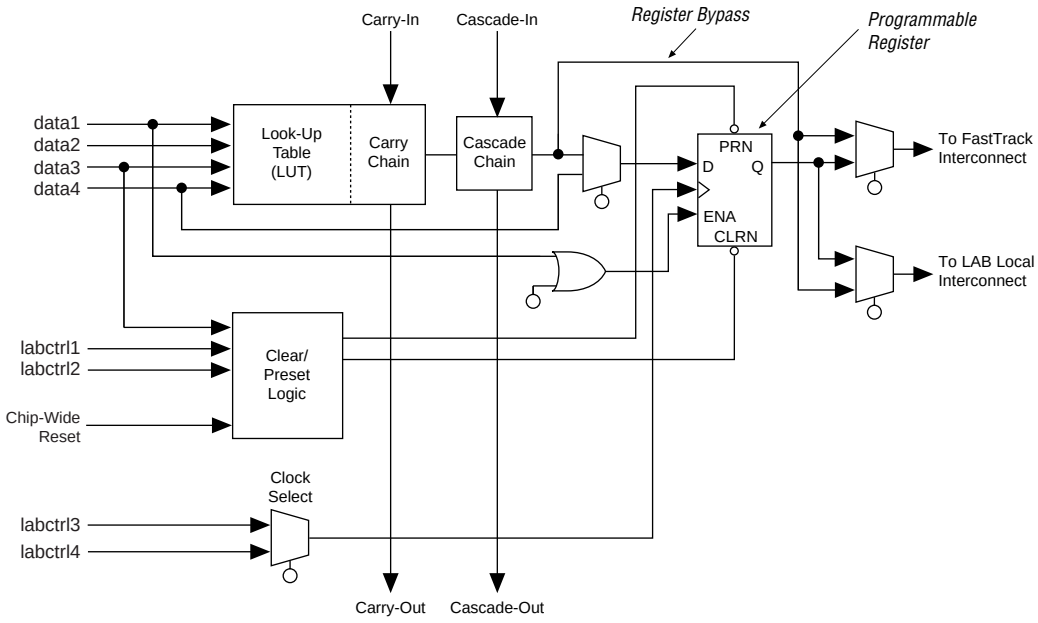
Notes:

(1) All registers can be asynchronously cleared by EAB local interconnect signals, global signals, or the chip-wide reset.

(2) EP1K10, EP1K30, and EP1K50 devices have 88 EAB local interconnect channels; EP1K100 devices have 104 EAB local interconnect channels.

The EAB can use Altera megafunctions to implement dual-port RAM applications where both ports can read or write, as shown in Figure 3. The ACEX 1K EAB can also be used in a single-port mode (see Figure 4).

Figure 8. ACEX 1K Logic Element



The programmable flipflop in the LE can be configured for D, T, JK, or SR operation. The clock, clear, and preset control signals on the flipflop can be driven by global signals, general-purpose I/O pins, or any internal logic. For combinatorial functions, the flipflop is bypassed and the LUT's output drives the LE's output.

The LE has two outputs that drive the interconnect: one drives the local interconnect, and the other drives either the row or column FastTrack Interconnect routing structure. The two outputs can be controlled independently. For example, the LUT can drive one output while the register drives the other output. This feature, called register packing, can improve LE utilization because the register and the LUT can be used for unrelated functions.

The ACEX 1K architecture provides two types of dedicated high-speed data paths that connect adjacent LEs without using local interconnect paths: carry chains and cascade chains. The carry chain supports high-speed counters and adders, and the cascade chain implements wide-input functions with minimum delay. Carry and cascade chains connect all LEs in a LAB and all LABs in the same row. Intensive use of carry and cascade chains can reduce routing flexibility. Therefore, the use of these chains should be limited to speed-critical portions of a design.

