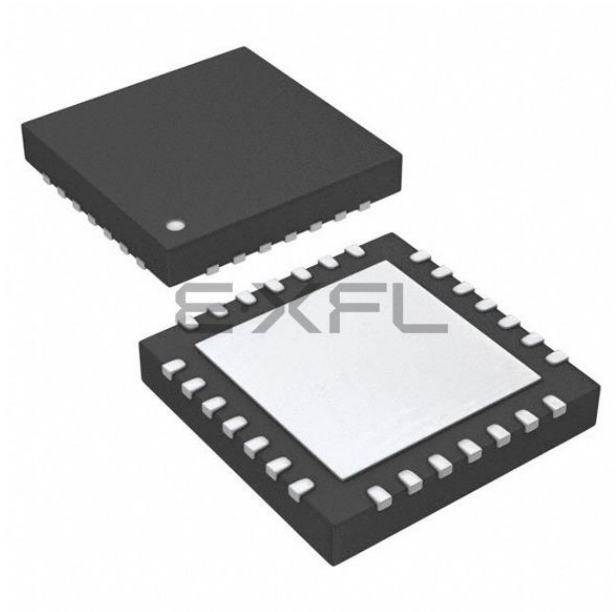

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## What is "Embedded - Microcontrollers"?

"Embedded - Microcontrollers" refer to small, integrated circuits designed to perform specific tasks within larger systems. These microcontrollers are essentially compact computers on a single chip, containing a processor core, memory, and programmable input/output peripherals. They are called "embedded" because they are embedded within electronic devices to control various functions, rather than serving as standalone computers. Microcontrollers are crucial in modern electronics, providing the intelligence and control needed for a wide range of applications.

## Applications of "Embedded - Microcontrollers"

| Details | |
|---|---|
| Product Status | Active |
| Core Processor | PIC |
| Core Size | 8-Bit |
| Speed | 64MHz |
| Connectivity | I²C, LINbus, SPI, UART/USART |
| Peripherals | Brown-out Detect/Reset, LVD, POR, PWM, WDT |
| Number of I/O | 25 |
| Program Memory Size | 32KB (16K x 16) |
| Program Memory Type | FLASH |
| EEPROM Size | 256 x 8 |
| RAM Size | 2K x 8 |
| Voltage - Supply (Vcc/Vdd) | 2.3V ~ 5.5V |
| Data Converters | A/D 35x10b; D/A 1x5b |
| Oscillator Type | Internal |
| Operating Temperature | -40°C ~ 85°C (TA) |
| Mounting Type | Surface Mount |
| Package / Case | 28-VQFN Exposed Pad |
| Supplier Device Package | 28-QFN (6x6) |
| Purchase URL | https://www.e-xfl.com/product-detail/microchip-technology/pic18f25k40t-i-ml |

**REGISTER 3-6: CONFIGURATION WORD 3H (30 0005h): WINDOWED WATCHDOG TIMER**

| U-1 | U-1 | R/W-1 | R/W-1 | R/W-1 | R/W-1 | R/W-1 | R/W-1 |
|---|---|---|---|---|---|---|---|
| — | — | WDTCCS<2:0> | | | WDTCWS<2:0> | | |
| bit 7 | | | | | | | bit 0 |

| **Legend:** | | | |
|---|---|---|---|
| R = Readable bit | W = Writable bit | U = Unimplemented bit, read as '1' | |
| -n = Value for blank device | '1' = Bit is set | '0' = Bit is cleared | x = Bit is unknown |

bit 7-6 **Unimplemented:** Read as '1'

bit 5-3 **WDTCCS<2:0>:** WDT Input Clock Selector bits

If WDTE<1:0> fuses = 2'b00

This bit is ignored.

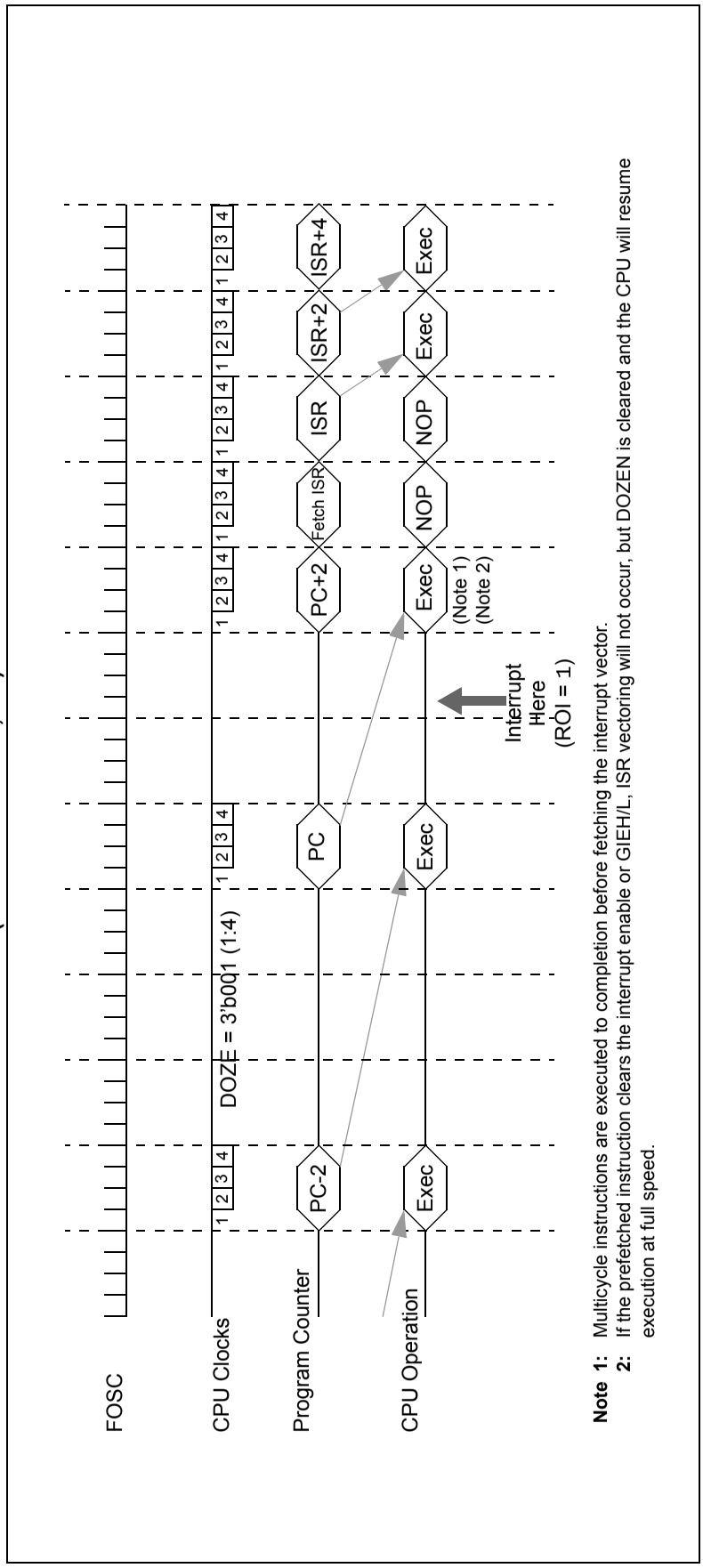
Otherwise:

111 = Software Control
110 = Reserved (Default to LFINTOSC)
.
.
.
010 = Reserved (Default to LFINTOSC)
001 = WDT reference clock is the 31.25 kHz MFINTOSC
000 = WDT reference clock is the 31.0 kHz LFINTOSC (default value)

bit 2-0 **WDTCWS<2:0>:** WDT Window Select bits

| WDTCWS | WINDOW at POR | | | Software control of WINDOW | Keyed access required? |
|---|---|---|---|---|---|
| | Value | Window delay Percent of time | Window opening Percent of time | | |
| 111 | 111 | n/a | 100 | Yes | No |
| 110 | 111 | n/a | 100 | No | Yes |
| 101 | 101 | 25 | 75 | | |
| 100 | 100 | 37.5 | 62.5 | | |
| 011 | 011 | 50 | 50 | | |
| 010 | 010 | 62.5 | 37.5 | | |
| 001 | 001 | 75 | 25 | | |
| 000 | 000 | 87.5 | 12.5 | | |

**FIGURE 6-1: DOZE MODE OPERATION EXAMPLE (DOZE<2:0> = 001, 1:4)**




**Note 1:** Multicycle instructions are executed to completion before fetching the interrupt vector.
**2:** If the prefetched instruction clears the interrupt enable or GIEH/L, ISR vectoring will not occur, but DOZEN is cleared and the CPU will resume execution at full speed.

