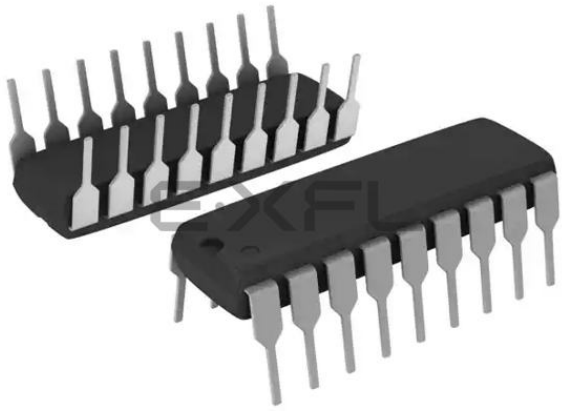

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"Embedded - Microcontrollers" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

Applications of "Embedded - Microcontrollers"

Details	
Product Status	Obsolete
Core Processor	PIC
Core Size	8-Bit
Speed	4MHz
Connectivity	-
Peripherals	POR, WDT
Number of I/O	13
Program Memory Size	896B (512 x 14)
Program Memory Type	FLASH
EEPROM Size	64 x 8
RAM Size	36 x 8
Voltage - Supply (Vcc/Vdd)	2V ~ 6V
Data Converters	-
Oscillator Type	External
Operating Temperature	0°C ~ 70°C (TA)
Mounting Type	Through Hole
Package / Case	18-DIP (0.300", 7.62mm)
Supplier Device Package	18-PDIP
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic16lf83-04-p

PIC16CXX devices contain an 8-bit ALU and working register. The ALU is a general purpose arithmetic unit. It performs arithmetic and Boolean functions between data in the working register and any register file.

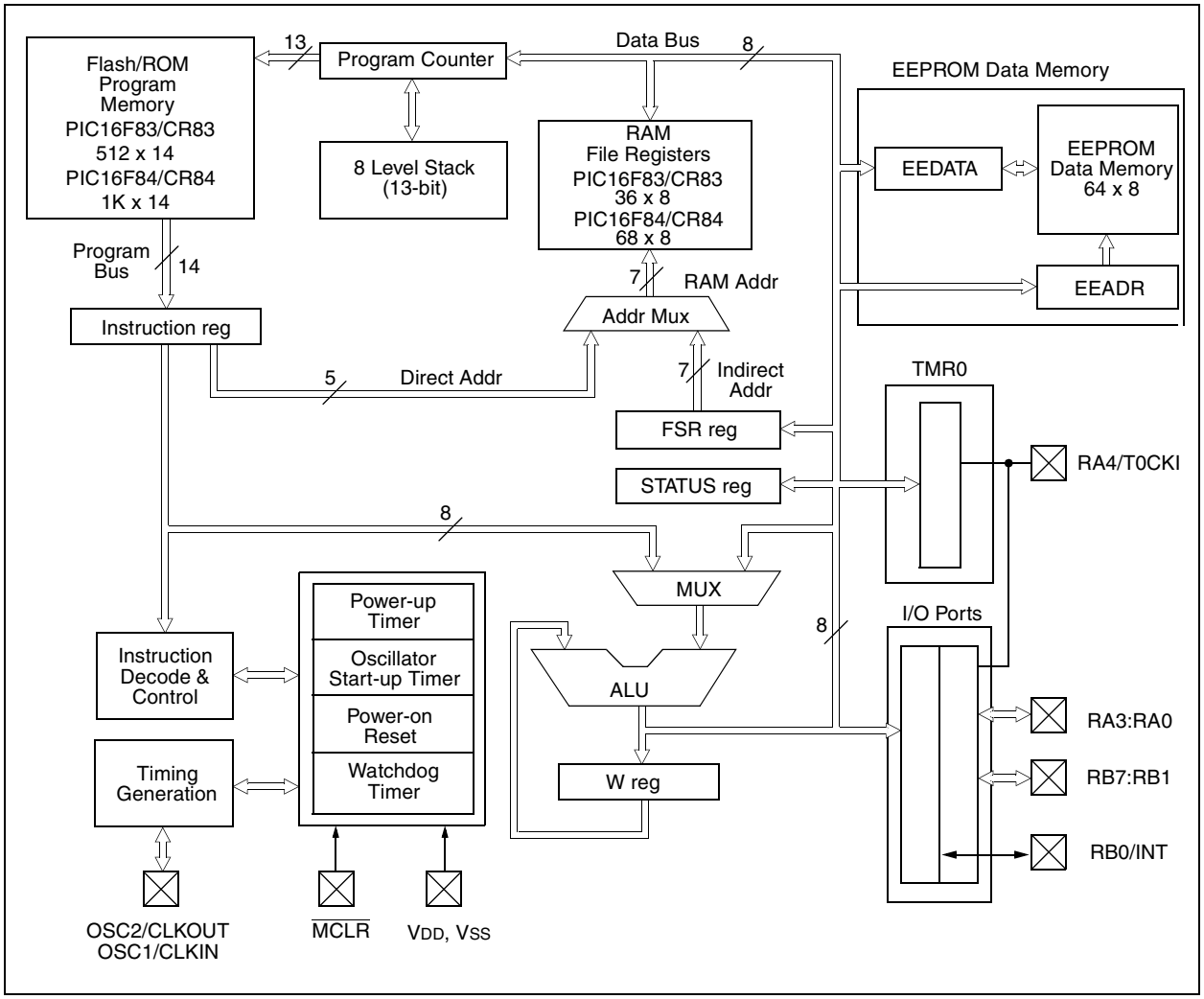
The ALU is 8-bits wide and capable of addition, subtraction, shift and logical operations. Unless otherwise mentioned, arithmetic operations are two's complement in nature. In two-operand instructions, typically one operand is the working register (W register), and the other operand is a file register or an immediate constant. In single operand instructions, the operand is either the W register or a file register.

The W register is an 8-bit working register used for ALU operations. It is not an addressable register.

Depending on the instruction executed, the ALU may affect the values of the Carry (C), Digit Carry (DC), and Zero (Z) bits in the STATUS register. The C and DC bits operate as a $\overline{\text{borrow}}$ and $\overline{\text{digit borrow}}$ out bit, respectively, in subtraction. See the SUBLW and SUBWF instructions for examples.

A simplified block diagram for the PIC16F8X is shown in Figure 3-1, its corresponding pin description is shown in Table 3-1.

