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Details

Product Status	Active
Core Processor	PIC
Core Size	8-Bit
Speed	20MHz
Connectivity	-
Peripherals	Brown-out Detect/Reset, POR, WDT
Number of I/O	13
Program Memory Size	3.5KB (2K x 14)
Program Memory Type	OTP
EEPROM Size	-
RAM Size	128 x 8
Voltage - Supply (Vcc/Vdd)	3V ~ 6V
Data Converters	-
Oscillator Type	External
Operating Temperature	-40°C ~ 85°C (TA)
Mounting Type	Surface Mount
Package / Case	18-SOIC (0.295", 7.50mm Width)
Supplier Device Package	18-SOIC
Purchase URL	https://www.e-xfl.com/product-detail/microchip-technology/pic16c622-20i-so

PIC16C62X

NOTES:

PIC16C62X

TABLE 1-1: PIC16C62X FAMILY OF DEVICES

		PIC16C620 ⁽³⁾	PIC16C620A ⁽¹⁾⁽⁴⁾	PIC16CR620A ⁽²⁾	PIC16C621 ⁽³⁾	PIC16C621A ⁽¹⁾⁽⁴⁾	PIC16C622 ⁽³⁾	PIC16C622A ⁽¹⁾⁽⁴⁾
Clock	Maximum Frequency of Operation (MHz)	20	40	20	20	40	20	40
Memory	EPROM Program Memory (x14 words)	512	512	512	1K	1K	2K	2K
	Data Memory (bytes)	80	96	96	80	96	128	128
Peripherals	Timer Module(s)	TMR0	TMR0	TMRO	TMR0	TMR0	TMR0	TMR0
	Comparators(s)	2	2	2	2	2	2	2
	Internal Reference Voltage	Yes	Yes	Yes	Yes	Yes	Yes	Yes
Features	Interrupt Sources	4	4	4	4	4	4	4
	I/O Pins	13	13	13	13	13	13	13
	Voltage Range (Volts)	2.5-6.0	2.7-5.5	2.5-5.5	2.5-6.0	2.7-5.5	2.5-6.0	2.7-5.5
	Brown-out Reset	Yes	Yes	Yes	Yes	Yes	Yes	Yes
	Packages	18-pin DIP, SOIC, 20-pin SSOP	18-pin DIP, SOIC, 20-pin SSOP	18-pin DIP, SOIC, 20-pin SSOP	18-pin DIP, SOIC, 20-pin SSOP	18-pin DIP, SOIC, 20-pin SSOP	18-pin DIP, SOIC, 20-pin SSOP	18-pin DIP, SOIC, 20-pin SSOP

All PICmicro[®] Family devices have Power-on Reset, selectable Watchdog Timer, selectable code protect and high I/O current capability. All PIC16C62X Family devices use serial programming with clock pin RB6 and data pin RB7.

- Note 1:** If you change from this device to another device, please verify oscillator characteristics in your application.
- 2:** For ROM parts, operation from 2.0V - 2.5V will require the PIC16LCR62XA parts.
- 3:** For OTP parts, operation from 2.5V - 3.0V will require the PIC16LC62X part.
- 4:** For OTP parts, operation from 2.7V - 3.0V will require the PIC16LC62XA part.

3.0 ARCHITECTURAL OVERVIEW

The high performance of the PIC16C62X family can be attributed to a number of architectural features commonly found in RISC microprocessors. To begin with, the PIC16C62X uses a Harvard architecture, in which, program and data are accessed from separate memories using separate busses. This improves bandwidth over traditional von Neumann architecture, where program and data are fetched from the same memory. Separating program and data memory further allows instructions to be sized differently than 8-bit wide data word. Instruction opcodes are 14-bits wide making it possible to have all single word instructions. A 14-bit wide program memory access bus fetches a 14-bit instruction in a single cycle. A two-stage pipeline overlaps fetch and execution of instructions. Consequently, all instructions (35) execute in a single cycle (200 ns @ 20 MHz) except for program branches.

The PIC16C620(A) and PIC16CR620A address 512 x 14 on-chip program memory. The PIC16C621(A) addresses 1K x 14 program memory. The PIC16C622(A) addresses 2K x 14 program memory. All program memory is internal.

The PIC16C62X can directly or indirectly address its register files or data memory. All special function registers including the program counter are mapped in the data memory. The PIC16C62X has an orthogonal (symmetrical) instruction set that makes it possible to carry out any operation on any register using any Addressing mode. This symmetrical nature and lack of 'special optimal situations' make programming with the PIC16C62X simple yet efficient. In addition, the learning curve is reduced significantly.

The PIC16C62X devices contain an 8-bit ALU and working register. The ALU is a general purpose arithmetic unit. It performs arithmetic and Boolean functions between data in the working register and any register file.

The ALU is 8-bits wide and capable of addition, subtraction, shift and logical operations. Unless otherwise mentioned, arithmetic operations are two's complement in nature. In two-operand instructions, typically one operand is the working register (W register). The other operand is a file register or an immediate constant. In single operand instructions, the operand is either the W register or a file register.

The W register is an 8-bit working register used for ALU operations. It is not an addressable register.

Depending on the instruction executed, the ALU may affect the values of the Carry (C), Digit Carry (DC), and Zero (Z) bits in the STATUS register. The C and DC bits operate as a Borrow and Digit Borrow out bit, respectively, bit in subtraction. See the `SUBLW` and `SUBWF` instructions for examples.

A simplified block diagram is shown in Figure 3-1, with a description of the device pins in Table 3-1.

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3.1 Clocking Scheme/Instruction Cycle

The clock input (OSC1/CLKIN pin) is internally divided by four to generate four non-overlapping quadrature clocks namely Q1, Q2, Q3 and Q4. Internally, the program counter (PC) is incremented every Q1, the instruction is fetched from the program memory and latched into the instruction register in Q4. The instruction is decoded and executed during the following Q1 through Q4. The clocks and instruction execution flow is shown in Figure 3-2.

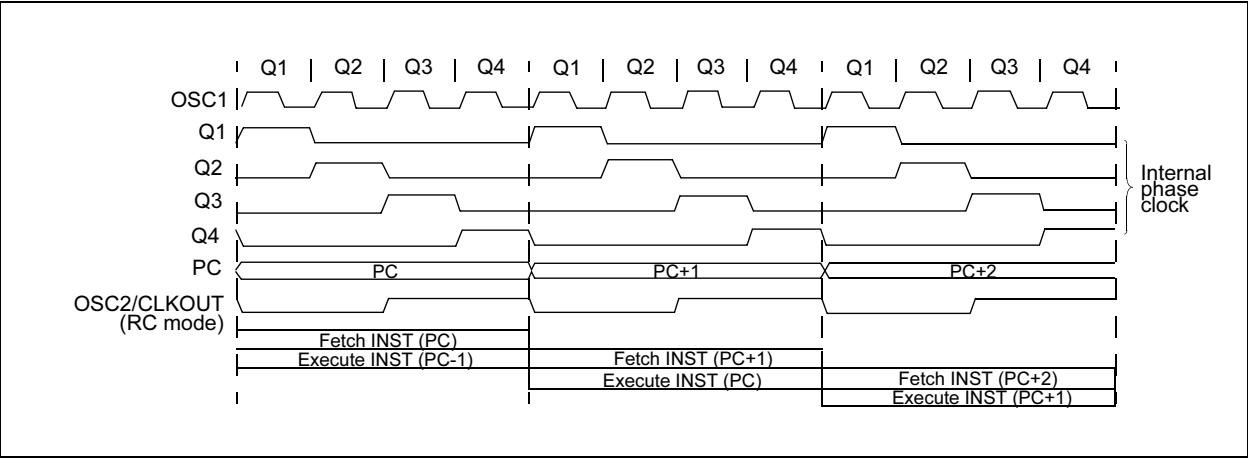
3.2 Instruction Flow/Pipelining

An "Instruction Cycle" consists of four Q cycles (Q1, Q2, Q3 and Q4). The instruction fetch and execute are pipelined such that fetch takes one instruction cycle while decode and execute takes another instruction cycle. However, due to the pipelining, each instruction effectively executes in one cycle. If an instruction causes the program counter to change (e.g., GOTO) then two cycles are required to complete the instruction (Example 3-1).