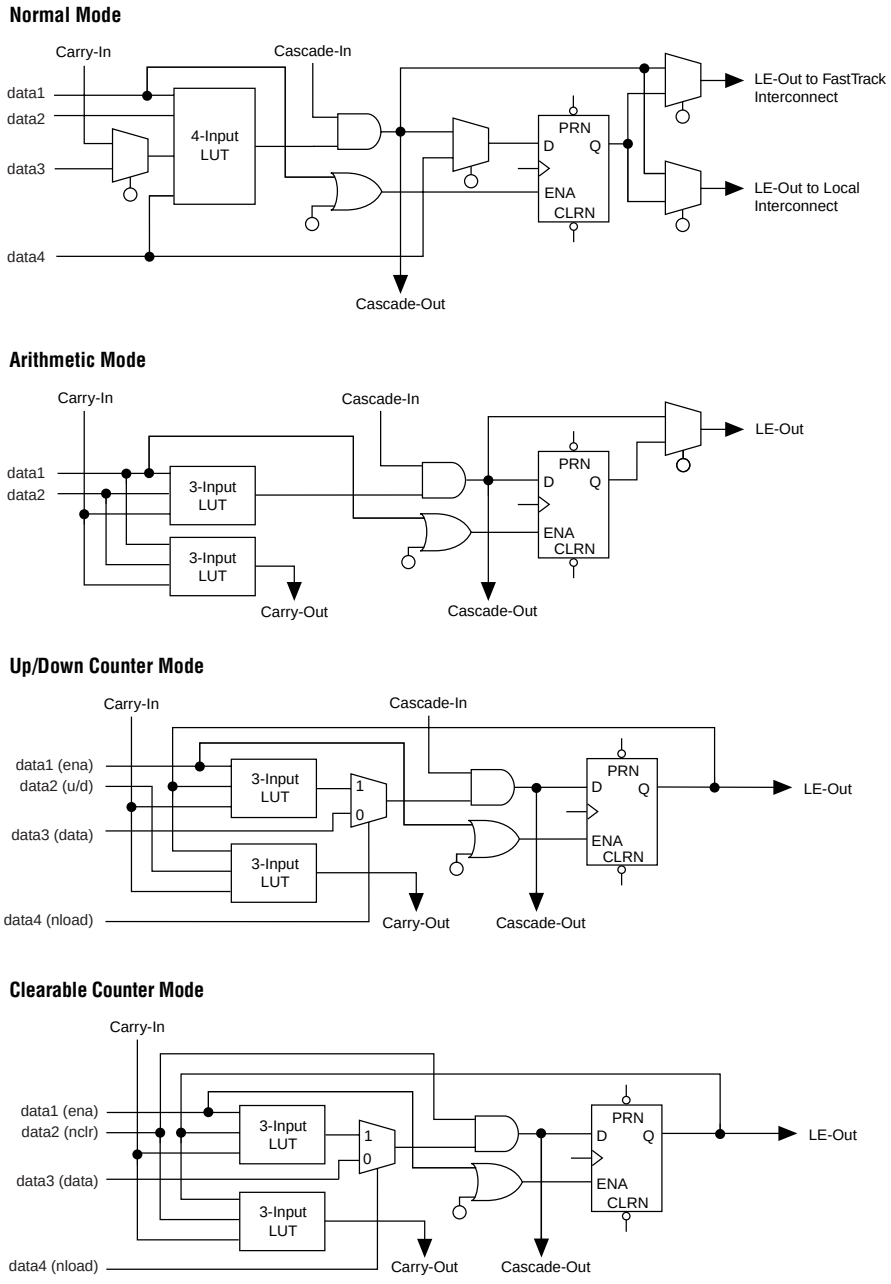
Carry Chain

The carry chain provides a very fast (as low as 0.2 ns) carry-forward function between LEs. The carry-in signal from a lower-order bit drives forward into the higher-order bit via the carry chain, and feeds into both the LUT and the next portion of the carry chain. This feature allows the ACEX 1K architecture to efficiently implement high-speed counters, adders, and comparators of arbitrary width. Carry chain logic can be created automatically by the compiler during design processing, or manually by the designer during design entry. Parameterized functions, such as LPM and DesignWare functions, automatically take advantage of carry chains.

Carry chains longer than eight LEs are automatically implemented by linking LABs together. For enhanced fitting, a long carry chain skips alternate LABs in a row. A carry chain longer than one LAB skips either from even-numbered LAB to even-numbered LAB, or from odd-numbered LAB to odd-numbered LAB. For example, the last LE of the first LAB in a row carries to the first LE of the third LAB in the row. The carry chain does not cross the EAB at the middle of the row. For instance, in the EP1K50 device, the carry chain stops at the eighteenth LAB, and a new carry chain begins at the nineteenth LAB.

Figure 9 shows how an *n*-bit full adder can be implemented in *n* + 1 LEs with the carry chain. One portion of the LUT generates the sum of two bits using the input signals and the carry-in signal; the sum is routed to the output of the LE. The register can be bypassed for simple adders or used for an accumulator function. Another portion of the LUT and the carry chain logic generates the carry-out signal, which is routed directly to the carry-in signal of the next-higher-order bit. The final carry-out signal is routed to an LE, where it can be used as a general-purpose signal.

Figure 11. ACEX 1K LE Operating Modes



Normal Mode

The normal mode is suitable for general logic applications and wide decoding functions that can take advantage of a cascade chain. In normal mode, four data inputs from the LAB local interconnect and the carry-in are inputs to a 4-input LUT. The compiler automatically selects the carry-in or the DATA3 signal as one of the inputs to the LUT. The LUT output can be combined with the cascade-in signal to form a cascade chain through the cascade-out signal. Either the register or the LUT can be used to drive both the local interconnect and the FastTrack Interconnect routing structure at the same time.

The LUT and the register in the LE can be used independently (register packing). To support register packing, the LE has two outputs; one drives the local interconnect, and the other drives the FastTrack Interconnect routing structure. The DATA4 signal can drive the register directly, allowing the LUT to compute a function that is independent of the registered signal; a 3-input function can be computed in the LUT, and a fourth independent signal can be registered. Alternatively, a 4-input function can be generated, and one of the inputs to this function can be used to drive the register. The register in a packed LE can still use the clock enable, clear, and preset signals in the LE. In a packed LE, the register can drive the FastTrack Interconnect routing structure while the LUT drives the local interconnect, or vice versa.