FIGURE 3-1: PIC16F8X BLOCK DIAGRAM



3.1 Clocking Scheme/Instruction Cycle

The clock input (from OSC1) is internally divided by four to generate four non-overlapping quadrature clocks namely Q1, Q2, Q3 and Q4. Internally, the program counter (PC) is incremented every Q1, the instruction is fetched from the program memory and latched into the instruction register in Q4. The instruction is decoded and executed during the following Q1 through Q4. The clocks and instruction execution flow is shown in Figure 3-2.

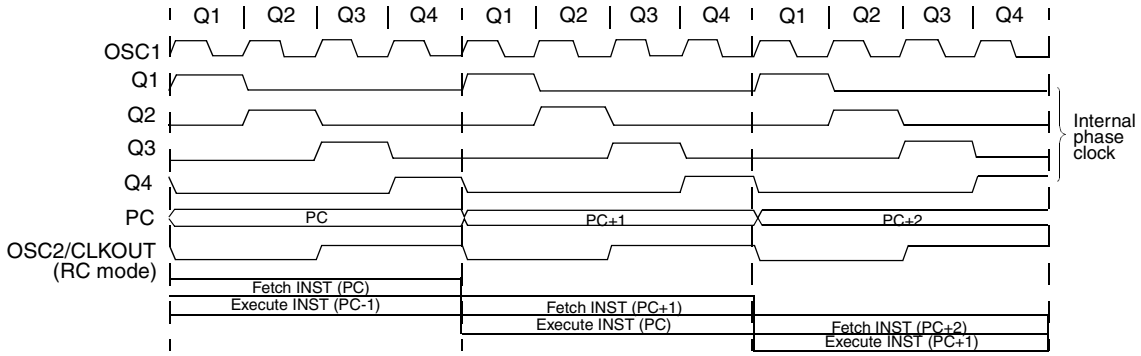
3.2 Instruction Flow/Pipelining

An "Instruction Cycle" consists of four Q cycles (Q1, Q2, Q3 and Q4). The instruction fetch and execute are pipelined such that fetch takes one instruction cycle while decode and execute takes another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the program counter to change (e.g., `GOTO`) then two cycles are required to complete the instruction (Example 3-1).

A fetch cycle begins with the Program Counter (PC) incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the "Instruction Register" in cycle Q1. This instruction is then decoded and executed during the Q2, Q3, and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).

FIGURE 3-2: CLOCK/INSTRUCTION CYCLE



EXAMPLE 3-1: INSTRUCTION PIPELINE FLOW

```
1. MOVLW 55h
2. MOVWF PORTB
3. CALL SUB_1
4. BSF   PORTA, BIT3
```

Fetch 1	Execute 1				
	Fetch 2	Execute 2			
		Fetch 3	Execute 3		
			Fetch 4	Flush	
				Fetch SUB_1	Execute SUB_1

All instructions are single cycle, except for any program branches. These take two cycles since the fetch instruction is "flushed" from the pipeline while the new instruction is being fetched and then executed.

4.5 Indirect Addressing; INDF and FSR Registers

The INDF register is not a physical register. Addressing INDF actually addresses the register whose address is contained in the FSR register (FSR is a *pointer*). This is indirect addressing.

EXAMPLE 4-1: INDIRECT ADDRESSING

- Register file 05 contains the value 10h
- Register file 06 contains the value 0Ah
- Load the value 05 into the FSR register
- A read of the INDF register will return the value of 10h
- Increment the value of the FSR register by one (FSR = 06)
- A read of the INDF register now will return the value of 0Ah.

Reading INDF itself indirectly (FSR = 0) will produce 00h. Writing to the INDF register indirectly results in a no-operation (although STATUS bits may be affected).

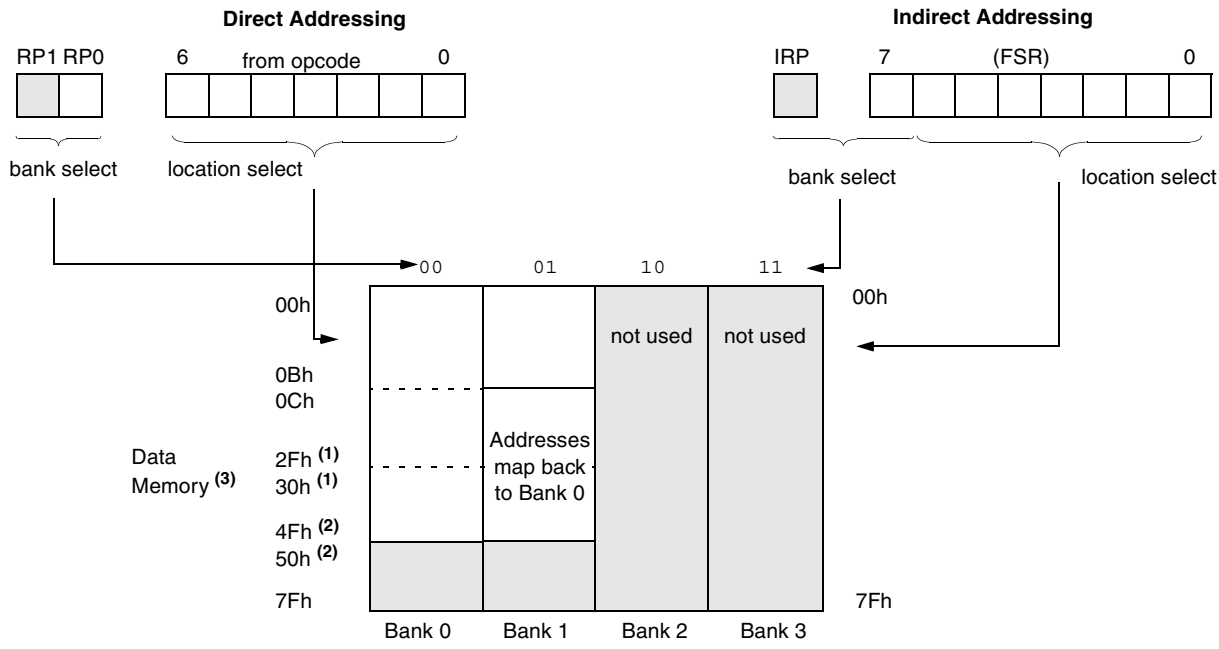
A simple program to clear RAM locations 20h-2Fh using indirect addressing is shown in Example 4-2.