## 8.0 RESETS

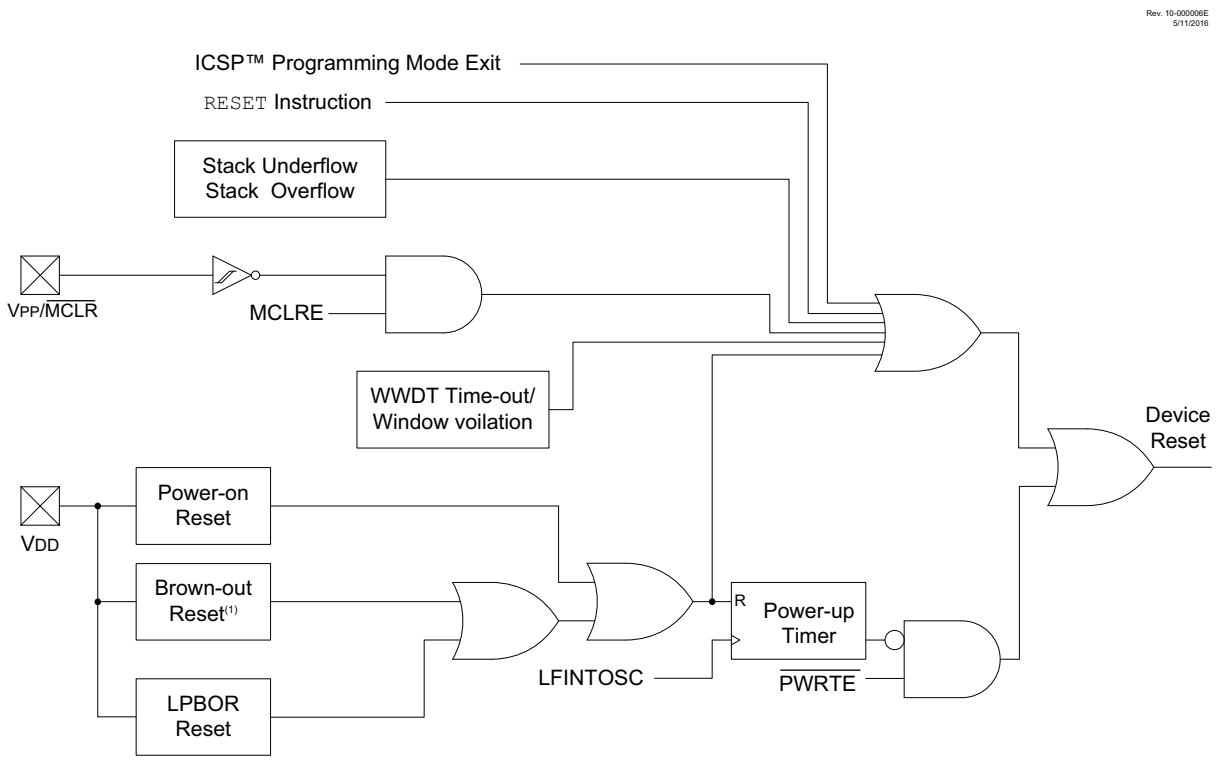
There are multiple ways to reset this device:

- Power-on Reset (POR)
- Brown-out Reset (BOR)
- Low-Power Brown-Out Reset (LPBOR)
- $\overline{\text{MCLR}}$ Reset
- WDT Reset
- RESET instruction
- Stack Overflow
- Stack Underflow
- Programming mode exit

To allow $V_{DD}$ to stabilize, an optional Power-up Timer can be enabled to extend the Reset time after a BOR or POR event.

A simplified block diagram of the On-Chip Reset Circuit is shown in Figure 8-1.