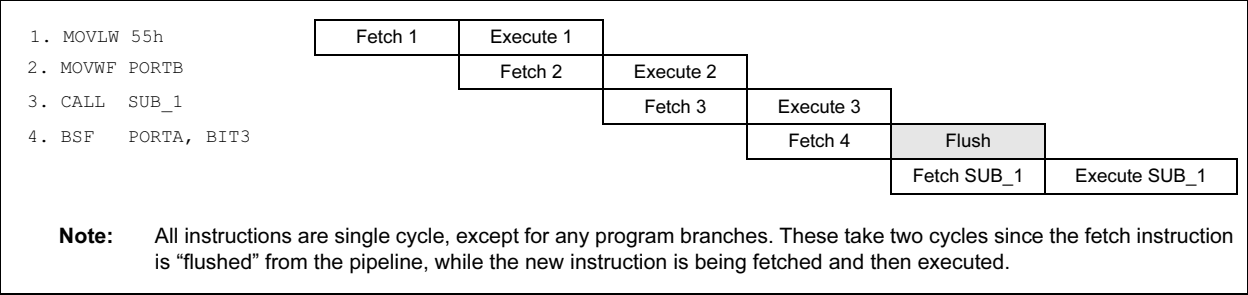
A fetch cycle begins with the program counter (PC) incrementing in Q1.

In the execution cycle, the fetched instruction is latched into the "Instruction Register (IR)" in cycle Q1. This instruction is then decoded and executed during the Q2, Q3 and Q4 cycles. Data memory is read during Q2 (operand read) and written during Q4 (destination write).

FIGURE 3-2: CLOCK/INSTRUCTION CYCLE



EXAMPLE 3-1: INSTRUCTION PIPELINE FLOW



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4.2 Data Memory Organization

The data memory (Figure 4-4, Figure 4-5, Figure 4-6 and Figure 4-7) is partitioned into two banks, which contain the General Purpose Registers and the Special Function Registers. Bank 0 is selected when the RP0 bit is cleared. Bank 1 is selected when the RP0 bit (STATUS <5>) is set. The Special Function Registers are located in the first 32 locations of each bank. Register locations 20-7Fh (Bank0) on the PIC16C620A/CR620A/621A and 20-7Fh (Bank0) and A0-BFh (Bank1) on the PIC16C622 and PIC16C622A are General Purpose Registers implemented as static RAM. Some Special Purpose Registers are mapped in Bank 1.

Addresses F0h-FFh of bank1 are implemented as common ram and mapped back to addresses 70h-7Fh in bank0 on the PIC16C620A/621A/622A/CR620A.

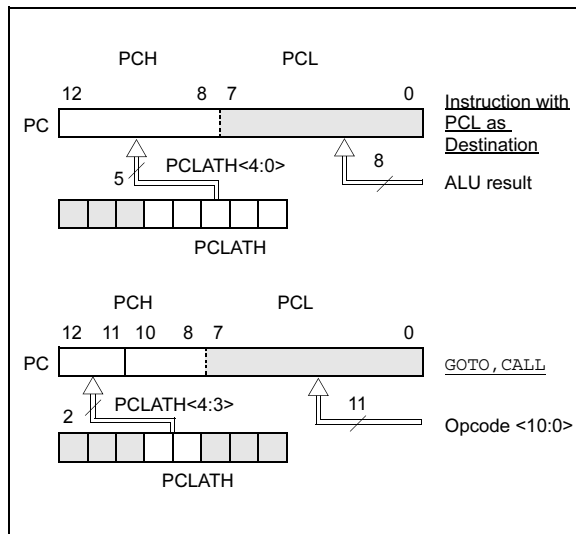
4.2.1 GENERAL PURPOSE REGISTER FILE

The register file is organized as 80 x 8 in the PIC16C620/621, 96 x 8 in the PIC16C620A/621A/CR620A and 128 x 8 in the PIC16C622(A). Each is accessed either directly or indirectly through the File Select Register FSR (Section 4.4).

4.3 PCL and PCLATH

The program counter (PC) is 13-bits wide. The low byte comes from the PCL register, which is a readable and writable register. The high byte (PC<12:8>) is not directly readable or writable and comes from PCLATH. On any RESET, the PC is cleared. Figure 4-8 shows the two situations for the loading of the PC. The upper example in the figure shows how the PC is loaded on a write to PCL (PCLATH<4:0> → PCH). The lower example in the figure shows how the PC is loaded during a CALL or GOTO instruction (PCLATH<4:3> → PCH).

FIGURE 4-8: LOADING OF PC IN DIFFERENT SITUATIONS



4.3.1 COMPUTED GOTO

A computed GOTO is accomplished by adding an offset to the program counter (ADDWF PCL). When doing a table read using a computed GOTO method, care should be exercised if the table location crosses a PCL memory boundary (each 256 byte block). Refer to the application note, "Implementing a Table Read" (AN556).

4.3.2 STACK

The PIC16C62X family has an 8-level deep x 13-bit wide hardware stack (Figure 4-2 and Figure 4-3). The stack space is not part of either program or data space and the stack pointer is not readable or writable. The PC is PUSHed onto the stack when a CALL instruction is executed or an interrupt causes a branch. The stack is POPed in the event of a RETURN, RETLW or a RETFIE instruction execution. PCLATH is not affected by a PUSH or POP operation.

The stack operates as a circular buffer. This means that after the stack has been PUSHed eight times, the ninth push overwrites the value that was stored from the first push. The tenth push overwrites the second push (and so on).

Note 1: There are no STATUS bits to indicate stack overflow or stack underflow conditions.

2: There are no instructions/mnemonics called PUSH or POP. These are actions that occur from the execution of the CALL, RETURN, RETLW and RETFIE instructions, or the vectoring to an interrupt address.

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FIGURE 5-3: BLOCK DIAGRAM OF RA3 PIN