EXAMPLE 4-2: HOW TO CLEAR RAM USING INDIRECT ADDRESSING

```
        movlw   0x20  ;initialize pointer
        movwf   FSR   ;  to RAM
NEXT    clrf    INDF  ;clear INDF register
        incf    FSR   ;inc pointer
        btfss   FSR,4 ;all done?
        goto    NEXT  ;NO, clear next
CONTINUE
        :             ;YES, continue
```

An effective 9-bit address is obtained by concatenating the 8-bit FSR register and the IRP bit (STATUS<7>), as shown in Figure 4-1. However, IRP is not used in the PIC16F8X.

FIGURE 4-1: DIRECT/INDIRECT ADDRESSING



Note 1: PIC16F83 and PIC16CR83 devices.
2: PIC16F84 and PIC16CR84 devices
3: For memory map detail see Figure 4-1.

NOTES:

8.1 Configuration Bits

The configuration bits can be programmed (read as '0') or left unprogrammed (read as '1') to select various device configurations. These bits are mapped in program memory location 2007h.

Address 2007h is beyond the user program memory space and it belongs to the special test/configuration memory space (2000h - 3FFFh). This space can only be accessed during programming.

To find out how to program the PIC16C84, refer to *PIC16C84 EEPROM Memory Programming Specification* (DS30189).

FIGURE 8-1: CONFIGURATION WORD - PIC16CR83 AND PIC16CR84

R-u	R-u	R-u	R-u	R-u	R-u	R/P-u	R-u	R-u	R-u	R-u	R-u	R-u	R-u
CP	CP	CP	CP	CP	CP	DP	CP	CP	CP	$\overline{\text{PWRTE}}$	WDTE	FOSC1	FOSC0
bit13													bit0

R = Readable bit
P = Programmable bit
- n = Value at POR reset
u = unchanged

bit 13:8 **CP**: Program Memory Code Protection bit
1 = Code protection off
0 = Program memory is code protected

bit 7 **DP**: Data Memory Code Protection bit
1 = Code protection off
0 = Data memory is code protected

bit 6:4 **CP**: Program Memory Code Protection bit
1 = Code protection off
0 = Program memory is code protected

bit 3 **$\overline{\text{PWRTE}}$**: Power-up Timer Enable bit
1 = Power-up timer is disabled
0 = Power-up timer is enabled

bit 2 **WDTE**: Watchdog Timer Enable bit
1 = WDT enabled
0 = WDT disabled

bit 1:0 **FOSC1:FOSC0**: Oscillator Selection bits
11 = RC oscillator
10 = HS oscillator
01 = XT oscillator
00 = LP oscillator

FIGURE 8-10: TIME-OUT SEQUENCE ON POWER-UP (MCLR NOT TIED TO VDD): CASE 1

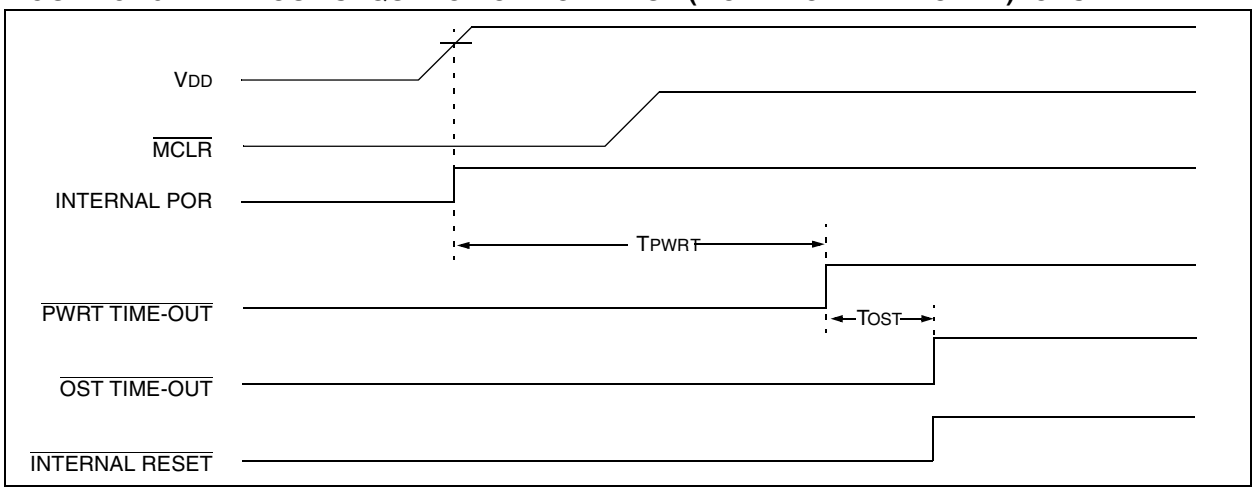
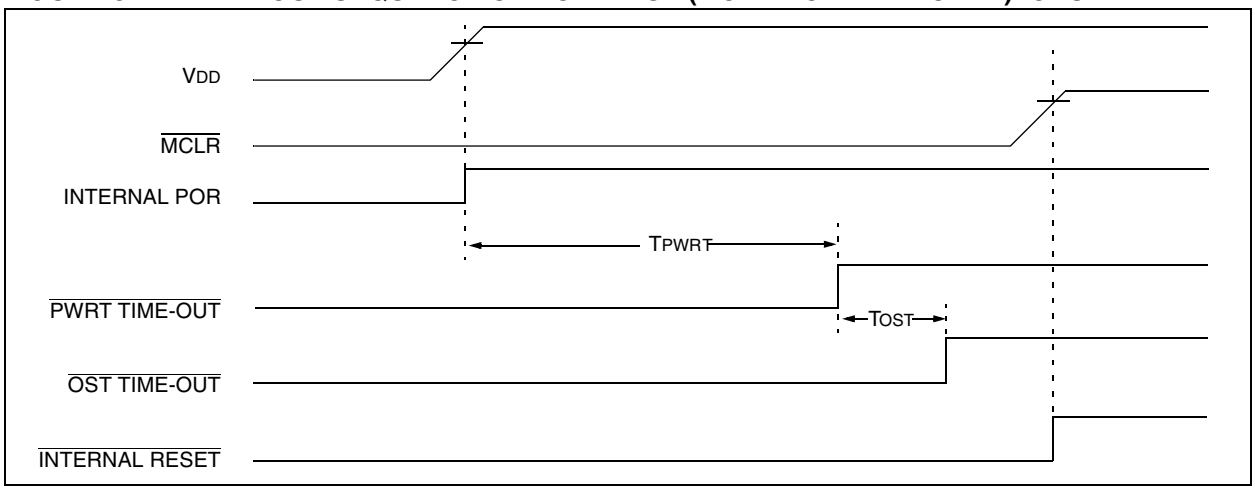


FIGURE 8-11: TIME-OUT SEQUENCE ON POWER-UP (MCLR NOT TIED TO VDD): CASE 2



8.9 Interrupts

The PIC16F8X has 4 sources of interrupt:

- External interrupt RB0/INT pin
- TMR0 overflow interrupt
- PORTB change interrupts (pins RB7:RB4)
- Data EEPROM write complete interrupt

The interrupt control register (INTCON) records individual interrupt requests in flag bits. It also contains the individual and global interrupt enable bits.