**FIGURE 8-1: SIMPLIFIED BLOCK DIAGRAM OF ON-CHIP RESET CIRCUIT**




**Note 1:** See Table 8-1 for BOR active conditions.

**TABLE 10-5: REGISTER FILE SUMMARY FOR PIC18(L)F24/25K40 DEVICES (CONTINUED)**

| Address | Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Value on POR, BOR |
|---|---|---|---|---|---|---|---|---|---|---|
| F26h to F22h | — | Unimplemented | | | | | | | | — |
| F21h | ANSELC | ANSELC7 | ANSELC6 | ANSELC5 | ANSELC4 | ANSELC3 | ANSELC2 | ANSELC1 | ANSELC0 | 11111111 |
| F20h | WPUC | WPUC7 | WPUC6 | WPUC5 | WPUC4 | WPUC3 | WPUC2 | WPUC1 | WPUC0 | 00000000 |
| F1Fh | ODCONC | ODCC7 | ODCC6 | ODCC5 | ODCC4 | ODCC3 | ODCC2 | ODCC1 | ODCC0 | 00000000 |
| F1Eh | SLRCONC | SLRC7 | SLRC6 | SLRC5 | SLRC4 | SLRC3 | SLRC2 | SLRC1 | SLRC0 | 11111111 |
| F1Dh | INLVLC | INLVLC7 | INLVLC6 | INLVLC5 | INLVLC4 | INLVLC3 | INLVLC2 | INLVLC1 | INLVLC0 | 11111111 |
| F1Ch | IOCCP | IOCCP7 | IOCCP6 | IOCCP5 | IOCCP4 | IOCCP3 | IOCCP2 | IOCCP1 | IOCCP0 | 00000000 |
| F1Bh | IOCCN | IOCCN7 | IOCCN6 | IOCCN5 | IOCCN4 | IOCCN3 | IOCCN2 | IOCCN1 | IOCCN0 | 00000000 |
| F1Ah | IOCCF | IOCCF7 | IOCCF6 | IOCCF5 | IOCCF4 | IOCCF3 | IOCCF2 | IOCCF1 | IOCCF0 | 00000000 |
| F19h | ANSELB | ANSELB7 | ANSELB6 | ANSELB5 | ANSELB4 | ANSELB3 | ANSELB2 | ANSELB1 | ANSELB0 | 11111111 |
| F18h | WPUB | WPUB7 | WPUB6 | WPUB5 | WPUB4 | WPUB3 | WPUB2 | WPUB1 | WPUB0 | 00000000 |
| F17h | ODCONB | ODCB7 | ODCB6 | ODCB5 | ODCB4 | ODCB3 | ODCB2 | ODCB1 | ODCB0 | 00000000 |
| F16h | SLRCONB | SLRB7 | SLRB6 | SLRB5 | SLRB4 | SLRB3 | SLRB2 | SLRB1 | SLRB0 | 11111111 |
| F15h | INLVLB | INLVLB7 | INLVLB6 | INLVLB5 | INLVLB4 | INLVLB3 | INLVLB2 | INLVLB1 | INLVLB0 | 11111111 |
| F14h | IOCBP | IOCBP7 | IOCBP6 | IOCBP5 | IOCBP4 | IOCBP3 | IOCBP2 | IOCBP1 | IOCBP0 | 00000000 |
| F13h | IOCBN | IOCBN7 | IOCBN6 | IOCBN5 | IOCBN4 | IOCBN3 | IOCBN2 | IOCBN1 | IOCBN0 | 00000000 |
| F12h | IOCBF | IOCBF7 | IOCBF6 | IOCBF5 | IOCBF4 | IOCBF3 | IOCBF2 | IOCBF1 | IOCBF0 | 00000000 |
| F11h | ANSELA | ANSELA7 | ANSELA6 | ANSELA5 | ANSELA4 | ANSELA3 | ANSELA2 | ANSELA1 | ANSELA0 | 11111111 |
| F10h | WPUA | WPUA7 | WPUA6 | WPUA5 | WPUA4 | WPUA3 | WPUA2 | WPUA1 | WPUA0 | 00000000 |
| F0Fh | ODCONA | ODCA7 | ODCA6 | ODCA5 | ODCA4 | ODCA3 | ODCA2 | ODCA1 | ODCA0 | 00000000 |
| F0Eh | SLRCONA | SLRA7 | SLRA6 | SLRA5 | SLRA4 | SLRA3 | SLRA2 | SLRA1 | SLRA0 | 11111111 |
| F0Dh | INLVLA | INLVLA7 | INLVLA6 | INLVLA5 | INLVLA4 | INLVLA3 | INLVLA2 | INLVLA1 | INLVLA0 | 11111111 |
| F0Ch | IOCAP | IOCAP7 | IOCAP6 | IOCAP5 | IOCAP4 | IOCAP3 | IOCAP2 | IOCAP1 | IOCAP0 | 00000000 |
| F0Bh | IOCAN | IOCAN7 | IOCAN6 | IOCAN5 | IOCAN4 | IOCAN3 | IOCAN2 | IOCAN1 | IOCAN0 | 00000000 |
| F0Ah | IOCAF | IOCAF7 | IOCAF6 | IOCAF5 | IOCAF4 | IOCAF3 | IOCAF2 | IOCAF1 | IOCAF0 | 00000000 |
| F09h to EFFh | — | Unimplemented | | | | | | | | — |
| EFEh | RC7PPS | — | — | — | RC7PPS<4:0> | | | | | ---00000 |
| EFDh | RC6PPS | — | — | — | RC6PPS<4:0> | | | | | ---00000 |
| EFCh | RC5PPS | — | — | — | RC5PPS<4:0> | | | | | ---00000 |
| EFBh | RC4PPS | — | — | — | RC4PPS<4:0> | | | | | ---00000 |
| EFAh | RC3PPS | — | — | — | RC3PPS<4:0> | | | | | ---00000 |
| EF9h | RC2PPS | — | — | — | RC2PPS<4:0> | | | | | ---00000 |
| EF8h | RC1PPS | — | — | — | RC1PPS<4:0> | | | | | ---00000 |
| EF7h | RC0PPS | — | — | — | RC0PPS<4:0> | | | | | ---00000 |
| EF6h | RB7PPS | — | — | — | RB7PPS<4:0> | | | | | ---00000 |
| EF5h | RB6PPS | — | — | — | RB6PPS<4:0> | | | | | ---00000 |
| EF4h | RB5PPS | — | — | — | RB5PPS<4:0> | | | | | ---00000 |
| EF3h | RB4PPS | — | — | — | RB4PPS<4:0> | | | | | ---00000 |

**Legend:** x = unknown, u = unchanged, — = unimplemented, q = value depends on condition
**Note 1:** Not available on LF devices.

Operations on the FSRs with POSTDEC, POSTINC and PREINC affect the entire register pair; that is, roll-overs of the FSRnL register from FFh to 00h carry over to the FSRnH register. On the other hand, results of these operations do not change the value of any flags in the STATUS register (e.g., Z, N, OV, etc.).

The PLUSW register can be used to implement a form of indexed addressing in the data memory space. By manipulating the value in the W register, users can reach addresses that are fixed offsets from pointer addresses. In some applications, this can be used to implement some powerful program control structure, such as software stacks, inside of data memory.

#### 10.6.3.3 Operations by FSRs on FSRs

Indirect addressing operations that target other FSRs or virtual registers represent special cases. For example, using an FSR to point to one of the virtual registers will not result in successful operations. As a specific case, assume that FSR0H:FSR0L contains FE7h, the address of INDF1. Attempts to read the value of the INDF1 using INDF0 as an operand will return 00h. Attempts to write to INDF1 using INDF0 as the operand will result in a NOP.

On the other hand, using the virtual registers to write to an FSR pair may not occur as planned. In these cases, the value will be written to the FSR pair but without any incrementing or decrementing. Thus, writing to either the INDF2 or POSTDEC2 register will write the same value to the FSR2H:FSR2L.

Since the FSRs are physical registers mapped in the SFR space, they can be manipulated through all direct operations. Users should proceed cautiously when working on these registers, particularly if their code uses indirect addressing.

Similarly, operations by indirect addressing are generally permitted on all other SFRs. Users should exercise the appropriate caution that they do not inadvertently change settings that might affect the operation of the device.

## 10.7 Data Memory and the Extended Instruction Set

Enabling the PIC18 extended instruction set (XINST Configuration bit = 1) significantly changes certain aspects of data memory and its addressing. Specifically, the use of the Access Bank for many of the core PIC18 instructions is different; this is due to the introduction of a new addressing mode for the data memory space.

What does not change is just as important. The size of the data memory space is unchanged, as well as its linear addressing. The SFR map remains the same. Core PIC18 instructions can still operate in both Direct and Indirect Addressing mode; inherent and literal instructions do not change at all. Indirect addressing with FSR0 and FSR1 also remain unchanged.

### 10.7.1 INDEXED ADDRESSING WITH LITERAL OFFSET

Enabling the PIC18 extended instruction set changes the behavior of indirect addressing using the FSR2 register pair within Access RAM. Under the proper conditions, instructions that use the Access Bank – that is, most bit-oriented and byte-oriented instructions – can invoke a form of indexed addressing using an offset specified in the instruction. This special addressing mode is known as Indexed Addressing with Literal Offset, or Indexed Literal Offset mode.

When using the extended instruction set, this addressing mode requires the following:

- The use of the Access Bank is forced (‘a’ = 0) and
- The file address argument is less than or equal to 5Fh.

Under these conditions, the file address of the instruction is not interpreted as the lower byte of an address (used with the BSR in direct addressing), or as an 8-bit address in the Access Bank. Instead, the value is interpreted as an offset value to an Address Pointer, specified by FSR2. The offset and the contents of FSR2 are added to obtain the target address of the operation.

### 10.7.2 INSTRUCTIONS AFFECTED BY INDEXED LITERAL OFFSET MODE

Any of the core PIC18 instructions that can use direct addressing are potentially affected by the Indexed Literal Offset Addressing mode. This includes all byte-oriented and bit-oriented instructions, or almost one-half of the standard PIC18 instruction set. Instructions that only use Inherent or Literal Addressing modes are unaffected.

Additionally, byte-oriented and bit-oriented instructions are not affected if they do not use the Access Bank (Access RAM bit is ‘1’), or include a file address of 60h or above. Instructions meeting these criteria will continue to execute as before. A comparison of the different possible addressing modes when the extended instruction set is enabled is shown in Figure 10-7.

Those who desire to use byte-oriented or bit-oriented instructions in the Indexed Literal Offset mode should note the changes to assembler syntax for this mode. This is described in more detail in **Section 35.2.1 “Extended Instruction Syntax”**.