Arithmetic Mode

The arithmetic mode offers two 3-input LUTs that are ideal for implementing adders, accumulators, and comparators. One LUT computes a 3-input function; the other generates a carry output. As shown in Figure 11, the first LUT uses the carry-in signal and two data inputs from the LAB local interconnect to generate a combinatorial or registered output. For example, in an adder, this output is the sum of three signals: a, b, and carry-in. The second LUT uses the same three signals to generate a carry-out signal, thereby creating a carry chain. The arithmetic mode also supports simultaneous use of the cascade chain.

Up/Down Counter Mode

The up / down counter mode offers counter enable, clock enable, synchronous up / down control, and data loading options. These control signals are generated by the data inputs from the LAB local interconnect, the carry-in signal, and output feedback from the programmable register. Two 3-input LUTs are used; one generates the counter data, and the other generates the fast carry bit. A 2-to-1 multiplexer provides synchronous loading. Data can also be loaded asynchronously with the clear and preset register control signals without using the LUT resources.

Clearable Counter Mode

The clearable counter mode is similar to the up / down counter mode, but it supports a synchronous clear instead of the up / down control. The clear function is substituted for the cascade-in signal in the up / down counter mode. Two 3-input LUTs are used; one generates the counter data, and the other generates the fast carry bit. Synchronous loading is provided by a 2-to-1 multiplexer. The output of this multiplexer is ANDed with a synchronous clear signal.

Internal Tri-State Emulation

Internal tri-state emulation provides internal tri-states without the limitations of a physical tri-state bus. In a physical tri-state bus, the tri-state buffers' output enable (OE) signals select which signal drives the bus. However, if multiple OE signals are active, contending signals can be driven onto the bus. Conversely, if no OE signals are active, the bus will float. Internal tri-state emulation resolves contending tri-state buffers to a low value and floating buses to a high value, thereby eliminating these problems. The Altera software automatically implements tri-state bus functionality with a multiplexer.

Clear & Preset Logic Control

Logic for the programmable register's clear and preset functions is controlled by the DATA3, LABCTRL1, and LABCTRL2 inputs to the LE. The clear and preset control structure of the LE asynchronously loads signals into a register. Either LABCTRL1 or LABCTRL2 can control the asynchronous clear. Alternatively, the register can be set up so that LABCTRL1 implements an asynchronous load. The data to be loaded is driven to DATA3; when LABCTRL1 is asserted, DATA3 is loaded into the register.

During compilation, the compiler automatically selects the best control signal implementation. Because the clear and preset functions are active-low, the Compiler automatically assigns a logic high to an unused clear or preset.

The clear and preset logic is implemented in one of the following six modes chosen during design entry:

- Asynchronous clear
- Asynchronous preset
- Asynchronous clear and preset
- Asynchronous load with clear
- Asynchronous load with preset
- Asynchronous load without clear or preset

Asynchronous Clear

The flipflop can be cleared by either LABCTRL1 or LABCTRL2. In this mode, the preset signal is tied to VCC to deactivate it.

Asynchronous Preset

An asynchronous preset is implemented as an asynchronous load, or with an asynchronous clear. If DATA3 is tied to VCC, asserting LABCTRL1 asynchronously loads a one into the register. Alternatively, the Altera software can provide preset control by using the clear and inverting the register's input and output. Inversion control is available for the inputs to both LEs and IOEs. Therefore, if a register is preset by only one of the two LABCTRL signals, the DATA3 input is not needed and can be used for one of the LE operating modes.

Asynchronous Preset & Clear

When implementing asynchronous clear and preset, LABCTRL1 controls the preset, and LABCTRL2 controls the clear. DATA3 is tied to VCC, so that asserting LABCTRL1 asynchronously loads a one into the register, effectively presetting the register. Asserting LABCTRL2 clears the register.

Asynchronous Load with Clear

When implementing an asynchronous load in conjunction with the clear, LABCTRL1 implements the asynchronous load of DATA3 by controlling the register preset and clear. LABCTRL2 implements the clear by controlling the register clear; LABCTRL2 does not have to feed the preset circuits.

Asynchronous Load with Preset

When implementing an asynchronous load in conjunction with preset, the Altera software provides preset control by using the clear and inverting the input and output of the register. Asserting LABCTRL2 presets the register, while asserting LABCTRL1 loads the register. The Altera software inverts the signal that drives DATA3 to account for the inversion of the register's output.

Asynchronous Load without Preset or Clear

When implementing an asynchronous load without preset or clear, LABCTRL1 implements the asynchronous load of DATA3 by controlling the register preset and clear.

For improved routing, the row interconnect consists of a combination of full-length and half-length channels. The full-length channels connect to all LABs in a row; the half-length channels connect to the LABs in half of the row. The EAB can be driven by the half-length channels in the left half of the row and by the full-length channels. The EAB drives out to the full-length channels. In addition to providing a predictable, row-wide interconnect, this architecture provides increased routing resources. Two neighboring LABs can be connected using a half-row channel, thereby saving the other half of the channel for the other half of the row.