The global interrupt enable bit, GIE (INTCON<7>) enables (if set) all un-masked interrupts or disables (if cleared) all interrupts. Individual interrupts can be disabled through their corresponding enable bits in INTCON register. Bit GIE is cleared on reset.

The "return from interrupt" instruction, `RETFIE`, exits interrupt routine as well as sets the GIE bit, which re-enable interrupts.

The RB0/INT pin interrupt, the RB port change interrupt and the TMR0 overflow interrupt flags are contained in the INTCON register.

When an interrupt is responded to; the GIE bit is cleared to disable any further interrupt, the return address is pushed onto the stack and the PC is loaded with 0004h. For external interrupt events, such as the RB0/INT pin or PORTB change interrupt, the interrupt latency will be three to four instruction cycles. The exact latency depends when the interrupt event occurs (Figure 8-17). The latency is the same for both one and two cycle instructions. Once in the interrupt service routine the source(s) of the interrupt can be determined by polling the interrupt flag bits. The interrupt flag bit(s) must be cleared in software before re-enabling interrupts to avoid infinite interrupt requests.

Note 1:	Individual interrupt flag bits are set regardless of the status of their corresponding mask bit or the GIE bit.

FIGURE 8-16: INTERRUPT LOGIC

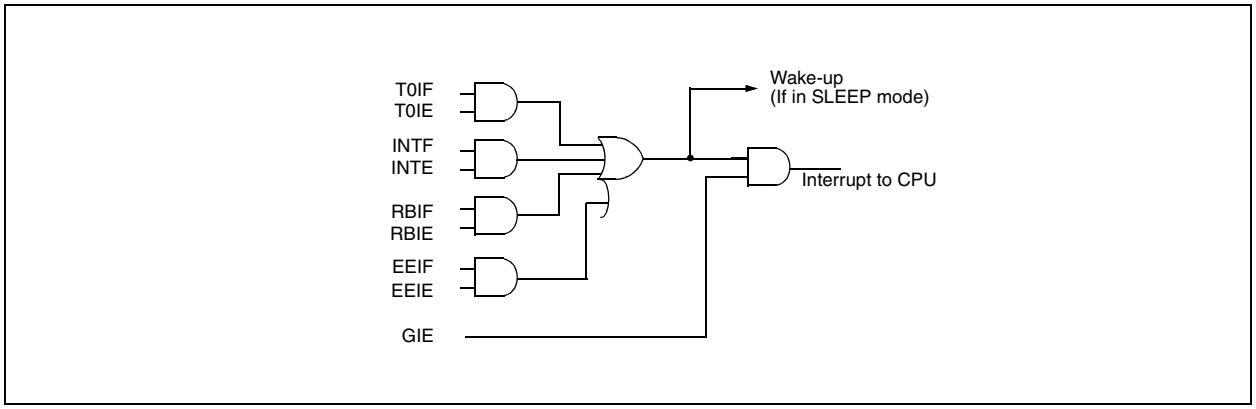
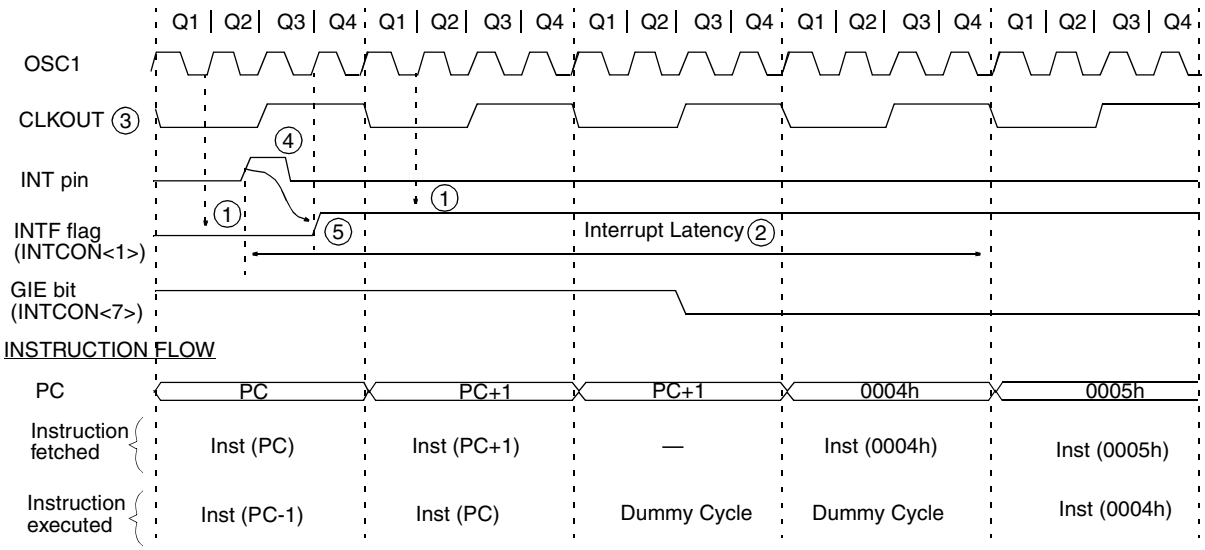


FIGURE 8-17: INT PIN INTERRUPT TIMING



Note 1: INTF flag is sampled here (every Q1).
2: Interrupt latency = 3-4Tcy where Tcy = instruction cycle time.
Latency is the same whether Inst (PC) is a single cycle or a 2-cycle instruction.
3: CLKOUT is available only in RC oscillator mode.
4: For minimum width of INT pulse, refer to AC specs.
5: INTF is enabled to be set anytime during the Q4-Q1 cycles.

8.9.1 INT INTERRUPT

External interrupt on RB0/INT pin is edge triggered: either rising if INTEDG bit (OPTION_REG<6>) is set, or falling, if INTEDG bit is clear. When a valid edge appears on the RB0/INT pin, the INTF bit (INTCON<1>) is set. This interrupt can be disabled by clearing control bit INTE (INTCON<4>). Flag bit INTF must be cleared in software via the interrupt service routine before re-enabling this interrupt. The INT interrupt can wake the processor from SLEEP (Section 8.12) only if the INTE bit was set prior to going into SLEEP. The status of the GIE bit decides whether the processor branches to the interrupt vector following wake-up.