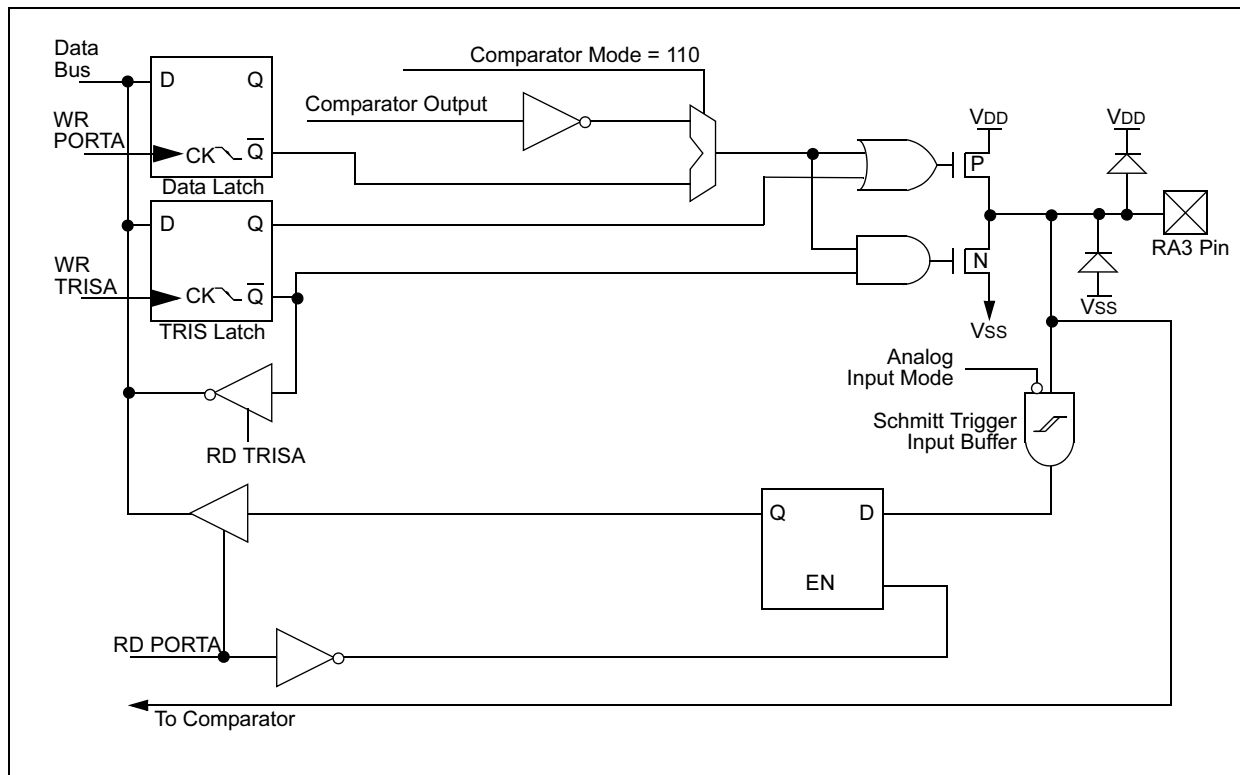
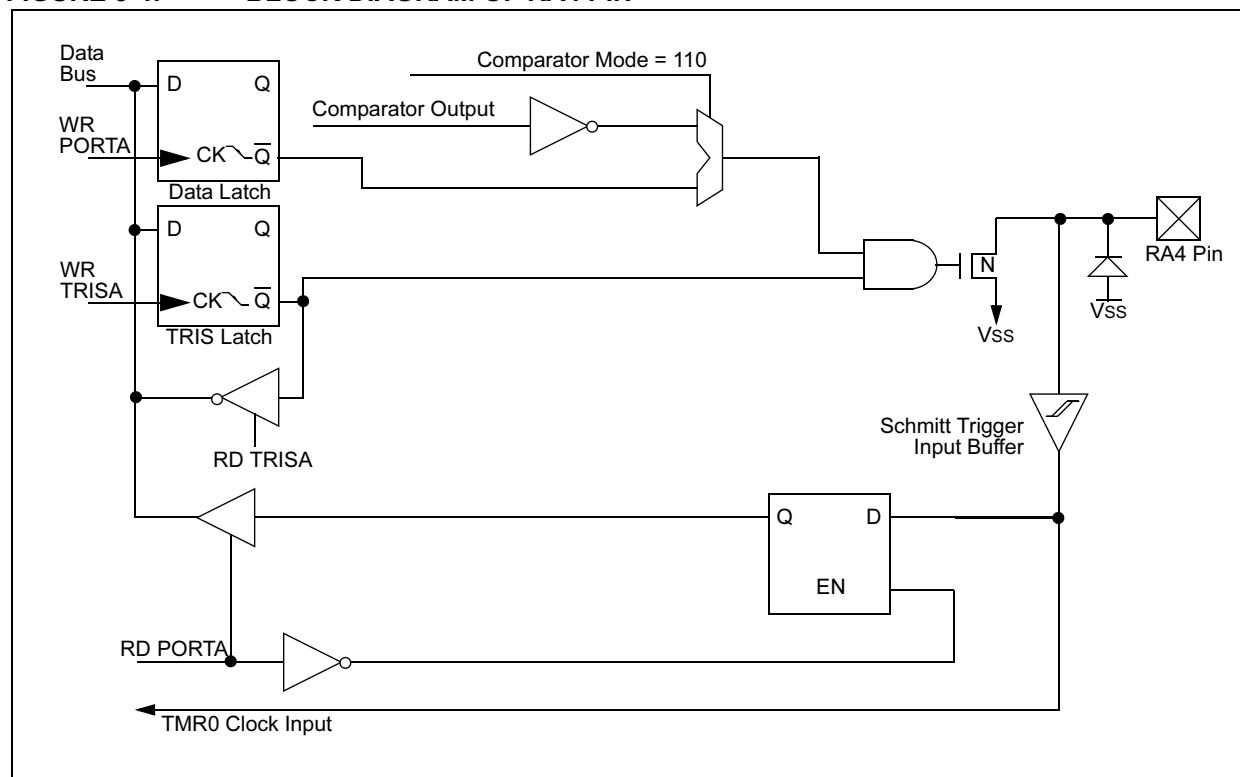


FIGURE 5-4: BLOCK DIAGRAM OF RA4 PIN



PIC16C62X

5.2 PORTB and TRISB Registers

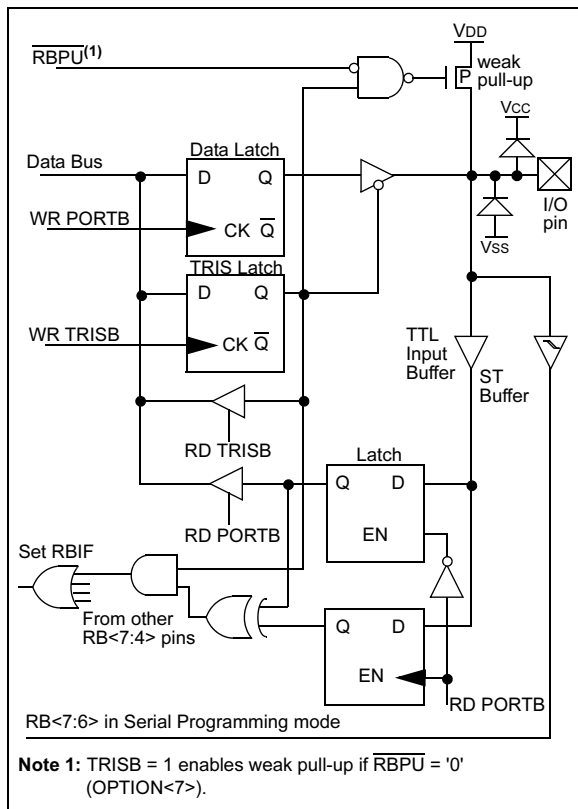
PORTB is an 8-bit wide, bi-directional port. The corresponding data direction register is TRISB. A '1' in the TRISB register puts the corresponding output driver in a High Impedance mode. A '0' in the TRISB register puts the contents of the output latch on the selected pin(s).

Reading PORTB register reads the status of the pins, whereas writing to it will write to the port latch. All write operations are read-modify-write operations. So a write to a port implies that the port pins are first read, then this value is modified and written to the port data latch.

Each of the PORTB pins has a weak internal pull-up ($\approx 200 \mu\text{A}$ typical). A single control bit can turn on all the pull-ups. This is done by clearing the RBPV (OPTION<7>) bit. The weak pull-up is automatically turned off when the port pin is configured as an output. The pull-ups are disabled on Power-on Reset.

Four of PORTB's pins, RB<7:4>, have an interrupt on change feature. Only pins configured as inputs can cause this interrupt to occur (e.g., any RB<7:4> pin configured as an output is excluded from the interrupt on change comparison). The input pins (of RB<7:4>) are compared with the old value latched on the last read of PORTB. The "mismatch" outputs of RB<7:4> are OR'ed together to generate the RBIF interrupt (flag latched in INTCON<0>).

FIGURE 5-5: BLOCK DIAGRAM OF RB<7:4> PINS



This interrupt can wake the device from SLEEP. The user, in the interrupt service routine, can clear the interrupt in the following manner:

- Any read or write of PORTB. This will end the mismatch condition.
- Clear flag bit RBIF.