Table 6 summarizes the FastTrack Interconnect routing structure resources available in each ACEX 1K device.

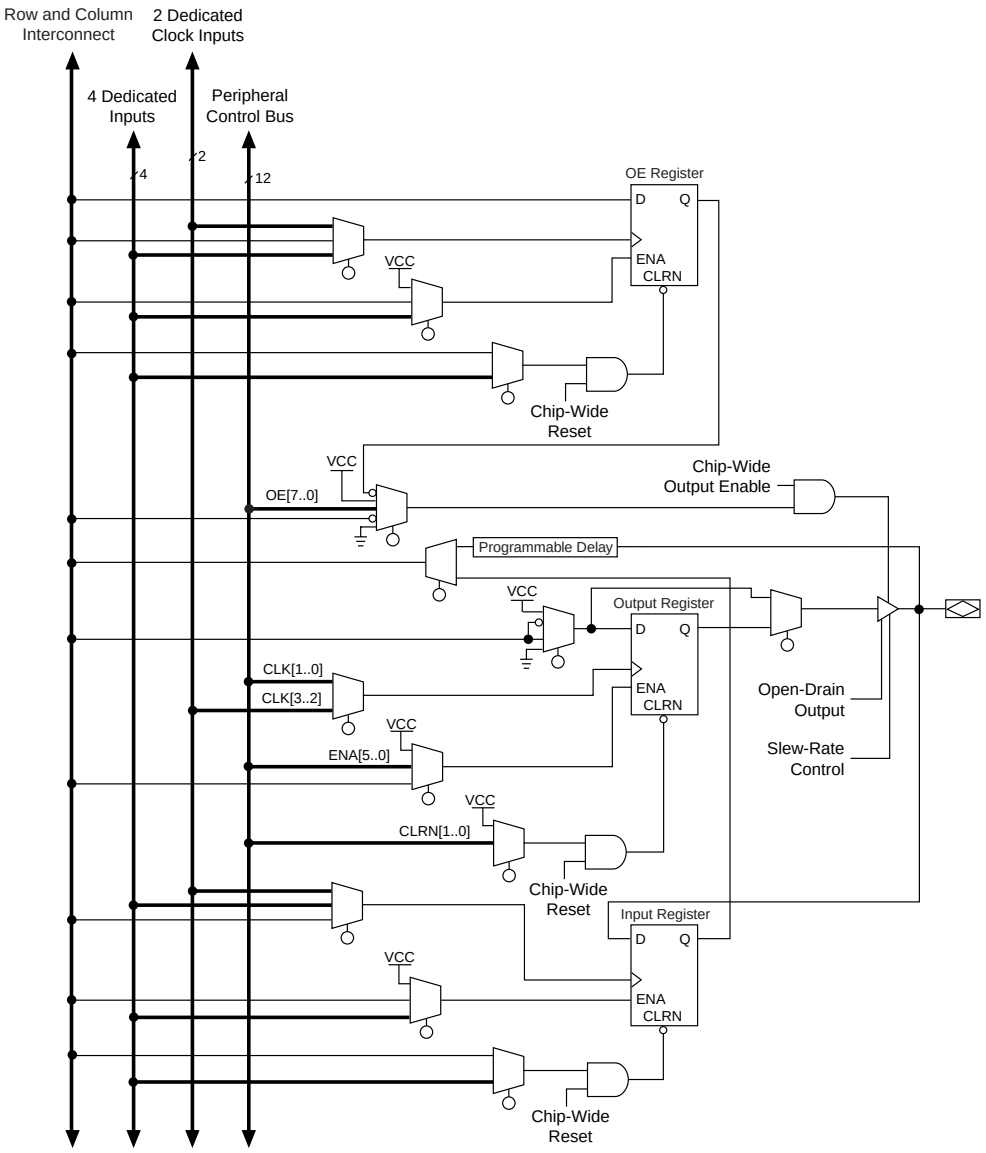
Table 6. ACEX 1K FastTrack Interconnect Resources

Device	Rows	Channels per Row	Columns	Channels per Column
EP1K10	3	144	24	24
EP1K30	6	216	36	24
EP1K50	10	216	36	24
EP1K100	12	312	52	24

In addition to general-purpose I/O pins, ACEX 1K devices have six dedicated input pins that provide low-skew signal distribution across the device. These six inputs can be used for global clock, clear, preset, and peripheral output-enable and clock-enable control signals. These signals are available as control signals for all LABs and IOEs in the device. The dedicated inputs can also be used as general-purpose data inputs because they can feed the local interconnect of each LAB in the device.

Figure 14 shows the interconnection of adjacent LABs and EABs, with row, column, and local interconnects, as well as the associated cascade and carry chains. Each LAB is labeled according to its location: a letter represents the row and a number represents the column. For example, LAB B3 is in row B, column 3.