8.9.2 TMR0 INTERRUPT

An overflow (FFh → 00h) in TMR0 will set flag bit T0IF (INTCON<2>). The interrupt can be enabled/disabled by setting/clearing enable bit T0IE (INTCON<5>) (Section 6.0).

8.9.3 PORT RB INTERRUPT

An input change on PORTB<7:4> sets flag bit RBIF (INTCON<0>). The interrupt can be enabled/disabled by setting/clearing enable bit RBIE (INTCON<3>) (Section 5.2).

Note 1: For a change on the I/O pin to be recognized, the pulse width must be at least TCY wide.

9.1 Instruction Descriptions

ADDLW **Add Literal and W**

Syntax: [*label*] ADDLW k

Operands: $0 \le k \le 255$

Operation: (W) + k → (W)

Status Affected: C, DC, Z

Encoding:

11	111x	kkkk	kkkk

Description: The contents of the W register are added to the eight bit literal 'k' and the result is placed in the W register.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process data	Write to W

Example:

```
ADDLW   0x15
```

Before Instruction
W = 0x10
After Instruction
W = 0x25

ANDLW **AND Literal with W**

Syntax: [*label*] ANDLW k

Operands: $0 \le k \le 255$

Operation: (W) .AND. (k) → (W)

Status Affected: Z

Encoding:

11	1001	kkkk	kkkk

Description: The contents of W register are AND'ed with the eight bit literal 'k'. The result is placed in the W register.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal "k"	Process data	Write to W

Example

```
ANDLW   0x5F
```

Before Instruction
W = 0xA3
After Instruction
W = 0x03

ADDWF **Add W and f**

Syntax: [*label*] ADDWF f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: (W) + (f) → (destination)

Status Affected: C, DC, Z

Encoding:

00	0111	dfff	ffff

Description: Add the contents of the W register with register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example

```
ADDWF   FSR, 0
```

Before Instruction
W = 0x17
FSR = 0xC2
After Instruction
W = 0xD9
FSR = 0xC2

ANDWF **AND W with f**

Syntax: [*label*] ANDWF f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: (W) .AND. (f) → (destination)

Status Affected: Z

Encoding:

00	0101	dfff	ffff

Description: AND the W register with register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example

```
ANDWF   FSR, 1
```

Before Instruction
W = 0x17
FSR = 0xC2
After Instruction
W = 0x17
FSR = 0x02

COMF **Complement f**

Syntax: [*label*] COMF f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: $(\bar{f}) \rightarrow$ (destination)

Status Affected: Z

Encoding:

00	1001	dfff	ffff

Description: The contents of register 'f' are complemented. If 'd' is 0 the result is stored in W. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example
```
COMF    REG1,0
```
Before Instruction
REG1 = 0x13
After Instruction
REG1 = 0x13
W = 0xEC

DECF **Decrement f**

Syntax: [*label*] DECF f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: (f) - 1 → (destination)

Status Affected: Z

Encoding:

00	0011	dfff	ffff

Description: Decrement register 'f'. If 'd' is 0 the result is stored in the W register. If 'd' is 1 the result is stored back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example
```
DECF    CNT, 1
```
Before Instruction
CNT = 0x01
Z = 0
After Instruction
CNT = 0x00
Z = 1

DECFSZ **Decrement f, Skip if 0**

Syntax: [*label*] DECFSZ f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: (f) - 1 → (destination);
skip if result = 0

Status Affected: None

Encoding:

00	1011	dfff	ffff

Description: The contents of register 'f' are decremented. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.
If the result is 1, the next instruction, is executed. If the result is 0, then a NOP is executed instead making it a 2TCY instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

If Skip: (2nd Cycle)

Q1	Q2	Q3	Q4
No-Operation	No-Operation	No-Operation	No-Operation

Example
```
HERE     DECFSZ   CNT, 1
         GOTO     LOOP
CONTINUE •
         •
         •
```
Before Instruction
PC = address HERE
After Instruction
CNT = CNT - 1
if CNT = 0,
PC = address CONTINUE
if CNT ≠ 0,
PC = address HERE+1

INCFSZ **Increment f, Skip if 0**

Syntax: [*label*] INCFSZ f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: (f) + 1 → (destination), skip if result = 0

Status Affected: None

Encoding:

00	1111	dfff	ffff

Description: The contents of register 'f' are incremented. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.
If the result is 1, the next instruction is executed. If the result is 0, a NOP is executed instead making it a 2TCY instruction.

Words: 1

Cycles: 1(2)

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

If Skip: (2nd Cycle)

Q1	Q2	Q3	Q4
No-Operation	No-Operation	No-Operation	No-Operation

Example

```
HERE     INCFSZ     CNT, 1
         GOTO       LOOP
CONTINUE •
         •
         •
```

Before Instruction
PC = address HERE
After Instruction
CNT = CNT + 1
if CNT= 0,
PC = address CONTINUE
if CNT≠ 0,
PC = address HERE +1

IORLW **Inclusive OR Literal with W**

Syntax: [*label*] IORLW k

Operands: $0 \le k \le 255$

Operation: (W) .OR. k → (W)

Status Affected: Z

Encoding:

11	1000	kkkk	kkkk

Description: The contents of the W register is OR'ed with the eight bit literal 'k'. The result is placed in the W register.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process data	Write to W

Example

```
IORLW   0x35
```

Before Instruction
W = 0x9A
After Instruction
W = 0xBF
Z = 1

IORWF **Inclusive OR W with f**

Syntax: [*label*] IORWF f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: (W) .OR. (f) → (destination)

Status Affected: $\overline{Z}$

Encoding:

00	0100	dfff	ffff

Description: Inclusive OR the W register with register 'f'. If 'd' is 0 the result is placed in the W register. If 'd' is 1 the result is placed back in register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example
```
IORWF          RESULT, 0
```
Before Instruction
RESULT = 0x13
W = 0x91
After Instruction
RESULT = 0x13
W = 0x93
Z = 1

MOVF **Move f**

Syntax: [*label*] MOVF f,d

Operands: $0 \le f \le 127$
$d \in [0,1]$

Operation: (f) → (destination)

Status Affected: Z

Encoding:

00	1000	dfff	ffff

Description: The contents of register f is moved to a destination dependant upon the status of d. If d = 0, destination is W register. If d = 1, the destination is file register f itself. d = 1 is useful to test a file register since status flag Z is affected.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write to destination

Example
```
MOVF    FSR, 0
```
After Instruction
W = value in FSR register
Z = 1

MOVLW **Move Literal to W**

Syntax: [*label*] MOVLW k

Operands: $0 \le k \le 255$

Operation: k → (W)

Status Affected: None

Encoding:

11	00xx	kkkk	kkkk

Description: The eight bit literal 'k' is loaded into W register. The don't cares will assemble as 0's.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read literal 'k'	Process data	Write to W

Example
```
MOVLW   0x5A
```
After Instruction
W = 0x5A

MOVWF **Move W to f**

Syntax: [*label*] MOVWF f

Operands: $0 \le f \le 127$

Operation: (W) → (f)

Status Affected: None

Encoding:

00	0000	1fff	ffff

Description: Move data from W register to register 'f'.