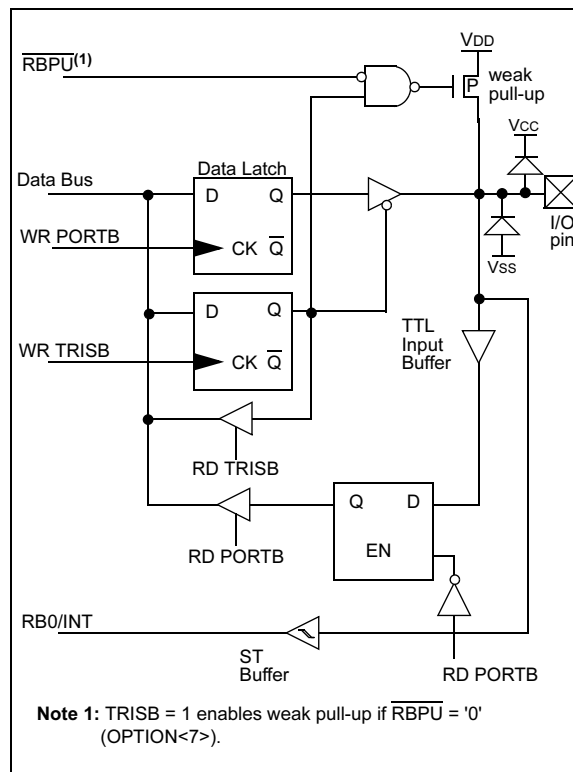
A mismatch condition will continue to set flag bit RBIF. Reading PORTB will end the mismatch condition and allow flag bit RBIF to be cleared.

This interrupt on mismatch feature, together with software configurable pull-ups on these four pins allow easy interface to a key pad and make it possible for wake-up on key-depression. (See AN552, "Implementing Wake-Up on Key Strokes.")

Note: If a change on the I/O pin should occur when the read operation is being executed (start of the Q2 cycle), then the RBIF interrupt flag may not get set.

The interrupt-on-change feature is recommended for wake-up on key depression operation and operations where PORTB is only used for the interrupt on change feature. Polling of PORTB is not recommended while using the interrupt-on-change feature.

FIGURE 5-6: BLOCK DIAGRAM OF RB<3:0> PINS



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7.4 Comparator Response Time

Response time is the minimum time, after selecting a new reference voltage or input source, before the comparator output has a valid level. If the internal reference is changed, the maximum delay of the internal voltage reference must be considered when using the comparator outputs. Otherwise the maximum delay of the comparators should be used (Table 12-2).

7.5 Comparator Outputs

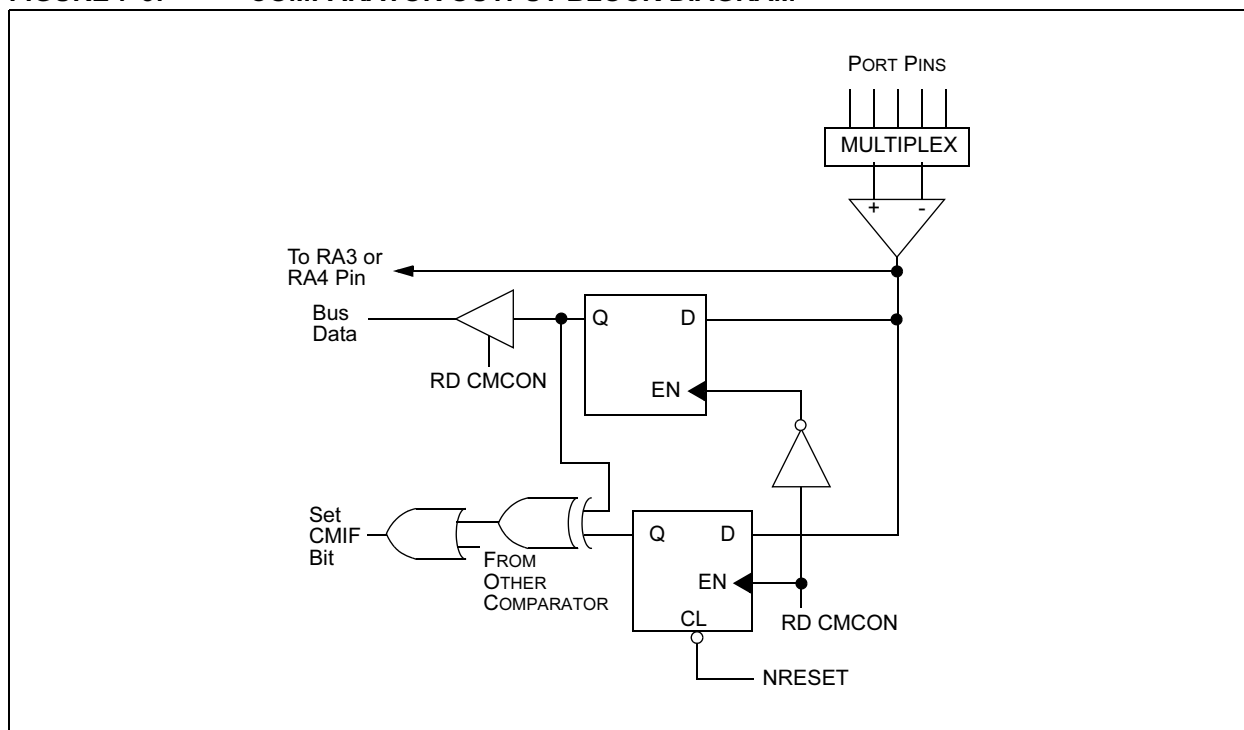
The comparator outputs are read through the CMCON register. These bits are read only. The comparator outputs may also be directly output to the RA3 and RA4 I/O pins. When the CM<2:0> = 110, multiplexors in the output path of the RA3 and RA4 pins will switch and the output of each pin will be the unsynchronized output of the comparator. The uncertainty of each of the comparators is related to the input offset voltage and the response time given in the specifications. Figure 7-3 shows the comparator output block diagram.

The TRISA bits will still function as an output enable/disable for the RA3 and RA4 pins while in this mode.

Note 1: When reading the PORT register, all pins configured as analog inputs will read as a '0'. Pins configured as digital inputs will convert an analog input according to the Schmitt Trigger input specification.

2: Analog levels on any pin that is defined as a digital input may cause the input buffer to consume more current than is specified.

FIGURE 7-3: COMPARATOR OUTPUT BLOCK DIAGRAM



7.6 Comparator Interrupts

The comparator interrupt flag is set whenever there is a change in the output value of either comparator. Software will need to maintain information about the status of the output bits, as read from CMCON<7:6>, to determine the actual change that has occurred. The CMIF bit, PIR1<6>, is the comparator interrupt flag. The CMIF bit must be RESET by clearing '0'. Since it is also possible to write a '1' to this register, a simulated interrupt may be initiated.

The CMIE bit (PIE1<6>) and the PEIE bit (INTCON<6>) must be set to enable the interrupt. In addition, the GIE bit must also be set. If any of these bits are clear, the interrupt is not enabled, though the CMIF bit will still be set if an interrupt condition occurs.

Note: If a change in the CMCON register (C1OUT or C2OUT) should occur when a read operation is being executed (start of the Q2 cycle), then the CMIF (PIR1<6>) interrupt flag may not get set.

The user, in the interrupt service routine, can clear the interrupt in the following manner:

- Any read or write of CMCON. This will end the mismatch condition.
- Clear flag bit CMIF.

A mismatch condition will continue to set flag bit CMIF. Reading CMCON will end the mismatch condition and allow flag bit CMIF to be cleared.