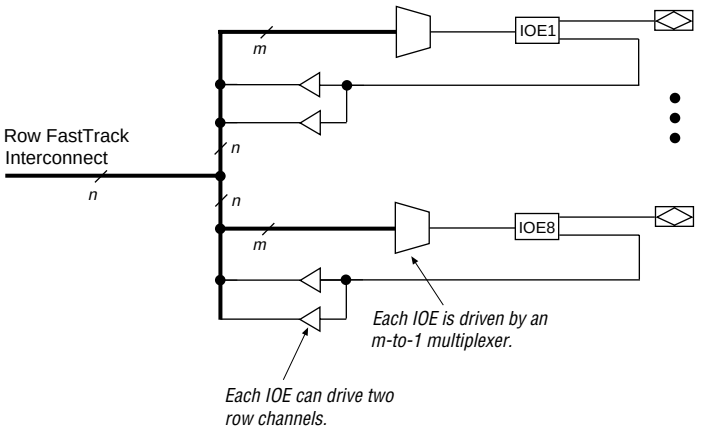
Figure 15. ACEX 1K Bidirectional I/O Registers



Row-to-IOE Connections

When an IOE is used as an input signal, it can drive two separate row channels. The signal is accessible by all LEs within that row. When an IOE is used as an output, the signal is driven by a multiplexer that selects a signal from the row channels. Up to eight IOEs connect to each side of each row channel (see Figure 16).

Figure 16. ACEX 1K Row-to-IOE Connections *Note (1)*



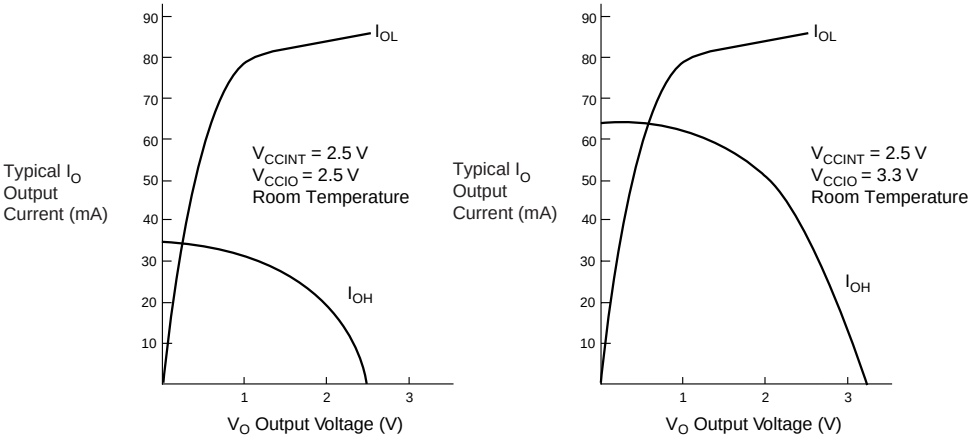
Note:
(1) The values for *m* and *n* are shown in Table 8.

Table 8 lists the ACEX 1K row-to-IOE interconnect resources.

Table 8. ACEX 1K Row-to-IOE Interconnect Resources

Device	Channels per Row *(n)*	Row Channels per Pin *(m)*
EP1K10	144	18
EP1K30	216	27
EP1K50	216	27
EP1K100	312	39

Figure 23. Output Drive Characteristics of ACEX 1K Devices



Timing Model

The continuous, high-performance FastTrack Interconnect routing resources ensure accurate simulation and timing analysis as well as predictable performance. This predictable performance contrasts with that of FPGAs, which use a segmented connection scheme and, therefore, have an unpredictable performance.

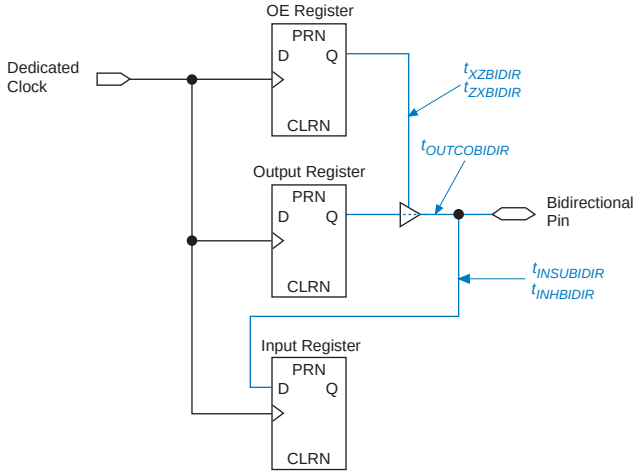
Device performance can be estimated by following the signal path from a source, through the interconnect, to the destination. For example, the registered performance between two LEs on the same row can be calculated by adding the following parameters:

- LE register clock-to-output delay (t_{CO})
- Interconnect delay ($t_{SAMEROW}$)
- LE look-up table delay (t_{LUT})
- LE register setup time (t_{SU})

The routing delay depends on the placement of the source and destination LEs. A more complex registered path may involve multiple combinatorial LEs between the source and destination LEs.

Timing simulation and delay prediction are available with the simulator and Timing Analyzer, or with industry-standard EDA tools. The Simulator offers both pre-synthesis functional simulation to evaluate logic design accuracy and post-synthesis timing simulation with 0.1-ns resolution. The Timing Analyzer provides point-to-point timing delay information, setup and hold time analysis, and device-wide performance analysis.