Words: 1

Cycles: 1

Q Cycle Activity:

Q1	Q2	Q3	Q4
Decode	Read register 'f'	Process data	Write register 'f'

Example
```
MOVWF   OPTION_REG
```
Before Instruction
OPTION = 0xFF
W = 0x4F
After Instruction
OPTION = 0x4F
W = 0x4F

TABLE 10-1: DEVELOPMENT TOOLS FROM MICROCHIP

		PIC12C5XX	PIC14000	PIC16C5X	PIC16CXXX	PIC16C6X	PIC16C7XX	PIC16C8X	PIC16C9XX	PIC17C4X	PIC17C75X	24CXX 25CXX 93CXX	HCS200 HCS300 HCS301
Emulator Products	PICMASTER®/ PICMASTER-CE In-Circuit Emulator	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓		
	ICEPIC™ Low-Cost In-Circuit Emulator	✓		✓	✓	✓	✓	✓	✓				
Software Tools	MPLAB™ Integrated Development Environment	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓		
	MPLAB™ C17 Compiler									✓	✓		
	*fuzzy*TECH®-MP Explorer/Edition Fuzzy Logic Dev. Tool	✓	✓	✓	✓	✓	✓	✓	✓	✓			
	MP-DriveWay™ Applications Code Generator			✓	✓	✓	✓	✓	✓	✓			
	Total Endurance™ Software Model											✓	
Programmers	PICSTART®Plus Low-Cost Universal Dev. Kit	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓		
	PRO MATE® II Universal Programmer	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓
	KEELOQ® Programmer												✓
	SEEVAL® Designers Kit											✓	
Demo Boards	PICDEM-1			✓	✓			✓		✓			
	PICDEM-2					✓	✓						
	PICDEM-3								✓				
	KEELOQ® Evaluation Kit												✓

FIGURE 10-5: RESET, WATCHDOG TIMER, OSCILLATOR START-UP TIMER AND POWER-UP TIMER TIMING

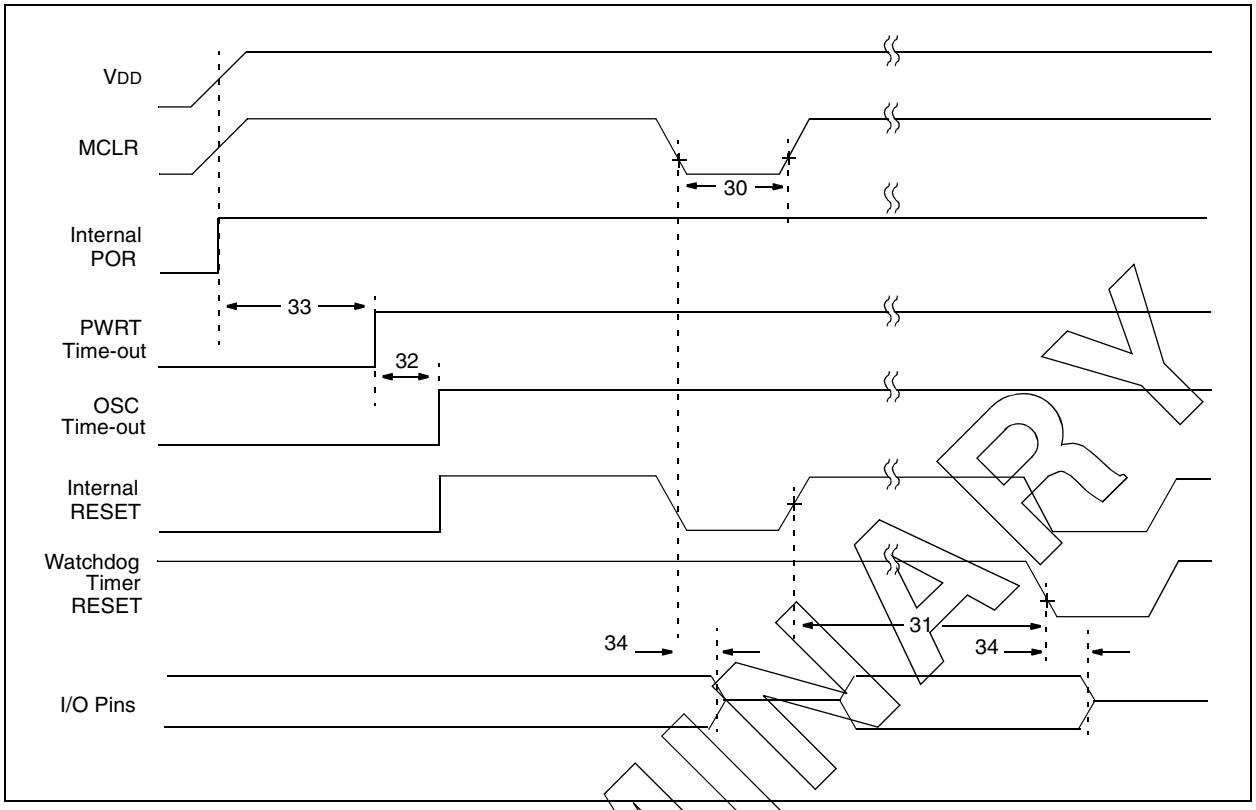


TABLE 10-5 RESET, WATCHDOG TIMER, OSCILLATOR START-UP TIMER AND POWER-UP TIMER REQUIREMENTS

Parameter No.	Sym	Characteristic	Min	Typ†	Max	Units	Conditions
30	TmcL	$\overline{MCLR}$ Pulse Width (low)	1000 *	—	—	ns	2.0V ≤ VDD ≤ 6.0V
31	Twdt	Watchdog Timer Time-out Period (No Prescaler)	7 *	18	33 *	ms	VDD = 5.0V
32	Tost	Oscillation Start-up Timer Period		1024TOSC		ms	TOSC = OSC1 period
33	Tpwrt	Power-up Timer Period	28 *	72	132 *	ms	VDD = 5.0V
34	TIOZ	I/O Hi-impedance from $\overline{MCLR}$ Low or reset	—	—	100 *	ns	

* These parameters are characterized but not tested.