7.7 Comparator Operation During SLEEP

When a comparator is active and the device is placed in SLEEP mode, the comparator remains active and the interrupt is functional if enabled. This interrupt will

wake up the device from SLEEP mode when enabled. While the comparator is powered-up, higher SLEEP currents than shown in the power-down current specification will occur. Each comparator that is operational will consume additional current as shown in the comparator specifications. To minimize power consumption while in SLEEP mode, turn off the comparators, CM<2:0> = 111, before entering SLEEP. If the device wakes up from SLEEP, the contents of the CMCON register are not affected.

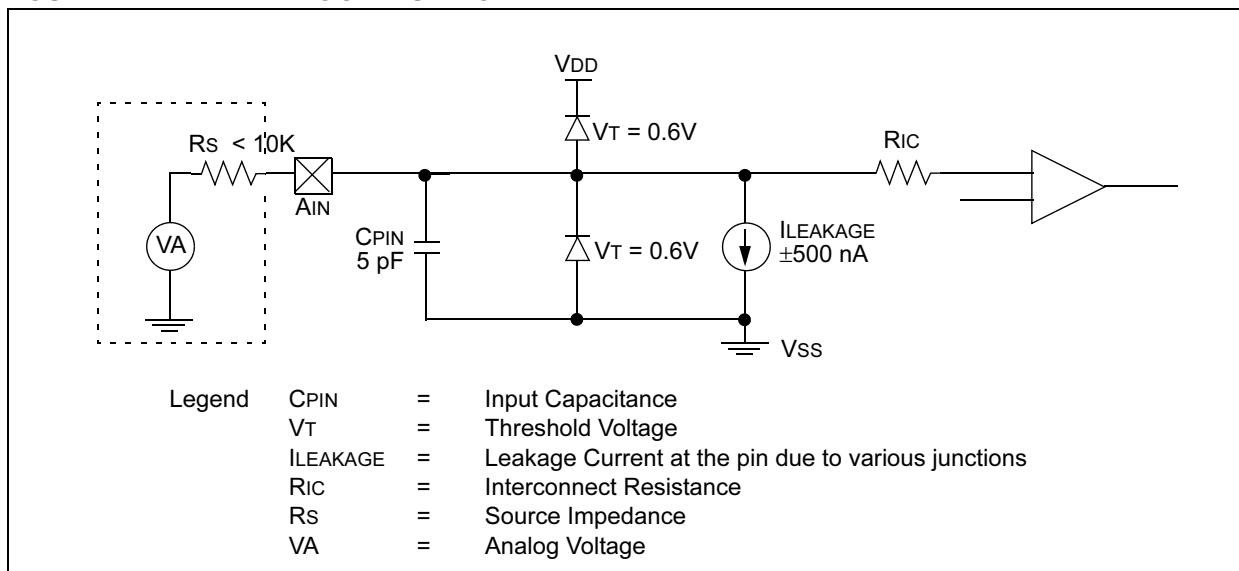
7.8 Effects of a RESET

A device RESET forces the CMCON register to its RESET state. This forces the comparator module to be in the comparator RESET mode, CM<2:0> = 000. This ensures that all potential inputs are analog inputs. Device current is minimized when analog inputs are present at RESET time. The comparators will be powered-down during the RESET interval.

7.9 Analog Input Connection Considerations

A simplified circuit for an analog input is shown in Figure 7-4. Since the analog pins are connected to a digital output, they have reverse biased diodes to VDD and VSS. The analog input therefore, must be between VSS and VDD. If the input voltage deviates from this range by more than 0.6V in either direction, one of the diodes is forward biased and a latchup may occur. A maximum source impedance of 10 kΩ is recommended for the analog sources. Any external component connected to an analog input pin, such as a capacitor or a Zener diode, should have very little leakage current.

FIGURE 7-4: ANALOG INPUT MODEL



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FIGURE 9-11: EXTERNAL POWER-ON RESET CIRCUIT (FOR SLOW V_{DD} POWER-UP)

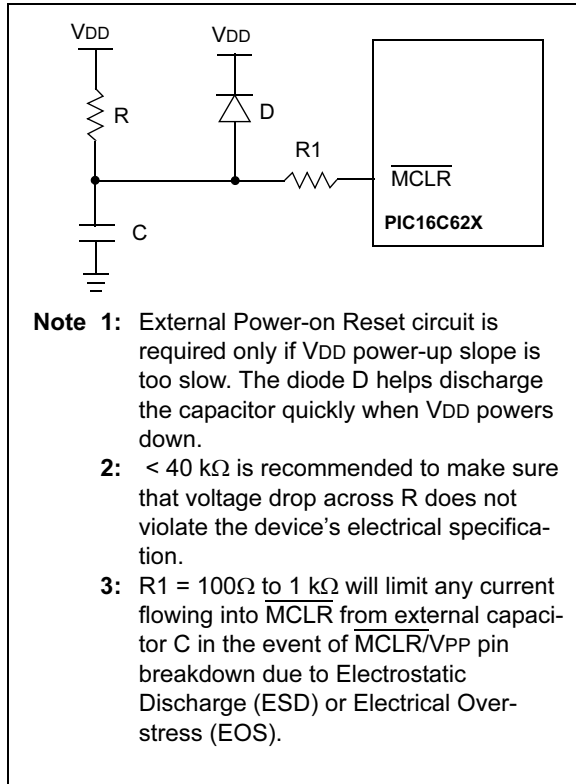


FIGURE 9-12: EXTERNAL BROWN-OUT PROTECTION CIRCUIT 1

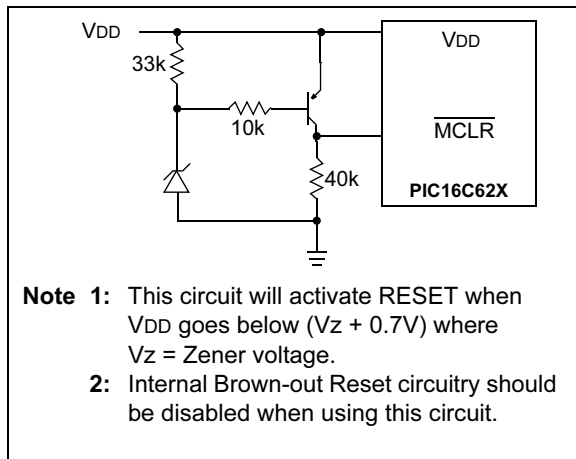


FIGURE 9-13: EXTERNAL BROWN-OUT PROTECTION CIRCUIT 2

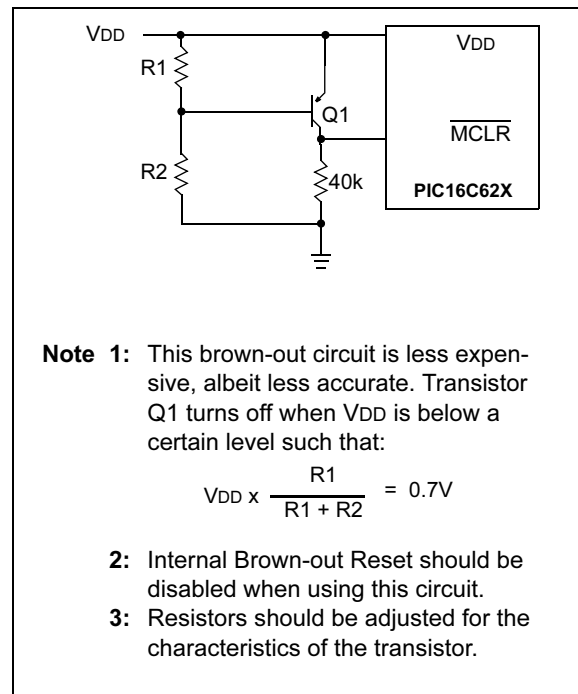
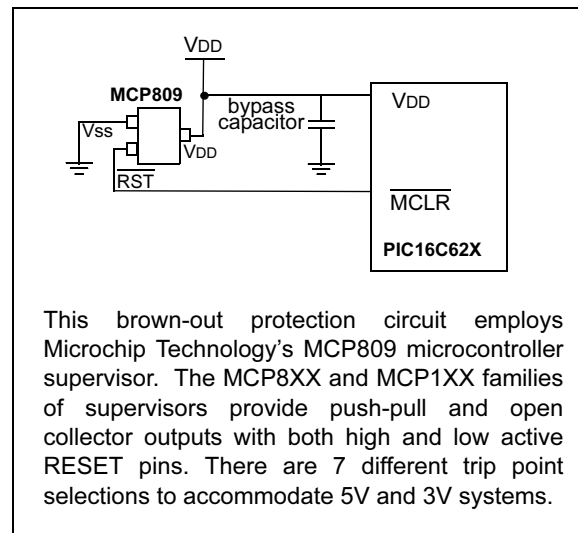