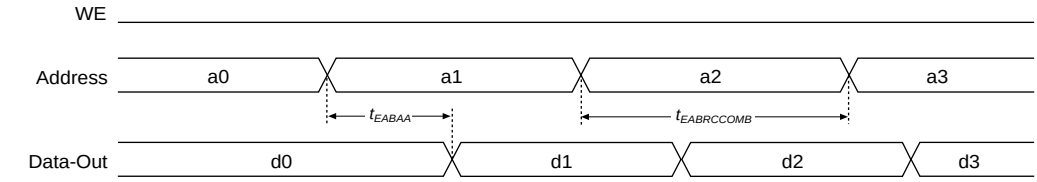
Figure 28. Synchronous Bidirectional Pin External Timing Model



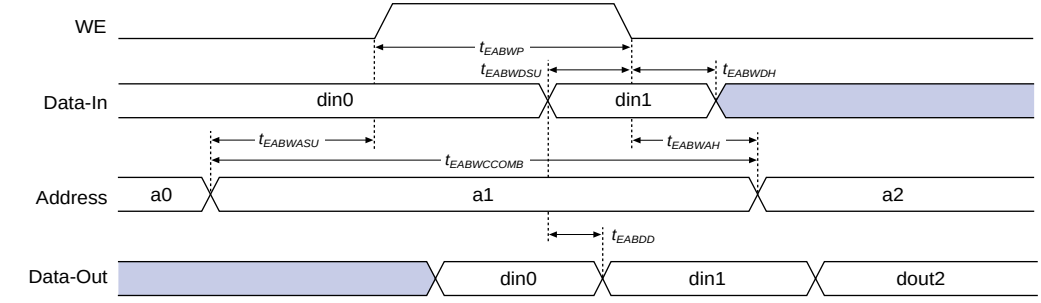
Tables 29 and 30 show the asynchronous and synchronous timing waveforms, respectively, for the EAB macroparameters in Table 24.

Figure 29. EAB Asynchronous Timing Waveforms

EAB Asynchronous Read



EAB Asynchronous Write



Tables 30 through 36 show EP1K10 device internal and external timing parameters.

Table 30. EP1K10 Device LE Timing Microparameters *Note (1)*

Symbol	Speed Grade						Unit
	-1		-2		-3		
	Min	Max	Min	Max	Min	Max	
t_{LUT}		0.7		0.8		1.1	ns
t_{CLUT}		0.5		0.6		0.8	ns
t_{RLUT}		0.6		0.7		1.0	ns
t_{PACKED}		0.4		0.4		0.5	ns
t_{EN}		0.9		1.0		1.3	ns
t_{CICO}		0.1		0.1		0.2	ns
t_{CGEN}		0.4		0.5		0.7	ns
t_{CGENR}		0.1		0.1		0.2	ns
t_{CASC}		0.7		0.9		1.1	ns
t_{C}		1.1		1.3		1.7	ns
t_{CO}		0.5		0.7		0.9	ns
t_{COMB}		0.4		0.5		0.7	ns
t_{SU}	0.7		0.8		1.0		ns
t_{H}	0.9		1.0		1.1		ns
t_{PRE}		0.8		1.0		1.4	ns
t_{CLR}		0.9		1.0		1.4	ns
t_{CH}	2.0		2.5		2.5		ns
t_{CL}	2.0		2.5		2.5		ns

Table 37. EP1K30 Device LE Timing Microparameters (Part 2 of 2) *Note (1)*

Symbol	Speed Grade						Unit
	-1		-2		-3		
	Min	Max	Min	Max	Min	Max	
t_{COMB}		0.4		0.4		0.6	ns
t_{SU}	0.4		0.6		0.6		ns
t_{H}	0.7		1.0		1.3		ns
t_{PRE}		0.8		0.9		1.2	ns
t_{CLR}		0.8		0.9		1.2	ns
t_{CH}	2.0		2.5		2.5		ns
t_{CL}	2.0		2.5		2.5		ns

Table 38. EP1K30 Device IOE Timing Microparameters *Note (1)*

Symbol	Speed Grade						Unit
	-1		-2		-3		
	Min	Max	Min	Max	Min	Max	
t_{IOD}		2.4		2.8		3.8	ns
t_{IOC}		0.3		0.4		0.5	ns
t_{IOCO}		1.0		1.1		1.6	ns
t_{IOCOMB}		0.0		0.0		0.0	ns
t_{IOSU}	1.2		1.4		1.9		ns
t_{IOH}	0.3		0.4		0.5		ns
t_{IOCLR}		1.0		1.1		1.6	ns
t_{OD1}		1.9		2.3		3.0	ns
t_{OD2}		1.4		1.8		2.5	ns
t_{OD3}		4.4		5.2		7.0	ns
t_{XZ}		2.7		3.1		4.3	ns
t_{ZX1}		2.7		3.1		4.3	ns
t_{ZX2}		2.2		2.6		3.8	ns
t_{ZX3}		5.2		6.0		8.3	ns
t_{INREG}		3.4		4.1		5.5	ns
t_{IOFD}		0.8		1.3		2.4	ns
t_{INCOMB}		0.8		1.3		2.4	ns