† Data in "Typ" column is at 5V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

NOTES:

TABLE 11-1 CROSS REFERENCE OF DEVICE SPECS FOR OSCILLATOR CONFIGURATIONS AND FREQUENCIES OF OPERATION (COMMERCIAL DEVICES)

OSC	PIC16CR84-04 PIC16CR83-04	PIC16CR84-10 PIC16CR83-10	PIC16LCR84-04 PIC16LCR83-04
RC	VDD: 4.0V to 6.0V IDD: 4.5 mA max. at 5.5V IPD: 14 μA max. at 4V WDT dis Freq: 4.0 MHz max.	VDD: 4.5V to 5.5V IDD: 1.8 mA typ. at 5.5V IPD: 1.0 μA typ. at 5.5V WDT dis Freq: 4..0 MHz max.	VDD: 2.0V to 6.0V IDD: 4.5 mA max. at 5.5V IPD: 5 μA max. at 2V WDT dis Freq: 2.0 MHz max.
XT	VDD: 4.0V to 6.0V IDD: 4.5 mA max. at 5.5V IPD: 14 μA max. at 4V WDT dis Freq: 4.0 MHz max.	VDD: 4.5V to 5.5V IDD: 1.8 mA typ. at 5.5V IPD: 1.0 μA typ. at 5.5V WDT dis Freq: 4.0 MHz max.	VDD: 2.0V to 6.0V IDD: 4.5 mA max. at 5.5V IPD: 5 μA max. at 2V WDT dis Freq: 2.0 MHz max.
HS	VDD: 4.5V to 5.5V IDD: 4.5 mA typ. at 5.5V IPD: 1.0 μA typ. at 4.5V WDT dis Freq: 4.0 MHz max.	VDD: 4.5V to 5.5V IDD: 10 mA max. at 5.5V typ. IPD: 1.0 μA typ. at 4.5V WDT dis Freq: 10 MHz max.	Do not use in HS mode
LP	VDD: 4.0V to 6.0V IDD: 48 μA typ. at 32 kHz, 2.0V IPD: 0.6 μA typ. at 3.0V WDT dis Freq: 200 kHz max.	Do not use in LP mode	VDD: 2.0V to 6.0V IDD: 45 μA max. at 32 kHz, 2.0V IPD: 5 μA max. at 2V WDT dis Freq: 200 kHz max.

The shaded sections indicate oscillator selections which are tested for functionality, but not for MIN/MAX specifications. It is recommended that the user select the device type that ensures the specifications required.

11.1 DC CHARACTERISTICS: PIC16CR84, PIC16CR83 (Commercial, Industrial)

DC Characteristics Power Supply Pins			Standard Operating Conditions (unless otherwise stated) Operating temperature 0°C ≤ TA ≤ +70°C (commercial) -40°C ≤ TA ≤ +85°C (industrial)				
Parameter No.	**Sym**	**Characteristic**	**Min**	**Typ†**	**Max**	**Units**	**Conditions**
D001	VDD	Supply Voltage	4.0	—	6.0	V	XT, RC and LP osc configuration
D001A			4.5	—	5.5	V	HS osc configuration
D002	VDR	RAM Data Retention Voltage(1)	1.5*	—	—	V	Device in SLEEP mode
D003	VPOR	VDD start voltage to ensure internal Power-on Reset signal	—	VSS	—	V	See section on Power-on Reset for details
D004	SVDD	VDD rise rate to ensure internal Power-on Reset signal	0.05*	—	—	V/ms	See section on Power-on Reset for details
	IDD	Supply Current(2)					RC and XT osc configuration(4)
D010			—	1.8	4.5	mA	FOSC = 4.0 MHz, VDD = 5.5V
D010A			—	7.3	10	mA	FOSC = 4.0 MHz, VDD = 5.5V (During EEPROM programming)
							HS OSC CONFIGURATION (PIC16CR84-10)
D013			—	5	10	mA	FOSC = 10 MHz, VDD = 5.5V
D020	IPD	Power-down Current(3)	—	7.0	28	µA	VDD = 4.0V, WDT enabled, industrial
D021			—	1.0	14	µA	VDD = 4.0V, WDT disabled, commercial
D021A			—	1.0	16	µA	VDD = 4.0V, WDT disabled, industrial

* These parameters are characterized but not tested.

† Data in "Typ" column is at 5.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

Note 1: This is the limit to which VDD can be lowered in SLEEP mode without losing RAM data.

2: The supply current is mainly a function of the operating voltage and frequency. Other factors such as I/O pin loading and switching rate, oscillator type, internal code execution pattern, and temperature also have an impact on the current consumption.
The test conditions for all IDD measurements in active operation mode are:
OSC1=external square wave, from rail to rail; all I/O pins tristated, pulled to VDD, T0CKI = VDD, $\overline{MCLR}$ = VDD; WDT enabled/disabled as specified.

3: The power down current in SLEEP mode does not depend on the oscillator type. Power-down current is measured with the part in SLEEP mode, with all I/O pins in hi-impedance state and tied to VDD and VSS.

4: For RC osc configuration, current through Rext is not included. The current through the resistor can be estimated by the formula IR = VDD/2Rext (mA) with Rext in kOhm.