**REGISTER 15-2: TRISx: TRI-STATE CONTROL REGISTER**

| R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 |
|---|---|---|---|---|---|---|---|
| TRISx7 | TRISx6 | TRISx5 | TRISx4 | TRISx3 | TRISx2 | TRISx1 | TRISx0 |
| bit 7 | | | | | | | bit 0 |

| **Legend:** | | |
|---|---|---|
| R = Readable bit | W = Writable bit | U = Unimplemented bit, read as '0' |
| '1' = Bit is set | '0' = Bit is cleared | x = Bit is unknown |
| -n/n = Value at POR and BOR/Value at all other Resets | | |

bit 7-0 **TRISx<7:0>:** TRISx Port I/O Tri-state Control bits
1 = Port output driver is disabled
0 = Port output driver is enabled

**TABLE 15-3: TRIS REGISTERS**

| Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 |
|---|---|---|---|---|---|---|---|---|
| TRISA | TRISA7 | TRISA6 | TRISA5 | TRISA4 | TRISA3 | TRISA2 | TRISA1 | TRISA0 |
| TRISB | TRISB7(1) | TRISB6(1) | TRISB5 | TRISB4 | TRISB3 | TRISB2 | TRISB1 | TRISB0 |
| TRISC | TRISC7 | TRISC6 | TRISC5 | TRISC4 | TRISC3 | TRISC2 | TRISC1 | TRISC0 |
| TRISE | — | — | — | — | —(2) | — | — | — |

**Note 1:** Bits RB6 and RB7 read '1' while in Debug mode.
**2:** TRISE3 bit is read-only, and will read '1'

**REGISTER 15-4: ANSELx: ANALOG SELECT REGISTER**

| R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 | R/W-1/1 |
|---|---|---|---|---|---|---|---|
| ANSELx7 | ANSELx6 | ANSELx5 | ANSELx4 | ANSELx3 | ANSELx2 | ANSELx1 | ANSELx0 |
| bit 7 | | | | | | | bit 0 |

| **Legend:** | | |
|---|---|---|
| R = Readable bit | W = Writable bit | U = Unimplemented bit, read as '0' |
| '1' = Bit is set | '0' = Bit is cleared | x = Bit is unknown |
| -n/n = Value at POR and BOR/Value at all other Resets | | |

bit 7-0 **ANSELx<7:0>**: Analog Select on Pins Rx<7:0>

1 = Digital Input buffers are disabled.

0 = ST and TTL input devices are enabled

**TABLE 15-5: ANALOG SELECT PORT REGISTERS**

| Name | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 |
|---|---|---|---|---|---|---|---|---|
| ANSELA | ANSELA7 | ANSELA6 | ANSELA5 | ANSELA4 | ANSELA3 | ANSELA2 | ANSELA1 | ANSELA0 |
| ANSELB | ANSELB7 | ANSELB6 | ANSELB5 | ANSELB4 | ANSELB3 | ANSELB2 | ANSELB1 | ANSELB0 |
| ANSELC | ANSELC7 | ANSELC6 | ANSELC5 | ANSELC4 | ANSELC3 | ANSELC2 | ANSELC1 | ANSELC0 |

## 19.4 Timer1/3/5 Prescaler

Timer1/3/5 has four prescaler options allowing 1, 2, 4 or 8 divisions of the clock input. The CKPS bits of the TxCON register control the prescale counter. The prescale counter is not directly readable or writable; however, the prescaler counter is cleared upon a write to TMRxH or TMRxL.

## 19.5 Secondary Oscillator

A secondary low-power 32.768 kHz oscillator circuit is built-in between pins SOSCI (input) and SOSCO (amplifier output). This internal circuit is to be used in conjunction with an external 32.768 kHz crystal. The secondary oscillator is not dedicated only to Timer1/3/5; it can also be used by other modules.

The oscillator circuit is enabled by setting the SOSCEN bit of the OSCEN register (Register 4-7). This can be used as the clock source to the Timer using the TMRxCLK bits.The oscillator will continue to run during Sleep.

| | |
|---|---|
| **Note:** | The oscillator requires a start-up and stabilization time before use. Thus, the SOSCEN bit of the OSCEN register should be set and a suitable delay observed prior to enabling Timer1/3/5. A software check can be performed to confirm if the secondary oscillator is enabled and ready to use. This is done by polling the SOR bit of the OSCSTAT (Register 4-4). |

## 19.6 Timer1/3/5 Operation in Asynchronous Counter Mode

If control bit $\overline{\text{SYNC}}$ of the TxCON register is set, the external clock input is not synchronized. The timer increments asynchronously to the internal phase clocks. If external clock source is selected then the timer will continue to run during Sleep and can generate an interrupt on overflow, which will wake-up the processor. However, special precautions in software are needed to read/write the timer (see **Section 19.6.1 "Reading and Writing Timer1/3/5 in Asynchronous Counter Mode"**).

| | |
|---|---|
| **Note:** | When switching from synchronous to asynchronous operation, it is possible to skip an increment. When switching from asynchronous to synchronous operation, it is possible to produce an additional increment. |

### 19.6.1 READING AND WRITING TIMER1/3/5 IN ASYNCHRONOUS COUNTER MODE

Reading TMRxH or TMRxL while the timer is running from an external asynchronous clock will ensure a valid read (taken care of in hardware). However, the user should keep in mind that reading the 16-bit timer in two 8-bit values itself, poses certain problems, since the timer may overflow between the reads. For writes, it is recommended that the user simply stop the timer and write the desired values. A write contention may occur by writing to the timer registers, while the register is incrementing. This may produce an unpredictable value in the TMRxH:TMRxL register pair.

## 20.0 TIMER2/4/6 MODULE

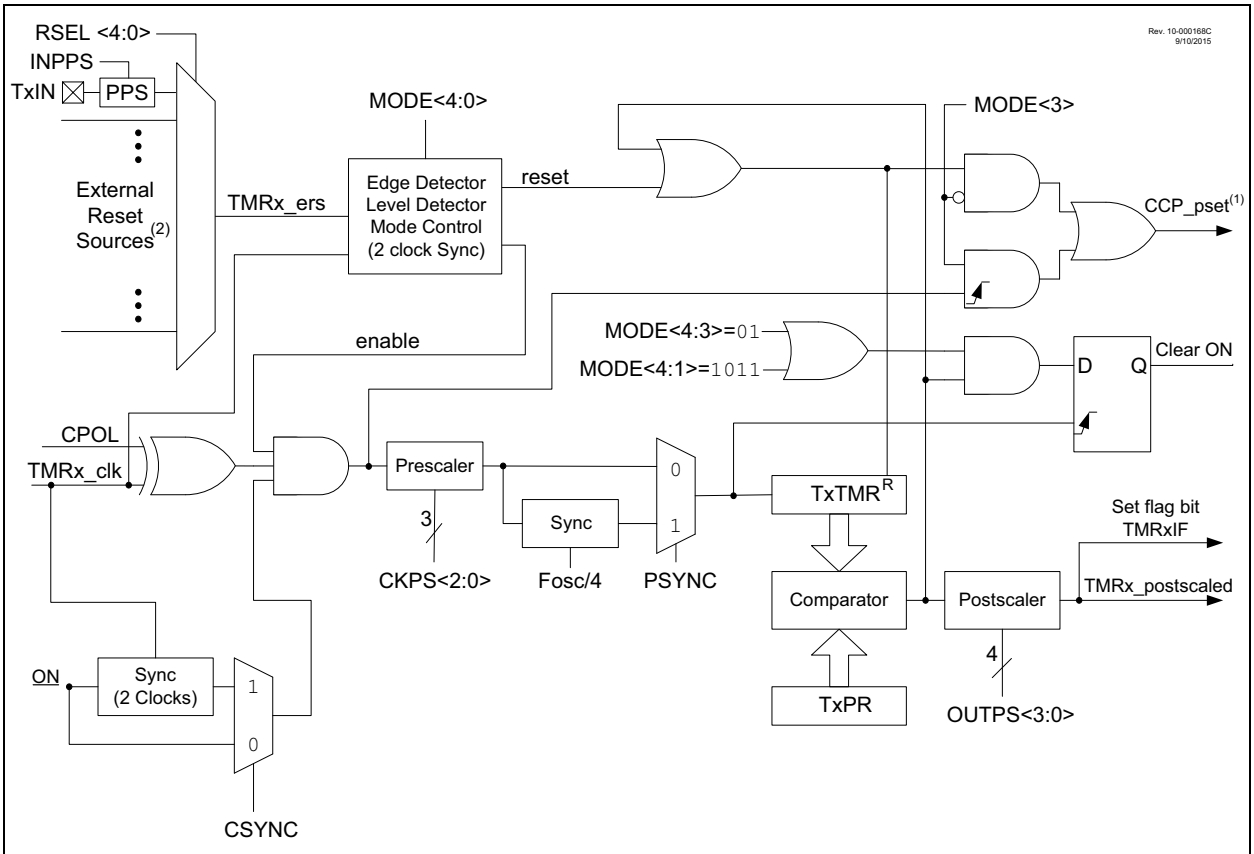
The Timer2/4/6 modules are 8-bit timers that can operate as free-running period counters or in conjunction with external signals that control start, run, freeze, and reset operation in One-Shot and Monostable modes of operation. Sophisticated waveform control such as pulse density modulation are possible by combining the operation of these timers with other internal peripherals such as the comparators and CCP modules. Features of the timer include:

- 8-bit timer register
- 8-bit period register
- Selectable external hardware timer Resets
- Programmable prescaler (1:1 to 1:128)
- Programmable postscaler (1:1 to 1:16)
- Selectable synchronous/asynchronous operation
- Alternate clock sources
- Interrupt-on-period
- Three modes of operation:
  - Free Running Period
  - One-shot
  - Monostable

See Figure 20-1 for a block diagram of Timer2. See Figure 20-2 for the clock source block diagram.

| **Note:** | Three identical Timer2 modules are implemented on this device. The timers are named Timer2, Timer4 and Timer6. All references to Timer2 apply as well to Timer4 and Timer6. All references to PR2 apply equally to other timers as well. |
|---|---|

**FIGURE 20-1: TIMER2 BLOCK DIAGRAM**




**Note 1:** Signal to the CCP to trigger the PWM pulse.
**2:** See Table 20.5 for description of CCP interaction in the different TMR modes.
**3:** See Table 20.5 for external Reset sources.

## 20.6 Timer2 Operation During Sleep

When PSYNC = 1, Timer2 cannot be operated while the processor is in Sleep mode. The contents of the TMR2 and T2PR registers will remain unchanged while processor is in Sleep mode.

When PSYNC = 0, Timer2 will operate in Sleep as long as the clock source selected is also still running. Selecting the LFINTOSC, MFINTOSC, or HFINTOSC oscillator as the timer clock source will keep the selected oscillator running during Sleep.