FIGURE 9-14: EXTERNAL BROWN-OUT PROTECTION CIRCUIT 3



9.5 Interrupts

The PIC16C62X has 4 sources of interrupt:

- External interrupt RB0/INT
- TMR0 overflow interrupt
- PORTB change interrupts (pins RB<7:4>)
- Comparator interrupt

The interrupt control register (INTCON) records individual interrupt requests in flag bits. It also has individual and global interrupt enable bits.

A global interrupt enable bit, GIE (INTCON<7>) enables (if set) all un-masked interrupts or disables (if cleared) all interrupts. Individual interrupts can be disabled through their corresponding enable bits in INTCON register. GIE is cleared on RESET.

The “return from interrupt” instruction, RETFIE, exits interrupt routine, as well as sets the GIE bit, which re-enable RB0/INT interrupts.

The INT pin interrupt, the RB port change interrupt and the TMR0 overflow interrupt flags are contained in the INTCON register.

The peripheral interrupt flag is contained in the special register PIR1. The corresponding interrupt enable bit is contained in special registers PIE1.

When an interrupt is responded to, the GIE is cleared to disable any further interrupt, the return address is pushed into the stack and the PC is loaded with 0004h.

Once in the interrupt service routine, the source(s) of the interrupt can be determined by polling the interrupt flag bits. The interrupt flag bit(s) must be cleared in software before re-enabling interrupts to avoid RB0/INT recursive interrupts.

For external interrupt events, such as the INT pin or PORTB change interrupt, the interrupt latency will be three or four instruction cycles. The exact latency depends when the interrupt event occurs (Figure 9-16). The latency is the same for one or two cycle instructions. Once in the interrupt service routine, the source(s) of the interrupt can be determined by polling the interrupt flag bits. The interrupt flag bit(s) must be cleared in software before re-enabling interrupts to avoid multiple interrupt requests.

Note 1: Individual interrupt flag bits are set regardless of the status of their corresponding mask bit or the GIE bit.

2: When an instruction that clears the GIE bit is executed, any interrupts that were pending for execution in the next cycle are ignored. The CPU will execute a NOP in the cycle immediately following the instruction which clears the GIE bit. The interrupts which were ignored are still pending to be serviced when the GIE bit is set again.

FIGURE 9-15: INTERRUPT LOGIC

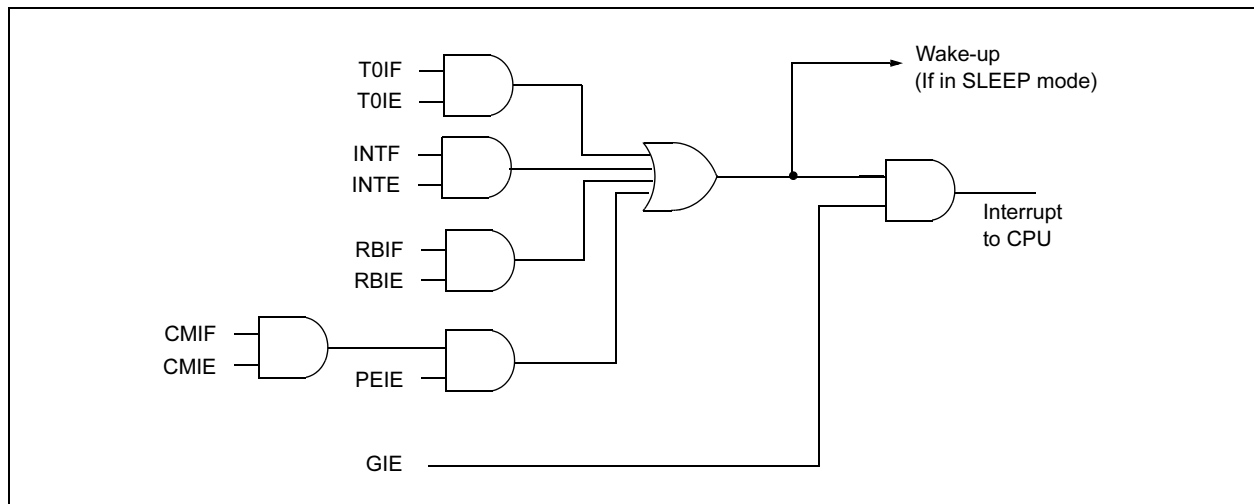


TABLE 9-6: SUMMARY OF INTERRUPT REGISTERS

Address	Name	Bit 7	Bit 6	Bit 5	Bit 4	Bit 3	Bit 2	Bit 1	Bit 0	Value on POR Reset	Value on all other RESETS ⁽¹⁾
0Bh	INTCON	GIE	PEIE	TOIE	INTE	RBIE	TOIF	INTF	RBIF	0000 000x	0000 000u
0Ch	PIR1	—	CMIF	—	—	—	—	—	—	-0-- ----	-0-- ----
8Ch	PIE1	—	CMIE	—	—	—	—	—	—	-0-- ----	-0-- ----

Note 1: Other (non Power-up) Resets include MCLR Reset, Brown-out Reset and Watchdog Timer Reset during normal operation.

9.6 Context Saving During Interrupts

During an interrupt, only the return PC value is saved on the stack. Typically, users may wish to save key registers during an interrupt (e.g., W register and STATUS register). This will have to be implemented in software.

Example 9-3 stores and restores the STATUS and W registers. The user register, W_TEMP, must be defined in both banks and must be defined at the same offset from the bank base address (i.e., W_TEMP is defined at 0x20 in Bank 0 and it must also be defined at 0xA0 in Bank 1). The user register, STATUS_TEMP, must be defined in Bank 0. The Example 9-3:

- Stores the W register
- Stores the STATUS register in Bank 0
- Executes the ISR code
- Restores the STATUS (and bank select bit register)
- Restores the W register

EXAMPLE 9-3: SAVING THE STATUS AND W REGISTERS IN RAM

```

MOVWF    W_TEMP        ;copy W to temp register,
                        ;could be in either bank
SWAPF    STATUS,W       ;swap status to be saved
                        ;into W
BCF       STATUS,RP0    ;change to bank 0 regardless
                        ;of current bank
MOVWF    STATUS_TEMP    ;save status to bank 0
                        ;register
:
:      (ISR)
:
SWAPF    STATUS_TEMP,   ;swap STATUS_TEMP register
W         W             ;into W, sets bank to origi-
                        ;nal
                        ;state
MOVWF    STATUS         ;move W into STATUS register
SWAPF    W_TEMP,F       ;swap W_TEMP
SWAPF    W_TEMP,W       ;swap W_TEMP into W