Table 44. EP1K50 Device LE Timing Microparameters (Part 2 of 2) Note (1)							
Symbol	Speed Grade						Unit
	-1		-2		-3		
	Min	Max	Min	Max	Min	Max	
t_{CO}		0.6		0.6		0.7	ns
t_{COMB}		0.3		0.4		0.5	ns
t_{SU}	0.5		0.6		0.7		ns
t_{H}	0.5		0.6		0.8		ns
t_{PRE}		0.4		0.5		0.7	ns
t_{CLR}		0.8		1.0		1.2	ns
t_{CH}	2.0		2.5		3.0		ns
t_{CL}	2.0		2.5		3.0		ns

Table 45. EP1K50 Device IOE Timing Microparameters Note (1)							
Symbol	Speed Grade						Unit
	-1		-2		-3		
	Min	Max	Min	Max	Min	Max	
t_{IOD}		1.3		1.3		1.9	ns
t_{IOC}		0.3		0.4		0.4	ns
t_{IOCO}		1.7		2.1		2.6	ns
t_{IOCOMB}		0.5		0.6		0.8	ns
t_{IOSU}	0.8		1.0		1.3		ns
t_{IOH}	0.4		0.5		0.6		ns
t_{IOCLR}		0.2		0.2		0.4	ns
t_{OD1}		1.2		1.2		1.9	ns
t_{OD2}		0.7		0.8		1.7	ns
t_{OD3}		2.7		3.0		4.3	ns
t_{XZ}		4.7		5.7		7.5	ns
t_{ZX1}		4.7		5.7		7.5	ns
t_{ZX2}		4.2		5.3		7.3	ns
t_{ZX3}		6.2		7.5		9.9	ns
t_{INREG}		3.5		4.2		5.6	ns
t_{IOFD}		1.1		1.3		1.8	ns
t_{INCOMB}		1.1		1.3		1.8	ns

Table 52. EP1K100 Device IOE Timing Microparameters *Note (1)*							
Symbol	**Speed Grade**						**Unit**
	-1		**-2**		**-3**		
	Min	**Max**	**Min**	**Max**	**Min**	**Max**	
t_{IOD}		1.7		2.0		2.6	ns
t_{IOC}		0.0		0.0		0.0	ns
t_{IOCO}		1.4		1.6		2.1	ns
t_{IOCOMB}		0.5		0.7		0.9	ns
t_{IOSU}	0.8		1.0		1.3		ns
t_{IOH}	0.7		0.9		1.2		ns
t_{IOCLR}		0.5		0.7		0.9	ns
t_{OD1}		3.0		4.2		5.6	ns
t_{OD2}		3.0		4.2		5.6	ns
t_{OD3}		4.0		5.5		7.3	ns
t_{XZ}		3.5		4.6		6.1	ns
t_{ZX1}		3.5		4.6		6.1	ns
t_{ZX2}		3.5		4.6		6.1	ns
t_{ZX3}		4.5		5.9		7.8	ns
t_{INREG}		2.0		2.6		3.5	ns
t_{IOFD}		0.5		0.8		1.2	ns
t_{INCOMB}		0.5		0.8		1.2	ns

Table 53. EP1K100 Device EAB Internal Microparameters *Note (1)*							
Symbol	**Speed Grade**						**Unit**
	-1		**-2**		**-3**		
	Min	**Max**	**Min**	**Max**	**Min**	**Max**	
$t_{EABDATA1}$		1.5		2.0		2.6	ns
$t_{EABDATA1}$		0.0		0.0		0.0	ns
t_{EABWE1}		1.5		2.0		2.6	ns
t_{EABWE2}		0.3		0.4		0.5	ns
t_{EABRE1}		0.3		0.4		0.5	ns
t_{EABRE2}		0.0		0.0		0.0	ns
t_{EABCLK}		0.0		0.0		0.0	ns
t_{EABCO}		0.3		0.4		0.5	ns
$t_{EABBYPASS}$		0.1		0.1		0.2	ns
t_{EABSU}	0.8		1.0		1.4		ns
t_{EABH}	0.1		0.1		0.2		ns
t_{EABCLR}	0.3		0.4		0.5		ns
t_{AA}		4.0		5.1		6.6	ns
t_{WP}	2.7		3.5		4.7		ns
t_{RP}	1.0		1.3		1.7		ns
t_{WDSU}	1.0		1.3		1.7		ns
t_{WDH}	0.2		0.2		0.3		ns
t_{WASU}	1.6		2.1		2.8		ns
t_{WAH}	1.6		2.1		2.8		ns
t_{RASU}	3.0		3.9		5.2		ns
t_{RAH}	0.1		0.1		0.2		ns
t_{WO}		1.5		2.0		2.6	ns
t_{DD}		1.5		2.0		2.6	ns
t_{EABOUT}		0.2		0.3		0.3	ns
t_{EABCH}	1.5		2.0		2.5		ns
t_{EABCL}	2.7		3.5		4.7		ns