**FIGURE 25-2: On Off Keying (OOK) Synchronization**

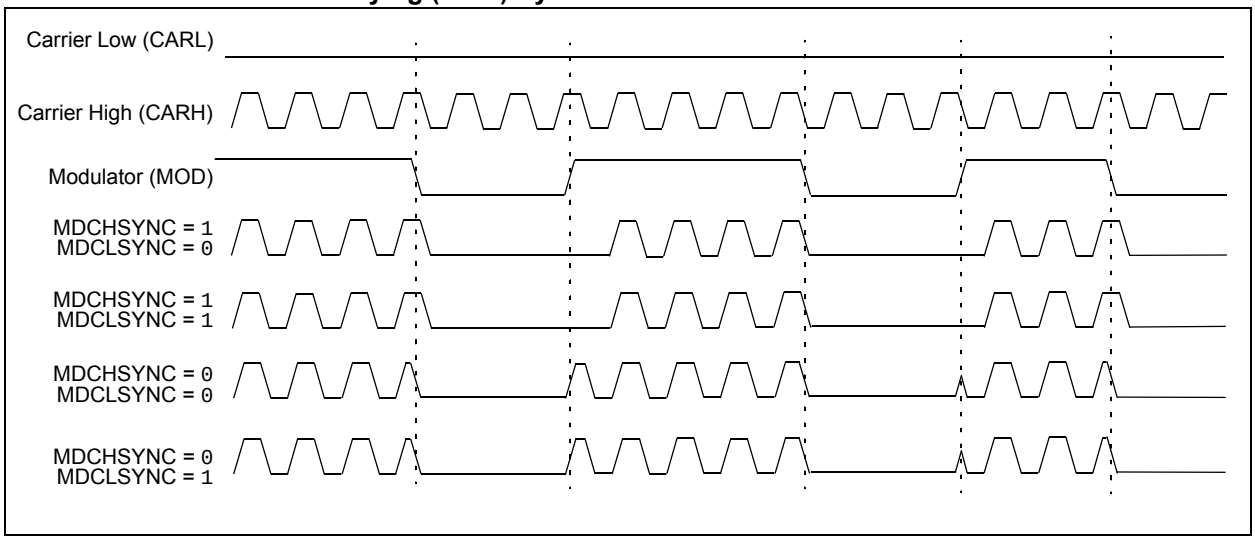



**FIGURE 25-3: No Synchronization (MDSHSYNC = 0, MDCLSYNC = 0)**

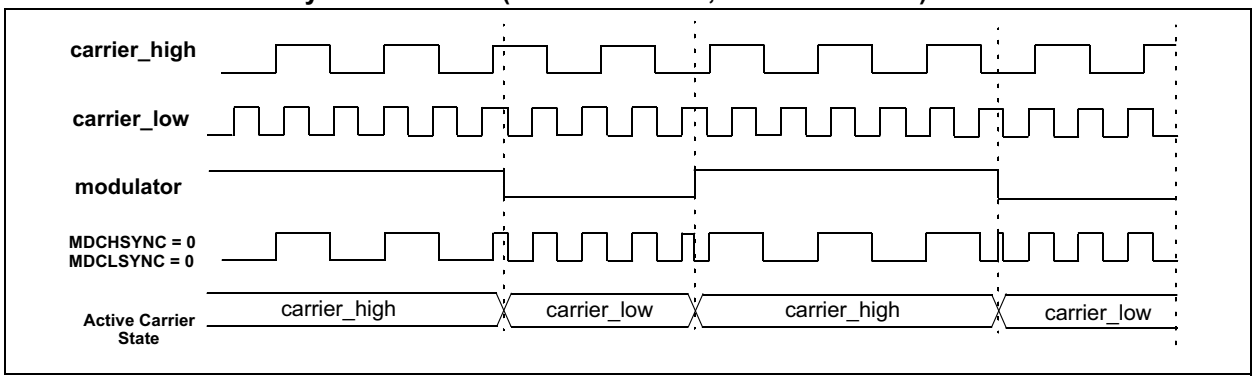



**FIGURE 25-4: Carrier High Synchronization (MDSHSYNC = 1, MDCLSYNC = 0)**

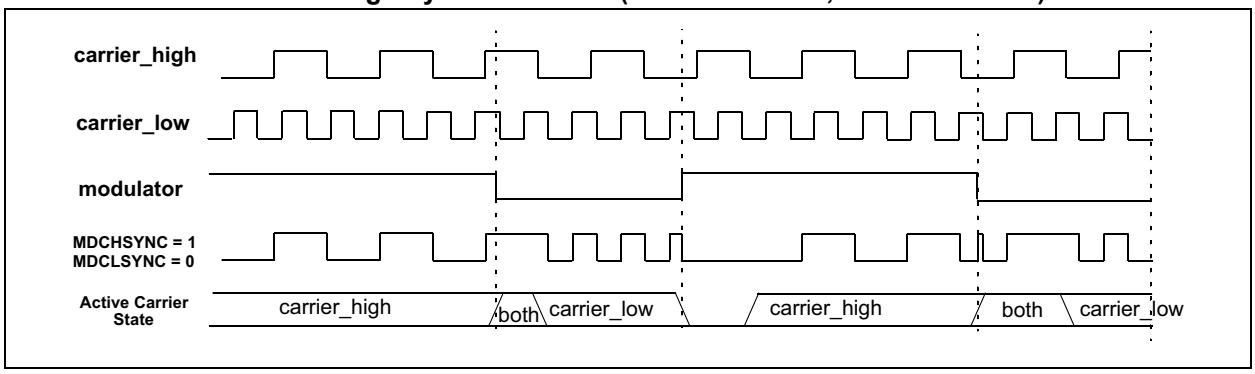
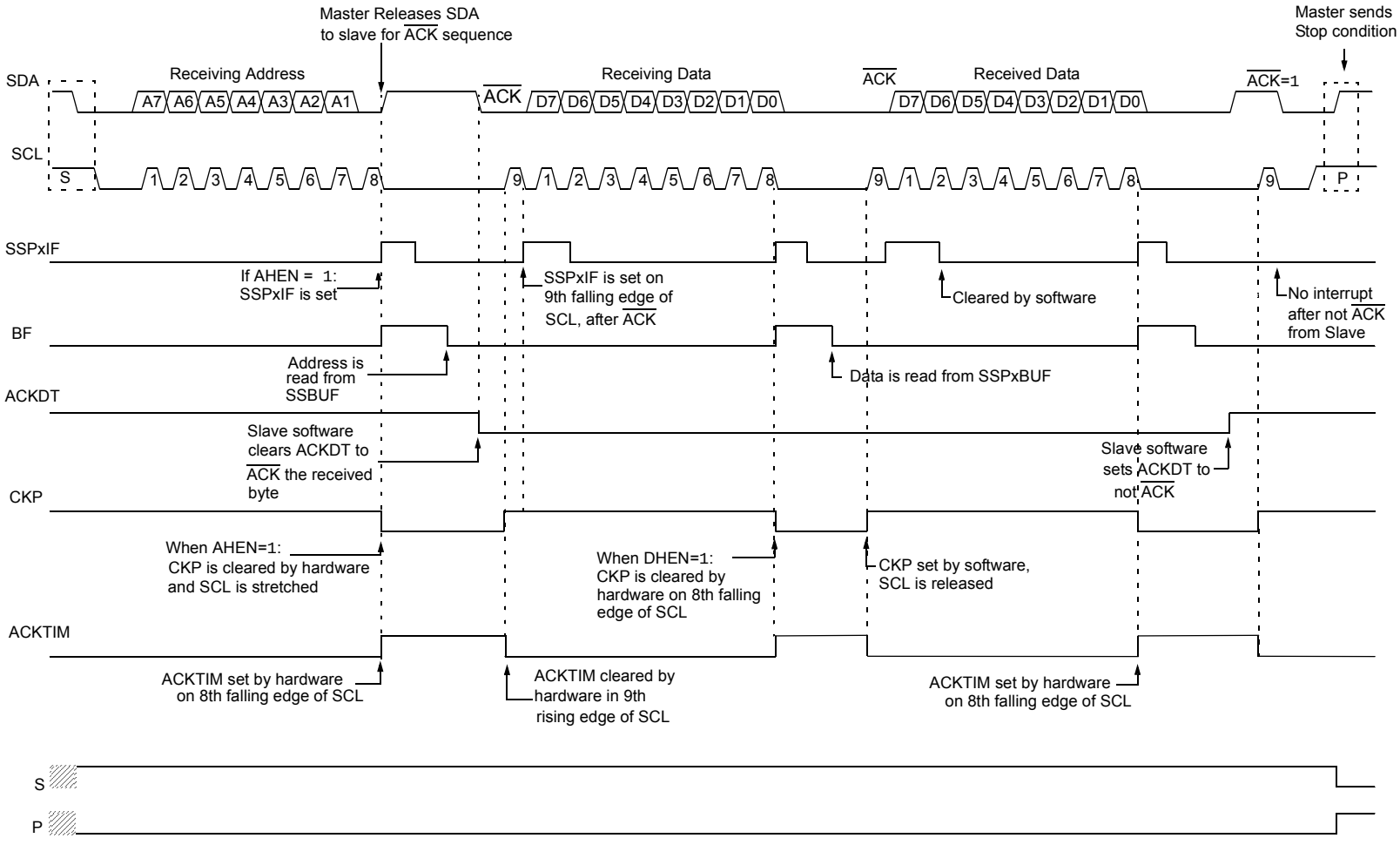


**FIGURE 26-16: I²C SLAVE, 7-BIT ADDRESS, RECEPTION (SEN = 0, AHEN = 1, DHEN = 1)**



#### 26.9.3.3 7-bit Transmission with Address Hold Enabled

Setting the AHEN bit of the SSPxCON3 register enables additional clock stretching and interrupt generation after the eighth falling edge of a received matching address. Once a matching address has been clocked in, CKP is cleared and the SSPxIF interrupt is set.

Figure 26-19 displays a standard waveform of a 7-bit address slave transmission with AHEN enabled.

1. Bus starts Idle.
2. Master sends Start condition; the S bit of SSPxSTAT is set; SSPxIF is set if interrupt on Start detect is enabled.
3. Master sends matching address with R/$\overline{W}$ bit set. After the eighth falling edge of the SCL line the CKP bit is cleared and SSPxIF interrupt is generated.
4. Slave software clears SSPxIF.
5. Slave software reads the ACKTIM bit of SSPxCON3 register, and R/$\overline{W}$ and D/$\overline{A}$ of the SSPxSTAT register to determine the source of the interrupt.
6. Slave reads the address value from the SSPxBUF register clearing the BF bit.
7. Slave software decides from this information if it wishes to $\overline{ACK}$ or not $\overline{ACK}$ and sets the ACKDT bit of the SSPxCON2 register accordingly.
8. Slave sets the CKP bit releasing SCL.
9. Master clocks in the $\overline{ACK}$ value from the slave.
10. Slave hardware automatically clears the CKP bit and sets SSPxIF after the $\overline{ACK}$ if the R/$\overline{W}$ bit is set.
11. Slave software clears SSPxIF.
12. Slave loads value to transmit to the master into SSPxBUF setting the BF bit.

**Note:** SSPxBUF cannot be loaded until after the $\overline{ACK}$.

13. Slave sets the CKP bit releasing the clock.
14. Master clocks out the data from the slave and sends an $\overline{ACK}$ value on the ninth SCL pulse.
15. Slave hardware copies the $\overline{ACK}$ value into the ACKSTAT bit of the SSPxCON2 register.
16. Steps 10-15 are repeated for each byte transmitted to the master from the slave.
17. If the master sends a not $\overline{ACK}$ the slave releases the bus allowing the master to send a Stop and end the communication.

**Note:** Master must send a not $\overline{ACK}$ on the last byte to ensure that the slave releases the SCL line to receive a Stop.

After the write to the SSPxBUF, each bit of the address will be shifted out on the falling edge of SCL until all seven address bits and the R/$\overline{W}$ bit are completed. On the falling edge of the eighth clock, the master will release the SDA pin, allowing the slave to respond with an Acknowledge. On the falling edge of the ninth clock, the master will sample the SDA pin to see if the address was recognized by a slave. The status of the $\overline{ACK}$ bit is loaded into the ACKSTAT Status bit of the SSPxCON2 register. Following the falling edge of the ninth clock transmission of the address, the SSPxIF is set, the BF flag is cleared and the Baud Rate Generator is turned off until another write to the SSPxBUF takes place, holding SCL low and allowing SDA to float.