FIGURE 12-13: WDT TIMER TIME-OUT PERIOD vs. VDD

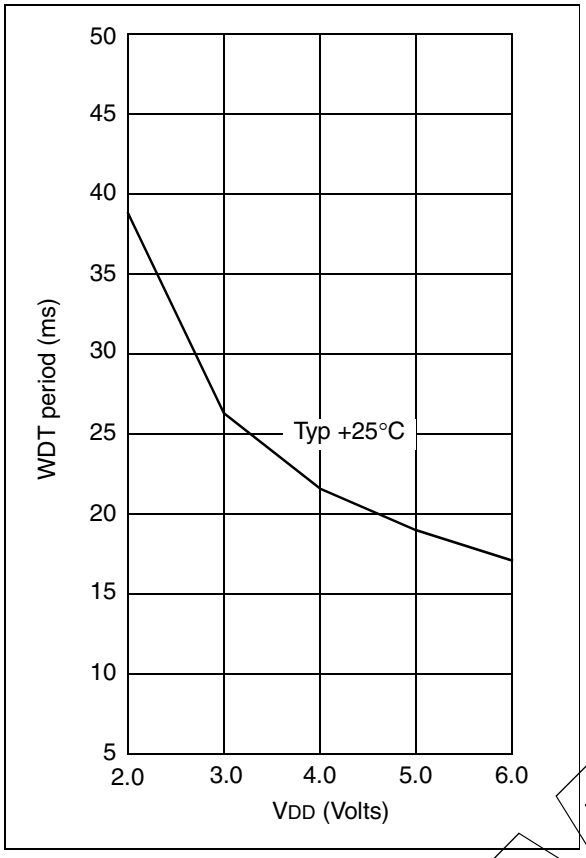


FIGURE 12-14: TRANSCONDUCTANCE (gm) OF HS OSCILLATOR vs. VDD

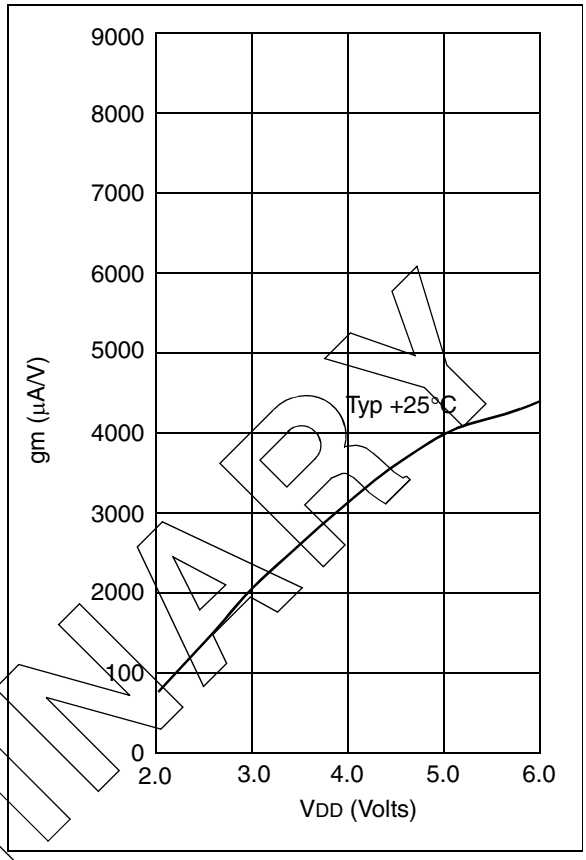


FIGURE 12-21: TYPICAL DATA MEMORY ERASE/WRITE CYCLE TIME VS. VDD

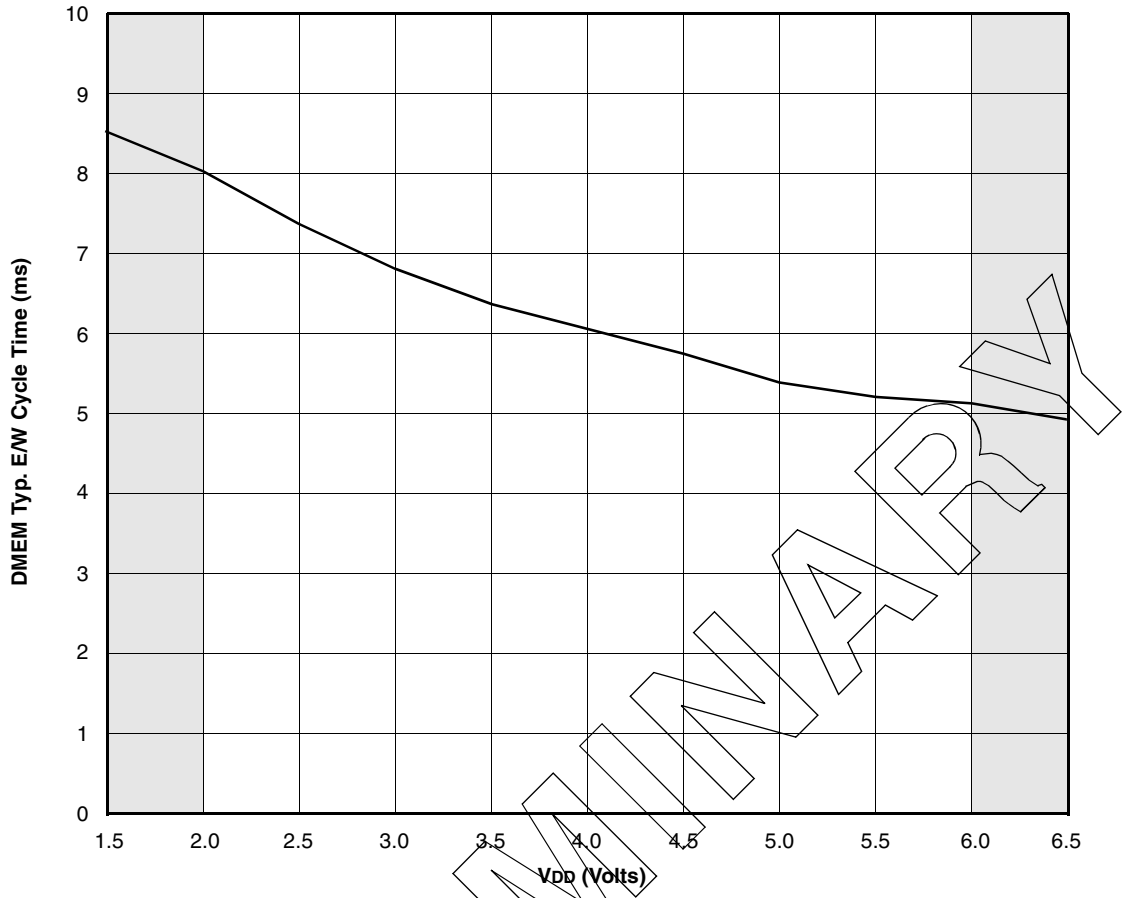


TABLE 12-2 INPUT CAPACITANCE*

Pin Name	Typical Capacitance (pF)	
	18L PDIP	18L SOIC
PORTA	5.0	4.3
PORTB	5.0	4.3
$\overline{MCLR}$	17.0	17.0
OSC1/CLKIN	4.0	3.5
OSC2/CLKOUT	4.3	3.5
T0CKI	3.2	2.8

* All capacitance values are typical at 25°C. A part to part variation of ±25% (three standard deviations) should be taken into account.