Table 57. EP1K100 External Bidirectional Timing Parameters *Notes (1), (2)*

Symbol	Speed Grade						Unit
	-1		-2		-3		
	Min	Max	Min	Max	Min	Max	
$t_{INSUBIDIR}$ (3)	1.7		2.5		3.3		ns
$t_{INHBIDIR}$ (3)	0.0		0.0		0.0		ns
$t_{INSUBIDIR}$ (4)	2.0		2.8		–		ns
$t_{INHBIDIR}$ (4)	0.0		0.0		–		ns
$t_{OUTCOBIDIR}$ (3)	2.0	5.2	2.0	6.9	2.0	9.1	ns
$t_{XZBIDIR}$ (3)		5.6		7.5		10.1	ns
$t_{ZXBIDIR}$ (3)		5.6		7.5		10.1	ns
$t_{OUTCOBIDIR}$ (4)	0.5	3.0	0.5	4.6	–	–	ns
$t_{XZBIDIR}$ (4)		4.6		6.5		–	ns
$t_{ZXBIDIR}$ (4)		4.6		6.5		–	ns

Notes to tables:

(1) All timing parameters are described in Tables 22 through 29 in this data sheet.
(2) These parameters are specified by characterization.
(3) This parameter is measured without the use of the ClockLock or ClockBoost circuits.
(4) This parameter is measured with the use of the ClockLock or ClockBoost circuits.

Power Consumption

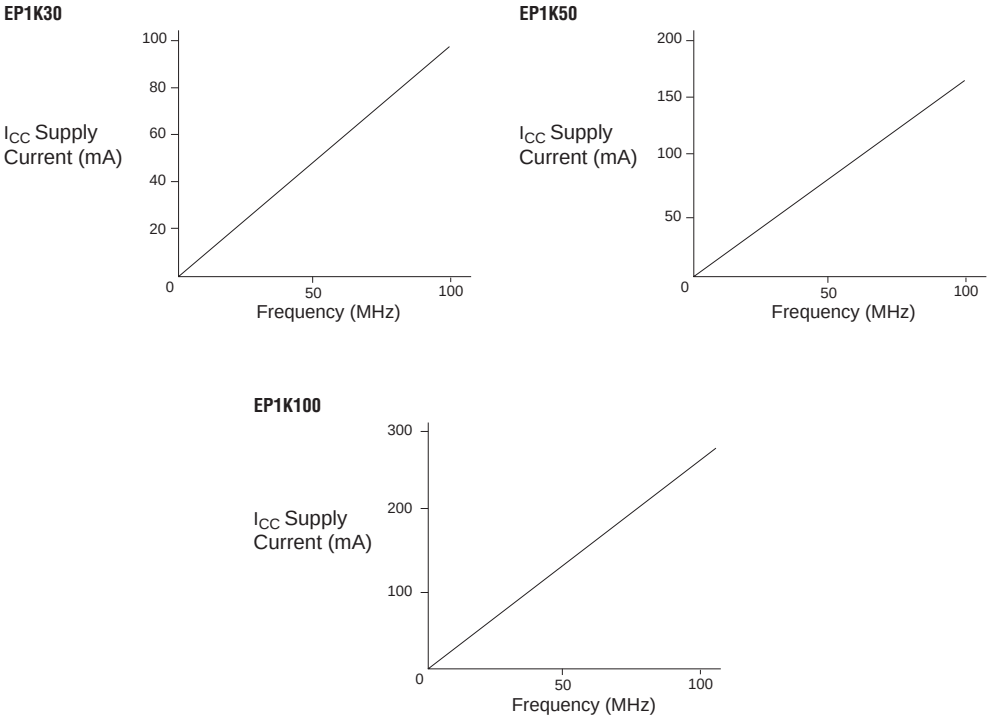
The supply power (P) for ACEX 1K devices can be calculated with the following equation:

$$P = P_{INT} + P_{IO} = (I_{CCSTANDBY} + I_{CCACTIVE}) \times V_{CC} + P_{IO}$$

The $I_{CCACTIVE}$ value depends on the switching frequency and the application logic. This value is calculated based on the amount of current that each LE typically consumes. The P_{IO} value, which depends on the device output load characteristics and switching frequency, can be calculated using the guidelines given in *Application Note 74 (Evaluating Power for Altera Devices).*

Compared to the rest of the device, the embedded array consumes a negligible amount of power. Therefore, the embedded array can be ignored when calculating supply current.

Figure 31. ACEX 1K $I_{CCACTIVE}$ vs. Operating Frequency



Configuration & Operation

The ACEX 1K architecture supports several configuration schemes. This section summarizes the device operating modes and available device configuration schemes.

Operating Modes

The ACEX 1K architecture uses SRAM configuration elements that require configuration data to be loaded every time the circuit powers up. The process of physically loading the SRAM data into the device is called *configuration*. Before configuration, as V_{CC} rises, the device initiates a Power-On Reset (POR). This POR event clears the device and prepares it for configuration. The ACEX 1K POR time does not exceed 50 µs.

When configuring with a configuration device, refer to the relevant configuration device data sheet for POR timing information.