#### 26.10.6.1 BF Status Flag

In Transmit mode, the BF bit of the SSPxSTAT register is set when the CPU writes to SSPxBUF and is cleared when all eight bits are shifted out.

#### 26.10.6.2 WCOL Status Flag

If the user writes the SSPxBUF when a transmit is already in progress (i.e., SSPSR is still shifting out a data byte), the WCOL bit is set and the contents of the buffer are unchanged (the write does not occur).

The WCOL bit must be cleared by software before the next transmission.

#### 26.10.6.3 ACKSTAT Status Flag

In Transmit mode, the ACKSTAT bit of the SSPxCON2 register is cleared when the slave has sent an Acknowledge ($\overline{ACK}$ = 0) and is set when the slave does not Acknowledge ($\overline{ACK}$ = 1). A slave sends an Acknowledge when it has recognized its address (including a general call), or when the slave has properly received its data.

#### 26.10.6.4 Typical transmit sequence:

1. The user generates a Start condition by setting the SEN bit of the SSPxCON2 register.
2. SSPxIF is set by hardware on completion of the Start.
3. SSPxIF is cleared by software.
4. The MSSP module will wait the required start time before any other operation takes place.
5. The user loads the SSPxBUF with the slave address to transmit.
6. Address is shifted out the SDA pin until all eight bits are transmitted. Transmission begins as soon as SSPxBUF is written to.
7. The MSSP module shifts in the $\overline{ACK}$ bit from the slave device and writes its value into the ACKSTAT bit of the SSPxCON2 register.
8. The MSSP module generates an interrupt at the end of the ninth clock cycle by setting the SSPxIF bit.
9. The user loads the SSPxBUF with eight bits of data.
10. Data is shifted out the SDA pin until all eight bits are transmitted.
11. The MSSP module shifts in the $\overline{ACK}$ bit from the slave device and writes its value into the ACKSTAT bit of the SSPxCON2 register.
12. Steps 8-11 are repeated for all transmitted data bytes.
13. The user generates a Stop or Restart condition by setting the PEN or RSEN bits of the SSPxCON2 register. Interrupt is generated once the Stop/Restart condition is complete.

## 27.4 EUSART Baud Rate Generator (BRG)

The Baud Rate Generator (BRG) is an 8-bit or 16-bit timer that is dedicated to the support of both the asynchronous and synchronous EUSART operation. By default, the BRG operates in 8-bit mode. Setting the BRG16 bit of the BAUDxCON register selects 16-bit mode.

The SPxBRGH, SPxBRGL register pair determines the period of the free running baud rate timer. In Asynchronous mode the multiplier of the baud rate period is determined by both the BRGH bit of the TXxSTA register and the BRG16 bit of the BAUDxCON register. In Synchronous mode, the BRGH bit is ignored.

Table 27-3 contains the formulas for determining the baud rate. Example 27-1 provides a sample calculation for determining the baud rate and baud rate error.

Typical baud rates and error values for various asynchronous modes have been computed for your convenience and are shown in Table 27-5. It may be advantageous to use the high baud rate (BRGH = 1), or the 16-bit BRG (BRG16 = 1) to reduce the baud rate error. The 16-bit BRG mode is used to achieve slow baud rates for fast oscillator frequencies.

Writing a new value to the SPxBRGH, SPxBRGL register pair causes the BRG timer to be reset (or cleared). This ensures that the BRG does not wait for a timer overflow before outputting the new baud rate.

If the system clock is changed during an active receive operation, a receive error or data loss may result. To avoid this problem, check the status of the RCIDL bit to make sure that the receive operation is idle before changing the system clock.

**EXAMPLE 27-1: CALCULATING BAUD RATE ERROR**

For a device with FOSC of 16 MHz, desired baud rate of 9600, Asynchronous mode, 8-bit BRG:

$$Desired\ Baud\ Rate = \frac{FOSC}{64([SPxBRGH:SPxBRGL] + 1)}$$

Solving for SPxBRGH:SPxBRGL:

$$SPBRGH:SPBRGL = \frac{\frac{FOSC}{Desired\ Baud\ Rate}}{64} - 1$$

$$= \frac{\frac{16000000}{9600}}{64} - 1$$

$$= [25.042] = 25$$

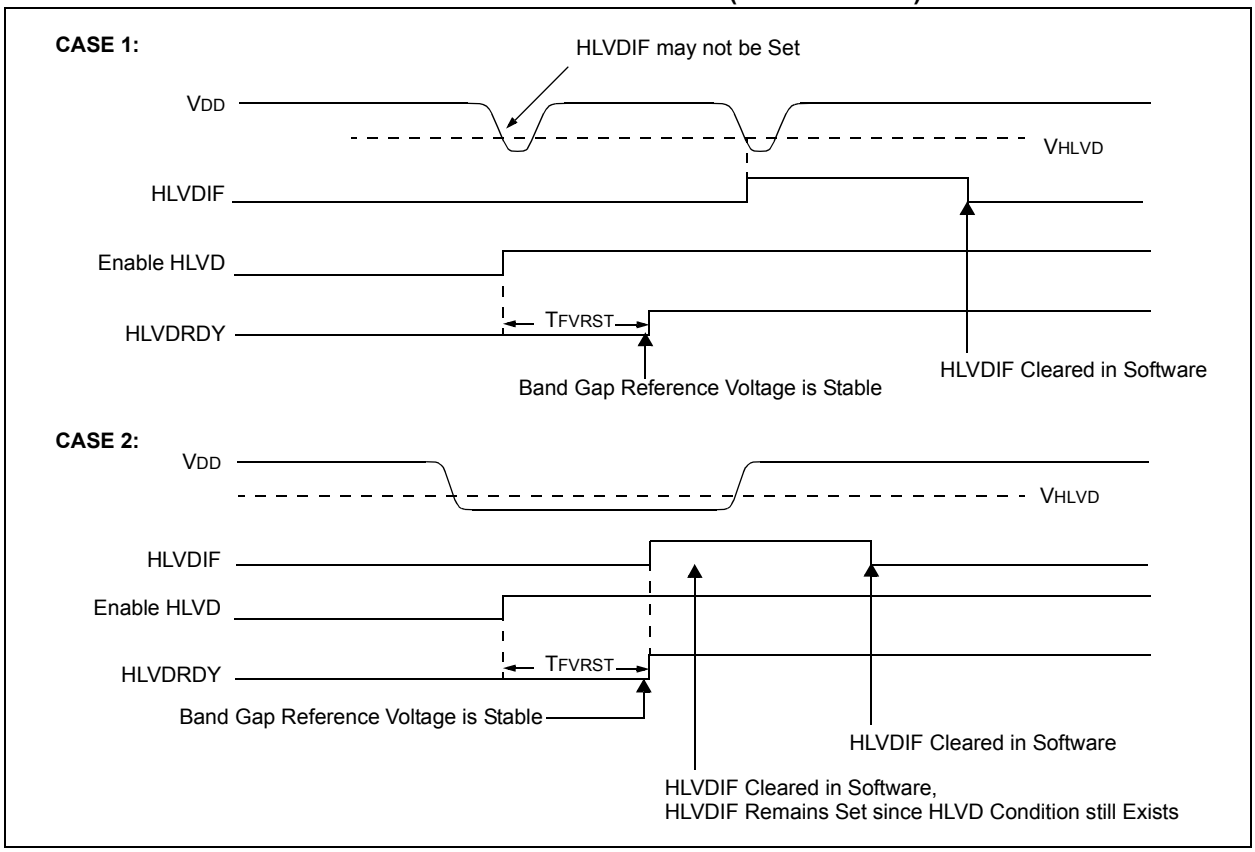
$$Calculated\ Baud\ Rate = \frac{16000000}{64(25 + 1)}$$

$$= 9615$$

$$Error = \frac{Calc.\ Baud\ Rate - Desired\ Baud\ Rate}{Desired\ Baud\ Rate}$$

$$= \frac{(9615 - 9600)}{9600} = 0.16\%$$

**FIGURE 33-2: LOW-VOLTAGE DETECT OPERATION (HLVDINTL = 1)**



**RRNCF** **Rotate Right f (No Carry)**

Syntax: RRNCF f {,d {,a}}

Operands: $0 \le f \le 255$
$d \in [0,1]$
$a \in [0,1]$

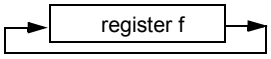
Operation: (f<n>) → dest<n – 1>,
(f<0>) → dest<7>