```

PIC16C62X

BCF		Bit Clear f						
Syntax:	[<i>label</i>] BCF f,b							
Operands:	$0 \leq f \leq 127$ $0 \leq b \leq 7$							
Operation:	$0 \rightarrow (f)$							
Status Affected:	None							
Encoding:	<table border="1"><tr><td>01</td><td>00bb</td><td>bfff</td><td>ffff</td></tr></table>				01	00bb	bfff	ffff
01	00bb	bfff	ffff					
Description:	Bit 'b' in register 'f' is cleared.							
Words:	1							
Cycles:	1							
Example	<div>BCF FLAG_REG, 7</div> <div>Before Instruction</div> <div>FLAG_REG = 0xC7</div> <div>After Instruction</div> <div>FLAG_REG = 0x47</div>							

BSF		Bit Set f			
Syntax:	[<i>label</i>] BSF f,b				
Operands:	$0 \leq f \leq 127$ $0 \leq b \leq 7$				
Operation:	$1 \rightarrow (f)$				
Status Affected:	None				
Encoding:	01	01bb	bfff	ffff	
Description:	Bit 'b' in register 'f' is set.				
Words:	1				
Cycles:	1				
Example	BSF FLAG_REG, 7				
	Before Instruction				
	FLAG_REG = 0x0A				
	After Instruction				
	FLAG_REG = 0x8A				

BTFSC		Bit Test, Skip if Clear							
Syntax:	[<i>label</i>] BTFSC f,b								
Operands:	$0 \leq f \leq 127$ $0 \leq b \leq 7$								
Operation:	skip if (f<b) = 0								
Status Affected:	None								
Encoding:	<table><tr><td>01</td><td>10bb</td><td>bfff</td><td>ffff</td></tr></table>					01	10bb	bfff	ffff
01	10bb	bfff	ffff						
Description:	<p>If bit 'b' in register 'f' is '0', then the next instruction is skipped.</p> <p>If bit 'b' is '0', then the next instruction fetched during the current instruction execution is discarded, and a NOP is executed instead, making this a two-cycle instruction.</p>								
Words:	1								
Cycles:	1(2)								
Example	HERE FALSE TRUE	BTFSC GOTO • • •	FLAG,1 PROCESS_CO DE						

Before Instruction
PC = address HERE

After Instruction
if FLAG<1> = 0,
PC = address TRUE
if FLAG<1> = 1,
PC = address FALSE

PIC16C62X

CLRW Clear W

Syntax: [*label*] CLRW

Operands: None

Operation: 00h → (W)
1 → Z

Status Affected: Z

Encoding:

00	0001	0000	0011
----	------	------	------

Description: W register is cleared. Zero bit (Z) is set.

Words: 1

Cycles: 1

Example CLRW

Before Instruction

W = 0x5A

After Instruction

W = 0x00

Z = 1

CLRWDT Clear Watchdog Timer

Syntax: [*label*] CLRWDT

Operands: None

Operation: 00h → WDT
0 → WDT prescaler,
1 → \overline{TO}
1 → \overline{PD}

Status Affected: \overline{TO} , \overline{PD}

Encoding:

00	0000	0110	0100
----	------	------	------

Description: CLRWDT instruction resets the Watchdog Timer. It also resets the prescaler of the WDT. STATUS bits \overline{TO} and \overline{PD} are set.

Words: 1

Cycles: 1

Example CLRWDT

Before Instruction

WDT counter = ?

After Instruction

WDT counter = 0x00

WDT prescaler = 0

\overline{TO} = 1

\overline{PD} = 1

COMF Complement f

Syntax: [*label*] COMF f,d

Operands: $0 \leq f \leq 127$
 $d \in [0,1]$

Operation: $(\bar{f}) \rightarrow (\text{dest})$

Status Affected: Z

Encoding:

00	1001	dfff	ffff
----	------	------	------

Description: The contents of register 'f' are complemented. If 'd' is 0, the result is stored in W. If 'd' is 1, the result is stored back in register 'f'.

Words: 1

Cycles: 1

Example COMF REG1,0

Before Instruction

REG1 = 0x13

After Instruction

REG1 = 0x13

W = 0xEC

DECF Decrement f

Syntax: [*label*] DECF f,d

Operands: $0 \leq f \leq 127$
 $d \in [0,1]$

Operation: $(f) - 1 \rightarrow (\text{dest})$

Status Affected: Z

Encoding:

00	0011	dfff	ffff
----	------	------	------

Description: Decrement register 'f'. If 'd' is 0, the result is stored in the W register. If 'd' is 1, the result is stored back in register 'f'.

Words: 1

Cycles: 1

Example DECF CNT, 1

Before Instruction

CNT = 0x01

Z = 0

After Instruction

CNT = 0x00

Z = 1

MOVF		Move f			
Syntax:	[<i>label</i>] MOVF f,d				
Operands:	$0 \leq f \leq 127$ $d \in [0,1]$				
Operation:	(f) \rightarrow (dest)				
Status Affected:	Z				
Encoding:	00	1000	dfff	ffff	
Description:	The contents of register f is moved to a destination dependent upon the status of d. If d = 0, destination is W register. If d = 1, the destination is file register f itself. d = 1 is useful to test a file register since status flag Z is affected.				
Words:	1				
Cycles:	1				
Example	MOVF FSR, 0				
	After Instruction				
	W	=	value in FSR		
	register				
	Z	=	1		

MOVWF	Move W to f				
Syntax:	[<i>label</i>] MOVWF f				
Operands:	0 ≤ f ≤ 127				
Operation:	(W) → (f)				
Status Affected:	None				
Encoding:	<table><tr><td>00</td><td>0000</td><td>1fff</td><td>ffff</td></tr></table>	00	0000	1fff	ffff
00	0000	1fff	ffff		
Description:	Move data from W register to register 'f'.				
Words:	1				
Cycles:	1				
Example	MOVWF OPTION Before Instruction OPTION = 0xFF W = 0x4F After Instruction OPTION = 0x4F W = 0x4F				

NOP	No Operation			
Syntax:	[<i>label</i>] NOP			
Operands:	None			
Operation:	No operation			
Status Affected:	None			
Encoding:	00	0000	0xx0	0000
Description:	No operation.			
Words:	1			
Cycles:	1			
Example	NOP			

OPTION	Load Option Register				
Syntax:	[<i>label</i>] OPTION				
Operands:	None				
Operation:	(W) → OPTION				
Status Affected:	None				
Encoding:	<table><tr><td>00</td><td>0000</td><td>0110</td><td>0010</td></tr></table>	00	0000	0110	0010
00	0000	0110	0010		
Description:	<p>The contents of the W register are loaded in the OPTION register.</p> <p>This instruction is supported for code compatibility with PIC16C5X products. Since OPTION is a readable/writable register, the user can directly address it.</p>				
Words:	1				
Cycles:	1				
Example	<div>To maintain upward compatibility with future PICmicro[®] products, do not use this instruction.</div>				

11.3 MPLAB C17 and MPLAB C18 C Compilers

The MPLAB C17 and MPLAB C18 Code Development Systems are complete ANSI C compilers for Microchip's PIC17CXXX and PIC18CXXX family of microcontrollers. These compilers provide powerful integration capabilities, superior code optimization and ease of use not found with other compilers.

For easy source level debugging, the compilers provide symbol information that is optimized to the MPLAB IDE debugger.

11.4 MPLINK Object Linker/ MPLIB Object Librarian

The MPLINK object linker combines relocatable objects created by the MPASM assembler and the MPLAB C17 and MPLAB C18 C compilers. It can link relocatable objects from pre-compiled libraries, using directives from a linker script.

The MPLIB object librarian manages the creation and modification of library files of pre-compiled code. When a routine from a library is called from a source file, only the modules that contain that routine will be linked in with the application. This allows large libraries to be used efficiently in many different applications.