Status Affected: N, Z

Encoding:

| 0100 | 00da | ffff | ffff |
|---|---|---|---|

Description: The contents of register 'f' are rotated one bit to the right. If 'd' is '0', the result is placed in W. If 'd' is '1', the result is placed back in register 'f' (default).
If 'a' is '0', the Access Bank will be selected (default), overriding the BSR value. If 'a' is '1', then the bank will be selected as per the BSR value.
If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever f ≤ 95 (5Fh). See **Section 35.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.




Words: 1

Cycles: 1

Q Cycle Activity:

| Q1 | Q2 | Q3 | Q4 |
|---|---|---|---|
| Decode | Read register 'f' | Process Data | Write to destination |

Example 1:
```
RRNCF   REG, 1, 0
Before Instruction
    REG     =   1101 0111
After Instruction
    REG     =   1110 1011
```

Example 2:
```
RRNCF   REG, 0, 0
Before Instruction
    W       =   ?
    REG     =   1101 0111
After Instruction
    W       =   1110 1011
    REG     =   1101 0111
```

**SETF** **Set f**

Syntax: SETF f {,a}

Operands: $0 \le f \le 255$
$a \in [0,1]$

Operation: FFh → f

Status Affected: None

Encoding:

| 0110 | 100a | ffff | ffff |
|---|---|---|---|

Description: The contents of the specified register are set to FFh.
If 'a' is '0', the Access Bank is selected.
If 'a' is '1', the BSR is used to select the GPR bank.
If 'a' is '0' and the extended instruction set is enabled, this instruction operates in Indexed Literal Offset Addressing mode whenever f ≤ 95 (5Fh). See **Section 35.2.3 "Byte-Oriented and Bit-Oriented Instructions in Indexed Literal Offset Mode"** for details.

Words: 1

Cycles: 1

Q Cycle Activity:

| Q1 | Q2 | Q3 | Q4 |
|---|---|---|---|
| Decode | Read register 'f' | Process Data | Write register 'f' |

Example:
```
SETF    REG, 1
Before Instruction
    REG     =   5Ah
After Instruction
    REG     =   FFh
```

**TBLWT** **Table Write**

Syntax: TBLWT ( *; *+; *-; +*)

Operands: None

Operation: if TBLWT*,
(TABLAT) → Holding Register;
TBLPTR – No Change;
if TBLWT*+,
(TABLAT) → Holding Register;
(TBLPTR) + 1 → TBLPTR;
if TBLWT*-,
(TABLAT) → Holding Register;
(TBLPTR) – 1 → TBLPTR;
if TBLWT+*,
(TBLPTR) + 1 → TBLPTR;
(TABLAT) → Holding Register;

Status Affected: None

Encoding:

| | | | |
|---|---|---|---|
| 0000 | 0000 | 0000 | 11nn<br>nn=0 *<br>=1 *+<br>=2 *-<br>=3 +* |

Description: This instruction uses the three LSBs of TBLPTR to determine which of the eight holding registers the TABLAT is written to. The holding registers are used to program the contents of Program Memory (P.M.). (Refer to **Section 11.1 "Program Flash Memory"** for additional details on programming Flash memory.)
The TBLPTR (a 21-bit pointer) points to each byte in the program memory. TBLPTR has a 2-MByte address range. The LSb of the TBLPTR selects which byte of the program memory location to access.

TBLPTR[0] = 0: Least Significant Byte of Program Memory Word
TBLPTR[0] = 1: Most Significant Byte of Program Memory Word

The TBLWT instruction can modify the value of TBLPTR as follows:

- no change
- post-increment
- post-decrement
- pre-increment

Words: 1

Cycles: 2

Q Cycle Activity:

| Q1 | Q2 | Q3 | Q4 |
|---|---|---|---|
| Decode | No operation | No operation | No operation |
| No operation | No operation (Read TABLAT) | No operation | No operation (Write to Holding Register ) |

**TBLWT** **Table Write (Continued)**

Example1:
```
TBLWT *+;
```
Before Instruction
TABLAT = 55h
TBLPTR = 00A356h
HOLDING REGISTER (00A356h) = FFh
After Instructions (table write completion)
TABLAT = 55h
TBLPTR = 00A357h
HOLDING REGISTER (00A356h) = 55h

Example 2:
```
TBLWT +*;
```
Before Instruction
TABLAT = 34h
TBLPTR = 01389Ah
HOLDING REGISTER (01389Ah) = FFh
HOLDING REGISTER (01389Bh) = FFh
After Instruction (table write completion)
TABLAT = 34h
TBLPTR = 01389Bh
HOLDING REGISTER (01389Ah) = FFh
HOLDING REGISTER (01389Bh) = 34h

**TABLE 37-5: MEMORY PROGRAMMING SPECIFICATIONS**

| Standard Operating Conditions (unless otherwise stated) | | | | | | | |
|---|---|---|---|---|---|---|---|
| **Param No.** | **Sym.** | **Characteristic** | **Min.** | **Typ†** | **Max.** | **Units** | **Conditions** |
| **Data EEPROM Memory Specifications** | | | | | | | |
| MEM20 | $E_D$ | DataEE Byte Endurance | 100k | — | — | E/W | -40°C ≤ TA ≤ +85°C |
| MEM21 | $T_{D\_RET}$ | Characteristic Retention | — | 40 | — | Year | Provided no other specifications are violated |
| MEM22 | $N_{D\_REF}$ | Total Erase/Write Cycles before Refresh | 1M<br>500k | 10M<br>— | —<br>— | E/W | -40°C ≤ TA ≤ +60°C<br>-40°C ≤ TA ≤ +85°C |
| MEM23 | $V_{D\_RW}$ | VDD for Read or Erase/Write operation | VDDMIN | — | VDDMAX | V | |
| MEM24 | $T_{D\_BEW}$ | Byte Erase and Write Cycle Time | — | 4.0 | 5.0 | ms | |
| **Program Flash Memory Specifications** | | | | | | | |
| MEM30 | $E_P$ | Flash Memory Cell Endurance | 10k | — | — | E/W | -40°C ≤ TA ≤ +85°C (**Note 1**) |
| MEM32 | $T_{P\_RET}$ | Characteristic Retention | — | 40 | — | Year | Provided no other specifications are violated |
| MEM33 | $V_{P\_RD}$ | VDD for Read operation | VDDMIN | — | VDDMAX | V | |
| MEM34 | $V_{P\_REW}$ | VDD for Row Erase or Write operation | VDDMIN | — | VDDMAX | V | |
| MEM35 | $T_{P\_REW}$ | Self-Timed Row Erase or Self-Timed Write | — | 2.0 | 2.5 | ms | |

† Data in “Typ” column is at 3.0V, 25°C unless otherwise stated. These parameters are for design guidance only and are not tested.

**Note 1:** Flash Memory Cell Endurance for the Flash memory is defined as: One Row Erase operation and one Self-Timed Write.

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