The object linker/librarian features include:

- Efficient linking of single libraries instead of many smaller files
- Enhanced code maintainability by grouping related modules together
- Flexible creation of libraries with easy module listing, replacement, deletion and extraction

11.5 MPLAB C30 C Compiler

The MPLAB C30 C compiler is a full-featured, ANSI compliant, optimizing compiler that translates standard ANSI C programs into dsPIC30F assembly language source. The compiler also supports many command-line options and language extensions to take full advantage of the dsPIC30F device hardware capabilities, and afford fine control of the compiler code generator.

MPLAB C30 is distributed with a complete ANSI C standard library. All library functions have been validated and conform to the ANSI C library standard. The library includes functions for string manipulation, dynamic memory allocation, data conversion, time-keeping, and math functions (trigonometric, exponential and hyperbolic). The compiler provides symbolic information for high level source debugging with the MPLAB IDE.

11.6 MPLAB ASM30 Assembler, Linker, and Librarian

MPLAB ASM30 assembler produces relocatable machine code from symbolic assembly language for dsPIC30F devices. MPLAB C30 compiler uses the assembler to produce its object file. The assembler generates relocatable object files that can then be archived or linked with other relocatable object files and archives to create an executable file. Notable features of the assembler include:

- Support for the entire dsPIC30F instruction set
- Support for fixed-point and floating-point data
- Command line interface
- Rich directive set
- Flexible macro language
- MPLAB IDE compatibility

11.7 MPLAB SIM Software Simulator

The MPLAB SIM software simulator allows code development in a PC hosted environment by simulating the PICmicro series microcontrollers on an instruction level. On any given instruction, the data areas can be examined or modified and stimuli can be applied from a file, or user defined key press, to any pin. The execution can be performed in Single-Step, Execute Until Break, or Trace mode.

The MPLAB SIM simulator fully supports symbolic debugging using the MPLAB C17 and MPLAB C18 C Compilers, as well as the MPASM assembler. The software simulator offers the flexibility to develop and debug code outside of the laboratory environment, making it an excellent, economical software development tool.

11.8 MPLAB SIM30 Software Simulator

The MPLAB SIM30 software simulator allows code development in a PC hosted environment by simulating the dsPIC30F series microcontrollers on an instruction level. On any given instruction, the data areas can be examined or modified and stimuli can be applied from a file, or user defined key press, to any of the pins.

The MPLAB SIM30 simulator fully supports symbolic debugging using the MPLAB C30 C Compiler and MPLAB ASM30 assembler. The simulator runs in either a Command Line mode for automated tasks, or from MPLAB IDE. This high speed simulator is designed to debug, analyze and optimize time intensive DSP routines.

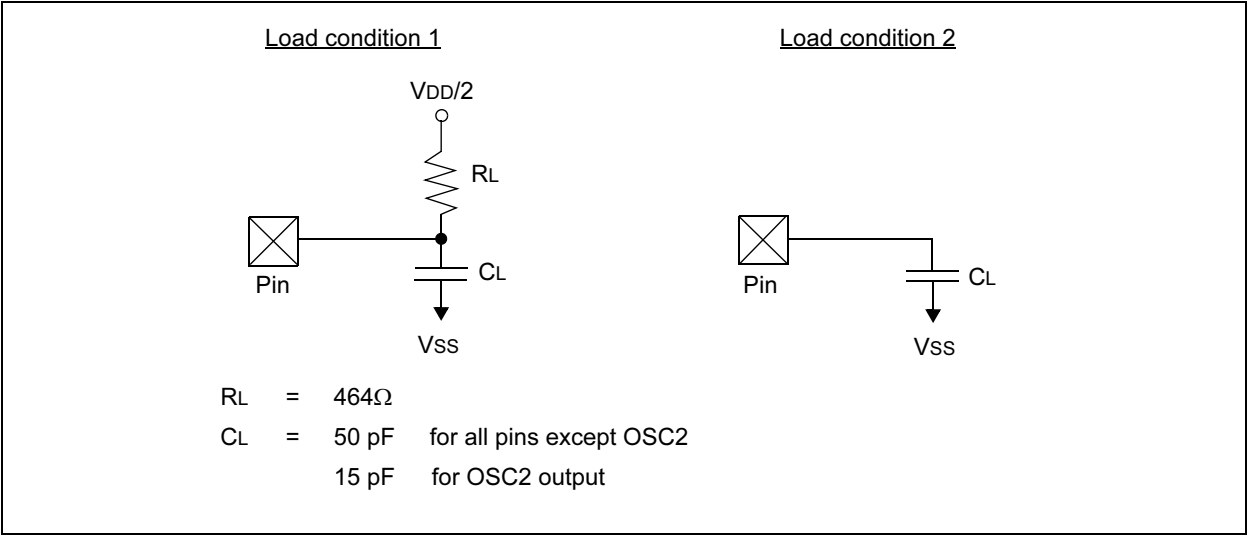
12.8 Timing Parameter Symbolology

The timing parameter symbols have been created with one of the following formats:

- 1. TppS2ppS
- 2. TppS

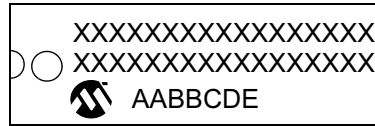
T			
F	Frequency	T	Time
Lowercase subscripts (pp) and their meanings:			
pp			
ck	CLKOUT	osc	OSC1
io	I/O port	t0	T0CKI
mc	MCLR		
Uppercase letters and their meanings:			
S			
F	Fall	P	Period
H	High	R	Rise
I	Invalid (Hi-impedance)	V	Valid
L	Low	Z	Hi-Impedance

FIGURE 12-11: LOAD CONDITIONS

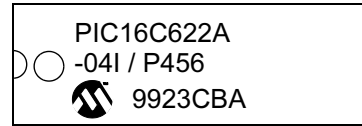


14.1 Package Marking Information

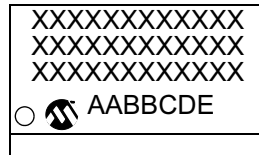
18-Lead PDIP



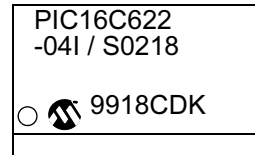
Example



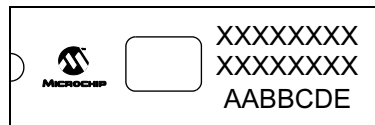
18-Lead SOIC (.300")



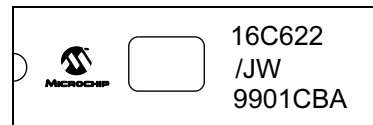
Example



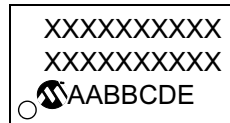
18-Lead Cerdip Windowed



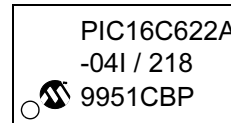
Example



20-Lead SSOP



Example



Legend: XX...X Customer specific information*
 Y Year code (last digit of calendar year)
 YY Year code (last 2 digits of calendar year)
 WW Week code (week of January 1 is week '01')
 NNN Alphanumeric traceability code

Note: In the event the full Microchip part number cannot be marked on one line, it will be carried over to the next line thus limiting the number of available characters for customer specific information.

* Standard PICmicro device marking consists of Microchip part number, year code, week code, and traceability code. For PICmicro device marking beyond this, certain price adders apply. Please check with your Microchip Sales Office. For QTP devices, any special marking adders are included in QTP price.

READER RESPONSE

Please list the following information, and use this outline to provide us with your comments about this document.

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Would you like a reply? Y N

Device: PIC16C62X Literature Number: DS30235J

Questions:

1. What are the best features of this document?

2. How does this document meet your hardware and software development needs?

3. Do you find the organization of this document easy to follow? If not, why?

4. What additions to the document do you think would enhance the structure and subject?

5. What deletions from the document could be made without affecting the overall usefulness?

6. Is there any incorrect or misleading information (what and where)?

7. How would you improve this document?
